0/23 Questions Answered

Exam 2

STUDENT NAME

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Q1 Excel with Integrity Pledge

0 Points

I will complete this exam in a fair, honest, respectful, responsible and trustworthy manner. This means that I will complete the exam as if the professor was watching my every action. I will act according to the professor's instructions, and I will neither give nor receive any aid or assistance other than what is authorized. I know that the integrity of this exam and this class is up to me, and I pledge to not take any action that would break the trust of my classmates or professor, or undermine the fairness of this class.

I promise to complete this exam in keeping with the Excel with Integrity Pledge.

Fill in your Name and today's Date below:

Enter your answer here

Save Answer

Q2 Instructions

0 Points

You have 90 minutes to complete the exam. Work to maximize points. If you get stuck, work through other problems and come back to it.

In general, if you think you've spotted a typo in the exam, do your best to answer in the spirit of the question. Keep in mind that some questions have interesting code examples with intentional bugs for you to find as part of the question. Questions asked during the exam will not be answered, so do your best to interpret the questions.

In general, assume that any necessary libraries (JUnit, ArrayList, List, and so on) have been imported.

You can use your notes and a compiler. However, the test is designed as if you were taking it during lecture without a compiler and it's highly suggested you not rely on your compiler and use pen & paper to figure out your answers.

Stay calm – you can do this!

Save Answer

Q3 Partition

6 Points

Consider this specification for partition:

int partition(int[] numbers)

Partition() method description:

Chooses a pivot value from the array. Then, changes the array so that all elements smaller than or equal to the pivot appear in indices less than the index of the pivot value, and all elements larger than the pivot appear in indices greater than the index of the pivot value. Returns the final index of the pivot value, which may be different than its starting index. All elements in the input should appear at some index in the output.

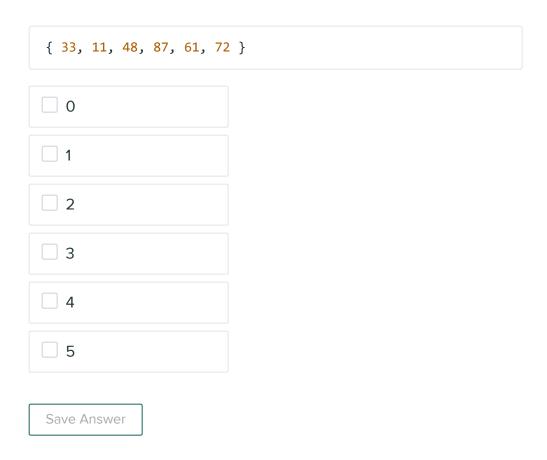
Q3.1 A

2 Points

Consider this *input* array:

```
{ 61, 33, 11, 87, 48, 72 }
```

Consider the following array *after* a call to **partition()**, for the input array above. What are **all** the indices that could have been **returned** from **partition()** for this to be a valid partition result? Select **all** possible indices.



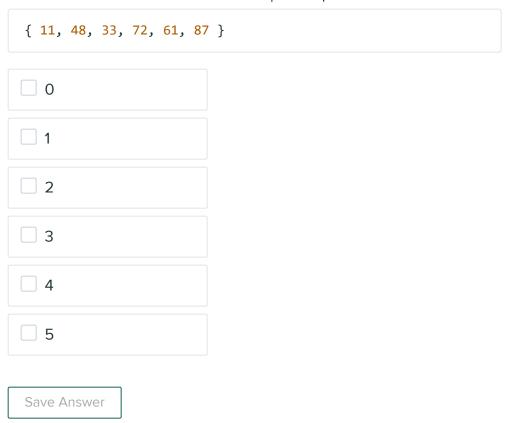
Q3.2 B

2 Points

Consider this *input* array:

```
{ 61, 33, 11, 87, 48, 72 }
```

Consider the following array *after* a call to **partition()**, for the input array above. What are **all** the indices that could have been **returned** from **partition()** for this to be a valid partition result? Select **all** possible indices.



Q3.3 C

2 Points

Consider this *input* array:

```
{ 61, 33, 11, 87, 48, 72 }
```

Consider the following array *after* a call to **partition()**, for the input array above. What are **all** the indices that could have been **returned** from **partition()** for this to be a valid partition result? Select **all** possible indices.

```
{ 33, 11, 48, 61, 87, 72 }
```

_ o	
<u> </u>	
_ 2	
_ 3	
_ 4	
<u> </u>	
Save Answer	

Q4 Run-time

7 Points

Recall that we say f is O(g) if there exist nO and C such that for all n > 1n0, f(n) < C * g(n)

Recall that Θ (big-Theta) represents a *tight* bound, O (big-O) an *upper* bound, and Ω (big-Omega) a *lower* bound. (We intentionally do not give definitions for Θ and Ω).

Q4.11

1 Point

Give a Θ bound for the number of steps the following program takes in terms of n, assuming doSomeWork() does a constant amount of work.

```
for(int i = 0; i <= n * n; i = i + 2) {
  for(int j = n; j > 0; j -= 1) {
    doSomeWork();
  }
}
```

○ Θ(1)
 ○ Θ(log(n))
 ○ Θ(n)
 ○ Θ(n*log(n))
 ○ Θ(n^2)
 ○ Θ(n^3)
 ○ Θ(2^n)
 ○ Θ(n!)

Save Answer

Q4.2 2

1 Point

Give a Θ bound for the number of steps the following program takes in terms of n, assuming **doSomeWork()** does a constant amount of work.

```
for(int i = 0; i < n / 2; i += 1) {
  for(int j = 1; j < n; j = j * 2) {
    doSomeWork();
  }
}</pre>
```

- **O** ⊝(1)
- $\bigcirc \Theta(\log(n))$
- **O** Θ(n)
- $O \Theta(n^*log(n))$
- **O** Θ(n^2)
- O Θ(n^3)
- **O** Θ(2^n)
- **O** Θ(n!)

Save Answer

Q4.3 3a

1 Point

Consider the following pair of functions. Indicate whether f is O(g)and/or f is $\Omega(g)$. Select one, both or the choice of Neither is correct as appropriate.



$$g(x) = 10 * x * log(x)$$

- fis O(g)
- Neither is correct

Save Answer

Q4.4 3b

1 Point

Consider the following pair of functions. Indicate whether f is O(g) and/or f is $\Omega(g)$. Select one, both or the choice of Neither is correct as appropriate.

$$f(x) = x^2 + x + \log(x)$$

$$g(x) = x * log(x) + x^2$$

- Neither is correct

Save Answer

Q4.5 3c

1 Point

Consider the following pair of functions. Indicate whether f is O(g) and/or f is $\Omega(g)$. Select one, both or the choice of Neither is correct as appropriate.



 \Box f is O(g)

 \Box f is $\Omega(g)$

Neither is correct

Save Answer

Q4.6 3d

1 Point

Consider the following pair of functions. Indicate whether f is O(g) and/or f is $\Omega(g)$. Select one, both or the choice of Neither is correct as appropriate.

$$f(x) = \frac{x * (x-1)}{2}$$

$$g(x) = 5 * x * log(x)$$

f is O(g)

Neither is correct

Save Answer

Q4.7 3e

1 Point

Consider the following pair of functions. Indicate whether f is O(g) and/or f is $\Omega(g)$. Select one, both or the choice of Neither is correct as

appropriate.

$$f(x) = \frac{x * (x-1)}{x}$$
 $g(x) = 3 * x * log(x)$

☐ f is O(g)

 \Box f is $\Omega(g)$

Neither is correct

Save Answer

Q5 HashTable

9 Points

Consider this implementation of a hash table based on one from the lecture notes, with some portions especially relevant for the question bolded.

interface Hasher< K > { int hash(K key); }

```
class HashTable< K, V > {
    class Entry {
        K k; V v;
        public Entry(K k, V v) { this.k = k; this.v = v; }
}
List< Entry >[] buckets; // An array of Lists of Entries
int size;
Hasher< K > hasher;

public HashTable(Hasher< K > h, int startCapacity) {
    this.size = 0;
    this.hasher = h;
    this.buckets = (List< Entry >[])(new List[startCapacity]);
}

public double loadFactor() { return (double)(this.size) /
this.buckets.length; }
```

```
public V get(K k) {
 int hashCode = this.hasher.hash(k);
 int index = hashCode % this.buckets.length;
 if(this.buckets[index] == null) {
  return null;
 }
 else {
  for(Entry e: this.buckets[index]) {
   if(e.k.equals(k)) { return e.v; }
  return null;
public void set(K k, V v) {
 int hashCode = this.hasher.hash(k);
 int index = hashCode % this.buckets.length;
 if(this.buckets[index] == null) {
  this.buckets[index] = new ArrayList< Entry >();
  this.buckets[index].add(new Entry(k, v));
 }
 else {
  for(Entry e: this.buckets[index]) {
   if(e.k.equals(k)) { e.v = v; return; }
  this.buckets[index].add(new Entry(k, v));
 }
 this.size += 1;
```

Consider these implementations of the **Hasher** interface:

}

```
class HasherA implements Hasher< String > {
  public int hash(String s) { return Character.codePointAt(s, 0); }
}
class HasherB implements Hasher< String > {
  int i;
  public HasherB() { this.i = 0; }
  public int hash(String s) { i += 1; return s.length() + i; }
}
```

Recall that Character.codePointAt retrieves the ASCII code of the character at the given index. Consider this sequence of operations:

- HashTable<String, Integer> ht = new HashTable<>(???, 5);
- ht.set("blue", 400);
- System.out.println(ht.size);
- ht.set("green", 400);
- System.out.println(ht.size);
- System.out.println(ht.get("blue"));
- ht.set("green", 300);
- System.out.println(ht.size);
- System.out.println(ht.get("green"));
- 10. ht.set("black", 200);
- 11. System.out.println(ht.size);

Consider using each of the two Hasher implementations to replace the ??? in line 1 above by filling it in with either **new HasherA()** or **new HasherB()**.

ASCII Table:

a	97	f	102
b	98	50	103
c	99	h	104
d	100	i	105
e	101	j	106

l 108 m 109 n 110	k	107
<u> </u>	1	108
n 110	m	109
	n	110
0 111	О	111

p	112
q	113
r	114
S	115
t	116

u	117
V	118
W	119
X	120
y	121

Z	122

Q5.1 HasherA - Output

3 Points

For HasherA, write the **output** below (6 lines of output):

Enter your answer here

Save Answer

Q5.2 HasherB - Output

3 Points

For HasherB, write the **output** below (6 lines of output):

Enter your answer here

Save Answer

Q5.3 HasherA - Length

1 Point

For HasherA, at the end of running all the statements, what would be the length of the longest List in the entries array? Write your answer as a number below:

Enter your answer here

Save Answer

Q5.4 HasherB - Length

1 Point

For HasherB, at the end of running all the statements, what would be the length of the longest List in the entries array? Write your answer as a number below:

Enter your answer here

Save Answer

Q5.5 Incorrect

1 Point

Which of the two hash implementations (**HasherA** or **HasherB**) produces *incorrect* results? Identify which one does, and **give a one-sentence answer why**:

Enter your answer here

Save Answer

Q6 Merge Sort

2 Points

Suppose we have a modified version of the merge sort called **threeMergeSort()**, and instead of cutting in the middle, it always cuts at the 1/3 location of the list (i.e. 1/3 of the list is in the front, and 2/3 are in the back). The **threeMergeSort()** method is given as follows and you can assume that **merge()** works well for this code.

```
public static void threeMergeSort(int[] in){
  if (in.length > 2){
    int mid = in.length/3;
    //left inclusive, right exclusive when using copyOfRange
    int[] front = Arrays.copyOfRange(in, 0, mid);
    int[] back = Arrays.copyOfRange(in, mid, in.length);
    threeMergeSort(front);
    threeMergeSort(back);
    merge(front, back, in);
  }
  else{
    if(in.length == 1){
      return;
    if(in.length == 2){
      Arrays.sort(in);
  }
}
```

Answer the following two questions about threeMergeSort()

Q6.1

1 Point

Please answer the following question about **threeMergeSort()** if we try to sort the following array

 $\{9, 4, 5, 2, 3, 6, 1\}$

What is the correct collection of lowest-level sub-arrays that have reached base cases?

- **O** {9}, {4, 5}, {2}, {3}, {6, 1}
- **O** {9}, {4}, {5}, {2}, {3}, {6}, {1}
- **O** {9, 4}, {5}, {2}, {3}, {6, 1}
- **O** {9}, {4, 5}, {2, 3}, {6, 1}

Save Answer

Q6.2

1 Point

Is it true that this **threeMergeSort()** has the same efficiency as the normal merge sort asymptotically?

- O true
- O false

Save Answer

Q7

4 Points

Consider this helper method:

```
public static void dosth(int loc, int[] arr) {
    int j = loc;
    while (j > 0 && arr[j] < arr[j-1]){
        int temp = arr[j];
        arr[j] = arr[j-1];
        arr[j-1] = temp;
        j--;
    }
}</pre>
```

Q7.1

1 Point

Consider this method skeleton:

```
void isort(int[] arr) {
  for(int i = 0; i < arr.length; i += 1) {
  }
}</pre>
```

Fill in a **single call to the helper method dosth()** that would complete the method and cause it to always change the input array to be in **increasing sorted order**. Write the body of the loop (a single call to dosth):

Enter your answer here

Save Answer

Q7.2

1 Point

Give a Θ bound (a tight bound) for the number of steps the **dosth()** method takes in terms of n in the **worst case** arrangement of the elements in arr, and assuming arr.length is at least as large as n.

O ⊝(1)
O ⊖(log(n))
O ⊖(n)
O ⊖(n*log(n))
O ⊝(n^2)
O ⊖(n^3)
O ⊖(2^n)
O ⊖(n!)
Save Answer
Q7.3 1 Point
Give a Θ bound (a tight bound) for the number of steps the dosth() method takes in terms of n in the best case arrangement of the elements in arr, and assuming arr.length is at least as large as n.
O ⊝(1)
O ⊖(log(n))
O ⊖(n)
O ⊖(n*log(n))
O Θ(n^2)
O ⊖(n^3)
O ⊖(2^n)
O ⊖(n!)
Save Answer
074

1 Point

Give a Θ bound (a tight bound) for the number of steps the completed isort() method you wrote takes in terms of arr.length in the worst case arrangement of the elements in arr, for input arrays of arbitrary size.
○ ⊖(1)
O ⊖(log(n))
O ⊖(n)
O ⊝(n*log(n))
O Θ(n^2)
O Θ(n^3)
O Θ(2^n)
O ⊖(n!)
Save Answer
Save All Answers Submit & View Submission >