

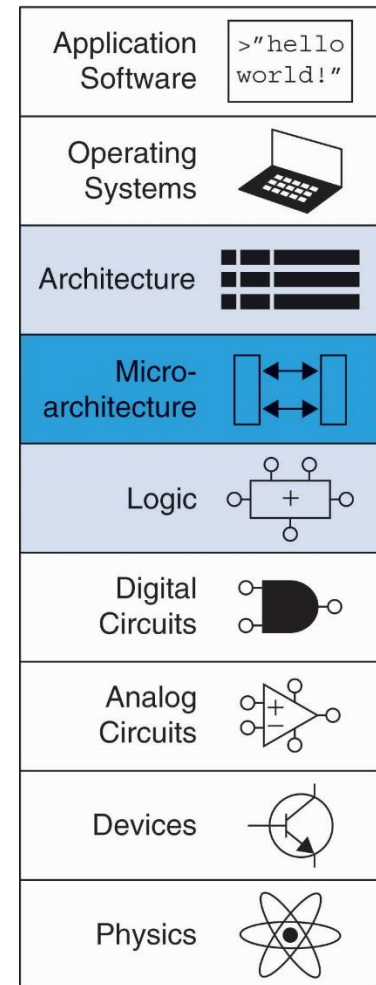
Digital Design & Computer Architecture

Sarah Harris & David Harris

Chapter 7: Microarchitecture

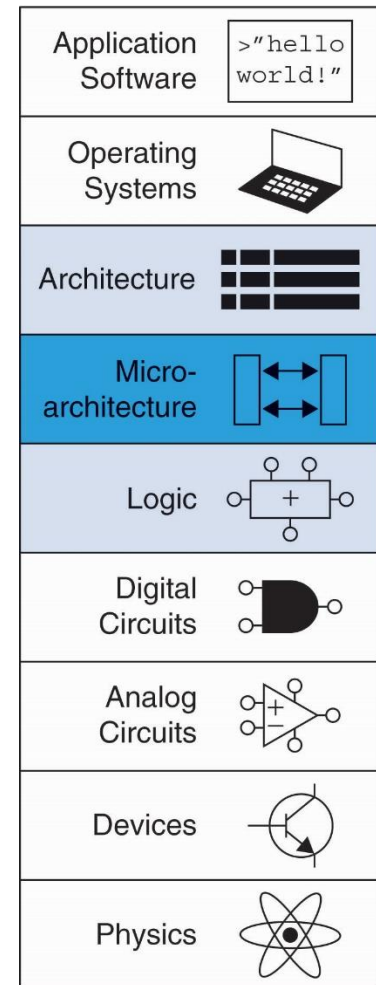
Chapter 7 :: Topics

- Introduction
- Performance Analysis
- Single-Cycle Processor
- Multicycle Processor
- Pipelined Processor
- Advanced Microarchitecture



Introduction

- **Microarchitecture:** how to implement an architecture in hardware
- Processor:
 - **Datapath:** functional blocks
 - **Control:** control signals



Microarchitecture

- **Multiple implementations** for a single architecture:
 - **Single-cycle:** Each instruction executes in a single cycle
 - **Multicycle:** Each instruction is broken up into series of shorter steps
 - **Pipelined:** Each instruction broken up into series of steps & multiple instructions execute at once

Processor Performance

- **Program execution time**

Execution Time = (#instructions)(cycles/instruction)(seconds/cycle)

- **Definitions:**

- CPI: Cycles/instruction
- clock period: seconds/cycle
- IPC: instructions/cycle = IPC

- **Challenge is to satisfy constraints of:**

- Cost
- Power
- Performance

RISC-V Processor

- Consider **subset** of RISC-V instructions:
 - R-type ALU instructions:
 - **add, sub, and, or, slt**
 - Memory instructions:
 - **lw, sw**
 - Branch instructions:
 - **beq**

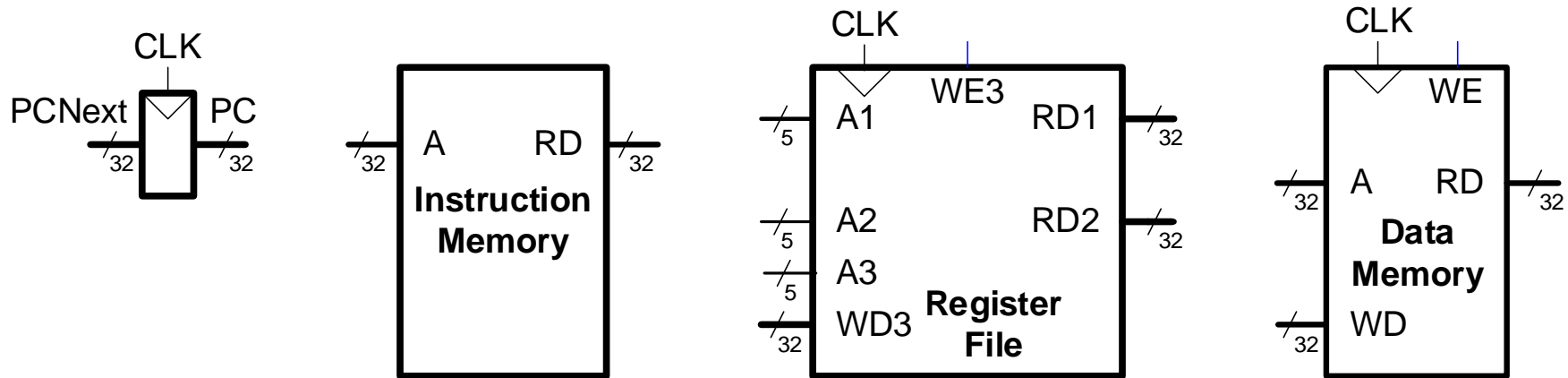
Architectural State Elements

Determines everything about a processor:

- **Architectural state:**

- 32 registers
- PC
- Memory

RISC-V Architectural State Elements



Chapter 7: Microarchitecture

Single-Cycle RISC-V Processor

Single-Cycle RISC-V Processor

- Datapath
- Control

Example Program

- Design datapath
- View example program executing

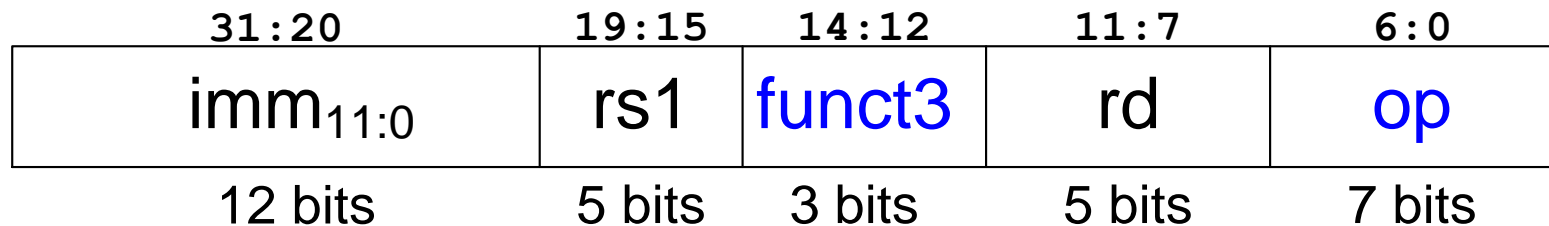
Example Program:

Address	Instruction	Type	Fields					Machine Language	
0x1000	L7: lw x6, -4(x9)	I	imm _{11:0}	rs1	f3	rd	op		
			1111111111100	01001	010	00110	0000011	FFC4A303	
0x1004	sw x6, 8(x9)	S	imm _{11:5}	rs2	rs1	f3	imm _{4:0}	op	
			0000000 00110	01001	010	01000	0100011	0064A423	
0x1008	or x4, x5, x6	R	funct7	rs2	rs1	f3	rd	op	
			0000000 00110	00101	110	00100	0110011	0062E233	
0x100C	beq x4, x4, L7	B	imm _{12,10:5}	rs2	rs1	f3	imm _{4:1,11}	op	
			1111111 00100	00100	000	10101	1100011	FE420AE3	

Single-Cycle RISC-V Processor

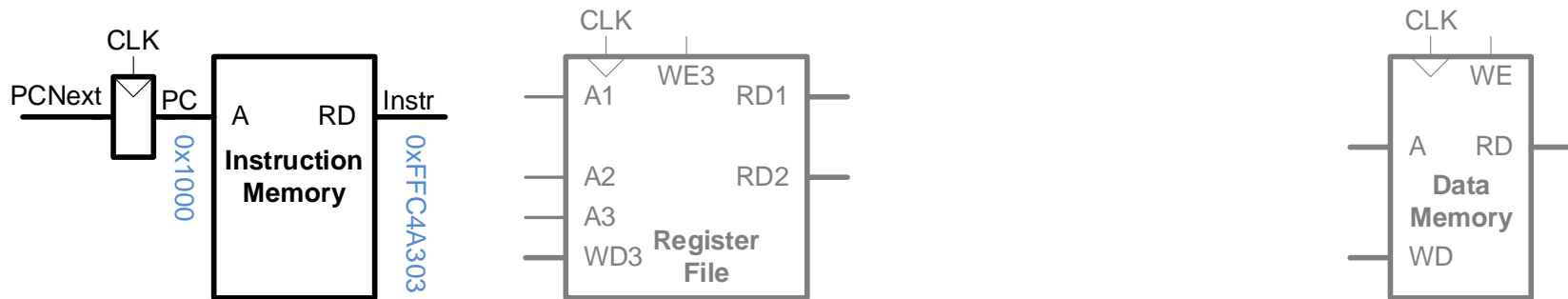
- **Datapath:** start with `lw` instruction
- **Example:** `lw x6, -4(x9)`
`lw rd, imm(rs1)`

I-Type



Single-Cycle Datapath: l_w fetch

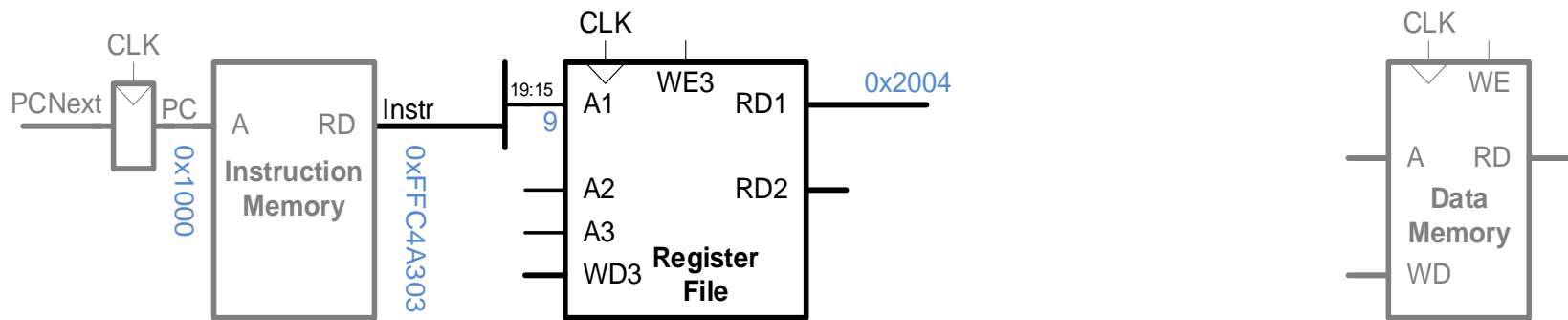
STEP 1: Fetch instruction



Address	Instruction	Type	Fields				Machine Language	
0x1000	L7: lw x6, -4(x9)	I	imm _{11:0} 111111111100	rs1 01001	f3 010	rd 00110	op 0000011	FFC4A303

Single-Cycle Datapath: lw Reg Read

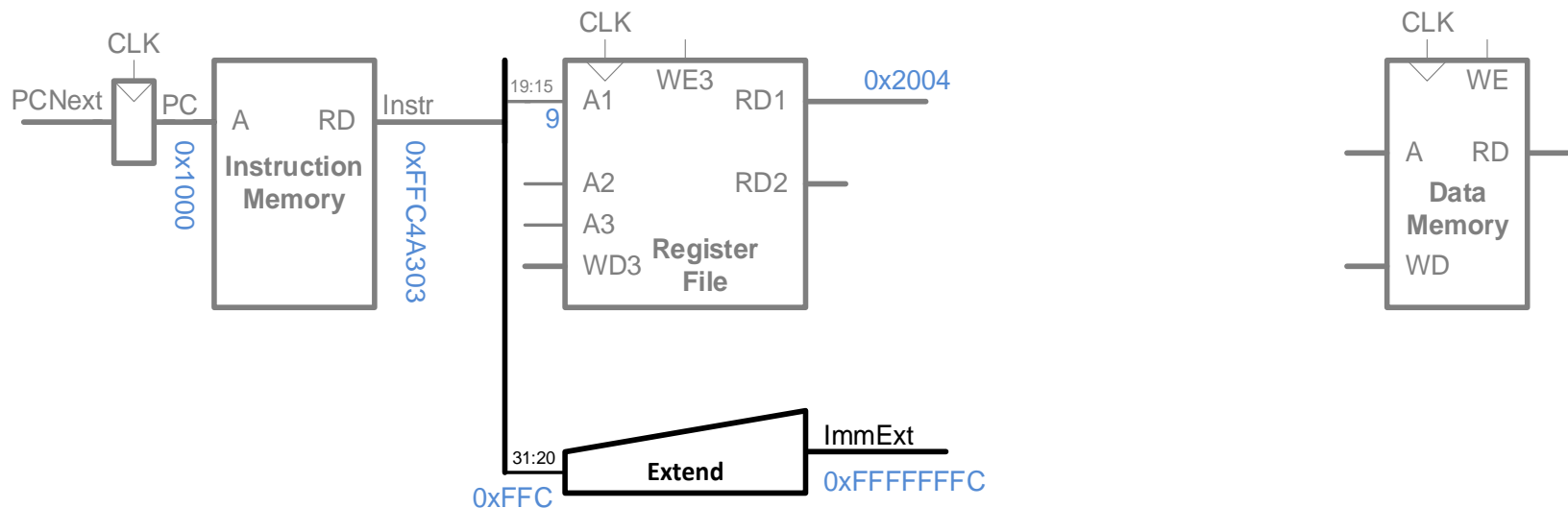
STEP 2: Read source operand (**rs1**) from RF



Address	Instruction	Type	Fields	Machine Language
0x1000	L7: lw x6, -4(x9)	I	<div><div>imm_{11:0}</div><div>111111111100</div><div>rs1</div><div>01001</div><div>f3</div><div>010</div><div>rd</div><div>00110</div><div>op</div><div>0000011</div></div>	FFC4A303

Single-Cycle Datapath: 1w Immediate

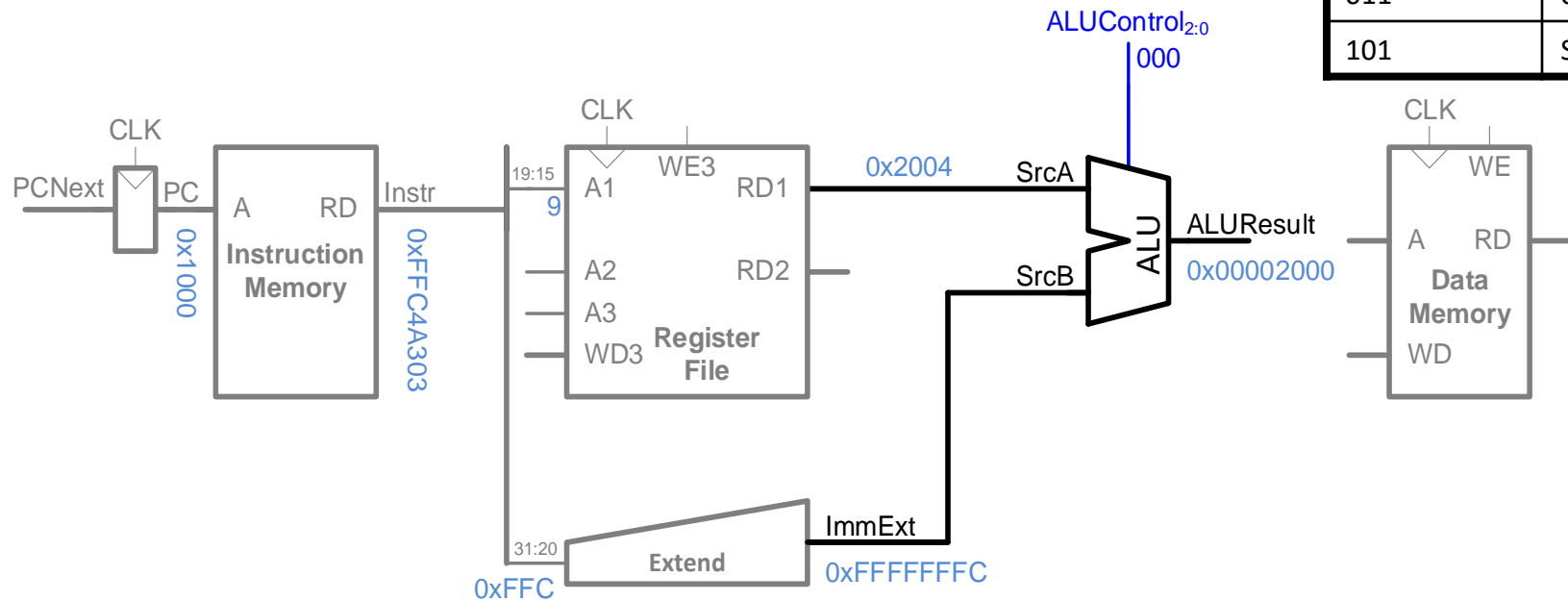
STEP 3: Extend the immediate



Address	Instruction	Type	Fields	Machine Language
0x1000	L7: lw x6, -4(x9)	I	<div><div>imm_{11:0}</div><div>1111111111100</div></div> <div><div>rs1</div><div>01001</div></div> <div><div>f3</div><div>010</div></div> <div><div>rd</div><div>00110</div></div> <div><div>op</div><div>0000011</div></div>	FFC4A303

Single-Cycle Datapath: lw Address

STEP 4: Compute the memory address

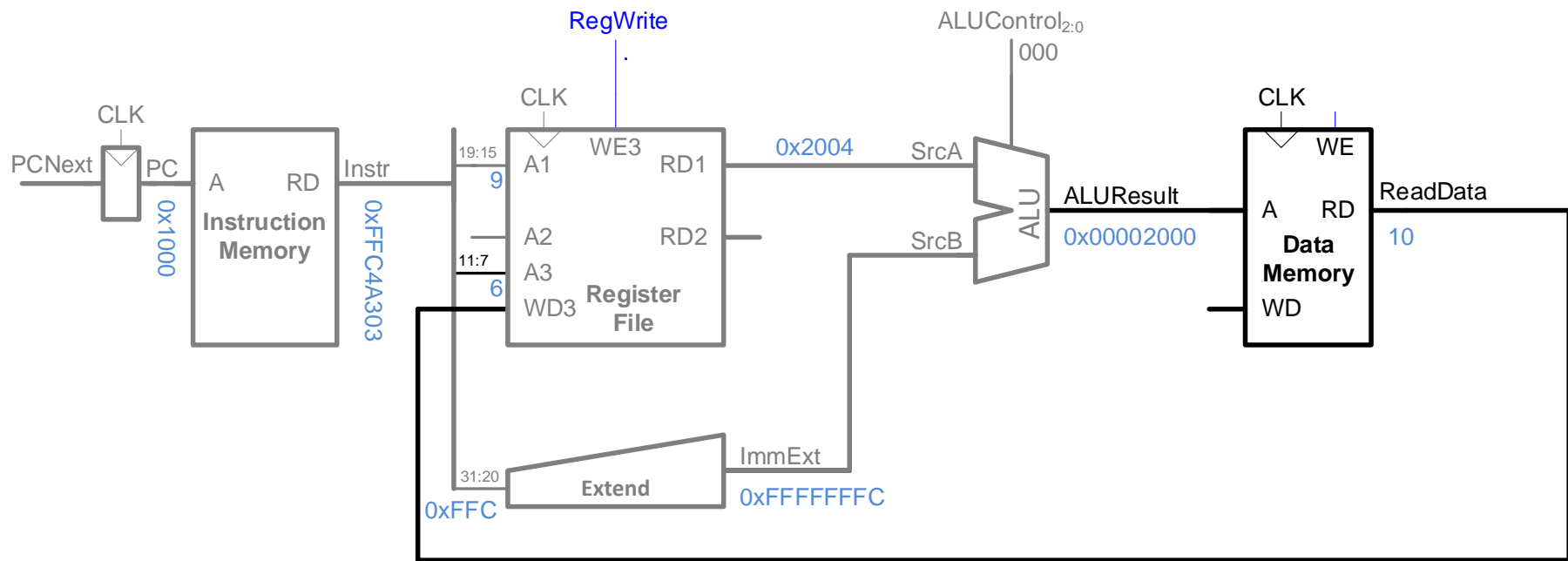


ALUControl _{2:0}	Function
000	add
001	subtract
010	and
011	or
101	SLT

Address	Instruction	Type	Fields	Machine Language										
0x1000	L7: lw x6, -4(x9)	I	<table><tr><th>imm_{11:0}</th><th>rs1</th><th>f3</th><th>rd</th><th>op</th></tr><tr><td>1111111111100</td><td>01001</td><td>010</td><td>00110</td><td>0000011</td></tr></table>	imm _{11:0}	rs1	f3	rd	op	1111111111100	01001	010	00110	0000011	FFC4A303
imm _{11:0}	rs1	f3	rd	op										
1111111111100	01001	010	00110	0000011										

Single-Cycle Datapath: lw Mem Read

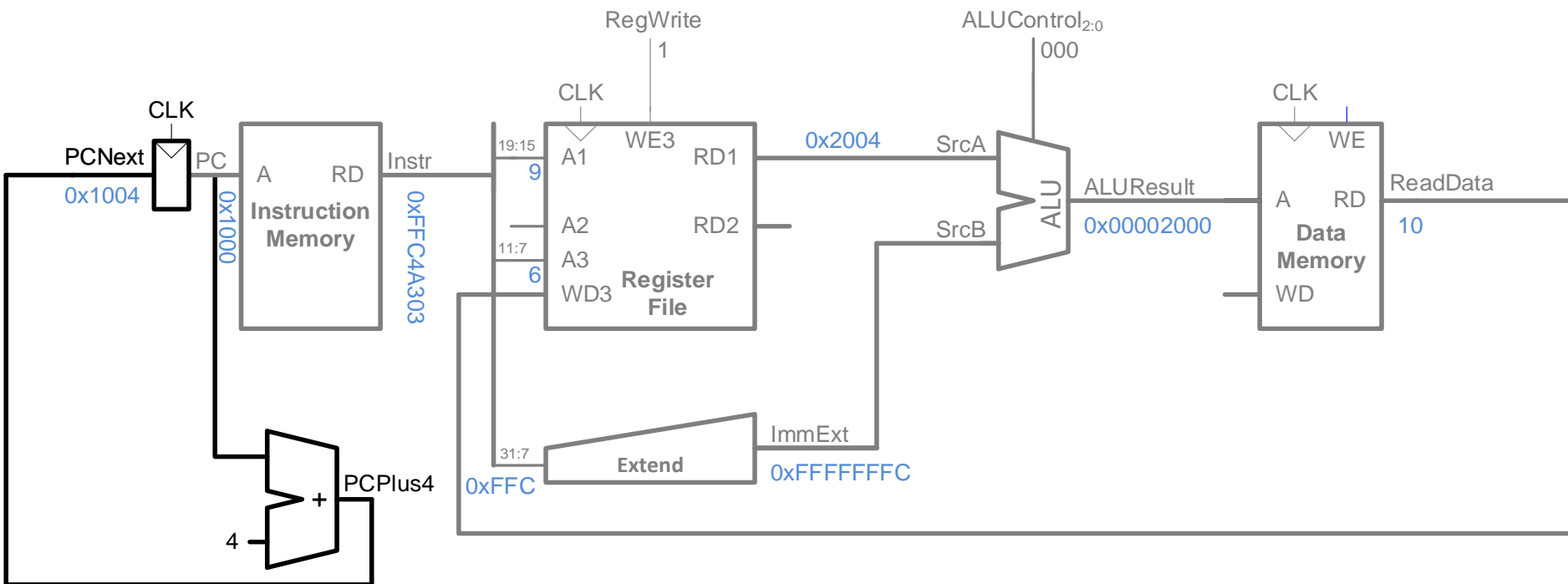
STEP 5: Read data from memory and write it back to register file



Address	Instruction	Type	Fields	Machine Language										
0x1000	L7: lw x6, -4(x9)	I	<table><tr><th>imm_{11:0}</th><th>rs1</th><th>f3</th><th>rd</th><th>op</th></tr><tr><td>111111111100</td><td>01001</td><td>010</td><td>00110</td><td>0000011</td></tr></table>	imm _{11:0}	rs1	f3	rd	op	111111111100	01001	010	00110	0000011	FFC4A303
imm _{11:0}	rs1	f3	rd	op										
111111111100	01001	010	00110	0000011										

Single-Cycle Datapath: PC Increment

STEP 6: Determine address of next instruction



Address	Instruction	Type	Fields			Machine Language	
			imm _{11:0}	rs1	f3	rd	op
0x1000	L7: lw x6, -4(x9)	I	111111111100	01001	010	00110	0000011
							FFC4A303

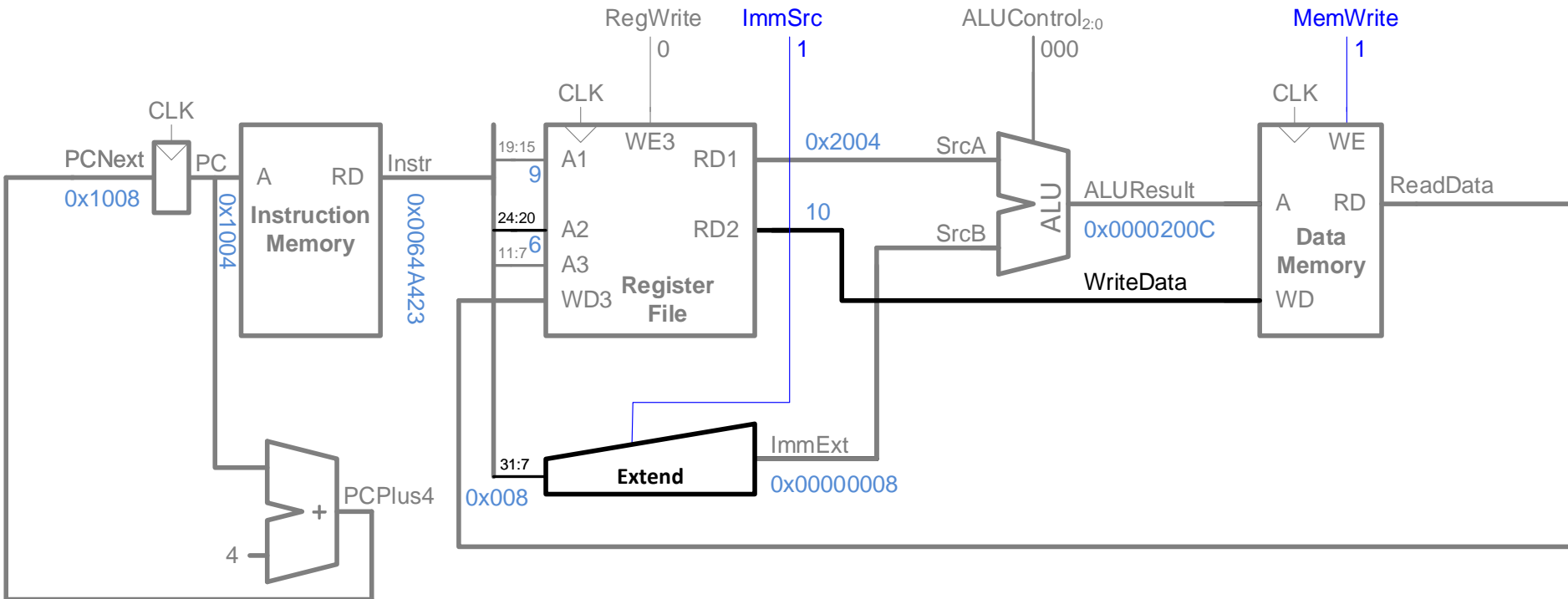
Chapter 7: Microarchitecture

Single-Cycle

**Datapath: Other
Instructions**

Single-Cycle Datapath: sw

- **Immediate:** now in {instr[31:25], instr[11:7]}
- **Add control signals:** ImmSrc, MemWrite

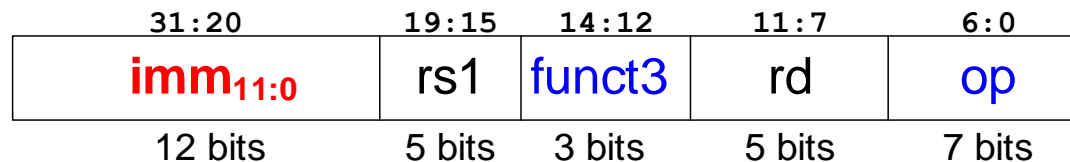


Address	Instruction	Type	Fields					Machine Language	
			imm _{11:5}	rs2	rs1	f3	imm _{4:0}	op	
0x1004	sw x6, 8(x9)	S	0000000	00110	01001	010	01000	0100011	0064A423

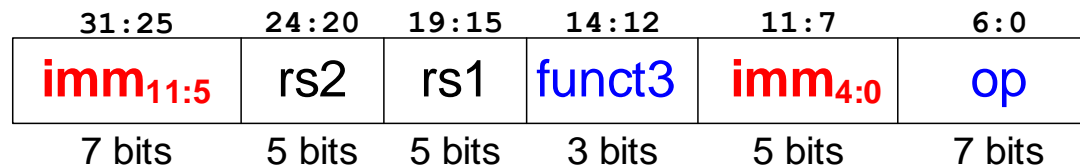
Single-Cycle Datapath: Immediate

ImmSrc	ImmExt	Instruction Type
0	{{20{instr[31]}}, instr[31:20] }	I-Type
1	{{20{instr[31]}}, instr[31:25], instr[11:7] }	S-Type

I-Type

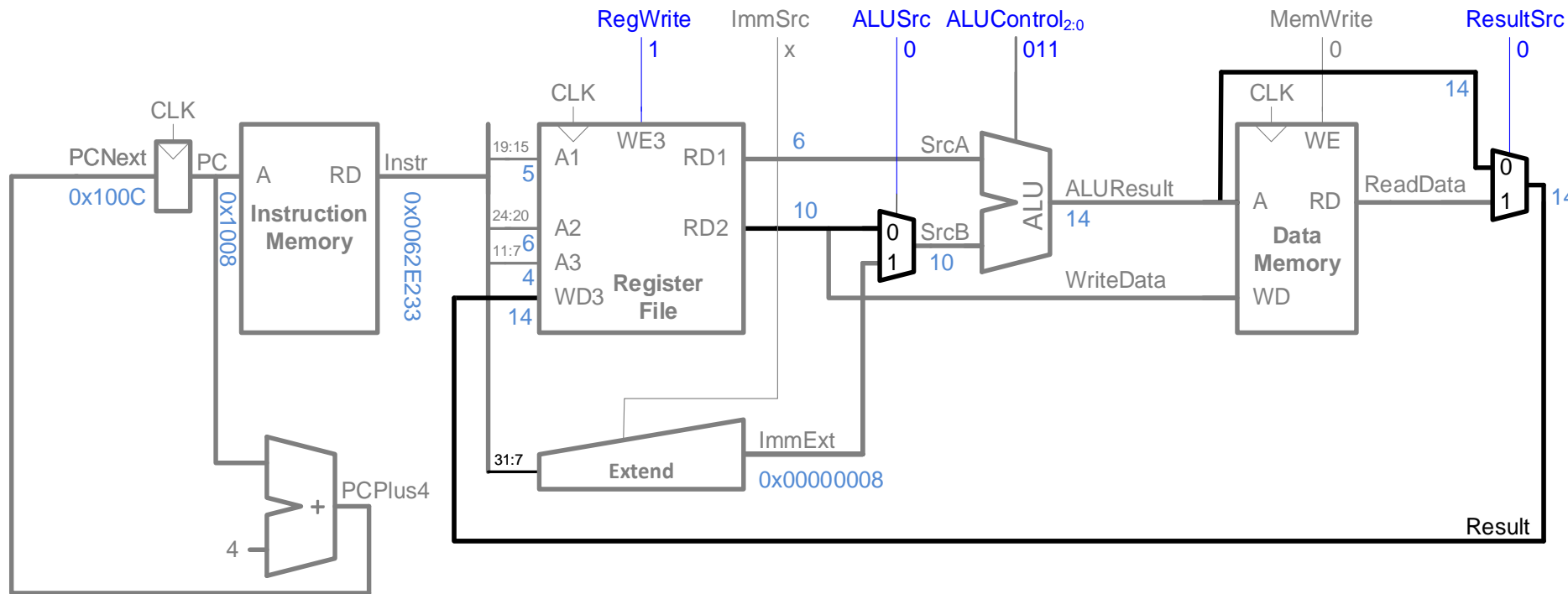


S-Type



Single-Cycle Datapath: R-type

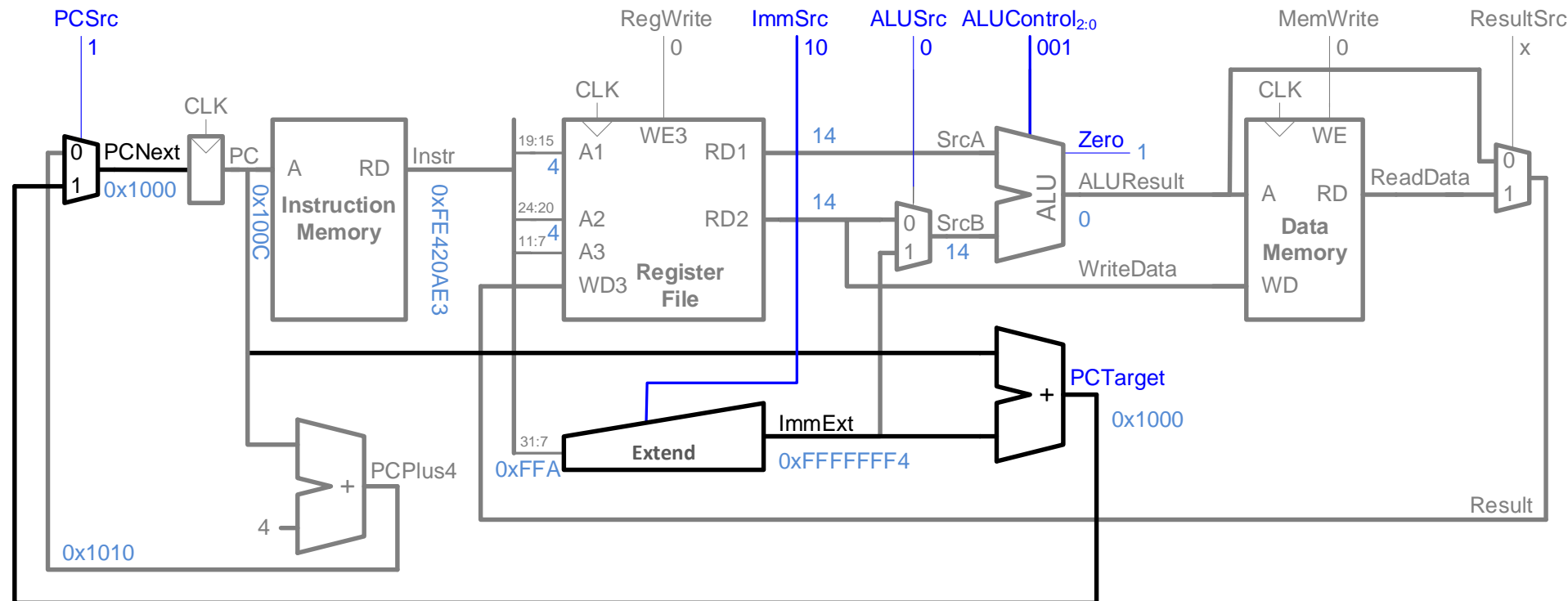
- Read from **rs1** and **rs2** (instead of **imm**)
- Write *ALUResult* to **rd**



Address	Instruction	Type	Fields					Machine Language	
0x1008	<code>or x4, x5, x6</code>	R	funct7	rs2	rs1	f3	rd	op	0062E233
			0000000	00110	00101	110	00100	0110011	

Single-Cycle Datapath: beq

Calculate **target address**: $PCTarget = PC + imm$

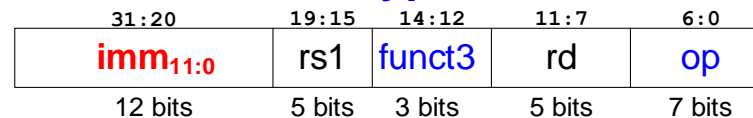


Address	Instruction	Type	Fields					Machine Language	
			imm _{12,10:5}	rs2	rs1	f3	imm _{4:1,11}	op	
0x100C	beq x4, x4, L7	B	1111111	00100	00100	000	10101	1100011	FE420AE3

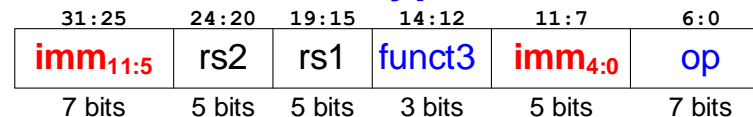
Single-Cycle Datapath: ImmExt

ImmSrc _{1:0}	ImmExt	Instruction Type
00	{{20{instr[31]}}, instr[31:20] }	I-Type
01	{{20{instr[31]}}, instr[31:25], instr[11:7] }	S-Type
10	{{19{instr[31]}}, instr[31], instr[7], instr[30:25], instr[11:8], 1'b0 }	B-Type

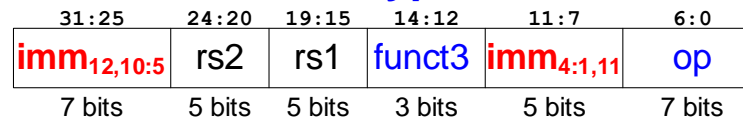
I-Type



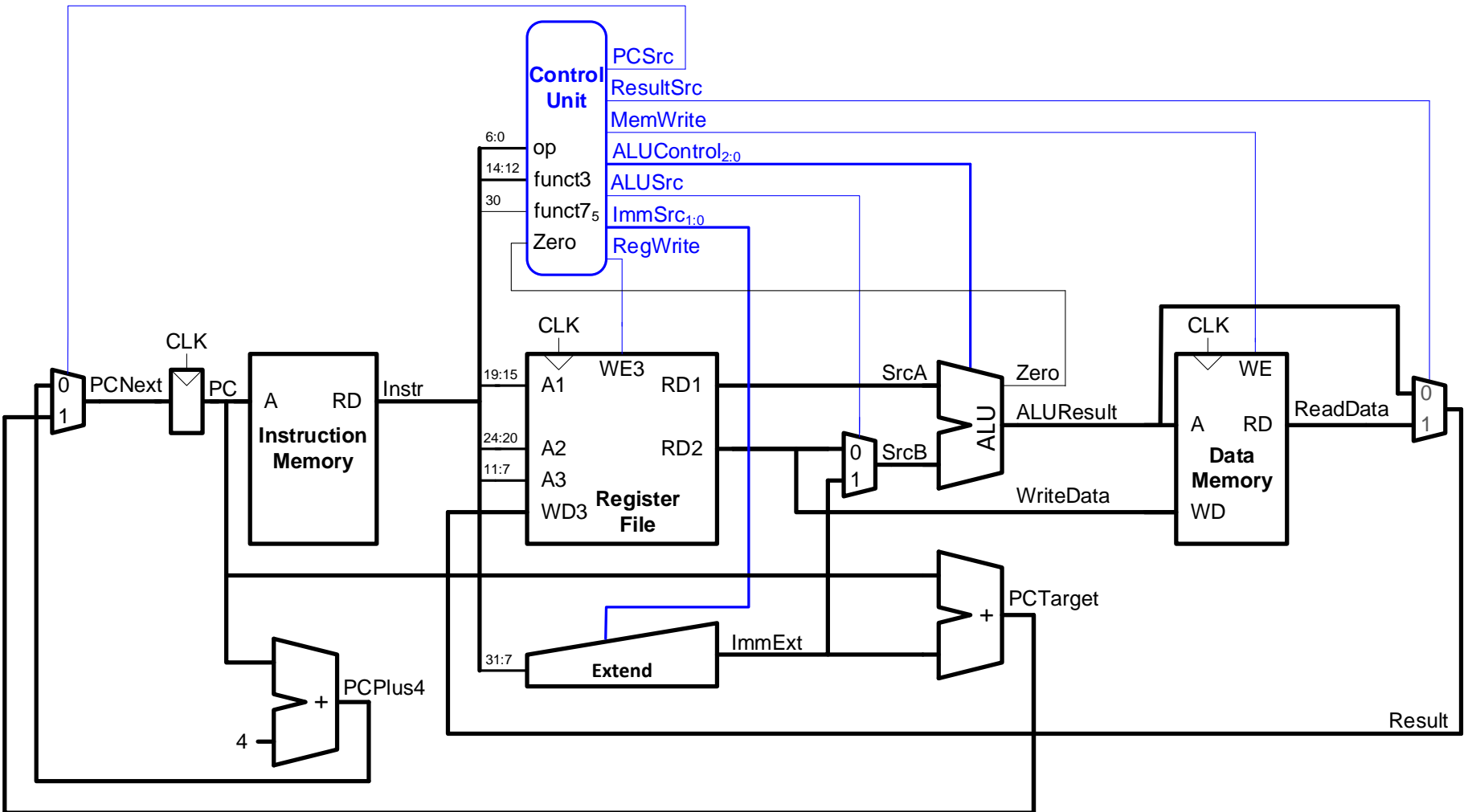
S-Type



B-Type



Single-Cycle RISC-V Processor

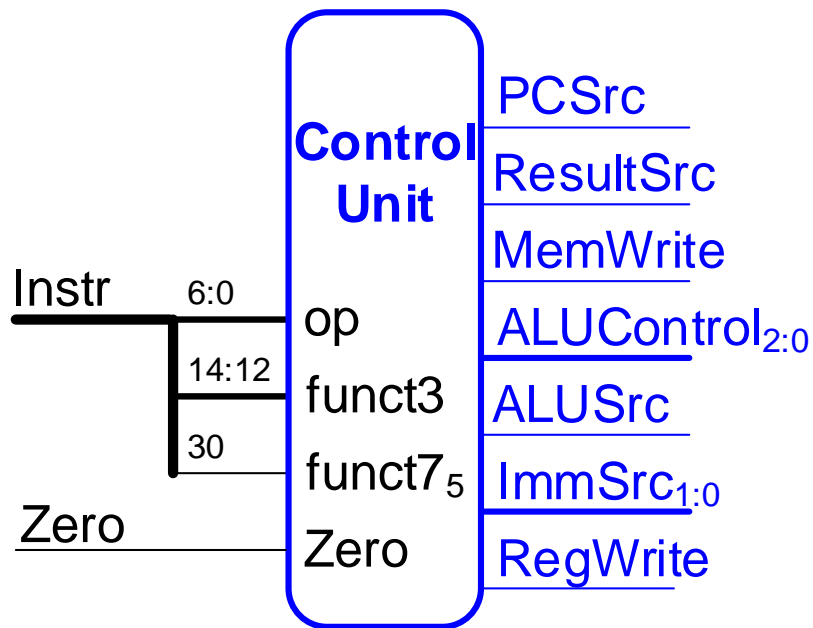


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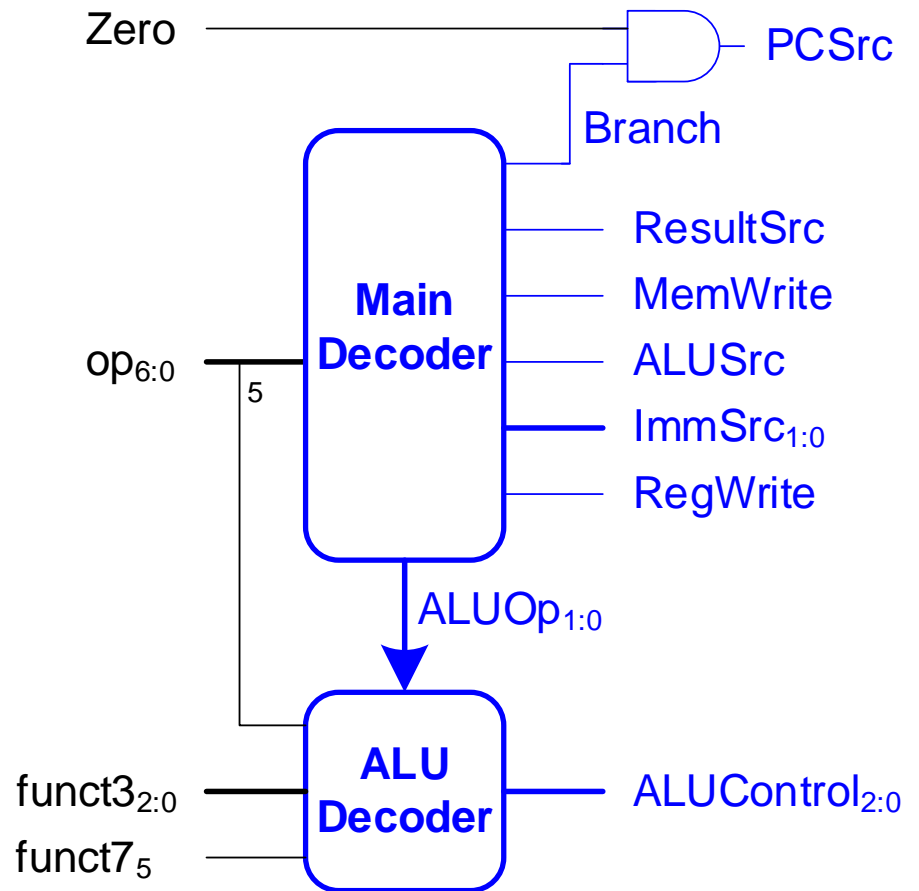
Single-Cycle Control

Single-Cycle Control

High-Level View

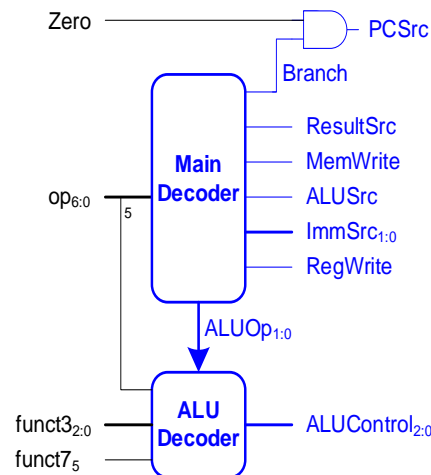


Low-Level View



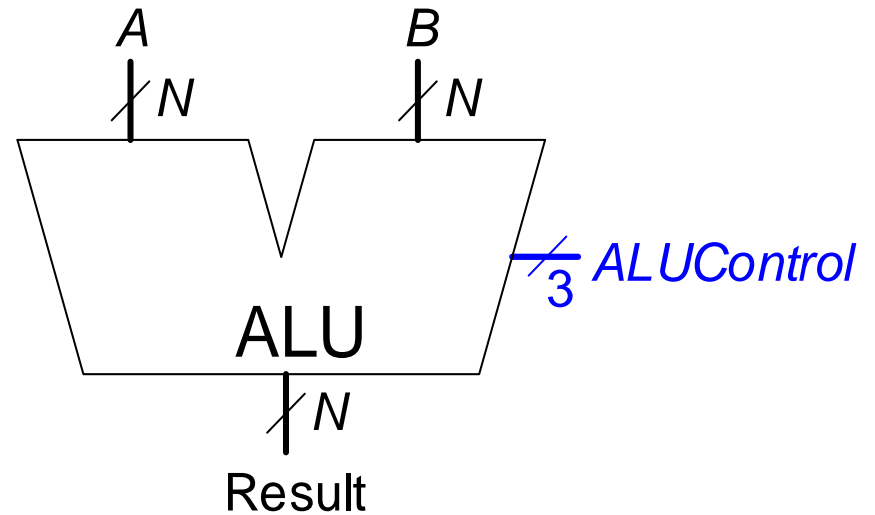
Single-Cycle Control: Main Decoder

op	Instr.	RegWrite	ImmSrc	ALUSrc	MemWrite	ResultSrc	Branch	ALUOp
3	lw							
35	sw							
51	R-type							
99	beq							



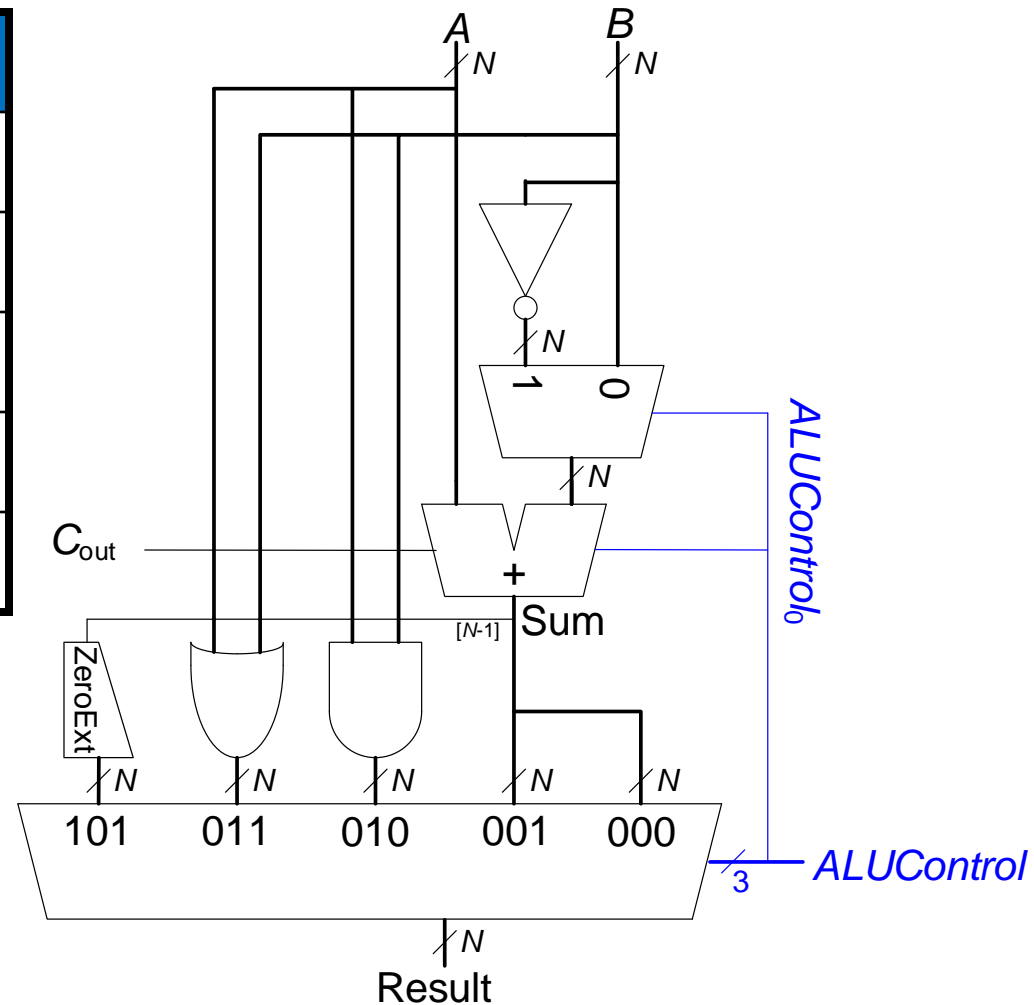
Review: ALU

ALUControl _{2:0}	Function
000	add
001	subtract
010	and
011	or
101	SLT

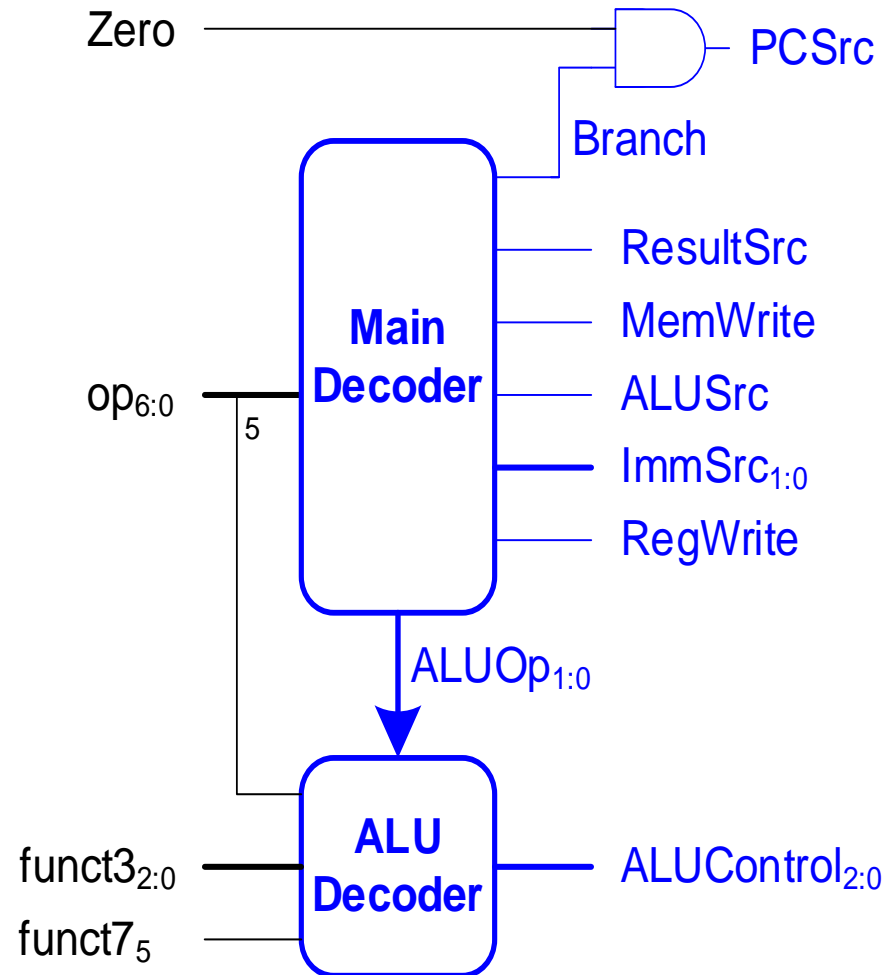


Review: ALU

ALUControl _{2:0}	Function
000	add
001	subtract
010	and
011	or
101	SLT

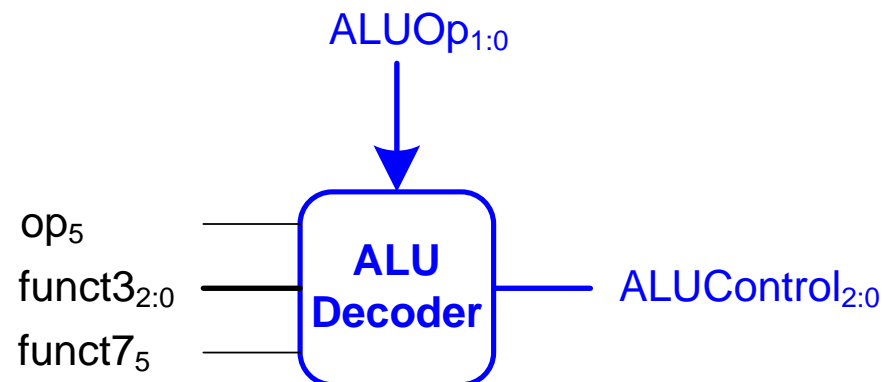


Single-Cycle Control: ALU Decoder



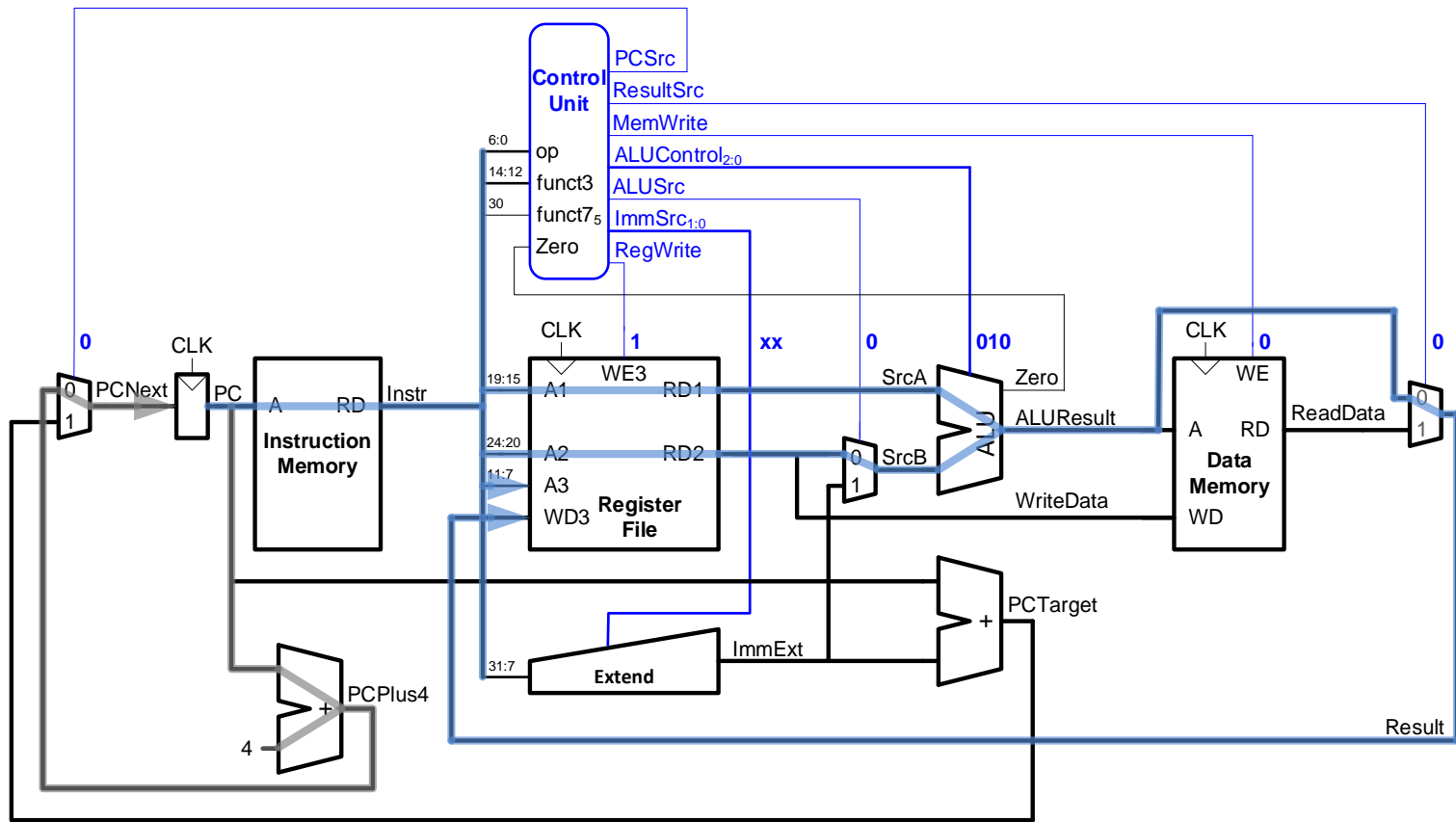
Single-Cycle Control: ALU Decoder

<i>ALUOp</i>	<i>funct3</i>	<i>op₅</i> , <i>funct7₅</i>	Instruction	<i>ALUControl_{2:0}</i>
00	x	x	lw, sw	000 (add)
01	x	x	beq	001 (subtract)
10	000	00, 01, 10	add	000 (add)
	000	11	sub	001 (subtract)
	010	x	slt	101 (set less than)
	110	x	or	011 (or)
	111	x	and	010 (and)



Example: and

op	Instruct	RegWrite	ImmSrc	ALUSrc	MemWrite	ResultSrc	Branch	ALUOp
51	R-type	1	XX	0	0	0	0	10



and x5, x6, x7

Chapter 7: Microarchitecture

Extending the Single-Cycle Processor

Extended Functionality: I-Type ALU

Enhance the single-cycle processor to handle **I-Type ALU instructions**: `addi`, `andi`, `ori`, and `slli`

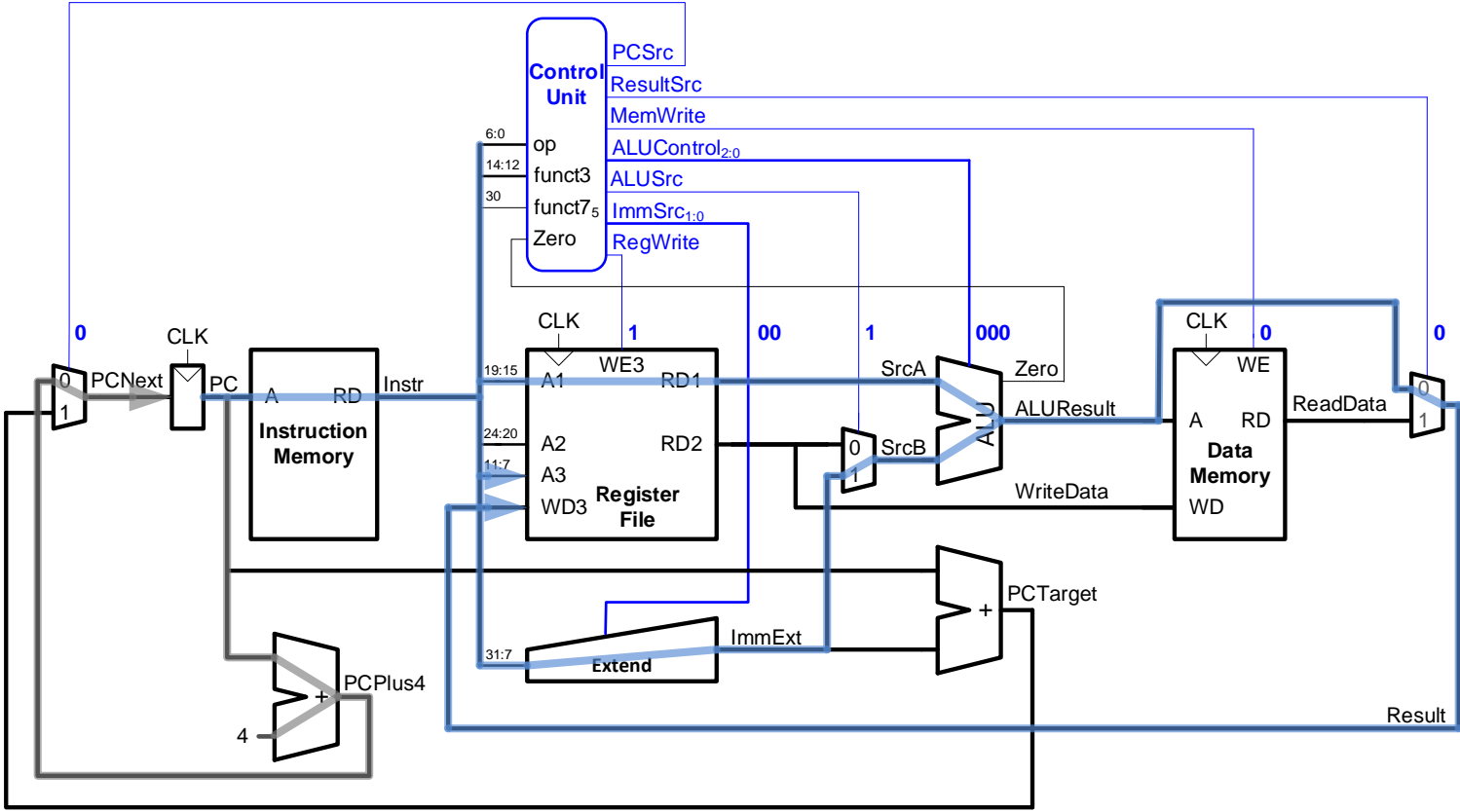
- **Similar to R-type** instructions
- But **second source** comes from **immediate**
- Change ***ALUSrc*** to select the immediate
- And ***ImmSrc*** to pick the correct immediate

Extended Functionality: I-Type ALU

op	Instruct.	RegWrite	ImmSrc	ALUSrc	MemWrite	ResultSrc	Branch	ALUOp
3	lw	1	00	1	0	1	0	00
35	sw	0	01	1	1	X	0	00
51	R-type	1	XX	0	0	0	0	10
99	beq	0	10	0	0	X	1	01
19	I-type	1	00	1	0	0	0	10

Extended Functionality: addi

op	Instruct.	RegWrite	ImmSrc	ALUSrc	MemWrite	ResultSrc	Branch	ALUOp
19	I-type	1	00	1	0	0	0	10



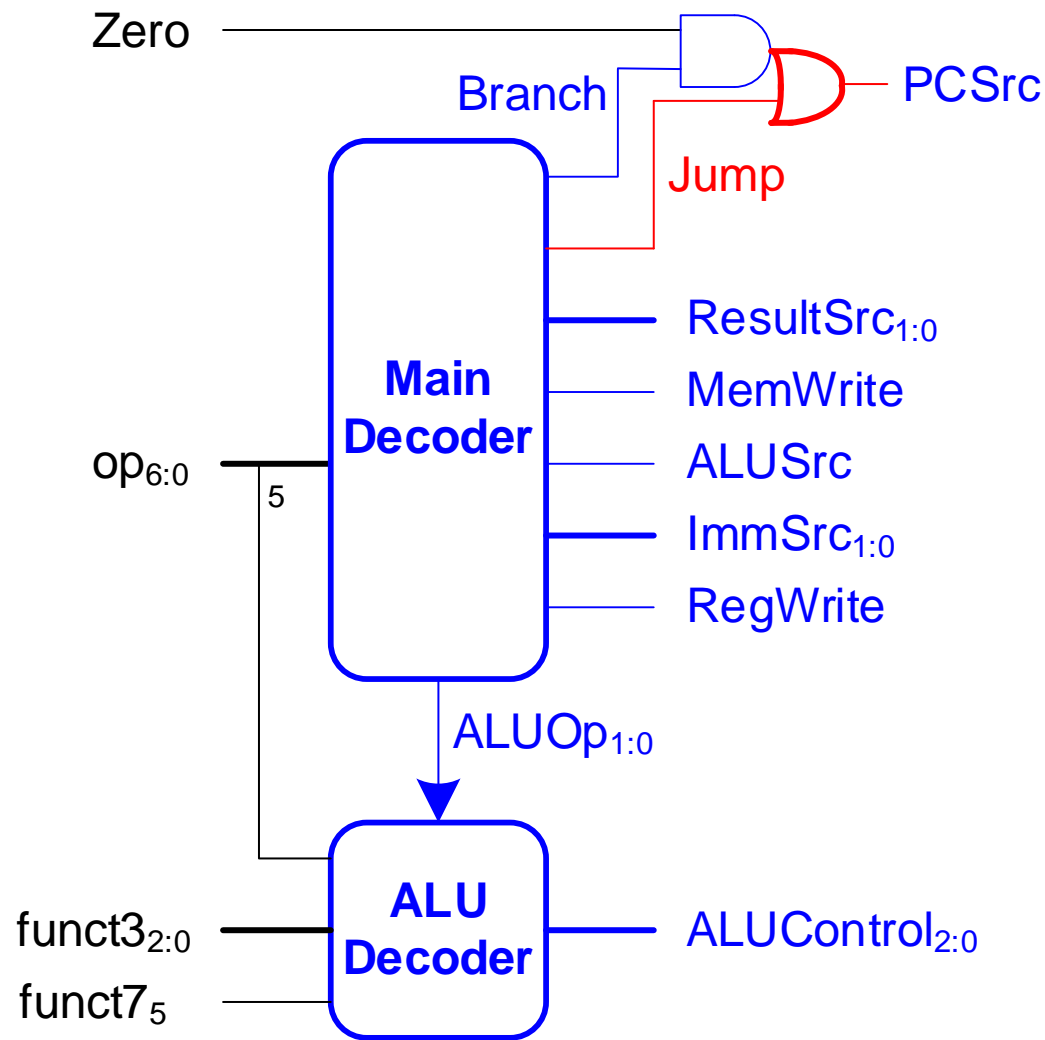
```
addi x5, x6, -33
```

Extended Functionality: `jal`

Enhance the single-cycle processor to handle `jal`

- **Similar to `beq`**
- But jump is **always taken**
 - *PCSrc* should be 1
- **Immediate format** is different
 - Need a new *ImmSrc* of 11
- And `jal` must **compute $PC+4$** and **store in `rd`**
 - Take $PC+4$ from adder through ResultMux

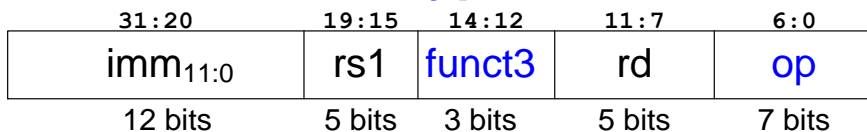
Extended Functionality: jal



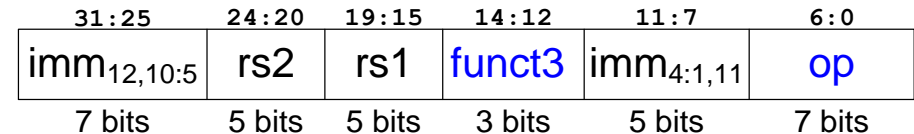
Extended Functionality: *ImmExt*

ImmSrc _{1:0}	ImmExt	Instruction Type
00	{{20{instr[31]}}, instr[31:20]}	I-Type
01	{{20{instr[31]}}, instr[31:25], instr[11:7]}	S-Type
10	{{19{instr[31]}}, instr[31], instr[7], instr[30:25], instr[11:8], 1'b0}	B-Type
11	{{12{instr[31]}}, instr[19:12], instr[20], instr[30:21], 1'b0}	J-Type

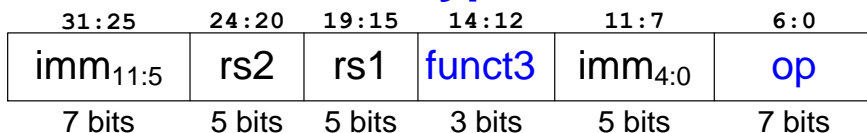
I-Type



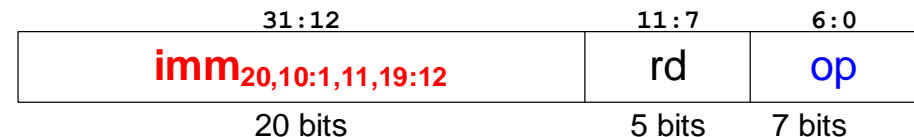
B-Type



S-Type



J-Type

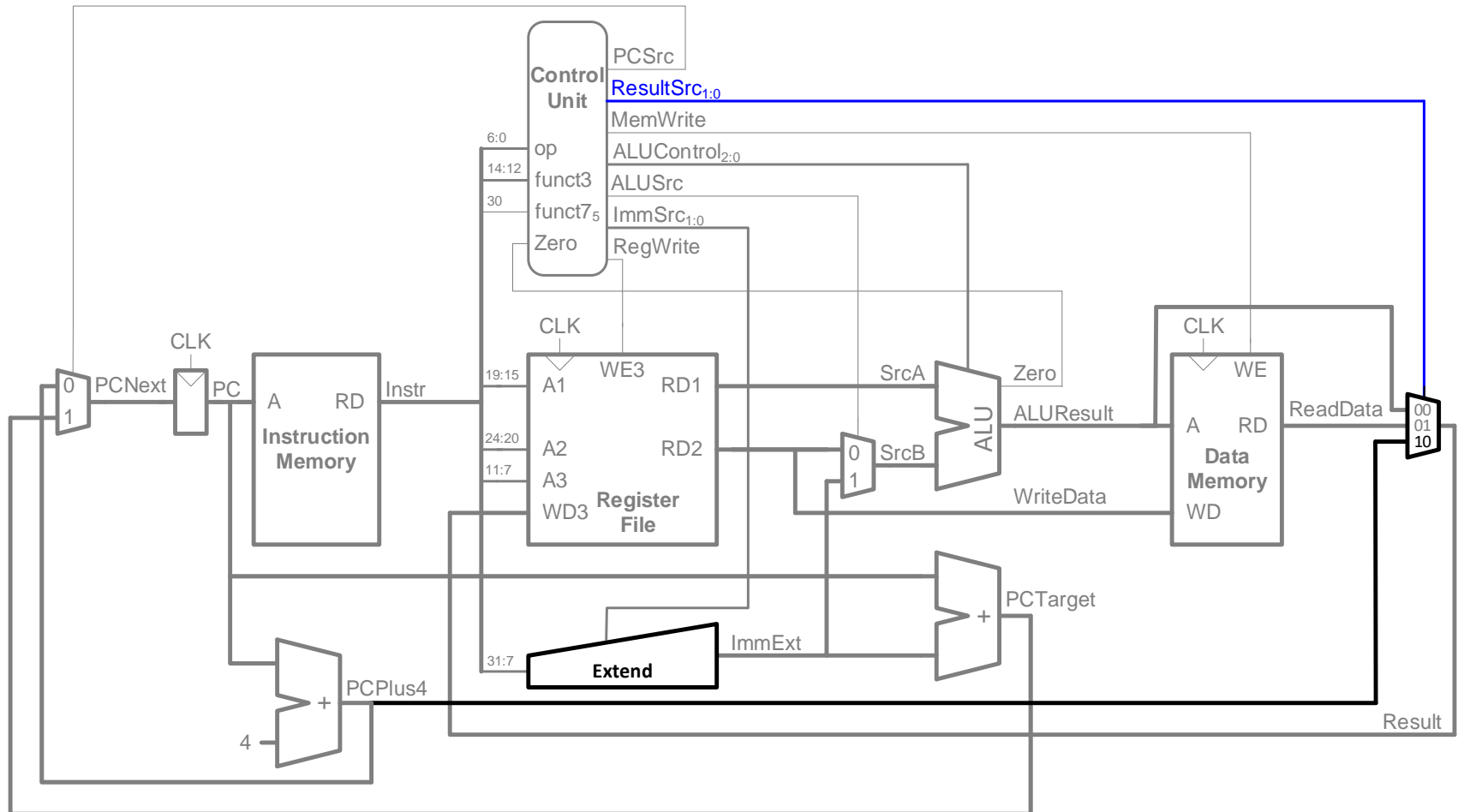


Extended Functionality: jal

op	Instruct.	RegWrite	ImmSrc	ALUSrc	MemWrite	ResultSrc	Branch	ALUOp	Jump
3	lw	1	00	1	0	01	0	00	0
35	sw	0	01	1	1	XX	0	00	0
51	R-type	1	XX	0	0	00	0	10	0
99	beq	0	10	0	0	XX	1	01	0
19	l-type	1	00	1	0	00	0	10	0
111	jal	1	11	X	0	10	0	XX	1

Extended Functionality: jal

op	Instruct.	RegWrite	ImmSrc	ALUSrc	MemWrite	ResultSrc	Branch	ALUOp	Jump
111	jal	1	11	X	0	10	0	XX	1



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Single-Cycle Performance

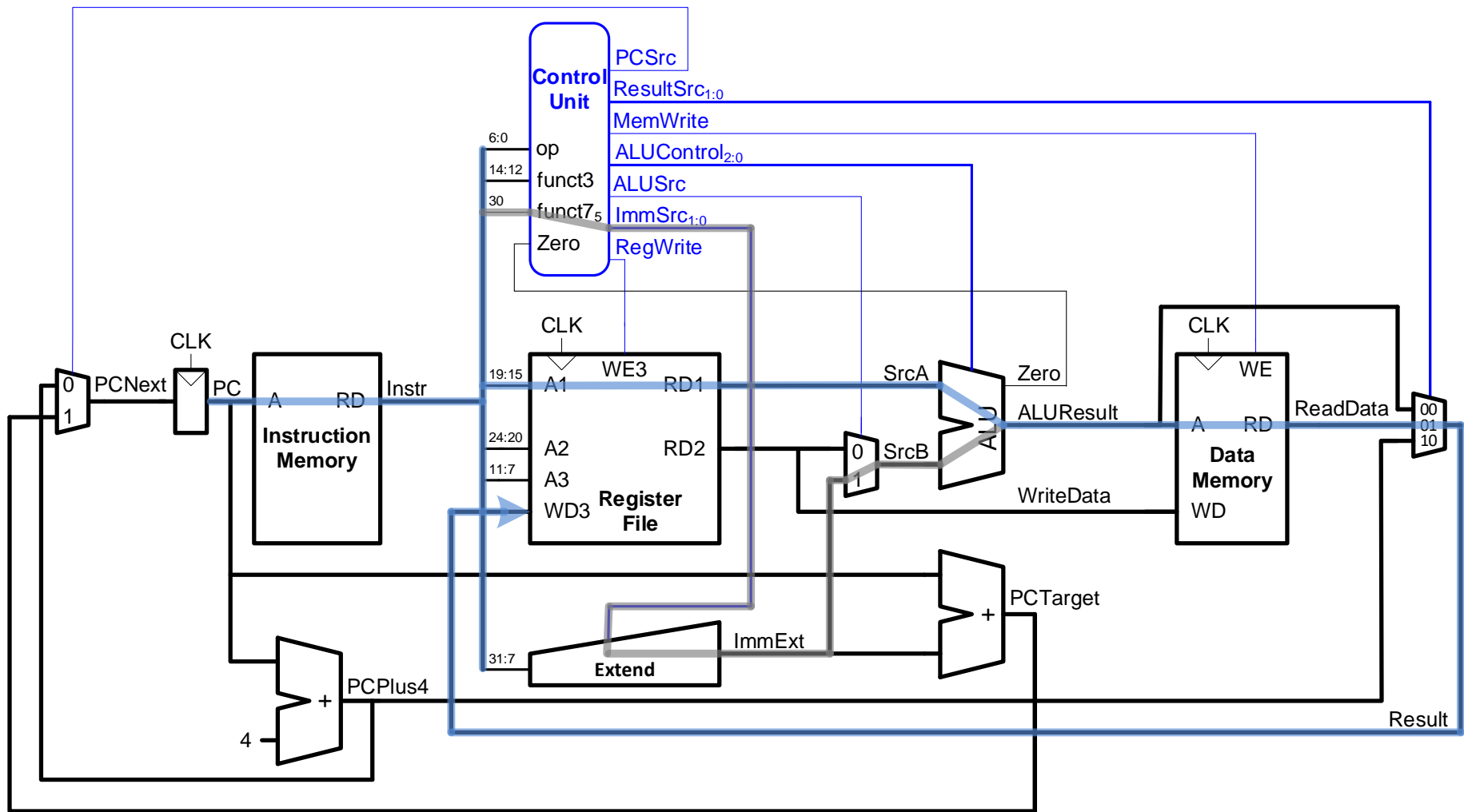
Processor Performance

Program Execution Time

= (#instructions)(cycles/instruction)(seconds/cycle)

= # instructions x CPI x T_c

Single-Cycle Processor Performance



T_c limited by critical path (1w)

Single-Cycle Processor Performance

- **Single-cycle critical path:**

$$T_{c_single} = t_{pcq_PC} + t_{mem} + \max[t_{RFread}, t_{dec} + t_{ext} + t_{mux}] + t_{ALU} + t_{mem} + t_{mux} + t_{RFsetup}$$

- **Typically, limiting paths are:**

- memory, ALU, register file

- So, $T_{c_single} = t_{pcq_PC} + t_{mem} + t_{RFread} + t_{ALU} + t_{mem} + t_{mux} + t_{RFsetup}$
 $= t_{pcq_PC} + 2t_{mem} + t_{RFread} + t_{ALU} + t_{mux} + t_{RFsetup}$

Single-Cycle Performance Example

Element	Parameter	Delay (ps)
Register clock-to-Q	t_{pcq_PC}	40
Register setup	t_{setup}	50
Multiplexer	t_{mux}	30
AND-OR gate	t_{AND-OR}	20
ALU	t_{ALU}	120
Decoder (Control Unit)	t_{dec}	25
Extend unit	t_{ext}	35
Memory read	t_{mem}	200
Register file read	t_{RFread}	100
Register file setup	$t_{RFsetup}$	60

$$T_{c_single} = t_{pcq_PC} + 2t_{mem} + t_{RFread} + t_{ALU} + t_{mux} + t_{RFsetup}$$
$$=$$

Single-Cycle Performance Example

Program with 100 billion instructions:

$$\text{Execution Time} = \# \text{ instructions} \times \text{CPI} \times T_c$$

Chapter 7: Microarchitecture

Multicycle RISC-V Processor

Single- vs. Multicycle Processor

- **Single-cycle:**
 - + simple
 - cycle time limited by longest instruction (1_w)
 - separate memories for instruction and data
 - 3 adders/ALUs
- **Multicycle processor** addresses these issues by breaking instruction into **shorter steps**
 - shorter instructions take fewer steps
 - can re-use hardware
 - cycle time is faster

Single- vs. Multicycle Processor

- **Single-cycle:**

- + simple
- cycle time limited by longest instruction (τ_w)
- separate memories for instruction and data
- 3 adders/ALUs

- **Multicycle:**

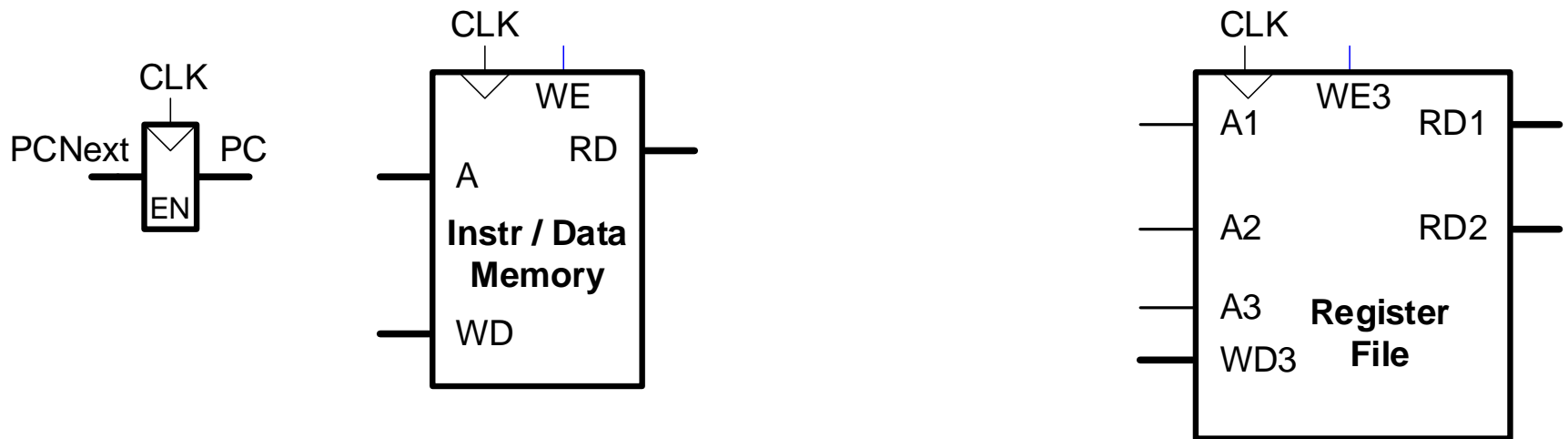
- + higher clock speed
- + simpler instructions run faster
- + reuse expensive hardware on multiple cycles
- sequencing overhead paid many times

**Same design steps
as single-cycle:**

- **first datapath**
- **then control**

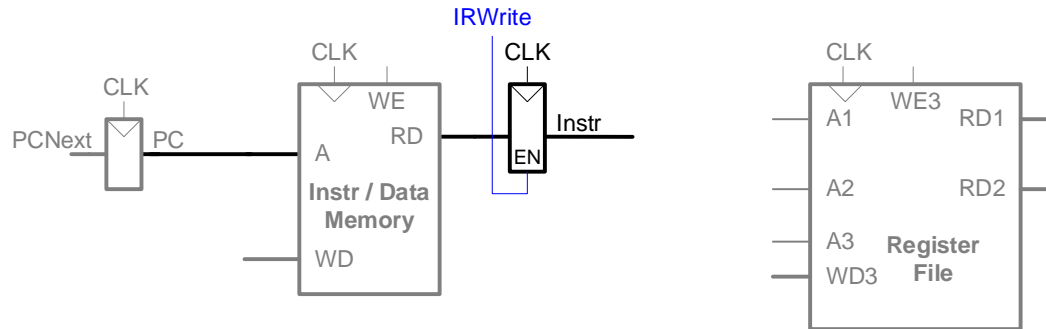
Multicycle State Elements

Replace separate Instruction and Data memories with a **single unified memory** – more realistic



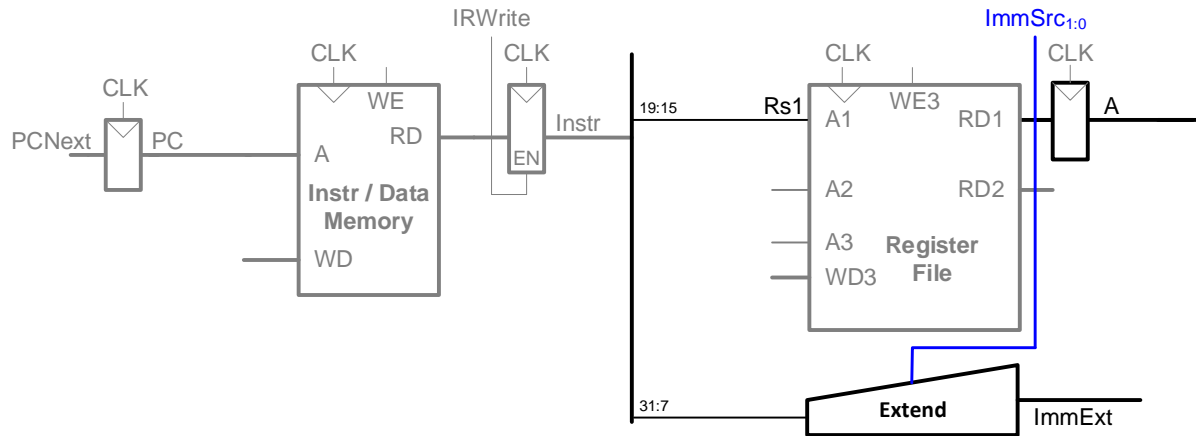
Multicycle Datapath: Instruction Fetch

STEP 1: Fetch instruction



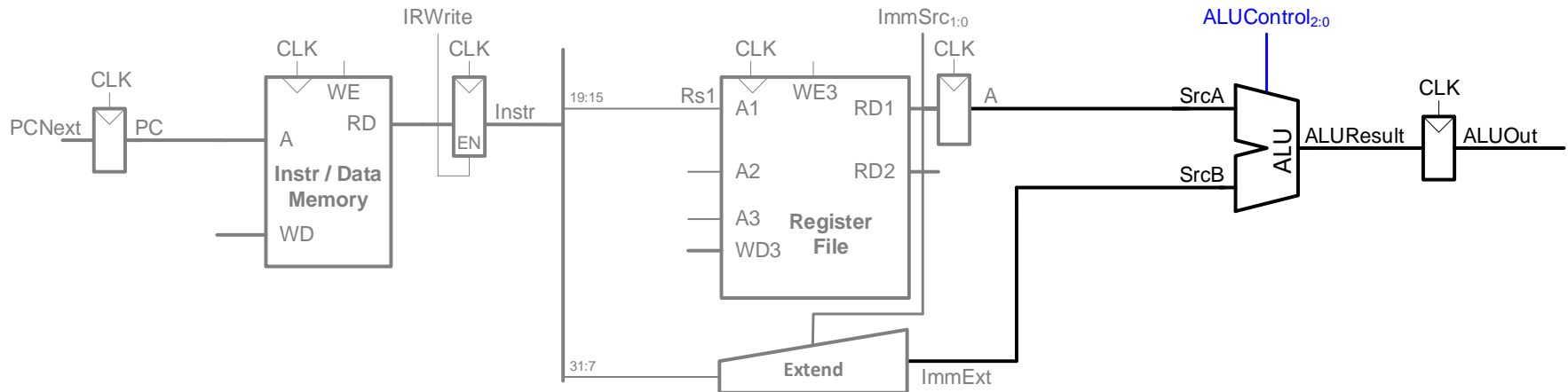
Multicycle Datapath: 1_w Get Sources

STEP 2: Read source operand from RF and extend immediate



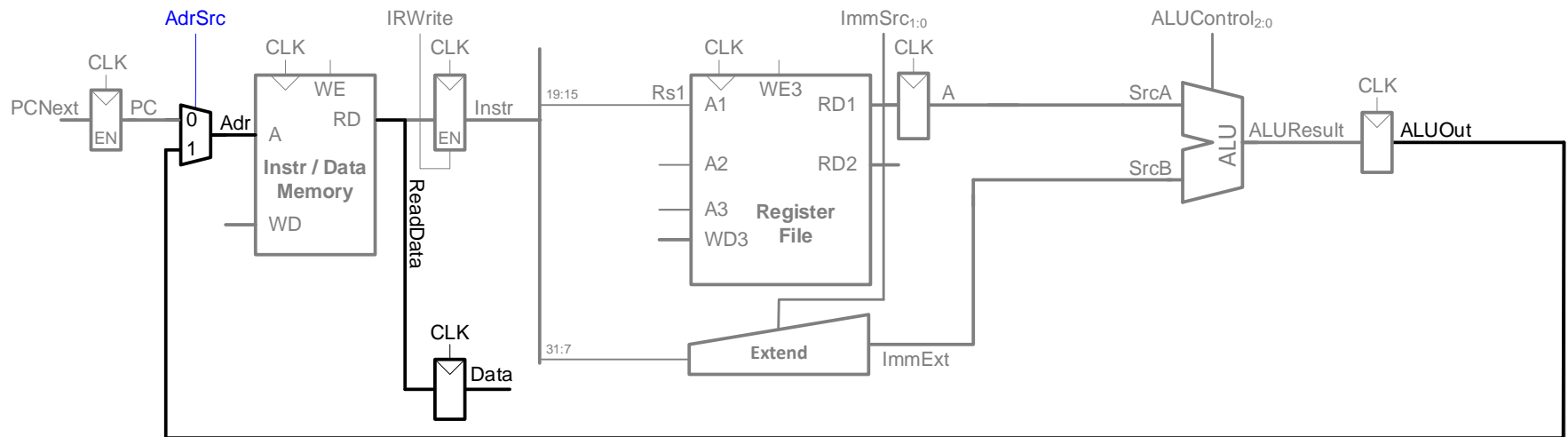
Multicycle Datapath: 1_w Address

STEP 3: Compute the memory address



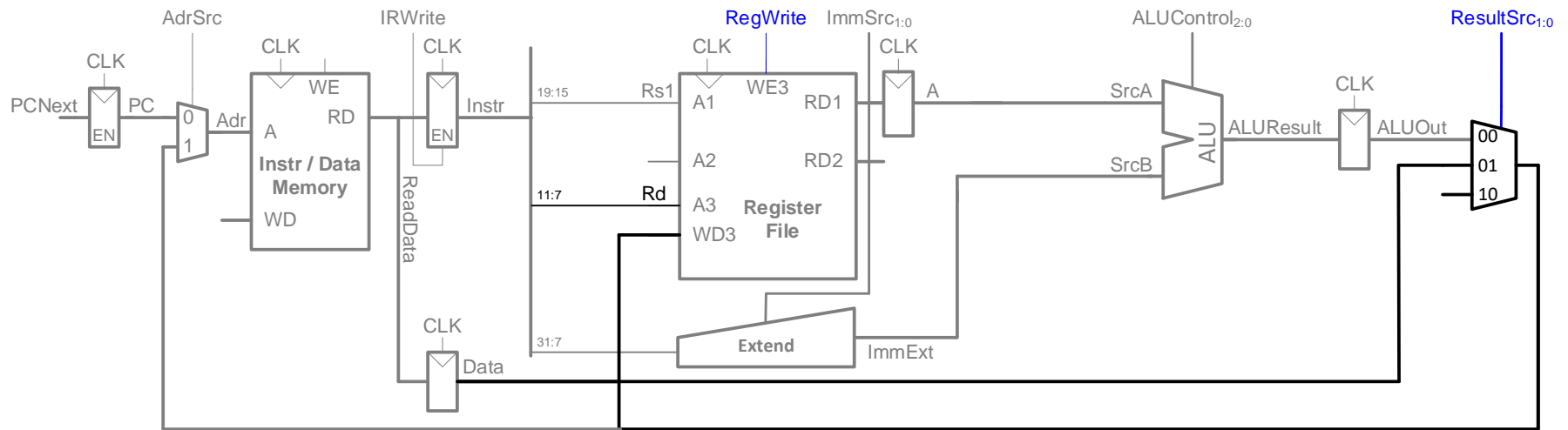
Multicycle Datapath: 1_W Memory Read

STEP 4: Read data from memory



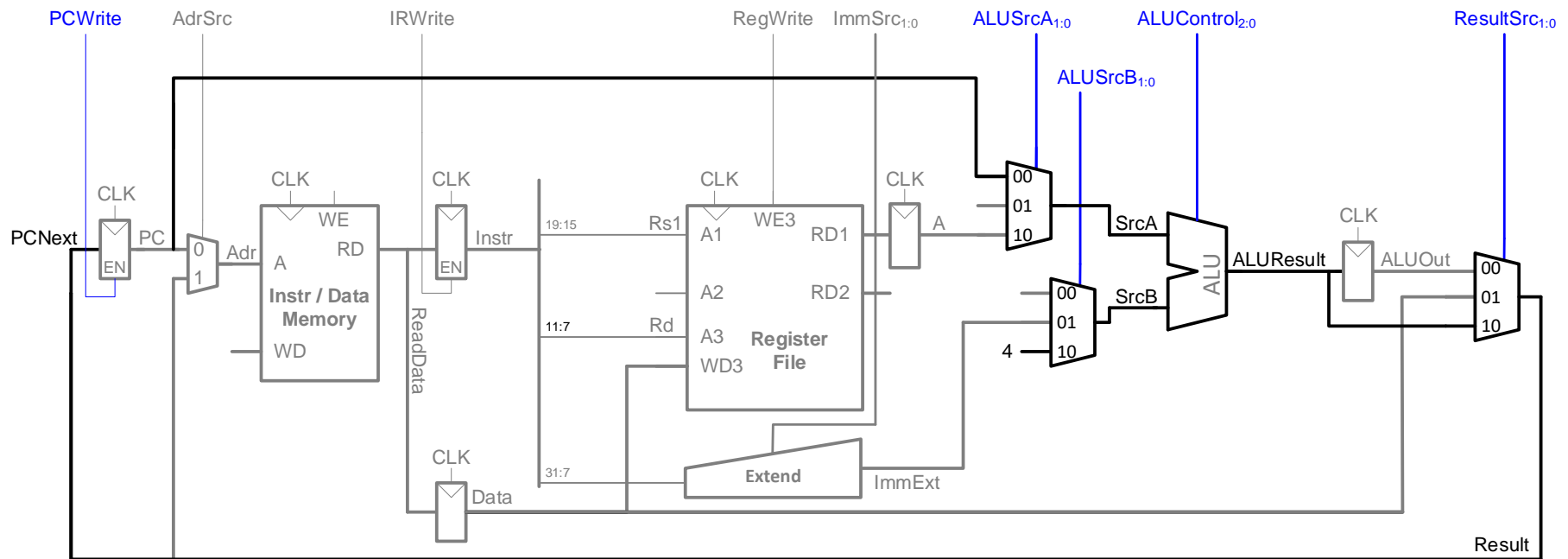
Multicycle Datapath: 1_W Write Register

STEP 5: Write data back to register file



Multicycle Datapath: Increment PC

STEP 6: Increment PC: $PC = PC + 4$

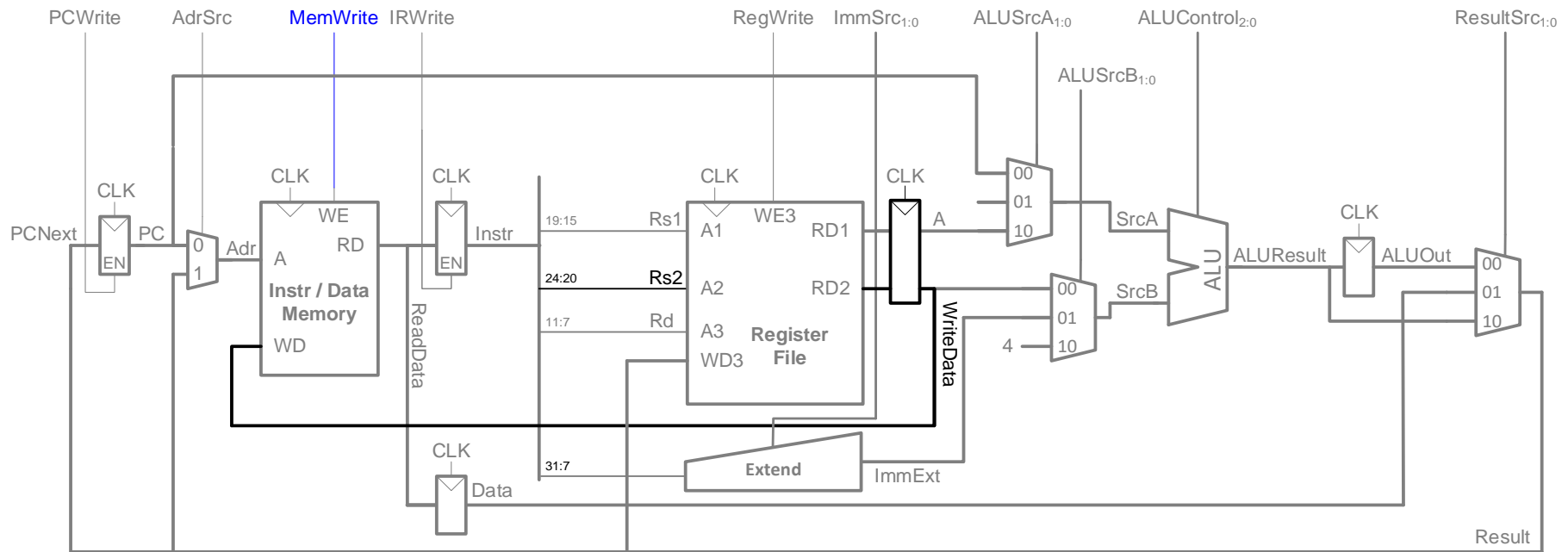


Chapter 7: Microarchitecture

Multicycle Datapath: Other Instructions

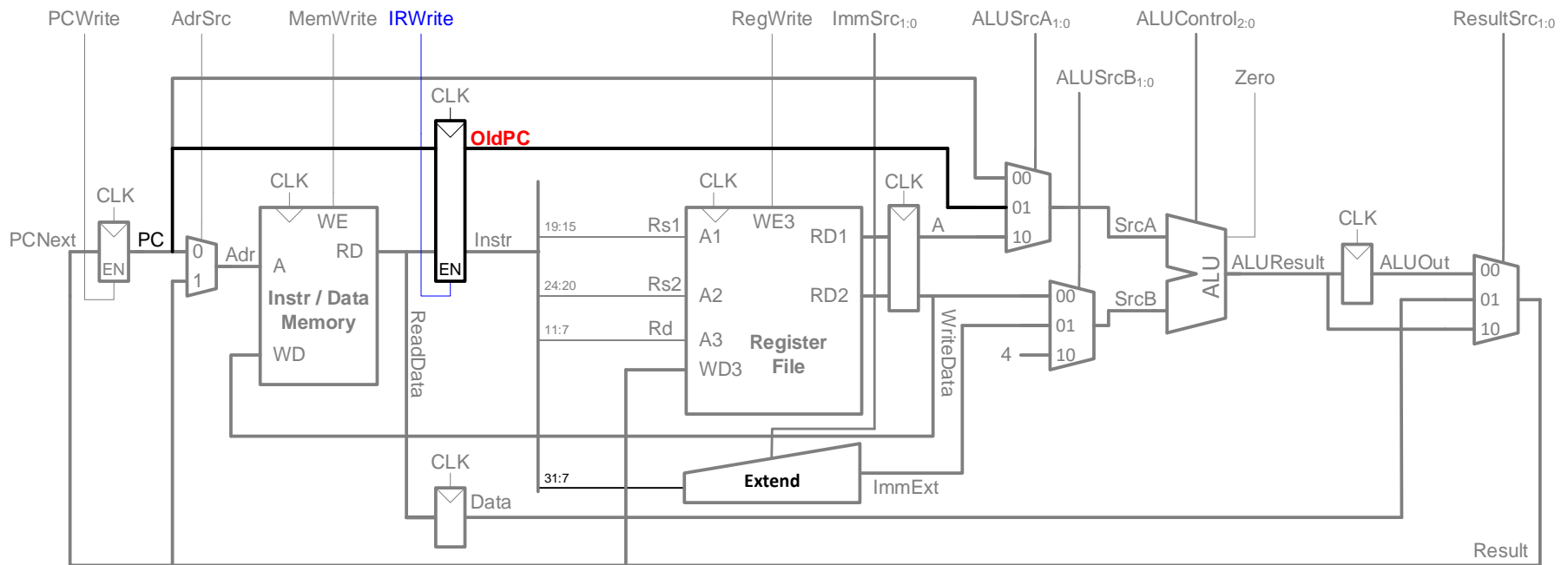
Multicycle Datapath: sw

Write data in **rs2** to memory



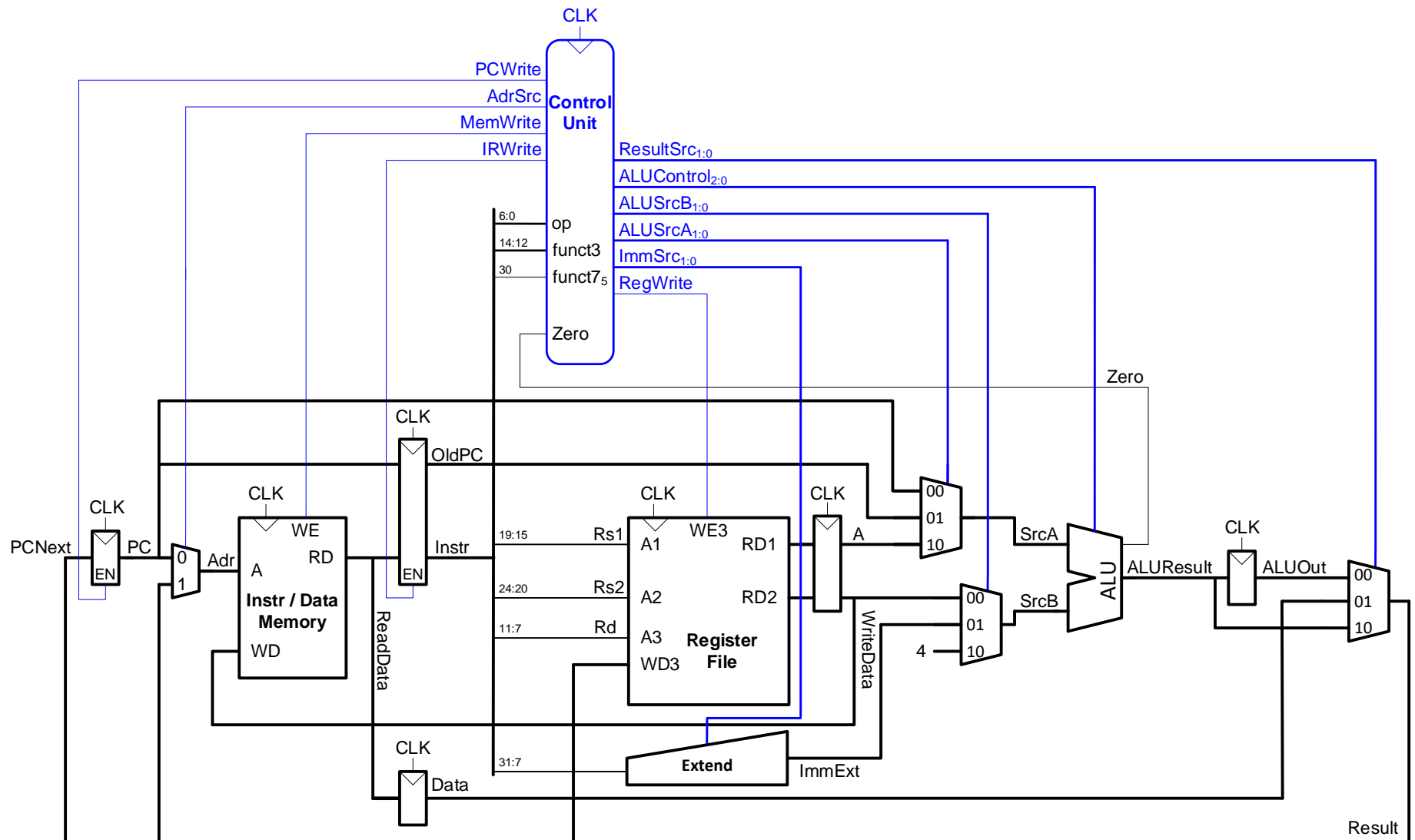
Multicycle Datapath: beq

Calculate branch target address:
 $BTA = PC + imm$



PC is updated in Fetch stage, so need to save **old (current) PC**

Multicycle RISC-V Processor

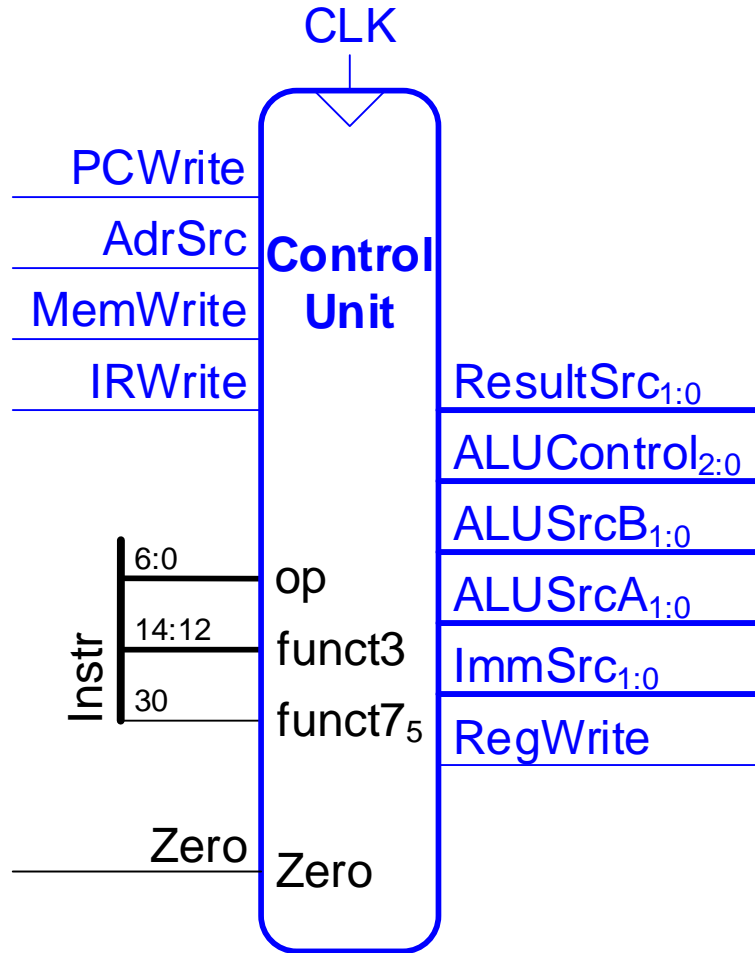


Chapter 7: Microarchitecture

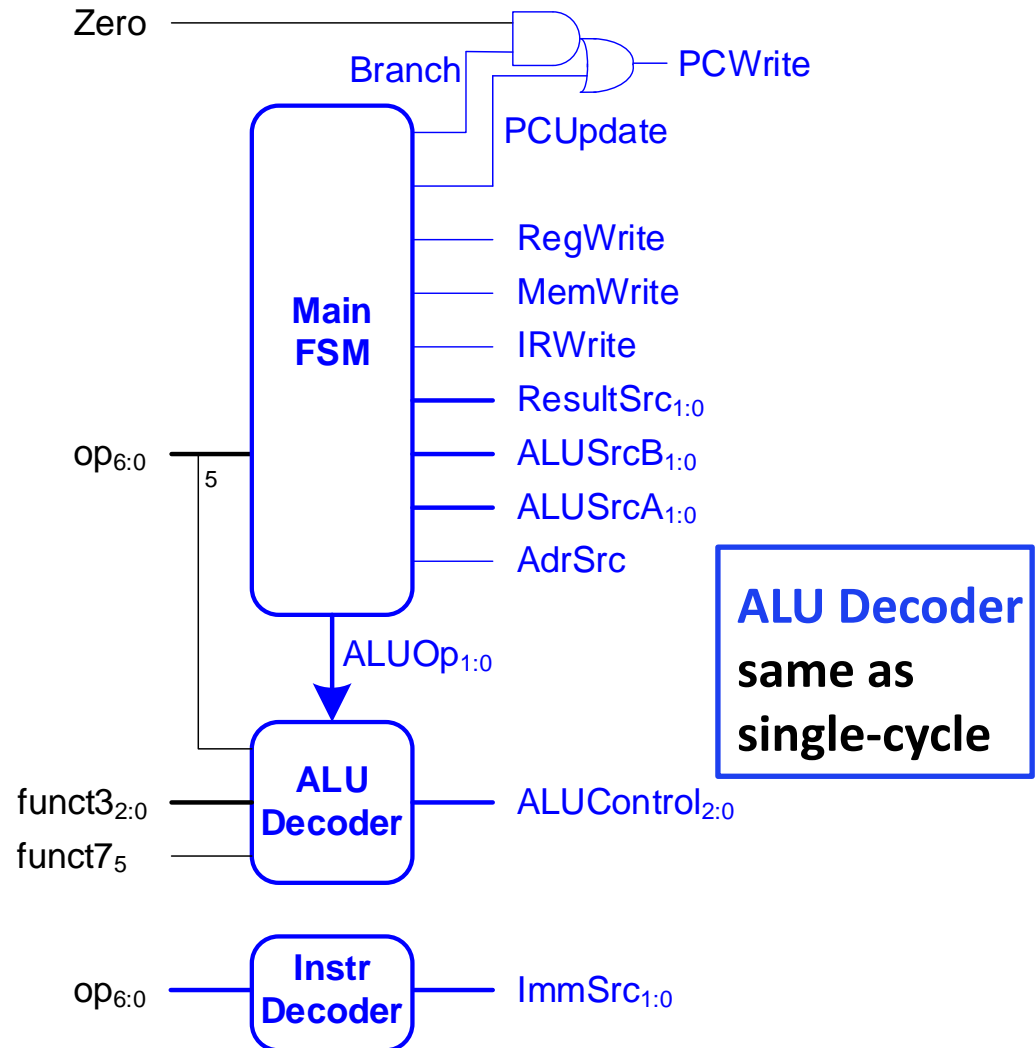
Multicycle Control

Multicycle Control

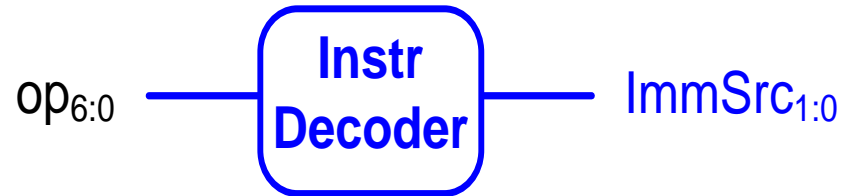
High-Level View



Low-Level View

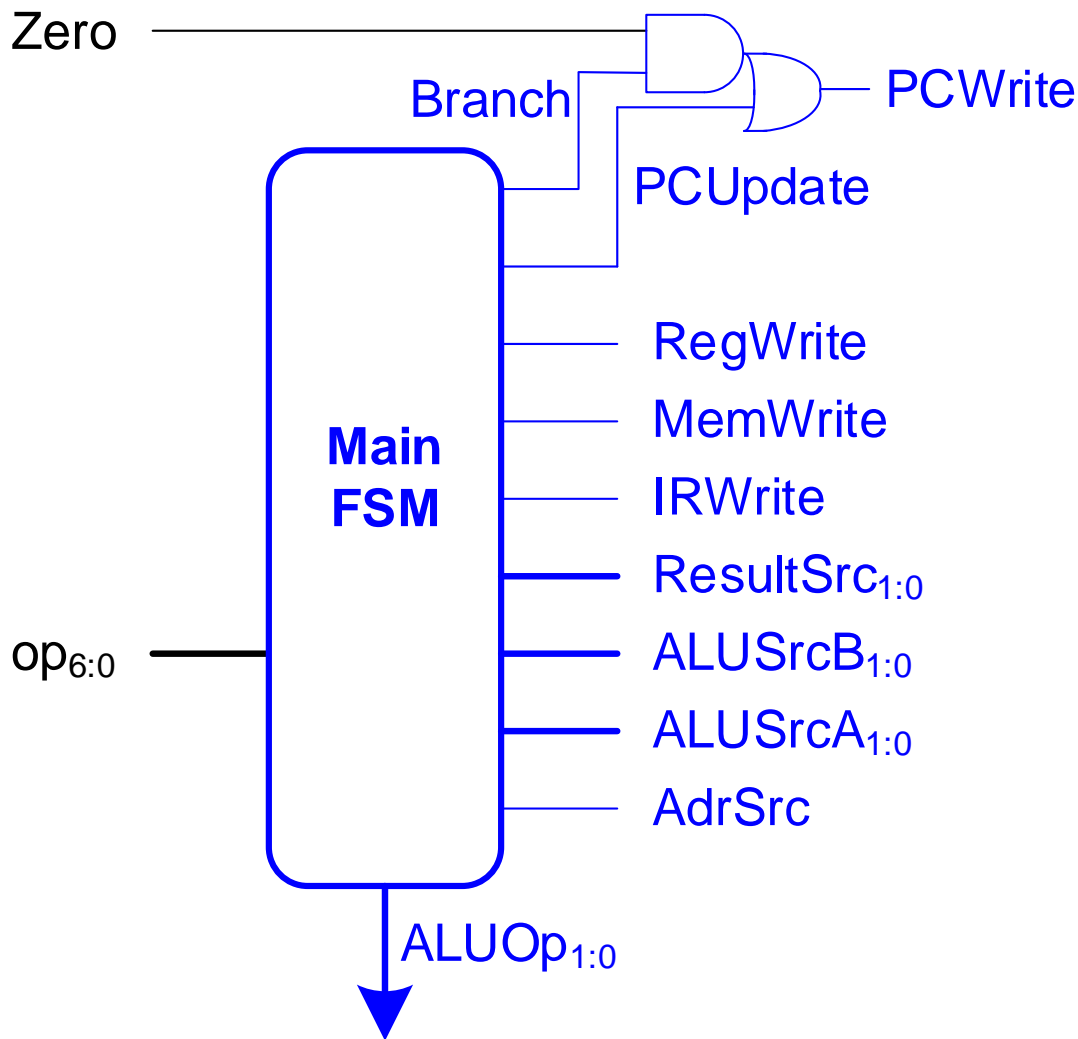


Multicycle Control: Instruction Decoder



op	Instruction	ImmSrc
3	lw	00
35	sw	01
51	R-type	XX
99	beq	10

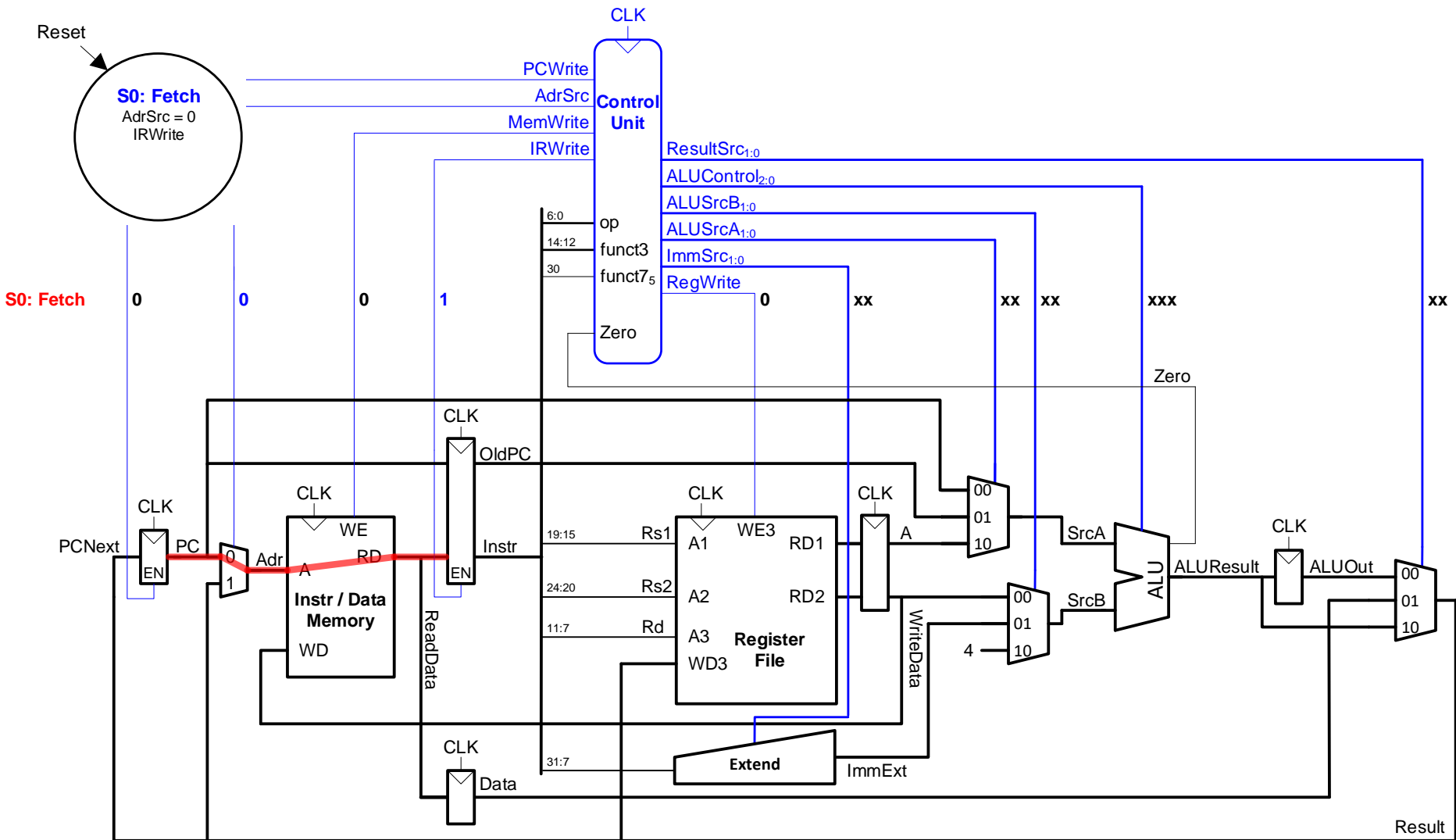
Multicycle Control: Main FSM



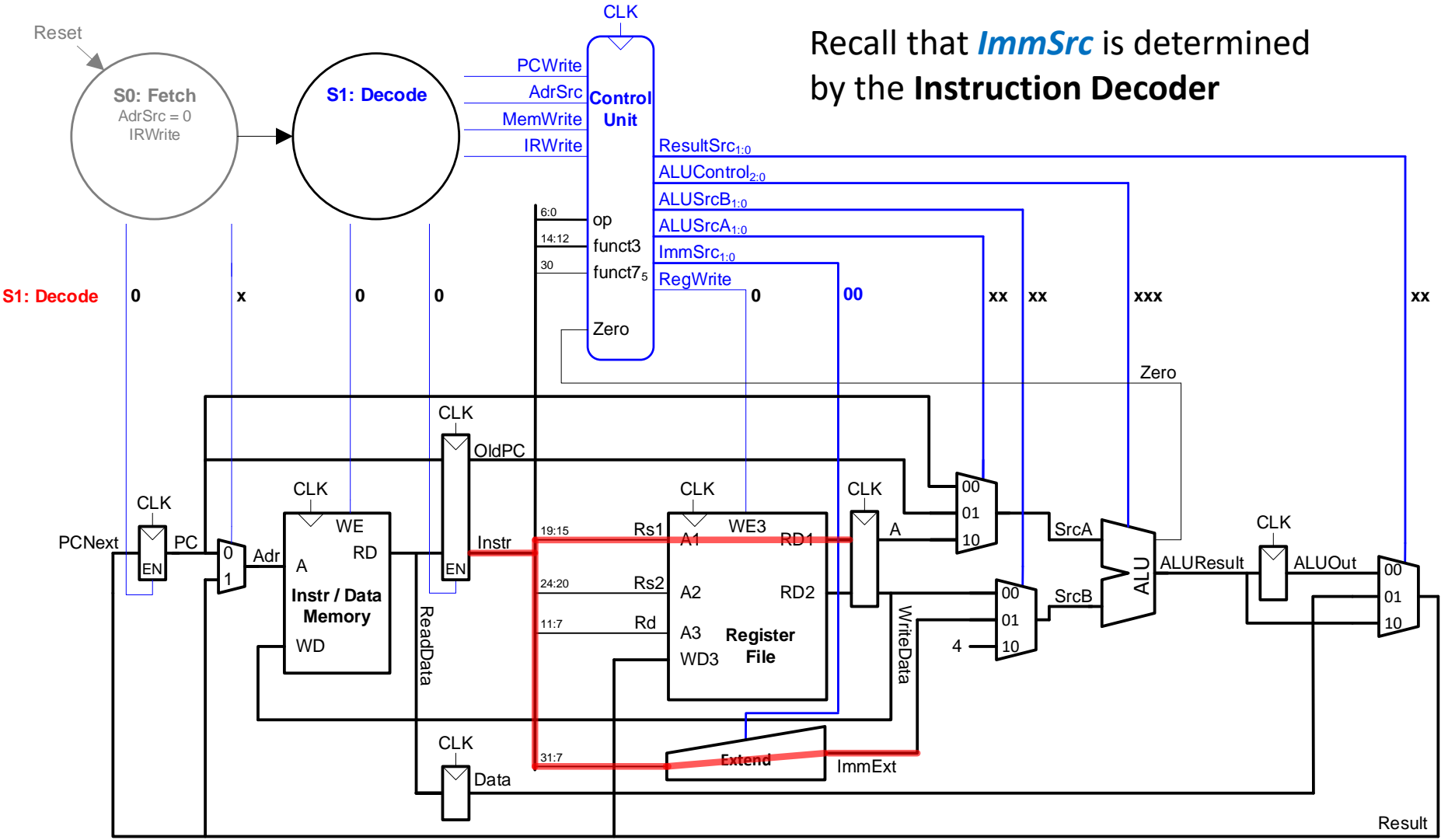
To declutter FSM:

- **Write enable signals** (RegWrite, MemWrite, IRWrite, PCUpdate, and Branch) are **0** if not listed in a state.
- **Other signals are don't care** if not listed in a state

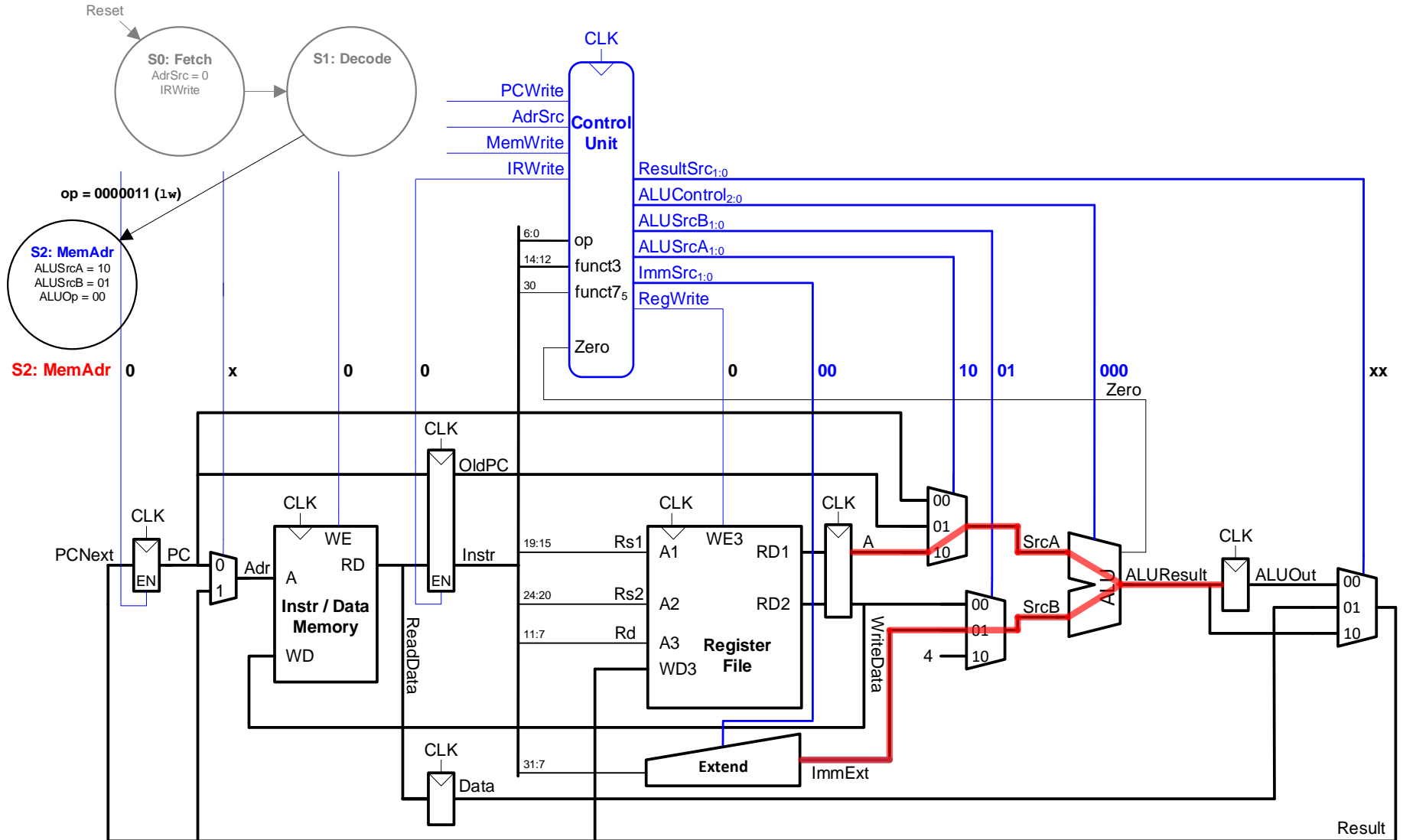
Main FSM: Fetch



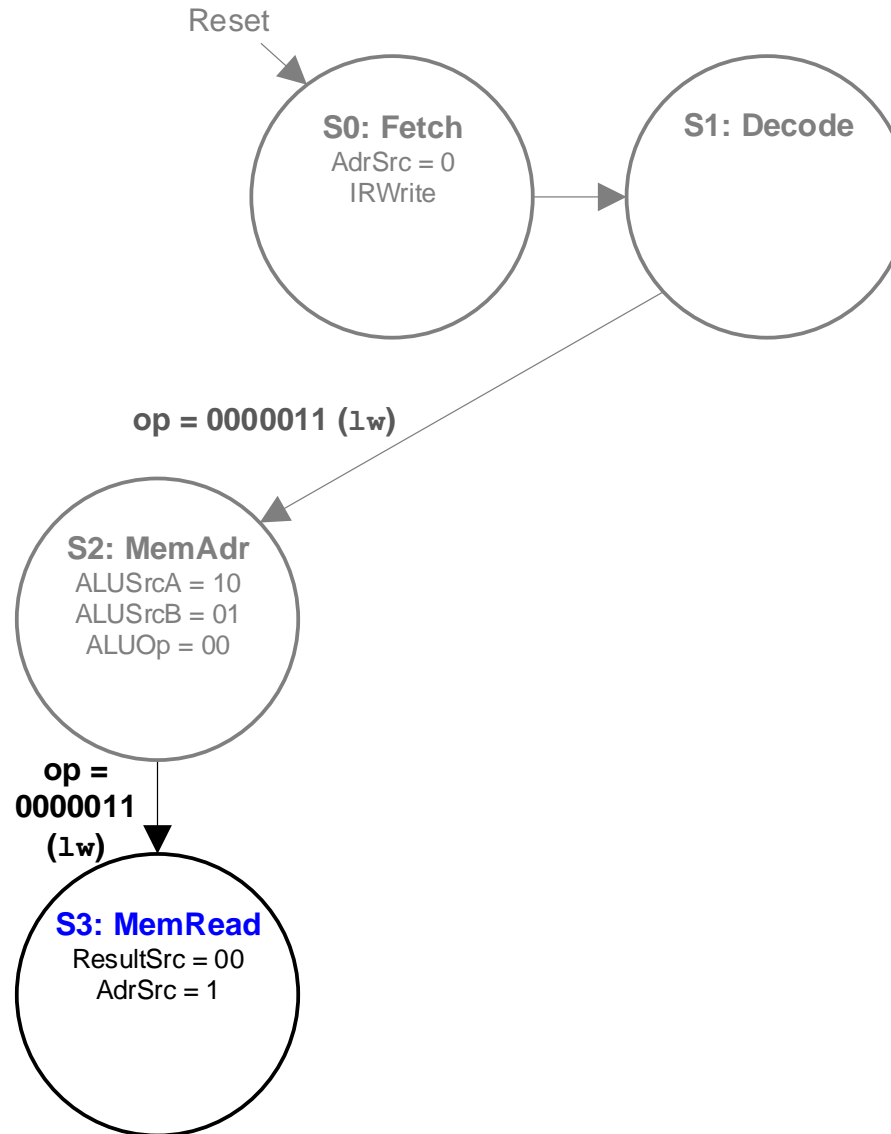
Main FSM: Decode



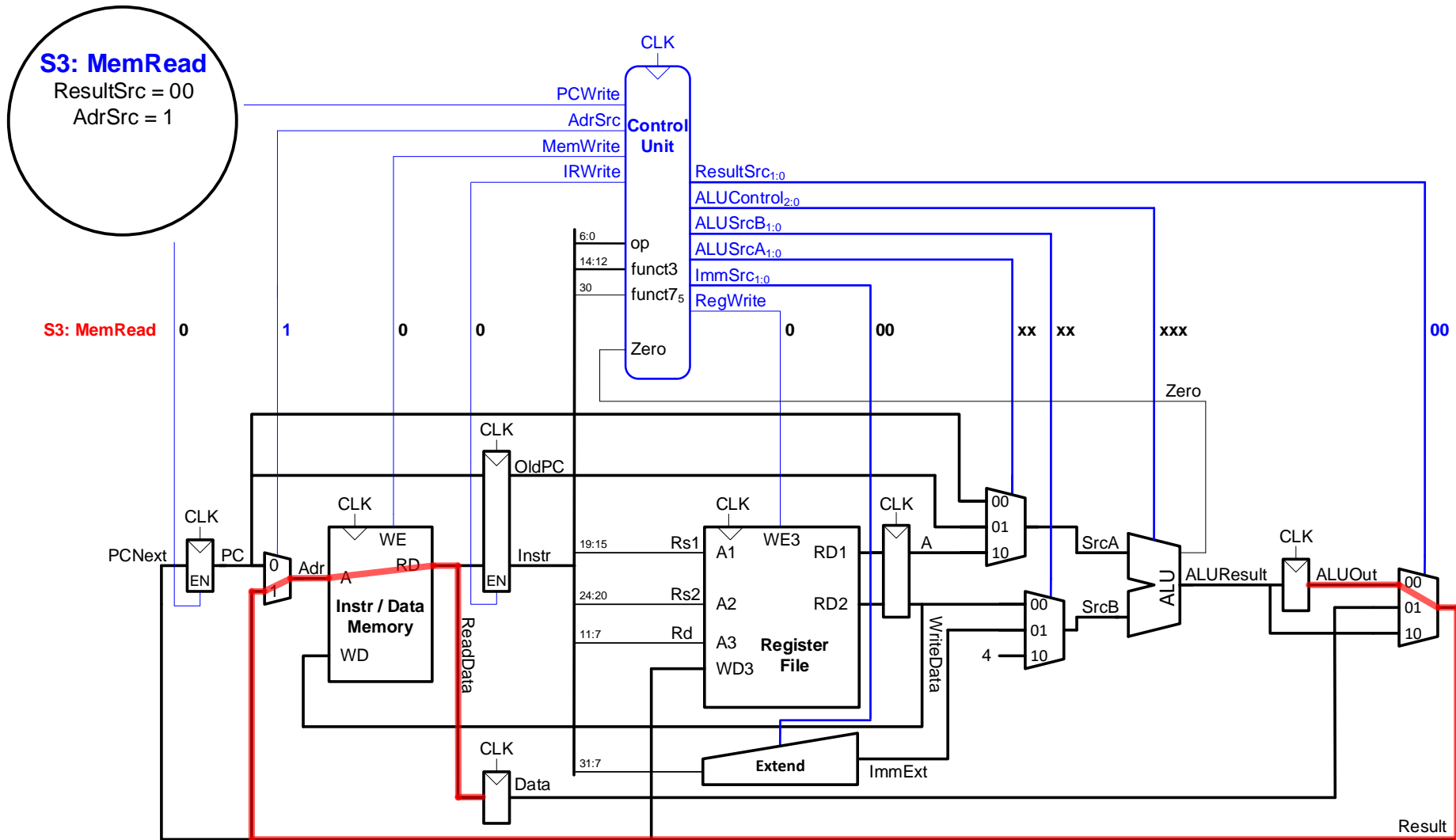
Main FSM: Address



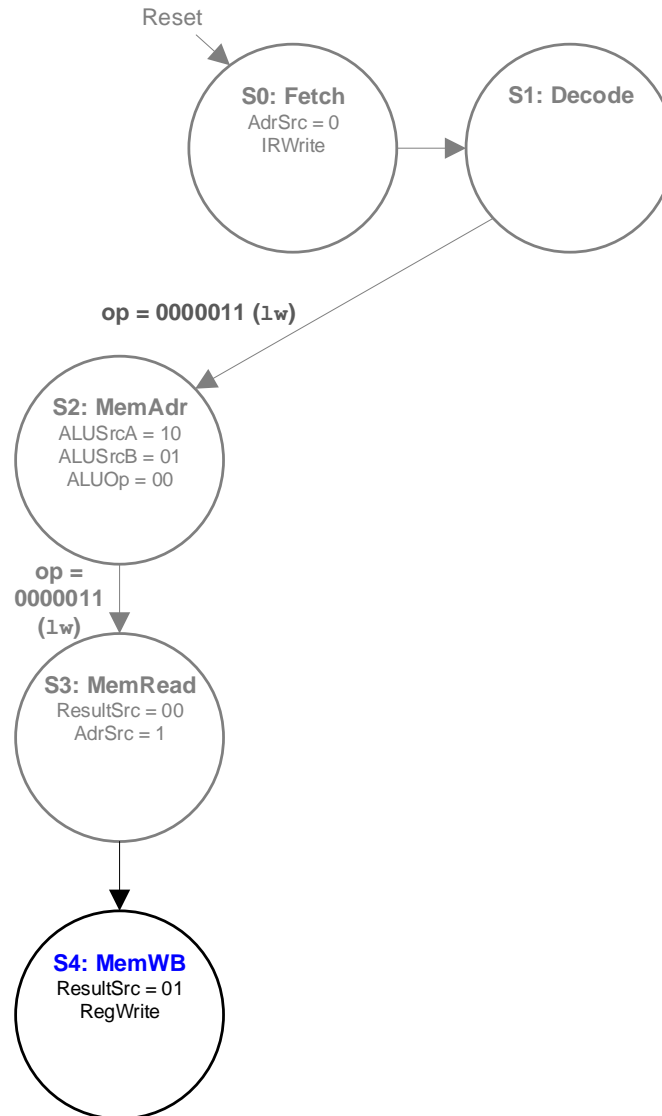
Main FSM: Read Memory



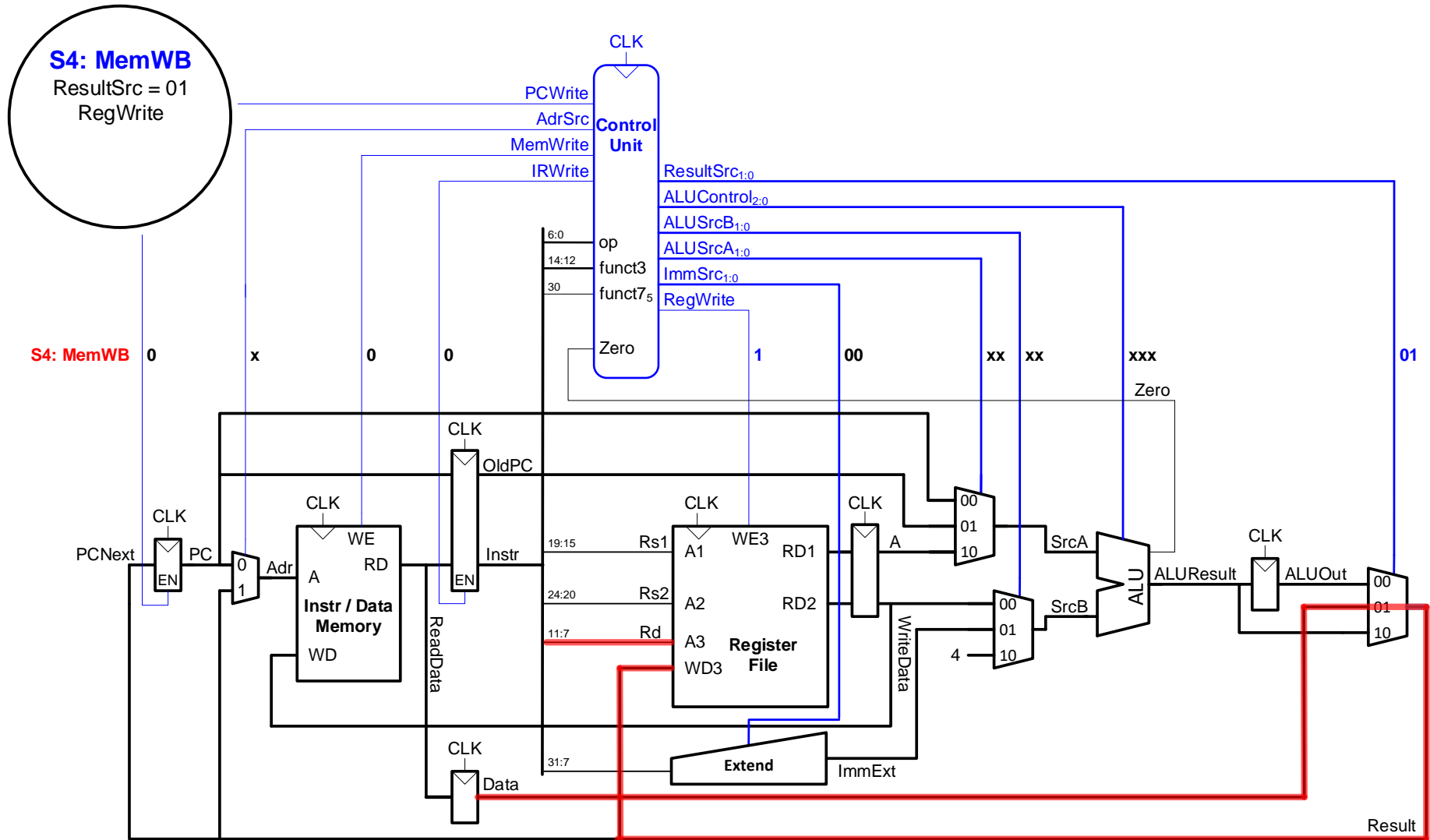
Main FSM: Read Memory Datapath



Main FSM: Write RF

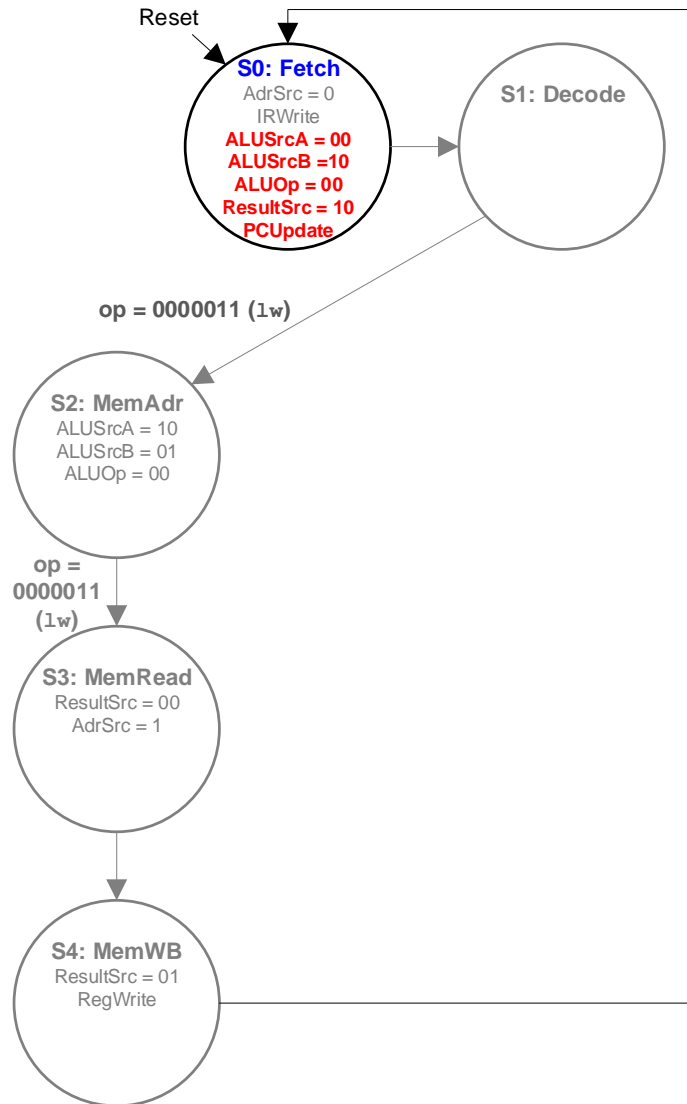


Main FSM: Write RF Datapath

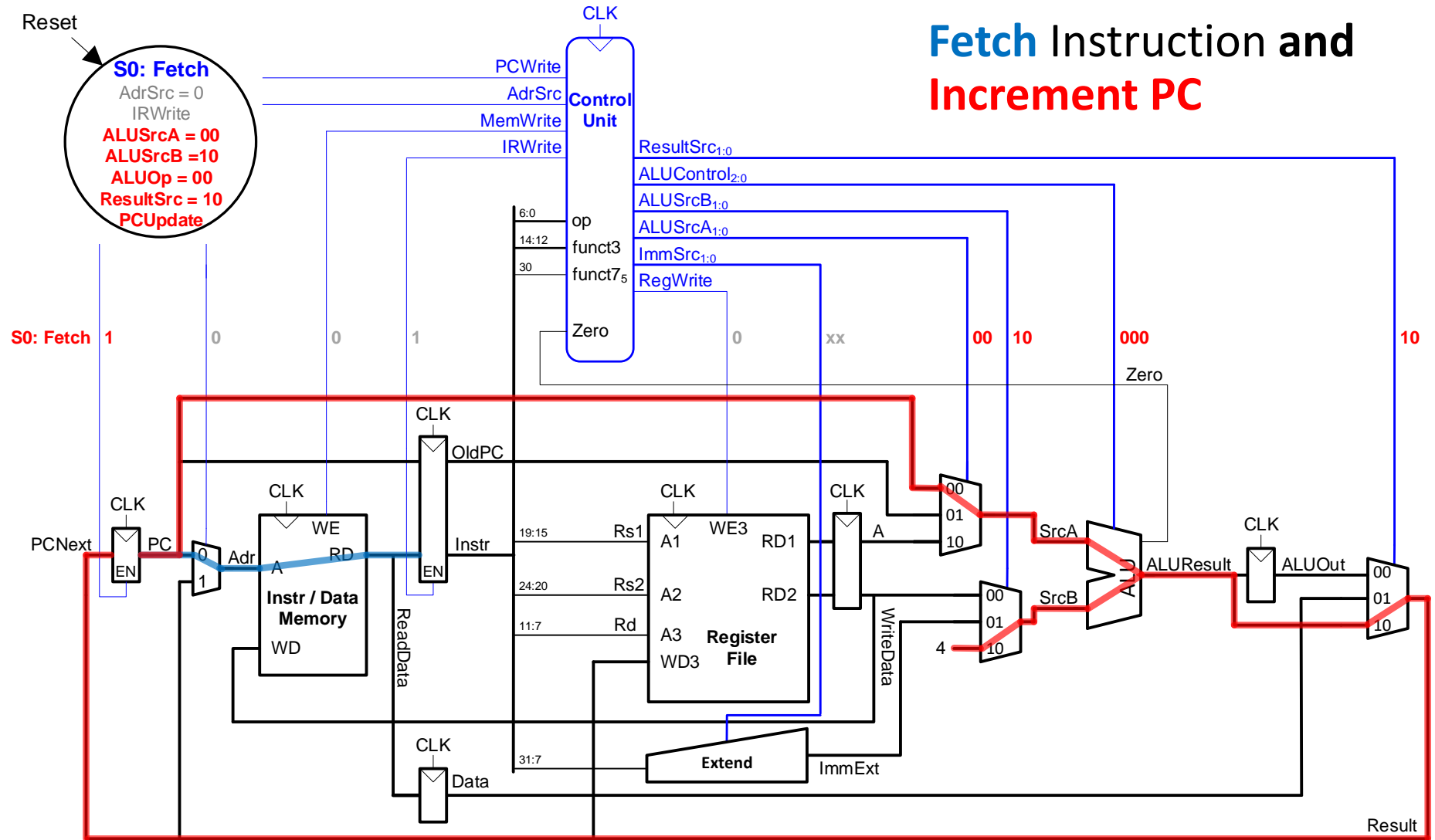


Main FSM: Fetch Revisited

Calculate **PC+4**
during Fetch stage
(ALU isn't being
used)



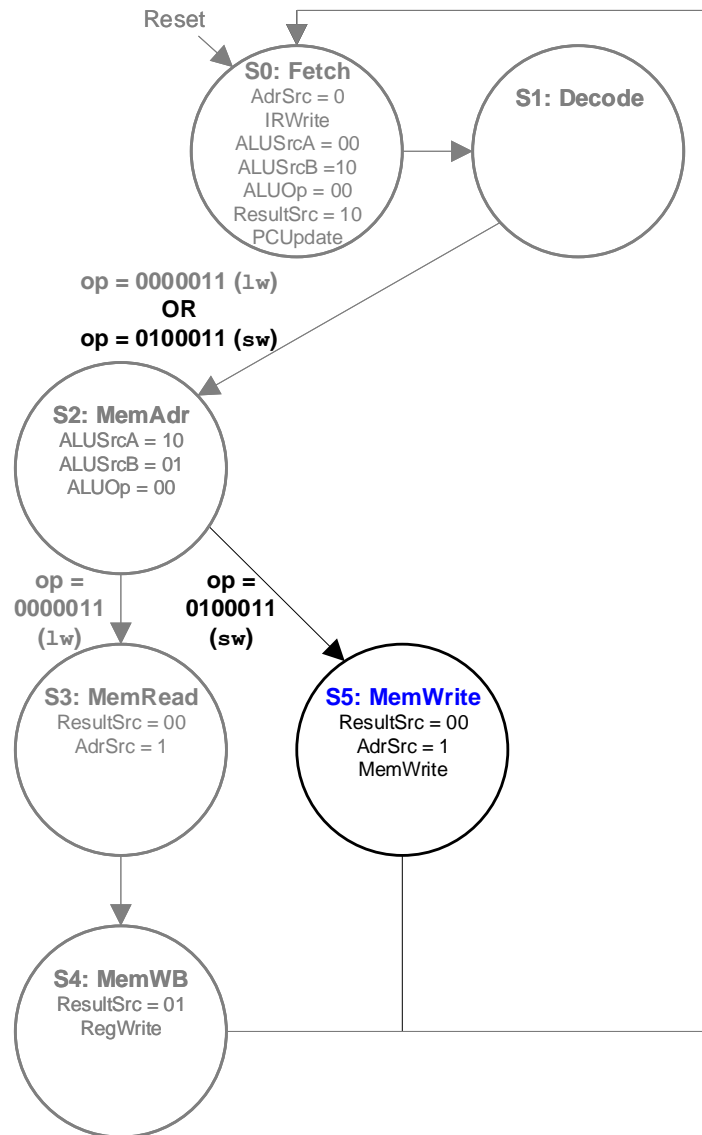
Main FSM: Fetch (PC+4) Datapath



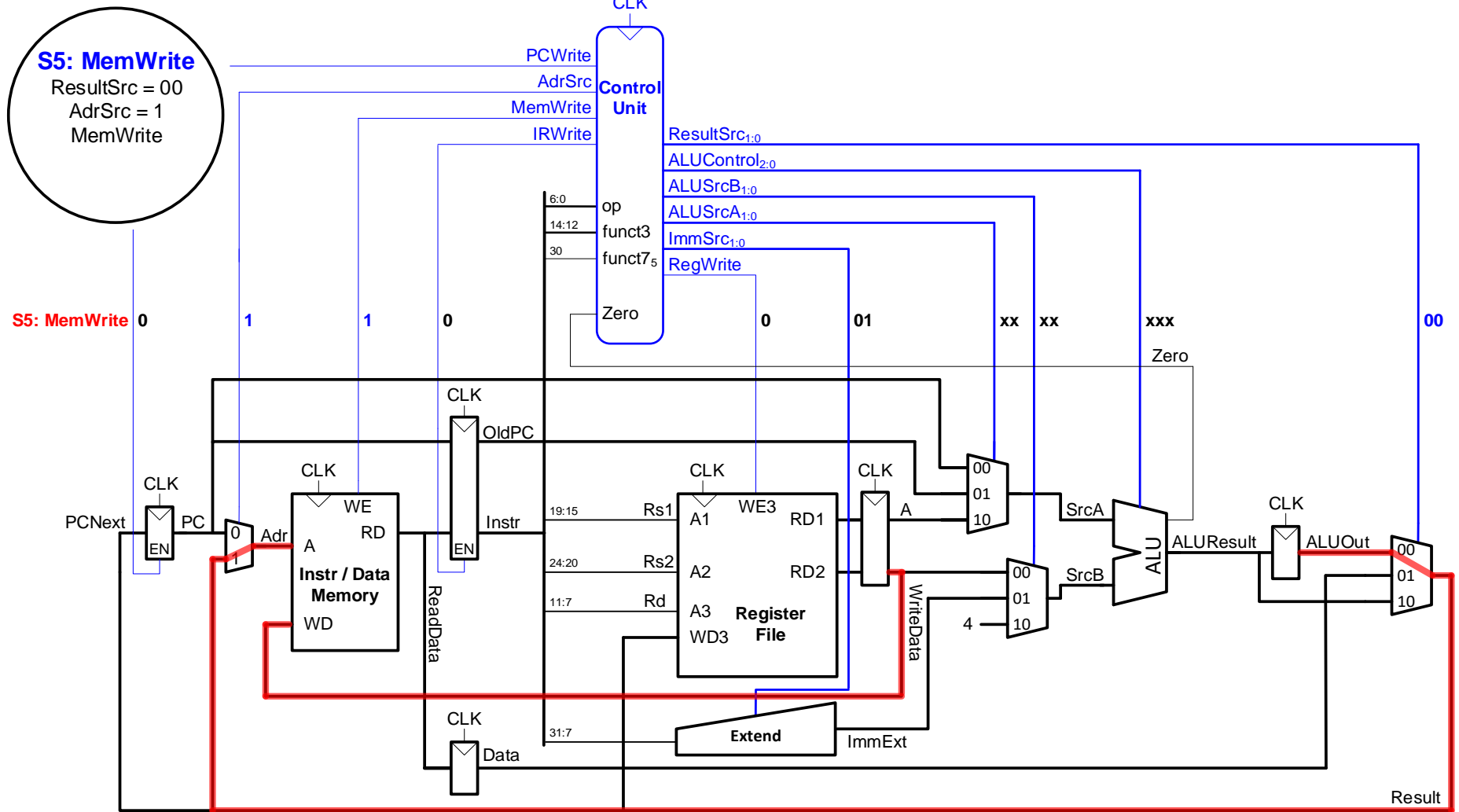
Chapter 7: Microarchitecture

Multicycle Control: Other Instructions

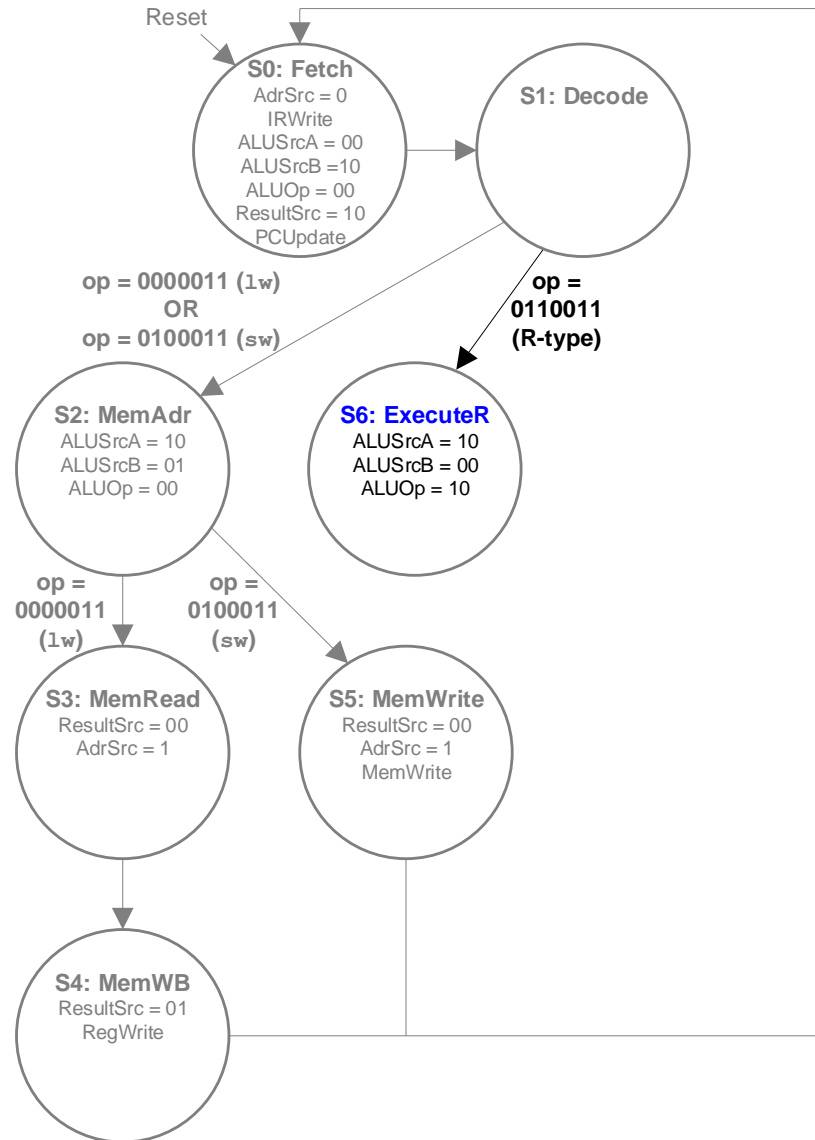
Main FSM: SW



Main FSM: sw Datapath



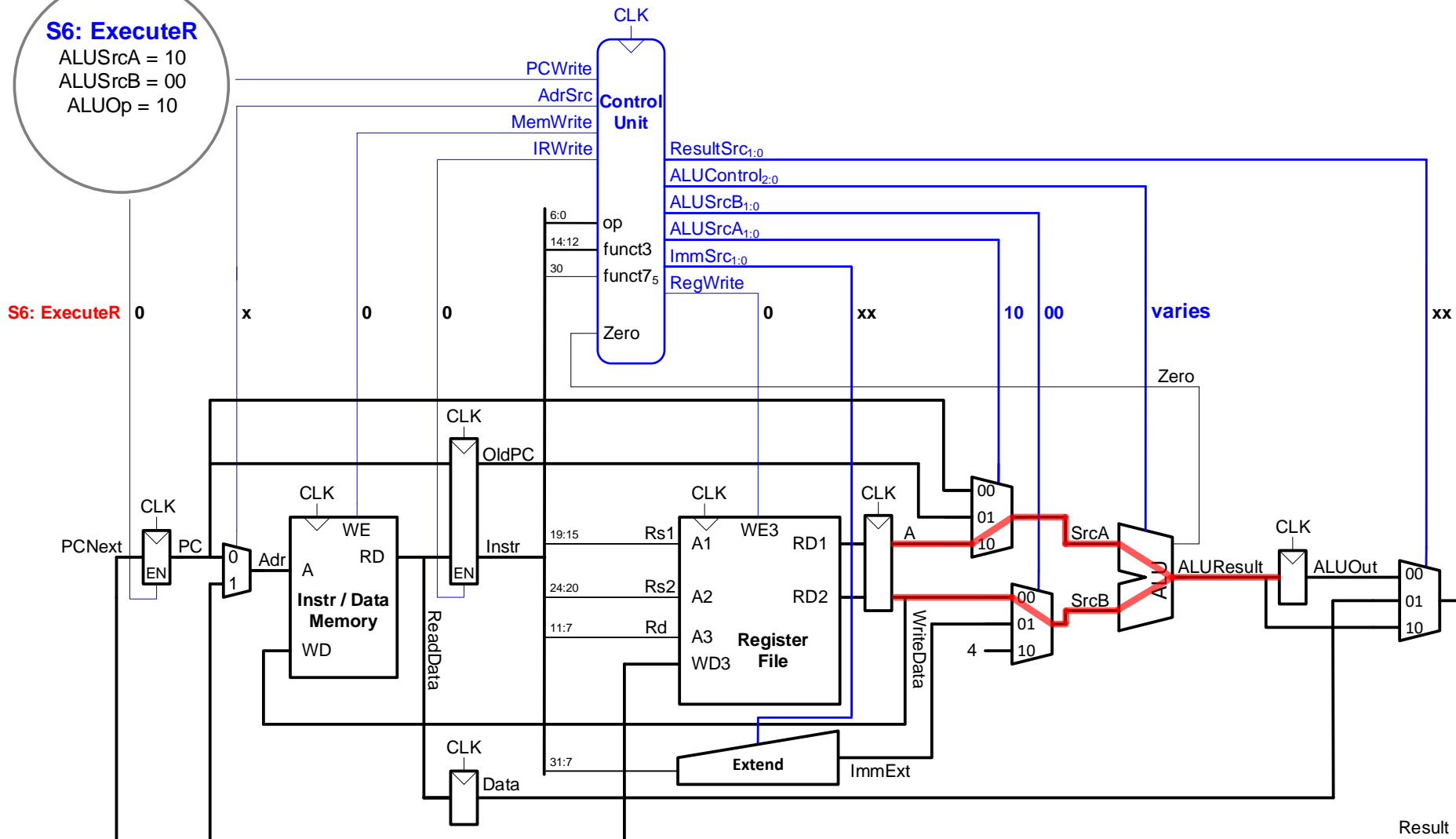
Main FSM: R-Type Execute



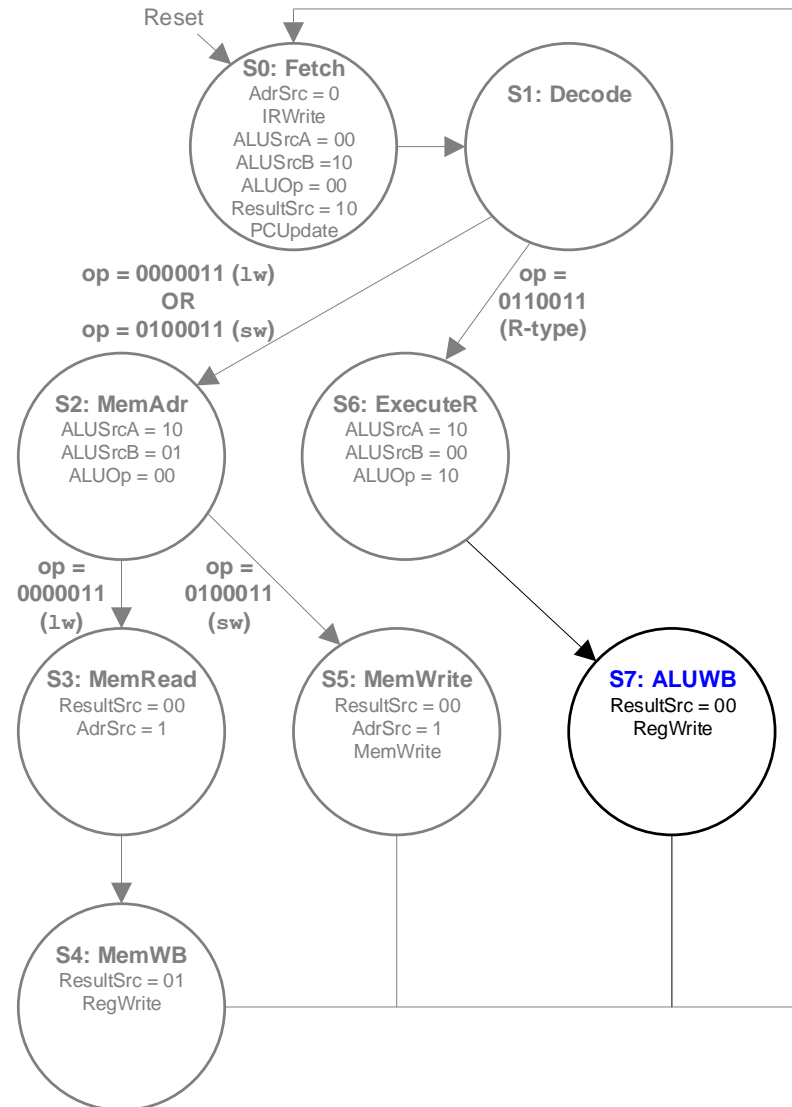
Main FSM: R-Type Execute Datapath

S6: ExecuteR

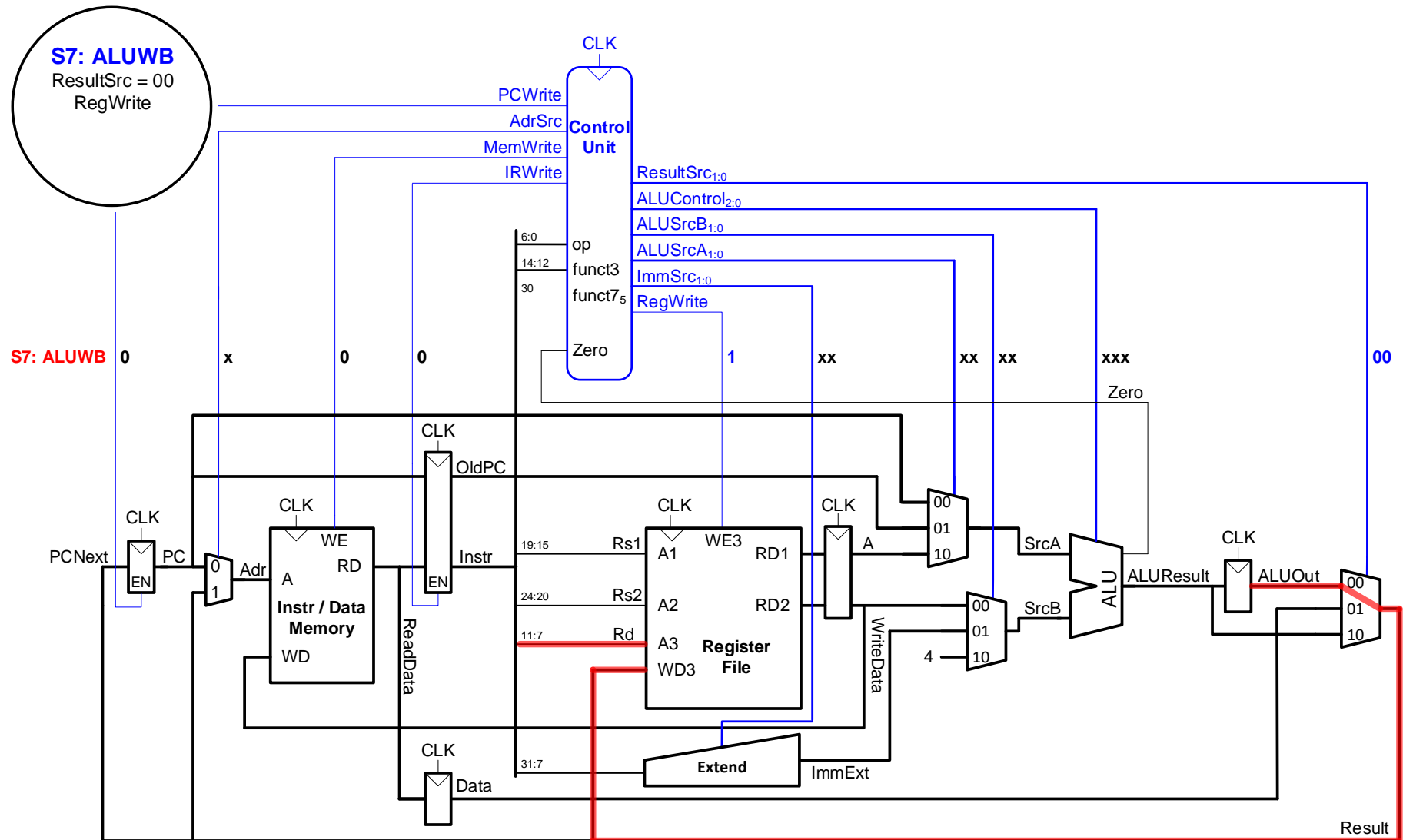
ALUSrcA = 10
ALUSrcB = 00
ALUOp = 10



Main FSM: R-Type ALU Write Back



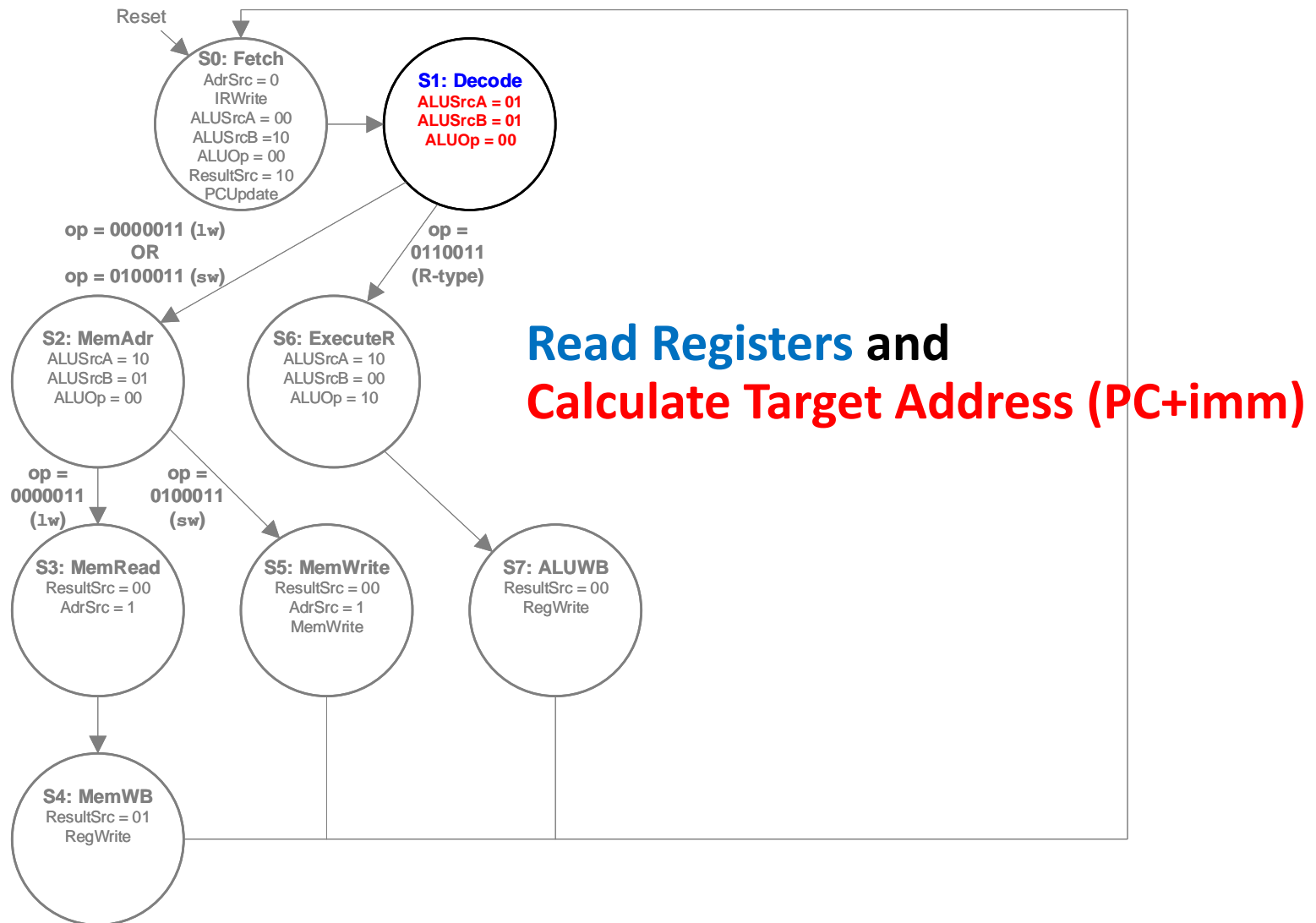
Main FSM: R-Type ALU Write Back




Main FSM: beq

- **Need to calculate:**
 - Branch Target Address
 - **rs1** - **rs2** (to see if equal)
- **ALU** isn't being used in Decode stage
 - Use it to calculate Target Address ($PC + imm$)

Main FSM: Decode Revisited

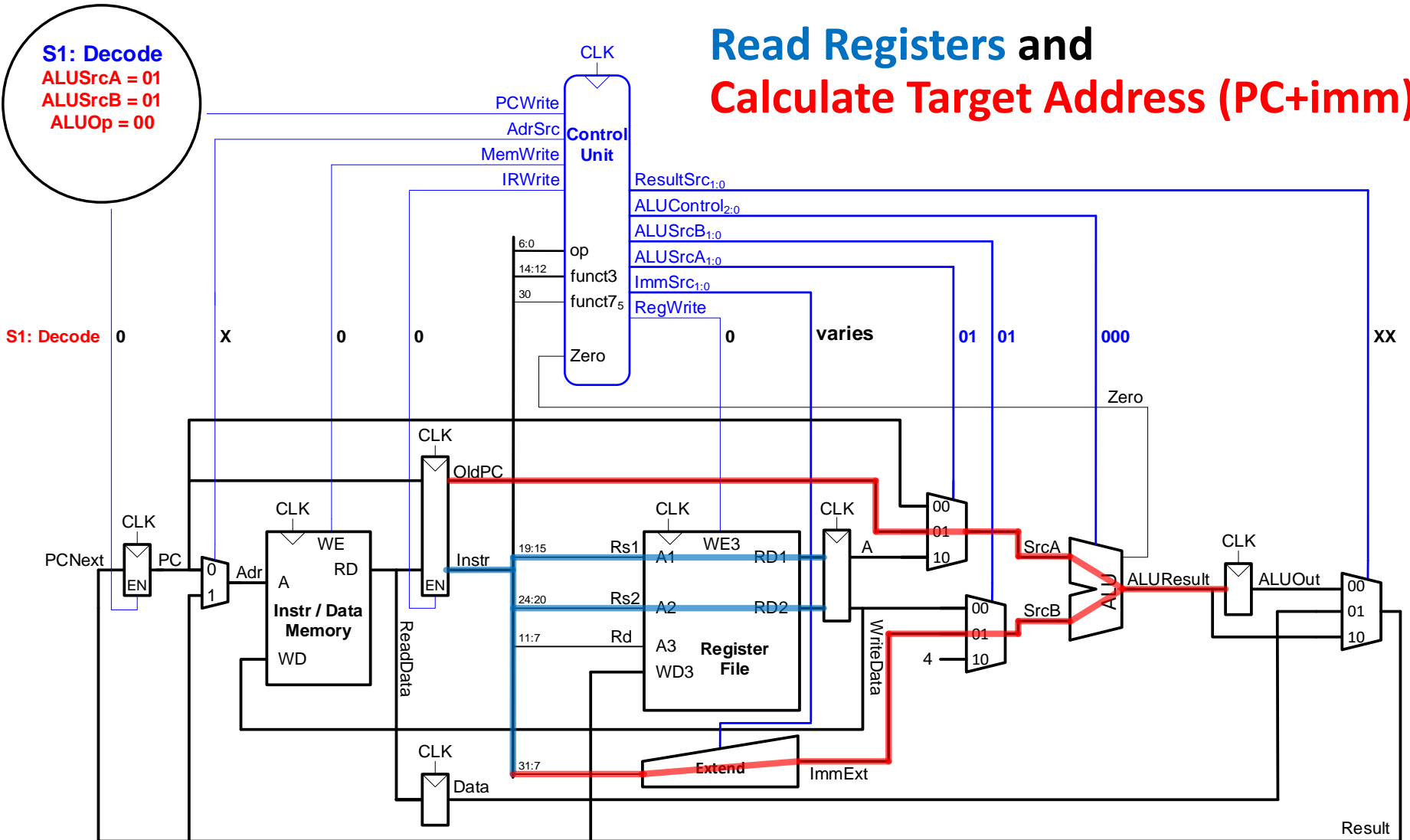


Main FSM: Decode (Target Address)

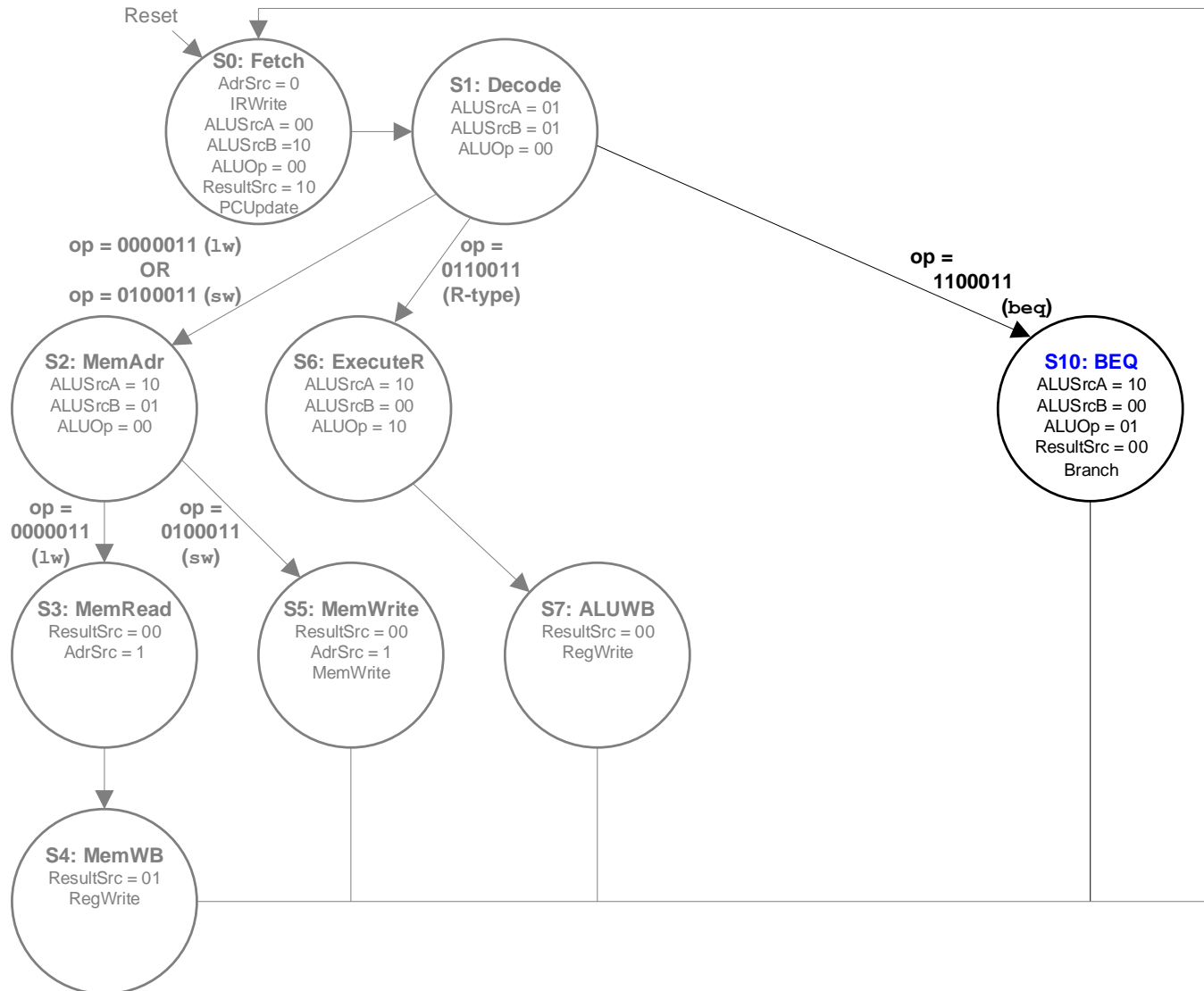


S1: Decode
 ALUSrcA = 01
 ALUSrcB = 01
 ALUOp = 00

Read Registers and Calculate Target Address (PC+imm)

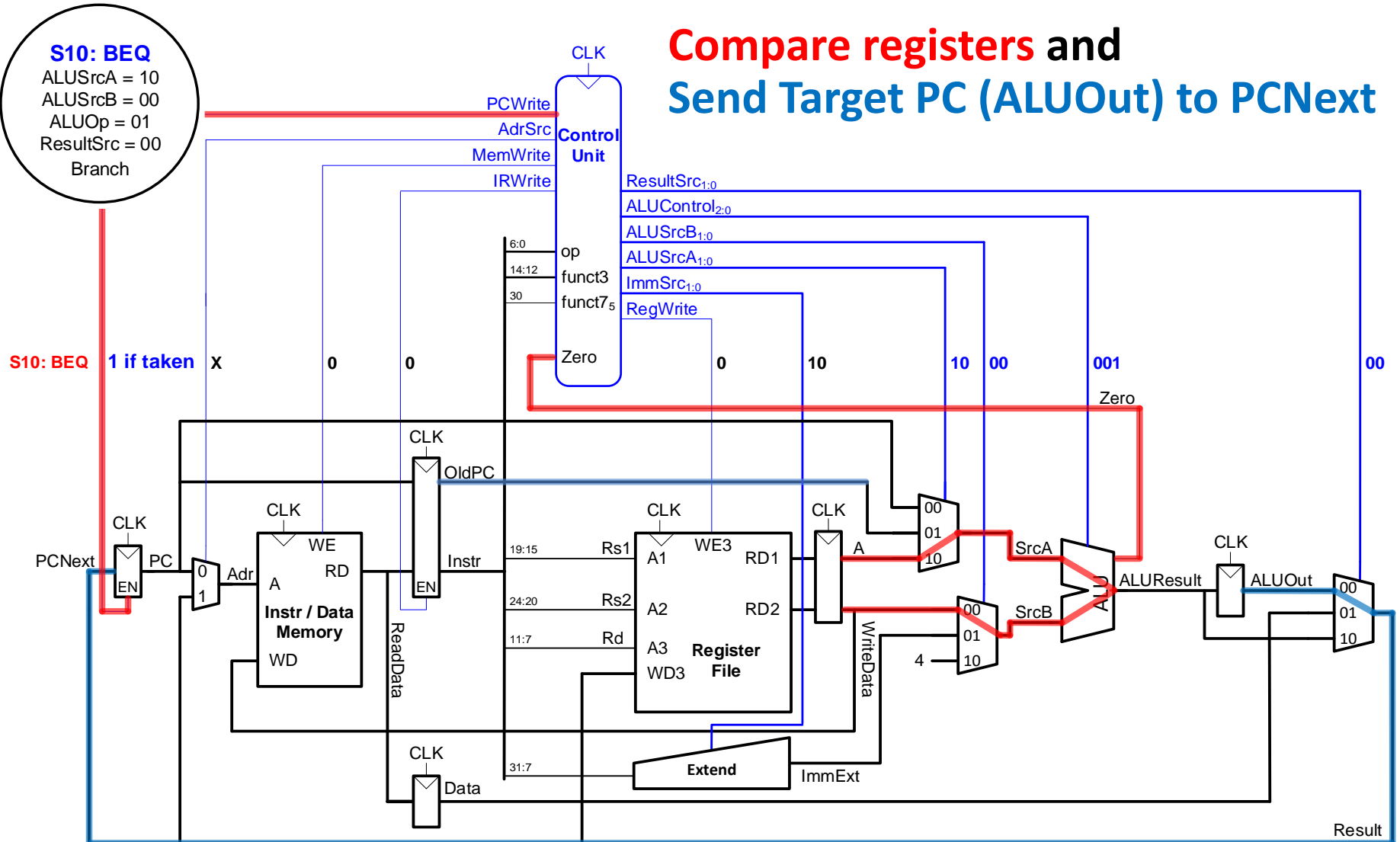


Main FSM: beq



Main FSM: beq Datapath

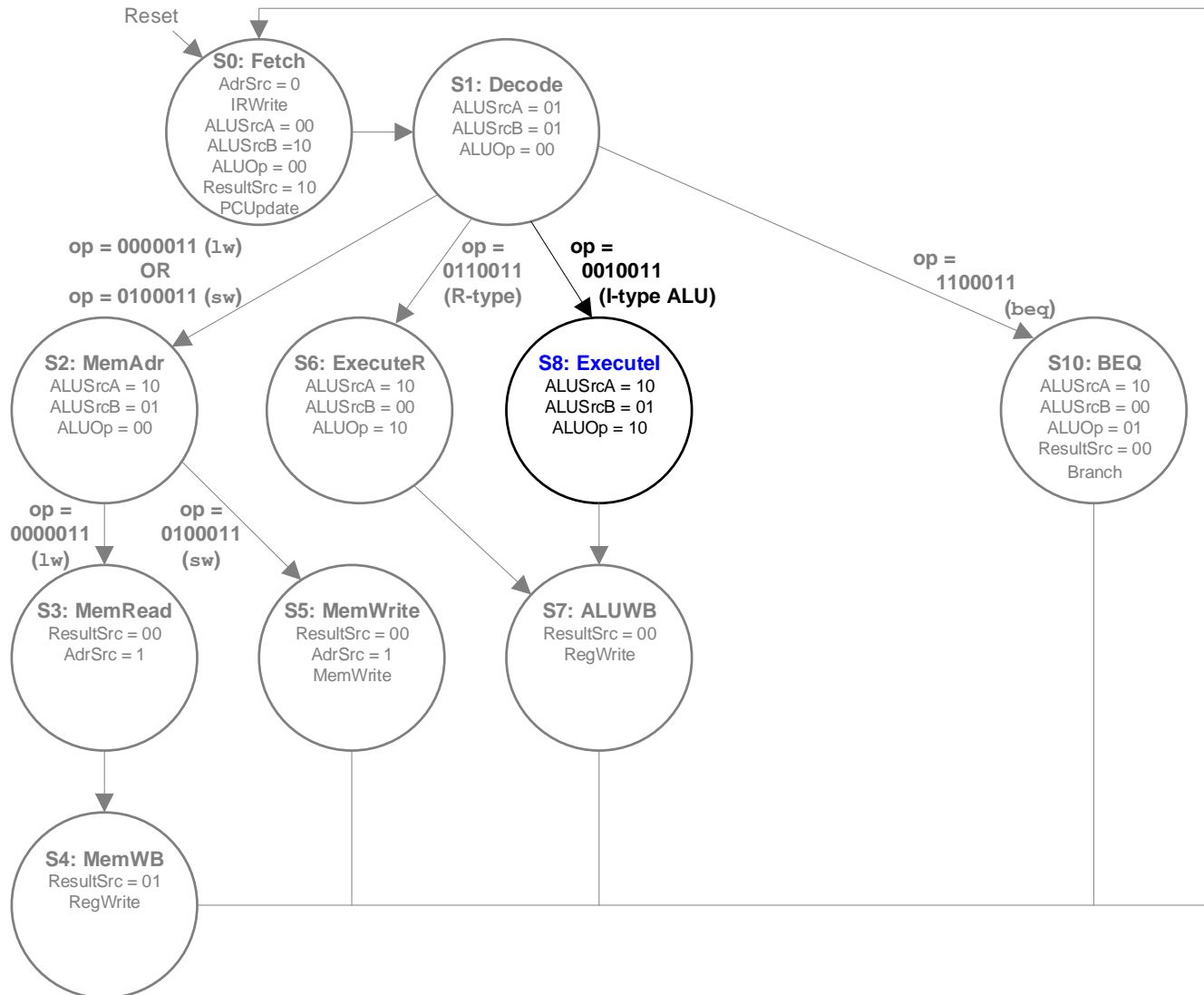
**Compare registers and
Send Target PC (ALUOut) to PCNext**



Chapter 7: Microarchitecture

Extending the RISC-V Multicycle Processor

Main FSM: I-Type ALU Execute

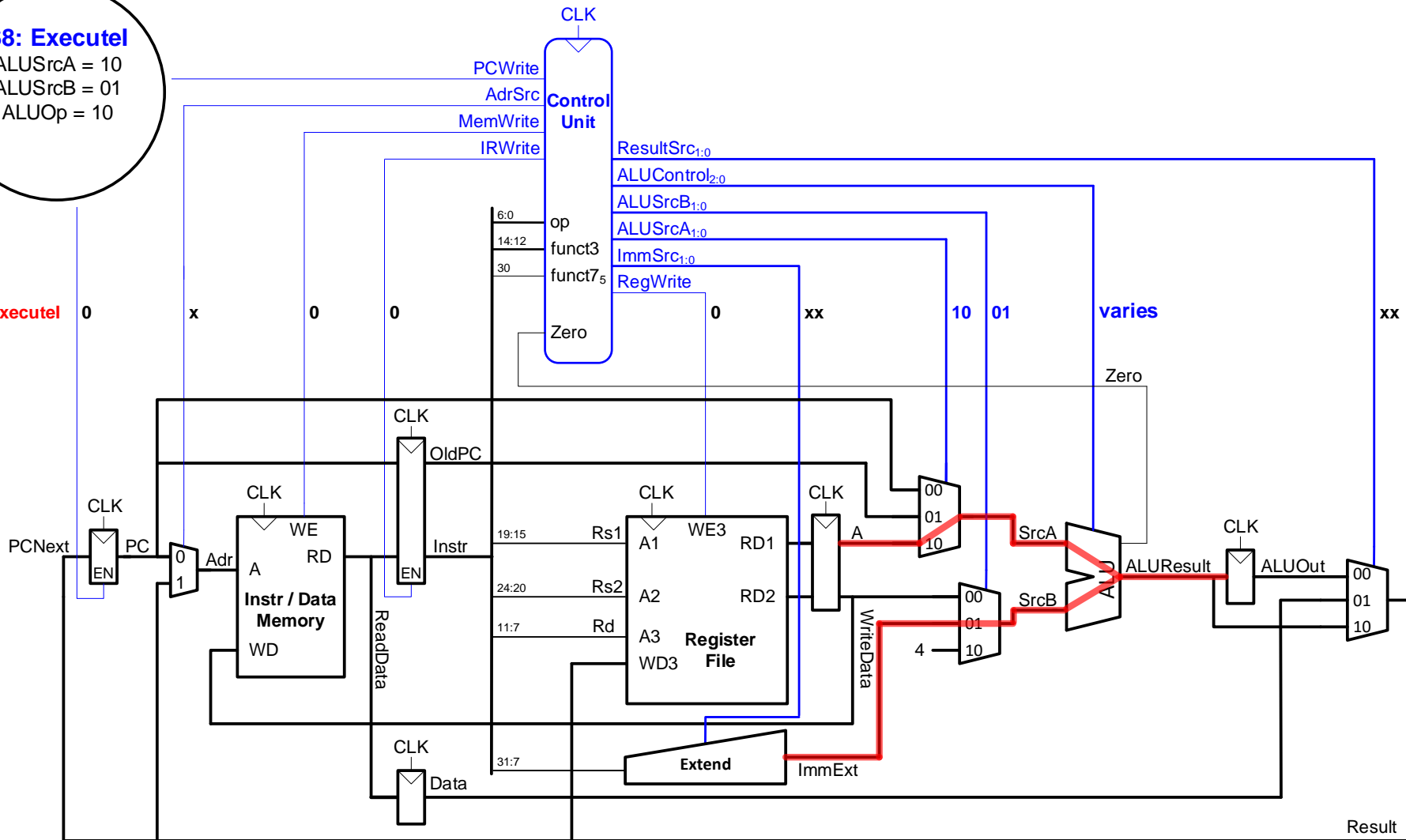


Main FSM: I-Type ALU Exec. Datapath

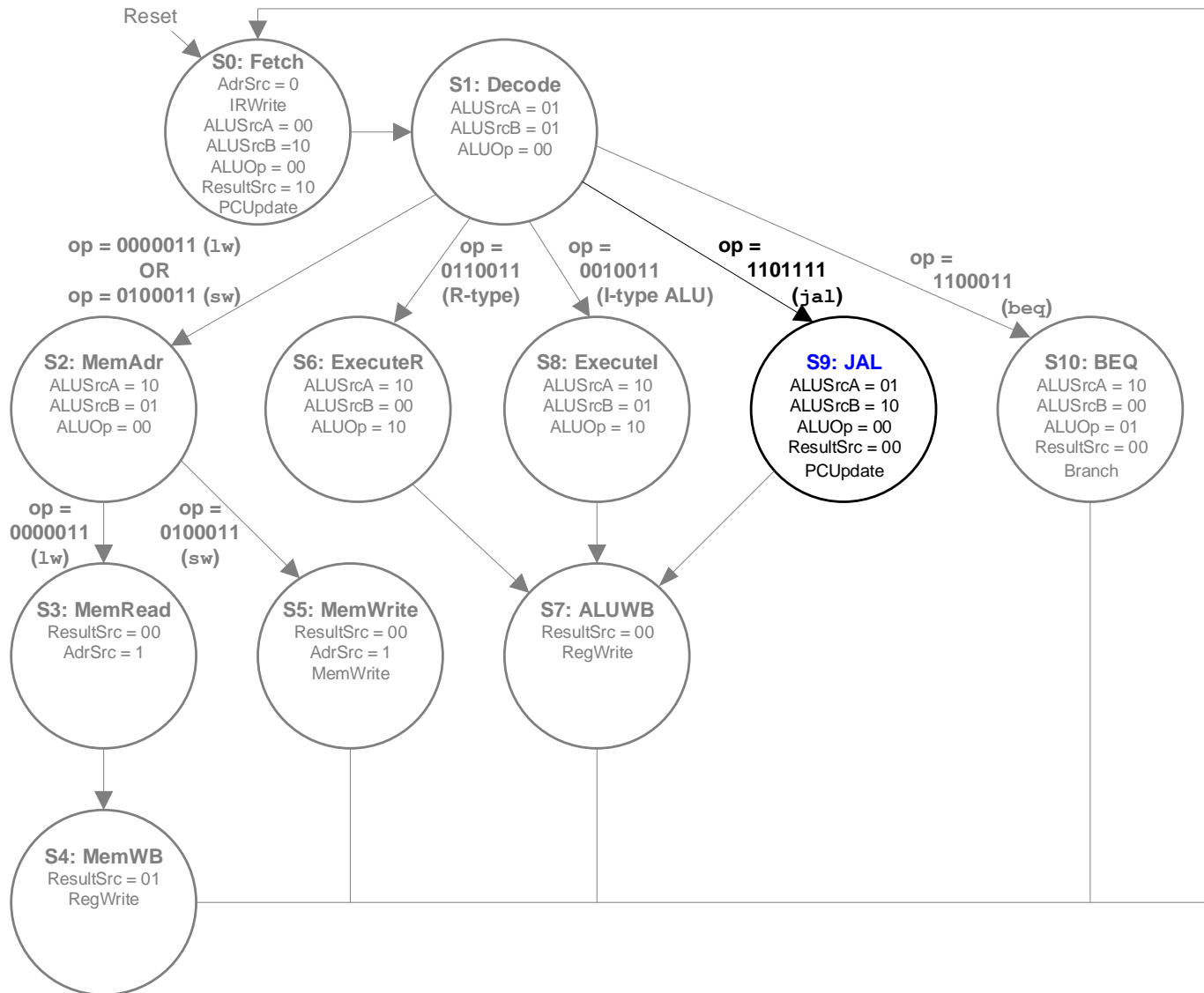
S8: Executel

ALUSrcA = 10
ALUSrcB = 01
ALUOp = 10

S8: Executel



Main FSM: jal

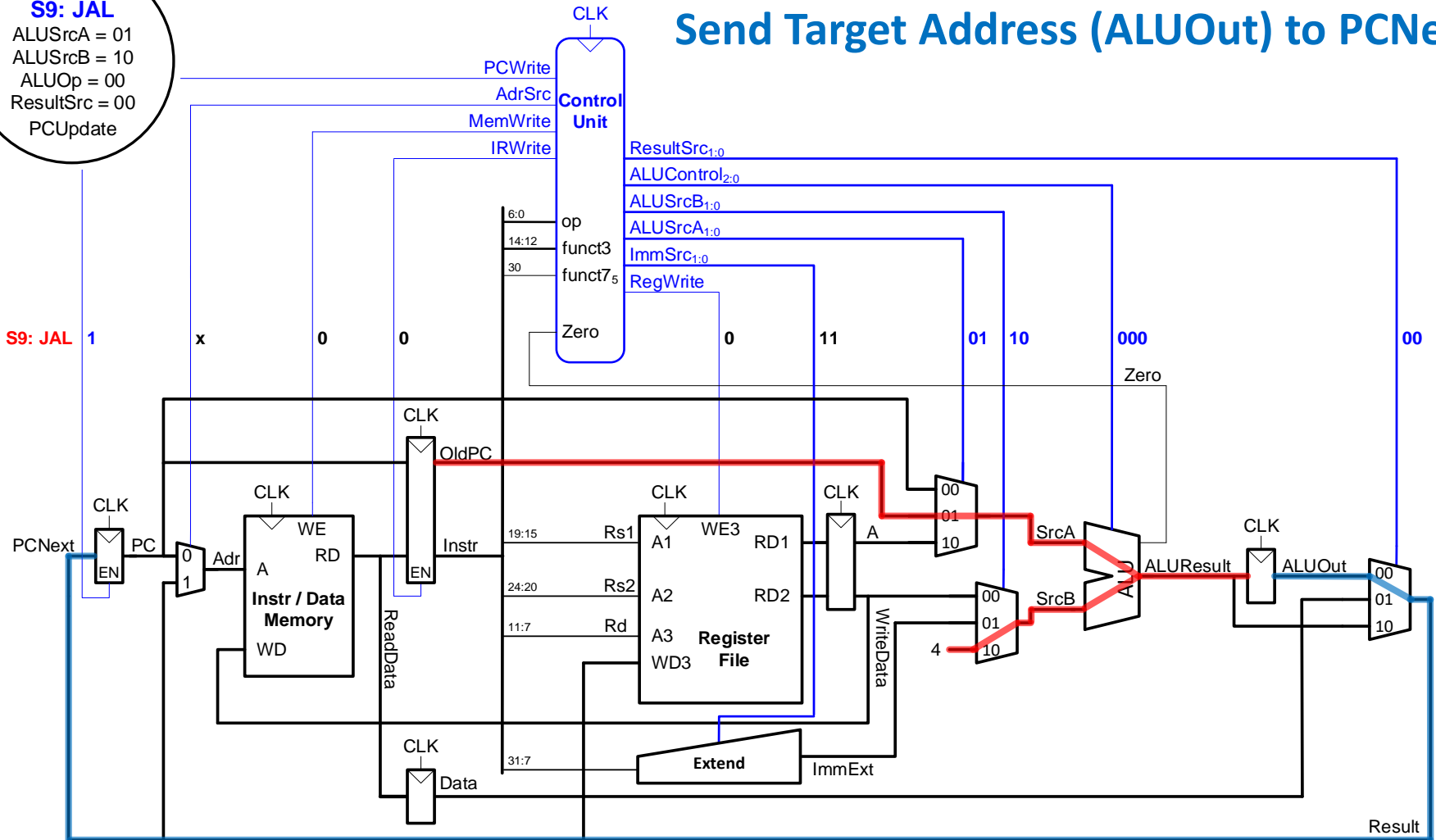


Main FSM: jal Datapath

Calculate PC + 4 and
Send Target Address (ALUOut) to PCNext

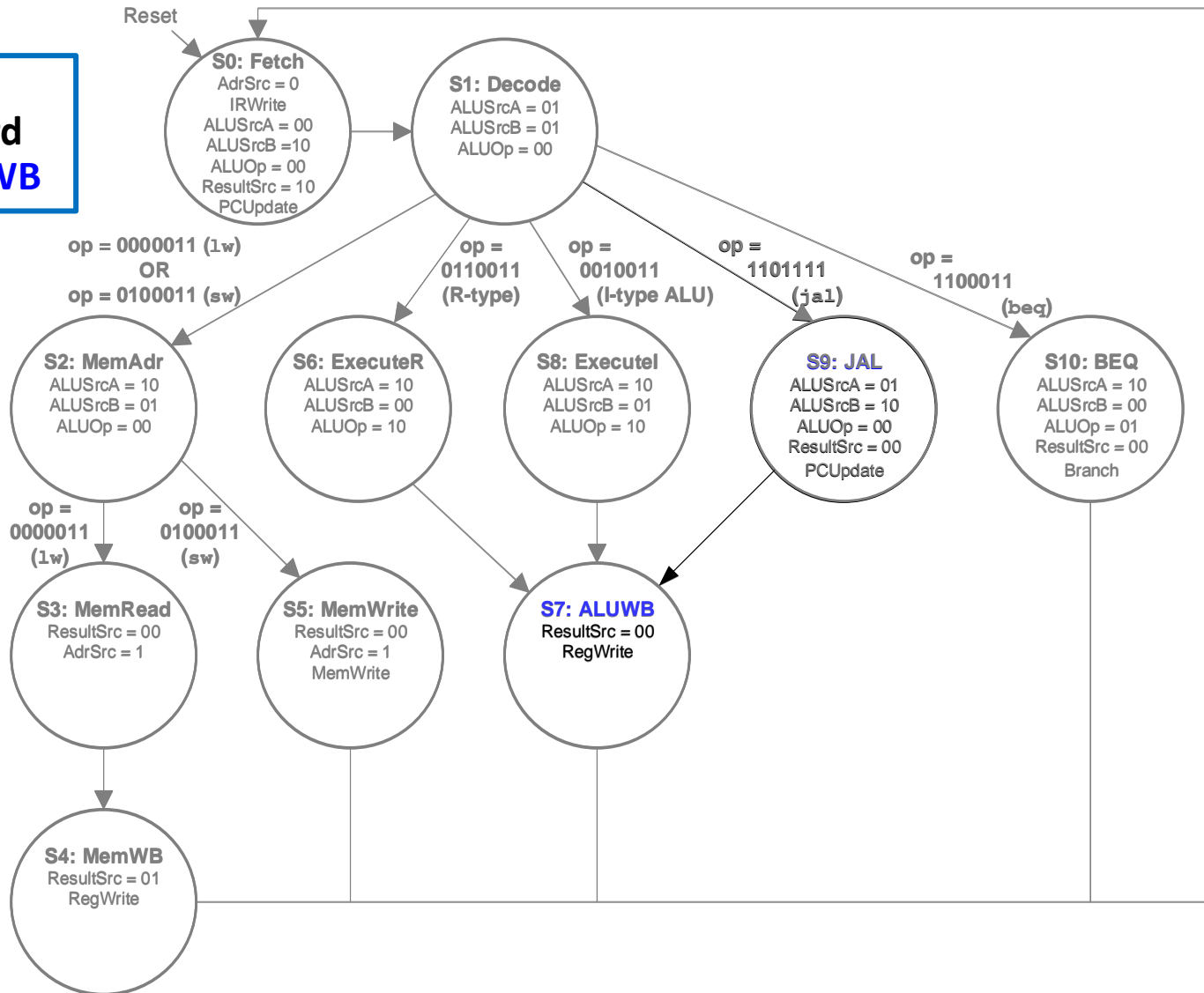
S9: JAL

ALUSrcA = 01
ALUSrcB = 10
ALUOp = 00
ResultSrc = 00
PCUpdate



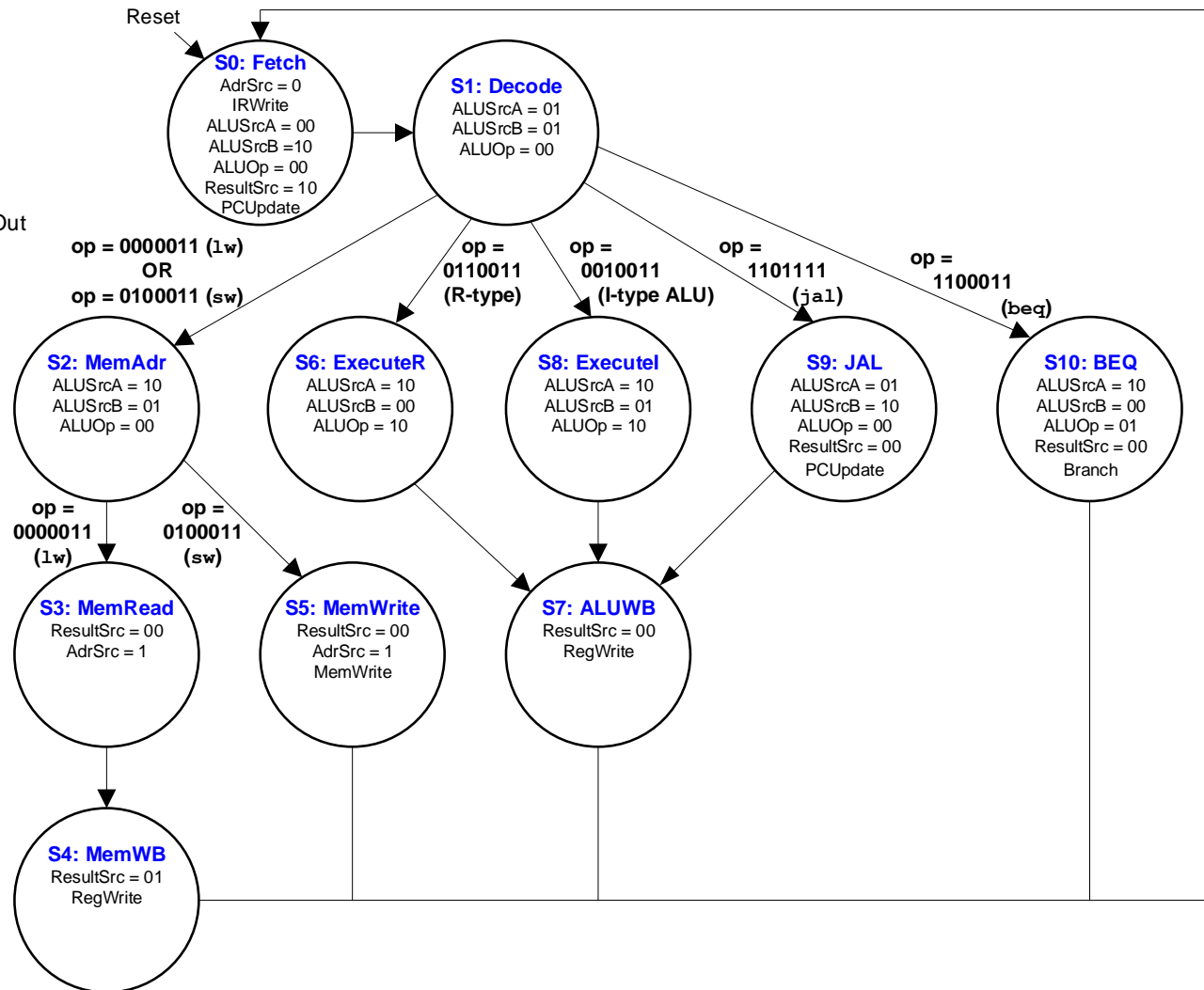
Main FSM: jal

PC + 4 is
written to **rd**
in **S7: ALUWB**



Multicycle Processor Main FSM

State	Datapath μ Op
Fetch	Instr \leftarrow Mem[PC]; PC \leftarrow PC+4
Decode	ALUOut \leftarrow PCTarget
MemAdr	ALUOut \leftarrow rs1 + imm
MemRead	Data \leftarrow Mem[ALUOut]
MemWB	rd \leftarrow Data
MemWrite	Mem[ALUOut] \leftarrow rd
ExecuteR	ALUOut \leftarrow rs1 op rs2
ExecuteI	ALUOut \leftarrow rs1 op imm
ALUWB	rd \leftarrow ALUOut
BEQ	ALUResult = rs1-rs2; if Zero, PC \leftarrow ALUOut
JAL	PC \leftarrow ALUOut; ALUOut \leftarrow PC+4



Chapter 7: Microarchitecture

Multicycle Performance

Multicycle Processor Performance

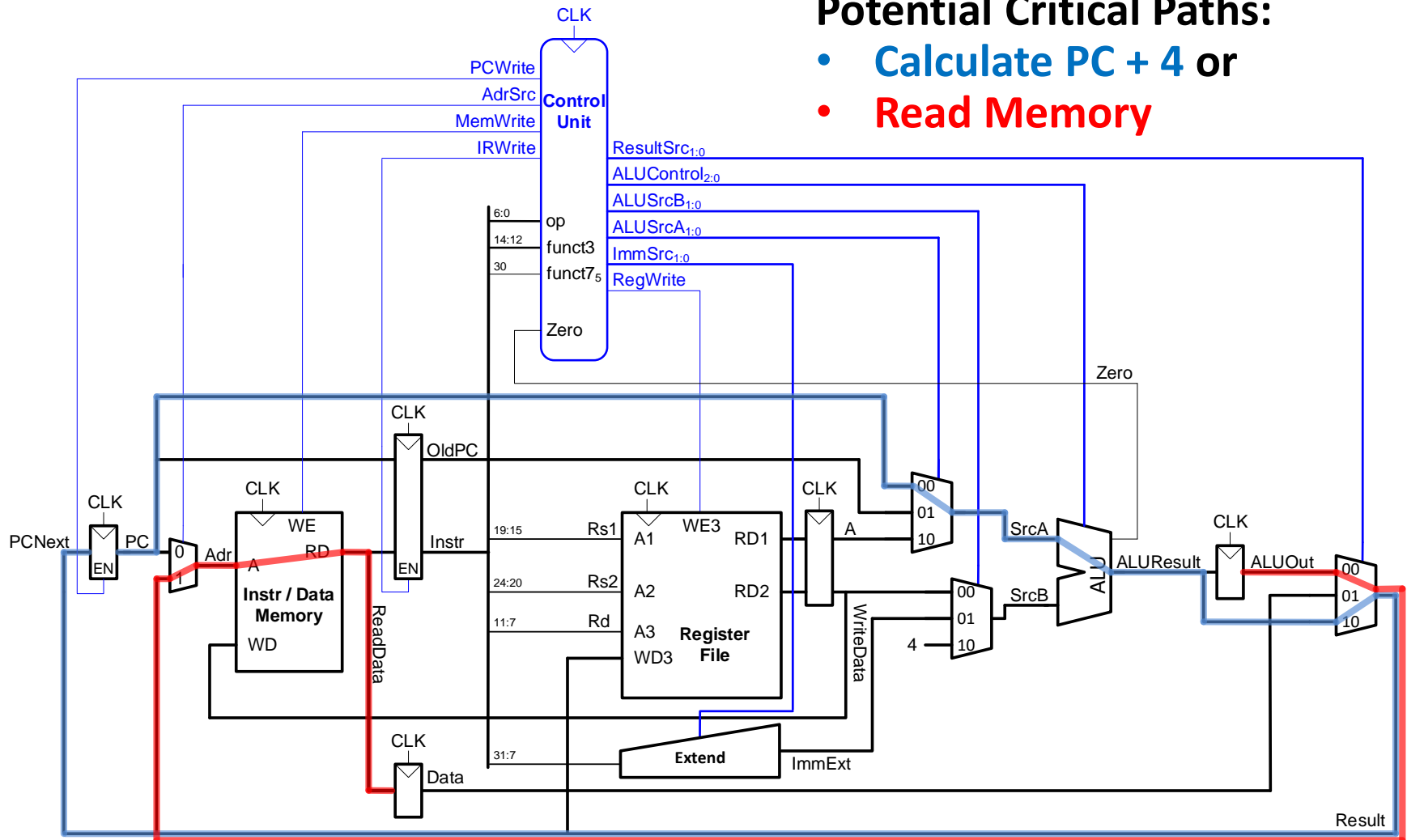
- Instructions take different number of cycles:
 - 3 cycles: `beq`
 - 4 cycles: `R-type`, `addi`, `sw`, `jal`
 - 5 cycles: `lw`
- CPI is weighted average
- SPECINT2000 benchmark:
 - 25% loads
 - 10% stores
 - 13% branches
 - 52% R-type

$$\text{Average CPI} = (0.13)(3) + (0.52 + 0.10)(4) + (0.25)(5) = 4.12$$

Multicycle Critical Path

Potential Critical Paths:

- Calculate PC + 4 or
- Read Memory



Multicycle Processor Performance

Multicycle critical path:

- **Assumptions:**
 - RF is faster than memory
 - Writing memory is faster than reading memory

$$T_{c_multi} = t_{pcq} + t_{dec} + 2t_{mux} + \max(t_{ALU}, t_{mem}) + t_{setup}$$

Multicycle Performance Example

Element	Parameter	Delay (ps)
Register clock-to-Q	t_{pcq_PC}	40
Register setup	t_{setup}	50
Multiplexer	t_{mux}	30
AND-OR gate	t_{AND-OR}	20
ALU	t_{ALU}	120
Decoder (Control Unit)	t_{dec}	25
Extend unit	t_{dec}	35
Memory read	t_{mem}	200
Register file read	t_{RFread}	100
Register file setup	$t_{RFsetup}$	60

$$T_{c_multi} = t_{pcq} + t_{dec} + 2t_{mux} + \max(t_{ALU}, t_{mem}) + t_{setup}$$

$$=$$

Multicycle Performance Example

For a program with **100 billion** instructions executing on a **multicycle** RISC-V processor

- **CPI** = 4.12 cycles/instruction
- **Clock cycle time:** $T_{c_multi} = 375 \text{ ps}$

Execution Time = (# instructions) \times CPI $\times T_c$

Chapter 7: Microarchitecture

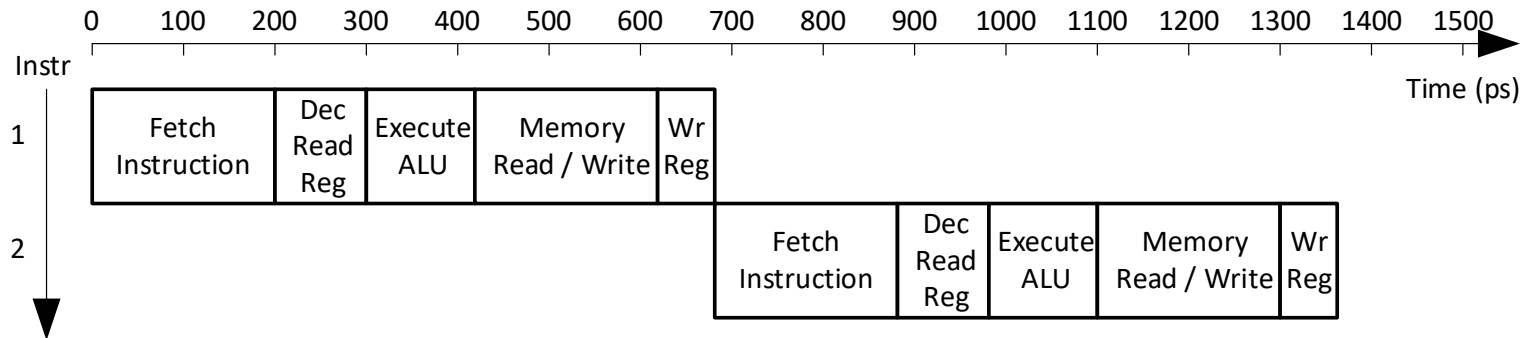
Pipelined RISC-V Processor

Pipelined RISC-V Processor

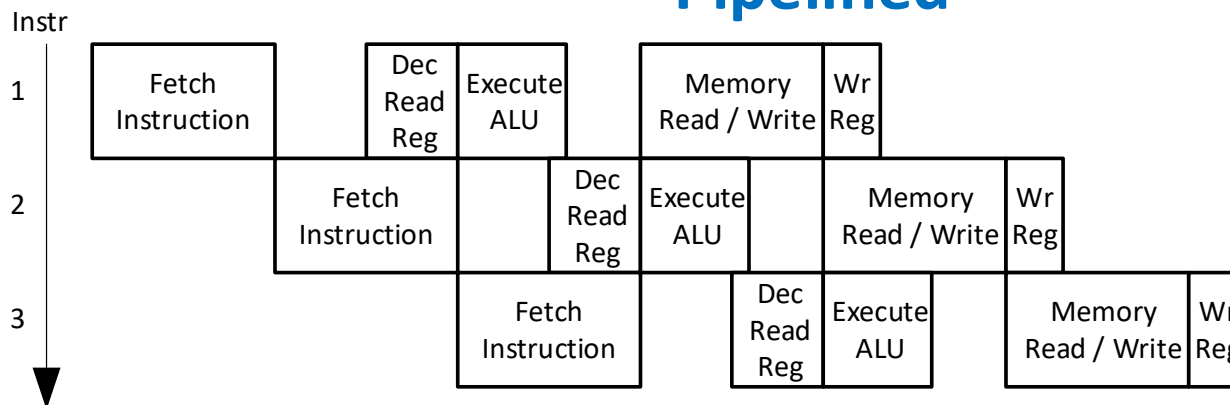
- **Temporal parallelism**
- Divide single-cycle processor into **5 stages**:
 - Fetch
 - Decode
 - Execute
 - Memory
 - Writeback
- Add **pipeline registers** between stages

Single-Cycle vs. Pipelined Processor

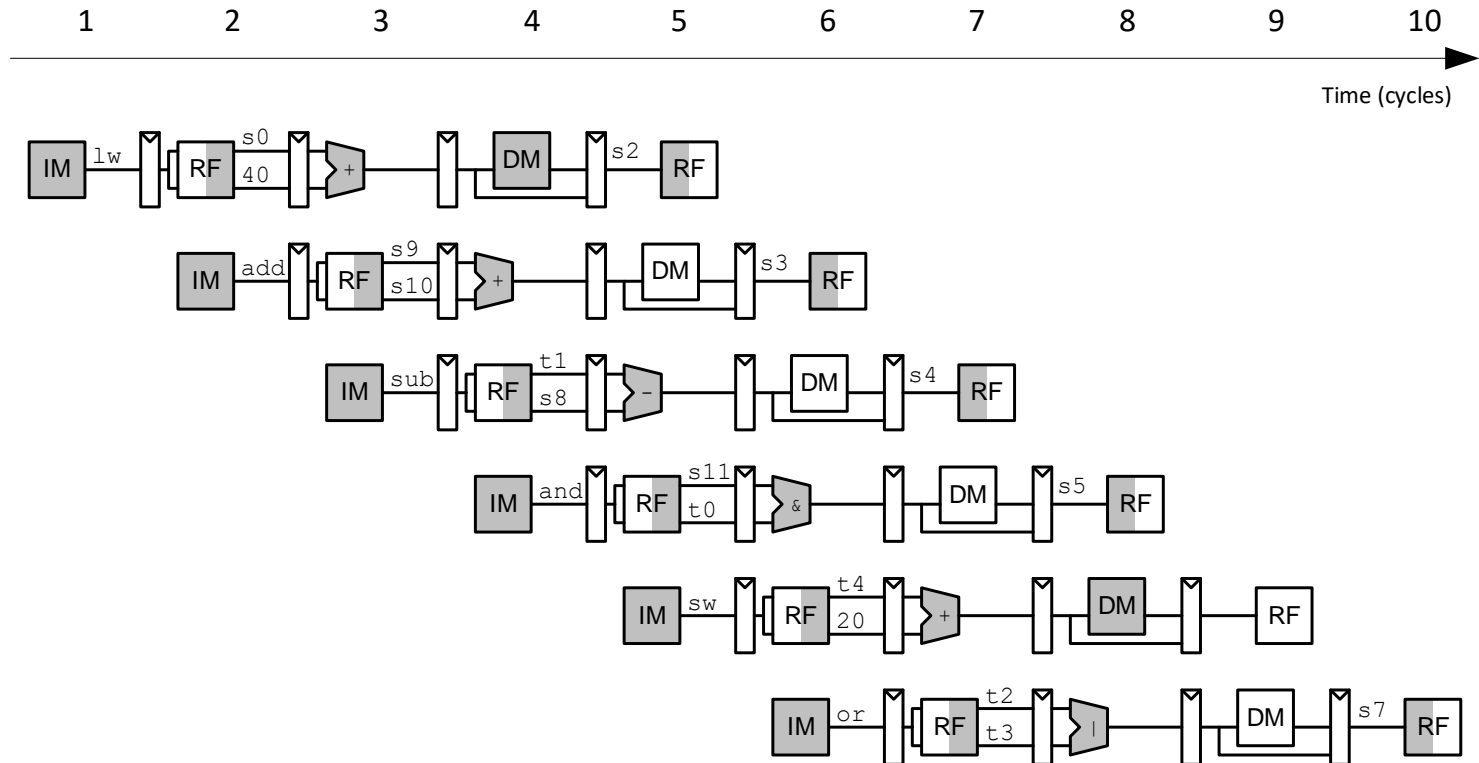
Single-Cycle



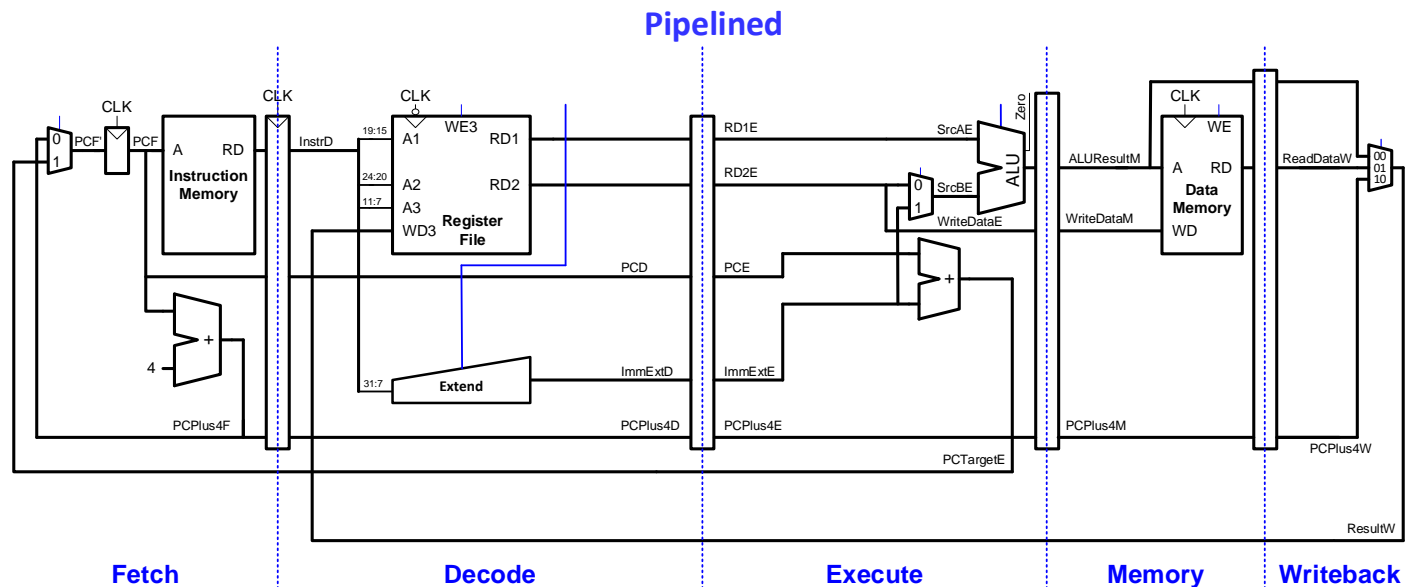
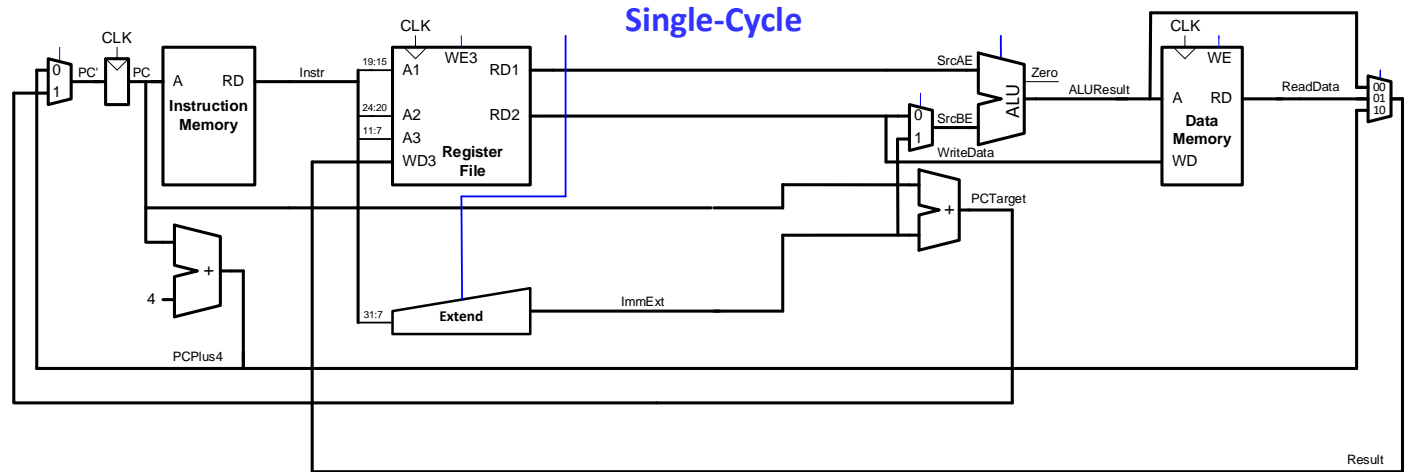
Pipelined



Pipelined Processor Abstraction

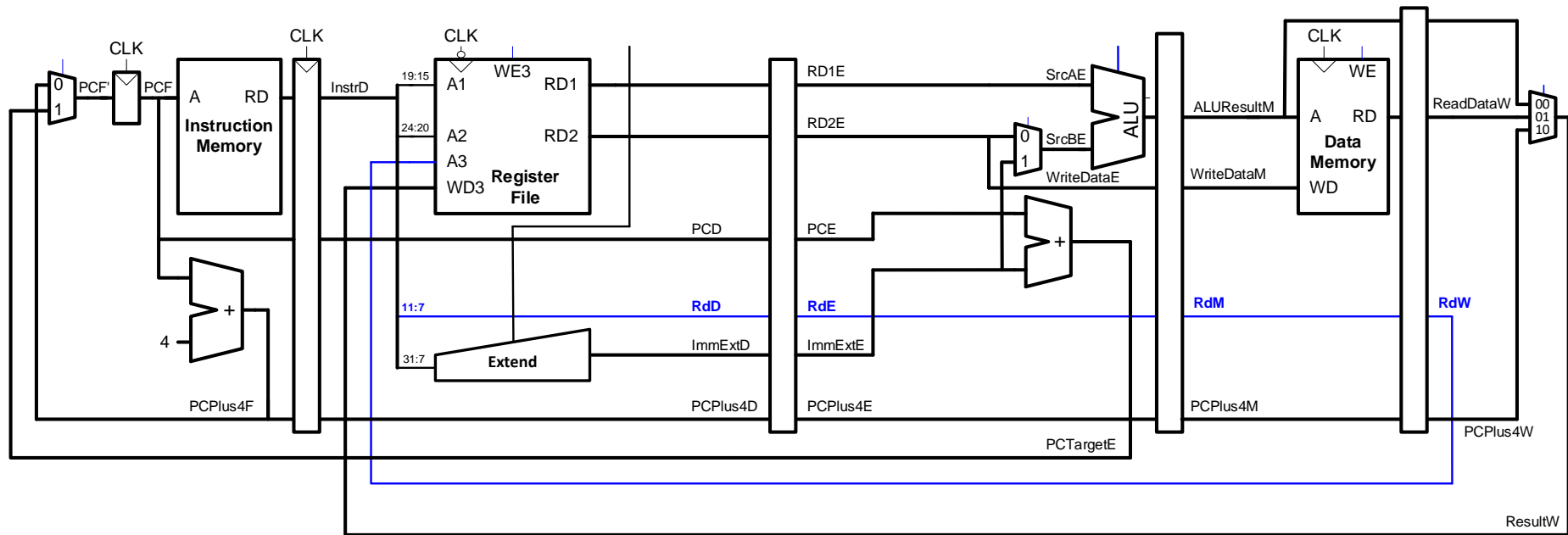


Single-Cycle & Pipelined Datapaths



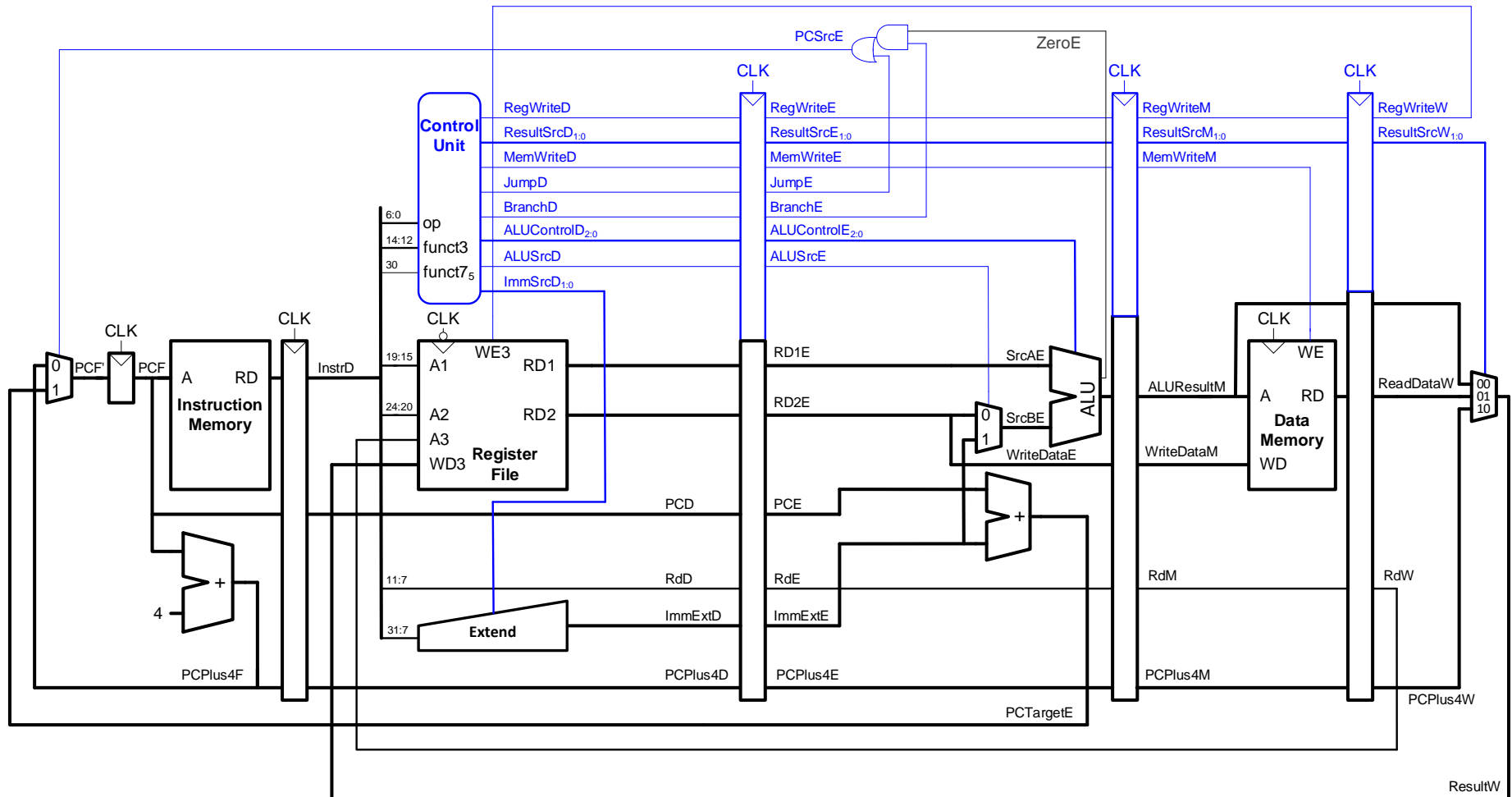
Signals in Pipelined Processor are appended with first letter of stage (i.e., PC**F**, PC**D**, PC**E**).

Corrected Pipelined Datapath



- ***Rd*** must arrive at same time as ***Result***
- Register file written on **falling edge** of *CLK*

Pipelined Processor with Control



- **Same control unit** as single-cycle processor
- **Control signals travel with** the instruction (drop off when used)

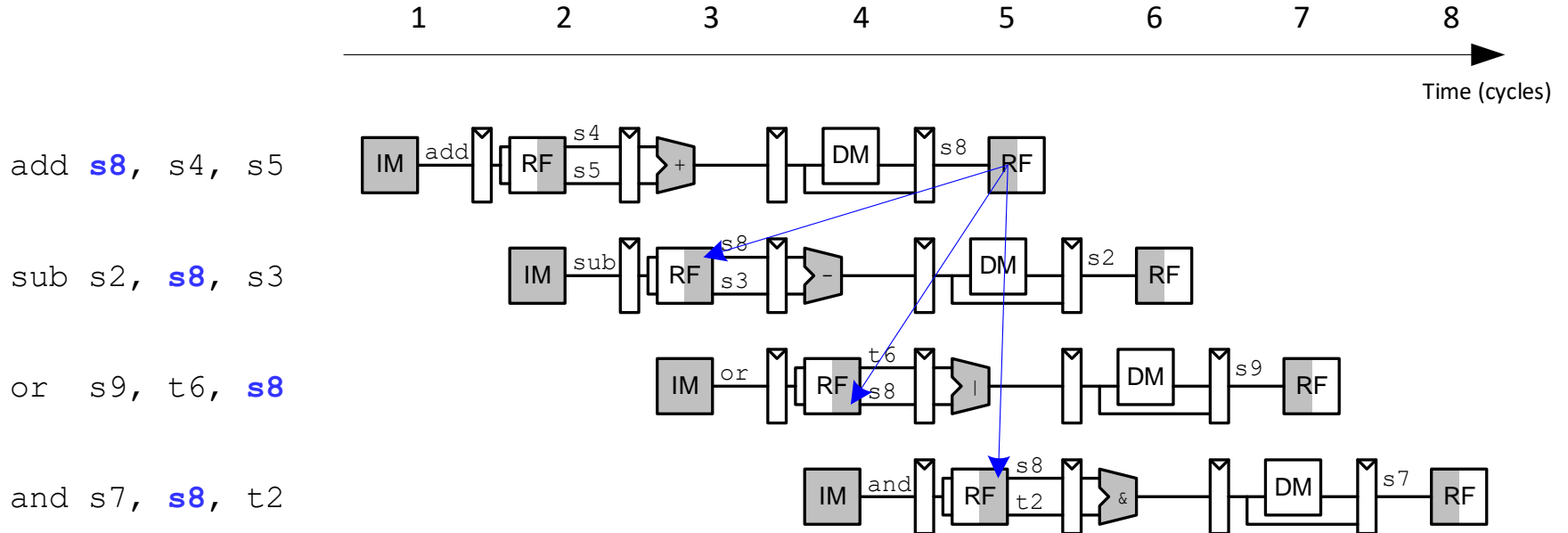
Chapter 7: Microarchitecture

Pipelined Processor Hazards

Pipelined Hazards

- When an instruction depends on result from instruction that hasn't completed
- Types:
 - **Data hazard:** register value not yet written back to register file
 - **Control hazard:** next instruction not decided yet (caused by branch)

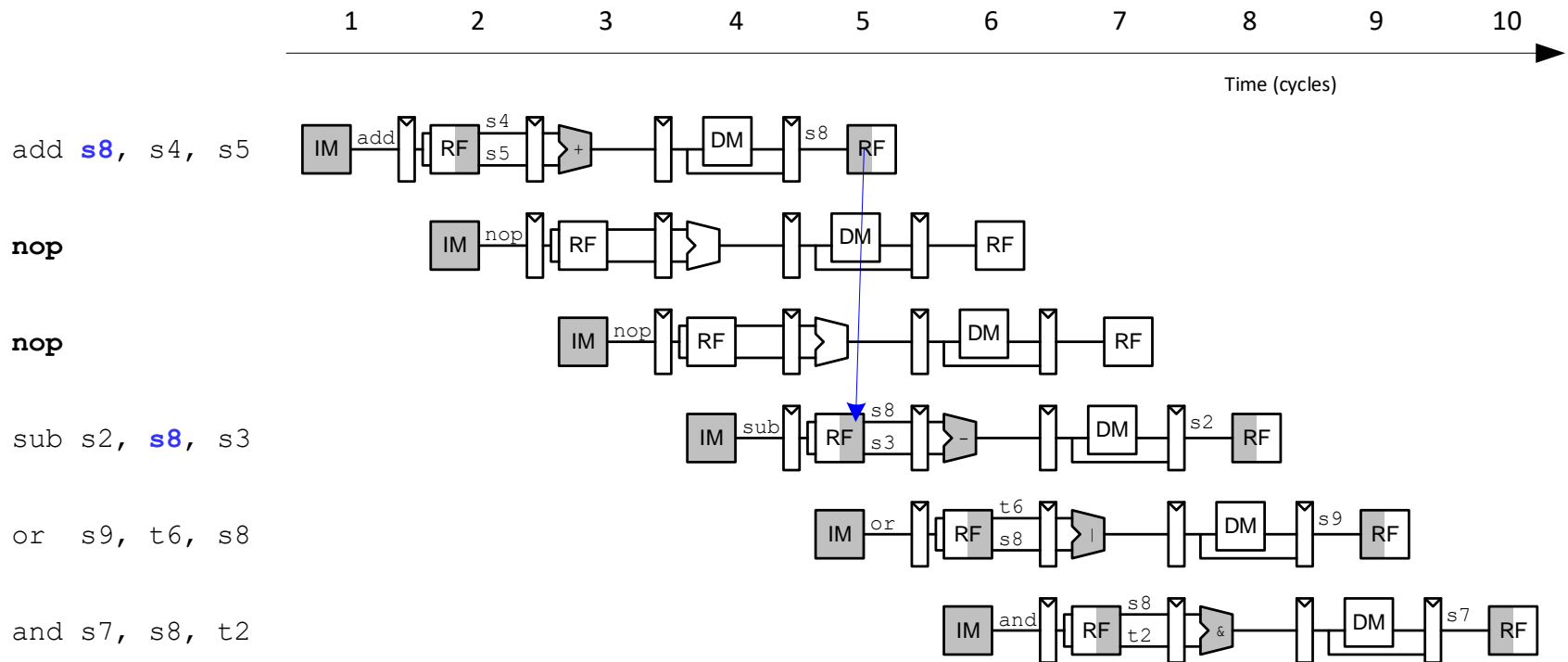
Data Hazard



Handling Data Hazards

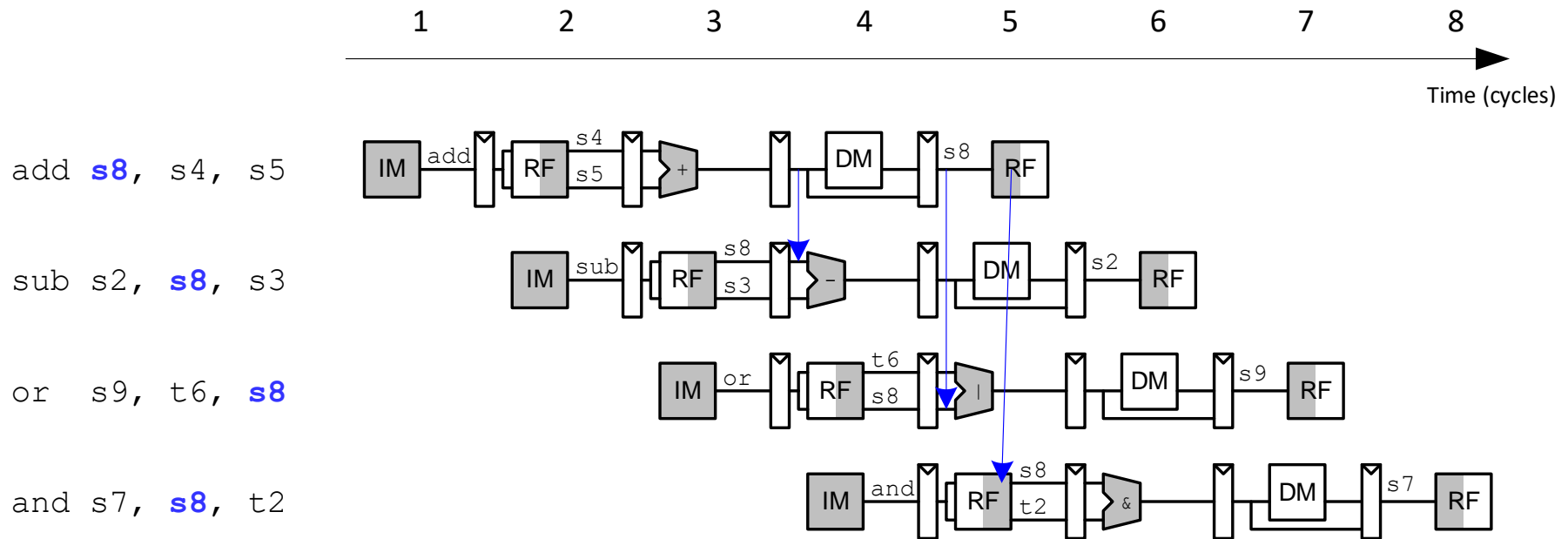
Handling Data Hazards

- **Insert** enough **nops** for result to be ready
- Or move independent useful instructions forward



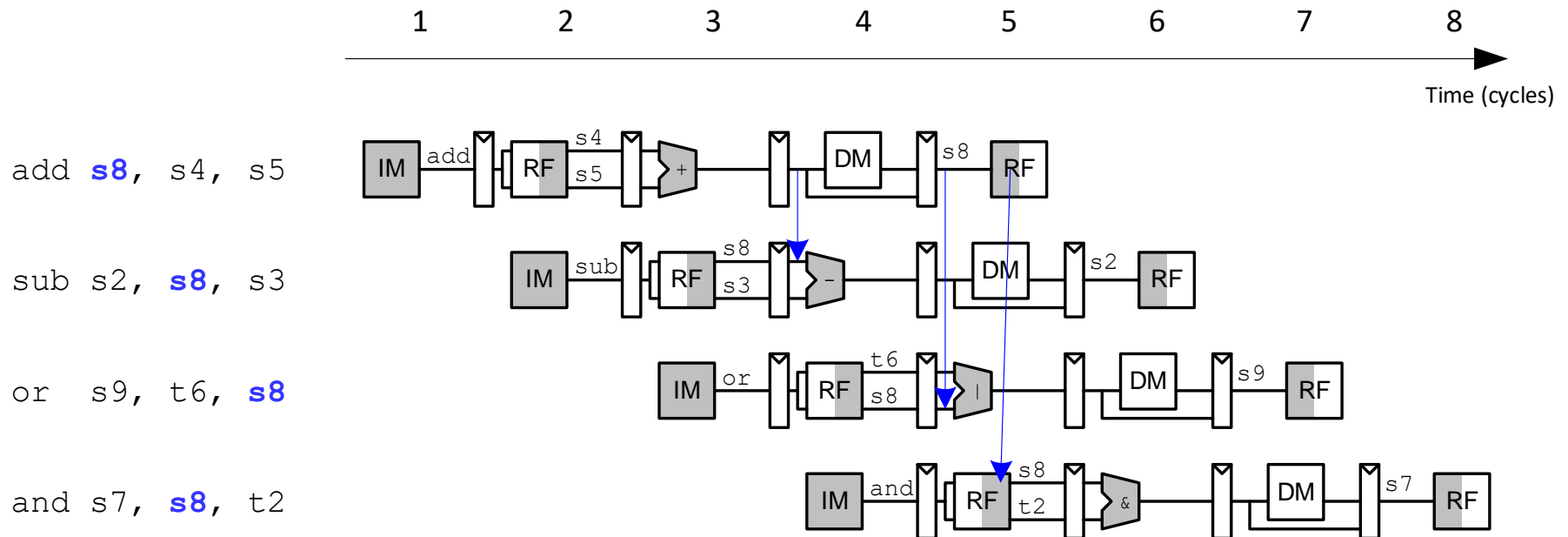
Data Forwarding

- Data is **available on internal busses** before it is written back to the register file (RF).
- **Forward data** from internal busses **to Execute stage**.

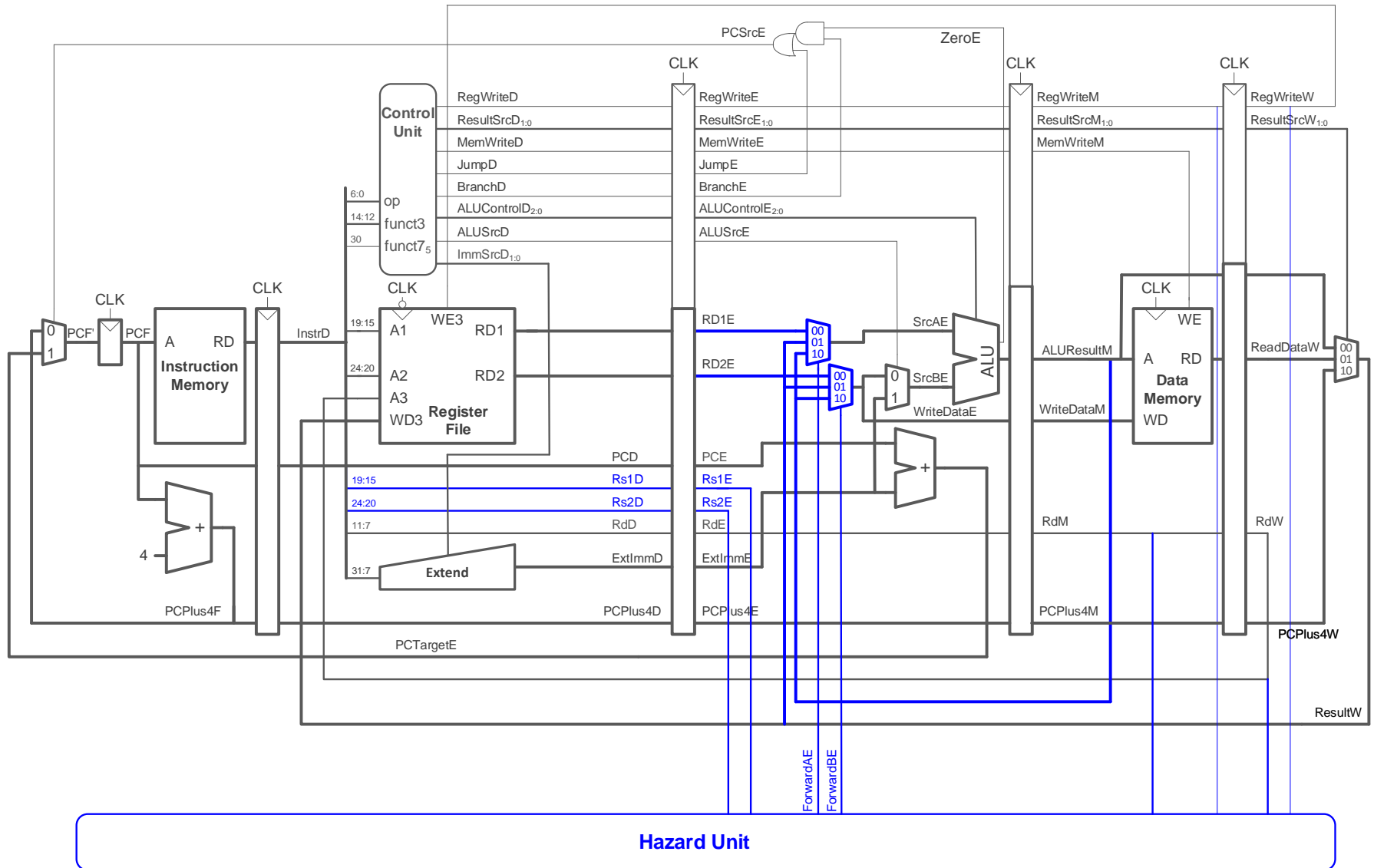


Data Forwarding

- Check if source register **in Execute stage matches** destination register of instruction **in Memory or Writeback stage**.
- If so, forward result.



Data Forwarding: Hazard Unit



Data Forwarding

- **Case 1:** **Execute** stage $Rs1$ or $Rs2$ matches **Memory** stage Rd ?
Forward from Memory stage
- **Case 2:** **Execute** stage $Rs1$ or $Rs2$ matches **Writeback** stage Rd ?
Forward from Writeback stage
- **Case 3:** Otherwise use value read from register file (as usual)

Equations for $Rs1$:

```
if      (( $Rs1E == RdM$ ) AND  $RegWriteM$ )           // Case 1
         $ForwardAE = 10$ 
else if (( $Rs1E == RdW$ ) AND  $RegWriteW$ )           // Case 2
         $ForwardAE = 01$ 
else     $ForwardAE = 00$                            // Case 3
```

***ForwardBE** equations are similar (replace $Rs1E$ with $Rs2E$)*

Data Forwarding

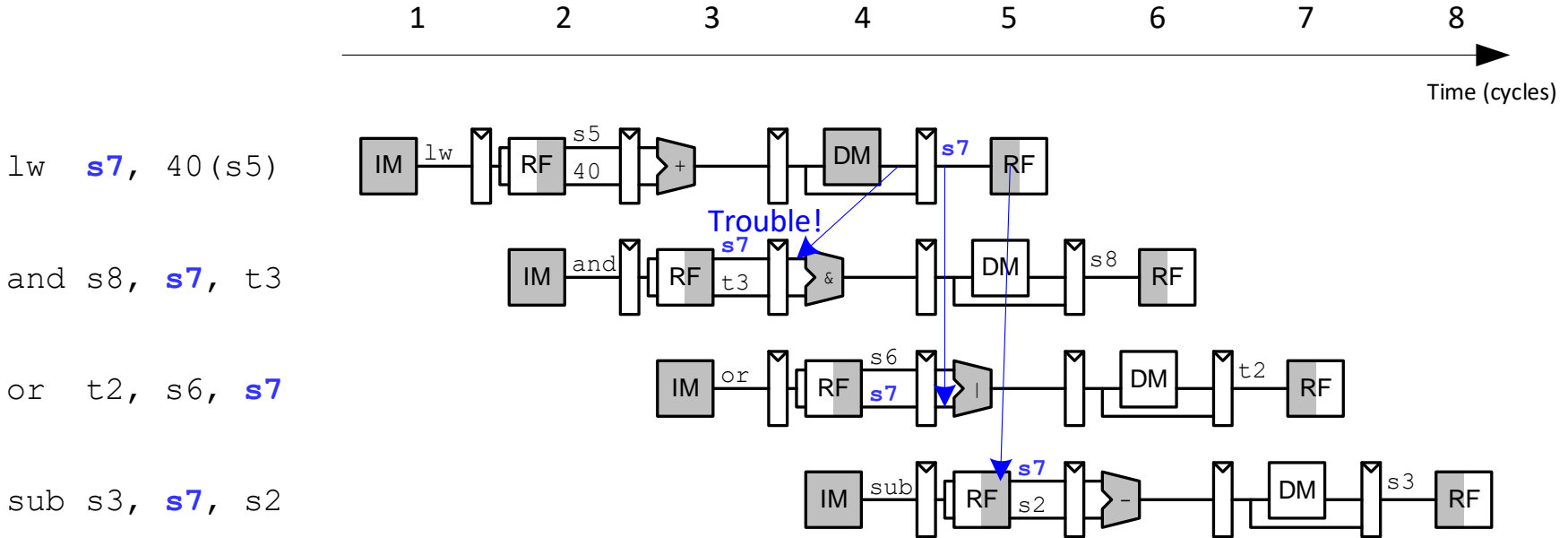
- **Case 1:** **Execute** stage $Rs1$ or $Rs2$ matches **Memory** stage Rd ?
Forward from Memory stage
- **Case 2:** **Execute** stage $Rs1$ or $Rs2$ matches **Writeback** stage Rd ?
Forward from Writeback stage
- **Case 3:** Otherwise use value read from register file (as usual)

Equations for $Rs1$:

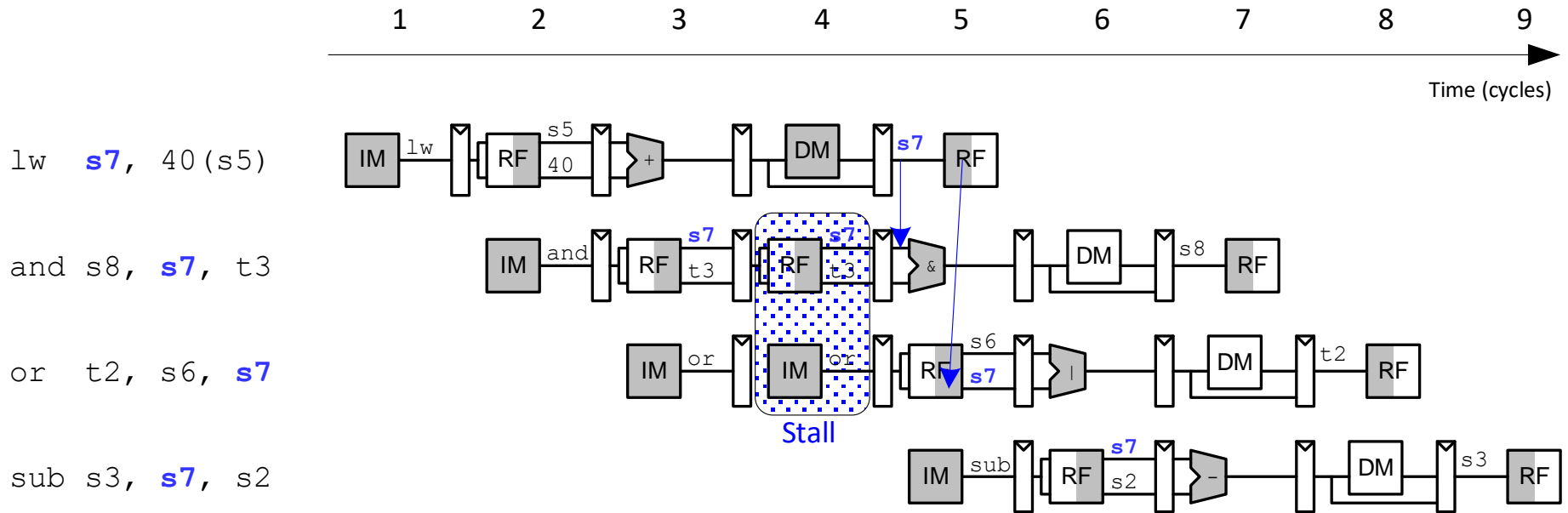
```
if      (( $Rs1E == RdM$ ) AND  $RegWriteM$ ) AND ( $Rs1E != 0$ ) // Case 1
        ForwardAE = 10
else if (( $Rs1E == RdW$ ) AND  $RegWriteW$ ) AND ( $Rs1E != 0$ ) // Case 2
        ForwardAE = 01
else    ForwardAE = 00                                     // Case 3
```

ForwardBE equations are similar (replace $Rs1E$ with $Rs2E$)

Data Hazard due to \perp_w Dependency



Stalling to solve 1w Data Dependency



Stalling Logic

- Is either **source register in the Decode stage** the same as the **destination register in the Execute stage**?

AND

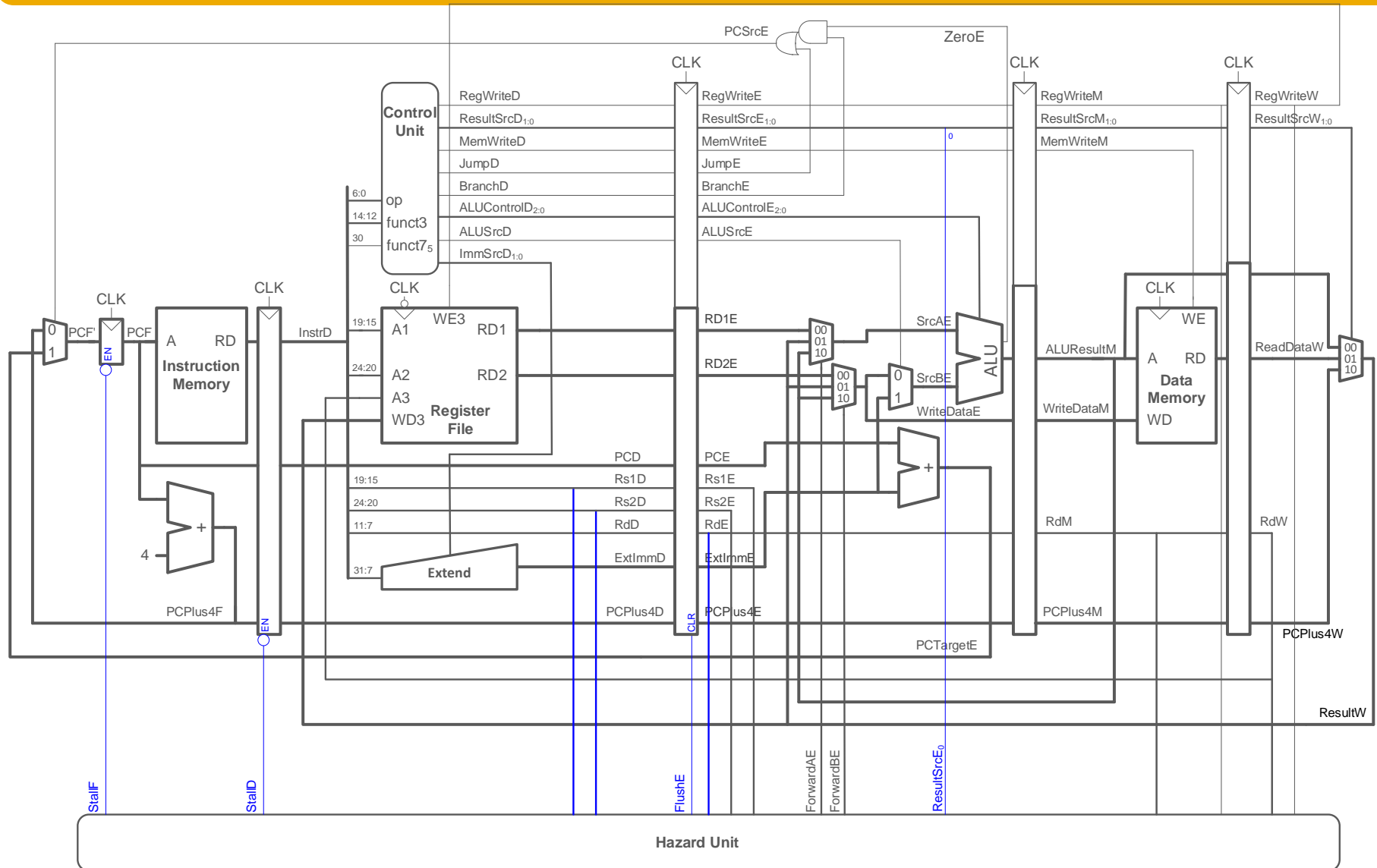
- Is the instruction in the **Execute stage a lw**?

$lwStall = ((Rs1D == RdE) \text{ OR } (Rs2D == RdE)) \text{ AND } ResultSrcE_0$

$StallF = StallD = FlushE = lwStall$

(Stall the Fetch and Decode stages, and flush the Execute stage.)

Stalling Hardware



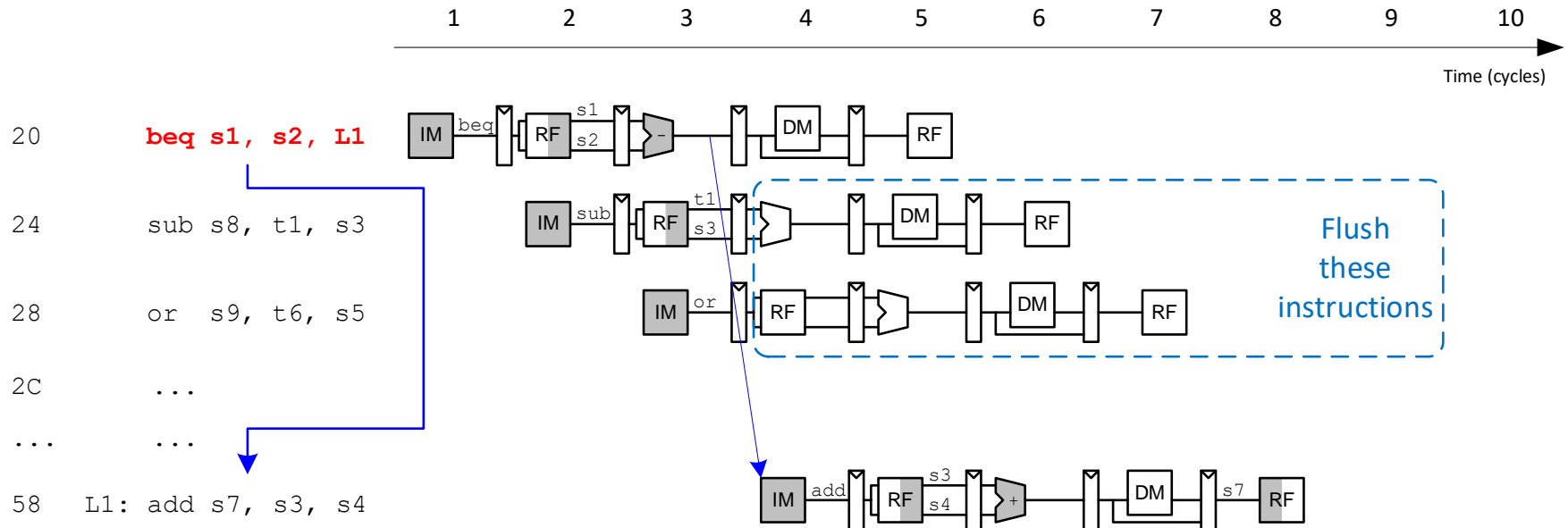
Chapter 7: Microarchitecture

Pipelined Processor Control Hazards

Control Hazards

- **beq:**
 - Branch **not determined** until the **Execute stage** of pipeline
 - **Instructions** after branch **fetch**ed before branch occurs
 - These **2 instructions must be flushed** if branch happens

Control Hazards



Branch misprediction penalty:

The number of instructions flushed when a branch is taken (in this case, 2 instructions)

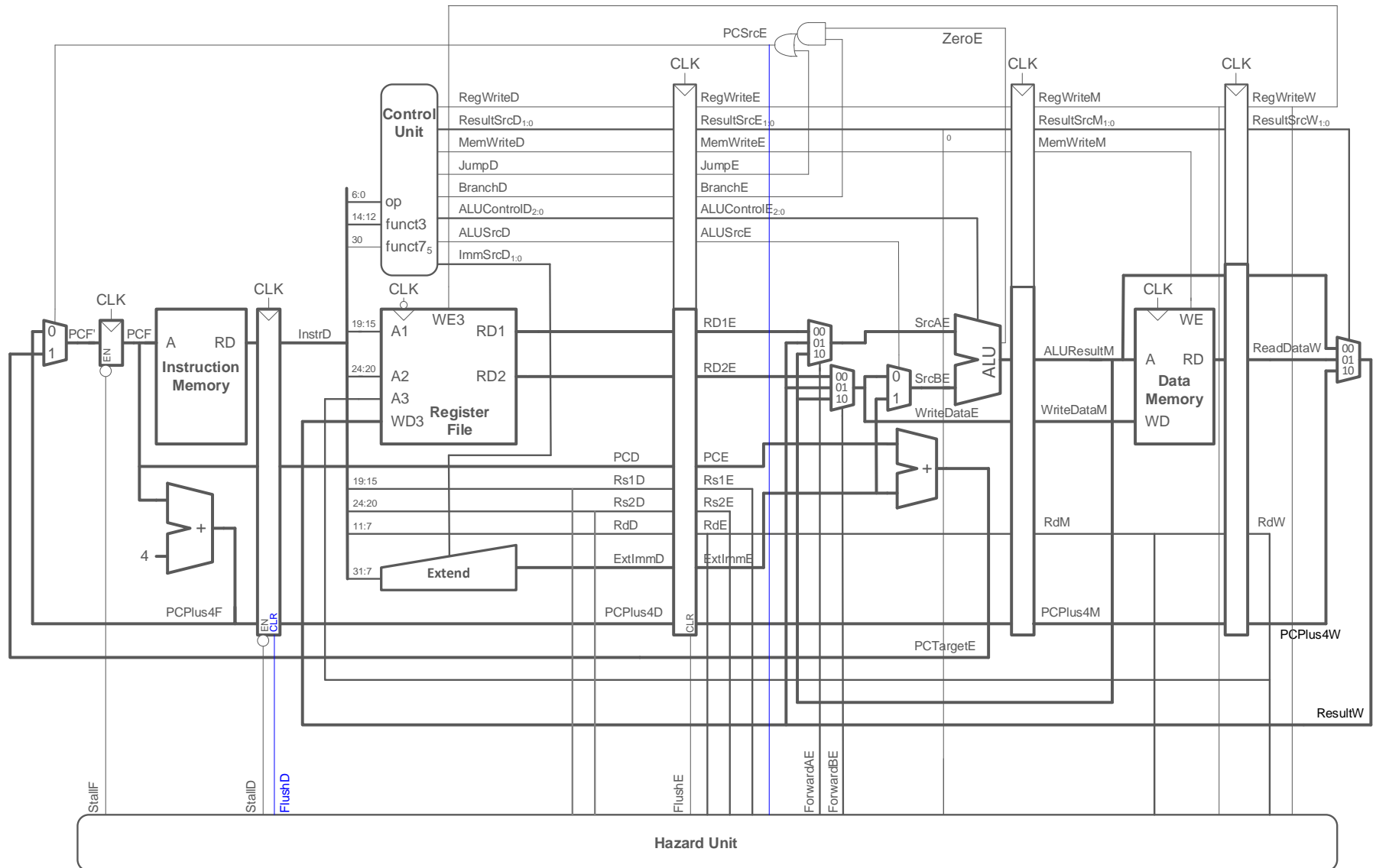
Control Hazards: Flushing Logic

- If branch is taken in execute stage, need to flush the instructions in the Fetch and Decode stages
 - Do this by clearing Decode and Execute Pipeline registers using *FlushD* and *FlushE*
- **Equations:**

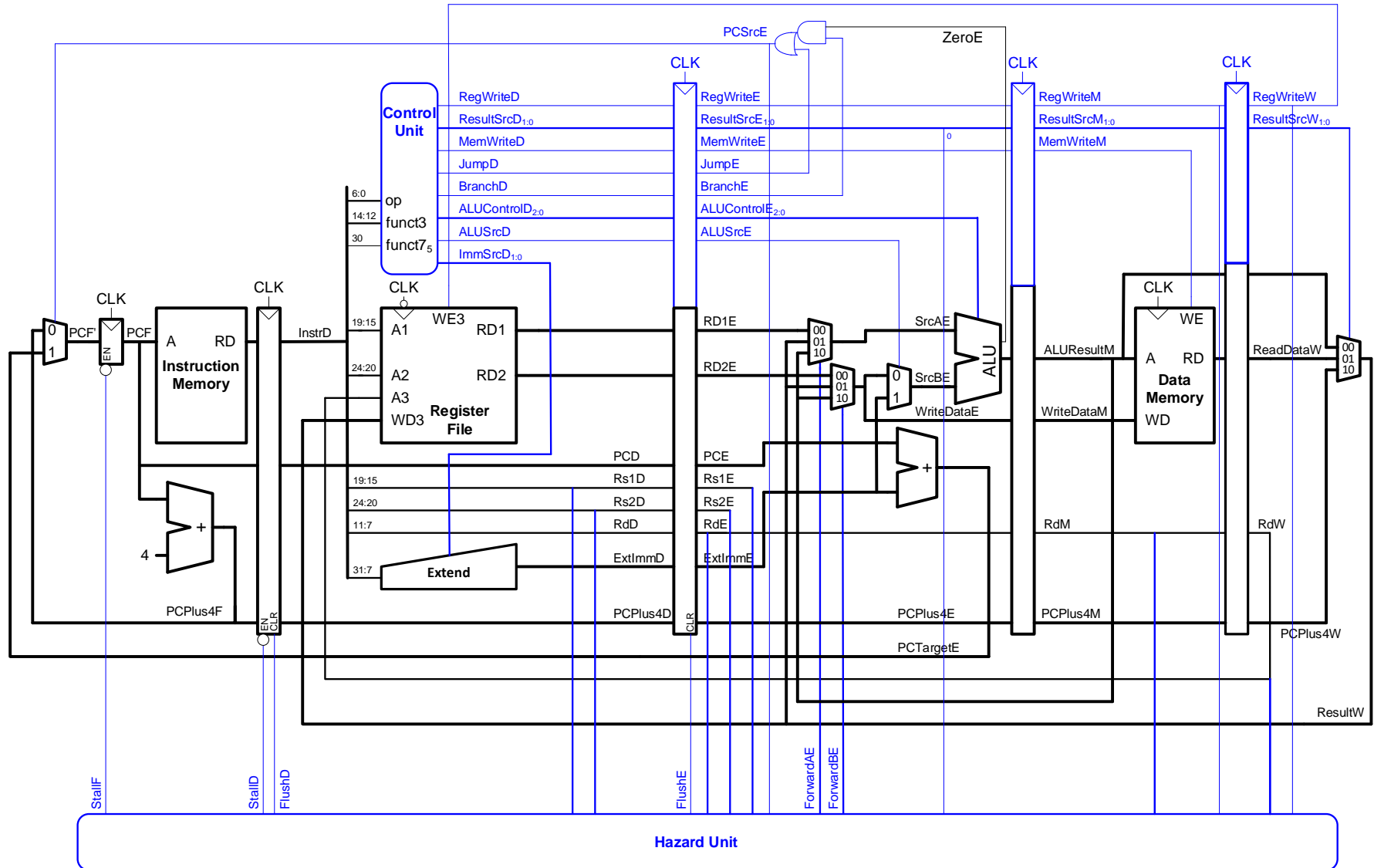
$$FlushD = PCSrcE$$

$$FlushE = lwStall \text{ OR } PCSrcE$$

Control Hazards: Flushing Hardware



RISC-V Pipelined Processor with Hazard Unit



Summary of Hazard Logic

Data hazard logic (shown for SrcA of ALU):

```
if      ((Rs1E == RdM) AND RegWriteM) AND (Rs1E != 0)    // Case 1
        ForwardAE = 10
else if ((Rs1E == RdW) AND RegWriteW) AND (Rs1E != 0)    // Case 2
        ForwardAE = 01
else      ForwardAE = 00                                // Case 3
```

Load word stall logic:

```
lwStall = ((Rs1D == RdE) OR (Rs2D == RdE)) AND ResultSrcE0
StallF = StallD = lwStall
```

Control hazard flush:

```
FlushD = PCSrcE
FlushE = lwStall OR PCSrcE
```


Chapter 7: Microarchitecture

Pipelined Performance

Pipelined Processor Performance Example

- **SPECINT2000 benchmark:**
 - 25% loads
 - 10% stores
 - 13% branches
 - 52% R-type
- **Suppose:**
 - 40% of loads used by next instruction
 - 50% of branches mispredicted
- **What is the average CPI?** (Ideally it's 1, but...)

Pipelined Processor Performance Example

Pipelined processor critical path:

$T_{c_pipelined} = \max \text{ of}$

$$t_{pcq} + t_{mem} + t_{setup}$$

$$2(t_{RRead} + t_{setup})$$

$$t_{pcq} + 4t_{mux} + t_{ALU} + t_{AND-OR} + t_{setup}$$

$$t_{pcq} + t_{mem} + t_{setup}$$

$$2(t_{pcq} + t_{mux} + t_{RFwrite})$$

Fetch

Decode

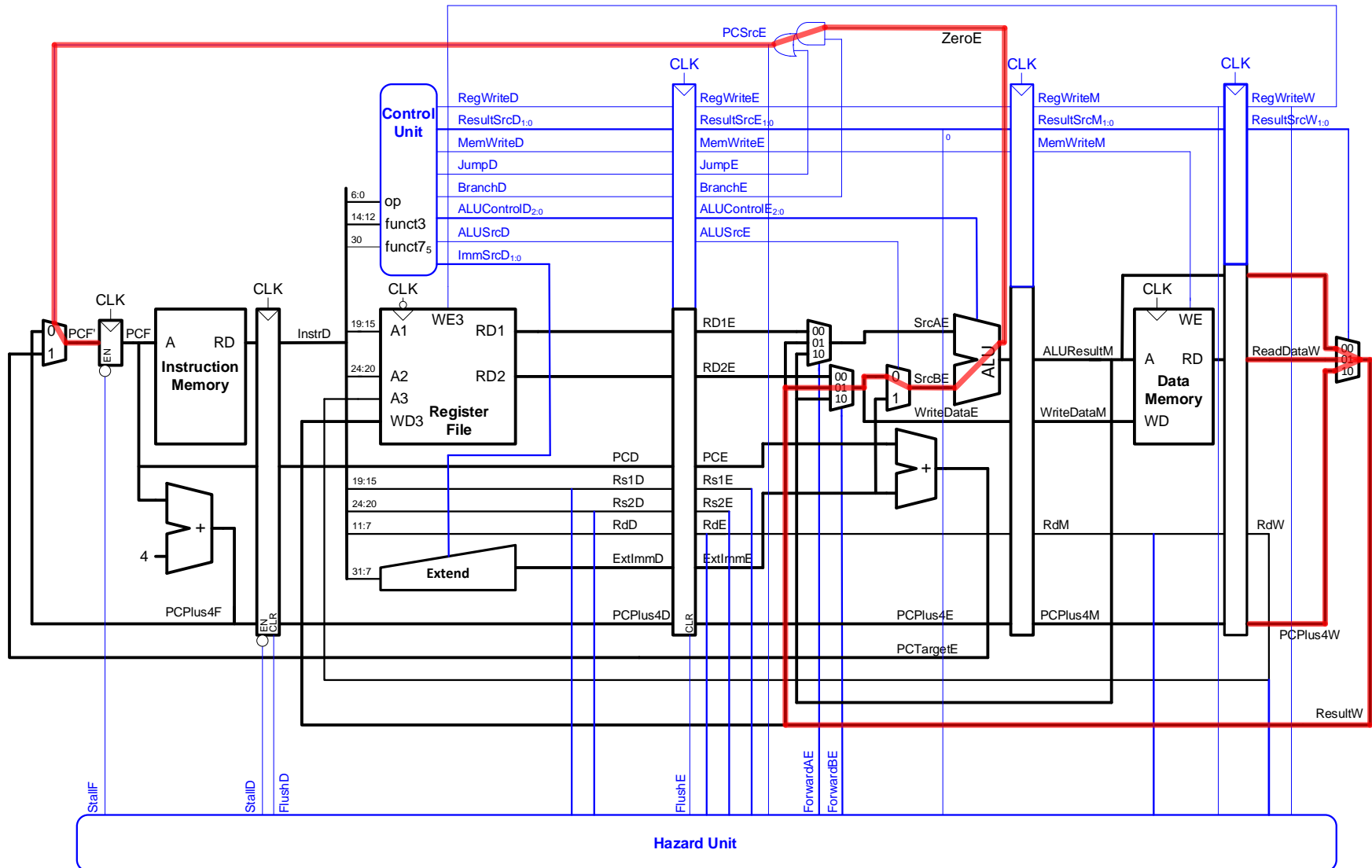
Execute

Memory

Writeback

- Decode and Writeback stages **both use the register file** in each cycle
- So each stage gets half of the cycle time ($T_c/2$) to do their work
- Or, stated a different way, **2x of their work** must fit in a cycle (T_c)

Pipelined Critical Path: Execute Stage



Pipelined Performance Example

Element	Parameter	Delay (ps)
Register clock-to-Q	t_{pcq_PC}	40
Register setup	t_{setup}	50
Multiplexer	t_{mux}	30
AND-OR gate	t_{AND-OR}	20
ALU	t_{ALU}	120
Decoder (Control Unit)	t_{dec}	25
Extend unit	t_{dec}	35
Memory read	t_{mem}	200
Register file read	t_{RFread}	100
Register file setup	$t_{RFsetup}$	60

$$T_{c_pipelined} = t_{pcq} + 4t_{mux} + t_{ALU} + t_{AND-OR} + t_{setup}$$
$$=$$

Pipelined Performance Example

Program with 100 billion instructions

$$\begin{aligned}\text{Execution Time} &= (\# \text{ instructions}) \times \text{CPI} \times T_c \\ &= (100 \times 10^9)(1.23)(350 \times 10^{-12}) \\ &= \mathbf{43 \text{ seconds}}\end{aligned}$$

Processor Performance Comparison

Processor	Execution Time (seconds)	Speedup (single-cycle as baseline)
Single-cycle	75	1
Multicycle	155	0.5
Pipelined	43	1.7

Chapter 7: Microarchitecture

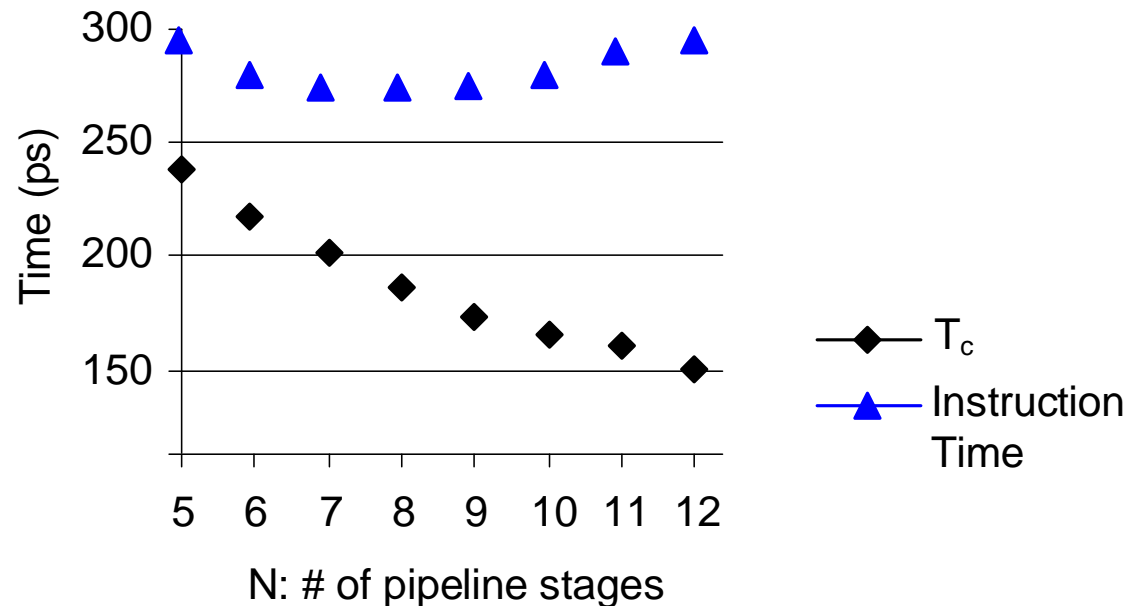
Advanced Microarchitecture

Advanced Microarchitecture

- Deep Pipelining
- Micro-operations
- Branch Prediction
- Superscalar Processors
- Out of Order Processors
- Register Renaming
- SIMD
- Multithreading
- Multiprocessors

Deep Pipelining

- **10-20 stages typical**
- Number of stages limited by:
 - Pipeline hazards
 - Sequencing overhead
 - Power
 - Cost



Micro-operations

- Decompose complex instructions into series of simple instructions called ***micro-operations*** (*micro-ops* or μ -ops)
- **At run-time**, complex instructions are decoded into one or more micro-ops
- Used heavily in **CISC** (complex instruction set computer) architectures (e.g., x86)

Complex Op

```
lw s1, 0(s2), postincr 4
```

Micro-op Sequence

```
lw    s1, 0(s2)
addi  s2, s2, 4
```

Without μ -ops, would need 2nd write port on the register file

Branch Prediction

- **Guess** whether branch will be taken
 - Backward branches are usually taken (loops)
 - Consider history to improve guess
- Good prediction **reduces fraction of branches requiring a flush**

Branch Prediction

- Ideal pipelined processor: $CPI = 1$
- Branch misprediction increases CPI
- **Static branch prediction:**
 - Check direction of branch (forward or backward)
 - If backward, predict taken
 - Else, predict not taken
- **Dynamic branch prediction:**
 - Keep **history** of last several hundred (or thousand) branches in *branch target buffer*, record:
 - Branch destination
 - Whether branch was taken

Dynamic Branch Prediction

- 1-bit branch predictor
- 2-bit branch predictor

Branch Prediction Example

```
addi s1, zero, 0      # s1 = sum
addi s0, zero, 0      # s0 = i
addi t0, zero, 10     # t0 = 10
```

```
For:                  # for (i=0; i<10; i=i+1)
    bge s0, t0, Done
    add s1, s1, s0     # sum = sum + i
    addi s0, s0, 1     # i = i + 1
    j For
```

Done:

1-Bit Branch Predictor

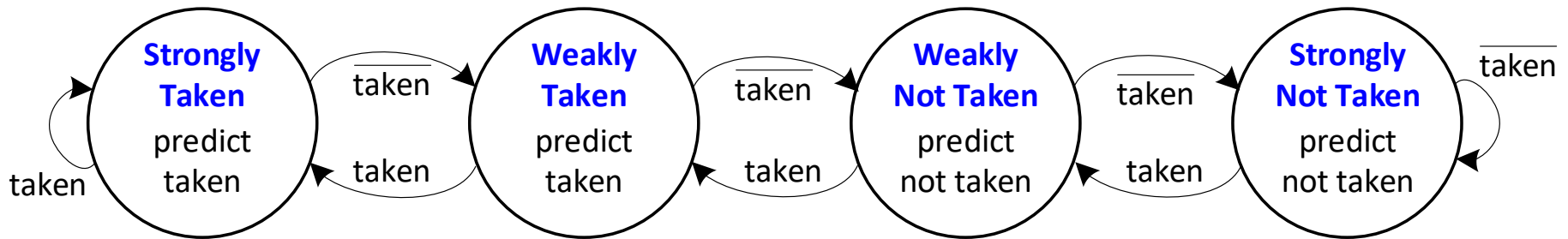
- **Remembers** whether branch was taken the last time and **does the same thing**
- Mispredicts first and last branch of loop

```
addi s1, zero, 0      # s1 = sum
addi s0, zero, 0      # s0 = i
addi t0, zero, 10     # t0 = 10
```

```
For:                  # for (i=0; i<10; i=i+1)
    bge s0, t0, Done
    add s1, s1, s0     # sum = sum + i
    addi s0, s0, 1     # i = i + 1
    j For
```

Done:

2-Bit Branch Predictor



```
addi s1, zero, 0      # s1 = sum
addi s0, zero, 0      # s0 = i
addi t0, zero, 10     # t0 = 10
```

```
For:                      # for (i=0; i<10; i=i+1)
    bge s0, t0, Done
    add s1, s1, s0        # sum = sum + i
    addi s0, s0, 1        # i = i + 1
    j    For
```

Done:

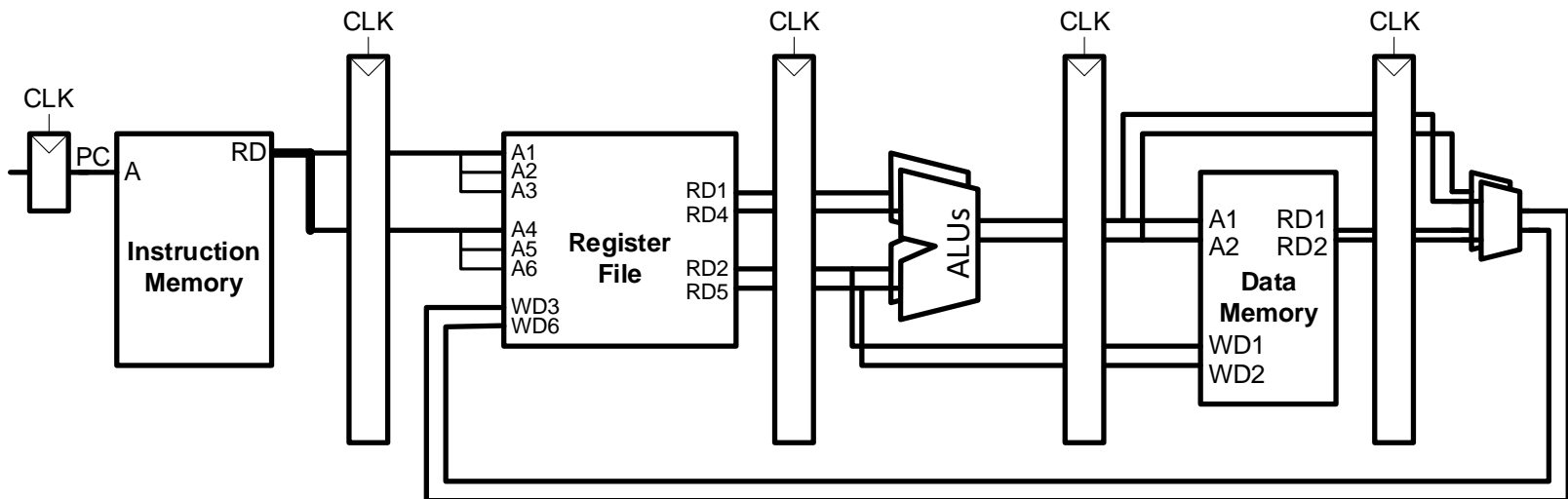
Only mispredicts **last branch** of loop

Chapter 7: Microarchitecture

Superscalar & Out of Order Processors

Superscalar Processors

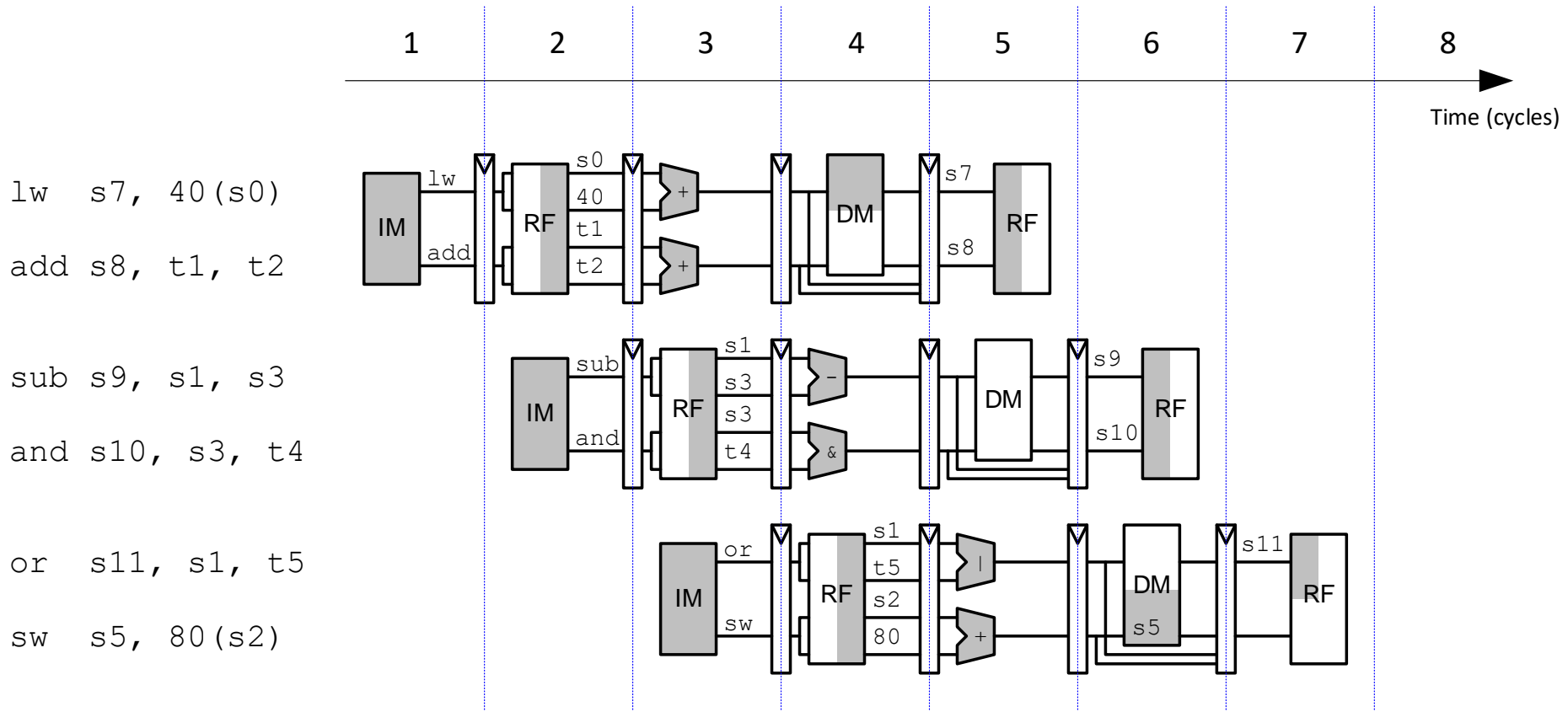
- Multiple copies of datapath execute multiple instructions at once
- Dependencies make it tricky to issue multiple instructions at once



Superscalar Example

Ideal IPC: 2

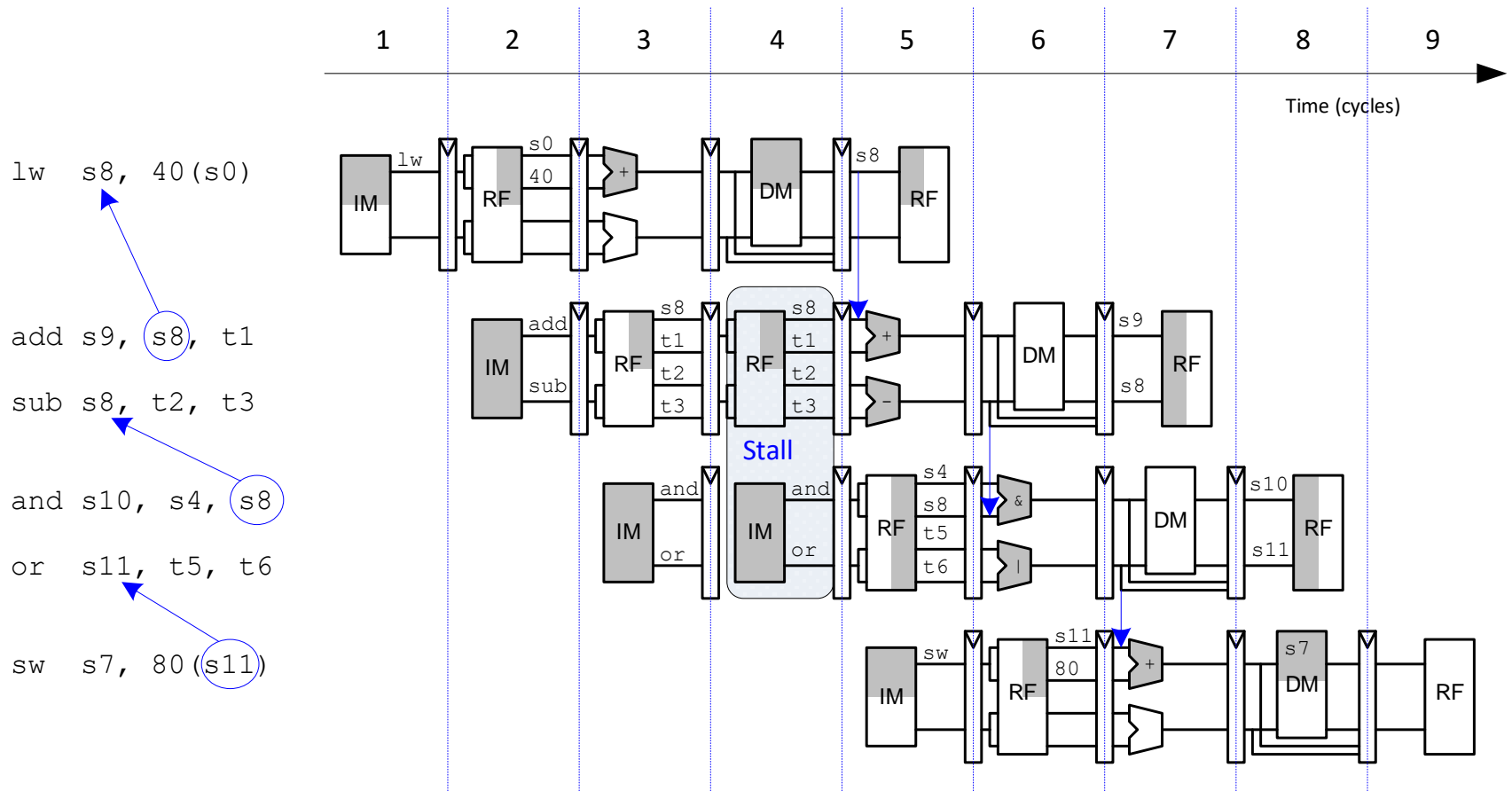
Actual IPC: 2



Superscalar with Dependencies

Ideal IPC: 2

Actual IPC: $6/5 = 1.2$



Out of Order (OOO) Processor

- Looks ahead across multiple instructions
- Issues as many instructions as possible at once
- Issues instructions out of order (as long as no dependencies)
- **Dependencies:**
 - **RAW** (read after write): one instruction writes, later instruction reads a register
 - **WAR** (write after read): one instruction reads, later instruction writes a register
 - **WAW** (write after write): one instruction writes, later instruction writes a register

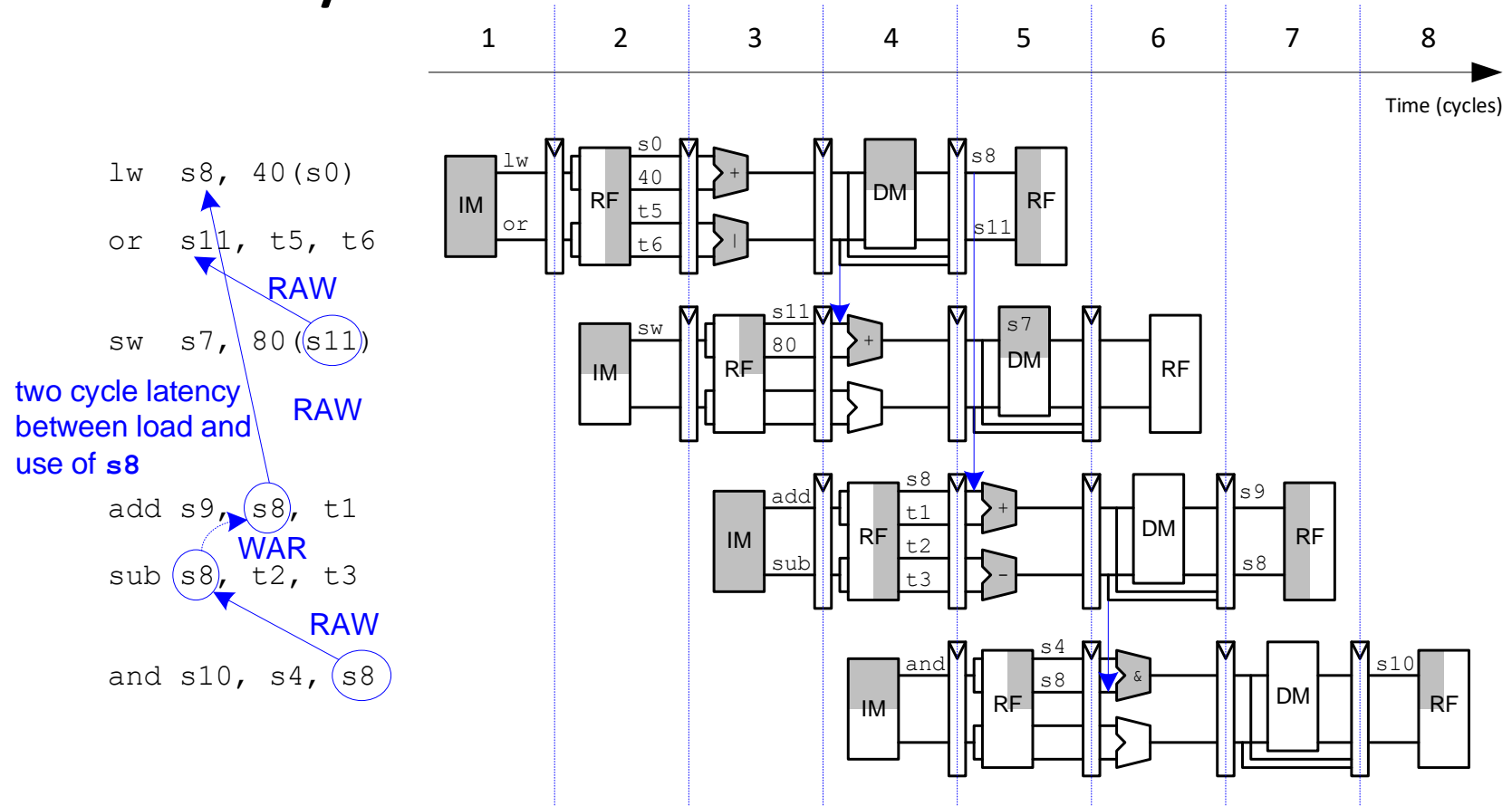
Out of Order (OOO) Processor

- **Instruction level parallelism (ILP):** number of instruction that can be issued simultaneously (average < 3)
- **Scoreboard:** table that keeps track of:
 - Instructions waiting to issue
 - Available functional units
 - Dependencies

Out of Order Processor Example

Ideal IPC: 2

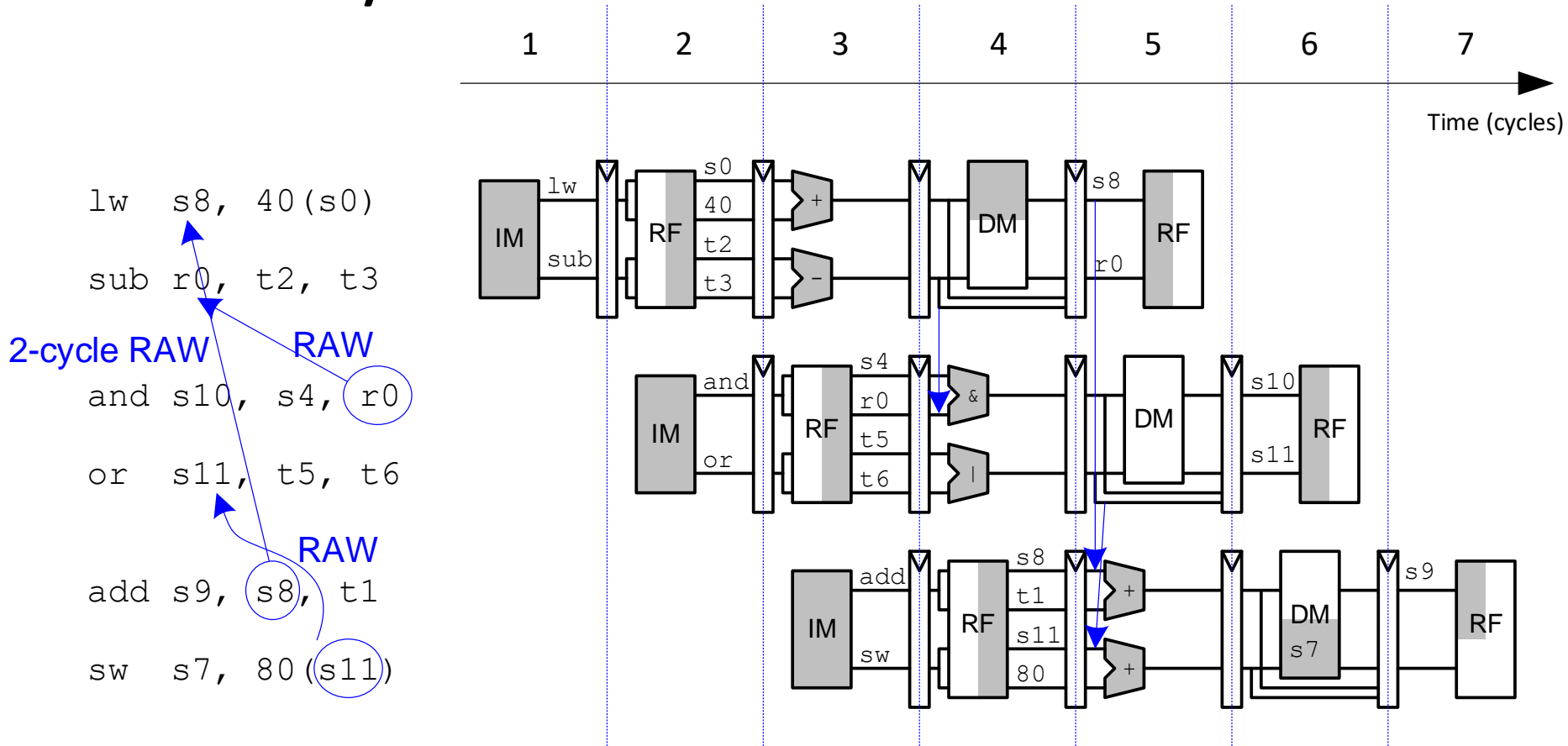
Actual IPC: $6/4 = 1.5$



Register Renaming

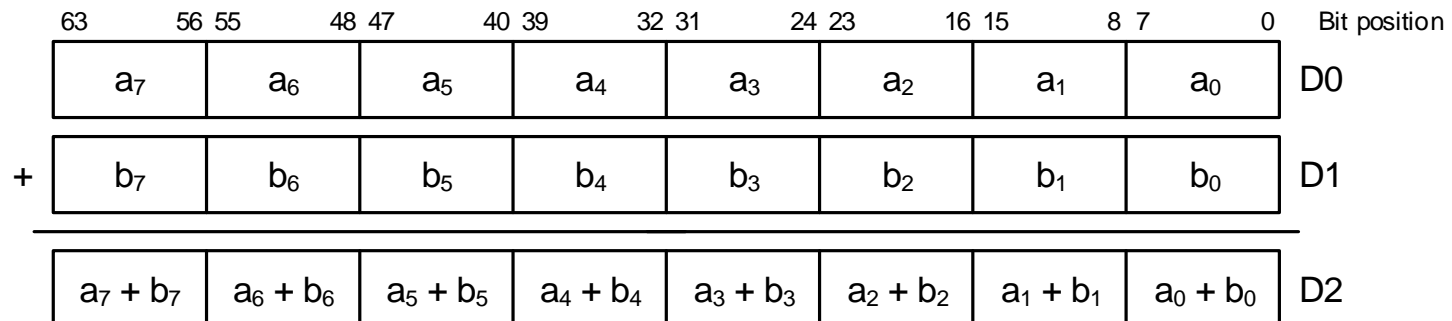
Ideal IPC: 2

Actual IPC: $6/3 = 2$



SIMD

- **Single Instruction Multiple Data (SIMD)**
 - Single instruction **acts on multiple pieces of data at once**
 - Common application: **graphics**
 - Can apply to short arithmetic operations (also called *packed arithmetic*)
- For example, add eight 8-bit elements



Chapter 7: Microarchitecture

Multithreading & Multiprocessors

Advanced Architecture Techniques

- **Multithreading**
 - Wordprocessor: thread for typing, spell checking, printing
- **Multiprocessors**
 - Multiple processors (cores) on a single chip

Threading: Definitions

- **Process:** program running on a computer
 - Multiple processes can run at once: e.g., surfing Web, playing music, writing a paper
- **Thread:** part of a program
 - Each process has multiple threads: e.g., a word processor may have threads for typing, spell checking, printing

Threads in a Conventional Processor

Single-core system:

- One thread runs at once
- When one thread stalls (for example, waiting for memory):
 - Architectural state of that thread stored
 - Architectural state of waiting thread loaded into processor and it runs
 - Called **context switching**
- Appears to user like all threads running simultaneously

Multithreading

- Multiple copies of architectural state
- Multiple threads **active** at once:
 - When one thread stalls, another runs immediately
 - If one thread can't keep all execution units busy, another thread can use them
- Does not increase instruction-level parallelism (ILP) of single thread, but increases throughput

Intel calls this “hyperthreading”

Multiprocessors

- Multiple processors (cores) with a method of communication between them
- Types:
 - **Homogeneous:** multiple cores with shared main memory
 - **Heterogeneous:** separate cores for different tasks (for example, DSP and CPU in cell phone)
 - **Clusters:** each core has own memory system

About these Notes

Digital Design and Computer Architecture Lecture Notes

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