## UNIVERSITY OF BRISTOL

**January 2021 Examination Period** 

## FACULTY OF ENGINEERING

Third Year Examination for the Degrees of Bachelor of Science Master of Engineering

COMS30039J Types and Lambda Calculus

## TIME ALLOWED: 2 Hours

This paper contains *two* questions, answer *both*. Each question is worth 25 marks. The maximum for this paper is *50 marks*. Credit will be given for partially correct answers.

PLEASE WRITE YOUR 7 DIGIT STUDENT NUMBER (NOT CANDIDATE NUMBER) ON THE ANSWER BOOKLET. YOUR STUDENT NUMBER CAN BE FOUND ON YOUR UCARD.

## **Other Instructions:**

You may use the lecture notes to help you, but you may not collaborate.

YOU MAY START IMMEDIATELY

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**Q1**. This question concerns the pure, untyped  $\lambda$ -calculus.

It will be useful to recall the combinators  $\underline{id} = \lambda x. x$  and  $\underline{const} = \lambda xy. x$ 

- (a) Give the set of free variables for each of the following terms:
  - i.  $\lambda xyz.xz(yz)$
  - ii.  $z(\lambda x. xy)$
  - iii.  $\lambda z.(\lambda x.y)(xz)$
  - iv.  $(\lambda x. x (\lambda y. xy))(\lambda x. y)$

[4 marks]

- (b) For each of the following, give a term M that satisfies the statement.
  - i.  $Mx \approx xx$
  - ii.  $M \underline{n} \approx \underline{2 * n}$
  - iii.  $M \approx MM$
  - iv.  $Mx \approx MM$

[8 marks]

(c) Define combinator sub as follows:

$$\underline{\operatorname{sub}} := \operatorname{fix} (\lambda f \, m n. \, \operatorname{ifz} n \, m \, (f \, (\operatorname{pred} m) \, (\operatorname{pred} n)))$$

Prove, by induction on n, that <u>sub</u> satisfies:

$$\underline{\mathsf{sub}} \ \underline{m} \ \underline{n} \ \approx \ \begin{cases} \underline{0} & \text{if } m \leq n \\ \underline{m-n} & \text{otherwise} \end{cases}$$

[6 marks]

(d) Prove that there does not exist a term M such that, for all terms N:

$$MN \approx \begin{cases} \underline{id} & \text{if } N \text{ is in normal form} \\ \underline{\text{const}} & \text{otherwise} \end{cases}$$

[3 marks]

(e) Find a finite sequence of *closed* terms  $M_1, M_2, \ldots, M_k$  for  $k \ge 0$  such that the following two equations are both satisfied:

$$(\lambda x. \times \underline{id} (x \underline{id} \underline{id})) M_1 M_2 \cdots M_k \times y \approx x (\lambda x. \times \underline{id} (x \times \underline{id})) M_1 M_2 \cdots M_k \times y \approx y$$

[3 marks]

- **Q2**. This question concerns type systems.
  - (a) Give a typing derivation for each of the following judgements:
    - i.  $\vdash \lambda xy.x: a \rightarrow b \rightarrow a$
    - ii.  $\vdash \lambda x. x(\lambda y. y) : ((a \rightarrow a) \rightarrow b) \rightarrow b$
    - iii.  $y : c \vdash (\lambda x. y)(\lambda xz. x) : c$

[6 marks]

- (b) For each of the following types, find a closed pure term that inhabits the type:
  - i.  $a \rightarrow b \rightarrow b$
  - ii.  $(a \rightarrow b) \rightarrow (a \rightarrow b \rightarrow c) \rightarrow a \rightarrow c$
  - iii.  $c \to ((a \to b \to a) \to c \to d) \to d$

[6 marks]

(c) Prove, by induction M, that: for all M, A,  $\Gamma$ ,  $\Gamma'$ , if  $\Gamma \vdash M : A$  and  $\Gamma \subseteq \Gamma'$  then  $\Gamma' \vdash M : A$ .

[7 marks]

- (d) i. Prove that the only closed normal form of type  $a \to b \to a$  is  $\lambda xy. x$ .
  - ii. Give a closed normal form of type  $a \rightarrow b \rightarrow a$  which is not itself an abstraction.

[7 marks]