CMSC 330 Quiz 4 Spring 2022 Solutions

01. Loco Lists

Q1.1. Write a CFG to represent a list consisting of positive integers and other lists. These lists (and sublists) can have arbitrarily large dimensions.

Notes:

- An empty list is also a valid list.
- You can use n to denote a positive integer in the CFG. You don't have to worry about representing multidigit numbers since n encapsulates them all.
- To represent ϵ in the CFG, you can either write the word epsilon or just type the letter e.

Examples of Valid Lists:

```
[]
[[[]]]
[1,[[1,2],[4],56],[[[564]]]]
[1,2,3]

S -> [S] | [T] | []
T -> U,T | U
U -> n | S
```

Q1.2. Is the language defined by the grammar given above regular?

Yes/No

Q2. Context-Free Grammars

My friend Hamza with a peculiar lexicon likes modifying common acronyms in his speech and texting as defined by the following CFG:

```
S -> A | B
A -> loL
L -> l | l out loud | A
B -> smH
H -> h | h my head | B
```

Can the grammar above be parsed by a LL1 recursive descent parser (like Project 4)? Justify your answer.

No. The first sets are not disjoint.

Q3. Ambiguity

Prove that the following grammar is ambiguous:

```
S -> bS | Sb | T
T -> Sa | Sb | Sc | \epsilon
S -> bS -> bT -> b
S -> Sb -> Tb -> b
```

Any other valid ambiguous productions i.e., same string generated with different paths are correct.

Q4. Operational Semantics

Using the gives rules, fill in the blanks the complete the derivation below:

$$\frac{A(x) = v}{A; \ n \Rightarrow n} \qquad \frac{A(x) = v}{A; \ x \Rightarrow v}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A, x : v_1; \ e_2 \Rightarrow v_2}{A; \ let \ x = e_1 \ in \ e_2 \Rightarrow v_2}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

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$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad v_3 \ is \ v_1 \wedge v_2}{A; \ e_1 \wedge e_2 \Rightarrow v_3}$$

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$$\frac{A; \ e_1 \Rightarrow v_1 \qquad A; \ e_2 \Rightarrow v_2 \qquad a; \ e_1 \wedge e_2 \Rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \Rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \Rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \Rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \Rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \Rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \Rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \Rightarrow v_3 \wedge e_1 \wedge e_2 \Rightarrow v_3 \wedge e_1 \wedge e_2 \rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \rightarrow v_3 \qquad a; \ e_1 \wedge e_1 \wedge e_2 \rightarrow v_3 \qquad a; \ e_1 \wedge e_2 \rightarrow v_3$$

Blank #1: let x = "cmsc"

Blank #2: A, x:"cmsc"; "330" ⇒ "330"

Blank #3: x ^ y

Blank #4: A, x:"cmsc", y:"330"(x) = "cmsc"

Blank #5: A, x:"cmsc", y:"330"(y) = "330"

Blank #6: "cmsc330" is "cmsc" ^ "330"