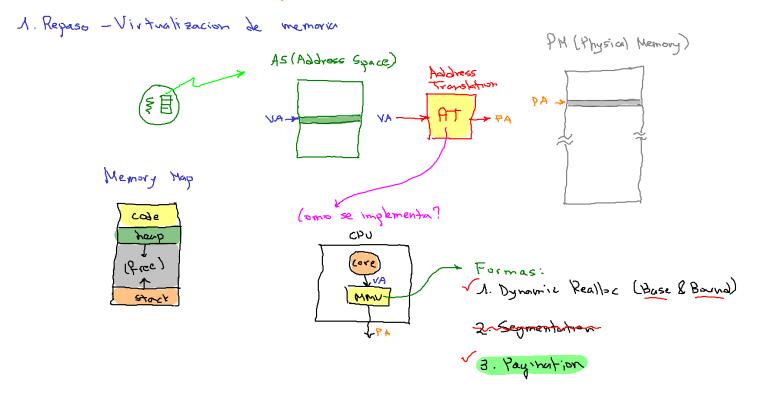
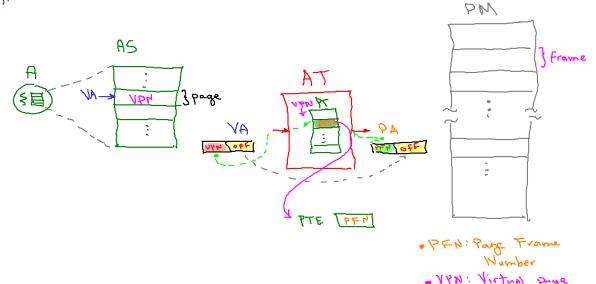
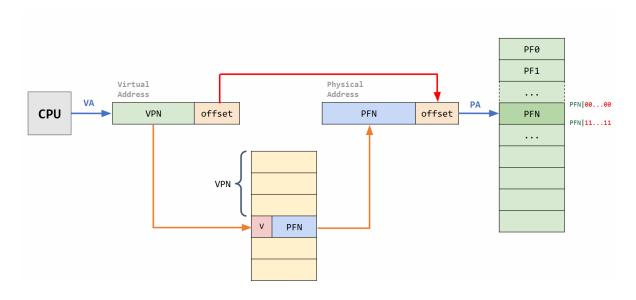
# 06/05/2025 - Sistemas Operativos - Ude@

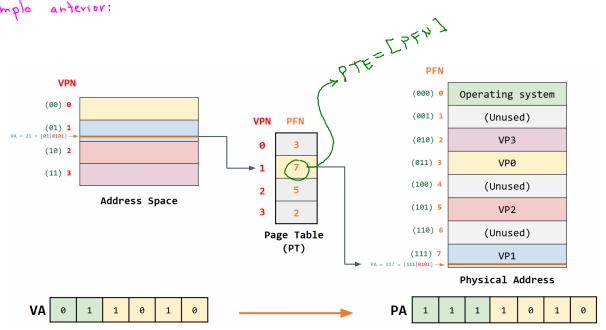




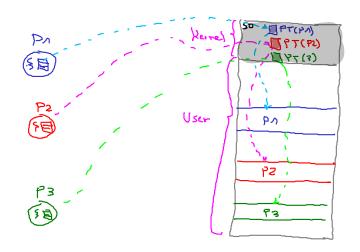


# Mecanisma de paginación:





## PT. -> Problema: Overhead de gasto de memoria por las 3.



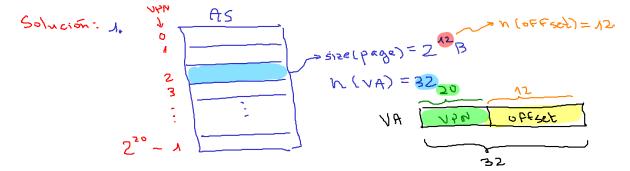
Suponga que se tiene una memoria virtual (Address Space) con las siguientes características:

gagnerion?

- 32 bits de direcciones.  $\rightarrow h(VA) = 32 bits$
- ~ size (Page) = 4KB = 4x 1024B = 22210B = 212B Páginas de 4KB
- Asuma que cada PTE tiene un tamaño de 4B 🦴 4 SIZe (PTE) = 48

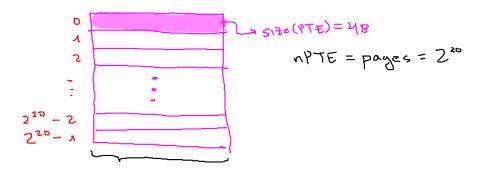
Responda las siguientes preguntas:

- VA VPW OFFSET ? 1. ¿Cual es el formato de las VA en este caso? —
- 2. ¿Cual es el tamaño de la tabla de página (PT)? 51ze (4T)=?

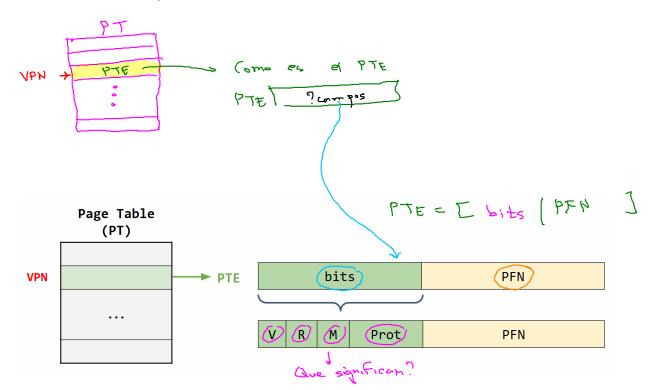


pages = 
$$\frac{5ize(AS)}{5ize(page)} = \frac{2^{32}}{2^{12}} = 2^{20}$$
 pages.

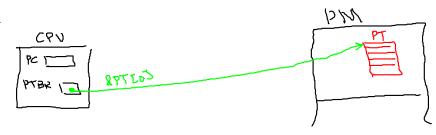
### 2. SIZe (PT)



## 4. Como es una PT?



## 5. Registro PTBR



```
// Extract the VPN from the virtual address
VPN = (VirtualAddress & VPN_MASK) >> SHIFT

// Form the address of the page-table entry (PTE)
PTEAddr = PTBR + (VPN * sizeof(PTE))

// Fetch the PTE
PTE = AccessMemory(PTEAddr)

// Check if process can access the page
if (PTE.Valid == False)
RaiseException(SEGMENTATION_FAULT)
else if (CanAccess(PTE.ProtectBits) == False)
RaiseException(PROTECTION_FAULT)
else
// Access is OK: form physical address and fetch it
offset = VirtualAddress & OFFSET_MASK
PhysAddr = (PTE.PFN << PFN_SHIFT) | offset
Register = AccessMemory(PhysAddr)
```

