CMSC 471

Constraint Satisfaction Problems



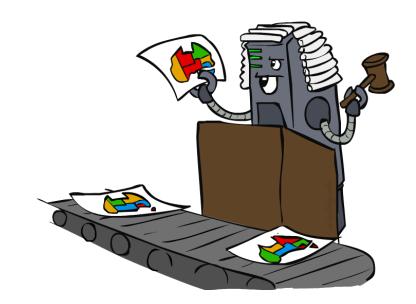


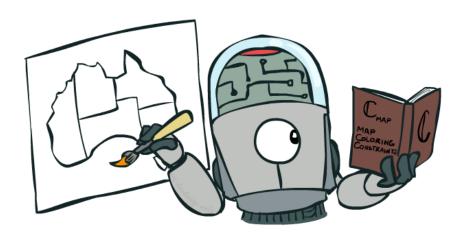
Instructor: KMA Solaiman

These slides were created by Dan Klein and Pieter Abbeel at UC Berkeley. [ai.berkeley.edu]

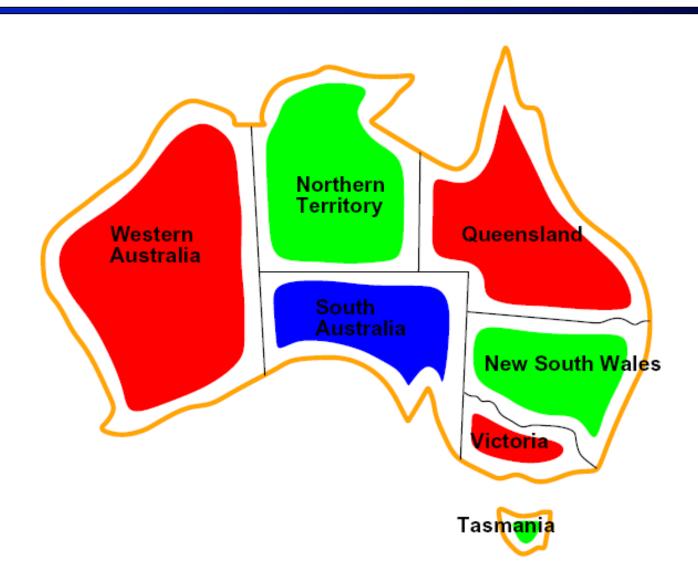
Constraint Satisfaction Problems

- Standard search problems:
 - State is a "black box": arbitrary data structure
 - Goal test can be any function over states
 - Successor function can also be anything
- Constraint satisfaction problems (CSPs):
 - A special subset of search problems
 - State is defined by variables X_i with values from a domain D (sometimes D depends on i)
 - Goal test is a set of constraints specifying allowable combinations of values for subsets of variables
- Simple example of a formal representation language
- Allows useful general-purpose algorithms with more power than standard search algorithms





CSP Examples



Example: Map Coloring

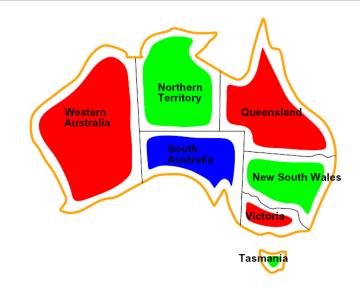
- Variables: WA, NT, Q, NSW, V, SA, T
- Domains: D = {red, green, blue}
- Constraints: adjacent regions must have different colors

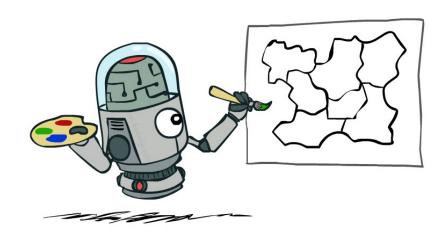
Implicit: $WA \neq NT$

Explicit: $(WA, NT) \in \{(red, green), (red, blue), \ldots\}$

Solutions are assignments satisfying all constraints, e.g.:

{WA=red, NT=green, Q=red, NSW=green, V=red, SA=blue, T=green}



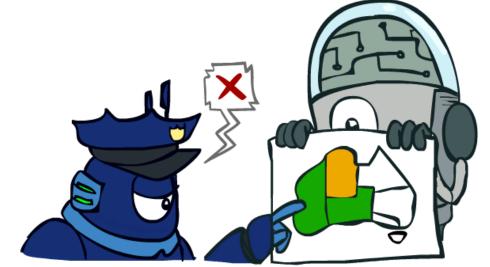


Solving CSPs

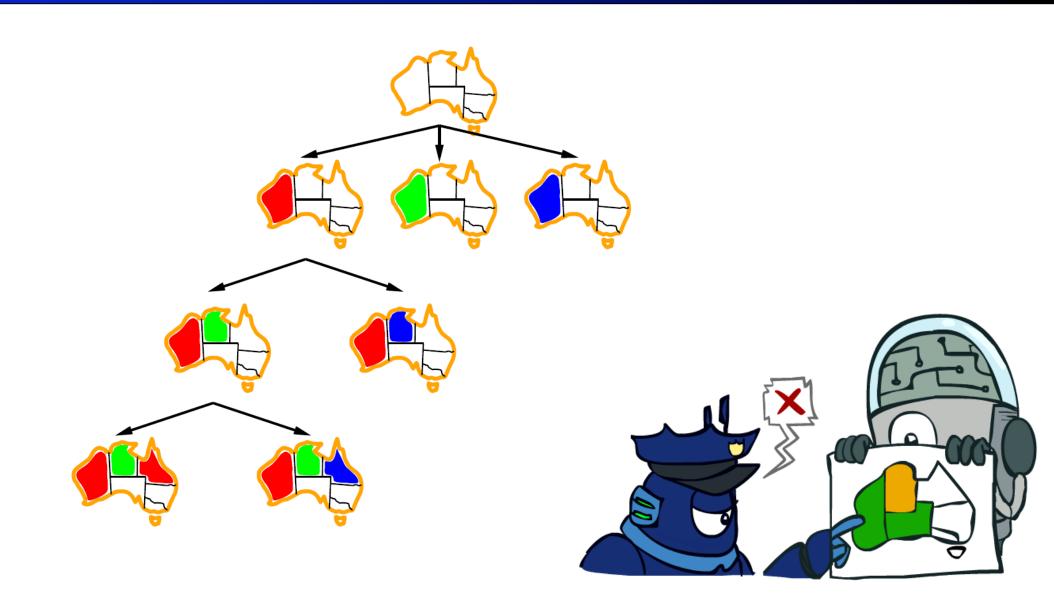


Backtracking Search

- Backtracking search is the basic uninformed algorithm for solving CSPs
- Idea 1: One variable at a time
 - Variable assignments are commutative, so fix ordering
 - I.e., [WA = red then NT = green] same as [NT = green then WA = red]
 - Only need to consider assignments to a single variable at each step
- Idea 2: Check constraints as you go
 - I.e. consider only values which do not conflict previous assignments
 - Might have to do some computation to check the constraint
 - "Incremental goal test"
- Depth-first search with these two improvements is called backtracking search (not the best name)



Backtracking Example



Improving Backtracking

- General-purpose ideas give huge gains in speed
- Ordering:
 - Which variable should be assigned next?
 - In what order should its values be tried?
- Filtering: Can we detect inevitable failure early?
- Structure: Can we exploit the problem structure?



Filtering

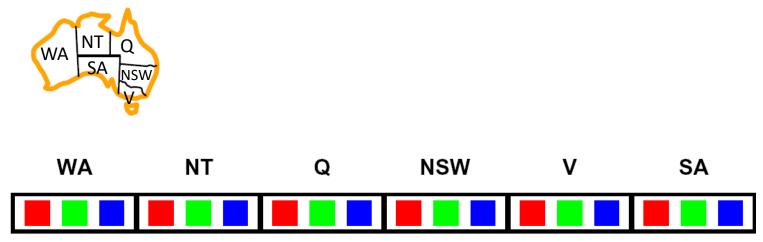


Filtering

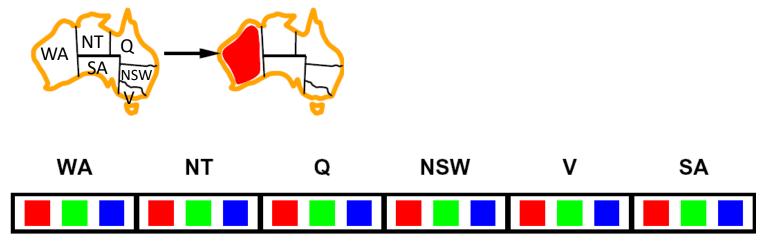


Keeping track of domains for unassigned variables and cross off bad options

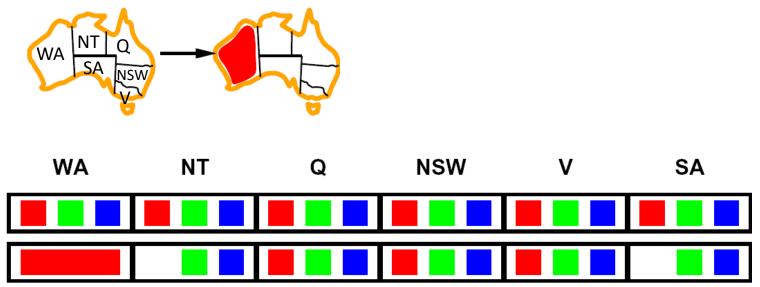
- Filtering: Keep track of domains for unassigned variables and cross off bad options
- Forward checking: Cross off values that violate a constraint when added to the existing assignment



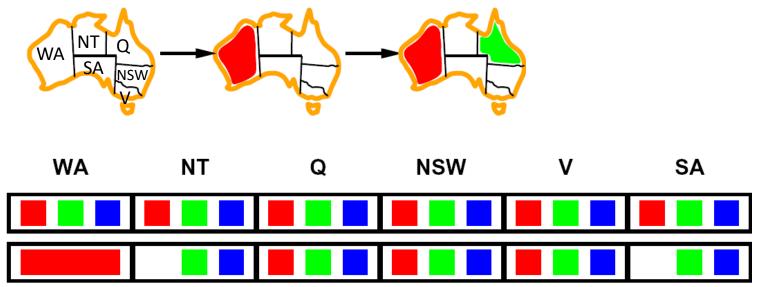
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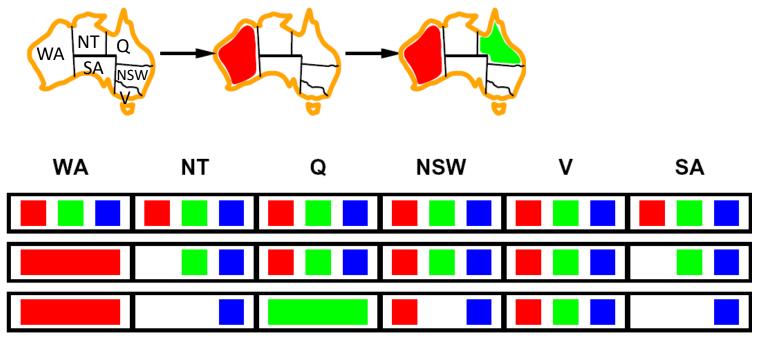
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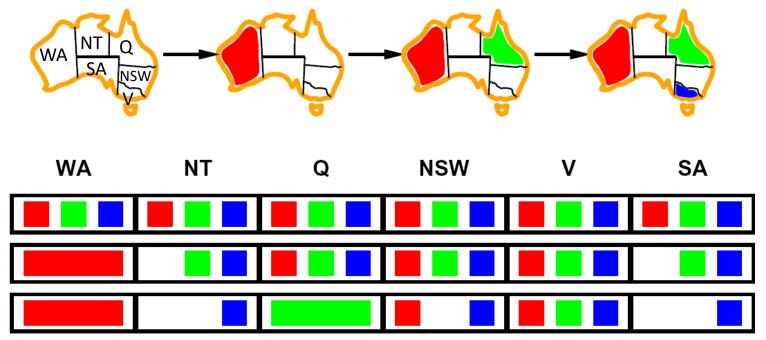
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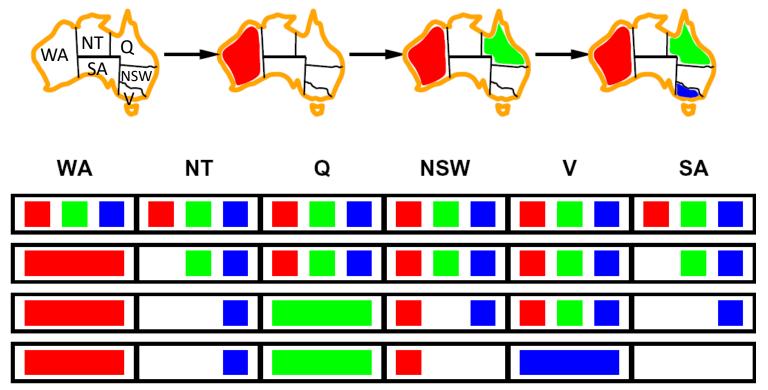
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Video of Demo Coloring – Backtracking with Forward Checking



 Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:





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- Why didn't we detect this yet?

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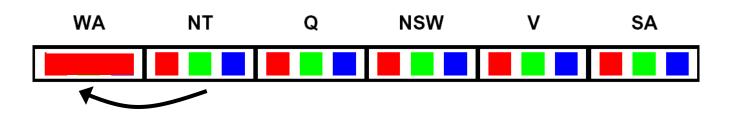


- NT and SA cannot both be blue!
- Why didn't we detect this yet?
- Constraint propagation: reason from constraint to constraint





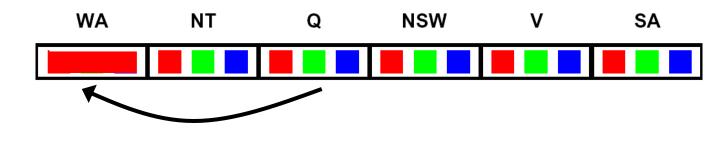






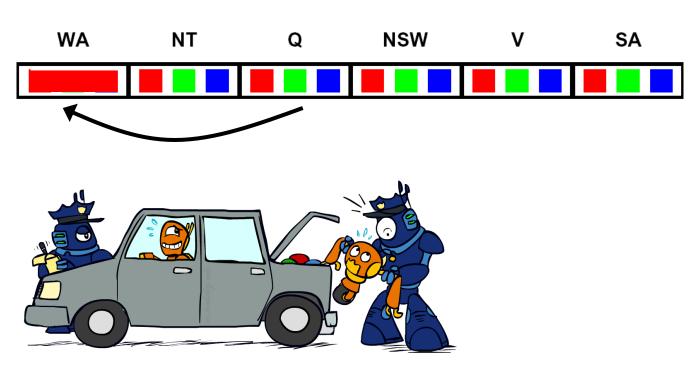






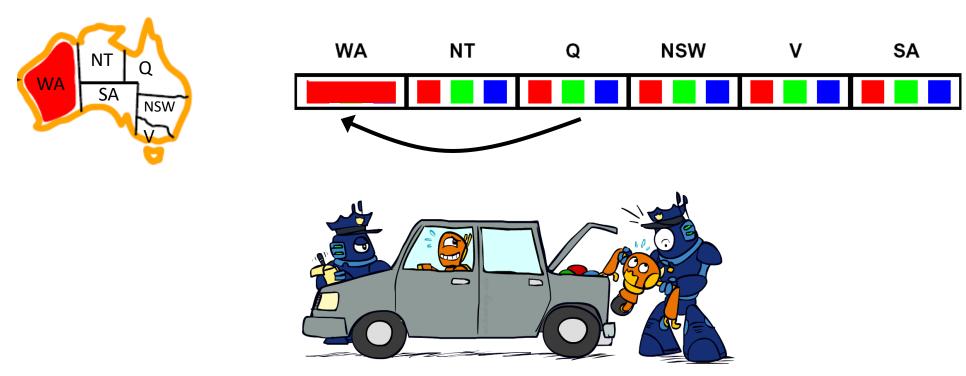






Delete from the tail!

An arc X → Y is consistent iff for every x in the tail there is some y in the head which could be assigned without violating a constraint



Delete from the tail!

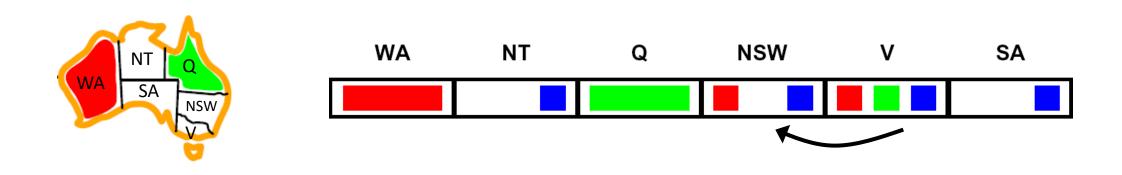
Forward checking: Enforcing consistency of arcs pointing to each new assignment

A simple form of propagation makes sure all arcs are consistent:

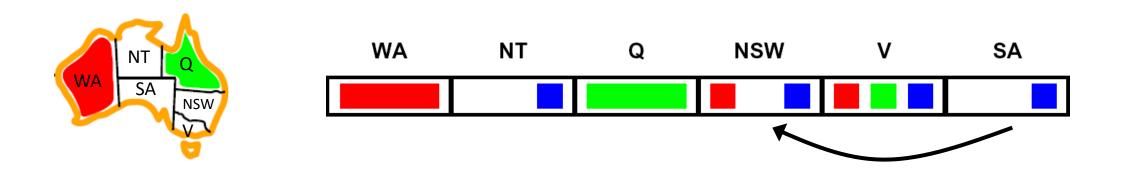




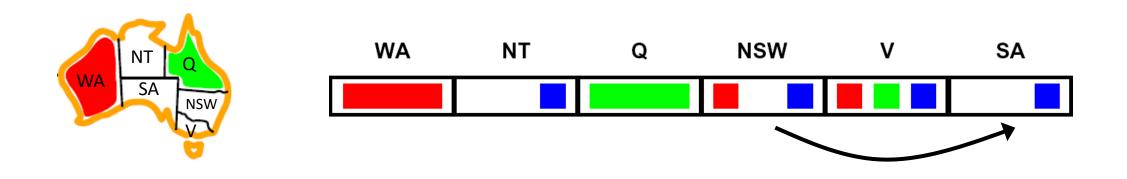
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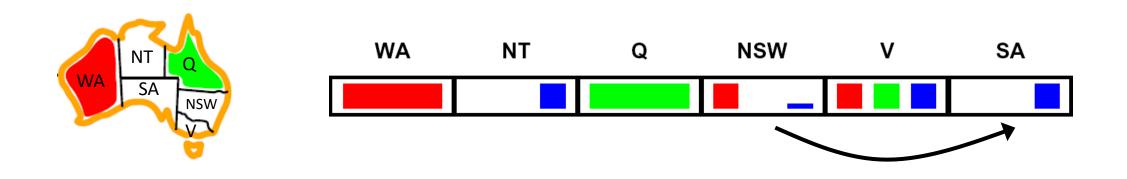
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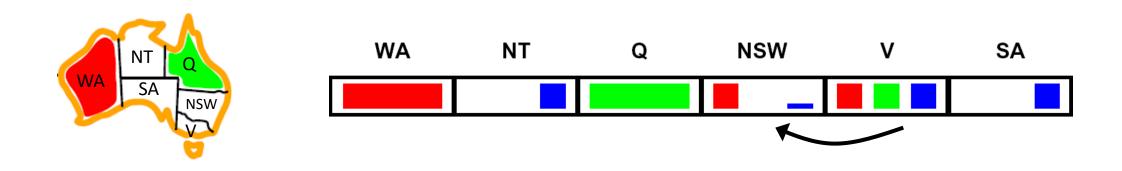
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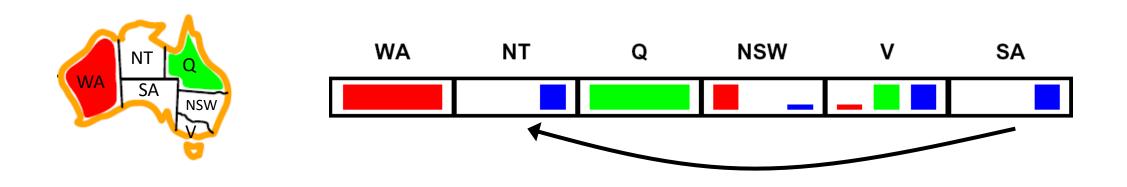
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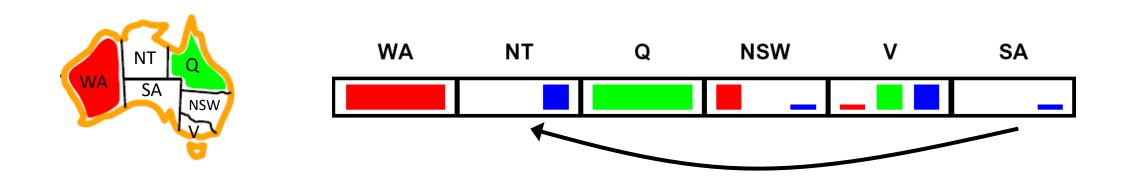


A simple form of propagation makes sure all arcs are consistent:



Arc Consistency of an Entire CSP

A simple form of propagation makes sure all arcs are consistent:



- Important: If X loses a value, neighbors of X need to be rechecked!
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment
- What's the downside of enforcing arc consistency?

Remember: Delete from the tail!

Enforcing Arc Consistency in a CSP

```
function AC-3(csp) returns the CSP, possibly with reduced domains
   inputs: csp, a binary CSP with variables \{X_1, X_2, \ldots, X_n\}
   local variables: queue, a queue of arcs, initially all the arcs in csp
   while queue is not empty do
      (X_i, X_i) \leftarrow \text{REMOVE-FIRST}(queue)
      if Remove-Inconsistent-Values(X_i, X_j) then
         for each X_k in Neighbors [X_i] do
            add (X_k, X_i) to queue
function Remove-Inconsistent-Values (X_i, X_i) returns true iff succeeds
   removed \leftarrow false
   for each x in Domain[X_i] do
      if no value y in DOMAIN[X<sub>i</sub>] allows (x,y) to satisfy the constraint X_i \leftrightarrow X_i
         then delete x from DOMAIN[X<sub>i</sub>]; removed \leftarrow true
   return removed
```

- Runtime: O(n²d³), can be reduced to O(n²d²)
- ... but detecting all possible future problems is NP-hard why?

Video of Demo Coloring – Backtracking with Forward Checking – Complex Graph

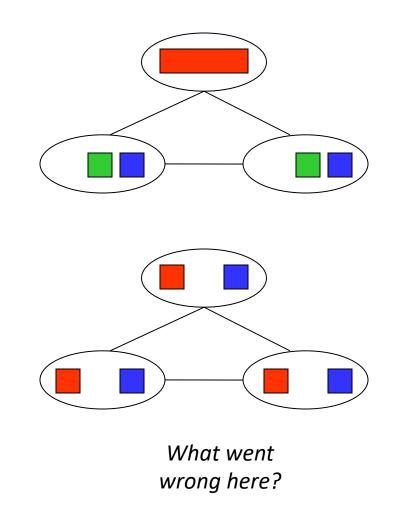


Video of Demo Coloring – Backtracking with Arc Consistency – Complex Graph



Limitations of Arc Consistency

- After enforcing arc consistency:
 - Can have one solution left
 - Can have multiple solutions left
 - Can have no solutions left (and not know it)



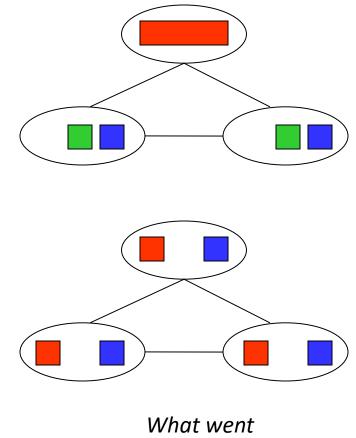
[Demo: coloring -- forward checking]

[Demo: coloring -- arc consistency]

Limitations of Arc Consistency

- After enforcing arc consistency:
 - Can have one solution left
 - Can have multiple solutions left
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• Arc consistency still runs inside a backtracking search!



What went wrong here?

[Demo: coloring -- forward checking]

[Demo: coloring -- arc consistency]

K-Consistency



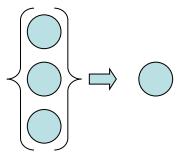
K-Consistency

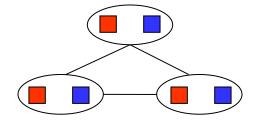
- Increasing degrees of consistency
 - 1-Consistency (Node Consistency): Each single node's domain has a value which meets that node's unary constraints
 - 2-Consistency (Arc Consistency): For each pair of nodes, any consistent assignment to one can be extended to the other
 - K-Consistency: For each k nodes, any consistent assignment to k-1 can be extended to the kth node.

- Higher k more expensive to compute
- (You need to know the k=2 case: arc consistency)









Ordering

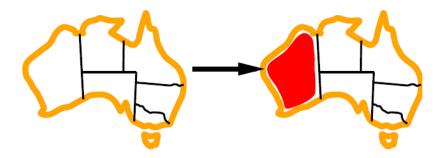


- Variable Ordering: Minimum remaining values (MRV):
 - Choose the variable with the fewest legal left values in its domain





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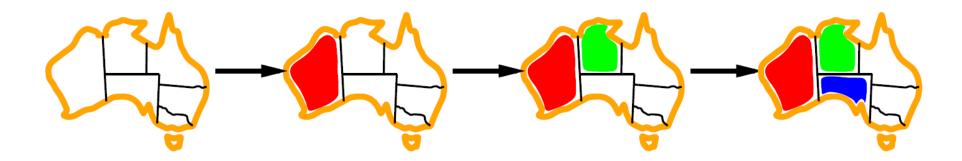


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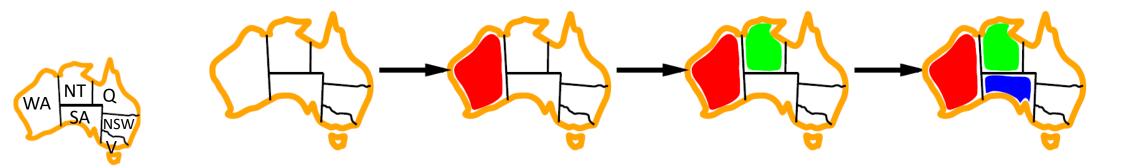


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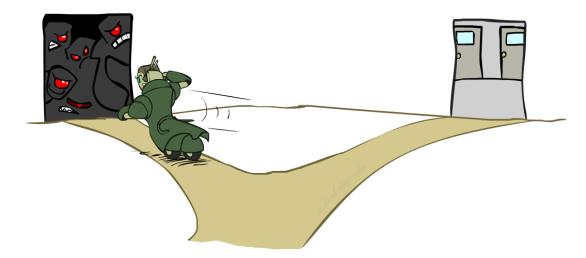


Why min rather than max?

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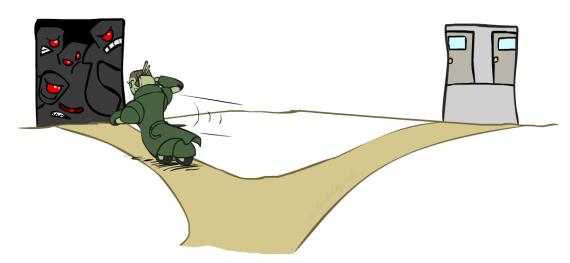
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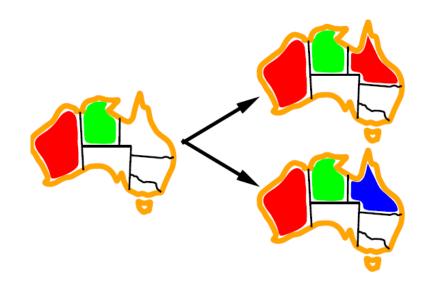
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- Why min rather than max?
- Also called "most constrained variable"
- "Fail-fast" ordering



- Value Ordering: Least Constraining Value
 - Given a choice of variable, choose the *least* constraining value
 - I.e., the one that rules out the fewest values in the remaining variables
 - Note that it may take some computation to determine this! (E.g., rerunning filtering)

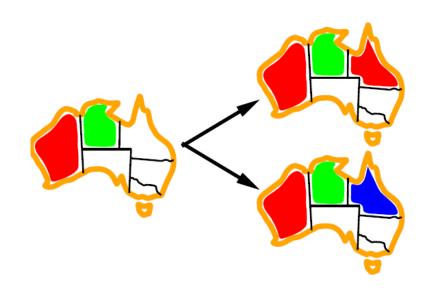




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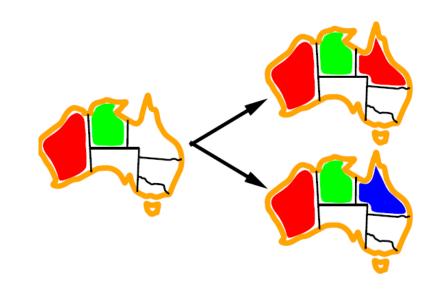




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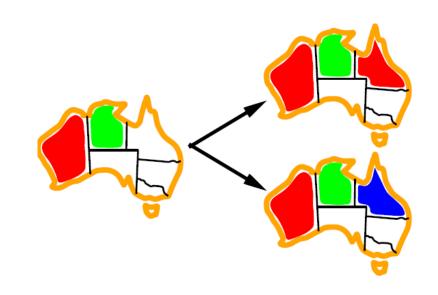


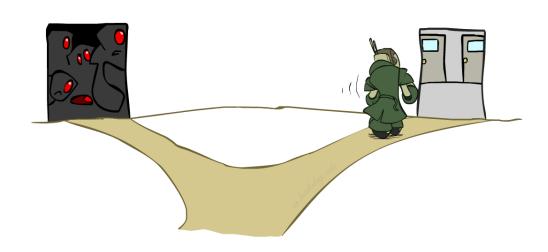


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Combining these ordering ideas makes
 1000 queens feasible





Demo: Coloring -- Backtracking + Forward Checking + Ordering