

Q4

- i) The Time complexity of the following program will be $O(\log(n))$ because in our For loop at each iteration our iterator will keep multiplying by 2. Thus will be $\log(n)$ base 2
- ii) In our lower bound case we will have $O(n)$ because if $n=1$, then it will begin the first for loop when $x = 0$ and when y becomes $x+1$ it won't fulfill the condition of the second loop that y which is $0+1$ is less than n which is 1. Hence $O(n)$ time complexity.

In our upper bound case we will have $O(n^3)$ because if our n is large then it will first run the first loop then run the second loop both of which are $O(n)$. So till here $O(n^2)$ and then in the function when y is $n-1$ and x is 0 then the loop in the function **func** will run the loop n times hence that is $O(n)$ so the total time complexity of the code fragment is $O(n^3)$.

- iii) For all values in our 1-D array $P[n]$, we add all value in array from i to j and store into the 2-D arrays $W[i][j]$. So for example for the case $W[x][y]$, we add all the values in $P[x]+P[x+1]+ \dots + P[y]$, where y begins from $x+1$ and iterates through to $n-1$.
- iv) **int x, y, n; ...**
/* P is a 1-D array of size n integers and W is a 2-D array of size n x n integers */

```
for(int x=0; x<n; x++)Tumharay  
{  
    W[x][x] =P[x];  
    for(int y=x+1; y<n;y++)  
    {  
        int s = P[y]+W[x][y-1];  
        W[x][y] = s ;  
    }  
}
```