

incremental A*

a path from start to end has been found by A*

but now the problem changes

in the example of trucks and parcels: a truck is out of order, a route is closed, a warehouse:

optimal plans may not still be optimal
may not be executable at all!

incremental A*: settings

A* has finished

but now:

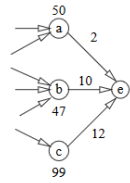
1. some $n \rightarrow m$ become unavailable
2. some new $n \rightarrow m$ are created
3. more generally, the cost of some $n \rightarrow m$ change

search from scratch

or use data obtained by the previous A* search

[italian] from scratch = dall'inizio, senza sfruttare quello che si era fatto in precedenza

obvious case

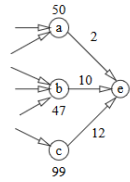


situation at end of A*

meaning of numbers: $d(a)=50$, $a \rightarrow e=2$, etc.

optimal path?

optimal path



best path: go to a, then e
cost is $d(a) + a \rightarrow e = 50 + 2$

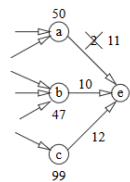
cost of the alternatives: $47 + 10 = 57$ and $99 + 12 = 111$
worse

how to get to a?

same method: choose cheapest predecessor + step

what if $a \rightarrow e$ increases?

increased cost at the end



cost of $a \rightarrow e$ increases

alternative paths to e:

- $d(a) + a \rightarrow e = 50 + 11 = 61$
- $d(b) + b \rightarrow e = 47 + 10 = 57$
- $d(c) + c \rightarrow e = 99 + 12 = 111$

go to b, then to e



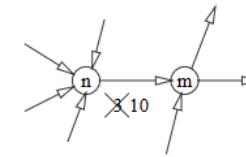
inner nodes?

data from previous search may be useful
proved when changing $a \rightarrow e$

e is the target state

what if an inner step changes cost?

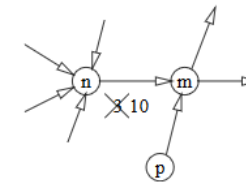
arbitrary change of cost



example: $n \rightarrow m$ changes from 3 to 10

both n and m are linked to other states!

effect of change



previously: $n \rightarrow m$ may be the optimal way to m

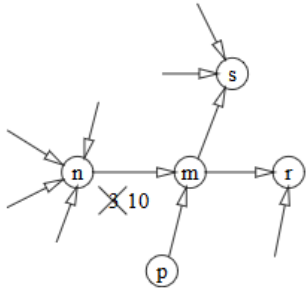
now: $p \rightarrow m$ may be

happens if: $d(p) + p \rightarrow m < d(n) + n \rightarrow m$

means: $start \Rightarrow p \rightarrow m$ cheaper than $start \Rightarrow n \rightarrow m$

not the only change!

chain reaction



$d(m)$ = best known $start \Rightarrow m$

previously $d(m) = d(n) + n \rightarrow m$ (with $n \rightarrow m = 3$)

now $d(m) = d(p) + p \rightarrow m$ (greater)

affects the successors of m

maybe $start \Rightarrow m \rightarrow r$ no longer the cheapest way to s

how to propagate the changes

seen before:

$d(_)$ correct \Rightarrow minimal path

$n \rightarrow m$ changed, $d(_)$ invalid

fix it

[note] As seen before, if $d(_)$ is valid an optimal path can be found by starting from the end and then repeatedly going back from the current node m to the predecessor n with a minimal $d(n) + n \rightarrow m$.

When some costs $n \rightarrow m$ change, $d(_)$ may not longer be correct. The previously optimal plan may no longer be optimal.

fix invalid distances

this is what A* does

invariant in A*:

$d(m)$ is the minimal cost of a path of closed states from the start to m

“path of closed states” = all states closed but possibly m

at each step, A* maintain this condition true

while attempting to enlarge the set of closed nodes

it “fixes” $d(_)$

[note] The invariant is valid for every state m , including the ones that have not not been reached at all. This is because the algorithm initializes $d(m) = \infty$, which is the correct distance for a state that is not reached by a path of closed states.

the invariant in A*

$d(m)$ is the minimal cost of a path of closed states from the start to m

when n enters the set of closed nodes,

if $start \Rightarrow n$ is a path of closed nodes and $n \rightarrow m$ exists:

- since the invariant was true before and $start \Rightarrow n$ is a path of closed nodes then $d(n)$ is correct
- path $start \Rightarrow n \rightarrow m$ is now a path of closed nodes to m
- its cost is $d(n) + n \rightarrow m$ since $d(n)$ is correct
- if the current value of $d(m)$ is less than or equal $d(n) + n \rightarrow m$,
this value is now wrong
(otherwise, the following steps are not done)
- set $d(m) = d(n) + n \rightarrow m$
- if m is closed, now $start \Rightarrow n \rightarrow m \rightarrow r$ is a new path of closed nodes to every successor r of m
- as a result, $d(r)$ may now be wrong
- to fix the problem, remove m from the closed states

A* tries to enlarge the set of closed nodes

ensuring that the invariant is always correct

fixing distances by A*

no need for a special algorithm for fixing distances
A* itself makes distances valid

works without changes if $n \rightarrow m$ may only decrease

just reopen m and continue

increasing costs

rule for always decreasing distances:

if $d(n) + n \rightarrow m < d(m)$, update $d(m)$ and reopen m

what it does:

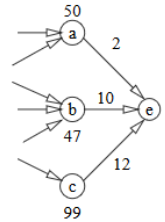
ensure that $d(_)$ is the minimal cost of a path of closed states

new path found \Rightarrow maybe better than the previous

in general?

minimal path

seen before:



if $a \rightarrow e$, $b \rightarrow e$, $c \rightarrow e$ exist

minimal cost of path to e is:

$$\min(d(a) + a \rightarrow e, d(b) + b \rightarrow e, d(c) + c \rightarrow e)$$

not only if e is the target

true in general

why: $d(a)$ is the minimal cost for reaching a , etc.



incremental A*

instead of: if $d(m) > d(n) + n \rightarrow m \dots$

use the new rule:

- $cp(m) = \{ n \mid n \rightarrow m, n \in \text{closed} \}$
 - $\text{temp} = \min_{n \in cp(m)} d(n) + n \rightarrow m$
 - if $\text{temp} \neq d(m)$ then:
 - set $d(m) = \text{temp}$
 - remove m from closed
-

reopening

A* tries to reach the target by closed states

by the invariant $d(\text{end})$ would be correct
allows going back by the cheapest path

removing states from the set of closed is in the opposite direction
BUT: necessary to maintain the invariant

actual algorithms

LPA*, D*, ...