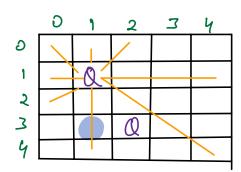
Backtracking 2

- N duceno
- Sudoku

Cuiven a N+N chemboard & boar of 2 queens. Check if they cannot attack each other.



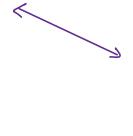
$$I/P \rightarrow (1,1) d (3,1) \Rightarrow false$$

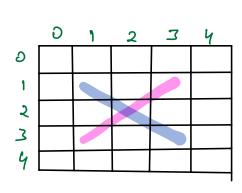
Direc



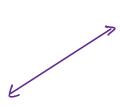


Same

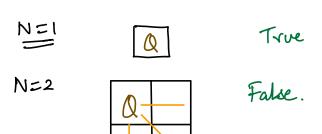


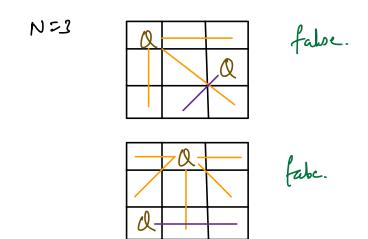


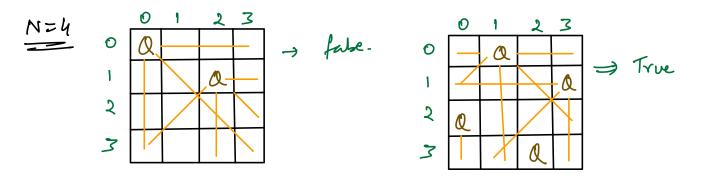
$$Y_1 - C_1 = = Y_2 - C_2$$



No liver a integer N, check if it is possible to place No green on a N*N chessboard s.t no greens attack each other.





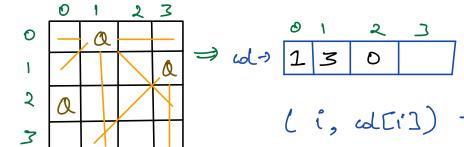


N lucero d N+N chestoard.

-> Every row should have enactly 1 queen.

-> Every column should have enactly 1 queen.

- a) Place queen row by row
- b) Place queen column by column.
- \Rightarrow We don't really need N*N entra space to keep track of placed queens. S.C \Rightarrow O(N)



(i, cd[i]) → locar of ith queen

boolean Nqueens (1°, N, col[])

if (r== N) return true; Il Pase care.

for (c > 0 to N-1) Il all possibilities.

lif (isValid (col [], r, c) Il valid possibilities

col[r] = c || DO

if (Nqueens (r+1, N, col[])) Il Recursion return true.

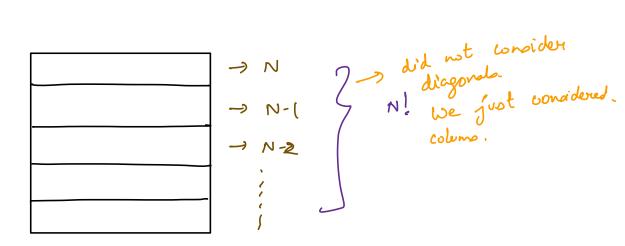
col [r] = -1 Il UNDD

return fabe.

boolean is Valid (col CI, Y, c) $Tc \Rightarrow O(N)$ {

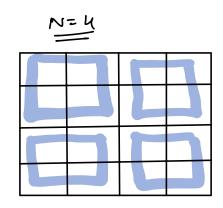
for $l i \Rightarrow o$ to Y-1)

{ j = col CiI j = col CiI



TC \Rightarrow O(N! \neq N) < O(N!) < O(N") SC \Rightarrow O(N + N) 10 orray recursion. Solve the given incomplete sudoku

Sudoku -> N * N grid where N is a perfect quere. Every row, every column of every block must have inique element.



2	=4		
1			
	Z	Ŋ	4
	2		
4		M	2

	٥	t	2	3
0	ん	4	1	3
> ₁	1	Z	Ŋ	4
2	3	2	4	1
	4	1	Μ	ત

$$\left(\frac{1}{2}\right) \times 2 = 0$$

$$\left(\frac{2}{2}\right) \approx 2 = 2$$

(3,2)

3-3%2=2

2-2%2=2

bool sudoku (A[][], N, T, E)

$$if(c = = N)$$

$$v++, c = 0$$

if
$$(r = = N)$$
 (
return true

```
A[~][c] = i // Do
valid case
              if ( check ( A, N, r, c) & & Sudoku (A, N, r, c+1))
Recursion
                   return True.
              ACYJ[c] = -1 11 UNOD
       return fale.
     bool cheek CACICI, N, r, c) T.C => O(N)
       for ( ° ) o to N-1) &
           ifl i!= c & A A[x][c] = = A[x][i]) (
               return fale.
          if [i!= x ll A[x](c] = = A[i](c]) 1
                return fale.
           sg = Sgrt(N)
           U = 8-8% Sq
           V = C - C% sq
           for ( i -> 0 to (sq-1))
           for ( 1 - 0 to (sq-1))
                 y = v + j = encluding correct cell.
                 if (n!= v II y!= e) dd A[n][y] == A[v][c]
                        return fabe.
           return true
```

ż

T.L > N +N ---- N2 time < O(N²)

S-C > O(N²)

L) recursive.