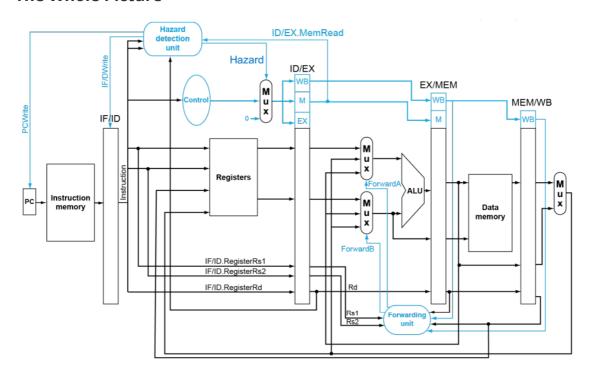
T7 Data Hazard

Different Type of Hazards

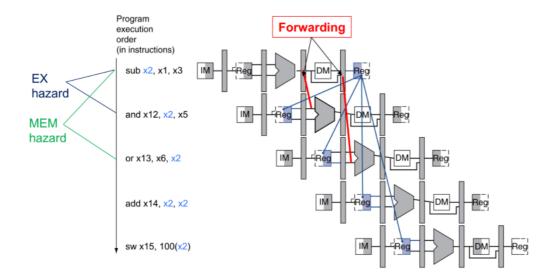
- Data hazards: Need to wait for previous instruction to complete its data read/write
- Control hazards: Decision on control action depends on previous instruction
- Structure hazards: A required resource is busy

The Whole Picture



Categories of Data Hazards

1. EX/MEM Data Hazards



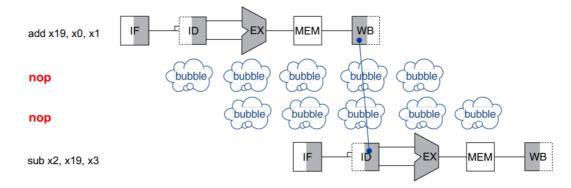
EX hazard:

- (1) Data hazard between two adjacent instructions
- (2) Forwarding path from EX/MEM pipeline register to ALU

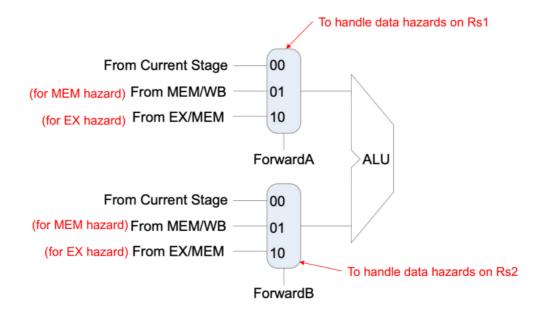
MEM hazard:

- (1) Data hazard between two instructions with one more instruction in between
- (2) Forwarding path from MEM/WB pipeline register to ALU

• Solution 1: Stall



• Solution 2: Forwarding



The control signal ForwardA and ForwardB is generated from **Forwarding unit**. Logic in Forwarding unit:

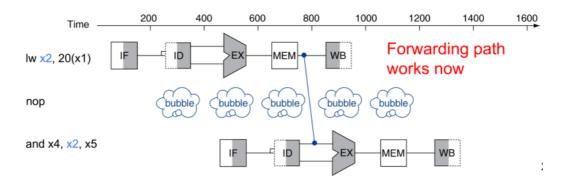
```
# EX hazard
if (EX/MEM.Regwrite and (EX/MEM.RegisterRd ≠ 0)
    and (EX/MEM.RegisterRd == ID/EX.RegisterRs1))
    ForwardA = 10
if (EX/MEM.Regwrite and (EX/MEM.RegisterRd ≠ 0)
    and (EX/MEM.RegisterRd == ID/EX.RegisterRs2))
    ForwardB = 10
# MEM hazard
```

```
if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
    and (MEM/WB.RegisterRd == ID/EX.RegisterRs1)
    and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
        and (EX/MEM.RegisterRd == ID/EX.RegisterRs1))
        ForwardA = 01

if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
        and (MEM/WB.RegisterRd == ID/EX.RegisterRs2)
    and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
        and (EX/MEM.RegisterRd == ID/EX.RegisterRs2)))
        ForwardB = 01
```

2. Load-Use Data Hazards

For load-use data hazards, forwarding cannot fix it completely. Need to stall for one cycle.



We will use hazard detection unit here to detect whether we need to insert a stall.

• Logic in hazard detection unit:

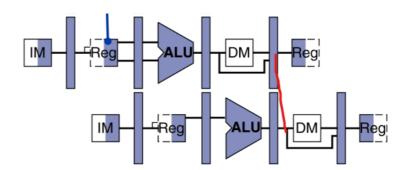
- Logic of stalls
 - Force control signals in ID/EX pipeline register to 0's
 - EX, MEM and WB in following cycles do no-operation (nop) with those 0 control signal
 - Prevent updates of PC and IF/ID registers
 - Instruction in IF/ID is held and decoded again
 - PC is held, so the same instruction is fetched again

3. Other Hazards:

Example:

```
Tw x10, 20(x1)
sw x10, 40(x1)
```

Can be solved by a new forwarding path:



Modification on hardware:

