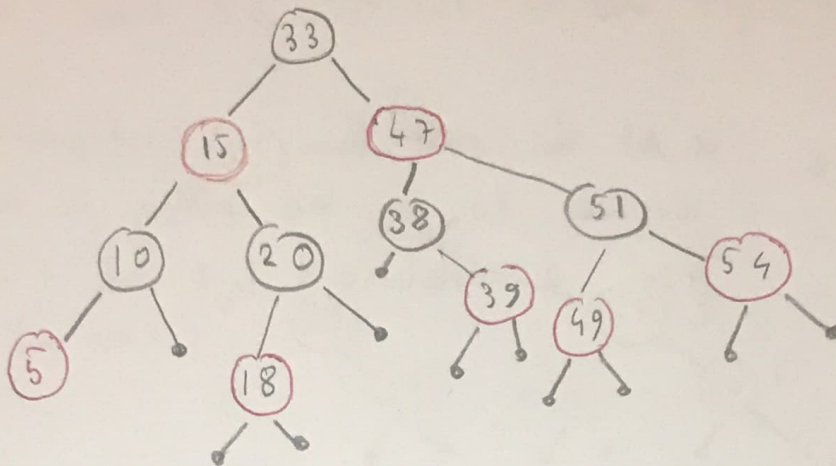
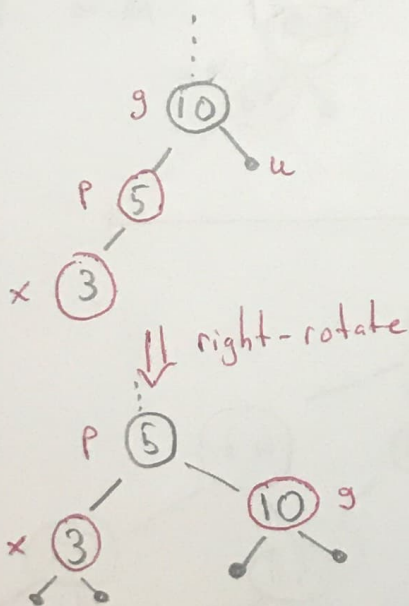


1-



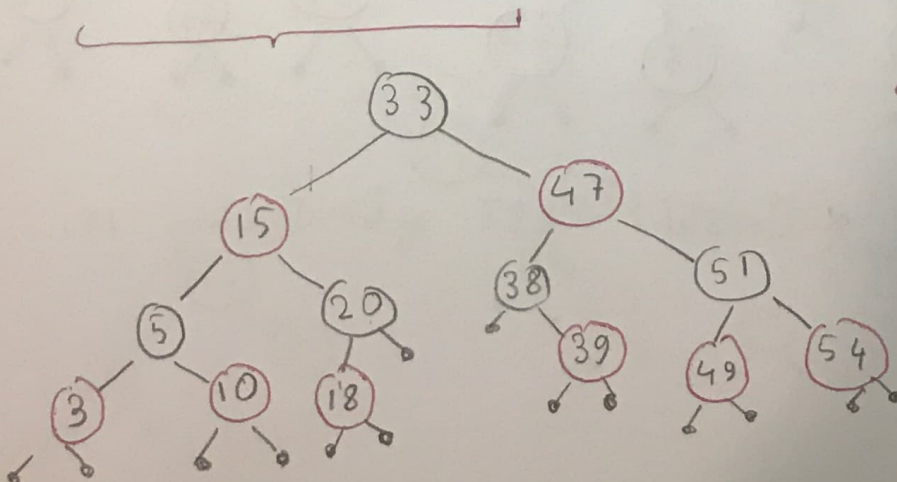
a) Insert 3

If I insert 3, it will be inserted as left child of 5.



At this point there is "red-parent-red-child" problem. To solve that we will apply right-rotate on 5 (p).

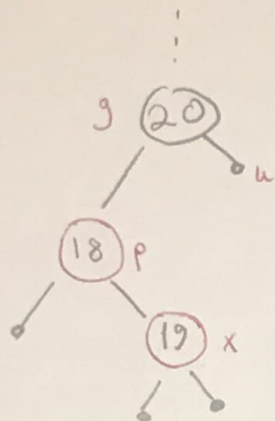
- α Fix color of p black.
- α Fix color of g red.



α Completed RBT after inserting 3.

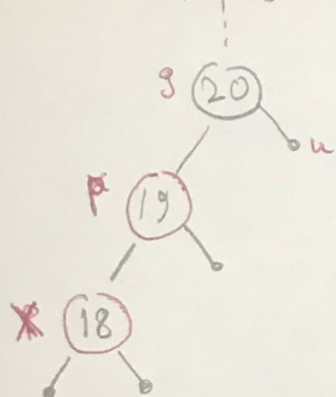
b) Insert 19

If I insert 19, it will be right child of 18.

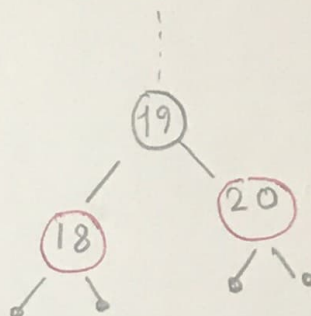


At this point, there is "red-parent-red-child" issue so to fix this problem, we may use 2 rotations. 1 left and 1 right. (Case 2)

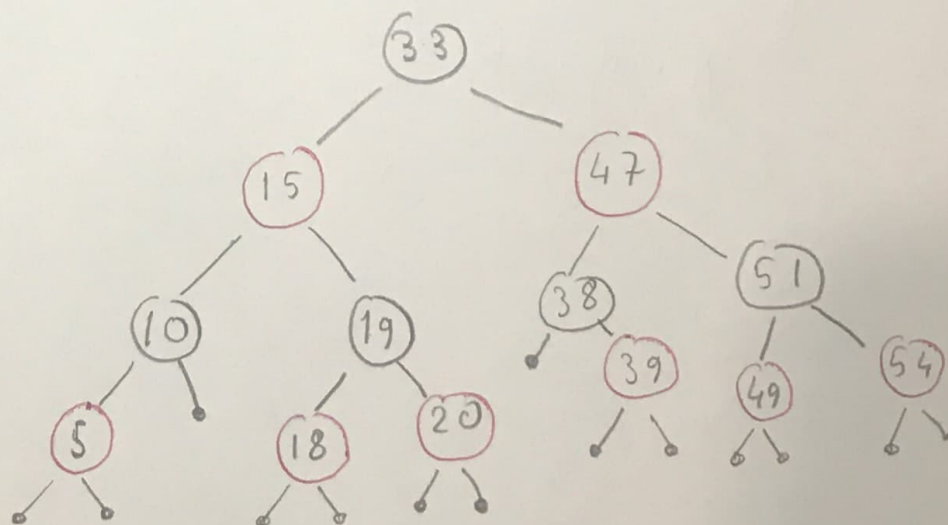
left-rotation



right-rotation
=>
to fix "red-parent-red-child" issue



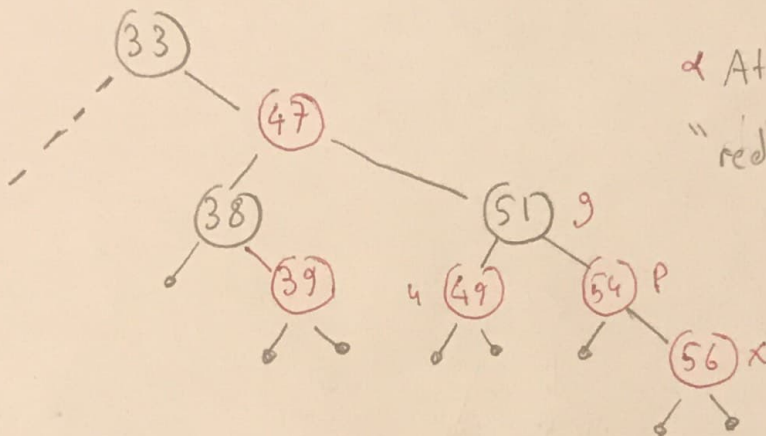
Make g red
Make p black



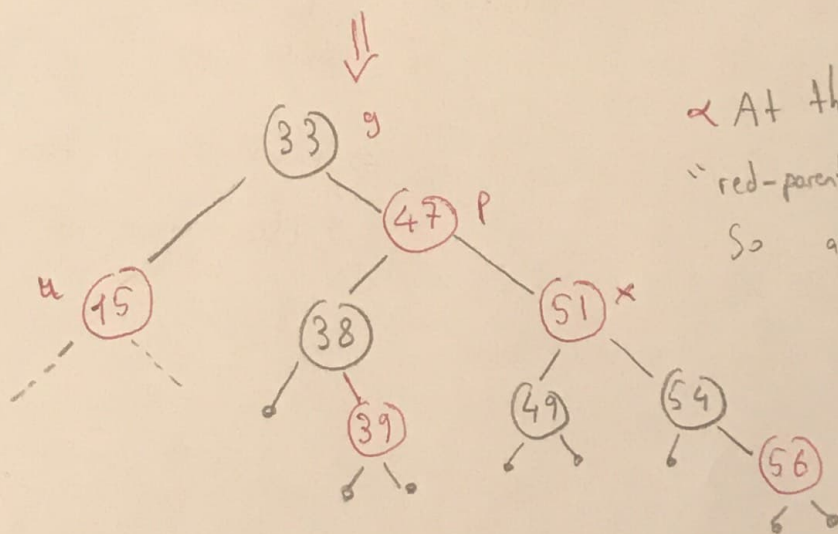
Completed RBT after inserting 19.

c) Insert 56

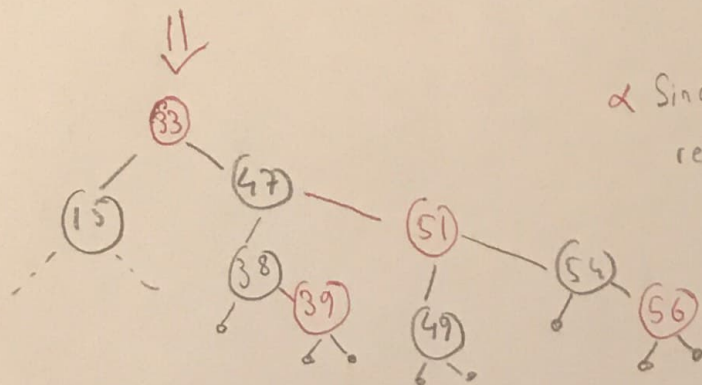
If I insert 56, it will be the right child of 54.



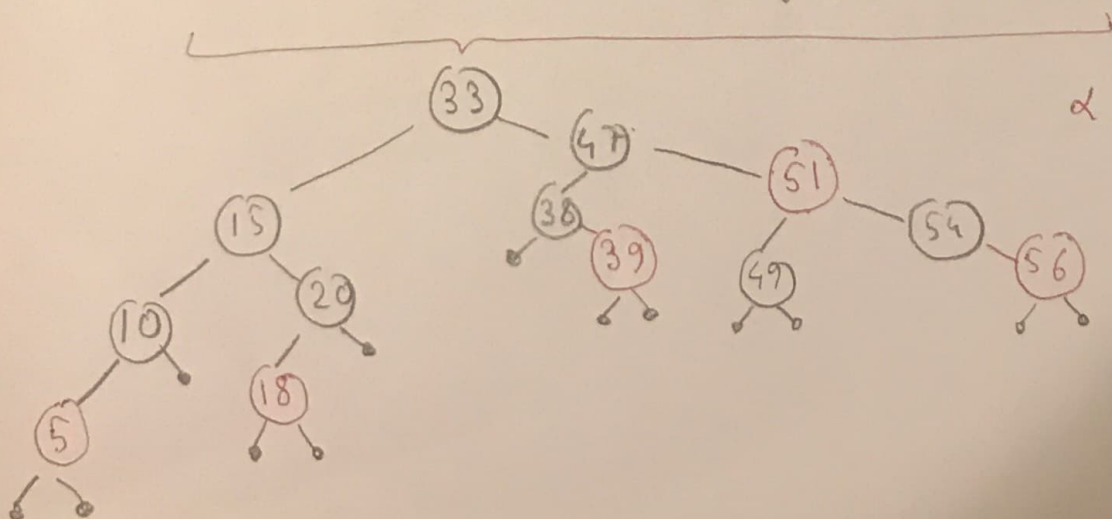
At this point we have "red-parent-red-child" problems. So to fix this, push color of parent and uncle to grandparent. (Case 1)



At this point, there is "red-parent-red-child" issue again. So apply the previous solution.



Since root can not be red, convert it to black.



Completed RBT after inserting 56.

2- a) To avoid from calculating same sub-problems, we can use something similar to recursive. This method traverse over the items and take the maximum value given items that are subject to weight constraint. We may utilize a 2-D matrix to store values.

So the algorithm is the following:

find Optimal Value ($W, n, \text{values-of-items}, \text{weight-of-items}$) {

↓

weight
constraint

↓

total
number of
items

↓

An array
that contains
values of items

↓

An array
that contains
weights of items

for ($i, i \leq n, i++$) \Rightarrow iterate over all items

for ($j, j \leq W, j++$) \Rightarrow iterate over each weight $\leq W$

if (i or j is 0)

put 0 to that cell
in 2-D matrix

else if weight of the item in cell $>$ Weight Constraint
put the value of item in the previous row and
same column

($T[i-1][j]$)

else

find max value of following and put it corresponding cell

$$\max \begin{cases} 1. \text{val}[i] + T[i-1, j - \text{weight of } i] \\ 2. T[i-1][j] \end{cases}$$

At the end, return the 2-D value matrix (T)

}

b) Since we have 2 for loops we should consider them.

Inner loop iterates W times.

Outer loop iterates n times.

$$O(W * n)$$

c) To find best solution, we can construct a matrix. In that matrix columns = W , rows = items

	value	weight	0	1	2	3	4	5.
1)	12	2	0	0	12	12	12	12
2)	10	1	0	10	12	22	22	22
3)	20	3	0	10	12	22	30	32
4)	15	2	0	10	15	25	30	37

- The max value that we can carry with 5 tons is $37 * 0.001$

- Items that we use are:
Item 1 - Item 2 - Item 4