

# *Bottling Computation Patterns*

## *Polymorphism and HOFs are the Secret Sauce*

**Refactor** arbitrary *repeated* code patterns ...

... into precisely *specified* and reusable **functions**

### *EXERCISE: Iteration*

Write a function that *squares* a list of `Int`

```
squares :: [Int] -> [Int]  
squares ns = ???
```

When you are done you should see

```
>>> squares [1,2,3,4,5]
[1,4,9,16,25]
```

## *Pattern: Iteration*

Next, lets write a function that converts a `String` to uppercase.

```
>>> shout "hello"
"HELLO"
```

Recall that in Haskell, a `String` is just a `[Char]`.

```
shout :: [Char] -> [Char]
shout = ???
```

Hoogle (<http://haskell.org/hoogle>) to see how to transform an individual `Char`

## *Iteration*

Common strategy: *iteratively* transform *each element* of input list

Like humans and monkeys, `shout` and `squares` share 93% of their DNA  
([http://www.livescience.com/health/070412\\_rhesus\\_monkeys.html](http://www.livescience.com/health/070412_rhesus_monkeys.html))

Super common *computation pattern*!

## *Abstract Iteration “Pattern” into Function*

Remember D.R.Y. (Don't repeat yourself)

**Step 1** Rename all variables to remove accidental *differences*

```
-- rename 'squares' to 'foo'
foo []      = []
foo (x:xs) = (x * x)      : foo xs
```

```
-- rename 'shout' to 'foo'
foo []      = []
foo (x:xs) = (toUpper x) : foo xs
```

### Step 2 Identify what is *different*

- In `squares` we *transform* `x` to `x * x`
- In `shout` we *transform* `x` to `Data.Char.toUpper x`

### Step 3 Make *differences* a parameter

- Make *transform* a parameter `f`

```
foo f []      = []
foo f (x:xs) = (f x) : foo f xs
```

**Done** We have *bottled* the computation pattern as `foo` (aka `map`)

```
map f []      = []
map f (x:xs) = (f x) : map f xs
```

`map` bottles the common pattern of iteratively transforming a list:



Fairy In a Bottle

# QUIZ

What is the type of `map`?

*aka filter*

*doTwice  $f x = f(f x)$*

```
map :: ???
map f [] = []
map f (x:xs) = (f x) : map f xs
```

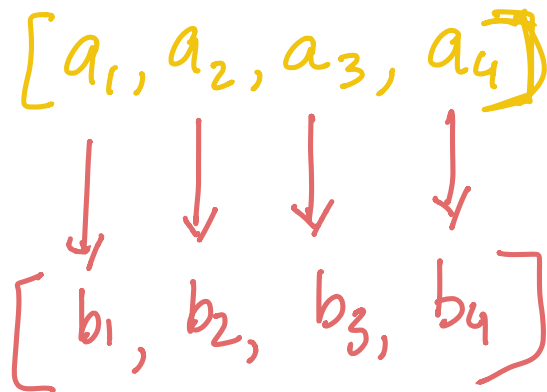
A.  $(\text{Int} \rightarrow \text{Int}) \rightarrow [\text{Int}] \rightarrow [\text{Int}]$

• B.  $(a \rightarrow a) \rightarrow [a] \rightarrow [a]$

C.  $[a] \rightarrow [b]$

• D.  $(a \rightarrow b) \rightarrow [a] \rightarrow [b]$

E.  $(a \rightarrow b) \rightarrow [a] \rightarrow [a]$



## *The type precisely describes `map`*

```
>>> :type map
map :: (a -> b) -> [a] -> [b]
```

That is, `map` takes two inputs

- a *transformer* of type `a -> b`
- a *list* of values `[a]`

and it returns as output

- a list of values `[b]`

that can only come by applying `f` to each element of the input list.

## *Reusing the Pattern*

Lets reuse the pattern by *instantiating* the transformer

# *shout*

*-- OLD with recursion*

```
shout :: [Char] -> [Char]
```

```
shout [] = []
```

```
shout (x:xs) = Char.toUpperCase x : shout xs
```

*-- NEW with map*

```
shout :: [Char] -> [Char]
```

```
shout xs = map (???) xs
```

# *squares*

*-- OLD with recursion*

```
squares :: [Int] -> [Int]
```

```
squares [] = []
```

```
squares (x:xs) = (x * x) : squares xs
```

*-- NEW with map*

```
squares :: [Int] -> [Int]
```

```
squares xs = map (???) xs
```

# EXERCISE

Suppose I have the following type

```
type Score = (Int, Int) -- pair of scores for Hw0, Hw1
```

Use `map` to write a function

```
total :: [Score] -> [Int]
total xs = map (???) xs
```

such that

```
>>> total [(10, 20), (15, 5), (21, 22), (14, 16)]
[30, 20, 43, 30]
```



# *The Case of the Missing Parameter*

Note that we can write `shout` like this

```
shout :: [Char] -> [Char]
shout = map Char.toUpper
```

Huh. No parameters? Can someone explain?

# *The Case of the Missing Parameter*

In Haskell, the following all mean the same thing

Suppose we define a function

```
add :: Int -> Int -> Int
add x y = x + y
```

Now the following all *mean the same thing*

```

plus x y = add x y
plus x   = add x
plus     = add

```

Why? *equational reasoning!* In general

```
foo x = e x
```

-- is equivalent to

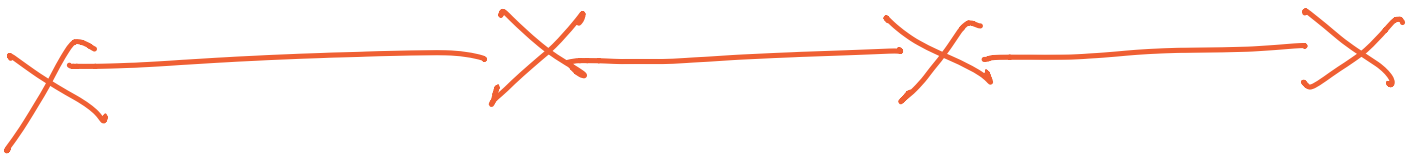
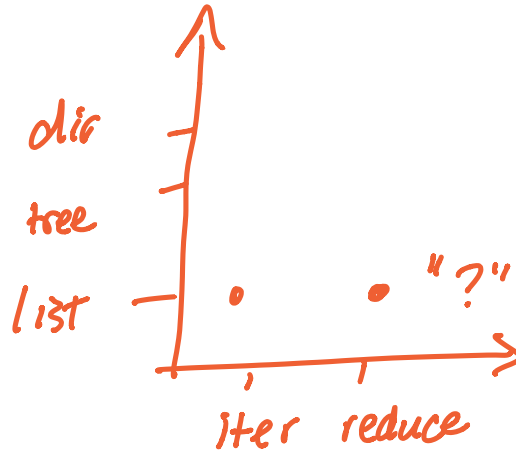
```
foo = e
```

as long as  $x$  doesn't appear in  $e$ .

Thus, to save some typing, we omit the extra parameter. *patterns*

HOFs & Bottling Patterns

Hello World! (10)



## Pattern: Reduction

Computation patterns are *everywhere* lets revisit our old `sumList`

```
sumList :: [Int] -> Int
sumList []      = 0
sumList (x:xs) = x + sumList xs
```

Next, a function that *concatenates* the `String` `s` in a list

```
catList :: [String] -> String
catList []      = ""
catList (x:xs) = x ++ (catList xs)
```

*Lets spot the pattern!*

**Step 1** Rename

```
foo []      = 0
foo (x:xs) = x + foo xs
```

```
foo []      = ""
foo (x:xs) = x ++ foo xs
```

**Step 2** Identify what is *different*

1. ???

2. ???

**Step 3** Make *differences* a parameter

foo p1 p2 [] = ???

foo p1 p2 (x:xs) = ???

## *EXERCISE: Reduction/Folding*

This pattern is commonly called *reducing* or *folding*

`foldr :: (a -> b -> b) -> b -> [a] -> b``foldr op base [] = base``foldr op base (x:xs) = op x (foldr op base xs)`

Can you figure out how `sumList` and `catList` are just *instances* of `foldr`?

*sumList = foldr (+) 0*

*catList = foldr (++) ""*

```
sumList :: [Int] -> Int  
sumList xs = foldr (?op) (?base) xs
```

```
catList :: [String] -> String  
catList xs = foldr (?op) (?base) xs
```

## *Executing foldr*

To develop some intuition about `foldr` lets “run” it a few times by hand.

```
foldr op b (a1:a2:a3:a4:[])
```

```
==>
```

```
  a1 `op` (foldr op b (a2:a3:a4:[]))
```

```
==>
```

```
  a1 `op` (a2 `op` (foldr op b (a3:a4:[])))
```

```
==>
```

```
  a1 `op` (a2 `op` (a3 `op` (foldr op b (a4:[]))))
```

```
==>
```

```
  a1 `op` (a2 `op` (a3 `op` (a4 `op` foldr op b [])))
```

```
==>
```

```
  a1 `op` (a2 `op` (a3 `op` (a4 `op` b)))
```

Look how it *mirrors* the structure of lists!

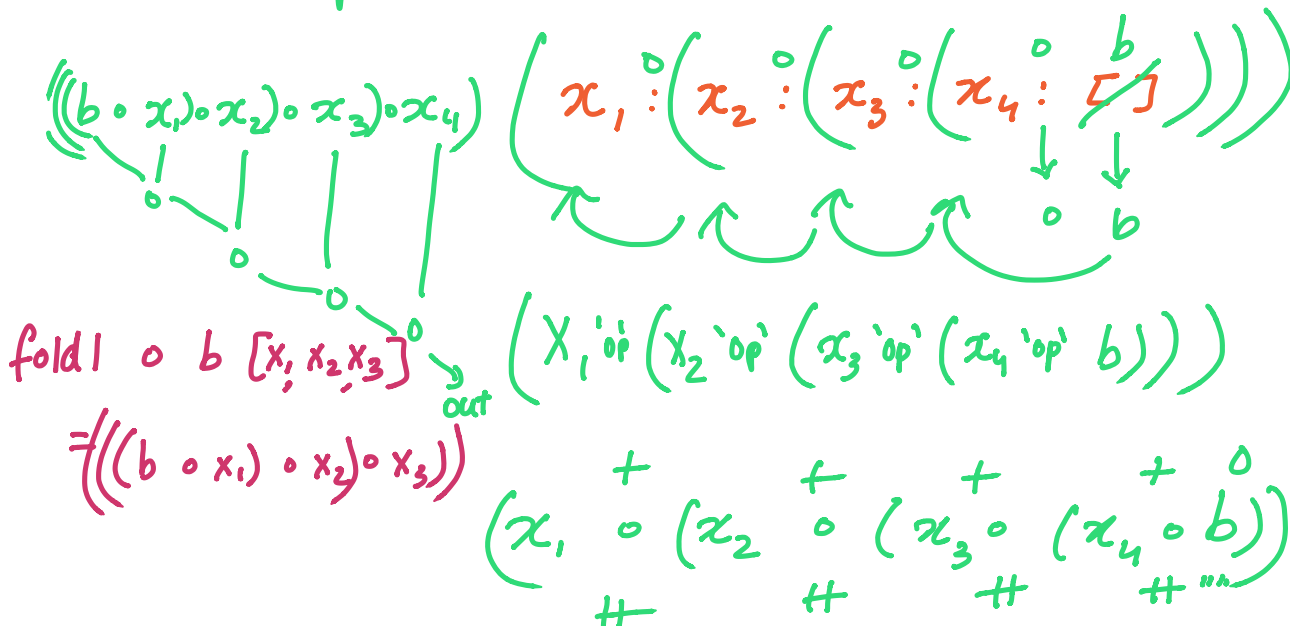
- (:) is replaced by op
- [] is replaced by base

So

```
foldr (+) 0 (x1:x2:x3:x4:[])
```

```
==> x1 + (x2 + (x3 + (x4 + 0)))
```

*foldl*



# Typing *foldr*

`foldr :: (a -> b -> b) -> b -> [a] -> b`

`foldr op base [] = base`

`foldr op base (x:xs) = op x (foldr op base xs)`

`foldr` takes as input

- a *reducer* function of type `a -> b -> b`
- a *base* value of type `b`
- a *list* of values to reduce `[a]`

and returns as output

- a *reduced* value `b`

$$(a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b$$

$$\text{foldr } op \ b \ (a_1 : a_2 : a_3 : a_4 : [])$$

$$(a_1 \circ (a_2 \circ (a_3 \circ (a_4 \circ \underline{b}))))$$

$$\underline{op} :: a \rightarrow b \rightarrow b$$

## QUIZ

Recall the function to compute the `len` of a list

```
len :: [a] -> Int
len []      = 0
len (x:xs) = 1 + len xs
```

Which of these is a valid implementation of ~~len~~ *len*

- ✗ A. `len = foldr (\n -> n + 1) 0`
- B. `len = foldr (\n m -> n + m) 0`
- C. `len = foldr (\_ n -> n + 1) 0`
- D. `len = foldr (\x xs -> 1 + len xs) 0`
- E. All of the above

$$\begin{array}{c}
 \underline{x_1} : x_2 : (x_3 : (\underline{x_4} : [])) \\
 \downarrow \\
 (1 + (1 + (1 + (1 + 0)))) \\
 \underbrace{\hspace{10em}} \\
 (\backslash x \underline{v} \rightarrow 1 + \underline{v})
 \end{array}$$

## The Missing Parameter Revisited

We wrote foldr as



```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr op base []      = base
foldr op base (x:xs) = op x (foldr op base xs)
```

but can also write this

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr op base = go
  where
    go []      = base
    go (x:xs) = op x (go xs)
```

Can someone explain where the `xs` went *missing*?

## *Trees*

Recall the `Tree a` type from last time

```

data Tree a
  = Leaf
  | Node a (Tree a) (Tree a)

```

For example here's a tree

```

tree2 :: Tree Int
tree2 = Node 2 Leaf Leaf

```

```

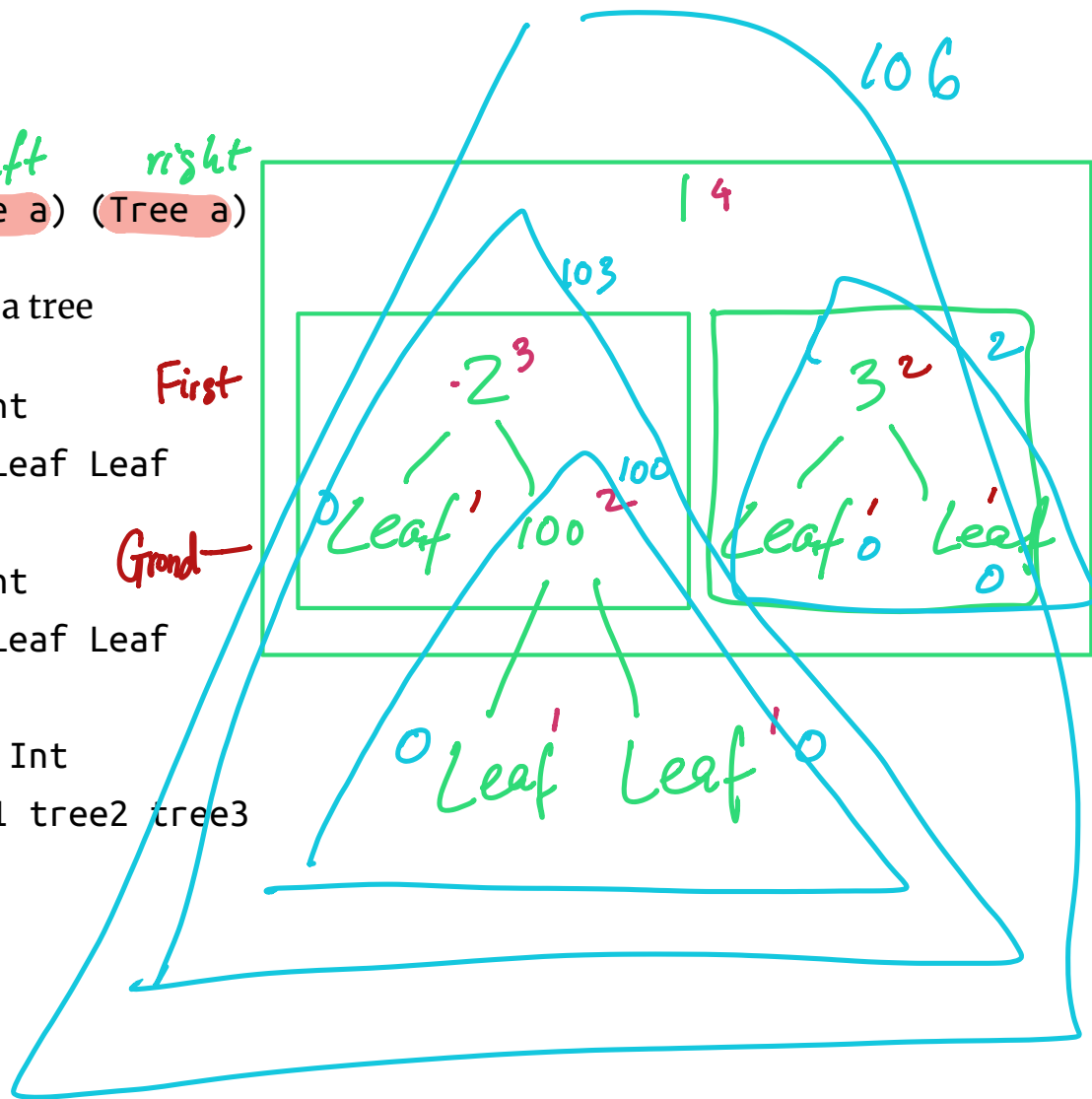
tree3 :: Tree Int
tree3 = Node 3 Leaf Leaf

```

```

tree123 :: Tree Int
tree123 = Node 1 tree2 tree3

```



## Some Functions on Trees

Lets write a function to compute the height of a tree

```
height :: Tree a -> Int
height Leaf      = 0
height (Node x l r) = 1 + max (height l) (height r)
```

Here's another to *sum* the leaves of a tree:

```
sumTree :: Tree Int -> Int
sumTree Leaf      = ???
sumTree (Node x l r) = ???
```

Gathers all the elements that occur as leaves of the tree:

```
toList :: Tree a -> [a]
toList Leaf      = ???
toList (Node x l r) = ???
```

Lets give it a whirl

```
>>> height tree123
2
```

```
>>> sumTree tree123
6
```

```
>>> toList tree123
[1,2,3]
```

## Pattern: Tree Fold

Can you spot the pattern? Those three functions are almost the same!

**Step 1:** Rename to maximize similarity

```
-- height
foo Leaf          = 0
foo (Node x l r) = 1 + max (foo l) (foo r)
```

```
-- sumTree
foo Leaf          = 0
foo (Node x l r) = foo l + foo r
```

```
-- toList
foo Leaf          = []
foo (Node x l r) = x : foo l ++ foo r
```

**Step 2:** Identify the differences

1. ???
2. ???

**Step 3** Make *differences* a parameter

```
foo p1 p2 Leaf          = ???
foo p1 p2 (Node x l r) = ???
```

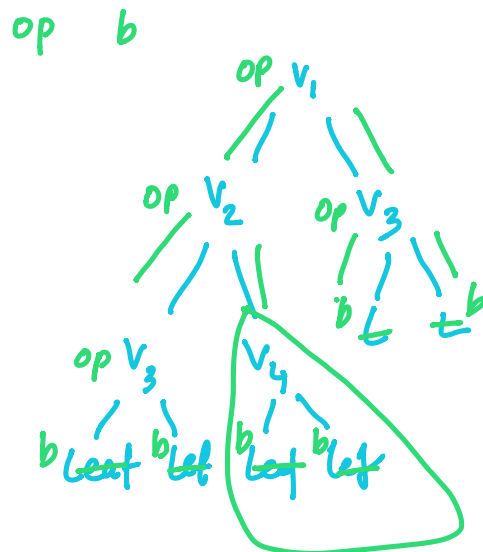
## *Pattern: Folding on Trees*

`tFold op b Leaf = b`

`tFold op b (Node x l r) = op x (tFold op b l) (tFold op b r)`

Lets try to work out the type of `tFold` !

`tFold :: t_op -> t_b -> Tree a -> t_out`



# QUIZ

Suppose that  $t :: \text{Tree Int}$ .

What does `tFold (\x y z -> y + z) 1 t` return?

- a. `0`
- b. the *largest* element in the tree `t`
- c. the *height* of the tree `t`
- d. the *number-of-leaves* of the tree `t`
- e. type error

# EXERCISE

Write a function to compute the *largest* element in a tree or `0` if tree is empty or all negative.

```
treeMax :: Tree Int -> Int
treeMax t = tFold f b t
  where
    f      = ???
    b      = ???
```

## *Map over Trees*

We can also write a `tmap` equivalent of `map` for `Tree` s

```
treeMap :: (a -> b) -> Tree a -> Tree b
treeMap f (Leaf x)    = Leaf (f x)
treeMap f (Node l r) = Node (treeMap f l) (treeMap f r)
```

which gives

```
>>> treeMap (\n -> n * n) tree123    -- square all elements
of tree
Node 1 (Node 4 Leaf Leaf) (Node 9 Leaf Leaf)
```

## *EXERCISE*

Recursion is **HARD TO READ** do we really have to use it ?

Lets rewrite `treeMap` using `tFold` !

```
treeMap :: (a -> b) -> Tree a -> Tree b
treeMap f t = tFold op base t
  where
    op      = ???
    base    = ???
```

When you are done, we should get

```
>>> animals = Node "cow" (Node "piglet" Leaf Leaf) (Leaf "hip
po" Leaf Leaf)
>>> treeMap reverse animals
Node "woc" (Node "telgip" Leaf Leaf) (Leaf "oppih" Leaf Leaf)
```



WED NOV 4

NEW DEADLINE  
for 01-TREES

## Examples: *foldDir*

```
data Dir a
  = Fil a           -- ^ A single file named `a`
  | Sub a [Dir a]   -- ^ A sub-directory name `a` with contents `[Dir a]`
```

```
data DirElem a
  = SubDir a        -- ^ A single Sub-Directory named `a`
  | File a          -- ^ A single File named `a`
```

```
foldDir :: ([a] -> r -> DirElem a -> r) -> r -> Dir a -> r
foldDir f r0 dir = go [] r0 dir
```

**where**

```
go stk r (Fil a)    = f stk r (File a)
go stk r (Sub a ds) = L.foldl' (go stk') r' ds
```

**where**

```
r'    = f stk r (SubDir a)
stk'  = a:stk
```

*foldDir* takes as input

- an *accumulator* *f* of type `[a] -> r -> DirElem a -> r`
  - takes as *input* the path `[a]` , the current result `r` , the next `DirElem [a]`
  - and returns as *output* the new result `r`

- an *initial* value of the result `r0` and
- directory to fold over `dir`

And returns the result of running the *accumulator* over the whole `dir` .

## *Examples: Spotting Patterns In The “Real” World*

These patterns in “toy” functions appear regularly in “real” code

1. Start with beginner’s version riddled with explicit recursion (swizzle-v0.html).
2. Spot the patterns and eliminate recursion using HOFs (swizzle-v1.html).
3. Finally refactor the code to “swizzle” and “unswizzle” without duplication (swizzle-v2.html).

**Try it yourself**

- Rewrite the code that swizzles `Char` to use the `Map k v` type in `Data.Map`

## Which is more readable? HOFs or Recursion

At first, *recursive* versions of `shout` and `squares` are easier to follow

- `fold` takes a bit of getting used to!

With practice, the *higher-order* versions become easier

- only have to understand specific operations
- recursion is lower-level & have to see “loop” structure
- worse, potential for making silly off-by-one errors → 2004

Indeed, HOFs were the basis of `map/reduce` and the big-data revolution (<http://en.wikipedia.org/wiki/MapReduce>)

- Can *parallelize* and *distribute* computation patterns just once ([https://www.usenix.org/event/osdio4/tech/full\\_papers/dean/dean.pdf](https://www.usenix.org/event/osdio4/tech/full_papers/dean/dean.pdf))
- Reuse (<http://en.wikipedia.org/wiki/MapReduce>) across hundreds or thousands of instances!

HOFs FTW!

0£ TREES Due  
Nov 4

10/27/20, 9:28 AM

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