

CSE12 - Lecture 14

Friday, May 3, 2024 10:00 AM

Lecture 14

Quicksort: Another magical (recursive) algorithm

<https://www.youtube.com/watch?v=ywWBy6J5gz8>

14	4	9	12	15	8	19	2
----	---	---	----	----	---	----	---

Select a **pivot** element:

14	4	9	12	15	8	19	2
----	---	---	----	----	---	----	---

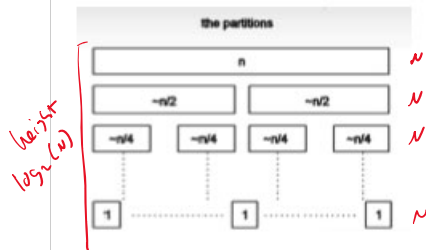
"Partition" the elements in the array (smaller or equal to pivot, larger or equal to pivot)

2	4	9	8	15	12	19	14
---	---	---	---	----	----	----	----

Magically sort the smaller elements and the larger elements (Quicksort)

2	4	8	9	12	15	19	21
---	---	---	---	----	----	----	----

Quick Sort: Using a "good" pivot
 < pivot value sorted >= pivot value median value
 How many levels will there be if you choose a pivot that divides the list in half?



- A. 1
- ☒ B. $\log(n)$
- C. n
- D. $n \cdot \log(n)$
- E. n^2

If the time to partition on each level takes N comparisons, how long does Quicksort take with a good partition?

- A. $O(1)$
- B. $O(\log(n))$
- C. $O(n)$
- ☒ D. $O(n \cdot \log(n))$
- E. $O(n^2)$

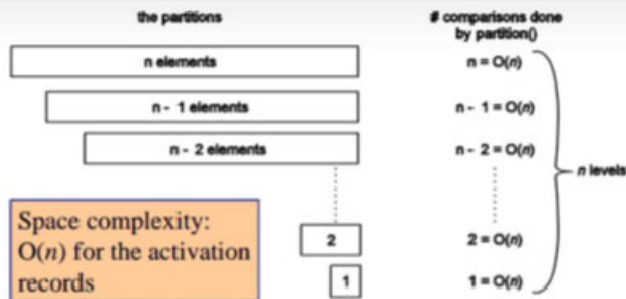
Which of these choices would be the worst choice for the pivot?

- ☒ A. The minimum element in the list
- B. The last element in the list
- C. The first element in the list
- ☒ D. A random element in the list

best case \rightarrow pivot = median value

Name: _____ PID: _____ Code: **2013**

Quick sort with a bad pivot



If the pivot always produces one empty partition and one with $n - 1$ elements, there will be n levels, each of which requires $O(n)$ comparisons: $O(n^2)$ time complexity

Which of these choices is a better choice for the pivot?

- A. The first element in the list
 B. A random element in the list → every partition randomly pick pivot value
 C. They are about the same

Shortcut → swap to visible
 [12, 4, 9, 3, 2, 8, 14, 15]
 pick the best partition

There are many ways to partition!

Quick sort - Middle Pivot

- We always pick the middle location as pivot
- The data we sort is [2, 3, 1, 5, 4, 6, 7]

After the first split, what is the order of elements in the list that was \leq pivot?

- A. 1 2 3 4
 B. 2 3 1 4
 C. 4 3 2 1
 D. 3 4 1 2
 E. None of the above

low 0 1 2 3 4 5 6 high
 sort {12, 4, 9, 3, 15, 8, 19, 2}

What is the first pivot? 15, index 4

How would you sort the data around the pivot?

lowIndex = 0

highIndex = length - 1 → 8 - 1 = 7

pivotIndex = (high - low) / 2 = (8 - 0) / 2 = 4

pivotValue = value[pivotIndex]; // 15

Return this at end

update this if it moves

check of left & middle

check of right & middle

if value[lowIndex] < pivotValue
 lowIndex++

else
 swap(lowIndex, highIndex)
 highIndex--

if value[highIndex] >= pivotValue
 highIndex--

else
 swap(lowIndex, highIndex)
 lowIndex++

exit condition [if (lowIndex > highIndex)
 return / break
 done = true;

```
import java.util.Arrays;
public class Sort {
    static void selectionSort(int[] arr) {
        for(int i = 0; i < arr.length; i++) {
            int minIndex = i;
            for(int j = i; j < arr.length; j++) {
                if(arr[minIndex] > arr[j]) { minIndex = j; }
            }
            int temp = arr[i];
            arr[i] = arr[minIndex];
            arr[minIndex] = temp;
        }
    }
    static void insertionSort(int[] arr) {
        for(int i = 0; i < arr.length; i++) {
            for(int j = i; j > 0; j--) {
                if(arr[j] < arr[j-1]) {
                    int temp = arr[j-1];
                    arr[j-1] = arr[j];
                    arr[j] = temp;
                }
            }
        }
    }
}
```

sorted array
 [1, 2, 3, 4, 5]

```
import java.util.Arrays;
public class SortFaster {
    static int[] combine(int[] p1, int[] p2) {...}
    static int[] mergeSort(int[] arr) {
        int len = arr.length;
        if(len <= 1) { return arr; }
        else {
            int[] p1 = Arrays.copyOfRange(arr, 0, len / 2);
            int[] p2 = Arrays.copyOfRange(arr, len / 2, len);
            int[] sortedPart1 = mergeSort(p1);
            int[] sortedPart2 = mergeSort(p2);
            int[] sorted = combine(sortedPart1, sortedPart2);
            return sorted;
        }
    }
}
```

```

int[] sortedPart2 = MergeSort(p2);
int[] sorted = combine(sortedPart1, sortedPart2);
return sorted;
}
}

static int partition(String[] array, int l, int h) {...}
static void qsort(String[] array, int low, int high) {
    if (high - low <= 1) { return; }
    int splitAt = partition(array, low, high);
    qsort(array, low, splitAt);
    qsort(array, splitAt + 1, high);
}
public static void sort(String[] array) {
    qsort(array, 0, array.length);
}
}

```

	Insertion	Selection	Merge	Quick
Best case time	Sorted array $O(N)$	$O(N^2)$	$O(N \log_2(N))$	Median value $O(N \log_2(N))$
Worst case time	Reverse sorted array $O(N^2)$	$O(N^2)$	$O(N \log_2(N))$	$O(N^2)$ Average case: $O(N \log_2(N))$
Key operations	swap(a, j, j-1) (until in the right place)	swap(a, i, indexOfMin) (after finding minimum value)	l = copy(a, 0, len/2) r = copy(a, len/2, len) ls = sort(l) rs = sort(r) merge(ls, rs)	p = partition(a, l, h) sort(a, l, p) sort(a, p+1, h)

Last note about sorting

Not only do we care about runtime, we also care about

- Space: do we need extra storage?
- Stable: if we have duplicates, do we maintain the same ordering?

Algorithm	Space	Stable
Bubble sort	$O(1)$	Yes
Selection sort	$O(1)$	No
Insertion sort	$O(1)$	Yes
Heap sort	$O(1)$	No
Merge sort	$O(n)$	Yes
Quick sort	$O(\log n)$	No

{ "I", "Greg" }

{ "I", "Jordan" }

["Greg", "Jordan"]

["Jordan", "Greg"]

Java Array List Sorting

↳ Quick sort → primitives

↳ Merge sort → objects