

# Instruction Cycle

## Lecture 04

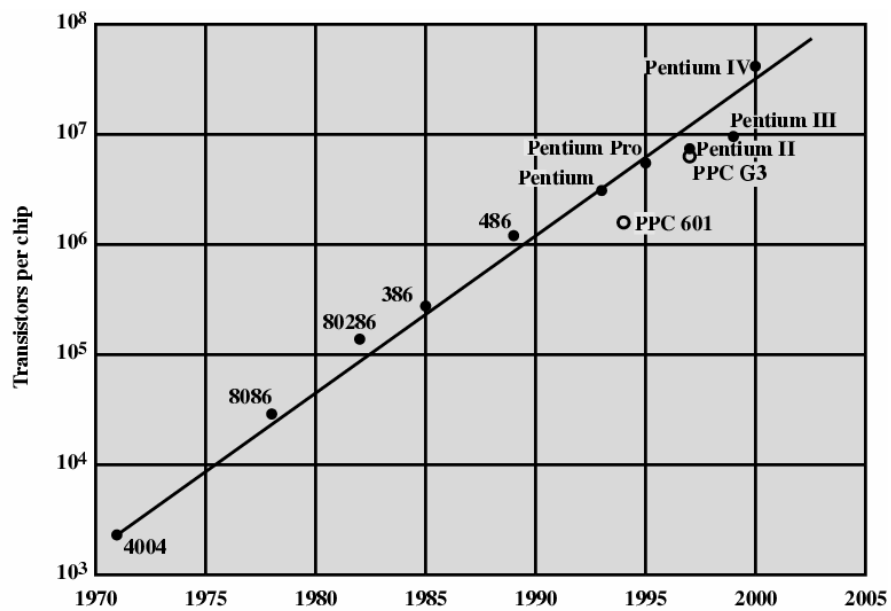
### Generations of Computer

- Vacuum tube - 1946-1957
- Transistor - 1958-1964
- Small scale integration - 1965 on
  - Up to 100 devices on a chip
- Medium scale integration - to 1971
  - 100-3,000 devices on a chip
- Large scale integration - 1971-1977
  - 3,000 - 100,000 devices on a chip
- Very large scale integration - 1978 -1991
  - 100,000 - 100,000,000 devices on a chip
- Ultra large scale integration – 1991 -
  - Over 100,000,000 devices on a chip

### Moore's Law

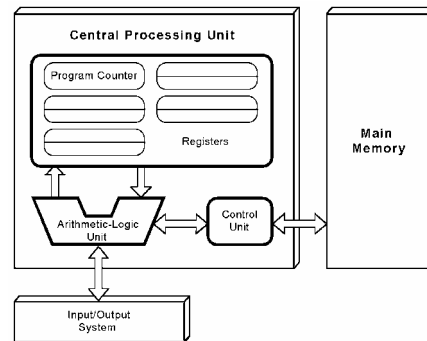
- Increased density of components on chip
- Moore's Law (1965) co-founder of Intel
  - Gordon Moore, Intel Co-founder
  - "The density of transistors in an integrated circuit will double every two year."
- Since 1970's development has slowed a little
  - Number of transistors doubles every 18 months
  - Cost of a chip has remained almost unchanged
- Higher packing density means shorter electrical paths, giving higher performance

### Growth in CPU Transistor Count



## The von Neumann Model

- The computers employ a fetch-decode-execute cycle to run programs as follows:
  - The control unit **fetches** the **next instruction** from **memory** using the **program counter** to determine where the instruction is located
  - The **instruction** is **decoded** into a language that the **ALU** can understand.
  - Any **data operands** required to execute the instruction are **fetches** from memory and placed into **registers** within the CPU
  - The ALU executes the instruction and places results in registers or memory



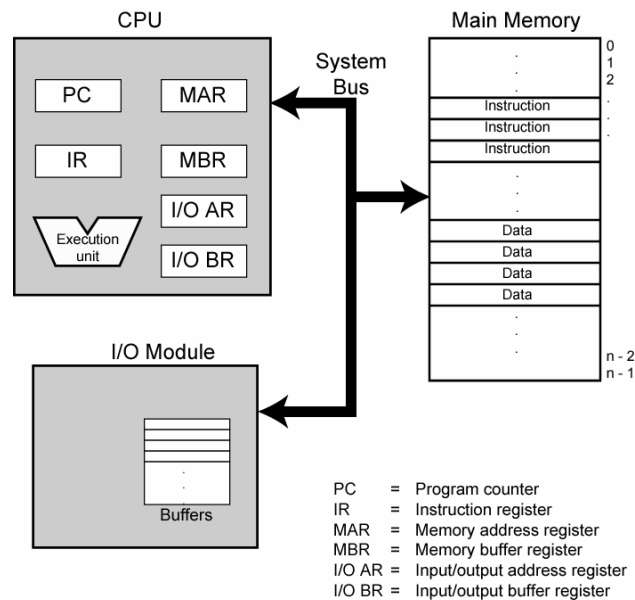
### Computer Components

- Almost all computer designs are based on concept developed by John von Neumann at the institute for Advanced Studies (IAS)
- The key concepts of von Neumann Architecture are
  - Data and instructions are stored in a single read-write memory
  - The contents of memory are addressable by location without regard the type of data contained therein
  - Execution occurs in a sequential fashion- unless modified from one instruction to the next
- A small set of basic logic components that can be combined in various ways to store binary data and to perform arithmetic and logic operation on that data

## What is a program?

- A general purpose configuration of hardware is to perform various functions on data depending on control signal applied to the hardware
  - The system accepts data and produces results
  - Rewiring of hardware is not required, but the set of control signal is applied “programming?”
- A sequence of steps
  - For each step, an arithmetic or logical operation is done on some data
  - For each operation, a different set of control signals is needed

## Computer Components: Top Level View

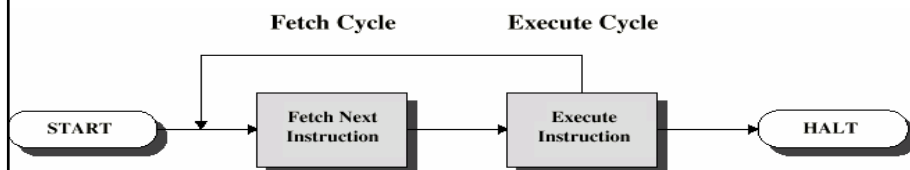


## Instruction Cycle

- Basic function performed by a computer is execution of a program
  - consists of set of instructions stored in the memory
  - Processor executes the instructions specified in the program
- Instruction processing consists of two steps:
  - Processor reads – Fetches from memory one at a time
  - Execute each instruction
- Program execution – repeating the steps over and over
  - Processing required for a single instruction is called “instruction cycle”

## Instruction fetch and execute

- At the beginning of each instruction cycle
  - Processor fetches an instruction from the memory
    - A register “PC”- program counter holds the address of the instruction to be fetched next
    - Processor always increment the PC after each instruction unless told otherwise
  - The fetched instruction is loaded into a register located in the processor – “the instruction Register (IR)”
  - The instruction contains bits – to specify the action the processor is to take



### Fetch Cycle

- Program Counter (PC) holds address of next instruction to fetch
- Processor fetches instruction from memory location pointed to by PC
- Increment PC
  - Unless told otherwise
- Instruction loaded into Instruction Register (IR)
- Processor interprets instruction and performs required actions

### Execute Cycle

- Processor-memory
  - data transfer between CPU and main memory
- Processor I/O
  - Data transfer between CPU and I/O module
- Data processing
  - Some arithmetic or logical operation on data
- Control
  - An instruction may specifies the alteration of sequence of operations
    - e.g. jump
- An instruction's execution may involve a Combination of these actions

### Example

Add the contents of the memory at address 940 to the contents of the memory at address 941 and store the result in the next location



### Partial List of Opcodes:

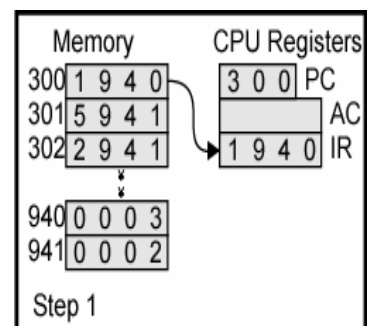
0001 = (1h) = Load AC from Memory

0010 = (2h) = Store AC to memory

0101 = (5h) = Add to AC from Memory

## 1. Program Execution

- PC contains 300 – the address of the first instruction
  - This instruction (the value 1940h) is loaded into the IR
  - PC is incremented (301)
- The process involve the use of
  - Memory Address Register (MAR)
  - Memory Buffer Register (MBR)

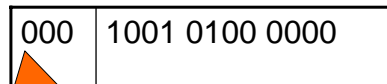


## 2. Program Execution

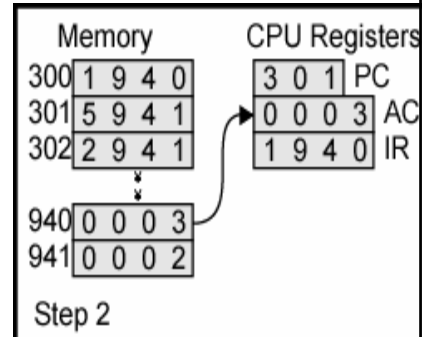
- The first nibble in the IR indicate that the AC is to be loaded

- Remaining 12 bits specifies the
  - Address 940 from which data are to be loaded

1940h = 0001 1001 0100 0000



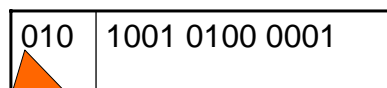
Load contents to AC



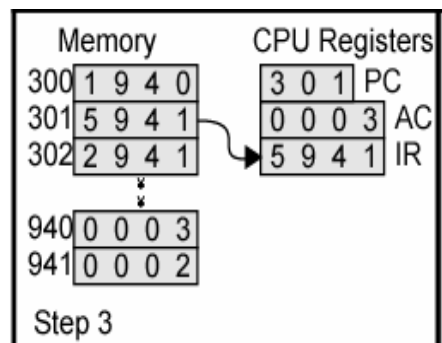
## Example of Program Execution

- The next instruction 5941 is fetched from the location 301
- PC is incremented

5941h = 0101 1001 0100 0001



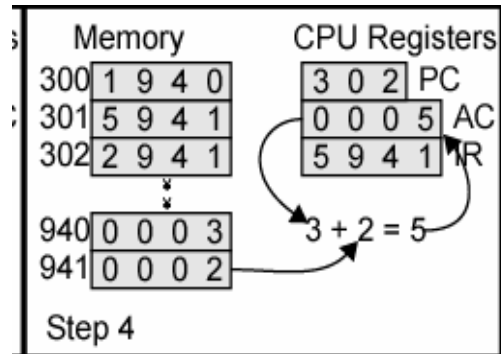
Add contents to AC





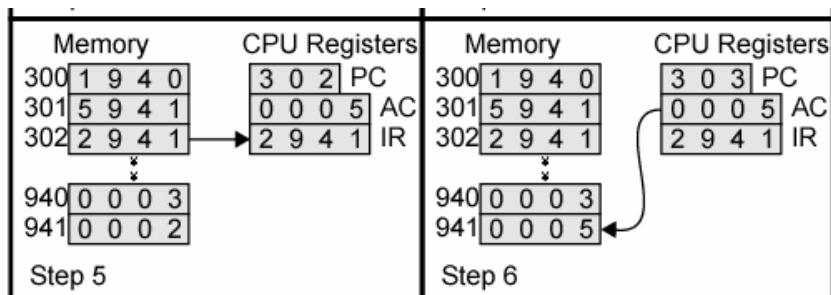
### Example of Program Execution

- The old contents of AC and the contents of the location 941 are added
- The result is stored in AC



### Example of Program Execution

- The next instruction 2941 is fetched from location 302
  - PC is incremented
- The contents of AC are stored in location 941



2941h= 0010 1001 0100 0001

0010	1001 0100 0001
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Store the contents of AC at address x

## Instruction Processing

- The *fetch-decode-execute cycle* is the series of steps that a computer carries out when it runs a program.
  - We first have to *fetch* an instruction from memory, and place it into the IR.
  - Once in the IR, it is *decoded* to determine what needs to be done next.
  - If a memory value (operand) is involved in the operation, it is retrieved and placed into the MBR.
- With everything in place, the instruction is *executed*.

## Instruction Processing

