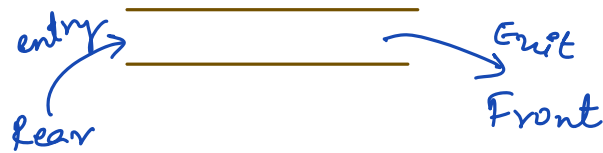


Welcome 😊

Agenda: Queues
Implementation
DA
1-2

Queue



FIFO \rightarrow First In First Out

Operations

- 1) Enqueue(n) \rightarrow Add n from the rear end.
 - 2) Dequeue() \rightarrow Remove element from front end.
 - 3) isEmpty() \rightarrow To check if queue is empty or not
 - 4) front()/rear() \rightarrow To get the element present at front/rear end
- T.C
O(1)

Implement Queue using dynamic arrays

enqueue(5)

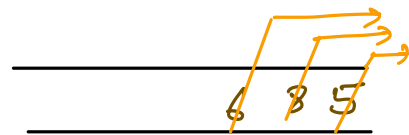
" (8)

" (6)

dequeue()

" ()

" ()

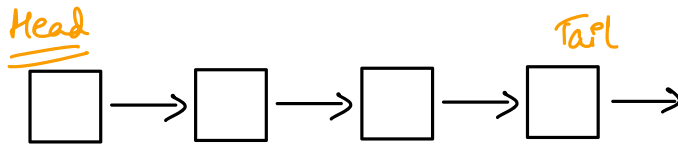


$r = \text{front} \neq 2$

$f = 0$

<pre> void enqueue (n) { r++ A[r] = n } </pre>	<pre> int dequeue () { if (isEmpty()) return -1 f++ return A[f-1] } </pre>	<pre> bool isEmpty () { return f > r } </pre>
--	--	--

Q Implement Queue using L.L



1) Insert at Head $\rightarrow O(1)$

1) Insert at Tail $\rightarrow O(1)$ ✓

2) Delete at Head $\rightarrow O(1)$ ✓

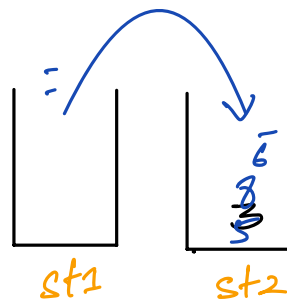
2) Delete at Tail $\rightarrow O(N)$

Q Implement Queue using Stacks..

```

enqueue (5)
" (8)
" (6)
dequeue ()
" ()
" ()

```



5
8
6
3
2
1

```

enqueue (7)
" (3)
dequeue ()
dequeue ()
dequeue ()

```

Insertion $\rightarrow O(1)$

Deletion $\rightarrow \cancel{O(N)} \rightarrow \underline{O(1)}$

$\left\{ \begin{array}{l} \underline{N} \text{ Insertions} \Rightarrow \underline{O(N)} \\ 1^{\text{st}} \text{ dequeue} \rightarrow \text{moved all st1} \rightarrow \text{st2} \\ \underline{N} \text{ Insertions} \quad \quad \quad \underline{O(N)} \\ \underline{N-1} \text{ dequeues} \quad \quad \quad \underline{O(1)} \end{array} \right.$

$\rightarrow \underline{O(N)}$ for all remaining $N-1$ dequeues.

```
void enqueue(n) {
    st1.push(n)
}
```

```
bool isEmpty() {
    return st1.empty()
    && st2.empty()
}
```

```
int dequeue() {
    if (isEmpty())
        return -1
    if (st2.empty()) {
        move()
    }
    return st2.pop()
}
```

```
void move() {
    while (!st1.empty()) {
        st2.push(st1.pop())
    }
}
```

T.C \Rightarrow $O(N)$

T.C \Rightarrow $O(1)$

Q Find N^{th} perfect number i.e. formed by digits 1 & 2

1 2 11 12 21 22 111 112 121 122
 1 2 3 4 5 6 7 8 9 10

```
int solve(int N)
```

```
{ if (N <= 2) return N
```

```
    // Queue
```

```
    q.enqueue(1)    q.enqueue(2)
```

```
    i = 3
```

```
    while (i <= N)
```

```
    { n = q.dequeue()
```

```
      a = n * 10 + 1
```

```
      b = n * 10 + 2
```

```
      if (i == N) return a
```

```
      if (i + 1 == N) return b
```

```
      q.enqueue(a)    q.enqueue(b)
```

```
      i = i + 2
```

```
    }
```

H.W \Rightarrow bit manip

122 121 112 111 22 21 12 11 2 1
 1 \rightarrow 1 \Rightarrow 11
 1 \rightarrow 2 \Rightarrow 12

i \rightarrow 2 1 1 1 1 1 1 1 1 1

n \rightarrow 1 2 1 1 1 1 1 1 1 1

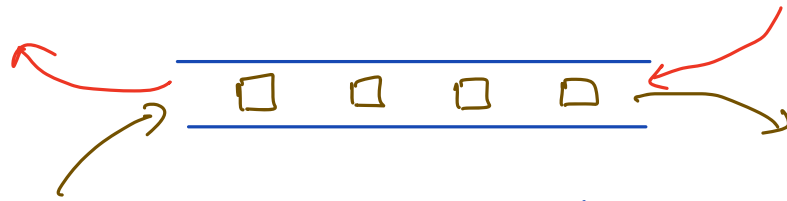
a \rightarrow 1 2 1 1 1 1 1 1 1 1

b \rightarrow 1 2 2 1 1 1 1 1 1 1

T.C $\rightarrow O(N)$

S.C $\rightarrow O(N)$

Doubly Ended Queue



D.S to use to implement \Rightarrow DLL \Rightarrow doubly L.L

Q Given an integer array A, \forall window of size K, find max. element.

eg: K=4

	1	8	5	6	7	4	2	0	3
	↓	↓	↓	↓	↓	↓			
	8	8	7	7	7	4			

1) Brute force \rightarrow \forall subarray of size K, find max.

T.C $\Rightarrow O(N * K)$ SC $\Rightarrow O(1)$

K=4

	1	8	5	6	7	4	2	0	3
	↓	↓	↓	↓	↓				
	X		X	X					

✓ stack + Queue \Rightarrow Doubly Ended Queue



Code

```
if ( f == i - K ) q.dequeueFront()
while ( !q.empty() && A[r] < A[i] ) q.dequeueRear()
q.enqueue(i)
print ( A[f] )
```

T.C $\rightarrow O(N)$

S.C $\rightarrow O(K)$

[1 8 5 6 7 4 2 0 3]
0 1 2 3 4 5 6 7 8

ans \Rightarrow 8 8 7

