# INDR371 - HW5

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# Question 1

### Part A

```
In [3]: | m.optimize() | | Gurobi Optimizer version 9.5.2 build v9.5.2rc0 (mac64[arm]) | Thread count: 8 physical cores, 8 logical processors, using up to 8 threads Optimize a model with 32 rows, 18 columns and 64 nonzeros | Model fingerprint: 0xcba38815 | Ocofficient statistics: | Matrix range | [1e+00, 1e+00] | Objective range | Se+00, 1e+01] | Objective range | O
```

Figure 1: Code and output of Question 1, Part A

In minisum algorithm, there are absolute valued terms. So I had to convert them into linear first. In order to do that, I added 2 variables r and s and added 2 constraints for each, to ensure that  $r_i = |x - a_i|$  and  $s_i = |y - b_i|$ . Then, I wrote r and s for absolute valued term in the objective function.

The objective function is 117 and the optimal location for the repair shop is (3,5).

## Part B

```
In [27]:
a.append(x.X)
b.append(y.X)
w.append(np.mean(w))
m.ModelSense = GRB.MINIMIZE
x = m.addVar(name="v",vtype=GRB.CONTINUOUS)
y = m.addVars(range(len(a)), name="r",vtype=GRB.CONTINUOUS)
r = m.addVars(range(len(a)), name="s",vtype=GRB.CONTINUOUS)
s = m.addVars(range(len(a)), name="s",vtype=GRB.CONTINUOUS)
m.addConstrs(r[i] >= x - a[i] for i in range(len(a)))
m.addConstrs(s[i] >= y - b[i] for i in range(len(b)))
m.addConstrs(s[i] >= b[i] - y for i in range(len(a)))
m.addConstrs(s[i] >= b[i] - y for i in range(len(a)))
m.setDbjective(c)

(a) Code

In [28]: m.optimize()

In [28]: m.optimize()

In [28]: m.optimize()

In [28]: b.optimize version 9.5.2 build v9.5.2rc0 (mac64[arm])
Thread count: 8 physical cores, 8 logical processors, using up to 8 threads Optimize a model with 45 rows, 21 columns and 99 nonzeros

Model fingerprint: 0xb5r01c80

Coefficient statistics:

Marix range
[le+00, le+01]
Objective range
[le+00, le+00]
Nour range
[le+00, le+01]
Objective range
[le+00, le+00]
Nour range
[le+00, le+00]
```

Figure 2: Code and output of Question 1, Part B

Also in minimax algorithm, there are absolute valued terms. So I had to convert them into linear first. In order to do that, I added 2 variables r and s and added 2 constraints for each, to ensure that  $r_i = |x - a_i|$  and  $s_i = |y - b_i|$ . Then, I defined a new variable c and added a constraint  $c \ge W_i * (r_i + s_i)$  for each point and set the objective function to be "c", so that we can minimize the maximum distance to a point. Additionally, I set the weight of the repair shop equal to the average of the weight of the facilities.

The objective function is 24.706 and the optimal location for the fire station is (2.53,5).

### Part C

- $z_{ij}$ : whether fire station i is assigned to facility j or not
- $x_i$ : x coordinate of fire station i
- $y_i$ : y coordinate of fire station i
- $a_i$ : x coordinate of facility j

- $b_j$ : y coordinate of facility j
- $r_{ij}$ :  $|x_i a_j|$
- $s_{ij}$ :  $|y_i b_j|$
- $W_i$ : weight of facility j

$$\min \quad c_1 + c_2 \tag{1a}$$

s.t. 
$$\sum_{i=1}^{2} z_{ij} = 1 \qquad \forall j$$
 (1b)

$$r_{ij} \ge x_i - a_j$$
  $\forall i, j$  (1c)

$$r_{ij} \ge a_j - x_i \qquad \forall i, j \tag{1d}$$

$$s_{ij} \ge y_i - b_j$$
  $\forall i, j$  (1e)

$$s_{ij} \ge b_j - y_i \qquad \forall i, j \tag{1f}$$

$$c_i \ge z_{ij} * W_j(r_{ij} + s_{ij})$$
  $\forall i, j$  (1g)

$$x_i, y_i \in \mathbb{Z}$$
  $\forall i$  (1h)

$$z_{ij} \in \{0, 1\} \qquad \forall i, j \tag{1i}$$

- Constraint 1b is for the assignment of 1 fire station to each facility.
- Constraints 1c, 1d, 1e and 1f are for the linearization of absolute valued terms.
- Constraint 1g is to keep the maximum distance between each fire station and the facilities assigned to it.
- Constraints 1h and 1i are for the domains of decision variables.

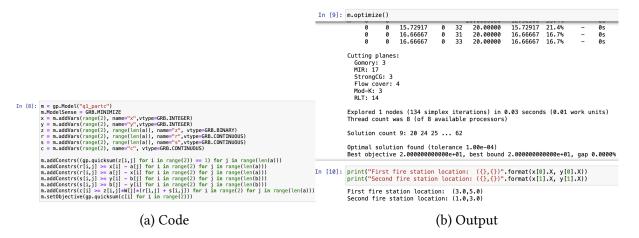


Figure 3: Code and output of Question 1, Part C

The objective function is 20 and the optimal locations for the fire stations are (1,3) and (3,5). The objective function in part C is expected to be less than part B since we do not consider the fixed cost of opening fire stations, therefore opening as much fire stations as possible would decrease the cost.

# **Question 2**

### Part A

- $x_{ir}$ : whether facility i is located at location r or not
- $r_{ij}$ : flow between facility i and j
- $d_{rs}$ : distance between location r and s

$$\min \sum_{i=1}^{5} \sum_{j=1, i \neq j}^{5} \sum_{r=1}^{12} \sum_{s=1, r \neq s}^{12} (r_{ij} * d_{rs} + r_{ji} * d_{sr}) * x_{ir} * x_{js}$$
 (2a)

s.t. 
$$\sum_{i=1}^{5} x_{ir} \le 1 \qquad \forall r$$
 (2b)

$$\sum_{r=1}^{12} x_{ir} = 1 \qquad \forall i \tag{2c}$$

$$\sum_{i=1}^{5} \sum_{j=1}^{5} d_{ij} * x_{(D_1,i)} * x_{(D_4,j)} \ge 2$$
(2d)

$$x_{ir} = 0$$
  $\forall i \in \{D_1, D_2, D_5\}, \forall r \in \{A, D, L, I\}$  (2e)

$$x_{(D_3),r} = 0 \qquad \forall r \notin A, D, L, I$$

$$x_{ir} \in \{0, 1\} \qquad \forall i, r$$

$$(2f)$$

$$(2g)$$

$$x_{ir} \in \{0, 1\} \qquad \forall i, r \tag{2g}$$

- Constraint 2b is for that at most one facility can be assigned to a single grid.
- Constraint 2c is for that a facility must be assigned to a location.
- Constraint 2d is for that facilities  $D_1$  and  $D_4$  must not share an edge i.e. the distances between locations of those facilities should be at least 2.
- Constraint 2e is for that  $D_1$ ,  $D_2$ , and  $D_5$  cannot be placed in grids A,D,L or I.
- Constraint 2f is for that  $D_3$  cannot be placed in grid locations other than A,D,L and I.
- Constraint 2g is for the binary domain of  $x_{ir}$ .

## Part B

Figure 4: Code of Question 2

Figure 5: Output of Question 2

The objective function is 1850 and the optimal grid locations are as follows:

•  $D_1 \to E$ •  $D_2 \to F$ •  $D_3 \to A$ 

D3	D5	D4				D4	D5	D3	
D1	D2						D2	D1	
			•		-				1
				<del>                                     </del>					
			1						1
D1	D2						D2	D1	
D1	D2 D5	D4				D4	D2 D5	D1	

Figure 6: Layout of locations with all mirrored versions