

Architecture of Grammar, day 5

DGfS Summerschool 2024 University of Göttingen

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DGfS Summerschool 2024, University of Göttingen, Aug. 12-16, 2024



Recap

- 'movement' relates multiple independent, semantically compatible DPs
- derivational view of 'movement' reliant on the Y-model creating identical copies / plain multidominance not tenable
- sketch of a model starting with planar LFs

Earlier argument for semantically contentful traces (Sauerland 2004):

- (2) a. *Polly visited every town that is near the lake Erik did Δ .
($\Delta = \text{visit } t$)
- b. Polly visited every town that is near the one Erik did Δ .
($\Delta = \text{visit } t$)

Effability & Economy

Can any conceptual representation that can be articulated in one language also articulated in another if the basic concepts are expressible in both languages?

Counterexample:

- (1) a. Der wievielte Tag des Monats ist heute? (GERMAN)
the how-many-th Tag des Monats is today
- b. *The how manyth day of the month is (it) today?
- c. Which day of the month is today?

But semantic and syntactic conditions exhibit more flexibility.

cases: Scope and Binding Economy

Wide scope blocked first, then becomes available:

- (2) Some boy admires every teacher. Every girl does too.
($\exists \gg \forall, * \forall \gg \exists$)
- (3) Some boy admires every teacher. Some girl does too.
($\exists \gg \forall, \forall \gg \exists$)

Long binding blocked first, then becomes available:

- (4) a. John said that he likes his mother. Bill does too.
b. **Bill said that John likes Bill's mother.*
- (5) a. John said that Mike likes his mother. Bill does too.
b. *Bill said that Mike likes Bill's mother.*

cases: Superiority and Weak Crossover

Pesetsky (1987): Object over subject blocked first, then becomes available.

- (6) a. Who invited who?
b. *Who did who invite?
- (7) a. Which girl invited which boy?
b. Which boy did which girl invite?

Weak crossover (Postal 1993): Binding blocked first, then becomes available.

- (8) *Which lawyer_i did his_i clients hate?
- (9) Which lawyer_i did only his_i clients hate?

cases: Inclusive Plurals

Φ on pronominals (after Sauerland 2002, Percus 2005): Inclusive plural blocked first, then becomes available.

(10) *Everyone who has only a single son invited them.

(11) Everyone who has one or more sons invited them.

Examples with duality:

(12) Everyone who has two sons invited *all/both of them.

(13) Everyone who has two or more sons invited all/*both of them.

Plural address of singular individual (Sauerland 2002, Wang 2023):

(14) [addressing friend] *Wie geht es Ihnen? (How are they?)

(15) [addressing boss] Wie geht es Ihnen?

Accounting for effability

Y-model:

- requires look-ahead to meaning

Meaning first model:

- preference for economical conceptual representations
- closely related to exhaustification

Binding Economy

Longer dependencies are less economical:

(14) $[[\text{the 'J}] [\lambda_x [\text{@} [x [\text{said} [\text{he}_x [\lambda_y [y [\text{like} [\text{the} [\text{his}_y \text{mother}]]]]]]]]]]]]]]]]]]]]]$

(15) $*[[\text{the 'J}] [\lambda_x [\text{@} [x [\text{said} [\text{he}_x [\text{like} [\text{the} [\text{his}_x \text{mother}]]]]]]]]]]]]]]]]]]]$

Relevant alternatives of p for economy calculation are structures q with:

- q must have same meaning as p
- q must only contain the same atoms p contains
- q can have a different pronunciation for p (contra Fox 1998)

Only the most economical structure (i.e. lowest dependency complexity) is licit.

Scope Economy

Fox (1998): lowering of the raised subject for narrow scope

- (16) a. $[[\text{every girl}] \lambda_x [[\text{every teacher}] \lambda_y [x \text{ likes } y]]]$
 b. $[[\text{every teacher}] \lambda_y [[\text{every girl}] [\text{likes } y]]]]$

Sauerland (2018): representations different

- (19) a. $[[\text{every girl}] [\lambda_x [[\text{every teacher}] [\lambda_y [\text{@} [x [\text{admire } y]]]]]]]$
 b. $*[[\text{every teacher}] [\lambda_y [[\text{every girl}] [\lambda_x [\text{@} [x [\text{admire } y]]]]]]]$

Dependency length exponentially contributes to complexity.

- (10) *Dependency Complexity (DC)* Let $\text{var}(\mathbf{A})$ be the set of occurrences of bound⁶ variables in \mathbf{A} and $\text{len}(x)$ be the number of complex concept units between a single occurrence $x \in \text{var}(\mathbf{A})$ and its binder λ_x within \mathbf{A} . Then we define the dependency complexity of \mathbf{A} as:

$$\text{DC}(\mathbf{A}) = \sum_{x \in \text{var}(\mathbf{A})} 2^{\text{len}(x)}$$

Account of superiority

- (17) a. Which girl invited which boy?
b. Which boy did which girl invite?

Singular *which* has uniqueness presupposition, *who* doesn't:

- (18) a. Which girl invited the teacher? – Mary / #None of them / Mary and Sue.
b. Who invited the teacher? – Mary / No one / Mary and Sue.

The uniqueness presuppositions project differently in (17a) and (17b):

- (19) a. Each girl invited exactly one boy.
b. Each boy invited exactly one girl.

(24)	a.			b.			c.		
	Abe Ben Cid			Abe Ben Cid			Abe Ben Cid		
	Ann	*		Ann			Ann	*	
	Bea	*		Bea	*	*	Bea	*	*
	Cel	*		Cel	*		Cel		

Weak Crossover

- (20) a. *?Which lawyer did his clients hate?
b. Which lawyer's clients hate him?
- (21) a. Which lawyer did only his clients hate?
Joe is the only lawyer x such that x 's clients hate Joe
b. Only which lawyer's clients hate him?
Joe is the only lawyer's x such that x 's clients hate x

Strong crossover doesn't improve:

- (22) *Which lawyer does only he hate?

Predicted if **only he** must bind the lower position.

Inclusive plurals

Strong plural interpretation derives from **exh**:

$$(23) \quad \mathbf{exh}_{\{SG(x)\}} PL(x) = \neg \text{atom}(x)$$

Pruning of SG allowed if singular and unpruned plural not possible:

(24) Everyone who has one or more sons invited them / *him.

Pruning of SG not allowed in Percus-style example:

(25) Everyone who has only a single son invited *them / him.

Additional economy measure: Pruning of an alternatives is uneconomical.
Also accounts for politeness plurals (and otherwise unmarked forms).

(26) [addressing boss] Wie geht es Ihnen / *[Dir]? (How are
they/*you)?

Relation to exhaustification

Basic exhaust operator (Fox 2008):

$$(27) \quad \mathbf{exh}_C p = p \wedge \bigwedge \{ \neg q \mid q \in \mathbf{Excl}(C, p) \}$$

$\mathbf{Excl}(C, p)$ is the intersection of all maximal subsets C' of C such that $\{ \neg q \mid q \in C' \}$ consistent with p .

Disjunction case:

$$(28) \quad A \vee B - \text{possible elements of } C: A, B, A \wedge B, \neg A \wedge B, A \wedge \neg B, \dots$$

Katzir's (2006) condition: Elements of C must not be structurally more complex than p .

Innocent exclusion: A not excluded, B not excluded.

$$(29) \quad \mathbf{Excl}(p \vee q, p) = \{ p \wedge q \}$$

Structural complexity (from Katzir 2007)

- (18) SUBSTITUTION SOURCE (first version, to be revised in 41):

Let ϕ be a parse tree. The *substitution source* for ϕ , written as $L(\phi)$, is the lexicon of the language.

- (19) STRUCTURAL COMPLEXITY:

Let ϕ, ψ be parse trees. If we can transform ϕ into ψ by a finite series of deletions, contractions, and replacements of constituents in ϕ with constituents of the same category taken from $L(\phi)$, we will write $\psi \lesssim \phi$. If $\psi \lesssim \phi$ and $\phi \lesssim \psi$ we will write $\phi \sim \psi$. If $\psi \lesssim \phi$ but not $\phi \lesssim \psi$ we will write $\psi < \phi$.

- (20) STRUCTURAL ALTERNATIVES:

Let ϕ be a parse tree. The set of *structural alternatives* for ϕ , written as $A_{str}(\phi)$, is defined as $A_{str}(\phi) := \{\phi' : \phi' \lesssim \phi\}$.

Blocking exclusion of simpler structures

Assume $\psi < \phi$ and $\psi \equiv \phi$: Blocking of too complex structures could be predicted:

(30) if $\psi \in \mathbf{Excl}(\phi, C)$, then $\mathbf{exh}(\phi) = \phi \wedge \neg\psi \wedge \dots = \#$

Magri's (2009) constraint: If an alternative ψ is contextually equivalent to the assertion ϕ it cannot be pruned [restated].

(31) $\#$ Some Italians live in a beautiful country.

Our situation: ϕ asserted, ψ equivalent:

(32) Magri's constraint applies also to exclusion of alternatives of ϕ that are structurally distinct.