

191-15-12485

Selection Sort

Implementation:

```
void selection sort (int arr[], int n)
```

```
{
    for (int i = 0; i < (n-1); i++)
```

```
{
    for (int j = i+1; j < n; j++)
```

```
{
    if (arr[i] > arr[j])
```

```
        swap (arr[i], arr[j]);
```

```
}
```

```
}
return;
```

Analysis:

```
for (int i = 0; i < (n-1); i++)
```

Cost · time

C_1 $(n-1)$

```
{
    for (int j = i+1; j < n; j++)
```

C_2 $\frac{n(n+1)}{2}$

```
{
    if (arr[i] > arr[j])
```

C_3 $\frac{n(n+1)}{2}$

```
        swap (arr[i], arr[j]);
```

C_4 $\frac{n(n+1)}{2}$

```
}
```

```
}
```

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∴ Time function $f(t) = c_1(n-1) + c_2 \frac{n(n+1)}{2} + c_3 \frac{n(n+1)}{2} + c_4 \frac{n(n+1)}{2}$
 Best case occurs when the array is already sorted

1	2	3	4	5
---	---	---	---	---

value of c_4 will be 0, so time function will be.

$$f(t) = \left(\frac{c_2 + c_3}{2} \right) * n^2 + \left(\frac{c_2 + c_3 + 2c_1}{2} \right) * n - c_1$$

This function look like $f(t) = an^2 + bn + c$, which is quadratic function

∴ Worst case will be occurred when the array is reversely sorted

5	4	3	2	1
---	---	---	---	---

$$f(t) = \left(\frac{c_4 + c_2 + c_3}{2} \right) * n^2 + \left(\frac{c_4 + c_2 + c_3 + 2c_1}{2} \right) * n - c_1$$

look like $f(t) = (an^2 + bn + c)$; which is quadratic function.

Time Complexity:

$$f(t) = an^2 + bn + c \text{ [Best Case]}$$

$$O(n^2)$$

Worst case:

$$f(t) = an^2 + bn + c$$

$$= O(n^2)$$

Implementation: Insertion sort

```
void insertion (int arr [], int n)
```

```
{
    int i, j, x;
```

```
    for (i = 0; i < n; i++)
```

```
{
```

```
        x = arr [i];
```

```
        j = i - 1;
```

```
        while (j > 0 && arr [j] > x)
```

```
            arr [j+1] = arr [j]; j--;
```

```
            arr [j+1] = x;
```

```
        }
```

Analysis:

```
void insertion (int arr [], int n)
```

```
{
    int i, j, x;
```

```
    for (i = 0; i < n; i++)
```

```
{
```

```
        x = arr [i];
```

```
        j = i - 1;
```

```
        while (j > 0 && arr [j] > x)
```

```
            arr [j+1] = arr [j]; j--;
```

Cost time

C_1

C_2

C_3

C_4

C_5

$n-1$

$n-1$

$\frac{n(n+1)}{2}$

$\frac{n(n+1)}{2}$

$\frac{n(n+1)}{2}$

$\frac{n(n+1)}{2}$


```

    arr[j+1] = x;
  }
}

```

∴ Time function:

$$f(t) = c_1 n + c_2(n-1) + c_3(n-1) + c_4 \frac{n(n+1)}{2} + c_5 \frac{n(n+1)}{2} + c_6(n-1)$$

Best case occurs when array is already sorted

1	2	3	5	7
---	---	---	---	---

When c_4 will execute for $(n-1)$ time and c_5 will execute for 0 time.

So function will be

$$\begin{aligned} f(t) &= c_1 n + c_2(n-1) + c_3(n-1) + c_4(n-1) + c_5 * 0 + c_6(n-1) \\ &= n(c_1 + c_2 + c_3 + c_4 + c_6) - (c_2 + c_3 + c_4 + c_6) \end{aligned}$$

∴ This look like $y = ax - b$. That's Linear eqⁿ worst case occurs when the array is reversely sorted,

7	5	3	2	1
---	---	---	---	---

∴ Time function will be

$$\begin{aligned} f(t) &= c_1 n + c_2 + c_3(n-1) + c_4 * \frac{n(n+1)}{2} + c_5 * \frac{n(n+1)}{2} + c_6(n-1) \\ &= \frac{c_1 + c_2}{2} n^2 + \frac{2(c_1 + c_2 + c_3 + c_6)}{2} n - (c_2 + c_3 + c_4 + c_6) \end{aligned}$$

\therefore look like $y = an^2 + bn + c$; This is a quadratic eq

Time Complexity :

Best Case:

We know, $y = an - b$ [\therefore Best case]

So, $\sim (n)$

worst case:

We know, $y = an^2 + bn + c$ [\therefore worst case]

So, $O(n^2)$

1	2	3	4	5
---	---	---	---	---

1	2	3	4	5
---	---	---	---	---