

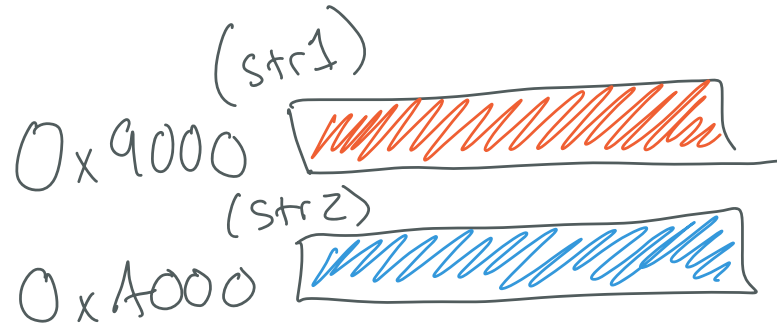
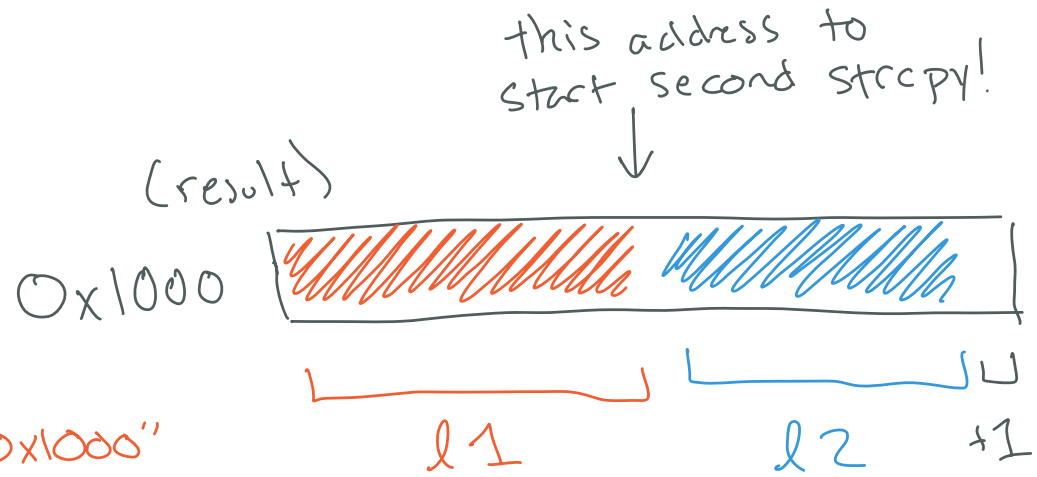
```
char* concat(char str1[], char str2[]) {
    int l1 = strlen(str1);
    int l2 = strlen(str2);
    char* result = malloc(l1 + l2 + 1);
}
```

`strcpy(result, str1)` "copy from  
0x9000 to 0x1000"

`strcpy(&result[l1], str2)` "copy from  
0x1000 to 0x1000 + l1 bytes"

`strcpy(result + l1, str2)`

both of the 1<sup>st</sup> args to  
strcpy compute "0x1000 + l1"



Imagine an example like this for `int32_t` data.

`&result[index]` must multiply index by 4 to get to  
the right address

`result + index` must also multiply by 4 (gcc does  
this!)

concat() from Tuesday:

```
char* concat(char str1[], char str2[]) {
    int l1 = strlen(str1);
    int l2 = strlen(str2);
    char* result = malloc(l1 + l2 + 1);
    for(int i = 0; i < l1; i += 1) {
        result[i] = str1[i];
    }
    for(int j = l1; j < l1 + l2; j += 1) {
        result[j] = str2[j - l1];
    }
    result[l1 + l2] = '\0';
    return result;
}
```

→ replace with strcpy(result, str1)

→ replace with strcpy(result, str2) X

would overwrite the beginning of result again

strcpy(result[l1], str2) X  
what about this option?

expected char\*, got char

## strcpy

<cstring>

```
char * strcpy ( char * destination, const char * source );
```

### Copy string

Copies the C string pointed by **source** into the array pointed by **destination**, including the terminating null character (and stopping at that point).

To avoid overflows, the size of the array pointed by **destination** shall be long enough to contain the same C string as **source** (including the terminating null character), and should not overlap in memory with **source**.

strcpy(&result[l1], str2)

strcpy(result + l1, str2)

Pointer arithmetic

Compute the address l1 bytes after result

← shift - 7

& expr

&

"ampersand"

How can result + 11 when they are two different type?

Will it give an error?

💬 😊 ...

Pointer Arithmetic

Not an error

$\text{ptr} + n$  where  $\text{ptr}$  is a pointer  $T^*$   
 $n$  is an integer

compute the address  $n * (\text{sizeof}(T))$  bytes  
from  $\text{ptr}$

$\text{ptr}$  of type  $T^*$

$\text{ptr}[\text{index}]$  access  $\text{sizeof}(T)$  bytes of memory  
at  $(\text{index} * \text{sizeof}(T))$  bytes after  $\text{ptr}$

## Structs in C

makes "Point" an abbrev for "struct Point"

"technically" the outline is a struct definition

```
typedef struct Point {  
    int x;  
    int y;  
} Point;
```

This is a struct declaration, typically at top level of file.

Look up  
"struct packing"  
online about  
how things stored  
in order in  
structs

// inside a function

```
Point p = {4, 5};
```

```
Point p2 = {22, 777};
```

variable declarations of  
a struct type allocate  
stack space for the struct

// field access (member access)

p.x

p.y

p2.x

p2.y

// field update

p.x = v

p.y = v

p2.x = v

p2.y = v

```
Point make-Point(int x, int y) { ... }
```

Struct variable decls. create space on the stack for that struct

```
typedef struct Point {
    int x, y;
} Point;

void example1() {
    Point p1 = { 4, 5 };
    Point p2 = { 200, 900 };
    printf("p1: %d, %d\tp2: %d %d\n", p1.x, p1.y, p2.x, p2.y);

    Point p3 = p2;
    p3.x = 777;
    printf("p2: %d, %d\tp3: %d %d\n", p2.x, p2.y, p3.x, p3.y);
}
```

Variable/Role

Address

Data

0/8 1/9 2/A 3/B 4/C 5/D 6/E 7/F

0x...00

0x...08

0x...10

0x...18

0x...20

0x...28

0x...30

0x...38

0x...40

0x...48

0x...50

0x...58

0x...60

0x...68

0x...70

0x...78

0x...80

0x...88

0x...90

0x...98

0x...A0

0x...A8

0x...B0

0x...B8

0x...C0

0x...C8

0x...D0

0x...D8

0x...E0

0x...E8

0x...F0

0x...F8

200 777 900  
200 900  
4 5

p3  
p2  
p1

```
[jpolitz@ieng6-203]:ss1-25-06-w3r-string-list:372$ ./point
p1: 4, 5      p2: 200 900
p2: 200, 900  p3: 777 900
```

```
Point make_Point(int x, int y) {
    Point p = { x, y };
    return p;
}

void example2() {
    Point p1 = make_Point(4, 5);
    Point p2 = make_Point(200, 900);
    printf("p1: %d, %d\tp2: %d %d\n", p1.x, p1.y, p2.x, p2.y);
}
```

make-Point

y  
x  
p

```
[jpolitz@ieng6-203]:ss1-25-06-w3r-string-list:374$ ./point
p1: 4, 5      p2: 200 900
```

Returning or assigning a struct main copies its members' values.

p2  
p1

5 900  
4 200  
4 200 5 900  
4 200 5 900

```

void update_X(Point p, int x) {
    p.x = x;
}
void example3() {
    Point p1 = { 4, 5 };
    update_X(p1, 333);
    printf("p1.x: %d\n", p1.x);
}

```

\$ ./point  
p1.x: 4

update-X | x  
          | P  
  
main | p1

When passing struct values as arguments, members values are copied

if update\_x were to return p, would p.x then be updated to 333?

If the only change was

```

Point update_X(Point p, int x) {
    p.x = x;
    return p;
}

```

This has no effect on p1

If we also changed main to  
p1 = update\_X(p1, 333)

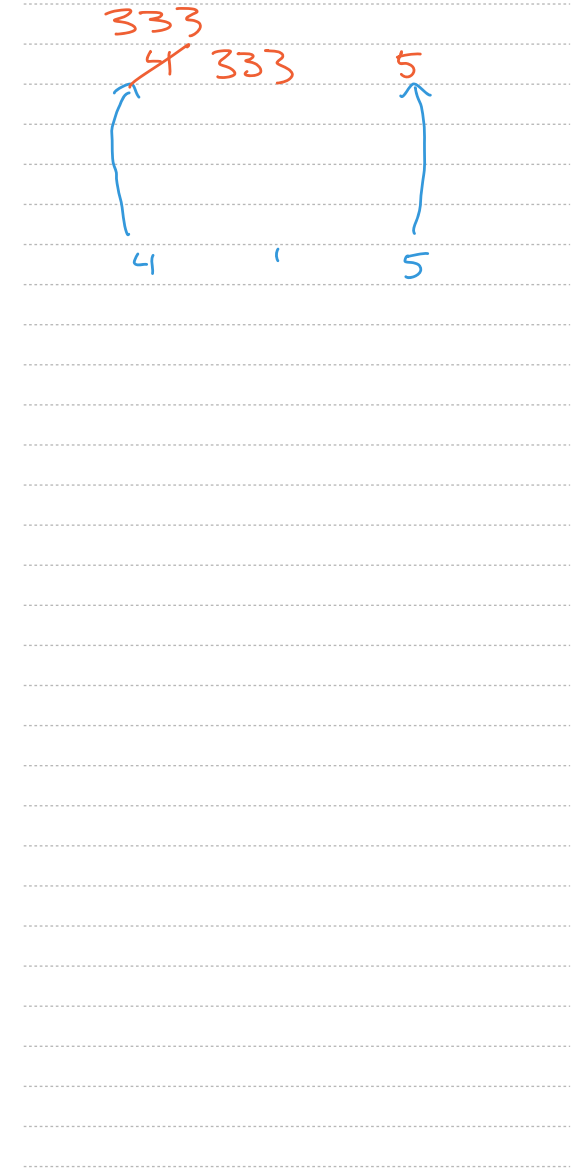
This does change p1 by copying return via assignment

Address

0x...00  
0x...08  
0x...10  
0x...18  
0x...20  
0x...28  
0x...30  
0x...38  
0x...40  
0x...48  
0x...50  
0x...58  
0x...60  
0x...68  
0x...70  
0x...78  
0x...80  
0x...88  
0x...90  
0x...98  
0x...A0  
0x...A8  
0x...B0  
0x...B8  
0x...C0  
0x...C8  
0x...D0  
0x...D8  
0x...E0  
0x...E8  
0x...F0  
0x...F8

Data

0/8 1/9 2/A 3/B 4/C 5/D 6/E 7/F



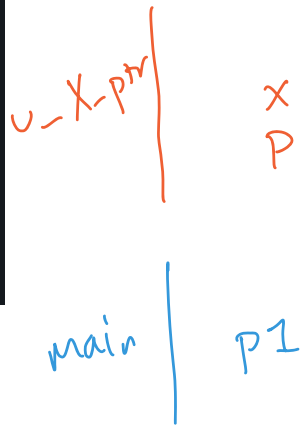
```

void update_X_ptr(Point* p, int x) {
    p->x = x;
}
void example4() {
    Point p1 = { 4, 5 };
    update_X_ptr(&p1, 444);
    printf("p1.x: %d\n", p1.x);
}

```

\$ ./point  
p1.x: 444

$p \rightarrow x$   
means at offset of x  
member from address stored in p



Variable/Role	Address	Data							
		0/8	1/9	2/A	3/B	4/C	5/D	6/E	7/F
	0x...88								
	0x...90								
	0x...98								
	0x...A0								
	0x...A8								
	0x...B0								
	0x...B8								
	0x...C0								
	0x...C8								
	0x...D0								
	0x...D8								
	0x...E0								
	0x...E8								
	0x...F0								
	0x...F8								
	0x...00								
	0x...08								
	0x...10								
	0x...18								
	0x...20								
	0x...28								
	0x...30								
	0x...38								
	0x...40								
	0x...48								
	0x...50								
	0x...58								
	0x...60								
	0x...68								
	0x...70								
	0x...78								
	0x...80								

what value?  
444  
0x...D0

4 444 5

```
void update_X_ptr(Point* p, int x) {
    p->x = x;
}
```

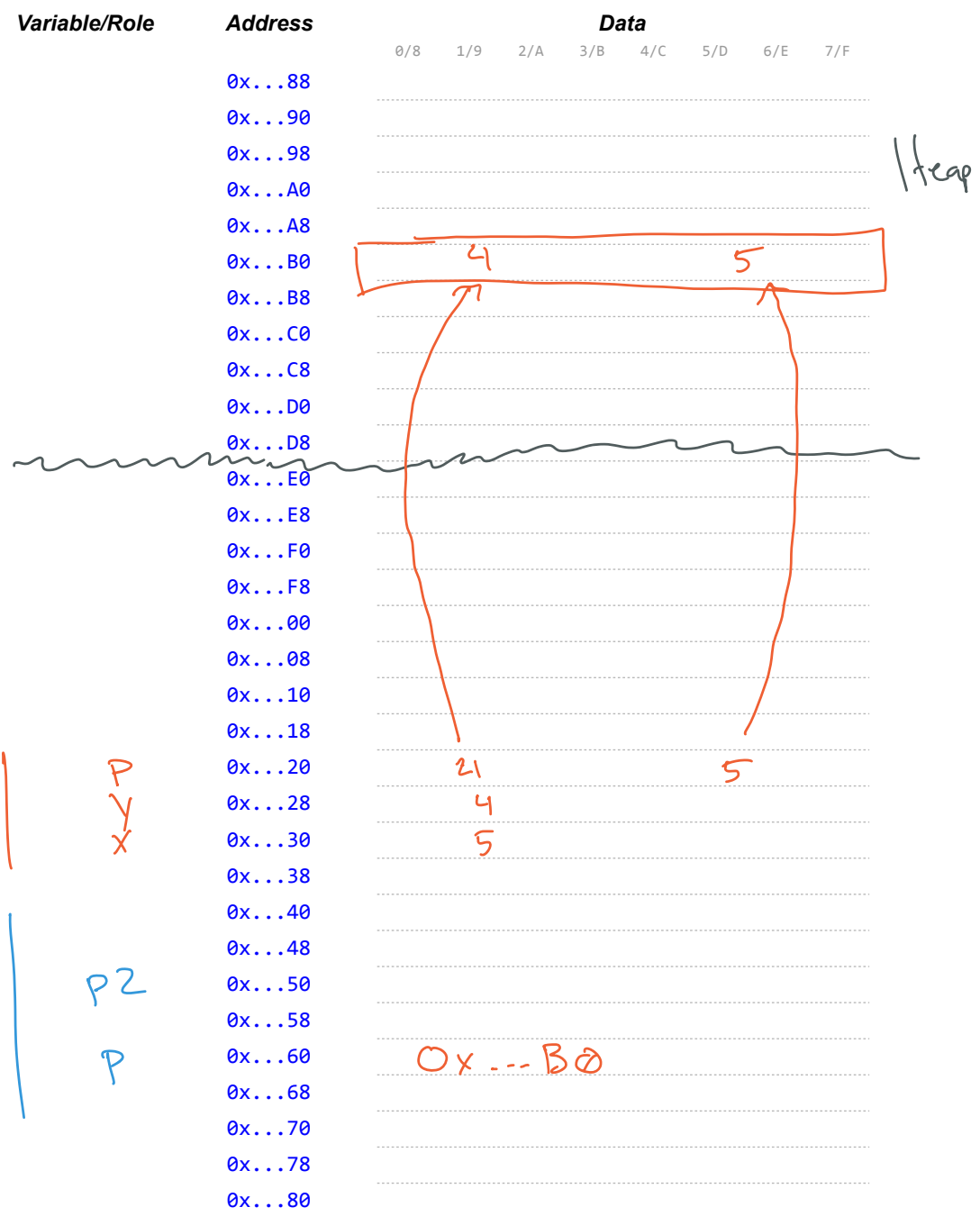
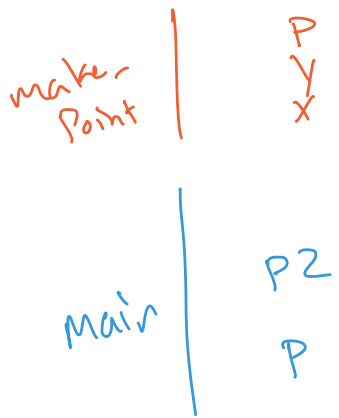
```
void example5() {
    printf("sizeof(Point): %ld\n", sizeof(Point));
    Point* p = malloc(sizeof(Point));
    *p = make_Point(4, 5);
    update_X_ptr(p, 555);
    printf("p->x: %d\n", p->x);

    Point* p2 = p;
    update_X_ptr(p2, 888);
    printf("p->x: %d, p2->x: %d\n", p->x, p2->x);
}
```

```
sizeof(Point): 8
p->x: 555
p->x: 888, p2->x: 888
```

\*p = val

assign into memory at  
the address stored in p  
the value val





```
void update_X_ptr(Point* p, int x) {
    p->x = x;
}
```

```
void example5() {
    printf("sizeof(Point): %ld\n", sizeof(Point));
    Point* p = malloc(sizeof(Point));
    *p = make_Point(4, 5);
    update_X_ptr(p, 555);
    printf("p->x: %d\n", p->x);

    Point* p2 = p;
    update_X_ptr(p2, 888);
    printf("p->x: %d, p2->x: %d\n", p->x, p2->x);
}
```

```
sizeof(Point): 8
p->x: 555
p->x: 888, p2->x: 888
```

\*p = val  
 assign into memory at  
 the address stored in p  
 the value val

update\_X\_ptr | x  
 | P  
 example5 | P2  
 | P

this assignment  
 just copies  
 address in p  
 into p2

Variable/Role	Address	Data
	0x...88	0/8 1/9 2/A 3/B 4/C 5/D 6/E 7/F
	0x...90	
	0x...98	
	0x...A0	
	0x...A8	
	0x...B0	4 555 5
	0x...B8	
	0x...C0	
	0x...C8	
	0x...D0	
	0x...D8	
	0x...E0	
	0x...E8	
	0x...F0	
	0x...F8	
	0x...00	
	0x...08	
	0x...10	
	0x...18	
	0x...20	555
	0x...28	
	0x...30	0x...B0
	0x...38	
	0x...40	
	0x...48	
	0x...50	0x...B0
	0x...58	
	0x...60	0x...B0
	0x...68	
	0x...70	
	0x...78	
	0x...80	

Heap

```
void update_X_ptr(Point* p, int x) {
    p->x = x;
}
```

```
void example5() {
    printf("sizeof(Point): %ld\n", sizeof(Point));
    Point* p = malloc(sizeof(Point));
    *p = make_Point(4, 5);
    update_X_ptr(p, 555);
    printf("p->x: %d\n", p->x);

    Point* p2 = p;
    update_X_ptr(p2, 888);
    printf("p->x: %d, p2->x: %d\n", p->x, p2->x);
}
```

```
sizeof(Point): 8
p->x: 555
p->x: 888, p2->x: 888
```

\*p = val

assign into memory at  
the address stored in p  
the value val

p and p2  
refer to same  
heap-allocated struct

update\_X\_ptr

x

P

example5

P2

P

this assignment  
just copies  
address in p  
into p2

Variable/Role

Address

Data

0/8 1/9 2/A 3/B 4/C 5/D 6/E 7/F

0x...88

0x...90

0x...98

0x...A0

0x...A8

0x...B0

0x...B8

0x...C0

0x...C8

0x...D0

0x...D8

0x...E0

0x...E8

0x...F0

0x...F8

0x...00

0x...08

0x...10

0x...18

0x...20

0x...28

0x...30

0x...38

0x...40

0x...48

0x...50

0x...58

0x...60

0x...68

0x...70

0x...78

0x...80

4 555 888 5

888

0x...B0

0x...B0

0x...B0

Heap

# 15 MIN BREAK RESUME 12:45

$\text{ptr} + n$  "pointer arithmetic" add  $n * \text{sizeof}(T)$  to  $T^* \text{ ptr}$   
Computes new address of type  $T^*$

$\text{ptr}[n]$  "indexing" look up  $\text{sizeof}(T)$  bytes of memory at  
offset  $n * \text{sizeof}(T)$  from  $\text{ptr}$   
Returns value of type  $T$

$\& x$  "addressof" computes address of variable  $x$   
For  $x$  of type  $T$ , returns  $T^*$

$\& \text{ptr}[n]$  "address of" same meaning as  $\text{ptr} + n$

$* \text{ptr}$  "dereference" Return  $\text{sizeof}(T)$  bytes of memory  
at address stored in  $\text{ptr}$  (Return type  $T$ )  
of type  $T^*$

$* \text{ptr} = \text{val}$   
 $\text{ptr}[n] = \text{val}$  "assignment" Compute addresses as above, but  
store  $\text{val}$  there rather than look up  
 $\text{ptr} : T^*$   $\text{val} : T$

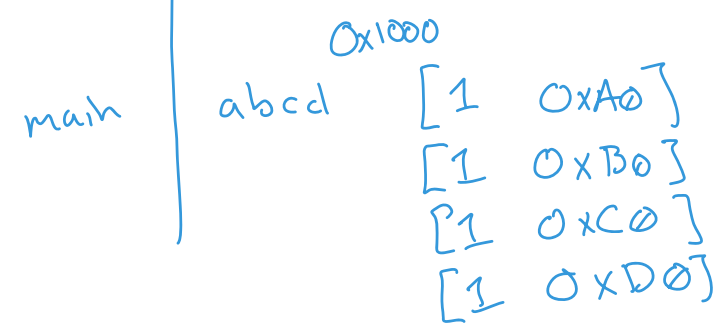
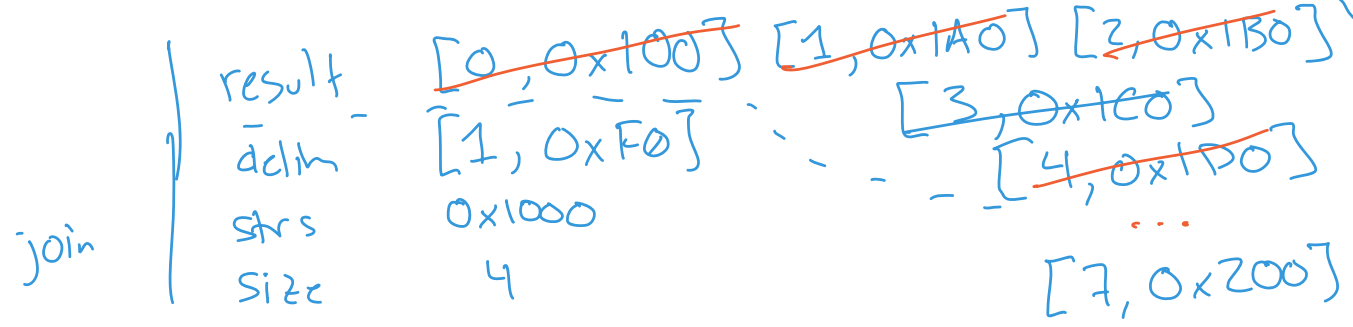
```
Str join(Str delim, Str strs[], int size) {
    Str result = str("");
    // high-level strategy: use concat() in a for loop
    for(int i = 0; i < size; i += 1) {
        result = concat(result, strs[i]);
        if(i < size - 1) {
            result = concat(result, delim);
        }
    }
    return result;
}

Str abcd[] = { str("a"), str("b"), str("c"), str("d") };
Str abcd_result = join(str("-"), abcd, 4);
printf("Expect a-b-c-d: %s %d\n", abcd_result.data, abcd_result.bytes)
```

Heap

- 0xA0 "a\0"
- 0xB0 "b\0"
- 0xC0 "c\0"
- 0xD0 "d\0"
- 0xF0 "-\0"
- 0x100 "\0"
- 0x1A0 "a\0"
- 0x1B0 "a-\0"
- 0x1C0 "a-b\0"
- 0x1D0 "a-b-\0"
- ...
- 0x200 "a-b-c-d\0"

made from concat



"garbage"

"memory leak"

none of these heap-allocated values are usable anymore!

This is why we have free(ptr) takes a malloc'd ptr and tells malloc the space can be re-used

As a programmer, find moments right before a heap-allocated value becomes unreachable, unusable, or otherwise not accessed again, and free at that point.

RUST (proglang)

```
[jpolitz@ieng6-203]:ss1-25-06-w3r-string-list:398$ valgrind ./str
==1193243== Memcheck, a memory error detector
==1193243== Copyright (C) 2002-2017, and GNU GPL'd, by Julian Seward et al.
==1193243== Using Valgrind-3.18.1 and LibVEX; rerun with -h for copyright info
==1193243== Command: ./str
==1193243==
abcdef 6
Should be hello,world: hello,world 11
Expect a-b-c-d: a-b-c-d 7
==1193243==
==1193243== HEAP SUMMARY:
==1193243==    in use at exit: 101 bytes in 23 blocks
==1193243== total heap usage: 24 allocs, 1 frees, 1,125 bytes allocated
==1193243==
==1193243== LEAK SUMMARY:
==1193243==    definitely lost: 101 bytes in 23 blocks
==1193243==    indirectly lost: 0 bytes in 0 blocks
==1193243==    possibly lost: 0 bytes in 0 blocks
==1193243==    still reachable: 0 bytes in 0 blocks
==1193243==    suppressed: 0 bytes in 0 blocks
==1193243== Rerun with --leak-check=full to see details of leaked memory
==1193243==
==1193243== For lists of detected and suppressed errors, rerun with: -s
==1193243== ERROR SUMMARY: 0 errors from 0 contexts (suppressed: 0 from 0)
```

```
// How to write a test for join()?
// join(str(", "), {str("hello"), str("world")}) -> str("hello, world")
// join(str(", "), {str("a"), str("b"), str("c")}) -> str("a,b,c")
Str strs[] = { str("hello"), str("world") };
Str result2 = join(str(", "), strs, 2);
printf("Should be hello,world: %s %d\n", result2.data, result2.bytes);

Str abcd[] = { str("a"), str("b"), str("c"), str("d") };
Str abcd_result = join(str("-"), abcd, 4);
printf("Expect a-b-c-d: %s %d\n", abcd_result.data, abcd_result.bytes);
```

"block" = malloc  
3 blocks

1  
"hello"  
2 3 4 5 6 7  
"hello, world"  
8 9 10 11 12 13 14 15  
"hello, world"

14 bytes

7 blocks

"\0"  
"a\0"  
"a-\0"  
"a-b\0"  
"a-b-\0"  
"a-b-c\0"  
"a-b-c-\0"

28 bytes

with a infinite loop with malloc and etc., and without freeing any of the heap memory, could we technically find the storage limit of the heap with valgrind?

what happens if you run out of space? i.e your memory leaks are larger than the space you have

malloc will return NULL!

← just count + check  
when malloc() = NULL

does free only deletes the data in malloc?

free tells malloc that space can be re-used

free can only be used on a ptr that was returned from malloc



Variable/Role	Address	Data							
		0/8	1/9	2/A	3/B	4/C	5/D	6/E	7/F
	0x...88								
	0x...90								
	0x...98								
	0x...A0								
	0x...A8								
	0x...B0								
	0x...B8								
	0x...C0								
	0x...C8								
	0x...D0								
	0x...D8								
	0x...E0								
	0x...E8								
	0x...F0								
	0x...F8								
	0x...00								
	0x...08								
	0x...10								
	0x...18								
	0x...20								
	0x...28								
	0x...30								
	0x...38								
	0x...40								
	0x...48								
	0x...50								
	0x...58								
	0x...60								
	0x...68								
	0x...70								
	0x...78								
	0x...80								