

# Welcome ☺

Agenda: Greedy Algo.  
3 ques<sup>ns</sup>.

Greedy  $\rightarrow$  Max Profits / Min Losses

Aman

Iphone  $\Rightarrow$  70K

Mukesh

Iphone  $\Rightarrow$  60K ✓

(10K Shoes 5K)  $\Rightarrow$  depends on multiple factors.

Q1 There is a limited time sale going on for toys.  
 $A[i]$   $\Rightarrow$  sale end time for  $i^{\text{th}}$  toy  
 $B[i]$   $\Rightarrow$  beauty of  $i^{\text{th}}$  toy.

It takes 1 unit of time to buy a toy. And toy can only be bought if current time of  $i^{\text{th}}$  toy  $< A[i]$

Buy toys s.t sum of beauty is maximized.

eg.  $A$  [ 3 1 3 2 3 ]  $t=0$  toy 4  $A[i]$  9  
 $B$  [ 6 5 3 1 9 ]  $t=1$  0 6  
 $t=2$  2 3  
18

$t=0$  toy 1  $A[i]$  5  
 $t=1$  0 6  
 $t=2$  4 9  
20

Q [ 1 2 ]  
[ 3 1500 ]  
Ans = 1503

⇒ Select toys not based just on the beauty of it  
but select them w.r.t sale end time

A	[	<sup>0</sup> 1	<sup>1</sup> 3	<sup>2</sup> 5	<sup>3</sup> 3	<sup>4</sup> 3	<sup>5</sup> 5	<sup>6</sup> 5	]
B	[	5	2	4	7	1	3	8	]

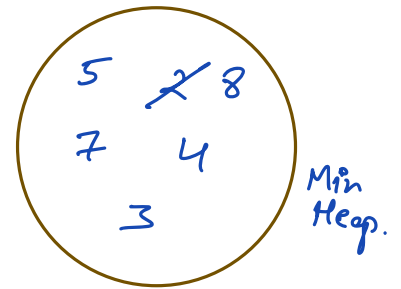
Step 1 ⇒ Sort array

A	[	<sup>0</sup> 1	<sup>1</sup> 3	<sup>2</sup> 3	<sup>3</sup> 3	<sup>4</sup> 5	<sup>5</sup> 5	<sup>6</sup> 5	]
B	[	5	2	7	(1)	4	3	(8)	]

very less beautiful.

high beauty toy  
⇒ purchase it  
⇒ remove smallest toy

t	toy	B[i]
t=0	0	5
t=1	1	2
t=2	2	7
t=3	4	4
t=4	5	3
t=5		



Code

```

1) sort w.r.t to A
t=0
for (i → 0 to N-1)
{
    if (A[i] > t) {
        Minheap.insert(B[i])
        t++
    }
    else { // sale end time > t
        if (Minheap.root < B[i]) {
            Minheap.getMin()
            Minheap.insert(B[i])
        }
    }
}

```

T.C →  $O(N \log N)$   
S.C →  $O(N)$

## Q2 Candy Distribution

There are  $N$  children with marks, teacher has to give them candies s.t

a) Every child has at least 1 candy.

b) Children with higher marks must have more candies than their neighbours.

$$\underline{i-1} < \overset{\downarrow}{i} > \underline{i+1}$$

Find min candies req. to do this.

eg.

A : [ 1 5 2 1 ]

C →     1   ~~1~~   ~~1~~   1  
             ~~1~~   2  
             2 3

- 1)  $\forall i, C[i] \geq 1$
- 2) if (  $A[i] > A[i-1]$  )  
     $C[i] = C[i-1] + 1$
- 3) if (  $A[i] > A[i+1]$  )  
     $C[i] = \max(C[i], C[i+1] + 1)$

Quiz eg: [ 4 4 4 4 4 ]

C     1     1     1     1     1

count = 5

eg: A : [ 1 6 3 1 10 12 20 5 2 ]

C :     1     ~~1~~   ~~1~~   1     ~~1~~   ~~1~~   ~~1~~   ~~1~~   1  
             2     2     2     2     3     4     2

cond-2 →

← cond-3

code

$\forall i, C[i] = 1$

for (  $i \rightarrow 1$  to  $N-1$  ) {

    if (  $A[i] > A[i-1]$  )  $C[i] = C[i-1] + 1$

}

ans = C[N-1]

for (i → N-2 to 0) {

if (A[i] > A[i+1]) C[i] = max(C[i], C[i+1] + 1)

ans += C[i]

}

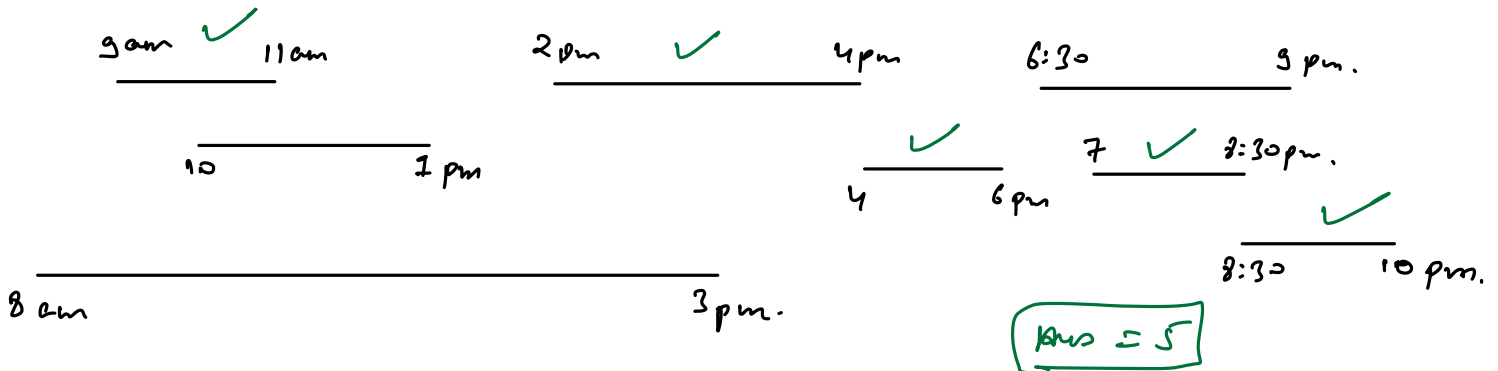
return ans;

T.C → O(N)

S.C → O(N)

Q Given N jobs with start & end time.

Find max. jobs that can be completed if only one job can be done at a time.

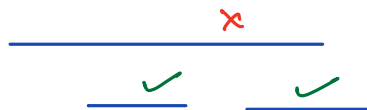


Diviz

S	[	1	5	8	7	12	13]
E	[	2	10	10	11	20	19]

Ans = 3

Start Time X



Duration X



End Time  $\rightarrow$  Early start time & less duration  
 $\Rightarrow$  Early End Time

Code

1. // Sort based on End Time.

ans = 1      e = E[0]

for (i  $\rightarrow$  1 to N-1)

{

    if (S[i]  $\geq$  e)

    {

        ans ++

        e = E[i]

    }

}

return ans;

T.C  $\rightarrow O(N \log N)$

S.C  $\rightarrow O(1)$