

# Bottling Computation Patterns

*Polymorphism and HOFs are the Secret Sauce*

**Refactor** arbitrary *repeated* code patterns ...

... into precisely *specified* and *reusable* functions

## EXERCISE: Iteration

Write a function that *squares* a list of `Int`

```
squares :: [Int] -> [Int]  
squares ns = ???
```

When you are done you should see

```
>>> squares [1,2,3,4,5]
[1,4,9,16,25]
```

## *Pattern: Iteration*

Next, let's write a function that converts a `String` to uppercase.

```
>>> shout "hello"
"HELLO"
```

Recall that in Haskell, a `String` is just a `[Char]`.

```
shout :: [Char] -> [Char]
shout = ???
```

Hoogle (<http://haskell.org/hoogle>) to see how to transform an individual `Char`

# *Iteration*

Common strategy: *iteratively transform each element of input list*

Like humans and monkeys, shout and squares share 93% of their DNA  
([http://www.livescience.com/health/070412\\_rhesus\\_monkeys.html](http://www.livescience.com/health/070412_rhesus_monkeys.html))

Super common *computation pattern!*

## *Abstract Iteration “Pattern” into Function*

Remember D.R.Y. (Don’t repeat yourself)

**Step 1** Rename all variables to remove accidental *differences*

```
-- rename 'squares' to 'foo'
foo []      = []
foo (x:xs) = (x * x)    : foo xs
```

```
-- rename 'shout' to 'foo'
foo []      = []
foo (x:xs) = (toUpper x) : foo xs
```

### Step 2 Identify what is *different*

- In `squares` we *transform* `x` to `x * x`
- In `shout` we *transform* `x` to `Data.Char.toUpper x`

### Step 3 Make *differences* a parameter

- Make *transform* a parameter `f`

```
foo f []      = []
foo f (x:xs) = (f x) : foo f xs
```

**Done** We have *bottled* the computation pattern as `foo` (aka `map`)

```
map f []      = []
map f (x:xs) = (f x) : map f xs
```

`map` bottles the common pattern of iteratively transforming a list:



Fairy In a Bottle

# QUIZ

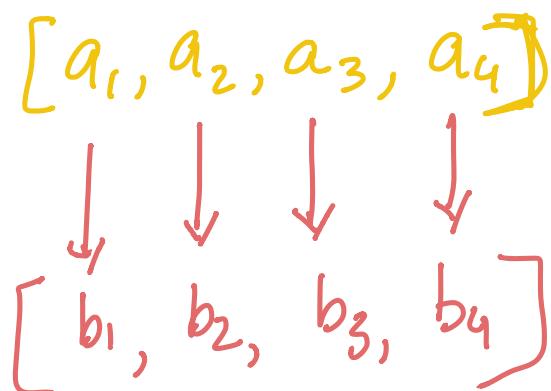
aka filter

What is the type of `map`?

```
map :: ???  
map f []      = []  
map f (x:xs) = (f x) : map f xs
```

doTwice  $f x = f(f x)$

- A.  $(\text{Int} \rightarrow \text{Int}) \rightarrow [\text{Int}] \rightarrow [\text{Int}]$
- B.  $(a \rightarrow a) \rightarrow [a] \rightarrow [a]$
- C.  $[a] \rightarrow [b]$
- D.  $(a \rightarrow b) \rightarrow [a] \rightarrow [b]$
- E.  $(a \rightarrow b) \rightarrow [a] \rightarrow [a]$



## *The type precisely describes map*

```
>>> :type map  
map :: (a -> b) -> [a] -> [b]
```

That is, `map` takes two inputs

- a *transformer* of type `a -> b`
- a *list of values* `[a]`

and it returns as output

- a *list of values* `[b]`

that can only come by applying `f` to each element of the input list.

## *Reusing the Pattern*

Lets reuse the pattern by *instantiating* the transformer

# shout

-- OLD with recursion

```
shout :: [Char] -> [Char]
shout []      = []
shout (x:xs) = Char.toUpperCase x : shout xs
```

-- NEW with map

```
shout :: [Char] -> [Char]
shout xs = map (????) xs
```

# squares

-- OLD with recursion

```
squares :: [Int] -> [Int]
squares []      = []
squares (x:xs) = (x * x) : squares xs
```

-- NEW with map

```
squares :: [Int] -> [Int]
squares xs = map (????) xs
```

# EXERCISE

Suppose I have the following type

```
type Score = (Int, Int) -- pair of scores for Hw0, Hw1
```

Use `map` to write a function

```
total :: [Score] -> [Int]
total xs = map (???) xs
```

such that

```
>>> total [(10, 20), (15, 5), (21, 22), (14, 16)]
[30, 20, 43, 30]
```

# *The Case of the Missing Parameter*

Note that we can write `shout` like this

```
shout :: [Char] -> [Char]  
shout = map Char.toUpperCase
```

Huh. No parameters? Can someone explain?

# *The Case of the Missing Parameter*

In Haskell, the following all mean the same thing

Suppose we define a function

```
add :: Int -> Int -> Int  
add x y = x + y
```

Now the following all *mean the same thing*

```
plus x y = add x y
plus x    = add x
plus      = add
```

Why? *equational reasoning!* In general

`foo x = e x`

-- is equivalent to

`foo = e`

as long as `x` doesn't appear in `e`.



Thus, to save some typing, we omit the extra parameter. *patterns*

## HOFs & Bottling Patterns

Hello World! (10)



## Pattern: Reduction

Computation patterns are *everywhere* lets revisit our old `sumList`

```
sumList :: [Int] -> Int
sumList []      = 0
sumList (x:xs) = x + sumList xs
```

Next, a function that *concatenates* the `String`s in a list

```
catList :: [String] -> String
catList []      = ""
catList (x:xs) = x ++ (catList xs)
```

*Lets spot the pattern!*

**Step 1** Rename

```
foo []      = 0
foo (x:xs) = x + foo xs
```

```
foo []      = ""
foo (x:xs) = x ++ foo xs
```

## Step 2 Identify what is *different*

1. ???

2. ???

## Step 3 Make *differences* a parameter

foo **p1** **p2** [] = ???

foo **p1** **p2** (x:xs) = ???

## EXERCISE: Reduction/Folding

This pattern is commonly called *reducing* or *folding*

`foldr :: (a -> b -> b) -> b -> [a] -> b`

`foldr op base [] = base`

`foldr op base (x:xs) = op x (foldr op base xs)`

Can you figure out how `sumList` and `catList` are just *instances* of `foldr`?

`sumList = foldr (+) 0`

`catList = foldr (++) ""`

```
sumList :: [Int] -> Int
sumList xs = foldr (?op) (?base) xs
```

```
catList :: [String] -> String
catList xs = foldr (?op) (?base) xs
```

## *Executing `foldr`*

To develop some intuition about `foldr` lets “run” it a few times by hand.

```

foldr op b (a1:a2:a3:a4:[])
==>
  a1 `op` (foldr op b (a2:a3:a4:[]))
==>
  a1 `op` (a2 `op` (foldr op b (a3:a4:[])))
==>
  a1 `op` (a2 `op` (a3 `op` (foldr op b (a4:[]))))
==>
  a1 `op` (a2 `op` (a3 `op` (a4 `op` foldr op b [])))
==>
  a1 `op` (a2 `op` (a3 `op` (a4 `op` b)))

```

Look how it *mirrors* the structure of lists!

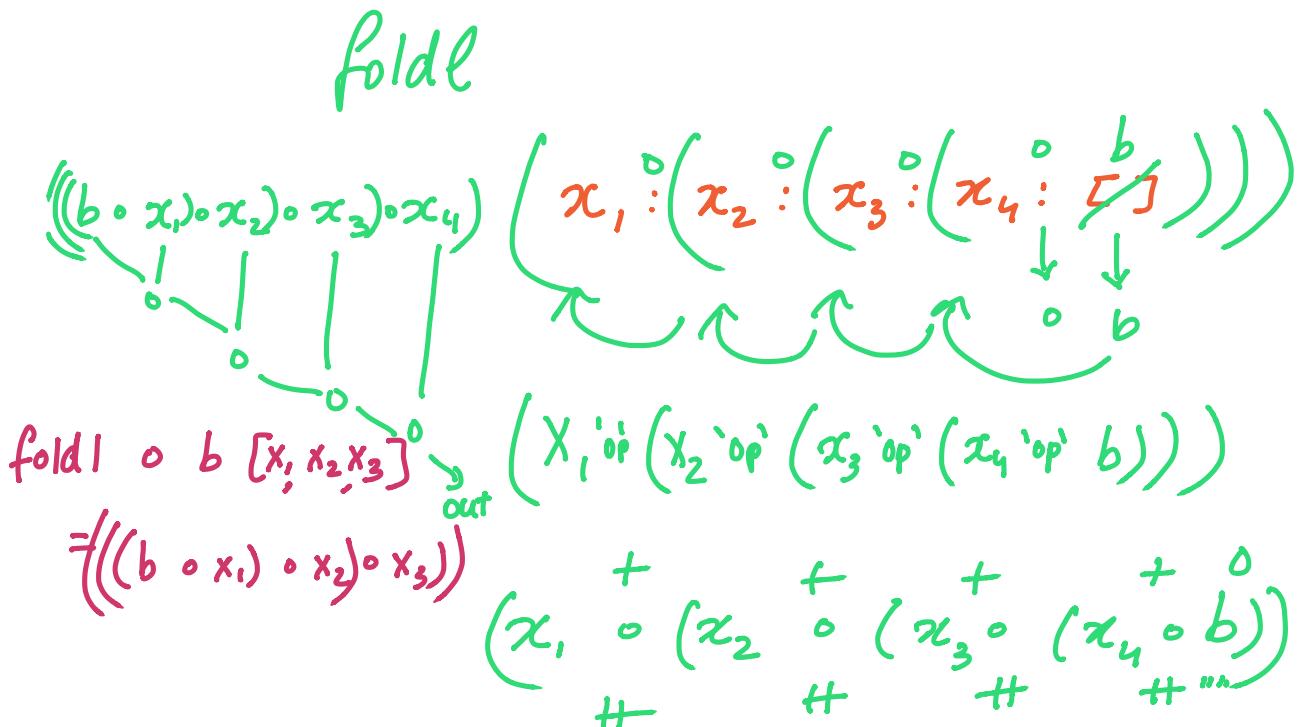
- `(:)` is replaced by `op`
- `[]` is replaced by `base`

So

```

foldr (+) 0 (x1:x2:x3:x4:[])
==> x1 + (x2 + (x3 + (x4 + 0)))

```



## *Typing foldr*

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr op base []      = base
foldr op base (x:xs) = op x (foldr op base xs)
```

`foldr` takes as input

- a *reducer* function of type  $a \rightarrow b \rightarrow b$
  - a *base* value of type  $b$
  - a *list* of values to reduce  $[a]$

and returns as output

- a *reduced* value b

$$\begin{aligned}
 & (\textcolor{red}{(a \rightarrow b \rightarrow b) \rightarrow b} \rightarrow [\textcolor{green}{a}]) \rightarrow b \\
 \text{foldr op } b \underset{=} {\overbrace {(\textcolor{green}{a_1 : a_2 : a_3 : a_4 : []})}} &= \\
 & \left( a_1 : \underset{\textcolor{red}{\curvearrowright}}{(a_2 : \underset{\textcolor{red}{\curvearrowright}}{(a_3 : \underset{\textcolor{red}{\curvearrowright}}{(a_4 : \underset{\textcolor{red}{\curvearrowright}}{b}))})} \right) \\
 & \circ :: \textcolor{green}{a \rightarrow b \rightarrow b}
 \end{aligned}$$

# QUIZ

Recall the function to compute the `len` of a list

```

len :: [a] -> Int
len []      = 0
len (x:xs) = 1 + len xs

```

Which of these is a valid implementation of ~~len~~ len

- A. `len = foldr (\n -> n + 1) 0`
- B. `len = foldr (\n m -> n + m) 0`
- C. `len = foldr (\_ n -> n + 1) 0`
- D. `len = foldr (\x xs -> 1 + len xs) 0`
- E. All of the above

$$\begin{array}{c}
\underline{x_1} : \underline{x_2} : \left( \underline{x_3} : \left( \underline{x_4} : [] \right) \right) \\
\downarrow \\
\left( 1 + \left( 1 + \left( 1 + \left( 1 + 0 \right) \right) \right) \right) \\
\swarrow \quad \searrow \\
(\lambda x \underline{v} \rightarrow 1 + \underline{v})
\end{array}$$

## The Missing Parameter Revisited

We wrote `foldr` as

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr op base []      = base
foldr op base (x:xs) = op x (foldr op base xs)
```

but can also write this

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr op base  = go
where
  go []      = base
  go (x:xs) = op x (go xs)
```

Can someone explain where the `xs` went *missing* ?

## Trees

Recall the `Tree a` type from last time

```
data Tree a
  = Leaf
  | Node a (Tree a) (Tree a)
```

For example here's a tree

```
tree2 :: Tree Int
```

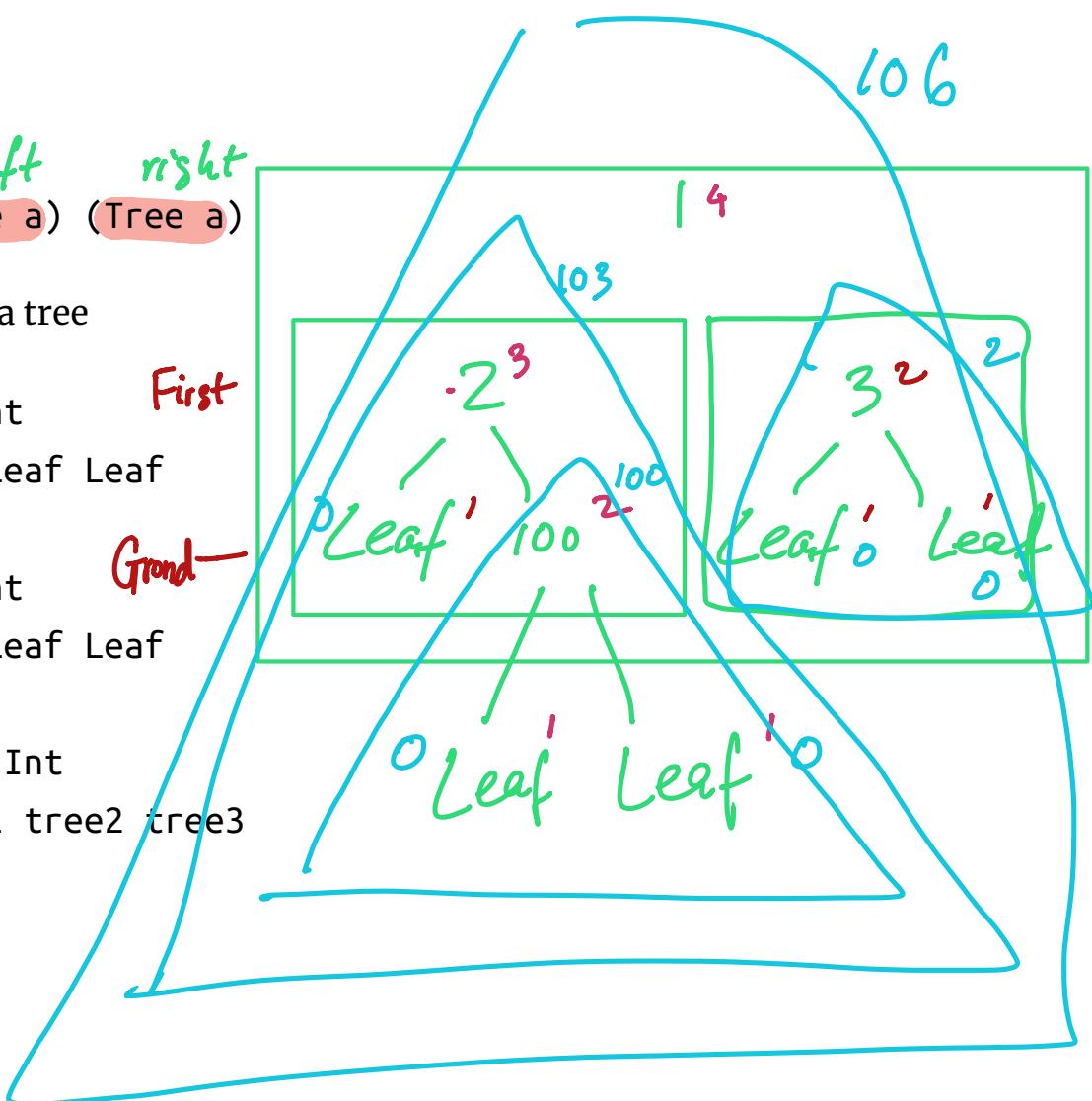
```
tree2 = Node 2 Leaf Leaf
```

```
tree3 :: Tree Int
```

```
tree3 = Node 3 Leaf Leaf
```

```
tree123 :: Tree Int
```

```
tree123 = Node 1 tree2 tree3
```



## Some Functions on Trees

Lets write a function to compute the `height` of a tree

```
height :: Tree a -> Int
height Leaf          = 0
height (Node x l r) = 1 + max (height l) (height r)
```

Here's another to *sum* the leaves of a tree:

```
sumTree :: Tree Int -> Int
sumTree Leaf          = ???
sumTree (Node x l r) = ???
```

Gathers all the elements that occur as leaves of the tree:

```
toList :: Tree a -> [a]
toList Leaf          = ???
toList (Node x l r) = ???
```

Lets give it a whirl

```
>>> height tree123
2
```

```
>>> sumTree tree123
6
```

```
>>> toList tree123
[1,2,3]
```

# Pattern: Tree Fold

Can you spot the pattern? Those three functions are almost the same!

**Step 1:** Rename to maximize similarity

```
-- height
foo Leaf      = 0
foo (Node x l r) = 1 + max (foo l) (foo r)
```

```
-- sumTree
foo Leaf      = 0
foo (Node x l r) = foo l + foo r
```

```
-- toList
foo Leaf      = []
foo (Node x l r) = x : foo l ++ foo r
```

**Step 2:** Identify the differences

1. ???
2. ???

**Step 3** Make *differences* a parameter

```
foo p1 p2 Leaf      = ???
foo p1 p2 (Node x l r) = ???
```

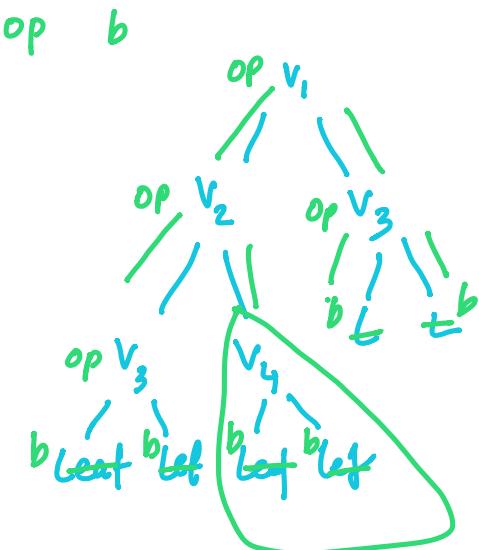
## Pattern: Folding on Trees

`tFold op b Leaf` = `b`

`tFold op b (Node x l r)` = `op x (tFold op b l) (tFold op b r)`

Lets try to work out the type of `tFold`!

`tFold :: t_op -> t_b -> Tree a -> t_out`



# QUIZ

Suppose that `t :: Tree Int`.

What does `tFold (\x y z -> y + z) 1 t` return?

- a. 0
- b. the *largest* element in the tree `t`
- c. the *height* of the tree `t`
- d. the *number-of-leaves* of the tree `t`
- e. type *error*

# EXERCISE

Write a function to compute the *largest* element in a tree or `0` if tree is empty or all negative.

```
treeMax :: Tree Int -> Int
treeMax t = tFold f b t
where
  f    = ???
  b    = ???
```

## *Map over Trees*

We can also write a `tmap` equivalent of `map` for `Trees`

```
treeMap :: (a -> b) -> Tree a -> Tree b
treeMap f (Leaf x)  = Leaf (f x)
treeMap f (Node l r) = Node (treeMap f l) (treeMap f r)
```

which gives

```
>>> treeMap (\n -> n * n) tree123      -- square all elements
of tree
Node 1 (Node 4 Leaf Leaf) (Node 9 Leaf Leaf)
```

## EXERCISE

Recursion is **HARD TO READ** do we really have to use it ?

Lets rewrite `treeMap` using `tFold` !

`treeMap :: (a -> b) -> Tree a -> Tree b`

`treeMap f t = tFold op base t`

**where**

`op = ???`

`base = ???`

When you are done, we should get

```
>>> animals = Node "cow" (Node "piglet" Leaf Leaf) (Leaf "hip  
po" Leaf Leaf)
```

```
>>> treeMap reverse animals
```

```
Node "woc" (Node "telgip" Leaf Leaf) (Leaf "oppih" Leaf Leaf)
```

**WED**

**NOV 4**

**NEW DEADLINE**  
for 01-TREES

## Examples: *foldDir*

```
data Dir a
  = Fil a          -- ^ A single file named `a`
  | Sub a [Dir a] -- ^ A sub-directory name `a` with contents `[Dir a]'
```

```
data DirElem a
  = SubDir a      -- ^ A single Sub-Directory named `a`
  | File a        -- ^ A single File named `a`
```

```
foldDir :: ([a] -> r -> DirElem a -> r) -> r -> Dir a -> r
foldDir f r0 dir = go [] r0 dir
```

**where**

```
go stk r (Fil a)    = f stk r (File a)
go stk r (Sub a ds) = L.foldl' (go stk') r' ds
```

**where**

```
r'    = f stk r (SubDir a)
stk' = a:stk
```

*foldDir* takes as input

- an *accumulator* *f* of type  $[a] \rightarrow r \rightarrow \text{DirElem } a \rightarrow r$ 
  - takes as *input* the path  $[a]$ , the current result  $r$ , the next  $\text{DirElem } [a]$
  - and returns as *output* the new result  $r$

- an *initial* value of the result `r0` and
- directory to fold over `dir`

And returns the result of running the *accumulator* over the whole `dir`.

## *Examples: Spotting Patterns In The “Real” World*

These patterns in “toy” functions appear regularly in “real” code

1. Start with beginner’s version riddled with explicit recursion (swizzle-v0.html).
2. Spot the patterns and eliminate recursion using HOFs (swizzle-v1.html).
3. Finally refactor the code to “swizzle” and “unswizzle” without duplication (swizzle-v2.html).

### Try it yourself

- Rewrite the code that swizzles `Char` to use the `Map k v` type in `Data.Map`

# Which is more readable? HOFs or Recursion

At first, recursive versions of `shout` and `squares` are easier to follow

- `fold` takes a bit of getting used to!

With practice, the *higher-order* versions become easier

- only have to understand specific operations
- recursion is lower-level & have to see “loop” structure
- worse, potential for making silly off-by-one errors

Indeed, HOFs were the basis of `map/reduce` and the big-data revolution  
(<http://en.wikipedia.org/wiki/MapReduce>)

- Can *parallelize* and *distribute* computation patterns just once  
([https://www.usenix.org/event/osdi04/tech/full\\_papers/dean/dean.pdf](https://www.usenix.org/event/osdi04/tech/full_papers/dean/dean.pdf))
- Reuse (<http://en.wikipedia.org/wiki/MapReduce>) across hundreds or thousands of instances!

HOFs FTW!

OF TRES DUE  
NOV 4

(<https://ucsd-cse230.github.io/fa20/feed.xml>)

(<https://twitter.com/ranjitjhala>)

(<https://plus.google.com/u/0/104385825850161331469>)

(<https://github.com/ranjitjhala>)

Generated by Hakyll (<http://jaspervdj.be/hakyll>), template by Armin Ronacher (<http://lucumr.pocoo.org>), suggest improvements here (<https://github.com/ucsd-progssys/liquidhaskell-blog/>).

01-trees due Friday 11/6

## Haskell Crash Course Part III

### *Writing Applications*

Lets write the classic “Hello world!” program.

For example, in Python you may write:

```
def main():
    print "hello, world!"

main()
```

and then you can run it:

```
$ python hello.py
hello world!
```

# Haskell is a *Pure* language.

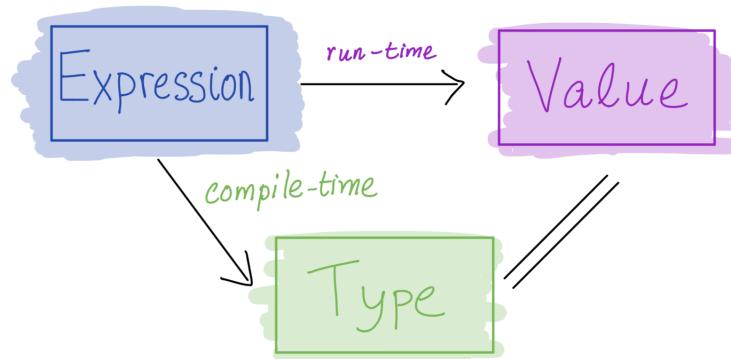
Not a *value* judgment, but a precise *technical* statement:

The “Immutability Principle”:

- A function must *always* return the same output for a given input
- A function’s behavior should *never change*

*foo :: In → Out*

## No Side Effects



Haskell’s most radical idea: *expression =\*> value*

- When you evaluate an expression you get a value and
- Nothing else happens

*“pure”*

Specifically, evaluation must not have an **side effects**

- change a global variable or
- print to screen or
- read a file or
- send an email or
- launch a missile.

OK  
in Python

anything OTHER THAN  
the OUTPUT VALUE

*But... how to write "Hello, world!"*

But, we want to ...

- print to screen
- read a file
- send an email

Thankfully, you can do all the above via a very clever idea: **Recipe**

# Recipes

This analogy is due to Joachim Breitner (<https://www.seas.upenn.edu/~cis194/fall16/lectures/06-io-and-monads.html>)

Haskell has a special type called `IO` – which you can think of as `Recipe`

`type Recipe a = IO a`

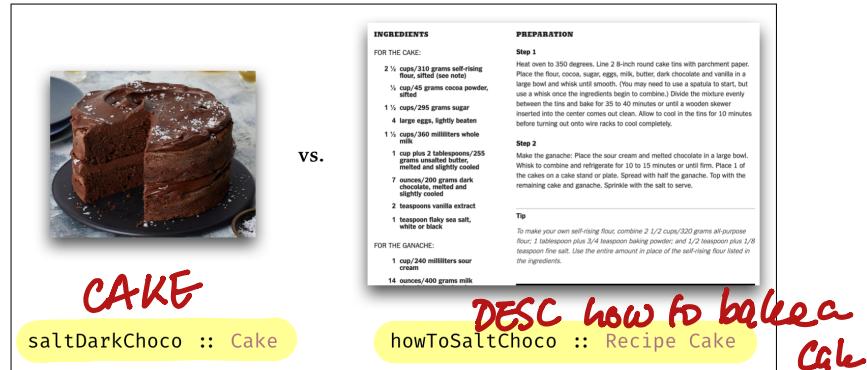
A `value` of type `Recipe a`

- is a `description` of a *computation* that can have *side-effects*
- which `when executed` performs some effectful I/O operations
- to produce a value of type `a`.

## Recipes have No Side Effects

A value of type `Recipe a` is

- A description of a computation that can have side-effects



Cake vs. Recipe

(L) chocolate *cake*, (R) a sequence of *instructions* on how to make a cake.

They are different (*Hint*: only one of them is delicious.)

Merely having a **Recipe Cake** has no effects! The recipe

- Does not make your oven *hot*
- Does not make your floor *dirty*

# Only One Way to Execute Recipes

Haskell looks for a special value

**main :: Recipe ()**

The value associated with `main` is handed to the **runtime system** and **executed**



Baker Aker

The Haskell runtime is a *master chef* who is the only one allowed to cook!

## *How to write an App in Haskell*

Make a `Recipe ()` that is handed off to the master chef `main`.

- `main` can be arbitrarily complicated
- composed of smaller sub-recipes

## *A Recipe to Print to Screen*

```
putStrLn :: String -> Recipe ()
```

The function `putStrLn`

- takes as input a `String`
- returns as output a `Recipe ()`

`putStrLn msg` is a `Recipe ()` – when executed prints out `msg` on the screen.

```
main :: Recipe ()  
main = putStrLn "Hello, world!"
```

... and we can compile and run it

```
$ ghc --make hello.hs  
$ ./hello  
Hello, world!
```

## QUIZ: How to Print Multiple Things?

Suppose I want to print two things e.g.

```
$ ghc --make hello.hs  
$ ./hello2  
Hello!  
World!
```

Can we try to compile and run this:

```
main = (putStrLn "Hello!", putStrLn "World!")
```

A. Yes!

*Recipe C)*

*Recipe C)*

B. No, there is a type error!

C. No, it compiles but produces a different result!

# A Collection of Recipes

Is just ... a *collection* of Recipes!

```
recPair :: (Recipe (), Recipe ())
recPair = (putStrLn "Hello!", putStrLn "World!")
```

```
recList :: [Recipe ()]
recList = [putStrLn "Hello!", putStrLn "World!"]
```

... we need a way to **combine** recipes!

## Combining? Just **do** it!

We can *combine* many recipes into a single one using a **do** block

```
foo :: Recipe a3
foo = do r1    -- r1 :: Recipe a1
        r2    -- r2 :: Recipe a2
        r3    -- r3 :: Recipe a3
```

(or if you *prefer* curly braces to indentation)

```
foo = do { r1;      -- r1 :: Recipe a1
          r2;      -- r2 :: Recipe a2
          r3 }     -- r3 :: Recipe a3
```

The **do** block combines sub-recipes `r1`, `r2` and `r3` into a *new* recipe that

- Will execute each sub-recipe in *sequence* and
- Return the value of type `a3` produced by the last recipe `r3`

## *Combining? Just **do** it!*

So we can write

```
main = do putStrLn "Hello!"
          putStrLn "World!"
```

or if you prefer

```
main = do { putStrLn "Hello!";
            putStrLn "World!"
          }
```

## EXERCISE: Combining Many Recipes

Write a function called `sequence` that

- Takes a *list* of recipes  $[r_1, \dots, r_n]$  as input and
- Returns a *single* recipe equivalent to `do {r1; ...; rn}`

`sequence :: [Recipe a] -> Recipe a`

`sequence rs = ???`

When you are done you should see the following behavior

-- *Hello.hs*

`main = sequence [putStrLn "Hello!", putStrLn "World!"]`

and then

\$ `ghc --make Hello.hs`

\$ `./hello`

Hello!

World!

# Using the Results of (Sub-) Recipes

Suppose we want a function that asks for the user's name

```
$ ./hello  
What is your name?  
Ranjit  
Hello Ranjit!
```

Annotations:

- A red arrow points from the text "Print / putStrLn" to the word "putStrLn" in the code.
- A green arrow points from the word "getline" to the function "getline" in the code.
- A red bracket labeled "# <<< user enters" spans the text "Ranjit".
- A red arrow points from the text "PutStrLn" to the word "putStrLn" in the code.

We can use the following sub-recipes

```
-- / read and return a line from stdin as String  
getLine :: Recipe String  
  
-- take a string s, return a recipe that prints s  
putStrLn :: String -> Recipe ()
```

But how to

- Combine the two sub-recipes while
- Passing the result of the first sub-recipe to the second.  
*'getline'*      *'putStrLn'*

## Naming Recipe Results via “Assignment”

You can write



```
x <- recipe
```

to *name* the result of executing `recipe`

- `x` can be used to refer to the result in later code

## Naming Recipe Results via “Assignment”

Lets, write a function that *asks* for the user's name

```
main = ask
```

```
ask :: Recipe ()  
ask = do name <- getLine;  
         putStrLn ("Hello " ++ name ++ "!")
```

Which produces the desired result

```
$ ./hello  
What is your name?  
Ranjit # user enters  
Hello Ranjit!
```

## EXERCISE

Modify the above code so that the program *repeatedly* asks for the users's name *until* they provide a *non-empty* string.

```
-- Hello.hs  
  
main = repeatAsk  
  
repeatAsk :: Recipe ()  
repeatAsk = _fill_this_in  
  
  
isEmpty :: String -> Bool  
isEmpty s = length s == 0
```

When you are done you should get the following behavior

```
$ ghc --make hello.hs
```

```
$ ./hello
```

What is your name?

# user hits return

What is your name?

# user hits return

What is your name?

# user hits return

What is your name?

Ranjit # user enters

Non-empty String

Hello Ranjit!

## EXERCISE

Modify your code to *also* print out a **count** in the prompt

```
$ ghc --make hello.hs
```

```
$ ./hello  
What is your name?  
What is your name?  
What is your name?  
What is your name?  
Ranjit  
Hello Ranjit!
```

(0) What is your name?  
(1) What is your name?  
(2) What is your name?  
(3) What is your name?  
Ranjit  
Hello Ranjit!

## *That's all about IO*

You should be able to implement build from `Directory.hs`

Using these library functions imported at the top of the file

```
import System.FilePath  (takeDirectory, takeFileName, (\/))  
import System.Directory (doesFileExist, listDirectory)
```

The functions are

- `takeDirectory`
- `takeFileName`
- `(\/)`
- `doesFileExist`
- `listDirectory`

Google the documentation to learn about how to use them.

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(<https://ucsd-cse230.github.io/fa20/feed.xml>)

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