



$S_i = F \checkmark$

maybe_ptr_ki := $cpk_i.load()$

maybe_ptr_ki ≠ $\phi \checkmark$

reload := $cpk_i.load()$

reload = maybe_ptr_ki ✓

$ptr_{ki} = maybe_ptr_{ki}$

$ptr_{ki} \neq \phi \checkmark$

* $ptr_{ki} = k_x \checkmark$ $\textcircled{1}$ ptr_{ki} could have been swapped here for ptr_{ki} but that's fine

$ptr_{ki} \neq \phi$ $\textcircled{2}$ ptr_{ki} swapped with $ptr_{kp} \rightarrow k' \checkmark$: $ptr_{ki} \neq \phi$ $\textcircled{3}$ swap results in this branch

cas(F ⇌ X) // no deletes of k_x

are possible now

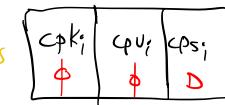
$ptr_{ki} \neq \phi$

retire($ptr_{ki} \cap ptr_{Vi}$)



retire[] is logical delete. retire(ptr_{ka}) means ptr_{ka} is not being held by any cpk_p . When ptr_{ka} is no longer hazardous or being protected, it will be deleted. Another thread can still deference ptr_{ka} to get k_A but $ptr_{ka} \neq \text{any } cpk_i \Rightarrow k_A \neq dB$ so returning k_A would be a correctness issue.

proto()



last delete results in this branch

cas[F ⇌ X] // no deletes of k_x are possible now

$ptr_{ki} \rightarrow \phi \checkmark$

$ptr_{ki} \neq cpk_i.load[] \checkmark$

$cps_i := F$



corrupted state!