Welcome 1

Aganda: Bridge construction

MST

Krusal's algo

Prim's algo

Dgickstra's algo.

Q Cuiven N islands & cost of construction of a bridge blue multiple pair of islands. (cost >0)

Find min-cost of construction reg. set it is possible to travel from one island to any other island. Via bridges.

If not possible return -1

3 9 6 4 5 7 6 3 7

4-2-6

3 2 5 7 7 7 3 3 3 6 3

⇒ pro = -1 → if graph is disconnected.

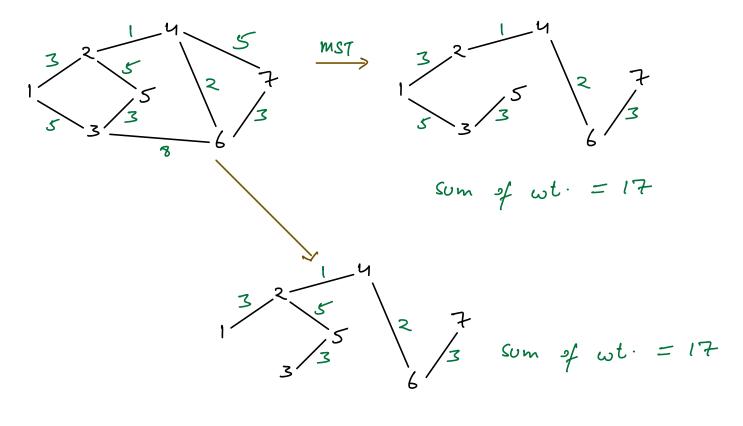
(thech using DSU)

Q In a wheeted graph with N nodes, what is nun # edges possible?

Min = N-1

(Trees)

Minimum Spanning Tree => Tree generated from a connected weighted graph s.t all nodes are connected & sum of weights of all selected edges is minimum.

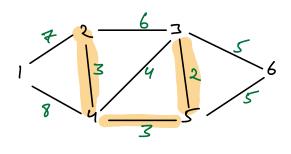


- => multiple MST possible for any graph.
- => braph with unique veights => Unique MST.

Algo. to find MST -> 1) Kruskal's algo. 2 greedy algo. 2) Prim's algo. I greedy algo.

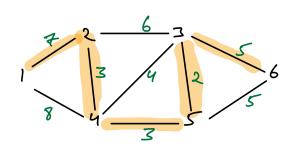
Kvuskal's algo

Select edges with min weight if it is not forming a cycle, till the complete graph is connected.



- Steps 1) Sort edges wirt weight Tic => O(Elog E)
 - 2) Consider each node as a set i.e parent[i] = i (DSU) (visited array won't work)

3) Travel all edges (U,V) & take union. If disjoint -> add in ans



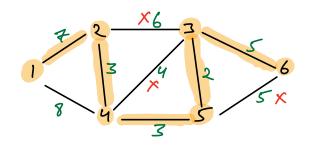
$$3\frac{2}{5}5$$

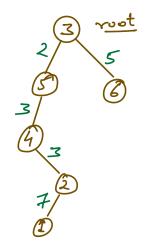
$$2\frac{3}{4}$$

T.C
$$\rightarrow$$
 O(Elog E + E)
S.C \rightarrow O(N)

Prim's algo

- i) Start with any node as root of MST of keep adding the other nodes as its children.
- 2) hive priority to nodes that have less edge weight.





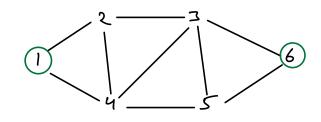
- 1) Start with root, insert its edges in a min heap (wiret to weight)
- 2) Pick min. weight edge from heap -> if it prims a upde i.e connect both visited nodes, repeat step 2 else add the other node as part of tree & insent its connected edges in min. heap.
- 3) Continue step 2 till complete tree is formed.

T.C O(Elog E)

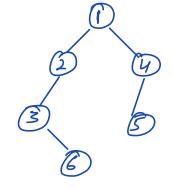
SC O(N+E)

VISC) Leap.

of Find min # edges to travel from u to v in undirected simple graph.



U=1 V=6 #=3

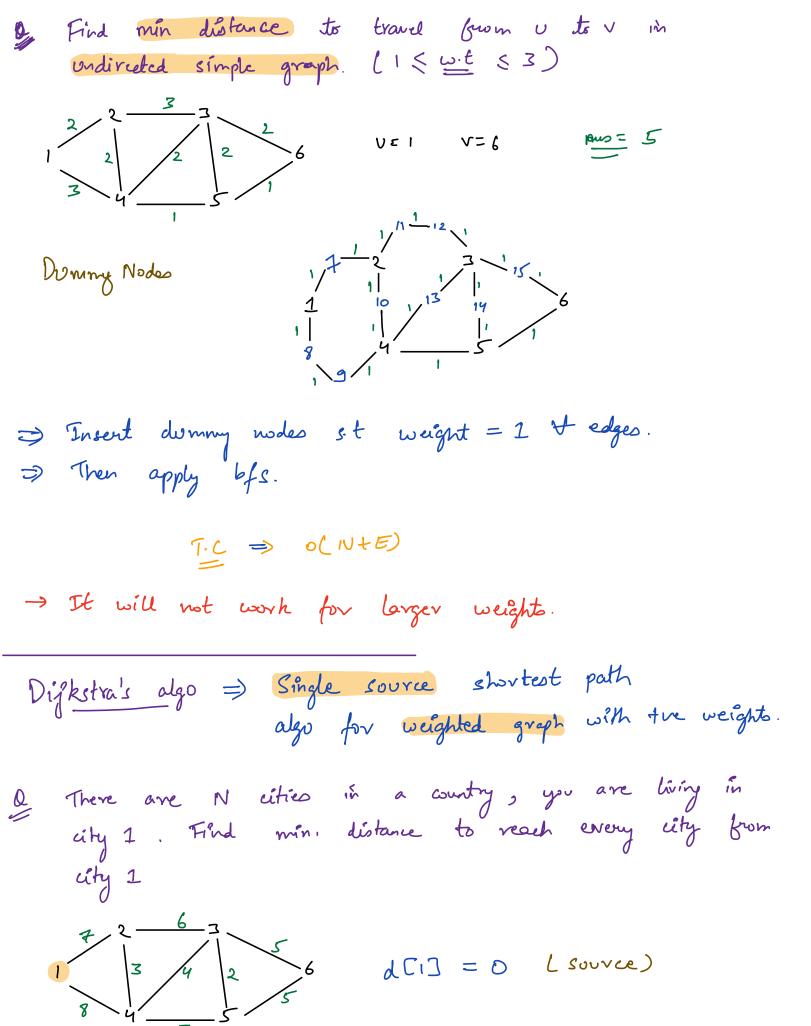


CXXXXXXX

Evel of V is ans.

T.C -> O(N+E)

s.c -> o(n)



$$dC1 = 0$$

$$dC2 = 2$$

$$dC3 = 2+2 = 4$$
Relaxing an edge.

To be continued