TDTS08: Lab Report

Lab 4: VLIW Processors

Name	PIN	Email
Alexander Yngve	930320-6651	aleyn573@student.liu.se
Pål Kastman	851212-7575	palka285@student.liu.se

Contents

1	Introduction	3
2	2.2 Dependencies	3 3 4
3	Result	4
4	Discussion	7
A	Graph file	8
В	VLIW files	9

1 Introduction

The purpose of this lab was to convert normal sequential code to VLIW instructions, so that we would get a greater performance. Basically we do what the VLIW compiler does during compilation time.

2 Method

The approach for this lab was the following:

- 1. Choose a basic block, and disassemble the block.
- 2. Find dependencies between the instruction in the block.
- 3. We pack the instruction into VLIWs.

2.1 Basic Block

To view all the basic blocks within the program (go.ss) the following command is issued:

```
vliwc /home/TDTS08/bin/go.ss | sort -nr +3 -4 | less
```

A basic block with at least 15 instructions is needed for the lab.

The next step is to view the disassemble the program to view the code for the basic block. This is done with the command below:

```
sslittle -na-sstrix-objdump -d /home/TDTS08/bin/go.ss | less
```

Which gives output similar to this:

```
lw $3,16($29)
41 b528:
                28 00 00 00
41 b52c:
                10 00 03 1d
41 b 530:
                28 00 00 00
                                   lw $4, -31444($28)
41 b 5 3 4:
                2c 85 04 1c
41 b 538:
                a2 00 00 00
                                   lui $6,4100
41 b53c:
                04 10 06
                           00
41 b 5 40:
                43 00 00
                           00
                                   addiu \$6,\$6,-6208
41 b 5 4 4:
                c0 e7 06 06
```

2.2 Dependencies

To pack the instructions into Very Long Instruction Words, the dependencies between the instructions must be resolved. The dependencies of interest are the true data dependencies, read after write, where the output of one instruction is required as an input to one of the following instructions.

Output dependencies (write after write) and anti-dependencies (write after read) are not considered since the are artificial dependencies which can be resolved in preprocessing.

The true data dependencies should be visualized in a directed graph where the address of each instruction is a node. The edges between the nodes symbolizes a dependency. This graph should also be described in a textual representation which will be used by the *vliwc* program. An example of the graph file looks like this:

```
\begin{array}{c} 0 \! \times \! 000000001 \\ 0 \! \times \! 000000001 & 0 \! \times \! 000000002 \end{array}
```

This file describes two instructions, 0x00000001 which is independent, and 0x00000002 which is dependent on 0x00000001.

2.3 VLIW

The last step is to pack the sequential instructions obtained in section 2.1 into VLIW format with the help of the dependency graph from section 2.2.

The text format for the VLIW file begins with a line which specifies alu_no , mul_no , fpu_no and bau_no - how many ALUs, MULs, FPUs and BAUs the VLIW processor will have.

The next lines are the Very Long Instruction Words in the form of addresses to the sequential instructions. The first alu_no instructions will be ALU instructions, the next mul_no instructions will be MUL instructions and so on.

An example VLIW file can look like this:

3 Result

The chosen basic block is number 2292 and consists of 16 instructions. The disassembly is shown below:

```
lw $3,16($29)
41 b528:
                28 00 00 00
41 \, b52 c:
                10 00 03 1d
                                    lw $4, -31444($28)
41 b530:
                28 00 00
                           00
41 b 5 3 4:
                2c 85 04
                           1c
41b538:
                a2 00 00 00
                                    lui $6,4100
41 b53c:
                04 10 06 00
41 b 540:
                43 00 00 00
                                    addiu \$6,\$6,-6208
41 b 5 4 4:
                c0 e7 06 06
                a2 00 00 00
41 b 5 48:
                                    lui $2,4097
41 \, b54 c:
                01 10 02 00
```

41 b550:	28	00	00	00	lw \$2,5960(\$2)
41b554:	48	17	02	02	
41 b558:	a2	00	00	00	lui \$5,4100
41 b55c:	04	10	05	00	
41b560:	43	00	00	00	addiu \$5,\$5,30352
41 b564:	90	76	05	05	
41b568:	55	00	00	00	sll \$3,\$3,0x2
41b56c:	02	03	03	00	, ,
41 b570:	42	00	00	00	addu \$6,\$3,\$6
41 b574:	00	06	06	03	
41b578:	28	00	00	00	lw \$7,0(\$6)
41 b57c:	00	00	07	06	, , ,
41b580:	42	00	00	00	addu \$5,\$3,\$5
41 b584:	00	05	05	03	, ,
41b588:	42	00	00	00	addu \$2,\$16,\$2
41b58c:	00	02	02	10	,
41b590:	42	00	00	00	addu \$2,\$2,\$7
41b594:	00	02	07	02	, ,
41b598:	34	00	00	00	sw \$2,0(\$6)
41 b59c:	00	00	02	06	, , ,
41 b5a0:	01	00	00	00	$j 41b758 < try_connect + 0x750 >$
41b5a4:	d6	6d	10	00	

Dependency analysis of the above code results in the graph shown in figure 1. The graph file is also included in appndix A. Note that the unconditional jump instruction at address 41b5a0 doesn't have any data dependencies but still needs to be run last due to procedural dependencies.

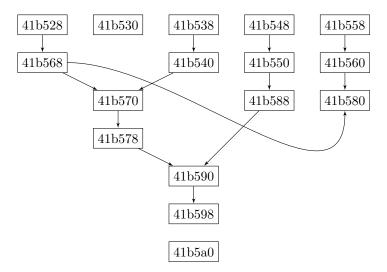


Figure 1 – Data dependencies between the instructions.

We first choose to make our architecture with no limit, this design can be seen in 1.

Table 1

	ALU		MUL	FPU			$\mathbf{B}\mathbf{A}\mathbf{U}$		
					41b528	41b530	41b538	41b548	41b558
41b540	41b568	41b560			41b560				
41b570	41b588	41b580							
					41b578				
41b590									
					41b598	41b5a0			

We then tried to reduce some of the units without loosing any performance, we first by tried by removing one of the BAUs (figure 2), since those are the most expensive units. In table 6 we can see that this design (VLIW2) doesn't perform any worse than the first design (VLIW1), so we continued to reduce more units in search of an even better design for this block.

Table 2

	ALU		MUL	FPU		BAU	
					41b528	41b558	41b538
41b568	41b540	41b560			41b548	41b530	
41b570	41b580				41b550		
41b588					41b578		
41b590							
					41b598	41b5a0	

Table 3

Al	ALU		FPU	BA	\U
				41b528	41b538
41b568	41b540			41b548	41b558
41b570	41b560			41b550	41b530
41b588	41b580			41b578	
41b590					
				41b598	41b5a0

Table 4

Al	LU	MUL	FPU	BAU
				41b528
				41b538
41b568	41b540			41b558
41b570	41b560			41b548
41b580				41b550
41b588				41b578
41b590				41b530
				41b598
				41b5a0

Table 5

Pricelis	t
ALU cost	2
MUL cost	16
FPU cost	32
BAU cost	100

Table 6

Design	VLIW1	VLIW2	VLIW3	VLIW4
No. ALU	3	3	2	2
No. MUL	1	1	1	1
No. FPU	1	1	1	1
No. BAU	5	3	2	1
Total Cost	554	354	252	152
No. Cycles	6	6	6	9
Cost per. ratio	3324	2124	1512	1368

4 Discussion

A Graph file

 $0 \times 0041 \text{b} 528$

 $0 \times 0041 \text{b} 530$

 $0 \times 0041 \text{b} 538$

 $0 \times 0041 \text{b} 548$

 $0\,\mathrm{x}0041\,\mathrm{b}558$

 $0\,x0041\,b528\ 0\,x0041\,b568$

 $0 \times 0041 \text{b} 538 \ 0 \times 0041 \text{b} 540$

 $0 \times 0041 b548 \quad 0 \times 0041 b550$

 $0\,x0041\,b558\ 0\,x0041\,b560$

 $0 \times 0041 b568 \ 0 \times 0041 b570$

 $0 \times 0041 b540 \ 0 \times 0041 b570$

 $0 \times 0041 b568 0 \times 0041 b580$

 $0 \times 0041 b560 0 \times 0041 b580$

 $0\,x0041\,b570\ 0\,x0041\,b578$

 $0\,x0041\,b550\ 0\,x0041\,b588$

 $0 \times 0041 b578 \quad 0 \times 0041 b590$

 $0 \times 0041 b588 \quad 0 \times 0041 b590$

 $0 \times 0041 b590 \quad 0 \times 0041 b598$

 $0 \times 0041 \text{b} 5 \text{a} 0$

B VLIW files

Listing 1 – VLIW 1

3			1	1	5				
nop	nop	nop	nop	nop	$0 \times 0041 b528$	$0 \times 0041 b530$	$0 \times 0041 b538$	$0 \times 0041 b548$	$0 \times 0041 b558$
$0 \times 0041 b540$	$0 \times 0041 b568$	$0 \times 0041 b560$	nop	nop	$0 \times 0041 b550$	nop	nop	nop	
nop									
$0 \times 0041 b570$	$0 \times 0041 b588$	$0 \times 0041 b580$	nop	nop	nop	nop	nop	nop	
nop									
nop	nop	nop	nop	nop	$0 \times 0041 b578$	nop	nop	nop	
nop									
$0 \times 0041 b590$	nop	nop	nop	nop	nop	nop	nop	nop	
nop									
nop	nop	nop	nop	nop	$0 \times 0041 b598$	$0 \times 0041 b5 a0$	nop	nop	
nop									

$\textbf{Listing 2} - VLIW \ 2$

3			1	1	3		
nop	nop	nop	nop	nop	$0 \times 0041 b528$	$0 \times 0041 b558$	$0 \times 0041 b538$
$0 \times 0041 b568$	$0 \times 0041 b540$	$0 \times 0041 b560$	nop	nop	$0 \times 0041 b548$	$0 \times 0041 b530$	nop
$0 \times 0041 b570$	nop	$0 \times 0041 b580$	nop	nop	$0 \times 0041 b550$	nop	nop
$0 \times 0041 b588$	nop	nop	nop	nop	$0 \times 0041 b578$	nop	nop
$0 \times 0041 b590$	nop	nop	nop	nop	nop	nop	nop
nop	nop	nop	nop	nop	$0 \times 0041 b598$	$0 \times 0041 b5a0$	nop

Listing 3 – VLIW 3

2		1	1	2	
nop	nop	nop	nop	$0 \times 0041 b528$	$0 \times 0041 b538$
$0 \times 0041 b568$	$0 \times 0041 b540$	nop	nop	$0 \times 0041 b548$	$0 \times 0041 \mathrm{b}558$
$0 \times 0041 b570$	$0 \times 0041 b560$	nop	nop	$0 \times 0041 b550$	$0 \times 0041 b530$
$0 \times 0041 b588$	$0 \times 0041 b580$	nop	nop	$0 \times 0041 b578$	nop
$0 \times 0041 b590$	nop	nop	nop	nop	nop
non	non	non	non	$0 \times 0041 b598$	$0 \times 0.041 \text{ b} 5 \text{ a} 0$

Listing 4 – VLIW 4