

CMPE 322/327 - Theory of Computation

Week 13: Rice's Theorem & Unrestricted Grammars

Burak Ekici

May 23-27, 2022

Outline

- 1 A Quick Recap
- 2 Rice's Theorem
- 3 Unrestricted Grammars

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

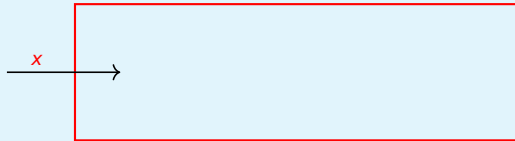
- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x
 - runs K on input $M_x \# x$

Theorem

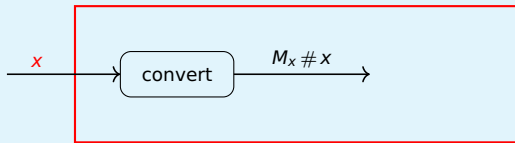
halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

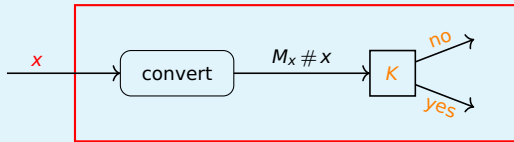
- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x
 - runs K on input $M_x \# x$
 - accepts if K rejects and loops if K accepts



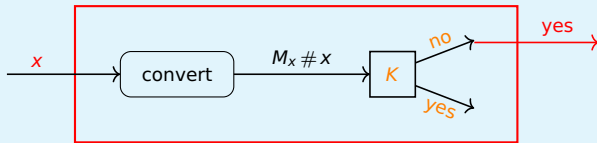
design of TM N



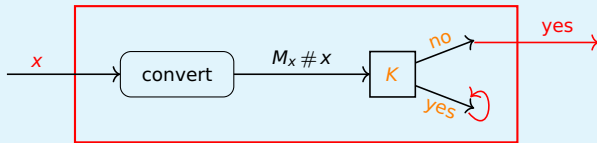
design of TM N



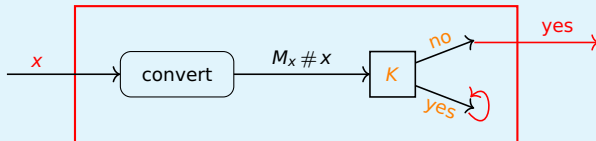
design of TM N



design of TM N



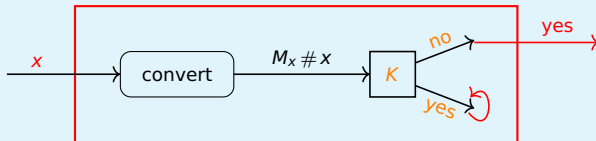
design of TM N

design of TM N

	ϵ	0	1	00	01	10	...
M_ϵ	x	✓	x	x	✓	x	...
M_0	✓	✓	x	x	x	✓	
M_1	x	x	✓	✓	✓	✓	
M_{00}	x	✓	✓	x	✓	x	
M_{01}	x	x	✓	x	✓	x	
M_{10}	✓	x	✓	x	x	x	
\vdots	\vdots						
\vdots	\vdots						

✓: halts

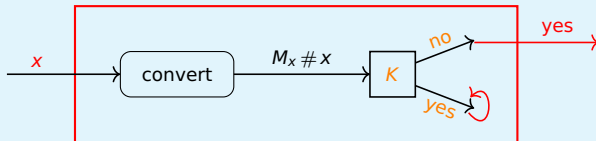
x: loops

design of TM N

	ϵ	0	1	00	01	10	...
M_ϵ	×	✓	×	×	✓	×	...
M_0	✓	✓	×	×	×	✓	
M_1	×	×	✓	✓	✓	✓	
M_{00}	×	✓	✓	×	✓	×	
M_{01}	×	×	✓	×	✓	×	
M_{10}	✓	×	✓	×	×	×	
\vdots	\vdots						
\vdots	\vdots						

✓: accepts

×: rejects

design of TM N

	ϵ	0	1	00	01	10	...
M_ϵ	✓	✓	x	x	✓	x	...
M_0	✓	x	x	x	x	✓	
M_1	x	x	x	✓	✓	✓	
M_{00}	x	✓	✓	✓	✓	x	
M_{01}	x	x	✓	x	x	x	
M_{10}	✓	x	✓	x	x	✓	
⋮	⋮						
⋮	⋮						

✓: halts

x: loops

$$L(N) = \{\epsilon, 00, 10, \dots\}$$

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x
 - runs K on input $M_x \# x$
 - accepts if K rejects and loops if K accepts
- for all inputs x N halts on x \iff K rejects $M_x \# x$

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x
 - runs K on input $M_x \# x$
 - accepts if K rejects and loops if K accepts
- for all inputs x N halts on $x \iff K$ rejects $M_x \# x \iff M_x$ does not halt on x

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x
 - runs K on input $M_x \# x$
 - accepts if K rejects and loops if K accepts
- for all inputs x N halts on $x \iff K$ rejects $M_x \# x \iff M_x$ does not halt on x
- N is different from all M_x

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. ()

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x
 - runs K on input $M_x \# x$
 - accepts if K rejects and loops if K accepts
- for all inputs x N halts on $x \iff K$ rejects $M_x \# x \iff M_x$ does not halt on x
- N is different from all M_x ⚡

Theorem

halting problem (HP) for TMs is undecidable (not recursive)

Proof. (diagonalization)

- suppose $HP = \{M \# x \mid \text{TM } M \text{ halts on input } x\}$ is recursive
- $HP = L(K)$ for some total TM K
- construct TM N that on input x
 - constructs M_x from x
 - runs K on input $M_x \# x$
 - accepts if K rejects and loops if K accepts
- for all inputs x N halts on x \iff K rejects $M_x \# x$ \iff M_x does not halt on x
- N is different from all M_x ⚡

All Together

1 MP is r.e.

(by corollary in w11.pdf on slide #29)

All Together

- 1 MP is r.e. (by corollary in w11.pdf on slide #29)
- 2 HP is **not recursive** (by theorem in w12.pdf on slide #9)

All Together

- 1 MP is r.e. (by corollary in w11.pdf on slide #29)
- 2 HP is not recursive (by theorem in w12.pdf on slide #9)
- 3 $A \leq_m B$ and B is r.e. $\Rightarrow A$ is r.e. (by the first theorem in w12.pdf on slides #14 – 15)

All Together

- ① MP is r.e. (by corollary in w11.pdf on slide #29)
- ② HP is not recursive (by theorem in w12.pdf on slide #9)
- ③ $A \leq_m B$ and B is r.e. \implies A is r.e. (by the first theorem in w12.pdf on slides #14 – 15)
- ④ $A \leq_m B$ and B is recursive \implies A is recursive (by the first theorem in w12.pdf on slide #16)

All Together

- ① MP is r.e. (by corollary in w11.pdf on slide #29)
- ② HP is not recursive (by theorem in w12.pdf on slide #9)
- ③ $A \leq_m B$ and B is r.e. \implies A is r.e. (by the first theorem in w12.pdf on slides #14 – 15)
- ④ $A \leq_m B$ and B is recursive \implies A is recursive (by the first theorem in w12.pdf on slide #16)
- ⑤ $HP \leq_m MP$ (by theorem in w12.pdf on slides #17 – 18)

All Together

- | | | |
|---|---|--|
| ① | MP is r.e. | (by corollary in w11.pdf on slide #29) |
| ② | HP is not recursive | (by theorem in w12.pdf on slide #9) |
| ③ | $A \leq_m B$ and B is r.e. \implies A is r.e. | (by the first theorem in w12.pdf on slides #14 – 15) |
| ④ | $A \leq_m B$ and B is recursive \implies A is recursive | (by the first theorem in w12.pdf on slide #16) |
| ⑤ | $HP \leq_m MP$ | (by theorem in w12.pdf on slides #17 – 18) |

Theorem

MP is not recursive

All Together

- | | | |
|---|---|--|
| ① | MP is r.e. | (by corollary in w11.pdf on slide #29) |
| ② | HP is not recursive | (by theorem in w12.pdf on slide #9) |
| ③ | $A \leq_m B$ and B is r.e. \implies A is r.e. | (by the first theorem in w12.pdf on slides #14 – 15) |
| ④ | $A \leq_m B$ and B is recursive \implies A is recursive | (by the first theorem in w12.pdf on slide #16) |
| ⑤ | $HP \leq_m MP$ | (by theorem in w12.pdf on slides #17 – 18) |

Theorem

MP is not recursive

Proof.

by statements 5, contrapositive of 4 and 2

All Together

- ① MP is r.e. (by corollary in w11.pdf on slide #29)
- ② HP is not recursive (by theorem in w12.pdf on slide #9)
- ③ $A \leq_m B$ and B is r.e. $\implies A$ is r.e. (by the first theorem in w12.pdf on slides #14 – 15)
- ④ $A \leq_m B$ and B is recursive $\implies A$ is recursive (by the first theorem in w12.pdf on slide #16)
- ⑤ $HP \leq_m MP$ (by theorem in w12.pdf on slides #17 – 18)

Theorem

MP is not recursive

Proof.

by statements 5, contrapositive of 4 and 2

Theorem

HP is r.e.

All Together

- ① MP is r.e. (by corollary in w11.pdf on slide #29)
- ② HP is not recursive (by theorem in w12.pdf on slide #9)
- ③ $A \leq_m B$ and B is r.e. $\implies A$ is r.e. (by the first theorem in w12.pdf on slides #14 – 15)
- ④ $A \leq_m B$ and B is recursive $\implies A$ is recursive (by the first theorem in w12.pdf on slide #16)
- ⑤ $HP \leq_m MP$ (by theorem in w12.pdf on slides #17 – 18)

Theorem

MP is not recursive

Proof.

by statements 5, contrapositive of 4 and 2

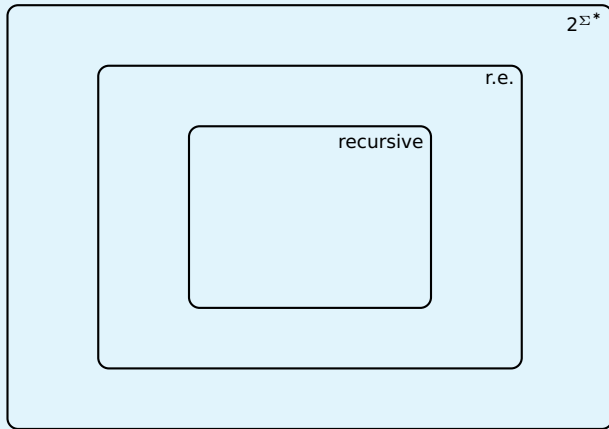
Theorem

HP is r.e.

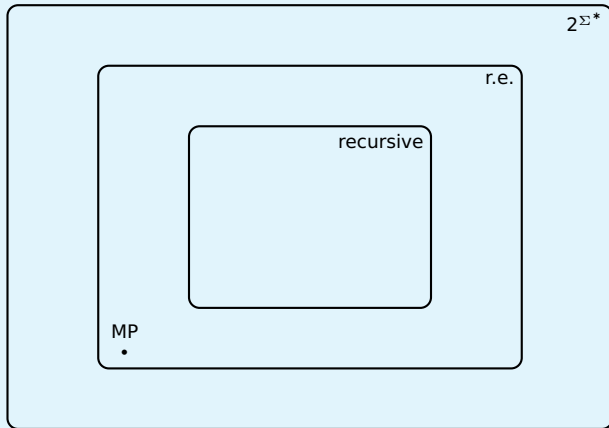
Proof.

by statements 5, 3 and 1

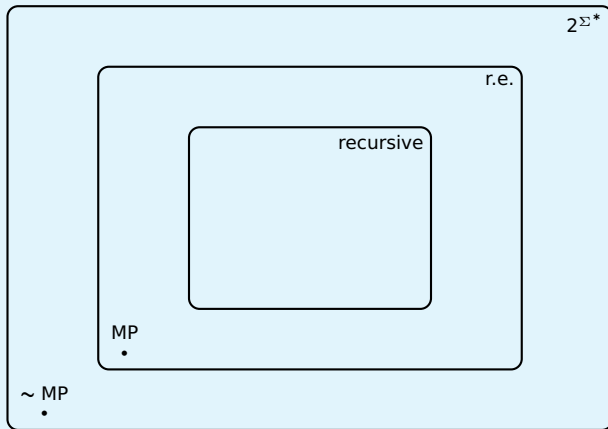
Summary



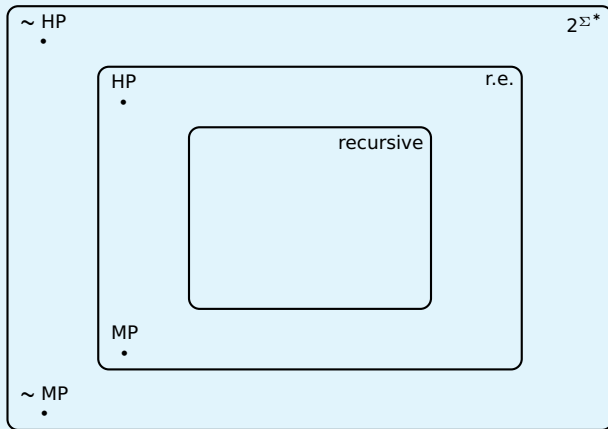
Summary



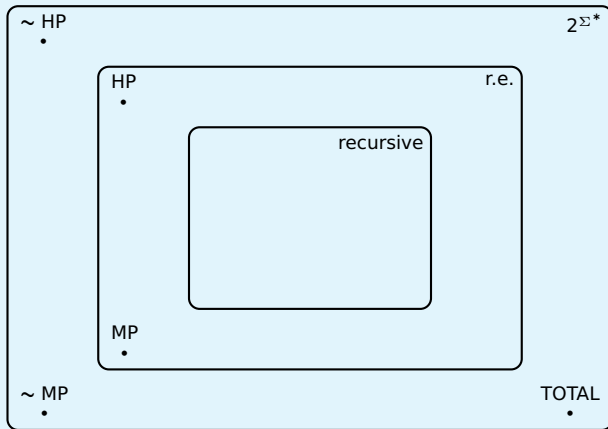
Summary



Summary



Summary



$TOTAL = \{M \mid \text{TM } M \text{ halts on all inputs}\}$

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M

question: is $L(M)$ finite ?

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M

question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M

question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)
- $A := \sim \text{HP}$
- $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M

question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$

- $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN}$ and $\sim \text{HP}$ is not r.e. $\implies \text{FIN}$ is not r.e.

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)
- $A := \sim \text{HP}$
- $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$
- $\sim \text{HP} \leq_m \text{FIN}$ and $\sim \text{HP}$ is not r.e. $\implies \text{FIN}$ is not r.e.

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)
- $A := \sim \text{HP}$
- $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$
- $\sim \text{HP} \leq_m \text{FIN}$ and **True** $\implies \text{FIN}$ is not r.e.

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)
- $A := \sim \text{HP}$
- $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$
- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN is not r.e.}$

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M

question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$

- $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N that
 - ① erases its input

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N that
 - ① erases its input
 - ② writes x on its tape

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N that
 - ① erases its input
 - ② writes x on its tape
 - ③ runs M on input x

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N that
 - ① erases its input
 - ② writes x on its tape
 - ③ runs M on input x
 - ④ accepts if M halts on x

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N that
 - ① erases its input
 - ② writes x on its tape
 - ③ runs M on input x
 - ④ accepts if M halts on x
- $M \# x \in \sim \text{HP} \iff M$ does not halt on x

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N that
 - ① erases its input
 - ② writes x on its tape
 - ③ runs M on input x
 - ④ accepts if M halts on x
- $M \# x \in \sim \text{HP} \iff M \text{ does not halt on } x \iff L(N) = \emptyset$

Theorem

finiteness problem for TMs is not semi-decidable (is not r.e.):

instance: TM M
question: is $L(M)$ finite ?

Proof.

$A \leq_m B$ and B is r.e. $\implies A$ is r.e.

- $A \leq_m B$ and A is not r.e. $\implies B$ is not r.e. (by the second theorem in w12.pdf on slides #14 – 15)

- $A := \sim \text{HP}$
• $B := \text{FIN} = \{M \mid L(M) \text{ is finite}\}$

- $\sim \text{HP} \leq_m \text{FIN} \implies \text{FIN}$ is not r.e.

reduction $\sigma: \sim \text{HP} \rightarrow_m \text{FIN}$

- σ transforms $M \# x$ into TM N that

- ① erases its input
- ② writes x on its tape
- ③ runs M on input x
- ④ accepts if M halts on x

- $M \# x \in \sim \text{HP} \iff M \text{ does not halt on } x \iff L(N) = \emptyset \iff N \in \text{FIN}$

Outline

- 1 A Quick Recap
- 2 Rice's Theorem
- 3 Unrestricted Grammars

Theorem (Rice's Theorem)

every **nontrivial** property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$

Theorem (Rice's Theorem)

every **nontrivial** property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A
- A is accepted by some TM K

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A
- A is accepted by some TM K
- reduction $\sigma: HP \rightarrow_m \{M \mid P(L(M)) = T\}$
 σ transforms $M \# x$ into N such that $L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$

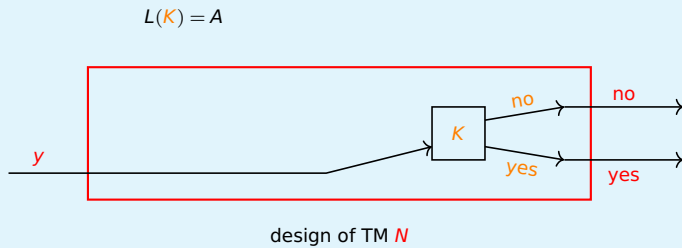
Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

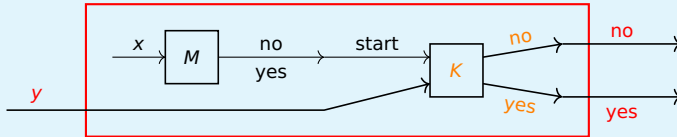
Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A
- A is accepted by some TM K
- reduction $\sigma: HP \rightarrow_m \{M \mid P(L(M)) = T\}$
 σ transforms $M \# x$ into N such that $L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$



$$L(K) = A \qquad L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$$



design of TM N

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A
- A is accepted by some TM K
- reduction $\sigma: HP \rightarrow_m \{M \mid P(L(M)) = T\}$
 σ transforms $M \# x$ into N such that $L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$
- $M \# x \in HP \implies P(L(N)) = P(A) = T \implies N \in \{M \mid P(L(M)) = T\}$

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A
- A is accepted by some TM K
- reduction $\sigma: HP \rightarrow_m \{M \mid P(L(M)) = T\}$

σ transforms $M \# x$ into N such that $L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$

- $M \# x \in HP \implies P(L(N)) = P(A) = T \implies N \in \{M \mid P(L(M)) = T\}$
 $M \# x \notin HP \implies P(L(N)) = P(\emptyset) = \perp \implies N \notin \{M \mid P(L(M)) = T\}$

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A
- A is accepted by some TM K
- reduction $\sigma: HP \rightarrow_m \{M \mid P(L(M)) = T\}$

σ transforms $M \# x$ into N such that $L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$

- $M \# x \in HP \implies P(L(N)) = P(A) = T \implies N \in \{M \mid P(L(M)) = T\}$
 $M \# x \notin HP \implies P(L(N)) = P(\emptyset) = \perp \implies N \notin \{M \mid P(L(M)) = T\}$
- $M \# x \in HP \iff N \in \{M \mid P(L(M)) = T\}$

Theorem (Rice's Theorem)

every nontrivial property P of r.e. sets is undecidable ($\{M \mid P(L(M)) = T\}$ is undecidable)

Proof.

reduction from HP

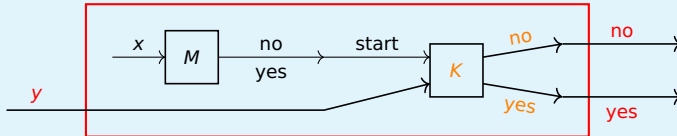
- without loss of generality: $P(\emptyset) = \perp$
- $P(A) = T$ for some r.e. set A
- A is accepted by some TM K
- reduction $\sigma: HP \rightarrow_m \{M \mid P(L(M)) = T\}$

σ transforms $M \# x$ into N such that $L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$

- $M \# x \in HP \implies P(L(N)) = P(A) = T \implies N \in \{M \mid P(L(M)) = T\}$
- $M \# x \notin HP \implies P(L(N)) = P(\emptyset) = \perp \implies N \notin \{M \mid P(L(M)) = T\}$
- $M \# x \in HP \iff N \in \{M \mid P(L(M)) = T\}$



$$L(K) = A \qquad L(N) = \begin{cases} A & \text{if } M \text{ halts on } x \\ \emptyset & \text{otherwise} \end{cases}$$



design of TM N

Corollary

emptiness, finiteness, regularity, context-freeness, recursiveness, ... are undecidable properties of r.e. sets

Some Decision Problems about TMs

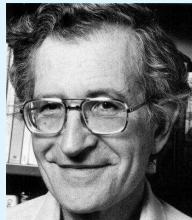
- instance: TM M , state q
question: does M enter state q on some input?
- instance: TM M
question: does M take more than 100 steps on some input?
- instance: TM M
question: does M take more than 100 steps on all inputs?
- instance: TM M
question: does M take less than 100 steps on all inputs?
- instance: TM M
question: does M accept all inputs?
- instance: TM M
question: is there TM N with less states and $L(M) = L(N)$?

Outline

- 1 A Quick Recap
- 2 Rice's Theorem
- 3 Unrestricted Grammars

recursive enumerable

Turing machines
unrestricted grammars



$$\{A \mid A \subseteq \Sigma^*\} = 2^{\Sigma^*}$$

Definitions

- **unrestricted grammar** is quadruple $G = (N, \Sigma, P, S)$

Definitions

- **unrestricted grammar** is quadruple $G = (N, \Sigma, P, S)$ with

① N : finite set of nonterminals

Definitions

- **unrestricted grammar** is quadruple $G = (N, \Sigma, P, S)$ with
 - 1 N : finite set of nonterminals
 - 2 Σ : finite set of terminals, disjoint from N

Definitions

- **unrestricted grammar** is quadruple $G = (N, \Sigma, P, S)$ with
 - 1 N : finite set of nonterminals
 - 2 Σ : finite set of terminals, disjoint from N
 - 3 P : finite set of productions $\alpha \rightarrow \beta$ with $\alpha \in (N \cup \Sigma)^* - \Sigma^*$ and $\beta \in (N \cup \Sigma)^*$

Definitions

- **unrestricted grammar** is quadruple $G = (N, \Sigma, P, S)$ with
 - 1 N : finite set of nonterminals
 - 2 Σ : finite set of terminals, disjoint from N
 - 3 P : finite set of productions $\alpha \rightarrow \beta$ with $\alpha \in (N \cup \Sigma)^* - \Sigma^*$ and $\beta \in (N \cup \Sigma)^*$
 - 4 $S \in N$: start symbol

Definitions

- unrestricted grammar is quadruple $G = (N, \Sigma, P, S)$ with
 - 1 N : finite set of nonterminals
 - 2 Σ : finite set of terminals, disjoint from N
 - 3 P : finite set of productions $\alpha \rightarrow \beta$ with $\alpha \in (N \cup \Sigma)^* - \Sigma^*$ and $\beta \in (N \cup \Sigma)^*$
 - 4 $S \in N$: start symbol
- one step derivation relation $\xrightarrow[G]{1}$ on $(N \cup \Sigma)^*$: $\gamma \alpha \delta \xrightarrow[G]{1} \gamma \beta \delta$ if $\alpha \rightarrow \beta \in P$ and $\gamma, \delta \in (N \cup \Sigma)^*$

Definitions

- unrestricted grammar is quadruple $G = (N, \Sigma, P, S)$ with
 - 1 N : finite set of nonterminals
 - 2 Σ : finite set of terminals, disjoint from N
 - 3 P : finite set of productions $\alpha \rightarrow \beta$ with $\alpha \in (N \cup \Sigma)^* - \Sigma^*$ and $\beta \in (N \cup \Sigma)^*$
 - 4 $S \in N$: start symbol
- one step derivation relation \xrightarrow{G} on $(N \cup \Sigma)^*$: $\gamma \alpha \delta \xrightarrow{G} \gamma \beta \delta$ if $\alpha \rightarrow \beta \in P$ and $\gamma, \delta \in (N \cup \Sigma)^*$
- $\xrightarrow[n]{G} = (\xrightarrow{G})^n \quad \forall n \geq 0 \quad \xrightarrow{*}{G} = \bigcup_{n \geq 0} \xrightarrow[n]{G}$

Definitions

- unrestricted grammar is quadruple $G = (N, \Sigma, P, S)$ with
 - 1 N : finite set of nonterminals
 - 2 Σ : finite set of terminals, disjoint from N
 - 3 P : finite set of productions $\alpha \rightarrow \beta$ with $\alpha \in (N \cup \Sigma)^* - \Sigma^*$ and $\beta \in (N \cup \Sigma)^*$
 - 4 $S \in N$: start symbol
- one step derivation relation \xrightarrow{G} on $(N \cup \Sigma)^*$: $\gamma \alpha \delta \xrightarrow{G} \gamma \beta \delta$ if $\alpha \rightarrow \beta \in P$ and $\gamma, \delta \in (N \cup \Sigma)^*$
- $\xrightarrow{G}^n = (\xrightarrow{G})^n \quad \forall n \geq 0 \quad \xrightarrow{G}^* = \bigcup_{n \geq 0} \xrightarrow{G}^n$

Definitions

- unrestricted grammar is quadruple $G = (N, \Sigma, P, S)$ with
 - 1 N : finite set of nonterminals
 - 2 Σ : finite set of terminals, disjoint from N
 - 3 P : finite set of productions $\alpha \rightarrow \beta$ with $\alpha \in (N \cup \Sigma)^* - \Sigma^*$ and $\beta \in (N \cup \Sigma)^*$
 - 4 $S \in N$: start symbol
- one step derivation relation \xrightarrow{G}_1 on $(N \cup \Sigma)^*$: $\gamma \alpha \delta \xrightarrow{G}_1 \gamma \beta \delta$ if $\alpha \rightarrow \beta \in P$ and $\gamma, \delta \in (N \cup \Sigma)^*$
- $\xrightarrow{G}_n = (\xrightarrow{G}_1)^n \quad \forall n \geq 0 \quad \xrightarrow{G}^* = \bigcup_{n \geq 0} \xrightarrow{G}_n$
- language generated by G : $L(G) = \{x \in \Sigma^* \mid S \xrightarrow{G}^* x\}$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

① $N = \{S, A, B, C, D, E\}$

② $\Sigma = \{a, b\}$

③ P consists of productions

$S \rightarrow ABC$	$DC \rightarrow B a C$	$D a \rightarrow a D$	$E a \rightarrow a E$	$a B \rightarrow B a$	$C \rightarrow \varepsilon$
$AB \rightarrow a A D \mid b A E \mid \varepsilon$	$EC \rightarrow B b C$	$D b \rightarrow b D$	$E b \rightarrow b E$	$b B \rightarrow B b$	

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- ① $N = \{S, A, B, C, D, E\}$
- ② $\Sigma = \{a, b\}$
- ③ P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow BaC & Da \rightarrow aD & Ea \rightarrow aE & aB \rightarrow Ba & C \rightarrow \epsilon \\
 AB \rightarrow aAD \mid bAE \mid \epsilon & EC \rightarrow BbC & Db \rightarrow bD & Eb \rightarrow bE & bB \rightarrow Bb &
 \end{array}$$

$abbabb \in L(G)$:

$$S \xrightarrow[G]{1} ABC$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

① $N = \{S, A, B, C, D, E\}$

② $\Sigma = \{a, b\}$

③ P consists of productions

$S \rightarrow ABC$	$DC \rightarrow B a C$	$D a \rightarrow a D$	$E a \rightarrow a E$	$a B \rightarrow B a$	$C \rightarrow \epsilon$
$AB \rightarrow a A D \mid b A E \mid \epsilon$	$EC \rightarrow B b C$	$D b \rightarrow b D$	$E b \rightarrow b E$	$b B \rightarrow B b$	

$abbabb \in L(G)$:

$$S \xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

① $N = \{S, A, B, C, D, E\}$

② $\Sigma = \{a, b\}$

③ P consists of productions

$S \rightarrow ABC$	$DC \rightarrow BaC$	$Da \rightarrow aD$	$Ea \rightarrow aE$	$aB \rightarrow Ba$	$C \rightarrow \epsilon$
$AB \rightarrow aAD \mid bAE \mid \epsilon$	$EC \rightarrow BbC$	$Db \rightarrow bD$	$Eb \rightarrow bE$	$bB \rightarrow Bb$	

$abbabb \in L(G)$:

$$S \xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC \xrightarrow[G]{1} aABaC$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \epsilon \\
 AB \rightarrow a AD \mid b A E \mid \epsilon & EC \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$S \xrightarrow[G]{1} ABC \xrightarrow[G]{1} a ADC \xrightarrow[G]{1} a AB a C \xrightarrow[G]{1} a b A E a C$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \varepsilon \\
 AB \rightarrow a A D \mid b A E \mid \varepsilon & EC \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$S \xrightarrow{G} ABC \xrightarrow{G} aADC \xrightarrow{G} aABaC \xrightarrow{G} abAEaC \xrightarrow{G} abAaEC$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \varepsilon \\
 AB \rightarrow a A D \mid b A E \mid \varepsilon & E C \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$S \xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC \xrightarrow[G]{1} aABaC \xrightarrow[G]{1} abAEaC \xrightarrow[G]{1} abAaEC \xrightarrow[G]{1} abAaBbC$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \epsilon \\
 AB \rightarrow a A D \mid b A E \mid \epsilon & EC \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$\begin{array}{l}
 S \xrightarrow[G]{1} ABC \xrightarrow[G]{1} a ADC \xrightarrow[G]{1} a A B a C \xrightarrow[G]{1} a b A E a C \xrightarrow[G]{1} a b A a E C \xrightarrow[G]{1} a b A a B b C \\
 \xrightarrow[G]{1} a b A B a b C
 \end{array}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

① $N = \{S, A, B, C, D, E\}$

② $\Sigma = \{a, b\}$

③ P consists of productions

$S \rightarrow ABC$	$DC \rightarrow B a C$	$D a \rightarrow a D$	$E a \rightarrow a E$	$a B \rightarrow B a$	$C \rightarrow \varepsilon$
$AB \rightarrow a A D \mid b A E \mid \varepsilon$	$EC \rightarrow B b C$	$D b \rightarrow b D$	$E b \rightarrow b E$	$b B \rightarrow B b$	

$abbabb \in L(G)$:

$$\begin{aligned}
 S &\xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC \xrightarrow[G]{1} aABaC \xrightarrow[G]{1} abAEaC \xrightarrow[G]{1} abAaEC \xrightarrow[G]{1} abAaBbC \\
 &\xrightarrow[G]{1} abABabC \xrightarrow[G]{1} abbAEabC
 \end{aligned}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \varepsilon \\
 AB \rightarrow a A D \mid b A E \mid \varepsilon & E C \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$\begin{array}{l}
 S \xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC \xrightarrow[G]{1} aABaC \xrightarrow[G]{1} abAEaC \xrightarrow[G]{1} abAaEC \xrightarrow[G]{1} abAaBbC \\
 \xrightarrow[G]{1} abABabC \xrightarrow[G]{1} abbAEabC \xrightarrow[G]{1} abbAaEbC
 \end{array}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \varepsilon \\
 AB \rightarrow a A D \mid b A E \mid \varepsilon & EC \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$\begin{array}{l}
 S \xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC \xrightarrow[G]{1} aABaC \xrightarrow[G]{1} abAEaC \xrightarrow[G]{1} abAaEC \xrightarrow[G]{1} abAaBbC \\
 \xrightarrow[G]{1} abABabC \xrightarrow[G]{1} abbAEabC \xrightarrow[G]{1} abbAaEbC \xrightarrow[G]{1} abbAabEC
 \end{array}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

① $N = \{S, A, B, C, D, E\}$

② $\Sigma = \{a, b\}$

③ P consists of productions

$S \rightarrow ABC$	$DC \rightarrow B a C$	$D a \rightarrow a D$	$E a \rightarrow a E$	$a B \rightarrow B a$	$C \rightarrow \varepsilon$
$AB \rightarrow a A D \mid b A E \mid \varepsilon$	$E C \rightarrow B b C$	$D b \rightarrow b D$	$E b \rightarrow b E$	$b B \rightarrow B b$	

$abbabb \in L(G)$:

$$\begin{aligned}
 S &\xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC \xrightarrow[G]{1} aABaC \xrightarrow[G]{1} abAEaC \xrightarrow[G]{1} abAaEC \xrightarrow[G]{1} abAaBbC \\
 &\xrightarrow[G]{1} abABabC \xrightarrow[G]{1} abbAEabC \xrightarrow[G]{1} abbAaE bC \xrightarrow[G]{1} abbAabEC \xrightarrow[G]{1} abbAabBbC
 \end{aligned}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

① $N = \{S, A, B, C, D, E\}$

② $\Sigma = \{a, b\}$

③ P consists of productions

$S \rightarrow ABC$	$DC \rightarrow B a C$	$D a \rightarrow a D$	$E a \rightarrow a E$	$a B \rightarrow B a$	$C \rightarrow \varepsilon$
$AB \rightarrow a A D \mid b A E \mid \varepsilon$	$EC \rightarrow B b C$	$D b \rightarrow b D$	$E b \rightarrow b E$	$b B \rightarrow B b$	

$abbabb \in L(G)$:

$$\begin{aligned}
 S &\xrightarrow[G]{1} ABC \xrightarrow[G]{1} aADC \xrightarrow[G]{1} aABaC \xrightarrow[G]{1} abAEaC \xrightarrow[G]{1} abAaEC \xrightarrow[G]{1} abAaBbC \\
 &\xrightarrow[G]{1} abABabC \xrightarrow[G]{1} abbAEabC \xrightarrow[G]{1} abbAaE bC \xrightarrow[G]{1} abbAabEC \xrightarrow[G]{1} abbAabBbC \\
 &\xrightarrow[G]{1} abbAaBbbC
 \end{aligned}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \varepsilon \\
 AB \rightarrow a A D \mid b A E \mid \varepsilon & EC \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$\begin{aligned}
 S &\xrightarrow[G]{1} ABC \xrightarrow[G]{1} a ADC \xrightarrow[G]{1} a AB a C \xrightarrow[G]{1} a b A E a C \xrightarrow[G]{1} a b A a E C \xrightarrow[G]{1} a b A a B b C \\
 &\xrightarrow[G]{1} a b A B a b C \xrightarrow[G]{1} a b b A E a b C \xrightarrow[G]{1} a b b A a E b C \xrightarrow[G]{1} a b b A a b E C \xrightarrow[G]{1} a b b A a b B b C \\
 &\xrightarrow[G]{1} a b b A a B b b C \xrightarrow[G]{1} a b b A B a b b C
 \end{aligned}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \epsilon \\
 AB \rightarrow a AD \mid b AE \mid \epsilon & EC \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$\begin{aligned}
 S &\xrightarrow[G]{1} ABC \xrightarrow[G]{1} a ADC \xrightarrow[G]{1} a AB a C \xrightarrow[G]{1} ab AE a C \xrightarrow[G]{1} ab A a EC \xrightarrow[G]{1} ab A a B b C \\
 &\xrightarrow[G]{1} ab AB ab C \xrightarrow[G]{1} abb AE ab C \xrightarrow[G]{1} abb A a E b C \xrightarrow[G]{1} abb A ab EC \xrightarrow[G]{1} abb A ab B b C \\
 &\xrightarrow[G]{1} abb A a B b b C \xrightarrow[G]{1} abb AB abb C \xrightarrow[G]{1} abbabb C
 \end{aligned}$$

Example

$\{xx \mid x \in \{a, b\}^*\} = L(G)$ for unrestricted grammar $G = (N, \Sigma, P, S)$ with

- 1 $N = \{S, A, B, C, D, E\}$
- 2 $\Sigma = \{a, b\}$
- 3 P consists of productions

$$\begin{array}{llllll}
 S \rightarrow ABC & DC \rightarrow B a C & D a \rightarrow a D & E a \rightarrow a E & a B \rightarrow B a & C \rightarrow \epsilon \\
 AB \rightarrow a A D \mid b A E \mid \epsilon & EC \rightarrow B b C & D b \rightarrow b D & E b \rightarrow b E & b B \rightarrow B b &
 \end{array}$$

$abbabb \in L(G)$:

$$\begin{aligned}
 S &\xrightarrow[G]{1} ABC \xrightarrow[G]{1} a ADC \xrightarrow[G]{1} a AB a C \xrightarrow[G]{1} ab AE a C \xrightarrow[G]{1} ab A a EC \xrightarrow[G]{1} ab A a B b C \\
 &\xrightarrow[G]{1} ab AB ab C \xrightarrow[G]{1} abb AE ab C \xrightarrow[G]{1} abb A a E b C \xrightarrow[G]{1} abb A ab EC \xrightarrow[G]{1} abb A ab B b C \\
 &\xrightarrow[G]{1} abb A a B b b C \xrightarrow[G]{1} abb AB abb C \xrightarrow[G]{1} abbabb C \xrightarrow[G]{1} abbabb
 \end{aligned}$$

Theorem

TMs and unrestricted grammars are **equivalent**:

- 1 $\forall \text{ TM } M \exists \text{ unrestricted grammar } G \text{ such that } L(G) = L(M)$
- 2 $\forall \text{ unrestricted grammar } G \exists \text{ TM } M \text{ such that } L(M) = L(G)$

Theorem

TMs and unrestricted grammars are equivalent:

- 1 $\forall \text{ TM } M \exists \text{ unrestricted grammar } G \text{ such that } L(G) = L(M)$
- 2 $\forall \text{ unrestricted grammar } G \exists \text{ TM } M \text{ such that } L(M) = L(G)$

Proof.

- 1 given $\text{TM } M = (Q, \Sigma, \Gamma, \vdash, \sqcup, \delta, s, t, r)$
construct unrestricted grammar $G = (N, \Sigma, P, S)$ with $N := ((\Sigma \cup \{\epsilon\}) \times \Gamma) \cup Q \cup \{S, T, U\}$

Theorem

TMs and unrestricted grammars are equivalent:

- 1 $\forall \text{ TM } M \exists \text{ unrestricted grammar } G \text{ such that } L(G) = L(M)$
- 2 $\forall \text{ unrestricted grammar } G \exists \text{ TM } M \text{ such that } L(M) = L(G)$

Proof.

- 1 given $\text{TM } M = (Q, \Sigma, \Gamma, \vdash, \sqcup, \delta, s, t, r)$
construct unrestricted grammar $G = (N, \Sigma, P, S)$ with $N := ((\Sigma \cup \{\epsilon\}) \times \Gamma) \cup Q \cup \{S, T, U\}$

Basic Idea

two tracks

Theorem

TMs and unrestricted grammars are equivalent:

- 1 $\forall \text{ TM } M \exists \text{ unrestricted grammar } G \text{ such that } L(G) = L(M)$
- 2 $\forall \text{ unrestricted grammar } G \exists \text{ TM } M \text{ such that } L(M) = L(G)$

Proof.

- 1 given $\text{TM } M = (Q, \Sigma, \Gamma, \vdash, \sqcup, \delta, s, t, r)$
construct unrestricted grammar $G = (N, \Sigma, P, S)$ with $N := ((\Sigma \cup \{\epsilon\}) \times \Gamma) \cup Q \cup \{S, T, U\}$

Basic Idea

two tracks

- **top track** for input string of TM

Theorem

TMs and unrestricted grammars are equivalent:

- 1 $\forall \text{ TM } M \exists \text{ unrestricted grammar } G \text{ such that } L(G) = L(M)$
- 2 $\forall \text{ unrestricted grammar } G \exists \text{ TM } M \text{ such that } L(M) = L(G)$

Proof.

- 1 given $\text{TM } M = (Q, \Sigma, \Gamma, \vdash, \sqcup, \delta, s, t, r)$
construct unrestricted grammar $G = (N, \Sigma, P, S)$ with $N := ((\Sigma \cup \{\varepsilon\}) \times \Gamma) \cup Q \cup \{S, T, U\}$

Basic Idea

two tracks

- top track for input string of TM
- **bottom track** for simulating tape of TM

Theorem

TMs and unrestricted grammars are equivalent:

- 1 $\forall \text{ TM } M \exists \text{ unrestricted grammar } G \text{ such that } L(G) = L(M)$
- 2 $\forall \text{ unrestricted grammar } G \exists \text{ TM } M \text{ such that } L(M) = L(G)$

Proof.

- 1 given $\text{TM } M = (Q, \Sigma, \Gamma, \vdash, \sqcup, \delta, s, t, r)$
construct unrestricted grammar $G = (N, \Sigma, P, S)$ with $N := ((\Sigma \cup \{\epsilon\}) \times \Gamma) \cup Q \cup \{S, T, U\}$

Basic Idea

two tracks

- top track for input string of TM
- bottom track for simulating tape of TM

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{array}{|c|} \hline \varepsilon \\ \hline \vdots \\ \hline \end{array} T$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{array}{|c|} \hline \varepsilon \\ \hline \vdots \\ \hline \end{array} T \quad T \rightarrow \begin{array}{|c|} \hline a \\ \hline a \\ \hline \end{array} T \quad \text{for all } a \in \Sigma$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{array}{|c|} \hline \varepsilon \\ \hline \vdots \\ \hline \end{array} T \quad T \rightarrow \begin{array}{|c|} \hline a \\ \hline a \\ \hline \end{array} T \quad \text{for all } a \in \Sigma \quad T \rightarrow U$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{array}{|c|} \hline \varepsilon \\ \hline \vdots \\ \hline \end{array} T \quad T \rightarrow \begin{array}{|c|} \hline a \\ \hline a \\ \hline \end{array} T \quad \text{for all } a \in \Sigma \quad T \rightarrow U \quad U \rightarrow \begin{array}{|c|} \hline \varepsilon \\ \hline \sqcup \\ \hline \end{array} U$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{array}{|c|} \hline \varepsilon \\ \hline \vdots \\ \hline \end{array} T \quad T \rightarrow \begin{array}{|c|} \hline a \\ \hline a \\ \hline \end{array} T \quad \text{for all } a \in \Sigma \quad T \rightarrow U \quad U \rightarrow \begin{array}{|c|} \hline \varepsilon \\ \hline \sqcup \\ \hline \end{array} U \quad U \rightarrow \varepsilon$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{bmatrix} \varepsilon \\ \vdots \end{bmatrix} T \quad T \rightarrow \begin{bmatrix} a \\ a \end{bmatrix} T \quad \text{for all } a \in \Sigma \quad T \rightarrow U \quad U \rightarrow \begin{bmatrix} \varepsilon \\ \sqcup \end{bmatrix} U \quad U \rightarrow \varepsilon$$

$$p \begin{bmatrix} c \\ a \end{bmatrix} \rightarrow \begin{bmatrix} c \\ b \end{bmatrix} q \quad \text{for all } \delta(p, a) = (q, b, R) \text{ and } c \in \Sigma \cup \{\varepsilon\}$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{bmatrix} \varepsilon \\ \vdots \end{bmatrix} T \quad T \rightarrow \begin{bmatrix} a \\ a \end{bmatrix} T \quad \text{for all } a \in \Sigma \quad T \rightarrow U \quad U \rightarrow \begin{bmatrix} \varepsilon \\ \sqcup \end{bmatrix} U \quad U \rightarrow \varepsilon$$

$$p \begin{bmatrix} c \\ a \end{bmatrix} \rightarrow \begin{bmatrix} c \\ b \end{bmatrix} q \quad \text{for all } \delta(p, a) = (q, b, R) \text{ and } c \in \Sigma \cup \{\varepsilon\}$$

$$\begin{bmatrix} e \\ d \end{bmatrix} p \begin{bmatrix} c \\ a \end{bmatrix} \rightarrow q \begin{bmatrix} e \\ d \end{bmatrix} \begin{bmatrix} c \\ b \end{bmatrix} \quad \text{for all } \delta(p, a) = (q, b, L) \text{ and } d \in \Gamma \text{ and } c, e \in \Sigma \cup \{\varepsilon\}$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{bmatrix} \varepsilon \\ \vdots \end{bmatrix} T \quad T \rightarrow \begin{bmatrix} a \\ a \end{bmatrix} T \quad \text{for all } a \in \Sigma \quad T \rightarrow U \quad U \rightarrow \begin{bmatrix} \varepsilon \\ \sqcup \end{bmatrix} U \quad U \rightarrow \varepsilon$$

$$p \begin{bmatrix} c \\ a \end{bmatrix} \rightarrow \begin{bmatrix} c \\ b \end{bmatrix} q \quad \text{for all } \delta(p, a) = (q, b, R) \text{ and } c \in \Sigma \cup \{\varepsilon\}$$

$$\begin{bmatrix} e \\ d \end{bmatrix} p \begin{bmatrix} c \\ a \end{bmatrix} \rightarrow q \begin{bmatrix} e \\ d \end{bmatrix} \begin{bmatrix} c \\ b \end{bmatrix} \quad \text{for all } \delta(p, a) = (q, b, L) \text{ and } d \in \Gamma \text{ and } c, e \in \Sigma \cup \{\varepsilon\}$$

$$\begin{bmatrix} c \\ d \end{bmatrix} t \rightarrow t c t \quad t \begin{bmatrix} c \\ d \end{bmatrix} \rightarrow t c t \quad t \rightarrow \varepsilon \quad \text{for all } d \in \Gamma \text{ and } c \in \Sigma \cup \{\varepsilon\}$$

Proof. (cont'd)

- P consists of

$$S \rightarrow s \begin{bmatrix} \varepsilon \\ \vdots \end{bmatrix} T \quad T \rightarrow \begin{bmatrix} a \\ a \end{bmatrix} T \quad \text{for all } a \in \Sigma \quad T \rightarrow U \quad U \rightarrow \begin{bmatrix} \varepsilon \\ \sqcup \end{bmatrix} U \quad U \rightarrow \varepsilon$$

$$p \begin{bmatrix} c \\ a \end{bmatrix} \rightarrow \begin{bmatrix} c \\ b \end{bmatrix} q \quad \text{for all } \delta(p, a) = (q, b, R) \text{ and } c \in \Sigma \cup \{\varepsilon\}$$

$$\begin{bmatrix} e \\ d \end{bmatrix} p \begin{bmatrix} c \\ a \end{bmatrix} \rightarrow q \begin{bmatrix} e \\ d \end{bmatrix} \begin{bmatrix} c \\ b \end{bmatrix} \quad \text{for all } \delta(p, a) = (q, b, L) \text{ and } d \in \Gamma \text{ and } c, e \in \Sigma \cup \{\varepsilon\}$$

$$\begin{bmatrix} c \\ d \end{bmatrix} t \rightarrow t c t \quad t \begin{bmatrix} c \\ d \end{bmatrix} \rightarrow t c t \quad t \rightarrow \varepsilon \quad \text{for all } d \in \Gamma \text{ and } c \in \Sigma \cup \{\varepsilon\}$$

- $x \in L(M) \iff x \in L(G)$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	$(1, \vdash, R)$	$(2, b, R)$	$(4, a, L)$	$(2, a, L)$
2	$(1, \vdash, R)$	$(3, a, L)$	$(1, a, R)$	$(4, a, L)$

$$\underset{1}{\vdash ab} \xrightarrow{1_M} \underset{1}{\vdash ab} \xrightarrow{1_M} \underset{2}{\vdash bb} \xrightarrow{1_M} \underset{1}{\vdash ba\sqcup} \xrightarrow{1_M} \underset{2}{\vdash baa} \xrightarrow{1_M} \underset{3}{\vdash baa}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	$(1, \vdash, R)$	$(2, b, R)$	$(4, a, L)$	$(2, a, L)$
2	$(1, \vdash, R)$	$(3, a, L)$	$(1, a, R)$	$(4, a, L)$

$$\underset{1}{\vdash} ab \xrightarrow[1]{M} \underset{1}{\vdash} ab \xrightarrow[2]{M} \underset{2}{\vdash} bb \xrightarrow[1]{M} \underset{1}{\vdash} ba \sqcup \xrightarrow[2]{M} \underset{2}{\vdash} baa \xrightarrow[3]{M} \underset{3}{\vdash} baa$$

$$S \rightarrow 1 \begin{array}{|c|} \hline \varepsilon \\ \hline \vdash \\ \hline \end{array} T$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	$(1, \vdash, R)$	$(2, b, R)$	$(4, a, L)$	$(2, a, L)$
2	$(1, \vdash, R)$	$(3, a, L)$	$(1, a, R)$	$(4, a, L)$

$$\underset{1}{\vdash} ab \xrightarrow[1]{M} \underset{1}{\vdash} ab \xrightarrow[1]{M} \underset{2}{\vdash} bb \xrightarrow[2]{M} \underset{1}{\vdash} ba \sqcup \xrightarrow[1]{M} \underset{2}{\vdash} baa \xrightarrow[2]{M} \underset{3}{\vdash} baa$$

$$S \rightarrow 1 \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow 1 \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	$(1, \vdash, R)$	$(2, b, R)$	$(4, a, L)$	$(2, a, L)$
2	$(1, \vdash, R)$	$(3, a, L)$	$(1, a, R)$	$(4, a, L)$

$$\underset{1}{\vdash} ab \xrightarrow[1]{M} \underset{1}{\vdash} ab \xrightarrow[2]{M} \underset{2}{\vdash} bb \xrightarrow[1]{M} \underset{1}{\vdash} ba \sqcup \xrightarrow[2]{M} \underset{2}{\vdash} baa \xrightarrow[3]{M} \underset{3}{\vdash} baa$$

$$S \rightarrow \underset{1}{\boxed{\begin{array}{c} \epsilon \\ \vdash \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{cc} \epsilon & a \\ \vdash & a \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{ccc} \epsilon & a & b \\ \vdash & a & b \end{array}}} T$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	$(1, \vdash, R)$	$(2, b, R)$	$(4, a, L)$	$(2, a, L)$
2	$(1, \vdash, R)$	$(3, a, L)$	$(1, a, R)$	$(4, a, L)$

$$\underset{1}{\vdash} ab \xrightarrow[1]{M} \vdash ab \xrightarrow[1]{M} \vdash bb \xrightarrow[2]{M} \vdash ba \sqcup \xrightarrow[1]{M} \vdash baa \xrightarrow[2]{M} \vdash baa \xrightarrow[3]{M}$$

$$S \rightarrow \underset{1}{\boxed{\begin{array}{c} \epsilon \\ \vdash \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{cc} \epsilon & a \\ \vdash & a \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{ccc} \epsilon & a & b \\ \vdash & a & b \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{ccc} \epsilon & a & b \\ \vdash & a & b \end{array}}} U$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\underset{1}{\vdash} ab \xrightarrow[1]{M} \vdash ab \xrightarrow[1]{M} \vdash bb \xrightarrow[2]{M} \vdash ba \sqcup \xrightarrow[1]{M} \vdash baa \xrightarrow[2]{M} \vdash baa \xrightarrow[3]{M}$$

$$S \rightarrow \underset{1}{\boxed{\begin{array}{c} \epsilon \\ \vdash \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{cc} \epsilon & a \\ \vdash & a \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{ccc} \epsilon & a & b \\ \vdash & a & b \end{array}}} T \rightarrow \underset{1}{\boxed{\begin{array}{ccc} \epsilon & a & b \\ \vdash & a & b \end{array}}} U \rightarrow \underset{1}{\boxed{\begin{array}{cccc} \epsilon & a & b & \epsilon \\ \vdash & a & b & \sqcup \end{array}}} U$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\underset{1}{\vdash} ab \xrightarrow[1]{M} \vdash ab \xrightarrow[1]{M} \vdash bb \xrightarrow[2]{M} \vdash ba \sqcup \xrightarrow[1]{M} \vdash baa \xrightarrow[2]{M} \vdash baa \xrightarrow[3]{M}$$

$$S \rightarrow \underset{1}{\begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array}} T \rightarrow \underset{1}{\begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array}} T \rightarrow \underset{1}{\begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array}} T \rightarrow \underset{1}{\begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array}} U \rightarrow \underset{1}{\begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array}} U \rightarrow \underset{1}{\begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array}}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned}
S &\rightarrow 1 \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow 1 \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow 1 \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow 1 \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow 1 \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow 1 \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} 1
\end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned}
S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \xrightarrow{1} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \xrightarrow{2} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array}
\end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned} S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\ &\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \xrightarrow{1} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \xrightarrow{2} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \xrightarrow{1} \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} \end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned}
S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \xrightarrow{1} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \xrightarrow{2} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \xrightarrow{1} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \xrightarrow{2} \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array}
\end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M} \vdash baa$$

$$\begin{aligned} S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\ &\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned}
S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|} \hline a & b & \epsilon \\ \hline b & a & a \\ \hline \end{array}
\end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	$(1, \vdash, R)$	$(2, b, R)$	$(4, a, L)$	$(2, a, L)$
2	$(1, \vdash, R)$	$(3, a, L)$	$(1, a, R)$	$(4, a, L)$

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned}
S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|c|} \hline a & b & \epsilon \\ \hline b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline a & b & \epsilon \\ \hline a & a & a \\ \hline \end{array}
\end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned} S &\rightarrow 1 \begin{bmatrix} \epsilon \\ \vdash \end{bmatrix} T \rightarrow 1 \begin{bmatrix} \epsilon & a \\ \vdash & a \end{bmatrix} T \rightarrow 1 \begin{bmatrix} \epsilon & a & b \\ \vdash & a & b \end{bmatrix} T \rightarrow 1 \begin{bmatrix} \epsilon & a & b \\ \vdash & a & b \end{bmatrix} U \rightarrow 1 \begin{bmatrix} \epsilon & a & b & \epsilon \\ \vdash & a & b & \sqcup \end{bmatrix} U \rightarrow 1 \begin{bmatrix} \epsilon & a & b & \epsilon \\ \vdash & a & b & \sqcup \end{bmatrix} \\ &\rightarrow \begin{bmatrix} \epsilon & 1 & a & b & \epsilon \\ \vdash & a & b & \sqcup \end{bmatrix} \rightarrow \begin{bmatrix} \epsilon & a & b & \epsilon \\ \vdash & b & b & \sqcup \end{bmatrix} \rightarrow \begin{bmatrix} \epsilon & a & b & 1 & \epsilon \\ \vdash & b & a & \sqcup \end{bmatrix} \rightarrow \begin{bmatrix} \epsilon & a & b & \epsilon \\ \vdash & b & a & a \end{bmatrix} \rightarrow \begin{bmatrix} \epsilon & 3 & a & b & \epsilon \\ \vdash & b & a & a \end{bmatrix} \\ &\rightarrow 33 \begin{bmatrix} a & b & \epsilon \\ b & a & a \end{bmatrix} \rightarrow 33a3 \begin{bmatrix} b & \epsilon \\ a & a \end{bmatrix} \rightarrow 33a3b3 \begin{bmatrix} \epsilon \\ a \end{bmatrix} \end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned}
S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \\
&\rightarrow \begin{array}{|c|c|c|} \hline a & b & \epsilon \\ \hline b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|} \hline a & b & \epsilon \\ \hline b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|} \hline a & b & \epsilon \\ \hline b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|} \hline a & b & \epsilon \\ \hline b & a & a \\ \hline \end{array}
\end{aligned}$$

Example

TM $M = (\{1, 2, 3, 4\}, \{a, b\}, \{a, b, \vdash, \sqcup\}, \vdash, \sqcup, \delta, 1, 3, 4)$ with

δ	\vdash	a	b	\sqcup
1	(1, \vdash , R)	(2, b , R)	(4, a , L)	(2, a , L)
2	(1, \vdash , R)	(3, a , L)	(1, a , R)	(4, a , L)

$$\vdash ab \xrightarrow[1]{1_M} \vdash ab \xrightarrow[1]{1_M} \vdash bb \xrightarrow[2]{1_M} \vdash ba \sqcup \xrightarrow[1]{1_M} \vdash baa \xrightarrow[2]{1_M} \vdash baa \xrightarrow[3]{1_M}$$

$$\begin{aligned} S &\rightarrow \begin{array}{|c|} \hline \epsilon \\ \hline \vdash \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|} \hline \epsilon & a \\ \hline \vdash & a \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} T \rightarrow \begin{array}{|c|c|c|} \hline \epsilon & a & b \\ \hline \vdash & a & b \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} U \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \\ &\rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & a & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & b & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & \sqcup \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline \epsilon & a & b & \epsilon \\ \hline \vdash & b & a & a \\ \hline \end{array} \\ &\rightarrow \begin{array}{|c|c|c|c|} \hline a & b & \epsilon \\ \hline b & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline a & b & \epsilon \\ \hline a & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline a & b & \epsilon \\ \hline a & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline a & b & \epsilon \\ \hline a & a & a \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline a & b & \epsilon \\ \hline a & a & a \\ \hline \end{array} \rightarrow^* ab \end{aligned}$$

Thanks! & Questions?