

# CSCI 6333/6315 Database Systems

Spring 2020

## ASSIGNMENT 4: Data Storage and Querying, Transaction Management

All answers shall be typed using a word processor and some drawing utilities. A cover page shall be prepared with course title, homework number, submitted date and time, and contact info including your email address.

**Due: Midnight, Tuesday, April 28, 2020**

The total score of this assignment is 110. Each part of the problem has 10 points.

Problem 1. This problem has two parts:

- a. Construct a  $B^+$ -tree for the following key values:  
2, 3, 5, 7, 11, 17, 19, 23, 29, 31.  
Assume that the number of pointers that will fit in one internal node is 4 and each leaf node can store 3 key values.
- b. After the  $B^+$ -tree is constructed for Part a, show the final tree after the following operations:
  - b.1. Insert 9
  - b.2. Insert 10
  - b.3. Insert 8
  - b.4. Delete 23
  - b.5. Delete 19

Problem 2. This problem has two parts:

- a. Suppose that we are using extendable hashing on a file that contains records with the following search key values:  
2, 3, 5, 7, 11, 17, 19, 23, 29, 31.  
Show the extendable hash structure for this file if the hash function is  $h(x) = x \% 8$ , and each bucket can hold three records.
- b. After Completing Part a, show the extendable hash structures after the following operations:
  - b.1. Delete 11
  - b.2. Delete 31
  - b.3. Insert 1
  - b.4. Insert 15

Problem 3. Consider the following SQL query for our University Database:

```
select T.dept_name
from department as T, department as S
where T.budget > S.budget and S.building = "MAG"
```

Write an efficient relational-algebra expression that is equivalent to this query. Justify your choice.

Problem 4. Let the relations  $r_1(A, B, C)$  and  $r_2(C, D, E)$  have the following properties:  $r_1$  has 20,000 tuples,  $r_2$  has 45,000 tuples, 25 tuples of  $r_1$  fit on one block, and 30 tuples of  $r_2$  fit on one block. Estimate the number of block accesses required, using each of the following join strategies for  $r_1 \bowtie r_2$ :

- a. Nested-loop join
- b. Block nested-loop join
- c. Merge join
- d. Hash join

Problem 5. Show that the following equivalences hold.

- c.  $E_1 \bowtie_{\theta} (E_2 - E_3) = (E_1 \bowtie_{\theta} E_2 - E_1 \bowtie_{\theta} E_3)$
- d.  $\sigma_{\theta_1 \wedge \theta_2} (E_1 \bowtie_{\theta_3} E_2) = \sigma_{\theta_1} (E_1 \bowtie_{\theta_3} (\sigma_{\theta_2} (E_2)))$ , where  $\theta_2$  involves only attributes from  $E_2$ .