Operating Systems

Timers

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In this episode

- •We will replace the delay function that is cpu intensive with a timer mechanism that puts a thread that wants to sleep for a certain number of cycles in the waiting status.
- When enough time is passed, the process is put again in the ready list
- We need
 - A structure representing the timer event
 - A timer counter incremented by the interrupt routine
 - A queue of timer events to store all timers installed by a process
 - A syscall sleep(num_cycles) that installs a timer and puts a process to sleep

Timer

- A timer event stores
 - the awakening time of a process. This is the absolute time when a process should be woken up
 - The pointer to the pcb of the process to wake up
- The list of timer events is ordered by awakening time

```
typedef struct {
  ListItem list;
  int awakening time;
  struct PCB* pcb;
} TimerItem;
//list of timers
typedef ListHead TimerHandler;
// initializes a timer
void Timer_init();
// initializes a timer list
void TimerList_init(ListHead* timers);
// inserts in the list a new timer event
// the list is ordered by awakening time
TimerItem* TimerList_add(ListHead* timers,
              int awakening_time,
              struct PCB* pcb);
// returns, if existing
// a timer matching with current_time
    TimerItem* TimerList_current(ListHead* timers,
    int current_time);
// removes the head
int TimerList removeCurrent(ListHead* timers);
void TimerList_print(ListHead* timers);
// deallocates a timer
int TimerItem_free(TimerItem* item);
```

Global Variables

We need a global variable to keep track of the progressing time.

We need a list of timers to keep track of the processes to wake up

This variable is incremented each time the timerInterrupt is invoked.

In the PID of each process we need to store a pointer (if any) to the timer event potentially installed by the process. This is just for bookkeeping.

```
FILE* log file=NULL;
PCB* init pcb;
PCB* running;
int last pid;
ListHead ready list;
ListHead waiting_list;
ListHead zombie list;
ListHead timer_list;
ListHead resources list;
SyscallFunctionType
syscall vector[DSOS MAX SYSCALLS];
int syscall_numarg[DSOS_MAX_SYSCALLS];
ucontext_t interrupt_context;
ucontext t trap context;
ucontext t main context;
ucontext t idle context;
int shutdown now=0; // used for termination
char system stack[STACK SIZE];
// process wide signal mask
sigset_t signal_set;
char signal stack[STACK SIZE];
volatile int disastrOS_time=0;
```

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```
void timerInterrupt(){
   ++disastr0S_time;
   internal_schedule();
   setcontext(&running->cpu_state);
}
```

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```
typedef struct PCB{
 ListItem list; // MUST BE THE FIRST!!!
 int pid;
 int return value; // ret value for the parent
 ProcessStatus status;
 int signals;
 int signals mask;
 ListHead descriptors;
  struct PCB* parent;
 ListHead children;
 ucontext t cpu state;
 struct TimerItem *timer;
 char stack[STACK_SIZE];
 int syscall_num;
 long int syscall_args[DSOS_MAX_SYSCALLS_ARGS];
 int syscall retvalue;
} PCB;
```

Scheduler

Each time the scheduler is called we need to scan the list of timers to see if the current time matches the head of the list.

If this occurs, it means that we need to wake up the corresponding process, by moving it from the waiting to the ready queue.

We finally delete the timer event, and clear its reference in the PID of the process

```
void internal schedule() {
  TimerItem* elapsed timer=0;
  PCB* previous pcb=0;
  while( (elapsed_timer
    =TimerList_current(&timer_list,
                       disastrOS_time)) ){
    PCB* pcb to wake=elapsed timer->pcb;
    List_detach(&waiting_list,
                (ListItem*) pcb to wake);
    pcb_to_wake->status=Ready;
    pcb to wake->timer=0;
    List_insert(&ready_list,
                (ListItem*) previous pcb,
                (ListItem*) pcb_to_wake);
    previous pcb=pcb to wake;
    TimerList_removeCurrent(&timer_list);
 // regular schedule
  if (ready list.first){
    PCB* next process=
         (PCB*) List_detach(&ready_list,
                             ready_list.first);
    running->status=Ready;
    List_insert(&ready_list,
                 ready list.last,
                (ListItem*) running);
    next process->status=Running;
    running=next_process;
```

sleep(int cycles)

- This new syscall will put a process to sleep for a certain number of cycles
- When a thread calls sleep, a new timer is created
- The awakening time of the timer is set as current_time+cycles
- The process is put in the waiting queue

```
void internal sleep(){
  if (running->timer) {
    printf("process has already a timer!!!\n");
    running->syscall retvalue=DSOS ESLEEP;
    return;
  int cycles_to_sleep=running->syscall_args[0];
  int wake time=disastrOS time+cycles to sleep;
  TimerItem* new timer=
    TimerList_add(&timer_list, wake_time, running);
  if (! new_timer) {
    printf("no new timer!!!\n");
    running->syscall retvalue=DSOS ESLEEP;
    return;
  running->status=Waiting;
  List_insert(&waiting_list,
    waiting_list.last,
    (ListItem*) running);
  if (ready_list.first)
    running=(PCB*)
    List_detach(&ready_list, ready_list.first);
 else {
    running=0;
    printf ("they are all sleeping\n");
    disastrOS_printStatus();
```

Timer test

We simply refer to the test program of the last episode, and we replace the waitABit with a call to sleep

Messages

When in kernel mode, we manage the system simply by changing the value of some data structures.

No matter what, when we exit kernel mode, the next process in running will be resumed.