- Logistic: quiz on Wednesday

Last Time: Security

- Security Properties: Integrity, Confidentiality, Authenticity

Authenticity is necessary for Confidentiality

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Threat Model	Mitigations / Techniques
Accidental corruption	- checksum/CRC
Adversarial modification	<ul><li>Secure hash</li><li>Message Authentication Code (keyed hash)</li></ul>
Replay	Idempotence of messages
Eavesdropping	encryption (AE AD)
	Authenticated encryption requires a pre-established shared secret. For communication with strangers:  - Trusted Third Party (Kerberos/Windows Active Directory): either relay the connection or Trent generates a new secret key and gives that to Alice and Bob  - AKE

peers: Alice + Bobeavesdropper: Eve

adversarial modification: Mallory

- trusted third party: Trent

- AKE:

- 1. Alice generates a key pair:  $(public_{Alice}, private_{Alice})$
- 2. Bob generates a key pair:  $(public_{Rob}, private_{Rob})$
- 3. Alice and Bob publishes their public keys
- 4. Alice sends some  $\boldsymbol{x}_1$  to Bob, and Bob sends some  $\boldsymbol{x}_2$  to Alice
- 5. Alice gets  $AKE(x_1, x_2, public_{Bob}, private_{Alice}) = secret key$  and Bob gets  $AKE(x_1, x_2, public_{Alice}, private_{Bob}) = secret key$ , and the secret keys that Alice gets and Bob gets are the same.
- 6. In addition, knowing  $public_{Alice}$ ,  $public_{Rob}$ ,  $x_1$ ,  $x_2$  does not reveal the secret key.
- But how do we know the public key of say Target?
  - Asking directly from Target does not work, since that message may be corrupted.
  - For a small number of entities, there could be a directory of public keys that were shared in a 100% secure way (e.g. an in-person meeting)
  - Or you could ask someone that you trust and you already know his/her public key
    - $sign(private_{s}, msg) \rightarrow signature$
    - $verify(public_{\downarrow}, msg, signature) \rightarrow bool$

- e.g. You are asking Keith for John's public key
  - $sign(private_{Keith},$  "John's public key is <x> according to Keith (expiring at time t)" )  $\rightarrow signature_{Keith}$
  - verify(public<sub>Keith</sub>, "John's public key is <x> according to Keith (expiring at time t)", signature<sub>Keith</sub>) → true. Then, this is a "certificate" that Keith verifies John's public key is <x>.
  - John can store this certificate, and show this to any person that trusts Keith to prove that John is actually John.
- Firefox —-----TCP----- Target @ "target.com"
  - Firefox trusts a list of certification authorities (whose public keys are programmed into Firefox)
  - When Firefox connects to "target.com", Target, to prove Target is actually Target, would provide:
    - "Hi, I'm target.com. My public key is <x>. Here is a certificate from a CA you trust".
    - And a certificate: "Target come's public key is <x> according to <CA>" + signature<sub>CA</sub> from private<sub>CA</sub>.
  - Firefox:
    - $verify(public_{CA}$ , "Target come's public key is <x> according to <CA>",  $signature_{CA}) \rightarrow true$
  - Then Firefox and Target does AKE to get a shared secret key. This shared secret key is used to do AEAD for all following communication in the current TCP connection.
  - These steps happen as part of the TLS layer. TLS translates between plaintext and ciphetext

## - Q&A

- A: This list of CAs is common across different browsers.
- Q: How does a CA decide to give the certificate to a specific entity? A: CA would have an intensive verification process (back in the days), but over the time the standard has been lowered. Now it's done via domain verification: if someone can put a provided verify.txt at URL/verify.txt within 5 minutes, a CA gives the certificate. This is indeed not secure, since DNS and routing are not secure.
- Q: What if CAs are forced to grant a certificate?
   A: Certificate Transparency Log: a log of all certificates granted by CAs. Big companies monitor this