

```

enum Type { Num, Bool, Any, Nothing }
struct TyEnv {

}

fn calc_type(e : &Expr<()>, type_env: &TyEnv) ->
    match e {
        Expr::Num(_) =>
            {
                 $\Gamma \langle \text{number} \rangle : \text{Num}$ 
            }

        Expr::True =>
            {
                 $\Gamma \text{ true} : \text{Bool}$ 
            }

        Expr::Add(e1, e2) =>
            {
                 $\Gamma (\text{op } e1 \text{ } e2) : \text{Num}$ 
                 $\text{when } \Gamma e1 \leq \text{Num} \text{ and } \Gamma e2 \leq \text{Num}$ 
                 $\text{and op is } +, -, *$ 
            }

        Expr::Let(x, ex, eb) =>
            {
                 $\Gamma x : T$ 
                 $\text{when } \Gamma(x) = T$ 
            }

        Expr::Id(x) =>
            {
                 $\Gamma (\text{let } (x \text{ } ex) \text{ } eb) : T$ 
                 $\text{when } \Gamma e : T1 \text{ and } \Gamma[x : T1] e : T$ 
            }

        Expr::Set(x, e) =>
            {
                 $\Gamma (\text{set! } x \text{ } e) : T$ 
                 $\text{when } e : T$ 
                 $\text{and } \Gamma(x) \leq T$ 
            }

        Expr::Call2(f, e1, e2) =>
            {
                 $\Gamma (f \text{ } e1 \text{ } e2 \dots) : T$ 
                 $\text{when (fun (f (x1 : T1) (x2 : T2) ...) -> T e)}$ 
                 $\text{and } e1 \leq T1, e2 \leq T2, \dots$ 
            }

        Expr::If(e1, e2, e3) =>
            {
                 $\Gamma (\text{if } e1 \text{ } e2 \text{ } e3) : T1 \cup T2$ 
                 $\text{when } \Gamma e2 : T1 \text{ and } \Gamma e3 : T2 \text{ and } \Gamma e1 : \text{Bool}$ 
            }

        Expr::Cast(t, e) =>
            {
                 $\Gamma (\text{cast } T \text{ } e) : T$ 
                 $\text{when } \Gamma e : T'$ 
            }
    }
}

```

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fn calc_type(e : &Expr<()>, type_env: &TyEnv) ->
...
Expr::Loop(e) =>
{
     $\Gamma$  (loop e) :  $T_1 \cup T_2 \cup \dots \cup T_n$ 
    when  $\Gamma_1 e_1 : T_1, \Gamma_2 e_2 : T_2, \dots, \Gamma_n e_n : T_n$ 
    and  $e_1, e_2, \dots, e_n$  are (break e)
        subexpressions of e not nested in another break
    and  $\Gamma_1, \Gamma_2, \dots, \Gamma_n$ 
        are the environments for the corresponding  $e_n$ 

     $\Gamma$  (break e) : Nothing
    when  $\Gamma e : T$ 

Expr::Break(e) =>

}
}

```