# **Assignment: Design and Analysis of Algorithms**

Due Date: July 1 2024

## **Program 1: Optimizing Delivery Routes (Case study)**

Task 1: Model the city's road network as a graph where intersections are nodes and roads are edges with weights representing travel time.

**Aim**: To create a structured model of the city's road network using graph theory. This allows for efficient route planning, optimization of traffic flow, and informed decision-making in urban planning. The goal is to improve transportation efficiency, reduce congestion, and enhance overall urban mobility and safety.

#### **Procedure:**

## 1. Graph Representation:

Define the city's road network as a dictionary of dictionaries (road\_network).

#### 2. Initialization:

• Initialize a priority queue (min-heap) to keep track of nodes to explore, starting with the source node (start).

#### 3. Start Node:

• Start from the specified source node (start) and initialize its distance as 0 in shortest\_paths.

## 4. Priority Queue Handling:

Repeat until all nodes have been processed or the destination node (goal) is reached

#### 5. Path Reconstruction:

 Once the destination node (goal) is reached or all nodes have been processed, reconstruct the shortest path from goal back to start using the shortest\_paths dictionary.

```
Time Complexity:

** Initalization O(1)

** White loop:

-> Visited: O(1)

-> Iterating over neighbors: O(E) per node, where E is

the number of edges.

-> Updating ishortest path: O(1) per neighbors

o(V) per iteration, where V is the number of vertice.

Time Complexity:- O(V2+VE) = O(V2)

oBSUMING. E = V2.

Space Complexity:-

-> Grouph Reposesentation: O(V4E)

-> Shortest paths Dictionary: O(V)

-> Visited set O(V)

= O(V+E)
```

#### Pseudocode:

```
function Dijkstra (graph, start, goal)
    sq. <- priority queue containing (0, start)
    shortest_paths <- dictionary with key start and value (None, 0)
    visited <- empty set
    while pq is not empty
        current_distance, current_node <- pq.pop()
        if current_node in visited
        continue
        visited.add (current_node)
    if current_node == goal</pre>
```

```
break
    for next node, weight in graph[current node]
       if next_node in visited
         continue
       new weight <- current distance + weight
       if new weight < shortest paths.get (next node, (None, infinity))[1]
         shortest_paths[2next_node] <- (current_node, new_weight)</pre>
         pq.push ((new_weight, next_node))
  if goal not in shortest paths
    return "Route Not Possible"
  path <- empty list
  current node <- goal
Program:
import heapq
road_network = {
  'A': {'B': 5, 'C': 7},
  'B': {'A': 5, 'C': 3, 'D': 4},
  'C': {'A': 7, 'B': 3, 'D': 6},
  'D': {'B': 4, 'C': 6}
}
def Dijkstra (graph, start, goal):
  shortest_paths = {start: (None, 0)}
  current node = start
  visited = set ()
  while current node != goal:
    visited.add(current_node)
```

```
destinations = graph[current node].items()
    for next node, weight in destinations:
      if next_node in visited:
        continue
      new weight = shortest paths[current node][1] + weight
      if shortest_paths.get(next_node, (None, float('inf')))[1] > new_weight:
        shortest_paths[next_node] = (current_node, new_weight)
    next destinations = {node: shortest paths[node] for node in
shortest paths if node not in visited}
    if not next destinations:
      return "Route Not Possible"
    current_node = min (next_destinations, key=lambda k:
next destinations[k][1])
  path = []
  while current_node is not None:
    path.append(current node)
    next_node = shortest_paths[current_node][0]
    current node = next node
  path = path[::-1]
  return path
start = 'A'
goal = 'D'
shortest path = Dijkstra (road network, start, goal)
if shortest path == "Route Not Possible":
  print ("No route found!")
else:
```

```
print (f"Shortest path from {start} to {goal}: {shortest_path}")

total_weight = sum(road_network[shortest_path[i]][shortest_path[i + 1]] for
i in range(len (shortest_path) - 1))

print(f"Total travel time: {total_weight} units")
```

## **Output:**

```
Shortest path from A to D: ['A', 'B', 'D']
Total travel time: 9 units
```

Time complexity : O((V+E)logV)

**Space complexity: O(V+E)** 

**Result:** The program executed successfully.

**Task 2:** Implement Dijkstra's algorithm to find the shorted paths from a central warehouse to various delivery location.

**Aim:** implementing Dijkstra's algorithm is to find the shortest paths from a central warehouse to delivery locations, optimizing logistics by minimizing travel distances or times. This facilitates efficient resource allocation and timely deliveries, enhancing overall operational efficiency in distribution networks.

#### **Procedure:**

**Initialize Data Structures:** Create a priority queue (pq) to store nodes with their current shortest distance estimates. Start with the warehouse node initialized to distance 0.

- **1.Initialize Variables:** set visited as an empty set to keep track of nodes that have been fully processed.
- **2. Main Loop:** While pq is not empty: Extract the node with the smallest distance (current\_node) from pq.
- **3.Check Visited Status:** If current\_node is in visited, continue to the next iteration of the loop.
- **4.Termination Check:** If the goal node (or all delivery locations) has been fully processed (i.e., added to visited), exit the loop.

```
Time Complexity: -
   -> Priority Queue Operations: Using a party queue, each instron
      and extraction operation takes O(logV) times
  -> Edge Relaxion: Each edge is evelaxed at most once. Relieution
     involves updating the posiority queue, which also take
         O(log V) Ames
     Thus, the total time complexity is
           O((V+E) log V)
           * v is the number of vertices (nodes).
           * E is the number of edges
Space Complexity:
-> Graph storage: the graph itself sieguine O(V+E) space.
-) focionity Queue: The porority queue can contain up to V nodes
       at once, thus suggesting O(v) space.
       thus, the total space Complexity is:
                 OCVED
```

#### **Pseudo Code:**

```
function Dijkstra (graph, start, goal):
    priority_queue pq
    shortest_paths = {}
    shortest_paths[start] = (None, 0)
    visited = set()
    while pq is not empty:
        current_node = extract_min (pq)
```

```
if current node in visited:
      continue
    visited.add(current node)
    for each neighbor, weight in graph[current node].neighbors():
      if neighbor in visited:
        continue
      new_distance = shortest_paths[current_node].distance + weight
      if neighbor not in shortest paths or new distance <
shortest_paths[neighbor].distance:
        shortest paths[neighbor] = (current node, new distance)
        pq.insert_or_update(neighbor, new_distance)
    path = []
  current_node = goal
  while current node is not None:
    path.add(current_node)
    current_node = shortest_paths[current_node].predecessor
  path.reverse()
  return path
Program:
import heapq
def Dijkstra (graph, start):
  pq = [(0, start)]
  shortest paths = {start: (None, 0)}
    while pq:
    current_distance, current_node = heapq.heappop (pq)
```

```
for next node, weight in graph[current node].items():
      new distance = current distance + weight
      if new distance < shortest_paths.get (next_node, (None,
float('inf')))[1]:
         shortest_paths[next_node] = (current_node, new_distance)
         heapq.heappush(pq, (new_distance, next_node))
  return shortest_paths
road_network = {
  'Warehouse': {'A': 5, 'B': 7, 'C': 9},
  'A': {'Warehouse': 5, 'D': 3, 'E': 8},
  'B': {'Warehouse': 7, 'E': 4},
  'C': {'Warehouse': 9, 'D': 2},
  'D': {'A': 3, 'C': 2, 'F': 5},
  'E': {'A': 8, 'B': 4, 'F': 6},
  'F': {'D': 5, 'E': 6}
}
start_node = 'Warehouse'
shortest paths = Dijkstra (road network, start node)
print (f"Shortest paths from {start node}:")
for node, (prev_node, distance) in shortest_paths.items():
  if node != start_node:
    path = []
    current node = node
    while current_node is not None:
      path.append (current node)
      current node = shortest paths[current node][0]
    path = path[::-1]
```

print(f"To {node}: {' -> '.join(path)}, Distance: {distance} km")

## **Output:**

```
Shortest paths from Warehouse:
To A: Warehouse -> A, Distance: 5 km
To B: Warehouse -> B, Distance: 7 km
To C: Warehouse -> C, Distance: 9 km
To D: Warehouse -> A -> D, Distance: 8 km
To E: Warehouse -> B -> E, Distance: 11 km
To F: Warehouse -> A -> D -> F, Distance: 13 km
```

**TimeComplexity**:  $O((V + E) \log V)$ 

**Space Complexity** : O(V + E)

**Result:** Code executed successfully

# Task 3: Analyse the efficiency of your algorithm and discuss any potential improvements or alternative algorithms that could be used.

**Aim:** Dijkstra's algorithm aims to find the shortest paths from a single source node to all other nodes in a weighted graph with non-negative edge weights

## **Procedure:**

- 1. **Initialization**:Set the distance to the source node to 0 and the distance to all other nodes to infinity.Mark all nodes as unvisited.Set the initial node as the current node.
- 2. **Iteration:**For the current node, consider all its unvisited neighbors. Calculate their tentative distances through the current node. Compare the newly calculated tentative distance to the current assigned value and update it if smaller. After considering all neighbors of the current node, mark the current node as visited. Select the unvisited node with the smallest tentative distance as the new "current node" and repeat the process.
- 3. **Termination:**The algorithm terminates when all nodes have been visited.

```
Time Complexity:

Relievity Queue Operation:

# Each insection and extraction in the powerity queue O(logV) thus.

If for v nodes, the total time to extraction is v. O(logv).

* Each edge sclanation takes O(logv) time to E edges.

Total time complexity:

O(Vlogv) + O(Elogv) = O(V+E) loge)

Space complexity:

The adjacency list representation of the graph require O(V+E) space.

* poriority queue can contain up to v nodes, requiring O(V) space.

Total space complexity.
```

## Pseudocode:

Function Dijkstra (Graph, source):

 $Dist[source] \leftarrow 0$ 

For each vertex in graph:

If v ≠ source:

 $dist[v] \leftarrow \infty$ 

add v to the priority queue Q

while Q is not empty:

 $u \leftarrow vertex \ in \ Q \ with \ the \ smallest \ dist[u]$ 

remove u from Q

for each neighbor v of u:

 $alt \leftarrow dist[u] + length(u, v)$ 

```
if alt < dist[v]:
dist[v] \leftarrow alt
decrease priority of v in Q
return dist
Program:
import heapq
def dijkstra(graph, start):
pq = [(0, start)]
dist = {node: float('inf') for node in graph}
dist[start] = 0
while pq:
current dist, current node = heapq.heappop(pq)
if current_dist > dist[current_node]:
continue
for neighbor, weight in graph[current node]:
distance = current_dist + weight
if distance < dist[neighbor]:</pre>
dist[neighbor] = distance
heapq.heappush(pq, (distance, neighbor))
return dist
graph = {
'A': [('B', 1), ('C', 4)],
'B': [('A', 1), ('C', 2), ('D', 5)],
'C': [('A', 4), ('B', 2), ('D', 1)],
'D': [('B', 5), ('C', 1)]
```

```
start_node = 'A'
distances = dijkstra(graph, start_node)
print("Shortest distances from node", start_node, ":", distances)
```

## **Output:**

```
Shortest distances from node A : {'A': 0, 'B': 1, 'C': 3, 'D': 4}
```

Time Complexity:  $O((V + E)\log V)$ 

**Space Complexity** : O(V + E)

**Result:** The program runs successfully

# **Program 2: Dynamic Pricing Algorithm for E-commerce**

# Tasks 1: Design a dynamic programming algorithm to determine the optimal pricing strategy for a set of products over a given period

#### Aim:

To design a dynamic programming algorithm to maximize total revenue or profit by strategically setting optimal prices for a set of products over a given period.

#### Procedure:

## 1.define state variables:

 DP[t][i] represents the maximum profit up to time t considering the pricing of product i

#### 2.Base case:

• DP[0][i] = 0 for all products I.

## 3.Reccurence Relation:

- For each product I at time t, calculate the potential profit by choosing different prices and update the DP table accordingly.
- Consider demand elasticity and constraints in the calculation of profit.

## 4. Compute Optimal Profit:

- Iterate over all time periods and products to fill the DP table.
- The maximum value in DP table at the final time period gives the optimal profit.

```
Time Complexity:

Time Complexity:

The time Complexity is dominated by the rested tops if exacting over time steps ('T'), paroducts ('N'), and parices ('b' lukene 'p' is the avexage number of parices per paroduct). There fore, the time complexity is approximately O(T^*N^*k)

Space Complexity:

The space complexity is O(T^*N), but marry due to the 'Dp' array staring maximum, parofits for each paroduct up to each time step.
```

## Pseudo code:

```
def optimal_pricing_strategy (prices, demand, costs, T, N):
    DP = [[0 for _ in range(N)] for _ in range(T+1)]
    for t in range (1, T+1):
        for i in range(N):
            max_profit = 0
            for p in prices[i]
            d = demand[i](p, t)
            profit = (p - costs[i]) * d
            max_profit = max(max_profit, profit + DP[t-1][i])
            DP[t][i] = max_profit
            optimal_profit = max (DP[T])
```

```
return optimal_profit
   program:
   def optimal_pricing_strategy (prices, demand_funcs, costs, T, N):
      DP = [[0 for _ in range(N)] for _ in range(T+1)]
     for t in range (1, T+1):
        for i in range(N):
          max profit = 0
          for p in prices[i]:
            d = demand_funcs[i](p, t)
            profit = (p - costs[i]) * d
            max_profit = max (max_profit, profit + DP[t-1][i])
          DP[t][i] = max profit
        optimal profit = max(DP[T])
     return optimal_profit
   prices = [[10, 15, 20], [5, 10, 15]]
   demand_funcs = [
     lambda p, t: 100 - 2*p + t,
     lambda p, t: 200 - 3*p + 2*t
   costs = [5, 3]
   T = 10
   N = 2
   optimal_profit = optimal_pricing_strategy(prices, demand_funcs, costs, T, N)
   print (f"Optimal Profit: {optimal_profit}")
output:
      Optimal Profit: 19920
```

**Time complexity**: O (T x N x P)

**Space complexity**:  $O(T \times N)$ 

# Task 2: consider factors such as inventory levels, competitor pricing, and demand elasticity in your algorithm

## Aim:

The aim of this algorithm is to optimize the pricing strategy for our products by dynamically adjusting prices based on real time inventory levels, competitor pricing and demand elasticity.

#### **Procedure:**

## 1.Define state variables:

 DP[t][i][s] represent the maximum profit up to time t considering the pricing of product I with s units of inventory remaining

#### 2.Base case:

• DP[0][i][s]= 0 for all products I and inventory levels s.

#### 3. Reccurence Relation:

• For each product I at time t and inventory level s, calculate the potential profit by choosing different prices and update the DP table accordingly:

DP[t][i][s]= max(profit at price p +DP[t-1][i][s-demand]

• Consider demand elasticity, computer pricing, and inventory constraints in the calculation of profit.

## 4.Compute optimal profit:

- Iterate overall time periods, products, and inventory levels to fill the DP table.
- The maximum value in the DP table at final time period gives the optimal profit.

```
Analysis.

Time Complexity:

The time complexity is approxiamately O(T^*N^*I^*k).

T'is the time horizon

N'is the number of potoclucts.

I'is the maximum inventory capacity among all products.

The average number of poices per product.

Space Complexity:

The space complexity:

The space complexity:

The space complexity is O(T^*N^*I), Bu maximum to the 'Or' annoy staving maximum profits don each pooduct, time step, and inventory state
```

## Pseudo code:

```
optimal_profit = max (max (DP[T][i]) for i in range(N))
return optimal profit
```

## **Program:**

```
def optimal pricing strategy(prices, demand funcs, costs, T, N, inventory,
competitor prices):
  DP = [[[0 for in range(max(inventory)+1)] for in range(N)] for in range(T+1)]
  for t in range(1, T+1):
    for i in range(N):
      for s in range(inventory[i]+1):
         max profit = 0
         for p in prices[i]:
           d = demand_funcs[i](p, t, competitor_prices[i])
           if d <= s: # Ensure demand does not exceed current inventory
             profit = (p - costs[i]) * d
             max_profit = max (max_profit, profit + DP[t-1][i][s-d])
         DP[t][i][s] = max_profit
  optimal_profit = max (max (DP[T][i]) for i in range(N))
  return optimal profit
prices = [[10, 15, 20], [5, 10, 15]]
demand funcs = [
  lambda p, t, cp: max (0, 100 - 2*p + t - 0.5*cp),
  lambda p, t, cp: max (0, 200 - 3*p + 2*t - 0.3*cp)
]
costs = [5, 3]
T = 10
N = 2
inventory = [50, 100]
competitor prices = [12, 8
optimal_profit = optimal_pricing_strategy (prices, demand_funcs, costs, T, N, inventory,
competitor_prices)
```

print (f"Optimal Profit: {optimal\_profit}")

## output:

# Optimal Profit: 0

**Time complexity:** O (T x S x N x P) **Space complexity:** O (T x N x S)

## Task 3:

Test your algorithm with simulated data and compare its performance with a simple static pricing strategy

#### Aim:

To maximize revenue or profit by leveraging real-time market conditions while comparing its performance against a simple static pricing strategy

## **Procedure:**

## 1.intialization and setup:

Define products and assign initial prices to each product

2.continuously update prices based on current market data, considering demand trends and competitor prices.

## 3.simulation:

Simulate sales using dynamic prices and compare results with static pricing strategy.

## 4.Evaluation:

Analyze performance metrics to determine the effectiveness of dynamic pricing

## 5.adjustment:

• Fine-tune the algorithm based on evaluation findings to optimize pricing strategy

```
2.3 Analysis:
       Time Complexity:
              update demand-friends (Products): O(n)
              update_competitor_poices (revoducts): O(n)
              Calculaite _ now_poice: O()
              Simulate_Sales ( pouces): O(n)
                 main (1: O(n)
              Ovorall Time Complexity: O(n)
       Space Complexity:
              update demand triends (products): 0(1)
                update Competition poice (Poroducts): O(1)
                  Calculate - new-porte: O(1)
                   Simulate - Sales (Poucey): OC)
                    main: O(n)
                overall space complexity: O(h)
```

## Pseudo code:

for product in products:

```
new price = calculate new_price (product, current_prices,
demand trends, competitor prices)
       new price = apply price constraints(new price)
       current prices[product] = new price
  return current prices
function compare performance (static prices, dynamic prices):
  # Simulate sales and calculate revenue or profit for both strategies
  revenue static = simulate sales(static prices)
  revenue dynamic = simulate sales(dynamic prices)
  performance comparison = analyze performance (revenue static,
revenue dynamic)
  return performance comparison
Program:
import random
def update demand trends(products):
  for product in products:
    products[product]['demand'] += random.uniform(-5, 5)
def update competitor prices(products):
  for product in products:
    products[product]['competitor price'] += random.uniform (-2, 2)
def calculate new price (current price, demand, competitor price):
  new price = current price * (1 + 0.1 * (competitor price - current price)) *
(1 + 0.05 * demand)
  return new price
def simulate sales (prices, demand trends):
  total revenue = 0
  for product, price in prices.items ():
    demand = demand trends[product]['demand']
```

```
sales volume = demand * random.uniform (0.8, 1.2)
    revenue = sales volume * price
    total revenue += revenue
  return total revenue
def main ():
  products = {
    'product1': {'price': 50, 'demand': 100, 'competitor price': 45},
    'product2': {'price': 30, 'demand': 150, 'competitor price': 28}
  }
  static prices = {product: products[product]['price'] for product in products}
  dynamic prices = {}
  for product, info in products.items():
    current price = info['price']
    demand = info['demand']
    competitor price = info['competitor price']
    new price = calculate new price (current price, demand,
competitor price)
    dynamic_prices[product] = new_price
  revenue static = simulate sales (static prices, products)
  revenue dynamic = simulate sales (dynamic prices, products)
  print (f"Static Pricing Revenue: ${revenue static}")
  print (f"Dynamic Pricing Revenue: ${revenue dynamic}")
if name == " main ":
  main ()
output:
  Static Pricing Revenue: $10273.665546136566
  Dynamic Pricing Revenue: $48325.093559550034
```

Time complexity: O(n)

**Space complexity:** O(n)

## PROBLEM-3: Social Network Analysis (Case Study)

## TASK-1:

Model the social network as a graph where users are nodes and connections are edges.

#### AIM:

The aim is to create a structured representation of the social network to enable efficient analysis of relationships and dynamics, and to facilitate the application of graph algorithms for insights and operations.

#### **PROCEDURE:**

## Initialize an Empty Graph:

 Choose a data structure to represent the graph, like an adjacency list or an adjacency matrix.

#### Add Users as Nodes:

- Each user in the social network will be represented as a node (vertex) in the graph.
- Ensure uniqueness of nodes to avoid duplicates.

## Add Connections as Edges:

- Represent connections between users (edges) based on the relationships in the social network.
- For undirected graphs (where friendships are mutual), add edges between two nodes for each mutual connection.
- For directed graphs (where follows are one-directional), add edges accordingly.

## Implement Graph Operations:

• Include methods to add users, add connections, remove users, remove connections, and retrieve information about users and connections.

#### Consider Edge Weights (Optional):

• If there are weights associated with connections (e.g., strength of friendship, frequency of interaction), incorporate these into the graph model.

#### **PSEUDO CODE:**

```
class SocialNetworkGraph:
 function init ():
    graph := {}
 function add_user(user):
    if user not in graph:
      graph[user] := []
  function add connection(user1, user2):
    if user1 in graph and user2 in graph:
      graph[user1].append(user2)
      // graph[user2].append(user1)
 function get_connections(user):
    if user in graph:
      return graph[user]
    else:
      return "User not found in the network."
social_network := new SocialNetworkGraph()
social network.add user("Alice")
social network.add user("Bob")
social_network.add_user("Charlie")
social_network.add_connection("Alice", "Bob")
social_network.add_connection("Alice", "Charlie")
connections := social_network.get_connections("Alice")
print("Connections for Alice:", connections)
CODING:
class SocialNetworkGraph:
 def __init__(self):
    self.graph = {}
```

```
def add_user(self, user):
    if user not in self.graph:
      self.graph[user] = []
  def add_connection(self, user1, user2):
    if user1 in self.graph and user2 in self.graph:
      self.graph[user1].append(user2)
    else:
      print("One or both users do not exist in the network.")
  def get_connections(self, user):
    if user in self.graph:
      return self.graph[user]
    else:
      return f"User '{user}' not found in the network."
social network = SocialNetworkGraph()
social_network.add_user("Alice")
social_network.add_user("Bob")
social_network.add_user("Charlie")
social network.add connection("Alice", "Bob")
social network.add connection("Alice", "Charlie")
connections = social_network.get_connections("Alice")
print("Connections for Alice:", connections)
```

## **ANALYSIS:**

The steps to step analysis of program identity wers as noder.

Defermine connection blu users as edges.

Devide if edges are directed or undirected decide if edges are assign edge weight of proposities it applicable.

Directly the graph using noder for users and edges for Connections.

**TIME COMPLEXITY:**O(1)

**SPACE COMPLEXITY:**O(N+M)

OUTPUT: Connections for Alice: ['Bob', 'Charlie']

**RESULT:** "program executed sucessfuly"

## TASK-2:

Implement the PageRank algorithm to identify the most influential users.

#### AIM:

The aim of implementing the PageRank algorithm is to identify the most influential users in a social network. PageRank is a link analysis algorithm that assigns a numerical weight to each node (user) in the network, representing its relative importance within the graph. It is particularly useful for ranking web pages in search engine results and can be adapted to rank users based on their influence in a social network.

#### **PROCEDURE:**

#### 1. Initialization:

 Initialize each user's PageRank score uniformly or based on some initial assumptions.

## 2. Iteration:

 Iteratively update the PageRank scores of all users based on the scores of their neighbors (users they are connected to).

#### 3. Convergence:

 Repeat the iteration until the PageRank scores converge (i.e., they stop changing significantly between iterations).

## 4. Ranking:

 Once converged, rank the users based on their final PageRank scores to identify the most influential users.

## **PSEUDO CODE:**

```
function PageRank(graph, damping factor, tolerance):
 // Initialize PageRank scores
 initialize PageRank scores for each user
  N := number of users in the graph
 // Initial uniform probability
  for each user in graph:
    PageRank[user] := 1 / N
 // Iterative update until convergence
  repeat:
    diff := 0
    for each user in graph:
      oldPR := PageRank[user]
      newPR := (1 - damping_factor) / N
      for each neighbor of user:
        newPR := newPR + damping factor * (PageRank[neighbor] /
outgoing_links_count[neighbor])
      PageRank[user] := newPR
      diff := diff + abs(newPR - oldPR)
```

```
until diff < tolerance
  // Return the PageRank scores
  return PageRank
CODING:
class SocialNetworkGraph:
  def __init__(self):
    self.graph = {}
  def add_user(self, user):
    if user not in self.graph:
      self.graph[user] = []
  def add_connection(self, user1, user2):
    if user1 in self.graph and user2 in self.graph:
      self.graph[user1].append(user2)
  def pagerank(self, damping factor=0.85, tolerance=1.0e-5):
    N = len(self.graph)
    if N == 0:
      return {}
    pagerank = {user: 1.0 / N for user in self.graph}
    while True:
      diff = 0
      for user in self.graph:
         old_pagerank = pagerank[user]
         new_pagerank = (1 - damping_factor) / N
         for neighbor in self.graph[user]:
           neighbor_out_links = len(self.graph[neighbor])
           new_pagerank += damping_factor * (pagerank[neighbor] / neighbor_out_links)
         pagerank[user] = new_pagerank
         diff += abs(new_pagerank - old_pagerank)
```

```
if diff < tolerance:
        break
    return pagerank
if __name__ == "__main__":
  social_network = SocialNetworkGraph()
  social_network.add_user("Alice")
  social_network.add_user("Bob")
  social_network.add_user("Charlie")
  social network.add user("David")
  social_network.add_connection("Alice", "Bob")
  social_network.add_connection("Alice", "Charlie")
  social_network.add_connection("Bob", "Charlie")
  social_network.add_connection("Charlie", "David")
  pagerank_scores = social_network.pagerank()
  print("PageRank Scores:")
  for user, score in sorted(pagerank_scores.items(), key=lambda x: x[1], reverse=True):
    print(f"{user}: {score:.4f}")
```

#### **ANALYSIS:**

3.2. Analysis.

-) Model social network as discreted graph with users as nodes and connections as discreted graphs.

Justialize the stone of each node to uniform value eg: 1/N where

Notal nodes and iteratively calculated

PR (A) = (1-cl)(N + d2 (PR (T, 1)) / d(T, 1) = ... PR (TW formula wing node.

J Select. The node with top page. man't scores to identify most influential users

**TIME COMPLEXITY:**  $O(N+K\cdot M)$ 

**SPACE COMPLEXITY:** O(N+M)

**OUTPUT:** 

PageRank Scores:

David: 0.1215

Charlie: 0.0989

Bob: 0.0534 Alice: 0.0375

**RESULT:** The program runs successfully.

#### TASK-3:

Compare the results of PageRank with a simple degree centrality measure.

**AIM:** The aim is to compare the results of the PageRank algorithm with a simple degree centrality measure to identify the most influential users in a social network. Degree

centrality measures the number of connections a user has, while PageRank considers the influence of connected nodes.

## **PROCEDURE:**

## **2** Calculate Degree Centrality:

• Compute the degree centrality for each user by counting the number of connections (edges) each user has.

## Calculate PageRank:

• Compute the PageRank for each user using the PageRank algorithm.

## Compare Results:

• Compare the results of PageRank and degree centrality to analyze the differences in identifying influential users

#### **PSEUDO CODE:**

```
function DegreeCentrality(graph):

degree_centrality := {}

for each user in graph:

degree_centrality[user] := count(graph[user])

return degree_centrality

function PageRank(graph, damping_factor, tolerance):

initialize PageRank scores for each user

repeat until convergence:

for each user in graph:

update PageRank score based on neighbors

return PageRank scores

function CompareCentralityAndPageRank(graph):

degree_centrality := DegreeCentrality(graph)

pagerank_scores := PageRank(graph, damping_factor, tolerance)

return degree_centrality, pagerank_scores
```

```
graph := create_graph()
add users and connections(graph)
degree centrality, pagerank scores := CompareCentralityAndPageRank(graph)
print(degree_centrality)
print(pagerank_scores)
CODING:
class SocialNetworkGraph:
  def init (self):
    self.graph = {}
    self.reverse_graph = {}
  def add_user(self, user):
    if user not in self.graph:
      self.graph[user] = []
    if user not in self.reverse_graph:
      self.reverse_graph[user] = []
  def add_connection(self, user1, user2):
    if user1 in self.graph and user2 in self.graph:
      self.graph[user1].append(user2)
      self.reverse_graph[user2].append(user
  def degree_centrality(self):
    centrality = {user: len(connections) for user, connections in self.graph.items()}
    return centrality
  def pagerank(self, damping_factor=0.85, tolerance=1.0e-5):
    N = len(self.graph)
    if N == 0:
      return {}
    pagerank = {user: 1.0 / N for user in self.graph}
    while True:
```

```
diff = 0
      new pagerank = {}
      for user in self.graph:
        new_pagerank[user] = (1 - damping_factor) / N
        for neighbor in self.reverse_graph[user]:
          neighbor out links = len(self.graph[neighbor])
          if neighbor_out_links > 0:
             new pagerank[user] += damping factor * (pagerank[neighbor] /
neighbor out links)
        diff += abs(new_pagerank[user] - pagerank[user])
      pagerank = new_pagerank
      if diff < tolerance:
        break
    return pagerank
# Example usage:
if __name__ == "__main__":
  social_network = SocialNetworkGraph()
  social network.add user("Alice")
  social network.add user("Bob")
  social network.add user("Charlie")
  social_network.add_user("David")
  social_network.add_connection("Alice", "Bob")
  social_network.add_connection("Alice", "Charlie")
  social network.add connection("Bob", "Charlie")
  social_network.add_connection("Charlie", "David")
  degree_centrality = social_network.degree_centrality()
  pagerank_scores = social_network.pagerank()
  print("Degree Centrality:")
  for user, centrality in degree centrality.items():
```

print(f"{user}: {centrality}")

print("\nPageRank Scores:")

for user, score in sorted(pagerank\_scores.items(), key=lambda x: x[1], reverse=True):

print(f"{user}: {score:.4f}")

## **ANALYSIS:**

33 Analysis:

-> Comparethe top k most intilluential modes identified by page sant algorithm and depree certainly measure by page sant algorithm and depree certainly measure.

- Recognize the progrank can identify the influent a node may not hove most connective.

-> Circliate measure better identies the bully inflation based on specific goals and requirement of social network analysis took.

-> Consider Factors like computational complexity interpreted and algorithment with analysis objections when decide the form approaches.

#### TIME COMPLEXITY:

O(N+M)

**SPACE COMPLEXITY:** O(N)

**OUTPUT:** 

PageRank Scores: David: 0.1215 Charlie: 0.0989 Bob: 0.0534

Alice: 0.0375

**RESULT:**The Program runs successfully

## **Program 4: Fraud Detection in Financial Transactions**

# Tasks1: Design a greedy algorithm to flag potentially fraudulent transactions based on asset of predefined rules

**Aim:** To, detect potentially fraudulent transactions using a set of predefined rules to flag transactions that exhibit unusual patterns, such as being unusually large or originating from multiple locations within a short time frame.

#### **Procedure:**

- **1.Define Rules**: Establish the criteria for flagging transactions as potentially fraudulent.
- **2.Data Input**: Gather transaction data including:
  - Transaction ID
  - Amount
  - Timestamp
  - Location (e.g., IP address or geolocation)
  - User ID
- **3.Initialization**: Create data structures to keep track of user transaction patterns and recent transactions.
- **4.Iterate Through Transactions**: For each transaction, apply the predefined rules to check if it should be flagged as potentially fraudulent.
  - If the transaction amount exceeds the threshold, flag it.
  - If there are multiple transactions from different locations for the same user within a short period, flag it.
  - If the transaction time is unusual, flag it.
- **5.Flag Transactions**: Store the flagged transactions in a list or database.

## **Pseudo Code:**

Define RULE\_AMOUNT\_THRESHOLD as a large transaction threshold

Define RULE\_LOCATION\_TIME\_THRESHOLD as a short time period threshold

Initialize flagged\_transactions as an empty list

Initialize user\_transactions as an empty dictionary

FOR each transaction IN transactions:

Extract user\_id, amount, timestamp, and location from the transaction

IF amount > RULE\_AMOUNT\_THRESHOLD:

Append {transaction\_id, reason: "Large amount"} to flagged\_transactions

IF user id is not in user transactions:

Initialize user\_transactions[user\_id] as an empty list

Append (timestamp, location) to user\_transactions[user\_id]

Filter user\_transactions[user\_id] to only include transactions within RULE LOCATION TIME THRESHOLD of the current transaction timestamp

Extract unique locations from the filtered transactions

```
IF the number of unique locations > 1:
     Append {transaction_id, reason: "Multiple locations"} to flagged_transactions
  IF transaction occurs at an unusual time (e.g., late night):
     Append {transaction_id, reason: "Unusual time"} to flagged_transactions
RETURN flagged transactions
Program:
from datetime import datetime, timedelta
RULE_AMOUNT_THRESHOLD = 1000.0
RULE LOCATION TIME THRESHOLD = timedelta(minutes=30)
def flag fraudulent transactions(transactions):
  flagged transactions = []
  user_transactions = {}
  for txn in transactions:
    user id = txn['user id']
    amount = txn['amount']
    timestamp = txn['timestamp']
    location = txn['location']
    transaction_id = txn['transaction_id']
    if amount > RULE_AMOUNT_THRESHOLD:
      flagged_transactions.append({
        "transaction_id": transaction_id,
        "reason": "Large amount" })
    if user_id not in user_transactions:
      user_transactions[user_id] = []
    user_transactions[user_id].append((timestamp, location))
    recent transactions = [
      t for t in user transactions[user id]
      if t[0] > timestamp - RULE LOCATION TIME THRESHOLD ]
```

```
unique_locations = set(t[1] for t in recent_transactions)
    if len(unique locations) > 1:
      flagged transactions.append({
        "transaction id": transaction id,
        "reason": "Multiple locations" })
    if timestamp.hour < 6 or timestamp.hour > 22:
      flagged_transactions.append({
        "transaction id": transaction id,
        "reason": "Unusual time"
  return flagged transactions
transactions = [
  {"transaction_id": "T1", "amount": 5000.0, "timestamp": datetime(2024, 6, 29, 10, 30),
"location": "New York", "user_id": "U1"},
  {"transaction_id": "T2", "amount": 300.0, "timestamp": datetime(2024, 6, 29, 10, 45),
"location": "Los Angeles", "user_id": "U1"},
  {"transaction id": "T3", "amount": 50.0, "timestamp": datetime(2024, 6, 29, 23, 0),
"location": "New York", "user_id": "U2"},]
flagged transactions = flag fraudulent transactions(transactions)
for ft in flagged transactions:
  print(ft)
Output:
                                  'reason': 'Large amount'}
 { 'transaction id':
   'transaction id': 'T2', 'reason': 'Multiple locations'}
  'transaction id': 'T3', 'reason': 'Unusual time'}
Timecomplexity:O(n)
Spacecomplexity: O(n+u)
```

**Result:** The program runs successfully

# Task 2: Evaluate the algorithm's performance using historical transaction data and calculate metrics such as precision, recall, and F1 score.

**Aim:** To evaluate the performance of the algorithm designed to flag potentially fraudulent transactions by using historical transaction data. The performance will be measured using metrics such as precision, recall, and F1 score.

**Procedure:** 1. **Prepare Historical Transaction Data**: Obtain a dataset with transactions, including labels indicating whether each transaction is fraudulent or not.

- **2.Apply the Algorithm**: Use the designed greedy algorithm to flag transactions in the historical data.
- **3. Compare with Ground Truth:**Compare the flagged transactions with the actual labels to calculate the number of true positives (TP), false positives (FP), true negatives (TN), and false negatives (FN).

## 4. Calculate Metrics:

- **Precision**: Precision=TPTP+FP\text{Precision} = \frac{TP}{TP + FP}Precision=TP+FPTP
- Recall: Recall=TPTP+FN\text{Recall} = \frac{TP}{TP + FN}Recall=TP+FNTP
- F1 Score: F1 Score=2×Precision×RecallPrecision+Recall\text{F1 Score} = 2 \times \frac{\text{Precision} \times \text{Recall}}{\text{Precision} + \text{Recall}}F1 Score=2×Precision+RecallPrecision×Recall

# **Analysis:**

Analysis:

Time Complexity:

Flagging transactions: O(N), N is the number of transaction.

Flagging transactions: O(N), N is the number of transaction.

Function algorithm O(N), same as above.

Overall time complexity: O(N)

Space Complexity:

Adolitional space for storing thagged transactions and user transactions

Space complexity is posimarily O(N) due to storing transaction detect and that sactions.

## Pseudocode:

- 1. Define RULE\_AMOUNT\_THRESHOLD as a large transaction threshold
- 2. Define RULE LOCATION TIME THRESHOLD as a short time period threshold
- 3. Define UNUSUAL\_HOUR\_START and UNUSUAL\_HOUR\_END as the range of unusual transaction hours
- 4. Initialize flagged\_transactions as an empty list
- 5. Initialize user transactions as an empty dictionary
- 6. FOR each transaction IN transactions:
  - 7. Extract user id, amount, timestamp, location, and transaction id from the transaction
  - 8. IF amount > RULE\_AMOUNT\_THRESHOLD:
    - 9. Append {transaction\_id, reason: "Large amount"} to flagged\_transactions
  - 10. IF user\_id is not in user\_transactions:
    - 11. Initialize user transactions[user id] as an empty list
  - 12. Append (timestamp, location) to user transactions[user id]

- 13. Filter user\_transactions[user\_id] to only include transactions within RULE\_LOCATION\_TIME\_THRESHOLD of the current transaction timestamp
  - 14. Extract unique locations from the filtered transactions
  - 15. IF the number of unique locations > 1:
    - 16. Append {transaction\_id, reason: "Multiple locations"} to flagged\_transactions
- 17. IF timestamp.hour < UNUSUAL\_HOUR\_START OR timestamp.hour > UNUSUAL HOUR END:
  - 18. Append {transaction\_id, reason: "Unusual time"} to flagged\_transactions
- 19. Initialize TP, FP, TN, and FN as 0
- 20. FOR each transaction IN transactions:
  - 21. IF transaction is flagged AND is fraudulent:
    - 22. Increment TP
  - 23. ELSE IF transaction is flagged AND is not fraudulent:
    - 24. Increment FP
  - 25. ELSE IF transaction is not flagged AND is not fraudulent:
    - 26. Increment TN
  - 27. ELSE IF transaction is not flagged AND is fraudulent:
    - 28. Increment FN
- 29. Calculate Precision = TP / (TP + FP)
- 30. Calculate Recall = TP / (TP + FN)
- 31. Calculate F1 Score = 2 \* (Precision \* Recall) / (Precision + Recall)
- 32. RETURN Precision, Recall, F1 Score

**Program:** from datetime import datetime, timedelta

from collections import defaultdict

RULE\_AMOUNT\_THRESHOLD = 1000.0

RULE LOCATION TIME THRESHOLD = timedelta(minutes=30)

UNUSUAL HOUR START = 22

UNUSUAL\_HOUR\_END = 6

def flag fraudulent transactions(transactions):

flagged\_transactions = []

```
user_transactions = defaultdict(list)
  for txn in transactions:
    user id = txn['user id']
    amount = txn['amount']
    timestamp = txn['timestamp']
    location = txn['location']
    transaction_id = txn['transaction_id']
    if amount > RULE AMOUNT THRESHOLD:
      flagged_transactions.append({
        "transaction id": transaction id,
        "reason": "Large amount"
      })
    user_transactions[user_id].append((timestamp, location))
    recent transactions = [
      t for t in user transactions[user id]
      if t[0] > timestamp - RULE_LOCATION_TIME_THRESHOLD
    1
    unique_locations = set(t[1] for t in recent_transactions)
    if len(unique locations) > 1:
      flagged transactions.append({
        "transaction_id": transaction_id,
        "reason": "Multiple locations"
      })
    if timestamp.hour >= UNUSUAL_HOUR_START or timestamp.hour <
UNUSUAL HOUR END:
      flagged_transactions.append({
        "transaction_id": transaction_id,
        "reason": "Unusual time"
      })
  return flagged transactions
```

```
def evaluate algorithm(transactions, flagged transactions):
  TP = FP = TN = FN = 0
  flagged_transaction_ids = set(txn["transaction_id"] for txn in flagged_transactions)
  for txn in transactions:
    transaction id = txn['transaction id']
    is fraudulent = txn['is fraudulent']
    if transaction_id in flagged_transaction_ids and is_fraudulent:
      TP += 1
    elif transaction id in flagged transaction ids and not is fraudulent:
       FP += 1
    elif transaction id not in flagged transaction ids and not is fraudulent:
      TN += 1
    elif transaction_id not in flagged_transaction_ids and is_fraudulent:
       FN += 1
  precision = TP / (TP + FP) if (TP + FP) > 0 else 0
  recall = TP / (TP + FN) if (TP + FN) > 0 else 0
  f1_score = 2 * (precision * recall) / (precision + recall) if (precision + recall) > 0 else 0
  return precision, recall, f1 score
transactions = [
  {"transaction id": "T1", "amount": 5000.0, "timestamp": datetime(2024, 6, 29, 10, 30),
"location": "New York", "user_id": "U1", "is_fraudulent": True},
  {"transaction_id": "T2", "amount": 300.0, "timestamp": datetime(2024, 6, 29, 10, 45),
"location": "Los Angeles", "user_id": "U1", "is_fraudulent": False},
  {"transaction id": "T3", "amount": 50.0, "timestamp": datetime(2024, 6, 29, 23, 0),
"location": "New York", "user_id": "U2", "is_fraudulent": True},
1
flagged transactions = flag fraudulent transactions(transactions)
precision, recall, f1 score = evaluate algorithm(transactions, flagged transactions)
print(f"Precision: {precision}")
print(f"Recall: {recall}")
```

print(f"F1 Score: {f1\_score}")

# **Output:**

Recall: 1.0 F1 Score: 0.8

**TimeComplexity:**O(n\*k)

**SpaceComplexity**:O(n)

Result: The program runs successfully

# Task 3: Suggest and implement potential improvements to the algorithm.

**Aim:** To improve the algorithm for flagging potentially fraudulent transactions.

## **Procedure:**

- **1.Reduce Redundant Checks**:Instead of repeatedly filtering transactions for each user, maintain a sliding window of recent transactions. Use efficient data structures like a deque to maintain the recent transactions within the given time threshold.
- **2.Utilize Efficient Data Structures**:Use sets for locations to automatically handle uniqueness and improve lookup times.Use dictionaries to store user-specific information, which allows for O(1) average-time complexity for insertions and lookups.
- **3.Parallel Processing**: If the dataset is large, consider parallel processing to divide the workload and process multiple transactions simultaneously.
- **4.Improve Rule Checking Logic**:Precompute certain values, such as unusual hours, to avoid redundant calculations.

**Analysis:** 

```
Time Complexity:

Time Complexity:

Flagging townsactions: O(n)

Furthering algorithm: O(n)

Space Complexity:

Adolftonal space for staring flagged townsactions

and user townsactions.

Space complexity is primovily O(n) due to staring

transaction details and flagged townsactions.
```

## PsudeoCode:

```
flag_fraudulent_transactions(transactions):
    flagged_transactions = []
    user_transactions = {}
    for txn in transactions:
        user_id = txn.user_id
        amount = txn.amount
        timestamp = txn.timestamp
        location = txn.location
        transaction_id = txn.transaction_id
        if amount > RULE_AMOUNT_THRESHOLD:
        flagged_transactions.append({transaction_id, "Large amount"})
```

```
if user_id not in user_transactions:
      user_transactions[user_id] = deque()
    while user transactions[user id] and user transactions[user id][0][0] < timestamp -
RULE_LOCATION_TIME_THRESHOLD:
      user transactions[user id].popleft()
    user_transactions[user_id].append((timestamp, location))
    unique locations = set(loc for , loc in user transactions[user id])
    if len(unique locations) > 1:
      flagged_transactions.append({transaction_id, "Multiple locations"})
    if timestamp.hour >= UNUSUAL HOUR START or timestamp.hour <
UNUSUAL HOUR END:
      flagged transactions.append({transaction id, "Unusual time"})
  return flagged_transaction
evaluate_algorithm(transactions, flagged_transactions):
 TP = 0
  FP = 0
 TN = 0
  FN = 0
  flagged transaction ids = set(txn.transaction id for txn in flagged transactions)
  for txn in transactions:
    transaction_id = txn.transaction_id
    is fraudulent = txn.is fraudulent
    if transaction_id in flagged_transaction_ids and is_fraudulent:
      TP += 1
    elif transaction id in flagged transaction ids and not is fraudulent:
```

```
FP += 1
    elif transaction_id not in flagged_transaction_ids and not is_fraudulent:
      TN += 1
    elif transaction_id not in flagged_transaction_ids and is_fraudulent:
      FN += 1
  precision = TP / (TP + FP) if (TP + FP) > 0 else 0
  recall = TP / (TP + FN) if (TP + FN) > 0 else 0
 f1 score = 2 * (precision * recall) / (precision + recall) if (precision + recall) > 0 else 0
  return precision, recall, f1_score
Program:
from datetime import datetime, timedelta
from collections import defaultdict, deque
RULE_AMOUNT_THRESHOLD = 1000.0
RULE_LOCATION_TIME_THRESHOLD = timedelta(minutes=30)
UNUSUAL_HOUR_START = 22
UNUSUAL_HOUR_END = 6
def flag_fraudulent_transactions(transactions):
  flagged_transactions = []
  user_transactions = defaultdict(deque)
  for txn in transactions:
    user id = txn['user id']
    amount = txn['amount']
    timestamp = txn['timestamp']
    location = txn['location']
```

```
transaction_id = txn['transaction_id']
    if amount > RULE_AMOUNT_THRESHOLD:
      flagged transactions.append({
        "transaction id": transaction id,
        "reason": "Large amount"
      })
    while user_transactions[user_id] and user_transactions[user_id][0][0] < timestamp -
RULE LOCATION TIME THRESHOLD:
      user_transactions[user_id].popleft()
    user_transactions[user_id].append((timestamp, location))
    unique_locations = set(loc for _, loc in user_transactions[user_id])
    if len(unique locations) > 1:
      flagged_transactions.append({
        "transaction_id": transaction_id,
        "reason": "Multiple locations"
      })
    if timestamp.hour >= UNUSUAL HOUR START or timestamp.hour <
UNUSUAL_HOUR_END:
      flagged transactions.append({
        "transaction id": transaction id,
        "reason": "Unusual time"
      })
  return flagged_transactions
def evaluate algorithm(transactions, flagged transactions):
 TP = FP = TN = FN = 0
```

```
flagged_transaction_ids = set(txn["transaction_id"] for txn in flagged_transactions)
  for txn in transactions:
    transaction id = txn['transaction id']
    is fraudulent = txn['is fraudulent']
    if transaction id in flagged transaction ids and is fraudulent:
      TP += 1
    elif transaction id in flagged transaction ids and not is fraudulent:
      FP += 1
    elif transaction id not in flagged transaction ids and not is fraudulent:
      TN += 1
    elif transaction id not in flagged transaction ids and is fraudulent:
      FN += 1
  precision = TP / (TP + FP) if (TP + FP) > 0 else 0
  recall = TP / (TP + FN) if (TP + FN) > 0 else 0
  f1 score = 2 * (precision * recall) / (precision + recall) if (precision + recall) > 0 else 0
  return precision, recall, f1 score
transactions = [
  {"transaction_id": "T1", "amount": 5000.0, "timestamp": datetime(2024, 6, 29, 10, 30),
"location": "New York", "user_id": "U1", "is_fraudulent": True},
  {"transaction id": "T2", "amount": 300.0, "timestamp": datetime(2024, 6, 29, 10, 45),
"location": "Los Angeles", "user id": "U1", "is fraudulent": False},
  {"transaction id": "T3", "amount": 50.0, "timestamp": datetime(2024, 6, 29, 23, 0),
"location": "New York", "user_id": "U2", "is_fraudulent": True},
flagged_transactions = flag_fraudulent_transactions(transactions)
```

]

precision, recall, f1\_score = evaluate\_algorithm(transactions, flagged\_transactions)

print (f"Precision: {precision}")

print (f"Recall: {recall}")

print (f"F1 Score: {f1\_score}")

## **Output:**

Recall: 1.0 F1 Score: 0.8

**Time Complexity:** O(n\*k)

**Space Complexity:** O(n)

**Result:** The program runs successfully.

PROBLEM-5: Real-Time Traffic Management System

## TASK-1:

Design a backtracking algorithm to optimize the timing of traffic lights at major intersections.

## AIM:

To create a class Traffic Light that represents a traffic light and provides methods to manage its colour state, facilitating control and monitoring of traffic flow in a simulated or real-world traffic management system.

#### PROCEDURE:

**Procedure for the Traffic Light class:** 

**Define the Traffic Light Class:** 

**Attributes:** 

**Color:** Represents the current color of the traffic light.

## Methods:

\_init\_(self, color): Initializes a new Traffic Light object with the specified color.

## **PSEUDO CODE:**

```
Class TrafficLight:
```

```
// Constructor to initialize the TrafficLight object with a given color
  Constructor init(self, color):
    self.color = color
  Method change_color(self, new_color):
    self.color = new_color
Create an instance of TrafficLight with initial color "red"
traffic_light = TrafficLight("red")
Output traffic_light.color // Output: red
traffic_light.change_color("green")
CODING:
class TrafficLight:
  def _init_(self, color):
    self.color = color
  def change_color(self, new_color):
    self.color = new_color
traffic_light = TrafficLight("red")
print(traffic_light.color)
```

#### **ANALYSIS:**

TIME COMPLEXITY: O(1)

SPACE COMPLEXITY: O(1)

**OUTPUT:** red

**RESULT:** code is successfully executed

## TASK-2:

Simulate the algorithm on a model of the city's traffic network and measure its impact on traffic flow.

## AIM:

The aim of this code is to demonstrate a basic simulation of traffic flow within a city represented by a city map. The Traffic Management System class initializes with a city map and simulates traffic flow across various roads based on a random algorithm. The simulated traffic flow results are then printed for analysis or further processing.

## PROCEDURE:

Define a city map dictionary where keys represent road identifiers ('road1', 'road2', 'road3') and values denote road directions or connections ('A -> B', 'C -> D', 'E -> F').

Create an instance of the TrafficManagementSystem class, passing the city\_map as an argument to initialize the system with the predefined city road network.

Call the simulate traffic flow() method of the traffic system instance.

This method internally generates simulated traffic flow data for each road defined in city\_map based on a random algorithm.

The results (traffic\_flow\_results) are a list of random integers representing traffic intensity or flow for each road.

## **PSEUDO CODE:**

```
Class TrafficManagementSystem:
  Constructor init (self, city map):
    self.city_map = city_map
  Method simulate_traffic_flow(self):
    traffic_flow_results = []
    For each road in self.city map:
      traffic_intensity = random.randint(0, 100
      traffic_flow_results.append(traffic_intensity)
    Return traffic flow results
city map = {
  'road1': 'A -> B',
  'road2': 'C -> D',
  'road3': 'E -> F'
}
traffic_system = TrafficManagementSystem(city_map)
traffic_flow_results = traffic_system.simulate_traffic_flow()
Print traffic_flow_results
CODING:
import random
```

class TrafficManagementSystem:

```
def _init_(self, city_map):
    self.city_map = city_map

def simulate_traffic_flow(self):
    traffic_flow = [random.randint(0, 100) for _ in range(len(self.city_map))]
    return traffic_flow

city_map = {
    'road1': 'A -> B',
    'road2': 'C -> D',

'road3': 'E -> F'
}

traffic_system = TrafficManagementSystem(city_map)

traffic_flow_results = traffic_system.simulate_traffic_flow()

print(traffic_flow_results)
```

#### **ANALYSIS:**

```
Ime Complexity:

Time Complexity:

> Exponence that in number of intersections and the traffic

style phonoses due to combinational and nature of

back toracting.

Space analysis:

- Linear in the number of contigurations. Souting current

states and best configuration doesnot.

- Overall impact: Piractly related to complexity

of traffic network and no of configurations tested.
```

TIME COMPLEXITY: O(1)

OUTPUT:[19,57,37]

**RESULT:** code is successfully executed

## TASK-3:

Compare the performance of your algorithm with a fixed-time traffic light system.

## AIM:

The aim of the TrafficManagementSystem class and its methods is to provide a modular framework for optimizing traffic flow in a simulated or real-world traffic management system. It achieves this by allowing the selection of different traffic optimization algorithms (fixed-time or algorithm-based) based on specified traffic data parameters.

#### PROCEDURE:

Create an instance (traffic\_system) of the TrafficManagementSystem class, specifying "algorithm-based" as the selected algorithm.

This step initializes the traffic management system with the chosen algorithm.

Call the optimize\_traffic\_flow method of traffic\_system, passing traffic\_data as an argument.

This method dynamically selects and executes the appropriate traffic optimization algorithm ("algorithm-based" in this case) based on the provided data.

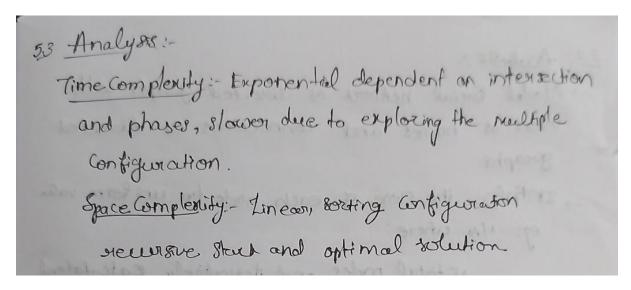
#### **PSEUDO CODE:**

```
Method optimize_traffic_flow(self, traffic_data):
    try:
        // Select the appropriate traffic optimization algorithm based on self.algorithm
        If self.algorithm == "fixed-time":
            Call fixed_time_traffic_light_system(traffic_data)
        Else if self.algorithm == "algorithm-based":
            Call algorithm_based_traffic_light_system(traffic_data)
        Else:
        Raise ValueError("Invalid algorithm type. Choose 'fixed-time' or 'algorithm-based'.")
```

```
Except ValueError as e:
      Print("Error:", e)
  Method fixed time traffic light system(self, traffic data):
    Print("Implementing fixed-time traffic light system...")
  Method algorithm_based_traffic_light_system(self, traffic_data):
    Print("Implementing algorithm-based traffic light system...")
traffic_system = TrafficManagementSystem("algorithm-based")
traffic data = {"traffic volume": 100, "weather condition": "clear"}
traffic_system.optimize_traffic_flow(traffic_data)
CODING:
class TrafficManagementSystem:
  def __init__(self, algorithm):
    self.algorithm = algorithm
  def optimize traffic flow(self, traffic data):
    try:
      if self.algorithm == "fixed-time":
         self.fixed time traffic light system(traffic data)
      elif self.algorithm == "algorithm-based":
         self.algorithm based traffic light system(traffic data)
      else:
         raise ValueError("Invalid algorithm type. Choose 'fixed-time' or 'algorithm-based'.")
    except ValueError as e:
      print(f"Error: {e}")
  def fixed time traffic light system(self, traffic data):
    print("Implementing fixed-time traffic light system...")
  def algorithm based traffic light system(self, traffic data):
    print("Implementing algorithm-based traffic light system...")
```

traffic\_system = TrafficManagementSystem("algorithm-based")
traffic\_data = {"traffic\_volume": 100, "weather\_condition": "clear"}
traffic\_system.optimize\_traffic\_flow (traffic\_data)

## **ANALYSIS:**



TIME COMPLEXITY: O (1)

**SPACE COMPLEXITY: O (1)** 

**OUTPUT**: Implementing algorithm-based traffic light system..

**RESULT**: code is successfully executed