



NANJING UNIVERSITY

ACM-ICPC Codebook 0
Miscellaneous

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1 General

1.1 Template

```
1 #include <iostream>
2 #include <cstdio>
3 #include <cstring>
4 #include <climits>
5 #include <vector>
6 #include <stack>
7 #include <queue>
8 #include <string>
9 #include <algorithm>
10 using namespace std;
11
12 #define rep(i, n) for (int i = 0; i < (n); i++)
13 #define Rep(i, n) for (int i = 1; i <= (n); i++)
14 typedef long long LL;
15
16 int main(){
17
18     return 0;
19 }
```

2 String

2.1 Knuth-Morris-Pratt algorithm

Single-pattern matching.

Usage:

construct(p)	Construct the failure table of pattern p.
match(t, p)	Match pattern p in text t.
found(pos)	Report the pattern found at pos.

Time complexity: $O(l)$.

```
1 const int SIZE = 10005;
2 int fail[SIZE];
3 int len;
4
5 void construct(const char* p){
6     len = strlen(p);
```

```

7   fail[0] = fail[1] = 0;
8   for (int i = 1; i < len; i++) {
9       int j = fail[i];
10      while (j && p[i] != p[j]) j = fail[j];
11      fail[i+1] = p[i] == p[j] ? j+1 : 0;
12  }
13 }
14
15 inline void found(int pos){
16     // ! add codes for having found at pos
17 }
18
19 void match(const char* t, const char* p){ // must be called after construct
20     int n = strlen(t);
21     int j = 0;
22     rep (i, n){
23         while (j && p[j] != t[i]) j = fail[j];
24         if (p[j] == t[i]) j++;
25         if (j == len) found(i - len + 1);
26     }
27 }

```

2.2 Trie

Support insertion and search for a set of words.

- △ If duplicate word exists, only the last one is preserved.
- △ The tag must not be 0, which is considered as not being a word.

Usage:

<code>id(c)</code>	Covert character to its id.
<code>add(s, t)</code>	Add word <i>s</i> into Trie, where <i>t</i> is the tag attached to <i>s</i> .
<code>search(s)</code>	Search for word <i>s</i> . Return the tag attached to <i>s</i> if found; otherwise return 0.

Time complexity: $O(l|\Sigma|)$ for insertion, $O(l)$ for search.

```

1  const int MAXN = 12000;
2  const int CHARN = 26;
3
4  inline int id(char c){
5      return c - 'a';
6  }
7
8  struct Trie{
9      int n;

```

```

10  int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
11  int tag[MAXN];
12
13  Trie(){
14      memset(tr[0], 0, sizeof(tr[0]));
15      tag[0] = 0;
16      n = 1;
17  }
18
19  // tag should not be 0
20  void add(const char* s, int t){
21      int p = 0, c, len = strlen(s);
22      rep (i, len){
23          c = id(s[i]);
24          if (!tr[p][c]){
25              memset(tr[n], 0, sizeof(tr[n]));
26              tag[n] = 0;
27              tr[p][c] = n++;
28          }
29          p = tr[p][c];
30      }
31      tag[p] = t;
32  }
33
34  // returns 0 if not found
35  // AC automaton does not need this function
36  int search(const char* s){
37      int p = 0, c, len = strlen(s);
38      rep (i, len){
39          c = id(s[i]);
40          if (!tr[p][c]) return 0;
41          p = tr[p][c];
42      }
43      return tag[p];
44  }
45  };

```

2.3 Aho-Corasick automaton

Automaton for multi-pattern matching.

△ See the warnings of Trie.

△ If a word has too many suffixes, the automaton might run slow.

Usage:

<code>add(s, t)</code>	Add word s into Trie, where t is the tag attached to s .
<code>construct()</code>	Construct the automaton after all words added.
<code>find(text)</code>	Find words in text.
<code>found(pos, j)</code>	Report a word found in node j , the last character of which is at pos .

Requirement:

2.2 Trie

Time complexity: $O(l|\Sigma|)$ for insertion and construction, $O(l)$ for finding, provided the number of suffixes of a word is constant.

```

1 struct AC : Trie{
2     int fail[MAXN];
3     int last[MAXN];
4
5     void construct(){
6         queue<int> q;
7         fail[0] = 0;
8         rep (c, CHARN){
9             if (int u = tr[0][c]){
10                 fail[u] = 0;
11                 q.push(u);
12                 last[u] = 0;
13             }
14         }
15         while (!q.empty()){
16             int r = q.front(); q.pop();
17             rep (c, CHARN){
18                 int u = tr[r][c];
19                 if (!u){
20                     tr[r][c] = tr[fail[r]][c];
21                     continue;
22                 }
23                 q.push(u);
24                 int v = fail[r];
25                 while (v && !tr[v][c]) v = fail[v];
26                 fail[u] = tr[v][c];
27                 last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
28             }
29         }
30     }
31
32     void found(int pos, int j){
33         if (j) {
34             // ! add codes for having found word with tag[j]
35             found(pos, last[j]);
36         }

```

```
37     }
38
39     void find(const char* text){ // must be called after construct()
40         int p = 0, c, len = strlen(text);
41         rep (i, len){
42             c = id(text[i]);
43             p = tr[p][c];
44             if (tag[p])
45                 found(i, p);
46             else if (last[p])
47                 found(i, last[p]);
48         }
49     }
50 };
```