# **Lesson 5. Introduction to Stochastic Processes**

#### 1 Overview

- A **stochastic process** is a sequence of random variables ordered by an index set
- Examples:

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○ {S_n; n = 0, 1, 2, ...} = {S_0, S_1, S_2, ...} with discrete index set {0, 1, 2, ...}

○ {Y_t; t \ge 0} with continuous index set {t \ge 0}
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- The indices *n* and *t* are often referred to as "time"
  - $\{S_n; n = 0, 1, 2, ...\}$  is a **discrete-time process**
  - ∘  $\{Y_t; t \ge 0\}$  is a **continuous-time process**
- The **state space** of a stochastic process is the range (possible values) of its random variables
  - State spaces can be discrete or continuous
     (i.e. the random variables of a stochastic process can be discrete or continuous)
- A stochastic process can be described by the joint distribution of its random variables
- Working with joint distributions can be difficult
- Instead, we can describe a stochastic process via an algorithm for generating its sample paths
- Recall: a **sample path** is a record of the time-dependent behavior of a system
- A stochastic process generates sample paths
  - $\circ$  e.g. a sequence of random variates of  $S_0, S_1, S_2, \dots$
- This lesson: an example

### 2 The Case of the Leaky Bit Bucket

Bit Bucket Computers specializes in installing and maintaining highly reliable computer systems. One of its standard configurations is to install a primary computer, an identical backup computer that is idle until needed, and provide a service contract that guarantees complete repair of a failed computer within 48 hours. If it has not fixed a computer within 48 hours, then it replaces the computer.

Computer systems are rated in terms of their "time to failure" (TTF). The engineers at Bit Bucket Computers have developed a probability distribution for the TTF of the individual computers and a probability distribution for the time required to complete repairs. They would like to have a TTF rating for the entire system. A failure of the system is when both computers are down simultaneously.

# Some additional details from the engineers:

- TTF for a computer:
  - Let  $X_i$  represent the TTF of the *i*th computer in service
  - $\circ X_1, X_2, \dots$  are independent and **time-stationary** (i.e. identically distributed)
    - ⇒ A new computer and a computer that has just been repaired have the same TTF
  - $\circ X_1, X_2, \dots$  have common cdf  $F_X$
  - $F_X$  is the Weibull distribution with parameters  $\alpha = 2$ ,  $\beta = 812$ 
    - ⇒ Expected TTF is 720 hours (30 days) with standard deviation 376 hours (16 days)
    - ♦ Due to their flexible nature, Weibull distributions are commonly used for failure times
- Service time:
  - $\circ$  Let  $R_i$  denote the time required to repair the *i*th computer failure
  - $\circ$   $R_1, R_2, \dots$  are independent and time-stationary
  - $\circ$   $R_1, R_2, \dots$  have common cdf  $F_R$
  - Based on service records,  $F_R$  is the uniform distribution on [4, 48]
- $X_1, X_2, \ldots$  and  $R_1, R_2, \ldots$  are independent
  - ⇒ Repair time of a computer is not affected by its TTF or the number of times it has been repaired
- System logic:
  - o After a system is installed, the primary computer is started
  - When it fails, the backup computer is immediately started and a service call is made to Bit Bucket
  - If the primary computer is repaired before the backup computer fails, then the primary computer becomes the backup computer, and the former backup computer remains the primary computer
  - o If at any time neither computer is available, the entire system fails
  - o Only one computer can be repaired at a time, and are repaired first-come-first-served

### 3 Simulating the Leaky Bit Bucket

- We're interested in D, the time the entire system fails
- *D* is a random variable
  - *D* is a (complex) function of random variables  $X_1, X_2, \ldots$  and  $R_1, R_2, \ldots$
- Let's generate values of  $X_1, X_2, \ldots$  and  $R_1, R_2, \ldots$  and use these to simulate values of D
- We can describe this simulation algorithmically as follows
- System events of interest

• <i>e</i> <sub>1</sub> =	
o <i>e</i> <sub>2</sub> =	

State of	the system: the critical variable that characterizes system status
State spa	.ce:
The <b>cloc</b>	<b>k time</b> $C_i$ of system event $e_i$ is the time the next system event of type $e_i$ occurs
o Wł	en no type $e_i$ event is pending, $C_i \leftarrow \infty$
The <i>n</i> th	<b>event epoch</b> $T_n$ is the time at which the $n$ th system event occurs
At $T_{n+1}$ ,	the time of the $(n + 1)$ st event, two things can happen:
	e system state can change e clocks can be reset
How exa	ctly? We need to describe <b>subroutines</b> for the system events
Let rand	om() be a function that generates variates for Uniform[0,1]
Subrouti	ne for system event $e_1$ :
Subrouti	ne for system event $e_2$ :

•	Let's also create an	"initial"	system event $e_0$	epresentin	g the inst	allation of	the com	puter s	ystem:
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$$e_0()$$
:  
1:  $S_0 \leftarrow 0$  (initially no computers down)  
2:  $C_1 \leftarrow F_X^{-1}(\mathsf{random}())$  (set clock for first computer TTF)  
3:  $C_2 \leftarrow +\infty$  (no pending repair)

# • Putting this all together:

algorithm BitBucketSimulation:

1: 
$$n \leftarrow 0$$
 (initialize system event counter)  
2:  $T_0 \leftarrow 0$  (initialize event epoch)  
3:  $e_0()$  (execute initial system event)  
4:  $T_{n+1} \leftarrow \min\{C_1, C_2\}$  (advance time to next pending system event)  
5:  $I \leftarrow \arg\min\{C_1, C_2\}$  (find index of next system event)  
6:  $C_I \leftarrow \infty$  (event  $I$  no longer pending)  
7:  $e_I()$  (execute system event  $I$ )  
8:  $n \leftarrow n + 1$  (update event counter)  
9: go to line 4

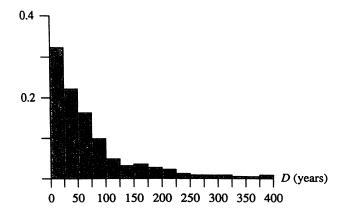
**Example 1.** Suppose the first three values generated by  $F_X^{-1}(random())$  are 612, 36, and 975. In addition, suppose the first two values generated by  $F_R^{-1}(random())$  are 39 and 9. Generate the sample path using the algorithm BitBucketSimulation.

Event counter	System event	Time	State	Failure clock	Repair clock $C_2$
n	I	$T_n$	$S_n$	$C_1$	$C_2$
0					
1					
2					
3					
4					

Let's combine these:	
	$Y_t$ = number of down computers at time $t$ for $t \ge 0$
or equivalently,	
The <b>time average</b> of $Y_t$ u	p to the $n$ th event epoch is
event epoch?	ulated sample path from Example 1, graph $Y_t$ . What is the time average of $Y_t$ up to t
$Y_t$ $2$ $1$	
	<b>→</b> t
Recall: we're interested $L$	D, the time of total system failure, which is:
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• Recall:

- To get information about the distribution of D, we run this simulation many times, say m = 500:
  - 1: **for** r = 1 to m **do**
  - 2: algorithm BitBucketSimulation
  - 3: end for
- Sample results:
  - ∘ Average of generated values of *D*: 551606 hours  $\approx$  63 years
  - $\circ$  Histogram of generated values of *D*:



- $\circ$  2% of the generated values of *D* are less than 2 years
- Is this acceptable or unacceptable?