

```
(def (g y)
      (+ y 1))
```

fun name to call | argument value

```
(def (f x)
      (+ (g (+ x 2)) 3))
```

fun name defined

function body

```
(def (main input)
      (f (+ y 4)))
```

argument name

Late May/Early June

This week

↓

```
type expr =
| ENum of int
| EBool of bool
...
| EDef of string * string * expr
| EApp of expr * expr
```

```
type expr =
| ENum of int
| EBool of bool
...
| EApp of string * expr
```

```
type def =
| Def of string * string * expr
```

```
type prog =
| Prog of def list * expr
```

Which representation do you want to implement first?

Some things to discuss:

- Compiling the new abstract syntax (getting its answer into EAX)
- How the environment works (a new kind of name)
- Are there programs we can represent with one but not the other?

```
type expr =  
| ENum of int  
| EBool of bool  
...  
| EApp of string * expr
```

```
type def =  
| Def of string * string * expr
```

```
type prog =  
| Prog of def list * expr
```

```
let rec compile_expr e si env =  
  match e with  
  ...  
  | EApp(fname, arg) -> ...
```

```
let compile_def d ... =  
  ...
```

```
let compile_prog p ... =  
  ...
```

```
let rec compile_expr e si env =  
  match e with  
  ...  
  | EApp(fname, arg) ->  
INSTRUCTIONS FOR FUNCTION CALL
```

```
let compile_def d ... =  
INSTRUCTIONS FOR FUNCTION BODY
```

```
(let (x 10)
  (let (z (g x))
    (+ 3 z)))
```

```
(def (g y)
      (+ y 1))
```

```
mov eax, 10
```

esp	0x34 0x2c
eax	10 10

g:

```
mov [esp-4], eax
```

```
mov eax, [esp-4]
```

```
mov [esp-8], after_call
```

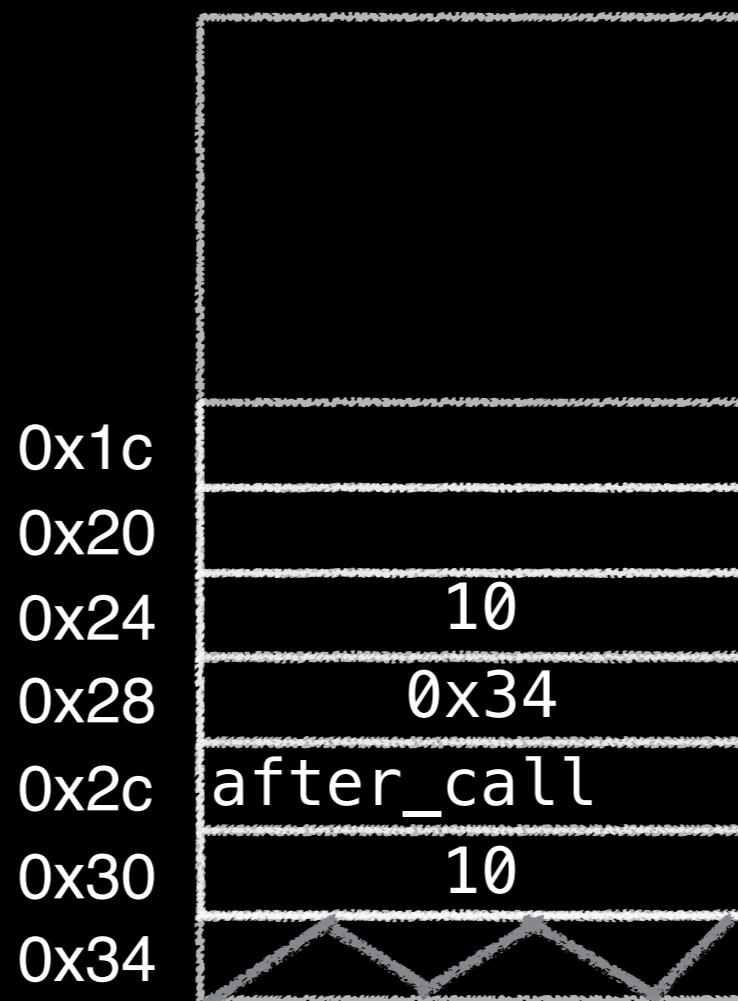
```
mov [esp-12], esp
```

```
mov [esp-16], eax
```

```
sub esp, 8
```

```
jmp g
```

```
after_call:
```



Where is the argument y in terms of the **current value** of esp?

- A: esp
- B: esp-4
- C: esp-8**
- D: esp+4
- E: esp-12

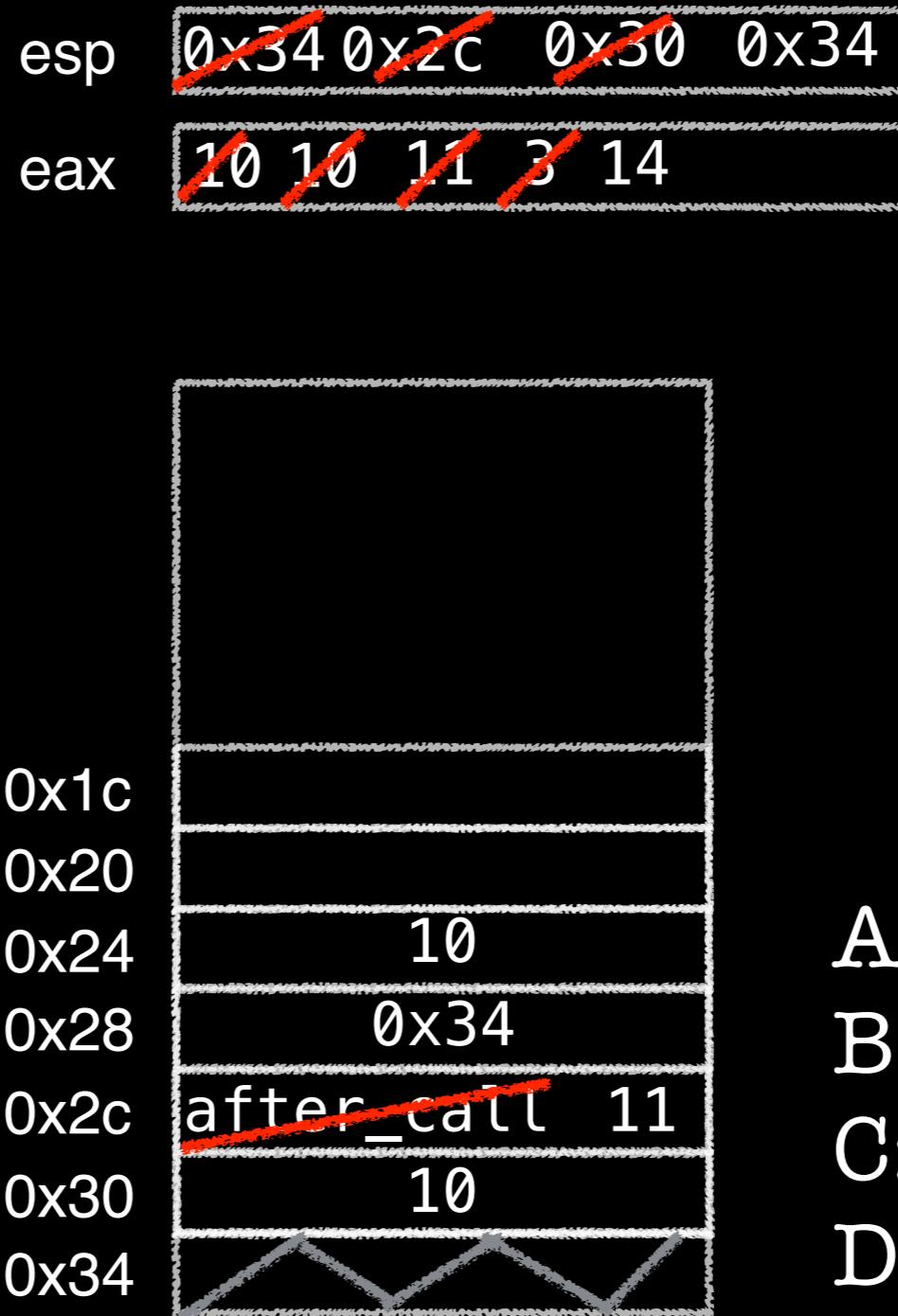
```
let rec compile_expr e si env =
  match e with
```

```
let compile_def d ... =
  INSTRUCTIONS FOR FUNCTION BODY
```

...
| EApp(fname, arg) ->
INSTRUCTIONS FOR FUNCTION CALL

```
(let (x 10)
  (let (z (g x))
    (+ 3 z)))
```

```
mov eax, 10
mov [esp-4], eax
mov eax, [esp-4]
mov [esp-8], after_call
mov [esp-12], esp
mov [esp-16], eax
sub esp, 8
jmp g
after_call:
  mov esp, [esp-8]
  mov [esp-8], eax
  mov eax, 3
  add eax, [esp-8]
```



```
(def (g y)
      (+ y 1))
```

```
g:
  mov eax, [esp-8]
  add eax, 1
  ret
```

What is the
current value
in [esp-8]?

- A: 10
- B: 0x34
- C: after_call
- D: the other 10

```
let rec compile_expr e si env =
  match e with
```

...
| EApp(fname, arg) ->
INSTRUCTIONS FOR FUNCTION CALL

```
let compile_def d ... =
INSTRUCTIONS FOR FUNCTION BODY
```

```
(let (x 10)
  (let (z (g x))
    (+ 3 z)))
```

```
mov eax, 10
mov [esp-4], eax
mov eax, [esp-4]
mov [esp-8], after_call
mov [esp-12], esp
mov [esp-16], eax
sub esp, 8
jmp g
after_call:
  mov esp, [esp-8]
  mov [esp-8], eax
  mov eax, 3
  add eax, [esp-8]
```

esp ~~0x34 0x2c 0x30 0x34~~

eax ~~10 10 11 3 14~~

```
(def (g y)
      (+ y 1))
```

g:

```
mov eax, [esp-8]
..
```

Call setup:

- Always these 3 values
- Always this order
- Always start at current si
- Always subtract to point esp at the return address

0x10	10
0x24	10
0x28	0x34
0x2c	after_call 11
0x30	10
0x34	

```
let rec compile_expr e si env =
  match e with
```

...
| EApp(fname, arg) ->
INSTRUCTIONS FOR FUNCTION CALL

```
let compile_def d ... =
INSTRUCTIONS FOR FUNCTION BODY
```

```
(let (x 10)
  (let (z (g x))
    (+ 3 z)))
```

```
mov eax, 10
mov [esp-4], eax
mov eax, [esp-4]
mov [esp-8], after_call
```

Callee has an easy job:

- Rely on (first) argument in [esp-8], so env starts with [(arg, 2)]
- Start at a “higher” si=3 for any local vars
- Expect [esp] to contain return pointer, use ret

```
mov [esp-8], eax
mov eax, 3
add eax, [esp-8]
```

0x2c	after_call	11
0x30		10
0x34		

```
let rec compile_expr e si env =
  match e with
```

...
| EApp(fname, arg) ->
INSTRUCTIONS FOR FUNCTION CALL

```
(def (g y)
      (+ y 1))
```

```
g:
  mov eax, [esp-8]
  add eax, 1
  ret
```

```
let compile_def d ... =
INSTRUCTIONS FOR FUNCTION BODY
```

```
(let (x 10)
  (let (z (g x))
    (+ 3 z)))
```

```
(def (g y)
      (+ y 1))
```

```
mov eax, 10
mov [esp-4], eax
mov eax, [esp-4]
mov [esp-8], after_call
```

esp	0x34	0x2c	0x30	0x34
eax	10	10	11	3

```
g:
mov eax, [esp-8]
add eax, 1
ret
```

After the call:

mov [esp-1] Rely on old esp at [esp-8] (always)
sub esp, 8 Expect answer to be in eax from callee

jmp g

after_call:

mov esp, [esp-8]

0x20	
0x24	10
0x28	0x34
0x2c	after_call 11
0x30	10
0x34	

```
mov [esp-8], eax
mov eax, 3
add eax, [esp-8]
```

```
let rec compile_expr e si env =
  match e with
```

...
| EApp(fname, arg) ->
INSTRUCTIONS FOR FUNCTION CALL

```
let compile_def d ... =
INSTRUCTIONS FOR FUNCTION BODY
```

```
(let (x 10)
  (let (z (g x))
    (+ 3 z)))
```

```
mov eax, 10
mov [esp-4], eax
mov eax, [esp-4]
mov [esp-8], after_call
mov [esp-12], esp
mov [esp-16], eax
sub esp, 8
jmp g
after_call:
  mov esp, [esp-8]
  mov [esp-8], eax
  mov eax, 3
  add eax, [esp-8]
```

esp	0x34 0x2c 0x30 0x34
eax	10 10 11 3 14
0x1c	
0x20	
0x24	10
0x28	0x34
0x2c	after_call 11
0x30	10
0x34	

```
(def (g y)
      (+ y 1))
```

```
g:
  mov eax, [esp-8]
  add eax, 1
  ret
```

```
let rec compile_expr e si env =
  match e with
```

```
... | EApp(fname, arg) ->
  INSTRUCTIONS FOR FUNCTION CALL
```

```
let compile_def d ... =
  INSTRUCTIONS FOR FUNCTION BODY
```

```
[cs131s@ieng6-201]:pa2-grading:509$ ./software/nonrdist/valgrind-3.10.0/bin/valgrind ./output/foo.run
==31414== Memcheck, a memory error detector
==31414== Copyright (C) 2002-2013, and GNU GPL'd, by Julian Seward et al.
==31414== Using Valgrind-3.10.0 and LibVEX; rerun with -h for copyright info
==31414== Command: ./output/foo.run
==31414==
==31414== Invalid write of size 4
==31414==    at 0x80488E9: ??? (in /home/linux/ieng6/cs131s/cs131s/pa2-grading/output/foo.run)
==31414==    by 0x408BA42: (below main) (in /usr/lib/libc-2.17.so)
==31414== Address 0xfe9e2318 is on thread 1's stack
==31414== 4 bytes below stack pointer
==31414==
==31414== Invalid write of size 4
==31414==    at 0x80488F1: ??? (in /home/linux/ieng6/cs131s/cs131s/pa2-grading/output/foo.run)
==31414==    by 0x408BA42: (below main) (in /usr/lib/libc-2.17.so)
==31414== Address 0xfe9e2314 is on thread 1's stack
==31414== 8 bytes below stack pointer
==31414==
```

How much stack space is used by this expression (what will be the largest N in an esp-N generated)?

(let (x 5) (+ x 1))

- A: 16 bytes (some instruction has stack index 4)
- B: 12 bytes (some instruction has stack index 3)
- C: 8 bytes (some instruction has stack index 2)
- D: 4 bytes (some instruction has stack index 1)

```
(let (x 10)
  (let (z (g x))
    (+ 3 z)))
```

```
mov eax, 10
mov [esp-4], eax
mov eax, [esp-4]
mov [esp-8], after_call
mov [esp-12], esp
mov [esp-16], eax
sub esp, 8
jmp g
after_call:
  mov esp, [esp-8]
  mov [esp-8], eax
  mov eax, 3
  add eax, [esp-8]
```

```
(def (twosqrt3x x)
  (let (ans (do_sqrt (* x 3)))
    (* 2 ans)))
(twosqrt3x 4)
```

twosqrt3x:

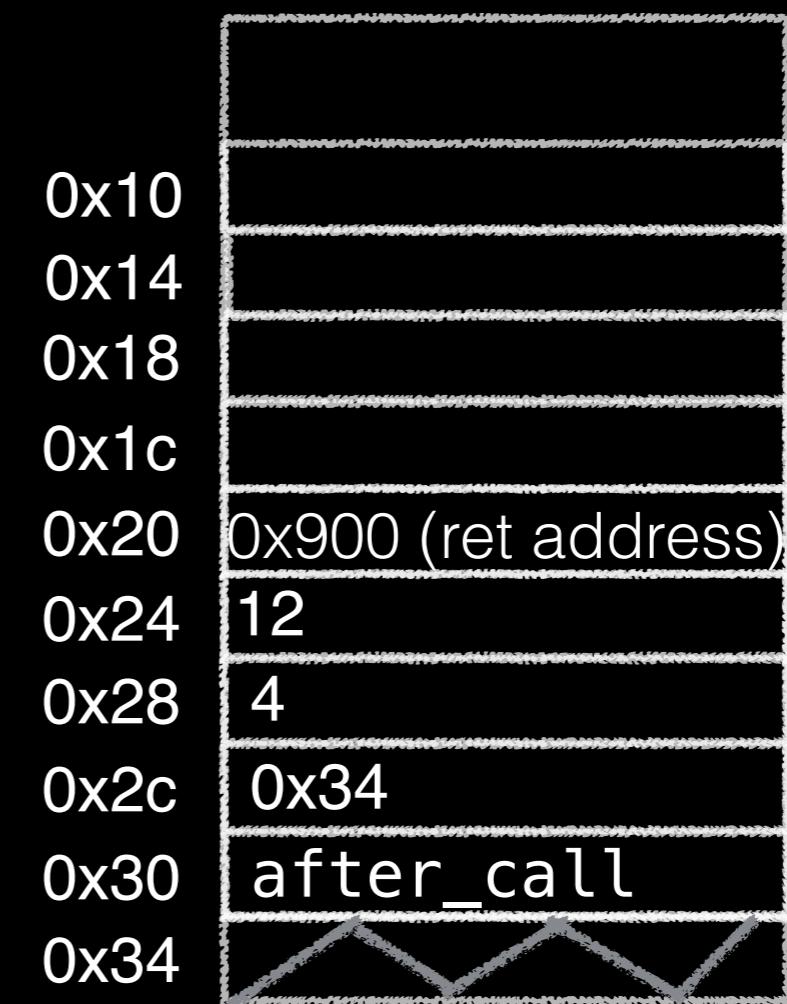
```
  mov eax, [esp-8]
  imul eax, 3
  sub esp, 8
  push eax
  call do_sqrt
```

0x900

```
  mov eax, 4
  mov [esp-4], after_call
  mov [esp-8], esp
  mov [esp-12], eax
  sub esp, 4
  jmp twosqrt3x
after_call:
```

```
int do_sqrt(int val) {
    float asF = (float)val;
    return (int)(sqrt(asF));
}
```

esp ~~0x34 0x30 0x28 0x24~~
eax ~~+ 12~~



```
(def (twosqrt3x x)
  (let (ans (do_sqrt (* x 3)))
    (* 2 ans)))
(twosqrt3x 4)
```

twosqrt3x:

```
  mov eax, [esp-8]
  imul eax, 3
  sub esp, 8
  push eax
  call do_sqrt
```

0x900

What will be in ESP when do_sqrt is complete?

- A: 0x24 (the value of ESP before call)
- B: 0x20 (do_sqrt's return address)
- C: 0x30 (the original ret address)
- D: We can't possibly know
- E: None of the above

after_call:

```
int do_sqrt(int val) {
    float asF = (float)val;
    return (int)(sqrt(asF));
}
```

esp C Happens
eax C Happens

0x10	
0x14	
0x18	
0x1c	C Happens
0x20	0x900 (ret address)
0x24	12
0x28	4
0x2c	0x34
0x30	after_call
0x34	

```
(def (twosqrt3x x)
  (let (ans (do_sqrt (* x 3)))
    (* 2 ans)))
(twosqrt3x 4)
```

twosqrt3x:

```
  mov eax, [esp-8]
  imul eax, 3
  sub esp, 8
  push eax
  call do_sqrt
```

0x900

```
int do_sqrt(int val) {
    float asF = (float)val;
    return (int)(sqrt(asF));
}
```

esp C Happens
eax C Happens

What will be in EAX when do_sqrt is complete?

- A: We can't possibly know
- B: 3 (the return value)
- C: 12
- D: 4
- E: None of the above

after_call:

0x10
0x14
0x18
0x1c
0x20 C Happens 0x900 (ret address)
0x24 12
0x28 4
0x2c 0x34
0x30 after_call
0x34

```
(def (twosqrt3x x)
  (let (ans (do_sqrt (* x 3)))
    (* 2 ans)))
(twosqrt3x 4)
```

twosqrt3x:

```
  mov eax, [esp-8]
  imul eax, 3
  sub esp, 8
  push eax
  call do_sqr
```

0x900

Just as it was
before call

```
  mov eax, 4
  mov [esp-4], after_call
  mov [esp-8], esp
  mov [esp-12], eax
  sub esp, 4
  jmp twosqrt3x
after_call:
```

```
int do_sqrt(int val) {
    float asF = (float)val;
    return (int)(sqrt(asF));
}
```

esp

0x34 0x30 0x28 0x24

eax

+ 12 3

0x10
0x14
0x18
0x1c
0x20
0x24
0x28
0x2c
0x30
0x34

0x900 (ret address)
12
4
0x34
after_call

Return
value

```
(def (twosqrt3x x)
  (let (ans (do_sqrt (* x 3)))
    (* 2 ans)))
(twosqrt3x 4)
```

twosqrt3x:

```
  mov eax, [esp-8]
  imul eax, 3
  sub esp, 8
  push eax
  call do_sqrt
0x900  add esp, 12
  mov [esp-12], eax
  mov eax, 2
  imul eax, [esp-12]
  ret
  mov eax, 4
  mov [esp-4], after_call
  mov [esp-8], esp
  mov [esp-12], eax
  sub esp, 4
  jmp twosqrt3x
after_call:
  mov esp, [esp-8]
  ret
```

```
int do_sqrt(int val) {
  float asF = (float)val;
  return (int)(sqrt(asF));
}
```

esp 
eax 

0x10
0x14
0x18
0x1c
0x20 0x900 (ret address)
0x24 12 3
0x28 4
0x2c 0x34
0x30 after_call
0x34



sum:

```
→ mov eax, [esp-8]
    cmp eax, 1
    jne false_branch
        mov eax, 1
        jmp after_if
false_branch:
    mov eax, [esp-8]
    sub eax, 1
    mov [esp-12], after_call
    mov [esp-16], esp
    mov [esp-20], eax
    sub esp, 12
    jmp sum
after_call:
    mov esp, [esp-8]
    mov [esp-12], eax
    mov eax, [esp-8]
    add eax, [esp-12]
after_if:
    ret
```

def sum(x):

```
if x = 1: 1
else:
    x + sum(x - 1)
```

eq	NO	NO	YES
esp	0x30	0x24	0x18
	0x24	0x28	0x30

eax	3	3	2	1	2	3	6
-----	---	---	---	---	---	---	---

What values will be stored in 0x10, 0x14, 0x18 by the next three instructions?

- A: 1, 0x20, after_call
- B: 2, 0x24, after_call
- C: 2, 0x20, after_call
- D: 1, 0x24, after_call

0x10	1
0x14	0x24
0x18	after_call 1
0x1c	2
0x20	0x30
0x24	after_call 3
0x28	3
0x2c	old_esp
0x30	return_ptr
0x34	


```

def twosqrt3x(x):
    let ans = do_sqrt(x * 3) in
        2 * ans
twosqrt3x(4)
twosqrt3x:
    mov eax, [esp-8]
    imul eax, 3
    sub esp, 8
    push eax
    call do_sqrt
0x900    add esp, 12
    mov [esp-12], eax
    mov eax, 2
    imul eax, [esp-12]
    ret
    mov eax, 4
    mov [esp-4], after_call
    mov [esp-8], esp
    mov [esp-12], eax
    sub esp, 4
    jmp twosqrt3x
after_call:
    mov esp, [esp-8]
    ret

```

```

int do_sqrt(int val) {
    float asD = (float)val;
    return (int)(sqrt(asD));
}

```

