# Bit Manipulations CS 491 – Competitive Programming

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### **Objectives**

- Compute binary representations of an integer
  - standard
  - one's compliment
  - two's compliment of arbitrary integers.
- ▶ Demonstrate the properties of boolean operations *and*, *or*, *not*, *xor*.
- Use shifting operations to test, set, and toggle arbitrary bits.
- Quickly determine if an integer is a power of 2.
- Quickly determine the number of set bits in an integer.
- Quickly determine the least significant set bit in an integer.

# Representation of a Positive Integer

- ► I think you know this very well by now....
  - ► Each digit is a successive power of 2
  - Let's use 6 bit integers for our examples.

```
2 = 000010
```

8 = 001000

10 = 001010

17 = 010001

## One's Compliment

- ► If you just "flip all the bits" you get one's compliment.
- ► In C++, the ~ operator will do this.

$$2 = 000010$$
  $\sim 2 = 111101$   
 $8 = 001000$   $\sim 8 = 110111$   
 $10 = 001010$   $\sim 10 = 110101$   
 $17 = 010001$   $\sim 17 = 101110$ 

We don't use one's compliment for negation though:

```
0 = 000000 ~0 = 111111
```

## Two's Compliment

- ► Take one's compliment (flip the bits) and then add one.
- ► In C++, regular old negation will do this.

```
000010
             \sim 2 = 111101
                              -2 = 111110
 001000
             ~8 = 110111
                              -8 = 111000
 001010
            ~10 =
                  110101
                             -10 = 110110
                             -17 = 101111
= 010001
            \sim 17 = 101110
  000000
             ~0 =
                  111111
                                   000000
                              -0 =
```

## Properties of And, Or, Not

#### Binary And &

- Commutative and associative.
- ► Identity is "all ones".

#### Binary Or |

- Commutative and associative.
- Identity is "all zeros."

#### Not ~

► Is its own inverse.  $\sim (\sim x) = x =$ 

## Example

```
a = 011001
b = 001010
c = 100110
a & b = 001000 a | b = 011011
b & c = 000010 b | c = 101110
a & c = 000000 a | c = 111111
```

#### Exclusive or

► Is true if bits are different.

$$0 \hat{\ } 0 = 0$$

$$0 \hat{1} = 1$$

$$1 \hat{1} = 0$$

- ► Is a good way to toggle bits:
- ▶ (a ^ b) ^ b == a

# **Shifting Operations**

Introduction and Objectives

- ► Use << to shift left, >> to shift right.
  - **▶** 001010 << 2 = 101000
  - ► 001010 >> 2 = 000010
- ► Allows easy multiplication and division by 2.
- ► Allows easy bit inspection and manipulations.

#### Check bit i

#### Set bit i

$$n = (1 << i)$$

#### Toggle bit *i*

# Some operations

Think about how you could do these operations.

- ightharpoonup Check if a number is divisible by 2.  $\mathcal{O}(1)$
- ► Clear lower *n* bits.  $\mathcal{O}(1)$
- ▶ Clear bits above  $n. \mathcal{O}(1)$
- ▶ Check if *n* is a power of 2.  $\mathcal{O}(1)$ 
  - $\blacktriangleright$  Hint: what is x & (x-1)?
- ▶ Count number of set bits in n.  $\mathcal{O}(b)$  b = number of bits.
- ▶ Get least significant set bit.  $\mathcal{O}(1)$ 
  - Hint: you need the two's compliment.

### Clear bits n and up

► To clear upper *n* bits, you need to create a bitmask that sets the lower bits.

```
mask = (1 << n) - 1:
x = x & mask;
 Example
x = 110011 -- lets clear bits 2 and up
mask = (1 << 2) -1;
     = 000100 - 1
     = 000011
x \& mask = 110011
         & 000011
           000011
```

#### Clear bits n and down

► To clear lower *n* bits, you need to create a bitmask that sets the lower bits, then compliment.

```
mask = \sim ((1 << (n+1)) - 1);
x = x & mask;
 Example
x = 110011 -- lets clear bits 2 and up
mask = \sim ((1 << 3) -1);
     = \sim (001000 - 1)
     = ~ 000111
          111000
x \& mask = 110011
          & 111000
            110000
```

# Check if n is odd or power of two

- ightharpoonup If x is odd, x | 1 is 1.
- ► Check x is power of 2, x & (x-1) will be zero.

▶ Use (x && !(x & (x-1))) to exclude when x is zero.



#### Check number of set bits in n

```
Consider n\mathcal{E}(n-1)...
        101100
n
n-1
      = 101011
        101000
so...
num = 0;
while (n>0) {
  num++;
  n = n & (n-1);
}
```

# Get least significant bit

```
n & (-n) will do this.

n = 101100
-n = 010100 (= 010011 + 1)
& -----
000100
```