

TYPES AND λ -CALCULUS

Problem Sheet 7

** 1.

- (a) Find an inhabitant of the type $((a \rightarrow b) \rightarrow b) \rightarrow a \rightarrow b$.
- (b) Give the corresponding proof of the corresponding formula.

** 2. Sketch an algorithm to decide the following problem and justify that it works:

Given: two typable terms M and N
Decide: if $M =_{\beta} N$

** 3. A term M is a *fixed point combinator* just if, for all terms P , $M P =_{\beta} P (M P)$. In other words, M computes a fixed point of its argument.

Prove that no fixed point combinator is typable, i.e. if closed term M is a fixed point combinator, then M is not typable.

Hint: Try to arrive at a contradiction by obtaining a β -equality in which the two sides are distinct normal forms.

** 4. Suppose we add a fixed point combinator fix to our lambda calculus as a new primitive. In other words, we extend the syntax of terms by the rule:

$$\frac{}{\text{fix} \in \Lambda}$$

and we extend the type system by the following rule:

$$\frac{}{\Gamma \vdash \text{fix} : (A \rightarrow A) \rightarrow A} \text{ (TFix)}$$

This type makes sense since `fix` takes a function as input and returns a fixed point of the function. (We should also extend the definition of β -reduction, but it is not important to this question.)

- (a) Show that every type in this extended system is inhabited.
- (b) What is the consequence for the Curry-Howard correspondence?