Subarrays → contiguous part of array

3	6	9	12	14	19	20	23	25	27

14 19 20 23	✓
9 6 12	✗
6 9 12	✓
20	✓
9 6 3	✗

complete array is also a subarray

uniquely identify the subarray

↓

[start, end] , start <= end
end < N

0	1	2	3	4	5	6
4	2	10	3	12	-2	15

start	end
1	1
1	2
1	3
1	4
1	5
1	6

ans = 6

count all subarrays

start	# count
0	7
1	6
2	5
3	4
4	3
5	2
6	1
	28

count of subarray in array of size N

start	# cont
0	N
1	N-1
2	N-2
3	.
.	.
.	.
.	.
N-1	1

N*(N+1)/2

```
printSubarray ( arr[], int start, int end)
{
    for ( i = start ; i <= end ; i++)
        print(arr[i]);
}
```

```
sumSubarray ( arr[], int start, int end)
{
    int sum = 0;
    for ( i = start ; i <= end ; i++)
        sum += arr[i];
}
```

print all subarrays

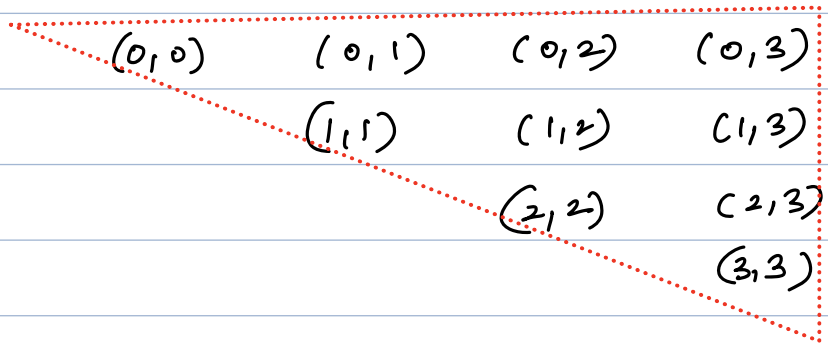
	0	1	2
All:	2	8	9

$\frac{N*(N+1)}{2}$

s	e	
0	0	[2]
0	1	[2, 8]
0	2	[2, 8, 9]
1	1	[8]
1	2	[8, 9]
2	2	[9]

To consider all subarrays, you have to consider all possible pairs of (s, e)

0	1	2	3
4	1	3	6




```
for( i=0; i<N; i++)  → fix start
{
    for( j=i; j<N; j++)  → fix end
    {
        [i, j]
        for( k=i; k<=j; k++)
            print(arr[k])
    }
}
```

T.C: $O(N^3)$

Find sum of each subarray

	0	1	2
All:	2	8	9

B.F:- consider all subarray & iterate.

S	e	
0	0	~~[2]~~ 2
0	1	~~[2, 8]~~ 10
0	2	~~[2, 8, 9]~~ 19
1	1	~~[8]~~ 8
1	2	~~[8, 9]~~ 17
2	2	~~[9]~~ 9

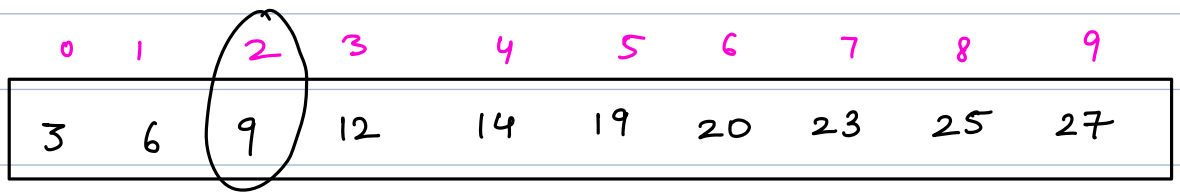
```
for( i=0; i<N; i++)  → fix start
{
    for( j=i; j<N; j++)  → fix end
    {
        [i, j]   int sum=0;
        for( k=i; k<=j; k++)
            sum += arr[k];
        print(sum);
    }
}
```

T.C: $O(N^3)$
S.C: $O(1)$

// build pf sum

```
for( i=0; i<N; i++)  → fix start
{
    for( j=i; j<N; j++)  → fix end
    {
        [i, j]
        if(i==0) sum = pf[j];
        else  sum = pf[j] - pf[i-1];
        print(sum);
    }
}
```

T.C: $O(N^2)$
S.C: $O(N)$



[2,2] → arr[2]

[2,3] → arr[2] + arr[3] = [2,2] + arr[3]

[2,4] → arr[2] + arr[3] + arr[4] = [2,3] + arr[4]

[i, j] → [i − j-1] + arr[j]

0	1	2	3
4	2	-1	3

```
for(int i=0; i<N; i++)
{
    sum=0;
    for(int j=i; j<N; j++)
    {
        sum += arr[j];
        print(sum);
    }
}
```

i	j	sum = 0
0	0	4
	1	4+2 = 6
	2	6+(-1) = 5
	3	5+3 = 8
		sum = 0
1	1	2
	2	2+(-1) = 1
	3	1+3 = 4

T.C: $O(N^2)$

S.C: $O(1)$

Find total sum of all subarray sums.

0	1	2	3
3	2	-1	5

s	e	sum
0	0	3
0	1	5
0	2	4
0	3	9
1	1	2
1	2	1
1	3	6
2	2	-1
2	3	4
3	3	5

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```
int ts = 0;
for(int i=0; i<N; i++)
{
    sum = 0;
    for(int j = i; j<N; j++)
    {
        sum += arr[j];
        ts += sum;
    }
}
```

T.C: $O(N^2)$

S.C: $O(1)$

	0	1	2
A:	-1	3	4

0,0	-1	arr[0]
0,1	2	arr[0] + arr[1]
0,2	6	arr[0] + arr[1] + arr[2]
1,1	3	arr[1]
1,2	7	arr[1] + arr[2]
2,2	4	arr[2]

21 = 3* arr[0] + 4* arr[1] + 3* arr[2]

3*(-1) + 4*3 + 3*4 = 21

contribution of i^{th} element = No of subarrays in which i^{th} element lies * arr[i]

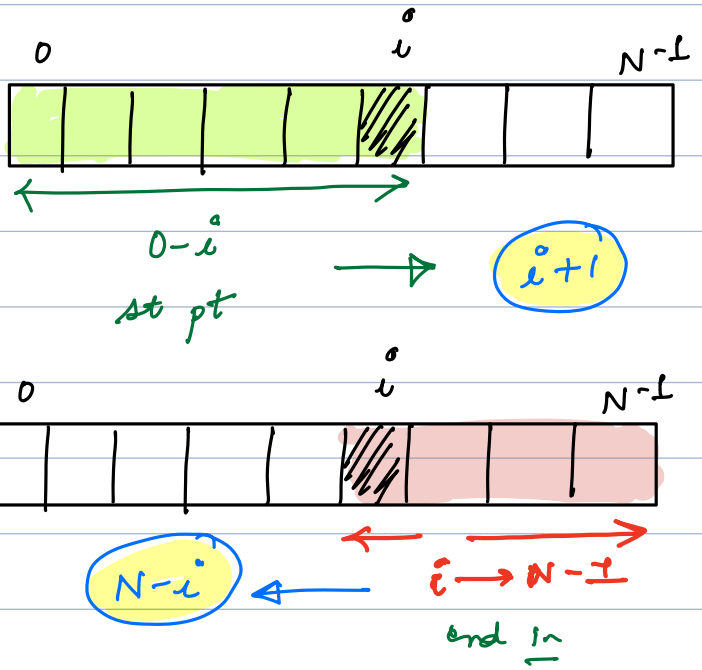
- In how many subarrays i^{th} element is coming?

	0	1	2	3	4	5
A:	3	-2	4	1	2	6

0,2	1,2	2,2
0,3	1,3	2,3
0,4	1,4	2,4
0,5	1,5	2,5

ans = 12

= no of st pt * no of end pts



$(N-1) - i + 1 = N - i$

No of subarrays = $(i+1) * (N-i)$

cont of i^{th} ele = $(i+1) * (N-i) * arr[i]$

T.C: O(N)
S.C: O(1)

```
int ans = 0;
for( i = 0; i < N; i++)
{
    ans += (i+1) * (N-i) * arr[i];
}
```

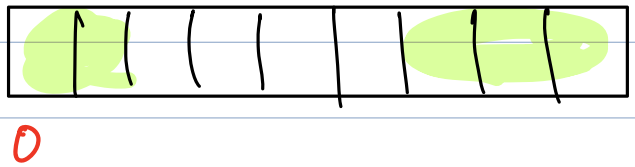
Pick both from sides!

3 elements

3 or 5

0

0



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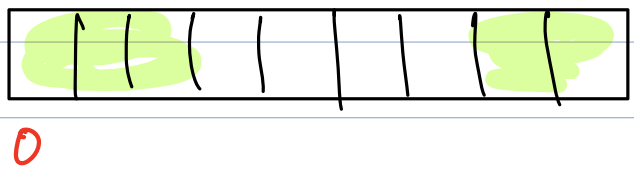

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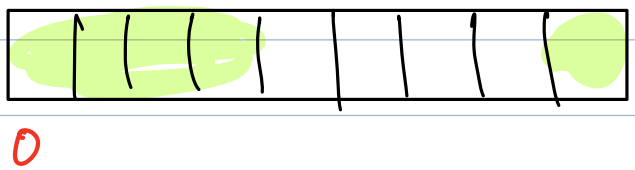

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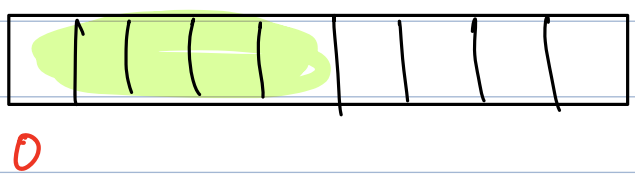

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