# The code library of SOWHAT

Version 1.2

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# Content

1	头文件	፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟፟	5
2	数论		5
	2.1	常用公式	5
	2.2	欧几里得算法	5
	2.3	快速欧几里得	5
	2.4	扩展欧几里得	6
	2.5	输出模线性方程 ax=b(modn)的最小解(所有解)	6
	2.6	中国剩余定理	6
	2.7	快速幂模	6
	2.8	欧拉函数	7
	2.9	质因子分解	7
	2.10	素数线性筛法+质因子个数和	8
	2.11	miller rabin 素数测试	8
	2.12	(待补充) 快速斐波那契数	9
3	图论-	——最短路径	9
	3.1	Dijkstra	9
	3.2	SPFA	15
	3.3	Floyd	17
	3.4	Bellman-Ford	19
4	图论—	——最小生成树	20
	4.1	Prim	20
	4.2	Kruskal	22
5	图论—	——二分图匹配	22
	5.1	匈牙利算法	22
	5.2	KM	24
6	图论—	——网络流	28
	6.1	最大流	28
	6	.1.1 EK	28
	6	.1.2 Dinic && SAP	28
	6.2	最小费用最大流	31
	6.3	有上下界的最大流	34
	6.4	混合图的欧拉回路	37
7	图论—	——连通性	40
	7.1	强连通分量	40

		7.1.1	Kosaraju	40
		7.1.2	Tarjan	41
	7.2	双连通	i分量	41
		7.2.1	Tarjan	41
8	数据	居结构		43
	8.1	(待补	充)并查集	43
	8.2	(待补	·充)AC 自动机	43
	8.3	二叉查	找树	43
	8.4	树状数	ɪ组	45
	8.5	线段树	†	46
		8.5.1	注意事项	46
		8.5.2	线段树结点定义	46
		8.5.3	扫描线定义	47
		8.5.4	线段树一般模板	47
		8.5.5	矩形面积交	50
		8.5.6	矩形面积并	51
		8.5.7	线段树求矩形覆盖 K 次交面积(可求面积并)	52
		8.5.8	周长轮廓并	54
		8.5.9	求区间连续递增最长子序列	56
		8.5.10	线段树求区间连续为 N 空白的左端点	59
		8.5.11	线段树区间加法乘法	61
		8.5.12	线段树区间异或,覆盖,最长序列	63
	8.6	划分树	†	66
	8.7	归并树	†	71
	8.8	Treap		73
	8.9	矩形切	]割	75
	8.10	最近公	·共祖先	77
9	计算	淳几何		79
	9.1	注意事	项	79
	9.2	各种公	式	79
	9.3	基本类	型定义	81
	9.4	基础函	j数	81
	9.5	各种极	始排序	81
	9.6	点、直	[线、线段	82
		9.6.1	两直线位置关系	82
		9.6.2	判断两线段是否相交	82
		9.6.3	判断两直线是否相交	83
		9.6.4	判断直线和线段是否相交	83
		9.6.5	判断两向量位置	83
		9.6.6	求两直线交点	83
		9.6.7	求两线段交点	84
		9.6.8	判断点是否在线段上	84

		9.6.9	点到线的距离,点	瓦到线段的最小距离.		84
		9.6.10	判断点是否在线段	注上		85
		9.6.11	求垂点 (不是垂足	크)		85
		9.6.12	通过两点求直线方	5程		85
		9.6.13	向量的旋转			85
		9.6.14	判断点是否在多边	5形内		85
	9.7	三角形				86
		9.7.1	计算三角形面积,	边长		86
		9.7.2	计算三角形的外接	接圆半径,内切圆半径	3 L	87
		9.7.3	各种心(外心,内	<b>內心,垂心,重心,</b> 费	曼马点)	87
	9.8	圆				89
		9.8.1	圆的位置关系			89
		9.8.2	两圆相交面积			90
		9.8.3	判断圆和矩形是否	5相交		90
		9.8.4	最小覆盖圆			91
		9.8.5	扇形重心距圆心路	三离		91
	9.9	多边形				92
		9.9.1	判断两个矩形是否	5相交		92
		9.9.2	多边形面积			92
		9.9.3	多边形的重心			92
	9.10	凸包相关.				92
		9.10.1	二维凸包			92
		9.10.2	三维凸包			94
	9.11	旋转卡壳.				97
	9.12	半平面交.				100
	9.13	整点相关.				01
	9.14	球面相关.				102
	9.15	模拟退火才	这多边形费马点			02
10	DP.					04
	10.1	各种背包.			1	04
	10.2	最长不下降	译子序列			106
	10.3	最长公共于	产序列			107
	10.4	最大子段和	I			108
	10.5	最大子矩阵	<b>F和</b>		1	109
11	杂 7	杂 8				110
	11.1	简单图判定	<u> </u>			110
	11.2	KMP			1	110
	11.3	等价表达式	<b>.</b>		1	111
	11.4	字串种类数	Д(сору)		1	113
	11.5	欧拉回路判	刂断		1	114
	11.6	大数				115

### 1 头文件 && 宏

```
#include <set>
#include <map>
#include <queue>
#include <stack>
#include <math.h>
#include <stdio.h>
#include <stdlib.h>
#include <iostream>
#include <limits.h>
#include <string.h>
#include <string>
#include <algorithm>
#define MID(x,y) ( ( x + y ) >> 1 )
#define L(x) ( x << 1 )
#define R(x) ( x << 1 | 1 )
#define FOR(i,s,t) for(int i=(s); i<(t); i++)
#define BUG puts("here!!!")
#define STOP system("pause")
using namespace std;
int main()
{
   return 0;
}
```

## 2 数论

typedef long long LL;// 可以根据题目要求调整 LL 的类型

#### 2.1 常用公式

```
n!=(n/e)*n * sqrt(2*pi*n)
```

### 2.2 欧几里得算法

```
LL gcd(LL a,LL b) {
    return b==0?a:gcd(b,a%b);
}
```

### 2.3 快速欧几里得

```
LL kgcd( LL a, LL b ){
   if(!a)return b;
   if(!b)return a;
   if( !(a&1) && !(b&1) ) return kgcd(a/2,b/2)<<1;
   else if( !(b&1) ) return kgcd( a,b>>1 );
   else if( !(a&1) ) return kgcd( a>>1,b );
   else return kgcd( abs(a-b),min(a,b) );
}
```

```
//计算机中的负数模(a%n)与数论中的不一样,需要转换
LL rightmod(LL a, LL n){
   return (n+a%n)%n;
2.4 扩展欧几里得
LL ext_euclid(LL a,LL b,LL &x,LL &y){
   LL t,d;
   if( b == 0 ){
      x = 1;
      y = 0;
      return a;
   d = ext_euclid(b,a%b,x,y);
   t = x;
   x = y;
   y = t - (a/b)*y;
   return d;
}
2.5 输出模线性方程 ax=b(modn)的最小解(所有解)
LL modf(LL a,LL b,LL n){
   LL x,y;
   LL d = ext_euclid(a,n,x,y);
   if( b%d ) cout << "FOREVER\n";</pre>
   else
   {
      x = rightmod(x*b/d,n);
      //for(i=0;i<d;i++) print x+i(n/d)所有解
      return rightmod(x,n/d);
   }return 0;
}
2.6 中国剩余定理
a 数组表示 ai, n 表示互质的 ni 数组, k 为个数
LL CRT(LL *a,LL *n,LL k){
   LL i,N = 1,ans=0;
   LL t;
   for(i = 0; i < k; i++)
      N *= n[i];
   for( i = 0; i < k; i++ )
      t = N/n[i]*modf(N/n[i],1,n[i]);
      ans = (ans + a[i]*t)%N;
   return ans;
}
2.7 快速幂模
求 b=a^k%M
LL fastmod(LL a, LL k, LL M){
   LL b = 1;
```

```
while( k ){
       if( k&1 )
           b = a*b%M;
       a = (a\%M)*(a\%M)\%M;
       k /= 2;
   return b;
}
2.8 欧拉函数
LL euler(LL n ){
   LL i,m = (int)sqrt( n + 0.5 ),ans = n;
   for( i = 2; i <= m; i++ )
       if( n%i == 0 )
           ans = ans/i*(i-1);
           while( n\%i == 0 ) n /= i;
   if( n > 1 ) ans = ans/n*(n-1);
   return ans;
}
2.9 质因子分解
count 表示质因子个数,d[i][0]和 d[i][1]分别表示第i个质因子和个数
LL d[100][2];
LL divprime(LL n){
   LL i,count=0,k=n;
   for( i = 2; i*i <= k; i++ )
  if( n%i == 0 ){
           d[count][0] = i;
           d[count][1] = 0;
           while( n%i == 0 )
              n/=i,d[count][1]++;
           count++;
   if(n > 1)
       d[count][0] = n, d[count][1] = 1, count++;
   return count;
}
//搜索质因子的各种组合,即找到 n 所有约数
//调用方式 finddiv(0,1,n,count)
//count 为因子个数,变量 ji 保存的是每次搜到的约数
void finddiv( LL ceng,LL ji,LL n,LL count ){
   if( ceng == count )
       cout<<ji<<endl;</pre>
       return;
   LL i,m;
   for( i = 0,m = 1; i <= d[ceng][1] ; i++ )
       finddiv( ceng+1, ji*m, n, count), m *= d[ceng][0];
}
```

#### 2.10 素数线性筛法+质因子个数和

```
const int SULEN = 1000000;
bool flag[SULEN+1];//是否为素数
int prime[1000001];
int sum[SULEN+1]; //n 的所有质因子的和
void xianxingshai(void)
   int count,i,j;
   fill( flag,flag+SULEN,true );
   fill( sum ,sum +SULEN, 0 );
   flag[0] = flag[1] = false;
   for( count = 0, i = 2; i <= SULEN; i++ )
       if( flag[i] )
       {
           prime[count++] = i;sum[i] = i;
       for(j = 0; j < count && i*prime[j] <= SULEN; <math>j++)
           flag[ i*prime[j] ] = false;
           if( i%prime[j] == 0 ){
               sum[ i*prime[j] ] = sum[i];
              break;
           }
           else
               sum[ i*prime[j] ] = sum[ i ] + prime[j];
       }
   }
//nlogn 素数筛法
int isprime[1000000];
void nlognprime( int n ){
   fill( isprime,isprime+1000000,1 );
   isprime[0] = isprime[1] = 0;
   int i,j,m= (int)sqrt( n+0.5);
   for( i= 2; i <= m; i++ )
       if( isprime[i] )
       for( j = i*i; j <= n; j+=i )
           isprime[j] = 0;
}
```

### 2.11 miller rabin 素数测试

```
LL witness( LL a, LL n ) {
    LL u = n-1,t = 0;
    while(!(u&1)) {u>>=1;t++;}
    int x1,x0 = fastmod(a,u,n);
    for( int i = 1; i <= t; i++ ) {
        x1 = x0*x0%n;
        if( x1 == 1 && x0 != 1 && x0 != n-1 )
            return 1;
        x0 = x1;
    }
    if( x0 != 1 ) return 1;
    else return 0;
}
LL miller_rabin(LL n ,LL s ) {//s 为测试次数
    if( n==2 ) return 1;
```

```
if( n%2==0 || n < 2 ) return 0;
int j,a;
for( j = 1; j <= s; j++ ){
    a = rand()*(n-1)/RAND_MAX + 1;
    if( witness(a,n) ) return 0;
}
return 1;
}

//n 的阶乘中素因子 a 的幂
int aofn( int a,int n ){
    int count = 0;
    while( n ){
        n/=a;
        count += n;
}
return count;
}</pre>
```

#### 2.12 (待补充) 快速斐波那契数

### 3 图论—最短路径

### 3.1 Dijkstra

```
/*
***************
最短路: dijkstra (邻接矩阵表示)
n 表示 点数, MAX 是 n 的上界+10
***************
*/
int map[MAX][MAX];
int Dijkstra(int from, int to, int n) // DIJ + 邻接矩阵
{
   int dis[MAX];
   bool used[MAX];
   memset(used,false,sizeof(used));
   for(int i=1; i<=n; i++)
      dis[i] = INT_MAX;
   dis[from] = 0;
   used[from] = true;
   int now = from;
   for(int i=1; i<n; i++)
      for(int k=1; k<=n; k++)</pre>
          if( map[now][k] && dis[k] > dis[now] + map[now][k] )
             dis[k] = dis[now] + map[now][k];
      int min = INT_MAX;
      for(int k=1; k<=n; k++)
          if( !used[k] && dis[k] < min )</pre>
             min = dis[now = k];
      used[now] = true;
```

```
return dis[to];
/*
***************
最短路: dijkstra (邻接表表示)
n 表示 点数, MAX 是 n 的上界+10
**************
typedef struct NODE{
   int to, len;
   NODE *next;
}NODE;
NODE *p[MAX], node[MAX*MAX];
void Add(int from,int to,int len)
   node[cou].next = p[from];
   node[cou].to = to;
   node[cou].len = len;
   p[from] = &node[cou++];
}
int Dijkstra_List(int from,int to,int n)
   int dis[MAX];
   bool used[MAX];
   memset(used, false, sizeof(used));
   for(int i=1; i<=n; i++)
       dis[i] = INT_MAX;
   dis[from] = 0;
   used[from] = true;
   int now = from;
   for(int i=1; i<=n; i++)
      NODE *head = p[now];
      while( head != NULL )
          int len = head->len;
          int v = head->to;
          if( dis[v] > dis[now] + len )
             dis[v] = dis[now] + len;
          head = head->next;
      int min = INT_MAX;
      for(int k=1; k<=n; k++)</pre>
          if( |used[k]| \&\& dis[k] < min )
             min = dis[now = k];
      used[now] = true;
   return dis[to];
}
最短路: dijkstra + priority(邻接矩阵表示)
记得加头文件!!!!! <queue>
```

```
n 表示 点数, MAX 是 n 的上界+10
**************
*/
int map[MAX][MAX];
typedef struct PRI{
   int u,len;
}RPI;
priority_queue<PRI> Q;
bool operator<(PRI a,PRI b)</pre>
                             // 大于号!!!!
   return a.len > b.len;
}
int Dijkstra_priority(int from,int to,int n)// 出发点,终止点,总节点
   while( !Q.empty() ) Q.pop(); //记得清空
   PRI dis[MAX];
   bool used[MAX];
   memset(used,false,sizeof(used));
   for(int i=1; i<=n; i++)
   {
      dis[i].len = INT MAX;
      dis[i].u = i;
   dis[from].len = 0;
   used[from] = 1;
   int now = from;
   for(int i=1; i<n; i++)</pre>
      for(int k=1; k<=n; k++)
         if( map[now][k] && !used[k] && dis[k].len > dis[now].len +
map[now][k] )
             dis[k].len = dis[now].len + map[now][k];
            Q.push(dis[k]);
         }
      now = Q.top().u;
      Q.pop();
      used[now] = 1;
   return dis[to].len;
}
****************
最短路: dijkstra + priority(邻接矩阵表示)
记得加头文件!!!!! <queue>
n 表示 点数, MAX 是 n 的上界+10
***************
*/
int map[MAX][MAX],dis[MAX];
struct cmp{
   bool operator()(int a,int b)
      return dis[a] > dis[b];
int Dijkstra_priority(int from,int to,int n)// 出发点,终止点,总节点
```

```
{
   priority_queue(int,vector<int>,cmp) q;
   bool used[MAX];
   memset(used,false,sizeof(used));
   for(int i=0; i<n; i++)
      dis[i] = INT_MAX;
   dis[from] = 0;
   while( !q.empty() )
       int now = q.top(); q.pop();
      if( used[now] ) continue;
      used[now] = true;
      for(int k=0; k< n; k++)
          if( map[now][k] && dis[k].len > dis[now].len + map[now][k] )
             dis[k] = dis[now] + map[now][k];
             q.push(k);
   return dis[to].len;
}
*****************
最短路: dijkstra + priority(邻接表表示)
记得加头文件!!!!! <queue>
n 表示 点数, MAX 是 n 的上界+10
***************
*/
typedef struct NODE{
   int to, len;
   NODE *next;
}NODE;
NODE *p[MAX],node[MAX*MAX];
int dis[MAX];
struct cmp{
   bool operator()(int a,int b)
      return dis[a] > dis[b];
};
void Add(int from,int to,int len)
   node[cou].next = p[from];
   node[cou].to = to;
   node[cou].len = len;
   p[from] = &node[cou++];
typedef pair<int,int> pii;
int Dijkstra_List_priority(int from,int to,int n) // 出发点,终止点,总节
点
{
   priority_queue<pii,vector<pii>,greater<pii> > q;
   int dis[MAX];
   bool used[MAX];
   memset(used, false, sizeof(used));
   for(int i=1; i<=n; i++)</pre>
```

```
dis[i] = INT_MAX;
   dis[from] = 0;
   q.push(make_pair(0, from));
   while( !q.empty() )
      int now = q.top().second; q.pop();
      if( used[now] ) continue;
      used[now] = true;
      NODE *head = p[now];
      while( head != NULL )
          int v = head->to;
          int len = head->len;
          if( dis[v] > dis[now] + len )
             dis[v] = dis[now] + len;
             q.push(make_pair(dis[v], v));
          head = head->next;
      }
   return dis[to];
}
****************
最短路: dijkstra + heap
n 表示 点数, MAX 是 n 的上界+10
***************
*/
typedef struct NODE{
   int to, len;
   NODE *next;
}NODE;
NODE *p[MAX], node[MAX*MAX];
void Add(int from,int to,int len)
   node[cou].next = p[from];
   node[cou].to = to;
   node[cou].len = len;
   p[from] = &node[cou++];
}
const int MAX = 1010;
                        // DIJ + HEAP
class HEAP{
   public:
      typedef struct PRIH{
          int u;
          int len;
       }PRIH;
       int heap_size;
      int pre[MAX*2];
      PRIH a[MAX*2];
      void Change(int x,int y)
          pre[a[x].u] = y;
          pre[a[y].u] = x;
          swap(a[x],a[y]);
```

```
void Heap(int i)
   int r,l,larg;
   while(1)
    {
       1 = (i << 1);
       r = 1+1;
       if( 1 <= heap_size && a[1].len < a[i].len )</pre>
           larg = 1;
       else
           larg = i;
       if( r <= heap_size && a[r].len < a[larg].len )</pre>
           larg = r;
       if( larg != i )
           Change(i,larg);
           i = larg;
       else
           break;
   }
void init(int s,int n)
                              // s 是 出发点, n 是点数
   for(int i=0; i<MAX*2; i++)</pre>
    {
       pre[i] = i;
       a[i].u = i;
       a[i].len = INT_MAX;
   a[s].len = 0;
   heap_size = n;
   Heap(1);
void Update(int i)
   int parent;
   while(i > 1)
       parent = (i>>1);
       if( a[i].len < a[parent].len )</pre>
           Change(i,parent);
           i = parent;
       else
           break;
   }
                                 // 更新点 x 的长为 key 的路径
void Modify(int x,int key)
   x = pre[x];
    if( a[x].len <= key )
       return ;
   a[x].len = key;
   Update(x);
                      // 删除最小值
void pop()
```

```
Change(1,heap_size);
          heap_size--;
          Heap(1);
                       // 取堆中最小值
      int Gettop()
      {
          return a[1].u;
      }
};
HEAP heap;
int Dijkstra Heap(int from,int to,int n)
   heap.init(from,n);
   int now = from;
   for(int i=1; i<=n; i++)</pre>
      now = heap.Gettop();
      int rulen = heap.a[1].len;
      heap.pop();
      NODE *head = p[now];
      while( head != NULL )
          int v = head->to;
          int len = head->len;
          int rv = heap.pre[v];
          if( heap.a[rv].len > rulen + len )
             heap.Modify(v,rulen+len);
          head = head->next;
      }
   int rto = heap.pre[to];
   return heap.a[rto].len;
}
3.2 SPFA
最短路: SPFA + 邻接矩阵
n 表示 点数, MAX 是 n 的上界+10
***************
*/
int map[MAX][MAX];
queue<int> q;
int SPFA(int from,int to,int n)
   while( !q.empty() ) q.pop();
   int dis[MAX];
   int cnt[MAX];
   bool inq[MAX];
   memset(cnt,0,sizeof(cnt));
   memset(inq,false,sizeof(inq));
   for(int i=1; i<=n; i++)
      dis[i] = INT MAX;
   dis[from] = 0;
   q.push(from);
   inq[from] = true;
   cnt[from]++;
```

```
int neg = 0;
   while( !q.empty() )
       type now = q.front();
       q.pop();
       inq[now] = false;
       for(int i=1; i<=n; i++)</pre>
          if( map[now][i] && dis[i] > dis[now] + map[now][i] )
              dis[i] = dis[now] + map[now][i];
              if( !inq[i] )
              {
                  inq[i] = true;
                  q.push(i);
                  ++cnt[i];
                  if( cnt[i] > n ) // 判断是否存在负环
                     neg = 1;
                     break;
                  }
              }
       if( neg )
          break;
   if( neg )
       cout << "SPFA say : Exist negative circle !!!!!" << endl;</pre>
       return -1;
   return dis[to];
}
/*
**************
最短路: SPFA + 邻接表
n 表示 点数, MAX 是 n 的上界+10
*/
typedef struct NODE{
   int to, len;
   NODE *next;
}NODE;
NODE *p[MAX],node[MAX*MAX];
void Add(int from,int to,int len)
{
   node[cou].next = p[from];
   node[cou].to = to;
   node[cou].len = len;
   p[from] = &node[cou++];
}
int SPFA_List(int from,int to,int n)
   while( !q.empty() ) q.pop();
   int dis[MAX],cnt[MAX],neg = 0;
   bool inq[MAX];
   for(int i=1; i<=n; i++)
       dis[i] = INT_MAX;
```

```
memset(cnt,0,sizeof(cnt));
   memset(inq,false,sizeof(inq));
   dis[from] = 0;
   q.push(from);
   inq[from] = cnt[from] = 1;
   while( !q.empty() )
      int now = q.front();
      q.pop();
      inq[now] = false;
      NODE *head = p[now];
      while( head != NULL )
          int v = head->to;
          int len = head->len;
          if( dis[v] > dis[now] + len )
             dis[v] = dis[now] + len;
             if(!inq[v])
             {
                inq[v] = true;
                q.push(v);
                cnt[v]++;
                if( cnt[v] > n )
                    neg = 1;
                    break;
             }
          }
          head = head->next;
      if( neg )
          break;
   if( neg )
      cout << "SPFA_List say : Exist negative circle !!!!!" << endl;</pre>
      return -1;
   return dis[to];
3.3 Floyd
最短路: Floyd 邻接矩阵
n 表示 点数, MAX 是 n 的上界+10
这个做法是 map 初始化为 INT MAX 了
***************
int map[max][max];
int Floyd(int from,int to,int n)
{
   for(int i=1; i<=n; i++)
      for(int k=1; k<=n; k++)
          for(int j=1; j<=n; j++)
             if( map[k][i] != INT_MAX && map[i][j] != INT_MAX && map[k][i]
```

```
+ map[i][j] < map[k][j] )
                  map[k][j] = map[k][i] + map[i][j];
   return map[from][to];
Floyd 记录路径
#define INF (1 << 29)
#define N 100
int g[N + 1][N + 1];
int d[N + 1][N + 1];
int n;
int path[N + 1][N + 1];
int path2[N + 1][N + 1];
int begin, end, mid;
int ans;
static void floyd_warshall(void);
static void output(int a, int b);
int
main(void)
   int m;
   int i, j, w;
   while (scanf("%d", &n), n != -1) {
       scanf("%d", &m);
       for (i = 1; i <= n; ++i)
       for (j = 1; j <= n; ++j)
           g[i][j] = INF;
       while (m--) {
           scanf("%d%d%d", &i, &j, &w);
           if (g[i][j] > w)
               g[i][j] = g[j][i] = w;
       }
       memset(path, -1, sizeof(path));
       floyd_warshall();
       if (ans == INF)
           printf("No solution.\n");
       else {
           output(begin, end);
           printf("%d %d\n", end, mid);
   }
   return 0;
}
static void
floyd_warshall(void)
{
   int i, j, k;
   for (i = 1; i <= n; ++i)
```

```
for (j = 1; j <= n; ++j)
      d[i][j] = g[i][j];
   ans = INF;
   for (k = 1; k \le n; ++k) {
      for (i = 1; i \le k - 1; ++i)
      for (j = i + 1; j \le k - 1; ++j) {
          if (ans > d[i][j] + g[i][k] + g[k][j]) { /* 引入 k 这个点构成环 */
             ans = d[i][j] + g[i][k] + g[k][j];
             begin = i, end = j, mid = k; /* i, j需要 output 输出, k单独
输出 */
             memcpy(path2, path, sizeof(path)); /* 保存当前最优的路径 */
          }
      }
      for (i = 1; i <= n; ++i)
      for (j = 1; j \le n; ++j) {
          if (d[i][j] > d[i][k] + d[k][j]) {
             d[i][j] = d[i][k] + d[k][j];
             path[i][j] = k;
          }
      }
   }
}
static void
output(int a, int b)
{
   if (path2[a][b] == -1) {
      printf("%d ", a);
      return;
   }
   output(a, path2[a][b]);
   output(path2[a][b], b);
}
3.4 Bellman-Ford
最短路: Bellman-ford
n 表示 点数, MAX 是 n 的上界+10
**************
*/
void BAdd(int from,int to,int len)// bellman
   e[coue].u = from;
   e[coue].v = to;
   e[coue++].w = len;
int Bellman(int from,int to,int n)
{
   int dis[MAX];
   bool used[MAX];
   int neg = 0;
   memset(used, false, sizeof(used));
   for(int i=1; i<=n; i++)
```

```
dis[i] = INT_MAX;
   dis[from] = 0;
   for(int i=1; i<n; i++)
       for(int k=0; k<coue; k++)</pre>
           int v = e[k].v;
           int u = e[k].u;
           int len = e[k].w;
           if( dis[u] != INT MAX && dis[v] > dis[u] + len )
               dis[v] = dis[u] + len;
       }
    for(int i=1; i<n && !neg; i++)// 判断是否是负环
       for(int k=0; k<coue && !neg; k++)</pre>
           if(dis[e[k].u] != INT_MAX && dis[e[k].v] > dis[e[k].u] + e[k].w)
               neg = 1;
   if( neg )
       cout << "BellmanFord say : Exist negative circle !!!!!" << endl;</pre>
       return -1;
   return dis[to];
}
```

### 4 图论—最小生成树

#### 4.1 **Prim**

```
/*
实例化: Prime< 100 , 10000 > Net;
初始化: Net.init( 30 );//点的数量
加边, Net.add( from, to, w );//directional..
求 W(mst); w = Net.prim();
*/
class Edge {
public:
   int to,w;
   Edge *next;
   void add( int t, int ww,Edge *&b){
       to = t;w = ww;next=b;b=this;
};
template<int Vsize, int Esize>
class Prim{
private:
   int countt, N;
   Edge *biao[ Vsize ],A[ 3*Esize ];
   int cnt[ Vsize ];// how many nodes refresh it
   int d[ Vsize ],flag[ Vsize ];
public:
   void init( int n ){
       countt = 0;N = n;
       for( int i = 0; i <= N; i++ )
           biao[i] = NULL;
   }
```

```
void add( int from,int to,int w ){
       A[ countt++ ].add( to, w, biao[from] );
   }
   int prim( void ){
       int i,j,minn,sum = 0,now = 1;
       memset( flag, 0, sizeof(flag) );
       memset( cnt, 0, sizeof(cnt) );
       fill( d, d+Vsize, INF );
       d[now] = 0;
       flag[ now ] = 1;
       for(i = 1; i < N; i++){
          for(Edge *p = biao[now]; p!=NULL; p = p->next )
              if( !flag[p->to] ){
                 if( d[p->to] > p->w ){
                     d[p->to] = p->w;
                     cnt[p->to] = 1;
                 else if( d[p->to] == p->w ){
                     cnt[p->to] ++;
              }
          minn = INF;
          for( j = 1; j <= N; j++ )
              if( !flag[j] && d[j] < minn )</pre>
                 minn = d[now = j];
          flag[ now ] = 1;
          //if( cnt[now] > 1 ) return -1;//返回-1说明 mst 不唯一
          sum += d[now];
       }
       return sum;
   }
};
最小生成树 Prim 算法 ,n 为点的个数,起始点从1开始
*/
int map[MAX][MAX],n;
int Prim( int map[][], int n )
   int used[MAX],dis[MAX],i,j,sum = 0,now,min;
   memset(used,0,sizeof(used));
   fill(dis,dis+MAX,INT_MAX);
   now = 1; dis[now] = 0; used[now] = 1;
   for(i=1; i<n; i++)
   {
       for(j=1; j<=n; j++)
          if( !used[j] && dis[j] > map[now][j] )
              dis[j] = map[now][j];
       min = INT MAX;
       for(j=1; j<=n; j++)
          if( !used[j] && dis[j] < min )
              min = dis[now = j];
       used[now] = 1;
       sum += min;
   }
```

```
return sum;
4.2 Kruskal
************************************
最小生成树 Kruskal 算法 ,n 为点的个数,e 为边数
记得头文件!!! <algorithm> !!!!!!!
***************
int pre[MAX]; // 记录父节点
typedef struct NODE
               // x, y 为边的两个顶点
   int x,y,len;
}NODE;
NODE edge[LALA]; // 边的条数根据题意
int find(int x)
   while( x != pre[x] )
      x = pre[x];
   return x;
int cmp ( NODE a ,NODE b)
   return a.len < b.len;
int Kruskal()
   int i, sum = 0, a, b;
   sort(edge,edge+e,cmp);
   for(i=0; i<n; i++)
      pre[i] = i;
   for(i=0; i<e; i++)
      a = find( edge[i].x );
      b = find( edge[i].y );
      if( a != b )
          sum += edge[i].len;
         pre[b] = a;
      }
   return sum;
}
```

## 5 图论—二分图匹配

### 5.1 匈牙利算法

```
int used[MAX],mat[MAX];
void init()
{
   memset(map,0,sizeof(map));
int Augment(int s,int n,int x)
   int i;
   for(i=s; i<=n; i++)</pre>
      if( !used[i] && map[x][i] )
          used[i] = 1;
          if( mat[i] == -1 || Augment(s,n,mat[i]) )
             mat[i] = x;
             return 1;
       }
   return 0;
}
int Hungary(int s,int n)
   int i, sum = 0;
   memset(mat,-1,sizeof(mat));
   for(i=s; i<=n; i++)</pre>
      memset(used,0,sizeof(used));
      if( Augment(s,n,i) )
          sum++;
   return sum;
}
二分图最大匹配 匈牙利算法 数组模拟邻接表
s 为数组起始点,n 为终止点,len[v] 存以 v 为起点的个数
建图用 map[x][++len[x]] = y;
*/
int map[MAX][MAX];
int used[MAX],mat[MAX],len[MAX];
void init()
{
   memset(len,0,sizeof(len));
   memset(map,0,sizeof(map));
int Augment(int x)
   int i,k;
   for(i=1; i<=len[x]; i++)
   {
      k = map[x][i];
      if( !used[k] )
          used[k] = 1;
          if( mat[k] == -1 || Augment(mat[k]) )
             mat[k] = x;
```

```
return 1;
            }
        }
    return 0;
}
int Hungary(int s,int n)
    memset(mat,-1,sizeof(mat));
    int i, sum = 0;
    for(i=s; i<=n; i++)
        memset(used,0,sizeof(used));
        if( Augment(s,i) )
            sum++;
    return sum;
}
5.2 KM
#define INF (1 << 29)
#define N 150
int g[N][N];
int n;
static int lx[N];
static int ly[N];
static char usedx[N];
static char usedy[N];
static int match[N];
static int slack[N];
static void km(void);
int
main(void)
    int i, j;
    int sum;
    scanf("%d", &n);
    for (i = 0; i < n; ++i)
for (j = 0; j < n; ++j)
    scanf("%d", &g[i][j]);</pre>
    km();
    sum = 0;
    for (i = 0; i < n; ++i)
        sum += g[i][match[i]];
    printf("%d\n", sum);
    return 0;
}
```

```
static int
find(int x)
   int y;
   int t;
   usedx[x] = 1;
   for (y = 0; y < n; ++y) {
       if (usedy[y])
           continue;
       t = 1x[x] + 1y[y] - g[x][y];
       if (t == 0) {
           usedy[y] = 1;
           if (match[y] == -1 || find(match[y])) {
               match[y] = x;
               return 1;
           }
       } else {
   if (slack[y] > t)
               slack[y] = t;
       }
   return 0;
}
static void
km(void)
{
   int x, y;
   int i, j;
   int d;
   memset(match, -1, sizeof(match));
   memset(lx, 0, sizeof(lx));
   memset(ly, 0, sizeof(ly));
   for (i = 0; i < n; ++i)
   for (j = 0; j < n; ++j)
if (g[i][j] > lx[i])
           lx[i] = g[i][j];
   for (x = 0; x < n; ++x) {
       for (i = 0; i < n; ++i)
           slack[i] = INF;
       while (1) {
           memset(usedx, 0, sizeof(usedx));
           memset(usedy, 0, sizeof(usedy));
           if (find(x))
               break;
           d = INF;
           for (y = 0; y < n; ++y) {
               if (!usedy[y] \&\& d > slack[y])
                   d = slack[y];
           for (i = 0; i < n; ++i) {
               if (usedx[i])
                   lx[i] -= d;
           }
```

```
for (i = 0; i < n; ++i) {
              if (usedy[i])
                  ly[i] += d;
              else
                  slack[i] -= d;
          }
       }
   }
}
KM(copy)
const int MAX = 1024;
int n;
                    // X 的大小
                             // X 到 Y 的映射(权重)
int weight [MAX] [MAX];
                             // 标号
int lx [MAX], ly [MAX];
                         // 是否被搜索过
bool sx [MAX], sy [MAX];
                    // Y(i) 与 X(match [i]) 匹配
int match [MAX];
// 初始化权重
void init (int size);
// 从 X(u) 寻找增广道路,找到则返回 true bool path (int u);
// 参数 maxsum 为 true , 返回最大权匹配, 否则最小权匹配
int bestmatch (bool maxsum = true);
void init (int size)
{
   // 根据实际情况,添加代码以初始化
   n = size;
   for (int i = 0; i < n; i ++)
       for (int j = 0; j < n; j ++)
           scanf ("%d", &weight [i] [j]);
}
bool path (int u)
   sx [u] = true;
   for (int v = 0; v < n; v ++)
       if (!sy [v] && 1x[u] + 1y [v] == weight [u] [v])
           sy [v] = true;
           if (match [v] == -1 || path (match [v]))
              match [v] = u;
              return true;
              }
   return false;
}
int bestmatch (bool maxsum)
   int i, j;
   if (!maxsum)
       for (i = 0; i < n; i ++)
           for (j = 0; j < n; j ++)
              weight [i] [j] = -weight [i] [j];
```

```
}
   // 初始化标号
   for (i = 0; i < n; i ++)
       lx[i] = -0x1FFFFFFF;
       ly [i] = 0;
       for (j = 0; j < n; j ++)
          if (lx [i] < weight [i] [j])</pre>
              lx [i] = weight [i] [j];
       }
   memset (match, -1, sizeof (match));
   for (int u = 0; u < n; u ++)
       while (1)
           {
          memset (sx, 0, sizeof (sx));
          memset (sy, 0, sizeof (sy));
           if (path (u))
              break;
           // 修改标号
           int dx = 0x7FFFFFFF;
           for (i = 0; i < n; i ++)
              if (sx [i])
                  for (j = 0; j < n; j ++)
                     if(!sy [j])
                         dx = min (lx[i] + ly [j] - weight [i] [j], dx);
          for (i = 0; i < n; i ++)
              if (sx [i])
                 lx [i] -= dx;
              if (sy [i])
                 ly [i] += dx;
          }
   int sum = 0;
   for (i = 0; i < n; i ++)
       sum += weight [match [i]] [i];
   if (!maxsum)
       {
           sum = -sum;
           for (i = 0; i < n; i ++)
              for (j = 0; j < n; j ++)
                 weight [i] [j] = -weight [i] [j];  // 如果需要保
持 weight [ ] [ ] 原来的值,这里需要将其还原
       }
   return sum;
```

}

### 6 图论—网络流

#### 6.1 最大流

#### 6.1.1 EK

```
/*
***************
最大流—EK 算法 (记得初始化 cap) 邻接矩阵形式
int cap[MAX][MAX];
                      //建图的时候用 cap[from][to] += 流量
                      //cap 存的是流量
int EKarp(int s,int t)
                      // 起点,终点
                         // a 存每次增广的流量
   queue<int> Q;
   int flow[MAX][MAX],a[MAX],u,v,f,pre[MAX];
                      // 增广的流量和
   memset(flow,0,sizeof(flow));
   while(1)
   {
      Q.push(s);
      memset(a,0,sizeof(a));
      a[s] = INT_MAX;
      while( !Q.empty() )
         u = Q.front();
         Q.pop();
         for(v=1; v<=m; v++)
            if( !a[v] && cap[u][v] > flow[u][v] )
                Q.push(v);
                a[v] = a[u] < cap[u][v] - flow[u][v] ? a[u] : cap[u][v]
- flow[u][v];
                pre[v] = u;
            }
      if( a[t] == 0 )
         break;
      for(u=t; u!=s; u=pre[u])
         flow[pre[u]][u] += a[t];
         flow[u][pre[u]] -= a[t];
      f += a[t];
   return f;
}
        6.1.2 Dinic && SAP
const long long INF = 210000000000056011;
typedef long long Tp;
using namespace std;
class Arc{
public:
   int to;
```

```
Tp w;
   Arc *next,*anti;
   Arc *ass( int tt,Tp ww,Arc* &b ){
       to = tt;w = ww;next = b;b = this;return this;
   }
};
template<class Tp,int Vsize, int Esize >
class Network{
private:
   Arc A[ 2*Esize+10 ],*biao[ Vsize ];
   int countt,d[ Vsize ],S,T,N;
   int bfs( void ){
       queue<int> q;
       fill( d,d+Vsize,-1 );
       d[S] = 0;
       q.push(S);
       while( !q.empty() ){
           int u = q.front();q.pop();
           for( Arc *p=biao[u]; p != NULL; p=p->next )
   if( d[p->to] == -1 && p->w > 0 )
                   d[p->to] = d[u]+1,q.push(p->to);
       return d[T] != -1;
   Tp dinic( int u,Tp sum ){
       Tp f,o;
       if( u == T ) return sum;
       o = sum;
       for( Arc *p=biao[u]; p != NULL && sum; p=p->next )
           if( d[p->to] == d[u]+1 && p->w > 0 ){
               f = dinic(p->to,min(p->w,sum));
               p->w-=f;p->anti->w+=f;
               sum -= f;
       return o-sum;
public:
   Network( void ) {}
   Network( int n ) { init(n) ;}
   void init( int n ){
       countt = 0;N = n;
       for( int i = 0; i <= n; i++ )
           biao[i] = NULL;
   void add( int from, int to, Tp w ){
       Arc *p1 = A[countt++].ass( to,w ,biao[from]);
       Arc *p2 = A[countt++].ass( from,0,biao[ to] );
       p1->anti = p2; p2->anti = p1;
   Tp MaxFlowDinic( int s, int t ){
       S = s;T = t;
       Tp total = 0;
       while( bfs() )
           total += dinic( S,INF );
       return total;
   Tp MaxFlowSap( int s, int t){
       S = s;T = t;
       int i,md;
```

```
Tp now,total;
       Arc *p,*locate,*ge[ Vsize ],*di[ Vsize ],*path[ Vsize ];
       int dist[ Vsize ], cd[ Vsize ];
       int his[ Vsize ], pre[ Vsize ];
       bool flag;
       memset( dist,0, sizeof(dist) );
       memset( cd, 0, sizeof(cd) );
       cd[0] = N;
       for(i = 0; i <= N; i++)
           di[i] = ge[i] = biao[i];
       for( total = 0, now = INF, i = S; dist[i] < N; ){
          his[i] = now;
           for( flag = false,p=di[i]; p!=NULL; p= p->next){
              if(p->w>0 \&\& dist[i] ==dist[p->to] + 1){
                  now = min(now,p->w);
                  pre[ p->to ] = i;
                  path[ p->to ] = p;
                  di[i] = p;
                  i = p->to;
                  if( i == T ){
                      for( total += now; i != S; i = pre[i] )
                         path[i]->w -= now, path[i]->anti->w += now;
                      now = INF;
                  flag = true;
                  break;
              }
           if(!flag){
              for( md = N-1,p=ge[i];p!= NULL;p=p->next )
                  if(p->w > 0 \&\& dist[p->to] < md)
                     md = dist[p->to],locate = p;
              di[i] = locate;
              if( !(--cd[ dist[i] ] ) ) break;
              cd[ dist[i] = md+1 ] ++;
              if( i != S )
                  i = pre[i],now = his[i];
           }
       }
       return total;
   int check( int u ){
       bool flag[Vsize];
       memset( flag, 0, sizeof(flag) );
       dfs( S,flag );
       return count( flag+1,flag+1+N,true);
   void dfs( int u ,bool *flag){
       flag[u] = true;
       for( Arc *p=biao[u];p != NULL; p = p->next ){
           if( p->w == 0 \mid | flag[p->to])
              continue;
           dfs( p->to ,flag);
       }
   }
//最大流题目构图:
//1.最大权闭合子图:
// 对于 DAG 图中的边 u->v, 建立 u 到 v 容量为 INF 的弧
// 正权点 U 连接 S->U,容量为 Weight (U)
```

```
// 负权点 V 连接 V->T,容量为-Weight (V)
// 闭合图的权 = 正权点权值和 - 最大流
// 隐含关系,一旦有 u,必定有 v。。
//2.二分图最小点权覆盖:
// 隐含关系,对于每条边的两个点,至少一个是在覆盖中的
// 对于每个点 x 输入 X,连接 S 到 x,容量为点权
// 对于每个点 y 属于 Y, 连接 y 到 T, 容来为点权
// 对于每条边 u->v,连接一条容量无限的弧
//3.二分图最大点权独立集:
// 隐含关系,在独立集中没有两个点之间有边
// 最大点权独立集 = 总的点权 - 最小点权覆盖
//4.二分加最大流,二分边的容量等判断最大流是否有答案
6.2 最小费用最大流
const int maxn = 110;
const int inf = 1008610086;
using namespace std;
class Arc{
public:
   int to,w,cost;
   Arc *next,*anti;
   Arc *add( int tt,int ww,int cc ,Arc *&b){
       to = tt; w = ww; cost = cc;
       next = b;b = this;return this;
};
class Network{
   public:
   Arc a[ 6000 ],*biao[maxn];
   int d[maxn],i,countt;
   Network() {}
   Network(int n){init(n);}
   void init(int n ){
       countt = 0;
       for( i = 0; i < n; i++)
          biao[i] = NULL;
   void add( int from, int to, int w, int c ){
       Arc *p1 = a[ countt++ ].add( to, w, c, biao[from] );
       Arc *p2 = a[ countt++ ].add( from, 0, -c, biao[to]);
       p1->anti = p2;
       p2->anti = p1;
   }
   int minimum cost maxflow(int s,int t){
       queue<int> q;
       int i,u,cost=0,flow=0;
       while(1){
          bool inq[maxn];
          int prep[maxn];
          int alow[maxn];
          Arc *prea[maxn];
          fill( d, d+maxn, inf );
          memset( inq, 0, sizeof(inq) );
```

```
alow[s] = inf; d[s] = 0;
          q.push( s );
          while( !q.empty() ){
             u = q.front();q.pop();
              inq[u] = false;
             for( Arc *p=biao[u];p!=NULL;p=p->next )
                 if( p->w && d[p->to] > d[u] + p->cost ){
                    d[p->to] = d[u] + p->cost;
                    prep[p->to] = u;
                    prea[p->to] = p;
                    alow[p->to] = min(alow[u], p->w);
                    if( !inq[p->to] ){inq[p->to]=true;q.push(p->to);}
                 }
          //bellmanford is upper
          if( d[t] == inf ) break;
          for( u = t; u != s; u = prep[u] ){
             prea[u]->w -= alow[t];
             prea[u]->anti->w += alow[t];
          cost += d[t]*alow[t];
          flow += alow[t];
      return cost;
   }
   bool full( int s ){
      for( Arc *p = biao[s]; p != NULL; p = p->next )
          if( p->w ) return false;
      return true;
   }
};
******************
最小费用最大流 数组模拟邻接表 点的起点为 1, 共 n 个点
****建图的时候,先初始化 init(),再调用
Add(from,to,c,pay); 如果是双向边还需建 Add(to,from,c,pay)
*/
typedef struct NODE{
   int from,to,cap,cost;
   int next;
}NODE;
NODE node[LALA];
                        // 根据题意
                        // 相当于指针数组
int p[MAX];
int cou,n,m;
void init()
{
   memset(p,-1,sizeof(p));
   memset(node,'\0',sizeof(node));
                        // 初始化为偶数
   cou = 2;
void Add(int u,int v,int cap,int cost)
   node[cou].from = u;
   node[cou].to = v;
   node[cou].cap = cap;
   node[cou].cost = cost;
   node[cou].next = p[u];
```

```
p[u] = cou++;
   node[cou].from = v;
   node[cou].to = u;
   node[cou].cap = 0;
   node[cou].cost = -cost;
   node[cou].next = p[v];
   p[v] = cou++;
int MincostMaxflow(int s,int t,int n )
   queue<int> q;
   int inq[MAX],pre[MAX],dis[MAX],re[MAX];
   int u,v,i,a,c = 0,ind,cost,cap;
   while(1)
   {
       memset(inq,0,sizeof(inq));
       fill(dis,dis+MAX,INT_MAX);
       dis[s] = 0;
       inq[s] = 1;
       pre[s] = s;
       q.push(s);
       while( !q.empty() )
           u = q.front();
           q.pop();
           inq[u] = 0;
           ind = p[u];
           while( ind != -1 )
           {
               u = node[ind].from;
               v = node[ind].to;
               cost = node[ind].cost;
               cap = node[ind].cap;
               if(cap > 0 \&\& dis[v] > dis[u] + cost)
                   dis[v] = dis[u] + cost;
                   pre[v] = u;
                   re[v] = ind;
                   if( !inq[v] )
                   {
                       q.push(v);
                       inq[v] = 1;
                   }
               ind = node[ind].next;
           }
       if( dis[t] >= 0 )
           break;
       a = INT MAX;
       for(u=t; u!=s; u=pre[u])
           if( node[re[u]].cap < a )</pre>
               a = node[re[u]].cap;
       for(u=t; u!=s; u=pre[u])
           node[re[u]^1].cap += a;
           node[re[u]].cap -= a;
       c += dis[t]*a;
```

```
return -c;
6.3 有上下界的最大流
const int inf = 2100000000;
class Arc{
public:
   int to,w;
   int 1,r;
   Arc *next,*anti;
   Arc *add( int tt,int ll,int rr,Arc *&b ){
       to = tt; l = ll; r = rr;
       w = rr-ll;next = b;b = this; return this;
   }
};
template<int Vsize, int Esize>
class BoundNetwork{
public:
   int d[ Vsize ],S,T,bound,countt,in[ Vsize ],out[ Vsize ];
   Arc a[ Esize ],*biao[ Vsize ];
   void init( int n ){
       bound = n;//S = n - 2;T = n - 1;
       for( int i = 0; i < n; i++ )
           biao[i] = NULL;
       memset( in, 0, sizeof(in) );
       memset( out, 0, sizeof(out) );
       countt = 0;
   }
   void add( int from,int to, int l,int r ){
       Arc *p1 = a[ countt++ ].add( to, l, r, biao[from] );
       Arc *p2 = a[ countt++ ].add( from, 0, 0, biao[to] );
       p1->anti = p2;
       p2->anti = p1;
   }
   int bfs( void ){
       queue<int> q;
       fill( d,d+Vsize,-1 );
       d[S] = 0;
       q.push(S);
       while( !q.empty() ){
           int u = q.front();q.pop();
           for( Arc *p=biao[u]; p!=NULL; p=p->next )
              if( d[ p->to ] == -1 \&\& p->w > 0 )
                  d[p->to] = d[u] + 1,q.push(p->to);
       return d[ T ] != -1;
   int dinic( int u ,int sum){
       int f,o = sum;
       if( u == T ) return sum;
       for( Arc *p = biao[u]; p!=NULL && sum; p=p->next )
           if(d[p->to] == d[u] + 1 && p->w > 0){
```

```
f = dinic(p->to, min(p->w,sum));
           p->w-=f;
           p->anti->w += f;
           sum -= f;
   return o - sum;
}
int max_flow1( int s,int t ){
   S = s;T = t;
   int total = 0;
   while( bfs() )
       total += dinic( S, 2100000000);
   return total;
}
int max_flow(int s,int t){
   S = s;T = t;
   int i,now_flow,total,md;
   Arc *ge[Vsize],*di[Vsize],*path[Vsize];
   int dist[Vsize],cd[Vsize],his[Vsize],pre[Vsize],n=bound;
   Arc *p,*locate;
   bool flag;
   memset( dist,0,sizeof(dist) );
   memset( cd,0,sizeof( cd) );
   cd[0] = n;
   for( i = 0; i < n ; i++ ) di[i] = ge[i] = biao[i];
   for( total = 0, now_flow = inf,i = S; dist[i] < n; ){</pre>
       his[i] = now flow;
       for( flag = false,p=di[i]; p!=NULL; p= p->next){
           if(p->w>0 \&\& dist[i] ==dist[p->to] + 1){
               now_flow = min( now_flow,p->w );
               pre[ p->to ] = i;
               path[ p->to ] = p;
               di[i] = p;
               i = p \rightarrow to;
               if( i == T ){
                  for( total += now_flow; i != S; i = pre[i] ){
                      path[i]->w -= now_flow;
                      path[i]->anti->w += now_flow;
                  now_flow = inf;
               flag = true;
               break;
       if( !flag ){
           for( md = n-1,p=ge[i];p!= NULL;p=p->next )
               if(p->w>0 \&\& dist[p->to] < md)
                  md = dist[p->to],locate = p;
           di[i] = locate;
           if( !(--cd[ dist[i] ] ) ) break;
           cd[ dist[i] = md+1 ] ++;
           if( i != S )
               i = pre[i],now_flow = his[i];
       }
   return total;
```

```
}
   void construct( int specialS,int specialT ){
       for( int i = 0; i < bound - 2; i++ )
           for( Arc *p=biao[i]; p!=NULL; p = p->next ){
              in[ p->to ] += p->l;
              out[ i ] += p->l;
       for( int i = 0; i < bound - 2; i++){
           add( specialS, i, 0, in[i] );
           add( i, specialT, 0, out[i] );
       }
   }
   bool full( void ){
       for( Arc *p = biao[S]; p != NULL; p = p->next )
           if( p->w > 0 ) return false;
       return true;
};
//示例程序 zoj 3229 Shoot the bullet ~
BoundNetwork< 1500, 1500*80 > Net;
int main(void){
   int n,m,from,to;
   int i,j,c,d,r,l;
   while( scanf("%d%d",&n,&m) != EOF ){
       Net.init(n+m+4);
       for( i = 0; i < m; i++){
           scanf("%d",&j);
           Net.add( n+i, n+m+1, j, inf );
       }
       queue<int> q;
       for( i = 0; i < n; i++){
           scanf("%d%d",&c,&d);
           Net.add( n+m, i, 0, d );
           while (c--){
              scanf("%d%d%d",&j,&r,&l);
               q.push( Net.countt );
              Net.add( i, n+j, r, 1 );
           }
       j = Net.countt;
       Net.add( n+m+1, n+m, 0, inf );
       Net.construct(n+m+2,n+m+3);
       Net.max_flow(n+m+2,n+m+3);
       if( !Net.full() ){
           puts("-1\n");continue;
       Net.a[j].w = Net.a[j+1].w = 0;
       Net.max flow1(n+m,n+m+1);
       int total = 0;
       for( Arc *p = Net.biao[Net.S]; p != NULL; p = p->next )
           if( p->to >= 0 \&\& p->to < n ) {
              total += p->r - p->w;
       printf("%d\n",total);
       while( !q.empty() ){
           int u = q.front();q.pop();
           printf("%d\n",Net.a[u].r - Net.a[u].w );
```

```
}
    printf("\n");
}
    return 0;
}
```

## 6.4 混合图的欧拉回路

大致就是, 先将无向边定向, 就是比如 1<->3, 可以定它的方向为 1->3, 1 的出度++, 3 的入度++即可。

读入的时候如果遇到无向边,把这条边加入待建的网络中,流量为 1。读入完后,然后用出度减入度得到 x,如果 x 为奇数,肯定不存在欧拉回路,如果没有奇数,就用最大流求解。

如果 x 大于 0,则建一条 s (源点)到当前点容量为 x/2 的边,如果 x 小于 0,建一条从当前点到 t (汇点)容量为|x/2|的边。

然后求最大流,如果是满流(即 s 出的流==t 入的流即可,s 指的是建图的时候连接 s 的边的容量和)就满足欧拉回路。

```
int pre[MAX],n;
int ind[MAX],outd[MAX];
int lev[MAX];
typedef struct MAP{
   int cap, to;
   int next;
}MAP;
MAP node[3000];
int head[MAX];
int cou;
void init()
{
   cou = 2;
   memset(node,'\0',sizeof(node));
   memset(head,-1,sizeof(head));
   memset(ind,0,sizeof(ind));
   memset(outd,0,sizeof(outd));
   for(int i=1; i<=n; i++)</pre>
       pre[i] = i;
int find(int x)
   while( x != pre[x] )
       x = pre[x];
   return x;
void Union(int x,int y)
   int a = find(x);
   int b = find(y);
   if( a == b )
       return ;
   int p = min(a,b);
   pre[a] = pre[b] = pre[x] = pre[y] = p;
int check()
```

```
{
   for(int i=1; i<=n; i++)</pre>
       if( find(i) != 1 )
           return 0;
   return 1;
void Add(int u,int v,int cap)
{
    node[cou].to = v;
    node[cou].cap = cap;
    node[cou].next = head[u];
    head[u] = cou++;
    node[cou].to = u;
    node[cou].cap = 0;
    node[cou].next = head[v];
    head[v] = cou++;
queue<int> q;
int BFS(int s,int t)
   int p,u,v,cap;
   memset(lev,-1,sizeof(lev));
   q.push(s);
   lev[s] = 0;
   while( !q.empty() )
   {
       u = q.front();
       q.pop();
       p = head[u];
       while (p!=-1)
       {
           v = node[p].to;
           cap = node[p].cap;
           if( cap > 0 \& lev[v] == -1 )
               lev[v] = lev[u] + 1;
               q.push(v);
           p = node[p].next;
       }
   return lev[t] != -1;
int Dinic(int k,int sum,int s,int t)
   int i,a,os;
   if(k == t)
       return sum;
   os = sum;
   int p = head[k];
   while(p != -1 \&\& sum)
   {
       int to = node[p].to;
       int cap = node[p].cap;
       if( lev[to] == lev[k] + 1 && cap > 0 )
           a = Dinic(to,min(sum,cap),s,t);
           node[p^1].cap += a;
           node[p].cap -= a;
```

```
sum -= a;
       }
       p = node[p].next;
   return os - sum;
int main()
   int m, from, to, s;
   int ncases;
   scanf("%d",&ncases);
   while( ncases-- )
       scanf("%d%d",&n,&m);
       init();
       while( m-- )
           scanf("%d%d%d",&from,&to,&s);
           if( from == to )
               continue;
           ind[to]++;
           outd[from]++;
           Union(from, to);
           if( s != 1 )
               Add(from, to, 1);
       if( !check() )
           printf("impossible\n");
           continue;
       int flag = 1;
       int sum = 0;
       for(int i=1; i<=n; i++)
           outd[i] -= ind[i];
           if( outd[i] % 2 == 1 )
               flag = 0;
               break;
           outd[i] /= 2;
           if(outd[i] > 0)
               Add(0,i,outd[i]);
               sum += outd[i];
           else
               Add(i,n+1,-outd[i]);
       if(!flag)
           printf("impossible\n");
           continue;
       int ans = 0;
       while( BFS(0,n+1) )
           ans += Dinic(0,INT_MAX,0,n+1);
       if( ans == sum )
           printf("possible\n");
```

## 7 图论—连通性

## 7.1 强连通分量

## 7.1.1 Kosaraju

```
class Edge{
public:
   int to;
   Edge *next;
   void add( int tt, Edge *&b ){to=tt;next=b;b=this;}
};
template< int Esize, int Vsize >
class Kosarajo{
public:
   Edge *biao[ Vsize ],a[ 2*Esize+10 ],*fan[ Vsize ];
   int flag[ Vsize ],countt,N,M;
   int TIME,GROUP,label[ Vsize ];
   stack<int> S;
   void init(void){
       for(int i = 0; i < N; i++){
           biao[i] = fan[i] = NULL;
           flag[i] = 0;
       GROUP = countt = 0;
   }
   void dfs1( int u ){
       flag[u] = 1;
       for( Edge *p=biao[u];p!=NULL;p=p->next )
           if( !flag[p->to] ) dfs1(p->to);
       S.push( u );
   void dfs2( int u,int group ){
       flag[u] = 1;
       for( Edge *p = fan[u]; p!=NULL; p=p->next )
           if( !flag[p->to] ) dfs2(p->to,group);
       label[u] = group;
   }
   int getscc(){
       for( int i = 0; i < N; i++ )
           if( !flag[i] ) dfs1(i);
       while( !S.empty() ){
           int u = S.top();S.pop();
           if( !flag[u] ) dfs2(u,GROUP++);
       return GROUP;
```

```
}
};
          7.1.2 Tarjan
class Edge{
public:
   int to;
   Edge *next;
   void add( int tt, Edge *&b ){to=tt;next=b;b=this;}
template<int Vsize, int Esize>
class TarjanScc{
public:
   Edge *biao[ Vsize ],a[ Esize*2+10];
   int flag[ Vsize ],countt,N,M,d[ Vsize ];
   int TIME,GROUP,label[ Vsize ],in[ Vsize ],low[ Vsize ];
   stack<int> S;
   void init(void){
       for(int i = 0; i < N; i++ )
           biao[i] = NULL,in[i]=flag[i]=0;
       TIME = GROUP = countt = 0;
   void add( int from, int to ){
       a[ countt++ ].add( to, biao[from] );
   void scc( int u ){
       low[u] = d[u] = TIME++;
       flag[u] = 1;
       S.push( u );
       in[u] = 1;
       for( Edge *p = biao[u]; p != NULL; p = p->next ){
           int v = p \rightarrow to;
           if( !flag[v] ) {
               scc(v);
               low[u] = min( low[u], low[v] );
           else if( d[v] < d[u] && in[v] == 1)
               low[u] = min(low[u], d[v]);
       }
       if( low[u] == d[u] ){
           while(!S.empty() && d[S.top()] >= d[u]){
               label[ S.top() ] = GROUP;
               in[ S.top() ] = 0;
               S.pop();
           GROUP++;
       }
   }
};
7.2 双连通分量
          7.2.1 Tarjan
const int MAXV=101;
struct edge{
   int to;
   struct edge *next;
```

```
void add(int t,struct edge *&b){to=t;next=b;b=this;}
}*biao[ MAXV ],a[ 5000 ];//分别是临界表和节点仓库
//如果有一条边 from-to,那么添加边的语句为 a[ count++ ].add( to,biao[from] );
//用 tarjan 算法求桥,割点,双连通分量
int d[ MAXV ];//发现时间
int low[ MAXV ];//最早可达祖先的发现时间
int flag[ MAXV ];//0 白, 1 灰, 2 黑 int out[ MAXV ];// 当前点的不相交子树, i 点去掉所得到的连通图个数为
out[i]+1(ROOT 除外)
bool is[ MAXV ];//是否是割点
int N;//点的个数
int NN;//割点个数
int TIME;//时间戳
int ROOT;//每次 dfs 都是从 ROOT 开始的
void bcc( int u ,int father){
   d[u] = ++TIME;
   low[u] = d[u];
   flag[u] = 1;
   for(edge *p = biao[u]; p != NULL; p=p->next ){
      int v = p->to;
      if( flag[ v ] == 1 && father != v )
          low[u] = min(low[u], d[v]);
      else if( !flag[ v ] ){
          bcc(v,u);
          low[u] = min( low[u], low[v] );
          if( d[u] \leftarrow low[v] \&\& u != ROOT )
             is[u] = true;
          if( d[u] <= low[v] ) out[u]++;</pre>
          if (d[u] < low[v]);
          //条件成立,则此边是桥,否则, uv 在同一个双连通分量中
      }
   }
   if( u == ROOT \&\& out[u] >= 2 ) is[u] =true;
   flag[u] = 2;
   if( is[u] ) NN++;
void init(void){
   fill(flag,flag+MAXV,0);
   fill(out,out+MAXV,0);
   fill( is,is+MAXV,0);
   for(int i = 0;i <= MAXV;i++)biao[i]=NULL;</pre>
   TIME = 0;
   ROOT = 1;
}
```

# 8 数据结构

- 8.1 (待补充)并查集
- 8.2 (待补充) AC 自动机
- 8.3 二叉查找树

```
typedef struct BST{
   int key;
   BST *lchild,*rchild,*parent; // 存左右孩子以及父节点
BST *head, *p, node[100];
int cou;
void init()
{
   head = p = NULL;
   cou = 0;
   memset(node,'\0',sizeof(node));
void BuildBST(BST *&head,int x) // 创建 BST, 即 BST 的插入~
{
   if( head == NULL )
       node[cou].key = x;
       node[cou].parent = p;
       head = &node[cou++];
       return ;
   p = head;
   if( head->key > x )
       BuildBST(head->lchild,x);
   else
       BuildBST(head->rchild,x);
}
void InorderTraver(BST *head) // 中序遍历
{
   if( head == NULL )
       return ;
   InorderTraver( head->lchild );
   cout << head->key << ' '
   InorderTraver( head->rchild );
}
BST* Search( BST *&head,int x )//在 BST 中查找 key 为 x 的值
   if( head->key == x )
       p = head;
       return head;
   if( head->key > x )
       Search(head->lchild,x);
```

```
else
       Search(head->rchild,x);
}
BST* Minmum(BST* head) // 返回 BST 中最小值
   while( head->lchild != NULL )
       head = head->lchild;
   return head;
}
BST *Maxmum( BST* head ) // 返回 BST 中最大值
   while( head->rchild != NULL )
       head = head->rchild;
   return head;
}
BST* Successor(BST *head) // 返回节点 head 的后继节点
   if( head->rchild != NULL )
       return Minmum(head->rchild);
   BST* y = head->parent;
   while( y != NULL && head == y->rchild )
       head = y;
       y = y->parent;
   return y;
}
BST* Predecessor(BST *head) // 返回节点 head 的前驱节点
{
   if( head->lchild != NULL )
       return Maxmum(head->lchild);
   BST* y = head->parent;
   while( y != NULL && head == y->lchild )
       head = y;
       y = y->parent;
   return y;
}
void Delet(BST *z) //删除节点 z
   BST *x,*y;
   if( z->lchild == NULL || z->rchild == NULL )
       y = z;
   else
       y = Successor(z);
   if( y->lchild != NULL )
       x = y \rightarrow lchild;
       x = y \rightarrow rchild;
   if( x != NULL )
       x->parent = y->parent;
   if( y->parent == NULL )
       head = x;
   else
```

```
if( y == y->parent->lchild )
          y->parent->lchild = x;
      else
          y->parent->rchild = x;
   if(y!=z)
      z \rightarrow key = y \rightarrow key;
}
int main()
{
   int x;
   while( cin >> x \&\& x)
      BuildBST(head,x);
   InorderTraver(head); //中序遍历输出
   cout << endl;</pre>
   while( cin >> x )
      BST *z;
      Search(head,x);
      z = p;
      Delet(z);
      if( head == NULL )
          break;
      InorderTraver(head); //删除一个 中序遍历一次
return 0;
8.4 树状数组
***************
树状数组
c ----> 树状数组
一维树状数组
**************
int c[MAX][MAX];
int Lowbit(int x)
   return x & (-x);
}
void Updata(int x,int num)// num 可能都为 1, 具体问题具体分析
   int i;
   for(i=x; i<MAX; i+=Lowbit(i))</pre>
      c[i] += num; // 若 num 为 1, c[i][k]++
int Getsum(int x)//sum 的返回值的类型可能是 long long ,根据题意判断
{
   int sum = 0,i;
   for(i=x; i>0; i-=Lowbit(i))
      sum += c[i];
   return sum;
}
/*
```

```
树状数组
c ----> 树状数组
 二维树状数组
int c[MAX][MAX];
int Lowbit(int x)
{
   return x & (-x);
}
void Updata(int x,int y, int num) // num 可能都为 1, 具体问题具体分析
   int i,k;
   for(i=x; i<MAX; i+=Lowbit(i))</pre>
       for(k=y; k<MAX; k+=Lowbit(k))</pre>
          c[i][k] += num; // 若 num 为 1, c[i][k]++
}
int Getsum(int x,int y)//sum 的返回值的类型可能是 long long , 根据题意判断
   int i,k,sum = 0;
   for(i=x; i>0; i-=Lowbit(i))
       for(k=y; k>0; k-=Lowbit(k))
          sum += c[i][k];
   return sum;
}
```

## 8.5 线段树

#### 8.5.1 注意事项

## 8.5.2 线段树结点定义

```
****************
   当前区间左端点的位置
1
   当前区间右端点的位置
r
   求得当前区间的长度,即r-1
len()
   (1 + r)/2
mid()
   判断一个区间是否覆盖当前区间
in()
   给当前区间左右端点位置赋值
1r()
// 一维线段树
struct Tnode{
 int 1,r;
```

```
int len() { return r - 1;}
int mid() { return MID(1,r);}
   bool in(int ll,int rr) { return l >= ll \&\& r <= rr; } void lr(int ll,int rr){ l = ll; r = rr; }
};
******
内容同上
Tnode son 为子树的结点
**************************
*******/
                      // 二维线段树(树套树)
struct T2node
{
   int l,r;
   Tnode son[MAXM<<2];</pre>
   int len() { return r - 1;}
   int mid() { return MID(1,r);}
   bool in(int ll,int rr) { return l >= ll && r <= rr; }
void lr(int ll,int rr){ l = ll; r = rr;}</pre>
};
         8.5.3 扫描线定义
                       // 扫描线定义
struct Sline
   double x,y1,y2;
   int flag;
};
         8.5.4 线段树一般模板
                   ****************
                一维线段树,例子是求区间乘积,Updata 修改的是点的值
                根一律为 1 , MAX 为区间总长度
******************************
                       // 一维线段树
struct Tnode{
   int l,r,val;
   long long sum;
   int len() { return r - 1;}
   int mid() { return MID(1,r);}
   bool in(int ll,int rr) { return l >= ll && r <= rr; }</pre>
   void lr(int ll,int rr){ l = ll; r = rr;}
};
Tnode node[MAX<<2];</pre>
int a[MAX];
void Updata_sum(int t)
   node[t].sum = (node[L(t)].sum % MOD * node[R(t)].sum % MOD) % MOD;
void Build(int t,int l,int r)
{
   node[t].lr(1,r);
   if(node[t].len() == 1)
      node[t].sum = node[t].val = a[1];
      return ;
   int mid = MID(1,r);
```

```
Build(L(t),1,mid);
   Build(R(t),mid,r);
   Updata_sum(t);
}
void Updata(int t,int l,int r,int xx)
   if( node[t].in(l,r) )
   {
      node[t].sum = node[t].val = val;
      return ;
   if( node[t].len() == 1 ) return ;
   int mid = node[t].mid();
   if( l < mid ) Updata(L(t),l,r,val);</pre>
   if( r > mid ) Updata(R(t),1,r,val);
   Updata_sum(t);
}
int Query(int t,int l,int r)
   if( node[t].in(l,r) ) return node[t].sum;
   if( node[t].len() == 1 ) return 0;
   int mid = node[t].mid();
   long long ans = 111;
   if( l < mid ) ans *= Query(L(t),l,r);
   ans %= MOD;
   if( r > mid ) ans *= Query(R(t),1,r);
   ans %= MOD;
   return ans;
}
二维线段树,例子是求区间最大值
const int MAX = 100;
                              //第二维
const int MAXM = 1010;
struct Tnode{
   int l,r,val;
   int len() { return r - 1;}
   int mid() { return MID(1,r);}
   bool in(int ll,int rr) { return l >= ll \&\& r <= rr; } void lr(int ll,int rr) { l = ll; r = rr; }
};
struct T2node
   int l,r;
   Tnode son[MAXM<<2];</pre>
   int len() { return r - 1;}
   int mid() { return MID(1,r);}
   bool in(int ll,int rr) { return l >= ll && r <= rr; }
   void lr(int 11,int rr){ 1 = 11; r = rr;}
T2node node[MAX<<2];
void init()
   memset(node,0,sizeof(node));
void sub_build(int f,int t,int l,int r)
```

```
{
   node[f].son[t].lr(l,r); node[f].son[t].val = -inf;
   if( l == r - 1 ) return;
   int mid = MID(1,r);
   sub_build(f,L(t),1,mid);
   sub_build(f,R(t),mid,r);
void Build(int t,int l1,int r1,int l2,int r2)
{
   node[t].lr(l1,r1);
                                     // 建立子树
   sub_build(t,1,12,r2);
   if( l1 == r1 - 1 ) return ;
   int mid = MID(l1,r1);
   Build(L(t), 11, mid, 12, r2);
   Build(R(t),mid,r1,l2,r2);
}
void sub_updata(int f,int t,int l,int r,int val)
   node[f].son[t].val = max(node[f].son[t].val,val);
   if( node[f].son[t].in(l,r) ) return ;
   if( node[f].son[t].len() == 1 ) return ;
   int mid = node[f].son[t].mid();
   if( l < mid ) sub_updata(f,L(t),l,r,val);</pre>
   if( r > mid ) sub_updata(f,R(t),l,r,val);
}
void Updata(int t,int l1,int r1,int l2,int r2,int val)
   sub updata(t,1,12,r2,val);
                                             // 更新子树
   if( node[t].in(l1,r1) ) return ;
   if( node[t].len() == 1 ) return ;
   int mid = node[t].mid();
   if( l1 < mid ) Updata(L(t),l1,r1,l2,r2,val);</pre>
   if( r1 > mid ) Updata(R(t),l1,r1,l2,r2,val);
}
int sub_query(Tnode *node,int t,int l,int r)
   if( node[t].in(l,r) ) return node[t].val;
   if( node[t].len() == 1 ) return -inf;
   int mid = node[t].mid();
   int ans = -inf;
   if( 1 < mid )
       ans = max(ans,sub_query(node,L(t),1,r));
   if(r > mid)
       ans = max(ans,sub_query(node,R(t),1,r));
   return ans;
}
int Query(int t,int l1,int r1,int l2,int r2)
{
   if( node[t].in(l1,r1) )
       return sub_query(node[t].son,1,12,r2);
   if( node[t].len() == 1 ) return -inf;
   int mid = node[t].mid();
   int ans = -inf;
   if( 11 < mid )
       ans = max(ans,Query(L(t),l1,r1,l2,r2));
```

```
if( r1 > mid )
       ans = max(ans,Query(R(t),l1,r1,l2,r2));
   return ans;
}
```

#### 8.5.5 矩形面积交

```
线段树求矩形交的面积(即每块区域被覆盖多于两次的面积)(hdu1255)
**********************************
const int MAX = 2010;
struct Tnode{int 1,r,cover;double once,len;};
Tnode node[MAX<<2];</pre>
struct Sline{double x,y1,y2;int flag;};
Sline 1[MAX];
double y[MAX];
void add line(double x1,double y1,double x2,double y2,int &cnt)
{
   1[cnt].x = x1; 1[cnt].y1 = y1; 1[cnt].y2 = y2;
   1[cnt].flag = 1;
   y[cnt++] = y1;
   1[cnt].x = x2; 1[cnt].y1 = y1; 1[cnt].y2 = y2;
   l[cnt].flag = -1;
   y[cnt++] = y2;
bool cmp(Sline a,Sline b)
   if(a.x == b.x)
      return a.flag > b.flag;
   return a.x < b.x;
}
void Build(int t,int l,int r)
{
   node[t].l = 1; node[t].r = r;
   node[t].cover = 0; node[t].once = node[t].len = 0.0;
   if( node[t].l == node[t].r - 1 ) return ;
   int mid = MID(1,r);
   Build(L(t),1,mid);
   Build(R(t),mid,r);
void Updata_len(int t)
   if( node[t].cover >= 2 )
      node[t].once = 0;
      node[t].len = y[node[t].r] - y[node[t].1];
      return ;
   if( node[t].cover == 1 )
      if( node[t].l == node[t].r - 1 )
          node[t].len = 0;
      else
          node[t].len = node[R(t)].len + node[L(t)].len + node[R(t)].once
+ node[L(t)].once;
      node[t].once = y[node[t].r] - y[node[t].l] - node[t].len;
      return ;
   if( node[t].cover == 0 )
```

```
if( node[t].1 == node[t].r - 1 )
           node[t].len = node[t].once = 0;
       else
       {
           node[t].len = node[R(t)].len + node[L(t)].len;
           node[t].once = node[R(t)].once + node[L(t)].once;
       return ;
   }
void Updata(int t, Sline p)
   if( y[node[t].1] >= p.y1 && y[node[t].r] <= p.y2 )
    {
       node[t].cover += p.flag;
       Updata_len(t);
       return ;
   if( node[t].l == node[t].r - 1 ) return ;
   int mid = MID(node[t].1,node[t].r);
   if( p.y1 < y[mid] ) Updata(L(t),p);
if( p.y2 > y[mid] ) Updata(R(t),p);
   Updata_len(t);
void solve(int n,int cnt)
   Build(1,0,cnt-1);
   double ans = 0.0;
   Updata(1,1[0]);
   for(int i=1; i<n; i++)
       ans += node[1].len*(l[i].x - l[i-1].x);
       Updata(1,1[i]);
   printf("%.21f\n",ans);
int main()
   cnt = 0;
   add_line(x1,y1,x2,y2,cnt);
   sort(y,y+cnt);
   sort(1,1+cnt,cmp);
   int t = unique(y,y+cnt) - y;
   solve(cnt,t);
return 0;
}
          8.5.6 矩形面积并
void Build(int t,int l,int r)
{
   node[t].l = 1; node[t].r = r;
   node[t].cover = 0;
    if( l == r - 1 ) return;
    int mid = MID(1,r);
   Build(R(t),mid,r);
   Build(L(t),1,mid);
}
```

```
void Updata_len(int t)
   if( node[t].cover > 0 )
      node[t].len = y[node[t].r] - y[node[t].1];
   else
      if( node[t].l == node[t].r - 1 )
         node[t].len = 0;
      else
         node[t].len = node[R(t)].len + node[L(t)].len;
}
void Updata(int t,Sline p)
   if( y[node[t].1] >= p.y1 && y[node[t].r] <= p.y2 )
      node[t].cover += p.flag;
      Updata_len(t);
      return ;
   if( node[t].l == node[t].r - 1 ) return ;
   int mid = MID(node[t].1,node[t].r);
   if( p.y1 < y[mid] ) Updata(L(t),p);
if( p.y2 > y[mid] ) Updata(R(t),p);
   Updata_len(t);
long long solve(int n,int cnt)
   Build(1,0,cnt-1);
   Updata(1,1[0]);
   long long ans = 011;
   for(int i=1; i<n; i++)
      ans += (1[i].x - 1[i-1].x)*node[1].len;
      Updata(1,1[i]);
   return ans;
}
        8.5.7 线段树求矩形覆盖 K 次交面积(可求面积并)
线段树求矩形交面积 (覆盖大于等于两次)
             矩形并面积 (覆盖大于等于一次)
             矩形覆盖 K 次交面积 (覆盖大于等于 N 次)
      注意:根据题意设定 slen 类型
      slen[i] 当前区间覆盖大于等于 i 次的长度
            当前区间覆盖次数, cover == 0 不代表子区间没有被覆盖
      cover
const int MAX = 60010;
                         // 矩形个数*2
const int K = 2;
typedef long long LL;
struct Sline{ int x,y1,y2,flag;};
                      // 一维线段树
struct Tnode{
                                //大小至少为 K+1,需要用到 slen[K]
   int l, r, cover, slen[K+1];
   int len() { return r - 1;}
   int mid() { return MID(1,r);}
   bool in(int ll,int rr) { return l >= ll \&\& r <= rr; } void lr(int ll,int rr) { <math>l = ll; r = rr; }
```

```
};
Tnode
       node[MAX<<2];
       1[MAX];
Sline
int
       y[MAX], cnty, cnt;
void add_line(int x1,int y1,int x2,int y2,int &cnt)
{
   l[cnt].x = x1; l[cnt].y1 = y1; l[cnt].y2 = y2;
   1[cnt].flag = 1;
   y[cnt++] = y1;
   1[cnt].x = x2; 1[cnt].y1 = y1; 1[cnt].y2 = y2;
   1[cnt].flag = -1;
   y[cnt++] = y2;
}
void Build(int t,int l,int r)
   node[t].lr(1,r);
   node[t].cover = 0;
   FOR(i, 0, K+1)
       node[t].slen[i] = 0;
   if( node[t].len() == 1 )
       return ;
   int mid = MID(1,r);
   Build(L(t),1,mid);
   Build(R(t),mid,r);
}
void Updata_len(int t)
   int cc = node[t].cover;
   FOR(i, 0, K+1)
   {
       if(cc >= i)
           node[t].slen[i] = y[node[t].r] - y[node[t].1];
       else
           if(node[t].len() == 1)
               node[t].slen[i] = 0;
           else
               node[t].slen[i]
                                             node[L(t)].slen[i-cc]
node[R(t)].slen[i-cc];
   }
void Updata(int t, Sline p)
   if( p.y1 \leftarrow y[node[t].1] && p.y2 \rightarrow y[node[t].r] )
       node[t].cover += p.flag;
       Updata_len(t);
       return ;
   if( node[t].len() == 1 ) return ;
   int mid = node[t].mid();
   if( p.y1 < y[mid] )</pre>
                              Updata(L(t), p);
   if(p.y2 > y[mid])
                          Updata(R(t), p);
   Updata_len(t);
}
```

```
bool cmp(Sline a, Sline b)
   if(a.x == b.x)
      return a.flag > b.flag;
   return a.x < b.x;
}
LL solve(int n)
   LL ans = 0;
   sort(y, y+n);
   cnty = unique(y, y+n) - y;
   sort(1, 1+n, cmp);
   Build(1, 0, cnty-1);
   Updata(1, 1[0]);
   FOR(i, 1, n)
      ans += node[1].slen[K] * 1ll * (l[i].x - l[i-1].x);
      Updata(1, 1[i]);
   return ans;
int main()
{
   int ind = 1, ncases, n, x1, y1, x2, y2;
   scanf("%d", &ncases);
   while( ncases-- )
      cnt = 0;
      scanf("%d", &n);
      FOR(i, 0, n)
         scanf("%d%d%d%d", &x1, &y1, &x2, &y2);
         x2++; y2++;
         add_line(x1, y1, x2, y2, cnt);
      }
      if(K > n)
         printf("Case %d: 0\n", ind++);
         continue;
      }
      LL ans = solve(cnt);
      printf("Case %d: %lld\n", ind++, ans);
   }
return 0;
        8.5.8 周长轮廓并
线段树矩形周长并 poj 1177 Picture
```

```
const int MAX = 5010;
struct Tnode{ int l,r,num,len,cover;bool lb,rb;};
struct Sline{ int x,y1,y2,flag;};
Tnode node[MAX<<2];</pre>
Sline l[MAX<<1];</pre>
int y[MAX<<1];</pre>
void add_line(int x1,int y1,int x2,int y2,int &cnt)
    1[cnt].x = x1; 1[cnt].y1 = y1; 1[cnt].y2 = y2; 1[cnt].flag = 1;
    y[cnt++] = y1;
    1[cnt].x = x2; 1[cnt].y1 = y1; 1[cnt].y2 = y2; 1[cnt].flag = -1;
   y[cnt++] = y2;
void init()
   memset(node,0,sizeof(node));
void Build(int t,int l,int r)
    node[t].l = 1; node[t].r = r;
    node[t].num = 0;
    if(l == r - 1) return;
    int mid = MID(1,r);
    Build(R(t),mid,r);
    Build(L(t),1,mid);
}
void Updata_len(int t)
{
    if( node[t].cover > 0 )
    {
       node[t].num = node[t].lb = node[t].rb = 1;
       node[t].len = y[node[t].r] - y[node[t].1];
    }
    else
        if( node[t].l == node[t].r - 1 )
           node[t].rb = node[t].lb = node[t].len = node[t].num = 0;
       else
        {
           node[t].rb = node[R(t)].rb; node[t].lb = node[L(t)].lb;
           node[t].len = node[L(t)].len + node[R(t)].len;
           node[t].num
                                node[L(t)].num
                                                        node[R(t)].num
node[R(t)].lb*node[L(t)].rb;
               //两线段重合的话,得减一下。。
void Updata(int t,Sline p)
    if( y[node[t].1] >= p.y1 && y[node[t].r] <= p.y2 )
    {
       node[t].cover += p.flag;
       Updata len(t);
       return ;
    if( node[t].l == node[t].r - 1 ) return ;
    int mid = MID(node[t].l ,node[t].r);
    if( p.y1 < y[mid] ) Updata(L(t),p);</pre>
    if( p.y2 > y[mid] ) Updata(R(t),p);
   Updata_len(t);
int solve(int n,int cnt,Sline *1)
```

```
{
   init();
   Build(1,0,cnt-1);
   int ans = 0,last = 0,lines = 0;
   for(int i=0; i<n; i++)
       Updata(1,1[i]);
       if( i >= 1 )
           ans += 2 * lines * (l[i].x - l[i-1].x);//计算平行于 x 轴的长度
       ans += abs(node[1].len - last);
                                             // 计算平行于 y 轴的长度
       last = node[1].len;
       lines = node[1].num;
   return ans;
}
bool cmp(Sline a,Sline b)
   if( a.x == b.x ) return a.flag > b.flag;
   return a.x < b.x;
int main()
   int n,x1,x2,y1,y2;
   while( ~scanf("%d",&n) )
       if( n == 0 )
           printf("0\n");
           continue;
       int cnt = 0;
       for(int i=0; i<n; i++)
           scanf("%d%d%d%d",&x1,&y1,&x2,&y2);
           add_line(x1,y1,x2,y2,cnt);
       sort(y,y+cnt);
       sort(1,1+cnt,cmp);
       int t = cnt;
       t = unique(y,y+cnt) - y;
       int ans = solve(cnt,t,1);
       printf("%d\n",ans);
   }
return 0;
```

#### 8.5.9 求区间连续递增最长子序列

```
当前区间右端点的值
*******************************
int a[MAX];
                        // 一维线段树
struct Tnode{
   int l,r,lval, rval, max, add, ll, rr;
   int len() { return r - 1;}
   int mid() { return MID(l,r);}
   bool in(int 11,int rr) { return 1 >= 11 && r <= rr; }</pre>
   void lr(int ll,int rr){ l = ll; r = rr;}
Tnode node[MAX<<2];</pre>
void Updata val(int t) //左右区间连接的时候更新
   node[t].ll = node[L(t)].ll;
   node[t].rr = node[R(t)].rr;
   node[t].lval = node[L(t)].lval;
   if( node[L(t)].lval == node[L(t)].len()
          && node[R(t)].11 > node[L(t)].rr)
       node[t].lval += node[R(t)].lval;
   node[t].rval = node[R(t)].rval;
   if( node[R(t)].rval == node[R(t)].len()
          && node[R(t)].11 > node[L(t)].rr)
       node[t].rval += node[L(t)].rval;
   int mval = 0;
   if( node[R(t)].11 > node[L(t)].rr )
       mval = node[L(t)].rval + node[R(t)].lval;
   node[t].max = max(mval, max(node[L(t)].max, node[R(t)].max));
   node[t].max = max(node[t].max, max(node[t].lval, node[t].rval));
void Build(int t,int l,int r)
   node[t].add = node[t].lval = node[t].rval = node[t].max = 0;
   node[t].lr(1,r);
   if( node[t].len() == 1 )
       node[t].lval = node[t].rval = node[t].max = 1;
       node[t].11 = node[t].rr = a[node[t].1];
       return ;
   int mid = MID(1,r);
   Build(L(t),1,mid);
   Build(R(t),mid,r);
   Updata_val(t);
}
void Updata llrr(int t, int add) // 更新区间左右端点的值
   node[t].ll += add;
   node[t].rr += add;
}
void Push down(int t) // 向下推进 add 值
   if(node[t].len() == 1)
       node[t].add = 0;
       return ;
```

```
if( node[t].add )
       node[L(t)].add += node[t].add;
       Updata_llrr( L(t), node[t].add );
       node[R(t)].add += node[t].add;
       Updata_llrr( R(t), node[t].add);
       node[t].add = 0;
   }
}
void Updata(int t,int l,int r,int val)//更新的是 add 值
   if( node[t].in(l,r) )
       node[t].add += val;
       Updata_llrr(t, val);
       Push_down(t);
       return ;
   if( node[t].len() == 1 ) return ;
   Push_down(t);
   int mid = node[t].mid();
   if( 1 < mid ) Updata(L(t),1,r,val);</pre>
   if( r > mid ) Updata(R(t),1,r,val);
   Updata_val(t);
}
int Query(int t,int l,int r)
   if( node[t].in(l,r) )
       return node[t].max;
   if( node[t].len() == 1 ) return 0;
   Push_down(t);
   int mid = node[t].mid();
   int ans = 0;
   if( r <= mid )
       ans = \max(ans, Query(L(t),1,r));
   else
       if( 1 >= mid )
           ans = \max(ans, Query(R(t),1,r));
       { //因为当前区间可能不完全包括连续值,所以需要以下操作
           ans = \max(ans, Query(L(t),l,r));
           ans = \max(ans, Query(R(t),1,r));
           int 11, rr;
           if( node[L(t)].r - node[L(t)].rval > 1 )
              11 = node[L(t)].rval;
           else
              ll = node[L(t)].r - l;
           // 找当前区间右边最大连续值
           if( node[R(t)].1 + node[R(t)].lval > r )
              rr = r - node[R(t)].1;
           else
              rr = node[R(t)].lval;
           // 找当前区间左边最大连续值
           ans = max(ans, 11);
```

```
ans = max(ans, rr);
          if( node[R(t)].11 > node[L(t)].rr )
              int mval = rr + 11;
             ans = max(ans, mval);
          }
   Updata_val(t);
   return ans;
}
         8.5.10线段树求区间连续为 N 空白的左端点
线段树求区间连续为 N 空白的左端点
Codeforces Beta Round #43 D. Parking Lot
// 一维线段树
struct Tnode{
   int l,r, lval, rval, mval, max, cover;
   int len() { return r - 1;}
int mid() { return MID(l,r);}
bool in(int ll,int rr) { return l >= ll && r <= rr; }
void lr(int ll,int rr){ l = ll; r = rr;}</pre>
Tnode node[MAX<<2];</pre>
void Build(int t,int l,int r)
   node[t].lr(1,r);
   node[t].lval = node[t].rval = node[t].mval = node[t].max = r - 1;
   if( node[t].len() == 1 ) return ;
   int mid = MID(1,r);
   Build(L(t),1,mid);
   Build(R(t),mid,r);
}
void Updata_val(int t)
   node[t].lval = node[L(t)].lval;
   if( node[L(t)].lval == node[L(t)].len() )
       node[t].lval += node[R(t)].lval;
   node[t].rval = node[R(t)].rval;
   if( node[R(t)].rval == node[R(t)].len() )
       node[t].rval += node[L(t)].rval;
   node[t].mval = node[L(t)].rval + node[R(t)].lval;
   node[t].max = max(node[L(t)].max, node[R(t)].max);
   node[t].max = max(node[t].max, max(node[t].lval, max(node[t].rval,
node[t].mval)));
```

node[t].lval = node[t].rval = node[t].mval = node[t].max = (val ?

void Updata(int t,int l,int r,int val)

if( node[t].len() == 1 ) return ;

if( node[t].in(l,r) )

int mid = node[t].mid();

0 : node[t].len());
 return ;

```
if( l < mid ) Updata(L(t),l,r,val);
if( r > mid ) Updata(R(t),l,r,val);
   Updata_val(t);
}
void Query(int t, int len, int &pos)
    if( node[t].max == node[t].len() && node[t].max >= len )
    {
        pos = node[t].1;
        return ;
    if( node[t].max < len ) return ;</pre>
    if( node[t].len() == 1 ) return ;
    if( node[L(t)].max >= len )
        Query(L(t), len, pos);
    else
        if( node[t].mval >= len )
            pos = node[L(t)].r - node[L(t)].rval;
            return ;
        else
            if( node[R(t)].max >= len )
                Query(R(t), len, pos);
struct NODE
{
    int len, pos;
NODE a[110];
int main()
{
    int 1, b, f, n, cnt, pos, op, len;
   while( ~scanf("%d%d%d", &l, &b, &f) )
        scanf("%d", &n);
        Build(1, -b, l + f);
        FOR(i, 1, n+1)
            scanf("%d%d", &op, &len);
            if( op == 1 )
            {
                a[i].len = len;
                pos = -1000;
                Query(1, len+b+f, pos);
                if(pos == -1000)
                   puts("-1");
a[i].pos = -1;
                }
                else
                    printf("%d\n", pos+b);
                    a[i].pos = pos + b;
                if( pos != -1000 )
                    Updata(1, pos + b, pos + b + len, 1);
            }
            else
            {
```

```
pos = a[len].pos;
              Updata(1, pos, pos + a[len].len, 0);
          }
       }
   }
return 0;
         8.5.11 线段树区间加法乘法
线段树区间加法乘法
BZOJ 1798: [Ahoi2009]Seq 维护序列 seq
Change operations:
1 a b c 区间[a , b]的每个元素都乘以 c
2 a b c 区间[a , b]的每个元素都加上 c
Output operations:
3 a b 输出区间[a , b]的和 mod d
                             ***************
const int MAX = 100010;
                           // 一维线段树
struct Tnode{
   int l,r;
   long long sum, mul, add;
   bool ind;
   int len() { return r - 1;}
int mid() { return MID(1,r);}
bool in(int ll,int rr) { return l >= ll && r <= rr; }
void lr(int ll,int rr){ l = ll; r = rr;}</pre>
Tnode node[MAX<<2];</pre>
int a[MAX],d;
void init()
{
   memset(node,0,sizeof(node));
void Build(int t,int l,int r)
   node[t].lr(l,r);
   node[t].mul = 1;
   node[t].add = 0;
   if( node[t].len() == 1 )
       node[t].sum = a[node[t].1];
       return ;
   int mid = MID(1, r);
   Build(L(t), 1, mid);
   Build(R(t), mid, r);
   node[t].sum = (node[L(t)].sum + node[R(t)].sum) % d;
void Mul(long long &sum,long long mul)
   sum *= mul; sum %= d;
void Add(long long &sum,long long add)
   sum += add; sum %= d;
void Updata_sub(int t,long long mul,long long add)
{
```

```
Mul(node[t].sum, mul);
   Add(node[t].sum, add*node[t].len());
   Mul(node[t].add, mul);
   Mul(node[t].mul, mul);
   Add(node[t].add, add);
}
void Pushdown(int t)
   if( node[t].mul == 1 && node[t].add == 0 ) return ;
   Updata_sub(L(t), node[t].mul, node[t].add);
   Updata_sub(R(t), node[t].mul, node[t].add);
   node[t].mul = 1; node[t].add = 0;
void Updata(int t,int l,int r,int op,long long val)
   if( op == 0 && val == 1 || op == 1 && val == 0 ) return ;
   if( node[t].in(l,r) )
       if( op == 0 )
           Mul(node[t].sum, val);
           Mul(node[t].mul, val);
Mul(node[t].add, val);
       else
           Add(node[t].sum, node[t].len()*val);
           Add(node[t].add, val);
       return ;
   if( node[t].len() == 1 ) return ;
   Pushdown(t);
   int mid = node[t].mid();
   if( l < mid ) Updata(L(t),l,r,op,val);</pre>
   if( r > mid ) Updata(R(t),1,r,op,val);
   node[t].sum = (node[L(t)].sum + node[R(t)].sum) % d;
}
long long Query(int t,int l,int r)
{
   if( node[t].in(l,r) ) return node[t].sum;
   if( node[t].len() == 1 ) return 0;
   Pushdown(t);
   int mid = node[t].mid();
   long long ans = 011;
   if( l < mid ) ans += Query(L(t),l,r);
   ans %= d;
   if(r > mid) ans += Query(R(t),1,r);
   ans %= d;
   return ans;
}
int main()
   int n,m,x,y,z,ind;
   while( ~scanf("%d%d",&n,&d) )
       init();
       for(int i=0; i<n; i++)
           scanf("%d",&a[i]);
```

```
Build(1,0,n);
        scanf("%d",&m);
while( m-- )
             scanf("%d",&ind);
             if( ind == 1 )
             {
                 scanf("%d%d%d",&x,&y,&z);
                 Updata(1,x-1,y,0,z);
             if( ind == 2 )
                 scanf("%d%d%d",&x,&y,&z);
                 Updata(1,x-1,y,1,z);
             if( ind == 3 )
                 scanf("%d%d",&x,&y);
                 long long ans = Query(1,x-1,y);
printf("%1ld\n",ans);
        }
    }
return 0;
```

#### 8.5.12 线段树区间异或,覆盖,最长序列

```
*******
线段树区间异或,覆盖,最长序列。hdu 3397 Sequence operation
Change operations:
0 a b change all characters into '0's in [a , b]
1 a b change all characters into '1's in [a , b]
2 a b change all '0's into '1's and change all '1's into '0's in [a, b]
Output operations:
3 a b output the number of '1's in [a, b]
4 a b output the length of the longest continuous '1' string in [a , b]
const int MAX = 100010;
struct Tnode{
   int l,r,lb,rb,lw,rw,sb,sw,sum;bool t;short cover;
   int mid() { return MID(1,r);}
   int len() { return r - 1;}
   bool in(int ll,int rr) { return l >= 11 \&\& r <= rr; } void lr(int ll,int rr) { l = 1l; r = rr; }
};
Tnode node[MAX*3];
bool aa[MAX];
void init()
   memset(aa,false,sizeof(aa));
   memset(node,0,sizeof(node));
void Updata_len(int t)
   node[t].sum = node[t].len() - node[t].sum;
void Swap(int t)
   swap(node[t].rw,node[t].rb);
   swap(node[t].lw,node[t].lb);
   swap(node[t].sw,node[t].sb);
```

```
}
void Updata_wb(int t,int v1,int v2)
   node[t].rb = node[t].lb = node[t].sb = v1;
   node[t].rw = node[t].lw = node[t].sw = v2;
   node[t].sum = v1;
void Updata cover(int t)
   if( node[t].cover == 1 )
       Updata_wb(L(t), 0, node[L(t)].len());
       Updata_wb(R(t), 0, node[R(t)].len());
   }
   else
   {
       Updata_wb(L(t), node[L(t)].len(), 0);
       Updata_wb(R(t), node[R(t)].len(), 0);
void Pushdown_len(int t)
   if( node[t].cover > 0 && node[t].len() != 1 )
       node[R(t)].cover = node[L(t)].cover = node[t].cover;
       Updata_cover(t);
       node[L(t)].t = node[R(t)].t = 0;
   node[t].cover = 0;
void Updata_sum(int t)
   node[t].lw = node[L(t)].lw + (node[L(t)].lw == node[L(t)].len() ?
node[R(t)].lw : 0 );
   node[t].rw = node[R(t)].rw + (node[R(t)].rw == node[R(t)].len() ?
node[L(t)].rw : 0);
   node[t].sw = max(node[R(t)].sw, max(node[L(t)].sw, node[L(t)].rw +
node[R(t)].lw));
   node[t].lb = node[L(t)].lb + (node[L(t)].lb == node[L(t)].len() ?
node[R(t)].lb : 0 );
   node[t].rb = node[R(t)].rb + (node[R(t)].rb == node[R(t)].len() ?
node[L(t)].rb : 0 );
   node[t].sb = max(node[R(t)].sb, max(node[L(t)].sb, node[L(t)].rb +
node[R(t)].1b));
   node[t].sum = node[L(t)].sum + node[R(t)].sum;
void Pushdown_xor(int t)
   if( node[t].len() == 1 )
       node[t].t = 0;
       return ;
   if( node[t].t )
       node[R(t)].t = !node[R(t)].t;
       node[L(t)].t = !node[L(t)].t;
       Swap(R(t)); Swap(L(t));
       Updata_len(R(t)); Updata_len(L(t));
       node[t].t = !node[t].t;
   }
void Build(int t,int l,int r)
```

```
{
   node[t].lr(l,r);
   if( node[t].len() == 1 )
       node[t].sum = node[t].lb = node[t].rb = node[t].sb = aa[1];
       node[t].lw = node[t].rw = node[t].sw = 1 - aa[1];
       return ;
   int mid = MID(1,r);
   Build(L(t),1,mid);
   Build(R(t),mid,r);
   Updata_sum(t);
void Updata(int t,int l,int r,int val)
   Pushdown_len(t);
   Pushdown_xor(t);
   if( node[t].in(l,r) )
       if( val != 3 )
           node[t].cover = val;
           if( node[t].cover == 1 )
               Updata_wb(t, 0, node[t].len());
           else
               Updata_wb(t, node[t].len(), 0);
           return ;
       else
           node[t].t = !node[t].t;
           Swap(t);
           Updata_len(t);
           return ;
       }
   if( node[t].len() == 1 ) return ;
   int mid = MID(node[t].1,node[t].r);
   if( l < mid ) Updata(L(t), l, r, val);</pre>
   if( r > mid ) Updata(R(t), l, r, val);
   Updata_sum(t);
}
int Query_nor(int t,int l,int r)
{
   Pushdown_len(t);
   Pushdown_xor(t);
   if( node[t].in(l,r) ) return node[t].sum;
   if( node[t].len() == 1 ) return 0;
   int mid = node[t].mid();
   int ans = 0;
   if( 1 < mid )
                  ans += Query_nor(L(t),1,r);
   if( r > mid )
                  ans += Query_nor(R(t),1,r);
   Updata_sum(t);
   return ans;
}
int Query_xor(int t,int l,int r)
   Pushdown len(t);
   Pushdown_xor(t);
   if( node[t].in(l,r) ) return node[t].sb;
   if( node[t].len() == 1 ) return 0;
```

```
int mid = node[t].mid();
   int ans = 0;
   if( 1 >= mid )
       ans = max(ans,Query_xor(R(t),1,r));
   else
       if( r <= mid )
          ans = max(ans,Query_xor(L(t),1,r));
       else
       {
           ans = max(ans, Query_xor(L(t), 1, mid));
           ans = max(ans, Query_xor(R(t), mid, r));
int a = ( node[L(t)].rb <= mid - 1 ? node[L(t)].rb : mid - 1);</pre>
           int b = (node[R(t)].lb \leftarrow r - mid ? node[R(t)].lb : r - mid);
           ans = max(ans, a+b);
   Updata_sum(t);
   return ans;
}
int main()
   int n,m,ncases,a,b,ind;
   scanf("%d",&ncases);
   while( ncases-- )
   {
       init();
scanf("%d%d",&n,&m);
       for(int i=0; i<n; i++)
           scanf("%d",&a);
                                                          //数列的初始值
           if(a) aa[i] = 1;
       Build(1,0,n);
       while( m-- )
           scanf("%d%d%d",&ind,&a,&b);
           if( ind == 0 ) Updata(1,a,b+1,1);
           if( ind == 1 ) Updata(1,a,b+1,2);
if( ind == 2 ) Updata(1,a,b+1,3);
           if( ind == 3 )
           {
              int ans = Query_nor(1,a,b+1);
              printf("%d\n",ans);
           if( ind == 4 )
              int ans = Query_xor(1,a,b+1);
              printf("%d\n",ans);
           }
       }
return 0;
8.6 划分树
                    *****************
                         划分树求区间第K数
```

```
const int MAX = 100010;
class Parti tree{
public :
                                       // 我的一维线段树定义
   class Tnode{
   public :
       int l,r;
       int len() { return r - 1;}
       int mid() { return MID(1,r);}
       bool in(int ll,int rr) { return l >= ll && r <= rr; }</pre>
       void lr(int 11,int rr){ 1 = 11; r = rr;}
   };
   Tnode node[MAX<<2];</pre>
   int num_left[20][MAX], seg[20][MAX],sa[MAX];//sa 数组存 sort 后的结果
                     //seg 数组存的是 d 层划分后的数字 (类似快排 Partation
 (d-1) 次后的结果)
                     //num left 存的是 d 层在 i 之前(包括 i)小于 sa[mid] 的数
的数目
   void init()
       memset(seg,0,sizeof(seg));
       memset(num_left,0,sizeof(num_left));
       memset(node,0,sizeof(node));
   void build(int s,int t)
   {
       sort(sa+s,sa+t+s);
       Parti_build(1,s,t,1);
   int query(int s,int t,int k)
   {
       return find rank(1,s,t,1,k);
   }
   void Parti build(int t,int l,int r,int d)
   {
       node[t].lr(l, r);
       if( node[t].len() == 0 ) return ;
       int mid = MID(l, r), lsame = mid - l + 1;
       for(int i=l; i<=r; i++)//首先确定分到每一侧的数的数目
          if( seg[d][i] < sa[mid] )//因为相同的元素可能被分到两侧
              lsame--;
       int lpos = 1,rpos = mid+1;
       for(int i=1; i<=r; i++)
       {
          if( i == 1 )
              num_left[d][i] = 0;
              num_left[d][i] = num_left[d][i-1];
          if( seg[d][i] < sa[mid] )</pre>
          {
              num_left[d][i]++;
              seg[d+1][lpos++] = seg[d][i];
          if( seg[d][i] > sa[mid] )
              seg[d+1][rpos++] = seg[d][i];
          if( seg[d][i] == sa[mid] )
              if( lsame > 0 ) // 如果大于 0, 说明左侧可以分和 sa[mid]相同
的数字
              {
                 lsame--;
                 num_left[d][i]++;
```

```
seg[d+1][lpos++] = seg[d][i];
             }
             else
                           //反之,说明左侧数字已经分满了,就分到右边去
                seg[d+1][rpos++] = seg[d][i];
      Parti_build(L(t), 1, mid, d+1);
      Parti_build(R(t), mid+1, r, d+1);
   int find_rank(int t,int l,int r,int d,int val)
   {
      if( node[t].len() == 0 ) return seg[d][1];
                           //s表示区间[1,r]有多少个小于sa[mid]的数被分到
      int s,ss;
左边
      if(1 == node[t].1)
          ss = 0;
      else
          ss = num_left[d][1-1];
      s = num_left[d][r] - ss;//ss表示从当前区间的L到1-1有多少个小于sa[mid]
的数被分到左边
      if(s >= val)
          return find_rank(L(t), node[t].l+ss, node[t].l+ss+s-1, d+1,
val);
      else
      {
          int mid = node[t].mid();
          int bb = 1 - node[t].1 - ss; //表示从当前区间 L 到 1-1 有多少个分
到右边
          int b = r - l + 1 - s;
                                        //表示[1,r]有多少个分到右边
          return find rank(R(t), mid+bb+1, mid+bb+b,d+1,val-s);
      }
   }
};
Parti tree t;
int main()
   int n,m,x,y,k;
   while( ~scanf("%d%d",&n,&m) )
      t.init();
                   //可以不要
      for(int i=1; i<=n; i++)
          scanf("%d",&t.sa[i]);
          t.seg[1][i] = t.sa[i];
      t.build(1,n);
      while( m-- )
          scanf("%d%d%d",&x,&y,&k);
          int ans = t.query(x, y, k);
          printf("%d\n",ans);
   }
return 0;
                          **************
```

划分树求区间中位数

```
const int MAX = 100010;
LL s[MAX];
class Parti_tree{
public :
   class Tnode{
   public :
       int l,r;
       int len() { return r - 1;}
       int mid() { return MID(1,r);}
       bool in(int 11,int rr) { return 1 >= 11 && r <= rr; }
       void lr(int ll,int rr){ l = ll; r = rr;}
   };
   Tnode node[MAX<<2];</pre>
   int num_left[20][MAX], seg[20][MAX],sa[MAX];
   LL less_mid, less_midt, less_num, sum_left[20][MAX];
                        //less mid 是区间内小于中位数的数的和 ,less num 是
   void init()
小于中位数的数量
                     // sum left 存的是 d 层 i 之前小于中位数的和
   {
       memset(num_left,0,sizeof(num_left));
       memset(node,0,sizeof(node));
   void build(int s,int t)
   {
       sort(sa+s,sa+t+s);
       Parti_build(1,s,t,1);
   int query(int s,int t,int k)
       less mid = less midt = less num = 0;
       return find rank(1,s,t,1,k);
   }
   void Parti build(int t,int l,int r,int d)
   {
       node[t].lr(l, r);
       if( node[t].len() == 0 ) return ;
       int mid = MID(l, r), lsame = mid - l + 1;
       for(int i=1; i<=r; i++)</pre>
          if( seg[d][i] < sa[mid] )</pre>
              lsame--;
       int lpos = 1,rpos = mid+1;
       for(int i=1; i<=r; i++)
       {
          if( i == 1 )
              num_left[d][i] = sum_left[d][i] = 0;
          else
          {
              num_left[d][i] = num_left[d][i-1];
              sum left[d][i] = sum left[d][i-1];
          if( seg[d][i] < sa[mid] )</pre>
              num_left[d][i]++;
              seg[d+1][lpos++] = seg[d][i];
              sum_left[d][i] += seg[d][i];
          if( seg[d][i] > sa[mid] )
              seg[d+1][rpos++] = seg[d][i];
          if( seg[d][i] == sa[mid] )
              if(lsame > 0)
```

```
lsame--;
                  num_left[d][i]++;
                  seg[d+1][lpos++] = seg[d][i];
                  sum_left[d][i] += seg[d][i];
               }
               else
                  seg[d+1][rpos++] = seg[d][i];
       Parti_build(L(t), 1, mid, d+1);
       Parti build(R(t), mid+1, r, d+1);
   int find_rank(int t,int l,int r,int d,int val)
       if( node[t].len() == 0 ) return seg[d][1];
       int s,ss;
       if(1 == node[t].1)
           s = num_left[d][r];
           less_midt = sum_left[d][r];
           ss = 0;
       }
       else
           s = num_left[d][r] - num_left[d][l-1];
           ss = num_left[d][1-1];
           less_midt = sum_left[d][r] - sum_left[d][1-1];
       if( s >= val )
           return find_rank(L(t), node[t].l+ss, node[t].l+ss+s-1, d+1,
val);
       else
       {
           int mid = node[t].mid();
           int bb = 1 - node[t].1 - ss;
           int b = r - l + 1 - s;
           less_mid += less_midt;
           less_num += s;
           return find_rank(R(t), mid+bb+1, mid+bb+b,d+1,val-s);
       }
   }
};
Parti_tree t;
int main()
   int n,m,x,y,k;
   int ncases,ind = 1;
   scanf("%d",&ncases);
   while( ncases-- )
       scanf("%d",&n);
       t.init();
       s[0] = 0;
       for(int i=1; i<=n; i++)
           scanf("%d",&t.sa[i]);
           s[i] = s[i-1] + t.sa[i];
           t.seg[1][i] = t.sa[i];
```

```
}
      scanf("%d",&m);
      t.build(1,n);
      printf("Case #%d:\n",ind++);
      while( m-- )
          scanf("%d%d",&x,&y);
          if(x == y)
          {
             printf("0\n");
             continue;
          x++; y++;//这题是从 0 开始输入的。= =。。 改成从 1 的
          int mid = (y-x)/2+1;
          LL ans = t.query(x, y, mid);
          LL sum = s[y] - s[x-1];
          LL suml = t.less_mid;
          LL sumr = sum - suml - ans;
          LL out = t.less_num*ans - suml + sumr - (y - x - t.less_num)*ans;
          printf("%I64d\n",out);
      printf("\n");
return 0;
    归并树
8.7
                        归并树求区间第 K 数 (效率没有划分树高)
const int MAX = 100010;
class Merger_tree{
public :
   class Tnode{
   public :
       int l,r;
       int len() { return r - 1;}
      int mid() { return MID(l,r);}
      bool in(int ll,int rr) { return l >= ll && r <= rr; }</pre>
      void lr(int ll,int rr){ l = ll; r = rr;}
   };
   Tnode node[MAX<<2];</pre>
   int seg[20][MAX],a[MAX],n;
   void init()
   {
      memset(seg,0,sizeof(seg));
      memset(node,0,sizeof(node));
   void build(int s,int t){ n = t; Merger_build(1,s,t,1); }
   int query(int x,int y,int k) { return find_rank(n,x,y,k); };
   void Merger_build(int t,int l,int r,int deep)
   {
      node[t].lr(l, r);
      if( node[t].len() == 0 )
          seg[deep][1] = a[1];
          return ;
```

```
int mid = MID(1, r);
       Merger_build(L(t), 1, mid, deep+1);
       Merger_build(R(t), mid+1, r, deep+1);
       int k = l, i = l, j = mid+1;
       while( i <= mid && j <= r )
           if( seg[deep+1][i] < seg[deep+1][j] )</pre>
               seg[deep][k++] = seg[deep+1][i++];
               seg[deep][k++] = seg[deep+1][j++];
       while( i <= mid )</pre>
           seg[deep][k++] = seg[deep+1][i++];
       while (j <= r)
           seg[deep][k++] = seg[deep+1][j++];
   int find_k(int t,int l,int r,int deep,int val)
       if( node[t].in(l,r) )
           int ll = node[t].1, rr = node[t].r;
           while( ll < rr )
           {
               int mid = MID(ll, rr);
               if( seg[deep][mid] < val )</pre>
                   11 = mid + 1;
               else
                   rr = mid;
           if( seg[deep][ll] <= val )</pre>
               return ll - node[t].l + 1;
               return ll - node[t].1;
       if( node[t].len() == 0 ) return 0;
       int ans = 0;
       int mid = node[t].mid();
       if( l \le mid ) ans += find_k(L(t), l, r, deep+1, val);
       if( r \ge mid ) ans += find_k(R(t), l, r, deep+1, val);
       return ans;
   int find_rank(int n,int x,int y,int k)
       int l = 1, r = n;
       while(l < r)
           int mid = MID(1, r);
           if( find_k(1, x, y, 1, seg[1][mid]) < k )</pre>
               1 = mid + 1;
           else
               r = mid;
       return seg[1][1];
   }
};
Merger_tree t;
int main()
{
```

```
int n,m,x,y,k;
   while( ~scanf("%d%d",&n,&m) )
       t.init();
       for(int i=1; i<=n; i++)
           scanf("%d",&t.a[i]);
       t.build(1,n);
       while( m-- )
       {
           scanf("%d%d%d",&x,&y,&k);
           int ans = t.query(x,y,k);
           printf("%d\n",ans);
   }
return 0;
}
8.8 Treap
int sum;
template<class Tp>
class Treap_Node{
public:
   Tp value;
   int time, fix;
   Treap_Node *left,*right;
   Treap_Node() { puts("you forget something!"); }
   Treap_Node(Tp v,int f){
       left=right=NULL;
       value=v;fix=f;time=1;
   }
};
template<class Tp>
class Treap{
private:
   Treap_Node<Tp> *Root;
   void Left_Rotate( Treap_Node<Tp>* &H ){
       Treap_Node<Tp> *temp = H->right;
       H->right = temp->left;
       temp->left = H;
       H = temp;
   void Right_Rotate( Treap_Node<Tp>* &H ){
       Treap_Node<Tp> *temp = H->left;
       H->left = temp->right;
       temp->right = H;
       H = temp;
   void insert( Treap_Node<Tp>* &H, Tp v ){
       if( H == NULL )
           H = new Treap_Node<Tp>(v,rand());
       else if( v < H->value ){
           insert( H->left, v );
           if( H->left->fix < H->fix )
               Right_Rotate( H );
       }
```

```
else if( v > H->value ){
           insert( H->right, v );
           if( H->right->fix < H->fix )
               Left_Rotate( H );
       else H->time++;
   bool del( Treap_Node<Tp>* &H, Tp v ){
       if( H == NULL ) return false;
       if( v == H->value ){
           if( H->time > 1 ) {H->time--; return true;}
           if( H->left == NULL || H->right == NULL ){
               Treap Node<Tp> *t = H;
               H = (H->right==NULL)?H->left:H->right;
               delete t;return true;
           }
           else{
               if( H->left->fix < H->right->fix ){
                   Right_Rotate( H );
                  return del( H->right, v );
               }else{
                   Left_Rotate( H );
                  return del( H->left, v );
               }
           }
       else if( v < H->value )
           return del( H->left, v );
       else
           return del( H->right, v );
   Treap Node<Tp>* find( Treap Node<Tp>* H, Tp v ){
       if( H == NULL \mid \mid v == H->value )
           return H;
       if( v < H->value ) return find( H->left,v );
       else return find( H->right, v );
   void Destroy( Treap_Node<Tp>* H ){
       if( H == NULL ) return;
       Destroy( H->left );
       Destroy( H->right );
       delete H;
   void out( Treap_Node<Tp> *H){
       if( H == NULL ) return;
       out( H->left );
       printf("%s %.4lf\n",H->value.c_str(),H->time*100.0/sum);
       out( H->right );
   Treap Node<Tp> *L(Treap Node<Tp>* &H ){
       if( H == NULL ) return H;
       return H->left!=NULL?L(H->left):H;
   Treap Node<Tp> *R(Treap Node<Tp>* &H ){
       if( H == NULL ) return H;
       return H->right!=NULL?R(H->right):H;
public:
   Treap() { Root = NULL; }
```

```
~Treap() { Destroy( Root );}
  void Insert( Tp v ){ insert( Root, v ); }
void Delete( Tp v ){    del( Root, v ); }
   Treap_Node<Tp> *Find( Tp v ){ return find(Root,v);}
   Treap_Node<Tp> *Min( void ){return L(Root);}
   Treap_Node<Tp> *Max( void ){return R(Root);}
  void print(void) { out( Root );}
int main(void){
  sum = 0;
   char a[100];
   Treap<string> T;
  while(gets(a) && a[0] != '\0'){
     string aa(a);
     T.Insert( aa );
     T.Delete( aa );
     T.Insert( aa );
     sum++;
   T.print();
  return 0;
8.9 矩形切割
矩形切割 v2.0
//
//
                                      //
               矩形类,包括矩形坐标和颜色;
// 1. class Rect:
                                       //
// 2. clear:
               清空所有矩形;
                                       //
               插入一个矩形;
// 3. insert:
                                       //
               获取各个颜色的面积统计;
// 4. get_color:
                                       //
// 5. total_area:
               获取面积总和;
                                       //
               get_color 简洁版;
// 6. color:
                                        //
//
2008 CopyRight(c) by elf
//
                                        //
2011 revised by fookwood
using namespace std;
// clr >= 0 为合法颜色, clr = -1 的方块为哨兵
template<class T>
class Rect {
public:
  T x1, y1, x2, y2;
  int clr;
  Rect() {}
   Rect( T a, T b, T c, T d, int clr ):
     x1(a), y1(b), x2(c), y2(d), clr(clr) {}
  T area() const { return (y2-y1)*(x2-x1); }
};
const Rect<int> SENTINEL( 0, 0, 0, 0, -1 );
template<class T>
```

```
class RectCut {
   deque< Rect<T> > Q;
   map<int, T>
                     Μ;
public:
   // 在队列中放置一个 sentinel
   void clear() { Q.assign( 1, SENTINEL ); M.clear(); }
   void insert( T a, T b, T c, T d, int clr ) {
       insert( Rect<T>( a, b, c, d, clr ) );
   void insert( Rect<T> r ) {
       M[r.clr] += r.area();
       Q.push_back( r );
       while( Q.front().clr != -1 ) {
           Rect<T> z = Q.front();
           Q.pop_front();
           if( z.x1 >= r.x2 || z.x2 <= r.x1 ||
              z.y1 >= r.y2 || z.y2 <= r.y1 ) {
Q.push_back( z );
               continue;
           if( z.x1 < r.x1 ) {
               Q.push_back( Rect<T>( z.x1, z.y1, r.x1, z.y2, z.clr ) );
               z.x1 = r.x1;
           if(z.x2 > r.x2) {
               Q.push back( Rect<T>( r.x2, z.y1, z.x2, z.y2, z.clr ) );
               z.x2 = r.x2;
           if( z.y1 < r.y1 ) {
               Q.push_back( Rect<T>( z.x1, z.y1, z.x2, r.y1, z.clr ) );
               z.y1 = r.y1;
           if(z.y2 > r.y2) {
               Q.push_back( Rect<T>( z.x1, r.y2, z.x2, z.y2, z.clr ) );
               z.y2 = r.y2;
           if( ( M[z.clr] -= z.area() ) == 0 ) M.erase( M.find( z.clr ) );
       Q.push_back( Q.front() );
       Q.pop_front();
   T color( int co ){
       return M[co];
   vector<pair<int, T> > get_color() const {
       return vector<pair<int, T> >( M.begin(), M.end() );
   T total area(){
       T sum = 0;
       vector< pair<int,T> > V = get_color();
       for( int i = 0; i < V.size(); i++ )
           sum += V[i].second;
       return sum;
   }
};
```

```
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//演示程序~zoj 1128~
double a, b, c, d;
int T = 1, N;
RectCut<double> RC;
int main() {
   while(cin >> N && N) {
       RC.clear();
       while(N--) {
          cin >> a >> b >> c >> d;
          RC.insert(a, b, c, d, 1);
       printf("Test case #%d\n", T++);
       printf("Total explored area: %.21f\n",
              RC.get_color().front().second);
       puts("");
   }
}
8.10 最近公共祖先
const int maxn = 50001;
using namespace std;
class Node{
public:
   int to,w;
   Node *next;
   void add( int tt,int ww,Node *&b ){
       to = tt;w= ww;next = b;b=this;
   }
};
int ab(int a,int b){return a>b?(a-b):(b-a);}
Node *biao[ maxn ],a[ 200010 ];
int TIME, countt, N;
int founder[ maxn ];
int timepoint[ maxn*2 ];
int timedeep [ maxn*2 ];
int opt[ maxn*2 ][ 20 ];
int flag[ maxn ];
int d[ maxn ],pre[maxn];
void get(int u,int deep){
   timedeep[ founder[u] = TIME++ ] = deep;
   timepoint[ TIME-1 ] = u;
   flag[u] = 1;
inline void dfs( int u,int deep ){
   timedeep[ founder[u] = TIME++ ] = deep;
   timepoint[ TIME-1 ] = u;
   flag[u] = 1;
   for( Node*p=biao[u]; p != NULL; p = p->next )
       if( !flag[p->to] ){
          d[p->to] = d[u] + p->w;
          dfs( p->to , deep +1 );
          timepoint[TIME++] = u;
```

```
timedeep[TIME-1]= deep;
       }
}
int main(void){
   int i,from,m,to,w;
    int j,tt,countt;
    int start, end, u, v, k, s = 0;
   while( scanf("%d",&N) != EOF ){
       if( s++ ) cout << endl;</pre>
       for(i = 0; i < N; i++)
           biao[i] = NULL;
       countt = 0;
       for( i = 1; i < N; i++){
           scanf("%d%d%d",&from,&to,&w);
           a[countt++].add( to,w,biao[from] );
           a[countt++].add( from,w,biao[to] );
       }
       for( i = 0; i < N; i++ )
           flag[i] = 0;
       d[3*N/4] = 0;
       TIME = 0;
       dfs(3*N/4,0);
       for(i = 0; i < 2*N - 1; i++)
           opt[i][0] = i;
       for( j = 1; (1<<j) <= 2*N-1; j++){
           tt = 1 << j;
           for( i = 0; i + tt -1 < 2*N -1; i++){
               if( timedeep[ opt[i][j-1] ] <</pre>
                   timedeep[ opt[i+tt/2][j-1] ] )
                   opt[i][j] = opt[i][j-1];
                   opt[i][j] = opt[i+tt/2][j-1];
           }
       }
       cin >> m;
       while( m-- ){
           scanf("%d%d",&u,&v);
start = min( founder[u],founder[v] );
           end = max( founder[u],founder[v] );
           k = (int)(log((double)(end-start+1))/log(2.0));
           int p1 = opt[start][k];
           int p2 = opt[end-(1<< k)+1][k];
           int p;
           if( timedeep[p1] > timedeep[p2] )
               p = timepoint[ p2 ];
           else p = timepoint[ p1 ];
           cout << d[u]-d[p]+d[v]-d[p]<<endl;
       }
    }
   return 0;
}
```

# 9 计算几何

## 9.1 注意事项

1、可能输出-0.00000;

# 9.2 各种公式

//[Pick 定理] 设以整数点为顶点的多边形的面积为 S, 多边形内部的整数点数为 N, 多边形边界上的整数点数为 L, 则

N + L/2 - 1 = S

//涉及乘法的时候 注意保证不越界 而且要保证精度

```
费马点
```

对于三角形: 当三个角都小于 **120** 度,则费马点在三角形内部与任意两定点的连线构成的 角都为 **120** 度;

//若存在一个角大于等于 120 度,则费马点为此角顶点。

//对于平面四边形: 若为凸四边形,则费马点为两对角线交点;对于凹四边形,为其凹顶点。

//在不是凸四边形的情况下,只要枚举四个顶点即可。

//计算四面体体积,欧拉四面体公式

//a=0A, b=0B, c=0C, l=AB, m=BC, n=CA

double v\_4mianti(double a,double b,double c,double 1,double m)
{

return sqrt(4\*a\*a\*b\*b\*c\*c-a\*a\*(b\*b+c\*c-m\*m)\*(b\*b+c\*c-m\*m)-

 $b*b*(c*c+a*a-n*n)*(c*c+a*a-n*n)-c*c*(a*a+b*b-l*1)*(a*a+b*b-l*1)\\ +(a*a+b*b-l*1)*(b*b+c*c-m*m)*(c*c+a*a-n*n))/12;$ 

### 三角形:

- 1. 半周长 P=(a+b+c)/2
- 2. 面积 S=aHa/2=absin(C)/2=sqrt(P(P-a)(P-b)(P-c))
- 3. 中线 Ma=sqrt(2(b^2+c^2)-a^2)/2=sqrt(b^2+c^2+2bccos(A))/2
- 4. 角平分线 Ta=sqrt(bc((b+c)^2-a^2))/(b+c)=2bccos(A/2)/(b+c)
- 5. 高线 Ha=bsin(C)=csin(B)=sqrt(b^2-((a^2+b^2-c^2)/(2a))^2)
- 6. 内切圆半径 r=S/P=asin(B/2)sin(C/2)/sin((B+C)/2)

 $=4R\sin(A/2)\sin(B/2)\sin(C/2)=\operatorname{sqrt}((P-a)(P-b)(P-c)/P)$   $=P\tan(A/2)\tan(B/2)\tan(C/2)$ 

7. 外接圆半径 R=abc/(4S)=a/(2sin(A))=b/(2sin(B))=c/(2sin(C))

#### 四边形:

D1,D2 为对角线,M 对角线中点连线,A 为对角线夹角

- 1. a^2+b^2+c^2+d^2=D1^2+D2^2+4M^2
- 2. S=D1D2sin(A)/2

(以下对圆的内接四边形)

- 3. ac+bd=D1D2
- 4. S=sqrt((P-a)(P-b)(P-c)(P-d)),P 为半周长

#### 正 n 边形:

R 为外接圆半径, r 为内切圆半径

- 1. 中心角 A=2PI/n
- 2. 内角 C=(n-2)PI/n
- 3. 边长 a=2sqrt(R^2-r^2)=2Rsin(A/2)=2rtan(A/2)
- 4. 面积 S=nar/2=nr^2tan(A/2)=nR^2sin(A)/2=na^2/(4tan(A/2))

### 圆:

- 1. 弧长 l=rA
- 2. 弦长 a=2sqrt(2hr-h^2)=2rsin(A/2)
- 3. 弓形高 h=r-sqrt(r^2-a^2/4)=r(1-cos(A/2))=atan(A/4)/2
- 4. 扇形面积 S1=r1/2=r^2A/2
- 5. 弓形面积 S2=(rl-a(r-h))/2=r^2(A-sin(A))/2

### 棱柱:

- 1. 体积 V=Ah,A 为底面积,h 为高
- 2. 侧面积 S=1p,1 为棱长,p 为直截面周长
- 3. 全面积 T=S+2A

### 棱锥:

- 1. 体积 V=Ah/3,A 为底面积,h 为高
- (以下对正棱锥)
- 2. 侧面积 S=1p/2,1 为斜高,p 为底面周长
- 3. 全面积 T=S+A

### 棱台:

- 1. 体积 V=(A1+A2+sqrt(A1A2))h/3,A1.A2 为上下底面积,h 为高(以下为正棱台)
- 2. 侧面积 S=(p1+p2)1/2,p1.p2 为上下底面周长,1 为斜高
- 3. 全面积 T=S+A1+A2

#### 圆柱:

- 1. 侧面积 S=2PIrh
- 2. 全面积 T=2PIr(h+r)
- 3. 体积 V=PIr^2h

### 圆锥:

- 1. 母线 l=sqrt(h^2+r^2)
- 2. 侧面积 S=PIrl
- 3. 全面积 T=PIr(1+r)
- 4. 体积 V=PIr^2h/3

### 圆台:

- 1. 母线 l=sqrt(h^2+(r1-r2)^2)
- 2. 侧面积 S=PI(r1+r2)1
- 3. 全面积 T=PIr1(1+r1)+PIr2(1+r2)
- 4. 体积 V=PI(r1^2+r2^2+r1r2)h/3

### 球:

- 1. 全面积 T=4PIr^2
- 2. 体积 V=4PIr^3/3

### 球台:

- 1. 侧面积 S=2PIrh
- 2. 全面积 T=PI(2rh+r1^2+r2^2)
- 3. 体积 V=PIh(3(r1^2+r2^2)+h^2)/6

### 球扇形:

1. 全面积 T=PIr(2h+r0),h 为球冠高,r0为球冠底面半径

 2. 体积 V=2PIn/2h/3
 スコース2
 スコース3
 スコース4

 Euler 的任意四面体体製公式(日知边长求体积)
 VI
 V2
 V3
 V4

已知 4 点坐标求体积(其中四个点的坐标分别为(31,**y1**,z1),(x2,y2,z2),(x3,y3,z3),(x4,y4,z4))

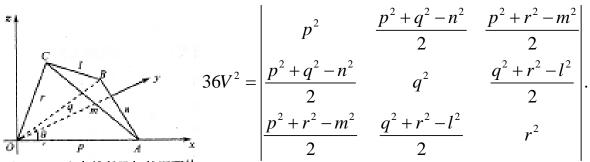


图 2.1 六条梭长已知的四面体

# 9.3 基本类型定义

```
//二维
struct point{ double x,y;}; //点
struct beeline{ double A,B,C;}; //直线 (直线方程)
struct line{ point a,b;}; //直线
                                        (两点式)
struct segment{ point a,b;}; // 线段
//三维
struct point3 {
    double x, y, z;
    point3(){};
    point3(double xx, double yy, double zz) : x(xx), y(yy), z(zz) {}
struct plane{ point3 a, b, c;};
9.4 基础函数
// 大小比较
const double eps = 1e-6;
bool dy(double x,double y) { return x > y + eps;} // x > y bool xy(double x,double y) { return x < y - eps;} // x < y bool dyd(double x,double y) { return x > y - eps;} // x >= y bool xyd(double x,double y) { return x < y + eps;} // x <= y bool dd(double x,double y) { return fabs( x - y ) < eps;} // x == y double disp2p(point a,point b) // a b 两点之间的距离
{
    return sqrt( (a.x - b.x ) * (a.x - b.x ) + (a.y - b.y ) * (a.y - b.y ));
}
// 叉积 (三点)
double crossProduct(point a, point b, point c)//向量 ac 在 ab 的方向 顺时针是
正
{
    return (c.x - a.x)*(b.y - a.y) - (b.x - a.x)*(c.y - a.y);
}
//叉积(三点坐标)
double crossProduct(double x0,double y0,double x1,double y1,double
x2, double y2)
{
    return (x2 - x0)*(y1 - y0) - (x1 - x0)*(y2 - x0);
}
9.5 各种极角排序
极角排序,以 C 为顶点,按象限极角排序,相同角度,距离近的排前面
```

int quad(point a)// 判断象限的函数,每个象限包括半个坐标轴

```
{
   if( dy(a.x,0) && xyd(a.y,0) ) return 1;
   if( xyd(a.x,0) && dy(a.y,0) ) return 2;
   if( xy(a.x,0) && xyd(a.y,0) ) return 3;
   if( dyd(a.x,0) && xy(a.y,0) ) return 4;
               // 和原点重合 , 一般不会有这种数据的。。
   return 0;
bool cmp(point& a,point& b)
   point p1 = a, p2 = b;
   p1.x -= C.x; p1.y -= C.y;
   p2.x -= C.x; p2.y -= C.y;
   int 11,12;
   11 = quad(p1); 12 = quad(p2);
   if( 11 == 12 )
      double c = crossProduct(C,a,b);
      return xy(c,0) \mid\mid dd(c,0.0) \&\& xy(fabs(a.x),fabs(b.x));
   return 11 < 12;
             极角排序,以C为顶点,按照向量排序,范围(-180°,+180°]
bool cmp(point& a,point& b)
{
   double t1 = atan2(a.y - C.y, a.x - C.x);
   double t2 = atan2(b.y - C.y, b.x - C.x);
   if( dd(t1, t2) ) return xy(fabs(a.x),fabs(b.x));
   return xy(a.t, b.t);
}
```

# 9.6 点、直线、线段

## 9.6.1 两直线位置关系

```
平行 垂直 相交即不平行
         线的位置关系:
//判断直线是否平行,平行返回1
bool parallel(point u1,point u2,point v1,point v2)
   return dd( (u1.x - u2.x)*(v1.y - v2.y) - (v1.x - v2.x)*(u1.y - u2.y),
0.0);
}
bool parallel(line u,line v)
   return dd( (u.a.x - u.b.x)*(v.a.y - v.b.y) - (v.a.x - v.b.x)*(u.a.y - v.b.y)
u.b.y), 0.0);
//判断两直线是否垂直,垂直返回1
bool perpendicular(point u1,point u2,point v1,point v2)
   return dd((u1.x - u2.x)*(v1.x - v2.x) + (u1.y - u2.y)*(v1.y - v2.y),0.0);
}
```

# 9.6.2 判断两线段是否相交

//判断线段 p1p2 与 p3p4 是否相交,包括端点与端点,端点与线段 const double eps = 1e-10;

```
struct point{ double x,y;
                                };
double crossProduct(point a, point b, point c )
bool onSegment(point a, point b, point c)
bool s2s_inst(point p1,point p2, point p3, point p4)
   double d1 = crossProduct(p3,p4,p1);
   double d2 = crossProduct(p3,p4,p2);
   double d3 = crossProduct(p1,p2,p3);
   double d4 = crossProduct(p1,p2,p4);
   if( xy(d1 * d2,0.0) && xy(d3 * d4,0.0) ) return true;
          //如果不判端点相交,则下面这句话不需要
                                                   Ш
   if(
          dd(d1,0.0)
                       &&
                            onSegment(p3,p4,p1)
                                                        dd(d2,0.0)
                                                                     &&
onSegment(p3,p4,p2)
         dd(d3,0.0)
                       &&
                            onSegment(p1,p2,p3)
                                                   Ш
                                                        dd(d4.0.0)
                                                                     &&
    Ш
onSegment(p1,p2,p4) )
       return true;
   return false;
}
```

### 9.6.3 判断两直线是否相交

不平行即相交

### 9.6.4 判断直线和线段是否相交

```
//判断直线和线段是否相交,共线的不算相交 double crossProduct(point a,point b,point c)//向量 ac 在 ab 的方向 bool s2l_inst(point s1,point s2,point l1,point l2)//s 是线段,1 是直线 { // xyd 包括端点在直线上。xy 是穿过 return xyd(crossProduct(l1,l2,s1) * crossProduct(l1,l2,s2),0.0); }
```

# 9.6.5 判断两向量位置

```
//判断向量 a1a2 与 b1b2 的位置
double rota_angle(point a1,point a2,point b1,point b2) //返回 b1b2 在 a1a2 的方向
{
    point t;
    t.x = b2.x - (b1.x - a1.x);
    t.y = b2.y - (b1.y - a1.y);
    return crossProduct(a1,a2,t);
}
```

### 9.6.6 求两直线交点

```
((11.a.x - 11.b.x)*(12.a.y - 12.b.y) - (11.a.y -
11.b.y)*(12.a.x - 12.b.x));
   ans.x += (11.b.x - 11.a.x)*t;
   ans.y += (11.b.y - 11.a.y)*t;
   return ans;
}
        9.6.7 求两线段交点
//求两线段的交点(先判断相交),然后按两直线求交点求得
point s2s_inst_p(point u1,point u2,point v1,point v2)
{
   point ans = u1;
   double t = ((u1.x - v1.x)*(v1.y - v2.y) - (u1.y - v1.y)*(v1.x - v2.x))/
            ((u1.x - u2.x)*(v1.y - v2.y) - (u1.y - u2.y)*(v1.x - v2.x));
   ans.x += (u2.x - u1.x)*t;
   ans.y += (u2.y - u1.y)*t;
   return ans;
}
        9.6.8 判断点是否在线段上
//判断点 c 是否在线段 ab 上。
double crossProduct(point a,point b,point c)
bool onSegment(point a, point b, point c)
{
   double maxx = max(a.x,b.x);
   double maxy = max(a.y,b.y);
   double minx = min(a.x,b.x);
   double miny = min(a.y,b.y);
   if( dd(crossProduct(a,b,c),0.0) && dyd(c.x,minx) && xyd(c.x,maxx) &&
dyd(c.y,miny) && xyd(c.y,maxy) )
      return true;
   return false;
}
        9.6.9 点到线的距离,点到线段的最小距离
                      ***************
         点到线的距离, 点到线段的最小距离
//点到直线的距离
double disp2l(point a,point l1,point l2)
{
   return fabs( crossProduct(a,11,12) )/disp2p(11,12);
}
//点到线段的最短距离
double crossProduct(point a,point b,point c)
double disp2p(point a,point b)
double disp2seg(point p,point 11,point 12)
{
   point t = p;
   t.x += 11.y - 12.y; t.y += 12.x - 11.x;
   if( dyd(crossProduct(l1,t,p)*crossProduct(l2,t,p),0.0) ) //包括点和线
段共线
      return xy(disp2p(p,11),disp2p(p,12))? disp2p(p,11):disp2p(p,12);
   return fabs( crossProduct(p,11,12) )/disp2p(11,12);
}
```

### 9.6.10判断点是否在线段上

```
判断点 c 是否在线段 ab 上。
//判断点 c 是否在线段 ab 上。 包括端点,如果不包括的话 把 xyd -> xy double crossProduct(point a,point b,point c)
bool onSegment(point a, point b, point c)
   if( dd(crossProduct(a,b,c),0.0) && dyd(c.x,min(a.x,b.x)) &&
       xyd(c.x,max(a.x,b.x))
                              &&
                                        dyd(c.y,min(a.y,b.y))
                                                                  &&
xyd(c.y,max(a.y,b.y)) )
       return true;
   return false;
}
         9.6.11 求垂点(不是垂足)
// 求一个点, 使得 ab 垂直于 1112
point foot line(point a, point 11, point 12) //ac 在 1112 的逆时针方向
{
   point c;
   c.x = a.x - 12.y + 11.y;
   c.y = a.y + 12.x - 11.x;
   return c;
}
         9.6.12 通过两点求直线方程
// 通过两点求直线 (Ax + By + C = 0)
                    double A,B,C; };
struct beeline{
beeline makeline(point a,point b)
{
   beeline line;
   line.A = a.y - b.y;
   line.B = a.x - b.x;
   line.C = -line.A * a.x - line.B*a.y;
   return line;
}
         9.6.13 向量的旋转
//向量的旋转 //底边线段 ab 绕 a 逆时针旋转角度 ang(弧度制),原 Whirl 函数
point Rotate(double ang, point a, point b)
{
   b.x -= a.x; b.y -= a.y;
   point c;
   c.x = b.x * cos(ang) - b.y * sin(ang) + a.x;
   c.y = b.x * sin(ang) + b.y * cos(ang) + a.y;
   return c;
}
         9.6.14 判断点是否在多边形内
//判断点是否在多边形内 凸凹均可
double crossProduct(point a,point b,point c)//向量 ac 在 ab 的方向
bool onSegment(point a, point b, point c)
bool s2s_inst(point p1,point p2, point p3, point p4)
// p 中按顺时针或者逆时针方向存多边形的点, n 为多边形的点数
bool point inPolygon(point pot,point p[],int n)
```

```
{
   int count = 0;
   point a = pot,b;
   b.x = 1e20; b.y = pot.y;
   p[n] = p[0];
   for(int i=0; i<n; i++)
       if( onSegment(p[i],p[i+1],pot) ) return true;
       if( !dd(p[i].y,p[i+1].y) )
       {
           int tmp = -1;
           if( onSegment(a,b,p[i]) )
              tmp = i;
           else
              if( onSegment(a,b,p[i+1]) )
                  tmp = i+1;
           if( tmp != -1 \& dd(p[tmp].y,max(p[i].y,p[i+1].y)) ||
              tmp == -1 \&\& s2s_inst(p[i],p[i+1],a,b))
              count++;
       }
   return count % 2 == 1;
}
//判断点 a 是否在凸多边形内,p 数组至少开[n + 2]个点
bool pin_convexh(point *p,int n,point a)
{
   p[n] = p[0]; p[n+1] = p[1];
   for(int i=0; i<n; i++)
       if( xy(crossProduct(p[i],p[i+1],a)*
           crossProduct(p[i+1],p[i+2],a),0.0)
           return false;
   return true;
}
```

# 9.7 三角形

# 9.7.1 计算三角形面积,边长

```
double area_triangle(double a,double b,double c)
   double p = (a+b+c)/2.0;
   return sqrt(p*(p-a)*(p-b)*(p-c));
//计算三角形面积(三点,坐标)
double area_triangle(double x0,double y0,double x1,double y1,double
x2, double y2)
{
   return fabs( crossProduct(x0,y0,x1,y1,x2,y2) )/2.0;
//计算三角形面积(已知三条中线长)
double area midline(double x,double y,double z)
{
   return 4.0/3*area triangle(x,y,z);
//公式 已知中线 x y z 求边长 a b c
void triangle_edge(double x,double y,double z,double &a,double &b,double
&c)
{
   a = 2.0/3.0*sqrt(2*x*x + 2*z*z - y*y);
   b = 2.0/3.0*sqrt(2*y*y + 2*z*z - x*x);
   c = 2.0/3.0*sqrt(2*x*x + 2*y*y - z*z);
}
        9.7.2 计算三角形的外接圆半径,内切圆半径
double circumcir_r(double a,double b,double c)//已知三边求外接圆半径
{
   return a*b*c/(4*area_triangle(a,b,c));
double circumcir r(point ap, point bp, point cp)//已知三点求外接圆半径
{
   double a,b,c;
   a = disp2p(ap,bp); b = disp2p(bp,cp); c = disp2p(cp,ap);
   return a*b*c/sqrt((a+b+c)*(a+b-c)*(a+c-b)*(b+c-a));
}
double incir r(double a,double b,double c)//已知三边求内切圆半径
   return 2*area_triangle(a,b,c)/(a+b+c);
double incir_r(point ap, point bp, point cp)//已知三点求内切圆半径
   double a,b,c;
   a = disp2p(ap,bp); b = disp2p(bp,cp); c = disp2p(cp,ap);
   return area_3p(ap,bp,cp)*2.0/(a+b+c);
}
         9.7.3 各种心(外心,内心,垂心,重心,费马点)
//计算三角形重心
//到三角形三点距离的平方和最小的点
//三角形内到三边距离之积最大的点
point 121 inst p(point u1,point u2,point v1,point v2)
point barycenter(point a, point b, point c)
{
   point u1,u2,v1,v2;
   u1.x = (a.x + b.x)/2;
   u1.y = (a.y + b.y)/2;
```

```
v1.x = (a.x + c.x)/2;
   v1.y = (a.y + c.y)/2;
   return 121_inst_p(u1,c,v1,b);
//三角形的内心
point incenter(point a,point b,point c)
   line u,v;
   double m,n;
   u.a = a;
   m = atan2(b.y - a.y, b.x - a.x);
   n = atan2(c.y - a.y, c.x - a.x);
   u.b.x = u.a.x + cos((m+n)/2);
   u.b.y = u.a.y + sin((m+n)/2);
   v.a = b;
   m = atan2(a.y - b.y, a.x - b.x);
   n = atan2(c.y - b.y, c.x - b.x);
   v.b.x = v.a.x + cos((m+n)/2);
   v.b.y = v.a.y + sin((m+n)/2);
   return 121_inst_p(u,v);
//计算三角形外心(外接圆圆心)
point 121 inst p(line u,line v)
point circumcenter(point a,point b,point c)
{
   point ua,ub,va,vb;
   ua.x = (a.x + b.x)/2;
   ua.y = (a.y + b.y)/2;
   ub.x = ua.x - a.y + b.y;//根据 垂直判断,两线段点积为 0
   ub.y = ua.y + a.x - b.x;
   va.x = (a.x + c.x)/2;
   va.y = (a.y + c.y)/2;
   vb.x = va.x - a.y + c.y;
   vb.y = va.y + a.x - c.x;
   return 121_inst_p(ua,ub,va,vb);
//这种不需要相交的那个函数 (外接圆圆心)
point circumcenter(point a,point b,point c)
   point ret;
   double a1 = b.x - a.x, b1 = b.y - a.y, c1 = (a1*a1 + b1*b1)/2;
   double a2 = c.x - a.x, b2 = c.y - a.y, c2 = (a2*a2 + b2*b2)/2;
   double d = a1 * b2 - a2 * b1;
   ret.x = a.x + (c1*b2 - c2*b1)/d;
   ret.y = a.y + (a1*c2 - a2*c1)/d;
return ret;
}
//三角形的垂心(垂线的交点)
point perpencenter(point a,point b,point c)
   point ua,ub,va,vb;
   ua = c;
   ub.x = ua.x - a.y + b.y;
   ub.y = ua.y + a.x - b.x;
   va = b;
   vb.x = va.x - a.y + c.y;
   vb.y = va.y + a.x - c.x;
   return 121_inst_p(ua,ub,va,vb);
//费马点
```

```
//到三角形三顶点距离之和最小的点
point fermentpoint(point a,point b,point c){
   point u,v;
   double
step=fabs(a.x)+fabs(a.y)+fabs(b.x)+fabs(b.y)+fabs(c.x)+fabs(c.y);
   int i,j,k;
   u.x=(a.x+b.x+c.x)/3;
   u.y=(a.y+b.y+c.y)/3;
   while (step>1e-10)
      for (k=0;k<10;step/=2,k++)
          for (i=-1;i<=1;i++)
             for (j=-1;j<=1;j++){
                v.x=u.x+step*i;
                v.y=u.y+step*j;
(distance(u,a)+distance(u,b)+distance(u,c)>distance(v,a)+distance(v,b)+
distance(v,c))
                    u=v;
   return u;
}
9.8 圆
         9.8.1 圆的位置关系
            ********************
          圆的位置关系: 判断 相交 相切 内含
                       求 交点 切点
          圆和矩形是否相交
                              *************************************
//判断两圆是否相交 (不包括相切)
bool c2c_inst(point a,double r1,point b,double r2)
{
   if(xy(disp2p(a,b),r1+r2) && dy(disp2p(a,b),fabs(r1 - r2)))
      return true:
   return false;
}
//求 直线与圆的交点(保证有交点)
// 也可计算 线段与圆的交点, 然后求得交点判断是否在线段上
void 12c_inst_p(point c,double r,point 11,point 12,point &p1,point &p2)
   point p = c;
   double t;
   p.x += 11.y - 12.y;
   p.y += 12.x - 11.x;
   p = 121_{inst_p(p,c,11,12)};
   t = sqrt(r*r - disp2p(p,c)*disp2p(p,c))/disp2p(11,12);
   p1.x = p.x + (12.x - 11.x)*t;
   p1.y = p.y + (12.y - 11.y)*t;
   p2.x = p.x - (12.x - 11.x)*t;
   p2.y = p.y - (12.y - 11.y)*t;
//求两圆交点(先判断相交)
//交得弧是弧 p1->p2
void c2c inst p(point c1,double r1,point c2,double r2,point &p1,point &p2)
{
   point u,v;
```

double t;

```
t = (1 + (r1*r1 - r2*r2)/disp2p(c1,c2)/disp2p(c1,c2))/2;
   u.x = c1.x + (c2.x - c1.x)*t;
   u.y = c1.y + (c2.y - c1.y)*t;
   v.x = u.x + c1.y - c2.y;
   v.y = u.y - c1.x + c2.x;
   12c_inst_p(c1,r1,u,v,p1,p2);
}
//判断两圆是否相切 (外切和内切)
bool c2c_tangent(point a,double r1,point b,double r2)
   if( dd(disp2p(a,b),r1+r2) \mid dd(disp2p(a,b),fabs(r1-r2)))
       return true;
   return false;
//求两圆切点(先判断相切,内外切交点不一样哎)
point c2c_tangent_p(point a,double r1,point b,double r2)
   point t;
   if( dd(disp2p(a,b),r1 + r2) ) // 外切交点
       t.x = (r1*b.x + r2*a.x)/(r1 + r2);
       t.y = (r1*b.y + r2*a.y)/(r1 + r2);
       return t;
   t.x = (r1*b.x - r2*a.x)/(r1 - r2);
   t.y = (r1*b.y - r2*a.y)/(r1 - r2);
   return t;
//判断两圆是否内含(a内含于b,不包括相切)
bool c2c ainb(point a, double r1, point b, double r2)
{
   return xy(disp2p(a,b),r2 - r1); //a 在 b 中,如果是包括内切,用 xyd
}
         9.8.2 两圆相交面积
double gongxing area(double r1, double r2, double 1) // 两圆半径以及连心线长
度
{
   double cosaa = (r1*r1 + 1*l - r2*r2)/(2*r1*l);
                                                      // cos 值
   double fcosaa = acos(cosaa)*2;
                                                  //夹角
                                                  // 扇形面积
   double Shu = fcosaa*r1*r1/2;
                                                  // 三角形面积
   double Ssan = r1*r1*sin(fcosaa)/2;
                                                  //扇形面积减去三角形面
   return (Shu - Ssan);
积
}
double c2c inst area(point a,double r1,point b,double r2)
   if( c2c_ainb(a,r1,b,r2) ) return pi*r1*r1; // a 内含于 b if( c2c_ainb(b,r2,a,r1) ) return pi*r2*r2; // b 内含于 a
   if( !c2c_inst(a,r1,b,r2) ) return 0;
                                       // 如果单纯求相交面积, 需判断相交,
   double l = disp2p(a,b);
从这儿到下面这段
   return gongxing_area(r1,r2,1) + gongxing_area(r2,r1,1);
}
```

### 9.8.3 判断圆和矩形是否相交

```
//判断圆和矩形是否相交,包括只交一个点
bool c2r_inst(point a,double r,point p[])
   bool flag = false;
   for(int i=0; i<4; i++)//如果严格相交,不要等号
      if( dyd(disp2p(a,p[i]),r) )
         flag = true;
   if( !flag ) return false;
   flag = false;
   for(int i=0; i<4; i++)
      if( xyd(disp2seg(a,p[i],p[(i+1)%4]),r) )
         flag = true;
   if( !flag ) return false;
   return true;
}
        9.8.4 最小覆盖圆
n 个点的最小覆盖圆 O(n)算法
                                     ************
point circumcenter(point a,point b,point c)
void min_cover_circle(point p[],int n,point &c,double &r)
   random_shuffle(p,p+n);// #include <algorithm>
   c = p[0]; r = 0;
   for(int i=1; i<n; i++)</pre>
      if( dy(disp2p(p[i],c),r) )
         c = p[i]; r = 0;
         for(int k=0; k<i; k++)</pre>
             if( dy(disp2p(p[k],c),r) )
                c.x = (p[i].x + p[k].x)/2;
                c.y = (p[i].y + p[k].y)/2;
                r = disp2p(p[k],c);
                for(int j=0; j<k; j++)</pre>
                   if( dy(disp2p(p[j],c),r) )
                                          // 求外接圆圆心,三点必不共
线
                      c = circumcenter(p[i],p[k],p[j]);
                      r = disp2p(p[i],c);
                   }
             }
      }
}
        9.8.5 扇形重心距圆心距离
// 扇形重心距圆心距离
//Xc = 2*R*sinA/3/A
//A 为圆心角的一半
double dis_z2c(double r, double angle)
   return 2 * r * sin(angle) / 3 / angle;
}
```

# 9.9 多边形

# 9.9.1 判断两个矩形是否相交

```
//判断两个矩形是否相交(即有公共面积),用左下角和右上角点表示矩形
bool r2r inst(point a1,point a2,point b1,point b2)
{
   if( b1.x >= a2.x \mid \mid b2.x <= a1.x \mid \mid b1.y >= a2.y \mid \mid b2.y <= a1.y )
      return false;
   return true;
}
         9.9.2 多边形面积
// 计算多边形的面积,可以根据面积判断是哪种方向
//按照我写的这个,顺时针的话,没有加绝对值的面积为负
double area_polygon(point p[],int n)
   if( n < 3 ) return 0.0;
   double s = 0.0;
   for(int i=0; i<n; i++)
      s += p[(i+1)\%n].y * p[i].x - p[(i+1)\%n].x * p[i].y;
   return fabs(s)/2.0;
//改变时针方向
void change_wise(point p[],int n)
   for(int i=0; i<n/2; i++)
      swap(p[i],p[n-i-1]);
}
         9.9.3 多边形的重心
//计算多边形的重心(凸凹均可)
double crossProduct(point a, point b, point c)//向量 ac 在 ab 的方向
point bary center(point p[],int n)
   point ans,t;
   double area = 0.0,t2;
   ans.x = 0.0; ans.y = 0.0;
   for(int i=1; i<n-1; i++)
      t2 = crossProduct(p[i],p[0],p[i+1])/2.0;
      ans.x += (p[0].x + p[i].x + p[i+1].x)*t2;
      ans.y += (p[0].y + p[i].y + p[i+1].y)*t2;
      area += t2;
   ans.x /= (3*area);
   ans.y /= (3*area);
   return ans;
}
9.10 凸包相关
```

# 9.10.1 二维凸包

//判断一个多边形是否是凸包,判断线旋转方向即可bool is convexhull(point p[],int n)

```
{
   for(int i=0; i<n; i++)</pre>
      if( xy( crossProduct(p[i],p[(i+1)%n],p[(i+2)%n]) *
         crossProduct(p[(i+1)\%n],p[(i+2)\%n],p[(i+3)\%n]),0.0))
         return false;
   return true;
}
//凸包无序排序 s 为内点
point s;
bool cmp(point a,point b)
   double l1 = crossProduct(s,c[0],a);
   double 12 = crossProduct(s,c[0],b);
   if( dyd(11,0.0) && dyd(12,0.0) )
      return xy(crossProduct(s,a,b),0.0);
   if( dyd(11,0.0) && xyd(12,0.0) )
      return 0;
   if( xyd(11,0.0) && xyd(12,0.0) )
      return xy(crossProduct(s,a,b),0.0);
   return 1;
void sort chull()
   s.x = (c[0].x + 2 * c[1].x + c[2].x)/4; // S 求法:线的中点,俩中点再求
中点
   s.y = (c[0].y + 2 * c[1].y + c[2].y)/4;
   int tmp = 0;
   for(int i=1; i<n; i++)</pre>
                                  11
      if(
             xy(c[i].x,c[tmp].x)
                                         dd(c[i].x,c[tmp].x)
                                                              &&
xy(c[i].y,c[tmp].y) )
         tmp = i;
   swap(c[tmp],c[0]);
   sort(c+1,c+n,cmp);
求凸包
double crossProduct(point a,point b,point c)
double disp2p(point a,point b)
bool cmp(point a, point b) // 排序
   double len = crossProduct(c[0],a,b);
   if( dd(len,0.0) )
      return xy(disp2p(c[0],a),disp2p(c[0],b));
   return xy(len,0.0);
}
int stk[MAX];
int top;
//形成凸包条件: 输入点个数大于等于三且至少有三个点不共线
//纯净凸包(即 stk 中没有三点共线)
void Graham(int n)
{
   int tmp = 0;
   for(int i=1; i<n; i++)
             xy(c[i].x,c[tmp].x)
                                 - 11
                                        dd(c[i].x,c[tmp].x)
      if(
                                                              &&
xy(c[i].y,c[tmp].y) )
         tmp = i;
   swap(c[0],c[tmp]);
   sort(c+1,c+n,cmp);
```

```
stk[0] = 0; stk[1] = 1;
   top = 1;
   for(int i=2; i<n; i++)
      while( xyd( crossProduct(c[stk[top]],c[stk[top-1]],c[i]), 0.0 ) &&
top >= 1)
          top--;
      stk[++top] = i;
   }
void Graham(int n)
{
   int tmp = 0;
   for(int i=1; i<n; i++)</pre>
      if(
              xy(c[i].x,c[tmp].x)
                                   dd(c[i].x,c[tmp].x)
                                                                  &&
xy(c[i].y,c[tmp].y) )
          tmp = i;
   swap(c[0],c[tmp]);
   sort(c+1,c+n,cmp);
   stk[0] = c[0]; stk[1] = c[1];
   top = 1;
   for(int i=2; i<n; i++)
      while( xyd( crossProduct(stk[top],stk[top-1],c[i]), 0.0 ) && top >=
1)
          top--;
      stk[++top] = c[i];
   }
}
         9.10.2 三维凸包
//三维凸包模板,可以求三维凸包表面积,体积,表面多边形数和表面三角形数
#define eps 1e-7
#define MAXV 310
//三维点
struct pt{
   double x, y, z;
   pt(){}
   pt(double _x, double _y, double _z): x(_x), y(_y), z(_z){}
   pt operator - (const pt p1){return pt(x - p1.x, y - p1.y, z - p1.z);}
   pt operator * (pt p){return pt(y*p.z-z*p.y, z*p.x-x*p.z, x*p.y-y*p.x);}
//叉乘
   double
                                                  x*p.x+y*p.y+z*p.z;}
                              (pt
                                     p){return
             operator
//点乘
};
struct _3DCH{
   struct fac{
                     //表示凸包一个面上三个点的编号
      int a, b, c;
                     //表示该面是否属于最终凸包中的面
      bool ok;
   };
           //初始点数
   int n;
   pt P[MAXV];
                 //初始点
             //凸包表面的三角形数
   fac F[MAXV*8]; //凸包表面的三角形
```

```
int to[MAXV][MAXV];
   double vlen(pt a){return sqrt(a.x*a.x+a.y*a.y+a.z*a.z);}
                                                              //向量长度
   double area(pt a, pt b, pt c){return vlen((b-a)*(c-a));}
                                                              //三角形面
积*2
   double volume(pt a, pt b, pt c, pt d){return (b-a)*(c-a)^(d-a);}
                                                                     //
四面体有向体积*6
   //正: 点在面同向
   double ptof(pt &p, fac &f){
       pt m = P[f.b]-P[f.a], n = P[f.c]-P[f.a], t = p-P[f.a];
       return (m * n) ^ t;
   void deal(int p, int a, int b){
       int f = to[a][b];
       fac add;
       if (F[f].ok){
           if (ptof(P[p], F[f]) > eps)
              dfs(p, f);
              add.a = b, add.b = a, add.c = p, add.ok = 1;
              to[p][b] = to[a][p] = to[b][a] = cnt;
              F[cnt++] = add;
           }
       }
   }
   void dfs(int p, int cur){
       F[cur].ok = 0;
       deal(p, F[cur].b, F[cur].a);
       deal(p, F[cur].c, F[cur].b);
       deal(p, F[cur].a, F[cur].c);
   }
   bool same(int s, int t){
       pt &a = P[F[s].a], &b = P[F[s].b], &c = P[F[s].c];
       return fabs(volume(a, b, c, P[F[t].a])) < eps && fabs(volume(a, b,
c, P[F[t].b])) < eps && fabs(volume(a, b, c, P[F[t].c])) < eps;</pre>
   //构建三维凸包
   void construct(){
       cnt = 0;
       if (n < 4)
           return;
       /*******此段是为了保证前四个点不公面,若已保证,可去掉******/
       bool sb = 1;
       //使前两点不公点
       for (int i = 1; i < n; i++){
           if (vlen(P[0] - P[i]) > eps){
              swap(P[1], P[i]);
              sb = 0;
              break;
           }
       if (sb)return;
       sb = 1;
```

```
//使前三点不公线
   for (int i = 2; i < n; i++){
       if (vlen((P[0] - P[1]) * (P[1] - P[i])) > eps){
          swap(P[2], P[i]);
          sb = 0;
          break;
       }
   if (sb)return;
   sb = 1;
   //使前四点不共面
   for (int i = 3; i < n; i++){
       if (fabs((P[0] - P[1]) * (P[1] - P[2]) ^ (P[0] - P[i])) > eps){
          swap(P[3], P[i]);
          sb = 0;
          break;
       }
   if (sb)return;
   /*******此段是为了保证前四个点不公面******/
   fac add;
   for (int i = 0; i < 4; i++){
       add.a = (i+1)%4, add.b = (i+2)%4, add.c = (i+3)%4, add.ok = 1;
       if (ptof(P[i], add) > 0)
           swap(add.b, add.c);
       to[add.a][add.b] = to[add.b][add.c] = to[add.c][add.a] = cnt;
       F[cnt++] = add;
   }
   for (int i = 4; i < n; i++){
       for (int j = 0; j < cnt; j++){
           if (F[j].ok && ptof(P[i], F[j]) > eps){
              dfs(i, j);
              break;
           }
       }
   int tmp = cnt;
   cnt = 0;
   for (int i = 0; i < tmp; i++){
       if (F[i].ok){
          F[cnt++] = F[i];
       }
   }
//表面积
double area(){
   double ret = 0.0;
   for (int i = 0; i < cnt; i++){
       ret += area(P[F[i].a], P[F[i].b], P[F[i].c]);
   return ret / 2.0;
//体积
double volume(){
```

}

}

```
pt 0(0, 0, 0);
       double ret = 0.0;
       for (int i = 0; i < cnt; i++){
           ret += volume(0, P[F[i].a], P[F[i].b], P[F[i].c]);
       return fabs(ret / 6.0);
   }
   //表面三角形数
   int facetCnt tri(){
       return cnt;
   }
   //表面多边形数
   int facetCnt(){
       int ans = 0;
       for (int i = 0; i < cnt; i++){
           bool nb = 1;
           for (int j = 0; j < i; j++){
   if (same(i, j)){</pre>
                  nb = 0;
                  break;
               }
           }
           ans += nb;
       return ans;
   }
};
              //内有大数组,不易放在函数内
3DCH hull;
int main()
{
   while (~scanf("%d", &hull.n)){
       for (int i = 0; i < hull.n; i++)</pre>
           scanf("%lf%lf%lf", &hull.P[i].x, &hull.P[i].y, &hull.P[i].z);
       hull.construct();
       printf("%d\n", hull.facetCnt());
   return 0;
}
9.11 旋转卡壳
//旋转卡壳 先求凸包(纯净),得出 stk 数组, n 为 top+1
// 返回最远点对距离
double crossProduct(Point a, Point b, Point c)//向量 ac 在 ab 的方向
double RC_maxdisp2p_inhull(int stk[],int n)
{
   stk[n] = stk[0];
   int q = 1;
   double ans = 0.0;
   for(int i=0; i<n; i++)</pre>
       while( xy(fabs(crossProduct(c[stk[i]],c[stk[i+1]],c[stk[q]])),
   fabs(crossProduct(c[stk[i]],c[stk[i+1]],c[stk[(q+1)%n]]))) )
           q = (q+1)%n;
```

```
ans = max(ans,disp2p(c[stk[i]],c[stk[q]]));
   return ans;
}
//旋转卡壳求两个凸包之间最短距离
double crossProduct(point a,point b,point c)
double disp2p(point a,point b)
bool cmp(point a,point b) // 凸包排序
void Graham(point c[],point stk[],int &top,int n) // 凸包
double disp2seg(point p,point l1,point l2)// 点到线段的最短距离
double rota angle(point a1,point a2,point b1,point b2) //返回 b1b2 在 a1a2
double RC mindish2h(point p[],point q[],int np,int nq) // 返回最短距离
{
   int sp = 0, sq = 0;
   for(int i=1; i<np; i++)</pre>
       if(
              xy(p[i].y,p[sp].y)
                                     Ш
                                             dd(p[sp].y,p[i].y)
                                                                   &&
xy(p[i].x,p[sp].x))
          sp = i;
                        //求得第一个凸包最左下角的点
   for(int i=1; i<nq; i++)</pre>
                                      Ш
       if(
               dy(q[i].y,q[sq].y)
                                             dd(q[sq].y,q[i].y)
                                                                   &&
dy(q[i].x,q[sq].x))
          sq = i;
                        //求得第二个凸包最右上角的点
   int tp = sp,tq = sq;
   double ans = disp2p(p[sp],q[sq]);
   do
   {
       double len = rota angle(p[sp],p[(sp+1)%np],q[(sq+1)%nq],q[sq]);
       if( dd(len,0.0) )
                                      // 卡壳正好和俩凸包的边重合
          ans = min(ans,disp2seg(p[sp],q[sq],q[(sq+1)%nq]));
          ans = min(ans,disp2seg(p[(sp+1)%np],q[sq],q[(sq+1)%nq]));
          ans = min(ans,disp2seg(q[sq],p[sp],p[(sp+1)%np]));
          ans = min(ans,disp2seg(q[(sq+1)%nq],p[sp],p[(sp+1)%np]));
          sp++; sp %= np; sq++; sq %= nq;
       }
       else
          if( xy(len,0.0) )
                                      // 卡壳和第一个凸包的边重合
              ans = min(ans,disp2seg(q[sq],p[sp],p[(sp+1)%np]));
              sp++; sp %= np;
          }
          else
          {
              ans = min(ans,disp2seg(p[sp],q[sq],q[(sq+1)%nq]));// 卡克
和第二个凸包的边重合
              sq++; sq %= nq;
   }while( !(tp == sp && tq == sq) );
   return ans;
}
//求凸包最小面积外接矩形
double crossProduct(point a, point b, point c)//向量 ac 在 ab 的方向
double disp2p(point a,point b)
bool cmp(point a,point b) // 凸包排序
double disp2seg(point a,point 11,point 12)
```

```
point foot_line(point a,point l1,point l2) // 求一个点,使得 ab 垂直于 l1l2
double rota angle(point a1,point a2,point b1,point b2) //判断向量 a1a2 与
b1b2 的位置
double Graham(int n)
double RC minareaRectangle(point p[],int n) //返回面积
{
                  // 0 == ymin, 1 == xmin, 2 == ymax ,3 == xmax;
   memset(r,0,sizeof(r));
   for(int i=0; i<n; i++)</pre>
   {
       if( xy(p[i].y,p[r[0]].y) )
                                    r[0] = i;
       if( xy(p[i].x,p[r[1]].x) )
                                    r[1] = i;
       if( dy(p[i].y,p[r[2]].y) )
                                    r[2] = i;
       if( dy(p[i].x,p[r[3]].x) )
                                    r[3] = i;
   int tp = r[0];
   double area = inf;
   do
   {
       point t = foot_line(p[r[0]], p[r[0]], p[(r[0]+1)%n]);
       while( dy(rota_angle(t,p[r[0]],p[r[1]],p[(r[1]+1)%n]),0.0) )
           r[1]++, r[1] \%= n;
       while( dy(rota_angle(p[r[0]],t,p[r[3]],p[(r[3]+1)%n]),0.0) )
           r[3]++, r[3] \% = n;
       while( dy(disp2seg(p[(r[2]+1)%n],p[r[0]],p[(r[0]+1)%n]),
              disp2seg(p[r[2]],p[r[0]],p[(r[0]+1)%n])))
           r[2]++, r[2] \%= n;
       double a = disp2seg(p[r[2]],p[r[0]],p[(r[0]+1)%n]);
       t = foot_line(p[r[3]],p[r[0]],p[(r[0]+1)%n]);
       double b = disp2seg(p[r[1]],p[r[3]],t);
       area = min( area, a*b );
       r[0]++; r[0] \% = n;
   }while( r[0] != tp );
   return area;
//卡壳求最大三角形面积,凸包预处理
double RC(point *s,int n) // 点集是凸包
   int p,q,r;
   p = 0; q = 1; r = 2;
   double area = area_triangle(s[p],s[q],s[r]);
   for(p=0; p<n; p++)
       q = (p+1)%n; r = (p+2)%n;
       area = max(area,area_triangle(s[p],s[q],s[r]));
       while( xy(fabs(crossProduct(s[p],s[q],s[r])),
              fabs(crossProduct(s[p],s[q],s[(r+1)%n]))) && r != p )
       {
           area = max(area, area triangle(s[p], s[q], s[(r+1)%n]));
           r = (r+1)%n;
       if( r == p ) continue;
       int rr = (r+1)%n;
       while( q != rr && r != p )
           area = area = max(area,area_triangle(s[p],s[q],s[r]));
           while( xy(fabs(crossProduct(s[p],s[q],s[r])),
```

```
fabs(crossProduct(s[p],s[q],s[(r+1)%n]))) && r != p )
                 r = (r+1)%n;
          q = (q+1)%n;
      }
   }
   return area;
}
9.12 半平面交
// 线段 ab 向左方向推进 h 后得到 cd
void changepoint(point a,point b,point &c,point &d,double h)
   double len = disp2p(a,b);
   double dx = h / len * (a.y - b.y);
   double dy = h / len * (-a.x + b.x);
   c.x = a.x + dx; c.y = a.y + dy;
   d.x = b.x + dx; d.y = b.y + dy;
}
//半平面交 N^2 算法,注意初始化有界平面(初始化成矩形,p[n]=p[0])
//每增加一条直线 ab 切割 调用这个函数。最终切点存在 s 中,长度为 len。
void cut_hp(point a,point b,point *s,int &len)
{
   int tc = 0;
   point tp[MAX];
   for(int i=0; i<=len; i++)</pre>
      tp[i] = s[i];
   for(int i=0; i<len; i++)</pre>
   {
      if(xyd(crossProduct(a,b,tp[i]),0.0))//右侧区域的话是dyd
          s[tc++] = tp[i];
      if( xy(crossProduct(a,b,tp[i])*
             crossProduct(a,b,tp[i+1]),0.0) )
          s[tc++] = 121_inst_p(a,b,tp[i],tp[i+1]);
   s[tc] = s[0];
   len = tc;
}
// 半平面交 N*LOGN 注意初始化有界平面
// 半平面返回点集 s, 点集点的个数 len
//ln 是 n 个半平面,用 line 的两个端点表示,一律考虑线段左边是半平面
//如果 1n 中有线段右边是半平面的,可以交换线段两点坐标
point l2l_inst_p(line l1,line l2)
bool parallel(line u, line v)
bool equal_ang(line a,line b)// 第一次 unique 的比较函数
{
   return dd(a.ang,b.ang);
bool cmphp(line a, line b) // 排序的比较函数
   if( dd(a.ang,b.ang) ) return xy(crossProduct(b.a,b.b,a.a),0.0);
   return xy(a.ang,b.ang);
bool equal p(point a,point b)//第二次 unique 的比较函数
{
   return dd(a.x,b.x) && dd(a.y,b.y);
void makeline hp(double x1,double y1,double x2,double y2,line &1)
```

```
{
   1.a.x = x1; 1.a.y = y1; 1.b.x = x2; 1.b.y = y2;
   1.ang = atan2(y2 - y1,x2 - x1);
void makeline_hp(point a,point b,line &1) // 线段(向量 ab)左侧侧区域有效
   1.a = a; 1.b = b;
   1.ang = atan2(b.y - a.y,b.x - a.x); // 如果是右侧区域, 改成 a.y -
b.y,a.x - b.x
void inst hp nlogn(line *ln,int n,point *s,int &len)
{
   len = 0;
   sort(ln,ln+n,cmphp);
   n = unique(ln,ln+n,equal_ang) - ln;
   int bot = 0, top = 1;
   deq[0] = ln[0]; deq[1] = ln[1];
   for(int i=2; i<n; i++)
                          parallel(deq[top],deq[top-1])
                                                                      Ш
parallel(deq[bot],deq[bot+1]) )
           return ;
       while( bot < top && dy(crossProduct(ln[i].a,ln[i].b,
           121_inst_p(deq[top],deq[top-1])),0.0) )
       while( bot < top && dy(crossProduct(ln[i].a,ln[i].b,</pre>
           121_inst_p(deq[bot],deq[bot+1])),0.0) )
           bot++;
       deq[++top] = ln[i];
   while( bot < top && dy(crossProduct(deq[bot].a,deq[bot].b,</pre>
       121 inst p(deg[top],deg[top-1])),0.0) ) top--;
   while( bot < top && dy(crossProduct(deq[top].a,deq[top].b,</pre>
       121_inst_p(deq[bot],deq[bot+1])),0.0) ) bot++;
   if( top <= bot + 1 ) return ;</pre>
   for(int i=bot; i<top; i++)</pre>
       s[len++] = 121_inst_p(deq[i],deq[i+1]);
   if( bot < top + 1 ) s[len++] = 121_inst_p(deq[bot],deq[top]);</pre>
   len = unique(s,s+len,equal_p) - s;
}
9.13 整点相关
//计算多边形边的整点(包括顶点)
int gcd(int n,int m)
int intp_insegment(point a, point b)
int intp_edge(point p[],int n)
{
   int ans = n;
   for(int i=0; i<n; i++)
       ans += intp_insegment(p[i], p[(i+1)%n]);
   return ans;
// 求多边形内的整点个数,多边形顶点都是整点&& dx dy 不能为 0
double area polygon(point p[],int n)
int intp_inpolygon(point p[],int n)
   double area = area_polygon(p,n);
   int pinedge = intp_edge(p,n);
```

```
return (int)(area) - pinedge/2 + 1;
}
//计算在线段上的整点个数 (还是这个比较好)
int intp insegment(point a, point b)
   int aa = abs(b.y - a.y), bb = abs(b.x - a.x);
   if(aa == 0 && bb == 0) return 0;
   if(aa == 0)
                 return bb - 1;
   if(bb == 0)
                 return aa - 1;
   return gcd(aa, bb) - 1;
9.14 球面相关
//求大圆角度,返回大圆的圆心角,输入为弧度
double angle_3d(double lng1, double lat1, double lng2, double lat2)
          //经度,纬度,经度,纬度
              acos(cos(lat1)*cos(lat2)*cos(lng1
                                                          lng2)
sin(lat1)*sin(lat2));
//求大圆劣弧长度,输入为弧度
double dis_3d(double lng1, double lat1, double lng2, double lat2)
          //经度,纬度,经度,纬度
   double dlon = lng2 - lng1;
double dlat = lat2 - lat1;
  double a = pow((sin(dlat/2)), 2) + cos(lat1) * cos(lat2) * pow(sin(dlon/2),
2);
   double c = 2 * atan2(sqrt(a), sqrt(1-a));
   double d = r * c;
   return d;
//大地坐标转空间坐标系
//纬度, 经度,单位坐标, 输入为角度
   Point(double la, double lo) :
       x(\cos(\log PI/180) * \cos(\log PI/180)),
       y(sin(lo*PI/180) * cos(la*PI/180)),
       z(sin(la*PI/180)) {}
9.15 模拟退火求多边形费马点
const int MAX = 1010;
const double inf = 1e30;
const double eps = 1e-8;
const double pi = acos(-1.0);
const int N = 15; // 设定随机点的个数
const int L = 40; // 设定随机方向次数
bool dy(double x,double y)
                            {
                                 return x > y + eps;
                                                        // x > y
bool xy(double x,double y)
                                                       // x < y
                                 return x < y - eps;}
                            {
                            {
                                                        // x >= y
bool dyd(double x,double y)
                                 return x > y - eps;}
bool xyd(double x,double y)
                                                         // x <= y
                                  return x < y + eps;}
                                 return fabs(x - y) < eps;} // x ==
bool dd(double x,double y)
struct point { double x,y;
   point (double x, double y):x(x),y(y){}
   point ():x(0),y(0){}
double disp2p(point a,point b) // a b 两点之间的距离
{
   return sqrt(( a.x - b.x ) * ( a.x - b.x ) + ( a.y - b.y ) * ( a.y - b.y ));
```

```
}
point p[MAX];
point rp[MAX];
double len[MAX];
double min_dis(point a,point *p,int n)
{
   double min = inf;
   for(int i=0; i<n; i++)</pre>
       double len = disp2p(a,p[i]);
       if( xy(len,min) )
           min = len;
   return min;
bool check(point a,double x,double y)
   return dyd(a.x,0.0) && dyd(a.y,0) && xyd(a.x,x) && xyd(a.y,y);
point Rand(double x,double y)
   point c;
   c.x = (rand()\%1000 + 1) / 1000.0 * x;
   c.y = (rand()%1000 + 1) / 1000.0 * y;
   return c;
int main()
{
   int n,m,ncases;
   double x,y;
                          // time.h!!
   srand(time(NULL));
   scanf("%d",&ncases);
   while( ncases-- )
       scanf("%lf%lf%d",&x,&y,&n);
       for(int i=0; i<n; i++)
       scanf("%lf%lf",&p[i].x,&p[i].y);
for(int i=0; i<N; i++)</pre>
           rp[i] = Rand(x,y);
           len[i] = min_dis(rp[i],p,n);
       }
       point st;
       double step = max(x,y)/2;
       while( step > 0.001 )
       {
           for(int k=0; k<N; k++)
           {
               st = rp[k];
               for(int i=0; i<L; i++)
                   double ang = (rand()%1000+1)/1000.0*10*pi;
                   double xx = st.x + step*cos(ang);
                   double yy = st.y + step*sin(ang);
                   point t = point(xx,yy);
                   if( !check(t,x,y) ) continue;
                   double dis = min_dis(t,p,n);
                   if( dy(dis,len[k]) )
```

```
rp[k] = t;
                        len[k] = dis;
                   }
                }
            }
            step *= 0.8;
       int ind = 0;
       for(int i=1; i<N; i++)</pre>
            if( len[i] > len[ind] )
               ind = i;
       printf("The
                                  safest
                                                        point
                                                                            is
(%.11f, %.11f).\n",rp[ind].x,rp[ind].y);
return 0;
```

# 10 **DP**

# 10.1 各种背包

```
///////////1111111111
#define zero_one_pack(d, cost, weight, limit, i)
do {
   for (i = limit; i >= cost; --i)
       if (d[i] < d[i-cost] + weight)</pre>
           d[i] = d[i-cost] + weight;
} while (0)
#define complete_pack(d, cost, weight, limit, i)
do {
   for (i = cost; i <= limit; ++i)</pre>
       if (d[i] < d[i-cost] + weight)</pre>
           d[i] = d[i-cost] + weight;
} while (0)
#define multiple_pack(d, cost, weight, amount, limit, i, k)
do {
   if (cost * amount >= limit)
       complete_pack(d, cost, weight, limit, i);
   else {
       for (k = 1; k \le amount; k \le 1) {
           zero_one_pack(d, k*cost, k*weight, limit, i);
           amount -= k;
       zero_one_pack(d, amount*cost, amount*weight, limit, i);
} while (0)
struct node {
   int cost;
   int weight;
   int amount;
};
#define LIMIT 20
```

```
#define N 3
int
main(void)
   struct node obj[N];
   int d[LIMIT + 1];
   int i, j, k;
   printf("LIMIT: %d\n", LIMIT);
   srand(time(NULL));
   for (i = 0; i < N; ++i) {
   obj[i].cost = rand() % 10 + 1;</pre>
       obj[i].weight = rand() % 10 + 1;
       obj[i].amount = rand() % 3 + 1;
       printf("object %d: cost %d weight %d amount %d\n",
               i, obj[i].cost, obj[i].weight, obj[i].amount);
   }
   memset(d, 0, sizeof(d));
   for (i = 0; i < N; ++i)
       zero_one_pack(d, obj[i].cost, obj[i].weight, LIMIT, j);
   printf("zero one pack, max is %d\n", d[LIMIT]);
   memset(d, 0, sizeof(d));
   for (i = 0; i < N; ++i)
       complete_pack(d, obj[i].cost, obj[i].weight, LIMIT, j);
   printf("complete pack, max is %d\n", d[LIMIT]);
   memset(d, 0, sizeof(d));
   for (i = 0; i < N; ++i)
       multiple_pack(d, obj[i].cost, obj[i].weight, obj[i].amount, LIMIT,
j, k);
   printf("multiple pack, max is %d\n", d[LIMIT]);
   return 0;
}
////////////////2222222222222
0-1 背包
//m 为容量, n 为东西个数
memset(bag,0,sizeof(bag));
for(int i=1; i<=n; i++)
   for(int k=m; k>=w[i]; k--)
       if( bag[k-w[i]]+v[i] > bag[k] )
           bag[k] = bag[k-w[i]] + v[i];
cout << bag[m] << endl;</pre>
完全背包
for(int i=1; i<=n; i++)</pre>
   for(int k=w[i]; k<=c; k++)
           money[k] = max(money[k], money[k-w[i]]+w[i]);
cout << money[c] << endl;</pre>
多重背包
count = 1;
memset(hp,0,sizeof(hp));
for(int i=1; i<=n; i++)
```

```
cin >> mp[i] >> hurt[i] >> num[i];
// mp 为每个东西体积, hurt 为价值, num 为数目
for(int i=1; i<=n; i++)
   int sum = 0;
   for(int k=0; ; k++)
       int x = (int)pow(2,k);
                              // <math.h>
       if(sum + x > num[i])
          break;
       sum += x;
       h[count] = x*hurt[i];
       mmp[count] = x*mp[i];
       count++;
   int x = num[i] - sum;
   if(x == 0)
       continue;
   h[count] = x*hurt[i];
   mmp[count] = x*mp[i];
   count++;
for(int i=1; i<count; i++)</pre>
   for(int k=m; k>=mmp[i]; k--)
       if(hp[k] < hp[k-mmp[i]] + h[i])
          hp[k] = hp[k-mmp[i]] + h[i];
cout << hp[m] << endl;</pre>
10.2 最长不下降子序列
int main(void)
{
   int n,geshu,i,j,nTmp,s=0;
   int a[1005];
   int flag[1005];
   scanf("%d",&n);
   while(n--)
   {
       scanf("%d",&geshu);
       for(i = 0; i < geshu; i++)
          scanf("%d",&a[i]);
       flag[0] = 1;
       for( i = 1; i < geshu; i++ )//以下是关键代码
       {
          nTmp = 0;
          for( j = 0; j < i; j++)
              if( flag[j] > nTmp && a[i] >= a[j])//修改这里的符号
                 nTmp = flag[j];
          flag[i] = nTmp + 1;
       }
                                           //判断最长的子序列
       nTmp = max_element(flag,flag+geshu);
       printf(s++?"\n%d\n":"%d\n",nTmp);
   }
   return 0;
/////22222222
```

```
最长上升子序列
for(int i=1; i<=n; i++)
   cin >> los[i];
max[1] = 1;
for(int i=2; i<=n; i++)
{
   int maxl = 0;
   for(int k=1; k<i; k++)</pre>
       if(los[i] > los[k])
           if( max[k] > maxl)
              maxl = max[k];
   max[i] = maxl+1;
int outputmax = 0;
for(int i=1; i<=n; i++)</pre>
   if( max[i] > outputmax )
       outputmax = max[i];
10.3 最长公共子序列
//zoj 1953 最长公共子序列 带路径记录。。。 before 数组为记录的前驱
//最后求出的 c 数组即为最长公共子序列(只是其中一种可能)
int main()
{
   char a[110],b[110],c[110],temp;
   int dp[110][110],before[110][110][2];
   int i,j,t,count;
   int lena, lenb;
   while( cin >> a+1 >> b+1 )
   {
       lena = strlen(a+1);
       lenb = strlen(b+1);
       memset( dp,0,sizeof(dp) );
       for( i = 1; i <= lena; i++ )
           for( j = 1; j <= lenb; j++ )
              if( a[i] == b[j] )
              {
                  dp[i][j] = dp[i-1][j-1] + 1;
                  before[i][j][0] = i - 1;
                  before[i][j][1] = j - 1;
              }
              else
              {
                  if( dp[i-1][j] > dp[i][j-1] ){
                     dp[i][j] = dp[i-1][j];
                     before[i][j][0] = i - 1;
                     before[i][j][1] = j;
                  }
                  else{
                     dp[i][j] = dp[i][j-1];
                     before[i][j][0] = i;
                     before[i][j][1] = j - 1;
                  }
       i=lena;j=lenb;
       count = 0;
       while( i && j )
```

```
if( a[i] == b[j] )
               c[count++] = a[i];
           t = before[i][j][0];
           j = before[i][j][1];
           i = t;
       c[count] = '\0';
       for( i = 0,count--; i < count; i++,count-- )</pre>
           temp=c[i],c[i]=c[count],c[count]=temp;
       for( i = j = 1, t = 0; c[t]!='\0';)
           if(a[i] == c[t] \&\& b[j] == c[t])
               cout<<c[t],t++,i++,j++;
           else if( a[i] == c[t] )
               cout<<b[j++];</pre>
           else if( b[j] == c[t] )
               cout<<a[i++];
           else
               cout<<a[i++]<<b[j++];
       cout << a+i << b+j << endl;</pre>
   return 0;
//党姐的太长了 ><><><
最长公共子序列
int maxlen[MAXLEN][MAXLEN];
int main(void)
{
    char str[MAXLEN],line[MAXLEN];
   int lens, len1;
   while(cin >> str)
    {
       memset(maxlen,0,sizeof(maxlen));
       cin >> line;
       lens = strlen(str);
       lenl = strlen(line);
       for(int i=0; i<lens; i++)</pre>
           for(int k=0; k<len1; k++)</pre>
               if( str[i] == line[k] )
                   \max len[i+1][k+1] = \max len[i][k] + 1;
               else
                   if( maxlen[i+1][k] > maxlen[i][k+1] )
                       \max len[i+1][k+1] = \max len[i+1][k];
                   else
                       \max len[i+1][k+1] = \max len[i][k+1];
       cout << maxlen[lens][len1] << endl;</pre>
return 0;
10.4 最大子段和
最大子串和
int Maxsum(int n, int *a)
{
```

```
int sum = 0,b = 0;
   for(int i=0; i<n; i++)
       if(b>0)
           b += a[i];
       else
           b = a[i];
       if( b > sum )
           sum = b;
   return sum;
10.5 最大子矩阵和
int DP(int a[],int n)
   int i,f[101];
   int max = -2000000000;
   for(i = 2, f[1] = a[1]; i <= n; i++)
       if (f[i - 1] > 0)
           f[i] = f[i - 1] + a[i];
       else
           f[i] = a[i];
       if (f[i] > max)
           max = f[i];
   }
   return max;
}
int main(void)
   int n,i,he,j;
   int a[105][105],k;
   int sum[105], max;
   scanf("%d",&n);
   for( i = 1 ; i <= n ; i++ )
       for( j = 1 ; j <= n ; j++ )
          scanf("%d",&a[i][j]);
   max = -2000000000;
   for( i = 1 ; i <= n ; i++ )
       memset(sum,0,sizeof(sum));
       for(j = i ; j <= n ; j++ )
       {
           for(k = 1; k <= n; k++)
              sum[k] += a[j][k];
           he = DP(sum,n);
           if(he>max) max = he;
   }
   printf("%d\n",max);
   return 0;
}
```

# 11 杂7杂8

# 11.1 简单图判定

```
// 图序列判定,给定图的节点数还有每个点的度数,判断是否能构成简单图
//sum 是度的总和
const int MAX = 1010;
int deg[MAX];
bool is_simpleG(int n, int sum, int *deg)
   if( sum % 2 == 1 ) return false;
   sort(deg+1, deg+n+1, greater<int>());
   for(int i=1; i<=n; i++)
       int s = 0;
       for(int k=1; k<=i; k++)</pre>
           s += deg[k];
       int s2 = 0;
       for(int k=i+1; k<=n; k++)</pre>
           s2 += min(i, deg[k]);
       if(s > i*(i-1) + s2)
           return false;
   return true;
}
11.2\,\mathrm{KMP}
#define MAX 1001
using namespace std;
int next[MAX];
char s[MAX],t[MAX];
void get_next(char *s)
   int len = strlen(s);
   int i = 0, j = -1;
   next[0] = -1;
   while( i < len )
       if(j == -1 || s[i] == s[j])
       {
           i++, j++;
           next[i] = s[i] == s[j] ? next[j] : j;
       else
           j = next[j];
int KMP(char *s,char *t)
   get_next(t);
   int lens = strlen(s);
   int lent = strlen(t);
   int i = 0, j = 0;
   while( i < lens && j < lent )</pre>
       if( j == -1 || s[i] == t[j] )
           i++, j++;
```

```
else
           j = next[j];
   if( j >= lent )
       return i - lent;
   return -1;
int main()
   int n,pos;
   while( scanf("%s %s",&s,&t)!=EOF )
       pos = KMP(s,t);
       if(pos == -1)
           printf("Can't match\n");
           printf("%d\n",pos);
return 0;
11.3等价表达式
int count1,count2;
int optr[100],opnd[100];
int h[7][7] = \{\{1,1,-1,-1,-1,1,1\},\{1,1,-1,-1,-1,1,1\},\{1,1,1,1,-1,1,1\},
\{1,1,1,1,-1,1,1\},\{-1,-1,-1,-1,-1,0,-2\},\{1,1,1,1,-2,1,1\},\{-1,-1,-1,-1,-1\}
,-2,0}};
void push1(int x)
{
   opnd[count1++] = x;
}
void push2(int x)
{
   optr[count2++] = x;
int empty2()
   if(count2 == 0)
       return 1;
   return 0;
int pop1()
   count1--;
   return opnd[count1];
char pop2()
   count2--;
   return optr[count2];
int computer(int x,int y, char ch)
   switch(ch)
       case '+':return x+y;
       case '-':return x-y;
       case '*':return x*y;
       case '/':return x/y;
```

```
}
}
char gettop()
{
   return optr[count2-1];
}
int number(char ch)
{
   switch(ch)
   {
       case '+':return 0;
       case '-':return 1;
       case '*':return 2;
       case '/':return 3;
       case '(':return 4;
       case ')':return 5;
       case '#':return 6;
   }
int record(char a, char b)
   return h[number(a)][number(b)];
int main(void)
{
   char c,fir[50];
   int x,y,temp,ch,length,i,plus,result1,result2;
   plus = 10;
   gets(fir);
   length = strlen(fir);
   fir[length] = '#';
   fir[length+1] = '(0';
   push2('#');
   i = 0;
           while( fir[i]!='#' || gettop()!='#')
               c = fir[i];
               if( isdigit(c) || isalpha(c) )
                   if(isalpha(c))
                      push1(c-'a'+plus);
                      push1(c-'0');
                   i++;
               }
               else
               {
                   if(c == '+' || c== '-'||c=='*' || c=='('||c==')'||c=='#')
                   switch( record(gettop(),c) )
                       case -1: push2(c); i++; break;
                      case 0 : pop2(); i++; break;
                      case 1 : ch = pop2();
                               y = pop1();
                               x = pop1();
                               push1(computer(x,y,ch));
                               break;
```

```
}
                  else
                     i++;
              }
           result1= opnd[count1-1];
   printf("%d\n",result1);
return 0;
11.4字串种类数(copy)
#define MAXN 50000
#define MAXM 128
                         // 字母个数
int head[MAXN], succ[MAXN], succRank[MAXN];
int sa[MAXN], rank[MAXN];
int letter[MAXM];
void da(char* str, int len) {// 倍增算法
   memset(letter, 0, sizeof(letter));
   for (int i = 0; i < len; ++i)
       if (!letter[str[i]])
           letter[str[i]] = 1;
   int total = -1;
   for (int i = 0; i < MAXM; ++i)
       if (letter[i])
           letter[i] = ++total;
   memset(head, 255, sizeof(head));
for (int i = len - 1; i >= 0; --i) { // 由于 head 头指针的性质,所以要倒序添加,下同
       rank[i] = letter[str[i]];
       succ[i] = head[rank[i]];
       head[rank[i]] = i;
   }
   int j = 0;
   for (int i = 0; i < len; ++i) {
       while (head[j] == -1) ++j;
       sa[i] = head[j];
       head[j] = succ[head[j]];
   // 到此初始化排序完毕
   for (int k = 1; k < len; k <<= 1) {
       // 以下两个 for 对新一轮子串进行排序,由于第二关键字的排序在上一次排序
中己完成
       // 故仅对第一关键字进行排序即可
       for (int i = len - 1; i >= 0; --i)
                                              // 倒序添加
           if (sa[i] - k >= 0) {
              succRank[sa[i] - k] = rank[sa[i]];
succ[sa[i] - k] = head[rank[sa[i] - k]];
              head[rank[sa[i] - k]] = sa[i] - k;
       for (int i = len - 1; i >= len - k; --i) { // 倒序添加
                                          // 另外添加的尾字符的 rank 为
          succRank[i] = -1;
-1
           succ[i] = head[rank[i]];
           head[rank[i]] = i;
       j = 0; total = -1;
```

```
int preSuccRank = 0, preRank = 0;
       for (int i = 0; i < len; ++i) {
          while (head[j] == -1) ++j;
          sa[i] = head[j];
if (i == 0 ||
                           preRank != rank[sa[i]] || preSuccRank !=
succRank[sa[i]]) {
              preRank = rank[sa[i]];
              rank[sa[i]] = ++total;
                                        // 链表中保证不递减, 所以可以这么做
           } else {
              preRank = rank[sa[i]];
              rank[sa[i]] = total;
           preSuccRank = succRank[sa[i]];
           head[j] = succ[head[j]];
   }
}
int height[MAXN];
void calcHeight(char* str, int len) {
   int i, j, k = 0;
for (i = 0; i < len; ++i) {</pre>
                                // h[i] >= h[i - 1] - 1
       if (k != 0)
           --k;
       if (rank[i] == 0)
          continue;
       j = sa[rank[i] - 1];
       // 求 suffix(sa[rank[i] - 1]) 和 suffix(sa[rank[i]]) 的最大公共前缀
       while (str[i + k] == str[j + k]) // C 字符串以 '\0' 结尾
       height[rank[i]] = k; // h[i] = k;
   }
}
char buf[MAXN + 1];
int len;
int main() {
   int t;
   scanf("%d", &t);
   for (int cas = 1; cas <= t; ++cas) {
       scanf("%s", buf);
       len = strlen(buf);
       da(buf, len);
       calcHeight(buf, len);
       long long ans = len - sa[0];
       for (int i = 1; i < len; ++i)
          ans += len - sa[i] - height[i];
       printf("%lld\n", ans);
   return 0;
}
11.5 欧拉回路判断
先判连通
//有向图
               //判断是否是欧拉路,包括欧拉回路和欧拉道路
int isoula()
{
```

```
int i;
   int sum = 0,out[26];
   for(i=0; i<26; i++)
       if( ind[i] != outd[i] )
           out[sum] = i;
           sum++;
   if( sum == 0 )
   {
       i = 0;
       while( !used[i] )
           i++;
       beg = i;
       return 1;
   if( sum == 2 )
       int x = out[0], y = out[1];
       if(ind[x] - outd[x] == 1 && outd[y] - ind[y] == 1)
           beg = y;
           return 1;
       if(ind[y] - outd[y] == 1 && outd[x] - ind[x] == 1)
           beg = x;
           return 1;
   return 0;
输出欧拉路
stack<int> s;
void output(int u) // 用栈存输出编号,递归。。。
   MAP *head = alp[u];
   while( head != NULL )
       if( head->flag == 0 )
       {
           head->flag = 1;
           output(head->v);
           s.push(head->to);
       head = head->next;
   }
}
11.6 大数
const int maxn = 2100;
struct bign{
 int len, s[maxn];
 bign() {
   memset(s, 0, sizeof(s));
   len = 1;
  }
```

```
bign(int num) {
  *this = num;
bign(const char* num) {
  *this = num;
bign operator = (int num) {
  char s[maxn];
  sprintf(s, "%d", num);
  *this = s;
  return *this;
}
bign operator = (const char* num) {
  len = strlen(num);
  for(int i = 0; i < len; i++) s[i] = num[len-i-1] - '0';
  return *this;
string str() const {
  string res = "";
  for(int i = 0; i < len; i++) res = (char)(s[i] + '0') + res;
  if(res == "") res = "0";
 return res;
bign operator + (const bign& b) const{
  bign c;
  c.len = 0;
  for(int i = 0, g = 0; g \mid \mid i < max(len, b.len); i++) {
   int x = g;
   if(i < len) x += s[i];
   if(i < b.len) x += b.s[i];
   c.s[c.len++] = x % 10;
   g = x / 10;
  }
 return c;
void clean() {
 while(len > 1 && !s[len-1]) len--;
bign operator * (const bign& b) {
  bign c; c.len = len + b.len;
  for(int i = 0; i < len; i++)
   for(int j = 0; j < b.len; j++)
     c.s[i+j] += s[i] * b.s[j];
  for(int i = 0; i < c.len-1; i++){
   c.s[i+1] += c.s[i] / 10;
   c.s[i] \% = 10;
  }
 c.clean();
  return c;
bign operator - (const bign& b) {
  bign c; c.len = 0;
```

```
for(int i = 0, g = 0; i < len; i++) {
   int x = s[i] - g;
   if(i < b.len) \times -= b.s[i];
   if(x >= 0) g = 0;
   else {
     g = 1;
     x += 10;
   c.s[c.len++] = x;
  }
 c.clean();
 return c;
bool operator < (const bign& b) const{</pre>
 if(len != b.len) return len < b.len;</pre>
 for(int i = len-1; i >= 0; i--)
   if(s[i] != b.s[i]) return s[i] < b.s[i];</pre>
 return false;
bool operator > (const bign& b) const{
 return b < *this;
bool operator <= (const bign& b) {return !(*this > b);}
bool operator >= (const bign& b) {return !(*this < b);}</pre>
bool operator == (const bign& b) {
 return !(b < *this) && !(*this < b);
bign operator += (const bign& b) {
  *this = *this + b;
 return *this;
}
bign operator << ( int& n )</pre>
 if( *this == 0 || n == 0 ) return *this;
 bign t = *this;
 int i = t.len-1;
 t.len += n;
 int *p = t.s;
 for(; i >= 0; i--)
     p[i+n]=p[i],p[i] = 0;
 for( i = 0; i < n; i++ )
     p[i] = 0;
 return t;
}
  bign operator / (bign b )
 bign ans=0,big=*this,add,chu;
 bign one = 1;
 int t;
 while( big >= b )
     t = big.len - b.len;
     chu = b << t;
     if( chu > big )
```

```
chu = b << (--t);
        add = one<<t;
        while( chu <= big )</pre>
            big = big - chu;
            ans += add;
        }
    return ans;
    }
};
istream& operator >> (istream &in, bign& x) {
  string s;
  in >> s;
  x = s.c_str();
  return \overline{i}n;
}
ostream& operator << (ostream &out, const bign& x) {</pre>
  out << x.str();
  return out;
int main() {
    bign a,b,two=2,one=1,zero=0;
    int n,i;
    cin >> n;
    while( n-- )
    {
        cin >> a;
        b = a/two;
        if( a.s[0]%2==1 ) cout << b << endl;
        else
        {
            b = b-one;
            if( b.s[0]\%2 == 0 ) b=b-one;
            cout << b << endl;</pre>
        }
    return 0;
}
```