

Let's add **print** and **definitions** and **calls** to our compiler

```
fn snek_print(val: i64) -> i64 {  
  if (val == 1) {  
    println!("false")  
  } else if (val == 3) {  
    println!("true")  
  } else {  
    println!("{}", val >> 1)  
  }  
  return val;  
}
```

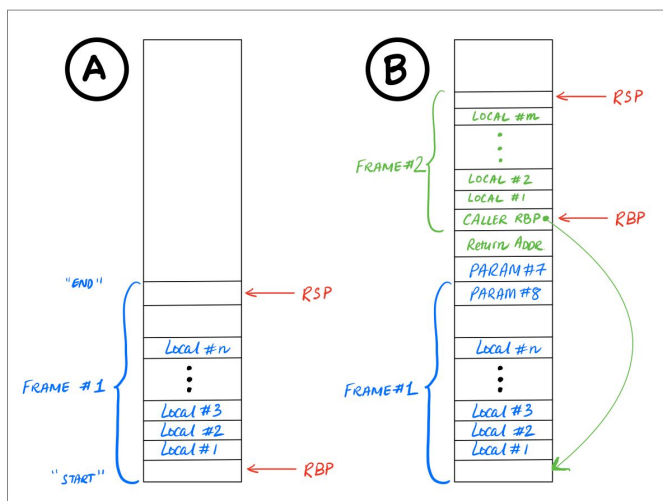
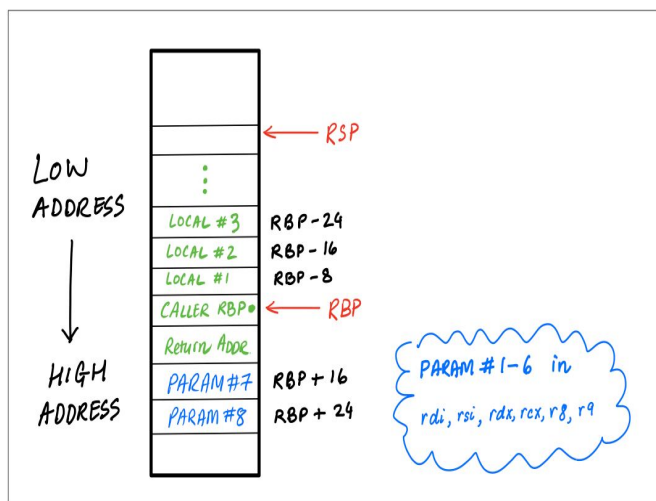
```
enum Expr {  
  ...  
  Print(Box<e>),  
}
```

## Example

## Generated code

```
(let (x (print input))  
  (+ x 1))
```

**print**



Let's add **print** and **definitions** and **calls** to our compiler

```
<prog> := <defn>+ <expr>

<defn> := (fun (<name> <name>) <expr>)
        | (fun (<name> <name> <name>) <expr>)

<expr> := ...
        | (<name> <expr>)
        | (<name> <expr> <expr>)
```

```
enum Defn {
  Fun1(String, String, Box<Expr>),
  Fun2(String, String, String, Box<Expr>),
}

enum Expr {
  ...
  Call1(String, Box<Expr>)
  Call2(String, Box<Expr>, Box<Expr>)
}
```

## Example

```
(fun (add x1 x2)
  (+ x1 x2))

(add input 10)
```

## Generated code

**def**

**call**

```

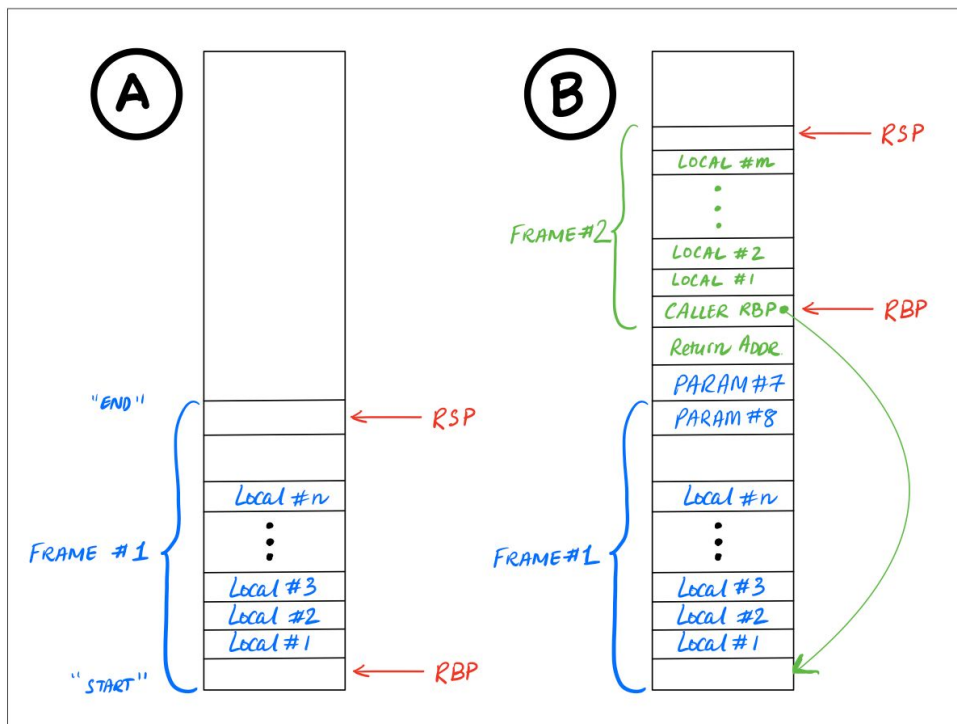
fn compile_def_body(args: &[String], sp: usize, body: &Expr, count: &mut i32) -> String
{
    let fun_entry = compile_entry(body, sp);
    let body_code = compile_expr(body, &init_env(args), sp, count, "time_to_exit");
    let fun_exit = compile_exit();

    format!("{fun_entry}
            {body_code}
            {fun_exit}")
}

fn compile_entry(e: &Expr, sp: usize) -> String {
    let vars = expr_vars(e) + sp;
    format!("push rbp
            mov rbp, rsp
            sub rsp, 8*{vars}")
}

fn compile_exit() -> String {
    format!("mov rsp, rbp
            pop rbp
            ret")
}

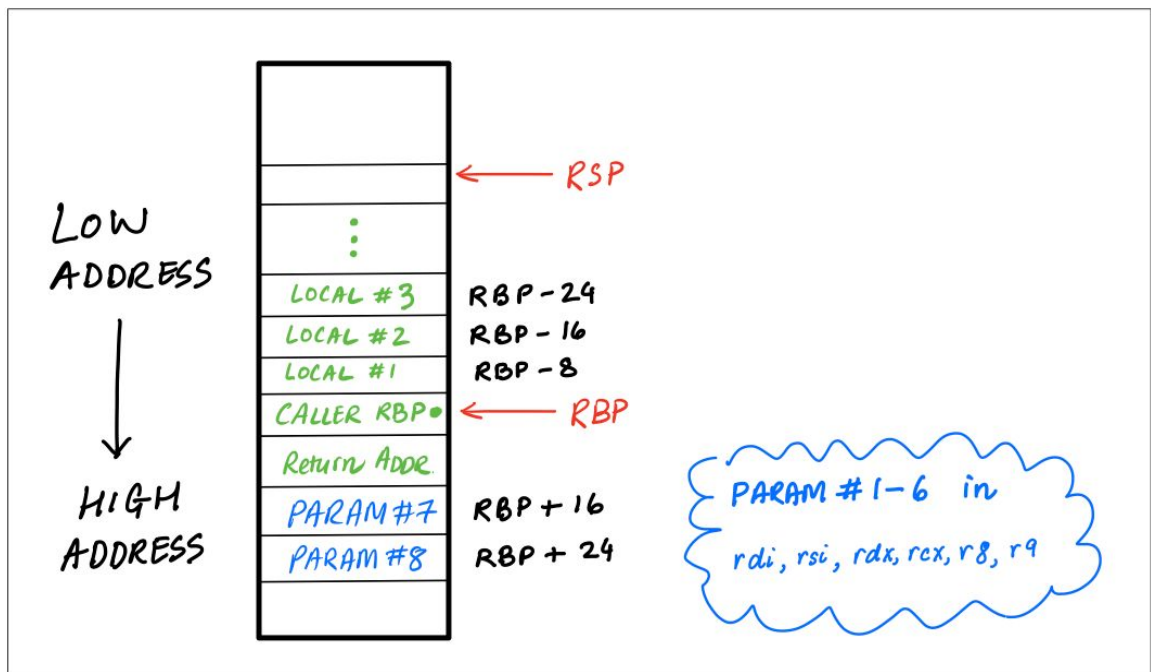
```



```

fn compile_expr(...) -> String {
  match e {
    ...
    Expr::Call2(f, e1, e2) => {
      let e1_code = compile_expr(e1, env, sp, count, brk);
      let e2_code = compile_expr(e2, env, sp + 1, count, brk);
      format!("{e1_code}
        mov [rbp - 8*{sp}], rax
        {e2_code}
        push rax
        mov rcx, [rbp - 8*{sp}]
        push rcx
        call fun_start_{f}
        add rsp, 8*2")
    }
  }
}

```



0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000

0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0101

[illegible][illegible]

### What does this mean for code generation?

*What should we do the next time we need a new type? (string, heap-allocated object, etc.)*

### Condition Codes (that matter for us): Overflow, Sign, Zero

many instructions set these; arithmetic, shifting, etc. mov does not

cmp <reg>, <val>      compute <reg> - <val> and set condition codes (value in <reg> does not change)  
some cases to think about:

<reg> = -2^64, <val> = 1      Overflow: \_\_\_\_\_ Sign: \_\_\_\_\_ Zero: \_\_\_\_\_

<reg> = 0, <val> = 1 Overflow: \_\_\_\_\_ Sign: \_\_\_\_\_ Zero: \_\_\_\_\_

```
<reg> = 1, <val> = 0 Overflow: ____ Sign: ____ Zero: ____
```

<reg> = -1, <val> = -2      Overflow: \_\_\_\_      Sign: \_\_\_\_      Zero: \_\_\_\_

`test <reg>, <val>` perform bitwise and on the two values, but don't change <reg>, and set condition codes as appropriate. Useful for mask checking. `test rax, 1` will set Z to true if and only if the LSB is 1

**<label>:**        *set this line as a label for jumping to later*

`jmp <label>`                *unconditionally jump to <label>*

`line <label>`                      *jump to <label> if Zero is not set (last cmped values not equal)*

**je <label>**      *jump to <label> if Zero is set (last cmped values are equal)*

*jump to <label> if Overflow is the same as Sign (which corresponds to  $\geq$  for last cmp)*

`je <label>`      *jump to <label> if Zero set or Overflow != Sign (which corresponds to <= for last cmp)*

**shl <reg>**      *shift <reg> to the left by 1, filling in least-significant bit with zero*

**sar <reg>**      *shift <reg> to the right by 1, filling in most-significant bit to preserve sign*

`shr <reg>`      *shift <reg> to the right by 1, filling in most-significant bit with zero*