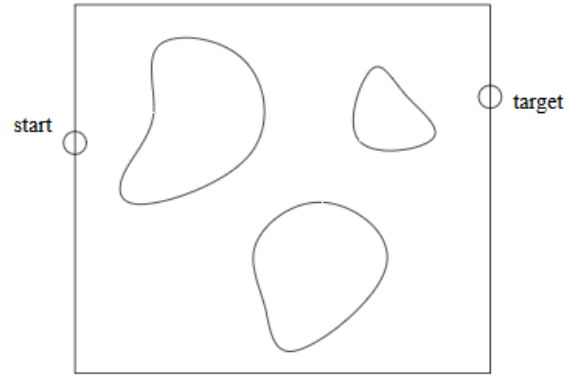


# geometric A\*



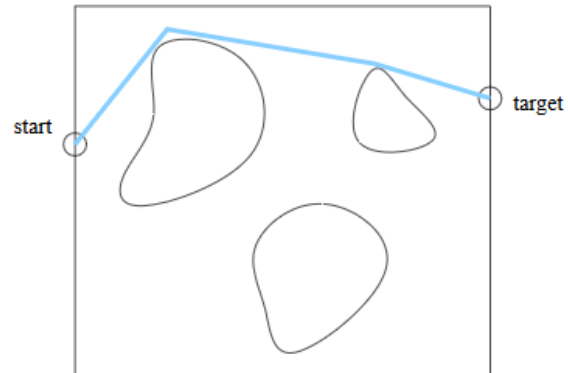
the problem: go from start to end

region of the plane

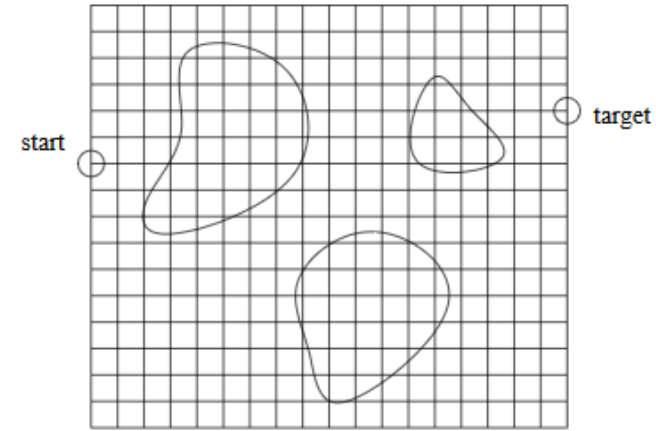
minimal path

applications:

- autonomous vehicles
- non-player characters in video games

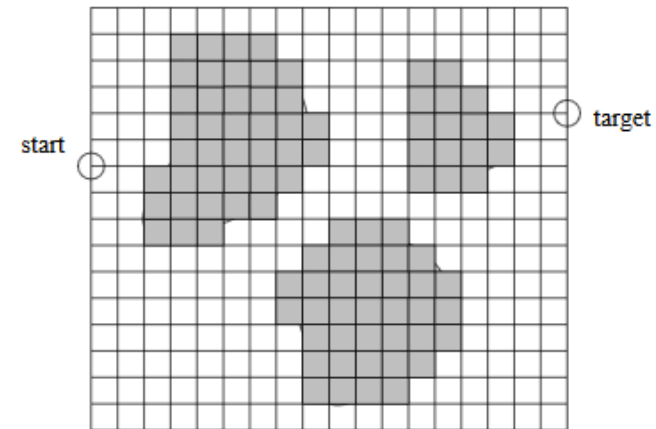


actual minimal path hard to find



approximation: grid

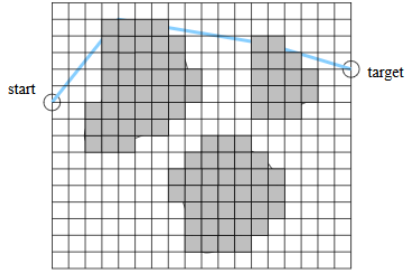
only moving on the grid + diagonals



occluded areas on the grid

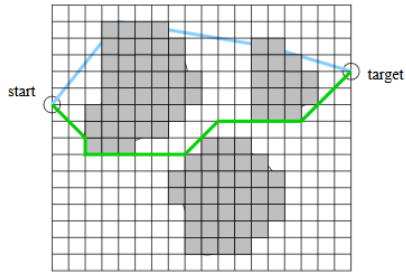
forbidden: diagonals inside a black box

horizontal or vertical line surrounded by two black boxes



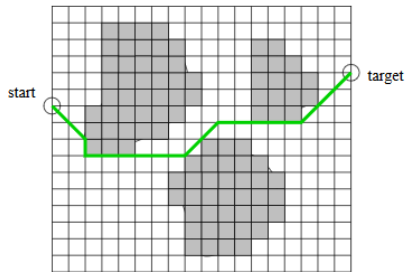
real optimal paths may become impossible as is

meaning: with the grid simplifications some optimal paths become impossible  
the simplification actually may make paths longer



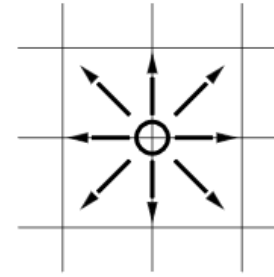
an optimal path in the grid

feasible also in the original problem  
but not optimal



finding optimal paths: use A\*

a state for each point in the grid  
allowed moves



allowed moves

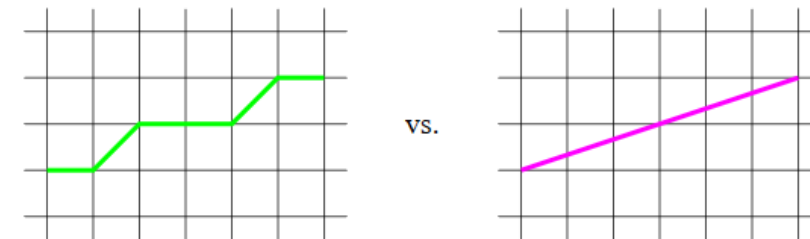
is a restriction

path not really optimal

suboptimality:

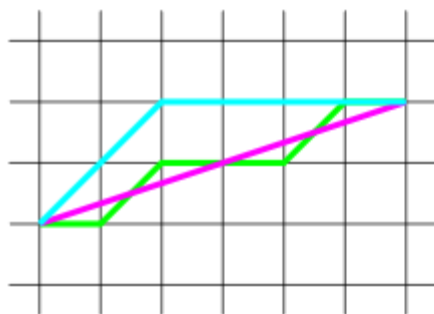
- the occluded areas may not be as in the grid
- the grid constraints the movements

increasing resolution attenuates the first  
not the second



optimal path vs. optimal in the grid

resolution irrelevant  
(why?)



every path in the grid is the same as:

- first all diagonal moves
- then all straight moves

increasing resolution is irrelevant

---

path sub-optimal: really a problem?

vehicles: large number of turns

npc: unnatural movements (drunken-like)

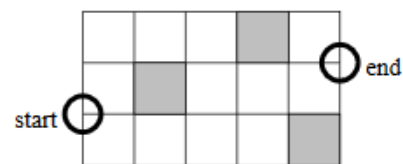
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some solutions:

- visibility graphs
- post smoothing
- Field D\*
- Theta\*



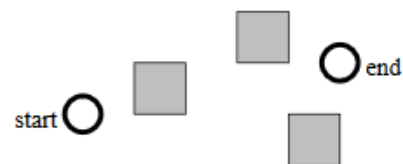
## visibility graph



simple problem

---

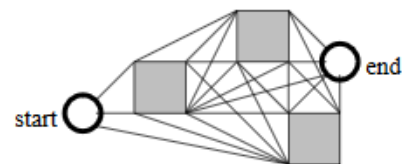
## visibility graph: no grid



consider only the occluded areas

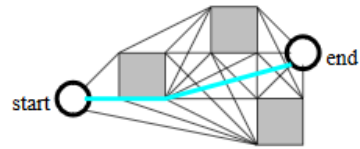
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## visibility graph: allowed moves



link every corner to every other  
also: start and end

## visibility graph: minimal path

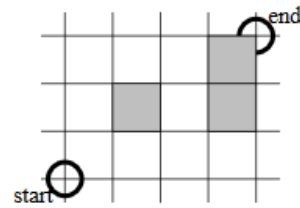


path in this new grid is optimal

but: many links  
on the grid: only from each state to neighbors

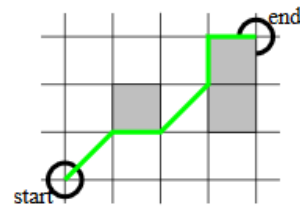
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## post-smoothing

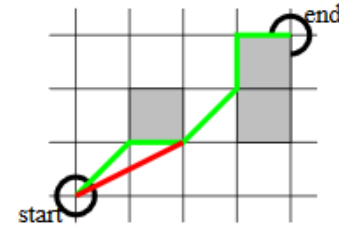


first find optimal path in the grid  
then smooth it

---

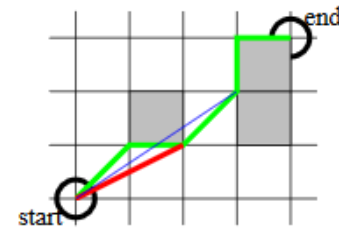


optimal path in the grid



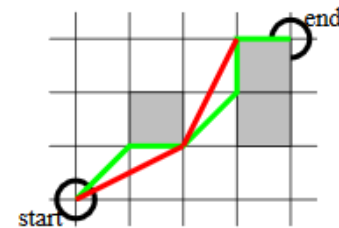
third is visible from one: skip second

---



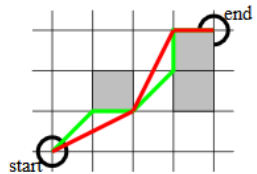
fourth is not visible from one: do not skip third  
repeat from third node

---



fifth node visible from third  
skip fourth

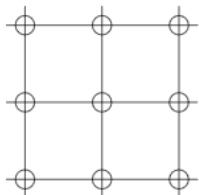
etc.



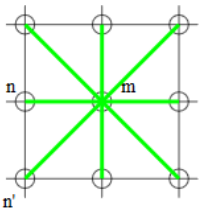
post-smoothing algorithm:

fix path by skipping nodes if successors are in line of sight of predecessor

# field D\*



the grid



A\* on grid: allow only moves on the grid

not only!

also:  $d(m) = \min d(n) + n \rightarrow m$  only for nodes  $n, n', \dots$  on the grid

[note] A\* allows only moves on the grid, but the grid matters not only to this. When determining the distance of a node from the start, only the predecessors on the grid are considered; this is the second aspect of the algorithm affected by the grid.

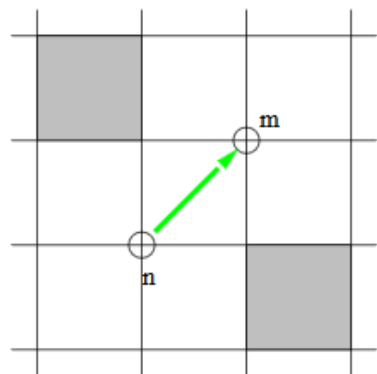
In the figure,  $m$  is a node for which field A\* is about to calculate the distance  $d(m)$ . Only nodes on the grid like  $n, n', \dots$  are considered as *predecessors* of  $m$ . Therefore,  $d(m)$  is calculated using them only.

## theta\*

incorporates smoothing into searching

---

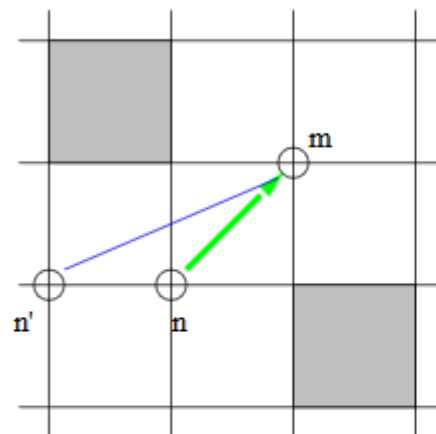
### allowed moves



A\* considers only moves like  $n \rightarrow m$   
 $n$  and  $m$  neighbors in the grid



### shortcuts

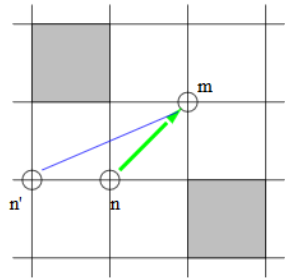


theta\* checks whether  $n$  has a predecessor  $n'$   
with a line-of-sight to  $m$

$n'$  becomes a new predecessor of  $m$

matters to?...

**shortcut = short distance**



theta\* checks whether  $n$  has a predecessor  $n'$  with a line-of-sight to  $m$

$n'$  becomes a new predecessor of  $m$

matters on:

- value of  $d(m)$   
 $n' \rightarrow m$  shorter than  $n \rightarrow m$
- optimal path:  
shorter  
smoothed (no turn on  $n$ )

[note] Precisely, theta\* stores the best predecessor of each node  $n$ , and only checks the line-of-sight from it to  $m$ . The new allowed move  $n' \rightarrow m$  is not really added as an allowed move, only  $n'$  stored as the new best predecessor of  $m$  if  $start \Rightarrow n' \rightarrow m$  is currently the shortest path to  $m$ .