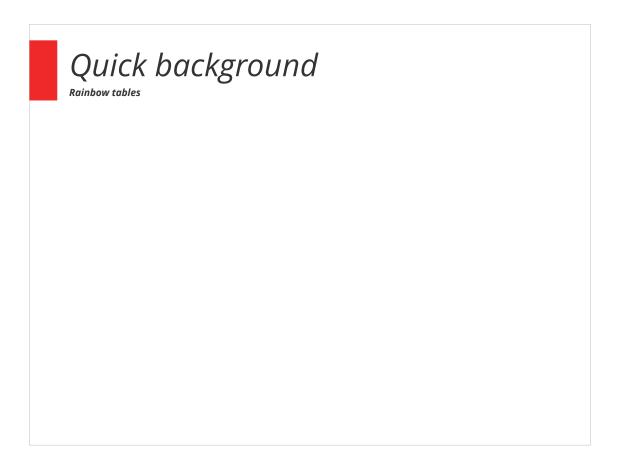
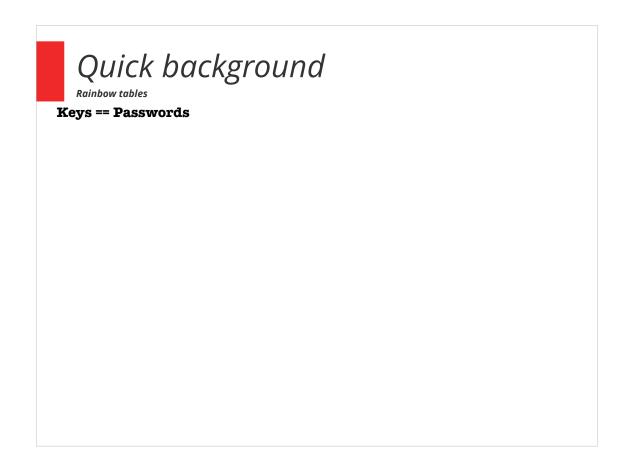


 Hi, I'm Chris. I'm going to try to tell you about rainbow table reduction functions in 5 minutes.



• Some quick background.



• First, whether I say "keys" or "passwords," I'm referring to the same thing.



- Plain-text, short, and user-generated
 - Example: monkey123

- They're plain-text strings.
- Short and user generated.



- Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of *all possible passwords* that could exist under certain constraints.

• Second, a "key space" is the set of all possible passwords that exist under certain constraints.



- Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of all possible passwords that could exist under certain constraints.

Constraints

• Key length - The number of characters that can exist in a key.

- Those constraints are:
 - "Key length" how many characters a key, can be...



- Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of all possible passwords that could exist under certain constraints.

Constraints:

- Key length The number of characters that can exist in a key.
- Allowable characters Permissible characters in a key.
 - Often a-z, A-Z, 0-9. Sometimes symbols.

• ... and, the "allowable characters" for a key.



- · Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of all possible passwords that could exist under certain constraints.

Constraints:

- Key length The number of characters that can exist in a key.
- Allowable characters Permissible characters in a key.
 - Often a-z, A-Z, 0-9. Sometimes symbols.

monkey123: Key length = 9 / allowable characters = 36 (a-z0-9)

- For example, all numbers and lower-case letters represent 36 characters.
- With a key length of 9 and 36 allowable characters, we could create the password 'monkey123'



- Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of *all possible passwords* that could exist under certain constraints.

Constraints:

- Key length The number of characters that can exist in a key.
- Allowable characters Permissible characters in a key.
 - Often a-z, A-Z, 0-9. Sometimes symbols.

 There are about 101 trillion possible passwords under these constraints.



- Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of *all possible passwords* that could exist under certain constraints.

Constraints:

- Key length The number of characters that can exist in a key.
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monkey123: Key length = 9 / allowable characters = 36 (a-z0-9) **Key space** = $^{\sim}36^{\circ}$ = 101 trillion

Cryptographic hash function - SHA-1

 Finally. A cryptographic hash function – SHA-1 in particular.



- · Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of *all possible passwords* that could exist under certain constraints.

Constraints:

- Key length The number of characters that can exist in a key.
- Allowable characters Permissible characters in a key.
 - Often a-z, A-Z, 0-9. Sometimes symbols.

monkey123: Key length = 9 / allowable characters = 36 (a-z0-9) **Key space** = $^{\sim}36^{\circ}$ = 101 trillion

Cryptographic hash function - SHA-1

• Variable length key in, 40-character hexadecimal string out

 With a hash function, any length of key in, and a 40-character hexadecimal string comes out.



- Plain-text, short, and user-generated
 - Example: monkey123

Key space

The set of all possible passwords that could exist under certain constraints.

Constraints:

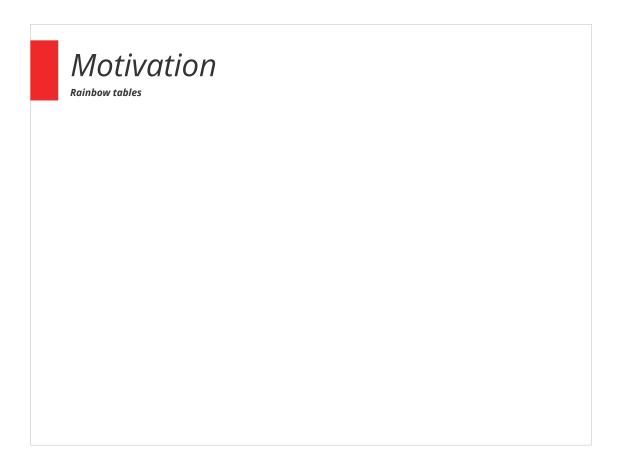
- Key length The number of characters that can exist in a key.
- Allowable characters Permissible characters in a key.
 - Often a-z, A-Z, 0-9. Sometimes symbols.

monkey123: Key length = 9 / allowable characters = 36 (a-z0-9)**Key space** = $^{\sim}36^9 = 101 trillion$

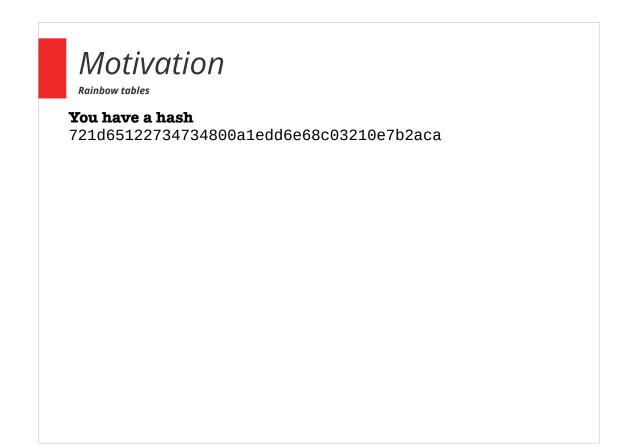
Cryptographic hash function - SHA-1

- Variable length key in, 40-character hexadecimal string out
- The same data in always results in the same hash out
- A second hash property is that the same key in always result in the same hash out.





• Now, the motivation.



• Say you have a hash.



721d65122734734800a1edd6e68c03210e7b2aca

What password generated this hash?

 You'd like to know what plain-text password produced your hash.



721d65122734734800a1edd6e68c03210e7b2aca

What password generated this hash?

• Most likely option: Brute force

- One way to answer this question is to use brute force.
- You would generate and hash every key in the key space, trying to find a hash that matches yours.



721d65122734734800a1edd6e68c03210e7b2aca

What password generated this hash?

- Most likely option: Brute force
- Surprisingly quick

- This can be done surprisingly quickly.
- In a key space of 101 trillion, 'monkey123' might have taken about 2 hours to find.



721d65122734734800a1edd6e68c03210e7b2aca

What password generated this hash?

- Most likely option: Brute force
- Surprisingly quick
- Not so quick when searching for multiple hashes

• But how long you were willing to wait for one hash, you may not be willing to wait for 100.



721d65122734734800a1edd6e68c03210e7b2aca

What password generated this hash?

- Most likely option: Brute force
- Surprisingly quick
- Not so quick when searching for multiple hashes
 - Regenerating keys, re-hashing them

 Because for each hash you search for, you're regenerating and rehashing the same keys each time.



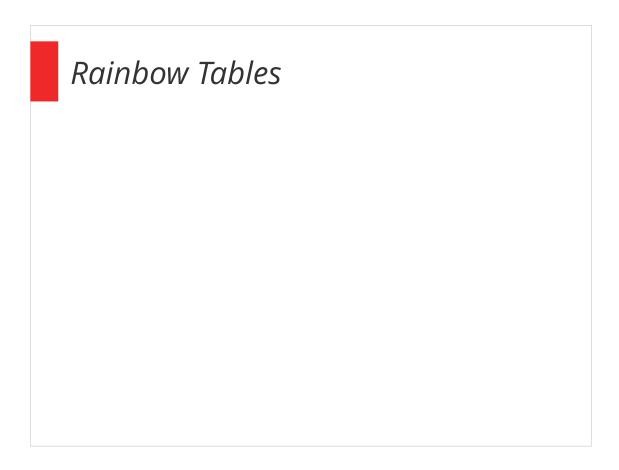
721d65122734734800a1edd6e68c03210e7b2aca

What password generated this hash?

- Most likely option: Brute force
- Surprisingly quick
- Not so quick when searching for multiple hashes
 - Regenerating keys, re-hashing them
 - Saving this work isn't feasible

• Due to storage constraints, saving all key:hash pair results isn't an option.

X



• This leads us to rainbow tables

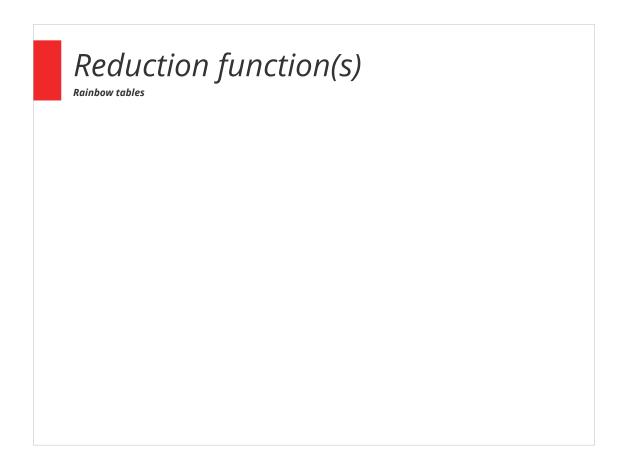
Rainbow Tables

A rainbow table offers a method by which one can search a very large key space multiple times without actually saving or recomputing the entire key space for each search.

Rainbow tables are a tradeoff between computation time and size.

- A rainbow table offers a method by which one can search a very large key space multiple times, without the need to save, or recompute, the entire key space for each search.
- It is a computation time/size trade off.





• Now, reduction functions!



• A reduction function takes as a pair, a hash, and a salt. It returns a plain-text key that exists within the target key space.



Consider it a "reverse hashing" function.

• It's sort of like a reverse hash function.



Consider it a "reverse hashing" function.

• The same (hash, salt) pair in always results in the same password out

 It's deterministic, in that the same "hash:salt pair in", always produces the same "plain-text key out".



Consider it a "reverse hashing" function.

- The same (hash, salt) pair in always results in the same password out
- · Output is evenly distributed across the key space

 And the entirety of the plain-text key space has equal chance of being produced by the reduction function.

Reduction function(s)

A reduction function takes a (hash, salt) pair and produces a plain-text key from the target key space.

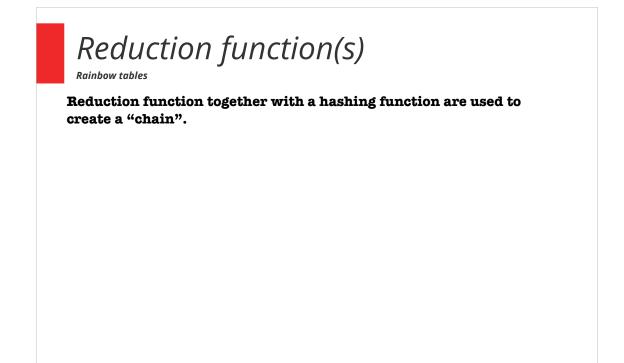
Consider it a "reverse hashing" function.

- The same (hash, salt) pair in always results in the same password out
- · Output is evenly distributed across the key space

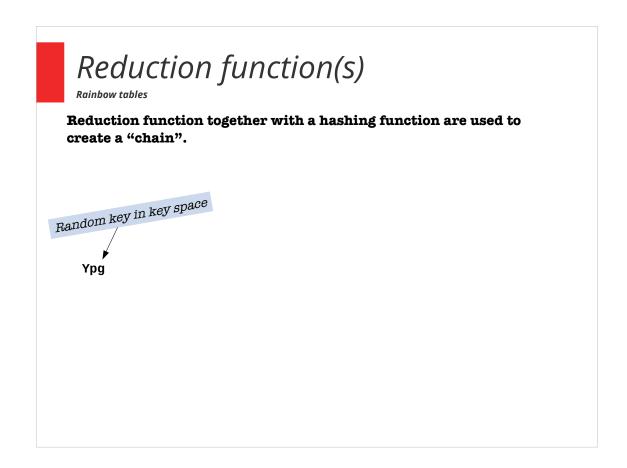
But a reduction function does not produce the password that generated a given hash!

• But the plain-text key responsible for the input hash is *not generated* by the reduction function!





• Together, the reduction and hash function are used to create what's called a "chain."



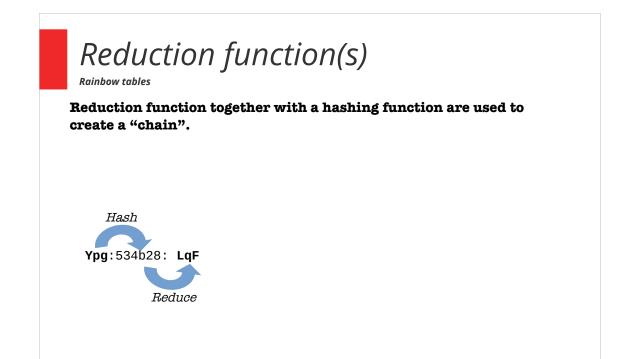
• To generate a chain, you start with a randomly generated plain-text key from the key space.



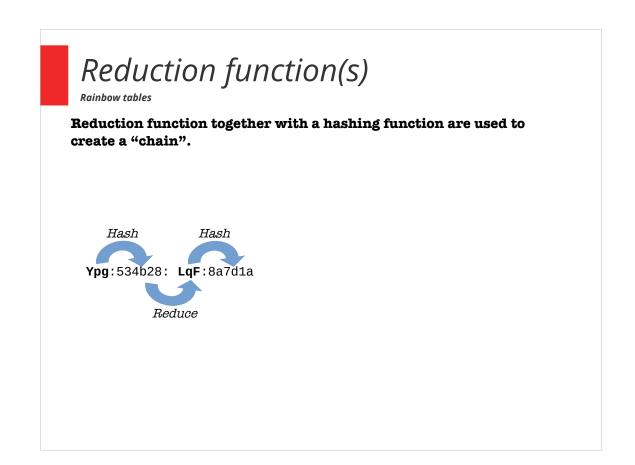
Reduction function together with a hashing function are used to create a "chain".



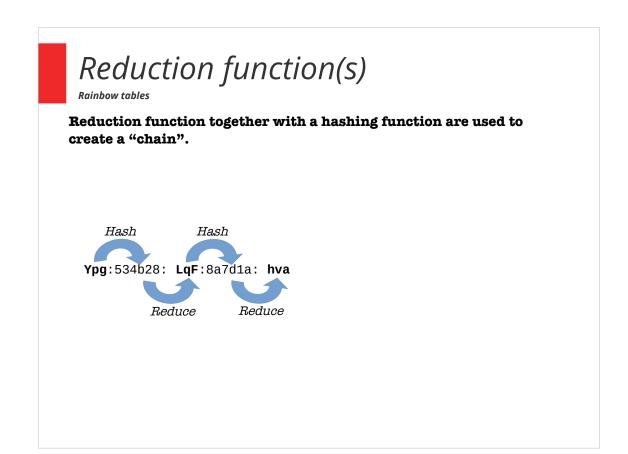
You hash that key.



• Then use the reduction function on the resulting hash.



• This gives you a new key, which you hash...



• And this carries on...

Reduction function(s) Rainbow tables

Reduction function together with a hashing function are used to create a "chain".

Hash Hash Hash

Ypg:534b28: LqF:8a7d1a: hva:5cd69f

Reduce Reduce

Reduction function(s) Rainbow tables

Reduction function together with a hashing function are used to create a "chain".

HashHashHashYpg:534b28:LqF:8a7d1a:hva:5cd69f:txPReduceReduceReduce



Hash Hash Hash Hash

Ypg:534b28: LqF:8a7d1a: hva:5cd69f: txP:1eba8d

Reduce Reduce Reduce

Reduction function(s) Rainbow tables

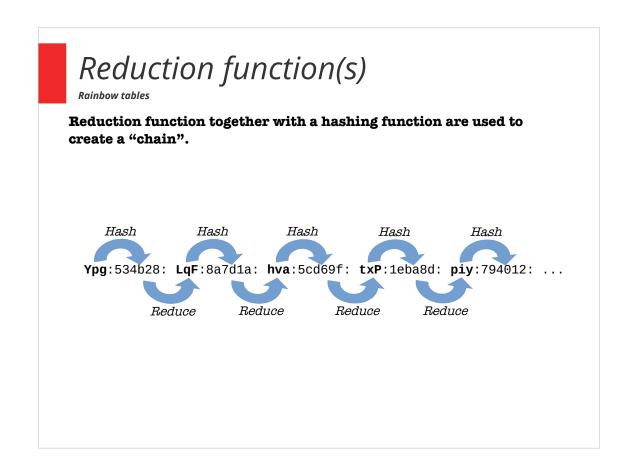
Reduction function together with a hashing function are used to create a "chain".

Hash Hash Hash

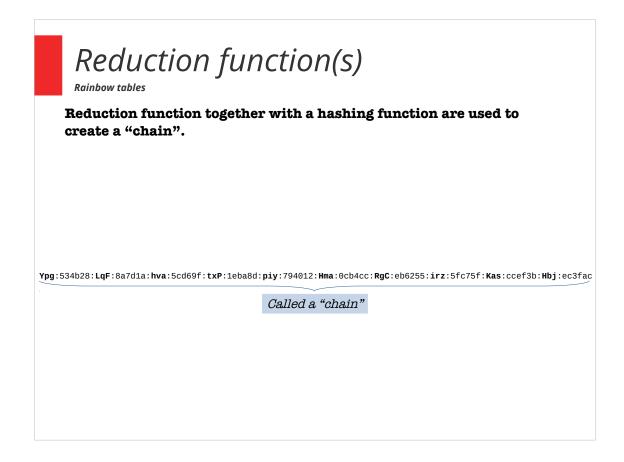
Ypg:534b28: LqF:8a7d1a: hva:5cd69f: txP:1eba8d: piy

Reduce Reduce Reduce Reduce

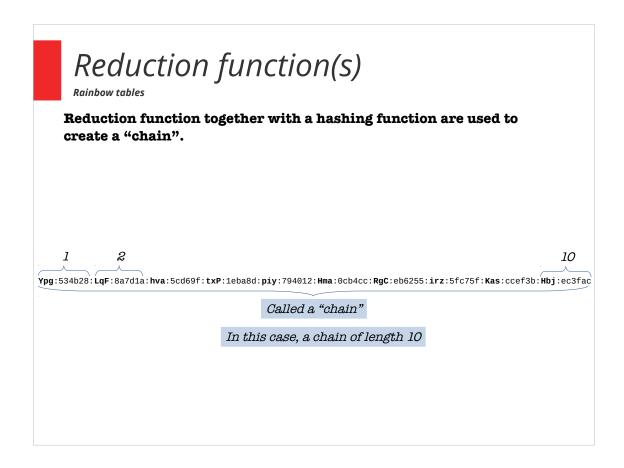




• For an arbitrary number of iterations



• What you'll end up with is a "chain" of plain-text keys and hashes.



• This particular chain would be considered of length 10.



In a Rainbow table, this chain constitutes one row of the table

• One chain is one row of a rainbow table.



In a Rainbow table, this chain constitutes one row of the table

 $\textbf{Ypg:} 534b28: \textbf{LqF:} 8a7d1a: \textbf{hva:} 5cd69f: \textbf{txP:} 1eba8d: \textbf{piy:} 794012: \textbf{Hma:} 0cb4cc: \textbf{RgC:} eb6255: \textbf{irz:} 5fc75f: \textbf{Kas:} ccef3b: \textbf{Hbj:} ec3fac \textbf{Juf:} 1eee70: \textbf{wwu:} 43db1e: \textbf{cbz:} 082829: \textbf{wfv:} ccbc9f: \textbf{nzg:} 629c4b: \textbf{szd:} 9f31f9: \textbf{cqy:} ef2dce: \textbf{utc:} d0206a: \textbf{mif:} 181072: \textbf{jkm:} e0f1c1 \textbf{mid:} 181072: \textbf{jk$

• Now ... a rainbow table will have millions of rows



In a Rainbow table, this chain constitutes one row of the table

```
Ypg:534b28:LqF:8a7d1a:hva:5cd69f:txP:1eba8d:piy:794012:Hma:0cb4cc:RgC:eb6255:irz:5fc75f:Kas:ccef3b:Hbj:ec3fac
Juf:1eee70:wwu:43db1e:cbz:082829:wfv:ccbc9f:nzg:629c4b:szd:9f31f9:cqy:ef2dce:utc:d0206a:mif:181072:jkm:e0f1c1
Axl:39f5ff:rdv:22efdd:zpy:18a63e:arm:322b66:pwd:6cf639:lfx:cc295a:qxn:a6d415:srr:6cda4d:smo:a280bd:qzx:bb8dd6
psy:b6741d:ryr:1834b0:rct:6229af:cod:67cdbf:qeu:745566:phg:58b0a2:sjp:e8dfb7:fbj:e8a8f1:upk:49ad89:otv:6bad6d
```

• Now a rainbow ... table will have millions of rows



In a Rainbow table, this chain constitutes one row of the table

```
Ypg:534b28:LqF:8a7d1a:hva:5cd69f:txP:1eba8d:piy:794012:Hma:0cb4cc:RgC:eb6255:irz:5fc75f:Kas:ccef3b:Hbj:ec3fac
Juf:1eee70:wwu:43db1e:cbz:082829:wfv:ccbc9f:nzg:629c4b:szd:9f31f9:cqy:ef2dce:utc:d0206a:mif:181072:jkm:e0f1c1
Axl:39f5ff:rdv:22efdd:zpy:18a63e:arm:322b66:pwd:6cf639:1fx:cc295a:qxn:a6d415:srr:6cda4d:smo:a280bd:qzx:bb8dd6
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```

Now a rainbow table will have millions ... of rows



In a Rainbow table, this chain constitutes one row of the table

```
Ypg:534b28:LqF:8a7d1a:hva:5cd69f:txP:1eba8d:piy:794012:Hma:0cb4cc:RgC:eb6255:irz:5fc75f:Kas:ccef3b:Hbj:ec3fac
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on:00d908:lia:1f2448:gwt:65db83:ohp:ae9a67:ihe:47199b:npp:e23cd4:owb:ef8a53:gln:8c7b56:bzz:0d748a:yyi:c86e68
viz:89fe8d:dvt:9fb4c5:lvo:0c1705:flh:e9f0e0:vuz:19b6f4:azo:7b76f4:buf:4106b0:boc:3026a9:xga:682bc4:lik:d73aa9
```

Now a rainbow table will have millions of rows



Absolutely no space savings at this point.

```
Ypg:534b28:LqF:8a7d1a:hva:5cd69f:txP:1eba8d:piy:794012:Hma:0cb4cc:RgC:eb6255:irz:5fc75f:Kas:ccef3b:Hbj:ec3fac
Juf:1eee70:wwu:43db1e:cbz:082829:wfv:ccbe9f:nzg:629c4b:szd:9f31f9:cqy:ef2dce:utc:d0206a:mif:181072:jkm:e0f1c1
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```

 But, as shown here, there is absolutely no space savings over just saving every key:hash pair from a key space.



Absolutely no space savings at this point.

Solution...

```
Ypg:534b28:LqF:8a7d1a:hva:5cd69f:txP:1eba8d:piy:794012:Hma:0cb4cc:RgC:eb6255:irz:5fc75f:Kas:ccef3b:Hbj:ec3fac
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Flp:e82c72:smo:a280bd:jsq:9e99a9:lrj:0e0e00:ndn:b92664:key:7377de:wcl:37e454:mie:22bc6a:ghj:377055:puv:09ce7e
xqL:6cee39:ldd:959267:kpw:0c6093:uln:5c064d:zzu:a041ef:jgp:224d8e:ybx:9028a9:ich:c70ed3:oao:6b4e18:abg:0f6a57
oko:4ab255:meb:70b6c1:ppu:ff081d:qkl:794b27:idu:833720:iuu:d57c7a:bnj:bc030d:woj:5af599:wzw:de16ea:yyk:07ee74
bfW:8250e2:cge:3e8184:jrz:987beb:kph:f4f7ed:tau:ccae45:bet:38ee18:pjr:3eaaf2:joi:892e76:srl:119545:ifl:291ceb
Dzm:4cda48:gea:7c82be:lmh:a3fe95:mpz:99a245:lls:bb60d7:qt;51731:mnb:98ac76:wwc:ac08f3:mzs:5bf966:coj:a4a52f
EST:b47641:ecl:22253a:xbb:566247:oxb:451b06:hgz:d9chaf:pht:16d0c5:oek:2c0aa0:plc:88f100:oyf:005e33:gyj:f1894b
con:00d908:lia:1f2448:gwt:65db83:ohp:ae9a67:ihe:47199b:npp:e23cd4:owb:efa8c3:gln:8c7b56:bzz:0d748a:yyi:c86e6
viz:89fe8d:dvt:9fb4c5:lvo:0c1705:flh:e9f0e0:vuz:19b6f4:azo:7b76f4:buf:4106b0:boc:3026a9:xga:682bc4:lik:d73aa9
```

The solution...



Absolutely no space savings at this point.

Solution...

Drop all entries from the middle of the table

```
Ypg:534b28:LqF:8a7d1a:hva:5cd69f:txP:1eba8d:piy:794012:Hma:0cb4cc:RgC:eb6255:irz:5fc75f:Kas:ccef3b:Hbj:ec3fac
Juf:1eee70:wwu:43db1e:cbz:082829:wfv:ccbc9f:nzg:629c4b:szd:9f31f9:cqy:ef2dce:utc:d0206a:mf:181072:jkm:e0f1c1
Axl:39f5ff:rdv:22efdd:zpy:18a63e:arm:322b66:pwd:6cf639:lfx:cc295a:qxn:a6d415:srr:6cda4d:smo:a280bd:qzx:bb8dd6
psy:b6741d:ryr:1834b0:rct:6229af:cod:67cdbf:qeu:745566:phg:58b0a2:sjp:e8dfb7:fbj:e8a8f1:upk:49ad89:otv:6bad6d
drP:3b50a4:adw:6bf788:jyf:846e10:zjs:e07110:izh:260933:jmd:866b85:pqm:2807f5:nlx:127202:ykv:45de73:dzf:8691c1
lQg:2ba3d4:eyc:08a9ff:row:6555a7:mcn:281c4d:svo:4c1a8b:tcp:64151f:jvu:89b3f3:jbb:80f6af:ksm:2127d9:adz:56c0f4
Eeg:acf56a:zxc:5ef6dd:hle:0e036c:hnu:2f1128:ztt:7def93:aww:8d2d0e:bmx:277b00:ffm:5894fd:vtg:1675c0:vqe:a201ad
rIc:bdaed9:jgh:1cc0d5:kxk:cb5d26:ssi:063f61:vlg:b86778:kke:600344:fhx:51cae4:zxm:07d38d:hbo:c910a4:rac:be1d88
Flp:e82c72:smo:a280bd:jsq:9e99a9:lrj:0e0e00:ndn:b92664:key:7377de:wcl:37e454:mie:22bc6a:ghj:377055:puv:09ce7e
xqL:6cea99:ldd:959267:kpw:0c6093:uln:5co64d:zzu:a041ef;jgp:224d8e:ybx:9028a9:ich:c70ed3:oao:6b4e18:abg:0f6a57
ok0:4ab255:meb:70b6c1:ppu:ff081d:qkl:794b27:idu:833720:iuu:d57c7a:bnj:bc030d:woj:5af599:wzw:de16ea:yyk:07ee74
bfW:8250e2:cge:3e8184:jrz:987beb:kph:f4f7ed:tau:ccae45:bet:38ee18:pjr:3eaaf2:joi:892e76:srl:119545:if1:291ceb
Dzm:4cda48:gea:7c82be:llm:ka3fe95:mpz:99245:lls:bb60d7:cgj:15731a:mnb:98ac76:wwc:ac088f3:mzs:5bf965:coj:a4a52f
EST:b47041:ecl:22253a:xbh:566247:oxb:451b06:hgz:d9cbaf:pht:16d9c5:oek:2c0aa0:plc:88f100:oyf:005e33:gyj:f1894b
con:00d908:lia:1f2448:gwt:65db83:ohp:ae9a67:ihe:47199b:npp:e23cd4:owb:efa8c3:gln:8c7b56:bzz:0d748a:yyi:c86e68
viz:89fe8d:dvt:9fb4c5:lvo:0c1705:flh:e9f0e0:vuz:19b6f4:azo:7b76f4:buf:4106b0:boc:3026a9:xga:682bc4:lik:d73aa9
```

Is to drop the entire middle of the table



Absolutely no space savings at this point.

Solution...

Drop all entries from the middle of the table

```
Ypg:534b28:LqF:8a7d1a:hva:5cd69f:txP:1eba8d:piy:794012:Hma:0cb4cc:RgC:eb6255:irz:5fc75f:Kas:ccef3b:Hbj:ec3fac
Juf:1eec70:wwu:43db1e:cbz:082829:wfv:ccbc9f:nzg:629c4b:szd:9f31f9:cqy:ef2dce:utc:d0206a:mif:181072:jkm:e0f1c1
Axl:39f5ff:rdv:22efdd:zpy:18a63e:arm:322b66:pwd:6cf639:lfx:cc295a:qxn:a6d415:srr:6cd4d:smo:a280bd:qzx:bb8dd6
pxy:b6741d:ryr:1834b0:rct:6229af:cod:67cdbf:qeu:745566:phg:58b0a2:sjp:e8dfb7:fbj:e8a8f1:upk:49ad89:otv:6bad6d
drP:3b50a4:adw:6bf788:jyf:846e10:zjs:e07110:izh:260933:jmd:866085:pqm:2807f5:nlx:127202:ykv:45de73:dzf:8691c1
lQg:2ba3d4:eyc:08a9ff:row:6555a7:mcn:281c4d:svo:4c1a8b:tcp:64151f:jvu:89b3f3:jhb:80f6af:ksm:2127d9:adz:56c0f4
Eeg:acf56a:zxc:5ef0dd:hle:0e036c:hnu:2f1128:ztt:7def93:aww:8d2d0e:bmx:277b00:ffm:5894fd:vtg:1675c0:vqe:a201ad
rIc:bdaed9:jgh:1cc0d5:kxk:cb5d26:ssi:063f61:vlg:b86778:kke:600344:fhx:51cae4:zxm:07d38d:hbo:c910a4:rac:be1d88
Flp:e82c72:smo:a280bd:jsq:9e99a9:lrj:0e0e00:ndn:b92664:key:7377de:wcl:37e454:mie:22bc6a:ghj:377055:puv:09ce7e
xqL:6ce39:ldd:959267:kpw:0c6093:uln:5co64d:zzu:a041ef;jgp:224d8e:ybx:9028a9:ich:c70ed3:oao:6b4e18:abg:0f6a57
ok0:4ab255:meb:70b6c1:ppu:ff081d:qkl:794b27:idu:833720:iuu:d57c7a:bnj:bc030d:woj:5af599:wzw:de16ea:yyk:07ee74
bfw:8250e2:cge:3e8184:jrz:987beb:kph:f4f7ed:tau:ccae45:bet:38ee18:pjr:3eaaf2:joi:892e76:rsl:1115945:ifl2:291ceb
Dzm:4cd48:gea:7c82be:llm:a3fe95:mpz:99245:1ls:bb60d7:cqj:51731a:mmb:98ac76:wwc:ac086f3:mzs:5bf965:coj:a4852f
EST:b47041:ecl:22253a:xbh:566247:oxb:451b06:hgz:d9cbaf:pht:16d9c5:oek:2c0aa0:plc:88f100:oyf:005e33:gyj:f1894b
con:00d908:lia:1f2448:gwt:65db83:ohp:ae9a67:ihe:47199b:npp:e23cd4:owb:efa8c3:gln:8c7b56:bzz:0d748a:yyi:c86e68
viz:89fe8d:dvt:9fb4c5:lvo:0c1705:f1h:e9f0e0:vuz:19b6f4:azo:7b76f4:buf:4106b0:boc:3026a9:xga:682bc4:lik:d73aa9
```

 From the first key on the left, to the final hash on the right.





Absolutely no space savings at this point.

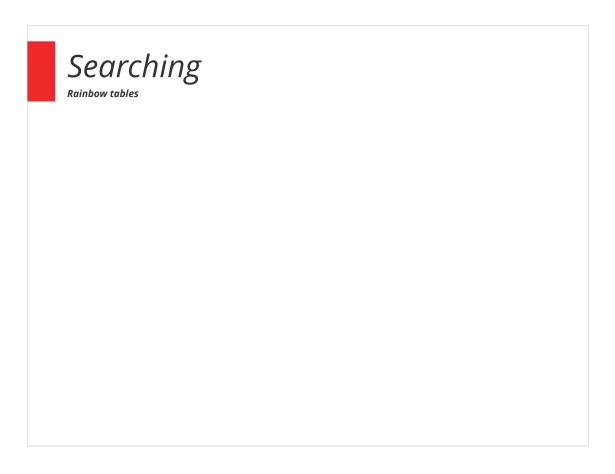
Solution...

Drop all entries from the middle of the table

Ypg:ec3fac
Juf:e0f1c1
Ax1:bb8dd6
psy:6bad6d
drP:8691c1
lQg:56c0f4
Eeg:a201ad
rIc:be1d88
Flp:09ce7e
xqL:0f6a57
ok0:07ce74
bfW:291ceb
Dzm:a4a52f
EST:f1894b
con:c86e68
viz:d73aa9

Keeping only the first key, and final hash.



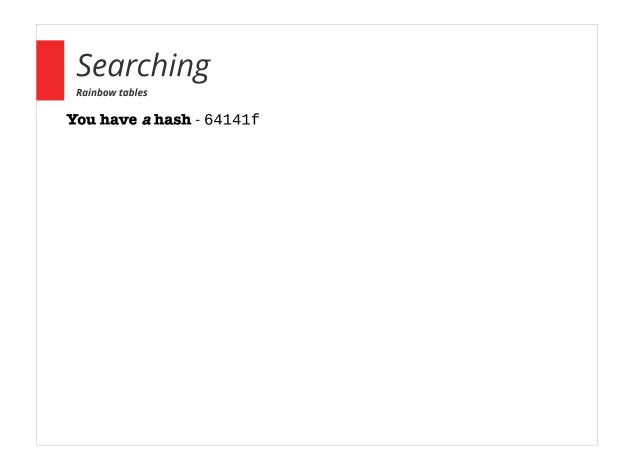


• There's an obvious question at this point...

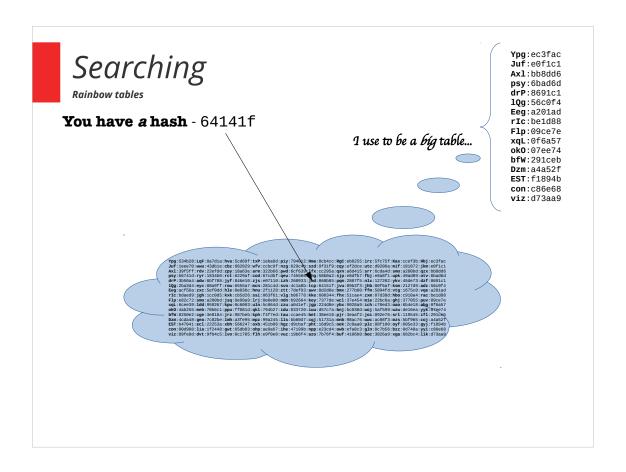


How do we search this table in such a way that we can find the any of the 144 key:hash pairs dropped from the full table? Ypg:ec3fac
Juf:e0f1c1
Ax1:bb8dd6
psy:6bad6d
drP:8691c1
lQg:56c0f4
Eeg:a201ad
rIc:be1d88
Flp:09ce7e
xqL:0f6a57
ok0:07ee74
bfW:291ceb
Dzm:a4a52f
EST:f1894b
con:c86e68
viz:d73aa9

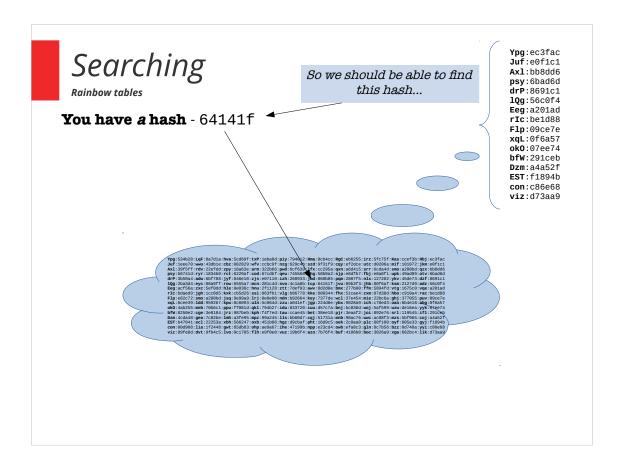
 How can we search this reduced table and still find any of the key:hash pairs that were dropped from the full table?



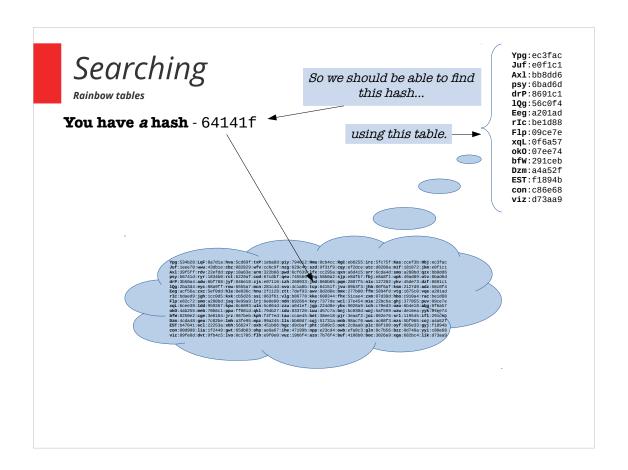
• Well lets see. Suppose you have a hash...



 A hash from a key:hash pair that was computed in the full-sized table, but was dropped.

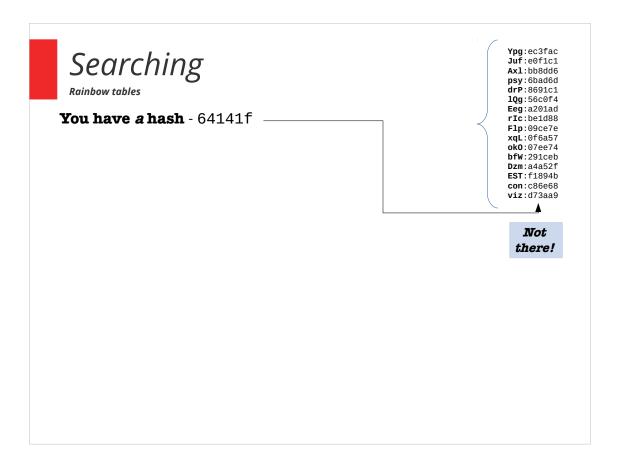


· We should be able to find this hash...

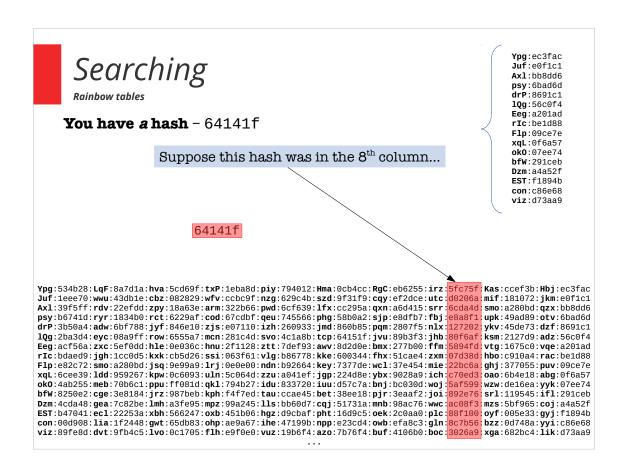


... in our reduced-size table.

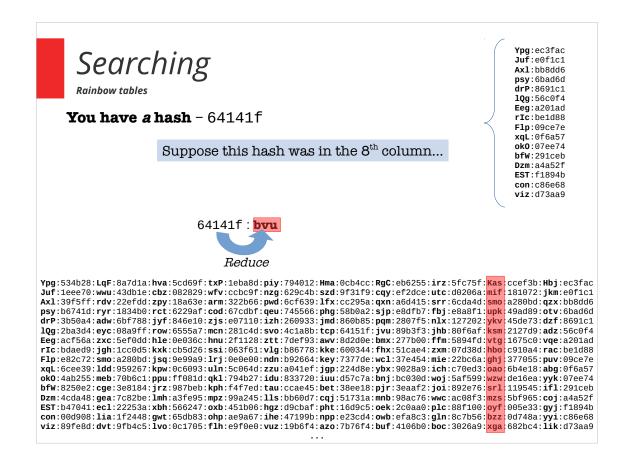
X



• But in this case, it's not in there.



 But suppose when the larger table was generated initially, that our hash was in, as an example, the 8th column



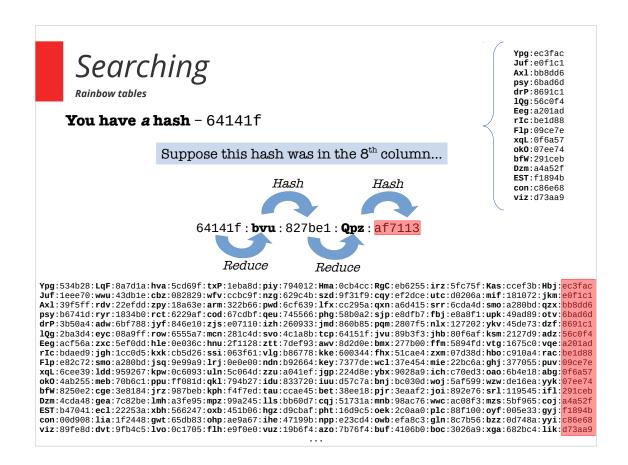
- Recalling the process used earlier to generate a chain by reducing a hash, and then hashing the key -
- Imagine this same reduce/hash process on our current hash.
- We first reduce our hash...



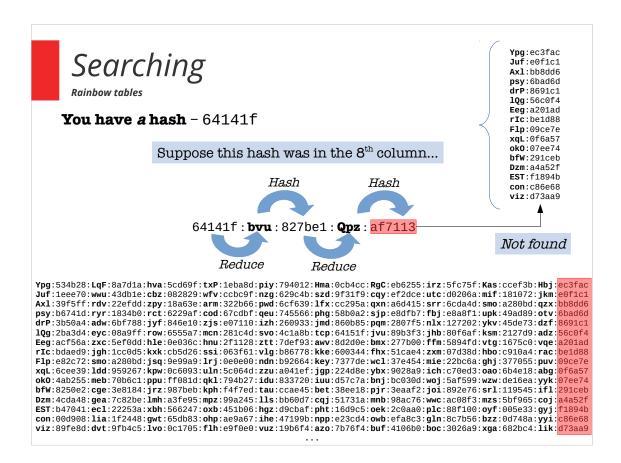
... and then hash the resulting key.



We reduce that new hash, and get a new key.

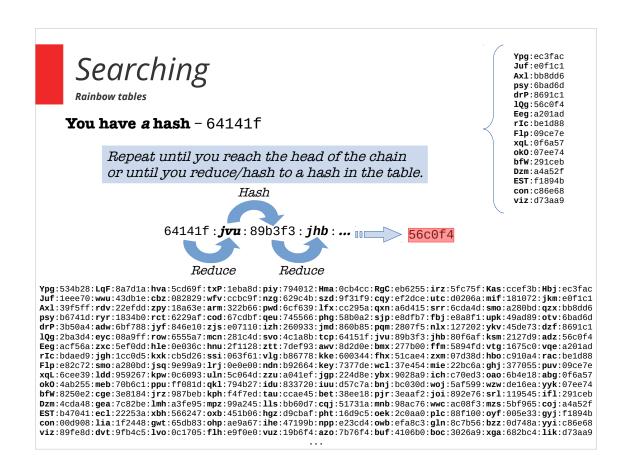


- Finally, we hash that key, getting a final hash.
- So if our hash existed in the 8th column of the larger table, through this process we should now have a hash that exists in the final column.

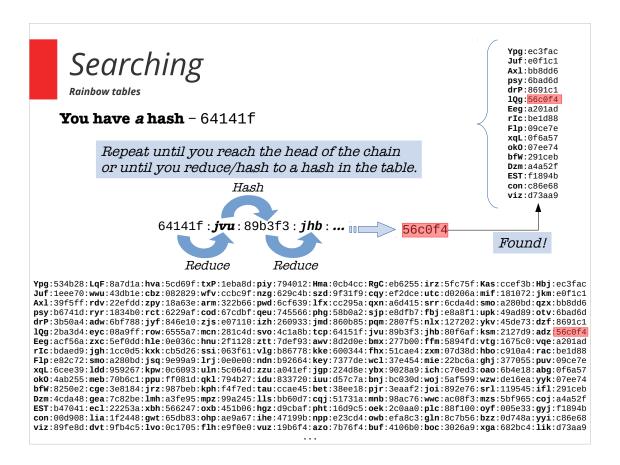


· But we don't.



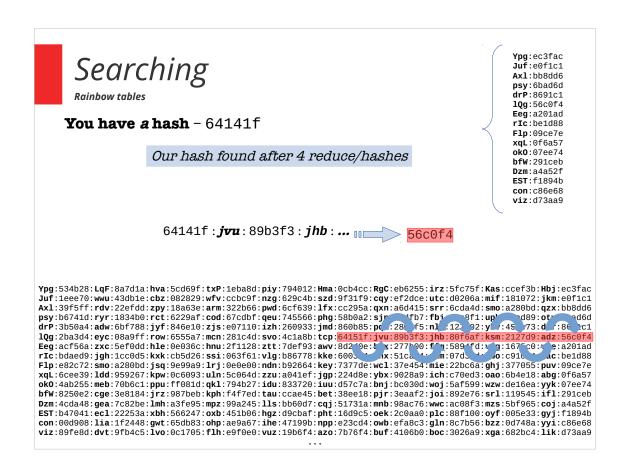


 This process is repeated, picking a column the hash you hold might have been in, and running it down a reduce/hash chain to see if the resulting hash exists in the reduced table.

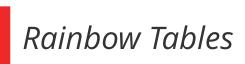


- You'll either reach the front of the table, not finding a match,
- Or you'll find a chain location that results in a hash that can be found in the reduced table.





 For our example hash, when 4 reduce/hash iterations are applied, we end up with a hash that is in the reduced table.



There's more to rainbow tables!

http://upnix.github.io/RainbowTables

- How do we get our key back?!
- What about this "salt" you mentioned for the reduction function?
- Why the name rainbow tables?
- What about chain collisions?

• And that's a very good thing...