

## TOP CRISP RELATIONS

26/9/24

①

The concept of relations between sets is build on the cartesian products operator of sets.

### CARTESIAN PRODUCT

of two sets  $A \& B$  denoted by  $A \times B$  is a set of all ordered pairs such that first element in the pair belongs to  $A$  & the second element belongs to  $B$ .

$$i.e. A \times B = \{ (a, b) \mid a \in A, b \in B \}$$

if  $A \neq B$  &  $A \& B$  are non-empty then  $A \times B \neq B \times A$

generalize to  $n$  number of sets

$$\underset{i=1}{\overset{n}{\times}} A_i = \{ (a_1, a_2, a_3, \dots, a_n) \mid a_i \in A_i \text{ for every } i=1, 2, \dots, n \}$$

observe  
Cardinality  $\left| \underset{i=1}{\overset{n}{\times}} A_i \right| = \prod_{i=1}^n |A_i|$

Example  $A_1 = \{a, b\}, A_2 = \{1, 2\}, A_3 = \{\alpha\}$

$$A_1 \times A_2 = \{(a, 1), (a, 2), (b, 1), (b, 2)\}$$

$$|A_1 \times A_2| = 4 \triangleq |A_1| = 2, |A_2| = 2$$

$$\therefore |A_1 \times A_2| = |A_1| \cdot |A_2|$$

From  $A_1 \times A_2 \times A_3$

(2)

### RELATIONS

An n-ary relation denoted as  $R(x_1, x_2, x_3, \dots, x_n)$

among crisp sets  $X_1, X_2, X_n$  is a subset of the Cartesian product and is indicator of an association or relate among the tuple elements

For  $n=2$  the Reln  $R(x_1, x_2)$  is termed binary.

3	ternary
4	quaternary
	quinary

$\therefore$  if the univer. of discourse or sets are finite, the 'n-ary' rel can be expr. as an  $n$ -dim relate matrx. In f brn Rel  $R(x, y)$

$$X = \{x_1, x_2, \dots, x_n\} \quad \& \quad Y = \{y_1, y_2, \dots, y_m\}$$

Relate matrx  $R$  is a two dimen. matrx where  $X$  repr. rows &  $Y$  repr. columns.

$$X = \{1, 2, 3, 4\}$$

$$\therefore X \times X = \{(1,1), (1,2), (1,3), (1,4), \dots, (4,4)\} \quad |X \times X| = 16$$

Let the relation  $R$  be defn as

$$R = \{(x,y) \mid y = x+1, x, y \in X\}$$

$$= \{(1,2), (2,3), (3,4)\}$$

The relat matrx  $R$  is given by

$$R = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

### Operations on Relations

③

Given two relation  $R$  &  $S$  defn on  $X \times Y$  represented by relation matrix  
the follow oper are support by  $R$  &  $S$

UNION  $R \cup S$

$$R \cup S(x, y) = \max(R(x, y), S(x, y))$$

Intersection  $R \cap S$

$$R \cap S(x, y) = \min(R(x, y), S(x, y))$$

Complm  $\bar{R}$

$$\bar{R}(x, y) = 1 - R(x, y)$$

### COMPOSITION OF RELATIONS $R \circ S$

Given  $R$  to be a relation on  $X, Y$  &  $S$  to be relation on  $Y, Z$

then  $R \circ S$  is a composition of relation on  $X, Z$  defined as.

$$R \circ S = \left\{ (x, z) \mid (x, y) \in X \times Z, \exists y \in Y \text{ such that } (x, y) \in R \text{ & } (y, z) \in S \right\}$$

↓                              ↓  
 refers to                      there  
 exist

### MIN MAX composition

given the set of relations of  $R$  &  $S$  the min-max composite  $T$  is

defined as  $T = R \circ S$

$$T(x, z) = \max_{y \in Y} (\min(R(x, y), S(y, z)))$$

at  $j$

### Example

Let  $R, S$  be defined on the set  $\{1, 3, 5\} \times \{1, 3, 5\}$ .

Let  $R: \{(x, y) \mid y = x+2\}$      $S: \{(x, y) \mid x < 4\}$

$$R = \{(1, 3), (3, 5)\}, \quad S = \{(1, 3), (1, 5), (3, 5)\}$$

Now

$$R: \begin{matrix} 1 & 3 & 5 \\ \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

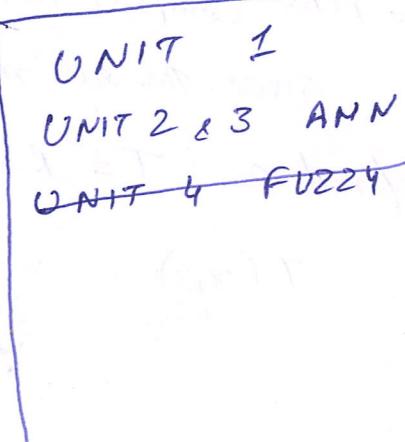
$$S: \begin{matrix} 1 & 3 & 5 \\ \begin{bmatrix} 1 & 1 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

Using min-max comp.

$$R \circ S = \begin{matrix} 1 & 3 & 5 \\ \begin{bmatrix} 0 & 0 & 1 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$$R \circ S = \{(1, 5)\}$$

$$\text{Also } S \circ R = \{(1, 5)\}$$



(1)

Fuzzy Relations is a fuzzy set defined on the Cartesian product of crisp set  $X_1, X_2 \dots X_n$  where the n-tuples  $(x_1, x_2, \dots x_n)$  may have a varying degree of membership within the relation. The membership value indicates the strength of the relation between the tuples.

Let  $R$  be the fuzzy relation between two sets  $X_1 \text{ & } X_2$

where  $X_1$  is the set of diseases &

$X_2$  is the set of symptoms.

$$X_1 = \{\text{Typhoid, viral fever, common cold}\} \quad \equiv X$$

$$X_2 = \{\text{runny nose, high temperature, shivering}\} \quad \equiv Y$$

The fuzzy relation  $R$  may be defined as

	Runny Nose	High Temp	Shivering
Typhoid	0.1	0.9	0.8
Viral fever	0.2	0.9	0.7
Common cold	0.9	0.4	0.6

## Fuzzy Cartesian Product

Let  $\tilde{A}$  be a fuzzy set defined on universe  $X$

$$\tilde{B} \quad " \quad " \quad " \quad " \quad Y$$

The cartesian product of  $\tilde{A} \times \tilde{B}$  under by  $\tilde{A} \times \tilde{B} \subseteq$  resulting a fuzzy relation  $R$  is given by

$$\tilde{R} = \tilde{A} \times \tilde{B} \subseteq X \times Y$$

where  $\tilde{R}$  has its membership function given by

$$\mu_{\tilde{R}}(x, y) = \mu_{\tilde{A} \times \tilde{B}}$$

Example page 206

Let  $\tilde{A} = \{(x_1, 0.2), (x_2, 0.7), (x_3, 0.4)\}$  &  $\tilde{B} = \{(y_1, 0.5), (y_2, 0.6)\}$ . Let  $\tilde{A}$  &  $\tilde{B}$  be the sets defined on the universes of discourse  $X = \{x_1, x_2, x_3\}$  &  $Y = \{y_1, y_2\}$ , respectively, then the fuzzy relation  $R$  as an element of the power set  $\mathcal{P}(X \times Y)$  is given by

$$\tilde{R} = \tilde{A} \times \tilde{B} = \begin{array}{cc|cc} & y_1 & y_2 \\ x_1 & 0.2 & 0.2 \\ x_2 & 0.5 & 0.6 \\ x_3 & 0.4 & 0.4 \end{array}$$

Since

$$\tilde{R}(x_1, y_1) = \min(\mu_{\tilde{A}}(x_1), \mu_{\tilde{B}}(y_1)) = \min(0.2, 0.5) = 0.2$$

$$\tilde{R}(x_1, y_2) = 0.2$$

Operations on Fuzzy Relation

Let  $R$  &  $S$  be two relations on  $X \times Y$

UNION

$$\mu_{R \cup S}(x, y) = \max(\mu_R(x, y), \mu_S(x, y))$$

INTERSECTION

$$\mu_{R \cap S}(x, y) = \min(\mu_R(x, y), \mu_S(x, y))$$

COMPLEMENT

$$\mu_{R^c}(x, y) = 1 - \mu_R(x, y)$$

Compos of relat = same as crisp set

MIN-MAX COMPOS

$$\mu_{R \circ S}(x, z) = \max_{y \in Y} (\min(\mu_R(x, y), \mu_S(y, z)))$$

Example

for  $X = \{x_1, x_2, x_3\}$ ,  $Y = \{y_1, y_2\}$ ,  $Z = \{z_1, z_2, z_3\}$ .

Let  $\tilde{R} = \begin{matrix} & y_1 & y_2 \\ x_1 & \begin{bmatrix} 0.5 & 0.1 \\ 0.2 & 0.9 \end{bmatrix} \\ x_2 & \\ x_3 & \begin{bmatrix} 0.3 & 0.6 \end{bmatrix} \end{matrix}$

$\tilde{S} = \begin{matrix} & z_1 & z_2 & z_3 \\ y_1 & \begin{bmatrix} 0.6 & 0.4 & 0.7 \end{bmatrix} \\ y_2 & \begin{bmatrix} 0.5 & 0.8 & 0.9 \end{bmatrix} \end{matrix}$

then  $R \circ S$  by max-min composition yields

$$R \circ S = \begin{matrix} & z_1 & z_2 & z_3 \\ x_1 & \begin{bmatrix} 0.5 & 0.4 & 0.5 \end{bmatrix} \\ x_2 & \begin{bmatrix} 0.5 & 0.8 & 0.9 \end{bmatrix} \\ x_3 & \begin{bmatrix} 0.6 & 0.6 & 0.7 \end{bmatrix} \end{matrix}$$

$$\begin{aligned} M_{R \circ S}^{nn}(x_1, z_1) &= \max \left( \min(R(x_1, y_1), S(y_1, z_1)), \min(R(x_1, y_2), S(y_2, z_1)) \right) \\ &= \max(\min(0.5, 0.6), \min(0.1, 0.5)) \\ &= \max(0.5, 0.1) \\ &= 0.5 \end{aligned}$$







# Fuzzy Logic & Inference

3/10/24 ①

## CRISP LOGIC (Proposition)

Consider Statement "Water boils at 90°C" True  
False  
 "Sky is blue" an agreement or disagreement  
can be made by "True" or "False"

a statement which is either "true" or "false" but not both is called proposition (It is denoted by upper case letter as P, Q, R, S...)

Examp P: Water boils at 90°C      }  
 Q: Sky is blue.      } Both are proposition (or atoms)

To represent real world complex inform several of propositions linked using  
connectives or operators      connectives or operators

### Proposition logic connectives

Symbol	Connective	Usage	Description
(may) $\wedge$	and	$P \wedge Q$	P and Q are true
$\vee$	or	$P \vee Q$	Either P or Q <del>is</del> is true
(may) $\neg$ or $\sim$	not	$\neg P \equiv \top P$	P is not true
$\Rightarrow$	implies	$P \Rightarrow Q$	P implies Q is true
=	equivalent	$P = Q$	P and Q are equal - same value (in truth value)

P	Q	$P \wedge Q$	$P \vee Q$	$\neg P$	$P \Rightarrow Q$	$P = Q$
T	T	T	T	F	T	T
T	F	F	T	F	F	F
F	T	F	T	T	T	F
F	F	F	F	T	T	T

## Interpretation Tautology contradiction

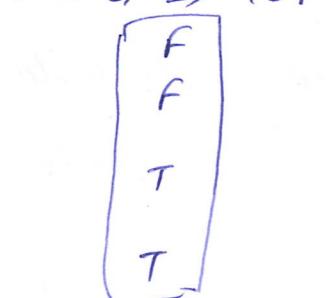
2

A logical formula comprising 'n' propositions will have  $2^n$  interpretations in its truth table (TT).

A formula which has all its entries records true is known as tautology & one which records false for all its entries is known as contradiction.

Q. Obtain a TT for the form  $(P \vee Q) \Rightarrow (\sim P)$ . Is it a tautology?

P	Q	$P \vee Q$	$\sim P$	$P \vee Q \Rightarrow \sim P$
T	T	T	F	
T	F	T	F	
F	T	T	T	
F	F	F	T	



not all are True

it is not a tautology

$$Q. H. \underbrace{((P \Rightarrow Q) \wedge (Q \Rightarrow P))}_{A:} = \underbrace{(P = Q)}_{B:} \text{ a tautology}$$

$$\begin{array}{ccccccc}
\cancel{\overline{P}} & Q & P \Rightarrow Q & Q \Rightarrow P & \underbrace{(P \Rightarrow Q) \wedge (Q \Rightarrow P)}_{A:} & \cancel{\overline{P=Q}} & A = B \\
& & & & & P = Q & T \\
& & & & & & T \\
& & & & & & T \\
& & & & & & T
\end{array}$$

TRY  
Yes it is a tautology

① <sup>Assumption</sup>  $(P \Rightarrow Q) = \neg P \vee Q$

?

ASSIGNMENT EXAMPLE 8.3 PG 214

## Inference in Propositional Logic

(3)

Inference is a technique by which, given a set of facts, or postulates, or axioms, or premises  $f_1, f_2, \dots, f_n$  a goal  $g$  is to be derived.

i.e. instead "where there is smoke there is fire" &

"There is smoke in the hill",

the statement "then the hill is on fire" can be easily deducted.

There are three rules for inferring facts (widely used in propositional logic)

- 1) Modus Ponens
- 2) Modus Tollens, and
- 3) Chain rule

1) MODUS PONENS (mod pon)

Given  $P \Rightarrow Q$  &  $P$  to be true,  $Q$  is true

$$\frac{P \Rightarrow Q \quad P}{Q}$$

} formed above one another the premises

} form below one & is the goal which can be inferred from the premises

2) Modus tollens

Given  $P \Rightarrow Q$  &  $\sim Q$  to be true  $\sim P$  is true

$$\frac{P \Rightarrow Q \quad \sim Q}{\sim P}$$

3) Chain rule

Given  $P \Rightarrow Q$  &  $Q \Rightarrow R$  to be true,  $P \Rightarrow R$  is true

$$\frac{P \Rightarrow Q \quad Q \Rightarrow R}{P \Rightarrow R}$$

it represents off the transitivity relation w.r.t " $\Rightarrow$ "

(4)

Q2 Given

Handwritten assign.

to be submitted on or before

4/10/24

i) CVD

ii)  $\sim H \Rightarrow A \wedge \sim B$ iii)  $(CVD) \Rightarrow \sim H$ (iv)  $(A \wedge \sim B) \Rightarrow (R \vee S)$ Can  $(R \vee S)$  be inferred from the above?

Mod Pon

$$\begin{array}{c} CVD \\ (CVD) \Rightarrow \sim H \\ \hline \sim H \quad (v) \end{array}$$

Chain n.

$$\begin{array}{c} \sim H \Rightarrow (A \wedge \sim B) \\ (A \wedge \sim B) \Rightarrow (R \vee S) \\ \hline \sim H \Rightarrow (R \vee S) \quad (vi) \end{array}$$

Mod Pov.

$$\begin{array}{c} \sim H \Rightarrow (R \vee S) \\ \sim H \\ \hline R \vee S \end{array}$$

Read Sect 8.1.2 pg 216

&amp; Example 8.6 pg 217

Vines for Plaem