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Introduction

Deal with images in

- Spatial domain Processing in pixel/spatial domain.
- Frequency Domain Processing in the transformed domain.

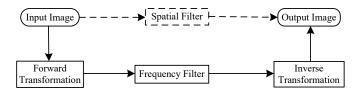


Figure: Frequency domain filtering





Frequency

Time series:

- In 1-D time series, frequency is related to $\frac{dA}{dt}$.
- If rate of change of amplitude higher then frequency is also higher

Images:

- For images, spatial frequency is related to $\frac{dl}{dx}$.
- For edges frequency is higher, smooth regions frequency is lower.



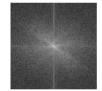


Figure: 1-D Time Signal

Figure: An Image.



Why Frequency Domain Processing?

- It is computationally faster to perform 2D Fourier transforms and a filter multiply than to perform a convolution in the image (spatial) domain.
- Frequency domain has a established suit of processes and tools that can be borrowed directly from signal processing in other domains.
- Easy to remove periodic noise in images.
- The parameters are easily manipulated.
- Very useful in compression few coefficients represent most of the information.



Outlines

Image Transforms

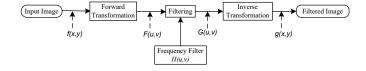


Figure: Frequency domain filtering

- $lue{}$ Convolution in spatial domain \leftrightarrow Multiplication in frequency domain.
- G(u,v) = H(u,v).F(u,v)





Image Transforms

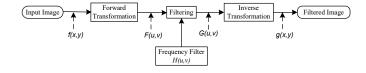


Figure: Frequency domain filtering

- $lue{}$ Convolution in spatial domain \leftrightarrow Multiplication in frequency domain.
- G(u,v) = H(u,v).F(u,v)

Standard Transforms - Kernel is image independent

- Discrete Fourier Transform (DFT).
- Discrete Cosine Transform (DCT).
- Discrete Sine Transform (DST).
- Discrete Hadamard Transform (DHT).

Image dependent Transforms - Kernel is image dependent

K-L Transform.



Discrete Fourier Transform

1-D Continuous Fourier Transform

$$F(u) = \int_{-\infty}^{\infty} f(x)e^{-j2\pi ux} dx \tag{1}$$

2-D Continuous Discrete Fourier Transform

$$F(u,v) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y)e^{-j2\pi(ux+vy)}dxdy$$
 (2)





Discrete Fourier Transform

1-D Continuous Fourier Transform

$$F(u) = \int_{-\infty}^{\infty} f(x) e^{-j2\pi ux} dx$$
 (1)

2-D Continuous Discrete Fourier Transform

$$F(u,v) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) e^{-j2\pi(ux+vy)} dxdy$$
 (2)

2-D Discrete Fourier Transform

$$F(u,v) = \frac{1}{N} \sum_{x=0}^{N-1} \sum_{y=0}^{N-1} f(x,y) e^{-j\frac{2\pi}{N}(ux+vy)}$$
(3)

2-D Inverse Discrete Fourier Transform

$$f(x,y) = \frac{1}{N} \sum_{u=0}^{N-1} \sum_{v=0}^{N-1} F(u,v) e^{j\frac{2\pi}{N}(ux+vy)}$$
(4)

For a 512 \times 512 image \rightarrow 512⁴ complex multiplications. Huge complexity!



Separability Property

2-D Discrete Fourier Transform

$$F(u,v) = \frac{1}{N} \sum_{x=0}^{N-1} \sum_{y=0}^{N-1} f(x,y) e^{-j\frac{2\pi}{N}(ux+vy)} = \frac{1}{N} \sum_{x=0}^{N-1} e^{-j\frac{2\pi}{N}ux} \cdot \sqrt{N} \cdot \frac{1}{\sqrt{N}} \sum_{y=0}^{N-1} f(x,y) e^{-j\frac{2\pi}{N}vy}$$

$$= \frac{1}{N} \sum_{x=0}^{N-1} e^{-j\frac{2\pi}{N}ux} \cdot \sqrt{N} \cdot F(x,v) = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} F(x,v) e^{-j\frac{2\pi}{N}ux} = F(u,v)$$





Separability Property

2-D Discrete Fourier Transform

$$F(u,v) = \frac{1}{N} \sum_{x=0}^{N-1} \sum_{y=0}^{N-1} f(x,y) e^{-j\frac{2\pi}{N}(ux+vy)} = \frac{1}{N} \sum_{x=0}^{N-1} e^{-j\frac{2\pi}{N}ux} \cdot \sqrt{N} \cdot \frac{1}{\sqrt{N}} \sum_{y=0}^{N-1} f(x,y) e^{-j\frac{2\pi}{N}vy}$$
$$= \frac{1}{N} \sum_{x=0}^{N-1} e^{-j\frac{2\pi}{N}ux} \cdot \sqrt{N} \cdot F(x,v) = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} F(x,v) e^{-j\frac{2\pi}{N}ux} = F(u,v)$$

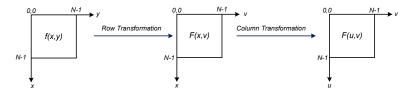


Figure: Separability - 2-D Transform as combination of 1-D Transform



Fast Fourier Transform

- Faster implementation of DFT.
- Exploiting Periodicity and Symmetry properties of DFT.

$$F(u) = \frac{1}{N} \sum_{x=0}^{N-1} f(x) e^{-j\frac{2\pi}{N}ux} = \frac{1}{N} \sum_{x=0}^{N-1} f(x) W_N^{ux}$$
 (5)

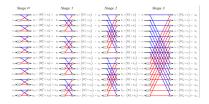


Figure: Butterfly structure.

Figure: 16 point FFT



- For an N-Point FFT, computational complexity would by O(NlogN).
- For an image we can have a computational complexity of $O(N^2 log N)$.
- \blacksquare For 512 \times 512 image 128 \times 512 \times 9 complex multiplications.
- Similarly we have faster implementation of IDFT.
- Implement 1-D FFT as an recursive function.



