

CS201
Fall 2019
Section 2
Homework Assignment 2
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I. Result Tables

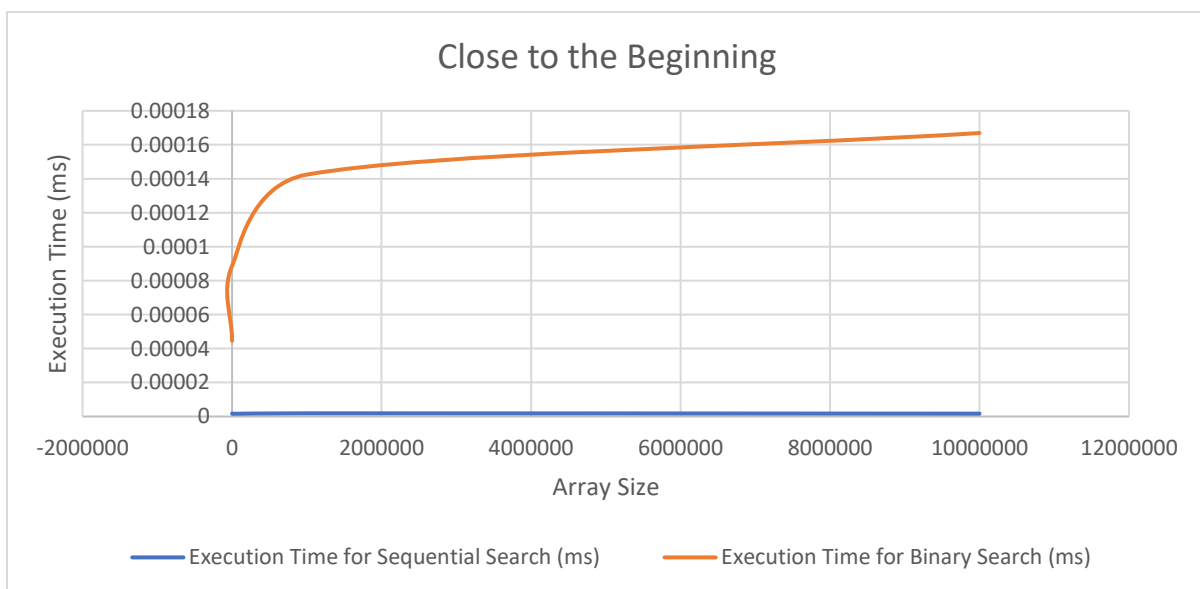
Close to the Beginning

Sequential Search

| Array Size | Time (ms) |
|------------|--------------|
| 101 | 0.0000017558 |
| 10001 | 0.0000016064 |
| 1000001 | 0.0000018875 |
| 10000001 | 0.0000016859 |

Binary Search

| Array Size | Time (ms) |
|------------|--------------|
| 101 | 0.0000445662 |
| 10001 | 0.0000897338 |
| 1000001 | 0.0001424410 |
| 10000001 | 0.0001669130 |



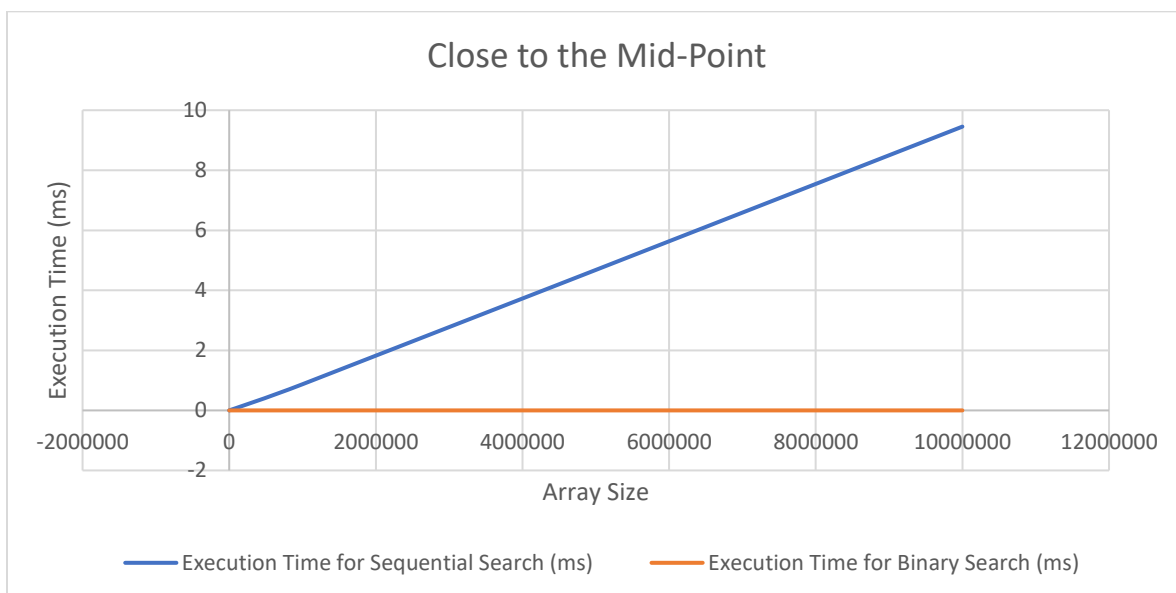
Close to the Mid-Point

Sequential Search

| Array Size | Time (ms) |
|------------|-----------|
| 101 | 0.000092 |
| 10001 | 0.008739 |
| 1000001 | 0.872176 |
| 10000001 | 9.453870 |

Binary Search

| Array Size | Time (ms) |
|------------|--------------|
| 101 | 0.0000080027 |
| 10001 | 0.0000080841 |
| 1000001 | 0.0000079643 |
| 10000001 | 0.0000080284 |



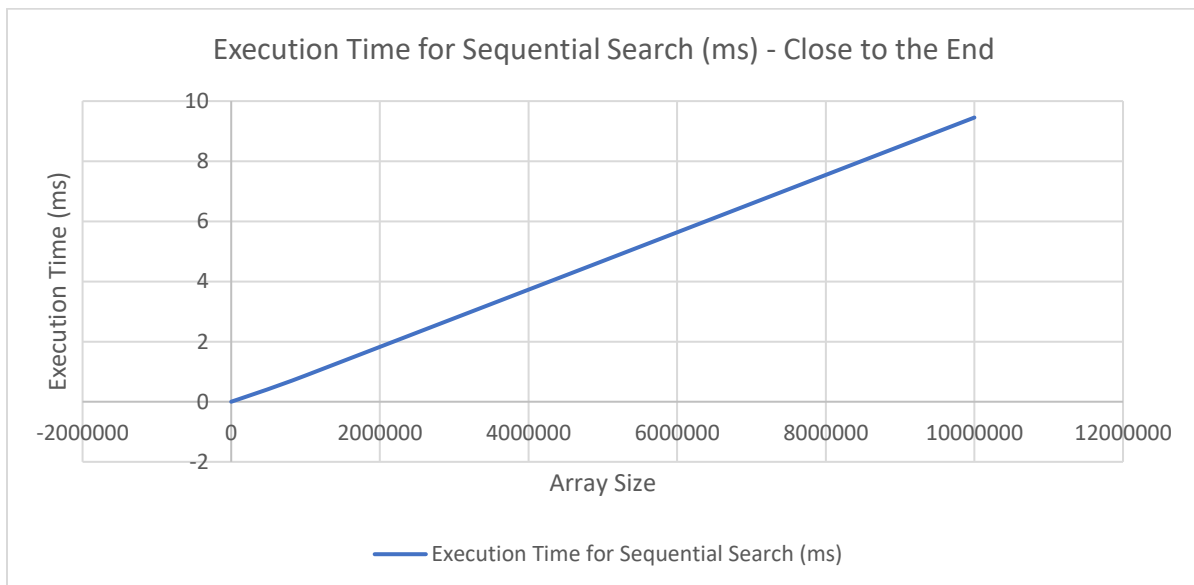
Close to the End

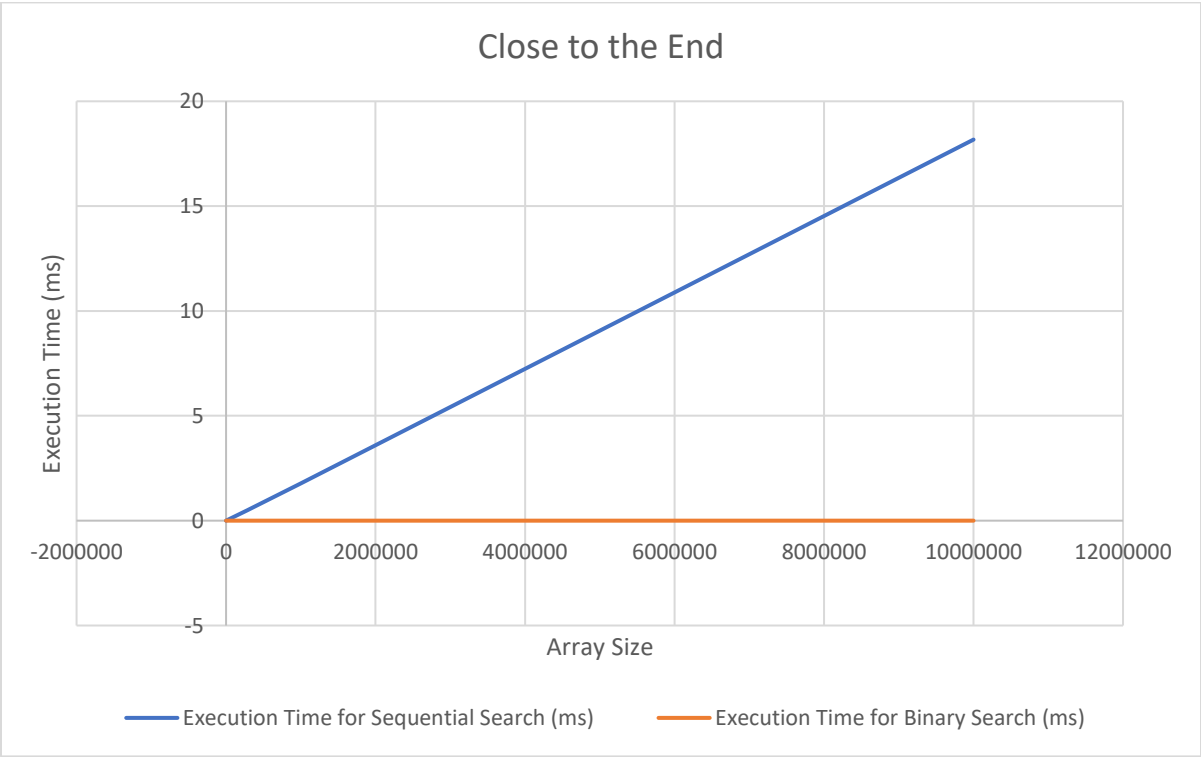
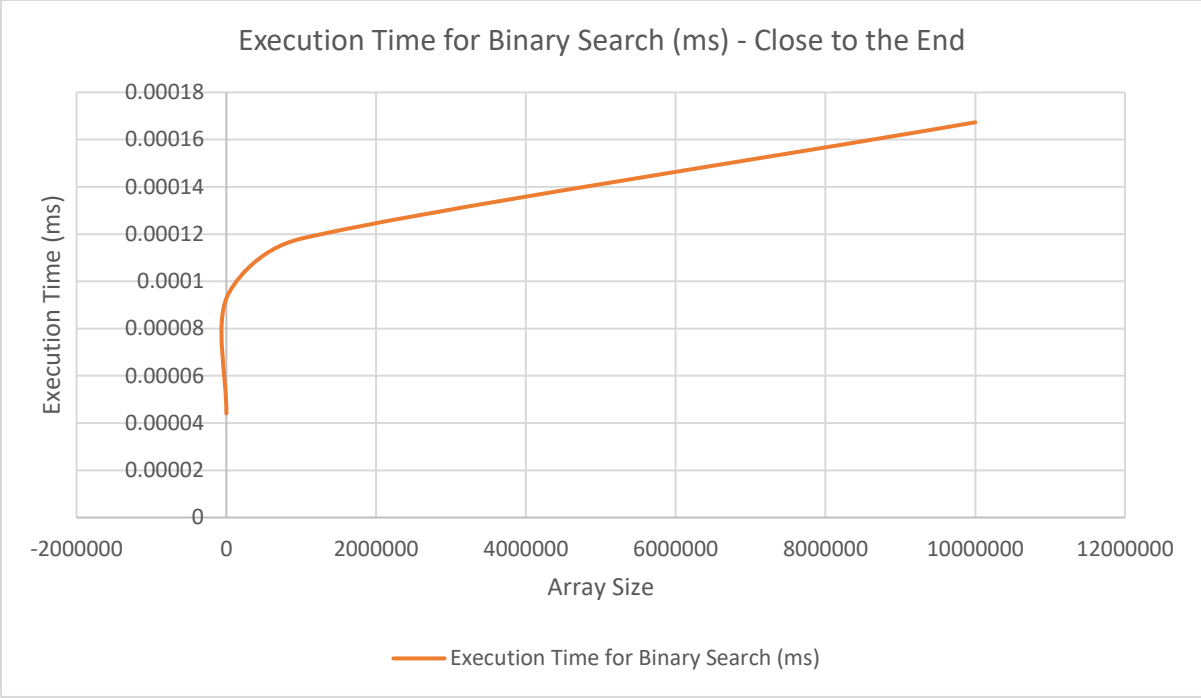
Sequential Search

| Array Size | Time (ms) |
|------------|-----------|
| 101 | 0.000185 |
| 10001 | 0.017436 |
| 1000001 | 1.7762 |
| 10000001 | 18.1729 |

Binary Search

| Array Size | Time (ms) |
|------------|--------------|
| 101 | 0.0000440448 |
| 10001 | 0.0000935630 |
| 1000001 | 0.0001180780 |
| 10000001 | 0.0001673350 |





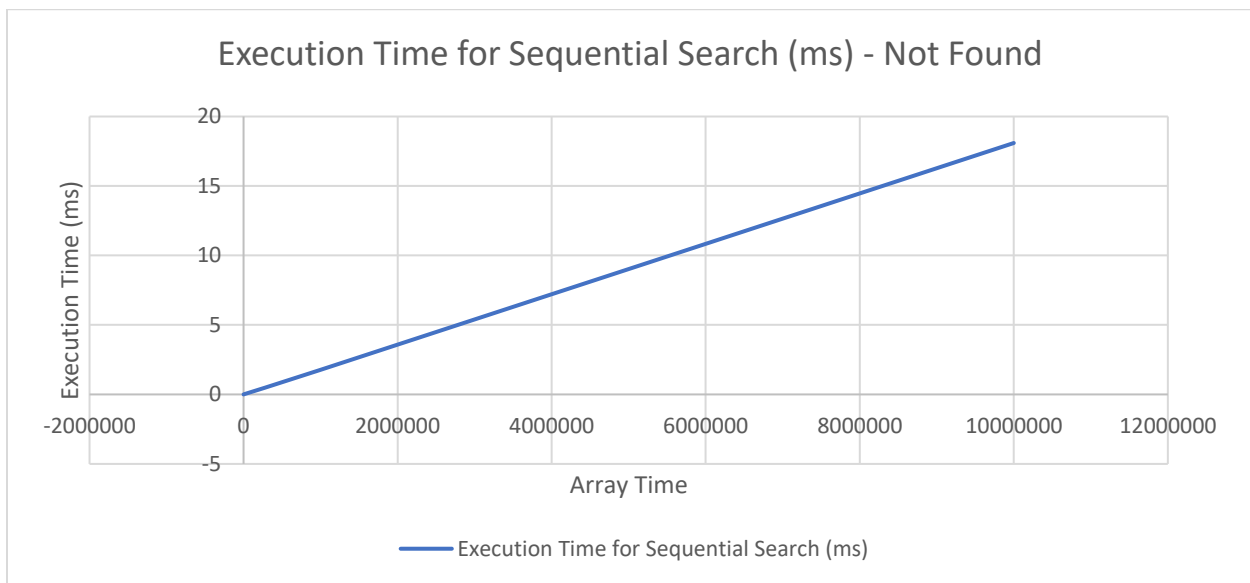
Not Found

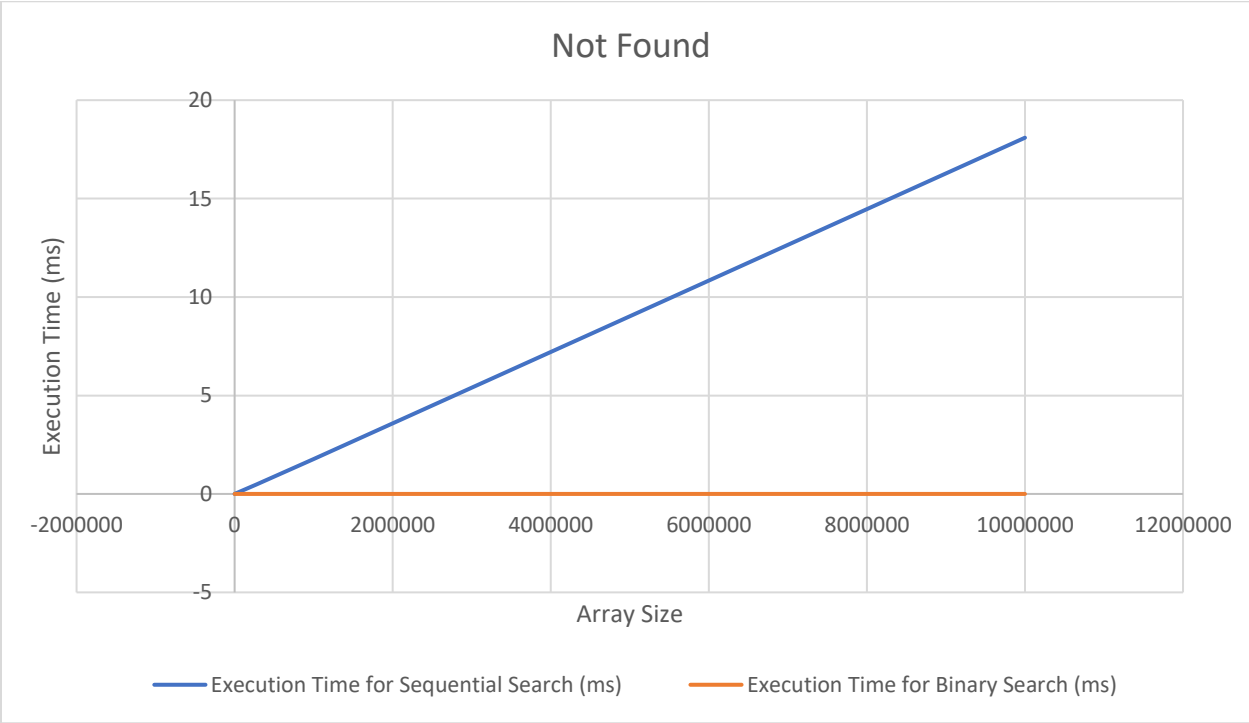
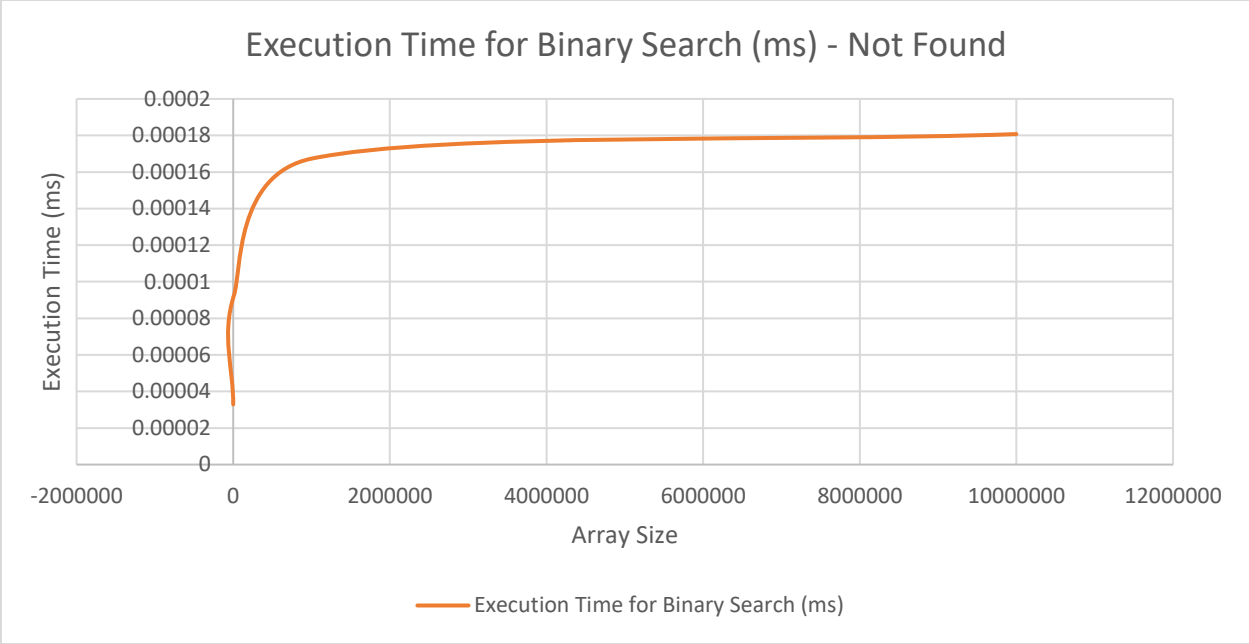
Sequential Search

| Array Size | Time (ms) |
|------------|-----------|
| 101 | 0.0001873 |
| 10001 | 0.018577 |
| 1000001 | 1.77191 |
| 10000001 | 18.0926 |

Binary Search

| Array Size | Time (ms) |
|------------|--------------|
| 101 | 0.0000329194 |
| 10001 | 0.0000926171 |
| 1000001 | 0.0001673350 |
| 10000001 | 0.0001807490 |





II. Big O Notations – Best and Worst Cases

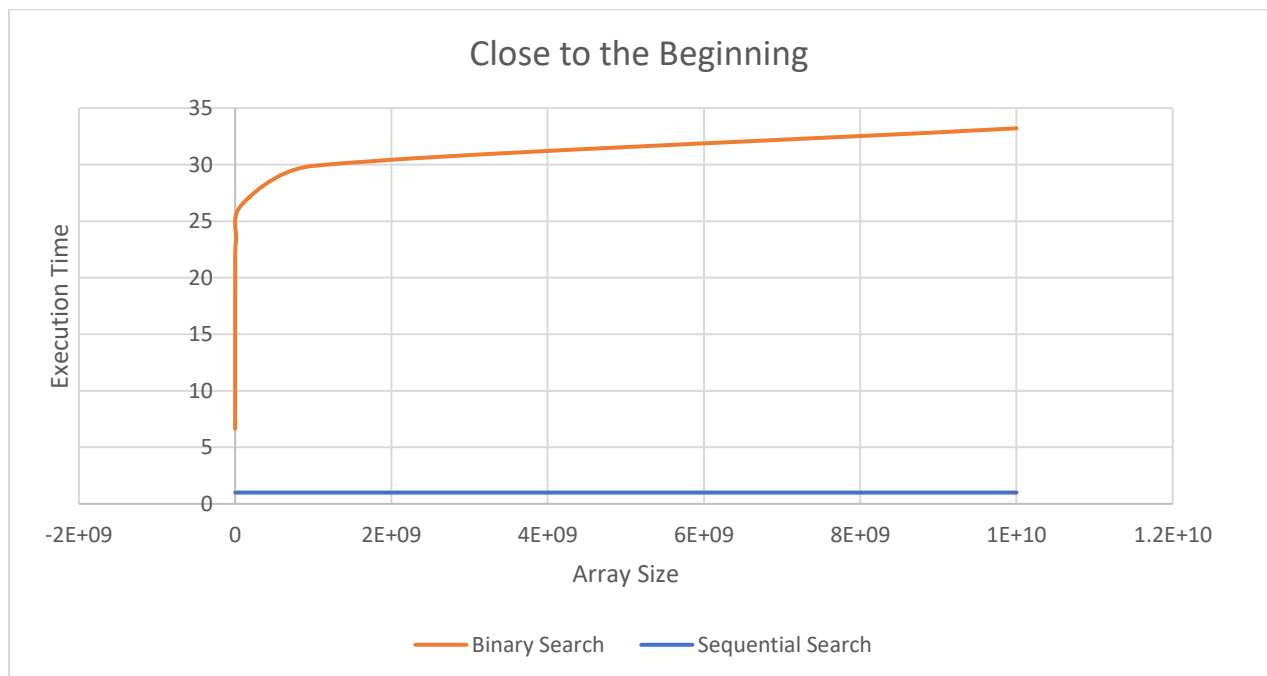
Sequential Search

- The best case arises where the key is close to the first index (close to the beginning) → **$O(1)$**
- The worst case arises where the key is close to the last index or not found in array (close to the end, not found) → **$O(N)$**
- Average Case → **$O(N) = \frac{O(N) + O(1)}{2}$**

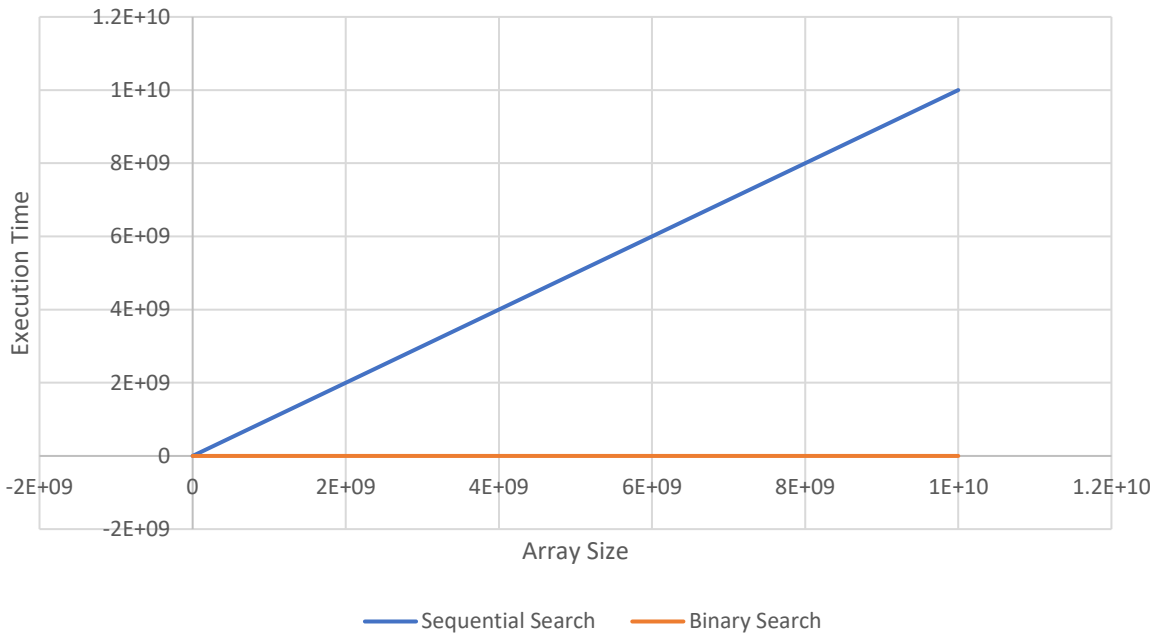
Binary Search

- The best case arises where the key is at the mid-point of the array (close to the mid-point) → **$O(1)$**
- The worst case arises where the key is close to the beginning or end or not found (close to the beginning, close to the end, not found) → **$O(\log(N))$**
- Average Case → **$O(\log(N)) = \frac{O(\log(N)) + O(1)}{2}$**

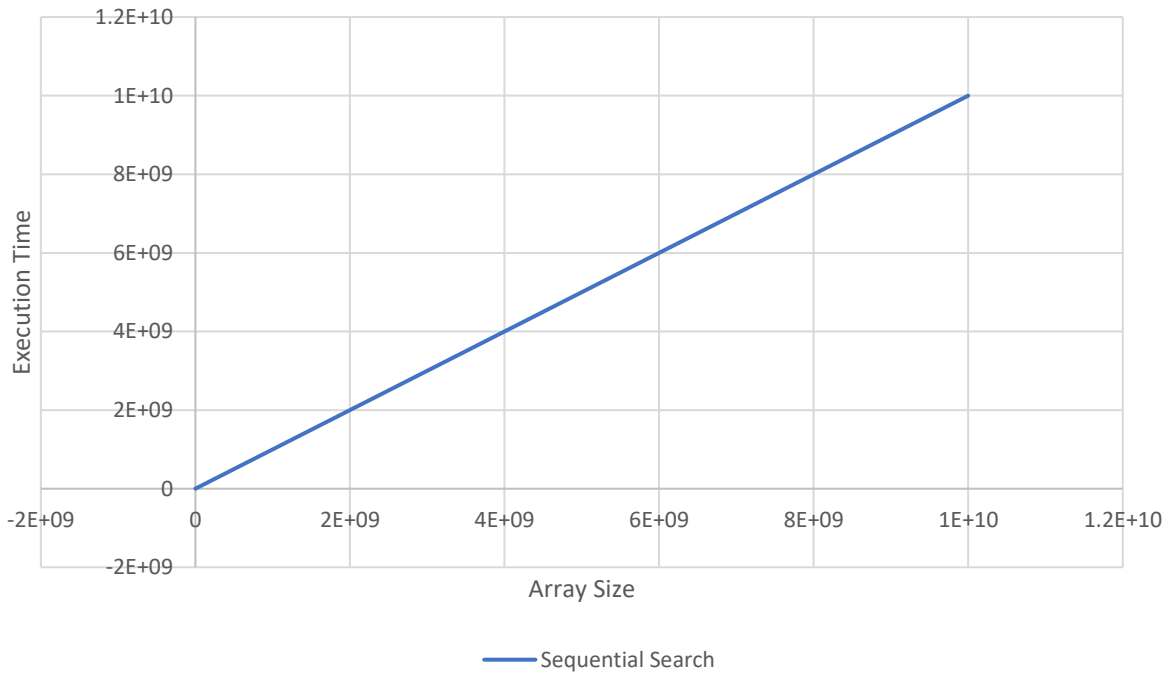
III. Theoretical Graphics



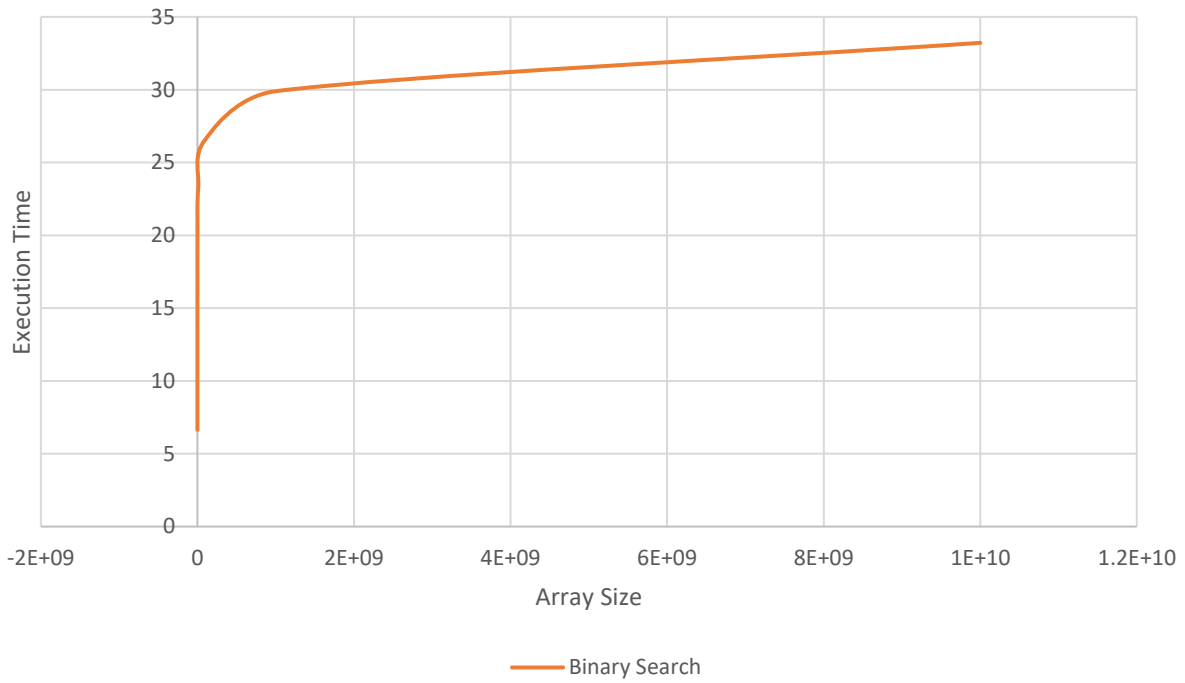
Close to the Mid-Point



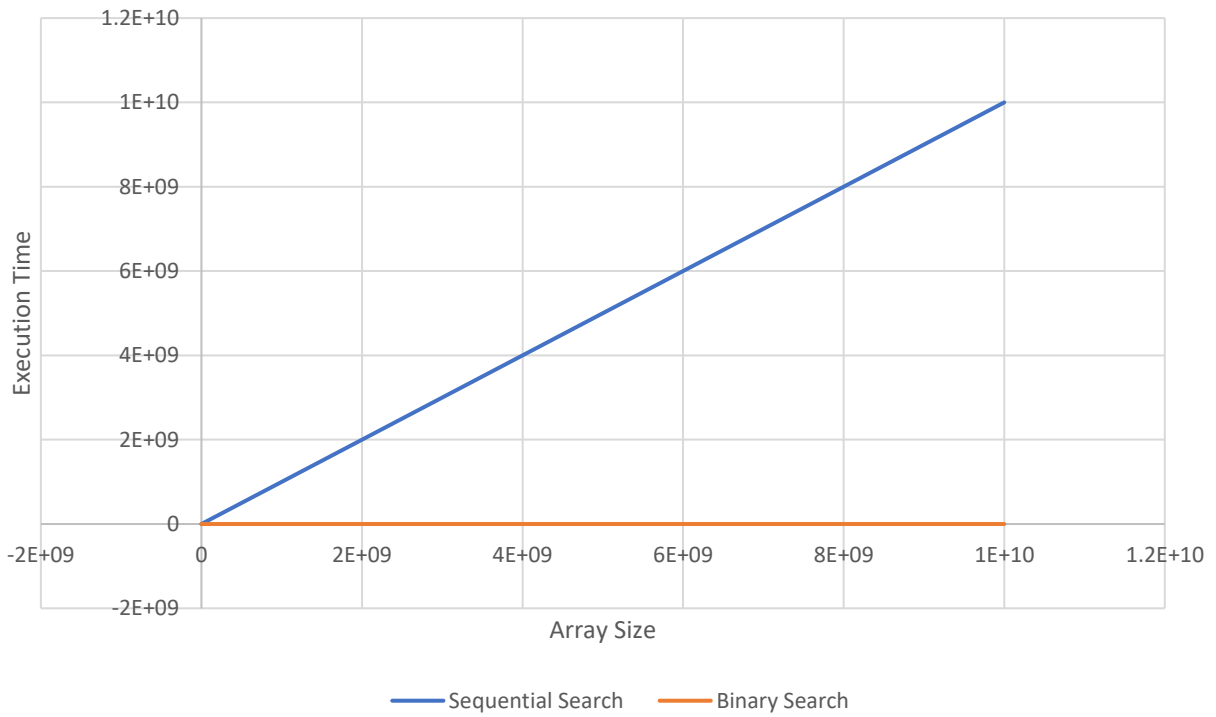
Execution Time for Sequential Search (ms) - Close to the End



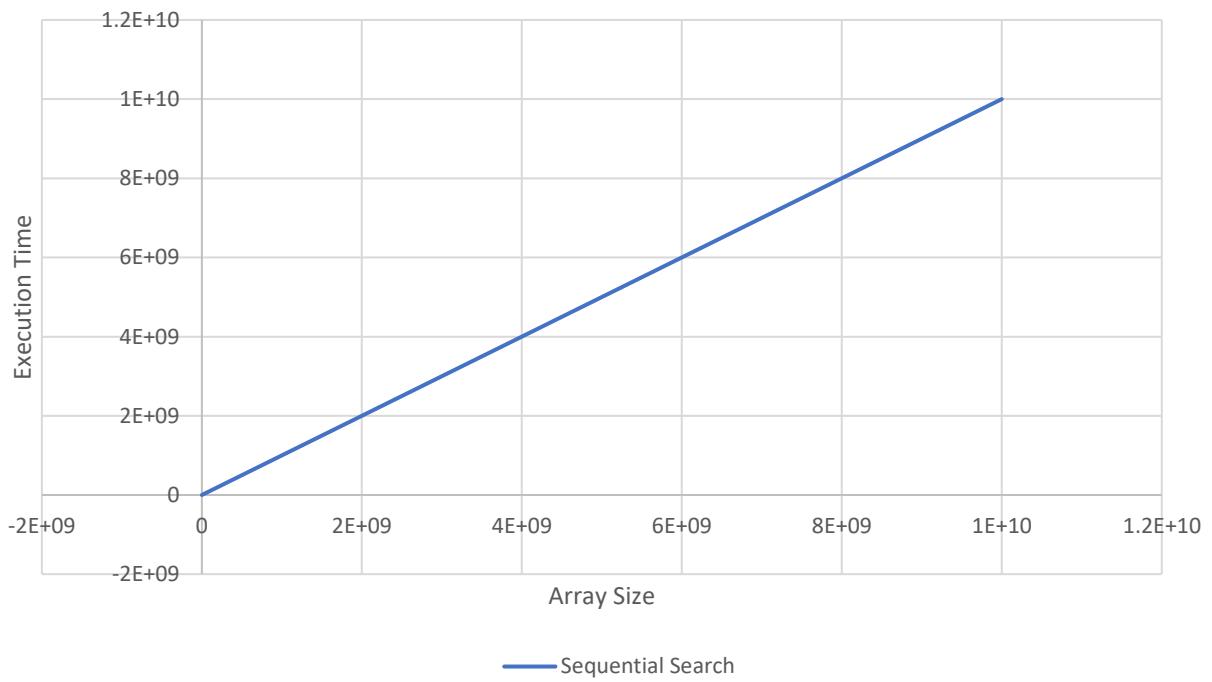
Execution Time for Binary Search - Close to the End



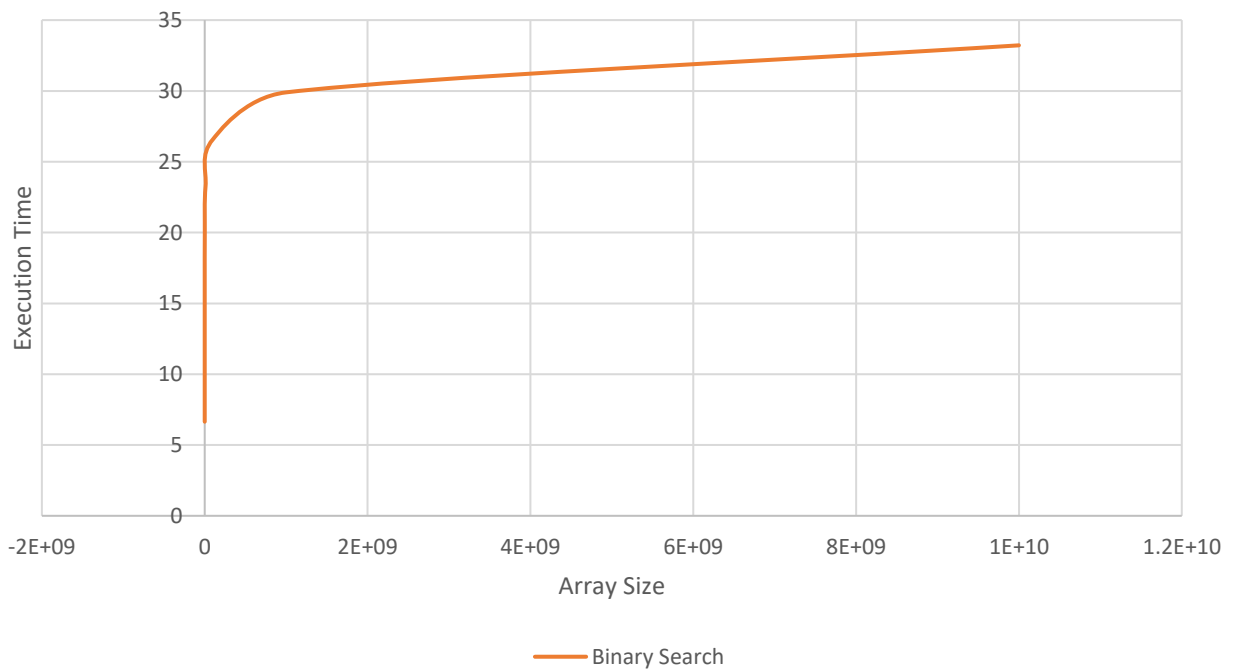
Close to the End

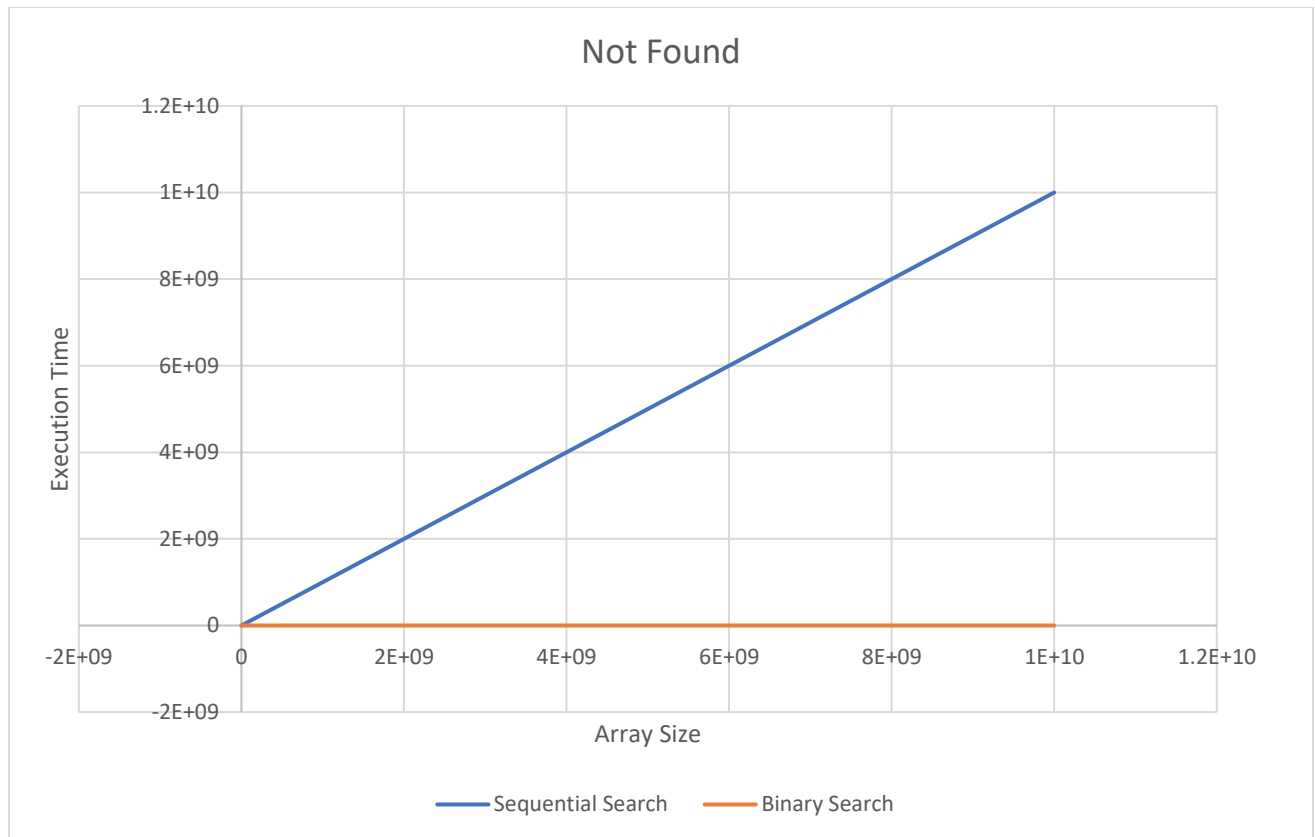


Execution Time for Sequential Search (ms) - Not Found



Execution Time for Binary Search (ms) - Not Found





IV. System Specifications

Vostro 5481

Device name DESKTOP-C4M6KNV
Processor Intel(R) Core(TM) i7-8565U CPU @ 1.80GHz 1.99 GHz
Installed RAM 8.00 GB (7.78 GB usable)

V. Conclusion

As we can see in the graphs that are drawn based on our data, although there are some errors based on my computer systems, the graphs are similar to expected graphs that are drawn based on mathematical functions. For sequential search this mathematical function is $O(N)$, so that our graphs should be linear. For binary search this mathematical function is $O(\log(N))$, so that our graphs should be logarithmic. As a result, we can see the similarities between graphs of our data and expected graphs.