# Lecture 24: Processes

CSE 29: Systems Programming and Software Tools

Aaron Schulman (Shalev)



### Process: A running program on an OS



- □ When you run a program on an OS (e.g., start in shell), it starts a "process"
- A process is an abstraction provided by the OS of a single isolated running program
- A process includes all of the context of the program
  - Memory: Stack/Heap and Code
  - Registers: Current state of execution (e.g., what instruction in the code is running)
  - Peripherals: Even disk access!
- Processes are completely isolated from each other:
  - You can not overwrite another processes' stack/heap or code!

## Process State: What does an OS maintain?



- Process ID: PID A unique number for a process
- The address space (range) in RAM for the process
- The execution state of the process (i.e., Registers)
- The hardware resources in use (e.g., open files on disk)

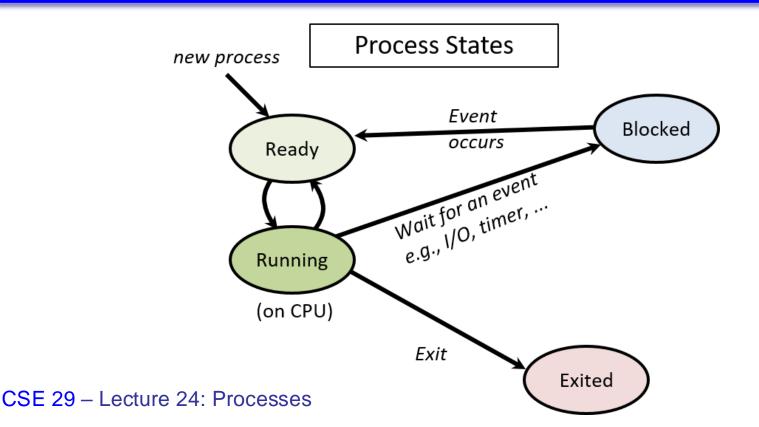
## Context Switching: Changing processes



- The Operating System decides when to switch between processes
  - This is known as "Context Switching"
  - The this is one of the big benefits of the process abstraction
- When an OS decides that a process had its time to execute it performs the context switch
  - The OS saves all context of the running process by copying all Registers to RAM
  - The OS then copies the saved context of another process from RAM to the Registers
  - The new context switched process begins!





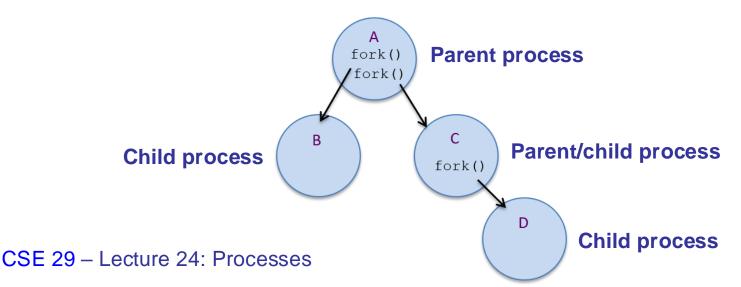






int fork(); Create a new process running the current code

- Creates a new process (new state for address space, PID, registers, hardware)
- Starts with a copy of the same memory of the running process
  - » Process starts with a *copy* of the stack, heap, and code



#### Fork example



```
pid_t pid;
pid = fork(); /* create a new child process
print("pid = %d\n", pid); /* both parent and child execute this */
                                                 Parent (pid 12)
   Parent: ____ B fork _____B
                                                  pid=fork()
                                                   pid:14
                                                              Child (pid 14)
                                                           Exact copy
                                                            of parent
                         Time
                                                            pid:0
```

CSE 29 – Lecture 24: Processes





int wait(pid t pid); Wait for the child process pid to end

- This is called only by the parent process
- NULL will wait for all children to finish





int execvp(char \*filename, char \*argv[]);

- Loads new code from a binary file (filename) into the running process
- Starts running the main() function in that code
- This function never returns to the calling processes' codebase

