

# 15. Dynamic Mechanism Design

---

ECON 7219 – Games With Incomplete Information

Benjamin Bernard

# Dynamic Revelation Principle

---

# Dynamic Selling Mechanism



## Before the auction:

- Auctioneer advertises the item he has for sale.
- Potential bidders form beliefs about their true valuation of the object.
- Potential bidders decide whether or not to attend the auction.

## At the auction:

- Bidders examine the item more closely and learn their true valuation.
- Bidders bid for the item and the highest bidder wins.

# Types and Utilities

## Type space:

- Player  $i$ 's utility depends only on their **ex-post** type  $\vartheta_i \in \Theta_i$ .
- At the time of entering the mechanism, player  $i$  has some information about  $\vartheta_i$  available, reflected in their **ex-ante** type  $\tau_i \in \mathcal{T}_i$ .
- We suppose types are one-dimensional:  $\Theta_i = [\underline{\vartheta}_i, \bar{\vartheta}_i]$  and  $\mathcal{T}_i = [\underline{\tau}_i, \bar{\tau}_i]$ .

## Quasi-linear utilities:

- The set of alternatives  $\mathcal{X}$  is  $\mathcal{Q} \times \mathbb{R}^n$ , where  $q \in \mathcal{Q}$  is the social state.
- We assume  $u_i(p, q, \vartheta_i) = v_i(q, \vartheta_i) - p_i$  such that  $0 \leq \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \leq K$ .
- Types are one-dimensional: there exists an order  $\succ_i$  of elements in  $\mathcal{Q}$  such that  $v_i$  has **increasing differences**, i.e.,  $q_H \succ q_L$  and  $\vartheta_i > \vartheta'_i$  imply

$$v_i(q_H, \vartheta_i) - v_i(q_L, \vartheta_i) \geq v_i(q_H, \vartheta'_i) - v_i(q_L, \vartheta'_i).$$

# Information of the Ex-Ante Type

## Independent types:

- Types  $(T_i, \theta_i)$  and  $(T_j, \theta_j)$  are independent for any  $i, j$ .
- Joint distribution  $F_i$  of  $(T_i, \theta_i)$  for any  $i$  is common knowledge.

## Information:

- Player  $i$ 's ex-ante information are reflected by his/her beliefs  $F_i(\vartheta_i | \tau_i)$  about his/her true valuation.
- We impose that  $F_i(\vartheta_i | \tau_i)$  is decreasing in  $\tau_i$  for any  $\vartheta_i \in (\underline{\vartheta}_i, \bar{\vartheta}_i)$ : higher values of  $\tau_i$  make high values of  $\vartheta_i$  more likely.
- The support of  $F_i(\vartheta_i | \tau_i)$  is independent of  $\tau_i$ : the ex-ante type does not provide any certainty about the payoff type.
- The partial derivative  $\partial F_i(\vartheta_i | \tau_i) / \partial \tau_i$  is bounded.

# Timing in Dynamic Mechanism Design



## Ex ante:

- Joint distribution of ex-ante and ex-post types is commonly known.
- Mechanism designer designs the mechanism.
- Players learn their ex-ante type and report an ex-ante type.

## Interim:

- Players observe their ex-post type (payoff type).
- Players decide what ex-post type to report.

## Ex post:

- Players' reports are publicly revealed.

## Individual rationality and incentive compatibility:

- Individual rationality has to be satisfied only at the ex-ante stage.
- Incentive compatibility has to hold at the ex-ante and interim stage.

# Direct Dynamic Mechanism

## Direct dynamic mechanism:

- A direct dynamic mechanism is a pair  $(p, q)$  with

$$q : \mathcal{T} \times \Theta \rightarrow \Delta(\mathcal{Q}), \quad p : \mathcal{T} \times \Theta \rightarrow \mathbb{R}^n.$$

- Denote by  $\alpha_q(\tau, \vartheta)$  the probability that  $q$  is chosen.

## Players' strategies:

- Players do not observe the ex-ante report of other players.
- We can write a strategy  $\sigma_i$  of player  $i$  using two maps

$$\sigma_i^0 : \mathcal{T}_i \rightarrow \mathcal{T}_i, \quad \sigma_i^1 : \mathcal{T}_i \times \Theta_i \times \mathcal{T}_i \rightarrow \Theta_i.$$

# Dynamic Revelation Principle

## Proposition 15.1

---

*For any indirect mechanism  $\Gamma$  and any PBE  $\sigma$  of that mechanism, there exists a direct dynamic mechanism  $\Gamma'$  and PBE  $\hat{\sigma}$  with*

$$\hat{\sigma}_i^0(\tau_i) = \tau_i, \quad \hat{\sigma}_i^1(\tau_i, \vartheta_i, \tau_i) = \vartheta_i,$$

*that induces the same distribution over outcomes in  $\Gamma'$  as  $\sigma$  does in  $\Gamma$ .*

---

### Remark:

- Dynamic revelation principle specifies truth-telling only on the path.
- Truth-telling off the path will follow from payoff independence of  $\tau_i$ .

**Proof:** Set  $\Gamma' = \Gamma \circ \sigma$ . Any deviation from  $\hat{\sigma}$  in  $\Gamma'$  has a corresponding deviation from  $\sigma$  in  $\Gamma$  and cannot be profitable.



# Interim and Ex-Ante Utilities

## Interim expected utility:

- Player  $i$ 's interim expected utility of reporting  $r_i \in \Theta_i$  is

$$u_i(R_i, r_i \mid \tau_i, \vartheta_i) := \mathbb{E}_{\tau_i, \vartheta_i} [v_i(q(R_i, T_{-i}, r_i, \theta_{-i}), \vartheta_i) - p_i(R_i, T_{-i}, r_i, \theta_{-i})]$$

conditional on ex-ante report  $R_i$  and truthful reporting by others.

## Ex-ante expected utility:

- Player  $i$ 's ex-ante utility for reporting  $R_i \in \mathcal{T}_i$  is

$$U_i(R_i, \sigma_i^1 \mid \tau_i) = \int_{\Theta_i} u_i(R_i, \sigma_i^1(\tau_i, \vartheta_i, R_i) \mid \tau_i, \vartheta_i) f_i(\vartheta_i \mid \tau_i) d\vartheta_i,$$

conditional on future report  $\sigma_i^1$  and truthful reporting by others.

- Denote by  $\hat{U}_i(R_i \mid \tau_i) = U_i(R_i, \hat{\sigma}_i^1 \mid \tau_i)$  the ex-ante utility of reporting  $R_i$ , conditional on truthful future report  $\hat{\sigma}_i^1(\tau_i, \vartheta_i, R_i) = \vartheta_i$ .

# Payoff Independence of the Ex-Ante Type

## Payoff independence:

- Conditional on truthful reporting by others, player  $i$ 's interim utility is

$$\begin{aligned}
 u_i(R_i, r_i \mid \tau_i, \vartheta_i) &= \mathbb{E}_{\tau_i, \vartheta_i} [v_i(q(R_i, T_{-i}, r_i, \theta_{-i}), \vartheta_i) - p_i(R_i, T_{-i}, r_i, \theta_{-i})] \\
 &= \sum_{q \in Q} v_i(q, \vartheta_i) \underbrace{\mathbb{E}_{\tau_i, \vartheta_i} [\alpha_q(R_i, T_{-i}, r_i, \theta_{-i})]}_{=: \bar{\alpha}_q(R_i, r_i)} - \underbrace{\mathbb{E}_{\tau_i, \vartheta_i} [p_i(R_i, T_{-i}, r_i, \theta_{-i})]}_{=: \bar{p}_i(R_i, r_i)} \\
 &=: u_i(r_i \mid R_i, \vartheta_i)
 \end{aligned}$$

- Knowing the ex-ante type is no longer valuable because:
  - At the interim stage, player  $i$  knows  $\vartheta_i$  already.
  - Types are independent, hence it does not help predict the others' types.
- It will be convenient to introduce the notation

$$\bar{v}_i(r_i \mid R_i, \vartheta_i) := \mathbb{E}_{R_i, \vartheta_i} [v_i(q(R_i, T_{-i}, r_i, \theta_{-i}), \vartheta_i)] = \sum_{q \in Q} v_i(q, \vartheta_i) \bar{\alpha}_q(R_i, r_i).$$

# Incentive Compatibility

## Definition 15.2

---

A direct mechanism is **incentive compatible** if:

1. It is incentive compatible with respect to the ex-post type, i.e., for every type  $(\tau_i, \vartheta_i)$ , and every ex-post report  $r_i \in \Theta_i$ ,

$$u_i(\vartheta_i \mid \tau_i, \vartheta_i) \geq u_i(r_i \mid \tau_i, \vartheta_i).$$

2. It is incentive-compatible with respect to the ex-ante type, i.e., for every  $\tau_i$ , every future report  $\sigma_i^1$ , and every ex-ante report  $R_i \in \mathcal{T}_i$

$$\hat{U}_i(\tau_i \mid \tau_i) \geq U_i(R_i, \sigma_i^1 \mid \tau_i).$$

---

# Incentive Compatibility

## Lemma 15.3

---

*A direct mechanism is incentive-compatible if and only if it satisfies*

1. *For every type  $(\tau_i, \vartheta_i)$ , and every ex-post report  $r_i \in \Theta_i$ ,*

$$u_i(\vartheta_i | \tau_i, \vartheta_i) \geq u_i(r_i | \tau_i, \vartheta_i).$$

2. *For every ex-ante type  $\tau_i$  and every ex-ante report  $R_i \in \mathcal{T}_i$*

$$\hat{U}_i(\tau_i | \tau_i) \geq \hat{U}_i(R_i, | \tau_i).$$

---

### Importance:

- Since ex-ante type  $\tau_i$  does not affect interim utilities, truth-telling on the path is sufficient to prevent deviations off the path as well.
- Thus, the revelation principle implies truth-telling also off the path.

# Proof of Lemma 15.3

## Proof of necessity:

- $\hat{U}_i(\tau_i | \tau_i) \geq U_i(R_i, \sigma_i^1 | \tau_i)$  for any  $\sigma_i^1$  implies  $\hat{U}_i(\tau_i | \tau_i) \geq \hat{U}_i(R_i | \tau_i)$ .

## Proof of sufficiency:

- Incentive-compatibility with respect to the ex-post type implies

$$\begin{aligned} U_i(R_i, \sigma_i^1 | \tau_i) &= \int_{\Theta_i} u_i(\sigma_i^1(\tau_i, \vartheta_i, R_i) | R_i, \vartheta_i) f_i(\vartheta_i | \tau_i) d\vartheta_i \\ &\leq \int_{\Theta_i} u_i(\vartheta_i | R_i, \vartheta_i) f_i(\vartheta_i | \tau_i) d\vartheta_i \\ &= \hat{U}_i(R_i | \tau_i) \leq \hat{U}_i(\tau_i | \tau_i). \end{aligned}$$

# Revenue Equivalence

---

# Characterizing Incentive Compatibility

## Static mechanism with one-dimensional types:

- Incentive compatibility = monotonicity + revenue equivalence.
- Selling mechanism:
  - Higher type has to receive the object with a higher likelihood.
  - Marginal increase in expected payments are equal to the marginal benefit of the increase in likelihood to obtain the item.
- Does a similar result hold for dynamic mechanisms?

## Interim incentive compatibility:

- For a given ex-ante type  $\tau_i$ , incentive compatibility with respect to the ex-post type is identical to the static case.

# Incentive Compatibility With Respect to the Ex-Post Type

## Lemma 15.4

*A direct dynamic mechanism is incentive compatible with respect to the ex-post type if and only if for every player  $i$  and every ex-ante type  $\tau_i$ :*

1. *“Monotonic” social state:  $h(\tau_i, \vartheta_i)$  is non-decreasing in  $\vartheta_i$ , where*

$$h(\tau_i, \vartheta_i) := \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(\tau_i, \vartheta_i).$$

2. *“Revenue equivalence” determines expected payments:*

$$\begin{aligned} \bar{p}_i(\tau_i, \vartheta_i) &= \bar{p}_i(\tau_i, \underline{\vartheta}) + \sum_{q \in Q} (v_i(q, \vartheta_i) \bar{\alpha}_q(\tau_i, \vartheta_i) - v_i(q, \underline{\vartheta}) \bar{\alpha}_q(\tau_i, \underline{\vartheta})) \\ &\quad - \sum_{q \in Q} \int_{\underline{\vartheta}}^{\vartheta_i} \frac{\partial v_i(q, x)}{\partial \vartheta_i} \bar{\alpha}_q(\tau_i, x) dx. \end{aligned}$$



# Incentive Compatibility With Respect to the Ex-Post Type

## Revenue equivalence:

- It turns out that a revenue-equivalence result holds.
- Step 1: show absolute continuity and differentiability of  $\hat{U}_i(\tau_i | \tau_i)$ .
- Step 2: use incentive compatibility to show the integral condition.

## Monotonicity:

- Monotonicity of the allocation function with respect to the ex-ante type is sufficient, but not necessary for incentive compatibility.
- Step 1: show a simple counterexample to necessity.
- Step 2: show sufficiency.

# Absolute Continuity

## Lemma 15.5

---

Suppose that there exist  $K_F$  and  $(K_q)_{q \in Q}$  such that

$$\left| \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \right| \leq K_q, \quad \left| \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} \right| \leq K_F.$$

Then for any incentive-compatible mechanism,  $\hat{U}_i(\tau_i | \tau_i)$  is increasing and absolutely continuous in  $\tau_i$ .

---

### Interpretation:

- Ex-ante expected utility is increasing in  $\tau_i$  under truth-telling.

### Selling mechanism:

- $K_q = 1$  since  $\frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} = 1_{\{q=\vartheta_i\}}$ .

# Proof of Monotonicity

## Proof of monotonicity:

- Integration by parts yields

$$\begin{aligned}
 \widehat{U}_i(R_i | \tau_i) &= \int_{\Theta_i} \underbrace{u_i(\vartheta_i | R_i, \vartheta_i)}_{\downarrow} \underbrace{f_i(\vartheta_i | \tau_i)}_{\uparrow} d\vartheta_i \\
 &= u_i(\bar{\vartheta} | R_i, \bar{\vartheta}) - \int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(R_i, \vartheta_i) F_i(\vartheta_i | \tau_i) d\vartheta_i.
 \end{aligned}$$

- For  $\tau_i^2 > \tau_i^1$ , ex-ante incentive compatibility implies

$$\begin{aligned}
 \widehat{U}_i(\tau_i^2 | \tau_i^2) - \widehat{U}_i(\tau_i^1 | \tau_i^1) &\geq \widehat{U}_i(\tau_i^1 | \tau_i^2) - \widehat{U}_i(\tau_i^1 | \tau_i^1) \\
 &= \int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(\tau_i^1, \vartheta_i) \underbrace{(F_i(\vartheta_i | \tau_i^1) - F_i(\vartheta_i | \tau_i^2))}_{\geq 0 \text{ by FOSD}} d\vartheta_i \geq 0.
 \end{aligned}$$

- This shows that  $\widehat{U}_i(\tau_i | \tau_i)$  is non-decreasing.

# Proof of Absolute Continuity

## Bounding the difference:

- Ex-ante incentive compatibility implies that for any  $\tau_i^2, \tau_i^1$ ,

$$\begin{aligned}\widehat{U}_i(\tau_i^2 | \tau_i^2) - \widehat{U}_i(\tau_i^1 | \tau_i^1) &\leq \widehat{U}_i(\tau_i^2 | \tau_i^2) - \widehat{U}_i(\tau_i^2 | \tau_i^1) \\ &\leq \sup_{R_i \in \mathcal{T}_i} \widehat{U}_i(R_i | \tau_i^2) - \widehat{U}_i(R_i | \tau_i^1).\end{aligned}$$

- Along the same lines, we obtain

$$\widehat{U}_i(\tau_i^2 | \tau_i^2) - \widehat{U}_i(\tau_i^1 | \tau_i^1) \geq \inf_{R_i \in \mathcal{T}_i} \widehat{U}_i(R_i | \tau_i^2) - \widehat{U}_i(R_i | \tau_i^1).$$

- Together, the two conditions yield

$$|\widehat{U}_i(\tau_i^2 | \tau_i^2) - \widehat{U}_i(\tau_i^1 | \tau_i^1)| \leq \sup_{R_i \in \mathcal{T}_i} |\widehat{U}_i(R_i | \tau_i^2) - \widehat{U}_i(R_i | \tau_i^1)|.$$

# Proof of Absolute Continuity

## Lipschitz continuity:

- For  $\tau_i^2 > \tau_i^1$ , we have  $F_i(\vartheta_i | \tau_i^1) - F_i(\vartheta_i | \tau_i^2) > 0$ , hence

$$\begin{aligned}
 |\widehat{U}_i(\tau_i^2 | \tau_i^2) - \widehat{U}_i(\tau_i^1 | \tau_i^1)| &\leq \sup_{R_i \in \mathcal{T}_i} |\widehat{U}_i(R_i | \tau_i^2) - \widehat{U}_i(R_i | \tau_i^1)| \\
 &\leq \int_{\Theta_i} \left| \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(R_i, \vartheta_i) \right| (F_i(\vartheta_i | \tau_i^1) - F_i(\vartheta_i | \tau_i^2)) \, d\vartheta_i \\
 &\leq \int_{\Theta_i} \max_q K_q \left| \frac{\partial F_i(\vartheta_i | \tau_i')}{\partial \tau_i} \right| |\tau_i^2 - \tau_i^1| \, d\vartheta_i \\
 &\leq \max_q K_q K_F(\bar{\vartheta} - \underline{\vartheta}) |\tau_i^2 - \tau_i^1|,
 \end{aligned}$$

where we have used the mean-value theorem.

Thus,  $\widehat{U}_i(\tau_i | \tau_i)$  is Lipschitz continuous and hence absolutely continuous.

# Revenue Equivalence for Ex-Ante Utility

## Proposition 15.6

Let  $U_i(\tau_i) := \hat{U}_i(\tau_i, \tau_i)$ . For any incentive-compatible direct mechanism:

1.  $U_i$  is differentiable everywhere except at most countably many points.

At any point of differentiability  $\tau_i$ , we have

$$U'_i(\tau_i) = - \int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(\tau_i, \vartheta_i) \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} d\vartheta_i.$$

2. For every ex-ante type  $\tau_i$ , we have

$$U_i(\tau_i) = U_i(\underline{\tau}) - \int_{\underline{\tau}}^{\tau_i} \int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(t, \vartheta_i) \frac{\partial F_i(\vartheta_i | t)}{\partial \tau_i} d\vartheta_i dt.$$

**Corollary:** An incentive-compatible mechanism is ex-ante individually rational if and only if  $U_i(\underline{\tau}) \geq 0$ .

# Proof of Proposition 15.6

## Derivative with respect to ex-ante type:

- Since  $\hat{U}_i(\tau_i | \tau_i)$  is monotonic, it is differentiable almost everywhere.
- Since  $F_i(\vartheta_i | \tau_i)$  is differentiable with respect to  $\tau_i$ ,

$$\hat{U}_i(R_i | \tau_i) = u_i(\bar{\vartheta} | R_i, \bar{\vartheta}) - \int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(R_i, \vartheta_i) F_i(\vartheta_i | \tau_i) d\vartheta_i$$

is differentiable as well, with derivative

$$\frac{\partial \hat{U}_i(R_i | \tau_i)}{\partial \tau_i} = - \int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(R_i, \vartheta_i) \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} d\vartheta_i.$$

# Proof of Proposition 15.6

## Revenue equivalence:

- Incentive-compatibility implies:

$$\frac{\hat{U}_i(\tau_i | \tau_i) - \hat{U}(\tau_i + \delta | \tau_i + \delta)}{\delta} \leq \frac{\hat{U}_i(\tau_i | \tau_i) - \hat{U}(\tau_i | \tau_i + \delta)}{\delta},$$

$$\frac{\hat{U}_i(\tau_i - \delta | \tau_i - \delta) - \hat{U}(\tau_i | \tau_i)}{\delta} \geq \frac{\hat{U}_i(\tau_i | \tau_i - \delta) - \hat{U}(\tau_i | \tau_i)}{\delta}.$$

- At any differentiability point of  $U_i(\tau_i) := \hat{U}_i(\tau_i | \tau_i)$ , we obtain

$$\left. \frac{\partial \hat{U}_i(R_i | \tau_i)}{\partial \tau_i} \right|_{R_i = \tau_i} \leq U'_i(\tau_i) \leq \left. \frac{\partial \hat{U}_i(R_i | \tau_i)}{\partial \tau_i} \right|_{R_i = \tau_i}.$$

- Therefore,

$$U'_i(\tau_i) = - \int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(R_i, \vartheta_i) \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} d\vartheta_i.$$



# Revenue Equivalence for Payments

## Proposition 15.7

---

*If a direct dynamic mechanism is incentive-compatible, then:*

$$\bar{p}_i(\tau_i, \vartheta_i) = \bar{p}_{i, \underline{\vartheta}}(\tau_i) + \sum_{q \in Q} v_i(q, \vartheta_i) \bar{\alpha}_q(\tau_i, \vartheta_i) - \int_{\underline{\vartheta}}^{\vartheta_i} h_i(\tau_i, x) dx,$$

where  $h_i(\tau_i, \vartheta_i) := \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(\tau_i, \vartheta_i)$  and

$$\begin{aligned} \bar{p}_{i, \underline{\vartheta}}(\tau_i) := & \bar{p}_i(\underline{\tau}, \underline{\vartheta}) - \sum_{q \in Q} v_i(q, \underline{\vartheta}) \bar{\alpha}_q(\underline{\tau}, \underline{\vartheta}) + \int_{\underline{\tau}}^{\tau_i} \int_{\underline{\vartheta}}^{\bar{\vartheta}} h_i(t, \vartheta) \frac{\partial F_i(\vartheta_i | t)}{\partial \tau_i} d\vartheta_i dt \\ & + \int_{\underline{\vartheta}}^{\bar{\vartheta}} \int_{\underline{\vartheta}}^{\vartheta_i} (h_i(\tau_i, x) f_i(\vartheta_i | \tau_i) - h_i(\underline{\tau}, x) f_i(\vartheta_i | \underline{\tau})) dx d\vartheta_i \end{aligned}$$


---

**Note:** Expected payments are determined uniquely up to  $\bar{p}_i(\underline{\tau}, \underline{\vartheta})$ .

# Proof of Proposition 15.7

## Use revenue equivalence for ex ante utility:

- Recall the definition of the ex-ante utility

$$U_i(\tau_i) = \int_{\Theta_i} \left( \sum_{q \in Q} v_i(q, \vartheta_i) \bar{\alpha}_q(\tau_i, \vartheta_i) - \bar{p}_i(\tau_i, \vartheta_i) \right) f_i(\vartheta_i | \tau_i) d\vartheta_i,$$

- From revenue equivalence for ex-ante utilities, we obtain

$$\begin{aligned} \int_{\Theta_i} \bar{p}_i(\tau_i, \vartheta_i) f(\vartheta_i | \tau_i) d\vartheta_i &= \sum_{q \in Q} \int_{\Theta_i} v_i(q, \vartheta_i) \bar{\alpha}_q(\tau_i, \vartheta_i) f_i(\vartheta_i | \tau_i) d\vartheta_i \\ &+ \int_{\Theta_i} \left( \bar{p}_i(\underline{\tau}, \vartheta_i) - \sum_{q \in Q} v_i(q, \vartheta_i) \bar{\alpha}_q(\underline{\tau}, \vartheta_i) \right) f_i(\vartheta_i | \underline{\tau}) d\vartheta_i \\ &+ \int_{\underline{\tau}}^{\tau} \int_{\Theta_i} h_i(t, \vartheta_i) \frac{\partial F_i(\vartheta_i | t)}{\partial \tau_i} d\vartheta_i dt \end{aligned}$$

# Proof of Proposition 15.7

## Use revenue equivalence for ex-post type:

- Replacing  $\bar{p}_i(\tau_i, \vartheta_i)$  and  $\bar{p}_i(\underline{\tau}, \vartheta_i)$  with the expression from revenue equivalence for the ex-post type yields

$$\begin{aligned} \bar{p}_i(\tau_i, \underline{\vartheta}) &= \bar{p}_i(\underline{\tau}, \underline{\vartheta}) + \sum_{q \in Q} v_i(q, \underline{\vartheta}) (\bar{\alpha}_q(\tau_i, \underline{\vartheta}) - \bar{\alpha}_q(\underline{\tau}, \underline{\vartheta})) \\ &\quad + \int_{\underline{\tau}}^{\tau_i} \int_{\underline{\vartheta}}^{\bar{\vartheta}} h_i(t, \vartheta_i) \frac{\partial F(\vartheta_i | t)}{\partial \tau_i} d\vartheta_i dt \\ &\quad + \int_{\underline{\vartheta}}^{\bar{\vartheta}} \int_{\underline{\tau}}^{\tau_i} (h_i(\tau_i, x) f(\vartheta_i | \tau_i) - h_i(\underline{\tau}, x) f(\vartheta_i | \underline{\tau})) dx d\vartheta_i. \end{aligned}$$

- Result now follows from revenue equivalence with respect to the ex-post type with the above expression for  $\bar{p}_i(\tau_i, \underline{\vartheta})$ .

# Monotonicity with Respect to Ex-Ante Type May Fail

## Selling to a single buyer:

- There are two states: buyer obtains the good ( $q = 1$ ) or not ( $q = 0$ ).
- Let  $q_i(R_i, \vartheta_i) = \bar{\alpha}_q(R_i, \vartheta_i)$  denote the probability of selling the good.
- Let  $\bar{p}_i(\tau_i, \vartheta_i) = p_i(\tau_i, \vartheta_i)$ , and suppose  $q_i$  and  $p_i$  are twice differentiable.

## Ex-ante utility:

- Since  $\frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} = 1_{\{q=1\}}$ , we obtain

$$\begin{aligned}\hat{U}_i(R_i | \tau_i) &= u_i(\bar{\vartheta} | R_i, \bar{\vartheta}) - \underbrace{\int_{\Theta_i} \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(R_i, \vartheta_i) F_i(\vartheta_i | \tau_i) d\vartheta_i}_{q_i(R_i, \vartheta_i)} \\ &= \bar{\vartheta} q_i(R_i, \bar{\vartheta}) - p_i(R_i, \bar{\vartheta}) - \int_{\Theta_i} q_i(R_i, \vartheta_i) F_i(\vartheta_i | \tau_i) d\vartheta_i.\end{aligned}$$

# Monotonicity with Respect to Ex-Ante Type May Fail

**First- and second-order constraints:**

$$\left. \frac{\partial \hat{U}_i(R_i | \tau_i)}{\partial R_i} \right|_{R_i = \tau_i} = \bar{\vartheta} q'_i(\tau_i, \bar{\vartheta}) - p'_i(\tau_i, \bar{\vartheta}) - \int_{\Theta_i} q'_i(\tau_i, \vartheta_i) F_i(\vartheta_i | \tau_i) d\vartheta_i = 0,$$

$$\left. \frac{\partial^2 \hat{U}_i(R_i | \tau_i)}{\partial R_i^2} \right|_{R_i = \tau_i} = \bar{\vartheta} q''_i(\tau_i, \bar{\vartheta}) - p''_i(\tau_i, \bar{\vartheta}) - \int_{\Theta_i} q''_i(\tau_i, \vartheta_i) F_i(\vartheta_i | \tau_i) d\vartheta_i \leq 0.$$

Differentiate first-order constraint with respect to  $\tau_i$ :

$$0 = \underbrace{\bar{\vartheta} q''_i(\tau_i, \bar{\vartheta}) - p''_i(\tau_i, \bar{\vartheta}) - \int_{\Theta_i} q''_i(\tau_i, \vartheta_i) F_i(\vartheta_i | \tau_i) d\vartheta_i}_{\leq 0 \text{ by SOSC}} - \int_{\Theta_i} q'_i(\tau_i, \vartheta_i) \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} d\vartheta_i.$$

Therefore,  $q_i(\tau_i, \vartheta_i)$  is increasing in  $\tau_i$  “on average”:

$$\int_{\Theta_i} q'_i(\tau_i, \vartheta_i) \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} d\vartheta_i \leq 0.$$

# Monotonicity

## Lemma 15.8

---

*Suppose  $q(\tau, \vartheta)$  is such that  $h_i(\tau_i, \vartheta_i) = \sum_{q \in Q} \frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} \bar{\alpha}_q(\tau_i, \vartheta_i)$  is non-decreasing in  $\tau_i$  and  $\vartheta_i$  for every player  $i$ . Then there exist payments  $p(\tau, \vartheta)$  such that the direct mechanism  $(q, p)$  is incentive-compatible.*

---

### Interpretation:

- The allocation of an incentive-compatible dynamic mechanism may not be monotonic in the ex-ante type.
- Monotonicity, however, is sufficient for incentive-compatibility.

# Proof of Lemma 15.8

## Incentive compatibility of ex-post type:

- Suppose that  $q(\tau, \vartheta)$  satisfies the monotonicity constraint.
- By Proposition 15.7, we must define  $\bar{p}_i(\tau_i, \vartheta_i)$  via the revenue equivalence for payments, choosing  $p_i(\underline{\tau}, \underline{\vartheta})$  such that  $U_i(\underline{\tau}) = 0$ .
- The mechanism  $(q, p)$  is incentive compatible with respect to ex-post type by Lemma 15.4
- By Lemma 15.3, it remains to show  $U_i(\tau_i) \geq \hat{U}_i(R_i | \tau_i)$ .

# Proof of Lemma 15.8

## Incentive compatibility of ex-post type:

- Recall that

$$\frac{\partial \hat{U}_i(R_i | \tau_i)}{\partial \tau_i} = - \int_{\Theta_i} h_i(R_i, \vartheta_i) \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} d\vartheta_i.$$

- The revenue equivalence for ex-ante utilities yields

$$\begin{aligned} U_i(\tau_i) - \hat{U}_i(R_i | \tau_i) &= \hat{U}_i(\tau_i | \tau_i) - \hat{U}_i(R_i | R_i) + \hat{U}_i(R_i | R_i) - \hat{U}_i(R_i | \tau_i) \\ &= \int_{R_i}^{\tau_i} U'(t) - \frac{\partial \hat{U}_i(R_i | t)}{\partial \tau_i} dt \\ &= \int_{R_i}^{\tau_i} \int_{\Theta_i} \underbrace{(h_i(R_i, \vartheta_i) - h_i(t, \vartheta_i))}_{\leq 0} \underbrace{\frac{\partial F_i(\vartheta_i | t)}{\partial \tau_i}}_{\leq 0 \text{ by FOSD}} d\vartheta_i dt \end{aligned}$$

- This shows that  $(q, p)$  is incentive-compatible.



# Dynamic vs. Static Mechanisms

## Revelation principle:

- The dynamic revelation principle gives us truth-telling only on the equilibrium path.
- If the ex-ante information is not directly payoff relevant, then truth-telling on the path is sufficient for off-path truth-telling.

## Revenue equivalence:

- Revenue equivalence for ex-post type holds as in the static case.
- If  $\frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i}$  is bounded, then it also holds for the ex-ante type.

## Monotonicity:

- Monotonicity for the ex-post type holds as in the static case.
- Monotonicity for the ex-ante type is sufficient for incentive-compatibility, but it is not necessary.

# Optimal Selling Mechanism

---

# Optimal Selling Mechanism

## Setup:

- There are  $i = 1, \dots, n$  potential buyers.
- There are  $n + 1$  social states  $q_i$ :  $i$  obtains the good and  $q_0$ : the seller keeps the good. Suppose the seller places no value on the item.
- Buyer  $i$  obtains the item with subjective probability

$$\bar{q}_i(\tau_i, \vartheta_i) := P_{\tau_i, \vartheta_i}(q(\tau, \vartheta) = q_i) = \bar{\alpha}_q(\tau_i, \vartheta_i) = h_i(\tau_i, \vartheta_i),$$

where we have used that  $\frac{\partial v_i(q, \vartheta_i)}{\partial \vartheta_i} = 1_{\{q=q_i\}}$ .

## Marginal distribution of ex-ante type:

- Let  $g_i(\tau_i) := \int_{\Theta_i} f_i(\tau_i, \vartheta_i) d\vartheta_i$  denote the marginal density function.
- Let  $G_i(\tau_i) := \int_{\underline{\tau}}^{\tau_i} g_i(t) dt$  denote the marginal distribution function.

# Optimal Selling Mechanism

## Expected revenue from a single buyer:

- By revenue equivalence, payments are determined by allocation rule.
- Seller's expected revenue from buyer  $i$  is

$$Rev_i = \int_{\mathcal{T}_i} \int_{\Theta_i} \vartheta_i \bar{q}_i(\tau_i, \vartheta_i) f_i(\vartheta_i | \tau_i) g_i(\tau_i) d\vartheta_i d\tau_i - \int_{\mathcal{T}_i} U_i(\tau_i) g_i(\tau_i) d\tau_i.$$

- Revenue is maximized if  $U_i(\underline{\tau}) = 0$ .
- Revenue equivalence for ex-ante utilities thus yields

$$\begin{aligned} \int_{\mathcal{T}_i} \underbrace{U_i(\tau_i)}_{\downarrow} \underbrace{(-g_i(\tau_i))}_{\uparrow} d\tau_i &= - \int_{\mathcal{T}_i} (1 - G_i(\tau_i)) U_i'(\tau_i) d\tau_i \\ &= \int_{\mathcal{T}_i} \int_{\Theta_i} (1 - G_i(\tau_i)) \bar{q}_i(\tau_i, \vartheta_i) \frac{\partial F_i(\vartheta_i | \tau_i)}{\partial \tau_i} d\vartheta_i d\tau_i. \end{aligned}$$

# Optimal Selling Mechanism

## Expected revenue from a single buyer:

- Define  $i$ 's **virtual valuation** as

$$\psi_i(\tau_i, \vartheta_i) := \vartheta_i + \frac{1 - G_i(\tau_i)}{g_i(\tau_i)} \frac{\partial F_i(\vartheta_i | \tau_i) / \partial \tau_i}{f_i(\vartheta_i | \tau_i)}.$$

- Then the seller's expected revenue from buyer  $i$  is

$$Rev_i = \int_{\mathcal{T}_i} \int_{\Theta_i} \psi_i(\tau_i, \vartheta_i) \bar{q}_i(\tau_i, \vartheta_i) f_i(\vartheta_i | \tau_i) g_i(\tau_i) d\vartheta_i d\tau_i.$$

- Since  $\bar{q}_i(\tau_i, \vartheta_i) = \int_{\mathcal{T}_{-i}} \int_{\Theta_{-i}} q_i(\tau, \vartheta) f_{-i}(\vartheta_{-i} | \tau_{-i}) g_{-i}(\tau_{-i})$ , we obtain

$$Rev_i = \int_{\mathcal{T}} \int_{\Theta} \psi_i(\tau_i, \vartheta_i) q_i(\tau, \vartheta) f(\vartheta | \tau) g(\tau) d\vartheta d\tau.$$

# Optimal Selling Mechanism

## Seller's expected revenue:

- Total revenue equals

$$Rev = \int_{\mathcal{T}} \int_{\Theta} \sum_{i=1}^n \psi_i(\tau_i, \vartheta_i) q_i(\tau, \vartheta) f(\vartheta | \tau) g(\tau) d\vartheta d\tau.$$

- Revenue is maximized for

$$q_i(\tau, \vartheta) = \begin{cases} 1 & \text{if } \psi_i(\tau_i, \vartheta_i) \geq \max(\max_j \psi_j(\tau_j, \vartheta_j), 0) \\ 0 & \text{otherwise.} \end{cases}$$

- If  $q_i$  is non-decreasing, then payments  $p$  given by revenue equivalence make the mechanism incentive-compatible.

# Optimal Selling Mechanism

## Proposition 15.9

---

*Suppose  $\psi_i(\tau_i, \vartheta_i)$  is non-decreasing in  $\tau_i$  and  $\vartheta_i$ . Then an incentive-compatible and individually rational direct dynamic mechanism  $(q, p)$  maximizes the seller's expected revenue if and only if*

$$q_i(\tau, \vartheta) = \begin{cases} 1 & \text{if } \psi_i(\tau_i, \vartheta_i) \geq \max(\max_j \psi_j(\tau_j, \vartheta_j), 0) \\ 0 & \text{otherwise,} \end{cases}$$

*and payments  $p_i(\tau_i, \vartheta_i)$  are determined from revenue equivalence such that  $U_i(\underline{\tau}) = 0$  for every buyer  $i$ .*

---

# Information Rent

## Information rent consists of two components:

- The hazard rate  $\frac{g_i(\tau_i)}{1-G_i(\tau_i)}$  of the ex-ante type.
- Informativeness measure  $\frac{\partial F_i(\vartheta_i | \tau_i) / \partial \tau_i}{f_i(\vartheta_i | \tau_i)}$  that captures how the buyers' knowledge of their valuation changes with ex-ante type.

## Reserve price:

- Since  $\psi_i$  is non-decreasing, there exists a cutoff

$$\hat{p}_i(\tau_i) := \min\{\vartheta_i \in \Theta_i \mid \psi_i(\vartheta_i, \tau_i) \geq 0\}$$

such that  $\psi(\vartheta_i, \tau_i) < 0$  if and only if  $\vartheta_i < \hat{p}_i(\tau_i)$ .

- Thus, the good is sold if and only if  $\vartheta_i \geq \hat{p}_i(\tau_i)$ .
- Since  $\psi_i$  is non-decreasing, the reserve price  $\hat{p}_i(\tau_i)$  is non-increasing in the announced ex-ante type  $\tau_i$ .



# Optimal Selling Mechanism With a Single Buyer

## Payments:

- For a pair  $(\tau_i, \vartheta_i)$  with  $\vartheta_i < \hat{p}_i(\tau_i)$ :

$$p_i(\tau_i, \vartheta_i) = p_{i, \underline{\vartheta}_i}(\tau_i) + \underbrace{\vartheta_i q_i(\tau_i, \vartheta_i)}_{=0} - \int_{\underline{\vartheta}_i}^{\vartheta_i} \underbrace{q_i(\tau_i, x)}_{=0} dx = p_{i, \underline{\vartheta}_i}(\tau_i).$$

- For a pair  $(\tau_i, \vartheta_i)$  with  $\vartheta_i \geq \hat{p}_i(\tau_i)$ :

$$\begin{aligned} p_i(\tau_i, \vartheta_i) &= p_{i, \underline{\vartheta}_i}(\tau_i) + \underbrace{\vartheta_i q_i(\tau_i, \vartheta_i)}_{=1} - \int_{\hat{p}_i(\tau_i)}^{\vartheta_i} \underbrace{q_i(\tau_i, x)}_{=1} dx \\ &= p_{i, \underline{\vartheta}_i}(\tau_i) + \hat{p}_i(\tau_i). \end{aligned}$$

- $p_{i, \underline{\vartheta}_i}(\tau_i)$  is determined by revenue equivalence and  $U_i(\underline{\tau}) = 0$ .

# Optimal Selling Mechanism With a Single Buyer

## Proposition 15.10

---

*Suppose  $\psi_i(\tau_i, \vartheta_i)$  is non-decreasing in  $\tau_i$  and  $\vartheta_i$ . There exist prices  $p_{i,\underline{\vartheta}_i}(\tau_i)$  and  $\hat{p}_i(\tau_i)$  such that the optimal selling mechanism takes the form*

$$q_i(\tau_i, \vartheta_i) = \begin{cases} 1 & \text{if } \vartheta_i \geq \hat{p}_i(\tau_i), \\ 0 & \text{otherwise,} \end{cases}$$
$$p_i(\tau_i, \vartheta_i) = \begin{cases} p_{i,\underline{\vartheta}_i}(\tau_i) + \hat{p}_i(\tau_i) & \text{if } \vartheta_i \geq \hat{p}_i(\tau_i), \\ p_{i,\underline{\vartheta}_i}(\tau_i) & \text{otherwise,} \end{cases}$$

---

## Implementation:

- Offer a menu of option contracts  $(p_{i,\underline{\vartheta}_i}(\tau_i), \hat{p}_i(\tau_i))_{\tau_i}$  to the buyer.
- In contract  $(p_{i,\underline{\vartheta}_i}(\tau_i), \hat{p}_i(\tau_i))$ , buyer buys a call option with exercise price  $\hat{p}_i(\tau_i)$  (= right to buy item at price  $\hat{p}_i(\tau_i)$ ) for the price  $p_{i,\underline{\vartheta}_i}(\tau_i)$ .

# Optimal Selling Mechanism With Multiple Buyers

If the informativeness measure is a function of only the ex-ante type

$$\frac{\partial F_i(\vartheta_i | \tau_i) / \partial \tau_i}{f_i(\vartheta_i | \tau_i)} = \phi_i(\tau_i),$$

then the optimal selling mechanism allows a similar interpretation. Set

$$\hat{p}_i(\tau_i) := -\frac{1 - G_i(\tau_i)}{g_i(\tau_i)} \phi_i(\tau_i).$$

## Implementation:

- At the beginning, buyers can buy a premium  $\hat{p}_i(\tau_i)$  for price  $p_{i,\underline{v}_i}(\tau_i)$ .
- After buyers learn their valuation, buyers bid in a second-price auction without reserve price, but in addition to the second-highest bid they also have to pay the acquired premium  $\hat{p}_i(\tau_i)$ .
- It is thus a weakly dominant strategy to bid their true valuation minus the premium, i.e., to bid their virtual valuation.

# **Value of Private Information**

---

# Decomposition of Information

## Decomposition into initial and additional information:

- First, every participant  $i$  observes realization  $\tau_i$  of  $T_i$ .
- Then, participant observes  $A_i = F(\theta_i | \tau_i)$ .
- Note that  $A_i$  is a transformation of the random variable  $\theta_i$  via  $F(\cdot | \tau_i)$ .
- It is distributed on  $[0, 1]$  and its conditional distribution is

$$\begin{aligned} P(A_i \leq \alpha_i | \tau_i) &= P(F(\theta_i | \tau_i) \leq \alpha_i | \tau_i) = P(\theta_i \leq F_i^{-1}(\alpha_i | \tau_i) | \tau_i) \\ &= F_i(F_i^{-1}(\alpha_i | \tau_i) | \tau_i) = \alpha_i. \end{aligned}$$

- $A_i$  is uniformly distributed and stochastically independent of  $T_i$ .

## Interpretation:

- $T_i$  is a noisy signal of  $\theta_i$ , where  $A_i$  is the noise.
- Upon learning the noise  $A_i = \alpha_i$ , buyer  $i$  learns  $\vartheta_i = F^{-1}(\alpha_i | \tau_i)$ .

# Value of Additional Information

## Consider two mechanisms:

- $A_i$  is learned privately by the buyer.
- $A_i$  is observed publicly, hence the seller does not need to elicit  $\alpha_i$ .
- Difference between seller's revenue is information rent for  $\alpha_i$ .

## $A_i$ is learned privately:

- This is isomorphic to the case we have studied.
- Define the virtual valuation

$$\begin{aligned}\psi_i(\tau_i, \alpha_i) &= F^{-1}(\alpha_i | \tau_i) + \frac{1 - G_i(\tau_i)}{g_i(\tau_i)} \frac{\partial F_i(F_i^{-1}(\alpha_i | \tau_i) | \tau_i) / \partial \tau_i}{f_i(F_i^{-1}(\alpha_i | \tau_i) | \tau_i)} \\ &= F^{-1}(\alpha_i | \tau_i) + \frac{1 - G_i(\tau_i)}{g_i(\tau_i)} \frac{\partial F_i^{-1}(\alpha_i | \tau_i)}{\partial \tau_i}.\end{aligned}$$

- Sell the item to buyer with the highest non-negative virtual valuation.

# Optimal Selling Mechanism for Publicly Observed Information

## Proposition 15.11

---

*Suppose  $\psi_i(\tau_i, \alpha_i)$  is non-decreasing in  $\tau_i$  and  $\alpha_i$  for every player  $i$ . Then the optimal selling mechanism when  $\alpha$  is privately observed is also optimal when  $\alpha$  is publicly observed.*

---

### Interpretation:

- Additional information can be elicited from the seller at no cost.
- Private information before entering the mechanism is more powerful than private information learned afterwards.
- The seller would like to contract early to minimize adverse selection.

### Note:

- The result does not hold if  $\psi_i$  is not monotonic.

# Revenue Equivalence for Publicly Observed Information

## Lemma 15.12

*For any incentive-compatible direct selling mechanism  $(q(\tau, \alpha), p(\tau, \alpha))$  with publicly observable  $\alpha$ , we have*

$$U_i(\tau_i) = U_i(\underline{\tau}) - \int_{\underline{\tau}}^{\tau_i} \int_0^1 \frac{\partial F_i^{-1}(\alpha_i | t)}{\partial \tau_i} \bar{q}_i(t, \alpha_i) d\alpha_i dt.$$

### Proof:

- Since  $U_i(\tau_i + \delta) \geq \hat{U}_i(\tau_i | \tau_i + \delta)$ , we obtain

$$\frac{U_i(\tau_i + \delta) - U_i(\tau_i)}{\delta} \geq \int_0^1 \frac{F_i^{-1}(\alpha_i | \tau_i + \delta) - F_i^{-1}(\alpha_i | \tau_i)}{\delta} \bar{q}_i(\tau_i, \alpha_i) d\alpha_i.$$

- Do the same for  $\tau_i - \delta$ , take limits, integrate.



# Proof of Proposition 15.11

## Expected revenue from a single buyer:

- The seller's expected revenue from buyer  $i$  is


$$Rev_i = \int_{\mathcal{T}_i} \int_0^1 F_i^{-1}(\alpha_i | \tau_i) \bar{q}_i(\tau_i, \alpha_i) g_i(\tau_i) d\alpha_i d\tau_i - \int_{\mathcal{T}_i} U_i(\tau_i) g(\tau_i) d\tau_i.$$

- Revenue equivalence for publicly observed  $\alpha$  thus yields

$$\begin{aligned} \int_{\mathcal{T}_i} \underbrace{U_i(\tau_i)}_{\downarrow} \underbrace{(-g(\tau_i))}_{\uparrow} d\tau_i &= - \int_{\mathcal{T}_i} (1 - G_i(\tau_i)) U_i'(\tau_i) d\tau_i \\ &= \int_{\mathcal{T}_i} \int_0^1 (1 - G_i(\tau_i)) \frac{\partial F_i^{-1}(\alpha_i | \tau_i)}{\partial \tau_i} \bar{q}_i(\tau_i, \vartheta_i) d\alpha_i d\tau_i. \end{aligned}$$

- Therefore, we obtain  $Rev_i = \int_{\mathcal{T}} \int_0^1 \psi_i(\tau_i, \alpha_i) q_i(\tau_i, \alpha_i) g(\tau_i) d\alpha_i d\tau_i$  for the same  $\psi_i$  as for privately observed  $\alpha$ .
- The remainder of proof works as before.

# Literature

 T. Börgers: *An Introduction to the Theory of Mechanism Design*, Chapter 11, Oxford University Press, 2015

