## 第一课 链表

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### 内容概述

### 1.8道经典链表常考题目

例1-a:链表逆序(easy)

例1-b:链表逆序2(medium)

例2:链表求交点 (easy)

例3:链表求环(medium)

例4:链表划分(medium)

例5:复杂链表的复制(hard)

例6-a:2个排序链表归并(easy)

例6-b:K个排序链表归并(hard)

- 2.详细讲解题目多种解题方法、代码实现
- 3.一些学习与找工作的建议



## 预备知识:链表基础

```
#include <stdio.h>
struct ListNode {
                                               1分钟时间填写代码!
   int val; //存储元素的数据域
   ListNode *next;
};
      //存储下一个节点地址的指针域
int main(){
   ListNode a;
   ListNode b;
                                      1.创建5个节点:
   ListNode C:
   ListNode d;
                                                                  40
   ListNode e;
   a.val = 10;
                                      2.将它们连接在一起:
   b.val = 20;
   c.val = 30;
   d.val = 40;
                                                       →| 30 | •
   e.val = 50:
   a.next = \&b;
                                      3.遍历它们,并打印节点的值:
    c.next = &d;
    d.next = &e;
   ListNode *head = &a;
   while (head) {
                                      head
       printf("%d\n", head->val);
   return 0;
```

## 预备知识:链表基础

next

```
#include <stdio.h>
struct ListNode {
    int val; //存储元素的数据域
                                  val
    ListNode *next;
} ;
       //存储下一个节点地址的指针域
int main(){
    ListNode a;
   ListNode b;
   ListNode C;
   ListNode d:
   ListNode e;
   a.val = 10;
   b.val = 20;
    c.val = 30;
   d.val = 40;
    e.val = 50;
    a.next = &b;
    b.next = &c;
    c.next = &d;
    d.next = &e;
    e.next = NULL;
    ListNode *head = &a;
    while(head) {
        printf("%d\n", head->val);
        head = head->next;
    return 0;
```

```
10
20
30
40
50
请按任意键继续. . .
```



### 例1-a:链表逆序-a

已知链表头节点指针head,将链表逆序。(不可申请额外空间)

#### 选自 LeetCode 206. Reverse Linked List

https://leetcode.com/problems/reverse-linked-list/description/

难度:Easy



### 例1-a:思路

依次遍历链表节点,每遍历一个节点即逆置一个节点

原始链表: 1 2 3 4 5 5 head

新链表:

(NULL)

Û

new\_head

循环1次:

**☆** head

新链表:

1

new\_head

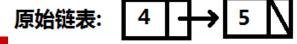
循环2次:

原始链表: 3 4 5 5

**☆** head

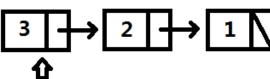
新链表: 2 → 1

**រេ** new\_head 循环3次:



1 head

新链表:



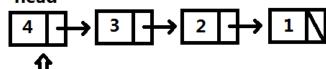
new\_head

循环4次:

原始链表: 5

1 head

新链表:



new\_head

循环5次:

原始链表: NULL)

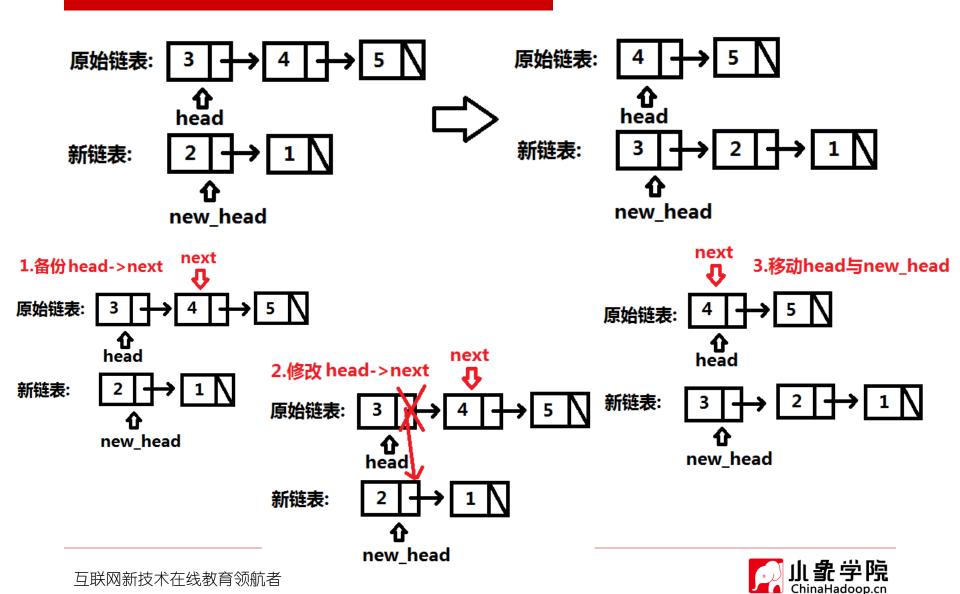
head

新链表:

 $5 \longrightarrow 4 \longrightarrow 3 \longrightarrow 2 \longrightarrow 1$ 

new\_head

## 例1-a:思路



### 例1-a:课堂练习

```
class Solution {
public:
   ListNode* reverseList(ListNode* head) {
       ListNode *new_head = NULL; //指向新链表头节点的指针
        while (head) {
            head = next; //遍历链表
        return new head; //返回新链表头节点
};
```

### 3分钟时间填写代码,有问题随时提出!



### 例1-a:实现

```
class Solution {
public:
    ListNode* reverseList(ListNode* head) {
        ListNode *new_head = NULL; //指向新链表头节点的指针
        while (head) {
              ListNode *next = head->next;
                                          //备份head->next
                                   //更新head->next
             head->next = new_head;
               new_head = head;
                                   //移动new head
            head = next; //遍历链表
        return new_head; //返回新链表头节点
};
```

## 例1-a:测试与leetcode提交结果

```
int main(){
    ListNode a(1);
    ListNode b(2);
    ListNode c(3);
    ListNode d(4);
    ListNode e(5);
    a.next = &b;
    b.next = &c; //将节点简单的链接,进行测试
    c.next = &d; //无需构造复杂的链表操作(插入,删除)
    d.next = &e;
    Solution solve;
    ListNode *head = &a;
    printf("Before reverse:\n");
    while(head){
        printf("%d\n", head->val);
        head = head->next;
    head = solve.reverseList(&a);
    printf("After reverse:\n");
    while (head) {
        printf("%d\n", head->val);
        head = head->next;
    return 0:
```

```
Before reverse:
1
2
3
4
5
After reverse:
5
4
3
2
1
请按任意键继续. . . _
```

#### Reverse Linked List

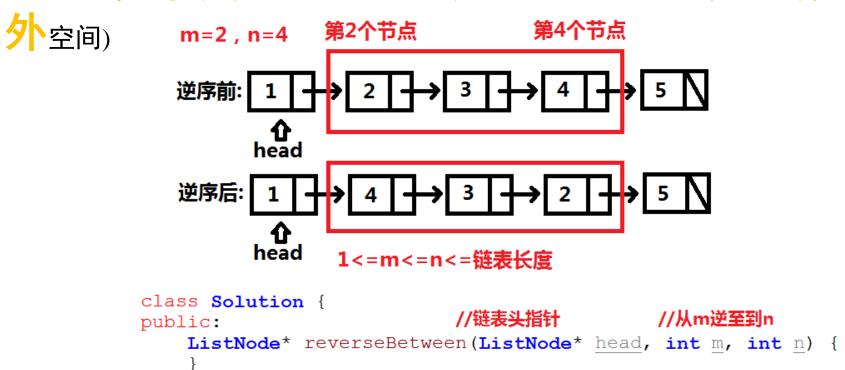
#### Submission Details

27 / 27 test cases passed. Status: Accepted
Runtime: 6 ms Submitted: 0 minutes ago



### 例1-b:链表逆序-b

### 已知链表头节点指针head,将链表从位置m到n逆序。(不申请额



#### 选自 LeetCode 92. Reverse Linked List II

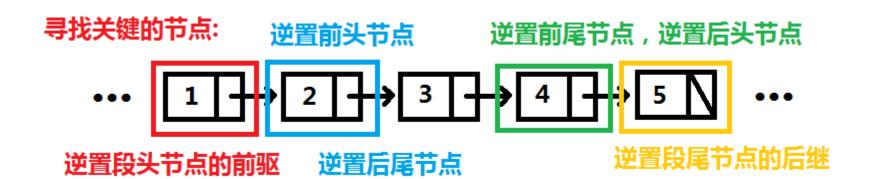
https://leetcode.com/problems/reverse-linked-list-ii/description/

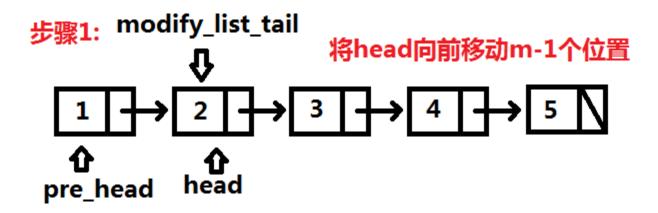
难度:Medium



};

### 例1-b:思路





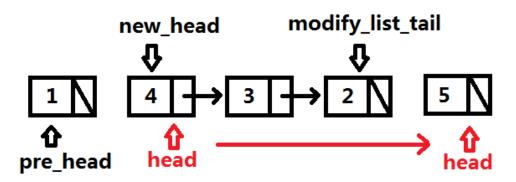
找到开始逆置的节点,记录该节点前驱、该节点



### 例1-b:思路

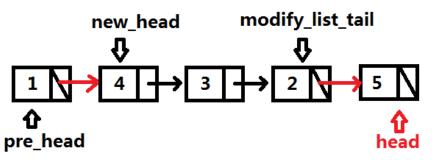
#### 步骤2:

从head开始, 逆置change\_len = n-m+1个节点



#### 步骤3:

将pre\_head与new\_head连接,modify\_list\_tail与head连接



### 思考:

- 1.最终结果应该<mark>返</mark>回哪 个节点?
- 2.如果m=1时,有什么 特殊的?



```
class Solution {
public:
   ListNode* reverseBetween(ListNode* head, int m, int n) {
       int change len = n - m + 1; //计算需要逆置的节点个数
                                                    例1-b:课堂练习
       ListNode *pre head = NULL; //初始化开始逆置的节点的前驱
       ListNode *result = head; //最终转换后的链表头节点, 非特殊情况即为head
       while (head && --m) {
                               //将head向前移动m-1个位置
                   1
          head = head->next;
                           //将modify_list_tail指向当前的head,即逆置后的链表尾
       ListNode *modify list tail = head;
       ListNode *new head = NULL;
                                //逆置change len个节点
       while (head && change len) {
          ListNode *next = head->next;
          head->next = new head;
          new head = head;
                                                 5分钟时间填写代码,
          head = next;
                                                 有问题随时提出!
                              //如果pre_head不空,说明不是从第一
       if (pre head) {
                               个节点开始逆置的m>1
       else{
                   5
       return result;
};
                                                                  ChinaHadoop.cn
```

```
class Solution {
                                                             例1-b:实现
public:
   ListNode* reverseBetween(ListNode* head, int m, int n) {
       int change len = n - m + 1; //计算需要逆置的节点个数
       ListNode *pre head = NULL; //初始化开始逆置的节点的前驱
       ListNode *result = head; //最终转换后的链表头节点,非特殊情况即为head
       while (head && --m) {
                                 //将head向前移动m-1个位置
              pre head = head;
                                 //记录head的前驱
           head = head->next;
                            //将modify_list_tail指向当前的head , 即逆置后的链表尾
       ListNode *modify list tail = head;
       ListNode *new head = NULL;
                                  //逆置change len个节点
       while (head && change len) {
           ListNode *next = head->next;
           head->next = new head;
           new head = head;
           head = next;
                                 //每完成一个节点逆置 , change len--;
             change_len--;
                                 //连接逆置后的链表尾与逆置段的后一个节点
         modify_list_tail->next = head;
                                //如果pre_head不空,说明不是从第一
       if (pre head) {
                                个节点开始逆置的m>1
            pre_head->next = new_head;
                                 //将逆置链表开始的节点前驱与逆置后的头节点链接
       else{
                                 //如果pre head为空,说明m==1从第一个节点开始逆置
            result = new head;
                                 结果即为逆置后的头节点
       return result;
```

ChinaHadoop.cn

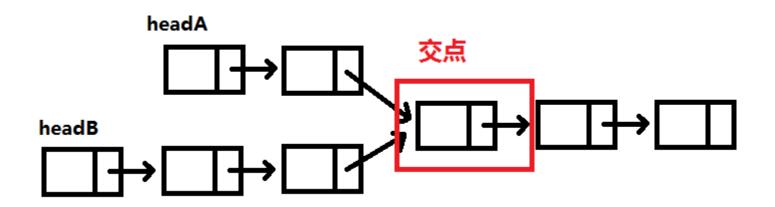
};

## 例1-b:测试与leetcode提交结果

```
int main(){
    ListNode a(1);
    ListNode b(2);
    ListNode c(3);
    ListNode d(4);
    ListNode e(5);
    a.next = \&b;
    b.next = &c;
    c.next = &d;
    d.next = &e;
    Solution solve;
    ListNode *head = solve.reverseBetween(&a, 2, 4);
    while (head) {
         printf("%d\n", head->val);
                                           Reverse Linked List II
         head = head->next;
                                           Submission Details
    return 0;
                                                                   Status: Accepted
                                              44 / 44 test cases passed.
                                              Runtime: 3 ms
                                                               Submitted: 0 minutes ago
```

### 例2:求两个链表的交点

已知链表A的头节点指针headA,链表B的头节点指针headB,两个链表相交,求两链表交点对应的节点



#### 选自 LeetCode 160. Intersection of Two Linked Lists

https://leetcode.com/problems/intersection-of-two-linked-lists/description/

难度:Easy



### 例2:题目要求

```
struct ListNode {
    int val; //数据域
    ListNode *next; //指针域
    ListNode(int x): val(x), next(NULL) {
    } //构造函数
};

class Solution {
    public: 链表A头节点指针 链表B头节点指针
    ListNode *getIntersectionNode(ListNode *headA, ListNode *headB) {
    } 求两个链表的交点
};
```

- 1.如果两个链表没有交点,则返回NULL
- 2.在求交点的过程中, 不可以破坏链表的结构或者修改链表的数据域
- 3.可以确保传入的链表A与链表B没有任何环
- 4\*.实现算法尽可能使时间复杂度O(n),空间复杂度O(1)

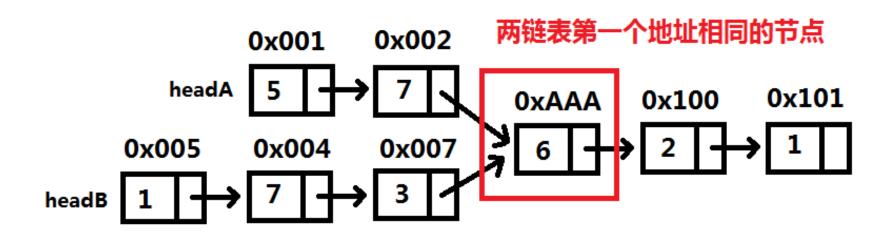


## 例2:方法1的必备知识(set的使用)

```
//STL set的使用
#include <set>
                                                                        //O(nlogn)或O(n)的方法判断两数组是否有相同元素
int main(){
                     std::set<int> test set; //STL set
                    const int A_LEN = 7; //测试数组A与B的长度
                                                                                                                                                                                                                                           b[3] = 3 in array A.
                    const int B LEN = 8;
                                                                                                                                                                                                                                           b[4] = 1 in array A.
                    int a[A LEN] = \{5, 1, 4, 8, 10, 1, 3\};
                                                                                                                                                                                                                                                                    = 1 in array A.
                    int b[B LEN] = \{2, 7, 6, 3, 1, 6, 0, 1\};
                    for (int i = 0; i < A LEN; i++) {
                                       test set.insert(a[i]); //将数组a的元素插入set
                    for (int i = 0; i < B LEN; i++) { // * hotelength{ hotelength} hotelength} // *hotelength{ hotelength} hotelength{ hotelengt
                                       if (test set.find(b[i]) != test set.end()) {
                                                          printf("b[%d] = %d in array A.\n", i, b[i]);
                   return 0:
```

## 例2:思路1,使用set求交集

- 1.遍历链表A,将A中节点对应的指针(地址),插入set
- 2.遍历链表B,将B中节点对应的指针(地址),在set中查找
- ,发现在set中的第一个节点地址,即是两个链表的<mark>交点</mark>。



## 例2:方法1,课堂练习

### 3分钟时间填写代码,有问题随时提出!

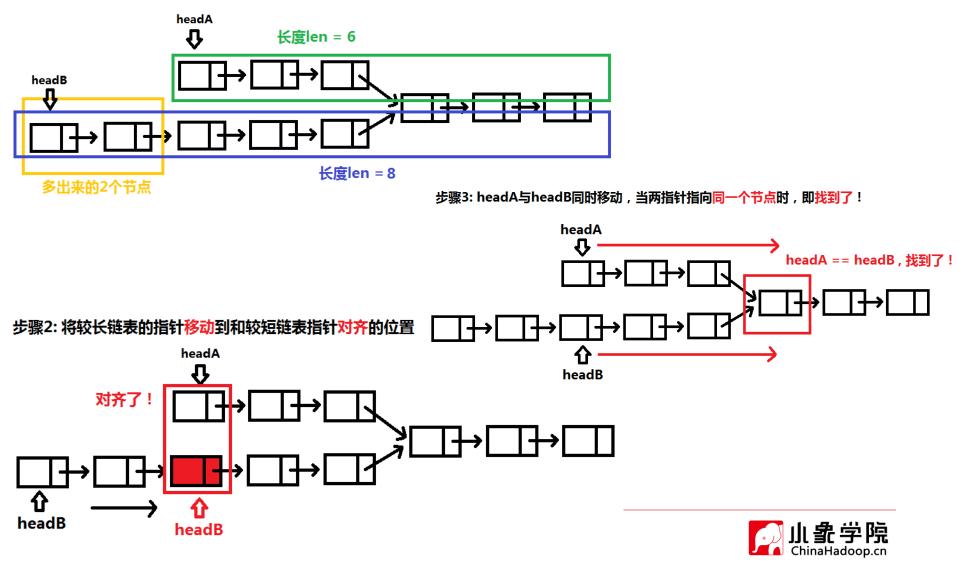


### 例2:方法1实现

```
class Solution {
public:
    ListNode *getIntersectionNode(ListNode *headA, ListNode *headB) {
       std::set<ListNode*> node_set;//设置查找集合node_set
              node set.insert(headA);
                                       //将链表A中的节点插入node_set
            headA = headA->next;
                                //遍历链表A
        while (headB) {
                 node_set.find(headB) != node_set.end()
            i f
                return headB;
                             //当在headB中找到第一个出现在node_set中的节点时
                headB = headB->next;
                                            //遍历链表B
        return NULL;
} ;
```

## 例2:思路2,空间复杂度O(1)

步骤1: 计算headA链表长度、计算headB链表长度, 较长的链表多出的长度



### 例2:方法2,课堂练习

};

```
ListNode *forward long list(int long len,
int get list length(ListNode *head) {
                                                    int short len, ListNode *head) {
   int len = 0;
   while (head) { //遍历链表,计算链表长度
                                          int delta = long len - short len;
                                         while(head && delta) {
                                                                 //将指针向前移动至多出
                                             head = head->next;
                                                                   节点个数后面的位置
       head = head->next;
   return len;
                                          return head;
class Solution {
public:
   ListNode *getIntersectionNode(ListNode *headA, ListNode *headB) {
       int list A len = get list length(headA);
       int list_B_len = get_list_length(headB); //求A、B两个链表长度
       if (list A len > list B len) {
          headA = forward long list(list A len, list B len, headA);
                  //如果链表A长,移动headA到对应位置
       else{
          headB = forward long list(list B len, list A len, headB);
                             //如果链表B长,移动headB到对应位置
       while (headA && headB) {
          if
                   3
              return headA;
                                        钟时间填写代码,有问题
          headA = headA->next;
          headB = headB->next;
       return NULL;
```

### 例2:方法2,实现

};

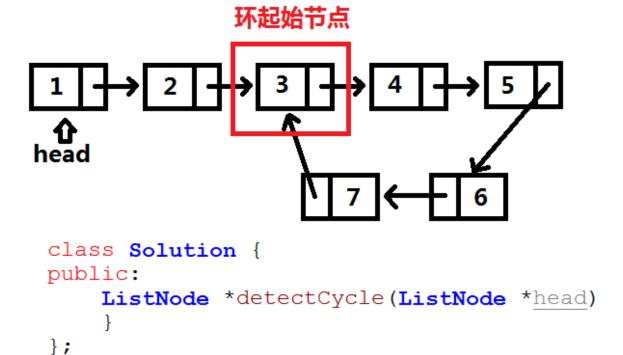
```
ListNode *forward long list(int long len,
int get list length(ListNode *head) {
                                                     int short len, ListNode *head) {
    int len = 0;
    while (head) { //遍历链表,计算链表长度
                                          int delta = long len - short len;
                                          while(head && delta) {
                                                                 //将指针向前移动至多出
            len++:
                                              head = head->next;
                                                                   节点个数后面的位置
        head = head->next;
                                                   delta--:
    return len;
                                          return head;
class Solution {
public:
   ListNode *getIntersectionNode(ListNode *headA, ListNode *headB) {
       int list A len = get list length(headA);
       int list_B_len = get_list_length(headB); //求A、B两个链表长度
       if (list A len > list B len) {
           headA = forward long list(list A len, list B len, headA);
                   //如果链表A长,移动headA到对应位置
       else{
           headB = forward long list(list B len, list A len, headB);
                               //如果链表B长,移动headB到对应位置
       while (headA && headB) {
           if
              headA == headB
                               //当两指针指向了同一个节点时,说明找到了!
               return headA;
           headA = headA->next;
           headB = headB->next;
       return NULL;
```

## 例2:测试与leetcode提交结果

```
int main(){
                       //将节点简单的链接一下,即可进行测试
    ListNode a1(1);
    ListNode a2(2);
                       //无需构造链表插入、删除等复杂方法
    ListNode b1(3);
    ListNode b2(4);
    ListNode b3(5);
    ListNode c1(6);
                                headA
    ListNode c2(7);
    ListNode c3(8);
    a1.next = &a2;
    a2.next = &c1;
                                                      c1 | -
                        headB
    c1.next = &c2;
    c2.next = &c3;
    b1.next = &b2;
    b2.next = &b3;
    b3.next = &c1;
    Solution solve;
    ListNode *result = solve.getIntersectionNode(&a1, &b1);
    printf("%d\n", result->val);
    return 0:
                                                 Intersection of Two Linked Lists
                                                 Submission Details
                                                                      Status: Accepted
                                                    42 / 42 test cases passed.
                                                                   Submitted: 0 minutes ago
                                                    Runtime: 63 ms
```

### 例3:链表求环

已知链表中可能存在环,若有环返回环起始节点,否则返回NULL。



#### 选自 LeetCode 141. Linked List Cycle 142. Linked List Cycle II

https://leetcode.com/problems/linked-list-cycle-ii/description/

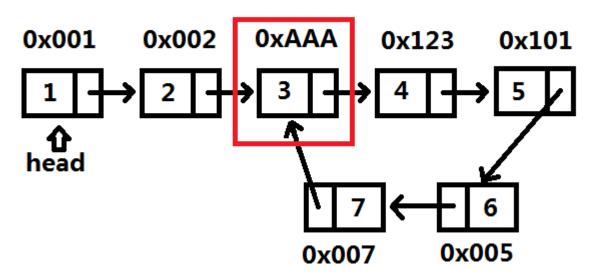
难度:Medium



### 例3:思路1,使用set求环起始节点

- 1.遍历链表,将链表中节点对应的指针(地址),插入set
- 2.在遍历时插入节点前,需要在set中查找,第一个在set中发现的节点地址,即是链表环的起点。

#### 该节点是遍历时,第一个在set中已出现的节点,即环的开始



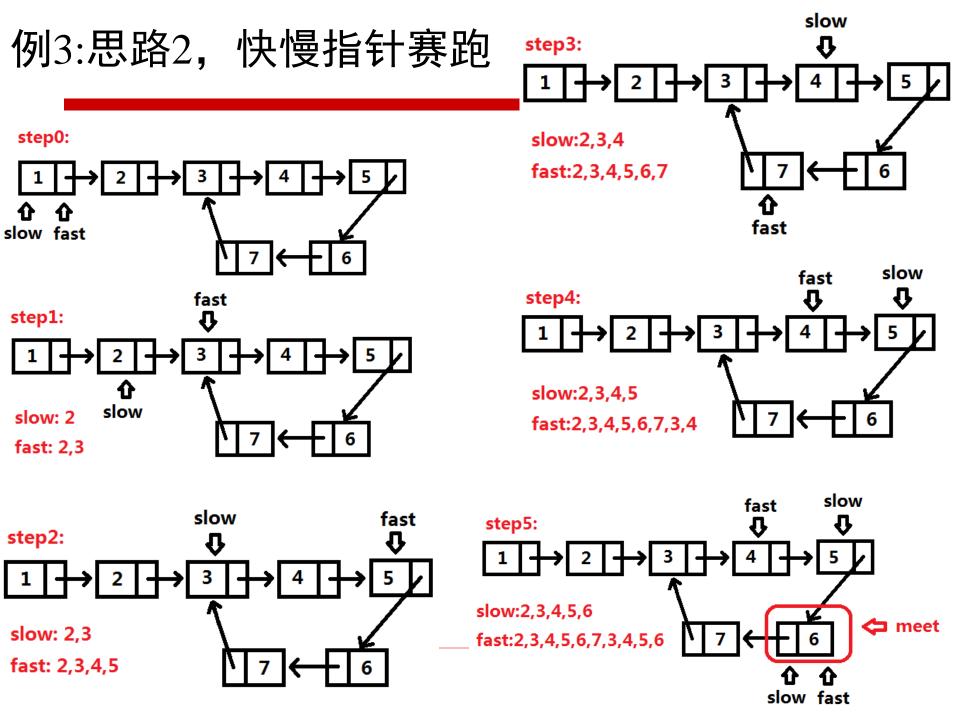


## 例3:方法1,课堂练习

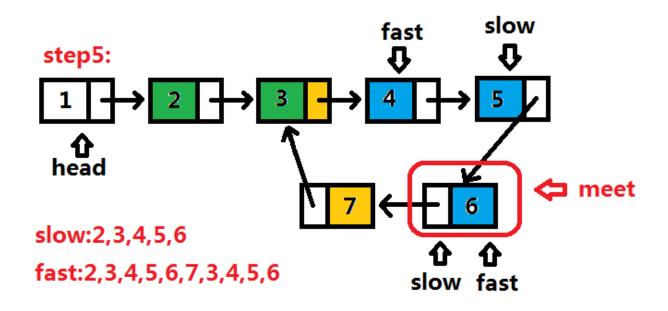
```
class Solution {
public:
   ListNode *detectCycle(ListNode *head) {
       std::set<ListNode *> node_set; //设置node_set
       while (head) { //遍历链表
          if
             return head;
                         //返回环的第一个节点
          head = head->next;
};
         3分钟时间填写代码,有问题
          随时提出!
```

### 例3:方法1实现

```
class Solution {
public:
   ListNode *detectCycle(ListNode *head) {
       std::set<ListNode *> node_set; //设置node_set
       while (head) { //遍历链表 //如果在node_set已经出现了
                node_set.find(head) != node_set.end()
           if
               return head;
                            //返回环的第一个节点
             node set.insert(head);
                                  //将节点插入node_set
           head = head->next;
          return NULL;
                           //没有遇到环,则返回NULL
};
```



### 例3:思路2,快慢指针



设:

1)绿色 2,3 段 为 a

2)蓝色 4,5,6 段为b

3)橙色 7,3 段为c

$$slow = a + b$$

$$fast = a + b + c + b$$

由于 fast 走的路程 是slow的两倍

故:

$$2*(a + b) = a + b + c + b$$

结论: 从head 与 meet 出发, 两指针速度一样, 相遇时即为环的起点



```
class Solution {
public:
                                       例3:方法2,课堂练习
   ListNode *detectCycle(ListNode *head)
       ListNode *fast = head; //快慢指针
       ListNode *slow = head;
       ListNode *meet = NULL; //相遇的节点
       while(fast){
          slow = slow->next; //slow与fast先各走一步 fast = fast->next;
          if (!fast) {
                                         分钟时间填写代码, 有问题
           if (fast == slow) {
                       3
                                       随时提出!
              break;
       if (meet == NULL) {
       while (head && meet) {
           if
              return head;
           head = head->next; //head与meet每次走1步
           meet = meet->next;
       return NULL;
};
```

```
class Solution {
                                            例3:方法2,实现
public:
   ListNode *detectCycle(ListNode *head) {
       ListNode *fast = head; //快慢指针
       ListNode *slow = head;
       ListNode *meet = NULL; //相遇的节点
       while(fast){
           slow = slow->next; //slow与fast先各走一步
           fast = fast->next;
           if (!fast) {
                                    //如果fast遇到链表尾,则返回NULL
                   return NULL;
              fast = fast->next;
                                  //fast再走1步
           if (fast == slow) {
                                     //fast与slow相遇 , 记录相遇位置
                 meet = fast:
               break;
        if (meet == NULL) {
                                 //如果没有相遇,则证明无环
              return NULL;
        while(head && meet) {
                             当head与meet相遇 , 说明遇到环的起始位置
            if
                head == meet
               return head;
            head = head->next; //head与meet每次走1步
            meet = meet->next;
        return NULL;
```

} **;** 

## 例3:测试与leetcode提交结果

```
int main(){
    ListNode a(1);
                                    Linked List Cycle II
    ListNode b(2);
    ListNode c(3);
                                    Submission Details
    ListNode d(4);
    ListNode e(5);
    ListNode f(6);
                                       16 / 16 test cases passed.
                                                                  Status: Accepted
    ListNode q(7);
    a.next = &b;
                                       Runtime: 9 ms
                                                              Submitted: 0 minutes ago
    b.next = &c;
    c.next = &d;
    d.next = &e;
    e.next = &f;
    f.next = &q;
    q.next = &c;
    Solution solve;
    ListNode *node = solve.detectCycle(&a);
    if (node) {
        printf("%d\n", node->val);
    else{
        printf("NULL\n");
    return 0:
}
```

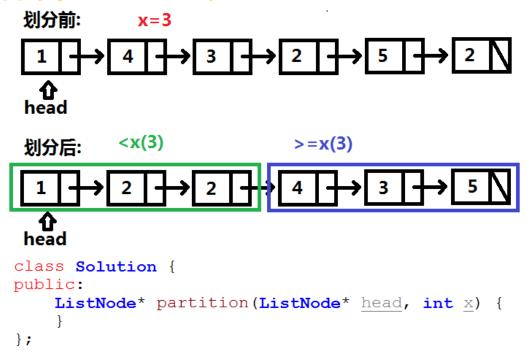
## 课间休息10分钟

# 课前小问题解答



## 例4:链表划分

已知链表头指针head与数值x,将所有小于x的节点放在大于或等于x的节点前,且保持这些节点的原来的相对位置。



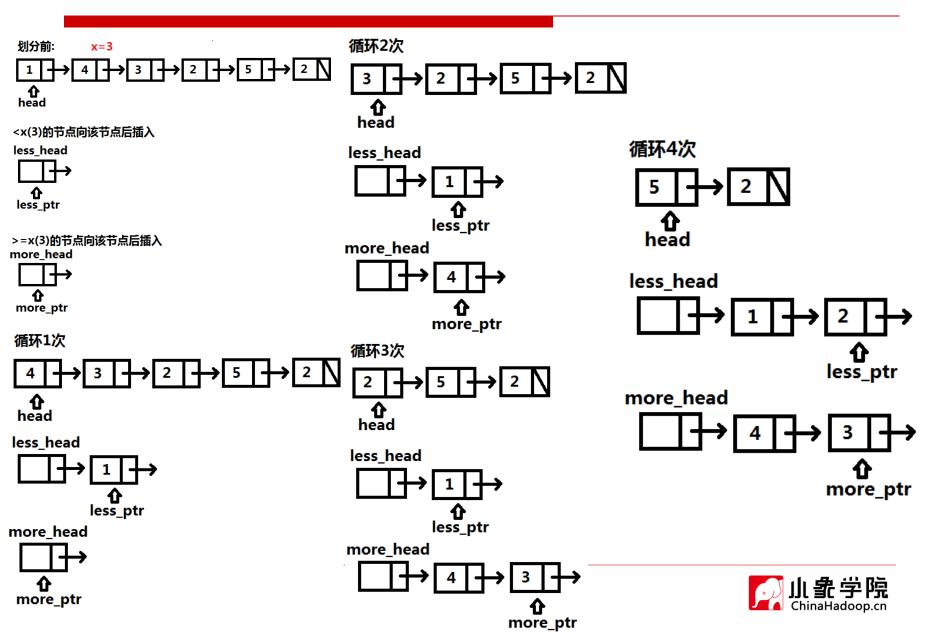
### 选自 LeetCode 86. Partition List

https://leetcode.com/problems/partition-list/description/

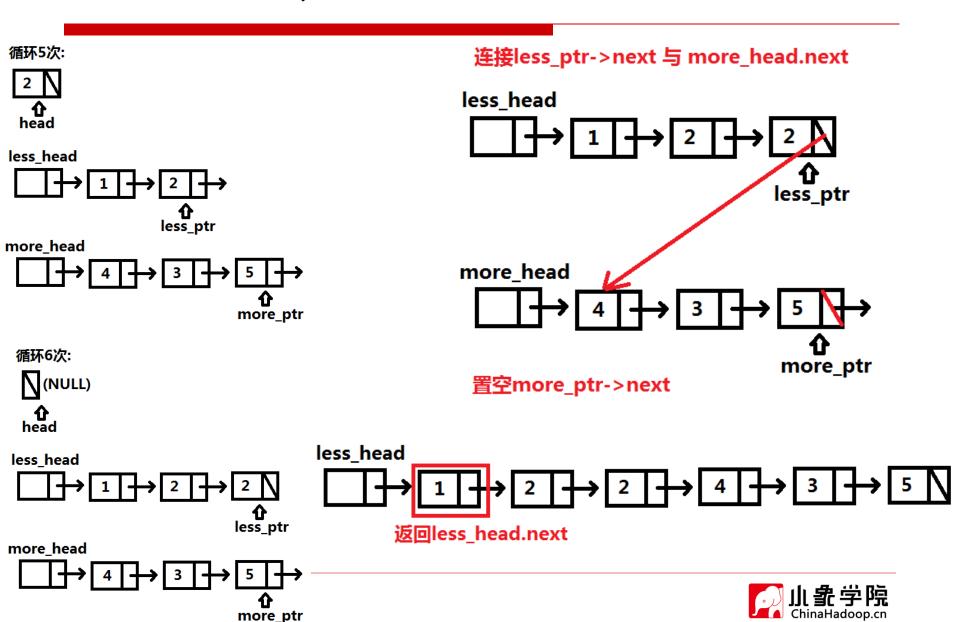
难度:Medium



## 例4:思路,巧用临时头节点



## 例4:思路,巧用临时头节点



## 例4:课堂练习

```
class Solution {
public:
   ListNode* partition(ListNode* head, int x) {
       ListNode less head(0);
                           //设置两个临时的头节点
       ListNode more head(0);
      ListNode *less ptr = &less head;
      ListNode *more_ptr = &more_head; //对应指针指向这两个头节点
       while (head) {
                            //如果节点值小于x , 则将该节点插入less_ptr后
          if (head->val < x) {
              less ptr->next = head;
                 //否则将该节点插入more_ptr后
          else
                                          3分钟时间填写代码,有
                                           问题随时提出!
              more ptr = head;
          head = head->next; //遍历链表
       more ptr->next = NULL;
                           //将more_ptr即链表尾节点next置空
       return less head.next;
              //less_head的next节点即为新链表头节点,返回
};
```

## 例4:实现

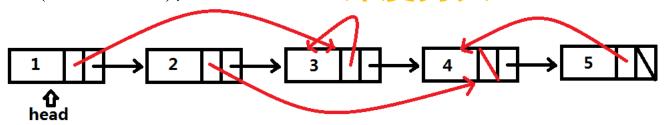
```
class Solution {
public:
   ListNode* partition(ListNode* head, int x) {
       ListNode less head(0);
                             //设置两个临时的头节点
       ListNode more head(0);
       ListNode *less ptr = &less head;
       ListNode *more_ptr = &more head; //对应指针指向这两个头节点
       while (head) {
                              //如果节点值小于x , 则将该节点插入less_ptr后
           if (head->val < x) {
               less ptr->next = head;
                                    //链接完成后,less_ptr向后移动,指向head
                 less ptr = head;
                   //否则将该节点插入more_ptr后
           else
               more ptr->next = head;
               more ptr = head;
           head = head->next; //遍历链表
                                        //将less链表尾,与more链表头相连
         less_ptr->next = more_head.next;
       more ptr->next = NULL;
                             //将more_ptr即链表尾节点next置空
       return less head.next;
               //less_head的next节点即为新链表头节点,返回
};
```

## 例4:测试与leetcode提交结果

```
int main(){
    ListNode a(1);
                             Partition List
    ListNode b(4);
                             Submission Details
    ListNode c(3);
    ListNode d(2);
    ListNode e(5);
                                                             Status: Accepted
                                166 / 166 test cases passed.
    ListNode f(2);
                               Runtime: 6 ms
                                                    Submitted: 1 hour, 9 minutes ago
    a.next = \&b;
    b.next = &c;
    c.next = &d;
    d.next = &e;
    e.next = &f;
    Solution solve;
    ListNode *head = solve.partition(&a, 3);
    while (head) {
         printf("%d\n", head->val);
         head = head->next;
    return 0;
```

## 例5:复杂的链表的深度拷贝

已知一个复杂的链表,节点中有一个指向本链表任意某个节点的随机指针(也可以为空),求这个链表的深度拷贝。



## 选自 LeetCode 138. Copy List with Random Pointer

https://leetcode.com/problems/copy-list-with-random-pointer/description/

难度:Hard

1分钟理解 题目,有问 题随时提 出!

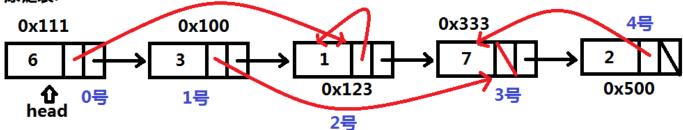


# 例5:必备知识(STL Map的使用)

```
a.random id = 3
#include <stdio.h>
                                                          b.random id = 1
#include <map> //STL map头文件
                                                          c.random id = 3
struct RandomListNode {
   int label;
   RandomListNode *next, *random;
   RandomListNode(int x) : label(x), next(NULL), random(NULL) {}
};
int main(){
   std::map<RandomListNode *, int> node map;
   RandomListNode a (5);
                          //设置一个节点map, key为节点地址, value为整型
   RandomListNode b(3);
   RandomListNode C(6);
   a.next = \&b;
   b.next = &c;
                           id = 1
                                            id = 2
    a.random = &c;
   b.random = &a:
   c.random = &c;
   node map[&a] = 1;
   node map[&b] = 2;
   node map[&c] = 3;
   printf("a.random id = %d\n", node map[a.random]);
   printf("b.random id = %d\n", node map[b.random]);
   printf("c.random id = %d\n", node map[c.random]);
   return 0:
```

## 例5:思路:节点地址与节点序号对应

### 原链表:



#### 新链表:



### Map1:

### 原链表节点地址 -> 节点位置(第几个节点)

0x111 -> 0	random 0x123指向 2号节点
0x100 -> 1	random 0x333指向 3号节点
0x123 -> 2	random 0x123指向 2号节点
0x333 -> 3	random NULL 指向 NULL
0x500 -> 4	random 0x333指向 3号节点

### Map2:

节点位置(第几个节点) -> 新链表节点地址

0 -> 0x200	random = 2号节点地址(0x202)
1 -> 0x201	random = 3号节点地址(0x210)
2 -> 0x202	random = 2号节点地址(0x202)
3 -> 0x210	random = NULL
4 -> 0x211	random = 3号节点地址(0x210)

## 例5:课堂练习

```
class Solution {
public:
   RandomListNode *copyRandomList(RandomListNode *head) {
       std::map<RandomListNode *, int> node_map; //地址到节点位置的map
       std::vector<RandomListNode *> node vec;
                                 //使用vector根据存储节点位置访问地址
       RandomListNode *ptr = head;
       int i = 0;
      while (ptr) { //将新链表节点push入node_vec,生成了新链表节点位置到地址的map
          node vec.push back(new RandomListNode(ptr->label));
                                                 3分钟时间填写代码,有
                                                 问题随时提出!
              //i记录节点位置
       node vec.push back(0);
       ptr = head;
                 //再次遍历原始列表 连接新链表的next指针、random指针
       i = 0;
       while(ptr){
           if (ptr->random) { //当random指针不空时
               int id = node map[ptr->random]; //根据node_map确认
                                                  原链表random指针
                                                  指向的 位置即id
          ptr = ptr->next;
          i++;
       return node vec[0];
} ;
```

## 例5:实现

};

```
class Solution {
public:
   RandomListNode *copyRandomList(RandomListNode *head) {
       std::map<RandomListNode *, int> node_map; //地址到节点位置的map
       std::vector<RandomListNode *> node vec;
       RandomListNode *ptr = head;
                                   //使用vector根据存储节点位置访问地址
       int i = 0;
       while (ptr) { //将新链表节点push入node_vec,生成了新链表节点位置到地址的map
           node vec.push back(new RandomListNode(ptr->label));
             node_map[ptr] = i;
                               //记录原始链表地址至节点位置的node_map
                           //遍历原始列表
           i++;
               //记录节点位置
       node vec.push back(0);
       ptr = head;
                  //再次遍历原始列表 连接新链表的next指针、random指针
       i = 0;
       while (ptr) {
                                           //连接新链表next指针
            node_vec[i]->next = node_vec[i+1];
           if (ptr->random) { //当random指针不空时
                int id = node map[ptr->random]; //根据node_map确认
                                                     原链表random指针
                  node vec[i]->random = node vec[id];
                                                     指向的 位置即id
                           //连接新链表random指针
           ptr = ptr->next;
           i++;
       return node vec[0];
```



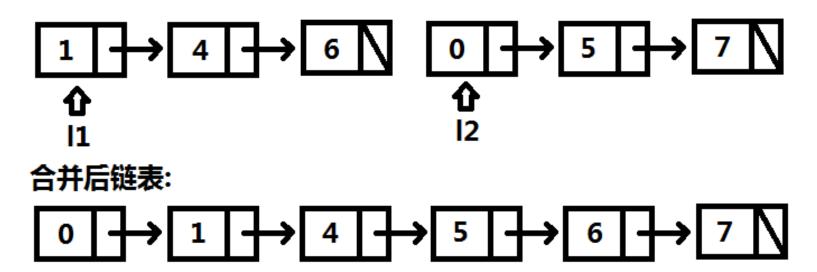
## 例5:测试与leetcode提交结果

return 0:

```
int main(){
                                     Copy List with Random Pointer
    RandomListNode a(1);
                                     Submission Details
    RandomListNode b(2);
    RandomListNode c(3);
    RandomListNode d(4);
                                                                 Status: Accepted
                                        12 / 12 test cases passed.
    RandomListNode e(5);
                                        Runtime: 66 ms
                                                              Submitted: 0 minutes ago
    a.next = \&b;
    b.next = &c;
    c.next = &d;
    d.next = &e;
    a.random = &c;
                         a(1)
    b.random = &d;
    c.random = &c;
    e.random = &d;
    Solution solve;
                                                             label = 1 rand = 3
    RandomListNode *head = solve.copyRandomList(&a);
                                                             label = 2 rand = 4
    while (head) {
                                                             label = 3 rand = 3
        printf("label = %d ", head->label);
                                                             label = 4 rand = NULL
        if (head->random) {
                                                             label = 5 rand = 4
             printf("rand = %d\n", head->random->label);
        else{
             printf("rand = NULL\n");
        head = head->next;
```

# 例6-a:排序链表的合并(2个)

已知两个已排序链表头节点指针11与12、将这两个链表合并,合并后仍为有序的,返回合并后的头节点。



## 选自 LeetCode 21. Merge Two Sorted Lists

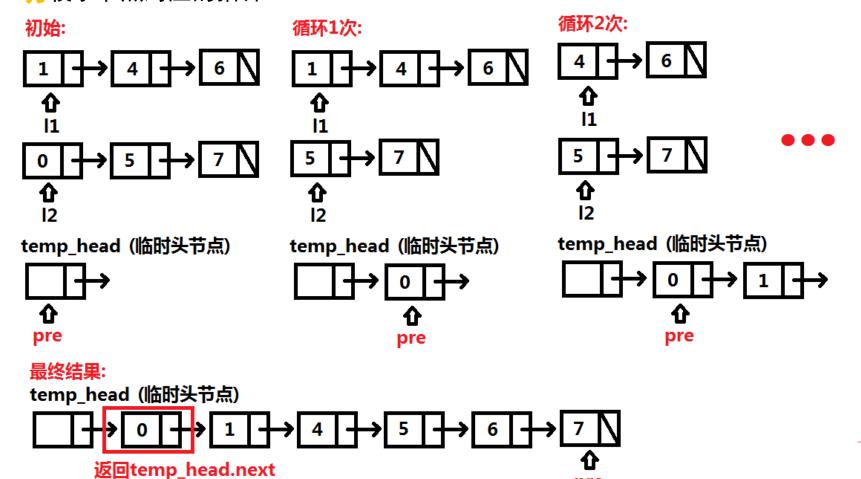
https://leetcode.com/problems/merge-two-sorted-lists/description/

难度:**Easy** 



## 例6-a:思路

比较11和12指向的节点,将较小的节点插入到**pre指针**后,并**向前移**动较小节点对应的指针。



pre

## 例6-a:课堂练习

```
class Solution {
public:
    ListNode* mergeTwoLists(ListNode* 11, ListNode* 12) {
        ListNode temp head(0);
                                     //设置临时头节点temp_head
        ListNode *pre = &temp head;
                                     //使用pre指针指向temp_head
        while (11 && 12) {
            if (11->val < 12->val) {
                                    //如果I1对应的节点小于I2对应的节点
                11 = 11 - \text{next};
            else{
                12 = 12 - \text{next};
                     //如果I1有剩余
        if (11) {
                     //如果I2有剩余
        if (12) {
                     5
        return temp head.next;
```

};

## 5分钟时间填写代码, 有 问题随时提出!

//11与12同时不空时,对它们进行比较



## 例6-a:实现

};

```
class Solution {
public:
   ListNode* mergeTwoLists(ListNode* 11, ListNode* 12) {
        ListNode temp head(0);
                                    //设置临时头节点temp_head
        ListNode *pre = &temp head;
                                    //使用pre指针指向temp_head
        while (11 && 12) {
            if (11->val < 12->val) {
                                    //11与12同时不空时,对它们进行比较
                 pre->next = 11;
                                   //如果I1对应的节点小于I2对应的节点
                11 = 11 - \text{next};
                                    //将pre与较小的节点进行连接
            else{
                 pre->next = 12;
                12 = 12 - \text{next};
                                 //pre指向新连接的节点
              pre = pre->next;
                    //如果I1有剩余
        if (11) {
                                //将I1接到pre后
              pre->next = 11;
                    //如果I2有剩余
        if (12) {
              pre->next = 12;
                                //将I2接到pre后
        return temp head.next;
```

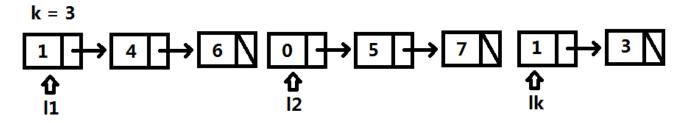


## 例6-a:测试与leetcode提交结果

```
int main(){
    ListNode a(1);
    ListNode b(4);
    ListNode c(6);
                        Merge Two Sorted Lists
                                                    青按仟意键绑
    ListNode d(0);
    ListNode e(5);
                        Submission Details
    ListNode f(7);
    a.next = &b;
                                                  Status: Accepted
                           208 / 208 test cases passed.
    b.next = &c;
                                              Submitted: 0 minutes ago
                           Runtime: 9 ms
    d.next = &e;
    e.next = &f;
    Solution solve;
    ListNode *head = solve.mergeTwoLists(&a, &d);
    while (head) {
         printf("%d\n", head->val);
         head = head->next;
    return 0:
```

# 例6-b:排序链表的合并(多个)

已知k个已排序链表头节点指针,将这k个链表合并,合并后仍为有序的,返回合并后的头节点。



### 合并后链表:



```
class Solution {
public: //链表头节点使用vector

ListNode* mergeKLists(std::vector<ListNode*>& lists) {
} //最终返回merge后的有序的链表头
};
```

### 选自 LeetCode 23. Merge k Sorted Lists

https://leetcode.com/problems/merge-k-sorted-lists/description/

难度:Hard

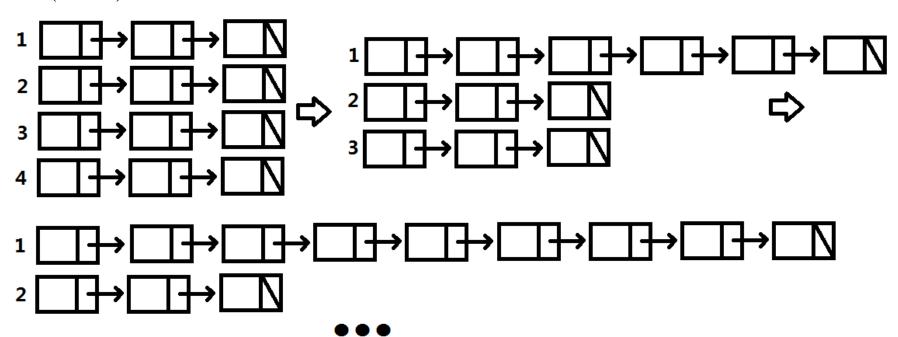


## 例6-b:方法1思考,暴力合并

方案1:最普通的方法, k个链表按顺序合并k-1次。

设有k个链表,平均每个链表有n个节点,时间复杂度:

 $(n+n)+(2n+n)+((k-1)n+n) = (1+2+...+k-1)n + (k-1)n = (1+2+...+k)n-n = (k^2+k-1)/2 * n = O(k^2*n)$ 



(3+3) + (6+3) + (9+3) = 27 次比较



## 例6-b:方法2思考,排序后相连

方案2:将k\*n个节点放到vector中,再将vector排序,再将节点顺序相连。

设有k个链表,平均每个链表有n个节点,时间复杂度:

```
kN*logkN + kN = O(kN*logkN) (比如k=100, n=10000) logkN = 20, k=100
```

```
#include <vector>
#include <algorithm> //STL 排序算法 std::sort
bool cmp (const ListNode *a, const ListNode *b) {
   return a->val < b->val;
               //比较函数,对节点进行从小到大的排序
int main(){
    ListNode a (3);
    ListNode b(2);
    ListNode C(5);
    ListNode d(0);
    std::vector<ListNode *> node vec;
    node vec.push back(&a);
    node vec.push back(&b);
    node vec.push back(&c);
    node_vec.push_back(&d); //调用排序函数
    std::sort(node vec.begin(), node vec.end(), cmp);
    for (int i = 0; i < node vec.size(); i++) {
       printf("%d\n", node vec[i]->val);
    return 0:
```

```
0
2
3
5
请按任意键继续..._
```



## 例6-b:方法3思考,分制后相连

方案3:对k个链表进行分制,两两进行合并。

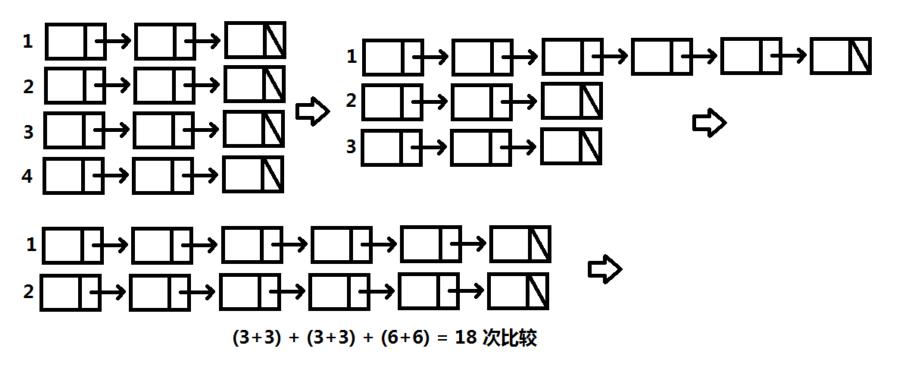
设有k个链表,平均每个链表有n个节点,时间复杂度:

第1轮,进行k/2次,每次处理2n个数字;第2轮,进行k/4次,每次处理4n个数字;…;

最后一次,进行k/(2^logk)次,每次处理2^logk\*N个值。

 $2N*k/2 + 4N*k/4 + 8N*k/8 + ... + 2^{\log k}*N*k/(2^{\log k})$ 

=Nk + Nk + ... + Nk = O(kNlogk)



# 例6-b:课堂练习(排序)

};

```
#include <vector>
#include <algorithm>
                                                 3分钟时间填写代码,有
bool cmp(const ListNode *a, const ListNode *b) {
    return a->val < b->val;
                                                 问题随时提出!
class Solution {
public:
   ListNode* mergeKLists(std::vector<ListNode*>& lists) {
       std::vector<ListNode *> node vec;
       for (int i = 0; i < lists.size(); i++){</pre>
                                         //遍历k个链表,将节点全部添加至
                                         node vec
           while (head) {
               head = head->next;
        if (node vec.size() == 0) {
           return NULL;
                         //根据节点数值进行排序
        std::sort(node vec.begin(), node vec.end(), cmp);
        for (int i = 1; i < node vec.size(); i++){</pre>
        node vec[node vec.size()-1]->next = NULL;
        return node vec[0];
```



## 例6-b:实现(排序)

};

```
#include <vector>
#include <algorithm>
bool cmp(const ListNode *a, const ListNode *b) {
    return a->val < b->val;
class Solution {
public:
    ListNode* mergeKLists(std::vector<ListNode*>& lists) {
        std::vector<ListNode *> node vec;
        for (int i = 0; i < lists.size(); i++){</pre>
                                            //遍历k个链表,将节点全部添加至
                ListNode *head = lists[i];
                                            node_vec
            while (head) {
                   node_vec.push_back(head);
                 head = head->next;
        if (node vec.size() == 0) {
             return NULL;
                           //根据节点数值进行排序
         std::sort(node vec.begin(), node vec.end(), cmp);
        for (int i = 1; i < node vec.size(); i++) {</pre>
                                                            //连接新的链表
                 node vec[i-1]->next = node vec[i];
         node vec[node vec.size()-1]->next = NULL;
         return node vec[0];
```

# 例6-b:课堂练习(分制)

ユルハグランスコンスパー くろうしし

```
ListNode* mergeKLists(std::vector<ListNode*>& lists) {
   if (lists.size() == 0) {
                          //如果lists为空,返回NULL
       return NULL:
   if (lists.size() == 1){
                           //如果只有两个list,调用两个list merge函数
   if (lists.size() == 2) {
       return mergeTwoLists(lists[0], lists[1]);
   int mid =
   std::vector<ListNode*> sub1 lists;
   std::vector<ListNode*> sub2 lists; //拆分lists为两个子lists
   for (int i = 0; i < mid; i++) {
       sub1 lists.push back(lists[i]);
                                               3分钟时间填写代码。有
   for (int i = mid; i < lists.size(); i++) {</pre>
       sub2 lists.push back(lists[i]);
                                               问题随时提出!
   ListNode *11 =
   ListNode *12 =
```

## 例6-b:实现(分制)

```
ListNode* mergeKLists(std::vector<ListNode*>& lists) {
    if (lists.size() == 0) {
                             //如果lists为空,返回NULL
        return NULL;
    if (lists.size() == 1){
                             //如果只有一个lists , 直接返回头指针
          return lists[0];
                             //如果只有两个list,调用两个list merge函数
      (lists.size() == 2){}
        return mergeTwoLists(lists[0], lists[1]);
                lists.size() / 2;
    int mid =
    std::vector<ListNode*> sub1 lists;
    std::vector<ListNode*> sub2 lists; //拆分lists为两个子lists
    for (int i = 0; i < mid; i++) {
        sub1 lists.push back(lists[i]);
    for (int i = mid; i < lists.size(); i++){</pre>
        sub2 lists.push back(lists[i]);
                    mergeKLists(sub1_lists);
    ListNode *11 =
    ListNode *12 =
                    mergeKLists(sub2_lists);
```

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## 例6-b:测试与leetcode提交结果

```
Merge k Sorted Lists
                                                              排序算法时间
int main(){
    ListNode a(1);
                                        Submission Details
    ListNode b(4);
    ListNode c(6);
    ListNode d(0);
                                           130 / 130 test cases passed.
                                                                       Status: Accepted
    ListNode e(5);
                                           Runtime: 42 ms
                                                                  Submitted: 0 minutes ago
    ListNode f(7);
    ListNode q(2);
    ListNode h(3);
                                         Merge k Sorted Lists
                                                              分制算法时间
    a.next = \&b;
    b.next = &c;
                                         Submission Details
    d.next = &e;
    e.next = &f;
    q.next = &h;
                                                                       Status: Accepted
                                            130 / 130 test cases passed.
    Solution solve;
                                            Runtime: 32 ms
                                                                   Submitted: 0 minutes ago
    std::vector<ListNode *> lists;
    lists.push back(&a);
    lists.push back(&d);
    lists.push back(&q);
    ListNode *head = solve.mergeKLists(lists);
    while (head) {
         printf("%d\n", head->val);
         head = head->next;
    return 0:
```

## 一些建议:关于课程复习与面试准备

- 1. 在不看回放与复习的前提下,重新编写这8道题目进行提交测试。
- 2. 尽量多的通过这些题目。
- 3. 再看一轮回放,再编写一遍题目并提交通过。
- 4. 纸上写代码, 反复练习。
- 5. 将leetcode其他链表相关题目完成。



一些建议:关于学习,如何成为一名优秀的研发工程师

- 1. 掌握一门编译语言
- 2. 掌握算法与数据结构
- 3. 掌握一门脚本语言
- 4. 掌握开发环境
- 5. 丰富其他前沿知识



## 一些建议:关于面试

- 1. 听清并搞懂问题
- 2. 冷静思考
- 3. 努力解决,不轻易放弃
- 4. 自信并谦虚



## 一些建议:关于实习

- 1. 实习是拿到BAT offer的捷径
- 2. 实习的岗位有许多:RD、QA、FE、OP、PM等
- 3. 尽量在一个实习岗位坚持半年
- 4. 实习还有不菲的薪水喔~
- 5. 关于学校课堂与上课, 自己把握



## 结束

# 非常感谢大家!

林沐

