

Fast Exponentiation

CS 491 – Competitive Programming

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Spring 2023

Objectives

- ▶ Use binary representation to compute exponentials in logarithmic time.
- ▶ Use a similar technique with matrices to compute Fibonacci numbers.

Naive Factorial

Remember the equation for n^m :

$$b^0 = 1$$

$$b^n = b * b^{n-1}$$

Recursive Implementation

```
1  int exponent(int base, int n) {  
2      if (n == 0)  
3          return 1;  
4      return n * exponent(base, n-1);  
5  }
```

Naive Factorial, II

Iterative Implementation

```
1  int fact(int n) {  
2      int out = 1;  
3      while (n>0) {  
4          out *= base;  
5          --n;  
6      }  
7      return out;  
8  }
```

- What is the time complexity?

Trick: Use a Binary Representation of the Exponent

► Q: What is 3^{22} ?

► Remember $a^x a^y = a^{x+y}$

► A: $3^{22} = 3^{10110_2} = 3^{10000_2} 3^{00100_2} 3^{00010_2} = 3^{16} 3^4 3^2$
 $= 1 \cdot 3^{10000_2} \times 0 \cdot 3^{01000_2} \times 1 \cdot 3^{00100_2} \times 1 \cdot 3^{00010_2} \times 0 \cdot 3^{00001_2}$

Implementation

Call with base=3 and n=22

```
1 int fexp(int b, int n) {
2     int out = 1;
3     int fact = b;
4
5     while (n>0) {
6         if (n & 1)
7             out *= fact;
8         fact *= fact;
9         n >>= 1;
10    }
11
12    return out;
13 }
```

► Initialize: out=1, fact=3, n=22

Implementation

Call with base=3 and n=22

```
1 int fexp(int b, int n) {  
2     int out = 1;  
3     int fact = b;  
4  
5     while (n>0) {  
6         if (n & 1)  
7             out *= fact;  
8         fact *= fact;  
9         n >>= 1;  
10    }  
11  
12    return out;  
13 }
```

- ▶ Initialize: out=1, fact=3, n=22
- ▶ n (22) & 1 is 0:
out=1, fact=9, n=11

Implementation

Call with base=3 and n=22

```
1 int fexp(int b, int n) {  
2     int out = 1;  
3     int fact = b;  
4  
5     while (n>0) {  
6         if (n & 1)  
7             out *= fact;  
8         fact *= fact;  
9         n >>= 1;  
10    }  
11  
12    return out;  
13 }
```

► Initialize: out=1, fact=3, n=22

► n (22) & 1 is 0:
out=1, fact=9, n=11

► n (11) & 1 is 1:
out=9, fact=81, n=5

Implementation

Call with base=3 and n=22

```
1 int fexp(int b, int n) {  
2     int out = 1;  
3     int fact = b;  
4  
5     while (n>0) {  
6         if (n & 1)  
7             out *= fact;  
8         fact *= fact;  
9         n >>= 1;  
10    }  
11  
12    return out;  
13 }
```

► Initialize: out=1, fact=3, n=22

► n (22) & 1 is 0:
out=1, fact=9, n=11

► n (11) & 1 is 1:
out=9, fact=81, n=5

► n (5) & 1 is 1:
out=729, fact=6561, n=2

Implementation

Call with base=3 and n=22

```
1 int fexp(int b, int n) {  
2     int out = 1;  
3     int fact = b;  
4  
5     while (n>0) {  
6         if (n & 1)  
7             out *= fact;  
8         fact *= fact;  
9         n >>= 1;  
10    }  
11  
12    return out;  
13 }
```

► Initialize: out=1, fact=3, n=22

► n (22) & 1 is 0:
out=1, fact=9, n=11

► n (11) & 1 is 1:
out=9, fact=81, n=5

► n (5) & 1 is 1:
out=729, fact=6561, n=2

► n (2) & 1 is 0:
out=729, fact=43046721, n=1

Implementation

Call with base=3 and n=22

```
1 int fexp(int b, int n) {  
2     int out = 1;  
3     int fact = b;  
4  
5     while (n>0) {  
6         if (n & 1)  
7             out *= fact;  
8         fact *= fact;  
9         n >>= 1;  
10    }  
11  
12    return out;  
13 }
```

► Initialize: out=1, fact=3, n=22

► n (22) & 1 is 0:
out=1, fact=9, n=11

► n (11) & 1 is 1:
out=9, fact=81, n=5

► n (5) & 1 is 1:
out=729, fact=6561, n=2

► n (2) & 1 is 0:
out=729, fact=43046721, n=1

► n (1) & 1 is 1:
out=31381059609, n=0

Calculating Fibonacci Numbers

Wrong Way

```
1  int fib(int n) {  
2      if (n<2) return 1;  
3      return fib(n-1) + fib(n-2)  
4  }
```

Reasonable way

```
1  int fib(int n, int a=1, int b=1) {  
2      if (n==1)  
3          return a;  
4      else  
5          return fib(n-1,b,a+b);  
6  }
```

Fast Fibonacci

- ▶ We can do even better! Consider this equation:

$$\begin{aligned}f'_1 &= f_1 + f_2 \\f'_2 &= f_1\end{aligned}$$

- ▶ We can represent this in matrix form.

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$$

- ▶ Repeating

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix}$$

- ▶ Again...!

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 3 \\ 2 \end{bmatrix} = \begin{bmatrix} 5 \\ 3 \end{bmatrix}$$

Square the Matrix

- ▶ Square the matrix to do multiple steps:

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 2 & 1 \\ 1 & 1 \end{bmatrix}$$

- ▶ Then...

$$\begin{bmatrix} 2 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} 5 \\ 3 \end{bmatrix}$$

- ▶ You can use the same technique as with fast exponents to “power up” this matrix and compute large Fibonacci numbers in logarithmic time.