

## Lesson 12. Dynamic Programming – Review

- Recall from Lessons 5-11:

- A **dynamic program** models situations where decisions are made in a sequential process in order to optimize some objective
- **Stages**  $t = 1, 2, \dots, T$ 
  - stage  $T \leftrightarrow$  end of decision process
- **States**  $n = 0, 1, \dots, N \leftarrow$  possible conditions of the system at each stage
- Two representations: **shortest/longest path** and **recursive**

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Shortest/longest path	Recursive
node $t_n$	$\leftrightarrow$ state $n$ at stage $t$
edge $(t_n, (t+1)_m)$	$\leftrightarrow$ allowable decision $x_t$ in state $n$ at stage $t$ that results in being in state $m$ at stage $t+1$
length of edge $(t_n, (t+1)_m)$	$\leftrightarrow$ contribution of decision $x_t$ in state $n$ at stage $t$ that results in being in state $m$ at stage $t+1$
length of shortest/longest path from node $t_n$ to end node	$\leftrightarrow$ value-to-go function $f_t(n)$
length of edges $(T_n, \text{end})$	$\leftrightarrow$ boundary conditions $f_T(n)$
shortest or longest path	$\leftrightarrow$ recursion is min or max:  $f_t(n) = \min_{x_t \text{ allowable}} \left\{ \begin{pmatrix} \text{contribution of} \\ \text{decision } x_t \end{pmatrix} + f_{t+1} \left( \begin{pmatrix} \text{new state} \\ \text{resulting} \\ \text{from } x_t \end{pmatrix} \right) \right\}$
source node $1_n$	$\leftrightarrow$ desired value-to-go function value $f_1(n)$

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**Example 1.** Simplexville Oil needs to build capacity to refine 1,000 barrels of oil and 2,000 barrels of gasoline per day. Simplexville can build a refinery at 2 locations. The cost of building a refinery is as follows:

Oil capacity per day	Gas capacity per day	Building cost (\$ millions)
0	0	0
1000	0	5
0	1000	7
1000	1000	14

The problem is to determine how much capacity should be built at each location in order to minimize the total building cost. Assume that the capacity requirements must be met exactly.

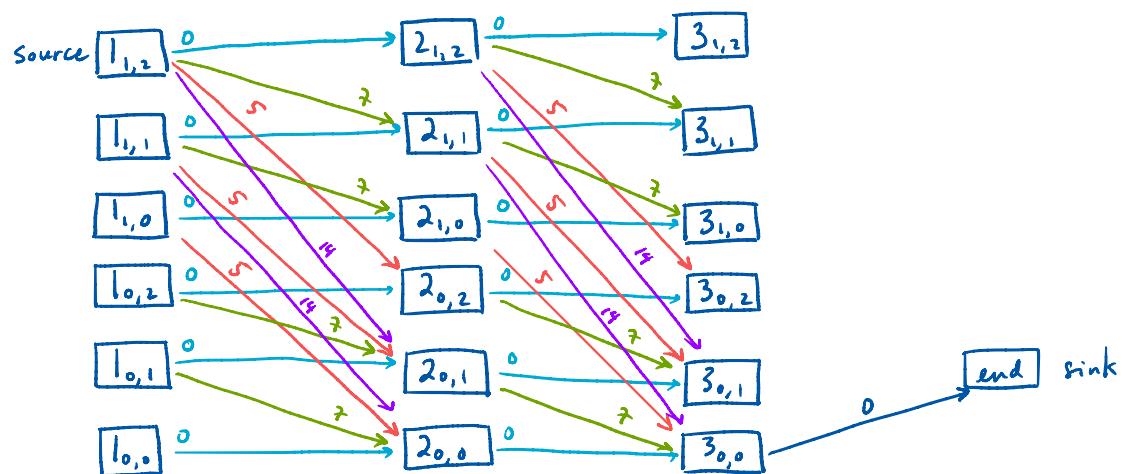
- Formulate this problem as a dynamic program by giving its shortest path representation.
- Formulate this problem as a dynamic program by giving its recursive representation. Solve the dynamic program.

$$\text{Stage } t \leftrightarrow \begin{cases} \text{deciding to build at location } t & t=1,2 \\ \text{end of decision-making process} & t=3 \end{cases}$$

$$\text{State } (n_1, n_2) \leftrightarrow n_1 \text{ oil capacity and } n_2 \text{ gas capacity still needed to be built}$$

$$n_1 = 0, 1 \\ n_2 = 0, 1, 2$$

Find shortest path.



## Recursive representation

- Stage  $t \leftrightarrow \begin{cases} \text{deciding to build at location } t & t=1, 2 \\ \text{end of decision-making process} & t=3 \end{cases}$
- State  $(n_1, n_2) \leftrightarrow n_1$  oil capacity and  $n_2$  gas capacity still needed to be built  $\begin{array}{l} n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{array}$
- Allowable decisions  $x_t$  at stage  $t$  and state  $(n_1, n_2)$ :
  - $x_t = (x_{t1}, x_{t2}) \leftrightarrow \text{build } x_{t1} \text{ oil capacity and } x_{t2} \text{ gas capacity at location } t$
  - $x_t \text{ must satisfy: } \begin{array}{l} x_{t1} \in \{0, 1\} \\ x_{t2} \in \{0, 1\} \\ x_{t1} \leq n_1 \\ x_{t2} \leq n_2 \end{array} \begin{array}{l} \text{can't overbuild capacity.} \\ \text{for } t=1, 2 \\ n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{array}$

- Contribution of  $x_t$  at stage  $t$  and state  $(n_1, n_2)$ :

$$c(x_{t1}, x_{t2}) = \begin{cases} 0 & \text{if } (x_{t1}, x_{t2}) = (0, 0) \\ 5 & \text{if } (x_{t1}, x_{t2}) = (1, 0) \\ 7 & \text{if } (x_{t1}, x_{t2}) = (0, 1) \\ 14 & \text{if } (x_{t1}, x_{t2}) = (1, 1) \end{cases} \begin{array}{l} \text{for } t=1, 2 \\ n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{array}$$

- Value-to-go function

$f_t(n_1, n_2) = \text{minimum total cost to build } n_1 \text{ oil capacity and } n_2 \text{ gas capacity}$   
                   with locations  $t, \dots, 2$  available  $\begin{array}{l} \text{for } t=1, 2, 3 \\ n_1 = 0, 1; n_2 = 0, 1, 2 \end{array}$

- Boundary conditions:  $f_3(n_1, n_2) = \begin{cases} 0 & \text{if } (n_1, n_2) = (0, 0) \\ +\infty & \text{o/w} \end{cases} \begin{array}{l} \text{for } n_1 = 0, 1; n_2 = 0, 1, 2. \end{array}$

- Recursion:

$$f_t(n_1, n_2) = \min_{\substack{x_{t1} \in \{0, 1\} \\ x_{t2} \in \{0, 1\} \\ x_{t1} \leq n_1 \\ x_{t2} \leq n_2}} \left\{ c(x_{t1}, x_{t2}) + f_{t+1}(n_1 - x_{t1}, n_2 - x_{t2}) \right\} \begin{array}{l} \text{for } t=1, 2 \\ n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{array}$$

oil  
 ↓  
 State  $n_1, n_2$   
 ↓  
 gas  
 ↓  
 decision  $x_{t1}, x_{t2}$   
 ↓  
 new state  $n_1 - x_{t1}, n_2 - x_{t2}$

- Desired value-to-go function value:  $f_1(1, 2)$

$$f_t(n_1, n_2) = \min_{\substack{x_{t1} \in \{0, 1\} \\ x_{t2} \in \{0, 1\} \\ x_{t1} \leq n_1 \\ x_{t2} \leq n_2}} \left\{ c(x_{t1}, x_{t2}) + f_{t+1}(n_1 - x_{t1}, n_2 - x_{t2}) \right\}$$

$$c(x_{t1}, x_{t2}) = \begin{cases} 0 & \text{if } (x_{t1}, x_{t2}) = (0, 0) \\ 5 & \text{if } (x_{t1}, x_{t2}) = (1, 0) \\ 7 & \text{if } (x_{t1}, x_{t2}) = (0, 1) \\ 14 & \text{if } (x_{t1}, x_{t2}) = (1, 1) \end{cases}$$

for  $t=1, 2$   
 $n_1 = 0, 1$   
 $n_2 = 0, 1, 2$

Solving backwards

Stage 3:  
(boundary conditions)

$$f_3(n_1, n_2) = \begin{cases} 0 & \text{if } (n_1, n_2) = (0, 0) \\ +\infty & \text{otherwise} \end{cases} \quad \text{for } n_1 = 0, 1 \\ n_2 = 0, 1, 2$$

Stage 2:

$$f_2(1, 2) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 1 \\ x_{22} \leq 2}} \left\{ c(x_{21}, x_{22}) + f_3(1 - x_{21}, 2 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(1, 2), c(1, 0) + f_3(0, 2) \\ 7 + \infty & c(0, 1) + f_3(1, 1), c(1, 1) + f_3(0, 1) \\ 14 + \infty & (0, 1) \quad (1, 1) \end{cases} = +\infty$$

$$f_2(1, 1) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 1 \\ x_{22} \leq 1}} \left\{ c(x_{21}, x_{22}) + f_3(1 - x_{21}, 1 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(1, 1), c(1, 0) + f_3(0, 1) \\ 7 + \infty & c(0, 1) + f_3(1, 0), c(1, 1) + f_3(0, 0) \\ 14 + 0 & (0, 1) \quad (1, 1) \end{cases} = 14$$

$$f_2(1, 0) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 1 \\ x_{22} \leq 0}} \left\{ c(x_{21}, x_{22}) + f_3(1 - x_{21}, 0 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(1, 0), c(1, 0) + f_3(0, 0) \\ 5 + 0 & (0, 0) \quad (1, 0) \end{cases} = 5$$

$$f_2(0, 2) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 0 \\ x_{22} \leq 2}} \left\{ c(x_{21}, x_{22}) + f_3(0 - x_{21}, 2 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(0, 2), c(0, 1) + f_3(0, 1) \\ 7 + \infty & (0, 0) \quad (0, 1) \end{cases} = +\infty$$

$$f_2(0, 1) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 0 \\ x_{22} \leq 1}} \left\{ c(x_{21}, x_{22}) + f_3(0 - x_{21}, 1 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(0, 1), c(0, 1) + f_3(0, 0) \\ 7 + 0 & (0, 0) \quad (0, 1) \end{cases} = 7$$

$$f_2(0, 0) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 0 \\ x_{22} \leq 0}} \left\{ c(x_{21}, x_{22}) + f_3(0 - x_{21}, 0 - x_{22}) \right\} = \min \begin{cases} 0 + 0 & c(0, 0) + f_3(0, 0) \\ 0 & (0, 0) \end{cases} = 0$$

Stage 1:

$$f_1(1, 2) = \min_{\substack{x_{11} \in \{0, 1\} \\ x_{12} \in \{0, 1\} \\ x_{11} \leq 1 \\ x_{12} \leq 2}} \left\{ c(x_{11}, x_{12}) + f_2(1 - x_{11}, 2 - x_{12}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_2(1, 2), c(1, 0) + f_2(0, 2) \\ 5 + \infty & c(0, 1) + f_2(1, 1), c(1, 1) + f_2(0, 1) \\ 14 + 7 & (0, 0) \quad (1, 1) \end{cases} = 21$$

Optimal value:  $f_1(1,2) = 21 \Rightarrow$  Minimum total cost of building  
1000 oil capacity + 2000 gas capacity = \$21 million

Optimal solution:  $(x_{11}, x_{12}) = (1,1) \Rightarrow$  At location 1, build 1000 oil capacity  
1000 gas capacity  
 $(x_{21}, x_{22}) = (0,1) \Rightarrow$  At location 2, build 1000 gas capacity