Find out the O-notation for the function. F(n) = 27 n + 16n

Let f(n) = 27 n + 16n

To obtain big Theta notation, we have to find out,

 $C_1 \times g(n) \leq F(n) \leq C_2 \times g(n)$ 

If we assume

C, = 27 and (2 = 43

For n 21 for q(n) = n we get

27 n = 27 n + 16n = 43 n

:. f(n) = 0 (n) for C1=27, C2=43 for n=1

Answer the following,

- (1) Find big theta (0) and big omega(2) notation
  - 1) f(n)=14\*7+83
  - 2) f(n)=83n"+84n.

(ii) Is 2" = O(2")? Explain.

To obtain big Omega -F(n) = 14 x 7 + 83 we can write it in polynomial form F(n) = 2n + 11n+6 where n=7

Now if g(n) = 7 (n) Then we get f(n) > 7(n) Hen4. 2n'+11n+6 = S(n)To obtain big theta notation. we have to find out  $C_1 \times g(n) \leq f(n) \leq C_2 \times g(n)$ If we assume CIZ7 and CZ= 26 Then with n=7, 7 n ≤ 2n + 11n+6 ≤ 26n. is true. :. f(n) = O(n) where n=7 C1=7 and Cz=26 2. f(n) = 83n + 84n & sl(n) If  $c_1 \times g_1(n) \leq c_1 \times g_2(n)$  Then  $f(n) \in O(g(n))$  $: F(n) = \theta(n^3) \text{ where } c_1 = 83 \text{ and } c_2 = 167 \text{ with } n \ge 1$ 

 $2^{n+1} = O(2^n)$  is True because  $2^{n+1} = 2 \cdot 2^n \le c \cdot 2^n$ , where  $c \ge 2$ 

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Check equalities (True / False):
  (1) 5n - 6n € O(n) + 1
  di, n! e o(m)
  (11), 2n2" + nlogn & D(n'z")
  (N) \(\frac{1}{2}\) \(\frac{1}{2}\) \(\frac{1}{2}\) \(\frac{1}{2}\) \(\frac{1}{2}\) \(\frac{1}{2}\)
  (v) n E O(n3) N 3 1 x N1 (0)0
  (vi) 2^n \in \Theta(2^{n+1})
  (vil) n! \in \Theta((n+1)!)
sol = (i) 5n - 6n & O(n) True beaux
    f(n)= 5n -6n
     g(n) = n~
 and f(n) < (x g(n) istrue.
 (11) f(n) = n! = 1 x 2 x 3 x 4 x .... h
       g(n)= n= n*n*n x .... *n
      F(n) < c * g(n) is true.
      Here ni & O(n) is false.
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 $g(n) = n^2 n$ 

(iii) 
$$f(n) = 2n^{2}2^{n} + n\log n \in \Theta(n^{2}n^{2}) \propto g(n) = n^{2}n^{2}$$

If  $n = 16$ , then

 $f(n) = 2(16)^{n}(2)^{16} + 16\log 16$ 
 $g(n) = (16)^{n}(2)^{16}$ 

This shows that,

$$f(n) \geq 2 \propto g(n)$$

If  $e(n) \geq 2 \propto g(n)$ 

$$f(n) \geq 2 \propto g(n)$$

$$f(n) = 2 \sim 2^{n} + n\log n \in \Theta(n^{2}n^{2}) \text{ is false}$$

(iv)

$$f(n) = 2 \sim 2^{n} + n\log n \in \Theta(n^{2}n^{2}) \text{ is false}$$

$$g(n) = 2 \sim 2^{n} + n\log n \in \Theta(n^{2}n^{2}) \text{ is false}$$

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$$f(n) =$$

(vii)  $n! \in \Theta((n+1)!)$   $\longrightarrow$  Falx n1 = 1×2×3 ... ×n  $(n+1)! = (n+1) \times n \times \dots 2 \times 1$ .. n! ∈ O(n+1)! is False Prove or disprove that F(n) = 1+2+3+ .... +nEO(n  $\frac{50}{2}$  let  $F(n) = 1 + 2 + 3 + ..., n = \frac{n(n+1)}{2} = \frac{1}{2}n^{2}$ ofg(n)= non a gold + cons andi By definition F(n) E O (g(n)) if  $c_1 * g(n) \leq f(n) \leq c_2 * g(n)$ If C1= = and n=2 Then  $\lceil c_1 \times g(n) \leq f(n) \rceil \Rightarrow \frac{1}{4}(2^n) \leq (2^n)$  is Trut If Cz=1 and n=2 (FCn) = (2\*g(n)) => \frac{1}{2}(2\*) \le 2\* is True As both the cases are true, it proves that 1+ 2+3+ .... + h & O(n).