

# SMART CONTRACT AUDIT REPORT

for

UpDeFi

Prepared By: Yiqun Chen

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## Contact

For more information about this document and its contents, please contact PeckShield Inc.

Name	Yiqun Chen	
Phone	+86 183 5897 7782	
Email	contact@peckshield.com	

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# 1 Introduction

Given the opportunity to review the design document and related smart contract source code of the UpDeFi, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

#### 1.1 About UpDeFi

UpDeFi is a yield farming aggregator running on Binance Smart Chain (BSC), with the goal of optimising DeFi users' yield farming at the lowest possible cost. The protocol has a number of built-in farming strategies and supports multiple farming pools (e.g., PancakeSwap, MarsSwap, etc). The protocol also has its utility token UP, which is distributed to protocol users according to their engagement or contribution.

ltem	Description
Target	UpDeFi
Туре	Solidity Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	January 28, 2022

Table 1.1: Basic Information of UpDeFi

In the following, we show the Git repository of reviewed files and the commit hash value used in this audit.

• https://github.com/up-defi/UpFarm-hardhat (d878a44)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

https://github.com/up-defi/UpFarm-hardhat (2adc051)

#### 1.2 About PeckShield

PeckShield Inc. [11] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

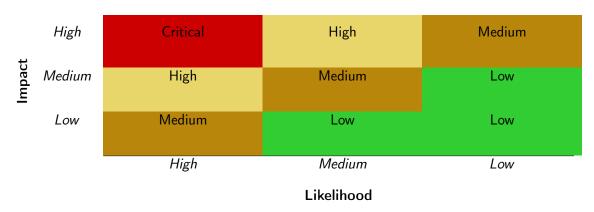


Table 1.2: Vulnerability Severity Classification

## 1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [10]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild:
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would

Table 1.3: The Full List of Check Items

Category	Check Item
	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
Basic Coding Bugs	Revert DoS
Dasic Coung Dugs	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
Advanced DeFi Scrutiny	Digital Asset Escrow
Advanced Berr Scrating	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
	Avoiding Use of Variadic Byte Array
Additional Recommendations	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [9], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings.

#### 1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during
	the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functional-
	ity that processes data.
Numeric Errors	Weaknesses in this category are related to improper calcula-
	tion or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like
	authentication, access control, confidentiality, cryptography,
	and privilege management. (Software security is not security
	software.)
Time and State	Weaknesses in this category are related to the improper man-
	agement of time and state in an environment that supports
	simultaneous or near-simultaneous computation by multiple
Forman Canadiai ana	systems, processes, or threads.
Error Conditions,	Weaknesses in this category include weaknesses that occur if
Return Values, Status Codes	a function does not generate the correct return/status code, or if the application does not handle all possible return/status
Status Codes	codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper manage-
Nesource Management	ment of system resources.
Behavioral Issues	Weaknesses in this category are related to unexpected behav-
Deliavioral issues	iors from code that an application uses.
Business Logics	Weaknesses in this category identify some of the underlying
Dusiness Togics	problems that commonly allow attackers to manipulate the
	business logic of an application. Errors in business logic can
	be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used
	for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of
	arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written
	expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices
	that are deemed unsafe and increase the chances that an ex-
	ploitable vulnerability will be present in the application. They
	may not directly introduce a vulnerability, but indicate the
	product has not been carefully developed or maintained.

# 2 | Findings

#### 2.1 Summary

Here is a summary of our findings after analyzing the UpDeFi implementation. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logic, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings
Critical	0
High	0
Medium	4
Low	3
Informational	0
Undetermined	0
Total	7

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities that need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

#### 2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 4 medium-severity vulnerabilities, and 3 low-severity vulnerabilities.

ID Title Severity Category **Status** PVE-001 Possible Costly share From Improper De-Time and State Medium Fixed posit Initialization In Strategy **PVE-002** Low Timely massUpdatePools In Up-**Business Logic** Fixed Farm::updateUPPerBlock() Potential Sandwich/MEV Attack For -**PVE-003** Confirmed Low Time and State safeSwap() **PVE-004** Medium Suggested Forbidden Transfer For Stak-Fixed Business Logic ingRewards Token Duplicate Pool Detection and Preven-**PVE-005** Low **Business Logic** Fixed tion **PVE-006** Medium Confirmed Trust Issue Of Admin Keys Security Features **PVE-007** Medium Minimum Delay Bypass in Timelock-Time and State Fixed Controller

Table 2.1: Key UpDeFi Audit Findings

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

# 3 Detailed Results

# 3.1 Possible Costly *share* From Improper Deposit Initialization In Strategy

• ID: PVE-001

Severity: MediumLikelihood: LowImpact: High

• Target: UpFarm/Strategy

Category: Time and State [6]CWE subcategory: CWE-362 [2]

#### Description

By design, the UpFarm contract is one of the main entries for interaction with users. It organizes a number of farming pools into which supported assets can be staked and implements an incentive mechanism that rewards the staking of different farming pools with the UP token. The Strategy contract implements the standard farming strategy, while the StrategyPCS and StrategyMars contracts inheriting from it implement the specific farming strategies.

In particular, one routine, i.e., Strategy::deposit(), is called inside the UpFarm::deposit() routine when the user stakes his assets, and its returned value representing the pool shares is internally recorded in the user.shares in the UpFarm contract. While examining the share calculation with the given stakes in the Strategy::deposit() routine, we notice an issue that may unnecessarily make the pool share extremely expensive and bring hurdles (or even causes loss) for later stakers. Specifically, the issue occurs when the pool is being initialized under the assumption that the current pool is empty.

```
93 function deposit( uint256 _wantAmt)
94 external
95 virtual
96 onlyOwner
97 nonReentrant
98 whenNotPaused
99 returns (uint256)
```

```
100
             uint256 beforeAmount = IERC20(wantAddress).balanceOf(address(this));
101
102
             IERC20(wantAddress).safeTransferFrom(
103
                 address (msg.sender),
104
                 address(this),
                 _{\tt wantAmt}
105
106
             );
107
             uint256 afterAmount = IERC20(wantAddress).balanceOf(address(this));
109
             uint256 realamount = afterAmount.sub(beforeAmount);
111
             uint256 sharesAdded = realamount;
112
             if (wantLockedTotal > 0 && sharesTotal > 0) {
113
                 sharesAdded = realamount
114
                     .mul(sharesTotal)
115
                     .mul(entranceFeeFactor)
116
                     .div(wantLockedTotal)
117
                     .div(entranceFeeFactorMax);
118
             }
119
             sharesTotal = sharesTotal.add(sharesAdded);
121
             if (isAutoComp) {
                 _farm();
122
123
             } else {
124
                 wantLockedTotal = wantLockedTotal.add(realamount);
125
126
             emit Deposit(msg.sender,realamount);
128
             return sharesAdded;
129
        }
131
         function farm() external virtual nonReentrant {
132
             _farm();
133
135
         function _farm() internal virtual {
136
             require(isAutoComp, "!isAutoComp");
137
             uint256 wantAmt = IERC20(wantAddress).balanceOf(address(this));
138
             wantLockedTotal = wantLockedTotal.add(wantAmt);
139
             IERC20(wantAddress).safeIncreaseAllowance(farmContractAddress, wantAmt);
141
             if (isCAKEStaking) {
142
                 IPancakeswapFarm(farmContractAddress).enterStaking(wantAmt);
143
             } else {
144
                 IPancakeswapFarm(farmContractAddress).deposit(pid, wantAmt);
145
             }
146
```

Listing 3.1: Strategy::deposit()&&farm()

Specifically, when the pool is being initialized, the share value directly takes the value of realamount (line 111), which is under control by the malicious actor. As this is the first stake, the current total

supply equals the calculated uint256 sharesAdded = realamount = 1WEI. With that, the actor can further transfer a huge amount of wantAddress tokens to Strategy contract with the goal of making per share extremely expensive.

An extremely expensive share can be very inconvenient to use as a small number of 1WEI may denote a large value. Furthermore, it can lead to precision issue in truncating the computed share value for staked assets. If truncated to be zero, the staked assets are essentially considered dust and kept by the pool with returning 0 that is internally recorded in the user.shares as stake credits.

This is a known issue that has been mitigated in popular Uniswap. When providing the initial liquidity to the contract (i.e. when totalSupply is 0), the liquidity provider must sacrifice 1000 LP tokens (by sending them to address(0)). By doing so, we can ensure the granularity of the LP tokens is always at least 1000 and the malicious actor is not the sole holder. This approach may bring an additional cost for the initial liquidity provider, but this cost is expected to be low and acceptable.

**Recommendation** Revise current execution logic of Strategy::deposit() to defensively calculate the share amount when the pool is being initialized. An alternative solution is to ensure guarded launch that safeguards the first stake to avoid being manipulated.

**Status** The issue has been addressed by the following commit: 909c967.

## 3.2 Timely massUpdatePools In UpFarm::updateUPPerBlock()

• ID: PVE-002

• Severity: Low

• Likelihood: Low

Impact: Low

• Target: UpFarm

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

The UpFarm contract implements an incentive mechanism that rewards the staking of supported assets with the UP token. The rewards are carried out by designating a number of staking pools. The staking users are rewarded in proportional to their staking assets in the pool.

The reward rate (per block) of the UP token can be adjusted via the updateUPPerBlock() routine. When analyzing its logic, we notice the lack of timely invoking massUpdatePools() to update the accUPPerShare and lastRewardBlock variables before the new reward-related configuration becomes effective. If the call to massUpdatePools() is not immediately invoked before updating the reward rate, certain situations may be crafted to create an unfair reward distribution.

```
function updateUPPerBlock(uint256 _upPerBlock) external onlyOwner {

UPPerBlock = _upPerBlock;
```

```
emit UpdateUPPerBlock(msg.sender,_upPerBlock);
325
}
```

Listing 3.2: UpFarm::updateUPPerBlock()

**Recommendation** Timely invoke massUpdatePools() in the updateUPPerBlock() routine. **Status** The issue has been addressed by the following commit: fbed970.

## 3.3 Potential Sandwich/MEV Attack For safeSwap()

• ID: PVE-003

Severity: Low

Likelihood: Low

Impact: Low

• Target: StrategyMars/StrategyPCS/Strategy

• Category: Time and State [8]

• CWE subcategory: CWE-682 [3]

#### Description

While examining the Strategy contract, we notice there is a routine (i.e., \_safeSwap()) that can be improved with effective slippage control. To elaborate, we show below the related code snippet of the Strategy contract. According to the design, the \_safeSwap() function is used to swap a certain token (specified by the input \_path) to another one. In the function, the swapExactTokensForTokensSupporting \_FeeOnTransferTokens() function of PancakeSwap/MarsSwap is called (lines 515 - 522) to swap the exact token to another one. However, we observe the second input amountOutMin parameter is calculated according to the current state of the PancakeSwap/MarsSwap pool (line 511), which may have been price-manipulated. In other words, the slippage control is ineffective and is therefore vulnerable to possible front-running attacks.

```
503
       function _safeSwap(
504
           address _uniRouterAddress,
505
           uint256 _amountIn,
506
           uint256 _slippageFactor,
           address[] memory _path,
507
508
           address _to,
509
           uint256 _deadline
510
       ) internal virtual {
511
           uint256[] memory amounts =
512
               IPancakeRouter02(_uniRouterAddress).getAmountsOut(_amountIn, _path);
513
           uint256 amountOut = amounts[amounts.length.sub(1)];
515
           IPancakeRouter02(_uniRouterAddress)
516
               517
```

```
518 amountOut.mul(_slippageFactor).div(1000),
519 _path,
520 _to,
521 _deadline
522 );
523 }
```

Listing 3.3: Strategy::\_safeSwap()

Note other routines, i.e, StrategyMars::earn() and Strategy::earn(), that use addLiquidity() also lack necessary slippage control.

Recommendation Improve the above-mentioned routine by adding effective slippage control.

**Status** The issue has been confirmed by the team. The team decides to leave it as is after the risk assessment and intends to use oracle to add effective slippage control in the future version.

#### 3.4 Suggested Forbidden Transfer For StakingRewards Token

• ID: PVE-004

Severity: MediumLikelihood: Medium

Impact: Medium

• Target: StakingRewards

Category: Business Logic [7]CWE subcategory: CWE-841 [4]

## Description

By design, the StakingRewards contract provides an incentive mechanism that rewards the staking of supported assets with the UP token. Specially, when the user deposits the UP token with the calling of deposit(), the LP token (i.e., "UP Farms Seed Token") will be minted to the user to represent the pool shares. Additionally, the StakingRewards contract supports all the standard ERC20 interfaces (including transfer()/transferFrom()) since it inherits from the standard ERC20 contract. In other words, the LP token can be transferred like the standard ERC20 token. Based on this, we notice there is a potential vulnerability that may result in the withdrawal failure.

To elaborate, we show below the related code snippet of the StakingRewards contract. In the deposit() function, the following statement is executed to record the user's deposit amount: user. amount = user.amount.add(realAmount), and at the same time the same amount of LP token will be minted. In the withdraw() function, the user.amount will be subtracted from the withdrawal amount of the token and the same withdrawal amount of LP token will be burned. This is reasonable under the assumption that the vault's internal asset balances (i.e., user.amount) are always consistent with actual token balances maintained in individual ERC20 token contracts. However, we notice the transfer() interface of the StakingRewards contract is inherited from the standard ERC20 contract,

which only maintains the LP token balances. If we assume Alice transfers the LP token to Bob, both Alice and Bob cannot withdraw the deposit UP token because of the inconsistency between the internal asset records (i.e., user.amount) and LP token balances maintained in ERC20 token contracts. We suggest to override the \_transfer() interface to forbid LP token transfer.

```
234
        function deposit(uint256 _pid, uint256 _amount) external override validatePid(_pid)
            nonReentrantwhenNotPaused {
235
236
             if (_amount > 0) {
237
                 uint256 beforeAmount = pool.lpToken.balanceOf(address(this));
238
                 pool.lpToken.safeTransferFrom(
239
                     address (msg.sender),
240
                     address(this),
                     _amount
241
242
                );
243
                 uint256 afterAmount = pool.lpToken.balanceOf(address(this));
244
                 uint256 realAmount = afterAmount.sub(beforeAmount);
245
                 if (address(uptoken) == address(pool.lpToken)) {
246
                     _mint(msg.sender, realAmount);
247
                 }
248
                 user.amount = user.amount.add(realAmount);
249
            }
250
             user.rewardDebt = user.amount.mul(pool.accUPPerShare).div(1e12);
251
             emit Deposit(msg.sender, _pid, _amount);
252
        }
254
        // Withdraw LP tokens from StakingReward.
255
        function withdraw(uint256 _pid, uint256 _amount) external override validatePid(_pid)
             nonReentrant {
256
257
            if (_amount > 0) {
258
                 user.amount = user.amount.sub(_amount);
259
                 if (address(uptoken) == address(pool.lpToken)) {
260
                     _burn(msg.sender, _amount);
261
                 pool.lpToken.safeTransfer(address(msg.sender), _amount);
262
263
            }
264
             user.rewardDebt = user.amount.mul(pool.accUPPerShare).div(1e12);
265
             emit Withdraw(msg.sender, _pid, _amount);
266
```

Listing 3.4: StakingRewards::deposit()&&withdraw()

Recommendation Suggest to override the \_transfer() interface as above-mentioned.

Status The issue has been addressed by the following commit: fbed970.

#### 3.5 Duplicate Pool Detection and Prevention

• ID: PVE-005

Severity: LowLikelihood: Low

• Impact: Low

• Target: UpFarm

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

The UpFarm contract provides an incentive mechanism that rewards the staking of supported assets with the UP token. The rewards are carried out by designating a number of staking pools into which supported assets can be staked. Each pool has its allocPoint\*multiplier/totalAllocPoint share of scheduled rewards and the rewards for stakers are proportional to their share of tokens in the pool.

In current implementation, there are a number of concurrent pools that share the rewarded tokens and more can be scheduled for addition (via a proper governance procedure or moderated by a privileged account). To accommodate these new pools, the design has the necessary mechanism in place that allows for dynamic additions of new staking pools that can participate in being incentivized as well.

The addition of a new pool is implemented in add(), whose code logic is shown below. It turns out it did not perform necessary sanity checks in preventing a new pool with a duplicate token from being added. Though it is a privileged interface (protected with the modifier onlyOwner), it is still desirable to enforce it at the smart contract code level, eliminating the concern of wrong pool introduction from human omissions.

```
80
        function add(
81
            uint256 _allocPoint,
82
            IERC20 _want,
83
            address _strat
84
        ) external onlyOwner {
85
            massUpdatePools();
86
            uint256 lastRewardBlock =
87
                block.number > startBlock ? block.number : startBlock;
88
            totalAllocPoint = totalAllocPoint.add(_allocPoint);
89
            poolInfo.push(
90
                PoolInfo({
91
                     want: _want,
92
                     allocPoint: _allocPoint,
                     lastRewardBlock: lastRewardBlock,
93
94
                     accUPPerShare: 0,
95
                     strat: _strat
96
                })
97
            );
٩R
            emit Add(_allocPoint,address(_want),_strat);
```

**Recommendation** Detect whether the given pool for addition is a duplicate of an existing pool. The pool addition is only successful when there is no duplicate.

Status The issue has been addressed by the following commit: fbed970.

#### 3.6 Trust Issue Of Admin Keys

• ID: PVE-006

• Severity: Medium

• Likelihood: Medium

• Impact: Medium

• Target: Multiple Contracts

• Category: Security Features [5]

• CWE subcategory: CWE-287 [1]

#### Description

In the UpDeFi protocol, there is a privileged account that plays a critical role in governing and regulating the protocol-wide operations (e.g., mint UpToken infinitely). In the following, we show the representative functions potentially affected by the privilege of the account.

```
19
        function mint(address _to, uint256 _amount) external {
20
            require(hasRole(MINTER_ROLE, msg.sender), "Caller is not a minter");
21
            require(totalSupply().add(_amount) <= cap(), "Out Of the Cap");</pre>
22
            _mint(_to, _amount);
23
       }
24
25
        function increaseCap(uint256 _increaseNum) external {
26
            require(hasRole(GOVERNOR_ROLE, msg.sender), "Caller is not a Governor");
27
            _cap = _cap.add(_increaseNum);
28
29
            emit IncreaseCap(msg.sender,_increaseNum);
```

Listing 3.6: UPToken::mint()&&increaseCap()

We emphasize that the privilege assignment may be necessary and consistent with the protocol design. However, it is worrisome if the privileged account is not governed by a DAO-like structure. Note that a compromised privileged account would allow the attacker to modify a number of sensitive system parameters, which directly undermines the assumption of the UpDeFi design.

**Recommendation** Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks.

Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

Status This issue has been confirmed by the team. The team intends to introduce multisig and timelock mechanisms to mitigate this issue when the protocol is deployed on the mainnet. Additionally, DAO governance will be applied once the governance tokens are sufficiently distributed to the public.

## 3.7 Minimum Delay Bypass In TimelockController

ID: PVE-007Severity: MediumLikelihood: LowImpact: High

Target: TimelockControllerCategory: Time and State [8]CWE subcategory: CWE-682 [3]

#### Description

The TimelockController, introduced in OpenZeppelin Contracts 3.3, is a smart contract that enforces a delay on all actions directed towards an owned contract. A typical setup is to position the TimelockController as the admin of an application smart contract so, whenever a privileged action is to be executed, it has to wait for a certain time specified by the TimelockController.

The security benefits of the TimelockController are twofold. Firstly, it provides an extra layer of security to a project's team by giving a heads up on every privileged action anticipated in the system. This allows the team to detect and react to malicious calls by compromised admin accounts. Secondly, it protects the community from the project's governance itself, allowing members to exit the protocol if they disagree with any impending changes.

In particular, scheduleBatch() (a schedule function) and executeBatch() (an execute function), allow the caller to enqueue and execute proposals that run multiple calls in sequence. However, there is a vulnerability in their implementation that can be exploited by the malicious EXECUTOR to execute arbitrary tasks bypassing the minimum delay protection.

To elaborate, we show below the related code snippet of the TimelockController contract. A malicious EXECUTOR could execute a batch with the calling of executeBatch(), including a set of calls, i.e., the call to the TimelockController itself to clear the minimum delay and grant PROPOSER and ADMIN rights to an address under their control, the call to scheduleBatch() to enqueue the batch by the controlled PROPOSER and the call to the arbitrary privileged functions under the TimelockController control. By doing so, the malicious EXECUTOR effectively takes full control of the TimelockController contract.

```
318
         function executeBatch(
319
             address[] calldata targets,
             uint256[] calldata values,
320
321
             bytes[] calldata datas,
322
             bytes32 predecessor,
323
             bytes32 salt
324
         ) external payable virtual onlyRole(EXECUTOR_ROLE) {
325
             require(
326
                 targets.length == values.length,
327
                 \verb|'TimelockController: length mismatch|'
328
             );
329
             require(
330
                 targets.length == datas.length,
331
                 "TimelockController: length mismatch"
332
             );
334
             bytes32 id =
335
                 hashOperationBatch(targets, values, datas, predecessor, salt);
336
             _beforeCall(predecessor);
337
             for (uint256 i = 0; i < targets.length; ++i) {</pre>
338
                 _call(id, i, targets[i], values[i], datas[i]);
339
             }
340
             _afterCall(id);
341
```

Listing 3.7: TimelockController::executeBatch()

```
224
         function scheduleBatch(
225
             address[] calldata targets,
226
             uint256[] calldata values,
227
             bytes[] calldata datas,
228
             bytes32 predecessor,
229
             bytes32 salt,
230
             uint256 delay
231
         ) external virtual onlyRole(PROPOSER_ROLE) {
232
             require(
233
                 targets.length == values.length,
234
                 "TimelockController: length mismatch"
235
             );
236
             require(
237
                 targets.length == datas.length,
238
                 "TimelockController: length mismatch"
239
             );
241
             bytes32 id =
242
                 hashOperationBatch(targets, values, datas, predecessor, salt);
243
             _schedule(id, delay);
244
             for (uint256 i = 0; i < targets.length; ++i) {</pre>
245
                 emit CallScheduled(
246
                     id,
247
                      i,
248
                      targets[i],
249
                      values[i],
```

```
250 datas[i],
251 predecessor,
252 delay
253 );
254 }
255 }
```

Listing 3.8: TimelockController::scheduleBatch()

**Recommendation** Considering the OpenZeppelin team has solved this vulnerability, we suggest to upgrade the TimelockController to the new version.

**Status** The issue has been addressed by the following commit: e698186.



# 4 Conclusion

In this audit, we have analyzed the UpDeFi design and implementation. UpDeFi is a yield farming aggregator running on Binance Smart Chain (BSC), with the goal of optimising DeFi users' yield farming at the lowest possible cost, which provides a number of built-in farming strategies and supports multiple farming pools (e.g., PancakeSwap, MarsSwap, etc). The current code base is well organized and those identified issues are promptly confirmed and fixed.

Meanwhile, we need to emphasize that smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



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