

NANJING UNIVERSITY

ACM-ICPC Codebook 0 Miscellaneous

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1 General

1.1 vimre

```
set nocompatible
 1
    syntax on
 2
    colorscheme slate
 3
    set number
 4
    set cursorline
 5
    set shiftwidth=2
 7
    set softtabstop=2
    set tabstop=2
 9
    set expandtab
    set magic
10
    set smartindent
11
12
    set backspace=indent,eol,start
    set cmdheight=1
13
14
    set laststatus=2
    set statusline=\ %<%F[\%1*\%\%*\%n\%R\%H]%=\ %y\ %0(%{&fileformat}\ %{&encoding}\ %c
15
       :%1/%L%)\
    set whichwrap=b,s,<,>,[,]
16
```

1.2 bashrc

```
mkdir -p ~/.trash
    alias rm=trash
2
3
   trash()
4
    {
5
      mv $@ ~/.trash/
6
7
    cleartrash()
8
9
       \rm -rvf ~/.trash
10
       mkdir -p ~/.trash
11
12
    }
```

1.3 runbash

```
1  if [ $# -ge 1 ]; then
2  fn=$1
3  echo ${fn} > .run.log
```

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```
4
    else
5
      fn=`cat .run.log`
   fi
6
7
8
   # cat $fn.cpp | xsel -ib
9
    if g++ $fn.cpp -std=c++11 -D LOCAL DEBUG -Wall -O2 -g -o $fn; then
10
      echo "****** Compilation Success! ****** [$fn]"
11
12
      if [ $# -ge 2 ]; then
13
       time -f "\n%U user, %S system, %e real" ./$fn < $2
14
      else
        time -f "\n%U user, %S system, %e real" ./$fn
15
      fi
16
17
    # cat $fn.cpp | xsel -ib
18
    else
      echo "****** Compilation Failed! ****** [$fn]"
19
20
   fi
```

1.4 Template

```
#include <bits/stdc++.h>
 1
    using namespace std;
 2
 3
 4
    #ifdef LOCAL DEBUG
    # define debug(fmt, ...) fprintf(stderr, "\033[94m%s:__" fmt "\n\033[0m", \
 5
         func , ## VA ARGS )
 6
 7
    #else
    # define debug(...) ((void) 0)
 8
    #endif
 9
    #define rep(i, n) for (int i=0; i<(n); i++)
10
    #define Rep(i, n) for (int i=1; i<=(n); i++)</pre>
11
12
    #define range(x) (x).begin(), (x).end()
13
    typedef long long LL;
14
    typedef unsigned long long ULL;
15
    template <unsigned p>
16
    struct Zp{
17
18
        unsigned x;
        Zp(unsigned x):x(x){}
19
20
        operator unsigned(){return x;}
        Zp operator ^ (ULL e) {
21
            Zp b=x, r=1;
22
23
            while (e) {
                if (e&1) r=r*b;
24
25
                b=b*b;
```

```
26
                e>>=1;
27
            }
28
            return r;
29
        Zp operator + (Zp rhs) {return (x+rhs)%p;}
30
        Zp operator - (Zp rhs) {return (x+p-rhs)%p;}
31
        Zp operator * (Zp rhs) {return x*rhs%p;}
32
33
        Zp operator / (Zp rhs) {return Zp(x)*(rhs^{(p-2)});}
34
    };
35
    typedef Zp<1000000007> zp;
36
37
    zp operator"" (ULL n){return n;}
38
```

2 String

2.1 Knuth-Morris-Pratt algorithm

Single-pattern matching.

Usage:

```
construct(p) Construct the failure table of pattern p.
match(t, p) Match pattern p in text t.
found(pos) Report the pattern found at pos.
```

Time complexity: O(l).

```
const int SIZE = 10005;
 1
    int fail[SIZE];
 2
 3
    int len;
 4
    void construct(const char* p){
 5
        len = strlen(p);
 6
 7
        fail[0] = fail[1] = 0;
 8
        for (int i = 1; i < len; i++) {</pre>
 9
             int j = fail[i];
10
            while (j && p[i] != p[j]) j = fail[j];
            fail[i+1] = p[i] == p[j] ? j+1 : 0;
11
12
        }
13
    }
14
15
    inline void found(int pos){
        //! add codes for having found at pos
16
17
    }
```

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```
18
19
    void match(const char* t, const char* p){ // must be called after construct
        int n = strlen(t);
20
        int j = 0;
21
        rep (i, n){
22
23
            while (j && p[j] != t[i]) j = fail[j];
            if (p[j] == t[i]) j++;
24
25
            if (j == len) found(i - len + 1);
26
        }
27
    }
```

2.2 Trie

Support insertion and search for a set of words.

- △ If duplicate word exists, only the last one is preserved.
- \triangle The tag must not be 0, which is considered as not being a word.

Usage:

```
    id(c) Covert character to its id.
    add(s, t) Add word s into Trie, where t is the tag attached to s.
    search(s) Search for word s. Return the tag attached to s if found; otherwise return 0.
```

Time complexity: $O(l|\Sigma|)$ for insertion, O(l) for search.

```
const int MAXN = 12000;
 1
 2
    const int CHARN = 26;
 3
 4
    inline int id(char c){
 5
        return c - 'a';
 6
    }
 7
 8
    struct Trie{
 9
        int n;
        int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
10
        int tag[MAXN];
11
12
        Trie(){
13
            memset(tr[0], 0, sizeof(tr[0]));
14
            tag[0] = 0; n = 1;
15
        }
16
17
18
        // tag should not be 0
19
        void add(const char* s, int t){
            int p = 0, c, len = strlen(s);
20
```

```
21
            rep (i, len){
22
                 c = id(s[i]);
23
                 if (!tr[p][c]){
                     memset(tr[n], 0, sizeof(tr[n]));
24
25
                     tag[n] = 0;
                     tr[p][c] = n++;
26
27
                 p = tr[p][c];
28
29
30
            tag[p] = t;
        }
31
32
        // returns 0 if not found
33
        // AC automaton does not need this function
34
        int search(const char* s){
35
            int p = 0, c, len = strlen(s);
36
            rep (i, len){
37
38
                 c = id(s[i]);
                 if (!tr[p][c]) return 0;
39
40
                 p = tr[p][c];
41
42
            return tag[p];
43
        }
44
    };
```

2.3 Aho-Corasick automaton

Automaton for multi-pattern matching.

 \triangle See the warnings of Trie.

△ If a word has too many suffixes, the automaton might run slow.

Usage:

```
add(s, t) Add word s into Trie, where t is the tag attached to s.

Construct() Construct the automaton after all words added.

Find words in text.

Found(pos, j) Report a word found in node j, the last character of which is at pos.
```

Requirement:

2.2 Trie

Time complexity: $O(l|\Sigma|)$ for insertion and construction, O(l) for finding, provided the number of suffixes of a word is constant.

```
1 | struct AC : Trie{
```

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```
int fail[MAXN];
 2
 3
        int last[MAXN];
 4
 5
        void construct(){
            queue<int> q;
 6
 7
            fail[0] = 0;
            rep (c, CHARN){
 8
 9
                 if (int u = tr[0][c]){
                     fail[u] = 0;
10
11
                     q.push(u);
                     last[u] = 0;
12
                 }
13
14
            }
15
            while (!q.empty()){
                 int r = q.front(); q.pop();
16
17
                 rep (c, CHARN){
                     int u = tr[r][c];
18
                     if (!u){
19
                         tr[r][c] = tr[fail[r]][c];
20
                         continue;
21
22
                     }
23
                     q.push(u);
24
                     int v = fail[r];
                     while (v && !tr[v][c]) v = fail[v];
25
                     fail[u] = tr[v][c];
26
                     last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
27
                 }
28
29
            }
        }
30
31
32
        void found(int pos, int j){
            if (j) {
33
                 //! add codes for having found word with tag[j]
34
35
                 found(pos, last[j]);
36
            }
37
        }
38
        void find(const char* text){ // must be called after construct()
39
40
            int p = 0, c, len = strlen(text);
            rep (i, len){
41
                 c = id(text[i]);
42
43
                 p = tr[p][c];
44
                 if (tag[p])
45
                     found(i, p);
                 else if (last[p])
46
                     found(i, last[p]);
47
            }
48
```

```
49 | }
50 |};
```

3 Game Theory

3.1 Nim games

以下游戏中,不能动的算输。

3.1.1 Bash game

有n个石子,每人最多拿m个,最少拿1个。 $n \mod (m+1) \neq 0$ 时先手必胜。

3.1.2 Fibonacci nim

有 n 个石子,第一轮可以拿不超过 n 个石子。此后,每次拿的石子数不超过前一次的 2 倍。当 n 是斐波那契数时先手必胜。

3.1.3 Wythoff's game

有 2 堆石子,分别有 a,b 个 $(a \le b)$,每人可以从一堆中拿任意多个,或从两堆中拿相同多个。当 $a = \lfloor (b-a) \frac{\sqrt{5}+1}{2} \rfloor$ 时先手必败。

4 Others

4.1 Fast Fourier transform

 \triangle The size of the sequence must be some power of 2.

 \triangle When performing convolution, the size of the sequence should be doubled. To compute k, one may call 32-__builtin_clz(a+b-1), where a and b are the lengths of two sequences.

Usage:

FFT(k) Initialize the structure with maximum sequence length 2^k .

fft(a) Perform Fourier transform on sequence a.

ifft(a) Perform inverse Fourier transform on sequence a.

conv(a, b) Convolve sequence a with b.

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Time complexity: $O(n \log n)$ for fft, ifft and conv.

```
const int NMAX = 1<<20;</pre>
 1
 2
    typedef complex<double> cplx;
 3
    const double PI = 2*acos(0.0);
    struct FFT{
 4
 5
        int rev[NMAX];
 6
        cplx omega[NMAX], oinv[NMAX];
 7
        int K, N;
 8
        FFT(int k){
 9
            K = k; N = 1 << k;
10
            rep (i, N){
11
                 rev[i] = (rev[i>1]>>1) | ((i&1)<<(K-1));
12
                omega[i] = polar(1.0, 2.0 * PI / N * i);
13
                oinv[i] = conj(omega[i]);
14
15
            }
16
        }
17
18
        void dft(cplx* a, cplx* w){
19
            rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
20
            for (int 1 = 2; 1 <= N; 1 *= 2)
                 int m = 1/2;
21
                 for (cplx* p = a; p != a + N; p += 1)
22
23
                     rep (k, m){
24
                         cplx t = w[N/1*k] * p[k+m];
                         p[k+m] = p[k] - t; p[k] += t;
25
26
                     }
27
            }
        }
28
29
        void fft(cplx* a){dft(a, omega);}
30
        void ifft(cplx* a){
31
32
            dft(a, oinv);
33
            rep (i, N) a[i] /= N;
        }
34
35
        void conv(cplx* a, cplx* b){
36
37
            fft(a); fft(b);
            rep (i, N) a[i] *= b[i];
38
39
            ifft(a);
40
        }
41
    };
```

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4.2 2-SAT

Usage:

```
init(n) Initialize the structure with at most n Boolean variables.

add_clause(x, Add clause: x = xval or y = yval.

xval, y, yval)

solve() Solve the 2-SAT problem. Return false if no solution.

value(i) Return the value of i-th variable in some solution, if exists.
```

Time complexity: O(m+n).

```
1
    const int MAXN = 100005;
 2
    struct twoSAT{
 3
        int n;
        vector<int> G[MAXN*2];
 4
 5
        bool mark[MAXN*2];
         int S[MAXN*2], c;
 6
 7
        void init(int n){
 8
 9
             this -> n = n;
             for (int i=0; i<n*2; i++) G[i].clear();</pre>
10
             memset(mark, 0, sizeof(mark));
11
         }
12
13
14
         bool dfs(int x){
15
             if (mark[x^1]) return false;
16
             if (mark[x]) return true;
             mark[x] = true;
17
             S[c++] = x;
18
             for (int i=0; i<G[x].size(); i++)</pre>
19
                 if (!dfs(G[x][i])) return false;
20
             return true;
21
         }
22
23
24
        void add clause(int x, bool xval, int y, bool yval){
             x = x * 2 + xval;
25
             y = y * 2 + yval;
26
             G[x^1].push_back(y);
27
28
             G[y^1].push back(x);
29
         }
30
         bool solve() {
31
             for (int i=0; i<n*2; i+=2){</pre>
32
                 if (!mark[i] && !mark[i+1]){
33
34
                      c = 0;
                      if (!dfs(i)){
35
                          while (c > 0) mark[S[--c]] = false;
36
```

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```
if (!dfs(i+1)) return false;
37
                    }
38
                }
39
            }
40
41
            return true;
42
        }
43
        inline bool value(unsigned i){return mark[2*i+1];}
44
45
    };
```