

Lesson 12. Dynamic Programming – Review

- Recall from Lessons 5-9:

- A **dynamic program** models situations where decisions are made in a sequential process in order to optimize some objective
- **Stages** $t = 1, 2, \dots, T$
 - stage $T \leftrightarrow$ end of decision process
- **States** $n = 0, 1, \dots, N \leftarrow$ possible conditions of the system at each stage
- Two representations: **shortest/longest path** and **recursive**

Shortest/longest path	Recursive
node t_n	\leftrightarrow state n at stage t
edge $(t_n, (t+1)_m)$	\leftrightarrow allowable decision x_t in state n at stage t that results in being in state m at stage $t+1$
length of edge $(t_n, (t+1)_m)$	\leftrightarrow contribution of decision x_t in state n at stage t that results in being in state m at stage $t+1$
length of shortest/longest path from node t_n to end node	\leftrightarrow value-to-go function $f_t(n)$
length of edges (T_n, end)	\leftrightarrow boundary conditions $f_T(n)$
shortest or longest path	\leftrightarrow recursion is min or max: $f_t(n) = \min_{x_t \text{ allowable}} \left\{ \begin{pmatrix} \text{contribution of} \\ \text{decision } x_t \end{pmatrix} + f_{t+1} \left(\begin{pmatrix} \text{new state} \\ \text{resulting} \\ \text{from } x_t \end{pmatrix} \right) \right\}$
source node 1_n	\leftrightarrow desired value-to-go function value $f_1(n)$

Example 1. Simplexville Oil needs to build capacity to refine 1,000 barrels of oil and 2,000 barrels of gasoline per day. Simplexville can build a refinery at 2 locations. The cost of building a refinery is as follows:

Oil capacity per day	Gas capacity per day	Building cost (\$ millions)
0	0	0
1000	0	5
0	1000	7
1000	1000	14

The problem is to determine how much capacity should be built at each location in order to minimize the total building cost. To make things a little simpler, assume that the capacity requirements must be met exactly.

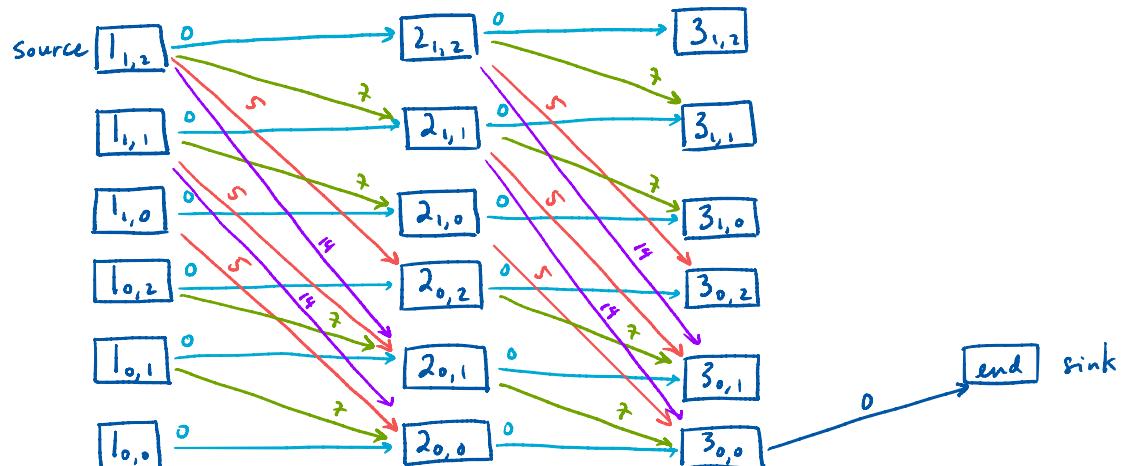
- Formulate this problem as a dynamic program by giving its shortest path representation.
- Formulate this problem as a dynamic program by giving its recursive representation. Solve the dynamic program.

Stage $t \leftrightarrow \begin{cases} \text{deciding to build at location } t & t=1,2 \\ \text{end of decision-making process} & t=3 \end{cases}$

State $(n_1, n_2) \leftrightarrow n_1 \text{ oil capacity and } n_2 \text{ gas capacity still needed to be built}$

$$\begin{aligned} n_1 &= 0, 1 \\ n_2 &= 0, 1, 2 \end{aligned}$$

Find shortest path:



Recursive representation

- Stage $t \leftrightarrow \begin{cases} \text{deciding to build at location } t & t=1, 2 \\ \text{end of decision-making process} & t=3 \end{cases}$
- State $(n_1, n_2) \leftrightarrow n_1$ oil capacity and n_2 gas capacity still needed to be built $\begin{matrix} n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{matrix}$
- Allowable decisions x_t at stage t and state (n_1, n_2) :
 - $x_t = (x_{t1}, x_{t2}) \leftrightarrow \text{build } x_{t1} \text{ oil capacity and } x_{t2} \text{ gas capacity at location } t$
 - $x_t \text{ must satisfy: } \begin{matrix} x_{t1} \in \{0, 1\} \\ x_{t2} \in \{0, 1\} \\ x_{t1} \leq n_1 \\ x_{t2} \leq n_2 \end{matrix} \text{ can't overbuild capacity.}$

- Contribution of x_t at stage t and state (n_1, n_2) :

$$c(x_{t1}, x_{t2}) = \begin{cases} 0 & \text{if } (x_{t1}, x_{t2}) = (0, 0) \\ 5 & \text{if } (x_{t1}, x_{t2}) = (1, 0) \\ 7 & \text{if } (x_{t1}, x_{t2}) = (0, 1) \\ 14 & \text{if } (x_{t1}, x_{t2}) = (1, 1) \end{cases} \quad \begin{matrix} \text{for } t=1, 2 \\ n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{matrix}$$

- Value-to-go function

$f_t(n_1, n_2) = \text{minimum total cost to build } n_1 \text{ oil capacity and } n_2 \text{ gas capacity}$
 with locations $t, \dots, 2$ available $\quad \begin{matrix} \text{for } t=1, 2, 3 \\ n_1 = 0, 1; n_2 = 0, 1, 2 \end{matrix}$

- Boundary conditions: $f_3(n_1, n_2) = \begin{cases} 0 & \text{if } (n_1, n_2) = (0, 0) \\ +\infty & \text{o/w} \end{cases} \quad \begin{matrix} \text{for } n_1 = 0, 1; n_2 = 0, 1, 2. \end{matrix}$

- Recursion:

$$f_t(n_1, n_2) = \min_{\substack{x_{t1} \in \{0, 1\} \\ x_{t2} \in \{0, 1\} \\ x_{t1} \leq n_1 \\ x_{t2} \leq n_2}} \left\{ c(x_{t1}, x_{t2}) + f_{t+1}(n_1 - x_{t1}, n_2 - x_{t2}) \right\} \quad \begin{matrix} \text{for } t=1, 2 \\ n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{matrix}$$

oil
 ↓
 State n_1, n_2
 ↓
 gas
 ↓
 decision x_{t1}, x_{t2}
 ↓
 new state $n_1 - x_{t1}, n_2 - x_{t2}$

- Desired value-to-go function value: $f_1(1, 2)$

$$f_t(n_1, n_2) = \min_{\substack{x_{t1} \in \{0, 1\} \\ x_{t2} \in \{0, 1\} \\ x_{t1} \leq n_1 \\ x_{t2} \leq n_2}} \left\{ c(x_{t1}, x_{t2}) + f_{t+1}(n_1 - x_{t1}, n_2 - x_{t2}) \right\}$$

$$c(x_{t1}, x_{t2}) = \begin{cases} 0 & \text{if } (x_{t1}, x_{t2}) = (0, 0) \\ 5 & \text{if } (x_{t1}, x_{t2}) = (1, 0) \\ 7 & \text{if } (x_{t1}, x_{t2}) = (0, 1) \\ 14 & \text{if } (x_{t1}, x_{t2}) = (1, 1) \end{cases}$$

for $t=1, 2$
 $n_1 = 0, 1$
 $n_2 = 0, 1, 2$

Solving backwards

Stage 3:
(boundary conditions)

$$f_3(n_1, n_2) = \begin{cases} 0 & \text{if } (n_1, n_2) = (0, 0) \\ +\infty & \text{otherwise} \end{cases} \quad \text{for } n_1 = 0, 1 \\ n_2 = 0, 1, 2$$

Stage 2:

$$f_2(1, 2) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 1 \\ x_{22} \leq 2}} \left\{ c(x_{21}, x_{22}) + f_3(1 - x_{21}, 2 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(1, 2), c(1, 0) + f_3(0, 2) \\ 7 + \infty & c(0, 1) + f_3(1, 1), c(1, 1) + f_3(0, 1) \\ 14 + \infty & (0, 1) \quad (1, 1) \end{cases} = +\infty$$

$$f_2(1, 1) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 1 \\ x_{22} \leq 1}} \left\{ c(x_{21}, x_{22}) + f_3(1 - x_{21}, 1 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(1, 1), c(1, 0) + f_3(0, 1) \\ 7 + \infty & c(0, 1) + f_3(1, 0), c(1, 1) + f_3(0, 0) \\ 14 + 0 & (0, 1) \quad (1, 1) \end{cases} = 14$$

$$f_2(1, 0) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 1 \\ x_{22} \leq 0}} \left\{ c(x_{21}, x_{22}) + f_3(1 - x_{21}, 0 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(1, 0), c(1, 0) + f_3(0, 0) \\ 5 + 0 & (0, 0) \quad (1, 0) \end{cases} = 5$$

$$f_2(0, 2) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 0 \\ x_{22} \leq 2}} \left\{ c(x_{21}, x_{22}) + f_3(0 - x_{21}, 2 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(0, 2), c(0, 1) + f_3(0, 1) \\ 7 + \infty & (0, 0) \quad (0, 1) \end{cases} = +\infty$$

$$f_2(0, 1) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 0 \\ x_{22} \leq 1}} \left\{ c(x_{21}, x_{22}) + f_3(0 - x_{21}, 1 - x_{22}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_3(0, 1), c(0, 1) + f_3(0, 0) \\ 7 + 0 & (0, 0) \quad (0, 1) \end{cases} = 7$$

$$f_2(0, 0) = \min_{\substack{x_{21} \in \{0, 1\} \\ x_{22} \in \{0, 1\} \\ x_{21} \leq 0 \\ x_{22} \leq 0}} \left\{ c(x_{21}, x_{22}) + f_3(0 - x_{21}, 0 - x_{22}) \right\} = \min \begin{cases} 0 + 0 & c(0, 0) + f_3(0, 0) \\ 0 & (0, 0) \end{cases} = 0$$

Stage 1:

$$f_1(1, 2) = \min_{\substack{x_{11} \in \{0, 1\} \\ x_{12} \in \{0, 1\} \\ x_{11} \leq 1 \\ x_{12} \leq 2}} \left\{ c(x_{11}, x_{12}) + f_2(1 - x_{11}, 2 - x_{12}) \right\} = \min \begin{cases} 0 + \infty & c(0, 0) + f_2(1, 2), c(1, 0) + f_2(0, 2) \\ 5 + \infty & c(0, 1) + f_2(1, 1), c(1, 1) + f_2(0, 1) \\ 14 + 7 & (0, 0) \quad (1, 1) \end{cases} = 21$$

Optimal value: $f_1(1,2) = 21 \Rightarrow$ Minimum total cost of building
1000 oil capacity + 2000 gas capacity = \$21 million

Optimal solution: $(x_{11}, x_{12}) = (1,1) \Rightarrow$ At location 1, build 1000 oil capacity
1000 gas capacity
 $(x_{21}, x_{22}) = (0,1) \Rightarrow$ At location 2, build 1000 gas capacity