# **Euclid's Algorithms**

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## **Objectives**

#### Your Objectives:

- ▶ Be able to calculate the GCD of two numbers using Euclid's algorithm.
- ▶ Use the extended Euclid's algorithm to solve Linear Diophantine equations.

- ▶ Let a > b > 0.
- ► Why?

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- ightharpoonup gcd(a,b) = gcd(b,mod(a,b))
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- ► That would be slow, so how about gcd(a,b) = gcd(b,a-nb), where n > 0 and a nb > 0 and minimal.
- ightharpoonup Easy! Just let n = mod(a, b)

### An example

$$\begin{aligned} \gcd(\mathsf{a}, \mathsf{b}) &= \gcd(\mathsf{b}, \mathsf{mod}(\mathsf{a}, \mathsf{b})) = \gcd(90, 25) \\ &= \gcd(25, 15) \\ &= \gcd(15, 10) \\ &= \gcd(10, 5) \\ &= \gcd(5, 0) \\ &= 5 \end{aligned}$$

# **Diophantine Equations**

- A Diophantine Equation is a polynomial equation where we are only interested in integer solutions.
- Linear Diophantine equation: ax + by = 1,
- It doesn't have to be 1....
- Running example: Suppose you go to the store. You buy x apples at 72 cents each and y oranges at 33 cents each. You spend \$5.85. How many of each did you buy?

Introduction

### How to do it

We want: ax + by = g, where g = gcd(a, b). We know a, b, and we calculate g. How can we get x and y?

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### JW 10 do 11

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▶ Then take  $a \mod b = a - \lfloor \frac{a}{b} \rfloor * b$  This gives:

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Rearrange a bit..

$$bx_1 + ay_1 - \left| \frac{a}{b} \right| by_1 = g \quad \Rightarrow \quad ay_1 + b(x_1 - \left| \frac{a}{b} \right| y_1) = g$$

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$$bx_1 + (a \mod b)y_1 = a$$

► Then take a mod  $b = a - \left| \frac{a}{b} \right| * b$  This gives:

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► This in turn gives us:

$$\begin{aligned}
 x &= y_1 \\
 y &= x_1 - \left| \frac{a}{b} \right| y_1
 \end{aligned}$$

#### The Code

```
x = y_1
                               y = x_1 - \left| \frac{a}{b} \right| y_1
// Stolen from Competitive Programming 3
// store x, y, and d as global variables
void extendedEuclid(int a, int b) {
   if (b == 0) \{ x = 1; y = 0; d = a; return; \}
   extendedEuclid(b, a % b):
   // similar as the original gcd
   int x1 = y;
   int v1 = x - (a / b) * v;
   x = x1;
   v = v1;
```

а	b	x	y	$a \times x + b \times y = 3$
72	33			
33	6			
6	3			
3	0			

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6	3	0	1	$6 \times 0 + 3 \times 1 = 3$
3	0	1	0	$3 \times 1 + 0 \times 0 = 3$

а	b	X	y	$a \times x + b \times y = 3$
72	33	-5	11	$72 \times -5 + 33 \times 11 = 3$
33	6	1	-5	$33 \times 1 + 6 \times -5 = 3$
6	3	0	1	$6 \times 0 + 3 \times 1 = 3$
3	0	1	Ο	$3 \times 1 + 0 \times 0 = 3$

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- Running the algorithm, we get...

$$72 \times -5 + 33 \times 11 = 3$$

lacktriangle We multiple both sides by 195 (since  $585=3\times195$ ) This gives us...

$$72 \times -975 + 33 \times 2145 = 585$$

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We can add  $72(\frac{33}{3})n$  to the 72 term and subtract  $33(\frac{72}{3})n$  from the second term and still have a valid equation.

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- ► Solve -975 + 11n > 0, this reduces to n > 88.6. So take n = 89.
- This gives us the final equation

$$72 \times 4 + 33 \times 9 = 585$$