Lesson 7. Arrival Counting Processes and the Poisson Arrival Process

1 Overview

- An arrival counting process is a stochastic process with one "arrival" system event type
- An "arrival" is broadly defined as any discrete unit that can be counted: for example,
 - o customer arrivals
 - o service requests
 - accidents in a factory

2 The Case of the Reckless Beehunter

Citizens of Beehunter have complained that a busy intersection has recently become more dangerous, and they are demanding that the city council take action to make the intersection safer. There have been 103 accidents at the intersection since record keeping began. The city council agrees to undertake a study of the intersection to determine if the accident rate has actually increased above the 1 per week average that is (unfortunately) considered normal. It hires a traffic engineer from nearby Vincennes, to perform the study.

The traffic engineer recommends that the number of accidents at the intersection be recorded for a 24-week period. If the number of accidents is significantly larger than expected, then she will declare that the intersection has indeed become more dangerous. During the study period, 36 accidents were observed.

- Our approach:
 - Model the time between accidents as random variables
 - \diamond Assume they are independent and time stationary with common cdf F_G (is this reasonable?)
 - $\Leftrightarrow E[G] = 1 \text{ week}$
 - Model the system as a stochastic process (algorithm that generates possible sample paths)
 - Using this model, determine if the probability that 36 accidents occur in a 24 week period is "small"

	System events:
•	State variables:

imulation algorithm – the same general algorithm from before: algorithm Simulation: 1: $n \leftarrow 0$ (initialize system event counter) $C_0 \leftarrow 0$ (initialize event epoch) $C_0 \leftarrow 0$ (execute initial system event) 2: $T_{n+1} \leftarrow \min\{C_1, \dots, C_k\}$ (advance time to next pending system event) $C_1 \leftarrow 0$ (find index of next system event) 3: $S_{n+1} \leftarrow S_n$ (temporarily maintain previous state) $C_1 \leftarrow 0$ (event I no longer pending) 4: $C_1 \leftarrow 0$ (execute system event I) (update event counter) 5: go to line 2 What does S_n equal for any n ?	System event subroutines:	
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	Output process: $Y_t \leftarrow S_n$ for all $t \in [T_n, T_{n+1})$,	or in words,
he stochastic process model defined above is called a renewal arrival-counting process		
C. LATIN, WOLLDINGS AND BUILDING	a renewal process for short	

 $\circ\;$ interarrival times are independent and time stationary

 $\circ \ \ arrivals \ occur \ one-at-a-time$

3 The Poisson arrival process

- Based on historical data and some statistical testing, the traffic engineer from Vincennes has determined that the time between accidents are in fact exponentially distributed with a mean of 1 week
- ⇒ Let's now assume that $G \sim \text{Exponential}(\lambda)$:

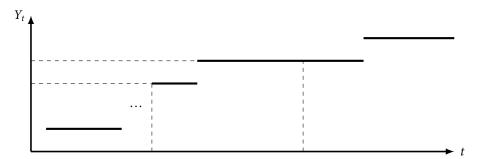
$$F_G(a) = 1 - e^{-\lambda a}$$
 for $a \ge 0$ $E[G] = \frac{1}{\lambda}$ $Var(G) = \frac{1}{\lambda^2}$

- This type of renewal process is known as a Poisson arrival process
- Let G_i be the interarrival time between arrivals i 1 and i
- We can directly write the event epoch T_n as a function of the interarrival times G_1, \ldots, G_n :

• Since G_i are exponential random variables with parameter λ , T_n is an **Erlang random variable with parameter** λ and n phases:

$$F_{T_n}(a) = 1 - \sum_{i=0}^{n-1} \frac{e^{-\lambda a} (\lambda a)^i}{j!} \quad \text{for } a \ge 0 \qquad E[T_n] = \frac{n}{\lambda} \qquad \text{Var}(T_n) = \frac{n}{\lambda^2}$$

• The output process Y_1, Y_2, \ldots and the event epoch process T_1, T_2, \ldots are fundamentally related:



• Therefore, we can get an explicit expression for the cdf of Y_t :

• And we can also get an expression for the pmf of Y_t :

mple 1. In the Beehunter case, at is the probability that the to ober of accidents? (Your calcula	otal number of accid	dents at week 24 is greater	<u>=</u>
operties of the Poisson pro	cess		
Let $\Delta t > 0$ be a time increment			
The forward-recurrence time F	D, is the time betwe	an t and the next arrival	
The independent-increments p dent random variables:	roperty: the numb	er of arrivals in nonoverla	pping time intervals are in
• As a consequence:			
The estation and in anomerate much		of arrivals in a time increm	ent of length Δt only depe
the length of the increment, not			
he length of the increment, not	e arrival rate of the	Poisson process	
he length of the increment, not As a consequence:		-	ribution as the interarriva

• The pmf and cdf may look familiar: Y_t is a **Poisson random variable** with parameter λt

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- Any arrival-counting process in which arrivals occur one-at-a-time and has independent and stationary increments must be a Poisson process
 - If you can justify your system having independent and stationary increments, then you can assume that interarrival times are exponentially distributed
 - o This is a very deep and powerful result

5 Why does the memoryless property hold?

- The memoryless property allows us to ignore when we start observing the Poisson process, since forward-recurrence times and interarrival times are distributed in the same way
- "Memoryless" ←→ how much time has passed doesn't matter
- Why is this true for Poisson processes?
- Let's consider G_n , the interarrival time between the (n-1)th and nth arrival (between T_{n-1} and T_n)
 - Recall that G_n ~ Exponential(λ)
- Pick some t between T_{n-1} and T_n
- We want to show that the forward-recurrence time $R_t \sim \text{Exponential}(\lambda)$
 - Equivalently, we show $F_{R_t}(a) = \Pr\{R_t \le a\} = 1 e^{-\lambda a}$



• Therefore:

$$\Pr\{R_t > a\} = \Pr\{G_n > t - T_{n-1} + a \mid G_n > t - T_{n-1}\}$$

$$= \frac{\Pr\{G_n > t - T_{n-1} + a \text{ and } G_n > t - T_{n-1}\}}{\Pr\{G_n > t - T_{n-1}\}}$$

$$= \frac{\Pr\{G_n > t - T_{n-1} + a\}}{\Pr\{G_n > t - T_{n-1}\}}$$

$$= \frac{e^{-\lambda(t - T_{n-1} + a)}}{e^{-\lambda(t - T_{n-1})}} = e^{-\lambda a}$$

$$\Rightarrow \Pr\{R_t \le a\} = 1 - e^{-\lambda a}$$

- Note: This "proof" is rough and sketchy we actually need to condition on T_{n-1} and Y_t
 - Repeated use of the law of total probability
- The independent-increments and stationary-increments properties follow from the memoryless property and the fundamental relationship between Y_t and T_n (see Nelson pp. 110-111)

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