

$$P_s[i] = P_s[i-1] + A[i]$$

$$\text{sum}(x \rightarrow y) = \text{pf}[y] - \text{pf}[x-1]$$

Given a string of lowercase English alphabets.

→ indexes

Find count of pairs (i, j) such that

$$i < j \text{ \& \& } \begin{aligned} \text{arr}[i] &= 'a' \\ \text{arr}[j] &= 'g' \end{aligned}$$

0 1 2 3 4 5
a b e g a g

$\left\{ \begin{array}{l} 0,3 \\ 0,5 \\ 4,5 \end{array} \right\}$

0 1 2 3 4 5 6
a c g d g a g

ans = 4

0 1 2 3 4 5 6 7
b c a g g a a g

ans = 5

Brute force → consider all pairs

T.C: $O(N^2)$

```
for ( i = 0; i < N; i++)  
{  
    if ( s[i] != 'a' ) continue;  
    for ( j = i+1; j < N; j++)  
    {  
        if ( s[j] == 'g' )  
            ans++;  
    }  
}
```

Traverse from end ← for every 'a' → looking for count of 'g's on right

OR

for every 'g' → calculate no of 'a's on left.

	<u>b</u>	<u>c</u>	<u>a</u>	<u>g</u>	<u>g</u>	<u>a</u>	<u>a</u>	<u>g</u>
cnt = 0	0	0	1	1	1	2	3	3
ans = 0	0	0	0	1	2	2	2	5

ans = 0, cnt = 0;

```
for ( i = 0; i < N; i++)  
{  
    if ( arr[i] == 'a' ) cnt++;  
    else if ( arr[i] == 'g' ) ans = ans + cnt;  
}
```

T.C: $O(N)$

↑
S.C: $O(1)$

Q Given an array of size N. Count the number of leaders in the array.

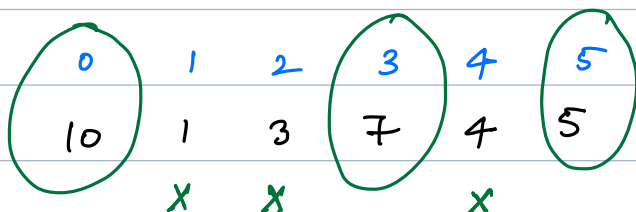
Leader: An element that is strictly greater than all the elements on its right side.

$$\boxed{A[i] > A[i+1 \dots n-1]}$$

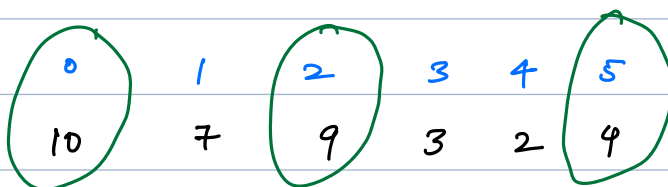
↓
leader

last element
is always a
leader.

Ex1:

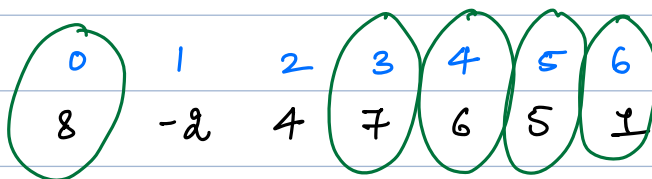


Ex2:-



ans = 3

Ex3:-



ans = 5

B-F For every index i , traverse on right

T.C: $O(N^2)$

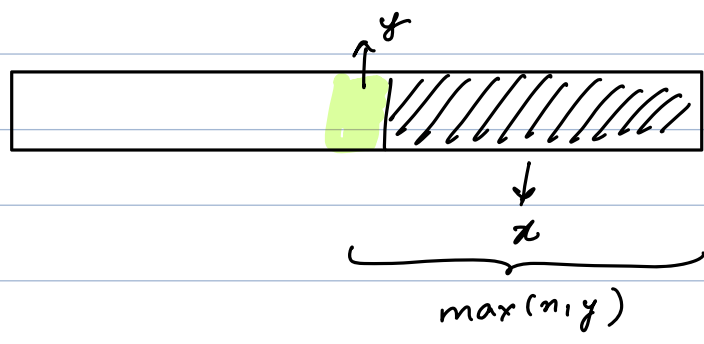
```

for(i=0; i<N; i++)
{
    bool flag = true;
    for(j=i+1; j<N; j++)
    {
        if(A[j] >= A[i])
        {
            flag = false; break;
        }
    }
    if(flag) { ans++; }
}

```

if $\text{element} > \max(i+1 \rightarrow n-1)$

↓
leader



0	1	2	3	4	5	
10	7	9	3	2	4	
3	2	2	1	1	1	ans
10	9	9	4	4	4	max

```
int ans = 1;
```

```
int maxe = arr[N-1];
```

```
for ( i = N-2; i >= 0; i-- )
```

T.C: $O(N)$

```
    if ( arr[i] > maxe )
```

```
    {
```

```
        ans ++;
```

```
        maxe = arr[i];
```

```
    }
```

```
}
```

Given an array of size N . Find the length of shortest range which contains both minimum and maximum element of the array.

0	1	2	3	4	5	6	7	8	9
1	2	3	1	3	4	6	4	6	3

min \rightarrow 1

max \rightarrow 6

len = 4

0	1	2	3	4	5	6	7	8	9	10
2	2	6	4	5	1	5	2	6	4	1

len = 3

min \rightarrow 1

max \rightarrow 6

0	1	2	3	4	5	6	7	8	9	10	11
1	6	4	2	7	7	5	1	3	1	1	5

min \rightarrow 1

max \rightarrow 7

$$\begin{array}{c} i \rightarrow j \\ \hline \downarrow \\ j - i + 1 \end{array}$$

0	1	2	3	4	5
8	8	8	8	8	8

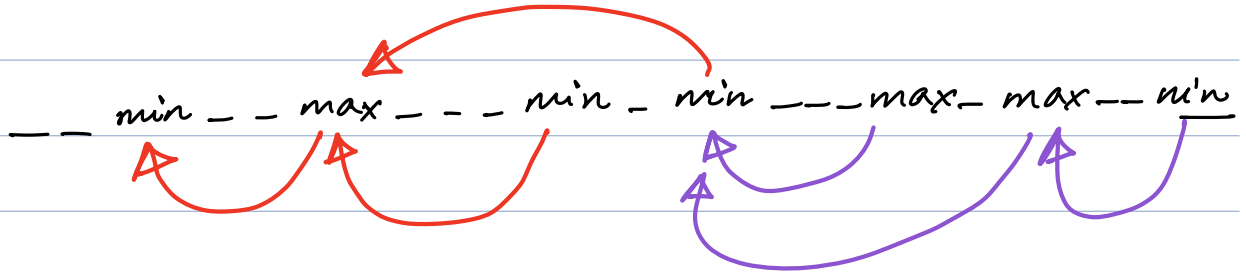
min \rightarrow 8

max \rightarrow 8

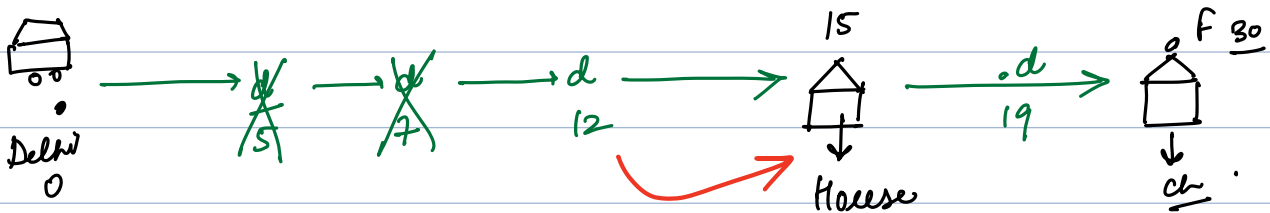
ans = 1

max & min → always on the border of the range

let's consider every possible pair of min & max



• For every max/min, find closest_min / closest_max on left



0	1	2	3	4	5	6	7	8	9
1	2	3	1	3	4	6	4	6	3

min = 1
max = 6

c_min = -1 0 0 0 3 3 3 3 3 3

c_max = -1 -1 -1 -1 -1 -1 -1 6 6 8

ans = N
10 4

0	1	2	3	4	5	6	7	8	9	10
2	2	6	4	5	1	5	2	6	4	1

min = 1
max = 6

c-min = -1 -1 -1 -1 -1 5 5 5 5 5 10

c-max = -1 -1 -1 2 2 2 2 2 2 8 8

ans = ~~11~~ 3

// find min & max element
if (min == max) return 1;

cmin = -1, cmax = -1, ans = N

for (i = 0; i < N; i++)

{

if (arr[i] == min)

{

c-min = i;

if (c-max != -1)

{

ans = min(ans, i - c-max + 1);

}

}

else if (arr[i] == max)

{

c-max = i;

if (c-min != -1)

{ ans = min(ans, i - c-min + 1); }

}

}

T.C: $O(N)$

S.C: $O(1)$

Best Time to Buy & sell stock - I



N days \rightarrow stock prices



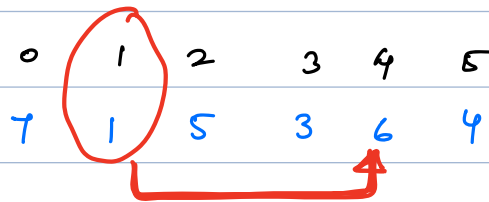
allowed to do only one transaction

Buy and sell



Buy before sell

Find Max Profit.



1) Consider all pairs
/ Buy & sell

$0 \rightarrow i-1$

min stock price

at i

selling day

\rightarrow leaders

OR

at $i+1$

buying

$i+1 \rightarrow n-1$

max stock price