## UNIVERSITY OF BRISTOL

**January 2021 Examination Period** 

## FACULTY OF ENGINEERING

Third Year Examination for the Degrees of Bachelor of Science Master of Engineering

COMS30039J Types and Lambda Calculus

## TIME ALLOWED: 2 Hours

This paper contains *two* questions, answer *both*. Each question is worth 25 marks. The maximum for this paper is *50 marks*. Credit will be given for partially correct answers.

PLEASE WRITE YOUR 7 DIGIT STUDENT NUMBER (NOT CANDIDATE NUMBER) ON THE ANSWER BOOKLET. YOUR STUDENT NUMBER CAN BE FOUND ON YOUR UCARD.

## **Other Instructions:**

You may use the lecture notes to help you, but you may not collaborate.

YOU MAY START IMMEDIATELY

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**Q1**. This question concerns the pure, untyped  $\lambda$ -calculus.

It will be useful to recall the combinators  $\mathbf{I} = \lambda x \cdot x$  and  $\mathbf{K} = \lambda x y \cdot x$ 

- (a) Give the set of free variables for each of the following terms:
  - i.  $\lambda xyz.xz(yz)$
  - ii.  $z(\lambda x. xy)$
  - iii.  $\lambda z.(\lambda x.y)(xz)$
  - iv.  $(\lambda x. x (\lambda y. xy))(\lambda x. y)$

[4 marks]

- (b) For each of the following, give a term M that satisfies the statement.
  - i.  $Mx =_{\beta} xx$
  - ii.  $M \lceil n \rceil =_{\beta} \lceil 2 * n \rceil$
  - iii.  $M =_{\beta} MM$
  - iv.  $Mx =_{\beta} MM$

[8 marks]

(c) Recall that there is a combinator **Pred** satisfying:

**Pred** 
$$\lceil 0 \rceil =_{\beta} \lceil 0 \rceil$$

**Pred** 
$$\lceil n+1 \rceil =_{\beta} \lceil n \rceil$$

Define combinators **Sub** and **Succ** as follows:

Succ 
$$= \lambda nfx f(nfx)$$

**Sub** 
$$= \lambda mn \ n \operatorname{Pred} m$$

Prove, by induction on *n*, that **Sub** satisfies:

**Sub** 
$$\lceil m \rceil \lceil n \rceil =_{\beta} \begin{cases} \lceil 0 \rceil & \text{if } m \leq n \\ \lceil m - n \rceil & \text{otherwise} \end{cases}$$

You may use the fact that, for all natural numbers p,  $\lceil p+1 \rceil =_{\beta} \mathbf{Succ} \lceil p \rceil$ .

[7 marks]

(d) Prove that there does not exist a term M such that, for all terms N:

$$MN =_{\beta} \begin{cases} \mathbf{I} & \text{if } N \text{ is in } \beta\text{-normal form} \\ \mathbf{K} & \text{otherwise} \end{cases}$$

[3 marks]

(e) Find a finite sequence of *closed* terms  $M_1, M_2, \ldots, M_k$  for  $k \ge 0$  such that the following two equations are both satisfied:

$$(\lambda x. x \mathbf{I} (x \mathbf{II})) M_1 M_2 \cdots M_k x y =_{\beta} x$$
$$(\lambda x. x \mathbf{I} (x x \mathbf{I})) M_1 M_2 \cdots M_k x y =_{\beta} y$$

[3 marks]

- **Q2**. This question concerns type systems.
  - (a) Give a typing derivation for each of the following judgements:
    - i.  $\vdash \lambda xy.x: a \rightarrow b \rightarrow a$
    - ii.  $\vdash \lambda x. x(\lambda y. y) : ((a \rightarrow a) \rightarrow b) \rightarrow b$
    - iii.  $y : c \vdash (\lambda x. y)(\lambda xz. x) : c$

[6 marks]

- (b) For each of the following types, find a closed term that inhabits the type:
  - i.  $a \rightarrow b \rightarrow b$
  - ii.  $(a \rightarrow b) \rightarrow (a \rightarrow b \rightarrow c) \rightarrow a \rightarrow c$
  - iii.  $c \to ((a \to b \to a) \to c \to d) \to d$

[6 marks]

(c) Prove, by induction on the Damas-Miler typing relation, that: if  $\Gamma \vdash M : A$  then, for all  $\Gamma'$ ,  $\Gamma \subseteq \Gamma'$  implies  $\Gamma' \vdash M : A$ .

[6 marks]

(d) Define the order, order(A) of a monotype A recursively on the structure of A:

$$order(a) = 0$$
  
 $order(B \rightarrow C) = max(1 + order(B), order(C))$ 

where  $\max(m, n)$  is the larger of the two natural numbers m and n. Prove that, if  $\operatorname{order}(A) \leq 2$  and  $\vdash M : A$ , then there is a term N such that:

- (i)  $M =_{\beta} \lambda x_1 \dots x_n$ . N
- (ii) and, moreover,  ${\it N}$  does not contain any abstraction as a subterm

[7 marks]