

Hw9 非编程题

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1. Maybe

```
fmap :: (a -> b) -> Maybe a -> Maybe b
fmap f Nothing  = Nothing          (fmap.1)
fmap f (Just a) = Just (f a)      (fmap.2)
```

1. $fmap\ id == id$

case 1:

```
fmap id Nothing === Nothing      (fmap.1)
                === id Nothing   (id)
```

case 2:

```
fmap id (Just a) === Just (id a) (fmap.2)
                === Just a       (id)
                === id (Just a)  (id)
```

Therefore $fmap\ id == id$ for all cases.

2. $fmap\ (f . g) == fmap\ f . fmap\ g$

case 1:

```
fmap (f . g) Nothing      === Nothing          (fmap.1)
(fmap f . fmap g) Nothing === fmap f (fmap g Nothing)  (.)
                        === fmap f Nothing      (fmap.1)
                        === Nothing              (fmap.1)
```

case 2:

```
fmap (f . g) (Just a)      === Just ((f . g) a)      (fmap.2)
(fmap f . fmap g) (Just a) === fmap f (fmap g (Just a)) (.)
                        === fmap f (Just (g a))      (fmap.2)
                        === Just (f (g a))            (fmap.2)
                        === Just ((f . g) a)          (.)
```

Therefore $fmap\ (f . g) == fmap\ f . fmap\ g$ for all cases.

2. []

```
fmap :: (a -> b) -> [a] -> [b]
fmap f []      = []                (fmap.1)
fmap f (a:as) = f a : fmap f as   (fmap.2)
```

1. $fmap\ id == id$

case 1:

```
fmap id [] === []      (fmap.1)
           === id []   (id)
```

case 2:

```
fmap id (x:xs) === id x : fmap id xs   (fmap.2)
               === id x : id xs        (IH)
               === x:xs                 (id)
               === id (x:xs)            (id)
```

Therefore $fmap\ id == id$ for all cases.

2. $fmap\ (f.\ g) == fmap\ f.\ fmap\ g$

case 1:

```
fmap (f . g) []      === []                (fmap.1)
(fmap f . fmap g) [] === fmap f (fmap g []) (.)
                     === fmap f []         (fmap.1)
                     === []                (fmap.1)
```

case 2:

```
fmap (f . g) (x:xs)   === (f . g) x : fmap (f . g) xs   (fmap.2)
(fmap f . fmap g) (x:xs) === fmap f (fmap g (x:xs))      (.)
                       === fmap f (g x : fmap g xs)       (fmap.2)
                       === (f . g) x : fmap f (fmap g xs) (fmap.2)
                       === (f . g) x : (fmap f . fmap g) xs (.)
                       === (f . g) x : fmap (f . g) xs    (IH)
```

Therefore $fmap\ (f.\ g) == fmap\ f.\ fmap\ g$ for all cases.

3. Either

```
fmap :: (a -> b) -> Either c a -> Either c b
fmap _ (Left c)  = Left c          (fmap.1)
fmap f (Right a) = Right (f a)     (fmap.2)
```

1. $fmap\ id == id$

case 1:

```
fmap id (Left a) === Left a          (fmap.1)
                  === id (Left a)     (id)
```

case 2:

```
fmap id (Right a) === Right (id a)    (fmap.2)
                  === Right a         (id)
                  === id (Right a)     (id)
```

Therefore $fmap\ id == id$ for all cases.

2. $fmap\ (f . g) == fmap\ f . fmap\ g$

case 1:

```
fmap (f . g) (Left a)      === Left a          (fmap.1)
(fmap f . fmap g) (Left a) === fmap f (fmap g (Left a)) (.)
                           === fmap f (Left a)     (fmap.1)
                           === Left a             (fmap.1)
```

case 2:

```
fmap (f . g) (Right a)     === Right ((f . g) a)    (fmap.2)
(fmap f . fmap g) (Right a) === fmap f (fmap g (Right a)) (.)
                           === fmap f (Right (g a))    (fmap.2)
                           === Right (f (g a))         (fmap.2)
                           === Right ((f . g) a)       (.)
```

Therefore $fmap\ (f . g) == fmap\ f . fmap\ g$ for all cases.