

duplicate key problem :

Hot keys are 98% of traffic. 1% cold keys are 2% traffic. Hot key op distribution: 95% **set** 5% **del** cold key op distribution: 20% **set** 10% **get** 70% **del**
agressive deletions on cold keys prevent needing to resize. lets say you have k_A k_B that hash to slot 10. lets say slot 10 is deleted. lets say k_A wins the cas request \rightarrow slot 10: p_{k_A} p_{v_A} F. thread trying to insert p_{k_B} p_{v_B} spins. a third thread comes along and deletes k_A . \rightarrow slot 10: ϕ ϕ D. the thread trying to insert p_{k_B} see's p_{k_A} was deleted during spin. it breaks the while loop. Then it probes and is placed in slot 11. \rightarrow slot 11: p_{k_B} p_{v_B} F. Then a fourth thread comes along trying to insert k_B which hashes to slot 10. Then the fourth thread claims the deleted slot 10 and inserts the key \rightarrow slot 10: p_{k_B} p_{v_B} F ~~slot 11: p_{k_B} p_{v_B} F~~