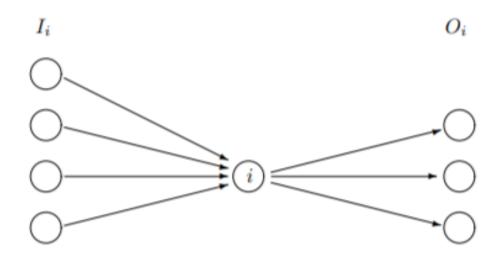
# Page Ranking for a Web Search Engine

• the number of links to and from a page give information about the importance of a page.



$$r_i = \sum_{j \in I_i} \frac{r_j}{N_j}.$$

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$$r_i^{(k+1)} = \sum_{j \in I_i} \frac{r_j^{(k)}}{N_j}, \quad k = 0, 1, \dots$$

page 1

page 2

page 4

page 5

page 3

page 6

page 5

page 3

page 4

page 6

page 2

page 1

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page 5

page 4

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#### page 1

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page 4 page 5

#### page 3

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#### page 5

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#### page 2

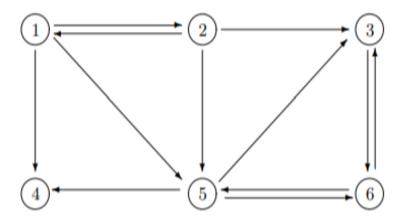
page 1

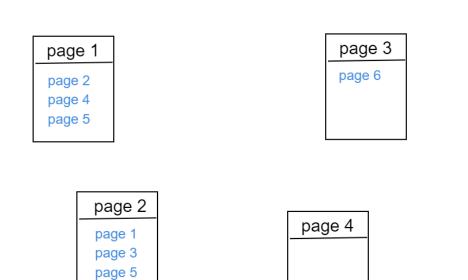
page 3 page 5 page 4

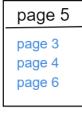
#### page 6

page 3

page 5



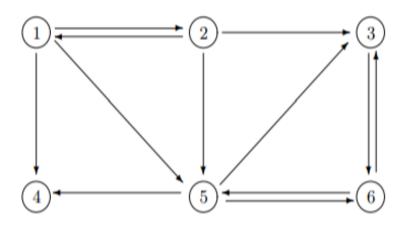




page 6

page 3

page 5



Corresponding Transition Probibility Matrix:

$$Q = \begin{pmatrix} 0 & \frac{1}{3} & 0 & 0 & 0 & 0 \\ \frac{1}{3} & 0 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{3} & 0 & 0 & \frac{1}{3} & \frac{1}{2} \\ \frac{1}{3} & 0 & 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & 0 & \frac{1}{2} \\ 0 & 0 & 1 & 0 & \frac{1}{3} & 0 \end{pmatrix}$$

$$Q = \begin{pmatrix} 0 & \frac{1}{3} & 0 & 0 & 0 & 0 \\ \frac{1}{3} & 0 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{3} & 0 & 0 & \frac{1}{3} & \frac{1}{2} \\ \frac{1}{3} & 0 & 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & 0 & \frac{1}{2} \\ 0 & 0 & 1 & 0 & \frac{1}{3} & 0 \end{pmatrix}.$$

$$r_i = \sum_{j \in I_i} \frac{r_j}{N_j}.$$

$$Q = \begin{pmatrix} 0 & \frac{1}{3} & 0 & 0 & 0 & 0 \\ \frac{1}{3} & 0 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{3} & 0 & 0 & \frac{1}{3} & \frac{1}{2} \\ \frac{1}{3} & 0 & 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & 0 & \frac{1}{2} \\ 0 & 0 & 1 & 0 & \frac{1}{3} & 0 \end{pmatrix}.$$

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is equivalent to the scalar product of row i and the vector r, which holds the ranks of all pages

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$$r_i = \sum_{j \in I_i} \frac{r_j}{N_j}.$$

is equivalent to the scalar product of row i and the vector r, which holds the ranks of all pages :

$$\lambda r = Qr, \qquad \lambda = 1,$$

i.e., r is an eigenvector of Q with eigenvalue  $\lambda = 1$ 

$$r_i^{(k+1)} = \sum_{j \in I_i} \frac{r_j^{(k)}}{N_j}, \quad k = 0, 1, \dots$$

$$r^{(k+1)} = Qr^{(k)}, \qquad k = 0, 1, \dots,$$

# Equivalent definition

$$r^{(k+1)} = Qr^{(k)}, \qquad k = 0, 1, \dots,$$

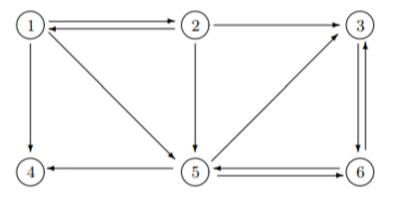
# Equivalent definition

$$r^{(k+1)} = Qr^{(k)}, \qquad k = 0, 1, \dots,$$

 If i were to randomly click around on this graph where would i end up spending most of my time?

# Trap

If i were to randomly click around on this graph where would i end up spending most of my time?



$$Q = \begin{pmatrix} 0 & \frac{1}{3} & 0 & 0 & 0 & 0 \\ \frac{1}{3} & 0 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{3} & 0 & 0 & \frac{1}{3} & \frac{1}{2} \\ \frac{1}{3} & 0 & 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & 0 & \frac{1}{2} \\ 0 & 0 & 1 & 0 & \frac{1}{3} & 0 \end{pmatrix}.$$

A square matrix A is called reducible if there is a permutation matrix P such that:

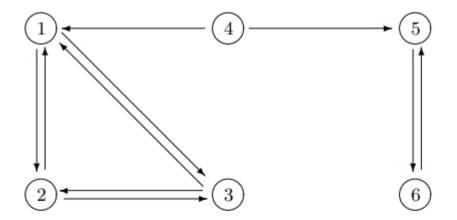
$$PAP^T = \begin{pmatrix} X & Y \\ 0 & Z \end{pmatrix}$$

where X and Z are both square. Otherwise the matrix is called irreducible

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Where is the trap of this graph?

# Eliminate traps:

$$Q' = (1-\alpha)Q + \frac{\alpha}{N} \; 1_{n \times n}$$

 $\frac{\alpha}{N}$ : random jump probability

# Q'is irreducible .why?

$$Q' = (1 - \alpha)Q + \frac{\alpha}{N} \, 1_{n \times n}$$

 $\frac{\alpha}{N}$ : random jump probability

We know the vector r, which holds the ranks of all pages r is an eigenvector of Q' with eigenvalue  $\lambda = 1$  but :

the existence of a unique eigenvalue with eigenvalue 1 is still not guaranteed

Let Q be an irreduible(Notraps), column − stochastic matrix with non negetive entries

⇒ Perron − Frobenius :

 $Largest\ eigenvalue = 1$ ,

$$eigenvector \ P = [p_1 \dots p_N]^T \ satisfies : p_i > 0 \ , \sum p_i = 1,$$

$$u = \frac{1}{N} [1, ..., 1]^T then \lim_{k \to \infty} Q^k u = P$$

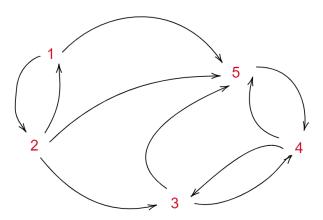
## Summary

- Adjancency matrix
- Normalize columns
- Eliminate trapes
- Use power method for computing the eigenvector

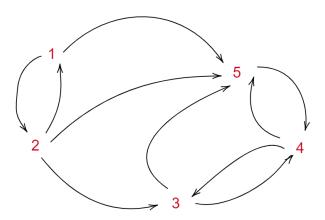
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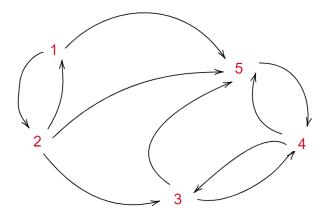
## Example:



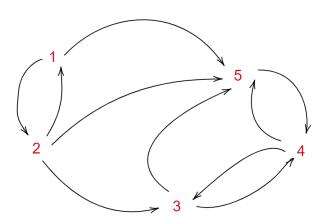
Adjancency matrix :



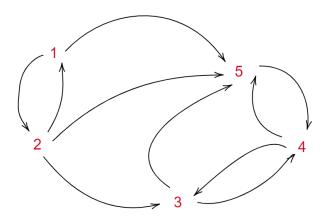
• Adjancency matrix :



• Normalize columns:



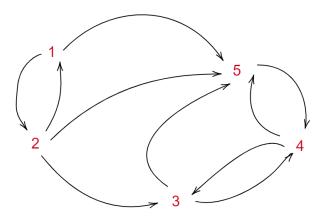
#### • Normalize columns:



## • Eliminate trapes:

$$Q' = (1 - \alpha)Q + \frac{\alpha}{N} \, 1_{n \times n}$$

 $\frac{\alpha}{N}$ : random jump probability

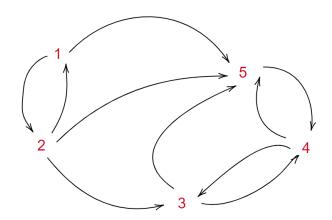


#### • Eliminate trapes:

$$Q' = (1-\alpha)Q + \frac{\alpha}{N} \ 1_{n \times n}$$

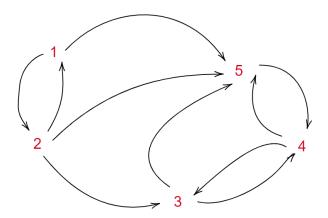
 $\frac{\alpha}{N}$ : random jump probability

$$\alpha = 0.01, Q' = \begin{bmatrix} 0.002 & 0.332 & 0.002 & 0.002 & 0.002 \\ 0.497 & 0.002 & 0.002 & 0.002 & 0.002 \\ 0.002 & 0.332 & 0.002 & 0.497 & 0.002 \\ 0.002 & 0.002 & 0.497 & 0.002 & 0.992 \\ 0.497 & 0.332 & 0.497 & 0.497 & 0.002 \end{bmatrix}$$



Use power method for computing the eigenvector

$$u = \begin{bmatrix} 0.2 \\ 0.2 \\ 0.2 \\ 0.2 \\ 0.2 \end{bmatrix}, Q'u = \begin{bmatrix} 0.07 \\ 0.0 \\ 0.7 \\ 0.30 \\ 0.36 \end{bmatrix}, Q'^2u = \begin{bmatrix} 0.003 \\ 0.004 \\ 0.8 \\ 0.45 \\ 0.30 \end{bmatrix}, ..., Q'^{10}u = \begin{bmatrix} 0.003 \\ 0.004 \\ 0.22 \\ 0.44 \\ 0.33 \end{bmatrix}, P = \begin{bmatrix} 0.0032 & 0.0036 & 0.2211 & 0.4401 & 0.3320 \end{bmatrix}$$



### Applications of PageRank Algorithm:

Search Engine Ranking
Internet Marketing
Social Network Analysis
Content Recommendation
Graph Analysis