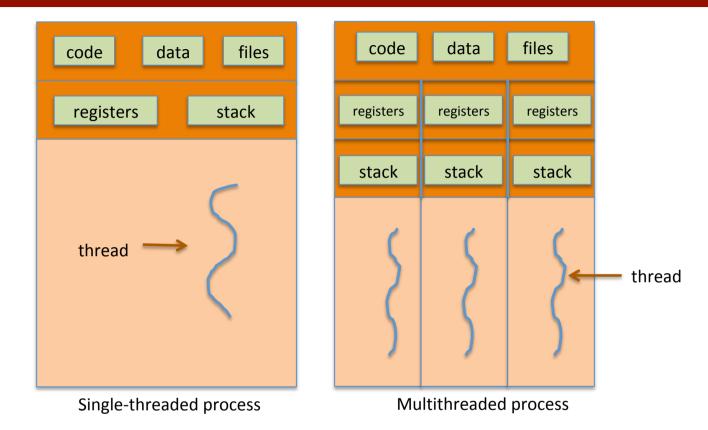


POSIX Threads

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Lecture 13
COMP304 - Operating Systems (OS)

Single vs Multithreaded Process



- A thread has an ID, a program counter, a register set, and a stack
- Shares the code section, data section and OS resources (e.g. files)
 with other threads within the same process

POSIX Threads API (Pthreads)

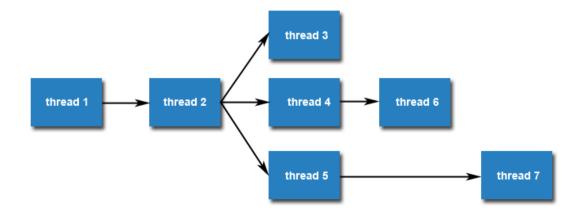
- Pthreads is the POSIX (Portable Operating System Interface for Unix) Thread Library
 - a POSIX standard (IEEE 1003.1c) API for thread creation and synchronization.
- API specifies behavior of the thread library, implementation is up to development of the library.
- Common in UNIX operating systems
 - Solaris, Linux, Mac OS X.
- Is this a Kernel or User Level Thread API?

POSIX Threads API

- Functions of pthreads API provide:
 - Thread management:
 - Creation/termination of threads
 - Set/Query thread attributes
 - Mutexes, semaphores
 - Condition variables
- All identifiers in the threads library begin with pthread

Creating Threads

- Initially, a main() program comprises a single, default thread. All
 other threads must be explicitly created by the programmer.
- pthread_create
 - creates a new thread and makes it executable.
- The maximum number of threads that may be created by a process is implementation dependent.
 - Programs that attempt to exceed the limit can fail or produce wrong results.
- Threads can create other threads (but there is no hierarchy)



Pthread_create()

Forking Pthreads

- thread_id is the thread id or handle (used to halt, etc.)
- thread_attribute various attributes
 - standard default values obtained by passing a NULL pointer
- thread_func the function to be run (takes and returns void*)
- func_arg an argument can be passed to thread_fun when it starts
- errorcode will be set to nonzero if the create operation fails

Thread Creation

- Each thread executes a specific function, thread func
 - For the program to perform different work in different threads, the arguments passed at thread creation distinguish the thread's "id" and any other unique features of the thread.
- After a thread is created, various attributes of it can be set
 - Priority of the thread
 - Stack size
 - Its scheduling policy

Simple Example

```
#include <stdio.h>
#include <stdib.h>
#include <pthread.h>

declares the various Pthreads

int main() {

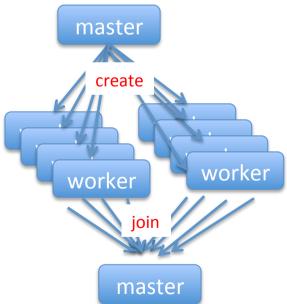
pthread_t threads[8];
int tn;

for(tn=0; tn<8; tn++) {
   pthread_create(&threads[tn], NULL, ParFun, NULL);
}

for(tn=0; tn<8; tn++) {
   pthread_join(threads[tn], NULL);
}

return 0;
}</pre>
```

- This code creates 8 threads that execute the function "ParFun".
- What happens to the master while workers are executing?
 - Does master become one of the workers?



Thread Termination

Pthread_exit()

 A thread returns from its starting routine by default, similar to a process terminating when it reaches to end of main

Pthread_cancel()

Thread is cancelled by another thread via this call

Pthread_join(...)

- From Unix specification: "suspends execution of the calling thread until the target thread terminates, unless the target thread has already terminated."
- After the last thread in a process terminates, the process terminates by calling exit()
- If the entire process is terminated, then all its threads will terminate

"Hello World"

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
int thread count=16; //accessible by all threads
void *hello(void* rank); //thread function
int main() {
 pthread t threads[16]; //thread handles
  int tn;
  for(tn=0; tn<16; tn++) {
    pthread create(&threads[tn], NULL, hello,(void*)tn);
  for(tn=0; tn<16; tn++) {
    pthread join(threads[tn], NULL);
                                                          Arguments to
                                        Start function
                                                          function
                                        to execute
  return 0;
                                                          e.g. Thread ID
        Wait for thread completion
```

"Hello World"

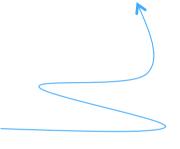
```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
int thread count=16; //global var, accessible by all threads
void *hello(void* rank); //thread function
int main() {
  pthread_t threads[16]; //thread handles
  int tn;
  for(tn=0; tn<16; tn++) {
    pthread create(&threads[tn], NULL, hello,(void*)tn);
  pthread exit(NULL);
                                    Prevents the process to die before its
                                    threads are done!
  return 0;
                                    main()/master will block and be kept
                                    alive to support the threads it created
```

"Hello World"

- By using thread id, different execution for each thread is possible
- How is possible for a thread to access thread_count?

Compiling a Pthread Program

gcc -o pth_hello pth_hello.c -lpthread



Link the Pthreads library

The Pthreads API is only available on POSIX systems — Linux, MacOS X, Solaris, HPUX, ...

Because threads share resources:

- Changes made by one thread to shared system resources (such as closing a file) will be seen by all other threads.
- Two pointers having the same value point to the same data.
 - Because of the shared address space
- Reading and writing to the same memory locations is possible, and therefore requires explicit synchronization by the programmer.

Shared Data

- Variables declared outside of 'main' are global variables
 - Those variables are shared by all threads
- Object allocated on the heap may be shared if pointer is passed as an argument to the thread function

```
char *message = "Hello World!\n";
pthread_create( &thread1, NULL, (void*)&print_fun, (void*)message);
```

- Variables on the stack are private (locally defined variables)
 - Passing pointer to these around to other threads can cause problems

Thread Synchronization

- Need to protect the shared data and synchronize threads
- Pthread provides several ways to synchronize threads:
- Mutexes (Locks)
- Semaphores
- Condition Variables
- Barriers
 - Synchronizing the threads to make sure that they all are at the same point in a program is called a barrier.

Mutex (Locks) in Pthreads

```
#include <pthread.h>
pthread mutex t mymutex = PTHREAD MUTEX INITIALIZER;
//create a mutex
pthread mutex init(&mymutex, NULL);
//use it
pthread mutex lock(mymutex);
                                   pthread mutex lock(&our lock);
//mutually excluded code region
                                   counter++
pthread mutex unlock(mymutex);
                                   pthread mutex unlock(&our lock);
//destroy a mutex
pthread mutex destroy(mymutex);
```

Condition Variables

- While mutexes implement synchronization by controlling thread access to data, condition variables allow threads to synchronize based upon the actual value of data.
 - Without condition variables, the programmer would need to have threads continually polling to check (possibly in a critical section) if the condition is met.
- A condition variable is always used in conjunction with a mutex lock.
- Set/query condition variable attributes

See an example usage:

https://computing.llnl.gov/tutorials/pthreads/#ConditionVariables

Condition Variable Code Example

See handout 13

- The main thread creates three threads.
- Two of those threads increment a "count" variable, while the third thread watches the value of "count".
- 3. When "count" reaches a predefined limit, the waiting thread is signaled by one of the incrementing threads.
- 4. The waiting thread "awakens" and then modifies count. The program continues until the incrementing threads reach TCOUNT.
- The main program prints the final value of count.

Reference link:

https://computing.llnl.gov/tutorials/pthreads/samples/condvar.c

Condition Variables (cont.)

Main Thread

- Declare and initialize global data/variables which require synchronization (such as "count")
- Declare and initialize a condition variable object
- Declare and initialize an associated mutex
- Create threads A and B to do work

Thread A

- Do work up to the point where a certain condition must occur (such as "count" must reach a specified value)
- Lock associated mutex and check value of a global variable
- Call pthread_cond_wait() to perform a blocking wait for signal from Thread-B. Note that a call to pthread_cond_wait() automatically and atomically unlocks the associated mutex variable so that it can be used by Thread-B.
- When signaled, wake up. Mutex is automatically and atomically locked.
- Explicitly unlock mutex
- Continue

Thread B

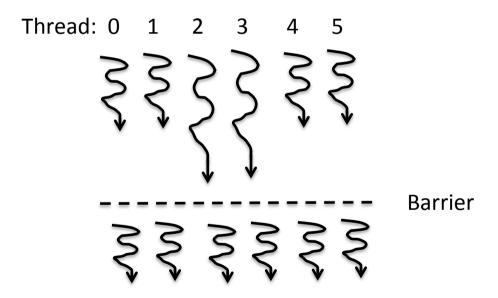
- Do work
- Lock associated mutex
- Change the value of the global variable that Thread-A is waiting upon.
- Check value of the global Thread-A wait variable. If it fulfills the desired condition, signal Thread-A.
- Unlock mutex.
- Continue

Semaphores in Pthreads

- Functions defined in semaphore.h:
- A semaphore is represented by a sem_t type.
- sem_init: for initializing semaphore
- sem_wait: for waiting on a semaphore
- sem_post: for signaling on a semaphore
- sem_destroy: for deallocating a semaphore if you no longer need it

Barriers

- Synchronizing the threads to make sure that they all are at the same point in a program is called a barrier.
- No thread can cross the barrier until all the threads have reached it.



Even though threads 2 and 3 reached barrier, they will wait for others to arrive.

Then all threads cross the barrier point together.

Barriers in Pthreads

• To (dynamically) initialize a barrier, use code similar to this (which sets the number of threads to 3):

```
pthread_barrier_t b;
pthread_barrier_init(&b,NULL,3);
```

- The second argument specifies an object attribute; using NULL yields the default attributes.
- To wait at a barrier, a process executes:

```
pthread_barrier_wait(&b);
```

To destroy a barrier

```
pthread_barrier_destroy(&b);
```

Thread Scheduling

- Threads can be scheduled by the operating system and run as independent entities
- Many-to-one and many-to-many models, thread library schedules user-level threads
 - Known as process-contention scope (PCS) since scheduling competition is within the process
- Kernel thread scheduled onto available CPU is systemcontention scope (SCS) – competition among all threads in system

Pthread Scheduling

- API allows specifying either PCS or SCS during thread creation
 - PTHREAD_SCOPE_PROCESS schedules threads using PCS scheduling
 - PTHREAD_SCOPE_SYSTEM schedules threads using SCS scheduling
- Can be limited by OS Linux and Mac OS X only allow PTHREAD SCOPE SYSTEM

Pthread Scheduling API

```
#include <pthread.h>
#include <stdio.h>
#define NUM THREADS 5
int main(int argc, char *argv[]) {
   int i, scope, policy;
  pthread t tid[NUM THREADS];
  pthread attr t attr;
   /* get the default attributes */
  pthread attr init(&attr);
   /* first inquire on the current scope */
   if (pthread attr getscope(&attr, &scope) != 0)
      fprintf(stderr, "Unable to get scheduling scope\n");
   else {
      if (scope == PTHREAD SCOPE PROCESS)
        printf("PTHREAD SCOPE PROCESS");
      else if (scope == PTHREAD SCOPE SYSTEM)
        printf("PTHREAD SCOPE SYSTEM");
      else
         fprintf(stderr, "Illegal scope value.\n");
```

Pthread Scheduling API (cont.)

```
/* get the current scheduling policy */
  if (pthread_attr_getschedpolicy(&attr, &policy) != 0)
    fprintf(stderr, "Unable to get policy.\n");
  else {
    if (policy == SCHED_OTHER) printf("SCHED OTHER\n");
    else if (policy == SCHED_RR) printf("SCHED RR\n");
    else if (policy == SCHED_FIFO) printf("SCHED FIFO\n");
  }

/* set the scheduling policy - FIFO, RR, or OTHER */
  if (pthread_attr_setschedpolicy(&attr, SCHED_FIFO) != 0)
    fprintf(stderr, "Unable to set policy.\n");
...
```

Pthread Scheduling API (cont.)

```
/* set the scheduling algorithm to PCS or SCS */
   pthread attr setscope(&attr, PTHREAD SCOPE SYSTEM);
   /* create the threads */
   for (i = 0; i < NUM THREADS; i++)
      pthread create(&tid[i],&attr,run method,NULL);
   /* now join on each thread */
   for (i = 0; i < NUM THREADS; i++)
      pthread join(tid[i], NULL);
/* Each thread will begin control in this function */
void *run method(void *param)
   /* do some work ... */
   pthread exit(0);
```

Summary of Pthreads

- Pthreads are user-level threads for POSIX systems
 - Familiar language for most programmers, particularly for systems people
 - Ability to shared data is convenient
 - Various supports for synchronization
- Pthread Tutorial
 - https://computing.llnl.gov/tutorials/pthreads/
- Reading from Book
 - **4.1, 4.3, 4.6, 4.7**
- Acknowledgements
 - These slides are adapted from
 - Öznur Özkasap (Koç University)
 - Operating System and Concepts (9th edition) Wiley
 - https://computing.llnl.gov/tutorials/pthreads/