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Quantum error correction beyond the surface code

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Joint work with Qian Xu*, Chris Pattison, Nithin Raveendran, Dolev Bluvstein, Jonathan Wurtz, Bane Vasic, Mikhail Lukin, Liang Jiang, Hengyun (Harry) Zhou. **Nat. Phys.** **20**, 1084-1090

Challenge of Large-Scale Quantum Computation

Quantum error correction will
bridge this gap!

Physical error
rates today

10^{-1} 10^{-2} 10^{-3}



What large-scale quantum
algorithms require



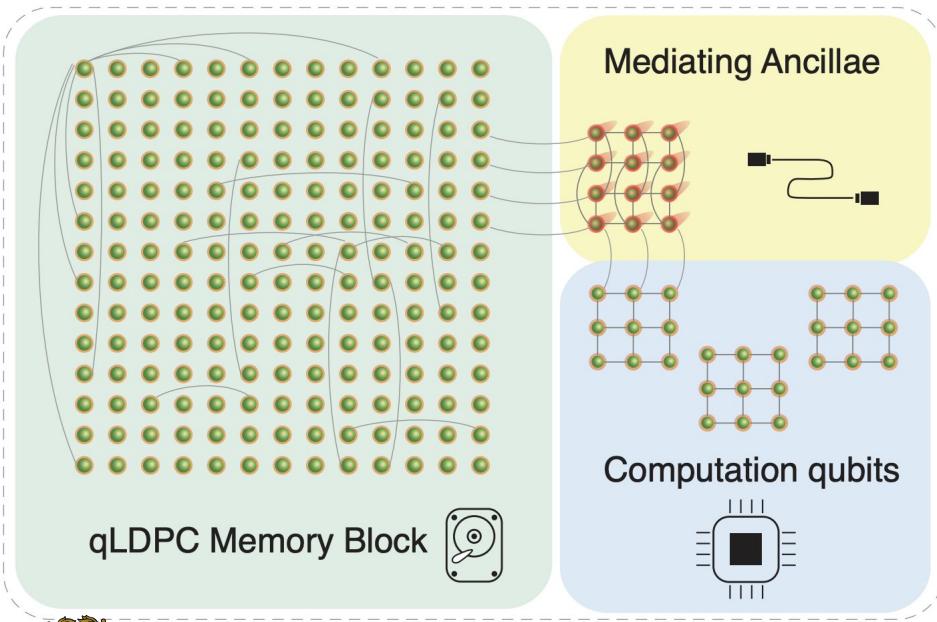
10^{-15}

Error rate of
encoded qubit

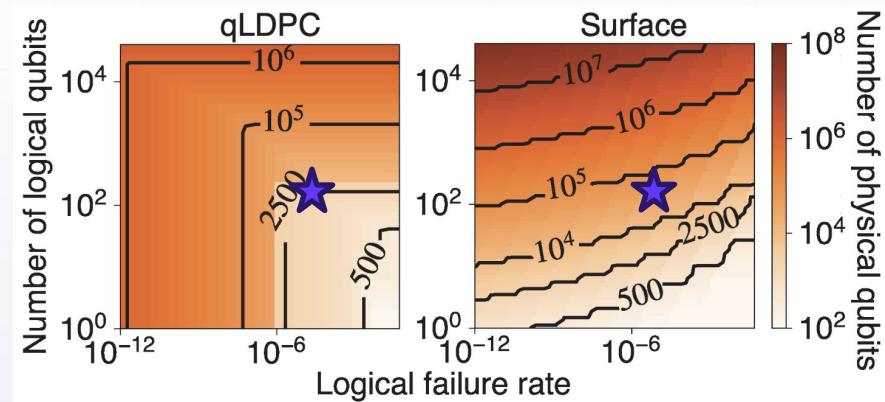


This Work: Bridging the Gap with qLDPC Codes

Implementation blueprint for memory & logical gates on a neutral atom quantum computer



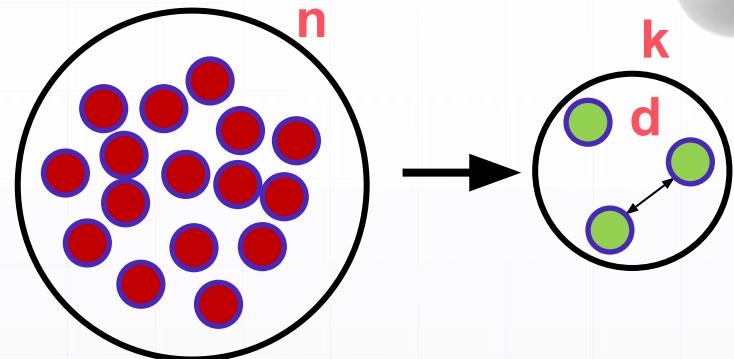
- Order of magnitude fewer physical qubits required, when replacing surface code with qLDPC code



Beyond the surface code?

- Good codes: $k = O(n), d = O(n)$
- Surface code: $k = O(1), d = O(\sqrt{n})$
- Best 2D-local codes: $kd^2 \leq O(n)$ (saturated by surface code)
- For improved properties: need non-locality!
- ... while preserving ldpc-ness (constant check weight + qubit degree)

The main message is that magic state distillation is *not* the dominant cost in a surface-code-based quantum computer. Rather, the large overhead of surface codes is due to their low encoding rate, which implies that a large number of qubits is required to simply store all data qubits of the computation. Litinski, Quantum (2019)



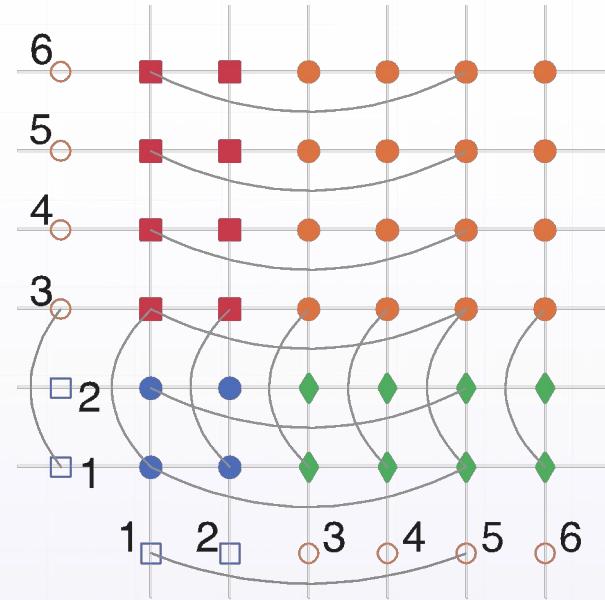
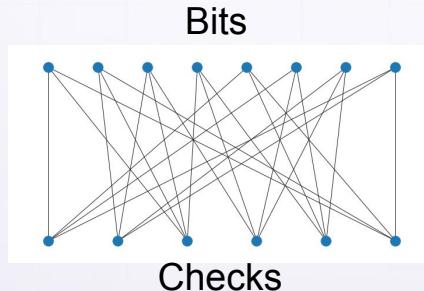
$[[n, k, d]]$ quantum code:

- n: number of physical qubits
- k: number of logical qubits
- d: code distance (error-correcting power)
- $R=k/n$: code rate



Product Structure of Some qLDPC Codes

- Hypergraph product code (HGP):
 - Construct quantum code by taking product between two classical codes
 - Example: surface code = HGP of two repetition codes
- Taking HGP of two classical codes with $[n, k = \Theta(n), d = \Theta(n)]$, we obtain a quantum code $[[n, k = \Theta(n), d = \Theta(\sqrt{n})]]$
- In this work: random (3,4)-biregular graphs

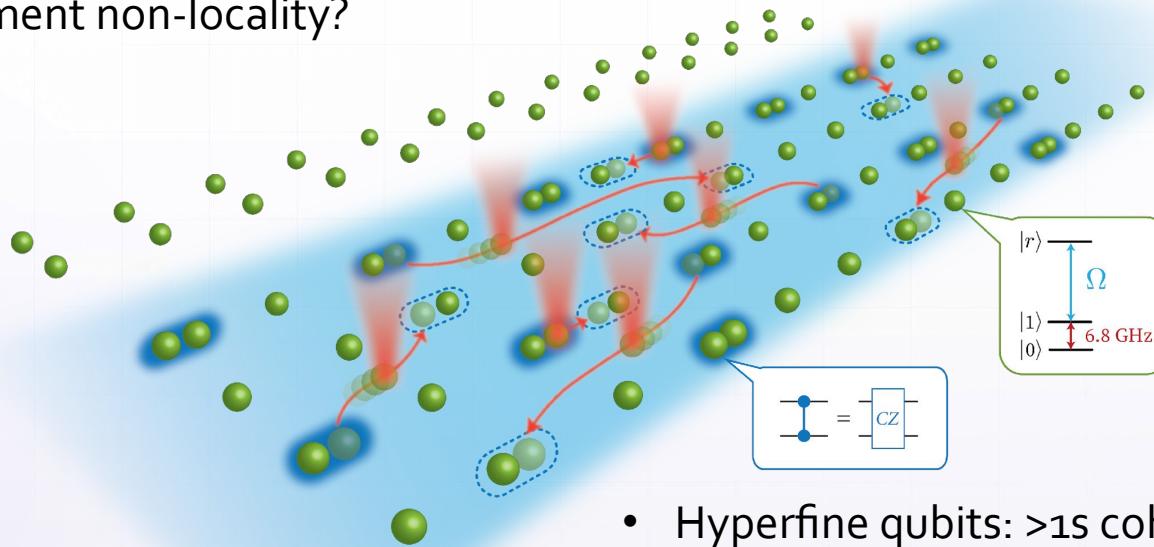


○ Classical bit □ Classical check
□ x □ = ● Data qubit ○ x □ = ● Z stabilizer
○ x ○ = ■ X stabilizer



Dynamically Reconfigurable Architecture with Neutral Atoms

How to implement non-locality?



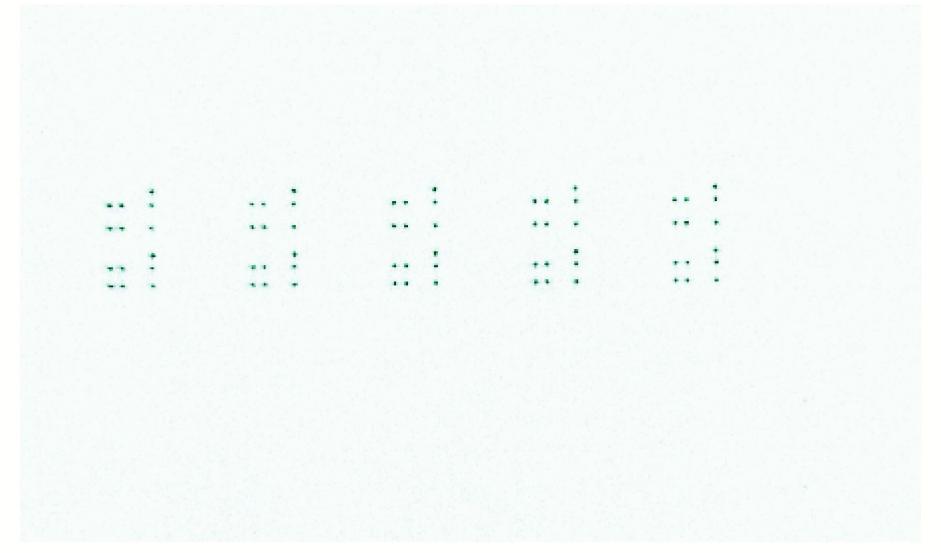
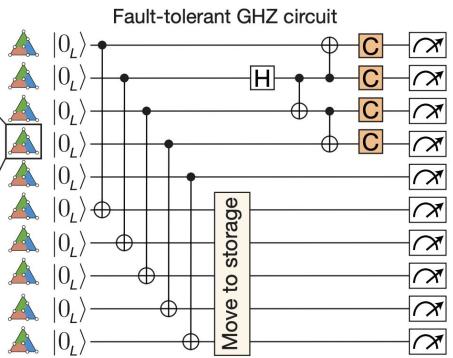
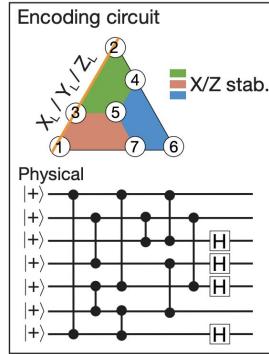
- Hyperfine qubits: $>1s$ coherence,
~99.98% global 1Q, ~99.9% local 1Q
- Rydberg-mediated two-qubit gates ~99.5%



Unique opportunities for error correction with reconfigurable atom arrays

Example: GHZ state with logical qubits

a

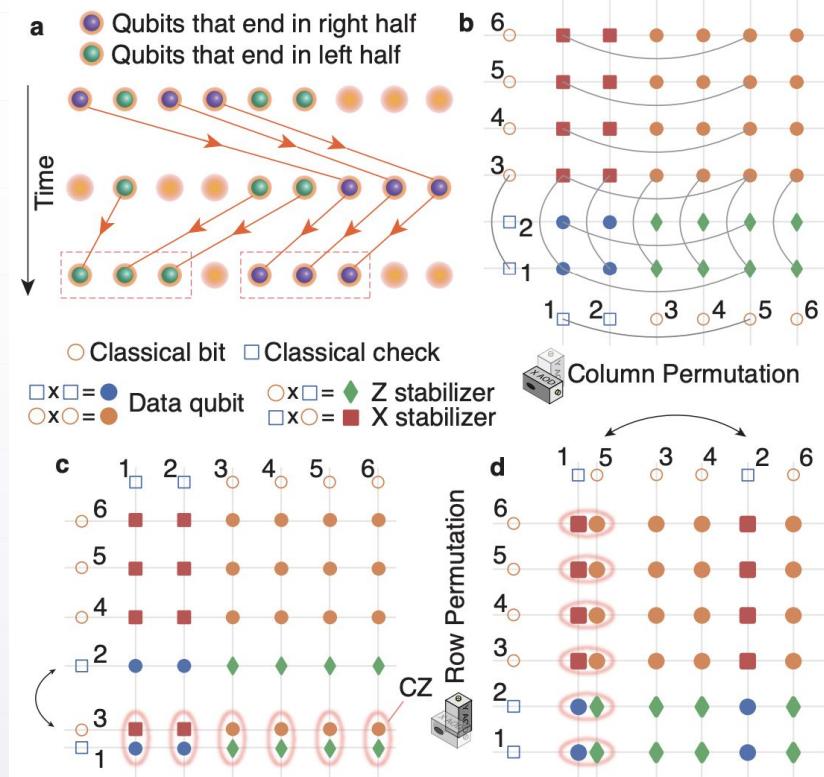


Enabling feature: efficient parallel classical control!



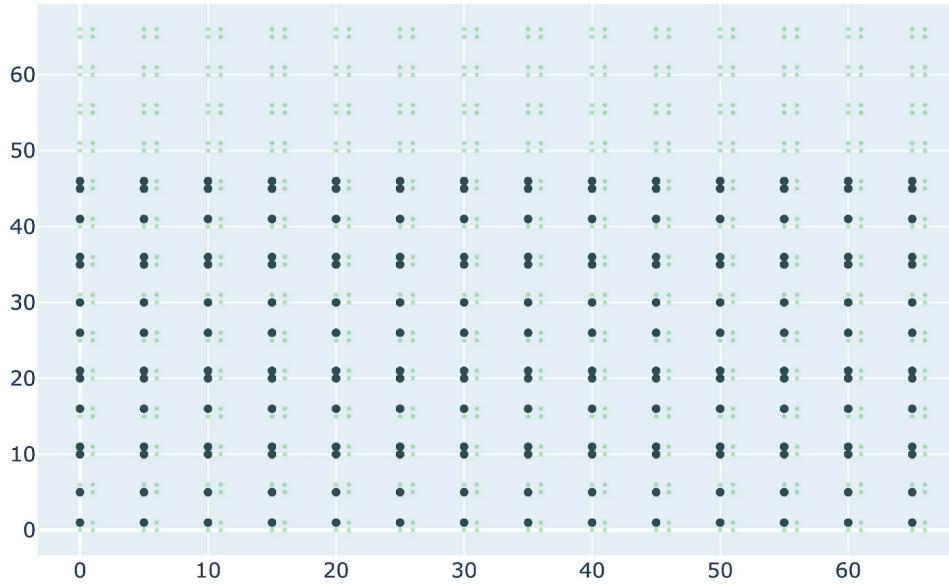
Implementing the Product Structure of HGP codes

- Key insight: product structure of codes matches well with product structure of optical tools (AODs)
- For each check:
 - Permute columns, do CZ gates
 - Permute rows, do CZ gates
- Efficient 1D rearrangement:
 - L qubits in a line
 - At each step:
 - Move to the right half qubits that in the end configuration appear in the right half
 - Repeat within each new half
 - #moves = $O(\log L)$



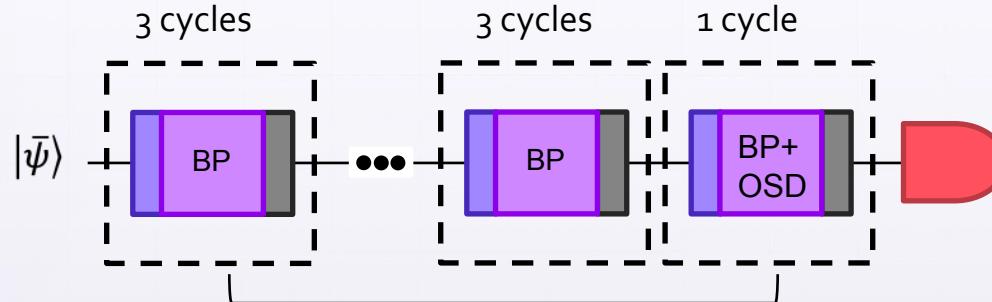
Full Layout for Hypergraph Product Code

- Very generic compilation and with current optical tools
- Movie generated using experimental software and commands!

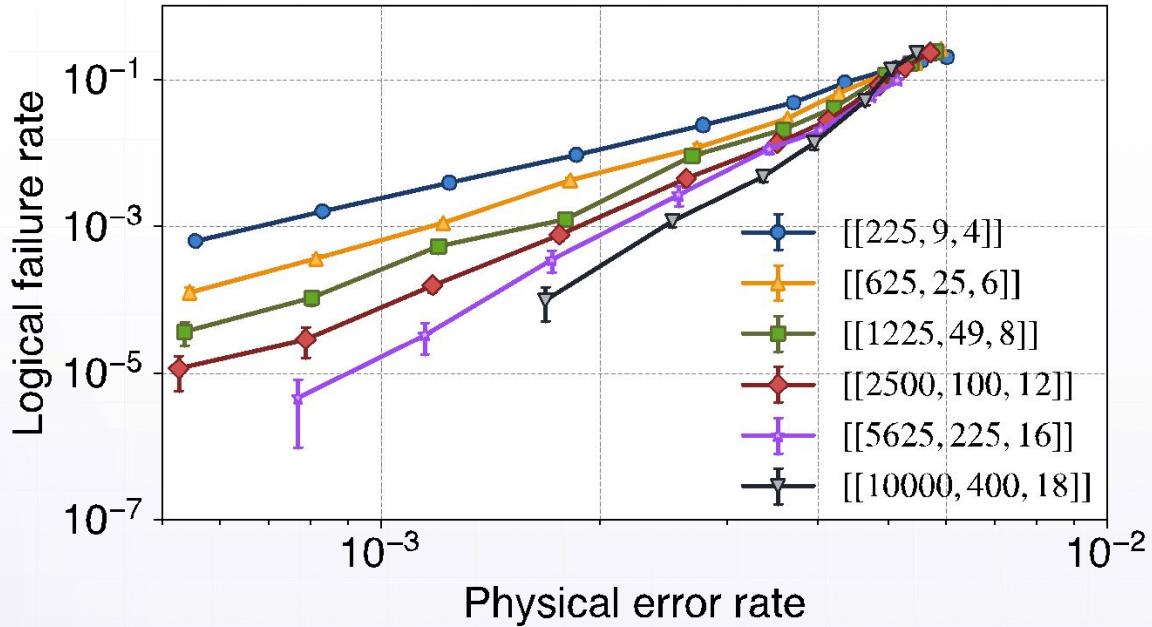


Circuit-Level Performance Evaluation

- Memory simulations: keep logical qubit alive for a long time
- HGP codes satisfy linear syndrome confinement:
 - “Qubits errors cannot grow without triggering more measurement syndromes”
- Confinement + bounded check weight:
 - Single-shot decoding
 - Single-ancilla syndrome extraction is fault-tolerant
 - We prove the existence of a threshold under these considerations!
- BP + BP-OSD decoding on circuit-level detector error model
 - Joint decoding of multiple rounds to improve performance



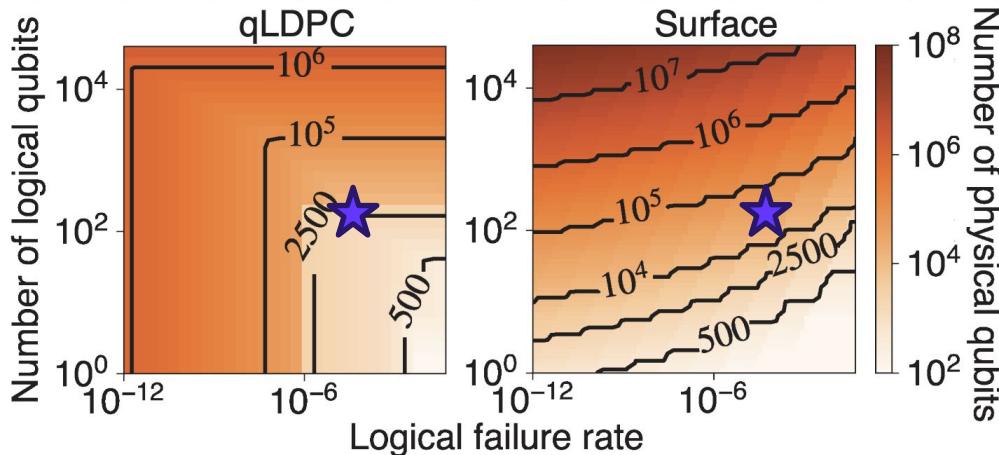
Competitive Memory Circuit-Level Performance



- Error threshold ~ 0.6% under circuit-level depolarizing noise without idling errors
- Long neutral atom coherence times -> adding in realistic idling errors (that scale with code size) produces minimal changes



Competitive Circuit-Level Performance



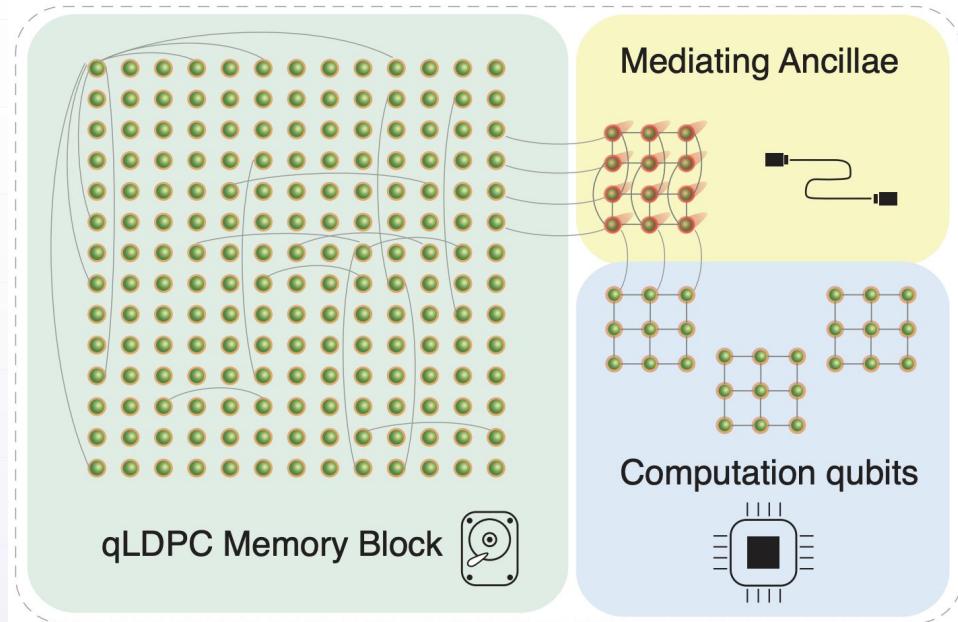
- qLDPC codes start out-performing surface codes at **several hundred physical qubits** and 0.1% error rates
- **Orders of magnitude** savings!
- <100k qubits enough for 1000 logical qubit computation!

Logical qubits	25	80	180	400
Logical failure rate	10^{-3}	10^{-4}	2×10^{-5}	6×10^{-6}
HGP code physical qubits (improvement over surface code)	1235 (1×)	4606 (2.8×)	10760 (4.0×)	19600 (6.9×)
LP code physical qubits (improvement over surface code)	851 (1.4×)	1367 (9.4×)	2670 (16.2×)	



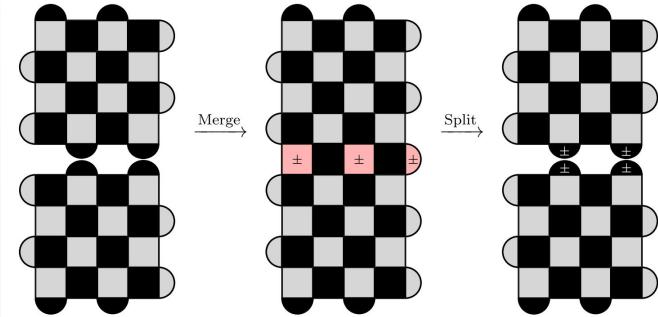
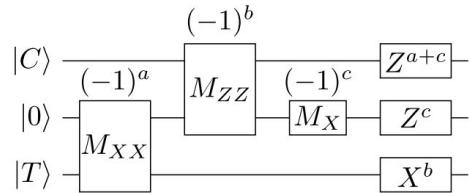
Computation with Logical Qubits

- qLDPC: memory
- Topological codes: processor
- Small enough number of computational qubits in parallel:
 - Can maintain constant overhead
- Our contribution:
 - Qubit efficient teleportation $q\text{LDPC} \leftrightarrow \text{surface}$
 - First numerical simulations of computation with qLDPC codes



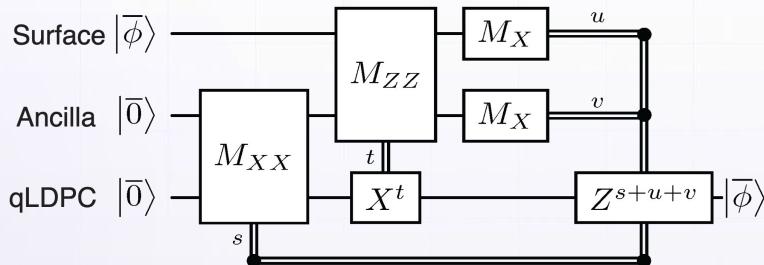
Lattice surgery

- Measurement-based gates: universal
 - Need: single and joint (logical) Clifford measurements
- Trick – for codes of similar boundary:
 - Merge codes
 - Logical measurement becomes stabilizer of new code
 - Measure stabilizer fault-tolerantly
 - Split codes
- Teleportation: also through measurement
 - Boundary of qLDPC and surface code can be very different
 - How to merge to perform joint measurement?

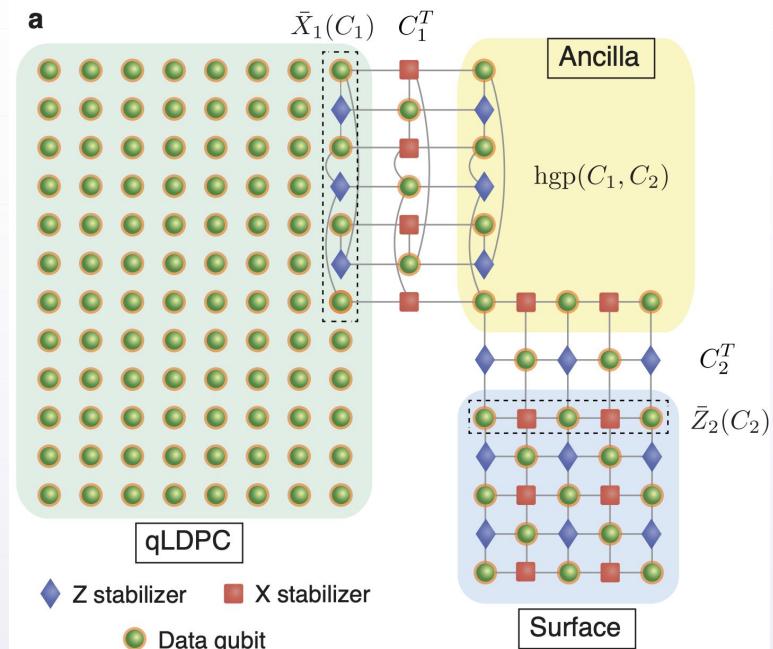


Teleportation through Lattice Surgery

- Mediate teleportation with ancilla block. Treat logicals as classical codes
- Ancilla block: HGP again!
- Logicals of same type:
 - Normal lattice surgery

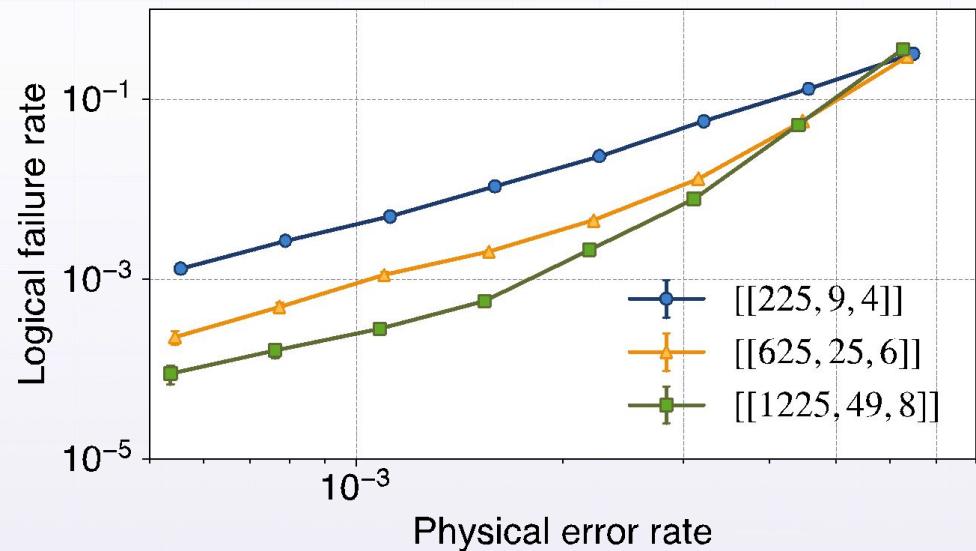
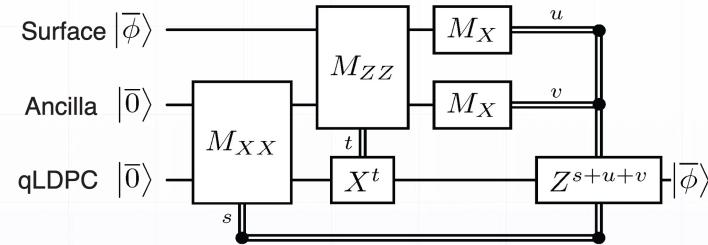


N.B. We prove this procedure is fault-tolerant



Lattice Surgery Performance

- Teleport logical qubits into surface code for computation
- Circuit-level simulations of teleportation process
- **Competitive threshold and error rates maintained!**

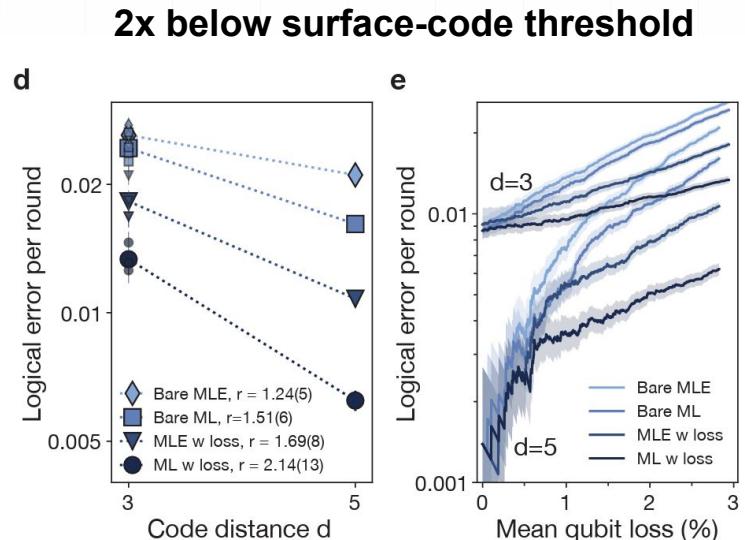


Outlook

Exploring LDPC challenges and opportunities on near-term hardware

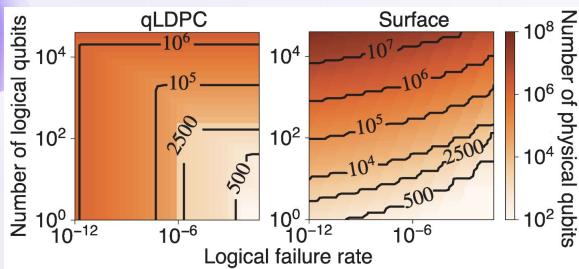
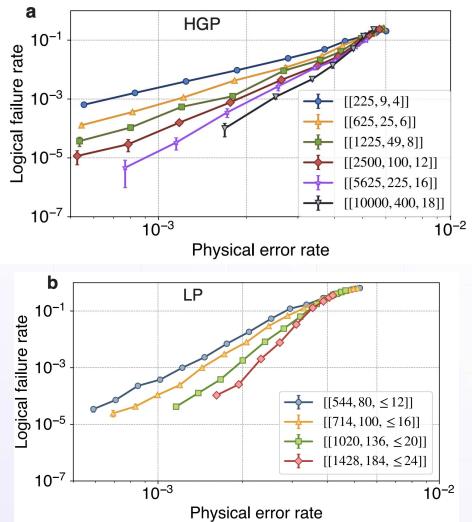
- Space-time trade-offs

- Improving qLDPC computation
- How to incorporate QEM into these
 - QEM largely unexplored in neutral atom
 - Some ideas: QEC naturally projects noise needed?), extrapolation based on posts

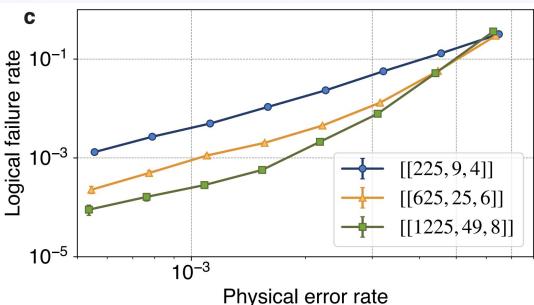
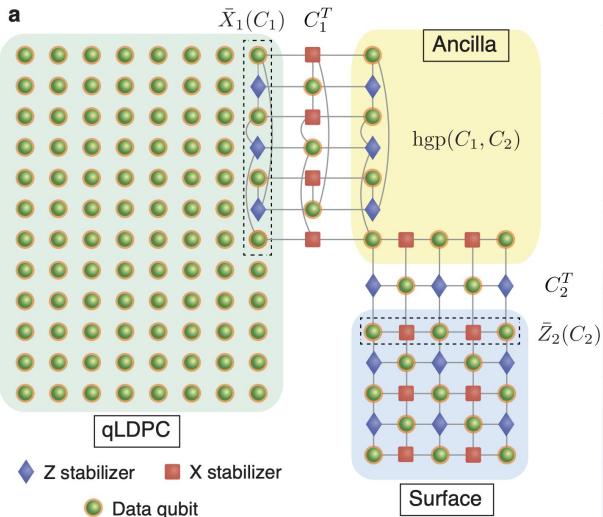


Summary

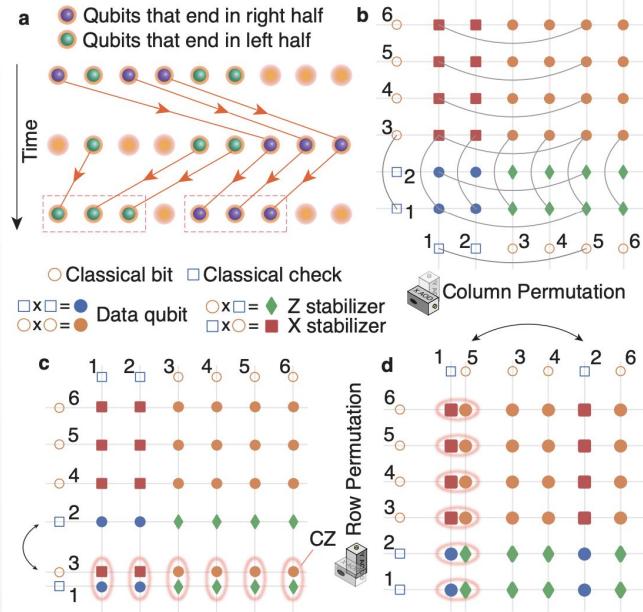
Memory



Gates



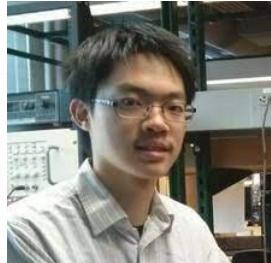
Practical implementation



Collaborators



Qian Xu
(U.Chicago)



Harry Zhou
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Chris Pattison
(Caltech)



Liang Jiang
(U.Chicago)



Mikhail Lukin
(Harvard)



Dolev
Bluvstein
(Harvard)



Nithin
Raveendran
(Arizona)



Jonathan Wurtz
(QuEra)



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(Arizona)



Thanks also to other members of Harvard atom array team.

