GraalVM Stamp Theory

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Abstract

The GraalVM compiler uses stamps to track type and range information during program analysis. Type information is recorded by using distinct subclasses of the abstract base class Stamp, i.e. IntegerStamp is used to represent an integer type. Each subclass introduces facilities for tracking range information. Every subclass of the Stamp class forms a lattice, together with an arbitrary top and bottom element each sublattice forms a lattice of all stamps. This Isabelle/HOL theory models stamps as instantiations of a lattice.

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1 Stamps: Type and Range Information

```
\begin{array}{c} \textbf{theory} \ StampLattice\\ \textbf{imports}\\ \ Values\\ \ HOL.Lattices\\ \textbf{begin} \end{array}
```

1.1 Void Stamp

The VoidStamp represents a type with no associated values. The VoidStamp lattice is therefore a simple single element lattice.

```
{\bf datatype}\ \mathit{void} =
  VoidStamp
instantiation \ void :: order
begin
definition less-eq\text{-}void :: void \Rightarrow void \Rightarrow bool where
  less-eq	ext{-}void\ a\ b=\mathit{True}
definition less\text{-}void :: void \Rightarrow void \Rightarrow bool where
  less-void\ a\ b=False
instance
  apply standard
  \mathbf{apply} \ (simp \ add: \ less-eq\text{-}void\text{-}def \ less-void\text{-}def) +
  by (metis (full-types) void.exhaust)
end
instantiation void :: semilattice-inf
begin
definition inf-void :: void \Rightarrow void \Rightarrow void where
  inf-void\ a\ b = VoidStamp
instance
  apply standard
  by (simp add: less-eq-void-def)+
\mathbf{end}
instantiation \ void :: semilattice-sup
begin
definition sup\text{-}void :: void \Rightarrow void \Rightarrow void where
  sup\text{-}void\ a\ b=VoidStamp
```

instance

 ${\bf apply} \ standard$

by $(simp\ add:\ less-eq\text{-}void\text{-}def)+$

end

 $\begin{array}{l} \textbf{instantiation} \ \textit{void} :: \textit{bounded-lattice} \\ \textbf{begin} \end{array}$

 $\begin{array}{ll} \textbf{definition} \ \textit{bot-void} :: \textit{void} \ \textbf{where} \\ \textit{bot-void} = \textit{VoidStamp} \end{array}$

 $\begin{array}{ll} \textbf{definition} \ top\text{-}void :: void \ \textbf{where} \\ top\text{-}void = VoidStamp \end{array}$

instance

apply standard

by $(simp\ add:\ less-eq\text{-}void\text{-}def)+$

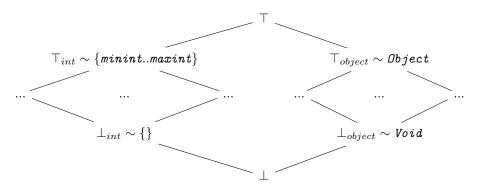
end

Definition of the stamp type

datatype stamp =

intstamp int64 int64 — Type: Integer; Range: Lower Bound & Upper Bound

1.2 Stamp Lattice



1.2.1 Stamp Order

Defines an ordering on the stamp type.

One stamp is less than another if the valid values for the stamp are a strict subset of the other stamp.

 $\begin{array}{l} \textbf{instantiation} \ stamp :: \ order \\ \textbf{begin} \end{array}$

```
fun less-eq\text{-}stamp :: stamp \Rightarrow stamp \Rightarrow bool where
  less-eq-stamp (intstamp l1\ u1) (intstamp l2\ u2) = (\{l1..u1\}\subseteq\{l2..u2\})
fun less-stamp :: stamp \Rightarrow stamp \Rightarrow bool where
 less-stamp (intstamp l1\ u1) (intstamp l2\ u2) = (\{l1..u1\} \subset \{l2..u2\})
lemma less-le-not-le:
 fixes x y :: stamp
 shows (x < y) = (x \le y \land \neg y \le x)
 \textbf{by} \ (\textit{metis subset-not-subset-eq stamp.exhaust less-stamp.simps less-eq-stamp.simps})
lemma order-refl:
 fixes x :: stamp
 shows x \leq x
 by (metis stamp.exhaust dual-order.refl less-eq-stamp.simps)
lemma order-trans:
 fixes x \ y \ z :: stamp
 shows x \leq y \Longrightarrow y \leq z \Longrightarrow x \leq z
proof -
 fix x :: stamp and y :: stamp and z :: stamp
 assume x \leq y
 assume y \leq z
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by auto
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by auto
 obtain l3\ u3 where zdef: z = intstamp\ l3\ u3
   using stamp.exhaust by auto
 have s1: \{l1..u1\} \le \{l2..u2\}
   using \langle x \leq y \rangle by (simp add: ydef xdef)
 have s2: \{l2..u2\} \le \{l3..u3\}
   using \langle y \leq z \rangle by (simp add: zdef ydef)
 from s1 \ s2 \ \text{have} \ \{l1..u1\} \le \{l3..u3\}
   by (meson dual-order.trans)
 then show x \le z
   by (simp add: zdef xdef)
qed
lemma antisym:
 fixes x y :: stamp
 shows x \leq y \Longrightarrow y \leq x \Longrightarrow x = y
proof -
 fix x :: stamp
 \mathbf{fix} \ y :: stamp
 assume xlessy: x \leq y
 assume ylessx: y \le x
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by auto
```

```
obtain l2\ u2 where ydef: y = intstamp\ l2\ u2
   using stamp.exhaust by auto
 from xlessy have s1: \{l1..u1\} \subseteq \{l2..u2\}
   by (simp add: ydef xdef)
 from ylessx have s2: \{l2..u2\} \subseteq \{l1..u1\}
   by (simp add: ydef xdef)
 have \{l1..u1\} \subseteq \{l2..u2\} \implies \{l2..u2\} \subseteq \{l1..u1\} \implies \{l1..u1\} = \{l2..u2\}
  then have s3: \{l1..u1\} = \{l2..u2\} \Longrightarrow (l1 = l2) \land (u1 = u2)
 then have (l1 = l2) \land (u1 = u2) \Longrightarrow x = y
   using xdef ydef by fastforce
 then show x = y
   using s1 s2 s3 by fastforce
\mathbf{qed}
instance
 apply standard
 by (simp add: antisym order-trans order-refl less-le-not-le)+
end
1.2.2
        Stamp Join
Defines the join operation for stamps.
For any two stamps, the join is defined as the intersection of the valid values
for the stamp.
instantiation stamp :: semilattice-inf
begin
notation inf (infix \sqcap 65)
fun inf-stamp :: stamp \Rightarrow stamp \Rightarrow stamp where
  inf-stamp (intstamp l1 u1) (intstamp l2 u2) = intstamp (max l1 l2) (min u1 u2)
lemma inf-le1:
 fixes x y :: stamp
 shows (x \sqcap y) \leq x
proof -
 fix x :: stamp
 \mathbf{fix} \ y :: stamp
 obtain l1 u1 where xdef: x = intstamp l1 u1
   using stamp.exhaust by auto
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by auto
 have joindef: x \sqcap y = intstamp (max l1 l2) (min u1 u2)
```

(is ?join = intstamp ?l3 ?u3)

```
by (simp add: ydef xdef)
 have leq: \{?l3..?u3\} \subseteq \{l1..u1\}
   \mathbf{by} \ simp
 have (x \sqcap y) \le x = (\{?l3..?u3\} \subseteq \{l1..u1\})
   using joindef by (simp add: xdef)
 then show (x \sqcap y) \leq x
   by (simp add: leq)
qed
lemma inf-le2:
 fixes x y :: stamp
 shows (x \sqcap y) \leq y
proof -
 \mathbf{fix}\ x::stamp
 \mathbf{fix} \ y :: stamp
 obtain l1 u1 where xdef: x = intstamp l1 u1
   using stamp.exhaust by auto
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by auto
 have joindef: x \sqcap y = intstamp \ (max \ l1 \ l2) \ (min \ u1 \ u2)
   (is ?join = intstamp ?l3 ?u3)
   by (simp add: ydef xdef)
 have leq: \{?l3..?u3\} \subseteq \{l2..u2\}
   by simp
 have (x \sqcap y) \le y = (\{?l3..?u3\} \subseteq \{l2..u2\})
   using joindef by (simp add: ydef)
  then show (x \sqcap y) \leq y
   by (simp add: leq)
qed
lemma inf-greatest:
 fixes x \ y \ z :: stamp
 shows x \le y \Longrightarrow x \le z \Longrightarrow x \le (y \sqcap z)
proof -
 fix x y z :: stamp
 assume xlessy: x \leq y
 assume xlessz: x \leq z
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by auto
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by auto
 obtain l3\ u3 where zdef: z = intstamp\ l3\ u3
   using stamp.exhaust by auto
 obtain l4 u4 where yzdef: y \sqcap z = intstamp l4 u4
   by (meson inf-stamp.elims)
 have max4: l4 = max l2 l3
   using yzdef by (simp add: zdef ydef)
 have min4: u4 = min u2 u3
   using yzdef by (simp add: zdef ydef)
```

```
have \{l1..u1\} \subseteq \{l2..u2\}

using xlessy by (simp \ add: \ ydef \ xdef)

have \{l1..u1\} \subseteq \{l3..u3\}

using xlessz by (simp \ add: \ zdef \ xdef)

have leq: \{l1..u1\} \subseteq \{l4..u4\}

using \langle \{l1..u1\} \subseteq \{l2..u2\} \rangle \langle \{l1..u1\} \subseteq \{l3..u3\} \rangle by (simp \ add: \ min4 \ max4)

have x \leq (y \cap z) = (\{l1..u1\} \subseteq \{l4..u4\})

by (simp \ add: \ xdef \ yzdef)

then show x \leq (y \cap z)

using leq by simp

qed

instance

apply standard

by (simp \ add: \ inf-greatest \ inf-le2 \ inf-le1)+

end
```

1.2.3 Stamp Meet

Defines the *meet* operation for stamps.

For any two stamps, the *meet* is defined as the union of the valid values for the stamp.

```
instantiation \ stamp :: semilattice-sup
begin
notation sup (infix \sqcup 65)
fun sup\text{-}stamp :: stamp \Rightarrow stamp \Rightarrow stamp where
 sup\text{-}stamp \ (intstamp \ l1 \ u1) \ (intstamp \ l2 \ u2) = intstamp \ (min \ l1 \ l2) \ (max \ u1 \ u2)
lemma sup-ge1:
 fixes x y :: stamp
 shows x \leq x \sqcup y
proof -
 \mathbf{fix} \ x :: stamp
 \mathbf{fix} \ y :: stamp
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by auto
 obtain l2\ u2 where ydef: y = intstamp\ l2\ u2
   using stamp.exhaust by auto
 have joindef: x \sqcup y = intstamp (min l1 l2) (max u1 u2)
   (is ?join = intstamp ?l3 ?u3)
   by (simp add: ydef xdef)
 have leq: \{l1..u1\} \subseteq \{?l3..?u3\}
   by simp
 have x \le x \sqcup y = (\{l1..u1\} \subseteq \{?l3..?u3\})
   using joindef by (simp add: xdef)
  then show x \leq x \sqcup y
```

```
by (simp add: leq)
qed
lemma sup-ge2:
 fixes x y :: stamp
 shows y \leq x \sqcup y
proof -
 \mathbf{fix} \ x :: stamp
 fix y :: stamp
 obtain l1 u1 where xdef: x = intstamp l1 u1
   using stamp.exhaust by auto
 obtain l2\ u2 where ydef: y = intstamp\ l2\ u2
   using stamp.exhaust by auto
 have joindef: x \sqcup y = intstamp (min l1 l2) (max u1 u2)
   (is ?join = intstamp ?l3 ?u3)
   by (simp add: ydef xdef)
 have leq: \{l2..u2\} \subseteq \{?l3..?u3\} (is ?subset-thesis)
   by simp
 have ?thesis = (?subset-thesis)
    by (metis StampLattice.sup-qe1 max.commute min.commute sup-stamp.elims
less-eq-stamp.simps
       sup-stamp.simps)
 then show ?thesis
   by simp
qed
lemma sup-least:
 fixes x y z :: stamp
 shows y \le x \Longrightarrow z \le x \Longrightarrow ((y \sqcup z) \le x)
proof -
 fix x y z :: stamp
 assume xlessy: y \leq x
 assume xlessz: z \leq x
 obtain l1 u1 where xdef: x = intstamp l1 u1
   using stamp.exhaust by auto
 obtain l2\ u2 where ydef: y = intstamp\ l2\ u2
   using stamp.exhaust by auto
 obtain l3\ u3 where zdef: z = intstamp\ l3\ u3
   using stamp.exhaust by auto
 have yzdef: y \sqcup z = intstamp (min l2 l3) (max u2 u3)
   (is ?meet = intstamp ?l4 ?u4)
   by (simp add: ydef zdef)
 have s1: \{l2..u2\} \subseteq \{l1..u1\}
   using xlessy by (simp add: ydef xdef)
 have s2: \{l3..u3\} \subseteq \{l1..u1\}
   using xlessz by (simp add: zdef xdef)
 have leq: \{?l4...?u4\} \subseteq \{l1..u1\} (is ?subset-thesis)
  by (metis (no-types, opaque-lifting) inf.orderE inf-stamp.simps max.bounded-iff
```

```
\begin{array}{l} max.cobounded2 \\ min.bounded\text{-}iff\ min.cobounded2\ stamp.inject\ xdef\ xlessy\ ydef\ zdef\ atLeastat-\\ Most-subset-iff \\ xlessz) \\ \text{have}\ (y \sqcup z \leq x) = ?subset\text{-}thesis \\ \text{by}\ (simp\ add:\ xdef\ yzdef) \\ \text{then show}\ (y \sqcup z \leq x) \\ \text{using}\ leq\ \text{by}\ simp \\ \text{qed} \\ \\ \text{instance} \\ \text{apply}\ standard \\ \text{by}\ (simp\ add:\ sup\text{-}least\ sup\text{-}ge2\ sup\text{-}ge1) + \\ \text{end} \\ \end{array}
```

1.2.4 Stamp Bounds

Defines the top and bottom elements of the stamp lattice.

This poses an interesting question as our stamp type is a union of the various Stamp subclasses, e.g. IntegerStamp, ObjectStamp, etc.

Each subclass should preferably have its own unique top and bottom element, i.e. An *IntegerStamp* would have the top element of the full range of integers allowed by the bit width and a bottom of a range with no integers. While the *ObjectStamp* should have *Object* as the top and *Void* as the bottom element.

```
instantiation stamp :: bounded\text{-}lattice
begin

notation bot \ (\bot 50)
notation top \ (\top 50)

definition width\text{-}min :: nat \Rightarrow int64 where width\text{-}min \ bits = -(2 \cap (bits-1))

definition width\text{-}max :: nat \Rightarrow int64 where width\text{-}max \ bits = (2 \cap (bits-1)) - 1

value (sint \ (width\text{-}min \ 64), \ sint \ (width\text{-}max \ 64))
value max\text{-}word::int64

lemma
assumes x = width\text{-}min \ 64
assumes y = width\text{-}max \ 64
shows sint \ x < sint \ y
by (simp \ add: \ assms \ width\text{-}max \ def \ width\text{-}min \ def)
```

Note that this definition is valid for unsigned integers only.

The bottom and top element for signed integers would be (-9223372036854775808, 9223372036854775807).

For unsigned we have (0, 18446744073709551615).

For Java we are likely to be more concerned with signed integers. To use the appropriate bottom and top for signed integers we would need to change our definition of less_eq from $11..u1 \le 12..u2$ to sint $11..sint u1 \le sint 12..sint u2$

We may still find an unsigned integer stamp useful. I plan to investigate the Java code to see if this is useful and then apply the changes to switch to signed integers.

```
definition bot-stamp = intstamp (-1) \theta
definition top-stamp = intstamp 0 (-1)
lemma bot-least:
 fixes a :: stamp
 shows (\bot) \le a
proof -
 obtain min max where bot-def: \bot = intstamp max min
   by (simp add: bot-stamp-def)
 have min < max
   using bot-def word-qt-0 unfolding bot-stamp-def by fastforce
 then have \{max..min\} = \{\}
   by (simp add: bot-def)
 then show ?thesis
   using less-eq-stamp.simps by (simp add: stamp.induct bot-stamp-def)
qed
lemma top-greatest:
 fixes a :: stamp
 shows a \leq (\top)
proof -
 obtain min max where top-def:\top = intstamp min max
   by (simp add: top-stamp-def)
 have max-is-max: \neg(\exists n. n > max)
   \mathbf{by}\ (\mathit{metis}\ \mathit{stamp.inject}\ \mathit{top-def}\ \mathit{top-stamp-def}\ \mathit{word-order.extremum-strict})
 have min-is-min: \neg(\exists n. n < min)
  by (metis not-less-iff-gr-or-eq stamp.inject top-def top-stamp-def word-coorder.not-eq-extremum)
 have \neg(\exists l u. \{min..max\} < \{l..u\})
   by (metis atLeastatMost-psubset-iff not-less min-is-min max-is-max)
 then show ?thesis
   unfolding top-stamp-def using less-eq-stamp.elims(3) by fastforce
qed
instance
 apply standard
 by (simp add: top-greatest bot-least)+
end
```

1.3 Java Stamp Methods

The following are methods from the Java Stamp class, they are the methods primarily used for optimizations.

```
definition is-unrestricted :: stamp \Rightarrow bool where is-unrestricted s = (T = s)

fun is-empty :: stamp \Rightarrow bool where is-empty s = (\bot = s)

fun as-constant :: stamp \Rightarrow Value option where as-constant (intstamp l u) = (if (card \{l..u\}) = 1 then Some (IntVal 64 (SOME x. x \in \{l..u\})) else None)

definition always-distinct :: stamp \Rightarrow stamp \Rightarrow bool where always-distinct stamp1 \Rightarrow stamp2 = (\bot = (stamp1 \sqcap stamp2))

definition never-distinct :: stamp \Rightarrow stamp \Rightarrow bool where never-distinct stamp1 \Rightarrow stamp2 \Rightarrow stamp \Rightarrow bool where stamp1 \Rightarrow stamp2 \Rightarrow stamp3 \Rightarrow stam
```

1.4 Mapping to Values

```
fun valid-value :: stamp => Value => bool where valid-value (intstamp l u) (IntVal b v) = (v \in \{l..u\}) | valid-value (intstamp l u) -= False
```

The *valid-value* function is used to map a stamp instance to the values that are allowed by the stamp.

It would be nice if there was a slightly more integrated way to perform this mapping as it requires some infrastructure to prove some fairly simple properties.

```
lemma bottom-range-empty: \neg(valid\text{-}value\ (\bot)\ v) unfolding bot-stamp-def using valid-value.elims(2) by fastforce lemma join-values: assumes joined = x\text{-}stamp\ \sqcap\ y\text{-}stamp shows valid-value joined x\longleftrightarrow (valid\text{-}value\ x\text{-}stamp\ x\land valid\text{-}value\ y\text{-}stamp\ x) proof (cases x) case UndefVal then show ?thesis using valid-value.elims(2) by auto next case (IntVal b x3) obtain lx\ ux where xdef: x\text{-}stamp\ = intstamp\ lx\ ux
```

```
using stamp.exhaust by auto
  obtain ly uy where ydef: y-stamp = intstamp ly uy
   using stamp.exhaust by auto
  obtain v where x = IntVal b v
   by (simp add: IntVal)
 have joined = intstamp (max lx ly) (min ux uy)
   (is joined = intstamp ?lj ?uj)
   by (simp add: xdef ydef assms)
  then have valid-value joined (IntVal b v) = (v \in \{?lj..?uj\})
   by simp
 then show ?thesis
   using \langle x = IntVal \ b \ v \rangle by (auto simp add: ydef xdef)
next
 case (ObjRef x5)
 then show ?thesis
   using valid-value. elims(2) by auto
next
 case (ObjStr\ x6)
 then show ?thesis
   using valid-value. elims(2) by auto
  case (Array Val x51 x52)
 then show ?thesis
   using valid-value.elims(2) by blast
\mathbf{qed}
lemma disjoint-empty:
 fixes x-stamp y-stamp :: stamp
 \mathbf{assumes} \perp = x\text{-}stamp \sqcap y\text{-}stamp
 shows \neg(valid\text{-}value x\text{-}stamp x \land valid\text{-}value y\text{-}stamp x)
 using bottom-range-empty by (simp add: join-values assms)
experiment begin
A possible equivalent alternative to the definition of less_eq
fun less-eq-alt :: 'a::ord \times 'a \Rightarrow 'a \times 'a \Rightarrow bool where
  less-eq-alt (l1, u1) (l2, u2) = ((\neg l1 \le u1) \lor l2 \le l1 \land u1 \le u2)
Proof equivalence
lemma
 fixes 11 12 u1 u2 :: int
 assumes l1 \le u1 \land l2 \le u2
 shows \{l1..u1\} \subseteq \{l2..u2\} = ((l1 \ge l2) \land (u1 \le u2))
 by (simp add: assms)
lemma
 fixes 11 12 u1 u2 :: int
 shows \{l1..u1\} \subseteq \{l2..u2\} = less-eq-alt\ (l1, u1)\ (l2, u2)
 by simp
```

1.5 Generic Integer Stamp

Experimental definition of integer stamps generically, restricting the datatype to only allow valid ranges and the bottom integer element (max_int..min_int).

```
lemma
 assumes (x::int) > 0
 shows (2 \hat{x})/2 = (2 \hat{x} - 1)
 sorry
definition max-signed-int :: 'a::len word where
 max-signed-int = (2 \cap (LENGTH('a) - 1)) - 1
definition min-signed-int :: 'a::len word where
 min-signed-int = -(2 \cap (LENGTH('a) - 1))
definition int-bottom :: 'a::len word \times 'a word where
 int-bottom = (max-signed-int, min-signed-int)
definition int-top :: 'a::len word \times 'a word where
 int-top = (min-signed-int, max-signed-int)
lemma
 fixes x :: 'a :: len word
 shows sint \ x \leq sint \ (((2 \ \widehat{\ } (LENGTH('a) - 1)) - 1)::'a \ word)
 using sint-greater-eq sorry
value sint (0::1 word)
value sint (1::1 word)
value sint (((2 \cap 0) - 1)::1 \ word)
value sint (((2 \hat{\ } 31) - 1)::32 \ word)
lemma max-signed:
 fixes a :: 'a::len word
 shows sint \ a \leq sint \ (max\text{-}signed\text{-}int::'a \ word)
proof (cases sint a = sint (max-signed-int::'a word))
 case True
 then show ?thesis
   by simp
next
 case False
 have sint\ a < sint\ (max-signed-int::'a\ word)
   using False unfolding max-signed-int-def sorry
 then show ?thesis
```

```
by simp
qed
lemma min-signed:
 fixes a :: 'a::len word
 shows sint \ a \ge sint \ (min\text{-}signed\text{-}int::'a \ word)
 sorry
value max-signed-int :: 32 word
value int-bottom::(32 word \times 32 word)
value sint (2147483647::32 word)
value sint (2147483648::32 word)
typedef (overloaded) ('a::len) intstamp =
  \{bounds :: ('a \ word, 'a \ word) \ prod . ((fst \ bounds) \le s \ (snd \ bounds) \lor bounds = \}
int-bottom)}
proof -
 show ?thesis
   by blast
qed
setup-lifting type-definition-intstamp
lift-definition lower :: ('a::len) intstamp \Rightarrow 'a word
 is prod.fst \circ Rep-intstamp.
lift-definition upper :: ('a::len) intstamp \Rightarrow 'a word
 is prod.snd \circ Rep-intstamp.
lift-definition lower-int :: ('a::len) intstamp \Rightarrow int
 is sint \circ prod.fst.
lift-definition upper-int :: ('a::len) intstamp \Rightarrow int
 is sint \circ prod.snd.
lift-definition range :: ('a::len) intstamp \Rightarrow int set
 is \lambda (l, u). \{sint \ l..sint \ u\}.
lift-definition bounds :: ('a::len) intstamp \Rightarrow ('a word \times 'a word)
 is Rep-intstamp.
lift-definition is-bottom :: ('a::len) intstamp \Rightarrow bool
 is \lambda x. x = int\text{-}bottom.
lift-definition from-bounds :: ('a::len word \times 'a word) \Rightarrow 'a intstamp
 is Abs-intstamp.
instantiation intstamp :: (len) order
begin
```

```
definition less-eq-intstamp :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow bool where
    less-eq-intstamp \ s1 \ s2 = (range \ s1 \subseteq range \ s2)
definition less-intstamp :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow bool where
    less-intstamp s1 \ s2 = (range \ s1 \subset range \ s2)
value int-bottom::(1 word \times 1 word)
value sint (0::1 word)
value sint (1::1 word)
value int-bottom::(2 word \times 2 word)
value sint (1::2 word)
value sint (2::2 word)
value sint ((2 ^(LENGTH(32) - 1) - 1)::32 word) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1) - 1)::32 word)) > sint ((-(2 ^(LENGTH(32) - 1) - 1) - 1) - 1) + (-(2 ^(LENGTH(32) - 1) + (-(2 ^(LENGTH(32) - 1) + (-(2 ^(LENGTH(32) - 1) - 1) + (-(2 ^(LENGT
- 1)))::32 word)
lemma bottom-is-bottom:
   assumes is-bottom s
    shows s \leq a
proof -
    have boundsdef: bounds s = int-bottom
        by (metis assms bounds.transfer is-bottom.rep-eq)
    obtain min \ max \ where \ bounds \ s = (max, \ min)
        by fastforce
    then have max \neq min
     by (metis boundsdef dual-order.eq-iff fst-conv int-bottom-def less-minus-one-simps(1)
max-signed
                min\text{-}signed\ not\text{-}less\ sint\text{-}0\ sint\text{-}n1\ snd\text{-}conv)
    then have sint min < sint max
     by (metis \land bounds \ s = (max, min) \land bounds def \ max-signed \ bounds def \ int-bottom-def
signed-word-eqI
                order.not-eq-order-implies-strict prod.sel(1))
    then have range s = \{\}
        by (simp add: \langle bounds \ s = (max, min) \rangle bounds.transfer range-def)
    then show ?thesis
        by (simp add: StampLattice.less-eq-intstamp-def)
qed
lemma bounds-has-value:
    fixes x y :: int
    assumes x < y
    shows card \{x..y\} > 0
    using assms by simp
\mathbf{lemma}\ bounds\text{-}has\text{-}no\text{-}value:
    fixes x y :: int
    assumes x < y
   shows card \{y..x\} = 0
```

```
by (simp add: assms)
{\bf lemma}\ bottom\text{-}unique:
  fixes a s :: 'a intstamp
  assumes is-bottom s
  shows a \leq s \longleftrightarrow is\text{-bottom } a
proof -
  have \forall x. \ sint \ (fst \ (bounds \ x)) \leq sint \ (snd \ (bounds \ x)) \lor is\text{-bottom} \ x
    using Rep-intstamp by (auto simp add: word-sle-eq is-bottom-def bounds-def)
  then have \forall x. (card (range x)) > 0 \lor is\text{-bottom } x
   by (simp add: bounds.transfer case-prod-beta range-def)
  obtain min max where bounds def: bounds s = (max, min)
   by fastforce
  have nooverlap: sint min < sint max
  \mathbf{by}\ (metis\ assms\ bounds. transfer\ bounds def\ fst\text{-}conv\ int\text{-}bottom\text{-}def\ is\text{-}bottom. rep\text{-}eq
min-signed
     order.not-eq-order-implies-strict\ signed-word-eq I\ sint-0\ snd-conv\ verit-la-disequality
       zero-neq-one max-signed)
  have range s = \{sint \ max..sint \ min\}
   by (simp add: bounds.transfer boundsdef range.rep-eq)
  then have card (range s) = 0
   by (simp add: nooverlap)
  then have \forall x. (card (range x)) > 0 \longrightarrow s < x
    by (auto simp add: less-intstamp-def \langle StampLattice.range\ s = \{sint\ max..sint\ max..sint\ max..sint\ max..sint\}
min\}\rangle)
  then show ?thesis
  by (meson \forall x. \ 0 < card \ (StampLattice.range \ x) \lor is-bottom \ x) \ bottom-is-bottom
less-intstamp-def
       less-eq-intstamp-def leD)
qed
lemma bottom-antisym:
 assumes is-bottom x
 shows x \leq y \Longrightarrow y \leq x \Longrightarrow x = y
  using assms proof (cases is-bottom y)
case True
  then show ?thesis
   by (metis Rep-intstamp-inverse assms is-bottom.rep-eq)
next
  {f case}\ {\it False}
  assume y \leq x
 have \neg (y \leq x)
   by (simp add: assms False bottom-unique)
  then show ?thesis
   by (simp\ add: \langle y \leq x \rangle)
qed
lemma int-antisym:
```

```
fixes x y :: 'a intstamp
 shows x \le y \Longrightarrow y \le x \Longrightarrow x = y
proof -
 \mathbf{fix} \ x :: 'a \ intstamp
 fix y :: 'a intstamp
 assume xlessy: x \leq y
 assume ylessx: y \le x
 obtain l1 u1 where xdef: bounds x = (l1, u1)
   by fastforce
 obtain l2 u2 where ydef: bounds y = (l2, u2)
   by fastforce
 from xlessy have s1: \{sint \ l1..sint \ u1\} \subseteq \{sint \ l2..sint \ u2\} (is ?xlessy)
   using xdef ydef less-eq-intstamp-def by (simp add: range-def bounds-def)
 from ylessx have s2: \{sint \ l2..sint \ u2\} \subseteq \{sint \ l1..sint \ u1\} (is ?ylessx)
   using xdef ydef less-eq-intstamp-def by (simp add: range-def bounds-def)
 show x = y proof (cases is-bottom x)
   case True
   then show ?thesis
     by (simp add: ylessx xlessy bottom-antisym)
  next
   case False
   then show ?thesis
     sorry
 qed
qed
instance
 apply standard
 using less-eq-intstamp-def less-intstamp-def apply (simp; blast)
 by (simp add: int-antisym less-eq-intstamp-def)+
end
value take-bit LENGTH(63) 20::int
value take-bit LENGTH(63) ((-20)::int)
value bit (20::int64) (63::nat)
value bit ((-20)::int64) (63::nat)
value ((-20)::int64) < (20::int64)
value take-bit LENGTH(63) ((-20)::int)
lift-definition smax :: 'a :: len \ word \Rightarrow 'a \ word \Rightarrow 'a \ word
 is \lambda a b. (if (sint a) \leq (sint b) then b else a).
lift-definition smin :: 'a::len \ word \Rightarrow 'a \ word \Rightarrow 'a \ word
 is \lambda a b. (if (sint a) \leq (sint b) then a else b).
instantiation intstamp :: (len) semilattice-inf
begin
```

```
notation inf (infix \sqcap 65)
definition join-bounds: 'a intstamp \Rightarrow 'a intstamp \Rightarrow ('a word \times 'a word) where
 join-bounds\ s1\ s2 = (smax\ (lower\ s1)\ (lower\ s2),\ smin\ (upper\ s1)\ (upper\ s2))
definition join-or-bottom :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow ('a word \times 'a word)
where
 join-or-bottom\ s1\ s2=(let\ bound=(join-bounds\ s1\ s2)\ in
   if \ sint \ (fst \ bound) \ge sint \ (snd \ bound) \ then \ int-bottom \ else \ bound)
definition inf-intstamp :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow 'a intstamp where
  inf-intstamp s1 s2 = from-bounds (join-or-bottom s1 s2)
lemma always-valid:
 fixes s1 s2 :: 'a intstamp
 shows Rep-intstamp (from-bounds (join-or-bottom s1 \ s2)) = join-or-bottom s1 \ s2
 by (smt (23) join-or-bottom-def from-bounds.transfer from-bounds-def mem-Collect-eq
word-sle-eq
     Abs\text{-}intstamp\text{-}inverse)
lemma invalid-join:
  fixes s1 s2 :: 'a intstamp
 assumes bound = join-bounds \ s1 \ s2
 assumes sint (fst bound) \ge sint (snd bound)
 shows from-bounds int-bottom = s1 \sqcap s2
 using assms by (simp add: join-or-bottom-def inf-intstamp-def)
lemma unfold-bounds:
  bounds \ x = (lower \ x, \ upper \ x)
 by (simp add: bounds.transfer lower.rep-eq upper.rep-eq)
lemma int-inf-le1:
 fixes x y :: 'a intstamp
 shows (x \sqcap y) \leq x
proof (cases is-bottom (x \sqcap y))
 case True
 then show ?thesis
   by (simp add: bottom-is-bottom)
next
  case False
 then show ?thesis
  using False proof -
 obtain l1 u1 where xdef: lower x = l1 \land upper x = u1
   by simp
  obtain l2 u2 where ydef: lower y = l2 \land upper y = u2
   by simp
 have joindef: x \sqcap y = from\text{-}bounds ((smax l1 l2, smin u1 u2))
   (is x \sqcap y = from\text{-}bounds (?l3, ?u3))
```

```
by (smt (z3) StampLattice.inf-intstamp-def StampLattice.join-bounds-def al-
ways-valid False
       is-bottom.rep-eq join-or-bottom-def xdef ydef)
 have leq: \{sint ?l3..sint ?u3\} \subseteq \{sint l1..sint u1\}
   bv (smt (z3) atLeastatMost-subset-iff smax.transfer smin.transfer)
 have (x \sqcap y) \le x = (\{sint ?l3..sint ?u3\} \subseteq \{sint l1..sint u1\})
   by (smt (z3) xdef less-eq-intstamp-def StampLattice.always-valid unfold-bounds
ydef range.rep-eq
     StampLattice.join-or-bottom-def bounds.abs-eq case-prod-conv inf-intstamp-def
False
       is-bottom.rep-eq join-bounds-def)
 then show (x \sqcap y) \leq x
   using leq by simp
qed
qed
lemma int-inf-le2:
 fixes x y :: 'a intstamp
 shows (x \sqcap y) \leq y
proof (cases is-bottom (x \sqcap y))
  case True
  then show ?thesis
   by (simp add: bottom-is-bottom)
next
 {f case}\ {\it False}
 then show ?thesis
 using False proof –
 obtain l1 u1 where xdef: lower x = l1 \land upper x = u1
   by simp
 obtain l2 u2 where ydef: lower y = l2 \land upper y = u2
   by simp
 have joindef: x \sqcap y = from\text{-}bounds ((smax l1 l2, smin u1 u2))
   (is x \sqcap y = from\text{-}bounds (?l3, ?u3))
   by (smt (z3) False StampLattice.inf-intstamp-def StampLattice.join-bounds-def
always-valid ydef
       is-bottom.rep-eq_join-or-bottom-def_xdef)
 have leq: \{sint ?l3..sint ?u3\} \subseteq \{sint l1..sint u1\}
   by (smt (z3) atLeastatMost-subset-iff smax.transfer smin.transfer)
 have (x \sqcap y) \leq y = (\{sint ? l3..sint ? u3\} \subseteq \{sint l2..sint u2\})
  by (smt (z3) less-eq-intstamp-def False StampLattice.always-valid unfold-bounds
range.rep-eq
     StampLattice.join-or-bottom-def\ bounds.abs-eq\ case-prod-conv\ inf-intstamp-def
xdef ydef
       is-bottom.rep-eq join-bounds-def)
 then show (x \sqcap y) \leq y
   by (smt (z3) atLeastatMost-subset-iff smax.transfer smin.transfer)
ged
qed
```

```
lemma
  assumes x \leq y
  assumes is-bottom y
  shows is-bottom x
  using assms by (auto simp add: bottom-unique bottom-is-bottom)
\mathbf{lemma} \ \mathit{int-inf-greatest} :
  fixes x y :: 'a intstamp
  \mathbf{shows}\ x \leq y \Longrightarrow x \leq z \Longrightarrow x \leq y \sqcap z
  sorry
instance
  apply standard
  by (simp add: local.int-inf-greatest local.int-inf-le2 local.int-inf-le1)+
end
instantiation intstamp :: (len) semilattice-sup
begin
notation sup (infix \sqcup 65)
instance apply standard sorry
\quad \text{end} \quad
instantiation intstamp :: (len) bounded-lattice
begin
notation bot (\perp 50)
notation top (\top 50)
\mathbf{definition}\ \mathit{bot\text{-}intstamp} = \mathit{int\text{-}bottom}
\mathbf{definition}\ to p\text{-}intstamp = int\text{-}top
instance apply standard sorry
end
value sint (0::1 word)
value sint (1::1 word)
datatype Stamp =
  BottomStamp \mid
  TopStamp \mid
  VoidStamp \mid
  Int8Stamp\ 8\ intstamp\ |
  Int16Stamp 16 intstamp |
```

```
Int64Stamp 64 intstamp
instantiation Stamp :: order
begin
fun less-eq-Stamp :: Stamp <math>\Rightarrow Stamp \Rightarrow bool where
 less-eq-Stamp \ BottomStamp - = True
 less-eq-Stamp - TopStamp = True \mid
 less-eq-Stamp\ VoidStamp\ VoidStamp\ =\ True\ |
 less-eq-Stamp (Int8Stamp v1) (Int8Stamp v2) = (v1 \le v2) |
 less-eq-Stamp (Int16Stamp v1) (Int16Stamp v2) = (v1 \le v2)
 less-eq-Stamp (Int32Stamp v1) (Int32Stamp v2) = (v1 \le v2)
 less-eq-Stamp (Int64Stamp v1) (Int64Stamp v2) = (v1 \le v2) |
 less-eq-Stamp - - = False
fun less-Stamp :: Stamp \Rightarrow Stamp \Rightarrow bool where
 less-Stamp\ BottomStamp\ BottomStamp\ =\ False\ |
 less-Stamp \ BottomStamp -= True \mid
 less-Stamp TopStamp TopStamp = False
 less-Stamp - TopStamp = True
 less-Stamp\ VoidStamp\ VoidStamp\ = False\ |
 less-Stamp (Int8Stamp v1) (Int8Stamp v2) = (v1 < v2)
 less-Stamp (Int16Stamp v1) (Int16Stamp v2) = (v1 < v2)
 less-Stamp (Int32Stamp v1) (Int32Stamp v2) = (v1 < v2)
 less-Stamp (Int64Stamp v1) (Int64Stamp v2) = (v1 < v2)
 less-Stamp - - = False
instance
 apply standard sorry
end
instantiation Stamp :: semilattice-inf
begin
notation inf (infix \sqcap 65)
fun inf-Stamp :: Stamp \Rightarrow Stamp \Rightarrow Stamp where
 inf-Stamp BottomStamp -= BottomStamp
 inf-Stamp - BottomStamp = BottomStamp
 inf-Stamp TopStamp - = TopStamp
 inf-Stamp - TopStamp = TopStamp
 inf-Stamp VoidStamp = VoidStamp |
 inf-Stamp (Int8Stamp v1) (Int8Stamp v2) = Int8Stamp (v1 \sqcap v2) |
 inf-Stamp (Int16Stamp v1) (Int16Stamp v2) = Int16Stamp (v1 \sqcap v2)
 inf-Stamp (Int32Stamp v1) (Int32Stamp v2) = Int32Stamp (v1 \sqcap v2) |
 inf-Stamp (Int64Stamp v1) (Int64Stamp v2) = Int64Stamp (v1 \sqcap v2)
```

Int32Stamp 32 intstamp

instance

```
apply standard sorry
instantiation Stamp :: semilattice-sup
begin
notation sup (infix \sqcup 65)
fun sup-Stamp :: Stamp \Rightarrow Stamp \Rightarrow Stamp where
 sup-Stamp\ BottomStamp\ -=\ BottomStamp
 sup\text{-}Stamp - BottomStamp = BottomStamp
 sup-Stamp TopStamp - = TopStamp
 sup\text{-}Stamp - TopStamp = TopStamp
 sup-Stamp\ VoidStamp\ VoidStamp\ =\ VoidStamp\ |
 sup-Stamp \ (Int8Stamp \ v1) \ (Int8Stamp \ v2) = Int8Stamp \ (v1 \sqcup v2) \mid
 sup\text{-}Stamp \ (Int16Stamp \ v1) \ (Int16Stamp \ v2) = Int16Stamp \ (v1 \sqcup v2) \mid
 sup-Stamp (Int32Stamp v1) (Int32Stamp v2) = Int32Stamp (v1 \sqcup v2) |
 sup\text{-}Stamp \ (Int64Stamp \ v1) \ (Int64Stamp \ v2) = Int64Stamp \ (v1 \sqcup v2)
instance
 apply standard sorry
\mathbf{end}
instantiation Stamp :: bounded-lattice
begin
notation bot (\perp 50)
notation top (\top 50)
definition top-Stamp :: Stamp where
 top	ext{-}Stamp = TopStamp
definition bot-Stamp :: Stamp where
 bot\text{-}Stamp = BottomStamp
instance
 apply standard apply (simp add: bot-Stamp-def)
 by (smt (verit, del-insts) less-eq-Stamp.simps(13) less-eq-Stamp.simps(2) sup.cobounded11
     sup-Stamp.simps(2))
end
lemma [code]: Rep-intstamp (from-bounds (l, u)) = (l, u)
 using Abs-intstamp-inverse from-bounds.rep-eq
 sorry
code-datatype Abs-intstamp
end
```