

Veriopt

April 12, 2023

Abstract

The Veriopt project aims to prove the optimization pass of the GraalVM compiler. The GraalVM compiler includes a sophisticated Intermediate Representation (IR) in the form of a sea-of-nodes based graph structure. We first define the IR graph structure in the Isabelle/HOL interactive theorem prover. We subsequently give the evaluation of the structure a semantics based on the current understanding of the purpose of each IR graph node. Optimization phases are then encoded including the static analysis passes required for an optimization. Each optimization phase is proved to be correct by proving that a bisimulation exists between the unoptimized and optimized graphs. The following document has been automatically generated from the Isabelle/HOL source to provide a very comprehensive definition of the semantics and optimizations introduced by the Veriopt project.

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1 Additional Theorems about Computer Words

theory *JavaWords*

imports

HOL-Library.Word

HOL-Library.Signed-Division

HOL-Library.Float

HOL-Library.LaTeXsugar

begin

Java supports 64, 32, 16, 8 signed ints, plus 1 bit (boolean) ints, and char is 16-bit unsigned. E.g. an 8-bit stamp has a default range of -128..+127. And a 1-bit stamp has a default range of -1..0, surprisingly.

During calculations the smaller sizes are sign-extended to 32 bits.

type-synonym *int64* = 64 word — long

type-synonym *int32* = 32 word — int

type-synonym *int16* = 16 word — short

type-synonym *int8* = 8 word — char

type-synonym *int1* = 1 word — boolean

abbreviation *valid-int-widths* :: nat set **where**

valid-int-widths $\equiv \{1, 8, 16, 32, 64\}$

type-synonym *iwidth* = nat

fun *bit-bounds* :: nat \Rightarrow (int \times int) **where**

bit-bounds bits = (((2 \wedge bits) div 2) * -1, ((2 \wedge bits) div 2) - 1)

definition *logic-negate* :: ('a::len) word \Rightarrow 'a word **where**

logic-negate x = (if x = 0 then 1 else 0)

fun *int-signed-value* :: iwidth \Rightarrow int64 \Rightarrow int **where**

int-signed-value b v = sint (signed-take-bit (b - 1) v)

fun *int-unsigned-value* :: iwidth \Rightarrow int64 \Rightarrow int **where**

int-unsigned-value b v = uint v

A convenience function for directly constructing -1 values of a given bit size.

fun *neg-one* :: iwidth \Rightarrow int64 **where**

neg-one b = mask b

1.1 Bit-Shifting Operators

definition *shiffl* (infix << 75) **where**

shiffl w n = (push-bit n) w

lemma *shiffl-power[simp]*: (x::('a::len) word) * (2 \wedge j) = x << j

unfolding *shiffl-def* **apply** (induction j)

apply simp unfolding funpow-Suc-right
by (*metis (no-types, opaque-lifting) push-bit-eq-mult*)

lemma ($x :: ('a :: \text{len}) \text{ word}$) * $((2^j) + 1) = x \ll j + x$
by (*simp add: distrib-left*)

lemma ($x :: ('a :: \text{len}) \text{ word}$) * $((2^j) - 1) = x \ll j - x$
by (*simp add: right-diff-distrib*)

lemma ($x :: ('a :: \text{len}) \text{ word}$) * $((2^j) + (2^k)) = x \ll j + x \ll k$
by (*simp add: distrib-left*)

lemma ($x :: ('a :: \text{len}) \text{ word}$) * $((2^j) - (2^k)) = x \ll j - x \ll k$
by (*simp add: right-diff-distrib*)

Unsigned shift right.

definition *shiftr* (**infix** $>>>$ 75) **where**
shiftr w n = drop-bit n w

corollary ($255 :: 8 \text{ word}$) $>>> (2 :: \text{nat}) = 63$ **by** *code-simp*

Signed shift right.

definition *sshiftr* :: $'a :: \text{len word} \Rightarrow \text{nat} \Rightarrow 'a :: \text{len word}$ (**infix** $>>$ 75) **where**
sshiftr w n = word-of-int ((sint w) div (2^n))

corollary ($128 :: 8 \text{ word}$) $>> 2 = 0xE0$ **by** *code-simp*

1.2 Fixed-width Word Theories

1.2.1 Support Lemmas for Upper/Lower Bounds

lemma *size32*: $\text{size } v = 32$ **for** $v :: 32 \text{ word}$
using *size-word.rep-eq*
using *One-nat-def add.right-neutral add-Suc-right len-of-numeral-defs(2) len-of-numeral-defs(3)*
mult.right-neutral mult-Suc-right numeral-2-eq-2 numeral-Bit0
by (*smt (verit, del-insts) mult.commute*)

lemma *size64*: $\text{size } v = 64$ **for** $v :: 64 \text{ word}$
using *size-word.rep-eq*
using *One-nat-def add.right-neutral add-Suc-right len-of-numeral-defs(2) len-of-numeral-defs(3)*
mult.right-neutral mult-Suc-right numeral-2-eq-2 numeral-Bit0
by (*smt (verit, del-insts) mult.commute*)

lemma *lower-bounds-equiv*:
assumes $0 < N$
shows $\neg(((2::\text{int}) \wedge (N-1))) = (2::\text{int}) \wedge N \text{ div } 2 * - 1$
by (*simp add: assms int-power-div-base*)

lemma *upper-bounds-equiv*:
assumes $0 < N$
shows $(2::int) \wedge (N-1) = (2::int) \wedge N \text{ div } 2$
by (*simp add: assms int-power-div-base*)

Some min/max bounds for 64-bit words

lemma *bit-bounds-min64*: $((fst (bit-bounds 64))) \leq (sint (v::int64))$
unfolding *bit-bounds.simps fst-def*
using *sint-ge[of v]* **by** *simp*

lemma *bit-bounds-max64*: $((snd (bit-bounds 64))) \geq (sint (v::int64))$
unfolding *bit-bounds.simps fst-def*
using *sint-lt[of v]* **by** *simp*

Extend these min/max bounds to extracting smaller signed words using *signed_take_bit*.

Note: we could use *signed* to convert between bit-widths, instead of *signed_take_bit*. But that would have to be done separately for each bit-width type.

corollary *sint(signed-take-bit 7 (128 :: int8)) = -128* **by** *code-simp*

ML-val $\langle @\{thm\ signed_take_bit_decr_length_iff\} \rangle$
declare $[[show_types=true]]$
ML-val $\langle @\{thm\ signed_take_bit_int_less_exp\} \rangle$

lemma *signed-take-bit-int-less-exp-word*:
fixes *ival* :: 'a :: len word
assumes $n < LENGTH('a)$
shows $sint(signed-take-bit\ n\ ival) < (2::int) \wedge n$
apply *transfer*
by (*smt (verit, best) not-take-bit-negative signed-take-bit-eq-take-bit-shift signed-take-bit-int-less-exp take-bit-int-greater-self-iff*)

lemma *signed-take-bit-int-greater-eq-minus-exp-word*:
fixes *ival* :: 'a :: len word
assumes $n < LENGTH('a)$
shows $-(2 \wedge n) \leq sint(signed-take-bit\ n\ ival)$
apply *transfer*
by (*smt (verit, best) signed-take-bit-int-greater-eq-minus-exp signed-take-bit-int-greater-eq-self-iff signed-take-bit-int-less-exp*)

lemma *signed-take-bit-range*:
fixes *ival* :: 'a :: len word
assumes $n < LENGTH('a)$
assumes $val = sint(signed-take-bit\ n\ ival)$
shows $-(2 \wedge n) \leq val \wedge val < 2 \wedge n$

using *signed-take-bit-int-greater-eq-minus-exp-word signed-take-bit-int-less-exp-word*
using *assms* **by** *blast*

A *bit_bounds* version of the above lemma.

lemma *signed-take-bit-bounds*:
fixes *ival* :: 'a :: len word
assumes $n \leq \text{LENGTH}('a)$
assumes $0 < n$
assumes $\text{val} = \text{sint}(\text{signed-take-bit } (n - 1) \text{ ival})$
shows $\text{fst } (\text{bit-bounds } n) \leq \text{val} \wedge \text{val} \leq \text{snd } (\text{bit-bounds } n)$
using *assms signed-take-bit-range lower-bounds-equiv upper-bounds-equiv*
by (*metis bit-bounds.simps fst-conv less-imp-diff-less nat-less-le sint-ge sint-lt*
snd-conv zle-diff1-eq)

lemma *signed-take-bit-bounds64*:
fixes *ival* :: int64
assumes $n \leq 64$
assumes $0 < n$
assumes $\text{val} = \text{sint}(\text{signed-take-bit } (n - 1) \text{ ival})$
shows $\text{fst } (\text{bit-bounds } n) \leq \text{val} \wedge \text{val} \leq \text{snd } (\text{bit-bounds } n)$
using *assms signed-take-bit-bounds*
by (*metis size64 word-size*)

lemma *int-signed-value-bounds*:
assumes $b1 \leq 64$
assumes $0 < b1$
shows $\text{fst } (\text{bit-bounds } b1) \leq \text{int-signed-value } b1 \text{ v2} \wedge$
 $\text{int-signed-value } b1 \text{ v2} \leq \text{snd } (\text{bit-bounds } b1)$
using *assms int-signed-value.simps signed-take-bit-bounds64* **by** *blast*

lemma *int-signed-value-range*:
fixes *ival* :: int64
assumes $\text{val} = \text{int-signed-value } n \text{ ival}$
shows $-(2 \wedge (n - 1)) \leq \text{val} \wedge \text{val} < 2 \wedge (n - 1)$
using *signed-take-bit-range assms*
by (*smt (verit, ccfv-SIG) One-nat-def diff-less int-signed-value.elims len-gt-0*
len-num1 power-less-imp-less-exp power-strict-increasing sint-greater-eq sint-less)

Some lemmas to relate (int) bit bounds to bit-shifting values.

lemma *bit-bounds-lower*:
assumes $0 < \text{bits}$
shows $\text{word-of-int } (\text{fst } (\text{bit-bounds } \text{bits})) = ((-1) << (\text{bits} - 1))$
unfolding *bit-bounds.simps fst-conv*
by (*metis (mono-tags, opaque-lifting) assms(1) mult-1 mult-minus1-right mult-minus-left*
of-int-minus of-int-power shiftl-power upper-bounds-equiv word-numeral-alt)

lemma *two-exp-div*:
assumes $0 < \text{bits}$


```

shows ((2::int) ^ bits div (2::int)) = (2::int) ^ (bits - Suc 0)
using assms by (auto simp: int-power-div-base)

```

```

declare [[show-types]]

```

Some lemmas about unsigned words smaller than 64-bit, for zero-extend operators.

```

lemma take-bit-smaller-range:
  fixes ival :: 'a :: len word
  assumes n < LENGTH('a)
  assumes val = sint(take-bit n ival)
  shows 0 ≤ val ∧ val < (2::int) ^ n
  by (simp add: assms signed-take-bit-eq)

```

```

lemma take-bit-same-size-nochange:
  fixes ival :: 'a :: len word
  assumes n = LENGTH('a)
  shows ival = take-bit n ival
  by (simp add: assms)

```

A simplification lemma for *new_int*, showing that upper bits can be ignored.

```

lemma take-bit-redundant[simp]:
  fixes ival :: 'a :: len word
  assumes 0 < n
  assumes n < LENGTH('a)
  shows signed-take-bit (n - 1) (take-bit n ival) = signed-take-bit (n - 1) ival
proof -
  have ¬ (n ≤ n - 1) using assms by arith
  then have ∧i . signed-take-bit (n - 1) (take-bit n i) = signed-take-bit (n - 1) i
    using signed-take-bit-take-bit by (metis (mono-tags))
  then show ?thesis
    by blast
qed

```

```

lemma take-bit-same-size-range:
  fixes ival :: 'a :: len word
  assumes n = LENGTH('a)
  assumes ival2 = take-bit n ival
  shows - (2 ^ n div 2) ≤ sint ival2 ∧ sint ival2 < 2 ^ n div 2
  using assms lower-bounds-equiv sint-ge sint-lt by auto

```

```

lemma take-bit-same-bounds:
  fixes ival :: 'a :: len word
  assumes n = LENGTH('a)
  assumes ival2 = take-bit n ival
  shows fst (bit-bounds n) ≤ sint ival2 ∧ sint ival2 ≤ snd (bit-bounds n)
  unfolding bit-bounds.simps
  using assms take-bit-same-size-range
  by force

```

Next we show that casting a word to a wider word preserves any upper/lower bounds. (These lemmas may not be needed any more, since we are not using `scast` now?)

lemma *scast-max-bound*:
assumes $\text{sint } (v :: 'a :: \text{len word}) < M$
assumes $\text{LENGTH}('a) < \text{LENGTH}('b)$
shows $\text{sint } ((\text{scast } v) :: 'b :: \text{len word}) < M$
unfolding $\text{Word.scast-eq Word.sint-sbintrunc'}$
using *Bit-Operations.signed-take-bit-int-eq-self-iff*
by (*smt (verit, best) One-nat-def assms(1) assms(2) decr-length-less-iff linorder-not-le power-strict-increasing-iff signed-take-bit-int-less-self-iff sint-greater-eq*)

lemma *scast-min-bound*:
assumes $M \leq \text{sint } (v :: 'a :: \text{len word})$
assumes $\text{LENGTH}('a) < \text{LENGTH}('b)$
shows $M \leq \text{sint } ((\text{scast } v) :: 'b :: \text{len word})$
unfolding $\text{Word.scast-eq Word.sint-sbintrunc'}$
using *Bit-Operations.signed-take-bit-int-eq-self-iff*
by (*smt (verit) One-nat-def Suc-pred assms(1) assms(2) len-gt-0 less-Suc-eq order-less-le order-less-le-trans power-le-imp-le-exp signed-take-bit-int-greater-eq-self-iff sint-lt*)

lemma *scast-bigger-max-bound*:
assumes $(\text{result} :: 'b :: \text{len word}) = \text{scast } (v :: 'a :: \text{len word})$
shows $\text{sint result} < 2 \wedge \text{LENGTH}('a) \text{ div } 2$
using *sint-lt upper-bounds-equiv scast-max-bound*
by (*smt (verit, best) assms(1) len-gt-0 signed-scast-eq signed-take-bit-int-greater-self-iff sint-ge sint-less upper-bounds-equiv*)

lemma *scast-bigger-min-bound*:
assumes $(\text{result} :: 'b :: \text{len word}) = \text{scast } (v :: 'a :: \text{len word})$
shows $-(2 \wedge \text{LENGTH}('a) \text{ div } 2) \leq \text{sint result}$
using *sint-ge lower-bounds-equiv scast-min-bound*
by (*smt (verit) assms len-gt-0 nat-less-le not-less scast-max-bound*)

lemma *scast-bigger-bit-bounds*:
assumes $(\text{result} :: 'b :: \text{len word}) = \text{scast } (v :: 'a :: \text{len word})$
shows $\text{fst } (\text{bit-bounds } (\text{LENGTH}('a))) \leq \text{sint result} \wedge \text{sint result} \leq \text{snd } (\text{bit-bounds } (\text{LENGTH}('a)))$
using *assms scast-bigger-min-bound scast-bigger-max-bound*
by *auto*

1.2.2 Support lemmas for take bit and signed take bit.

Lemmas for removing redundant `take_bit` wrappers.

lemma *take-bit-dist-addL[simp]*:
fixes $x :: 'a :: \text{len word}$
shows $\text{take-bit } b (\text{take-bit } b x + y) = \text{take-bit } b (x + y)$

```

proof (induction b)
  case 0
  then show ?case
    by simp
next
  case (Suc b)
  then show ?case
    by (simp add: add.commute mask-eqs(2) take-bit-eq-mask)
qed

```

```

lemma take-bit-dist-addR[simp]:
  fixes x :: 'a :: len word
  shows take-bit b (x + take-bit b y) = take-bit b (x + y)
  using take-bit-dist-addL by (metis add.commute)

```

```

lemma take-bit-dist-subL[simp]:
  fixes x :: 'a :: len word
  shows take-bit b (take-bit b x - y) = take-bit b (x - y)
  by (metis take-bit-dist-addR uminus-add-conv-diff)

```

```

lemma take-bit-dist-subR[simp]:
  fixes x :: 'a :: len word
  shows take-bit b (x - take-bit b y) = take-bit b (x - y)
  using take-bit-dist-subL
  by (metis (no-types, opaque-lifting) diff-add-cancel diff-right-commute diff-self)

```

```

lemma take-bit-dist-neg[simp]:
  fixes ix :: 'a :: len word
  shows take-bit b (- take-bit b (ix)) = take-bit b (- ix)
  by (metis diff-0 take-bit-dist-subR)

```

```

lemma signed-take-take-bit[simp]:
  fixes x :: 'a :: len word
  assumes  $0 < b$ 
  shows signed-take-bit (b - 1) (take-bit b x) = signed-take-bit (b - 1) x
  by (smt (verit, best) Suc-diff-1 assms lessI linorder-not-less signed-take-bit-take-bit)

```

```

lemma mod-larger-ignore:
  fixes a :: int
  fixes m n :: nat
  assumes  $n < m$ 
  shows  $(a \bmod 2^m) \bmod 2^n = a \bmod 2^n$ 
  by (smt (verit, del-insts) assms exp-mod-exp linorder-not-le mod-0-imp-dvd mod-mod-cancel mod-self order-less-imp-le)

```

```

lemma mod-dist-over-add:
  fixes a b c :: int64
  fixes n :: nat
  assumes 1:  $0 < n$ 
  assumes 2:  $n < 64$ 
  shows  $(a \bmod 2^n + b) \bmod 2^n = (a + b) \bmod 2^n$ 
proof –
  have 3:  $(0 :: \text{int64}) < 2^n$ 
    using assms by (simp add: size64 word-2p-lem)
  then show ?thesis
    unfolding word-mod-2p-is-mask[OF 3]
    apply transfer
    by (metis (no-types, opaque-lifting) and.right-idem take-bit-add take-bit-eq-mask)
qed

end

```

2 java.lang.Long

Utility functions from the Java Long class that Graal occasionally makes use of.

```

theory JavaLong
  imports JavaWords
           HOL-Library.FSet
begin

lemma negative-all-set-32:
   $n < 32 \implies \text{bit } (-1 :: \text{int32}) \ n$ 
  apply transfer by auto

```

```

definition MaxOrNeg :: nat set  $\Rightarrow$  int where
  MaxOrNeg s = (if s = {} then -1 else Max s)

```

```

definition MinOrHighest :: nat set  $\Rightarrow$  nat  $\Rightarrow$  nat where
  MinOrHighest s m = (if s = {} then m else Min s)

```

```

lemma MaxOrNegEmpty:
   $\text{MaxOrNeg } s = -1 \iff s = \{\}$ 
  unfolding MaxOrNeg-def by auto

```

2.1 Long.highestOneBit

```

definition highestOneBit :: (a::len) word  $\Rightarrow$  int where
  highestOneBit v = MaxOrNeg {n. bit v n}

```

```

lemma highestOneBitInvar:

```

highestOneBit $v = j \implies (\forall i::\text{nat}. (\text{int } i > j \longrightarrow \neg (\text{bit } v \ i)))$
apply (induction size v)
apply simp
by (smt (verit) MaxOrNeg-def Max-ge empty-iff finite-bit-word highestOneBit-def
 mem-Collect-eq of-nat-mono)

lemma highestOneBitNeg:
 highestOneBit $v = -1 \longleftrightarrow v = 0$
unfolding highestOneBit-def MaxOrNeg-def
by (metis Collect-empty-eq-bot bit-0-eq bit-word-eqI int-ops(2) negative-eq-positive
 one-neq-zero)

lemma higherBitsFalse:
fixes $v :: 'a :: \text{len word}$
shows $i > \text{size } v \implies \neg (\text{bit } v \ i)$
by (simp add: bit-word.rep-eq size-word.rep-eq)

lemma highestOneBitN:
assumes bit v n
assumes $\forall i::\text{nat}. (\text{int } i > n \longrightarrow \neg (\text{bit } v \ i))$
shows highestOneBit $v = n$
unfolding highestOneBit-def MaxOrNeg-def
by (metis Max-ge Max-in all-not-in-conv assms(1) assms(2) finite-bit-word mem-Collect-eq
 of-nat-less-iff order-less-le)

lemma highestOneBitSize:
assumes bit v n
assumes $n = \text{size } v$
shows highestOneBit $v = n$
by (metis assms(1) assms(2) not-bit-length wsst-TYs(3))

lemma highestOneBitMax:
 highestOneBit $v < \text{size } v$
unfolding highestOneBit-def MaxOrNeg-def
using higherBitsFalse
by (simp add: bit-imp-le-length size-word.rep-eq)

lemma highestOneBitAtLeast:
assumes bit v n
shows highestOneBit $v \geq n$
proof (induction size v)
case 0
then show ?case **by** simp
next
case (Suc x)
then have $\forall i. \text{bit } v \ i \longrightarrow i < \text{Suc } x$
by (simp add: bit-imp-le-length wsst-TYs(3))

```

then show ?case
  unfolding highestOneBit-def MaxOrNeg-def
  using assms by auto
qed

```

```

lemma highestOneBitElim:
  highestOneBit v = n
   $\implies ((n = -1 \wedge v = 0) \vee (n \geq 0 \wedge \text{bit } v \ n))$ 
  unfolding highestOneBit-def MaxOrNeg-def
  by (metis Max-in finite-bit-word le0 le-minus-one-simps(3) mem-Collect-eq of-nat-0-le-iff
of-nat-eq-iff)

```

A recursive implementation of highestOneBit that is suitable for code generation.

```

fun highestOneBitRec :: nat  $\Rightarrow$  ('a::len) word  $\Rightarrow$  int where
  highestOneBitRec n v =
    (if bit v n then n
     else if n = 0 then -1
     else highestOneBitRec (n - 1) v)

```

```

lemma highestOneBitRecTrue:
  highestOneBitRec n v = j  $\implies j \geq 0 \implies \text{bit } v \ j$ 
proof (induction n)
  case 0
  then show ?case
  by (metis diff-0 highestOneBitRec.simps leD of-nat-0-eq-iff of-nat-0-le-iff zle-diff1-eq)

```

```

next
  case (Suc n)
  then show ?case
  by (metis diff-Suc-1 highestOneBitRec.elims nat.discI nat-int)
qed

```

```

lemma highestOneBitRecN:
  assumes bit v n
  shows highestOneBitRec n v = n
  by (simp add: assms)

```

```

lemma highestOneBitRecMax:
  highestOneBitRec n v  $\leq$  n
  by (induction n; simp)

```

```

lemma highestOneBitRecElim:
  assumes highestOneBitRec n v = j
  shows  $((j = -1 \wedge v = 0) \vee (j \geq 0 \wedge \text{bit } v \ j))$ 
  using assms highestOneBitRecTrue by blast

```

```

lemma highestOneBitRecZero:
  v = 0  $\implies$  highestOneBitRec (size v) v = -1

```

by (*induction rule: highestOneBitRec.induct; simp*)

lemma *highestOneBitRecLess*:

assumes $\neg \text{bit } v \ n$

shows $\text{highestOneBitRec } n \ v = \text{highestOneBitRec } (n - 1) \ v$

using *assms* **by** *force*

Some lemmas that use masks to restrict `highestOneBit` and relate it to `highestOneBitRec`.

lemma *highestOneBitMask*:

assumes $\text{size } v = n$

shows $\text{highestOneBit } v = \text{highestOneBit } (\text{and } v \ (\text{mask } n))$

by (*metis assms dual-order.refl lt2p-lem mask-eq-iff size-word.rep-eq*)

lemma *maskSmaller*:

fixes $v :: 'a :: \text{len word}$

assumes $\neg \text{bit } v \ n$

shows $\text{and } v \ (\text{mask } (\text{Suc } n)) = \text{and } v \ (\text{mask } n)$

unfolding *bit-eq-iff*

by (*metis assms bit-and-iff bit-mask-iff less-Suc-eq*)

lemma *highestOneBitSmaller*:

assumes $\text{size } v = \text{Suc } n$

assumes $\neg \text{bit } v \ n$

shows $\text{highestOneBit } v = \text{highestOneBit } (\text{and } v \ (\text{mask } n))$

by (*metis assms highestOneBitMask maskSmaller*)

lemma *highestOneBitRecMask*:

shows $\text{highestOneBit } (\text{and } v \ (\text{mask } (\text{Suc } n))) = \text{highestOneBitRec } n \ v$

proof (*induction n*)

case *0*

then show *?case*

by (*smt (verit, ccfv-SIG) Word.mask-Suc-0 and-mask-lt-2p and-nonnegative-int-iff*

bit-1-iff bit-and-iff highestOneBitN highestOneBitNeg highestOneBitRec.simps mask-eq-exp-minus-1 of-int-0 uint-1-eq uint-and word-and-def)

next

case $(\text{Suc } n)$

then show *?case*

proof (*cases bit v (Suc n)*)

case *True*

have *1: highestOneBitRec (Suc n) v = Suc n*

by (*simp add: True*)

have $\forall i::\text{nat. } (\text{int } i > (\text{Suc } n) \longrightarrow \neg (\text{bit } (\text{and } v \ (\text{mask } (\text{Suc } (\text{Suc } n)))) \ i))$

by (*simp add: bit-and-iff bit-mask-iff*)

then have *2: highestOneBit (and v (mask (Suc (Suc n)))) = Suc n*

using *True highestOneBitN*

by (*metis bit-take-bit-iff lessI take-bit-eq-mask*)

then show *?thesis*

using *1 2* **by** *auto*

```

next
  case False
  then show ?thesis
    by (simp add: Suc maskSmaller)
qed
qed

```

Finally - we can use the mask lemmas to relate `highestOneBitRec` to its spec.

```

lemma highestOneBitImpl[code]:
  highestOneBit v = highestOneBitRec (size v) v
  by (metis highestOneBitMask highestOneBitRecMask maskSmaller not-bit-length
    wsst-TYs(ℐ))

```

```

lemma highestOneBit (0x5 :: int8) = 2 by code-simp

```

2.2 Long.lowestOneBit

```

definition lowestOneBit :: ('a::len) word ⇒ nat where
  lowestOneBit v = MinOrHighest {n . bit v n} (size v)

```

```

lemma max-bit: bit (v::('a::len) word) n ⇒ n < size v
  by (simp add: bit-imp-le-length size-word.rep-eq)

```

```

lemma max-set-bit: MaxOrNeg {n . bit (v::('a::len) word) n} < Nat.size v
  using max-bit unfolding MaxOrNeg-def
  by force

```

2.3 Long.numberOfLeadingZeros

```

definition numberOfLeadingZeros :: ('a::len) word ⇒ nat where
  numberOfLeadingZeros v = nat (Nat.size v - highestOneBit v - 1)

```

```

lemma MaxOrNeg-neg: MaxOrNeg {} = -1
  by (simp add: MaxOrNeg-def)

```

```

lemma MaxOrNeg-max: s ≠ {} ⇒ MaxOrNeg s = Max s
  by (simp add: MaxOrNeg-def)

```

```

lemma zero-no-bits:
  {n . bit 0 n} = {}
  by simp

```

```

lemma highestOneBit (0::64 word) = -1
  by (simp add: MaxOrNeg-neg highestOneBit-def)

```

```

lemma numberOfLeadingZeros (0::64 word) = 64
  unfolding numberOfLeadingZeros-def using MaxOrNeg-neg highestOneBit-def
  size64

```


by (*smt* (*verit*) *nat-int zero-no-bits*)

lemma *highestOneBit-top*: $\text{Max } \{\text{highestOneBit } (v::64 \text{ word})\} < 64$
unfolding *highestOneBit-def*
by (*metis* *Max-singleton int-eq-iff-numeral max-set-bit size64*)

lemma *numberOfLeadingZeros-top*: $\text{Max } \{\text{numberOfLeadingZeros } (v::64 \text{ word})\} \leq 64$
unfolding *numberOfLeadingZeros-def*
using *size64*
by (*simp* *add: MaxOrNeg-def highestOneBit-def nat-le-iff*)

lemma *numberOfLeadingZeros-range*: $0 \leq \text{numberOfLeadingZeros } a \wedge \text{numberOfLeadingZeros } a \leq \text{Nat.size } a$
unfolding *numberOfLeadingZeros-def*
using *MaxOrNeg-def highestOneBit-def nat-le-iff*
by (*smt* (*verit*) *bot-nat-0.extremum int-eq-iff*)

lemma *leadingZerosAddHighestOne*: $\text{numberOfLeadingZeros } v + \text{highestOneBit } v = \text{Nat.size } v - 1$
unfolding *numberOfLeadingZeros-def highestOneBit-def*
using *MaxOrNeg-def int-nat-eq int-ops(6) max-bit order-less-irrefl* **by** *fastforce*

2.4 Long.numberOfTrailingZeros

definition *numberOfTrailingZeros* :: (*'a::len*) *word* \Rightarrow *nat* **where**
numberOfTrailingZeros *v* = *lowestOneBit* *v*

lemma *lowestOneBit-bot*: $\text{lowestOneBit } (0::64 \text{ word}) = 64$
unfolding *lowestOneBit-def MinOrHighest-def*
by (*simp* *add: size64*)

lemma *bit-zero-set-in-top*: $\text{bit } (-1::'a::len \text{ word}) \ 0$
by *auto*

lemma *nat-bot-set*: $(0::nat) \in xs \longrightarrow (\forall x \in xs . 0 \leq x)$
by *fastforce*

lemma *numberOfTrailingZeros* ($0::64 \text{ word}$) = 64
unfolding *numberOfTrailingZeros-def*
using *lowestOneBit-bot* **by** *simp*

2.5 Long.bitCount

definition *bitCount* :: (*'a::len*) *word* \Rightarrow *nat* **where**
bitCount *v* = *card* {*n* . *bit* *v* *n*}

lemma *bitCount* 0 = 0
unfolding *bitCount-def*
by (*metis* *card.empty zero-no-bits*)

2.6 Long.zeroCount

definition *zeroCount* :: ('a::len) word \Rightarrow nat **where**
zeroCount v = card {n. n < Nat.size v \wedge \neg (bit v n)}

lemma *zeroCount-finite*: finite {n. n < Nat.size v \wedge \neg (bit v n)}
using *finite-nat-set-iff-bounded* **by** *blast*

lemma *negone-set*:
bit (-1::('a::len) word) n \longleftrightarrow n < LENGTH('a)
by *simp*

lemma *negone-all-bits*:
{n . bit (-1::('a::len) word) n} = {n . 0 \leq n \wedge n < LENGTH('a)}
using *negone-set*
by *auto*

lemma *bitCount-finite*:
finite {n . bit (v::('a::len) word) n}
by *simp*

lemma *card-of-range*:
x = card {n . 0 \leq n \wedge n < x}
by *simp*

lemma *range-of-nat*:
{(n::nat) . 0 \leq n \wedge n < x} = {n . n < x}
by *simp*

lemma *finite-range*:
finite {n::nat . n < x}
by *simp*

lemma *range-eq*:
fixes x y :: nat
shows card {y.. x } = card {y<.. x }
using *card-atLeastLessThan* *card-greaterThanAtMost* **by** *presburger*

lemma *card-of-range-bound*:
fixes x y :: nat
assumes x > y
shows x - y = card {n . y < n \wedge n \leq x}
proof –
have *finite*: finite {n . y \leq n \wedge n < x}
by *auto*
have *nonempty*: {n . y \leq n \wedge n < x} \neq {}
using *assms* **by** *blast*
have *simprep*: {n . y < n \wedge n \leq x} = {y<.. x }
by *auto*

```

have  $x - y = \text{card } \{y < ..x\}$ 
  by auto
then show ?thesis
  unfolding simpref by blast
qed

```

```

lemma bitCount (-1::('a::len) word) = LENGTH('a)
  unfolding bitCount-def using card-of-range
  by (metis (no-types, lifting) Collect-cong negone-all-bits)

```

```

lemma bitCount-range:
  fixes  $n :: ('a::len) \text{ word}$ 
  shows  $0 \leq \text{bitCount } n \wedge \text{bitCount } n \leq \text{Nat.size } n$ 
  unfolding bitCount-def
  by (metis atLeastLessThan-iff bot-nat-0.extremum max-bit mem-Collect-eq subsetI
  subset-eq-atLeast0-lessThan-card)

```

```

lemma zerosAboveHighestOne:
   $n > \text{highestOneBit } a \implies \neg(\text{bit } a \ n)$ 
  unfolding highestOneBit-def MaxOrNeg-def
  by (metis (mono-tags, opaque-lifting) Collect-empty-eq Max-ge finite-bit-word
  less-le-not-le mem-Collect-eq of-nat-le-iff)

```

```

lemma zerosBelowLowestOne:
  assumes  $n < \text{lowestOneBit } a$ 
  shows  $\neg(\text{bit } a \ n)$ 
proof (cases  $\{i. \text{bit } a \ i\} = \{\}$ )
  case True
  then show ?thesis by simp
next
  case False
  have  $n < \text{Min } (\text{Collect } (\text{bit } a)) \implies \neg \text{bit } a \ n$ 
    using False by auto
  then show ?thesis
    by (metis False MinOrHighest-def assms lowestOneBit-def)
qed

```

```

lemma union-bit-sets:
  fixes  $a :: ('a::len) \text{ word}$ 
  shows  $\{n . n < \text{Nat.size } a \wedge \text{bit } a \ n\} \cup \{n . n < \text{Nat.size } a \wedge \neg(\text{bit } a \ n)\} = \{n . n < \text{Nat.size } a\}$ 
  by fastforce

```

```

lemma disjoint-bit-sets:
  fixes  $a :: ('a::len) \text{ word}$ 
  shows  $\{n . n < \text{Nat.size } a \wedge \text{bit } a \ n\} \cap \{n . n < \text{Nat.size } a \wedge \neg(\text{bit } a \ n)\} = \{\}$ 
  by blast

```

```

lemma qualified-bitCount:

```

$bitCount\ v = card\ \{n . n < Nat.size\ v \wedge bit\ v\ n\}$
by (*metis* (*no-types*, *lifting*) *Collect-cong bitCount-def max-bit*)

lemma *card-eq*:
assumes *finite x* \wedge *finite y* \wedge *finite z*
assumes $x \cup y = z$
assumes $y \cap x = \{\}$
shows $card\ z - card\ y = card\ x$
using *assms add-diff-cancel-right' card-Un-disjoint*
by (*metis inf.commute*)

lemma *card-add*:
assumes *finite x* \wedge *finite y* \wedge *finite z*
assumes $x \cup y = z$
assumes $y \cap x = \{\}$
shows $card\ x + card\ y = card\ z$
using *assms card-Un-disjoint*
by (*metis inf.commute*)

lemma *card-add-inverses*:
assumes *finite* $\{n. Q\ n \wedge \neg(P\ n)\}$ \wedge *finite* $\{n. Q\ n \wedge P\ n\}$ \wedge *finite* $\{n. Q\ n\}$
shows $card\ \{n. Q\ n \wedge P\ n\} + card\ \{n. Q\ n \wedge \neg(P\ n)\} = card\ \{n. Q\ n\}$
apply (*rule card-add*)
using *assms apply simp*
apply *auto[1]*
by *auto*

lemma *ones-zero-sum-to-width*:
 $bitCount\ a + zeroCount\ a = Nat.size\ a$
proof –
have *add-cards*: $card\ \{n. (\lambda n. n < size\ a)\ n \wedge (bit\ a\ n)\} + card\ \{n. (\lambda n. n < size\ a)\ n \wedge \neg(bit\ a\ n)\} = card\ \{n. (\lambda n. n < size\ a)\ n\}$
apply (*rule card-add-inverses*) **by** *simp*
then have $\dots = Nat.size\ a$
by *auto*
then show *?thesis*
unfolding *bitCount-def zeroCount-def* **using** *max-bit*
by (*metis* (*mono-tags*, *lifting*) *Collect-cong add-cards*)
qed

lemma *intersect-bitCount-helper*:
 $card\ \{n . n < Nat.size\ a\} - bitCount\ a = card\ \{n . n < Nat.size\ a \wedge \neg(bit\ a\ n)\}$
proof –
have *size-def*: $Nat.size\ a = card\ \{n . n < Nat.size\ a\}$
using *card-of-range* **by** *simp*
have *bitCount-def*: $bitCount\ a = card\ \{n . n < Nat.size\ a \wedge bit\ a\ n\}$
using *qualified-bitCount* **by** *auto*
have *disjoint*: $\{n . n < Nat.size\ a \wedge bit\ a\ n\} \cap \{n . n < Nat.size\ a \wedge \neg(bit\ a\ n)\} = \{\}$

```

n))} = {}
  using disjoint-bit-sets by auto
  have union: {n . n < Nat.size a ∧ bit a n} ∪ {n . n < Nat.size a ∧ ¬(bit a n)}
= {n . n < Nat.size a}
  using union-bit-sets by auto
  show ?thesis
  unfolding bitCount-def
  apply (rule card-eq)
  using finite-range apply simp
  using union apply blast
  using disjoint by simp
qed

```

```

lemma intersect-bitCount:
  Nat.size a - bitCount a = card {n . n < Nat.size a ∧ ¬(bit a n)}
  using card-of-range intersect-bitCount-helper by auto

```

```

hide-fact intersect-bitCount-helper

```

```

end

```

3 Operator Semantics

```

theory Values
  imports
    JavaWords
begin

```

In order to properly implement the IR semantics we first introduce a type that represents runtime values. These runtime values represent the full range of primitive types currently allowed by our semantics, ranging from basic integer types to object references and arrays.

Note that Java supports 64, 32, 16, 8 signed ints, plus 1 bit (boolean) ints, and char is 16-bit unsigned. E.g. an 8-bit stamp has a default range of -128..+127. And a 1-bit stamp has a default range of -1..0, surprisingly.

During calculations the smaller sizes are sign-extended to 32 bits, but explicit widening nodes will do that, so most binary calculations should see equal input sizes.

An object reference is an option type where the *None* object reference points to the static fields. This is examined more closely in our definition of the heap.

```

type-synonym objref = nat option

```

```

datatype (discs-sels) Value =
  UndefVal |

```

IntVal iwidth int64 |

ObjRef objref |
ObjStr string

fun *intval-bits* :: *Value* \Rightarrow *nat* **where**
 intval-bits (*IntVal* *b* *v*) = *b*

fun *intval-word* :: *Value* \Rightarrow *int64* **where**
 intval-word (*IntVal* *b* *v*) = *v*

Converts an integer word into a Java value.

fun *new-int* :: *iwidth* \Rightarrow *int64* \Rightarrow *Value* **where**
 new-int *b* *w* = *IntVal* *b* (*take-bit* *b* *w*)

Converts an integer word into a Java value, iff the two types are equal.

fun *new-int-bin* :: *iwidth* \Rightarrow *iwidth* \Rightarrow *int64* \Rightarrow *Value* **where**
 new-int-bin *b1* *b2* *w* = (*if* *b1*=*b2* *then* *new-int* *b1* *w* *else* *UndefVal*)

fun *wf-bool* :: *Value* \Rightarrow *bool* **where**
 wf-bool (*IntVal* *b* *w*) = (*b* = 1) |
 wf-bool - = *False*

fun *val-to-bool* :: *Value* \Rightarrow *bool* **where**
 val-to-bool (*IntVal* *b* *val*) = (*if* *val* = 0 *then* *False* *else* *True*) |
 val-to-bool *val* = *False*

fun *bool-to-val* :: *bool* \Rightarrow *Value* **where**
 bool-to-val *True* = (*IntVal* 32 1) |
 bool-to-val *False* = (*IntVal* 32 0)

Converts an Isabelle bool into a Java value, iff the two types are equal.

fun *bool-to-val-bin* :: *iwidth* \Rightarrow *iwidth* \Rightarrow *bool* \Rightarrow *Value* **where**
 bool-to-val-bin *t1* *t2* *b* = (*if* *t1* = *t2* *then* *bool-to-val* *b* *else* *UndefVal*)

fun *is-int-val* :: *Value* \Rightarrow *bool* **where**
 is-int-val *v* = *is-IntVal* *v*

lemma *neg-one-value[simp]*: *new-int* *b* (*neg-one* *b*) = *IntVal* *b* (*mask* *b*)

by *simp*

lemma *neg-one-signed*[*simp*]:

assumes $0 < b$

shows *int-signed-value* b (*neg-one* b) = -1

by (*smt* (*verit*, *best*) *assms* *diff-le-self* *diff-less* *int-signed-value.simps* *less-one* *mask-eq-take-bit-minus-one* *neg-one.simps* *nle-le* *signed-minus-1* *signed-take-bit-of-minus-1* *signed-take-bit-take-bit* *verit-comp-simplify1*(1))

3.1 Arithmetic Operators

We need to introduce arithmetic operations which agree with the JVM.

Within the JVM, bytecode arithmetic operations are performed on 32 or 64 bit integers, unboxing where appropriate.

The following collection of *intval* functions correspond to the JVM arithmetic operations. We merge the 32 and 64 bit operations into a single function, even though the stamp of each *IRNode* tells us exactly what the bit widths will be. These merged functions make it easier to do the instantiation of *Value* as 'plus', etc. It might be worse for reasoning, because it could cause more case analysis, but this does not seem to be a problem in practice.

fun *intval-add* :: *Value* \Rightarrow *Value* \Rightarrow *Value* **where**

intval-add (*IntVal* $b1$ $v1$) (*IntVal* $b2$ $v2$) =
 (if $b1 = b2$ then *IntVal* $b1$ (*take-bit* $b1$ ($v1+v2$))) else *UndefVal*) |
intval-add - - = *UndefVal*

fun *intval-sub* :: *Value* \Rightarrow *Value* \Rightarrow *Value* **where**

intval-sub (*IntVal* $b1$ $v1$) (*IntVal* $b2$ $v2$) = *new-int-bin* $b1$ $b2$ ($v1-v2$) |
intval-sub - - = *UndefVal*

fun *intval-mul* :: *Value* \Rightarrow *Value* \Rightarrow *Value* **where**

intval-mul (*IntVal* $b1$ $v1$) (*IntVal* $b2$ $v2$) = *new-int-bin* $b1$ $b2$ ($v1*v2$) |
intval-mul - - = *UndefVal*

fun *intval-div* :: *Value* \Rightarrow *Value* \Rightarrow *Value* **where**

intval-div (*IntVal* $b1$ $v1$) (*IntVal* $b2$ $v2$) =
new-int-bin $b1$ $b2$ (*word-of-int*
 ((*int-signed-value* $b1$ $v1$) *sdiv* (*int-signed-value* $b2$ $v2$))) |
intval-div - - = *UndefVal*

fun *intval-mod* :: *Value* \Rightarrow *Value* \Rightarrow *Value* **where**

intval-mod (*IntVal* $b1$ $v1$) (*IntVal* $b2$ $v2$) =

```

    new-int-bin b1 b2 (word-of-int
      ((int-signed-value b1 v1) smod (int-signed-value b2 v2))) |
  intval-mod - - = UndefVal

```

```

fun intval-negate :: Value ⇒ Value where
  intval-negate (IntVal t v) = new-int t (- v) |
  intval-negate - = UndefVal

```

```

fun intval-abs :: Value ⇒ Value where
  intval-abs (IntVal t v) = new-int t (if int-signed-value t v < 0 then - v else v) |
  intval-abs - = UndefVal

```

TODO: clarify which widths this should work on: just 1-bit or all?

```

fun intval-logic-negation :: Value ⇒ Value where
  intval-logic-negation (IntVal b v) = new-int b (logic-negate v) |
  intval-logic-negation - = UndefVal

```

3.2 Bitwise Operators

```

fun intval-and :: Value ⇒ Value ⇒ Value where
  intval-and (IntVal b1 v1) (IntVal b2 v2) = new-int-bin b1 b2 (and v1 v2) |
  intval-and - - = UndefVal

```

```

fun intval-or :: Value ⇒ Value ⇒ Value where
  intval-or (IntVal b1 v1) (IntVal b2 v2) = new-int-bin b1 b2 (or v1 v2) |
  intval-or - - = UndefVal

```

```

fun intval-xor :: Value ⇒ Value ⇒ Value where
  intval-xor (IntVal b1 v1) (IntVal b2 v2) = new-int-bin b1 b2 (xor v1 v2) |
  intval-xor - - = UndefVal

```

```

fun intval-not :: Value ⇒ Value where
  intval-not (IntVal t v) = new-int t (not v) |
  intval-not - = UndefVal

```

3.3 Comparison Operators

```

fun intval-short-circuit-or :: Value ⇒ Value ⇒ Value where
  intval-short-circuit-or (IntVal b1 v1) (IntVal b2 v2) = bool-to-val-bin b1 b2 (((v1
    ≠ 0) ∨ (v2 ≠ 0))) |
  intval-short-circuit-or - - = UndefVal

```

```

fun intval-equals :: Value ⇒ Value ⇒ Value where
  intval-equals (IntVal b1 v1) (IntVal b2 v2) = bool-to-val-bin b1 b2 (v1 = v2) |
  intval-equals - - = UndefVal

```

```

fun intval-less-than :: Value ⇒ Value ⇒ Value where
  intval-less-than (IntVal b1 v1) (IntVal b2 v2) =
    bool-to-val-bin b1 b2 (int-signed-value b1 v1 < int-signed-value b2 v2) |

```


intval-less-than - - = *UndefVal*

fun *intval-below* :: *Value* \Rightarrow *Value* \Rightarrow *Value* **where**
intval-below (*IntVal* *b1* *v1*) (*IntVal* *b2* *v2*) = *bool-to-val-bin* *b1* *b2* (*v1* < *v2*) |
intval-below - - = *UndefVal*

fun *intval-conditional* :: *Value* \Rightarrow *Value* \Rightarrow *Value* \Rightarrow *Value* **where**
intval-conditional *cond* *tv* *fv* = (if (*val-to-bool* *cond*) then *tv* else *fv*)

3.4 Narrowing and Widening Operators

Note: we allow these operators to have *inBits*=*outBits*, because the Graal compiler also seems to allow that case, even though it should rarely / never arise in practice.

Some sanity checks that *take_bitN* and *signed_take_bit*(*N* - 1) match up as expected.

corollary *sint* (*signed-take-bit* 0 (1 :: *int32*)) = -1 **by** *code-simp*
corollary *sint* (*signed-take-bit* 7 ((256 + 128) :: *int64*)) = -128 **by** *code-simp*
corollary *sint* (*take-bit* 7 ((256 + 128 + 64) :: *int64*)) = 64 **by** *code-simp*
corollary *sint* (*take-bit* 8 ((256 + 128 + 64) :: *int64*)) = 128 + 64 **by** *code-simp*

fun *intval-narrow* :: *nat* \Rightarrow *nat* \Rightarrow *Value* \Rightarrow *Value* **where**
intval-narrow *inBits* *outBits* (*IntVal* *b* *v*) =
 (if *inBits* = *b* \wedge 0 < *outBits* \wedge *outBits* \leq *inBits* \wedge *inBits* \leq 64
 then *new-int* *outBits* *v*
 else *UndefVal*) |
intval-narrow - - = *UndefVal*

fun *intval-sign-extend* :: *nat* \Rightarrow *nat* \Rightarrow *Value* \Rightarrow *Value* **where**
intval-sign-extend *inBits* *outBits* (*IntVal* *b* *v*) =
 (if *inBits* = *b* \wedge 0 < *inBits* \wedge *inBits* \leq *outBits* \wedge *outBits* \leq 64
 then *new-int* *outBits* (*signed-take-bit* (*inBits* - 1) *v*)
 else *UndefVal*) |
intval-sign-extend - - = *UndefVal*

fun *intval-zero-extend* :: *nat* \Rightarrow *nat* \Rightarrow *Value* \Rightarrow *Value* **where**
intval-zero-extend *inBits* *outBits* (*IntVal* *b* *v*) =
 (if *inBits* = *b* \wedge 0 < *inBits* \wedge *inBits* \leq *outBits* \wedge *outBits* \leq 64
 then *new-int* *outBits* (*take-bit* *inBits* *v*)
 else *UndefVal*) |
intval-zero-extend - - = *UndefVal*

Some well-formedness results to help reasoning about narrowing and widening operators

lemma *intval-narrow-ok*:

```

assumes intval-narrow inBits outBits val  $\neq$  UndefVal
shows  $0 < \text{outBits} \wedge \text{outBits} \leq \text{inBits} \wedge \text{inBits} \leq 64 \wedge \text{outBits} \leq 64 \wedge$ 
        $\text{is-IntVal } \text{val} \wedge$ 
        $\text{intval-bits } \text{val} = \text{inBits}$ 
using assms intval-narrow.simps neq0-conv intval-bits.simps
by (metis Value.disc(2) intval-narrow.elims le-trans)

```

```

lemma intval-sign-extend-ok:
assumes intval-sign-extend inBits outBits val  $\neq$  UndefVal
shows  $0 < \text{inBits} \wedge$ 
        $\text{inBits} \leq \text{outBits} \wedge \text{outBits} \leq 64 \wedge$ 
        $\text{is-IntVal } \text{val} \wedge$ 
        $\text{intval-bits } \text{val} = \text{inBits}$ 
using assms intval-sign-extend.simps neq0-conv
by (metis intval-bits.simps intval-sign-extend.elims is-IntVal-def)

```

```

lemma intval-zero-extend-ok:
assumes intval-zero-extend inBits outBits val  $\neq$  UndefVal
shows  $0 < \text{inBits} \wedge$ 
        $\text{inBits} \leq \text{outBits} \wedge \text{outBits} \leq 64 \wedge$ 
        $\text{is-IntVal } \text{val} \wedge$ 
        $\text{intval-bits } \text{val} = \text{inBits}$ 
using assms intval-sign-extend.simps neq0-conv
by (metis intval-bits.simps intval-zero-extend.elims is-IntVal-def)

```

3.5 Bit-Shifting Operators

Note that Java shift operators use unary numeric promotion, unlike other binary operators, which use binary numeric promotion (see the Java language reference manual). This means that the left-hand input determines the output size, while the right-hand input can be any size.

```

fun shift-amount :: iwidth  $\Rightarrow$  int64  $\Rightarrow$  nat where
  shift-amount b val = unat (and val (if b = 64 then 0x3F else 0x1f))

```

```

fun intval-left-shift :: Value  $\Rightarrow$  Value  $\Rightarrow$  Value where
  intval-left-shift (IntVal b1 v1) (IntVal b2 v2) = new-int b1 (v1 << shift-amount
b1 v2) |
  intval-left-shift - - = UndefVal

```

Signed shift is more complex, because we sometimes have to insert 1 bits at the correct point, which is at *b1* bits.

```

fun intval-right-shift :: Value  $\Rightarrow$  Value  $\Rightarrow$  Value where
  intval-right-shift (IntVal b1 v1) (IntVal b2 v2) =
    (let shift = shift-amount b1 v2 in
     let ones = and (mask b1) (not (mask (b1 - shift) :: int64)) in
     (if int-signed-value b1 v1 < 0
      then new-int b1 (or ones (v1 >>> shift))
      else new-int b1 (v1 >>> shift))) |

```

intval-right-shift - - = *UndefVal*

fun *intval-uright-shift* :: *Value* \Rightarrow *Value* \Rightarrow *Value* **where**
intval-uright-shift (*IntVal* *b1* *v1*) (*IntVal* *b2* *v2*) = *new-int* *b1* (*v1* >>> *shift-amount*
b1 *v2*) |
intval-uright-shift - - = *UndefVal*

3.5.1 Examples of Narrowing / Widening Functions

experiment begin

corollary *intval-narrow* 32 8 (*IntVal* 32 (256 + 128)) = *IntVal* 8 128 **by** *simp*
corollary *intval-narrow* 32 8 (*IntVal* 32 (-2)) = *IntVal* 8 254 **by** *simp*
corollary *intval-narrow* 32 1 (*IntVal* 32 (-2)) = *IntVal* 1 0 **by** *simp*
corollary *intval-narrow* 32 1 (*IntVal* 32 (-3)) = *IntVal* 1 1 **by** *simp*

corollary *intval-narrow* 32 8 (*IntVal* 64 (-2)) = *UndefVal* **by** *simp*
corollary *intval-narrow* 64 8 (*IntVal* 32 (-2)) = *UndefVal* **by** *simp*
corollary *intval-narrow* 64 8 (*IntVal* 64 254) = *IntVal* 8 254 **by** *simp*
corollary *intval-narrow* 64 8 (*IntVal* 64 (256+127)) = *IntVal* 8 127 **by** *simp*
corollary *intval-narrow* 64 64 (*IntVal* 64 (-2)) = *IntVal* 64 (-2) **by** *simp*
end

experiment begin

corollary *intval-sign-extend* 8 32 (*IntVal* 8 (256 + 128)) = *IntVal* 32 ($2^{32} - 128$) **by** *simp*
corollary *intval-sign-extend* 8 32 (*IntVal* 8 (-2)) = *IntVal* 32 ($2^{32} - 2$) **by** *simp*
corollary *intval-sign-extend* 1 32 (*IntVal* 1 (-2)) = *IntVal* 32 0 **by** *simp*
corollary *intval-sign-extend* 1 32 (*IntVal* 1 (-3)) = *IntVal* 32 (*mask* 32) **by** *simp*

corollary *intval-sign-extend* 8 32 (*IntVal* 64 254) = *UndefVal* **by** *simp*
corollary *intval-sign-extend* 8 64 (*IntVal* 32 254) = *UndefVal* **by** *simp*
corollary *intval-sign-extend* 8 64 (*IntVal* 8 254) = *IntVal* 64 (-2) **by** *simp*
corollary *intval-sign-extend* 32 64 (*IntVal* 32 ($2^{32} - 2$)) = *IntVal* 64 (-2) **by** *simp*
corollary *intval-sign-extend* 64 64 (*IntVal* 64 (-2)) = *IntVal* 64 (-2) **by** *simp*
end

experiment begin

corollary *intval-zero-extend* 8 32 (*IntVal* 8 (256 + 128)) = *IntVal* 32 128 **by** *simp*
corollary *intval-zero-extend* 8 32 (*IntVal* 8 (-2)) = *IntVal* 32 254 **by** *simp*
corollary *intval-zero-extend* 1 32 (*IntVal* 1 (-1)) = *IntVal* 32 1 **by** *simp*
corollary *intval-zero-extend* 1 32 (*IntVal* 1 (-2)) = *IntVal* 32 0 **by** *simp*

corollary *intval-zero-extend* 8 32 (*IntVal* 64 (-2)) = *UndefVal* **by** *simp*
corollary *intval-zero-extend* 8 64 (*IntVal* 64 (-2)) = *UndefVal* **by** *simp*

```

corollary intval-zero-extend 8 64 (IntVal 8 254) = IntVal 64 254 by simp
corollary intval-zero-extend 32 64 (IntVal 32 ( $2^{32} - 2$ )) = IntVal 64 ( $2^{32} - 2$ ) by simp
corollary intval-zero-extend 64 64 (IntVal 64 (-2)) = IntVal 64 (-2) by simp
end

```

experiment begin

```

corollary intval-right-shift (IntVal 8 128) (IntVal 8 0) = IntVal 8 128 by eval
corollary intval-right-shift (IntVal 8 128) (IntVal 8 1) = IntVal 8 192 by eval
corollary intval-right-shift (IntVal 8 128) (IntVal 8 2) = IntVal 8 224 by eval
corollary intval-right-shift (IntVal 8 128) (IntVal 8 8) = IntVal 8 255 by eval
corollary intval-right-shift (IntVal 8 128) (IntVal 8 31) = IntVal 8 255 by eval
end

```

lemma *intval-add-sym*:

```

shows intval-add a b = intval-add b a
by (induction a; induction b; auto simp: add.commute)

```

```

lemma intval-add (IntVal 32 ( $2^{31}-1$ )) (IntVal 32 ( $2^{31}-1$ )) = IntVal 32 ( $2^{32} - 2$ )
by eval
lemma intval-add (IntVal 64 ( $2^{31}-1$ )) (IntVal 64 ( $2^{31}-1$ )) = IntVal 64 4294967294
by eval

```

end

3.6 Fixed-width Word Theories

```

theory ValueThms
imports Values
begin

```

3.6.1 Support Lemmas for Upper/Lower Bounds

```

lemma size32: size v = 32 for v :: 32 word
using size-word.rep-eq
using One-nat-def add.right-neutral add-Suc-right len-of-numeral-defs(2) len-of-numeral-defs(3)
mult.right-neutral mult-Suc-right numeral-2-eq-2 numeral-Bit0
by (smt (verit, del-insts) mult.commute)

lemma size64: size v = 64 for v :: 64 word
using size-word.rep-eq

```

```

using One-nat-def add.right-neutral add-Suc-right len-of-numeral-defs(2) len-of-numeral-defs(3)
mult.right-neutral mult-Suc-right numeral-2-eq-2 numeral-Bit0
by (smt (verit, del-ists) mult commute)

```

```

lemma lower-bounds-equiv:
  assumes 0 < N
  shows -(((2::int) ^ (N-1))) = (2::int) ^ N div 2 * - 1
  by (simp add: asms int-power-div-base)

```

```

lemma upper-bounds-equiv:
  assumes 0 < N
  shows (2::int) ^ (N-1) = (2::int) ^ N div 2
  by (simp add: asms int-power-div-base)

```

Some min/max bounds for 64-bit words

```

lemma bit-bounds-min64: ((fst (bit-bounds 64))) ≤ (sint (v::int64))
  unfolding bit-bounds.simps fst-def
  using sint-ge[of v] by simp

```

```

lemma bit-bounds-max64: ((snd (bit-bounds 64))) ≥ (sint (v::int64))
  unfolding bit-bounds.simps fst-def
  using sint-lt[of v] by simp

```

Extend these min/max bounds to extracting smaller signed words using `signed_take_bit`.

Note: we could use `signed` to convert between bit-widths, instead of `signed_take_bit`. But that would have to be done separately for each bit-width type.

```

value sint(signed-take-bit 7 (128 :: int8))

```

```

ML-val <@{thm signed-take-bit-decr-length-iff}>
declare [[show-types=true]]
ML-val <@{thm signed-take-bit-int-less-exp}>

```

```

lemma signed-take-bit-int-less-exp-word:
  fixes ival :: 'a :: len word
  assumes n < LENGTH('a)
  shows sint(signed-take-bit n ival) < (2::int) ^ n
  apply transfer
  by (smt (verit, best) not-take-bit-negative signed-take-bit-eq-take-bit-shift
    signed-take-bit-int-less-exp take-bit-int-greater-self-iff)

```

```

lemma signed-take-bit-int-greater-eq-minus-exp-word:
  fixes ival :: 'a :: len word
  assumes n < LENGTH('a)

```

```

shows  $-(2 \wedge n) \leq \text{sint}(\text{signed-take-bit } n \text{ ival})$ 
apply transfer
by (smt (verit, best) signed-take-bit-int-greater-eq-minus-exp
      signed-take-bit-int-greater-eq-self-iff signed-take-bit-int-less-exp)

```

```

lemma signed-take-bit-range:
  fixes ival :: 'a :: len word
  assumes  $n < \text{LENGTH}('a)$ 
  assumes  $\text{val} = \text{sint}(\text{signed-take-bit } n \text{ ival})$ 
  shows  $-(2 \wedge n) \leq \text{val} \wedge \text{val} < 2 \wedge n$ 
  using signed-take-bit-int-greater-eq-minus-exp-word signed-take-bit-int-less-exp-word
  using assms by blast

```

A *bit_bounds* version of the above lemma.

```

lemma signed-take-bit-bounds:
  fixes ival :: 'a :: len word
  assumes  $n \leq \text{LENGTH}('a)$ 
  assumes  $0 < n$ 
  assumes  $\text{val} = \text{sint}(\text{signed-take-bit } (n - 1) \text{ ival})$ 
  shows  $\text{fst } (\text{bit-bounds } n) \leq \text{val} \wedge \text{val} \leq \text{snd } (\text{bit-bounds } n)$ 
  using assms signed-take-bit-range lower-bounds-equiv upper-bounds-equiv
  by (metis bit-bounds.simps fst-conv less-imp-diff-less nat-less-le sint-ge sint-lt
      snd-conv zle-diff1-eq)

```

```

lemma signed-take-bit-bounds64:
  fixes ival :: int64
  assumes  $n \leq 64$ 
  assumes  $0 < n$ 
  assumes  $\text{val} = \text{sint}(\text{signed-take-bit } (n - 1) \text{ ival})$ 
  shows  $\text{fst } (\text{bit-bounds } n) \leq \text{val} \wedge \text{val} \leq \text{snd } (\text{bit-bounds } n)$ 
  using assms signed-take-bit-bounds
  by (metis size64 word-size)

```

```

lemma int-signed-value-bounds:
  assumes  $b1 \leq 64$ 
  assumes  $0 < b1$ 
  shows  $\text{fst } (\text{bit-bounds } b1) \leq \text{int-signed-value } b1 \text{ v2} \wedge$ 
       $\text{int-signed-value } b1 \text{ v2} \leq \text{snd } (\text{bit-bounds } b1)$ 
  using assms int-signed-value.simps signed-take-bit-bounds64 by blast

```

```

lemma int-signed-value-range:
  fixes ival :: int64
  assumes  $\text{val} = \text{int-signed-value } n \text{ ival}$ 
  shows  $-(2 \wedge (n - 1)) \leq \text{val} \wedge \text{val} < 2 \wedge (n - 1)$ 
  using signed-take-bit-range assms
  by (smt (verit, ccfv-SIG) One-nat-def diff-less int-signed-value.elims len-gt-0
      len-num1 power-less-imp-less-exp power-strict-increasing sint-greater-eq sint-less)

```

Some lemmas about unsigned words smaller than 64-bit, for zero-extend

operators.

```
lemma take-bit-smaller-range:
  fixes ival :: 'a :: len word
  assumes n < LENGTH('a)
  assumes val = sint(take-bit n ival)
  shows 0 ≤ val ∧ val < (2::int) ^ n
  by (simp add: assms signed-take-bit-eq)
```

```
lemma take-bit-same-size-nochange:
  fixes ival :: 'a :: len word
  assumes n = LENGTH('a)
  shows ival = take-bit n ival
  by (simp add: assms)
```

A simplification lemma for `new_int`, showing that upper bits can be ignored.

```
lemma take-bit-redundant[simp]:
  fixes ival :: 'a :: len word
  assumes 0 < n
  assumes n < LENGTH('a)
  shows signed-take-bit (n - 1) (take-bit n ival) = signed-take-bit (n - 1) ival
proof -
  have ¬ (n ≤ n - 1) using assms by arith
  then have ∧i . signed-take-bit (n - 1) (take-bit n i) = signed-take-bit (n - 1) i
    using signed-take-bit-take-bit by (metis (mono-tags))
  then show ?thesis
    by blast
qed
```

```
lemma take-bit-same-size-range:
  fixes ival :: 'a :: len word
  assumes n = LENGTH('a)
  assumes ival2 = take-bit n ival
  shows - (2 ^ n div 2) ≤ sint ival2 ∧ sint ival2 < 2 ^ n div 2
  using assms lower-bounds-equiv sint-ge sint-lt by auto
```

```
lemma take-bit-same-bounds:
  fixes ival :: 'a :: len word
  assumes n = LENGTH('a)
  assumes ival2 = take-bit n ival
  shows fst (bit-bounds n) ≤ sint ival2 ∧ sint ival2 ≤ snd (bit-bounds n)
  unfolding bit-bounds.simps
  using assms take-bit-same-size-range
  by force
```

Next we show that casting a word to a wider word preserves any upper/lower bounds. (These lemmas may not be needed any more, since we are not using `scast` now?)

```
lemma scast-max-bound:
```

```

assumes sint (v :: 'a :: len word) < M
assumes LENGTH('a) < LENGTH('b)
shows sint ((scast v) :: 'b :: len word) < M
unfolding Word.scast-eq Word.sint-sbintrunc'
using Bit-Operations.signed-take-bit-int-eq-self-iff
by (smt (verit, best) One-nat-def assms(1) assms(2) decr-length-less-iff linorder-not-le
power-strict-increasing-iff signed-take-bit-int-less-self-iff sint-greater-eq)

```

lemma *scast-min-bound*:

```

assumes M ≤ sint (v :: 'a :: len word)
assumes LENGTH('a) < LENGTH('b)
shows M ≤ sint ((scast v) :: 'b :: len word)
unfolding Word.scast-eq Word.sint-sbintrunc'
using Bit-Operations.signed-take-bit-int-eq-self-iff
by (smt (verit) One-nat-def Suc-pred assms(1) assms(2) len-gt-0 less-Suc-eq or-
der-less-le order-less-le-trans power-le-imp-le-exp signed-take-bit-int-greater-eq-self-iff
sint-lt)

```

lemma *scast-bigger-max-bound*:

```

assumes (result :: 'b :: len word) = scast (v :: 'a :: len word)
shows sint result < 2 ^ LENGTH('a) div 2
using sint-lt upper-bounds-equiv scast-max-bound
by (smt (verit, best) assms(1) len-gt-0 signed-scast-eq signed-take-bit-int-greater-self-iff
sint-ge sint-less upper-bounds-equiv)

```

lemma *scast-bigger-min-bound*:

```

assumes (result :: 'b :: len word) = scast (v :: 'a :: len word)
shows − (2 ^ LENGTH('a) div 2) ≤ sint result
using sint-ge lower-bounds-equiv scast-min-bound
by (smt (verit) assms len-gt-0 nat-less-le not-less scast-max-bound)

```

lemma *scast-bigger-bit-bounds*:

```

assumes (result :: 'b :: len word) = scast (v :: 'a :: len word)
shows fst (bit-bounds (LENGTH('a))) ≤ sint result ∧ sint result ≤ snd (bit-bounds
(LENGTH('a)))
using assms scast-bigger-min-bound scast-bigger-max-bound
by auto

```

Results about *new_int*.

lemma *new-int-take-bits*:

```

assumes IntVal b val = new-int b ival
shows take-bit b val = val
using assms by force

```

3.6.2 Support lemmas for take bit and signed take bit.

Lemmas for removing redundant *take_bit* wrappers.

lemma *take-bit-dist-addL[simp]*:


```

fixes  $x :: 'a :: \text{len word}$ 
shows  $\text{take-bit } b (\text{take-bit } b \ x + y) = \text{take-bit } b (x + y)$ 
proof (induction  $b$ )
  case  $0$ 
  then show  $?case$ 
    by simp
next
  case (Suc  $b$ )
  then show  $?case$ 
    by (simp add: add.commute mask-eqs(2) take-bit-eq-mask)
qed

```

```

lemma take-bit-dist-addR[simp]:
  fixes  $x :: 'a :: \text{len word}$ 
  shows  $\text{take-bit } b (x + \text{take-bit } b \ y) = \text{take-bit } b (x + y)$ 
  using take-bit-dist-addL by (metis add.commute)

```

```

lemma take-bit-dist-subL[simp]:
  fixes  $x :: 'a :: \text{len word}$ 
  shows  $\text{take-bit } b (\text{take-bit } b \ x - y) = \text{take-bit } b (x - y)$ 
  by (metis take-bit-dist-addR uminus-add-conv-diff)

```

```

lemma take-bit-dist-subR[simp]:
  fixes  $x :: 'a :: \text{len word}$ 
  shows  $\text{take-bit } b (x - \text{take-bit } b \ y) = \text{take-bit } b (x - y)$ 
  using take-bit-dist-subL
  by (metis (no-types, opaque-lifting) diff-add-cancel diff-right-commute diff-self)

```

```

lemma take-bit-dist-neg[simp]:
  fixes  $ix :: 'a :: \text{len word}$ 
  shows  $\text{take-bit } b (- \text{take-bit } b \ (ix)) = \text{take-bit } b (- ix)$ 
  by (metis diff-0 take-bit-dist-subR)

```

```

lemma signed-take-bit[simp]:
  fixes  $x :: 'a :: \text{len word}$ 
  assumes  $0 < b$ 
  shows  $\text{signed-take-bit } (b - 1) (\text{take-bit } b \ x) = \text{signed-take-bit } (b - 1) \ x$ 
  by (smt (verit, best) Suc-diff-1 assms lessI linorder-not-less signed-take-bit-take-bit)

```

```

lemma mod-larger-ignore:
  fixes  $a :: \text{int}$ 
  fixes  $m \ n :: \text{nat}$ 
  assumes  $n < m$ 
  shows  $(a \bmod 2^m) \bmod 2^n = a \bmod 2^n$ 
  by (smt (verit, del-insts) assms exp-mod-exp linorder-not-le mod-0-imp-dvd mod-mod-cancel)

```

mod-self order-less-imp-le)

```

lemma mod-dist-over-add:
  fixes a b c :: int64
  fixes n :: nat
  assumes 1:  $0 < n$ 
  assumes 2:  $n < 64$ 
  shows  $(a \bmod 2^n + b) \bmod 2^n = (a + b) \bmod 2^n$ 
proof –
  have 3:  $(0 :: \text{int64}) < 2^n$ 
    using assms by (simp add: size64 word-2p-lem)
  then show ?thesis
    unfolding word-mod-2p-is-mask[OF 3]
    apply transfer
    by (metis (no-types, opaque-lifting) and.right-idem take-bit-add take-bit-eq-mask)
qed

end

```

4 Stamp Typing

```

theory Stamp
  imports Values
begin

```

The GraalVM compiler uses the Stamp class to store range and type information for a given node in the IR graph. We model the Stamp class as a datatype, Stamp, and provide a number of functions on the datatype which correspond to the class methods within the compiler.

Stamp information is used in a variety of ways in optimizations, and so, we additionally provide a number of lemmas which help to prove future optimizations.

```

datatype Stamp =
  VoidStamp
  | IntegerStamp (stp-bits: nat) (stp-lower: int) (stp-upper: int)

  | KlassPointerStamp (stp-nonNull: bool) (stp-alwaysNull: bool)
  | MethodCountersPointerStamp (stp-nonNull: bool) (stp-alwaysNull: bool)
  | MethodPointersStamp (stp-nonNull: bool) (stp-alwaysNull: bool)
  | ObjectStamp (stp-type: string) (stp-exactType: bool) (stp-nonNull: bool) (stp-alwaysNull:
bool)
  | RawPointerStamp (stp-nonNull: bool) (stp-alwaysNull: bool)
  | IllegalStamp

```

```

fun is-stamp-empty :: Stamp  $\Rightarrow$  bool where
  is-stamp-empty (IntegerStamp b lower upper) = (upper < lower) |

```

is-stamp-empty $x = \text{False}$

Just like the `IntegerStamp` class, we need to know that our lo/hi bounds fit into the given number of bits (either signed or unsigned). Our integer stamps have infinite lo/hi bounds, so if the lower bound is non-negative, we can assume that all values are positive, and the integer bits of a related value can be interpreted as unsigned. This is similar (but slightly more general) to what `IntegerStamp.java` does with its test: `if (sameSignBounds())` in the `unsignedUpperBound()` method.

Note that this is a bit different and more accurate than what `StampFactory.forUnsignedInteger` does (it widens large unsigned ranges to the max signed range to allow all bit patterns) because its lo/hi values are only 64-bit.

```
fun valid-stamp :: Stamp  $\Rightarrow$  bool where
  valid-stamp (IntegerStamp bits lo hi) =
    ( $0 < \text{bits} \wedge \text{bits} \leq 64 \wedge$ 
      $\text{fst}(\text{bit-bounds } \text{bits}) \leq \text{lo} \wedge \text{lo} \leq \text{snd}(\text{bit-bounds } \text{bits}) \wedge$ 
      $\text{fst}(\text{bit-bounds } \text{bits}) \leq \text{hi} \wedge \text{hi} \leq \text{snd}(\text{bit-bounds } \text{bits})) \mid$ 
    valid-stamp s = True
```

experiment begin

corollary *bit-bounds* $1 = (-1, 0)$ **by** *simp*

end

— A stamp which includes the full range of the type

```
fun unrestricted-stamp :: Stamp  $\Rightarrow$  Stamp where
  unrestricted-stamp VoidStamp = VoidStamp |
  unrestricted-stamp (IntegerStamp bits lower upper) = (IntegerStamp bits ( $\text{fst}$ 
    ( $\text{bit-bounds } \text{bits}$ ) ( $\text{snd}(\text{bit-bounds } \text{bits})$ ))) |

  unrestricted-stamp (KlassPointerStamp nonNull alwaysNull) = (KlassPointerStamp
    False False) |
  unrestricted-stamp (MethodCountersPointerStamp nonNull alwaysNull) = (MethodCountersPointerStamp
    False False) |
  unrestricted-stamp (MethodPointersStamp nonNull alwaysNull) = (MethodPointersStamp
    False False) |
  unrestricted-stamp (ObjectStamp type exactType nonNull alwaysNull) = (ObjectStamp
    "" False False False) |
  unrestricted-stamp - = IllegalStamp
```

fun *is-stamp-unrestricted* :: *Stamp* \Rightarrow *bool* **where**
is-stamp-unrestricted *s* = (*s* = *unrestricted-stamp* *s*)

— A stamp which provides type information but has an empty range of values

fun *empty-stamp* :: *Stamp* \Rightarrow *Stamp* **where**
empty-stamp *VoidStamp* = *VoidStamp* |
empty-stamp (*IntegerStamp* *bits* *lower* *upper*) = (*IntegerStamp* *bits* (*snd* (*bit-bounds* *bits*)) (*fst* (*bit-bounds* *bits*))) |
empty-stamp (*KlassPointerStamp* *nonNull* *alwaysNull*) = (*KlassPointerStamp* *nonNull* *alwaysNull*) |
empty-stamp (*MethodCountersPointerStamp* *nonNull* *alwaysNull*) = (*MethodCountersPointerStamp* *nonNull* *alwaysNull*) |
empty-stamp (*MethodPointersStamp* *nonNull* *alwaysNull*) = (*MethodPointersStamp* *nonNull* *alwaysNull*) |
empty-stamp (*ObjectStamp* *type* *exactType* *nonNull* *alwaysNull*) = (*ObjectStamp* *type* *exactType* *nonNull* *alwaysNull*) |
empty-stamp *stamp* = *IllegalStamp*

— Calculate the meet stamp of two stamps

fun *meet* :: *Stamp* \Rightarrow *Stamp* \Rightarrow *Stamp* **where**
meet *VoidStamp* *VoidStamp* = *VoidStamp* |
meet (*IntegerStamp* *b1* *l1* *u1*) (*IntegerStamp* *b2* *l2* *u2*) = (
 if *b1* \neq *b2* *then* *IllegalStamp* *else*
 (*IntegerStamp* *b1* (*min* *l1* *l2*) (*max* *u1* *u2*))
) |
meet (*KlassPointerStamp* *nn1* *an1*) (*KlassPointerStamp* *nn2* *an2*) = (
 KlassPointerStamp (*nn1* \wedge *nn2*) (*an1* \wedge *an2*)
) |
meet (*MethodCountersPointerStamp* *nn1* *an1*) (*MethodCountersPointerStamp* *nn2* *an2*) = (
 MethodCountersPointerStamp (*nn1* \wedge *nn2*) (*an1* \wedge *an2*)
) |
meet (*MethodPointersStamp* *nn1* *an1*) (*MethodPointersStamp* *nn2* *an2*) = (
 MethodPointersStamp (*nn1* \wedge *nn2*) (*an1* \wedge *an2*)
) |
meet *s1* *s2* = *IllegalStamp*

— Calculate the join stamp of two stamps

fun *join* :: *Stamp* \Rightarrow *Stamp* \Rightarrow *Stamp* **where**
join *VoidStamp* *VoidStamp* = *VoidStamp* |
join (*IntegerStamp* *b1* *l1* *u1*) (*IntegerStamp* *b2* *l2* *u2*) = (
 if *b1* \neq *b2* *then* *IllegalStamp* *else*
 (*IntegerStamp* *b1* (*max* *l1* *l2*) (*min* *u1* *u2*))
) |
join (*KlassPointerStamp* *nn1* *an1*) (*KlassPointerStamp* *nn2* *an2*) = (
 KlassPointerStamp (*nn1* \vee *nn2*) (*an1* \vee *an2*)
) |
join *s1* *s2* = *IllegalStamp*

```

    if ((nn1 ∨ nn2) ∧ (an1 ∨ an2))
    then (empty-stamp (KlassPointerStamp nn1 an1))
    else (KlassPointerStamp (nn1 ∨ nn2) (an1 ∨ an2))
  ) |
  join (MethodCountersPointerStamp nn1 an1) (MethodCountersPointerStamp nn2
an2) = (
    if ((nn1 ∨ nn2) ∧ (an1 ∨ an2))
    then (empty-stamp (MethodCountersPointerStamp nn1 an1))
    else (MethodCountersPointerStamp (nn1 ∨ nn2) (an1 ∨ an2))
  ) |
  join (MethodPointersStamp nn1 an1) (MethodPointersStamp nn2 an2) = (
    if ((nn1 ∨ nn2) ∧ (an1 ∨ an2))
    then (empty-stamp (MethodPointersStamp nn1 an1))
    else (MethodPointersStamp (nn1 ∨ nn2) (an1 ∨ an2))
  ) |
  join s1 s2 = IllegalStamp

```

— In certain circumstances a stamp provides enough information to evaluate a value as a stamp, the `asConstant` function converts the stamp to a value where one can be inferred.

```

fun asConstant :: Stamp ⇒ Value where
  asConstant (IntegerStamp b l h) = (if l = h then IntVal b (word-of-int l) else
UndefVal) |
  asConstant - = UndefVal

```

— Determine if two stamps never have value overlaps i.e. their join is empty

```

fun alwaysDistinct :: Stamp ⇒ Stamp ⇒ bool where
  alwaysDistinct stamp1 stamp2 = is-stamp-empty (join stamp1 stamp2)

```

— Determine if two stamps must always be the same value i.e. two equal constants

```

fun neverDistinct :: Stamp ⇒ Stamp ⇒ bool where
  neverDistinct stamp1 stamp2 = (asConstant stamp1 = asConstant stamp2 ∧
asConstant stamp1 ≠ UndefVal)

```

```

fun constantAsStamp :: Value ⇒ Stamp where
  constantAsStamp (IntVal b v) = (IntegerStamp b (int-signed-value b v) (int-signed-value
b v)) |

```

```

  constantAsStamp - = IllegalStamp

```

— Define when a runtime value is valid for a stamp. The stamp bounds must be valid, and val must be zero-extended.

```

fun valid-value :: Value ⇒ Stamp ⇒ bool where
  valid-value (IntVal b1 val) (IntegerStamp b l h) =
    (if b1 = b then
      valid-stamp (IntegerStamp b l h) ∧
      take-bit b val = val ∧
      l ≤ int-signed-value b val ∧ int-signed-value b val ≤ h

```

```

    else False) |

valid-value (ObjRef ref) (ObjectStamp klass exact nonNull alwaysNull) =
  ((alwaysNull  $\longrightarrow$  ref = None)  $\wedge$  (ref=None  $\longrightarrow$   $\neg$  nonNull)) |
valid-value stamp val = False

definition wf-value :: Value  $\Rightarrow$  bool where
  wf-value v = valid-value v (constantAsStamp v)

lemma unfold-wf-value[simp]:
  wf-value v  $\implies$  valid-value v (constantAsStamp v)
using wf-value-def by auto

fun compatible :: Stamp  $\Rightarrow$  Stamp  $\Rightarrow$  bool where
  compatible (IntegerStamp b1 lo1 hi1) (IntegerStamp b2 lo2 hi2) =
    (b1 = b2  $\wedge$  valid-stamp (IntegerStamp b1 lo1 hi1)  $\wedge$  valid-stamp (IntegerStamp
b2 lo2 hi2)) |
  compatible (VoidStamp) (VoidStamp) = True |
  compatible - - = False

fun stamp-under :: Stamp  $\Rightarrow$  Stamp  $\Rightarrow$  bool where
  stamp-under (IntegerStamp b1 lo1 hi1) (IntegerStamp b2 lo2 hi2) = (hi1 < lo2) |
  stamp-under - - = False

— The most common type of stamp within the compiler (apart from the Void-
Stamp) is a 32 bit integer stamp with an unrestricted range. We use default-stamp
as it is a frequently used stamp.
definition default-stamp :: Stamp where
  default-stamp = (unrestricted-stamp (IntegerStamp 32 0 0))

value valid-value (IntVal 8 (255)) (IntegerStamp 8 (-128) 127)
end

```

5 Graph Representation

5.1 IR Graph Nodes

```

theory IRNodes
  imports
    Values
begin

```

The GraalVM IR is represented using a graph data structure. Here we define the nodes that are contained within the graph. Each node represents a Node subclass in the GraalVM compiler, the node classes have annotated fields to

indicate input and successor edges.

We represent these classes with each IRNode constructor explicitly labelling a reference to the node IDs that it stores as inputs and successors.

The `inputs_of` and `successors_of` functions partition those labelled references into input edges and successor edges of a node.

To identify each Node, we use a simple natural number index. Zero is always the start node in a graph. For human readability, within nodes we write `INPUT` (or special case thereof) instead of `ID` for input edges, and `SUCC` instead of `ID` for control-flow successor edges. Optional edges are handled as "INPUT option" etc.

```
type-synonym ID = nat
type-synonym INPUT = ID
type-synonym INPUT-ASSOC = ID
type-synonym INPUT-STATE = ID
type-synonym INPUT-GUARD = ID
type-synonym INPUT-COND = ID
type-synonym INPUT-EXT = ID
type-synonym SUCC = ID
```

```
datatype (discs-sels) IRNode =
  AbsNode (ir-value: INPUT)
  | AddNode (ir-x: INPUT) (ir-y: INPUT)
  | AndNode (ir-x: INPUT) (ir-y: INPUT)
  | BeginNode (ir-next: SUCC)
  | BytecodeExceptionNode (ir-arguments: INPUT list) (ir-stateAfter-opt: INPUT-STATE
option) (ir-next: SUCC)
  | ConditionalNode (ir-condition: INPUT-COND) (ir-trueValue: INPUT) (ir-falseValue:
INPUT)
  | ConstantNode (ir-const: Value)
  | DynamicNewArrayNode (ir-elementType: INPUT) (ir-length: INPUT) (ir-voidClass-opt:
INPUT option) (ir-stateBefore-opt: INPUT-STATE option) (ir-next: SUCC)
  | EndNode
  | ExceptionObjectNode (ir-stateAfter-opt: INPUT-STATE option) (ir-next: SUCC)

  | FrameState (ir-monitorIds: INPUT-ASSOC list) (ir-outerFrameState-opt: IN-
PUT-STATE option) (ir-values-opt: INPUT list option) (ir-virtualObjectMappings-opt:
INPUT-STATE list option)
  | IfNode (ir-condition: INPUT-COND) (ir-trueSuccessor: SUCC) (ir-falseSuccessor:
SUCC)
  | IntegerBelowNode (ir-x: INPUT) (ir-y: INPUT)
  | IntegerEqualsNode (ir-x: INPUT) (ir-y: INPUT)
  | IntegerLessThanNode (ir-x: INPUT) (ir-y: INPUT)
  | InvokeNode (ir-nid: ID) (ir-callTarget: INPUT-EXT) (ir-classInit-opt: IN-
PUT option) (ir-stateDuring-opt: INPUT-STATE option) (ir-stateAfter-opt: IN-
PUT-STATE option) (ir-next: SUCC)
  | InvokeWithExceptionNode (ir-nid: ID) (ir-callTarget: INPUT-EXT) (ir-classInit-opt:
```

INPUT option) (*ir-stateDuring-opt: INPUT-STATE option*) (*ir-stateAfter-opt: INPUT-STATE option*) (*ir-next: SUCC*) (*ir-exceptionEdge: SUCC*)
 | *IsNullNode* (*ir-value: INPUT*)
 | *KillingBeginNode* (*ir-next: SUCC*)
 | *LeftShiftNode* (*ir-x: INPUT*) (*ir-y: INPUT*)
 | *LoadFieldNode* (*ir-nid: ID*) (*ir-field: string*) (*ir-object-opt: INPUT option*) (*ir-next: SUCC*)
 | *LogicNegationNode* (*ir-value: INPUT-COND*)
 | *LoopBeginNode* (*ir-ends: INPUT-ASSOC list*) (*ir-overflowGuard-opt: INPUT-GUARD option*) (*ir-stateAfter-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *LoopEndNode* (*ir-loopBegin: INPUT-ASSOC*)
 | *LoopExitNode* (*ir-loopBegin: INPUT-ASSOC*) (*ir-stateAfter-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *MergeNode* (*ir-ends: INPUT-ASSOC list*) (*ir-stateAfter-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *MethodCallTargetNode* (*ir-targetMethod: string*) (*ir-arguments: INPUT list*)
 | *MulNode* (*ir-x: INPUT*) (*ir-y: INPUT*)
 | *NarrowNode* (*ir-inputBits: nat*) (*ir-resultBits: nat*) (*ir-value: INPUT*)
 | *NegateNode* (*ir-value: INPUT*)
 | *NewArrayNode* (*ir-length: INPUT*) (*ir-stateBefore-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *NewInstanceNode* (*ir-nid: ID*) (*ir-instanceClass: string*) (*ir-stateBefore-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *NotNode* (*ir-value: INPUT*)
 | *OrNode* (*ir-x: INPUT*) (*ir-y: INPUT*)
 | *ParameterNode* (*ir-index: nat*)
 | *PiNode* (*ir-object: INPUT*) (*ir-guard-opt: INPUT-GUARD option*)
 | *ReturnNode* (*ir-result-opt: INPUT option*) (*ir-memoryMap-opt: INPUT-EXT option*)
 | *RightShiftNode* (*ir-x: INPUT*) (*ir-y: INPUT*)
 | *ShortCircuitOrNode* (*ir-x: INPUT-COND*) (*ir-y: INPUT-COND*)
 | *SignExtendNode* (*ir-inputBits: nat*) (*ir-resultBits: nat*) (*ir-value: INPUT*)
 | *SignedDivNode* (*ir-nid: ID*) (*ir-x: INPUT*) (*ir-y: INPUT*) (*ir-zeroCheck-opt: INPUT-GUARD option*) (*ir-stateBefore-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *SignedRemNode* (*ir-nid: ID*) (*ir-x: INPUT*) (*ir-y: INPUT*) (*ir-zeroCheck-opt: INPUT-GUARD option*) (*ir-stateBefore-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *StartNode* (*ir-stateAfter-opt: INPUT-STATE option*) (*ir-next: SUCC*)
 | *StoreFieldNode* (*ir-nid: ID*) (*ir-field: string*) (*ir-value: INPUT*) (*ir-stateAfter-opt: INPUT-STATE option*) (*ir-object-opt: INPUT option*) (*ir-next: SUCC*)
 | *SubNode* (*ir-x: INPUT*) (*ir-y: INPUT*)
 | *UnsignedRightShiftNode* (*ir-x: INPUT*) (*ir-y: INPUT*)
 | *UnwindNode* (*ir-exception: INPUT*)
 | *ValuePhiNode* (*ir-nid: ID*) (*ir-values: INPUT list*) (*ir-merge: INPUT-ASSOC*)
 | *ValueProxyNode* (*ir-value: INPUT*) (*ir-loopExit: INPUT-ASSOC*)
 | *XorNode* (*ir-x: INPUT*) (*ir-y: INPUT*)
 | *ZeroExtendNode* (*ir-inputBits: nat*) (*ir-resultBits: nat*) (*ir-value: INPUT*)
 | *NoNode*

| *RefNode* (*ir-ref:ID*)

```
fun opt-to-list :: 'a option ⇒ 'a list where
  opt-to-list None = [] |
  opt-to-list (Some v) = [v]
```

```
fun opt-list-to-list :: 'a list option ⇒ 'a list where
  opt-list-to-list None = [] |
  opt-list-to-list (Some x) = x
```

The following functions, *inputs_of* and *successors_of*, are automatically generated from the GraalVM compiler. Their purpose is to partition the node edges into input or successor edges.

```
fun inputs-of :: IRNode ⇒ ID list where
  inputs-of-AbsNode:
  inputs-of (AbsNode value) = [value] |
  inputs-of-AddNode:
  inputs-of (AddNode x y) = [x, y] |
  inputs-of-AndNode:
  inputs-of (AndNode x y) = [x, y] |
  inputs-of-BeginNode:
  inputs-of (BeginNode next) = [] |
  inputs-of-BytecodeExceptionNode:
  inputs-of (BytecodeExceptionNode arguments stateAfter next) = arguments @
  (opt-to-list stateAfter) |
  inputs-of-ConditionalNode:
  inputs-of (ConditionalNode condition trueValue falseValue) = [condition, true-
  Value, falseValue] |
  inputs-of-ConstantNode:
  inputs-of (ConstantNode const) = [] |
  inputs-of-DynamicNewArrayNode:
  inputs-of (DynamicNewArrayNode elementType length0 voidClass stateBefore
  next) = [elementType, length0] @ (opt-to-list voidClass) @ (opt-to-list stateBefore)
  |
  inputs-of-EndNode:
  inputs-of (EndNode) = [] |
  inputs-of-ExceptionObjectNode:
  inputs-of (ExceptionObjectNode stateAfter next) = (opt-to-list stateAfter) |
  inputs-of-FrameState:
  inputs-of (FrameState monitorIds outerFrameState values virtualObjectMappings)
  = monitorIds @ (opt-to-list outerFrameState) @ (opt-list-to-list values) @ (opt-list-to-list
  virtualObjectMappings) |
  inputs-of-IfNode:
  inputs-of (IfNode condition trueSuccessor falseSuccessor) = [condition] |
```

inputs-of-IntegerBelowNode:
inputs-of (IntegerBelowNode x y) = [x, y] |
inputs-of-IntegerEqualsNode:
inputs-of (IntegerEqualsNode x y) = [x, y] |
inputs-of-IntegerLessThanNode:
inputs-of (IntegerLessThanNode x y) = [x, y] |
inputs-of-InvokeNode:
inputs-of (InvokeNode nid0 callTarget classInit stateDuring stateAfter next)
= callTarget # (opt-to-list classInit) @ (opt-to-list stateDuring) @ (opt-to-list
stateAfter) |
inputs-of-InvokeWithExceptionNode:
inputs-of (InvokeWithExceptionNode nid0 callTarget classInit stateDuring stateAfter
next exceptionEdge) = callTarget # (opt-to-list classInit) @ (opt-to-list stateDur-
ing) @ (opt-to-list stateAfter) |
inputs-of-IsNullNode:
inputs-of (IsNullNode value) = [value] |
inputs-of-KillingBeginNode:
inputs-of (KillingBeginNode next) = [] |
inputs-of-LeftShiftNode:
inputs-of (LeftShiftNode x y) = [x, y] |
inputs-of-LoadFieldNode:
inputs-of (LoadFieldNode nid0 field object next) = (opt-to-list object) |
inputs-of-LogicNegationNode:
inputs-of (LogicNegationNode value) = [value] |
inputs-of-LoopBeginNode:
inputs-of (LoopBeginNode ends overflowGuard stateAfter next) = ends @ (opt-to-list
overflowGuard) @ (opt-to-list stateAfter) |
inputs-of-LoopEndNode:
inputs-of (LoopEndNode loopBegin) = [loopBegin] |
inputs-of-LoopExitNode:
inputs-of (LoopExitNode loopBegin stateAfter next) = loopBegin # (opt-to-list
stateAfter) |
inputs-of-MergeNode:
inputs-of (MergeNode ends stateAfter next) = ends @ (opt-to-list stateAfter) |
inputs-of-MethodCallTargetNode:
inputs-of (MethodCallTargetNode targetMethod arguments) = arguments |
inputs-of-MulNode:
inputs-of (MulNode x y) = [x, y] |
inputs-of-NarrowNode:
inputs-of (NarrowNode inputBits resultBits value) = [value] |
inputs-of-NegateNode:
inputs-of (NegateNode value) = [value] |
inputs-of-NewArrayNode:
inputs-of (NewArrayNode length0 stateBefore next) = length0 # (opt-to-list state-
Before) |
inputs-of-NewInstanceNode:
inputs-of (NewInstanceNode nid0 instanceClass stateBefore next) = (opt-to-list
stateBefore) |
inputs-of-NotNode:

inputs-of (*NotNode value*) = [*value*] |
inputs-of-OrNode:
inputs-of (*OrNode x y*) = [*x*, *y*] |
inputs-of-ParameterNode:
inputs-of (*ParameterNode index*) = [] |
inputs-of-PiNode:
inputs-of (*PiNode object guard*) = *object* # (*opt-to-list guard*) |
inputs-of-ReturnNode:
inputs-of (*ReturnNode result memoryMap*) = (*opt-to-list result*) @ (*opt-to-list memoryMap*) |
inputs-of-RightShiftNode:
inputs-of (*RightShiftNode x y*) = [*x*, *y*] |
inputs-of-ShortCircuitOrNode:
inputs-of (*ShortCircuitOrNode x y*) = [*x*, *y*] |
inputs-of-SignExtendNode:
inputs-of (*SignExtendNode inputBits resultBits value*) = [*value*] |
inputs-of-SignedDivNode:
inputs-of (*SignedDivNode nid0 x y zeroCheck stateBefore next*) = [*x*, *y*] @ (*opt-to-list zeroCheck*) @ (*opt-to-list stateBefore*) |
inputs-of-SignedRemNode:
inputs-of (*SignedRemNode nid0 x y zeroCheck stateBefore next*) = [*x*, *y*] @ (*opt-to-list zeroCheck*) @ (*opt-to-list stateBefore*) |
inputs-of-StartNode:
inputs-of (*StartNode stateAfter next*) = (*opt-to-list stateAfter*) |
inputs-of-StoreFieldNode:
inputs-of (*StoreFieldNode nid0 field value stateAfter object next*) = *value* # (*opt-to-list stateAfter*) @ (*opt-to-list object*) |
inputs-of-SubNode:
inputs-of (*SubNode x y*) = [*x*, *y*] |
inputs-of-UnsignedRightShiftNode:
inputs-of (*UnsignedRightShiftNode x y*) = [*x*, *y*] |
inputs-of-UnwindNode:
inputs-of (*UnwindNode exception*) = [*exception*] |
inputs-of-ValuePhiNode:
inputs-of (*ValuePhiNode nid0 values merge*) = *merge* # *values* |
inputs-of-ValueProxyNode:
inputs-of (*ValueProxyNode value loopExit*) = [*value*, *loopExit*] |
inputs-of-XorNode:
inputs-of (*XorNode x y*) = [*x*, *y*] |
inputs-of-ZeroExtendNode:
inputs-of (*ZeroExtendNode inputBits resultBits value*) = [*value*] |
inputs-of-NoNode: *inputs-of* (*NoNode*) = [] |

inputs-of-RefNode: *inputs-of* (*RefNode ref*) = [*ref*]

fun *successors-of* :: *IRNode* ⇒ *ID list* **where**
successors-of-AbsNode:

successors-of (*AbsNode value*) = [] |
successors-of-AddNode:
successors-of (*AddNode x y*) = [] |
successors-of-AndNode:
successors-of (*AndNode x y*) = [] |
successors-of-BeginNode:
successors-of (*BeginNode next*) = [next] |
successors-of-BytecodeExceptionNode:
successors-of (*BytecodeExceptionNode arguments stateAfter next*) = [next] |
successors-of-ConditionalNode:
successors-of (*ConditionalNode condition trueValue falseValue*) = [] |
successors-of-ConstantNode:
successors-of (*ConstantNode const*) = [] |
successors-of-DynamicNewArrayNode:
successors-of (*DynamicNewArrayNode elementType length0 voidClass stateBefore next*) = [next] |
successors-of-EndNode:
successors-of (*EndNode*) = [] |
successors-of-ExceptionObjectNode:
successors-of (*ExceptionObjectNode stateAfter next*) = [next] |
successors-of-FrameState:
successors-of (*FrameState monitorIds outerFrameState values virtualObjectMappings*) = [] |
successors-of-IfNode:
successors-of (*IfNode condition trueSuccessor falseSuccessor*) = [trueSuccessor, falseSuccessor] |
successors-of-IntegerBelowNode:
successors-of (*IntegerBelowNode x y*) = [] |
successors-of-IntegerEqualsNode:
successors-of (*IntegerEqualsNode x y*) = [] |
successors-of-IntegerLessThanNode:
successors-of (*IntegerLessThanNode x y*) = [] |
successors-of-InvokeNode:
successors-of (*InvokeNode nid0 callTarget classInit stateDuring stateAfter next*) = [next] |
successors-of-InvokeWithExceptionNode:
successors-of (*InvokeWithExceptionNode nid0 callTarget classInit stateDuring stateAfter next exceptionEdge*) = [next, exceptionEdge] |
successors-of-IsNullNode:
successors-of (*IsNullNode value*) = [] |
successors-of-KillingBeginNode:
successors-of (*KillingBeginNode next*) = [next] |
successors-of-LeftShiftNode:
successors-of (*LeftShiftNode x y*) = [] |
successors-of-LoadFieldNode:
successors-of (*LoadFieldNode nid0 field object next*) = [next] |
successors-of-LogicNegationNode:
successors-of (*LogicNegationNode value*) = [] |
successors-of-LoopBeginNode:

successors-of (*LoopBeginNode ends overflowGuard stateAfter next*) = [*next*] |
successors-of-LoopEndNode:
successors-of (*LoopEndNode loopBegin*) = [] |
successors-of-LoopExitNode:
successors-of (*LoopExitNode loopBegin stateAfter next*) = [*next*] |
successors-of-MergeNode:
successors-of (*MergeNode ends stateAfter next*) = [*next*] |
successors-of-MethodCallTargetNode:
successors-of (*MethodCallTargetNode targetMethod arguments*) = [] |
successors-of-MulNode:
successors-of (*MulNode x y*) = [] |
successors-of-NarrowNode:
successors-of (*NarrowNode inputBits resultBits value*) = [] |
successors-of-NegateNode:
successors-of (*NegateNode value*) = [] |
successors-of-NewArrayNode:
successors-of (*NewArrayNode length0 stateBefore next*) = [*next*] |
successors-of-NewInstanceNode:
successors-of (*NewInstanceNode nid0 instanceClass stateBefore next*) = [*next*] |
successors-of-NotNode:
successors-of (*NotNode value*) = [] |
successors-of-OrNode:
successors-of (*OrNode x y*) = [] |
successors-of-ParameterNode:
successors-of (*ParameterNode index*) = [] |
successors-of-PiNode:
successors-of (*PiNode object guard*) = [] |
successors-of-ReturnNode:
successors-of (*ReturnNode result memoryMap*) = [] |
successors-of-RightShiftNode:
successors-of (*RightShiftNode x y*) = [] |
successors-of-ShortCircuitOrNode:
successors-of (*ShortCircuitOrNode x y*) = [] |
successors-of-SignExtendNode:
successors-of (*SignExtendNode inputBits resultBits value*) = [] |
successors-of-SignedDivNode:
successors-of (*SignedDivNode nid0 x y zeroCheck stateBefore next*) = [*next*] |
successors-of-SignedRemNode:
successors-of (*SignedRemNode nid0 x y zeroCheck stateBefore next*) = [*next*] |
successors-of-StartNode:
successors-of (*StartNode stateAfter next*) = [*next*] |
successors-of-StoreFieldNode:
successors-of (*StoreFieldNode nid0 field value stateAfter object next*) = [*next*] |
successors-of-SubNode:
successors-of (*SubNode x y*) = [] |
successors-of-UnsignedRightShiftNode:
successors-of (*UnsignedRightShiftNode x y*) = [] |
successors-of-UnwindNode:
successors-of (*UnwindNode exception*) = [] |

```

successors-of-ValuePhiNode:
successors-of (ValuePhiNode nid0 values merge) = [] |
successors-of-ValueProxyNode:
successors-of (ValueProxyNode value loopExit) = [] |
successors-of-XorNode:
successors-of (XorNode x y) = [] |
successors-of-ZeroExtendNode:
successors-of (ZeroExtendNode inputBits resultBits value) = [] |
successors-of-NoNode: successors-of (NoNode) = [] |

```

```

successors-of-RefNode: successors-of (RefNode ref) = [ref]

```

```

lemma inputs-of (FrameState x (Some y) (Some z) None) = x @ [y] @ z
unfolding inputs-of-FrameState by simp
lemma successors-of (FrameState x (Some y) (Some z) None) = []
unfolding inputs-of-FrameState by simp

```

```

lemma inputs-of (IfNode c t f) = [c]
unfolding inputs-of-IfNode by simp
lemma successors-of (IfNode c t f) = [t, f]
unfolding successors-of-IfNode by simp

```

```

lemma inputs-of (EndNode) = [] ∧ successors-of (EndNode) = []
unfolding inputs-of-EndNode successors-of-EndNode by simp

```

end

5.2 IR Graph Node Hierarchy

```

theory IRNodeHierarchy
imports IRNodes
begin

```

It is helpful to introduce a node hierarchy into our formalization. Often the GraalVM compiler relies on explicit type checks to determine which operations to perform on a given node, we try to mimic the same functionality by using a suite of predicate functions over the `IRNode` class to determine inheritance.

As one would expect, the function `is<ClassName>Type` will be true if the node parameter is a subclass of the `ClassName` within the GraalVM compiler.

These functions have been automatically generated from the compiler.

```

fun is-EndNode :: IRNode ⇒ bool where
  is-EndNode EndNode = True |
  is-EndNode _ = False

```

```

fun is-VirtualState :: IRNode  $\Rightarrow$  bool where
  is-VirtualState n = ((is-FrameState n))

fun is-BinaryArithmeticNode :: IRNode  $\Rightarrow$  bool where
  is-BinaryArithmeticNode n = ((is-AddNode n)  $\vee$  (is-AndNode n)  $\vee$  (is-MulNode
n)  $\vee$  (is-OrNode n)  $\vee$  (is-SubNode n)  $\vee$  (is-XorNode n))

fun is-ShiftNode :: IRNode  $\Rightarrow$  bool where
  is-ShiftNode n = ((is-LeftShiftNode n)  $\vee$  (is-RightShiftNode n)  $\vee$  (is-UnsignedRightShiftNode
n))

fun is-BinaryNode :: IRNode  $\Rightarrow$  bool where
  is-BinaryNode n = ((is-BinaryArithmeticNode n)  $\vee$  (is-ShiftNode n))

fun is-AbstractLocalNode :: IRNode  $\Rightarrow$  bool where
  is-AbstractLocalNode n = ((is-ParameterNode n))

fun is-IntegerConvertNode :: IRNode  $\Rightarrow$  bool where
  is-IntegerConvertNode n = ((is-NarrowNode n)  $\vee$  (is-SignExtendNode n)  $\vee$ 
(is-ZeroExtendNode n))

fun is-UnaryArithmeticNode :: IRNode  $\Rightarrow$  bool where
  is-UnaryArithmeticNode n = ((is-AbsNode n)  $\vee$  (is-NegateNode n)  $\vee$  (is-NotNode
n))

fun is-UnaryNode :: IRNode  $\Rightarrow$  bool where
  is-UnaryNode n = ((is-IntegerConvertNode n)  $\vee$  (is-UnaryArithmeticNode n))

fun is-PhiNode :: IRNode  $\Rightarrow$  bool where
  is-PhiNode n = ((is-ValuePhiNode n))

fun is-FloatingGuardedNode :: IRNode  $\Rightarrow$  bool where
  is-FloatingGuardedNode n = ((is-PiNode n))

fun is-UnaryOpLogicNode :: IRNode  $\Rightarrow$  bool where
  is-UnaryOpLogicNode n = ((is-IsNullNode n))

fun is-IntegerLowerThanNode :: IRNode  $\Rightarrow$  bool where
  is-IntegerLowerThanNode n = ((is-IntegerBelowNode n)  $\vee$  (is-IntegerLessThanNode
n))

fun is-CompareNode :: IRNode  $\Rightarrow$  bool where
  is-CompareNode n = ((is-IntegerEqualsNode n)  $\vee$  (is-IntegerLowerThanNode n))

fun is-BinaryOpLogicNode :: IRNode  $\Rightarrow$  bool where
  is-BinaryOpLogicNode n = ((is-CompareNode n))

```

```

fun is-LogicNode :: IRNode  $\Rightarrow$  bool where
  is-LogicNode n = ((is-BinaryOpLogicNode n)  $\vee$  (is-LogicNegationNode n)  $\vee$ 
    (is-ShortCircuitOrNode n)  $\vee$  (is-UnaryOpLogicNode n))

fun is-ProxyNode :: IRNode  $\Rightarrow$  bool where
  is-ProxyNode n = ((is-ValueProxyNode n))

fun is-FloatingNode :: IRNode  $\Rightarrow$  bool where
  is-FloatingNode n = ((is-AbstractLocalNode n)  $\vee$  (is-BinaryNode n)  $\vee$  (is-ConditionalNode
    n)  $\vee$  (is-ConstantNode n)  $\vee$  (is-FloatingGuardedNode n)  $\vee$  (is-LogicNode n)  $\vee$ 
    (is-PhiNode n)  $\vee$  (is-ProxyNode n)  $\vee$  (is-UnaryNode n))

fun is-AccessFieldNode :: IRNode  $\Rightarrow$  bool where
  is-AccessFieldNode n = ((is-LoadFieldNode n)  $\vee$  (is-StoreFieldNode n))

fun is-AbstractNewArrayNode :: IRNode  $\Rightarrow$  bool where
  is-AbstractNewArrayNode n = ((is-DynamicNewArrayNode n)  $\vee$  (is-NewArrayNode
    n))

fun is-AbstractNewObjectNode :: IRNode  $\Rightarrow$  bool where
  is-AbstractNewObjectNode n = ((is-AbstractNewArrayNode n)  $\vee$  (is-NewInstanceNode
    n))

fun is-IntegerDivRemNode :: IRNode  $\Rightarrow$  bool where
  is-IntegerDivRemNode n = ((is-SignedDivNode n)  $\vee$  (is-SignedRemNode n))

fun is-FixedBinaryNode :: IRNode  $\Rightarrow$  bool where
  is-FixedBinaryNode n = ((is-IntegerDivRemNode n))

fun is-DeoptimizingFixedWithNextNode :: IRNode  $\Rightarrow$  bool where
  is-DeoptimizingFixedWithNextNode n = ((is-AbstractNewObjectNode n)  $\vee$  (is-FixedBinaryNode
    n))

fun is-AbstractMemoryCheckpoint :: IRNode  $\Rightarrow$  bool where
  is-AbstractMemoryCheckpoint n = ((is-BytecodeExceptionNode n)  $\vee$  (is-InvokeNode
    n))

fun is-AbstractStateSplit :: IRNode  $\Rightarrow$  bool where
  is-AbstractStateSplit n = ((is-AbstractMemoryCheckpoint n))

fun is-AbstractMergeNode :: IRNode  $\Rightarrow$  bool where
  is-AbstractMergeNode n = ((is-LoopBeginNode n)  $\vee$  (is-MergeNode n))

fun is-BeginStateSplitNode :: IRNode  $\Rightarrow$  bool where
  is-BeginStateSplitNode n = ((is-AbstractMergeNode n)  $\vee$  (is-ExceptionObjectNode
    n)  $\vee$  (is-LoopExitNode n)  $\vee$  (is-StartNode n))

fun is-AbstractBeginNode :: IRNode  $\Rightarrow$  bool where
  is-AbstractBeginNode n = ((is-BeginNode n)  $\vee$  (is-BeginStateSplitNode n)  $\vee$ 

```


(is-KillingBeginNode n))

fun *is-FixedWithNextNode* :: *IRNode* \Rightarrow *bool* **where**
 is-FixedWithNextNode *n* = ((*is-AbstractBeginNode* *n*) \vee (*is-AbstractStateSplit* *n*)
 \vee (*is-AccessFieldNode* *n*) \vee (*is-DeoptimizingFixedWithNextNode* *n*))

fun *is-WithExceptionNode* :: *IRNode* \Rightarrow *bool* **where**
 is-WithExceptionNode *n* = ((*is-InvokeWithExceptionNode* *n*))

fun *is-ControlSplitNode* :: *IRNode* \Rightarrow *bool* **where**
 is-ControlSplitNode *n* = ((*is-IfNode* *n*) \vee (*is-WithExceptionNode* *n*))

fun *is-ControlSinkNode* :: *IRNode* \Rightarrow *bool* **where**
 is-ControlSinkNode *n* = ((*is-ReturnNode* *n*) \vee (*is-UnwindNode* *n*))

fun *is-AbstractEndNode* :: *IRNode* \Rightarrow *bool* **where**
 is-AbstractEndNode *n* = ((*is-EndNode* *n*) \vee (*is-LoopEndNode* *n*))

fun *is-FixedNode* :: *IRNode* \Rightarrow *bool* **where**
 is-FixedNode *n* = ((*is-AbstractEndNode* *n*) \vee (*is-ControlSinkNode* *n*) \vee (*is-ControlSplitNode*
 n) \vee (*is-FixedWithNextNode* *n*))

fun *is-CallTargetNode* :: *IRNode* \Rightarrow *bool* **where**
 is-CallTargetNode *n* = ((*is-MethodCallTargetNode* *n*))

fun *is-ValueNode* :: *IRNode* \Rightarrow *bool* **where**
 is-ValueNode *n* = ((*is-CallTargetNode* *n*) \vee (*is-FixedNode* *n*) \vee (*is-FloatingNode*
 n))

fun *is-Node* :: *IRNode* \Rightarrow *bool* **where**
 is-Node *n* = ((*is-ValueNode* *n*) \vee (*is-VirtualState* *n*))

fun *is-MemoryKill* :: *IRNode* \Rightarrow *bool* **where**
 is-MemoryKill *n* = ((*is-AbstractMemoryCheckpoint* *n*))

fun *is-NarrowableArithmeticNode* :: *IRNode* \Rightarrow *bool* **where**
 is-NarrowableArithmeticNode *n* = ((*is-AbsNode* *n*) \vee (*is-AddNode* *n*) \vee (*is-AndNode*
 n) \vee (*is-MulNode* *n*) \vee (*is-NegateNode* *n*) \vee (*is-NotNode* *n*) \vee (*is-OrNode* *n*) \vee
 (*is-ShiftNode* *n*) \vee (*is-SubNode* *n*) \vee (*is-XorNode* *n*))

fun *is-AnchoringNode* :: *IRNode* \Rightarrow *bool* **where**
 is-AnchoringNode *n* = ((*is-AbstractBeginNode* *n*))

fun *is-DeoptBefore* :: *IRNode* \Rightarrow *bool* **where**
 is-DeoptBefore *n* = ((*is-DeoptimizingFixedWithNextNode* *n*))

fun *is-IndirectCanonicalization* :: *IRNode* \Rightarrow *bool* **where**
 is-IndirectCanonicalization *n* = ((*is-LogicNode* *n*))

```

fun is-IterableNodeType :: IRNode  $\Rightarrow$  bool where
  is-IterableNodeType n = ((is-AbstractBeginNode n)  $\vee$  (is-AbstractMergeNode n)  $\vee$ 
    (is-FrameState n)  $\vee$  (is-IfNode n)  $\vee$  (is-IntegerDivRemNode n)  $\vee$  (is-InvokeWithExceptionNode
      n)  $\vee$  (is-LoopBeginNode n)  $\vee$  (is-LoopExitNode n)  $\vee$  (is-MethodCallTargetNode n)
       $\vee$  (is-ParameterNode n)  $\vee$  (is-ReturnNode n)  $\vee$  (is-ShortCircuitOrNode n))

fun is-Invoke :: IRNode  $\Rightarrow$  bool where
  is-Invoke n = ((is-InvokeNode n)  $\vee$  (is-InvokeWithExceptionNode n))

fun is-Proxy :: IRNode  $\Rightarrow$  bool where
  is-Proxy n = ((is-ProxyNode n))

fun is-ValueProxy :: IRNode  $\Rightarrow$  bool where
  is-ValueProxy n = ((is-PiNode n)  $\vee$  (is-ValueProxyNode n))

fun is-ValueNodeInterface :: IRNode  $\Rightarrow$  bool where
  is-ValueNodeInterface n = ((is-ValueNode n))

fun is-ArrayLengthProvider :: IRNode  $\Rightarrow$  bool where
  is-ArrayLengthProvider n = ((is-AbstractNewArrayNode n)  $\vee$  (is-ConstantNode
    n))

fun is-StampInverter :: IRNode  $\Rightarrow$  bool where
  is-StampInverter n = ((is-IntegerConvertNode n)  $\vee$  (is-NegateNode n)  $\vee$  (is-NotNode
    n))

fun is-GuardingNode :: IRNode  $\Rightarrow$  bool where
  is-GuardingNode n = ((is-AbstractBeginNode n))

fun is-SingleMemoryKill :: IRNode  $\Rightarrow$  bool where
  is-SingleMemoryKill n = ((is-BytecodeExceptionNode n)  $\vee$  (is-ExceptionObjectNode
    n)  $\vee$  (is-InvokeNode n)  $\vee$  (is-InvokeWithExceptionNode n)  $\vee$  (is-KillingBeginNode
    n)  $\vee$  (is-StartNode n))

fun is-LIRLowerable :: IRNode  $\Rightarrow$  bool where
  is-LIRLowerable n = ((is-AbstractBeginNode n)  $\vee$  (is-AbstractEndNode n)  $\vee$ 
    (is-AbstractMergeNode n)  $\vee$  (is-BinaryOpLogicNode n)  $\vee$  (is-CallTargetNode n)  $\vee$ 
    (is-ConditionalNode n)  $\vee$  (is-ConstantNode n)  $\vee$  (is-IfNode n)  $\vee$  (is-InvokeNode n)
     $\vee$  (is-InvokeWithExceptionNode n)  $\vee$  (is-IsNullNode n)  $\vee$  (is-LoopBeginNode n)  $\vee$ 
    (is-PiNode n)  $\vee$  (is-ReturnNode n)  $\vee$  (is-SignedDivNode n)  $\vee$  (is-SignedRemNode
    n)  $\vee$  (is-UnaryOpLogicNode n)  $\vee$  (is-UnwindNode n))

fun is-GuardedNode :: IRNode  $\Rightarrow$  bool where
  is-GuardedNode n = ((is-FloatingGuardedNode n))

fun is-ArithmeticLIRLowerable :: IRNode  $\Rightarrow$  bool where
  is-ArithmeticLIRLowerable n = ((is-AbsNode n)  $\vee$  (is-BinaryArithmeticNode n)  $\vee$ 
    (is-IntegerConvertNode n)  $\vee$  (is-NotNode n)  $\vee$  (is-ShiftNode n)  $\vee$  (is-UnaryArithmeticNode
    n))

```

```

fun is-SwitchFoldable :: IRNode  $\Rightarrow$  bool where
  is-SwitchFoldable n = ((is-IfNode n))

fun is-VirtualizableAllocation :: IRNode  $\Rightarrow$  bool where
  is-VirtualizableAllocation n = ((is-NewArrayNode n)  $\vee$  (is-NewInstanceNode n))

fun is-Unary :: IRNode  $\Rightarrow$  bool where
  is-Unary n = ((is-LoadFieldNode n)  $\vee$  (is-LogicNegationNode n)  $\vee$  (is-UnaryNode
n)  $\vee$  (is-UnaryOpLogicNode n))

fun is-FixedNodeInterface :: IRNode  $\Rightarrow$  bool where
  is-FixedNodeInterface n = ((is-FixedNode n))

fun is-BinaryCommutative :: IRNode  $\Rightarrow$  bool where
  is-BinaryCommutative n = ((is-AddNode n)  $\vee$  (is-AndNode n)  $\vee$  (is-IntegerEqualsNode
n)  $\vee$  (is-MulNode n)  $\vee$  (is-OrNode n)  $\vee$  (is-XorNode n))

fun is-Canonicalizable :: IRNode  $\Rightarrow$  bool where
  is-Canonicalizable n = ((is-BytecodeExceptionNode n)  $\vee$  (is-ConditionalNode n)  $\vee$ 
(is-DynamicNewArrayNode n)  $\vee$  (is-PhiNode n)  $\vee$  (is-PiNode n)  $\vee$  (is-ProxyNode
n)  $\vee$  (is-StoreFieldNode n)  $\vee$  (is-ValueProxyNode n))

fun is-UncheckedInterfaceProvider :: IRNode  $\Rightarrow$  bool where
  is-UncheckedInterfaceProvider n = ((is-InvokeNode n)  $\vee$  (is-InvokeWithExceptionNode
n)  $\vee$  (is-LoadFieldNode n)  $\vee$  (is-ParameterNode n))

fun is-Binary :: IRNode  $\Rightarrow$  bool where
  is-Binary n = ((is-BinaryArithmeticNode n)  $\vee$  (is-BinaryNode n)  $\vee$  (is-BinaryOpLogicNode
n)  $\vee$  (is-CompareNode n)  $\vee$  (is-FixedBinaryNode n)  $\vee$  (is-ShortCircuitOrNode n))

fun is-ArithmeticOperation :: IRNode  $\Rightarrow$  bool where
  is-ArithmeticOperation n = ((is-BinaryArithmeticNode n)  $\vee$  (is-IntegerConvertNode
n)  $\vee$  (is-ShiftNode n)  $\vee$  (is-UnaryArithmeticNode n))

fun is-ValueNumberable :: IRNode  $\Rightarrow$  bool where
  is-ValueNumberable n = ((is-FloatingNode n)  $\vee$  (is-ProxyNode n))

fun is-Lowerable :: IRNode  $\Rightarrow$  bool where
  is-Lowerable n = ((is-AbstractNewObjectNode n)  $\vee$  (is-AccessFieldNode n)  $\vee$ 
(is-BytecodeExceptionNode n)  $\vee$  (is-ExceptionObjectNode n)  $\vee$  (is-IntegerDivRemNode
n)  $\vee$  (is-UnwindNode n))

fun is-Virtualizable :: IRNode  $\Rightarrow$  bool where
  is-Virtualizable n = ((is-IsNullNode n)  $\vee$  (is-LoadFieldNode n)  $\vee$  (is-PiNode n)
 $\vee$  (is-StoreFieldNode n)  $\vee$  (is-ValueProxyNode n))

fun is-Simplifiable :: IRNode  $\Rightarrow$  bool where
  is-Simplifiable n = ((is-AbstractMergeNode n)  $\vee$  (is-BeginNode n)  $\vee$  (is-IfNode

```

$n) \vee (is-LoopExitNode\ n) \vee (is-MethodCallTargetNode\ n) \vee (is-NewArrayNode\ n))$

fun *is-StateSplit* :: *IRNode* \Rightarrow *bool* **where**
is-StateSplit *n* = ((*is-AbstractStateSplit* *n*) \vee (*is-BeginStateSplitNode* *n*) \vee (*is-StoreFieldNode* *n*))

fun *is-ConvertNode* :: *IRNode* \Rightarrow *bool* **where**
is-ConvertNode *n* = ((*is-IntegerConvertNode* *n*))

fun *is-sequential-node* :: *IRNode* \Rightarrow *bool* **where**
is-sequential-node (*StartNode* -) = *True* |
is-sequential-node (*BeginNode* -) = *True* |
is-sequential-node (*KillingBeginNode* -) = *True* |
is-sequential-node (*LoopBeginNode* - - -) = *True* |
is-sequential-node (*LoopExitNode* - - -) = *True* |
is-sequential-node (*MergeNode* - - -) = *True* |
is-sequential-node (*RefNode* -) = *True* |
is-sequential-node - = *False*

The following convenience function is useful in determining if two *IRNodes* are of the same type irregardless of their edges. It will return true if both the node parameters are the same node class.

fun *is-same-ir-node-type* :: *IRNode* \Rightarrow *IRNode* \Rightarrow *bool* **where**
is-same-ir-node-type *n1* *n2* = (
 ((*is-AbsNode* *n1*) \wedge (*is-AbsNode* *n2*)) \vee
 ((*is-AddNode* *n1*) \wedge (*is-AddNode* *n2*)) \vee
 ((*is-AndNode* *n1*) \wedge (*is-AndNode* *n2*)) \vee
 ((*is-BeginNode* *n1*) \wedge (*is-BeginNode* *n2*)) \vee
 ((*is-BytecodeExceptionNode* *n1*) \wedge (*is-BytecodeExceptionNode* *n2*)) \vee
 ((*is-ConditionalNode* *n1*) \wedge (*is-ConditionalNode* *n2*)) \vee
 ((*is-ConstantNode* *n1*) \wedge (*is-ConstantNode* *n2*)) \vee
 ((*is-DynamicNewArrayNode* *n1*) \wedge (*is-DynamicNewArrayNode* *n2*)) \vee
 ((*is-EndNode* *n1*) \wedge (*is-EndNode* *n2*)) \vee
 ((*is-ExceptionObjectNode* *n1*) \wedge (*is-ExceptionObjectNode* *n2*)) \vee
 ((*is-FrameState* *n1*) \wedge (*is-FrameState* *n2*)) \vee
 ((*is-IfNode* *n1*) \wedge (*is-IfNode* *n2*)) \vee
 ((*is-IntegerBelowNode* *n1*) \wedge (*is-IntegerBelowNode* *n2*)) \vee
 ((*is-IntegerEqualsNode* *n1*) \wedge (*is-IntegerEqualsNode* *n2*)) \vee
 ((*is-IntegerLessThanNode* *n1*) \wedge (*is-IntegerLessThanNode* *n2*)) \vee
 ((*is-InvokeNode* *n1*) \wedge (*is-InvokeNode* *n2*)) \vee
 ((*is-InvokeWithExceptionNode* *n1*) \wedge (*is-InvokeWithExceptionNode* *n2*)) \vee
 ((*is-IsNullNode* *n1*) \wedge (*is-IsNullNode* *n2*)) \vee
 ((*is-KillingBeginNode* *n1*) \wedge (*is-KillingBeginNode* *n2*)) \vee
 ((*is-LeftShiftNode* *n1*) \wedge (*is-LeftShiftNode* *n2*)) \vee
 ((*is-LoadFieldNode* *n1*) \wedge (*is-LoadFieldNode* *n2*)) \vee
 ((*is-LogicNegationNode* *n1*) \wedge (*is-LogicNegationNode* *n2*)) \vee
 ((*is-LoopBeginNode* *n1*) \wedge (*is-LoopBeginNode* *n2*)) \vee
 ((*is-LoopEndNode* *n1*) \wedge (*is-LoopEndNode* *n2*)) \vee

```

((is-LoopExitNode n1) ∧ (is-LoopExitNode n2)) ∨
((is-MergeNode n1) ∧ (is-MergeNode n2)) ∨
((is-MethodCallTargetNode n1) ∧ (is-MethodCallTargetNode n2)) ∨
((is-MulNode n1) ∧ (is-MulNode n2)) ∨
((is-NarrowNode n1) ∧ (is-NarrowNode n2)) ∨
((is-NegateNode n1) ∧ (is-NegateNode n2)) ∨
((is-NewArrayNode n1) ∧ (is-NewArrayNode n2)) ∨
((is-NewInstanceNode n1) ∧ (is-NewInstanceNode n2)) ∨
((is-NotNode n1) ∧ (is-NotNode n2)) ∨
((is-OrNode n1) ∧ (is-OrNode n2)) ∨
((is-ParameterNode n1) ∧ (is-ParameterNode n2)) ∨
((is-PiNode n1) ∧ (is-PiNode n2)) ∨
((is-ReturnNode n1) ∧ (is-ReturnNode n2)) ∨
((is-RightShiftNode n1) ∧ (is-RightShiftNode n2)) ∨
((is-ShortCircuitOrNode n1) ∧ (is-ShortCircuitOrNode n2)) ∨
((is-SignedDivNode n1) ∧ (is-SignedDivNode n2)) ∨
((is-SignedRemNode n1) ∧ (is-SignedRemNode n2)) ∨
((is-SignExtendNode n1) ∧ (is-SignExtendNode n2)) ∨
((is-StartNode n1) ∧ (is-StartNode n2)) ∨
((is-StoreFieldNode n1) ∧ (is-StoreFieldNode n2)) ∨
((is-SubNode n1) ∧ (is-SubNode n2)) ∨
((is-UnsignedRightShiftNode n1) ∧ (is-UnsignedRightShiftNode n2)) ∨
((is-UnwindNode n1) ∧ (is-UnwindNode n2)) ∨
((is-ValuePhiNode n1) ∧ (is-ValuePhiNode n2)) ∨
((is-ValueProxyNode n1) ∧ (is-ValueProxyNode n2)) ∨
((is-XorNode n1) ∧ (is-XorNode n2)) ∨
((is-ZeroExtendNode n1) ∧ (is-ZeroExtendNode n2)))

```

end

5.3 IR Graph Type

```

theory IRGraph
  imports
    IRNodeHierarchy
    Stamp
    HOL-Library.FSet
    HOL.Relation
begin

```

This theory defines the main Graal data structure - an entire IR Graph.

IRGraph is defined as a partial map with a finite domain. The finite domain is required to be able to generate code and produce an interpreter.

```

typedef IRGraph = {g :: ID → (IRNode × Stamp) . finite (dom g)}
proof -
  have finite(dom(Map.empty)) ∧ ran Map.empty = {} by auto
  then show ?thesis
    by fastforce

```

qed

setup-lifting *type-definition-IRGraph*

lift-definition *ids* :: *IRGraph* \Rightarrow *ID set*
is $\lambda g. \{nid \in dom\ g \cdot \nexists s. g\ nid = (Some\ (NoNode,\ s))\}$.

fun *with-default* :: '*c* \Rightarrow ('*b* \Rightarrow '*c*) \Rightarrow (('a \rightarrow '*b*) \Rightarrow '*a* \Rightarrow '*c*) **where**
with-default *def conv* = ($\lambda m\ k.$
(*case* *m* *k* *of* *None* \Rightarrow *def* | *Some* *v* \Rightarrow *conv* *v*))

lift-definition *kind* :: *IRGraph* \Rightarrow (*ID* \Rightarrow *IRNode*)
is *with-default* *NoNode* *fst* .

lift-definition *stamp* :: *IRGraph* \Rightarrow *ID* \Rightarrow *Stamp*
is *with-default* *IllegalStamp* *snd* .

lift-definition *add-node* :: *ID* \Rightarrow (*IRNode* \times *Stamp*) \Rightarrow *IRGraph* \Rightarrow *IRGraph*
is $\lambda nid\ k\ g.$ *if* *fst* *k* = *NoNode* *then* *g* *else* *g*(*nid* \mapsto *k*) **by** *simp*

lift-definition *remove-node* :: *ID* \Rightarrow *IRGraph* \Rightarrow *IRGraph*
is $\lambda nid\ g.$ *g*(*nid* := *None*) **by** *simp*

lift-definition *replace-node* :: *ID* \Rightarrow (*IRNode* \times *Stamp*) \Rightarrow *IRGraph* \Rightarrow *IRGraph*
is $\lambda nid\ k\ g.$ *if* *fst* *k* = *NoNode* *then* *g* *else* *g*(*nid* \mapsto *k*) **by** *simp*

lift-definition *as-list* :: *IRGraph* \Rightarrow (*ID* \times *IRNode* \times *Stamp*) *list*
is $\lambda g.$ *map* ($\lambda k. (k, the\ (g\ k))$) (*sorted-list-of-set* (*dom* *g*)) .

fun *no-node* :: (*ID* \times (*IRNode* \times *Stamp*)) *list* \Rightarrow (*ID* \times (*IRNode* \times *Stamp*)) *list*
where
no-node *g* = *filter* ($\lambda n. fst\ (snd\ n) \neq NoNode$) *g*

lift-definition *irgraph* :: (*ID* \times (*IRNode* \times *Stamp*)) *list* \Rightarrow *IRGraph*
is *map-of* \circ *no-node*
by (*simp* *add: finite-dom-map-of*)

definition *as-set* :: *IRGraph* \Rightarrow (*ID* \times (*IRNode* \times *Stamp*)) *set* **where**
as-set *g* = $\{(n, kind\ g\ n, stamp\ g\ n) \mid n \cdot n \in ids\ g\}$

definition *true-ids* :: *IRGraph* \Rightarrow *ID set* **where**
true-ids *g* = *ids* *g* - $\{n \in ids\ g. \exists n'. kind\ g\ n = RefNode\ n'\}$

definition *domain-subtraction* :: '*a set* \Rightarrow ('*a* \times '*b*) *set* \Rightarrow ('*a* \times '*b*) *set*
(**infix** ≤ 30) **where**
domain-subtraction *s* *r* = $\{(x, y) \cdot (x, y) \in r \wedge x \notin s\}$

notation (*latex*)
domain-subtraction ($\cdot \Leftarrow \cdot$)

code-datatype *irgraph*

fun *filter-none* **where**

filter-none *g* = {*nid* ∈ *dom g* . $\nexists s. g \text{ nid} = (\text{Some } (\text{NoNode}, s))$ }

lemma *no-node-clears*:

res = *no-node xs* $\longrightarrow (\forall x \in \text{set } res. \text{fst } (\text{snd } x) \neq \text{NoNode})$

by *simp*

lemma *dom-eq*:

assumes $\forall x \in \text{set } xs. \text{fst } (\text{snd } x) \neq \text{NoNode}$

shows *filter-none* (*map-of xs*) = *dom* (*map-of xs*)

unfolding *filter-none.simps* **using** *assms map-of-SomeD*

by *fastforce*

lemma *fil-eq*:

filter-none (*map-of* (*no-node xs*)) = *set* (*map fst* (*no-node xs*))

using *no-node-clears*

by (*metis dom-eq dom-map-of-conv-image-fst list.set-map*)

lemma *irgraph[code]*: *ids* (*irgraph m*) = *set* (*map fst* (*no-node m*))

unfolding *irgraph-def ids-def* **using** *fil-eq*

by (*smt Rep-IRGraph comp-apply eq-onp-same-args filter-none.simps ids.abs-eq ids-def irgraph.abs-eq irgraph.rep-eq irgraph-def mem-Collect-eq*)

lemma [*code*]: *Rep-IRGraph* (*irgraph m*) = *map-of* (*no-node m*)

using *Abs-IRGraph-inverse*

by (*simp add: irgraph.rep-eq*)

— Get the inputs set of a given node ID

fun *inputs* :: *IRGraph* \Rightarrow *ID* \Rightarrow *ID set* **where**

inputs *g nid* = *set* (*inputs-of* (*kind g nid*))

— Get the successor set of a given node ID

fun *succ* :: *IRGraph* \Rightarrow *ID* \Rightarrow *ID set* **where**

succ *g nid* = *set* (*successors-of* (*kind g nid*))

— Gives a relation between node IDs - between a node and its input nodes

fun *input-edges* :: *IRGraph* \Rightarrow *ID rel* **where**

input-edges *g* = ($\bigcup i \in \text{ids } g. \{(i,j) | j. j \in (\text{inputs } g i)\}$)

— Find all the nodes in the graph that have *nid* as an input - the usages of *nid*

fun *usages* :: *IRGraph* \Rightarrow *ID* \Rightarrow *ID set* **where**

usages *g nid* = {*i*. *i* ∈ *ids g* ∧ *nid* ∈ *inputs g i*}

fun *successor-edges* :: *IRGraph* \Rightarrow *ID rel* **where**

successor-edges *g* = ($\bigcup i \in \text{ids } g. \{(i,j) | j. j \in (\text{succ } g i)\}$)

fun *predecessors* :: *IRGraph* \Rightarrow *ID* \Rightarrow *ID set* **where**

predecessors *g nid* = {*i*. *i* ∈ *ids g* ∧ *nid* ∈ *succ g i*}

fun *nodes-of* :: *IRGraph* \Rightarrow (*IRNode* \Rightarrow *bool*) \Rightarrow *ID set* **where**

```

  nodes-of g sel = {nid ∈ ids g . sel (kind g nid)}
fun edge :: (IRNode ⇒ 'a) ⇒ ID ⇒ IRGraph ⇒ 'a where
  edge sel nid g = sel (kind g nid)

fun filtered-inputs :: IRGraph ⇒ ID ⇒ (IRNode ⇒ bool) ⇒ ID list where
  filtered-inputs g nid f = filter (f ∘ (kind g)) (inputs-of (kind g nid))
fun filtered-successors :: IRGraph ⇒ ID ⇒ (IRNode ⇒ bool) ⇒ ID list where
  filtered-successors g nid f = filter (f ∘ (kind g)) (successors-of (kind g nid))
fun filtered-usages :: IRGraph ⇒ ID ⇒ (IRNode ⇒ bool) ⇒ ID set where
  filtered-usages g nid f = {n ∈ (usages g nid). f (kind g n)}

fun is-empty :: IRGraph ⇒ bool where
  is-empty g = (ids g = {})

fun any-usage :: IRGraph ⇒ ID ⇒ ID where
  any-usage g nid = hd (sorted-list-of-set (usages g nid))

lemma ids-some[simp]: x ∈ ids g ⟷ kind g x ≠ NoNode
proof –
  have that: x ∈ ids g ⟶ kind g x ≠ NoNode
    using ids.rep-eq kind.rep-eq by force
  have kind g x ≠ NoNode ⟶ x ∈ ids g
    unfolding with-default.simps kind-def ids-def
    by (cases Rep-IRGraph g x = None; auto)
  from this that show ?thesis by auto
qed

lemma not-in-g:
  assumes nid ∉ ids g
  shows kind g nid = NoNode
  using assms ids-some by blast

lemma valid-creation[simp]:
  finite (dom g) ⟷ Rep-IRGraph (Abs-IRGraph g) = g
  using Abs-IRGraph-inverse by (metis Rep-IRGraph mem-Collect-eq)

lemma [simp]: finite (ids g)
  using Rep-IRGraph ids.rep-eq by simp

lemma [simp]: finite (ids (irgraph g))
  by (simp add: finite-dom-map-of)

lemma [simp]: finite (dom g) ⟶ ids (Abs-IRGraph g) = {nid ∈ dom g . ∄ s. g
  nid = Some (NoNode, s)}
  using ids.rep-eq by simp

lemma [simp]: finite (dom g) ⟶ kind (Abs-IRGraph g) = (λx . (case g x of None
  ⇒ NoNode | Some n ⇒ fst n))
  by (simp add: kind.rep-eq)

```



```

lemma [simp]: finite (dom g)  $\longrightarrow$  stamp (Abs-IRGraph g) = ( $\lambda x$  . (case g x of
```

None \Rightarrow *IllegalStamp* | *Some* *n* \Rightarrow *snd* *n*))

```

  using stamp.abs-eq stamp.rep-eq by auto
```

lemma [simp]: *ids* (*irgraph* *g*) = *set* (*map fst* (*no-node* *g*))

```

  using irgraph by auto
```

lemma [simp]: *kind* (*irgraph* *g*) = ($\lambda \textit{nid}$. (*case* (*map-of* (*no-node* *g*)) *nid* of

None \Rightarrow *NoNode* | *Some* *n* \Rightarrow *fst* *n*))

```

  using irgraph.rep-eq kind.transfer kind.rep-eq by auto
```

lemma [simp]: *stamp* (*irgraph* *g*) = ($\lambda \textit{nid}$. (*case* (*map-of* (*no-node* *g*)) *nid* of

None \Rightarrow *IllegalStamp* | *Some* *n* \Rightarrow *snd* *n*))

```

  using irgraph.rep-eq stamp.transfer stamp.rep-eq by auto
```

lemma *map-of-upd*: (*map-of* *g*)(*k* \mapsto *v*) = (*map-of* ((*k*, *v*) $\#$ *g*))

```

  by simp
```

lemma [code]: *replace-node* *nid* *k* (*irgraph* *g*) = (*irgraph* (((*nid*, *k*) $\#$ *g*)))

proof (*cases* *fst* *k* = *NoNode*)

```

  case True
    then show ?thesis
      by (metis (mono-tags, lifting) Rep-IRGraph-inject filter.simps(2) irgraph.abs-eq
        no-node.simps replace-node.rep-eq snd-conv)
  next
    case False
      then show ?thesis unfolding irgraph-def replace-node-def no-node.simps
        by (smt (verit, best) Rep-IRGraph comp-apply eq-onp-same-args filter.simps(2)
          id-def irgraph.rep-eq map-fun-apply map-of-upd mem-Collect-eq no-node.elims re-
          place-node.abs-eq replace-node-def snd-eqD)
  qed
```

lemma [code]: *add-node* *nid* *k* (*irgraph* *g*) = (*irgraph* (((*nid*, *k*) $\#$ *g*)))

```

  by (smt (z3) Rep-IRGraph-inject add-node.rep-eq filter.simps(2) irgraph.rep-eq
    map-of-upd no-node.simps snd-conv)
```

lemma *add-node-lookup*:

```

  gup = add-node nid (k, s) g  $\longrightarrow$ 
    (if k  $\neq$  NoNode then kind gup nid = k  $\wedge$  stamp gup nid = s else kind gup nid
    = kind g nid)
proof (cases k = NoNode)
  case True
    then show ?thesis
      by (simp add: add-node.rep-eq kind.rep-eq)
  next
    case False
      then show ?thesis
```

```

    by (simp add: kind.rep-eq add-node.rep-eq stamp.rep-eq)
qed

```

lemma *remove-node-lookup*:

```

  gup = remove-node nid g  $\longrightarrow$  kind gup nid = NoNode  $\wedge$  stamp gup nid =
  IllegalStamp
  by (simp add: kind.rep-eq remove-node.rep-eq stamp.rep-eq)

```

lemma *replace-node-lookup*[simp]:

```

  gup = replace-node nid (k, s) g  $\wedge$  k  $\neq$  NoNode  $\longrightarrow$  kind gup nid = k  $\wedge$  stamp
  gup nid = s
  by (simp add: replace-node.rep-eq kind.rep-eq stamp.rep-eq)

```

lemma *replace-node-unchanged*:

```

  gup = replace-node nid (k, s) g  $\longrightarrow$  ( $\forall$  n  $\in$  (ids g - {nid}) . n  $\in$  ids g  $\wedge$  n  $\in$  ids
  gup  $\wedge$  kind g n = kind gup n)
  by (simp add: kind.rep-eq replace-node.rep-eq)

```

5.3.1 Example Graphs

Example 1: empty graph (just a start and end node)

definition *start-end-graph*:: *IRGraph* **where**

```

  start-end-graph = irgraph [(0, StartNode None 1, VoidStamp), (1, ReturnNode
  None None, VoidStamp)]

```

Example 2: public static int sq(int x) return x * x;

```

[1 P(0)] / [0 Start] [4 *] | / V / [5 Return]

```

definition *eg2-sq*:: *IRGraph* **where**

```

  eg2-sq = irgraph [
    (0, StartNode None 5, VoidStamp),
    (1, ParameterNode 0, default-stamp),
    (4, MulNode 1 1, default-stamp),
    (5, ReturnNode (Some 4) None, default-stamp)
  ]

```

```

value input-edges eg2-sq

```

```

value usages eg2-sq 1

```

```

end

```

5.4 Structural Graph Comparison

theory

```

  Comparison

```

imports

IRGraph
begin

We introduce a form of structural graph comparison that is able to assert structural equivalence of graphs which differ in zero or more reference node chains for any given nodes.

fun *find-ref-nodes* :: *IRGraph* \Rightarrow (*ID* \rightarrow *ID*) **where**
find-ref-nodes *g* = *map-of*
 (*map* ($\lambda n. (n, \text{ir-ref } (\text{kind } g \ n))$) (*filter* ($\lambda id. \text{is-RefNode } (\text{kind } g \ id)$) (*sorted-list-of-set* (*ids* *g*))))

fun *replace-ref-nodes* :: *IRGraph* \Rightarrow (*ID* \rightarrow *ID*) \Rightarrow *ID list* \Rightarrow *ID list* **where**
replace-ref-nodes *g m xs* = *map* ($\lambda id. (\text{case } (m \ id) \text{ of } \text{Some } other \Rightarrow other \mid \text{None} \Rightarrow id)$) *xs*

fun *find-next* :: *ID list* \Rightarrow *ID set* \Rightarrow *ID option* **where**
find-next to-see seen = (*let* *l* = (*filter* ($\lambda nid. nid \notin seen$) *to-see*)
 in (*case* *l* of [] \Rightarrow *None* | *xs* \Rightarrow *Some* (*hd* *xs*)))

inductive *reachables* :: *IRGraph* \Rightarrow *ID list* \Rightarrow *ID set* \Rightarrow *ID set* \Rightarrow *bool* **where**
reachables *g* [] {} {} |
 [None = *find-next to-see seen*] \implies *reachables* *g to-see seen seen* |
 [Some *n* = *find-next to-see seen*;
 node = *kind* *g n*;
 new = (*inputs-of* *node*) @ (*successors-of* *node*);
 reachables *g (to-see @ new) ({n} \cup *seen*) seen'*] \implies *reachables* *g to-see seen*
seen'

code-pred (*modes*: *i* \Rightarrow *i* \Rightarrow *i* \Rightarrow *o* \Rightarrow *bool*) [*show-steps*,*show-mode-inference*,*show-intermediate-results*]
reachables .

inductive *nodeEq* :: (*ID* \rightarrow *ID*) \Rightarrow *IRGraph* \Rightarrow *ID* \Rightarrow *IRGraph* \Rightarrow *ID* \Rightarrow *bool*
where
 [*kind* *g1 n1* = *RefNode ref*; *nodeEq m g1 ref g2 n2*] \implies *nodeEq m g1 n1 g2 n2* |
 [*x* = *kind* *g1 n1*;
 y = *kind* *g2 n2*;
 is-same-ir-node-type *x y*;
 replace-ref-nodes g1 m (successors-of x) = *successors-of y*;
 replace-ref-nodes g1 m (inputs-of x) = *inputs-of y*]
 \implies *nodeEq m g1 n1 g2 n2*

code-pred [*show-modes*] *nodeEq* .

fun *diffNodesGraph* :: *IRGraph* \Rightarrow *IRGraph* \Rightarrow *ID set* **where**
diffNodesGraph g1 g2 = (*let* *refNodes* = *find-ref-nodes g1 in*
 { *n* . *n* \in *Predicate.the (reachables-i-i-i-o g1 [0] {})* } \wedge (*case* *refNodes n* of *Some*
 - \Rightarrow *False* | - \Rightarrow *True*) \wedge \neg (*nodeEq refNodes g1 n g2 n*))

```

fun diffNodesInfo :: IRGraph ⇒ IRGraph ⇒ (ID × IRNode × IRNode) set (infix
  ∩s 20)
  where
    diffNodesInfo g1 g2 = {(nid, kind g1 nid, kind g2 nid) | nid . nid ∈ diffNodesGraph
      g1 g2}

fun eqGraph :: IRGraph ⇒ IRGraph ⇒ bool (infix ≈s 20)
  where
    eqGraph isabelle-graph graal-graph = ((diffNodesGraph isabelle-graph graal-graph)
      = {})

end

```

5.5 Control-flow Graph Traversal

```

theory
  Traversal
imports
  IRGraph
begin

```

```

type-synonym Seen = ID set

```

nextEdge helps determine which node to traverse next by returning the first successor edge that isn't in the set of already visited nodes. If there is not an appropriate successor, None is returned instead.

```

fun nextEdge :: Seen ⇒ ID ⇒ IRGraph ⇒ ID option where
  nextEdge seen nid g =
    (let nids = (filter (λnid'. nid' ∉ seen) (successors-of (kind g nid))) in
      (if length nids > 0 then Some (hd nids) else None))

```

pred determines which node, if any, acts as the predecessor of another.

Merge nodes represent a special case where-in the predecessor exists as an input edge of the merge node, to simplify the traversal we treat only the first input end node as the predecessor, ignoring that multiple nodes may act as a successor.

For all other nodes, the predecessor is the first element of the predecessors set. Note that in a well-formed graph there should only be one element in the predecessor set.

```

fun pred :: IRGraph ⇒ ID ⇒ ID option where
  pred g nid = (case kind g nid of
    (MergeNode ends -) ⇒ Some (hd ends) |
    - ⇒
      (if IRGraph.predecessors g nid = {}
        then None else
        Some (hd (sorted-list-of-set (IRGraph.predecessors g nid)))
      )

```

)

Here we try to implement a generic fork of the control-flow traversal algorithm that was initially implemented for the ConditionalElimination phase

type-synonym $'a \text{ TraversalState} = (ID \times Seen \times 'a)$

inductive Step

$:: ('a \text{ TraversalState} \Rightarrow 'a) \Rightarrow IRGraph \Rightarrow 'a \text{ TraversalState} \Rightarrow 'a \text{ TraversalState option} \Rightarrow bool$

for $sa \ g$ **where**

— Hit a BeginNode with an IfNode predecessor which represents the start of a basic block for the IfNode. 1. nid' will be the successor of the begin node. 2. Find the first and only predecessor. 3. Extract condition from the preceding IfNode. 4. Negate condition if the begin node is second branch (we've taken the else branch of the condition) 5. Add the condition or the negated condition to stack 6. Perform any stamp updates based on the condition using the registerNewCondition function and place them on the top of the stack of stamp information

$\llbracket kind \ g \ nid = \text{BeginNode } nid';$

$nid \notin seen;$
 $seen' = \{nid\} \cup seen;$

$\text{Some } ifcond = pred \ g \ nid;$
 $kind \ g \ ifcond = \text{IfNode } cond \ t \ f;$

$analysis' = sa \ (nid, seen, analysis)$
 $\implies Step \ sa \ g \ (nid, seen, analysis) \ (\text{Some } (nid', seen', analysis')) \mid$

— Hit an EndNode 1. nid' will be the usage of EndNode 2. pop the conditions and stamp stack

$\llbracket kind \ g \ nid = \text{EndNode};$

$nid \notin seen;$
 $seen' = \{nid\} \cup seen;$

$nid' = any\text{-usage } g \ nid;$

$analysis' = sa \ (nid, seen, analysis)$
 $\implies Step \ sa \ g \ (nid, seen, analysis) \ (\text{Some } (nid', seen', analysis')) \mid$

— We can find a successor edge that is not in seen, go there

$\llbracket \neg(is\text{-EndNode } (kind \ g \ nid));$
 $\neg(is\text{-BeginNode } (kind \ g \ nid));$

$nid \notin seen;$
 $seen' = \{nid\} \cup seen;$

$\text{Some } nid' = nextEdge \ seen' \ nid \ g;$

```

    analysis' = sa (nid, seen, analysis)
  => Step sa g (nid, seen, analysis) (Some (nid', seen', analysis')) |

  — We cannot find a successor edge that is not in seen, give back None
  [[¬(is-EndNode (kind g nid));
    ¬(is-BEGINNode (kind g nid));

    nid ∉ seen;
    seen' = {nid} ∪ seen;

    None = nextEdge seen' nid g]]
  => Step sa g (nid, seen, analysis) None |

  — We've already seen this node, give back None
  [[nid ∈ seen]] => Step sa g (nid, seen, analysis) None

code-pred (modes: i ⇒ i ⇒ i ⇒ o ⇒ bool) Step .

end

```

6 Data-flow Semantics

```

theory IRTreeEval
  imports
    Graph.Stamp
  begin

```

We define a tree representation of data-flow nodes, as an abstraction of the graph view.

Data-flow trees are evaluated in the context of a method state (currently called MapState in the theories for historical reasons).

The method state consists of the values for each method parameter, references to method parameters use an index of the parameter within the parameter list, as such we store a list of parameter values which are looked up at parameter references.

The method state also stores a mapping of node ids to values. The contents of this mapping is calculated during the traversal of the control flow graph.

As a concrete example, as the *SignedDivNode::'a* can have side-effects (during division by zero), it is treated as part of the control-flow, since the data-flow phase is specified to be side-effect free. As a result, the control-flow semantics for *SignedDivNode::'a* calculates the value of a node and maps the node identifier to the value within the method state. The data-flow semantics then just reads the value stored in the method state for the node.

```

type-synonym ID = nat
type-synonym MapState = ID ⇒ Value
type-synonym Params = Value list

```

definition *new-map-state* :: *MapState* **where**

new-map-state = ($\lambda x.$ *UndefVal*)

6.1 Data-flow Tree Representation

datatype *IRUnaryOp* =

UnaryAbs
 | *UnaryNeg*
 | *UnaryNot*
 | *UnaryLogicNegation*
 | *UnaryNarrow* (*ir-inputBits*: *nat*) (*ir-resultBits*: *nat*)
 | *UnarySignExtend* (*ir-inputBits*: *nat*) (*ir-resultBits*: *nat*)
 | *UnaryZeroExtend* (*ir-inputBits*: *nat*) (*ir-resultBits*: *nat*)

datatype *IRBinaryOp* =

BinAdd
 | *BinMul*
 | *BinSub*
 | *BinAnd*
 | *BinOr*
 | *BinXor*
 | *BinShortCircuitOr*
 | *BinLeftShift*
 | *BinRightShift*
 | *BinURightShift*
 | *BinIntegerEquals*
 | *BinIntegerLessThan*
 | *BinIntegerBelow*

datatype (*discs-sels*) *IRExpr* =

UnaryExpr (*ir-uop*: *IRUnaryOp*) (*ir-value*: *IRExpr*)
 | *BinaryExpr* (*ir-op*: *IRBinaryOp*) (*ir-x*: *IRExpr*) (*ir-y*: *IRExpr*)
 | *ConditionalExpr* (*ir-condition*: *IRExpr*) (*ir-trueValue*: *IRExpr*) (*ir-falseValue*:
IRExpr)

| *ParameterExpr* (*ir-index*: *nat*) (*ir-stamp*: *Stamp*)

| *LeafExpr* (*ir-nid*: *ID*) (*ir-stamp*: *Stamp*)

| *ConstantExpr* (*ir-const*: *Value*)

| *ConstantVar* (*ir-name*: *string*)

| *VariableExpr* (*ir-name*: *string*) (*ir-stamp*: *Stamp*)

fun *is-ground* :: *IRExpr* \Rightarrow *bool* **where**

is-ground (*UnaryExpr* *op* *e*) = *is-ground* *e* |

is-ground (*BinaryExpr* *op* *e1* *e2*) = (*is-ground* *e1* \wedge *is-ground* *e2*) |

is-ground (*ConditionalExpr* *b* *e1* *e2*) = (*is-ground* *b* \wedge *is-ground* *e1* \wedge *is-ground*

```

e2) |
  is-ground (ParameterExpr i s) = True |
  is-ground (LeafExpr n s) = True |
  is-ground (ConstantExpr v) = True |
  is-ground (ConstantVar name) = False |
  is-ground (VariableExpr name s) = False

```

```

typedef GroundExpr = { e :: IRExpr . is-ground e }
using is-ground.simps(6) by blast

```

6.2 Functions for re-calculating stamps

Note: in Java all integer calculations are done as 32 or 64 bit calculations. However, here we generalise the operators to allow any size calculations. Many operators have the same output bits as their inputs. However, the unary integer operators that are not *normal_unary* are narrowing or widening operators, so the result bits is specified by the operator. The binary integer operators are divided into three groups: (1) *binary_fixed_32* operators always output 32 bits, (2) *binary_shift_ops* operators output size is determined by their left argument, and (3) other operators output the same number of bits as both their inputs.

abbreviation *binary-fixed-32-ops* :: IRBinaryOp set **where**
binary-fixed-32-ops ≡ {BinShortCircuitOr, BinIntegerEquals, BinIntegerLessThan, BinIntegerBelow}

abbreviation *binary-shift-ops* :: IRBinaryOp set **where**
binary-shift-ops ≡ {BinLeftShift, BinRightShift, BinURightShift}

abbreviation *normal-unary* :: IRUnaryOp set **where**
normal-unary ≡ {UnaryAbs, UnaryNeg, UnaryNot, UnaryLogicNegation}

fun stamp-unary :: IRUnaryOp ⇒ Stamp ⇒ Stamp **where**

```

stamp-unary op (IntegerStamp b lo hi) =
  unrestricted-stamp (IntegerStamp (if op ∈ normal-unary then b else (ir-resultBits
op)) lo hi) |

```

```

stamp-unary op - = IllegalStamp

```

fun stamp-binary :: IRBinaryOp ⇒ Stamp ⇒ Stamp ⇒ Stamp **where**
stamp-binary op (IntegerStamp b1 lo1 hi1) (IntegerStamp b2 lo2 hi2) =
 (if op ∈ binary-shift-ops then unrestricted-stamp (IntegerStamp b1 lo1 hi1)
 else if b1 ≠ b2 then IllegalStamp else
 (if op ∈ binary-fixed-32-ops
 then unrestricted-stamp (IntegerStamp 32 lo1 hi1)
 else unrestricted-stamp (IntegerStamp b1 lo1 hi1))) |

```

stamp-binary op - - = IllegalStamp

```



```

fun stamp-expr :: IRExpr ⇒ Stamp where
  stamp-expr (UnaryExpr op x) = stamp-unary op (stamp-expr x) |
  stamp-expr (BinaryExpr bop x y) = stamp-binary bop (stamp-expr x) (stamp-expr
y) |
  stamp-expr (ConstantExpr val) = constantAsStamp val |
  stamp-expr (LeafExpr i s) = s |
  stamp-expr (ParameterExpr i s) = s |
  stamp-expr (ConditionalExpr c t f) = meet (stamp-expr t) (stamp-expr f)

export-code stamp-unary stamp-binary stamp-expr

```

6.3 Data-flow Tree Evaluation

```

fun unary-eval :: IRUnaryOp ⇒ Value ⇒ Value where
  unary-eval UnaryAbs v = intval-abs v |
  unary-eval UnaryNeg v = intval-negate v |
  unary-eval UnaryNot v = intval-not v |
  unary-eval UnaryLogicNegation v = intval-logic-negation v |
  unary-eval (UnaryNarrow inBits outBits) v = intval-narrow inBits outBits v |
  unary-eval (UnarySignExtend inBits outBits) v = intval-sign-extend inBits outBits
v |
  unary-eval (UnaryZeroExtend inBits outBits) v = intval-zero-extend inBits outBits
v

```

```

fun bin-eval :: IRBinaryOp ⇒ Value ⇒ Value ⇒ Value where
  bin-eval BinAdd v1 v2 = intval-add v1 v2 |
  bin-eval BinMul v1 v2 = intval-mul v1 v2 |
  bin-eval BinSub v1 v2 = intval-sub v1 v2 |
  bin-eval BinAnd v1 v2 = intval-and v1 v2 |
  bin-eval BinOr v1 v2 = intval-or v1 v2 |
  bin-eval BinXor v1 v2 = intval-xor v1 v2 |
  bin-eval BinShortCircuitOr v1 v2 = intval-short-circuit-or v1 v2 |
  bin-eval BinLeftShift v1 v2 = intval-left-shift v1 v2 |
  bin-eval BinRightShift v1 v2 = intval-right-shift v1 v2 |
  bin-eval BinURightShift v1 v2 = intval-uright-shift v1 v2 |
  bin-eval BinIntegerEquals v1 v2 = intval-equals v1 v2 |
  bin-eval BinIntegerLessThan v1 v2 = intval-less-than v1 v2 |
  bin-eval BinIntegerBelow v1 v2 = intval-below v1 v2

```

```

lemmas eval-thms =
  intval-abs.simps intval-negate.simps intval-not.simps
  intval-logic-negation.simps intval-narrow.simps
  intval-sign-extend.simps intval-zero-extend.simps
  intval-add.simps intval-mul.simps intval-sub.simps
  intval-and.simps intval-or.simps intval-xor.simps
  intval-left-shift.simps intval-right-shift.simps

```

intval-uright-shift.simps intval-equals.simps
intval-less-than.simps intval-below.simps

inductive *not-undef-or-fail* :: *Value* \Rightarrow *Value* \Rightarrow *bool* **where**
 $\llbracket \text{value} \neq \text{UndefVal} \rrbracket \implies \text{not-undef-or-fail value value}$

notation (*latex output*)
not-undef-or-fail (- = -)

inductive
evaltree :: *MapState* \Rightarrow *Params* \Rightarrow *IRExpr* \Rightarrow *Value* \Rightarrow *bool* ($[-,-] \vdash - \mapsto -$ 55)
for *m p* **where**

ConstantExpr:
 $\llbracket \text{wf-value } c \rrbracket$
 $\implies [m,p] \vdash (\text{ConstantExpr } c) \mapsto c \mid$

ParameterExpr:
 $\llbracket i < \text{length } p; \text{valid-value } (p!i) \ s \rrbracket$
 $\implies [m,p] \vdash (\text{ParameterExpr } i \ s) \mapsto p!i \mid$

ConditionalExpr:
 $\llbracket [m,p] \vdash ce \mapsto cond;$
 $cond \neq \text{UndefVal};$
 $branch = (\text{if val-to-bool } cond \text{ then } te \text{ else } fe);$
 $[m,p] \vdash branch \mapsto result;$
 $result \neq \text{UndefVal} \rrbracket$
 $\implies [m,p] \vdash (\text{ConditionalExpr } ce \ te \ fe) \mapsto result \mid$

UnaryExpr:
 $\llbracket [m,p] \vdash xe \mapsto x;$
 $result = (\text{unary-eval } op \ x);$
 $result \neq \text{UndefVal} \rrbracket$
 $\implies [m,p] \vdash (\text{UnaryExpr } op \ xe) \mapsto result \mid$

BinaryExpr:
 $\llbracket [m,p] \vdash xe \mapsto x;$
 $[m,p] \vdash ye \mapsto y;$
 $result = (\text{bin-eval } op \ x \ y);$
 $result \neq \text{UndefVal} \rrbracket$
 $\implies [m,p] \vdash (\text{BinaryExpr } op \ xe \ ye) \mapsto result \mid$

LeafExpr:
 $\llbracket val = m \ n;$
 $\text{valid-value } val \ s \rrbracket$
 $\implies [m,p] \vdash \text{LeafExpr } n \ s \mapsto val$

code-pred (*modes*: $i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow \text{bool}$ as *evalT*)

```

[show-steps,show-mode-inference,show-intermediate-results]
evaltree .

inductive
  evaltrees :: MapState  $\Rightarrow$  Params  $\Rightarrow$  IRExp list  $\Rightarrow$  Value list  $\Rightarrow$  bool ( $[-,] \vdash - \mapsto_L$ 
- 55)
  for m p where

    EvalNil:
    [m,p]  $\vdash [] \mapsto_L []$  |

    EvalCons:
    [[m,p]  $\vdash x \mapsto xval$ ;
    [m,p]  $\vdash yy \mapsto_L yyval$ ]
     $\Rightarrow$  [m,p]  $\vdash (x\#yy) \mapsto_L (xval\#yyval)$ 

code-pred (modes: i  $\Rightarrow$  i  $\Rightarrow$  i  $\Rightarrow$  o  $\Rightarrow$  bool as evalTs)
  evaltrees .

definition sq-param0 :: IRExp where
  sq-param0 = BinaryExpr BinMul
    (ParameterExpr 0 (IntegerStamp 32 (- 2147483648) 2147483647))
    (ParameterExpr 0 (IntegerStamp 32 (- 2147483648) 2147483647))

values {v. evaltree new-map-state [IntVal 32 5] sq-param0 v}

declare evaltree.intros [intro]
declare evaltrees.intros [intro]

```

6.4 Data-flow Tree Refinement

We define the induced semantic equivalence relation between expressions. Note that syntactic equality implies semantic equivalence, but not vice versa.

definition *equiv-exprs* :: IRExp \Rightarrow IRExp \Rightarrow bool ($- \doteq -$ 55) **where**
 $(e1 \doteq e2) = (\forall m p v. ([m,p] \vdash e1 \mapsto v) \longleftrightarrow ([m,p] \vdash e2 \mapsto v))$

We also prove that this is a total equivalence relation (*equivp equiv-exprs*) (HOL.Equiv_Relations), so that we can reuse standard results about equivalence relations.

lemma *equivp equiv-exprs*
apply (auto simp add: equivp-def equiv-exprs-def)
by (metis equiv-exprs-def)+

We define a refinement ordering over IRExp and show that it is a preorder. Note that it is asymmetric because e2 may refer to fewer variables than e1.

instantiation *IRExpr* :: *preorder* **begin**

notation *less-eq* (**infix** \sqsubseteq 65)

definition

le-expr-def [*simp*]:

$$(e_2 \leq e_1) \longleftrightarrow (\forall m p v. (([m,p] \vdash e_1 \mapsto v) \longrightarrow ([m,p] \vdash e_2 \mapsto v)))$$

definition

lt-expr-def [*simp*]:

$$(e_1 < e_2) \longleftrightarrow (e_1 \leq e_2 \wedge \neg (e_1 \dot{=} e_2))$$

instance proof

fix *x y z* :: *IRExpr*

show $x < y \longleftrightarrow x \leq y \wedge \neg (y \leq x)$ **by** (*simp add: equiv-exprs-def; auto*)

show $x \leq x$ **by** *simp*

show $x \leq y \implies y \leq z \implies x \leq z$ **by** *simp*

qed

end

abbreviation (**output**) *Refines* :: *IRExpr* \Rightarrow *IRExpr* \Rightarrow *bool* (**infix** \sqsupseteq 64)

where $e_1 \sqsupseteq e_2 \equiv (e_2 \leq e_1)$

6.5 Stamp Masks

A stamp can contain additional range information in the form of masks. A stamp has an up mask and a down mask, corresponding to the bits that may be set and the bits that must be set.

Examples: A stamp where no range information is known will have; an up mask of -1 as all bits may be set, and a down mask of 0 as no bits must be set.

A stamp known to be one should have; an up mask of 1 as only the first bit may be set, no others, and a down mask of 1 as the first bit must be set and no others.

We currently don't carry mask information in stamps, and instead assume correct masks to prove optimizations.

locale *stamp-mask* =

fixes *up* :: *IRExpr* \Rightarrow *int64* (\uparrow)

fixes *down* :: *IRExpr* \Rightarrow *int64* (\downarrow)

assumes *up-spec*: $[m, p] \vdash e \mapsto \text{IntVal } b \ v \implies (\text{and } v \ (\text{not } ((\text{ucast } (\uparrow e)))) = 0$

and *down-spec*: $[m, p] \vdash e \mapsto \text{IntVal } b \ v \implies (\text{and } (\text{not } v) \ (\text{ucast } (\downarrow e))) = 0$

begin

lemma *may-implies-either*:

$[m, p] \vdash e \mapsto \text{IntVal } b \ v \implies \text{bit } (\uparrow e) \ n \implies \text{bit } v \ n = \text{False} \vee \text{bit } v \ n = \text{True}$

by *simp*

lemma *not-may-implies-false*:

$[m, p] \vdash e \mapsto \text{IntVal } b \ v \implies \neg(\text{bit } (\uparrow e) \ n) \implies \text{bit } v \ n = \text{False}$
using *up-spec*
using *bit-and-iff bit-eq-iff bit-not-iff bit-unsigned-iff down-spec*
by (*smt (verit, best) bit.double-compl*)

lemma *must-implies-true*:

$[m, p] \vdash e \mapsto \text{IntVal } b \ v \implies \text{bit } (\downarrow e) \ n \implies \text{bit } v \ n = \text{True}$
using *down-spec*
by (*metis bit.compl-one bit-and-iff bit-minus-1-iff bit-not-iff impossible-bit ucast-id*)

lemma *not-must-implies-either*:

$[m, p] \vdash e \mapsto \text{IntVal } b \ v \implies \neg(\text{bit } (\downarrow e) \ n) \implies \text{bit } v \ n = \text{False} \vee \text{bit } v \ n = \text{True}$
by *simp*

lemma *must-implies-may*:

$[m, p] \vdash e \mapsto \text{IntVal } b \ v \implies n < 32 \implies \text{bit } (\downarrow e) \ n \implies \text{bit } (\uparrow e) \ n$
by (*meson must-implies-true not-may-implies-false*)

lemma *up-mask-and-zero-implies-zero*:

assumes *and* $(\uparrow x) (\uparrow y) = 0$
assumes $[m, p] \vdash x \mapsto \text{IntVal } b \ xv$
assumes $[m, p] \vdash y \mapsto \text{IntVal } b \ yv$
shows *and* $xv \ yv = 0$
using *assms*
by (*smt (z3) and.commute and.right-neutral and-zero-eq bit.compl-zero bit.conj-cancel-right bit.conj-disj-distrib(1) ucast-id up-spec word-bw-assocs(1) word-not-dist(2)*)

lemma *not-down-up-mask-and-zero-implies-zero*:

assumes *and* $(\text{not } (\downarrow x)) (\uparrow y) = 0$
assumes $[m, p] \vdash x \mapsto \text{IntVal } b \ xv$
assumes $[m, p] \vdash y \mapsto \text{IntVal } b \ yv$
shows *and* $xv \ yv = yv$
using *assms*
by (*smt (z3) and-zero-eq bit.conj-cancel-left bit.conj-disj-distrib(1) bit.conj-disj-distrib(2) bit.de-Morgan-disj down-spec or-eq-not-not-and ucast-id up-spec word-ao-absorbs(2) word-ao-absorbs(8) word-bw-lcs(1) word-not-dist(2)*)

end

definition *IRExpr-up* :: *IRExpr* \Rightarrow *int64* **where**

IRExpr-up *e* = *not 0*

definition *IRExpr-down* :: *IRExpr* \Rightarrow *int64* **where**

IRExpr-down *e* = *0*

lemma *ucast-zero*: $(\text{ucast } (0::\text{int64})::\text{int32}) = 0$

```

    by simp

lemma ucast-minus-one: (ucast (-1::int64)::int32) = -1
  apply transfer by auto

interpretation simple-mask: stamp-mask
  IRExp-up :: IRExp ⇒ int64
  IRExp-down :: IRExp ⇒ int64
  unfolding IRExp-up-def IRExp-down-def
  apply unfold-locales
  by (simp add: ucast-minus-one)+

end

```

6.6 Data-flow Tree Theorems

```

theory IRTreeEvalThms
  imports
    Graph.ValueThms
    IRTreeEval
begin

```

6.6.1 Deterministic Data-flow Evaluation

```

lemma evalDet:
  [m,p] ⊢ e ↦ v₁ ⇒
  [m,p] ⊢ e ↦ v₂ ⇒
  v₁ = v₂
  apply (induction arbitrary: v₂ rule: evaltree.induct)
  by (elim EvalTreeE; auto)+

```

```

lemma evalAllDet:
  [m,p] ⊢ e ↦L v1 ⇒
  [m,p] ⊢ e ↦L v2 ⇒
  v1 = v2
  apply (induction arbitrary: v2 rule: evaltrees.induct)
  apply (elim EvalTreeE; auto)
  using evalDet by force

```

6.6.2 Typing Properties for Integer Evaluation Functions

We use three simple typing properties on integer values: *isIntVal32*, *isIntVal64* and the more general *isIntVal*.

```

lemma unary-eval-not-obj-ref:
  shows unary-eval op x ≠ ObjRef v
  by (cases op; cases x; auto)

```

```

lemma unary-eval-not-obj-str:
  shows unary-eval op x ≠ ObjStr v

```

by (*cases op; cases x; auto*)

lemma *unary-eval-int*:

assumes *def*: *unary-eval op x* \neq *UndefVal*
shows *is-IntVal* (*unary-eval op x*)
unfolding *is-IntVal-def* **using** *def*
apply (*cases unary-eval op x; auto*)
using *unary-eval-not-obj-ref unary-eval-not-obj-str* **by** *simp+*

lemma *bin-eval-int*:

assumes *def*: *bin-eval op x y* \neq *UndefVal*
shows *is-IntVal* (*bin-eval op x y*)
apply (*cases op; cases x; cases y*)
unfolding *is-IntVal-def* **using** *def* **apply** *auto*
apply *presburger+*
apply (*meson bool-to-val.elims*)
apply (*meson bool-to-val.elims*)
apply (*smt (verit) new-int.simps*)
by (*meson bool-to-val.elims*)**+**

lemma *IntVal0*:

(*IntVal 32 0*) = (*new-int 32 0*)
unfolding *new-int.simps*
by *auto*

lemma *IntVal1*:

(*IntVal 32 1*) = (*new-int 32 1*)
unfolding *new-int.simps*
by *auto*

lemma *bin-eval-new-int*:

assumes *def*: *bin-eval op x y* \neq *UndefVal*
shows $\exists b v. (bin-eval op x y) = new-int b v \wedge$
 $b = (if op \in binary-fixed-32-ops then 32 else intval-bits x)$
apply (*cases op; cases x; cases y*)
unfolding *is-IntVal-def* **using** *def* **apply** *auto*
apply *presburger+*
apply (*metis take-bit-and*)
apply *presburger*
apply (*metis take-bit-or*)
apply *presburger*
apply (*metis take-bit-xor*)

```

apply presburger
using IntVal0 IntVal1
apply (metis bool-to-val.elims new-int.simps)
apply presburger
apply (smt (verit) new-int.elims)
apply (smt (verit, best) new-int.elims)
apply (metis IntVal0 IntVal1 bool-to-val.elims new-int.simps)
apply presburger
apply (metis IntVal0 IntVal1 bool-to-val.elims new-int.simps)
apply presburger
apply (metis IntVal0 IntVal1 bool-to-val.elims new-int.simps)
by meson

lemma int-stamp:
  assumes i: is-IntVal v
  shows is-IntegerStamp (constantAsStamp v)
  using i unfolding is-IntegerStamp-def is-IntVal-def by auto

lemma validStampIntConst:
  assumes v = IntVal b ival
  assumes  $0 < b \wedge b \leq 64$ 
  shows valid-stamp (constantAsStamp v)
proof –
  have bnds: fst (bit-bounds b) ≤ int-signed-value b ival ∧ int-signed-value b ival
   $\leq \text{snd (bit-bounds b)}$ 
  using assms int-signed-value-bounds
  by presburger
  have s: constantAsStamp v = IntegerStamp b (int-signed-value b ival) (int-signed-value
b ival)
  using assms(1) constantAsStamp.simps(1) by blast
  then show ?thesis
  unfolding s valid-stamp.simps
  using assms(2) assms bnds by linarith
qed

lemma validDefIntConst:
  assumes v: v = IntVal b ival
  assumes  $0 < b \wedge b \leq 64$ 
  assumes take-bit b ival = ival
  shows valid-value v (constantAsStamp v)
proof –
  have bnds: fst (bit-bounds b) ≤ int-signed-value b ival ∧ int-signed-value b ival
   $\leq \text{snd (bit-bounds b)}$ 
  using assms int-signed-value-bounds
  by presburger
  have s: constantAsStamp v = IntegerStamp b (int-signed-value b ival) (int-signed-value
b ival)

```



```

    using assms(1) constantAsStamp.simps(1) by blast
  then show ?thesis
    unfolding s unfolding v unfolding valid-value.simps
    using assms validStampIntConst
    by simp
qed

```

6.6.3 Evaluation Results are Valid

A valid value cannot be *UndefVal*.

```

lemma valid-not-undef:
  assumes a1: valid-value val s
  assumes a2: s ≠ VoidStamp
  shows val ≠ UndefVal
  apply (rule valid-value.elims(1)[of val s True])
  using a1 a2 by auto

```

```

lemma valid-VoidStamp[elim]:
  shows valid-value val VoidStamp ⇒
    val = UndefVal
  using valid-value.simps by metis

```

```

lemma valid-ObjStamp[elim]:
  shows valid-value val (ObjectStamp klass exact nonNull alwaysNull) ⇒
    (∃ v. val = ObjRef v)
  using valid-value.simps by (metis val-to-bool.cases)

```

```

lemma valid-int[elim]:
  shows valid-value val (IntegerStamp b lo hi) ⇒
    (∃ v. val = IntVal b v)
  using valid-value.elims(2) by fastforce

```

```

lemmas valid-value-elim =
  valid-VoidStamp
  valid-ObjStamp
  valid-int

```

```

lemma evaltree-not-undef:
  fixes m p e v
  shows ([m,p] ⊢ e ↦ v) ⇒ v ≠ UndefVal
  apply (induction rule: evaltree.induct)
  using valid-not-undef wf-value-def by auto

```

```

lemma leafint:
  assumes ev: [m,p] ⊢ LeafExpr i (IntegerStamp b lo hi) ↦ val
  shows ∃ b v. val = (IntVal b v)

```

```

proof –
  have valid-value val (IntegerStamp b lo hi)
    using ev by (rule LeafExprE; simp)
  then show ?thesis by auto
qed

```

```

lemma default-stamp [simp]: default-stamp = IntegerStamp 32 (–2147483648)
2147483647
  using default-stamp-def by auto

```

```

lemma valid-value-signed-int-range [simp]:
  assumes valid-value val (IntegerStamp b lo hi)
  assumes lo < 0
  shows  $\exists v. (val = \text{IntVal } b \ v \wedge$ 
     $lo \leq \text{int-signed-value } b \ v \wedge$ 
     $\text{int-signed-value } b \ v \leq hi)$ 
  using assms valid-int
  by (metis valid-value.simps(1))

```

6.6.4 Example Data-flow Optimisations

6.6.5 Monotonicity of Expression Refinement

We prove that each subexpression position is monotonic. That is, optimizing a subexpression anywhere deep inside a top-level expression also optimizes that top-level expression.

Note that we might also be able to do this via reusing Isabelle’s *mono* operator (HOL.Orderings theory), proving instantiations like *mono(UnaryExprop)*, but it is not obvious how to do this for both arguments of the binary expressions.

```

lemma mono-unary:
  assumes  $x \geq x'$ 
  shows  $(\text{UnaryExpr op } x) \geq (\text{UnaryExpr op } x')$ 
  using UnaryExpr assms by auto

```

```

lemma mono-binary:
  assumes  $x \geq x'$ 
  assumes  $y \geq y'$ 
  shows  $(\text{BinaryExpr op } x \ y) \geq (\text{BinaryExpr op } x' \ y')$ 
  using BinaryExpr assms by auto

```

```

lemma never-void:
  assumes  $[m, p] \vdash x \mapsto xv$ 
  assumes valid-value xv (stamp-expr xe)
  shows stamp-expr xe  $\neq \text{VoidStamp}$ 
  using valid-value.simps

```

using *assms*(2) **by** *force*

lemma *compatible-trans*:

compatible x y \wedge *compatible y z* \implies *compatible x z*
by (*cases x*; *cases y*; *cases z*; *simp del: valid-stamp.simps*)

lemma *compatible-reft*:

compatible x y \implies *compatible y x*
using *compatible.elims*(2) **by** *fastforce*

lemma *mono-conditional*:

assumes $c \geq c'$
assumes $t \geq t'$
assumes $f \geq f'$
shows $(\text{ConditionalExpr } c \ t \ f) \geq (\text{ConditionalExpr } c' \ t' \ f')$
proof (*simp only: le-expr-def; (rule allI)+; rule impI*)
fix $m \ p \ v$
assume $a: [m, p] \vdash \text{ConditionalExpr } c \ t \ f \mapsto v$
then obtain *cond* **where** $c: [m, p] \vdash c \mapsto \text{cond}$ **by** *auto*
then have $c': [m, p] \vdash c' \mapsto \text{cond}$ **using** *assms* **by** *auto*

define *branch* **where** $b: \text{branch} = (\text{if val-to-bool cond then } t \text{ else } f)$
define *branch'* **where** $b': \text{branch}' = (\text{if val-to-bool cond then } t' \text{ else } f')$
then have $\text{beval}: [m, p] \vdash \text{branch} \mapsto v$ **using** $a \ b \ c \ \text{evalDet}$ **by** *blast*

from *beval* **have** $[m, p] \vdash \text{branch}' \mapsto v$ **using** *assms b b'* **by** *auto*
then show $[m, p] \vdash \text{ConditionalExpr } c' \ t' \ f' \mapsto v$
using *ConditionalExpr c' b'*
by (*simp add: evaltree-not-undef*)
qed

6.7 Unfolding rules for evaltree quadruples down to bin-eval level

These rewrite rules can be useful when proving optimizations. They support top-down rewriting of each level of the tree into the lower-level *bin_eval* / *unary_eval* level, simply by saying *unfoldingunfold_evaltree*.

lemma *unfold-const*:

shows $([m, p] \vdash \text{ConstantExpr } c \mapsto v) = (\text{wf-value } v \wedge v = c)$
by *blast*

```

lemma unfold-binary:
  shows ( $[m,p] \vdash \text{BinaryExpr } op \ x \ y \mapsto val$ ) = ( $\exists \ x \ y.$ 
    ( $[m,p] \vdash x \mapsto x$ )  $\wedge$ 
    ( $[m,p] \vdash y \mapsto y$ )  $\wedge$ 
    ( $val = \text{bin-eval } op \ x \ y$ )  $\wedge$ 
    ( $val \neq \text{UndefVal}$ )
  ) (is ?L = ?R)
proof (intro iffI)
  assume  $\mathcal{B}$ : ?L
  show ?R by (rule evaltree.cases[OF  $\mathcal{B}$ ]; blast+)
next
  assume ?R
  then obtain  $x \ y$  where  $[m,p] \vdash x \mapsto x$ 
    and  $[m,p] \vdash y \mapsto y$ 
    and  $val = \text{bin-eval } op \ x \ y$ 
    and  $val \neq \text{UndefVal}$ 
  by auto
  then show ?L
    by (rule BinaryExpr)
qed

```

```

lemma unfold-unary:
  shows ( $[m,p] \vdash \text{UnaryExpr } op \ x \mapsto val$ )
    = ( $\exists \ x.$ 
    ( $[m,p] \vdash x \mapsto x$ )  $\wedge$ 
    ( $val = \text{unary-eval } op \ x$ )  $\wedge$ 
    ( $val \neq \text{UndefVal}$ )
  ) (is ?L = ?R)
by auto

```

```

lemmas unfold-evaltree =
  unfold-binary
  unfold-unary

```

6.8 Lemmas about *new__int* and integer eval results.

```

lemma unary-eval-new-int:
  assumes def:  $\text{unary-eval } op \ x \neq \text{UndefVal}$ 
  shows  $\exists \ b \ v. \text{unary-eval } op \ x = \text{new-int } b \ v \wedge$ 
     $b = (\text{if } op \in \text{normal-unary} \text{ then } \text{intval-bits } x \text{ else } \text{ir-resultBits } op)$ 
proof (cases op ∈ normal-unary)
  case True
  then show ?thesis
    by (metis def empty-iff insert-iff intval-abs.elims intval-bits.simps intval-logic-negation.elims
      intval-negate.elims intval-not.elims unary-eval.simps(1) unary-eval.simps(2) unary-eval.simps(3)
      unary-eval.simps(4))

```

```

next
  case False
  consider ib ob where op = UnaryNarrow ib ob |
    ib ob where op = UnaryZeroExtend ib ob |
    ib ob where op = UnarySignExtend ib ob
  by (metis False IRUnaryOp.exhaust insert-iff)
  then show ?thesis
  proof (cases)
    case 1
    then show ?thesis
    by (metis False IRUnaryOp.sel(4) def intval-narrow.elims unary-eval.simps(5))
  next
    case 2
    then show ?thesis
    by (metis False IRUnaryOp.sel(6) def intval-zero-extend.elims unary-eval.simps(7))
  next
    case 3
    then show ?thesis
    by (metis False IRUnaryOp.sel(5) def intval-sign-extend.elims unary-eval.simps(6))
  qed
qed

```

```

lemma new-int-unused-bits-zero:
  assumes IntVal b ival = new-int b ival0
  shows take-bit b ival = ival
  using assms(1) new-int-take-bits by blast

```

```

lemma unary-eval-unused-bits-zero:
  assumes unary-eval op x = IntVal b ival
  shows take-bit b ival = ival
  using assms unary-eval-new-int
  by (metis Value.inject(1) Value.simps(5) new-int.elims new-int-unused-bits-zero)

```

```

lemma bin-eval-unused-bits-zero:
  assumes bin-eval op x y = (IntVal b ival)
  shows take-bit b ival = ival
  using assms bin-eval-new-int
  by (metis Value.distinct(1) Value.inject(1) new-int.elims new-int-take-bits)

```

```

lemma eval-unused-bits-zero:
   $[m,p] \vdash xe \mapsto (IntVal\ b\ ix) \implies take-bit\ b\ ix = ix$ 
  proof (induction xe)
    case (UnaryExpr x1 xe)
    then show ?case
    using unary-eval-unused-bits-zero by force
  next
    case (BinaryExpr x1 xe1 xe2)
    then show ?case
    using bin-eval-unused-bits-zero by force

```

```

next
  case (ConditionalExpr xe1 xe2 xe3)
  then show ?case
    by (metis (full-types) EvalTreeE(3))
next
  case (ParameterExpr i s)
  then have valid-value (p!i) s
    by fastforce
  then show ?case
    by (metis ParameterExprE Value.distinct(7) intval-bits.simps intval-word.simps
local.ParameterExpr valid-value.elims(2))
next
  case (LeafExpr x1 x2)
  then show ?case
    by (smt (z3) EvalTreeE(6) Value.simps(11) valid-value.elims(1) valid-value.simps(1))

next
  case (ConstantExpr x)
  then show ?case using wf-value-def
    by (metis EvalTreeE(1) constantAsStamp.simps(1) valid-value.simps(1))
next
  case (ConstantVar x)
  then show ?case
    by fastforce
next
  case (VariableExpr x1 x2)
  then show ?case
    by fastforce
qed

```

```

lemma unary-normal-bitsize:
  assumes unary-eval op x = IntVal b ival
  assumes op ∈ normal-unary
  shows ∃ ix. x = IntVal b ix
  apply (cases op)
    prefer 7 using assms apply blast
    prefer 6 using assms apply blast
    prefer 5 using assms apply blast
  using Value.distinct(1) Value.sel(1) assms(1) new-int.simps unary-eval.simps
    intval-abs.elims intval-negate.elims intval-not.elims intval-logic-negation.elims
  apply metis+
done

```

```

lemma unary-not-normal-bitsize:
  assumes unary-eval op x = IntVal b ival
  assumes op ∉ normal-unary
  shows b = ir-resultBits op ∧ 0 < b ∧ b ≤ 64
  apply (cases op)

```

```

using assms apply blast+
apply (metis IRUnaryOp.sel(4) Value.distinct(1) Value.sel(1) assms(1) intval-narrow.elims
intval-narrow-ok new-int.simps unary-eval.simps(5))
apply (smt (verit) IRUnaryOp.sel(5) Value.distinct(1) Value.sel(1) assms(1)
intval-sign-extend.elims new-int.simps order-less-le-trans unary-eval.simps(6))
apply (metis IRUnaryOp.sel(6) Value.distinct(1) assms(1) intval-bits.simps int-
val-zero-extend.elims linorder-not-less neq0-conv new-int.simps unary-eval.simps(7))
done

```

```

lemma unary-eval-bitsize:
  assumes unary-eval op x = IntVal b ival
  assumes  $2: x = \text{IntVal } bx \ ix$ 
  assumes  $0 < bx \wedge bx \leq 64$ 
  shows  $0 < b \wedge b \leq 64$ 
proof (cases op ∈ normal-unary)
  case True
  then obtain tmp where unary-eval op x = new-int bx tmp
  by (cases op; simp; auto simp: 2)
  then show ?thesis
  using assms by simp
next
  case False
  then obtain tmp where unary-eval op x = new-int b tmp ∧ 0 < b ∧ b ≤ 64
  apply (cases op; simp; auto simp: 2)
  apply (metis 2 Value.inject(1) Value.simps(5) assms(1) intval-narrow.simps(1)
intval-narrow-ok new-int.simps unary-eval.simps(5))
  apply (metis 2 Value.distinct(1) Value.inject(1) assms(1) bot-nat-0.not-eq-extremum
diff-is-0-eq intval-sign-extend.elims new-int.simps unary-eval.simps(6) zero-less-diff)
  by (smt (verit, del-insts) 2 Value.simps(5) assms(1) intval-bits.simps int-
val-zero-extend.simps(1) new-int.simps order-less-le-trans unary-eval.simps(7))
  then show ?thesis
  by blast
qed

```

```

lemma bin-eval-inputs-are-ints:
  assumes bin-eval op x y = IntVal b ix
  obtains xb yb xi yi where  $x = \text{IntVal } xb \ xi \wedge y = \text{IntVal } yb \ yi$ 
proof –
  have bin-eval op x y ≠ UndefVal
  by (simp add: assms)
  then show ?thesis
  using assms apply (cases op; cases x; cases y; simp)
  using that by blast+
qed

```

```

lemma eval-bits-1-64:
   $[m,p] \vdash xe \mapsto (\text{IntVal } b \text{ } ix) \implies 0 < b \wedge b \leq 64$ 
proof (induction xe arbitrary: b ix)
  case (UnaryExpr op x2)
  then obtain xv where
    xv:  $([m,p] \vdash x2 \mapsto xv) \wedge$ 
    IntVal b ix = unary-eval op xv
  using unfold-binary by auto
  then have b = (if op ∈ normal-unary then intval-bits xv else ir-resultBits op)
  using unary-eval-new-int
  by (metis Value.disc(1) Value.discI(1) Value.sel(1) new-int.simps)
  then show ?case
  by (metis xv UnaryExpr.IH unary-normal-bitsize unary-not-normal-bitsize)
next
  case (BinaryExpr op x y)
  then obtain xv yv where
    xy:  $([m,p] \vdash x \mapsto xv) \wedge$ 
     $([m,p] \vdash y \mapsto yv) \wedge$ 
    IntVal b ix = bin-eval op xv yv
  using unfold-binary by auto
  then have def: bin-eval op xv yv ≠ UndefVal and xv: xv ≠ UndefVal and yv ≠
  UndefVal
  using evaltree-not-undef xy by (force, blast, blast)
  then have b = (if op ∈ binary-fixed-32-ops then 32 else intval-bits xv)
  by (metis xy intval-bits.simps new-int.simps bin-eval-new-int)
  then show ?case
  by (metis BinaryExpr.IH(1) Value.distinct(7) Value.distinct(9) xv bin-eval-inputs-are-ints
  intval-bits.elims le-add-same-cancel1 less-or-eq-imp-le numeral-Bit0 xy zero-less-numeral)
next
  case (ConditionalExpr xe1 xe2 xe3)
  then show ?case
  by (metis (full-types) EvalTreeE(3))
next
  case (ParameterExpr x1 x2)
  then show ?case
  using ParameterExprE intval-bits.simps valid-stamp.simps(1) valid-value.elims(2)
  valid-value.simps(17)
  by (metis (no-types, lifting))
next
  case (LeafExpr x1 x2)
  then show ?case
  by (smt (z3) EvalTreeE(6) Value.distinct(7) Value.inject(1) valid-stamp.simps(1)
  valid-value.elims(1))
next
  case (ConstantExpr x)
  then show ?case using wf-value-def
  by (metis EvalTreeE(1) constantAsStamp.simps(1) valid-stamp.simps(1) valid-value.simps(1))
next
  case (ConstantVar x)

```



```

    then show ?case
      by blast
  next
    case (VariableExpr x1 x2)
    then show ?case
      by blast
  qed

lemma unfold-binary-width:
  assumes  $op \notin \text{binary-fixed-32-ops} \wedge op \notin \text{binary-shift-ops}$ 
  shows  $([m,p] \vdash \text{BinaryExpr } op \ xe \ ye \mapsto \text{IntVal } b \ val) = (\exists \ x \ y. \ ([m,p] \vdash \ xe \mapsto \text{IntVal } b \ x) \wedge \ ([m,p] \vdash \ ye \mapsto \text{IntVal } b \ y) \wedge \ (\text{IntVal } b \ val = \text{bin-eval } op \ (\text{IntVal } b \ x) \ (\text{IntVal } b \ y)) \wedge \ (\text{IntVal } b \ val \neq \text{UndefVal})) \text{ (is } ?L = ?R)$ 
proof (intro iffI)
  assume  $\exists: ?L$ 
  show  $?R$  apply (rule evaltree.cases[OF  $\exists$ ])
    apply force+ apply auto[1]
    using assms apply (cases op; auto)
    apply (smt (verit) intval-add.elims Value.inject(1))
    using intval-mul.elims Value.inject(1)
    apply (smt (verit) new-int.simps new-int-bin.simps)
    using intval-sub.elims Value.inject(1)
    apply (smt (verit) new-int.simps new-int-bin.simps)
    using intval-and.elims Value.inject(1)
    apply (smt (verit) new-int.simps new-int-bin.simps take-bit-and)
    using intval-or.elims Value.inject(1)
    apply (smt (verit) new-int.simps new-int-bin.simps take-bit-or)
    using intval-xor.elims Value.inject(1)
    apply (smt (verit) new-int.simps new-int-bin.simps take-bit-xor)
  by blast

next
  assume  $R: ?R$ 
  then obtain  $x \ y$  where  $[m,p] \vdash \ xe \mapsto \text{IntVal } b \ x$ 
    and  $[m,p] \vdash \ ye \mapsto \text{IntVal } b \ y$ 
    and  $\text{new-int } b \ val = \text{bin-eval } op \ (\text{IntVal } b \ x) \ (\text{IntVal } b \ y)$ 
    and  $\text{new-int } b \ val \neq \text{UndefVal}$ 
    using bin-eval-unused-bits-zero by force
  then show ?L
    using  $R$  by blast
  qed
end

```

7 Tree to Graph

```

theory TreeToGraph
  imports
    Semantics.IRTreeEval
    Graph.IRGraph
begin

```

7.1 Subgraph to Data-flow Tree

```

fun find-node-and-stamp :: IRGraph  $\Rightarrow$  (IRNode  $\times$  Stamp)  $\Rightarrow$  ID option where
  find-node-and-stamp g (n,s) =
    find ( $\lambda i.$  kind g i = n  $\wedge$  stamp g i = s) (sorted-list-of-set(ids g))

export-code find-node-and-stamp

```

```

fun is-preevaluated :: IRNode  $\Rightarrow$  bool where
  is-preevaluated (InvokeNode n - - - -) = True |
  is-preevaluated (InvokeWithExceptionNode n - - - -) = True |
  is-preevaluated (NewInstanceNode n - -) = True |
  is-preevaluated (LoadFieldNode n - -) = True |
  is-preevaluated (SignedDivNode n - - - -) = True |
  is-preevaluated (SignedRemNode n - - - -) = True |
  is-preevaluated (ValuePhiNode n -) = True |
  is-preevaluated - = False

```

```

inductive
  rep :: IRGraph  $\Rightarrow$  ID  $\Rightarrow$  IRExpr  $\Rightarrow$  bool (-  $\vdash$  -  $\simeq$  - 55)
  for g where

```

```

  ConstantNode:
  [[kind g n = ConstantNode c]
    $\implies$  g  $\vdash$  n  $\simeq$  (ConstantExpr c) |

```

```

  ParameterNode:
  [[kind g n = ParameterNode i;
   stamp g n = s]
    $\implies$  g  $\vdash$  n  $\simeq$  (ParameterExpr i s) |

```

```

  ConditionalNode:
  [[kind g n = ConditionalNode c t f;
   g  $\vdash$  c  $\simeq$  ce;
   g  $\vdash$  t  $\simeq$  te;
   g  $\vdash$  f  $\simeq$  fe]
    $\implies$  g  $\vdash$  n  $\simeq$  (ConditionalExpr ce te fe) |

```

```

  AbsNode:

```

$\llbracket \text{kind } g \ n = \text{AbsNode } x;$
 $g \vdash x \simeq xe \rrbracket$
 $\implies g \vdash n \simeq (\text{UnaryExpr } \text{UnaryAbs } xe) \mid$

NotNode:
 $\llbracket \text{kind } g \ n = \text{NotNode } x;$
 $g \vdash x \simeq xe \rrbracket$
 $\implies g \vdash n \simeq (\text{UnaryExpr } \text{UnaryNot } xe) \mid$

NegateNode:
 $\llbracket \text{kind } g \ n = \text{NegateNode } x;$
 $g \vdash x \simeq xe \rrbracket$
 $\implies g \vdash n \simeq (\text{UnaryExpr } \text{UnaryNeg } xe) \mid$

LogicNegationNode:
 $\llbracket \text{kind } g \ n = \text{LogicNegationNode } x;$
 $g \vdash x \simeq xe \rrbracket$
 $\implies g \vdash n \simeq (\text{UnaryExpr } \text{UnaryLogicNegation } xe) \mid$

AddNode:
 $\llbracket \text{kind } g \ n = \text{AddNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinAdd } xe \ ye) \mid$

MulNode:
 $\llbracket \text{kind } g \ n = \text{MulNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinMul } xe \ ye) \mid$

SubNode:
 $\llbracket \text{kind } g \ n = \text{SubNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinSub } xe \ ye) \mid$

AndNode:
 $\llbracket \text{kind } g \ n = \text{AndNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinAnd } xe \ ye) \mid$

OrNode:
 $\llbracket \text{kind } g \ n = \text{OrNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinOr } xe \ ye) \mid$

XorNode:

$\llbracket \text{kind } g \ n = \text{XorNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinXor } xe \ ye) \mid$

ShortCircuitOrNode:

$\llbracket \text{kind } g \ n = \text{ShortCircuitOrNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinShortCircuitOr } xe \ ye) \mid$

LeftShiftNode:

$\llbracket \text{kind } g \ n = \text{LeftShiftNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinLeftShift } xe \ ye) \mid$

RightShiftNode:

$\llbracket \text{kind } g \ n = \text{RightShiftNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinRightShift } xe \ ye) \mid$

UnsignedRightShiftNode:

$\llbracket \text{kind } g \ n = \text{UnsignedRightShiftNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinURightShift } xe \ ye) \mid$

IntegerBelowNode:

$\llbracket \text{kind } g \ n = \text{IntegerBelowNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinIntegerBelow } xe \ ye) \mid$

IntegerEqualsNode:

$\llbracket \text{kind } g \ n = \text{IntegerEqualsNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinIntegerEquals } xe \ ye) \mid$

IntegerLessThanNode:

$\llbracket \text{kind } g \ n = \text{IntegerLessThanNode } x \ y;$
 $g \vdash x \simeq xe;$
 $g \vdash y \simeq ye \rrbracket$
 $\implies g \vdash n \simeq (\text{BinaryExpr } \text{BinIntegerLessThan } xe \ ye) \mid$

NarrowNode:

$\llbracket \text{kind } g \ n = \text{NarrowNode } \text{inputBits } \text{resultBits } x; \\ g \vdash x \simeq xe \rrbracket \\ \implies g \vdash n \simeq (\text{UnaryExpr } (\text{UnaryNarrow } \text{inputBits } \text{resultBits}) \ xe) \mid$

SignExtendNode:

$\llbracket \text{kind } g \ n = \text{SignExtendNode } \text{inputBits } \text{resultBits } x; \\ g \vdash x \simeq xe \rrbracket \\ \implies g \vdash n \simeq (\text{UnaryExpr } (\text{UnarySignExtend } \text{inputBits } \text{resultBits}) \ xe) \mid$

ZeroExtendNode:

$\llbracket \text{kind } g \ n = \text{ZeroExtendNode } \text{inputBits } \text{resultBits } x; \\ g \vdash x \simeq xe \rrbracket \\ \implies g \vdash n \simeq (\text{UnaryExpr } (\text{UnaryZeroExtend } \text{inputBits } \text{resultBits}) \ xe) \mid$

LeafNode:

$\llbracket \text{is-preevaluated } (\text{kind } g \ n); \\ \text{stamp } g \ n = s \rrbracket \\ \implies g \vdash n \simeq (\text{LeafExpr } n \ s) \mid$

RefNode:

$\llbracket \text{kind } g \ n = \text{RefNode } n'; \\ g \vdash n' \simeq e \rrbracket \\ \implies g \vdash n \simeq e$

code-pred (*modes*: $i \Rightarrow i \Rightarrow o \Rightarrow \text{bool}$ as *exprE*) *rep* .

inductive

replist :: *IRGraph* \Rightarrow *ID* *list* \Rightarrow *IRExpr* *list* \Rightarrow *bool* ($- \vdash - \simeq_L -$ 55)
for *g* **where**

RepNil:

$g \vdash [] \simeq_L [] \mid$

RepCons:

$\llbracket g \vdash x \simeq xe; \\ g \vdash xs \simeq_L xse \rrbracket \\ \implies g \vdash x \# xs \simeq_L xe \# xse$

code-pred (*modes*: $i \Rightarrow i \Rightarrow o \Rightarrow \text{bool}$ as *exprListE*) *replist* .

definition *wf-term-graph* :: *MapState* \Rightarrow *Params* \Rightarrow *IRGraph* \Rightarrow *ID* \Rightarrow *bool* **where**

wf-term-graph *m* *p* *g* *n* = $(\exists \ e. (g \vdash n \simeq e) \wedge (\exists \ v. ([m, p] \vdash e \mapsto v)))$

values $\{t. \text{eg2-sq} \vdash 4 \simeq t\}$

7.2 Data-flow Tree to Subgraph

```
fun unary-node :: IRUnaryOp  $\Rightarrow$  ID  $\Rightarrow$  IRNode where
  unary-node UnaryAbs v = AbsNode v |
  unary-node UnaryNot v = NotNode v |
  unary-node UnaryNeg v = NegateNode v |
  unary-node UnaryLogicNegation v = LogicNegationNode v |
  unary-node (UnaryNarrow ib rb) v = NarrowNode ib rb v |
  unary-node (UnarySignExtend ib rb) v = SignExtendNode ib rb v |
  unary-node (UnaryZeroExtend ib rb) v = ZeroExtendNode ib rb v
```

```
fun bin-node :: IRBinaryOp  $\Rightarrow$  ID  $\Rightarrow$  ID  $\Rightarrow$  IRNode where
  bin-node BinAdd x y = AddNode x y |
  bin-node BinMul x y = MulNode x y |
  bin-node BinSub x y = SubNode x y |
  bin-node BinAnd x y = AndNode x y |
  bin-node BinOr x y = OrNode x y |
  bin-node BinXor x y = XorNode x y |
  bin-node BinShortCircuitOr x y = ShortCircuitOrNode x y |
  bin-node BinLeftShift x y = LeftShiftNode x y |
  bin-node BinRightShift x y = RightShiftNode x y |
  bin-node BinURightShift x y = UnsignedRightShiftNode x y |
  bin-node BinIntegerEquals x y = IntegerEqualsNode x y |
  bin-node BinIntegerLessThan x y = IntegerLessThanNode x y |
  bin-node BinIntegerBelow x y = IntegerBelowNode x y
```

```
inductive fresh-id :: IRGraph  $\Rightarrow$  ID  $\Rightarrow$  bool where
  n  $\notin$  ids g  $\Longrightarrow$  fresh-id g n
```

```
code-pred fresh-id .
```

```
fun get-fresh-id :: IRGraph  $\Rightarrow$  ID where
```

```
  get-fresh-id g = last(sorted-list-of-set(ids g)) + 1
```

```
export-code get-fresh-id
```

```
value get-fresh-id eg2-sq
```

```
value get-fresh-id (add-node 6 (ParameterNode 2, default-stamp) eg2-sq)
```

```
inductive
```

unrep :: *IRGraph* \Rightarrow *IRExpr* \Rightarrow (*IRGraph* \times *ID*) \Rightarrow *bool* (- \oplus - \rightsquigarrow - 55)
where

ConstantNodeSame:

$\llbracket \text{find-node-and-stamp } g \text{ (ConstantNode } c, \text{ constantAsStamp } c) = \text{Some } n \rrbracket$
 $\implies g \oplus (\text{ConstantExpr } c) \rightsquigarrow (g, n) \mid$

ConstantNodeNew:

$\llbracket \text{find-node-and-stamp } g \text{ (ConstantNode } c, \text{ constantAsStamp } c) = \text{None};$
 $n = \text{get-fresh-id } g;$
 $g' = \text{add-node } n \text{ (ConstantNode } c, \text{ constantAsStamp } c) \text{ } g \rrbracket$
 $\implies g \oplus (\text{ConstantExpr } c) \rightsquigarrow (g', n) \mid$

ParameterNodeSame:

$\llbracket \text{find-node-and-stamp } g \text{ (ParameterNode } i, s) = \text{Some } n \rrbracket$
 $\implies g \oplus (\text{ParameterExpr } i \text{ } s) \rightsquigarrow (g, n) \mid$

ParameterNodeNew:

$\llbracket \text{find-node-and-stamp } g \text{ (ParameterNode } i, s) = \text{None};$
 $n = \text{get-fresh-id } g;$
 $g' = \text{add-node } n \text{ (ParameterNode } i, s) \text{ } g \rrbracket$
 $\implies g \oplus (\text{ParameterExpr } i \text{ } s) \rightsquigarrow (g', n) \mid$

ConditionalNodeSame:

$\llbracket \text{find-node-and-stamp } g4 \text{ (ConditionalNode } c \text{ } t \text{ } f, s') = \text{Some } n;$
 $g \oplus ce \rightsquigarrow (g2, c);$
 $g2 \oplus te \rightsquigarrow (g3, t);$
 $g3 \oplus fe \rightsquigarrow (g4, f);$
 $s' = \text{meet (stamp } g4 \text{ } t) \text{ (stamp } g4 \text{ } f) \rrbracket$
 $\implies g \oplus (\text{ConditionalExpr } ce \text{ } te \text{ } fe) \rightsquigarrow (g4, n) \mid$

ConditionalNodeNew:

$\llbracket \text{find-node-and-stamp } g4 \text{ (ConditionalNode } c \text{ } t \text{ } f, s') = \text{None};$
 $g \oplus ce \rightsquigarrow (g2, c);$
 $g2 \oplus te \rightsquigarrow (g3, t);$
 $g3 \oplus fe \rightsquigarrow (g4, f);$
 $s' = \text{meet (stamp } g4 \text{ } t) \text{ (stamp } g4 \text{ } f);$
 $n = \text{get-fresh-id } g4;$
 $g' = \text{add-node } n \text{ (ConditionalNode } c \text{ } t \text{ } f, s') \text{ } g4 \rrbracket$
 $\implies g \oplus (\text{ConditionalExpr } ce \text{ } te \text{ } fe) \rightsquigarrow (g', n) \mid$

UnaryNodeSame:

$\llbracket \text{find-node-and-stamp } g2 \text{ (unary-node } op \text{ } x, s') = \text{Some } n;$
 $g \oplus xe \rightsquigarrow (g2, x);$
 $s' = \text{stamp-unary } op \text{ (stamp } g2 \text{ } x) \rrbracket$
 $\implies g \oplus (\text{UnaryExpr } op \text{ } xe) \rightsquigarrow (g2, n) \mid$

UnaryNodeNew:

$\llbracket \text{find-node-and-stamp } g2 \text{ (unary-node } op \text{ } x, s') = \text{None};$

$g \oplus xe \rightsquigarrow (g2, x);$
 $s' = \text{stamp-unary } op \ (\text{stamp } g2 \ x);$
 $n = \text{get-fresh-id } g2;$
 $g' = \text{add-node } n \ (\text{unary-node } op \ x, \ s') \ g2]$
 $\implies g \oplus (\text{UnaryExpr } op \ xe) \rightsquigarrow (g', n) \mid$

BinaryNodeSame:

$\llbracket \text{find-node-and-stamp } g3 \ (\text{bin-node } op \ x \ y, \ s') = \text{Some } n;$
 $g \oplus xe \rightsquigarrow (g2, x);$
 $g2 \oplus ye \rightsquigarrow (g3, y);$
 $s' = \text{stamp-binary } op \ (\text{stamp } g3 \ x) \ (\text{stamp } g3 \ y)]$
 $\implies g \oplus (\text{BinaryExpr } op \ xe \ ye) \rightsquigarrow (g3, n) \mid$

BinaryNodeNew:

$\llbracket \text{find-node-and-stamp } g3 \ (\text{bin-node } op \ x \ y, \ s') = \text{None};$
 $g \oplus xe \rightsquigarrow (g2, x);$
 $g2 \oplus ye \rightsquigarrow (g3, y);$
 $s' = \text{stamp-binary } op \ (\text{stamp } g3 \ x) \ (\text{stamp } g3 \ y);$
 $n = \text{get-fresh-id } g3;$
 $g' = \text{add-node } n \ (\text{bin-node } op \ x \ y, \ s') \ g3]$
 $\implies g \oplus (\text{BinaryExpr } op \ xe \ ye) \rightsquigarrow (g', n) \mid$

AllLeafNodes:

$\llbracket \text{stamp } g \ n = s;$
 $\text{is-preevaluated } (\text{kind } g \ n)]$
 $\implies g \oplus (\text{LeafExpr } n \ s) \rightsquigarrow (g, n)$

code-pred (*modes*: $i \Rightarrow i \Rightarrow o \Rightarrow \text{bool}$ as *unrepE*)
unrep .

unrepRules

$$\frac{\text{find-node-and-stamp } (g::\text{IRGraph}) \text{ (ConstantNode } (c::\text{Value}), \text{ constantAsStamp } c) = \text{Some } (n::\text{nat})}{g \oplus \text{ConstantExpr } c \rightsquigarrow (g, n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g::\text{IRGraph}) \text{ (ConstantNode } (c::\text{Value}), \text{ constantAsStamp } c) = \text{None} \\ (n::\text{nat}) = \text{get-fresh-id } g \\ (g'::\text{IRGraph}) = \text{add-node } n \text{ (ConstantNode } c, \text{ constantAsStamp } c) \end{array}}{g \oplus \text{ConstantExpr } c \rightsquigarrow (g', n)}$$

$$\frac{\text{find-node-and-stamp } (g::\text{IRGraph}) \text{ (ParameterNode } (i::\text{nat}), s::\text{Stamp}) = \text{Some } (n::\text{nat})}{g \oplus \text{ParameterExpr } i \ s \rightsquigarrow (g, n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g::\text{IRGraph}) \text{ (ParameterNode } (i::\text{nat}), s::\text{Stamp}) = \text{None} \\ (n::\text{nat}) = \text{get-fresh-id } g \\ (g'::\text{IRGraph}) = \text{add-node } n \text{ (ParameterNode } i, s) \end{array}}{g \oplus \text{ParameterExpr } i \ s \rightsquigarrow (g', n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g4::\text{IRGraph}) \text{ (ConditionalNode } (c::\text{nat}) \ (t::\text{nat}) \ (f::\text{nat}), s'::\text{Stamp}) = \text{Some } (n::\text{nat}) \\ g::\text{IRGraph} \oplus ce::\text{IExpr} \rightsquigarrow (g2::\text{IRGraph}, c) \\ g2 \oplus te::\text{IExpr} \rightsquigarrow (g3::\text{IRGraph}, t) \\ g3 \oplus fe::\text{IExpr} \rightsquigarrow (g4, f) \quad s' = \text{meet } (\text{stamp } g4 \ t) \ (\text{stamp } g4 \ f) \end{array}}{g \oplus \text{ConditionalExpr } ce \ te \ fe \rightsquigarrow (g4, n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g4::\text{IRGraph}) \text{ (ConditionalNode } (c::\text{nat}) \ (t::\text{nat}) \ (f::\text{nat}), s'::\text{Stamp}) = \text{None} \\ g::\text{IRGraph} \oplus ce::\text{IExpr} \rightsquigarrow (g2::\text{IRGraph}, c) \\ g2 \oplus te::\text{IExpr} \rightsquigarrow (g3::\text{IRGraph}, t) \quad g3 \oplus fe::\text{IExpr} \rightsquigarrow (g4, f) \\ s' = \text{meet } (\text{stamp } g4 \ t) \ (\text{stamp } g4 \ f) \quad (n::\text{nat}) = \text{get-fresh-id } g4 \\ (g'::\text{IRGraph}) = \text{add-node } n \text{ (ConditionalNode } c \ t \ f, s') \end{array}}{g \oplus \text{ConditionalExpr } ce \ te \ fe \rightsquigarrow (g', n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g3::\text{IRGraph}) \text{ (bin-node } (op::\text{IRBinaryOp}) \ (x::\text{nat}) \ (y::\text{nat}), s'::\text{Stamp}) = \text{Some } (n::\text{nat}) \\ g::\text{IRGraph} \oplus xe::\text{IExpr} \rightsquigarrow (g2::\text{IRGraph}, x) \\ g2 \oplus ye::\text{IExpr} \rightsquigarrow (g3, y) \\ s' = \text{stamp-binary } op \ (\text{stamp } g3 \ x) \ (\text{stamp } g3 \ y) \end{array}}{g \oplus \text{BinaryExpr } op \ xe \ ye \rightsquigarrow (g3, n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g3::\text{IRGraph}) \text{ (bin-node } (op::\text{IRBinaryOp}) \ (x::\text{nat}) \ (y::\text{nat}), s'::\text{Stamp}) = \text{None} \\ g::\text{IRGraph} \oplus xe::\text{IExpr} \rightsquigarrow (g2::\text{IRGraph}, x) \\ g2 \oplus ye::\text{IExpr} \rightsquigarrow (g3, y) \\ s' = \text{stamp-binary } op \ (\text{stamp } g3 \ x) \ (\text{stamp } g3 \ y) \\ (n::\text{nat}) = \text{get-fresh-id } g3 \\ (g'::\text{IRGraph}) = \text{add-node } n \text{ (bin-node } op \ x \ y, s') \end{array}}{g \oplus \text{BinaryExpr } op \ xe \ ye \rightsquigarrow (g', n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g2::\text{IRGraph}) \text{ (unary-node } (op::\text{IRUnaryOp}) \ (x::\text{nat}), s'::\text{Stamp}) = \text{Some } (n::\text{nat}) \\ g::\text{IRGraph} \oplus xe::\text{IExpr} \rightsquigarrow (g2, x) \\ s' = \text{stamp-unary } op \ (\text{stamp } g2 \ x) \end{array}}{g \oplus \text{UnaryExpr } op \ xe \rightsquigarrow (g2, n)}$$

$$\frac{\begin{array}{l} \text{find-node-and-stamp } (g2::\text{IRGraph}) \text{ (unary-node } (op::\text{IRUnaryOp}) \ (x::\text{nat}), s'::\text{Stamp}) = \text{None} \\ g::\text{IRGraph} \oplus xe::\text{IExpr} \rightsquigarrow (g2, x) \\ s' = \text{stamp-unary } op \ (\text{stamp } g2 \ x) \quad (n::\text{nat}) = \text{get-fresh-id } g2 \\ (g'::\text{IRGraph}) = \text{add-node } n \text{ (unary-node } op \ x, s') \end{array}}{g \oplus \text{UnaryExpr } op \ xe \rightsquigarrow (g', n)}$$

$$\frac{\text{stamp } (g::\text{IRGraph}) \ (n::\text{nat}) = (s::\text{Stamp}) \quad \text{is-preevaluated } (\text{kind } g \ n)}{g \oplus \text{LeafExpr } n \ s \rightsquigarrow (g, n)}$$

values $\{(n, g) . (eg2\text{-}sq \oplus sq\text{-}param0 \rightsquigarrow (g, n))\}$

7.3 Lift Data-flow Tree Semantics

definition *encodeeval* :: *IRGraph* \Rightarrow *MapState* \Rightarrow *Params* \Rightarrow *ID* \Rightarrow *Value* \Rightarrow *bool*
 $([\cdot, \cdot, \cdot] \vdash \cdot \mapsto \cdot \ 50)$
where
encodeeval *g m p n v* = $(\exists e. (g \vdash n \simeq e) \wedge ([m, p] \vdash e \mapsto v))$

7.4 Graph Refinement

definition *graph-represents-expression* :: *IRGraph* \Rightarrow *ID* \Rightarrow *IRExpr* \Rightarrow *bool*
 $(\vdash \cdot \preceq \cdot \ 50)$
where
 $(g \vdash n \preceq e) = (\exists e'. (g \vdash n \simeq e') \wedge (e' \leq e))$

definition *graph-refinement* :: *IRGraph* \Rightarrow *IRGraph* \Rightarrow *bool* **where**
graph-refinement *g1 g2* =
 $((ids\ g_1 \subseteq ids\ g_2) \wedge$
 $(\forall n. n \in ids\ g_1 \longrightarrow (\forall e. (g_1 \vdash n \simeq e) \longrightarrow (g_2 \vdash n \preceq e))))$

lemma *graph-refinement*:

graph-refinement *g1 g2* $\implies (\forall n\ m\ p\ v. n \in ids\ g1 \longrightarrow ([g1, m, p] \vdash n \mapsto v) \longrightarrow$
 $([g2, m, p] \vdash n \mapsto v))$
by (*meson encodeeval-def graph-refinement-def graph-represents-expression-def le-expr-def*)

7.5 Maximal Sharing

definition *maximal-sharing*:

maximal-sharing *g* = $(\forall n_1\ n_2. n_1 \in true\text{-}ids\ g \wedge n_2 \in true\text{-}ids\ g \longrightarrow$
 $(\forall e. (g \vdash n_1 \simeq e) \wedge (g \vdash n_2 \simeq e) \wedge (stamp\ g\ n_1 = stamp\ g\ n_2) \longrightarrow n_1 =$
 $n_2))$

end

7.6 Formedness Properties

theory *Form*

imports

Semantics.TreeToGraph

begin

definition *wf-start* **where**

wf-start *g* = $(0 \in ids\ g \wedge$
 $is\text{-}StartNode\ (kind\ g\ 0))$

definition *wf-closed* **where**

wf-closed *g* =
 $(\forall n \in ids\ g .$

$$\begin{aligned} &inputs\ g\ n \subseteq ids\ g \wedge \\ &succ\ g\ n \subseteq ids\ g \wedge \\ &kind\ g\ n \neq NoNode) \end{aligned}$$

definition *wf-phs* **where**

$$\begin{aligned} wf-phs\ g = & \\ &(\forall\ n \in ids\ g. \\ &\quad is-PhiNode\ (kind\ g\ n) \longrightarrow \\ &\quad length\ (ir-values\ (kind\ g\ n)) \\ &= length\ (ir-ends \\ &\quad (kind\ g\ (ir-merge\ (kind\ g\ n)))))) \end{aligned}$$

definition *wf-ends* **where**

$$\begin{aligned} wf-ends\ g = & \\ &(\forall\ n \in ids\ g . \\ &\quad is-AbstractEndNode\ (kind\ g\ n) \longrightarrow \\ &\quad card\ (usages\ g\ n) > 0) \end{aligned}$$

fun *wf-graph* :: *IRGraph* \Rightarrow *bool* **where**

$$wf-graph\ g = (wf-start\ g \wedge wf-closed\ g \wedge wf-phs\ g \wedge wf-ends\ g)$$

lemmas *wf-folds* =

$$\begin{aligned} &wf-graph.simps \\ &wf-start-def \\ &wf-closed-def \\ &wf-phs-def \\ &wf-ends-def \end{aligned}$$

fun *wf-stamps* :: *IRGraph* \Rightarrow *bool* **where**

$$\begin{aligned} wf-stamps\ g = &(\forall\ n \in ids\ g . \\ &(\forall\ v\ m\ p\ e . (g \vdash n \simeq e) \wedge ([m, p] \vdash e \mapsto v) \longrightarrow valid-value\ v\ (stamp-expr\ e))) \end{aligned}$$

fun *wf-stamp* :: *IRGraph* \Rightarrow (*ID* \Rightarrow *Stamp*) \Rightarrow *bool* **where**

$$\begin{aligned} wf-stamp\ g\ s = &(\forall\ n \in ids\ g . \\ &(\forall\ v\ m\ p\ e . (g \vdash n \simeq e) \wedge ([m, p] \vdash e \mapsto v) \longrightarrow valid-value\ v\ (s\ n))) \end{aligned}$$

lemma *wf-empty*: *wf-graph start-end-graph*

unfolding *start-end-graph-def wf-folds by simp*

lemma *wf-eg2-sq*: *wf-graph eg2-sq*

unfolding *eg2-sq-def wf-folds by simp*

fun *wf-logic-node-inputs* :: *IRGraph* \Rightarrow *ID* \Rightarrow *bool* **where**

$$\begin{aligned} wf-logic-node-inputs\ g\ n = & \\ &(\forall\ inp \in set\ (inputs-of\ (kind\ g\ n)) . (\forall\ v\ m\ p . ([g, m, p] \vdash inp \mapsto v) \longrightarrow wf-bool \\ &v)) \end{aligned}$$

fun *wf-values* :: *IRGraph* \Rightarrow *bool* **where**

$$wf-values\ g = (\forall\ n \in ids\ g .$$

$$\begin{aligned}
& (\forall v m p . ([g, m, p] \vdash n \mapsto v) \longrightarrow \\
& \quad (is-LogicNode (kind g n) \longrightarrow \\
& \quad \quad wf-bool v \wedge wf-logic-node-inputs g n)))
\end{aligned}$$

end

7.7 Dynamic Frames

This theory defines two operators, 'unchanged' and 'changeonly', that are useful for specifying which nodes in an IRGraph can change. The dynamic framing idea originates from 'Dynamic Frames' in software verification, started by Ioannis T. Kassios in "Dynamic frames: Support for framing, dependencies and sharing without restrictions", In FM 2006.

theory *IRGraphFrames*

imports

Form

begin

fun *unchanged* :: *ID set* \Rightarrow *IRGraph* \Rightarrow *IRGraph* \Rightarrow *bool* **where**

unchanged ns g1 g2 = $(\forall n . n \in ns \longrightarrow$
 $(n \in ids\ g1 \wedge n \in ids\ g2 \wedge kind\ g1\ n = kind\ g2\ n \wedge stamp\ g1\ n = stamp\ g2\ n))$

fun *changeonly* :: *ID set* \Rightarrow *IRGraph* \Rightarrow *IRGraph* \Rightarrow *bool* **where**

changeonly ns g1 g2 = $(\forall n . n \in ids\ g1 \wedge n \notin ns \longrightarrow$
 $(n \in ids\ g1 \wedge n \in ids\ g2 \wedge kind\ g1\ n = kind\ g2\ n \wedge stamp\ g1\ n = stamp\ g2\ n))$

lemma *node-unchanged*:

assumes *unchanged ns g1 g2*

assumes *nid* \in *ns*

shows *kind g1 nid* = *kind g2 nid*

using *assms* **by** *auto*

lemma *other-node-unchanged*:

assumes *changeonly ns g1 g2*

assumes *nid* \in *ids g1*

assumes *nid* \notin *ns*

shows *kind g1 nid* = *kind g2 nid*

using *assms*

using *changeonly.simps* **by** *blast*

Some notation for input nodes used

inductive *eval-uses*:: *IRGraph* \Rightarrow *ID* \Rightarrow *ID* \Rightarrow *bool*

for *g* **where**

use0: *nid* \in *ids g*

$\implies eval-uses\ g\ nid\ nid \mid$

```

use-inp:  $nid' \in inputs\ g\ n$ 
 $\implies eval\text{-}uses\ g\ nid\ nid' \mid$ 

use-trans:  $\llbracket eval\text{-}uses\ g\ nid\ nid';$ 
 $eval\text{-}uses\ g\ nid'\ nid'' \rrbracket$ 
 $\implies eval\text{-}uses\ g\ nid\ nid''$ 

fun eval-usages :: IRGraph  $\Rightarrow$  ID  $\Rightarrow$  ID set where
  eval-usages g nid = {n  $\in$  ids g . eval-uses g nid n}

lemma eval-usages-self:
  assumes nid  $\in$  ids g
  shows nid  $\in$  eval-usages g nid
  using assms eval-usages.simps eval-uses.intros(1)
  by (simp add: ids.rep-eq)

lemma not-in-g-inputs:
  assumes nid  $\notin$  ids g
  shows inputs g nid = {}
proof –
  have k: kind g nid = NoNode using assms not-in-g by blast
  then show ?thesis by (simp add: k)
qed

lemma child-member:
  assumes n = kind g nid
  assumes n  $\neq$  NoNode
  assumes List.member (inputs-of n) child
  shows child  $\in$  inputs g nid
  unfolding inputs.simps using assms
  by (metis in-set-member)

lemma child-member-in:
  assumes nid  $\in$  ids g
  assumes List.member (inputs-of (kind g nid)) child
  shows child  $\in$  inputs g nid
  unfolding inputs.simps using assms
  by (metis child-member ids-some inputs.elims)

lemma inp-in-g:
  assumes n  $\in$  inputs g nid
  shows nid  $\in$  ids g
proof –
  have inputs g nid  $\neq$  {}
  using assms
  by (metis empty-iff empty-set)

```

```

then have kind g nid  $\neq$  NoNode
  using not-in-g-inputs
  using ids-some by blast
then show ?thesis
  using not-in-g
  by metis
qed

```

```

lemma inp-in-g-wf:
  assumes wf-graph g
  assumes  $n \in \text{inputs } g \text{ nid}$ 
  shows  $n \in \text{ids } g$ 
  using assms unfolding wf-folds
  using inp-in-g by blast

```

```

lemma kind-unchanged:
  assumes  $\text{nid} \in \text{ids } g1$ 
  assumes unchanged (eval-usages g1 nid) g1 g2
  shows  $\text{kind } g1 \text{ nid} = \text{kind } g2 \text{ nid}$ 
proof –
  show ?thesis
    using assms eval-usages-self
    using unchanged.simps by blast
qed

```

```

lemma stamp-unchanged:
  assumes  $\text{nid} \in \text{ids } g1$ 
  assumes unchanged (eval-usages g1 nid) g1 g2
  shows  $\text{stamp } g1 \text{ nid} = \text{stamp } g2 \text{ nid}$ 
  by (meson assms(1) assms(2) eval-usages-self unchanged.elims(2))

```

```

lemma child-unchanged:
  assumes  $\text{child} \in \text{inputs } g1 \text{ nid}$ 
  assumes unchanged (eval-usages g1 nid) g1 g2
  shows unchanged (eval-usages g1 child) g1 g2
  by (smt assms(1) assms(2) eval-usages.simps mem-Collect-eq
    unchanged.simps use-inp use-trans)

```

```

lemma eval-usages:
  assumes  $us = \text{eval-usages } g \text{ nid}$ 
  assumes  $\text{nid}' \in \text{ids } g$ 
  shows  $\text{eval-uses } g \text{ nid nid}' \longleftrightarrow \text{nid}' \in us$  (is ?P  $\longleftrightarrow$  ?Q)
  using assms eval-usages.simps
  by (simp add: ids.rep-eq)

```

```

lemma inputs-are-uses:
  assumes  $\text{nid}' \in \text{inputs } g \text{ nid}$ 

```

```

shows eval-uses g nid nid'
by (metis assms use-inp)

lemma inputs-are-usages:
  assumes nid' ∈ inputs g nid
  assumes nid' ∈ ids g
  shows nid' ∈ eval-usages g nid
  using assms(1) assms(2) eval-usages inputs-are-uses by blast

lemma inputs-of-are-usages:
  assumes List.member (inputs-of (kind g nid)) nid'
  assumes nid' ∈ ids g
  shows nid' ∈ eval-usages g nid
  by (metis assms(1) assms(2) in-set-member inputs.elims inputs-are-usages)

lemma usage-includes-inputs:
  assumes us = eval-usages g nid
  assumes ls = inputs g nid
  assumes ls ⊆ ids g
  shows ls ⊆ us
  using inputs-are-usages eval-usages
  using assms(1) assms(2) assms(3) by blast

lemma elim-inp-set:
  assumes k = kind g nid
  assumes k ≠ NoNode
  assumes child ∈ set (inputs-of k)
  shows child ∈ inputs g nid
  using assms by auto

lemma encode-in-ids:
  assumes g ⊢ nid ≃ e
  shows nid ∈ ids g
  using assms
  apply (induction rule: rep.induct)
  apply simp+
  by fastforce+

lemma eval-in-ids:
  assumes [g, m, p] ⊢ nid ↦ v
  shows nid ∈ ids g
  using assms using encodeeval-def encode-in-ids
  by auto

lemma transitive-kind-same:
  assumes unchanged (eval-usages g1 nid) g1 g2
  shows  $\forall \text{nid}' \in (\text{eval-usages } g1 \text{ nid}) . \text{kind } g1 \text{ nid}' = \text{kind } g2 \text{ nid}'$ 
  using assms
  by (meson unchanged.elims(1))

```

```

theorem stay-same-encoding:
  assumes nc: unchanged (eval-usages g1 nid) g1 g2
  assumes g1: g1  $\vdash$  nid  $\simeq$  e
  assumes wf: wf-graph g1
  shows g2  $\vdash$  nid  $\simeq$  e
proof –
  have dom: nid  $\in$  ids g1
  using g1 encode-in-ids by simp
  show ?thesis
using g1 nc wf dom proof (induction e rule: rep.induct)
  case (ConstantNode n c)
  then have kind g2 n = ConstantNode c
  using dom nc kind-unchanged
  by metis
  then show ?case using rep.ConstantNode
  by presburger
next
  case (ParameterNode n i s)
  then have kind g2 n = ParameterNode i
  by (metis kind-unchanged)
  then show ?case
  by (metis ParameterNode.hyps(2) ParameterNode.premis(1) ParameterNode.premis(3)
  rep.ParameterNode stamp-unchanged)
next
  case (ConditionalNode n c t f ce te fe)
  then have kind g2 n = ConditionalNode c t f
  by (metis kind-unchanged)
  have c  $\in$  eval-usages g1 n  $\wedge$  t  $\in$  eval-usages g1 n  $\wedge$  f  $\in$  eval-usages g1 n
  using inputs-of-ConditionalNode
  by (metis ConditionalNode.hyps(1) ConditionalNode.hyps(2) ConditionalNode.hyps(3)
  ConditionalNode.hyps(4) encode-in-ids inputs.simps inputs-are-usages
  list.set-intros(1) set-subset-Cons subset-code(1))
  then show ?case using transitive-kind-same
  by (metis ConditionalNode.hyps(1) ConditionalNode.premis(1) IRNodes.inputs-of-ConditionalNode
   $\langle$ kind g2 n = ConditionalNode c t f $\rangle$  child-unchanged inputs.simps list.set-intros(1)
  local.ConditionalNode(5) local.ConditionalNode(6) local.ConditionalNode(7) local.ConditionalNode(9)
  rep.ConditionalNode set-subset-Cons subset-code(1) unchanged.elims(2))
next
  case (AbsNode n x xe)
  then have kind g2 n = AbsNode x
  using kind-unchanged
  by metis
  then have x  $\in$  eval-usages g1 n
  using inputs-of-AbsNode
  by (metis AbsNode.hyps(1) AbsNode.hyps(2) encode-in-ids inputs.simps inputs-are-usages
  list.set-intros(1))
  then show ?case
  by (metis AbsNode.IH AbsNode.hyps(1) AbsNode.premis(1) AbsNode.premis(3))

```



```

IRNodes.inputs-of-AbsNode ⟨kind g2 n = AbsNode x⟩ child-member-in child-unchanged
local.wf member-rec(1) rep.AbsNode unchanged.simps)
next
  case (NotNode n x xe)
  then have kind g2 n = NotNode x
    using kind-unchanged
    by metis
  then have x ∈ eval-usages g1 n
    using inputs-of-NotNode
    by (metis NotNode.hyps(1) NotNode.hyps(2) encode-in-ids inputs.simps in-
      puts-are-usages list.set-intros(1))
  then show ?case
    by (metis NotNode.IH NotNode.hyps(1) NotNode.prem(1) NotNode.prem(3)
      IRNodes.inputs-of-NotNode ⟨kind g2 n = NotNode x⟩ child-member-in child-unchanged
      local.wf member-rec(1) rep.NotNode unchanged.simps)
next
  case (NegateNode n x xe)
  then have kind g2 n = NegateNode x
    using kind-unchanged by metis
  then have x ∈ eval-usages g1 n
    using inputs-of-NegateNode
    by (metis NegateNode.hyps(1) NegateNode.hyps(2) encode-in-ids inputs.simps
      inputs-are-usages list.set-intros(1))
  then show ?case
    by (metis IRNodes.inputs-of-NegateNode NegateNode.IH NegateNode.hyps(1)
      NegateNode.prem(1) NegateNode.prem(3) ⟨kind g2 n = NegateNode x⟩ child-member-in
      child-unchanged local.wf member-rec(1) rep.NegateNode unchanged.elims(1))
next
  case (LogicNegationNode n x xe)
  then have kind g2 n = LogicNegationNode x
    using kind-unchanged by metis
  then have x ∈ eval-usages g1 n
    using inputs-of-LogicNegationNode inputs-of-are-usages
    by (metis LogicNegationNode.hyps(1) LogicNegationNode.hyps(2) encode-in-ids
      member-rec(1))
  then show ?case
    by (metis IRNodes.inputs-of-LogicNegationNode LogicNegationNode.IH Logic-
      NegationNode.hyps(1) LogicNegationNode.hyps(2) LogicNegationNode.prem(1) ⟨kind
      g2 n = LogicNegationNode x⟩ child-unchanged encode-in-ids inputs.simps list.set-intros(1)
      local.wf rep.LogicNegationNode)
next
  case (AddNode n x y xe ye)
  then have kind g2 n = AddNode x y
    using kind-unchanged by metis
  then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
    using inputs-of-LogicNegationNode inputs-of-are-usages
    by (metis AddNode.hyps(1) AddNode.hyps(2) AddNode.hyps(3) IRNodes.inputs-of-AddNode
      encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case

```

```

    by (metis AddNode.IH(1) AddNode.IH(2) AddNode.hyps(1) AddNode.hyps(2)
AddNode.hyps(3) AddNode.premis(1) IRNodes.inputs-of-AddNode ⟨kind g2 n = AddNode
x y⟩ child-unchanged encode-in-ids in-set-member inputs.simps local.wf member-rec(1)
rep.AddNode)
next
  case (MulNode n x y xe ye)
  then have kind g2 n = MulNode x y
    using kind-unchanged by metis
  then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
    using inputs-of-LogicNegationNode inputs-of-are-usages
  by (metis MulNode.hyps(1) MulNode.hyps(2) MulNode.hyps(3) IRNodes.inputs-of-MulNode
encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using MulNode inputs-of-MulNode
  by (metis ⟨kind g2 n = MulNode x y⟩ child-unchanged inputs.simps list.set-intros(1)
rep.MulNode set-subset-Cons subset-iff unchanged.elims(2))
next
  case (SubNode n x y xe ye)
  then have kind g2 n = SubNode x y
    using kind-unchanged by metis
  then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
    using inputs-of-LogicNegationNode inputs-of-are-usages
  by (metis SubNode.hyps(1) SubNode.hyps(2) SubNode.hyps(3) IRNodes.inputs-of-SubNode
encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using SubNode inputs-of-SubNode
  by (metis ⟨kind g2 n = SubNode x y⟩ child-member child-unchanged encode-in-ids
ids-some member-rec(1) rep.SubNode)
next
  case (AndNode n x y xe ye)
  then have kind g2 n = AndNode x y
    using kind-unchanged by metis
  then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
    using inputs-of-LogicNegationNode inputs-of-are-usages
  by (metis AndNode.hyps(1) AndNode.hyps(2) AndNode.hyps(3) IRNodes.inputs-of-AndNode
encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using AndNode inputs-of-AndNode
  by (metis ⟨kind g2 n = AndNode x y⟩ child-unchanged inputs.simps list.set-intros(1)
rep.AndNode set-subset-Cons subset-iff unchanged.elims(2))
next
  case (OrNode n x y xe ye)
  then have kind g2 n = OrNode x y
    using kind-unchanged by metis
  then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
    using inputs-of-OrNode inputs-of-are-usages
  by (metis OrNode.hyps(1) OrNode.hyps(2) OrNode.hyps(3) IRNodes.inputs-of-OrNode
encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using OrNode inputs-of-OrNode
  by (metis ⟨kind g2 n = OrNode x y⟩ child-member child-unchanged encode-in-ids
ids-some member-rec(1) rep.OrNode)
next

```

```

case (XorNode n x y xe ye)
then have kind g2 n = XorNode x y
  using kind-unchanged by metis
then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
  using inputs-of-XorNode inputs-of-are-usages
  by (metis XorNode.hyps(1) XorNode.hyps(2) XorNode.hyps(3) IRNodes.inputs-of-XorNode
encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using XorNode inputs-of-XorNode
  by (metis ⟨kind g2 n = XorNode x y⟩ child-member child-unchanged encode-in-ids
ids-some member-rec(1) rep.XorNode)
next
case (ShortCircuitOrNode n x y xe ye)
then have kind g2 n = ShortCircuitOrNode x y
  using kind-unchanged by metis
then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
  using inputs-of-XorNode inputs-of-are-usages
  by (metis ShortCircuitOrNode.hyps(1) ShortCircuitOrNode.hyps(2) ShortCir-
cuitOrNode.hyps(3) IRNodes.inputs-of-ShortCircuitOrNode encode-in-ids in-mono
inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using ShortCircuitOrNode inputs-of-ShortCircuitOrNode
  by (metis ⟨kind g2 n = ShortCircuitOrNode x y⟩ child-member child-unchanged
encode-in-ids ids-some member-rec(1) rep.ShortCircuitOrNode)
next
case (LeftShiftNode n x y xe ye)
then have kind g2 n = LeftShiftNode x y
  using kind-unchanged by metis
then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
  using inputs-of-XorNode inputs-of-are-usages
  by (metis LeftShiftNode.hyps(1) LeftShiftNode.hyps(2) LeftShiftNode.hyps(3)
IRNodes.inputs-of-LeftShiftNode encode-in-ids in-mono inputs.simps inputs-are-usages
list.set-intros(1) set-subset-Cons)
  then show ?case using LeftShiftNode inputs-of-LeftShiftNode
  by (metis ⟨kind g2 n = LeftShiftNode x y⟩ child-member child-unchanged
encode-in-ids ids-some member-rec(1) rep.LeftShiftNode)
next
case (RightShiftNode n x y xe ye)
then have kind g2 n = RightShiftNode x y
  using kind-unchanged by metis
then have x ∈ eval-usages g1 n ∧ y ∈ eval-usages g1 n
  using inputs-of-RightShiftNode inputs-of-are-usages
  by (metis RightShiftNode.hyps(1) RightShiftNode.hyps(2) RightShiftNode.hyps(3)
IRNodes.inputs-of-RightShiftNode encode-in-ids in-mono inputs.simps inputs-are-usages
list.set-intros(1) set-subset-Cons)
  then show ?case using RightShiftNode inputs-of-RightShiftNode
  by (metis ⟨kind g2 n = RightShiftNode x y⟩ child-member child-unchanged
encode-in-ids ids-some member-rec(1) rep.RightShiftNode)
next
case (UnsignedRightShiftNode n x y xe ye)
then have kind g2 n = UnsignedRightShiftNode x y

```

```

    using kind-unchanged by metis
  then have  $x \in \text{eval-usages } g1 \ n \wedge y \in \text{eval-usages } g1 \ n$ 
    using inputs-of-UnsignedRightShiftNode inputs-of-are-usages
    by (metis UnsignedRightShiftNode.hyps(1) UnsignedRightShiftNode.hyps(2) UnsignedRightShiftNode.hyps(3) IRNodes.inputs-of-UnsignedRightShiftNode encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using UnsignedRightShiftNode inputs-of-UnsignedRightShiftNode
    by (metis <kind  $g2 \ n = \text{UnsignedRightShiftNode } x \ y$ > child-member child-unchanged encode-in-ids ids-some member-rec(1) rep.UnsignedRightShiftNode)
next
  case (IntegerBelowNode  $n \ x \ y \ x e \ y e$ )
  then have kind  $g2 \ n = \text{IntegerBelowNode } x \ y$ 
    using kind-unchanged by metis
  then have  $x \in \text{eval-usages } g1 \ n \wedge y \in \text{eval-usages } g1 \ n$ 
    using inputs-of-IntegerBelowNode inputs-of-are-usages
    by (metis IntegerBelowNode.hyps(1) IntegerBelowNode.hyps(2) IntegerBelowNode.hyps(3) IRNodes.inputs-of-IntegerBelowNode encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using IntegerBelowNode inputs-of-IntegerBelowNode
    by (metis <kind  $g2 \ n = \text{IntegerBelowNode } x \ y$ > child-member child-unchanged encode-in-ids ids-some member-rec(1) rep.IntegerBelowNode)
next
  case (IntegerEqualsNode  $n \ x \ y \ x e \ y e$ )
  then have kind  $g2 \ n = \text{IntegerEqualsNode } x \ y$ 
    using kind-unchanged by metis
  then have  $x \in \text{eval-usages } g1 \ n \wedge y \in \text{eval-usages } g1 \ n$ 
    using inputs-of-IntegerEqualsNode inputs-of-are-usages
    by (metis IntegerEqualsNode.hyps(1) IntegerEqualsNode.hyps(2) IntegerEqualsNode.hyps(3) IRNodes.inputs-of-IntegerEqualsNode encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using IntegerEqualsNode inputs-of-IntegerEqualsNode
    by (metis <kind  $g2 \ n = \text{IntegerEqualsNode } x \ y$ > child-member child-unchanged encode-in-ids ids-some member-rec(1) rep.IntegerEqualsNode)
next
  case (IntegerLessThanNode  $n \ x \ y \ x e \ y e$ )
  then have kind  $g2 \ n = \text{IntegerLessThanNode } x \ y$ 
    using kind-unchanged by metis
  then have  $x \in \text{eval-usages } g1 \ n \wedge y \in \text{eval-usages } g1 \ n$ 
    using inputs-of-IntegerLessThanNode inputs-of-are-usages
    by (metis IntegerLessThanNode.hyps(1) IntegerLessThanNode.hyps(2) IntegerLessThanNode.hyps(3) IRNodes.inputs-of-IntegerLessThanNode encode-in-ids in-mono inputs.simps inputs-are-usages list.set-intros(1) set-subset-Cons)
  then show ?case using IntegerLessThanNode inputs-of-IntegerLessThanNode
    by (metis <kind  $g2 \ n = \text{IntegerLessThanNode } x \ y$ > child-member child-unchanged encode-in-ids ids-some member-rec(1) rep.IntegerLessThanNode)
next
  case (NarrowNode  $n \ ib \ rb \ x \ x e$ )
  then have kind  $g2 \ n = \text{NarrowNode } ib \ rb \ x$ 
    using kind-unchanged by metis

```

```

then have  $x \in \text{eval-usages } g1 \ n$ 
  using inputs-of-NarrowNode inputs-of-are-usages
  by (metis NarrowNode.hyps(1) NarrowNode.hyps(2) IRNodes.inputs-of-NarrowNode
encode-in-ids inputs.simps inputs-are-usages list.set-intros(1))
  then show ?case using NarrowNode inputs-of-NarrowNode
    by (metis  $\langle \text{kind } g2 \ n = \text{NarrowNode } ib \ rb \ x \rangle$  child-unchanged inputs.elims
list.set-intros(1) rep.NarrowNode unchanged.simps)
next
  case (SignExtendNode n ib rb x xe)
  then have  $\text{kind } g2 \ n = \text{SignExtendNode } ib \ rb \ x$ 
    using kind-unchanged by metis
  then have  $x \in \text{eval-usages } g1 \ n$ 
    using inputs-of-SignExtendNode inputs-of-are-usages
    by (metis SignExtendNode.hyps(1) SignExtendNode.hyps(2) encode-in-ids in-
puts.simps inputs-are-usages list.set-intros(1))
    then show ?case using SignExtendNode inputs-of-SignExtendNode
      by (metis  $\langle \text{kind } g2 \ n = \text{SignExtendNode } ib \ rb \ x \rangle$  child-member-in child-unchanged
in-set-member list.set-intros(1) rep.SignExtendNode unchanged.elims(2))
  next
    case (ZeroExtendNode n ib rb x xe)
    then have  $\text{kind } g2 \ n = \text{ZeroExtendNode } ib \ rb \ x$ 
      using kind-unchanged by metis
    then have  $x \in \text{eval-usages } g1 \ n$ 
      using inputs-of-ZeroExtendNode inputs-of-are-usages
      by (metis ZeroExtendNode.hyps(1) ZeroExtendNode.hyps(2) IRNodes.inputs-of-ZeroExtendNode
encode-in-ids inputs.simps inputs-are-usages list.set-intros(1))
      then show ?case using ZeroExtendNode inputs-of-ZeroExtendNode
        by (metis  $\langle \text{kind } g2 \ n = \text{ZeroExtendNode } ib \ rb \ x \rangle$  child-member-in child-unchanged
member-rec(1) rep.ZeroExtendNode unchanged.simps)
    next
      case (LeafNode n s)
      then show ?case
        by (metis kind-unchanged rep.LeafNode stamp-unchanged)
    next
      case (RefNode n n')
      then have  $\text{kind } g2 \ n = \text{RefNode } n'$ 
        using kind-unchanged by metis
      then have  $n' \in \text{eval-usages } g1 \ n$ 
        by (metis IRNodes.inputs-of-RefNode RefNode.hyps(1) RefNode.hyps(2) en-
code-in-ids inputs.elims inputs-are-usages list.set-intros(1))
        then show ?case
          by (metis IRNodes.inputs-of-RefNode RefNode.IH RefNode.hyps(1) RefNode.hyps(2)
RefNode.prem(1)  $\langle \text{kind } g2 \ n = \text{RefNode } n' \rangle$  child-unchanged encode-in-ids in-
puts.elims list.set-intros(1) local.wf rep.RefNode)
qed
qed

```

```

theorem stay-same:
  assumes nc: unchanged (eval-usages g1 nid) g1 g2
  assumes g1:  $[g1, m, p] \vdash \text{id} \mapsto v1$ 
  assumes wf: wf-graph g1
  shows  $[g2, m, p] \vdash \text{id} \mapsto v1$ 
proof –
  have nid:  $\text{id} \in \text{ids } g1$ 
    using g1 eval-in-ids by simp
  then have  $\text{id} \in \text{eval-usages } g1 \text{ nid}$ 
    using eval-usages-self by blast
  then have kind-same:  $\text{kind } g1 \text{ nid} = \text{kind } g2 \text{ nid}$ 
    using nc node-unchanged by blast
  obtain e where  $e: (g1 \vdash \text{id} \simeq e) \wedge ([m,p] \vdash e \mapsto v1)$ 
    using encodeeval-def g1
    by auto
  then have val:  $[m,p] \vdash e \mapsto v1$ 
    using g1 encodeeval-def
    by simp
  then show ?thesis using e nid nc
    unfolding encodeeval-def
proof (induct e v1 arbitrary: nid rule: evaltree.induct)
  case (ConstantExpr c)
    then show ?case
      by (meson local.wf stay-same-encoding)
  next
    case (ParameterExpr i s)
    have  $g2 \vdash \text{id} \simeq \text{ParameterExpr } i \text{ s}$ 
      using stay-same-encoding ParameterExpr
      by (meson local.wf)
    then show ?case using evaltree.ParameterExpr
      by (meson ParameterExpr.hyps)
  next
    case (ConditionalExpr ce cond branch te fe v)
    then have  $g2 \vdash \text{id} \simeq \text{ConditionalExpr } ce \text{ te } fe$ 
      using ConditionalExpr.premis(1) ConditionalExpr.premis(3) local.wf stay-same-encoding
      by presburger
    then show ?case
      by (meson ConditionalExpr.premis(1) ConditionalExpr.premis(3) local.wf
stay-same-encoding)
  next
    case (UnaryExpr xe v op)
    then show ?case
      using local.wf stay-same-encoding by blast
  next
    case (BinaryExpr xe x ye y op)
    then show ?case
      using local.wf stay-same-encoding by blast
  next
    case (LeafExpr val nid s)

```

```

    then show ?case
    by (metis local.wf stay-same-encoding)
qed
qed

```

```

lemma add-changed:
  assumes gup = add-node new k g
  shows changeonly {new} g gup
  using assms unfolding add-node-def changeonly.simps
  using add-node.rep-eq add-node-def kind.rep-eq stamp.rep-eq by simp

```

```

lemma disjoint-change:
  assumes changeonly change g gup
  assumes nochange = ids g - change
  shows unchanged nochange g gup
  using assms unfolding changeonly.simps unchanged.simps
  by blast

```

```

lemma add-node-unchanged:
  assumes new  $\notin$  ids g
  assumes nid  $\in$  ids g
  assumes gup = add-node new k g
  assumes wf-graph g
  shows unchanged (eval-usages g nid) g gup
proof -
  have new  $\notin$  (eval-usages g nid) using assms
  using eval-usages.simps by blast
  then have changeonly {new} g gup
  using assms add-changed by blast
  then show ?thesis using assms add-node-def disjoint-change
  using Diff-insert-absorb by auto
qed

```

```

lemma eval-uses-imp:
  ((nid'  $\in$  ids g  $\wedge$  nid = nid')
   $\vee$  nid'  $\in$  inputs g nid
   $\vee$  ( $\exists$  nid'' . eval-uses g nid nid''  $\wedge$  eval-uses g nid'' nid'))
 $\longleftrightarrow$  eval-uses g nid nid'
  using use0 use-inp use-trans
  by (meson eval-uses.simps)

```

```

lemma wf-use-ids:
  assumes wf-graph g
  assumes nid  $\in$  ids g
  assumes eval-uses g nid nid'
  shows nid'  $\in$  ids g
  using assms(3)
proof (induction rule: eval-uses.induct)

```

```

    case use0
    then show ?case by simp
next
    case use-inp
    then show ?case
        using assms(1) inp-in-g-wf by blast
next
    case use-trans
    then show ?case by blast
qed

lemma no-external-use:
  assumes wf-graph g
  assumes nid'  $\notin$  ids g
  assumes nid  $\in$  ids g
  shows  $\neg$ (eval-uses g nid nid')
proof -
  have 0: nid  $\neq$  nid'
    using assms by blast
  have inp: nid'  $\notin$  inputs g nid
    using assms
    using inp-in-g-wf by blast
  have rec-0:  $\nexists n . n \in$  ids g  $\wedge$  n = nid'
    using assms by blast
  have rec-inp:  $\nexists n . n \in$  ids g  $\wedge$  n  $\in$  inputs g nid'
    using assms(2) inp-in-g by blast
  have rec:  $\nexists$  nid'' . eval-uses g nid nid''  $\wedge$  eval-uses g nid'' nid'
    using wf-use-ids assms(1) assms(2) assms(3) by blast
  from inp 0 rec show ?thesis
    using eval-uses-imp by blast
qed

end

```

7.8 Tree to Graph Theorems

```

theory TreeToGraphThms
imports
  IRTreeEvalThms
  IRGraphFrames
  HOL-Eisbach.Eisbach
  HOL-Eisbach.Eisbach-Tools
begin

```

7.8.1 Extraction and Evaluation of Expression Trees is Deterministic.

First, we prove some extra rules that relate each type of IRNode to the corresponding IRExpr type that 'rep' will produce. These are very helpful

for proving that 'rep' is deterministic.

named-theorems *rep*

lemma *rep-constant* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = ConstantNode\ c \implies$
 $e = ConstantExpr\ c$
by (*induction rule: rep.induct; auto*)

lemma *rep-parameter* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = ParameterNode\ i \implies$
 $(\exists\ s.\ e = ParameterExpr\ i\ s)$
by (*induction rule: rep.induct; auto*)

lemma *rep-conditional* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = ConditionalNode\ c\ t\ f \implies$
 $(\exists\ ce\ te\ fe.\ e = ConditionalExpr\ ce\ te\ fe)$
by (*induction rule: rep.induct; auto*)

lemma *rep-abs* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = AbsNode\ x \implies$
 $(\exists\ xe.\ e = UnaryExpr\ UnaryAbs\ xe)$
by (*induction rule: rep.induct; auto*)

lemma *rep-not* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = NotNode\ x \implies$
 $(\exists\ xe.\ e = UnaryExpr\ UnaryNot\ xe)$
by (*induction rule: rep.induct; auto*)

lemma *rep-negate* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = NegateNode\ x \implies$
 $(\exists\ xe.\ e = UnaryExpr\ UnaryNeg\ xe)$
by (*induction rule: rep.induct; auto*)

lemma *rep-logicnegation* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = LogicNegationNode\ x \implies$
 $(\exists\ xe.\ e = UnaryExpr\ UnaryLogicNegation\ xe)$
by (*induction rule: rep.induct; auto*)

lemma *rep-add* [*rep*]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = AddNode\ x\ y \implies$
 $(\exists\ xe\ ye.\ e = BinaryExpr\ BinAdd\ xe\ ye)$

by (*induction rule: rep.induct; auto*)

lemma *rep-sub* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = SubNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinSub\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-mul* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = MulNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinMul\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-and* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = AndNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinAnd\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-or* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = OrNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinOr\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-xor* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = XorNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinXor\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-short-circuit-or* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = ShortCircuitOrNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinShortCircuitOr\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-left-shift* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = LeftShiftNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinLeftShift\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-right-shift* [*rep*]:
 $g \vdash n \simeq e \implies$
 $kind\ g\ n = RightShiftNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinRightShift\ xe\ ye)$
by (*induction rule: rep.induct; auto*)

lemma *rep-unsigned-right-shift* [rep]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = UnsignedRightShiftNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinURightShift\ xe\ ye)$
by (induction rule: *rep.induct*; auto)

lemma *rep-integer-below* [rep]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = IntegerBelowNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinIntegerBelow\ xe\ ye)$
by (induction rule: *rep.induct*; auto)

lemma *rep-integer-equals* [rep]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = IntegerEqualsNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinIntegerEquals\ xe\ ye)$
by (induction rule: *rep.induct*; auto)

lemma *rep-integer-less-than* [rep]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = IntegerLessThanNode\ x\ y \implies$
 $(\exists\ xe\ ye. e = BinaryExpr\ BinIntegerLessThan\ xe\ ye)$
by (induction rule: *rep.induct*; auto)

lemma *rep-narrow* [rep]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = NarrowNode\ ib\ rb\ x \implies$
 $(\exists\ x. e = UnaryExpr\ (UnaryNarrow\ ib\ rb)\ x)$
by (induction rule: *rep.induct*; auto)

lemma *rep-sign-extend* [rep]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = SignExtendNode\ ib\ rb\ x \implies$
 $(\exists\ x. e = UnaryExpr\ (UnarySignExtend\ ib\ rb)\ x)$
by (induction rule: *rep.induct*; auto)

lemma *rep-zero-extend* [rep]:

$g \vdash n \simeq e \implies$
 $kind\ g\ n = ZeroExtendNode\ ib\ rb\ x \implies$
 $(\exists\ x. e = UnaryExpr\ (UnaryZeroExtend\ ib\ rb)\ x)$
by (induction rule: *rep.induct*; auto)

lemma *rep-load-field* [rep]:

$g \vdash n \simeq e \implies$
 $is\ preevaluated\ (kind\ g\ n) \implies$
 $(\exists\ s. e = LeafExpr\ n\ s)$
by (induction rule: *rep.induct*; auto)

```

lemma rep-ref [rep]:
  g ⊢ n ≃ e ⇒
    kind g n = RefNode n' ⇒
      g ⊢ n' ≃ e
  by (induction rule: rep.induct; auto)

```

```

method solve-det uses node =
  (match node in kind - - = node - for node ⇒
    ⟨match rep in r: - ⇒ - = node - ⇒ - ⇒
      ⟨match IRNode.inject in i: (node - = node -) = - ⇒
        ⟨match RepE in e: - ⇒ (∧x. - = node x ⇒ -) ⇒ - ⇒
          ⟨match IRNode.distinct in d: node - ≠ RefNode - ⇒
            ⟨metis i e r d⟩⟩⟩⟩ |
    match node in kind - - = node - - for node ⇒
      ⟨match rep in r: - ⇒ - = node - - ⇒ - ⇒
        ⟨match IRNode.inject in i: (node - - = node - -) = - ⇒
          ⟨match RepE in e: - ⇒ (∧x y. - = node x y ⇒ -) ⇒ - ⇒
            ⟨match IRNode.distinct in d: node - - ≠ RefNode - ⇒
              ⟨metis i e r d⟩⟩⟩⟩ |
    match node in kind - - = node - - - for node ⇒
      ⟨match rep in r: - ⇒ - = node - - - ⇒ - ⇒
        ⟨match IRNode.inject in i: (node - - - = node - - -) = - ⇒
          ⟨match RepE in e: - ⇒ (∧x y z. - = node x y z ⇒ -) ⇒ - ⇒
            ⟨match IRNode.distinct in d: node - - - ≠ RefNode - ⇒
              ⟨metis i e r d⟩⟩⟩⟩ |
    match node in kind - - = node - - - for node ⇒
      ⟨match rep in r: - ⇒ - = node - - - ⇒ - ⇒
        ⟨match IRNode.inject in i: (node - - - = node - - -) = - ⇒
          ⟨match RepE in e: - ⇒ (∧x. - = node - - x ⇒ -) ⇒ - ⇒
            ⟨match IRNode.distinct in d: node - - - ≠ RefNode - ⇒
              ⟨metis i e r d⟩⟩⟩⟩)

```

Now we can prove that 'rep' and 'eval', and their list versions, are deterministic.

```

lemma repDet:
  shows (g ⊢ n ≃ e1) ⇒ (g ⊢ n ≃ e2) ⇒ e1 = e2
proof (induction arbitrary: e2 rule: rep.induct)
  case (ConstantNode n c)
  then show ?case using rep-constant by auto
next
  case (ParameterNode n i s)
  then show ?case
    by (metis IRNode.disc(2685) ParameterNodeE is-RefNode-def rep-parameter)
next
  case (ConditionalNode n c t f ce te fe)
  then show ?case
    using IRNode.distinct(593)
    using IRNode.inject(6) ConditionalNodeE rep-conditional

```

```

      by metis
next
  case (AbsNode n x xe)
  then show ?case
    by (solve-det node: AbsNode)
next
  case (NotNode n x xe)
  then show ?case
    by (solve-det node: NotNode)
next
  case (NegateNode n x xe)
  then show ?case
    by (solve-det node: NegateNode)
next
  case (LogicNegationNode n x xe)
  then show ?case
    by (solve-det node: LogicNegationNode)
next
  case (AddNode n x y xe ye)
  then show ?case
    by (solve-det node: AddNode)
next
  case (MulNode n x y xe ye)
  then show ?case
    by (solve-det node: MulNode)
next
  case (SubNode n x y xe ye)
  then show ?case
    by (solve-det node: SubNode)
next
  case (AndNode n x y xe ye)
  then show ?case
    by (solve-det node: AndNode)
next
  case (OrNode n x y xe ye)
  then show ?case
    by (solve-det node: OrNode)
next
  case (XorNode n x y xe ye)
  then show ?case
    by (solve-det node: XorNode)
next
  case (ShortCircuitOrNode n x y xe ye)
  then show ?case
    by (solve-det node: ShortCircuitOrNode)
next
  case (LeftShiftNode n x y xe ye)
  then show ?case
    by (solve-det node: LeftShiftNode)

```

```

next
  case (RightShiftNode n x y xe ye)
  then show ?case
    by (solve-det node: RightShiftNode)
next
  case (UnsignedRightShiftNode n x y xe ye)
  then show ?case
    by (solve-det node: UnsignedRightShiftNode)
next
  case (IntegerBelowNode n x y xe ye)
  then show ?case
    by (solve-det node: IntegerBelowNode)
next
  case (IntegerEqualsNode n x y xe ye)
  then show ?case
    by (solve-det node: IntegerEqualsNode)
next
  case (IntegerLessThanNode n x y xe ye)
  then show ?case
    by (solve-det node: IntegerLessThanNode)
next
  case (NarrowNode n x xe)
  then show ?case
    by (metis IRNode.distinct(2203) IRNode.inject(28) NarrowNodeE rep-narrow)
next
  case (SignExtendNode n x xe)
  then show ?case
    by (metis IRNode.distinct(2599) IRNode.inject(39) SignExtendNodeE rep-sign-extend)
next
  case (ZeroExtendNode n x xe)
  then show ?case
    by (metis IRNode.distinct(2753) IRNode.inject(50) ZeroExtendNodeE rep-zero-extend)
next
  case (LeafNode n s)
  then show ?case using rep-load-field LeafNodeE
    by (metis is-preevaluated.simps(53))
next
  case (RefNode n')
  then show ?case
    using rep-ref by blast
qed

lemma repAllDet:
   $g \vdash xs \simeq_L e1 \implies$ 
   $g \vdash xs \simeq_L e2 \implies$ 
   $e1 = e2$ 
proof (induction arbitrary: e2 rule: replist.induct)
  case RepNil
  then show ?case

```

```

    using replist.cases by auto
next
  case (RepCons x xe xs xse)
  then show ?case
    by (metis list.distinct(1) list.sel(1) list.sel(3) repDet replist.cases)
qed

```

```

lemma encodeEvalDet:
  [g,m,p] ⊢ e ↦ v1 ⟹
  [g,m,p] ⊢ e ↦ v2 ⟹
  v1 = v2
by (metis encodeeval-def evalDet repDet)

```

```

lemma graphDet: ([g,m,p] ⊢ n ↦ v1) ∧ ([g,m,p] ⊢ n ↦ v2) ⟹ v1 = v2
using encodeEvalDet by blast

```

7.8.2 Monotonicity of Graph Refinement

Lift refinement monotonicity to graph level. Hopefully these shouldn't really be required.

```

lemma mono-abs:
  assumes kind g1 n = AbsNode x ∧ kind g2 n = AbsNode x
  assumes (g1 ⊢ x ≃ xe1) ∧ (g2 ⊢ x ≃ xe2)
  assumes xe1 ≥ xe2
  assumes (g1 ⊢ n ≃ e1) ∧ (g2 ⊢ n ≃ e2)
  shows e1 ≥ e2
by (metis AbsNode assms(1) assms(2) assms(3) assms(4) mono-unary repDet)

```

```

lemma mono-not:
  assumes kind g1 n = NotNode x ∧ kind g2 n = NotNode x
  assumes (g1 ⊢ x ≃ xe1) ∧ (g2 ⊢ x ≃ xe2)
  assumes xe1 ≥ xe2
  assumes (g1 ⊢ n ≃ e1) ∧ (g2 ⊢ n ≃ e2)
  shows e1 ≥ e2
by (metis NotNode assms(1) assms(2) assms(3) assms(4) mono-unary repDet)

```

```

lemma mono-negate:
  assumes kind g1 n = NegateNode x ∧ kind g2 n = NegateNode x
  assumes (g1 ⊢ x ≃ xe1) ∧ (g2 ⊢ x ≃ xe2)
  assumes xe1 ≥ xe2
  assumes (g1 ⊢ n ≃ e1) ∧ (g2 ⊢ n ≃ e2)
  shows e1 ≥ e2
by (metis NegateNode assms(1) assms(2) assms(3) assms(4) mono-unary repDet)

```

```

lemma mono-logic-negation:
  assumes kind g1 n = LogicNegationNode x ∧ kind g2 n = LogicNegationNode x
  assumes (g1 ⊢ x ≃ xe1) ∧ (g2 ⊢ x ≃ xe2)
  assumes xe1 ≥ xe2
  assumes (g1 ⊢ n ≃ e1) ∧ (g2 ⊢ n ≃ e2)

```

shows $e1 \geq e2$
by (*metis* *LogicNegationNode* *assms*(1) *assms*(2) *assms*(3) *assms*(4) *mono-unary* *repDet*)

lemma *mono-narrow*:

assumes $\text{kind } g1 \ n = \text{NarrowNode } ib \ rb \ x \wedge \text{kind } g2 \ n = \text{NarrowNode } ib \ rb \ x$
assumes $(g1 \vdash x \simeq xe1) \wedge (g2 \vdash x \simeq xe2)$
assumes $xe1 \geq xe2$
assumes $(g1 \vdash n \simeq e1) \wedge (g2 \vdash n \simeq e2)$
shows $e1 \geq e2$
using *assms* *mono-unary* *repDet* *NarrowNode*
by *metis*

lemma *mono-sign-extend*:

assumes $\text{kind } g1 \ n = \text{SignExtendNode } ib \ rb \ x \wedge \text{kind } g2 \ n = \text{SignExtendNode } ib \ rb \ x$
assumes $(g1 \vdash x \simeq xe1) \wedge (g2 \vdash x \simeq xe2)$
assumes $xe1 \geq xe2$
assumes $(g1 \vdash n \simeq e1) \wedge (g2 \vdash n \simeq e2)$
shows $e1 \geq e2$
by (*metis* *SignExtendNode* *assms*(1) *assms*(2) *assms*(3) *assms*(4) *mono-unary* *repDet*)

lemma *mono-zero-extend*:

assumes $\text{kind } g1 \ n = \text{ZeroExtendNode } ib \ rb \ x \wedge \text{kind } g2 \ n = \text{ZeroExtendNode } ib \ rb \ x$
assumes $(g1 \vdash x \simeq xe1) \wedge (g2 \vdash x \simeq xe2)$
assumes $xe1 \geq xe2$
assumes $(g1 \vdash n \simeq e1) \wedge (g2 \vdash n \simeq e2)$
shows $e1 \geq e2$
using *assms* *mono-unary* *repDet* *ZeroExtendNode*
by *metis*

lemma *mono-conditional-graph*:

assumes $\text{kind } g1 \ n = \text{ConditionalNode } c \ t \ f \wedge \text{kind } g2 \ n = \text{ConditionalNode } c \ t \ f$
assumes $(g1 \vdash c \simeq ce1) \wedge (g2 \vdash c \simeq ce2)$
assumes $(g1 \vdash t \simeq te1) \wedge (g2 \vdash t \simeq te2)$
assumes $(g1 \vdash f \simeq fe1) \wedge (g2 \vdash f \simeq fe2)$
assumes $ce1 \geq ce2 \wedge te1 \geq te2 \wedge fe1 \geq fe2$
assumes $(g1 \vdash n \simeq e1) \wedge (g2 \vdash n \simeq e2)$
shows $e1 \geq e2$
using *ConditionalNodeE* *IRNode.inject*(6) *assms*(1) *assms*(2) *assms*(3) *assms*(4) *assms*(5) *assms*(6) *mono-conditional* *repDet* *rep-conditional*
by (*smt* (*verit*, *best*) *ConditionalNode*)

lemma *mono-add*:

assumes $\text{kind } g1 \ n = \text{AddNode } x \ y \wedge \text{kind } g2 \ n = \text{AddNode } x \ y$
assumes $(g1 \vdash x \simeq xe1) \wedge (g2 \vdash x \simeq xe2)$
assumes $(g1 \vdash y \simeq ye1) \wedge (g2 \vdash y \simeq ye2)$


```

assumes  $xe1 \geq xe2 \wedge ye1 \geq ye2$ 
assumes  $(g1 \vdash n \simeq e1) \wedge (g2 \vdash n \simeq e2)$ 
shows  $e1 \geq e2$ 
using mono-binary assms AddNodeE IRNode.inject(2) repDet rep-add
by (metis IRNode.distinct(205))

```

lemma *mono-mul*:

```

assumes  $kind\ g1\ n = MulNode\ x\ y \wedge kind\ g2\ n = MulNode\ x\ y$ 
assumes  $(g1 \vdash x \simeq xe1) \wedge (g2 \vdash x \simeq xe2)$ 
assumes  $(g1 \vdash y \simeq ye1) \wedge (g2 \vdash y \simeq ye2)$ 
assumes  $xe1 \geq xe2 \wedge ye1 \geq ye2$ 
assumes  $(g1 \vdash n \simeq e1) \wedge (g2 \vdash n \simeq e2)$ 
shows  $e1 \geq e2$ 
using mono-binary assms IRNode.inject(27) MulNodeE repDet rep-mul
by (smt (verit, best) MulNode)

```

lemma *term-graph-evaluation*:

```

 $(g \vdash n \sqsubseteq e) \implies (\forall\ m\ p\ v.\ ([m,p] \vdash e \mapsto v) \longrightarrow ([g,m,p] \vdash n \mapsto v))$ 
unfolding graph-represents-expression-def apply auto
by (meson encodeeval-def)

```

lemma *encodes-contains*:

```

 $g \vdash n \simeq e \implies$ 
 $kind\ g\ n \neq NoNode$ 
apply (induction rule: rep.induct)
apply (match IRNode.distinct in e: ?n ≠ NoNode ⇒
 $\langle presburger\ add: e \rangle +$ 
apply force
by fastforce)

```

lemma *no-encoding*:

```

assumes  $n \notin ids\ g$ 
shows  $\neg(g \vdash n \simeq e)$ 
using assms apply simp apply (rule notI) by (induction e; simp add: en-
codes-contains)

```

lemma *not-excluded-keep-type*:

```

assumes  $n \in ids\ g1$ 
assumes  $n \notin excluded$ 
assumes  $(excluded \sqsubseteq as-set\ g1) \subseteq as-set\ g2$ 
shows  $kind\ g1\ n = kind\ g2\ n \wedge stamp\ g1\ n = stamp\ g2\ n$ 
using assms unfolding as-set-def domain-subtraction-def by blast

```

method *metis-node-eq-unary* **for** $node :: 'a \Rightarrow IRNode =$

```

 $(match\ IRNode.inject\ in\ i: (node\ - = node\ -) = - \Rightarrow$ 
 $\langle metis\ i \rangle)$ 

```

method *metis-node-eq-binary* **for** $node :: 'a \Rightarrow 'a \Rightarrow IRNode =$

```

 $(match\ IRNode.inject\ in\ i: (node\ -\ - = node\ -\ -) = - \Rightarrow$ 

```

```

    ⟨metis i⟩
method metis-node-eq-ternary for node :: 'a ⇒ 'a ⇒ 'a ⇒ IRNode =
  (match IRNode.inject in i: (node - - - = node - - -) = - ⇒
    ⟨metis i⟩)

```

7.8.3 Lift Data-flow Tree Refinement to Graph Refinement

```

theorem graph-semantic-preservation:
  assumes a:  $e1' \geq e2'$ 
  assumes b:  $(\{n'\} \trianglelefteq \text{as-set } g1) \subseteq \text{as-set } g2$ 
  assumes c:  $g1 \vdash n' \simeq e1'$ 
  assumes d:  $g2 \vdash n' \simeq e2'$ 
  shows graph-refinement g1 g2
  unfolding graph-refinement-def apply rule
  apply (metis b d ids-some no-encoding not-excluded-keep-type singleton-iff sub-
    setI)
  apply (rule allI) apply (rule impI) apply (rule allI) apply (rule impI)
  unfolding graph-represents-expression-def
proof –
  fix n e1
  assume e:  $n \in \text{ids } g1$ 
  assume f:  $(g1 \vdash n \simeq e1)$ 

  show  $\exists e2. (g2 \vdash n \simeq e2) \wedge e1 \geq e2$ 
proof (cases n = n')
  case True
  have g:  $e1 = e1'$  using c f True repDet by simp
  have h:  $(g2 \vdash n \simeq e2') \wedge e1' \geq e2'$ 
    using True a d by blast
  then show ?thesis
    using g by blast
next
  case False
  have n  $\notin \{n'\}$ 
    using False by simp
  then have i:  $\text{kind } g1 \ n = \text{kind } g2 \ n \wedge \text{stamp } g1 \ n = \text{stamp } g2 \ n$ 
    using not-excluded-keep-type
    using b e by presburger
  show ?thesis using f i
proof (induction e1)
  case (ConstantNode n c)
  then show ?case
    by (metis eq-refl rep.ConstantNode)
next
  case (ParameterNode n i s)
  then show ?case
    by (metis eq-refl rep.ParameterNode)
next
  case (ConditionalNode n c t f ce1 te1 fe1)

```

```

have k: g1 ⊢ n ≃ ConditionalExpr ce1 te1 fe1 using f ConditionalNode
  by (simp add: ConditionalNode.hyps(2) rep.ConditionalNode)
obtain cn tn fn where l: kind g1 n = ConditionalNode cn tn fn
  using ConditionalNode.hyps(1) by blast
then have mc: g1 ⊢ cn ≃ ce1
  using ConditionalNode.hyps(1) ConditionalNode.hyps(2) by fastforce
from l have mt: g1 ⊢ tn ≃ te1
  using ConditionalNode.hyps(1) ConditionalNode.hyps(3) by fastforce
from l have mf: g1 ⊢ fn ≃ fe1
  using ConditionalNode.hyps(1) ConditionalNode.hyps(4) by fastforce
then show ?case
proof -
  have g1 ⊢ cn ≃ ce1 using mc by simp
  have g1 ⊢ tn ≃ te1 using mt by simp
  have g1 ⊢ fn ≃ fe1 using mf by simp
  have cer: ∃ ce2. (g2 ⊢ cn ≃ ce2) ∧ ce1 ≥ ce2
    using ConditionalNode
    using a b c d l no-encoding not-excluded-keep-type repDet singletonD
    by (metis-node-eq-ternary ConditionalNode)
  have ter: ∃ te2. (g2 ⊢ tn ≃ te2) ∧ te1 ≥ te2
    using ConditionalNode a b c d l no-encoding not-excluded-keep-type repDet
    singletonD
    by (metis-node-eq-ternary ConditionalNode)
  have ∃ fe2. (g2 ⊢ fn ≃ fe2) ∧ fe1 ≥ fe2
    using ConditionalNode a b c d l no-encoding not-excluded-keep-type repDet
    singletonD
    by (metis-node-eq-ternary ConditionalNode)
  then have ∃ ce2 te2 fe2. (g2 ⊢ n ≃ ConditionalExpr ce2 te2 fe2) ∧
    ConditionalExpr ce1 te1 fe1 ≥ ConditionalExpr ce2 te2 fe2
    using ConditionalNode.premis l rep.ConditionalNode cer ter
    by (smt (verit) mono-conditional)
  then show ?thesis
    by meson
qed
next
case (AbsNode n x xe1)
have k: g1 ⊢ n ≃ UnaryExpr UnaryAbs xe1 using f AbsNode
  by (simp add: AbsNode.hyps(2) rep.AbsNode)
obtain xn where l: kind g1 n = AbsNode xn
  using AbsNode.hyps(1) by blast
then have m: g1 ⊢ xn ≃ xe1
  using AbsNode.hyps(1) AbsNode.hyps(2) by fastforce
then show ?case
proof (cases xn = n')
case True
  then have n: xe1 = e1' using c m repDet by simp
  then have ev: g2 ⊢ n ≃ UnaryExpr UnaryAbs e2' using AbsNode.hyps(1)
    l m n
    using AbsNode.premis True d rep.AbsNode by simp

```

```

    then have r: UnaryExpr UnaryAbs e1' ≥ UnaryExpr UnaryAbs e2'
      by (meson a mono-unary)
    then show ?thesis using ev r
      by (metis n)
  next
    case False
    have g1 ⊢ xn ≃ xe1 using m by simp
    have ∃ xe2. (g2 ⊢ xn ≃ xe2) ∧ xe1 ≥ xe2
      using AbsNode
    using False b encodes-contains l not-excluded-keep-type not-in-g singleton-iff
      by (metis-node-eq-unary AbsNode)
    then have ∃ xe2. (g2 ⊢ n ≃ UnaryExpr UnaryAbs xe2) ∧ UnaryExpr
UnaryAbs xe1 ≥ UnaryExpr UnaryAbs xe2
      by (metis AbsNode.premis l mono-unary rep.AbsNode)
    then show ?thesis
      by meson
  qed
next
  case (NotNode n x xe1)
  have k: g1 ⊢ n ≃ UnaryExpr UnaryNot xe1 using f NotNode
    by (simp add: NotNode.hyps(2) rep.NotNode)
  obtain xn where l: kind g1 n = NotNode xn
    using NotNode.hyps(1) by blast
  then have m: g1 ⊢ xn ≃ xe1
    using NotNode.hyps(1) NotNode.hyps(2) by fastforce
  then show ?case
  proof (cases xn = n')
    case True
    then have n: xe1 = e1' using c m repDet by simp
    then have ev: g2 ⊢ n ≃ UnaryExpr UnaryNot e2' using NotNode.hyps(1)
l m n
      using NotNode.premis True d rep.NotNode by simp
    then have r: UnaryExpr UnaryNot e1' ≥ UnaryExpr UnaryNot e2'
      by (meson a mono-unary)
    then show ?thesis using ev r
      by (metis n)
  next
    case False
    have g1 ⊢ xn ≃ xe1 using m by simp
    have ∃ xe2. (g2 ⊢ xn ≃ xe2) ∧ xe1 ≥ xe2
      using NotNode
    using False i b l not-excluded-keep-type singletonD no-encoding
      by (metis-node-eq-unary NotNode)
    then have ∃ xe2. (g2 ⊢ n ≃ UnaryExpr UnaryNot xe2) ∧ UnaryExpr
UnaryNot xe1 ≥ UnaryExpr UnaryNot xe2
      by (metis NotNode.premis l mono-unary rep.NotNode)
    then show ?thesis
      by meson
  qed

```

```

next
  case (NegateNode n x xe1)
  have k:  $g1 \vdash n \simeq \text{UnaryExpr UnaryNeg xe1}$  using f NegateNode
    by (simp add: NegateNode.hyps(2) rep.NegateNode)
  obtain xn where l: kind g1 n = NegateNode xn
    using NegateNode.hyps(1) by blast
  then have m:  $g1 \vdash xn \simeq xe1$ 
    using NegateNode.hyps(1) NegateNode.hyps(2) by fastforce
  then show ?case
  proof (cases xn = n')
    case True
    then have n:  $xe1 = e1'$  using c m repDet by simp
    then have ev:  $g2 \vdash n \simeq \text{UnaryExpr UnaryNeg e2'}$  using NegateNode.hyps(1)
  l m n
    using NegateNode.premis True d rep.NegateNode by simp
  then have r:  $\text{UnaryExpr UnaryNeg e1'} \geq \text{UnaryExpr UnaryNeg e2'}$ 
    by (meson a mono-unary)
  then show ?thesis using ev r
    by (metis n)
  next
  case False
  have  $g1 \vdash xn \simeq xe1$  using m by simp
  have  $\exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
    using NegateNode
    using False i b l not-excluded-keep-type singletonD no-encoding
    by (metis-node-eq-unary NegateNode)
  then have  $\exists xe2. (g2 \vdash n \simeq \text{UnaryExpr UnaryNeg xe2}) \wedge \text{UnaryExpr}$ 
  UnaryNeg xe1  $\geq \text{UnaryExpr UnaryNeg xe2}$ 
    by (metis NegateNode.premis l mono-unary rep.NegateNode)
  then show ?thesis
    by meson
  qed
  next
  case (LogicNegationNode n x xe1)
  have k:  $g1 \vdash n \simeq \text{UnaryExpr UnaryLogicNegation xe1}$  using f LogicNegationNode
  by (simp add: LogicNegationNode.hyps(2) rep.LogicNegationNode)
  obtain xn where l: kind g1 n = LogicNegationNode xn
    using LogicNegationNode.hyps(1) by blast
  then have m:  $g1 \vdash xn \simeq xe1$ 
    using LogicNegationNode.hyps(1) LogicNegationNode.hyps(2) by fastforce
  then show ?case
  proof (cases xn = n')
    case True
    then have n:  $xe1 = e1'$  using c m repDet by simp
    then have ev:  $g2 \vdash n \simeq \text{UnaryExpr UnaryLogicNegation e2'}$  using
  LogicNegationNode.hyps(1) l m n
    using LogicNegationNode.premis True d rep.LogicNegationNode by simp
  then have r:  $\text{UnaryExpr UnaryLogicNegation e1'} \geq \text{UnaryExpr UnaryLog-}$ 

```

```

icNegation e2'
  by (meson a mono-unary)
  then show ?thesis using ev r
  by (metis n)
next
case False
have g1 ⊢ xn ≃ xe1 using m by simp
have ∃ xe2. (g2 ⊢ xn ≃ xe2) ∧ xe1 ≥ xe2
  using LogicNegationNode
  using False i b l not-excluded-keep-type singletonD no-encoding
  by (metis-node-eq-unary LogicNegationNode)
  then have ∃ xe2. (g2 ⊢ n ≃ UnaryExpr UnaryLogicNegation xe2) ∧
UnaryExpr UnaryLogicNegation xe1 ≥ UnaryExpr UnaryLogicNegation xe2
  by (metis LogicNegationNode.prem1 mono-unary rep.LogicNegationNode)
  then show ?thesis
  by meson
qed
next
case (AddNode n x y xe1 ye1)
have k: g1 ⊢ n ≃ BinaryExpr BinAdd xe1 ye1 using f AddNode
  by (simp add: AddNode.hyps(2) rep.AddNode)
obtain xn yn where l: kind g1 n = AddNode xn yn
  using AddNode.hyps(1) by blast
then have mx: g1 ⊢ xn ≃ xe1
  using AddNode.hyps(1) AddNode.hyps(2) by fastforce
from l have my: g1 ⊢ yn ≃ ye1
  using AddNode.hyps(1) AddNode.hyps(3) by fastforce
then show ?case
proof -
  have g1 ⊢ xn ≃ xe1 using mx by simp
  have g1 ⊢ yn ≃ ye1 using my by simp
  have xer: ∃ xe2. (g2 ⊢ xn ≃ xe2) ∧ xe1 ≥ xe2
    using AddNode
    using a b c d l no-encoding not-excluded-keep-type repDet singletonD
    by (metis-node-eq-binary AddNode)
  have ∃ ye2. (g2 ⊢ yn ≃ ye2) ∧ ye1 ≥ ye2
    using AddNode
    using a b c d l no-encoding not-excluded-keep-type repDet singletonD
    by (metis-node-eq-binary AddNode)
  then have ∃ xe2 ye2. (g2 ⊢ n ≃ BinaryExpr BinAdd xe2 ye2) ∧ BinaryExpr
BinAdd xe1 ye1 ≥ BinaryExpr BinAdd xe2 ye2
    by (metis AddNode.prem1 mono-binary rep.AddNode xer)
  then show ?thesis
  by meson
qed
next
case (MulNode n x y xe1 ye1)
have k: g1 ⊢ n ≃ BinaryExpr BinMul xe1 ye1 using f MulNode
  by (simp add: MulNode.hyps(2) rep.MulNode)

```

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obtain  $xn\ yn$  where  $l$ :  $kind\ g1\ n = MulNode\ xn\ yn$ 
  using  $MulNode.hyps(1)$  by blast
then have  $mx$ :  $g1 \vdash xn \simeq xe1$ 
  using  $MulNode.hyps(1)\ MulNode.hyps(2)$  by fastforce
from  $l$  have  $my$ :  $g1 \vdash yn \simeq ye1$ 
  using  $MulNode.hyps(1)\ MulNode.hyps(3)$  by fastforce
then show ?case
proof –
  have  $g1 \vdash xn \simeq xe1$  using  $mx$  by simp
  have  $g1 \vdash yn \simeq ye1$  using  $my$  by simp
  have  $xer$ :  $\exists\ xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
    using  $MulNode$ 
    using  $a\ b\ c\ d\ l\ no-encoding\ not-excluded-keep-type\ repDet\ singletonD$ 
    by (metis-node-eq-binary  $MulNode$ )
  have  $\exists\ ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
    using  $MulNode$ 
    using  $a\ b\ c\ d\ l\ no-encoding\ not-excluded-keep-type\ repDet\ singletonD$ 
    by (metis-node-eq-binary  $MulNode$ )
  then have  $\exists\ xe2\ ye2. (g2 \vdash n \simeq BinaryExpr\ BinMul\ xe2\ ye2) \wedge BinaryExpr$ 
     $BinMul\ xe1\ ye1 \geq BinaryExpr\ BinMul\ xe2\ ye2$ 
    by (metis  $MulNode.premis\ l\ mono-binary\ rep.MulNode\ xer$ )
  then show ?thesis
    by meson
qed
next
case ( $SubNode\ n\ x\ y\ xe1\ ye1$ )
have  $k$ :  $g1 \vdash n \simeq BinaryExpr\ BinSub\ xe1\ ye1$  using  $f\ SubNode$ 
  by (simp  $add$ :  $SubNode.hyps(2)\ rep.SubNode$ )
obtain  $xn\ yn$  where  $l$ :  $kind\ g1\ n = SubNode\ xn\ yn$ 
  using  $SubNode.hyps(1)$  by blast
then have  $mx$ :  $g1 \vdash xn \simeq xe1$ 
  using  $SubNode.hyps(1)\ SubNode.hyps(2)$  by fastforce
from  $l$  have  $my$ :  $g1 \vdash yn \simeq ye1$ 
  using  $SubNode.hyps(1)\ SubNode.hyps(3)$  by fastforce
then show ?case
proof –
  have  $g1 \vdash xn \simeq xe1$  using  $mx$  by simp
  have  $g1 \vdash yn \simeq ye1$  using  $my$  by simp
  have  $xer$ :  $\exists\ xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
    using  $SubNode$ 
    using  $a\ b\ c\ d\ l\ no-encoding\ not-excluded-keep-type\ repDet\ singletonD$ 
    by (metis-node-eq-binary  $SubNode$ )
  have  $\exists\ ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
using  $SubNode\ a\ b\ c\ d\ l\ no-encoding\ not-excluded-keep-type\ repDet\ singletonD$ 
  by (metis-node-eq-binary  $SubNode$ )
  then have  $\exists\ xe2\ ye2. (g2 \vdash n \simeq BinaryExpr\ BinSub\ xe2\ ye2) \wedge BinaryExpr$ 
     $BinSub\ xe1\ ye1 \geq BinaryExpr\ BinSub\ xe2\ ye2$ 
    by (metis  $SubNode.premis\ l\ mono-binary\ rep.SubNode\ xer$ )
  then show ?thesis

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```

      by meson
    qed
  next
    case (AndNode n x y xe1 ye1)
    have k:  $g1 \vdash n \simeq \text{BinaryExpr BinAnd } xe1 \ ye1$  using f AndNode
      by (simp add: AndNode.hyps(2) rep.AndNode)
    obtain xn yn where l: kind g1 n = AndNode xn yn
      using AndNode.hyps(1) by blast
    then have mx:  $g1 \vdash xn \simeq xe1$ 
      using AndNode.hyps(1) AndNode.hyps(2) by fastforce
    from l have my:  $g1 \vdash yn \simeq ye1$ 
      using AndNode.hyps(1) AndNode.hyps(3) by fastforce
    then show ?case
    proof -
      have  $g1 \vdash xn \simeq xe1$  using mx by simp
      have  $g1 \vdash yn \simeq ye1$  using my by simp
      have xer:  $\exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
        using AndNode
        using a b c d l no-encoding not-excluded-keep-type repDet singletonD
        by (metis-node-eq-binary AndNode)
      have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
        using AndNode a b c d l no-encoding not-excluded-keep-type repDet
        singletonD
        by (metis-node-eq-binary AndNode)
      then have  $\exists xe2 \ ye2. (g2 \vdash n \simeq \text{BinaryExpr BinAnd } xe2 \ ye2) \wedge \text{BinaryExpr}$ 
        BinAnd xe1 ye1  $\geq \text{BinaryExpr BinAnd } xe2 \ ye2$ 
        by (metis AndNode.premis l mono-binary rep.AndNode xer)
      then show ?thesis
        by meson
    qed
  next
    case (OrNode n x y xe1 ye1)
    have k:  $g1 \vdash n \simeq \text{BinaryExpr BinOr } xe1 \ ye1$  using f OrNode
      by (simp add: OrNode.hyps(2) rep.OrNode)
    obtain xn yn where l: kind g1 n = OrNode xn yn
      using OrNode.hyps(1) by blast
    then have mx:  $g1 \vdash xn \simeq xe1$ 
      using OrNode.hyps(1) OrNode.hyps(2) by fastforce
    from l have my:  $g1 \vdash yn \simeq ye1$ 
      using OrNode.hyps(1) OrNode.hyps(3) by fastforce
    then show ?case
    proof -
      have  $g1 \vdash xn \simeq xe1$  using mx by simp
      have  $g1 \vdash yn \simeq ye1$  using my by simp
      have xer:  $\exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
        using OrNode
        using a b c d l no-encoding not-excluded-keep-type repDet singletonD
        by (metis-node-eq-binary OrNode)
      have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 

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    using OrNode a b c d l no-encoding not-excluded-keep-type repDet singletonD
    by (metis-node-eq-binary OrNode)
    then have  $\exists xe2 ye2. (g2 \vdash n \simeq BinaryExpr BinOr xe2 ye2) \wedge BinaryExpr$ 
    BinOr xe1 ye1  $\geq BinaryExpr BinOr xe2 ye2$ 
    by (metis OrNode.premis l mono-binary rep.OrNode xer)
    then show ?thesis
    by meson
  qed
next
case (XorNode n x y xe1 ye1)
have k:  $g1 \vdash n \simeq BinaryExpr BinXor xe1 ye1$  using f XorNode
by (simp add: XorNode.hyps(2) rep.XorNode)
obtain xn yn where l: kind g1 n = XorNode xn yn
using XorNode.hyps(1) by blast
then have mx:  $g1 \vdash xn \simeq xe1$ 
using XorNode.hyps(1) XorNode.hyps(2) by fastforce
from l have my:  $g1 \vdash yn \simeq ye1$ 
using XorNode.hyps(1) XorNode.hyps(3) by fastforce
then show ?case
proof -
  have  $g1 \vdash xn \simeq xe1$  using mx by simp
  have  $g1 \vdash yn \simeq ye1$  using my by simp
  have xer:  $\exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
  using XorNode
  using a b c d l no-encoding not-excluded-keep-type repDet singletonD
  by (metis-node-eq-binary XorNode)
  have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
  using XorNode a b c d l no-encoding not-excluded-keep-type repDet
  singletonD
  by (metis-node-eq-binary XorNode)
  then have  $\exists xe2 ye2. (g2 \vdash n \simeq BinaryExpr BinXor xe2 ye2) \wedge BinaryExpr$ 
  BinXor xe1 ye1  $\geq BinaryExpr BinXor xe2 ye2$ 
  by (metis XorNode.premis l mono-binary rep.XorNode xer)
  then show ?thesis
  by meson
qed
next
case (ShortCircuitOrNode n x y xe1 ye1)
have k:  $g1 \vdash n \simeq BinaryExpr BinShortCircuitOr xe1 ye1$  using f ShortCir-
cuitOrNode
by (simp add: ShortCircuitOrNode.hyps(2) rep.ShortCircuitOrNode)
obtain xn yn where l: kind g1 n = ShortCircuitOrNode xn yn
using ShortCircuitOrNode.hyps(1) by blast
then have mx:  $g1 \vdash xn \simeq xe1$ 
using ShortCircuitOrNode.hyps(1) ShortCircuitOrNode.hyps(2) by fastforce
from l have my:  $g1 \vdash yn \simeq ye1$ 
using ShortCircuitOrNode.hyps(1) ShortCircuitOrNode.hyps(3) by fastforce
then show ?case
proof -

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    have  $g1 \vdash xn \simeq xe1$  using  $mx$  by simp
    have  $g1 \vdash yn \simeq ye1$  using  $my$  by simp
    have  $xer: \exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
      using ShortCircuitOrNode
      using  $a\ b\ c\ d\ l$  no-encoding not-excluded-keep-type repDet singletonD
      by (metis-node-eq-binary ShortCircuitOrNode)
    have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
      using ShortCircuitOrNode a b c d l no-encoding not-excluded-keep-type
repDet singletonD
      by (metis-node-eq-binary ShortCircuitOrNode)
    then have  $\exists xe2\ ye2. (g2 \vdash n \simeq BinaryExpr\ BinShortCircuitOr\ xe2\ ye2) \wedge$ 
BinaryExpr BinShortCircuitOr xe1 ye1  $\geq BinaryExpr BinShortCircuitOr xe2 ye2$ 
      by (metis ShortCircuitOrNode.premis l mono-binary rep.ShortCircuitOrNode
xer)
    then show ?thesis
      by meson
  qed
next
case (LeftShiftNode n x y xe1 ye1)
have  $k: g1 \vdash n \simeq BinaryExpr\ BinLeftShift\ xe1\ ye1$  using  $f$  LeftShiftNode
  by (simp add: LeftShiftNode.hyps(2) rep.LeftShiftNode)
obtain  $xn\ yn$  where  $l: kind\ g1\ n = LeftShiftNode\ xn\ yn$ 
  using LeftShiftNode.hyps(1) by blast
then have  $mx: g1 \vdash xn \simeq xe1$ 
  using LeftShiftNode.hyps(1) LeftShiftNode.hyps(2) by fastforce
from  $l$  have  $my: g1 \vdash yn \simeq ye1$ 
  using LeftShiftNode.hyps(1) LeftShiftNode.hyps(3) by fastforce
then show ?case
proof -
  have  $g1 \vdash xn \simeq xe1$  using  $mx$  by simp
  have  $g1 \vdash yn \simeq ye1$  using  $my$  by simp
  have  $xer: \exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
    using LeftShiftNode
    using  $a\ b\ c\ d\ l$  no-encoding not-excluded-keep-type repDet singletonD
    by (metis-node-eq-binary LeftShiftNode)
  have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
    using LeftShiftNode a b c d l no-encoding not-excluded-keep-type repDet
singletonD
    by (metis-node-eq-binary LeftShiftNode)
  then have  $\exists xe2\ ye2. (g2 \vdash n \simeq BinaryExpr\ BinLeftShift\ xe2\ ye2) \wedge$ 
BinaryExpr BinLeftShift xe1 ye1  $\geq BinaryExpr BinLeftShift xe2 ye2$ 
    by (metis LeftShiftNode.premis l mono-binary rep.LeftShiftNode xer)
  then show ?thesis
    by meson
  qed
next
case (RightShiftNode n x y xe1 ye1)
have  $k: g1 \vdash n \simeq BinaryExpr\ BinRightShift\ xe1\ ye1$  using  $f$  RightShiftNode
  by (simp add: RightShiftNode.hyps(2) rep.RightShiftNode)

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obtain  $xn\ yn$  where  $l: \text{kind } g1\ n = \text{RightShiftNode } xn\ yn$ 
  using  $\text{RightShiftNode.hyps}(1)$  by blast
then have  $m\!x: g1 \vdash xn \simeq xe1$ 
  using  $\text{RightShiftNode.hyps}(1)\ \text{RightShiftNode.hyps}(2)$  by fastforce
from  $l$  have  $m\!y: g1 \vdash yn \simeq ye1$ 
  using  $\text{RightShiftNode.hyps}(1)\ \text{RightShiftNode.hyps}(3)$  by fastforce
then show ?case
proof –
  have  $g1 \vdash xn \simeq xe1$  using  $m\!x$  by simp
  have  $g1 \vdash yn \simeq ye1$  using  $m\!y$  by simp
  have  $x\!e\!r: \exists\ xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
    using  $\text{RightShiftNode}$ 
    using  $a\ b\ c\ d\ l\ \text{no-encoding not-excluded-keep-type repDet singletonD}$ 
    by  $(\text{metis-node-eq-binary } \text{RightShiftNode})$ 
  have  $\exists\ ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
    using  $\text{RightShiftNode } a\ b\ c\ d\ l\ \text{no-encoding not-excluded-keep-type repDet}$ 
    singletonD
    by  $(\text{metis-node-eq-binary } \text{RightShiftNode})$ 
  then have  $\exists\ xe2\ ye2. (g2 \vdash n \simeq \text{BinaryExpr BinRightShift } xe2\ ye2) \wedge$ 
BinaryExpr BinRightShift  $xe1\ ye1 \geq \text{BinaryExpr BinRightShift } xe2\ ye2$ 
    by  $(\text{metis } \text{RightShiftNode.premis } l\ \text{mono-binary rep.RightShiftNode } x\!e\!r)$ 
  then show ?thesis
    by meson
  qed
next
  case  $(\text{UnsignedRightShiftNode } n\ x\ y\ xe1\ ye1)$ 
  have  $k: g1 \vdash n \simeq \text{BinaryExpr BinURightShift } xe1\ ye1$  using  $f\ \text{UnsignedRightShiftNode}$ 
    by  $(\text{simp add: UnsignedRightShiftNode.hyps}(2)\ \text{rep.UnsignedRightShiftNode})$ 
  obtain  $xn\ yn$  where  $l: \text{kind } g1\ n = \text{UnsignedRightShiftNode } xn\ yn$ 
    using  $\text{UnsignedRightShiftNode.hyps}(1)$  by blast
  then have  $m\!x: g1 \vdash xn \simeq xe1$ 
    using  $\text{UnsignedRightShiftNode.hyps}(1)\ \text{UnsignedRightShiftNode.hyps}(2)$  by
fastforce
  from  $l$  have  $m\!y: g1 \vdash yn \simeq ye1$ 
    using  $\text{UnsignedRightShiftNode.hyps}(1)\ \text{UnsignedRightShiftNode.hyps}(3)$  by
fastforce
  then show ?case
  proof –
    have  $g1 \vdash xn \simeq xe1$  using  $m\!x$  by simp
    have  $g1 \vdash yn \simeq ye1$  using  $m\!y$  by simp
    have  $x\!e\!r: \exists\ xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
      using  $\text{UnsignedRightShiftNode}$ 
      using  $a\ b\ c\ d\ l\ \text{no-encoding not-excluded-keep-type repDet singletonD}$ 
      by  $(\text{metis-node-eq-binary } \text{UnsignedRightShiftNode})$ 
    have  $\exists\ ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
      using  $\text{UnsignedRightShiftNode } a\ b\ c\ d\ l\ \text{no-encoding not-excluded-keep-type}$ 
      repDet singletonD
      by  $(\text{metis-node-eq-binary } \text{UnsignedRightShiftNode})$ 

```

```

      then have  $\exists xe2 ye2. (g2 \vdash n \simeq BinaryExpr BinURightShift xe2 ye2) \wedge$ 
BinaryExpr BinURightShift xe1 ye1  $\geq BinaryExpr BinURightShift xe2 ye2$ 
      by (metis UnsignedRightShiftNode.premis l mono-binary rep.UnsignedRightShiftNode
xer)
      then show ?thesis
      by meson
    qed
  next
  case (IntegerBelowNode n x y xe1 ye1)
  have k:  $g1 \vdash n \simeq BinaryExpr BinIntegerBelow xe1 ye1$  using f IntegerBe-
lowNode
  by (simp add: IntegerBelowNode.hyps(2) rep.IntegerBelowNode)
  obtain xn yn where l: kind g1 n = IntegerBelowNode xn yn
  using IntegerBelowNode.hyps(1) by blast
  then have mx:  $g1 \vdash xn \simeq xe1$ 
  using IntegerBelowNode.hyps(1) IntegerBelowNode.hyps(2) by fastforce
  from l have my:  $g1 \vdash yn \simeq ye1$ 
  using IntegerBelowNode.hyps(1) IntegerBelowNode.hyps(3) by fastforce
  then show ?case
  proof -
    have  $g1 \vdash xn \simeq xe1$  using mx by simp
    have  $g1 \vdash yn \simeq ye1$  using my by simp
    have xer:  $\exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
    using IntegerBelowNode
    using a b c d l no-encoding not-excluded-keep-type repDet singletonD
    by (metis-node-eq-binary IntegerBelowNode)
    have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
    using IntegerBelowNode a b c d l no-encoding not-excluded-keep-type repDet
singletonD
    by (metis-node-eq-binary IntegerBelowNode)
    then have  $\exists xe2 ye2. (g2 \vdash n \simeq BinaryExpr BinIntegerBelow xe2 ye2) \wedge$ 
BinaryExpr BinIntegerBelow xe1 ye1  $\geq BinaryExpr BinIntegerBelow xe2 ye2$ 
    by (metis IntegerBelowNode.premis l mono-binary rep.IntegerBelowNode
xer)
    then show ?thesis
    by meson
  qed
next
case (IntegerEqualsNode n x y xe1 ye1)
have k:  $g1 \vdash n \simeq BinaryExpr BinIntegerEquals xe1 ye1$  using f IntegerEqual-
sNode
by (simp add: IntegerEqualsNode.hyps(2) rep.IntegerEqualsNode)
obtain xn yn where l: kind g1 n = IntegerEqualsNode xn yn
using IntegerEqualsNode.hyps(1) by blast
then have mx:  $g1 \vdash xn \simeq xe1$ 
using IntegerEqualsNode.hyps(1) IntegerEqualsNode.hyps(2) by fastforce
from l have my:  $g1 \vdash yn \simeq ye1$ 
using IntegerEqualsNode.hyps(1) IntegerEqualsNode.hyps(3) by fastforce
then show ?case

```

```

proof –
  have  $g1 \vdash xn \simeq xe1$  using  $mx$  by simp
  have  $g1 \vdash yn \simeq ye1$  using  $my$  by simp
  have  $xer: \exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
    using IntegerEqualsNode
    using  $a\ b\ c\ d\ l$  no-encoding not-excluded-keep-type repDet singletonD
    by (metis-node-eq-binary IntegerEqualsNode)
  have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
    using IntegerEqualsNode  $a\ b\ c\ d\ l$  no-encoding not-excluded-keep-type
repDet singletonD
    by (metis-node-eq-binary IntegerEqualsNode)
  then have  $\exists xe2\ ye2. (g2 \vdash n \simeq BinaryExpr\ BinIntegerEquals\ xe2\ ye2) \wedge$ 
BinaryExpr BinIntegerEquals xe1 ye1  $\geq$  BinaryExpr BinIntegerEquals xe2 ye2
    by (metis IntegerEqualsNode.premis l mono-binary rep.IntegerEqualsNode
xer)
    then show ?thesis
      by meson
  qed
next
  case (IntegerLessThanNode  $n\ x\ y\ xe1\ ye1$ )
    have  $k: g1 \vdash n \simeq BinaryExpr\ BinIntegerLessThan\ xe1\ ye1$  using  $f$  IntegerLessThanNode
    by (simp add: IntegerLessThanNode.hyps(2) rep.IntegerLessThanNode)
    obtain  $xn\ yn$  where  $l: kind\ g1\ n = IntegerLessThanNode\ xn\ yn$ 
    using IntegerLessThanNode.hyps(1) by blast
    then have  $mx: g1 \vdash xn \simeq xe1$ 
    using IntegerLessThanNode.hyps(1) IntegerLessThanNode.hyps(2) by fast-force
    from  $l$  have  $my: g1 \vdash yn \simeq ye1$ 
    using IntegerLessThanNode.hyps(1) IntegerLessThanNode.hyps(3) by fast-force
    then show ?case
      proof –
        have  $g1 \vdash xn \simeq xe1$  using  $mx$  by simp
        have  $g1 \vdash yn \simeq ye1$  using  $my$  by simp
        have  $xer: \exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
          using IntegerLessThanNode
          using  $a\ b\ c\ d\ l$  no-encoding not-excluded-keep-type repDet singletonD
          by (metis-node-eq-binary IntegerLessThanNode)
        have  $\exists ye2. (g2 \vdash yn \simeq ye2) \wedge ye1 \geq ye2$ 
          using IntegerLessThanNode  $a\ b\ c\ d\ l$  no-encoding not-excluded-keep-type
repDet singletonD
          by (metis-node-eq-binary IntegerLessThanNode)
        then have  $\exists xe2\ ye2. (g2 \vdash n \simeq BinaryExpr\ BinIntegerLessThan\ xe2\ ye2)$ 
 $\wedge BinaryExpr\ BinIntegerLessThan\ xe1\ ye1 \geq BinaryExpr\ BinIntegerLessThan\ xe2\ ye2$ 
          by (metis IntegerLessThanNode.premis l mono-binary rep.IntegerLessThanNode
xer)
        then show ?thesis

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      by meson
    qed
  next
    case (NarrowNode n inputBits resultBits x xe1)
    have k:  $g1 \vdash n \simeq \text{UnaryExpr } (\text{UnaryNarrow inputBits resultBits}) xe1$  using
    f NarrowNode
      by (simp add: NarrowNode.hyps(2) rep.NarrowNode)
    obtain xn where l: kind g1 n = NarrowNode inputBits resultBits xn
      using NarrowNode.hyps(1) by blast
    then have m:  $g1 \vdash xn \simeq xe1$ 
      using NarrowNode.hyps(1) NarrowNode.hyps(2)
      by auto
    then show ?case
    proof (cases xn = n')
      case True
        then have n:  $xe1 = e1'$  using c m repDet by simp
        then have ev:  $g2 \vdash n \simeq \text{UnaryExpr } (\text{UnaryNarrow inputBits resultBits}) e2'$ 
        using NarrowNode.hyps(1) l m n
          using NarrowNode.premis True d rep.NarrowNode by simp
        then have r:  $\text{UnaryExpr } (\text{UnaryNarrow inputBits resultBits}) e1' \geq \text{UnaryExpr } (\text{UnaryNarrow inputBits resultBits}) e2'$ 
          by (meson a mono-unary)
        then show ?thesis using ev r
          by (metis n)
      next
        case False
        have  $g1 \vdash xn \simeq xe1$  using m by simp
        have  $\exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
          using NarrowNode
          using False b encodes-contains l not-excluded-keep-type not-in-g singleton-iff
          by (metis node-eq-ternary NarrowNode)
        then have  $\exists xe2. (g2 \vdash n \simeq \text{UnaryExpr } (\text{UnaryNarrow inputBits resultBits}) xe2) \wedge \text{UnaryExpr } (\text{UnaryNarrow inputBits resultBits}) xe1 \geq \text{UnaryExpr } (\text{UnaryNarrow inputBits resultBits}) xe2$ 
          by (metis NarrowNode.premis l mono-unary rep.NarrowNode)
        then show ?thesis
          by meson
    qed
  next
    case (SignExtendNode n inputBits resultBits x xe1)
    have k:  $g1 \vdash n \simeq \text{UnaryExpr } (\text{UnarySignExtend inputBits resultBits}) xe1$ 
    using f SignExtendNode
      by (simp add: SignExtendNode.hyps(2) rep.SignExtendNode)
    obtain xn where l: kind g1 n = SignExtendNode inputBits resultBits xn
      using SignExtendNode.hyps(1) by blast
    then have m:  $g1 \vdash xn \simeq xe1$ 
      using SignExtendNode.hyps(1) SignExtendNode.hyps(2)
      by auto
    then show ?case

```

```

proof (cases  $xn = n'$ )
  case True
    then have  $n: xe1 = e1'$  using  $c\ m\ repDet$  by simp
    then have  $ev: g2 \vdash n \simeq UnaryExpr\ (UnarySignExtend\ inputBits\ resultBits)$ 
     $e2'$  using  $SignExtendNode.hyps(1)\ l\ m\ n$ 
      using  $SignExtendNode.premis\ True\ d\ rep.SignExtendNode$  by simp
      then have  $r: UnaryExpr\ (UnarySignExtend\ inputBits\ resultBits)\ e1' \geq$ 
 $UnaryExpr\ (UnarySignExtend\ inputBits\ resultBits)\ e2'$ 
        by (meson a mono-unary)
      then show ?thesis using  $ev\ r$ 
        by (metis n)
    next
      case False
        have  $g1 \vdash xn \simeq xe1$  using  $m$  by simp
        have  $\exists\ xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
          using  $SignExtendNode$ 
          using  $False\ b\ encodes-contains\ l\ not-excluded-keep-type\ not-in-g\ singleton-iff$ 
          by (metis-node-eq-ternary SignExtendNode)
        then have  $\exists\ xe2. (g2 \vdash n \simeq UnaryExpr\ (UnarySignExtend\ inputBits\ result-$ 
 $Bits)\ xe2) \wedge UnaryExpr\ (UnarySignExtend\ inputBits\ resultBits)\ xe1 \geq UnaryExpr$ 
 $(UnarySignExtend\ inputBits\ resultBits)\ xe2$ 
          by (metis SignExtendNode.premis l mono-unary rep.SignExtendNode)
        then show ?thesis
          by meson
      qed
    next
      case ( $ZeroExtendNode\ n\ inputBits\ resultBits\ x\ xe1$ )
        have  $k: g1 \vdash n \simeq UnaryExpr\ (UnaryZeroExtend\ inputBits\ resultBits)\ xe1$ 
using  $f\ ZeroExtendNode$ 
          by (simp add: ZeroExtendNode.hyps(2) rep.ZeroExtendNode)
        obtain  $xn$  where  $l: kind\ g1\ n = ZeroExtendNode\ inputBits\ resultBits\ xn$ 
          using  $ZeroExtendNode.hyps(1)$  by blast
        then have  $m: g1 \vdash xn \simeq xe1$ 
          using  $ZeroExtendNode.hyps(1)\ ZeroExtendNode.hyps(2)$ 
          by auto
        then show ?case
          proof (cases  $xn = n'$ )
            case True
              then have  $n: xe1 = e1'$  using  $c\ m\ repDet$  by simp
              then have  $ev: g2 \vdash n \simeq UnaryExpr\ (UnaryZeroExtend\ inputBits\ resultBits)$ 
               $e2'$  using  $ZeroExtendNode.hyps(1)\ l\ m\ n$ 
                using  $ZeroExtendNode.premis\ True\ d\ rep.ZeroExtendNode$  by simp
                then have  $r: UnaryExpr\ (UnaryZeroExtend\ inputBits\ resultBits)\ e1' \geq$ 
 $UnaryExpr\ (UnaryZeroExtend\ inputBits\ resultBits)\ e2'$ 
                  by (meson a mono-unary)
                then show ?thesis using  $ev\ r$ 
                  by (metis n)
            next
              case False

```

```

    have  $g1 \vdash xn \simeq xe1$  using  $m$  by simp
    have  $\exists xe2. (g2 \vdash xn \simeq xe2) \wedge xe1 \geq xe2$ 
      using ZeroExtendNode
      using False b encodes-contains l not-excluded-keep-type not-in-g singleton-iff
      by (metis-node-eq-ternary ZeroExtendNode)
    then have  $\exists xe2. (g2 \vdash n \simeq \text{UnaryExpr } (\text{UnaryZeroExtend inputBits resultBits}) xe2) \wedge \text{UnaryExpr } (\text{UnaryZeroExtend inputBits resultBits}) xe1 \geq \text{UnaryExpr } (\text{UnaryZeroExtend inputBits resultBits}) xe2$ 
      by (metis ZeroExtendNode.premis l mono-unary rep.ZeroExtendNode)
    then show ?thesis
      by meson
  qed
next
  case (LeafNode n s)
  then show ?case
    by (metis eq-refl rep.LeafNode)
next
  case (RefNode n')
  then show ?case
    by (metis a b c d no-encoding not-excluded-keep-type rep.RefNode repDet singletonD)
  qed
qed
qed

```

lemma *graph-antics-preservation-subscript*:

```

  assumes  $a: e_1' \geq e_2'$ 
  assumes  $b: (\{n\} \trianglelefteq \text{as-set } g_1) \subseteq \text{as-set } g_2$ 
  assumes  $c: g_1 \vdash n \simeq e_1'$ 
  assumes  $d: g_2 \vdash n \simeq e_2'$ 
  shows graph-refinement  $g_1 g_2$ 
  using graph-antics-preservation assms by simp

```

lemma *tree-to-graph-rewriting*:

```

   $e_1 \geq e_2$ 
   $\wedge (g_1 \vdash n \simeq e_1) \wedge \text{maximal-sharing } g_1$ 
   $\wedge (\{n\} \trianglelefteq \text{as-set } g_1) \subseteq \text{as-set } g_2$ 
   $\wedge (g_2 \vdash n \simeq e_2) \wedge \text{maximal-sharing } g_2$ 
 $\implies \text{graph-refinement } g_1 g_2$ 
  using graph-antics-preservation
  by auto

```

```

declare [[simp-trace]]
lemma equal-refines:
  fixes  $e1 e2 :: \text{IRExpr}$ 
  assumes  $e1 = e2$ 
  shows  $e1 \geq e2$ 
  using assms

```



```

  by simp
declare [[simp-trace=false]]

```

```

lemma eval-contains-id[simp]:  $g1 \vdash n \simeq e \implies n \in \text{ids } g1$ 
  using no-encoding by blast

```

```

lemma subset-kind[simp]:  $\text{as-set } g1 \subseteq \text{as-set } g2 \implies g1 \vdash n \simeq e \implies \text{kind } g1 \ n = \text{kind } g2 \ n$ 
  using eval-contains-id unfolding as-set-def
  by blast

```

```

lemma subset-stamp[simp]:  $\text{as-set } g1 \subseteq \text{as-set } g2 \implies g1 \vdash n \simeq e \implies \text{stamp } g1 \ n = \text{stamp } g2 \ n$ 
  using eval-contains-id unfolding as-set-def
  by blast

```

```

method solve-subset-eval uses as-set eval =
  (metis eval as-set subset-kind subset-stamp |
   metis eval as-set subset-kind)

```

```

lemma subset-implies-evals:
  assumes  $\text{as-set } g1 \subseteq \text{as-set } g2$ 
  assumes  $(g1 \vdash n \simeq e)$ 
  shows  $(g2 \vdash n \simeq e)$ 
  using assms(2)
  apply (induction e)
    apply (solve-subset-eval as-set: assms(1) eval: ConstantNode)
    apply (solve-subset-eval as-set: assms(1) eval: ParameterNode)
    apply (solve-subset-eval as-set: assms(1) eval: ConditionalNode)
    apply (solve-subset-eval as-set: assms(1) eval: AbsNode)
    apply (solve-subset-eval as-set: assms(1) eval: NotNode)
    apply (solve-subset-eval as-set: assms(1) eval: NegateNode)
    apply (solve-subset-eval as-set: assms(1) eval: LogicNegationNode)
    apply (solve-subset-eval as-set: assms(1) eval: AddNode)
    apply (solve-subset-eval as-set: assms(1) eval: MulNode)
    apply (solve-subset-eval as-set: assms(1) eval: SubNode)
    apply (solve-subset-eval as-set: assms(1) eval: AndNode)
    apply (solve-subset-eval as-set: assms(1) eval: OrNode)
    apply (solve-subset-eval as-set: assms(1) eval: XorNode)
    apply (solve-subset-eval as-set: assms(1) eval: ShortCircuitOrNode)
    apply (solve-subset-eval as-set: assms(1) eval: LeftShiftNode)
    apply (solve-subset-eval as-set: assms(1) eval: RightShiftNode)
    apply (solve-subset-eval as-set: assms(1) eval: UnsignedRightShiftNode)
    apply (solve-subset-eval as-set: assms(1) eval: IntegerBelowNode)
    apply (solve-subset-eval as-set: assms(1) eval: IntegerEqualsNode)
    apply (solve-subset-eval as-set: assms(1) eval: IntegerLessThanNode)

```

```

    apply (solve-subset-eval as-set: assms(1) eval: NarrowNode)
    apply (solve-subset-eval as-set: assms(1) eval: SignExtendNode)
    apply (solve-subset-eval as-set: assms(1) eval: ZeroExtendNode)
    apply (solve-subset-eval as-set: assms(1) eval: LeafNode)
    by (solve-subset-eval as-set: assms(1) eval: RefNode)

lemma subset-refines:
  assumes as-set g1  $\subseteq$  as-set g2
  shows graph-refinement g1 g2
proof -
  have ids g1  $\subseteq$  ids g2 using assms unfolding as-set-def
  by blast
  then show ?thesis unfolding graph-refinement-def apply rule
  apply (rule allI) apply (rule impI) apply (rule allI) apply (rule impI)
  unfolding graph-represents-expression-def
  proof -
    fix n e1
    assume 1:n  $\in$  ids g1
    assume 2:g1  $\vdash$  n  $\simeq$  e1

    show  $\exists e2. (g2 \vdash n \simeq e2) \wedge e1 \geq e2$ 
    using assms 1 2 using subset-implies-evals
    by (meson equal-refines)
  qed
qed

```

```

lemma graph-construction:
  e1  $\geq$  e2
   $\wedge$  as-set g1  $\subseteq$  as-set g2
   $\wedge$  (g2  $\vdash$  n  $\simeq$  e2)
   $\implies$  (g2  $\vdash$  n  $\trianglelefteq$  e1)  $\wedge$  graph-refinement g1 g2
  using subset-refines
  by (meson encodeeval-def graph-represents-expression-def le-expr-def)

```

7.8.4 Term Graph Reconstruction

```

lemma find-exists-kind:
  assumes find-node-and-stamp g (node, s) = Some nid
  shows kind g nid = node
  using assms unfolding find-node-and-stamp.simps
  by (metis (mono-tags, lifting) find-Some-iff)

```

```

lemma find-exists-stamp:
  assumes find-node-and-stamp g (node, s) = Some nid
  shows stamp g nid = s
  using assms unfolding find-node-and-stamp.simps
  by (metis (mono-tags, lifting) find-Some-iff)

```

```

lemma find-new-kind:

```

```

assumes  $g' = \text{add-node } \text{nid} \ (\text{node}, s) \ g$ 
assumes  $\text{node} \neq \text{NoNode}$ 
shows  $\text{kind } g' \ \text{nid} = \text{node}$ 
using assms
using add-node-lookup by presburger

lemma find-new-stamp:
assumes  $g' = \text{add-node } \text{nid} \ (\text{node}, s) \ g$ 
assumes  $\text{node} \neq \text{NoNode}$ 
shows  $\text{stamp } g' \ \text{nid} = s$ 
using assms
using add-node-lookup by presburger

lemma sorted-bottom:
assumes finite xs
assumes  $x \in xs$ 
shows  $x \leq \text{last}(\text{sorted-list-of-set}(xs::\text{nat set}))$ 
using assms
using sorted2-simps(2) sorted-list-of-set(2)
by (smt (verit, del-insts) Diff-iff Max-ge Max-in empty-iff list.set(1) snoc-eq-iff-butlast
sorted-insort-is-snoc sorted-list-of-set(1) sorted-list-of-set.fold-insort-key.infinite sorted-list-of-set.fold-insort-ke)

lemma fresh:  $\text{finite } xs \implies \text{last}(\text{sorted-list-of-set}(xs::\text{nat set})) + 1 \notin xs$ 
using sorted-bottom
using not-le by auto

lemma fresh-ids:
assumes  $n = \text{get-fresh-id } g$ 
shows  $n \notin \text{ids } g$ 
proof –
  have finite (ids g) using Rep-IRGraph by auto
  then show ?thesis
    using assms fresh unfolding get-fresh-id.simps
    by blast
qed

lemma graph-unchanged-rep-unchanged:
assumes  $\forall n \in \text{ids } g. \text{kind } g \ n = \text{kind } g' \ n$ 
assumes  $\forall n \in \text{ids } g. \text{stamp } g \ n = \text{stamp } g' \ n$ 
shows  $(g \vdash n \simeq e) \longrightarrow (g' \vdash n \simeq e)$ 
apply (rule impI) subgoal premises e using e assms
  apply (induction n e)
    apply (metis no-encoding rep.ConstantNode)
    apply (metis no-encoding rep.ParameterNode)
    apply (metis no-encoding rep.ConditionalNode)
    apply (metis no-encoding rep.AbsNode)
    apply (metis no-encoding rep.NotNode)
    apply (metis no-encoding rep.NegateNode)
    apply (metis no-encoding rep.LogicNegationNode)

```

```

      apply (metis no-encoding rep.AddNode)
      apply (metis no-encoding rep.MulNode)
      apply (metis no-encoding rep.SubNode)
      apply (metis no-encoding rep.AndNode)
      apply (metis no-encoding rep.OrNode)
      apply (metis no-encoding rep.XorNode)
      apply (metis no-encoding rep.ShortCircuitOrNode)
      apply (metis no-encoding rep.LeftShiftNode)
      apply (metis no-encoding rep.RightShiftNode)
      apply (metis no-encoding rep.UnsignedRightShiftNode)
      apply (metis no-encoding rep.IntegerBelowNode)
      apply (metis no-encoding rep.IntegerEqualsNode)
      apply (metis no-encoding rep.IntegerLessThanNode)
      apply (metis no-encoding rep.NarrowNode)
      apply (metis no-encoding rep.SignExtendNode)
      apply (metis no-encoding rep.ZeroExtendNode)
      apply (metis no-encoding rep.LeafNode)
    by (metis no-encoding rep.RefNode)
  done

```

lemma *fresh-node-subset*:

```

  assumes  $n \notin \text{ids } g$ 
  assumes  $g' = \text{add-node } n (k, s) g$ 
  shows  $\text{as-set } g \subseteq \text{as-set } g'$ 
  using assms
  by (smt (verit, del-Insts) Collect-mono-iff Diff-idemp Diff-insert-absorb add-changed
    as-set-def disjoint-change unchanged.simps)

```

lemma *unrep-subset*:

```

  assumes  $(g \oplus e \rightsquigarrow (g', n))$ 
  shows  $\text{as-set } g \subseteq \text{as-set } g'$ 
  using assms proof (induction  $g e (g', n)$  arbitrary:  $g' n$ )
  case (ConstantNodeSame  $g c n$ )
  then show ?case by blast
next
  case (ConstantNodeNew  $g c n g'$ )
  then show ?case using fresh-ids fresh-node-subset
    by presburger
next
  case (ParameterNodeSame  $g i s n$ )
  then show ?case by blast
next
  case (ParameterNodeNew  $g i s n g'$ )
  then show ?case using fresh-ids fresh-node-subset
    by presburger
next
  case (ConditionalNodeSame  $g ce g2 c te g3 t fe g4 f s' n$ )
  then show ?case by blast
next

```

```

    case (ConditionalNodeNew g ce g2 c te g3 t fe g4 f s' n g')
    then show ?case using fresh-ids fresh-node-subset
      by (meson subset-trans)
next
    case (UnaryNodeSame g xe g2 x s' op n)
    then show ?case by blast
next
    case (UnaryNodeNew g xe g2 x s' op n g')
    then show ?case using fresh-ids fresh-node-subset
      by (meson subset-trans)
next
    case (BinaryNodeSame g xe g2 x ye g3 y s' op n)
    then show ?case by blast
next
    case (BinaryNodeNew g xe g2 x ye g3 y s' op n g')
    then show ?case using fresh-ids fresh-node-subset
      by (meson subset-trans)
next
    case (AllLeafNodes g n s)
    then show ?case by blast
qed

```

lemma *fresh-node-preserves-other-nodes*:

```

  assumes n' = get-fresh-id g
  assumes g' = add-node n' (k, s) g
  shows  $\forall n \in \text{ids } g. (g \vdash n \simeq e) \longrightarrow (g' \vdash n \simeq e)$ 
  using assms
  by (smt (verit, ccfv-SIG) Diff-idemp Diff-insert-absorb add-changed disjoint-change
    fresh-ids graph-unchanged-rep-unchanged unchanged.elims(2))

```

lemma *found-node-preserves-other-nodes*:

```

  assumes find-node-and-stamp g (k, s) = Some n
  shows  $\forall n \in \text{ids } g. (g \vdash n \simeq e) \longleftrightarrow (g \vdash n \simeq e)$ 
  using assms
  by blast

```

lemma *unrep-ids-subset[simp]*:

```

  assumes  $g \oplus e \rightsquigarrow (g', n)$ 
  shows  $\text{ids } g \subseteq \text{ids } g'$ 
  using assms unrep-subset
  by (meson graph-refinement-def subset-refines)

```

lemma *unrep-unchanged*:

```

  assumes  $g \oplus e \rightsquigarrow (g', n)$ 
  shows  $\forall n \in \text{ids } g. \forall e. (g \vdash n \simeq e) \longrightarrow (g' \vdash n \simeq e)$ 
  using assms unrep-subset fresh-node-preserves-other-nodes
  by (meson subset-implies-evals)

```

theorem *term-graph-reconstruction*:

$g \oplus e \rightsquigarrow (g', n) \implies (g' \vdash n \simeq e) \wedge \text{as-set } g \subseteq \text{as-set } g'$
subgoal premises e **apply** (rule *conjI*) **defer**
 using e *unrep-subset* **apply** *blast* **using** e
proof (induction g e (g', n) arbitrary: $g' n$)
 case (*ConstantNodeSame* $g' c n$)
 then have $\text{kind } g' n = \text{ConstantNode } c$
 using *find-exists-kind local.ConstantNodeSame* **by** *blast*
 then show $?case$ **using** *ConstantNode* **by** *blast*
next
 case (*ConstantNodeNew* $g c$)
 then show $?case$
 using *ConstantNode IRNode.distinct(683) add-node-lookup* **by** *presburger*
next
 case (*ParameterNodeSame* $i s$)
 then show $?case$
 by (*metis ParameterNode find-exists-kind find-exists-stamp*)
next
 case (*ParameterNodeNew* $g i s$)
 then show $?case$
 by (*metis IRNode.distinct(2447) ParameterNode add-node-lookup*)
next
 case (*ConditionalNodeSame* $g_4 c t f s' n g ce g_2 te g_3 fe$)
 then have $k: \text{kind } g_4 n = \text{ConditionalNode } c t f$
 using *find-exists-kind* **by** *blast*
 have $c: g_4 \vdash c \simeq ce$ **using** *local.ConditionalNodeSame unrep-unchanged*
 using *no-encoding* **by** *blast*
 have $t: g_4 \vdash t \simeq te$ **using** *local.ConditionalNodeSame unrep-unchanged*
 using *no-encoding* **by** *blast*
 have $f: g_4 \vdash f \simeq fe$ **using** *local.ConditionalNodeSame unrep-unchanged*
 using *no-encoding* **by** *blast*
 then show $?case$ **using** $c t f$
 using *ConditionalNode k* **by** *blast*
next
 case (*ConditionalNodeNew* $g_4 c t f s' g ce g_2 te g_3 fe n g'$)
 moreover have *ConditionalNode c t f* $\neq \text{NoNode}$
 using *unary-node.elims* **by** *blast*
 ultimately have $k: \text{kind } g' n = \text{ConditionalNode } c t f$
 using *find-new-kind local.ConditionalNodeNew*
 by *presburger*
 then have $c: g' \vdash c \simeq ce$ **using** *local.ConditionalNodeNew unrep-unchanged*
 using *no-encoding*
 by (*metis ConditionalNodeNew.hyps(9) fresh-node-preserves-other-nodes*)
 then have $t: g' \vdash t \simeq te$ **using** *local.ConditionalNodeNew unrep-unchanged*
 using *no-encoding fresh-node-preserves-other-nodes*
 by *metis*
 then have $f: g' \vdash f \simeq fe$ **using** *local.ConditionalNodeNew unrep-unchanged*
 using *no-encoding fresh-node-preserves-other-nodes*
 by *metis*
 then show $?case$ **using** $c t f$

```

    using ConditionalNode k by blast
next
case (UnaryNodeSame g' op x s' n g xe)
then have k: kind g' n = unary-node op x
    using find-exists-kind local.UnaryNodeSame by blast
then have g' ⊢ x ≃ xe using local.UnaryNodeSame by blast
then show ?case using k
    apply (cases op)
    using AbsNode unary-node.simps(1) apply presburger
    using NegateNode unary-node.simps(3) apply presburger
    using NotNode unary-node.simps(2) apply presburger
    using LogicNegationNode unary-node.simps(4) apply presburger
    using NarrowNode unary-node.simps(5) apply presburger
    using SignExtendNode unary-node.simps(6) apply presburger
    using ZeroExtendNode unary-node.simps(7) by presburger
next
case (UnaryNodeNew g2 op x s' g xe n g')
moreover have unary-node op x ≠ NoNode
    using unary-node.elims by blast
ultimately have k: kind g' n = unary-node op x
    using find-new-kind local.UnaryNodeNew
    by presburger
have x ∈ ids g2 using local.UnaryNodeNew
    using eval-contains-id by blast
then have x ≠ n using local.UnaryNodeNew(5) fresh-ids by blast
have g' ⊢ x ≃ xe using local.UnaryNodeNew fresh-node-preserves-other-nodes
    using ⟨x ∈ ids g2⟩ by blast
then show ?case using k
    apply (cases op)
    using AbsNode unary-node.simps(1) apply presburger
    using NegateNode unary-node.simps(3) apply presburger
    using NotNode unary-node.simps(2) apply presburger
    using LogicNegationNode unary-node.simps(4) apply presburger
    using NarrowNode unary-node.simps(5) apply presburger
    using SignExtendNode unary-node.simps(6) apply presburger
    using ZeroExtendNode unary-node.simps(7) by presburger
next
case (BinaryNodeSame g3 op x y s' n g xe g2 ye)
then have k: kind g3 n = bin-node op x y
    using find-exists-kind by blast
have x: g3 ⊢ x ≃ xe using local.BinaryNodeSame unrep-unchanged
    using no-encoding by blast
have y: g3 ⊢ y ≃ ye using local.BinaryNodeSame unrep-unchanged
    using no-encoding by blast
then show ?case using x y k apply (cases op)
    using AddNode bin-node.simps(1) apply presburger
    using MulNode bin-node.simps(2) apply presburger
    using SubNode bin-node.simps(3) apply presburger
    using AndNode bin-node.simps(4) apply presburger

```

```

    using OrNode bin-node.simps(5) apply presburger
    using XorNode bin-node.simps(6) apply presburger
    using ShortCircuitOrNode bin-node.simps(7) apply presburger
    using LeftShiftNode bin-node.simps(8) apply presburger
    using RightShiftNode bin-node.simps(9) apply presburger
    using UnsignedRightShiftNode bin-node.simps(10) apply presburger
    using IntegerEqualsNode bin-node.simps(11) apply presburger
    using IntegerLessThanNode bin-node.simps(12) apply presburger
    using IntegerBelowNode bin-node.simps(13) by presburger
next
case (BinaryNodeNew g3 op x y s' g xe g2 ye n g')
moreover have bin-node op x y  $\neq$  NoNode
  using bin-node.elims by blast
ultimately have k: kind g' n = bin-node op x y
  using find-new-kind local.BinaryNodeNew
  by presburger
then have k: kind g' n = bin-node op x y
  using find-exists-kind by blast
have x: g'  $\vdash$  x  $\simeq$  xe using local.BinaryNodeNew unrep-unchanged
  using no-encoding
  by (meson fresh-node-preserves-other-nodes)
have y: g'  $\vdash$  y  $\simeq$  ye using local.BinaryNodeNew unrep-unchanged
  using no-encoding
  by (meson fresh-node-preserves-other-nodes)
then show ?case using x y k apply (cases op)
  using AddNode bin-node.simps(1) apply presburger
  using MulNode bin-node.simps(2) apply presburger
  using SubNode bin-node.simps(3) apply presburger
  using AndNode bin-node.simps(4) apply presburger
  using OrNode bin-node.simps(5) apply presburger
  using XorNode bin-node.simps(6) apply presburger
  using ShortCircuitOrNode bin-node.simps(7) apply presburger
  using LeftShiftNode bin-node.simps(8) apply presburger
  using RightShiftNode bin-node.simps(9) apply presburger
  using UnsignedRightShiftNode bin-node.simps(10) apply presburger
  using IntegerEqualsNode bin-node.simps(11) apply presburger
  using IntegerLessThanNode bin-node.simps(12) apply presburger
  using IntegerBelowNode bin-node.simps(13) by presburger
next
case (AllLeafNodes g n s)
then show ?case using rep.LeafNode by blast
qed
done

```

lemma *ref-refinement*:

```

assumes g  $\vdash$  n  $\simeq$  e1
assumes kind g n' = RefNode n
shows g  $\vdash$  n'  $\trianglelefteq$  e1
using assms RefNode

```



```

by (meson equal-refines graph-represents-expression-def)

lemma unrep-refines:
  assumes  $g \oplus e \rightsquigarrow (g', n)$ 
  shows graph-refinement  $g$   $g'$ 
  using assms
  using graph-refinement-def subset-refines unrep-subset by blast

lemma add-new-node-refines:
  assumes  $n \notin \text{ids } g$ 
  assumes  $g' = \text{add-node } n \ (k, s) \ g$ 
  shows graph-refinement  $g$   $g'$ 
  using assms unfolding graph-refinement
  using fresh-node-subset subset-refines by presburger

lemma add-node-as-set:
  assumes  $g' = \text{add-node } n \ (k, s) \ g$ 
  shows  $(\{n\} \sqsubseteq \text{as-set } g) \subseteq \text{as-set } g'$ 
  using assms unfolding as-set-def domain-subtraction-def
  using add-changed
  by (smt (z3) case-prodE changeonly.simps mem-Collect-eq prod.sel(1) subsetI)

theorem refined-insert:
  assumes  $e_1 \geq e_2$ 
  assumes  $g_1 \oplus e_2 \rightsquigarrow (g_2, n')$ 
  shows  $(g_2 \vdash n' \sqsubseteq e_1) \wedge \text{graph-refinement } g_1 \ g_2$ 
  using assms
  using graph-construction term-graph-reconstruction by blast

lemma ids-finite: finite (ids  $g$ )
  using Rep-IRGraph ids.rep-eq by simp

lemma unwrap-sorted: set (sorted-list-of-set (ids  $g$ )) = ids  $g$ 
  using Rep-IRGraph set-sorted-list-of-set ids-finite
  by blast

lemma find-none:
  assumes find-node-and-stamp  $g \ (k, s) = \text{None}$ 
  shows  $\forall n \in \text{ids } g. \text{kind } g \ n \neq k \vee \text{stamp } g \ n \neq s$ 
proof -
  have  $(\nexists n. n \in \text{ids } g \wedge (\text{kind } g \ n = k \wedge \text{stamp } g \ n = s))$ 
    using assms unfolding find-node-and-stamp.simps using find-None-iff un-
  wrap-sorted
    by (metis (mono-tags, lifting))
  then show ?thesis
    by blast
qed

```

```

method ref-represents uses node =
  (metis IRNode.distinct(2755) RefNode dual-order.refl find-new-kind fresh-node-subset
node subset-implies-evals)

```

7.8.5 Data-flow Tree to Subgraph Preserves Maximal Sharing

```

lemma same-kind-stamp-encodes-equal:
  assumes kind g n = kind g n'
  assumes stamp g n = stamp g n'
  assumes  $\neg(\text{is-preevaluated } (\text{kind } g \ n))$ 
  shows  $\forall \ e. (g \vdash n \simeq e) \longrightarrow (g \vdash n' \simeq e)$ 
  apply (rule allI)
  subgoal for e
    apply (rule impI)
    subgoal premises eval using eval assms
      apply (induction e)
    using ConstantNode apply presburger
    using ParameterNode apply presburger
      apply (metis ConditionalNode)
      apply (metis AbsNode)
      apply (metis NotNode)
      apply (metis NegateNode)
      apply (metis LogicNegationNode)
      apply (metis AddNode)
      apply (metis MulNode)
      apply (metis SubNode)
      apply (metis AndNode)
      apply (metis OrNode)
      apply (metis XorNode)
      apply (metis ShortCircuitOrNode)
      apply (metis LeftShiftNode)
      apply (metis RightShiftNode)
      apply (metis UnsignedRightShiftNode)
      apply (metis IntegerBelowNode)
      apply (metis IntegerEqualsNode)
      apply (metis IntegerLessThanNode)
      apply (metis NarrowNode)

```

```

    apply (metis SignExtendNode)
    apply (metis ZeroExtendNode)
  defer
    apply (metis RefNode)
  by blast
done
done

```

lemma *new-node-not-present*:

```

assumes find-node-and-stamp g (node, s) = None
assumes n = get-fresh-id g
assumes g' = add-node n (node, s) g
shows  $\forall n' \in \text{true-ids } g. (\forall e. ((g \vdash n \simeq e) \wedge (g \vdash n' \simeq e)) \longrightarrow n = n')$ 
using assms
using encode-in-ids fresh-ids by blast

```

lemma *true-ids-def*:

```

true-ids g = {n ∈ ids g. ¬(is-RefNode (kind g n)) ∧ ((kind g n) ≠ NoNode)}
unfolding true-ids-def ids-def
using ids-def is-RefNode-def by fastforce

```

lemma *add-node-some-node-def*:

```

assumes k ≠ NoNode
assumes g' = add-node nid (k, s) g
shows g' = Abs-IRGraph ((Rep-IRGraph g)(nid ↦ (k, s)))
using assms
by (metis Rep-IRGraph-inverse add-node.rep-eq fst-conv)

```

lemma *ids-add-update-v1*:

```

assumes g' = add-node nid (k, s) g
assumes k ≠ NoNode
shows dom (Rep-IRGraph g') = dom (Rep-IRGraph g) ∪ {nid}
using assms ids.rep-eq add-node-some-node-def
by (simp add: add-node.rep-eq)

```

lemma *ids-add-update-v2*:

```

assumes g' = add-node nid (k, s) g
assumes k ≠ NoNode
shows nid ∈ ids g'
using assms
using find-new-kind ids-some by presburger

```

lemma *add-node-ids-subset*:

```

assumes n ∈ ids g
assumes g' = add-node n node g
shows ids g' = ids g ∪ {n}
using assms unfolding add-node-def
apply (cases fst node = NoNode)
using ids.rep-eq replace-node.rep-eq replace-node-def apply auto[1]

```

unfolding *ids-def*
by (*smt* (*verit*, *best*) *Collect-cong Un-insert-right dom-fun-upd fst-conv fun-upd-apply*
ids.rep-eq ids-def insert-absorb mem-Collect-eq option.inject option.simps(3) re-
place-node.rep-eq replace-node-def sup-bot.right-neutral)

lemma *convert-maximal:*

assumes $\forall n n'. n \in \text{true-ids } g \wedge n' \in \text{true-ids } g \longrightarrow (\forall e e'. (g \vdash n \simeq e) \wedge (g \vdash n' \simeq e') \longrightarrow e \neq e')$
shows *maximal-sharing* *g*
using *assms*
using *maximal-sharing* **by** *blast*

lemma *add-node-set-eq:*

assumes $k \neq \text{NoNode}$
assumes $n \notin \text{ids } g$
shows $\text{as-set } (\text{add-node } n (k, s) g) = \text{as-set } g \cup \{(n, (k, s))\}$
using *assms* **unfolding** *as-set-def add-node-def* **apply** *transfer* **apply** *simp*
by *blast*

lemma *add-node-as-set-eq:*

assumes $g' = \text{add-node } n (k, s) g$
assumes $n \notin \text{ids } g$
shows $\{n\} \sqsubseteq \text{as-set } g' = \text{as-set } g$
using *assms* **unfolding** *domain-subtraction-def*
using *add-node-set-eq*
by (*smt* (*z3*) *Collect-cong Rep-IRGraph-inverse UnCI UnE add-node.rep-eq as-set-def*
case-prodE2 case-prodI2 le-boolE le-boolI' mem-Collect-eq prod.sel(1) singletonD
singletonI)

lemma *true-ids:*

$\text{true-ids } g = \text{ids } g - \{n \in \text{ids } g. \text{is-RefNode } (\text{kind } g \ n)\}$
unfolding *true-ids-def*
by *fastforce*

lemma *as-set-ids:*

assumes $\text{as-set } g = \text{as-set } g'$
shows $\text{ids } g = \text{ids } g'$
using *assms*
by (*metis antisym equalityD1 graph-refinement-def subset-refines*)

lemma *ids-add-update:*

assumes $k \neq \text{NoNode}$
assumes $n \notin \text{ids } g$
assumes $g' = \text{add-node } n (k, s) g$
shows $\text{ids } g' = \text{ids } g \cup \{n\}$
using *assms* **apply** (*subst* *assms(3)*) **using** *add-node-set-eq as-set-ids*
by (*smt* (*verit*, *del-insts*) *Collect-cong Diff-idemp Diff-insert-absorb Un-commute*
add-node.rep-eq add-node-def ids.rep-eq ids-add-update-v1 ids-add-update-v2 insertE
insert-Collect insert-is-Un map-upd-Some-unfold mem-Collect-eq replace-node-def)

replace-node-unchanged)

lemma *true-ids-add-update*:

assumes $k \neq \text{NoNode}$
assumes $n \notin \text{ids } g$
assumes $g' = \text{add-node } n \ (k, s) \ g$
assumes $\neg(\text{is-RefNode } k)$
shows $\text{true-ids } g' = \text{true-ids } g \cup \{n\}$
using *assms* **using** *true-ids ids-add-update*
by (*smt* (*z3*) *Collect-cong Diff-iff Diff-insert-absorb Un-commute add-node-def*
find-new-kind insert-Diff-if insert-is-Un mem-Collect-eq replace-node-def replace-node-unchanged)

lemma *new-def*:

assumes $(\text{new} \sqsubseteq \text{as-set } g') = \text{as-set } g$
shows $n \in \text{ids } g \longrightarrow n \notin \text{new}$
using *assms*
by (*smt* (*z3*) *as-set-def case-prodD domain-subtraction-def mem-Collect-eq*)

lemma *add-preserves-rep*:

assumes *unchanged*: $(\text{new} \sqsubseteq \text{as-set } g') = \text{as-set } g$
assumes *closed*: *wf-closed* *g*
assumes *existed*: $n \in \text{ids } g$
assumes $g' \vdash n \simeq e$
shows $g \vdash n \simeq e$
proof (*cases* $n \in \text{new}$)
case *True*
have $n \notin \text{ids } g$
using *unchanged True unfolding as-set-def domain-subtraction-def*
by *blast*
then show *?thesis* **using** *existed by simp*
next
case *False*
then have *kind-eq*: $\forall n'. n' \notin \text{new} \longrightarrow \text{kind } g \ n' = \text{kind } g' \ n'$
— can be more general than *stamp_eq* because *NoNode* default is equal
using *unchanged not-excluded-keep-type*
by (*smt* (*z3*) *case-prodE domain-subtraction-def ids-some mem-Collect-eq subsetI*)
from *False* **have** *stamp-eq*: $\forall n' \in \text{ids } g'. n' \notin \text{new} \longrightarrow \text{stamp } g \ n' = \text{stamp } g' \ n'$
using *unchanged not-excluded-keep-type*
by (*metis equalityE*)
show *?thesis* **using** *assms(4) kind-eq stamp-eq False*
proof (*induction* *n e* *rule: rep.induct*)
case (*ConstantNode* *n c*)
then show *?case*
using *rep.ConstantNode kind-eq by presburger*
next

```

    case (ParameterNode n i s)
    then show ?case
      using rep.ParameterNode
      by (metis no-encoding)
  next
    case (ConditionalNode n c t f ce te fe)
    have kind: kind g n = ConditionalNode c t f
      using ConditionalNode.hyps(1) ConditionalNode.prem(3) kind-eq by pres-
    burger
    then have isin: n ∈ ids g
      by simp
    have inputs: {c, t, f} = inputs g n
      using kind unfolding inputs.simps using inputs-of-ConditionalNode by simp
    have c ∈ ids g ∧ t ∈ ids g ∧ f ∈ ids g
      using closed unfolding wf-closed-def
      using isin inputs by blast
    then have c ∉ new ∧ t ∉ new ∧ f ∉ new
      using new-def unchanged by blast
    then show ?case using ConditionalNode apply simp
      using rep.ConditionalNode by presburger
  next
    case (AbsNode n x xe)
    then have kind: kind g n = AbsNode x
      by simp
    then have isin: n ∈ ids g
      by simp
    have inputs: {x} = inputs g n
      using kind unfolding inputs.simps by simp
    have x ∈ ids g
      using closed unfolding wf-closed-def
      using isin inputs by blast
    then have x ∉ new
      using new-def unchanged by blast
    then show ?case
      using AbsNode
      using rep.AbsNode by presburger
  next
    case (NotNode n x xe)
    then have kind: kind g n = NotNode x
      by simp
    then have isin: n ∈ ids g
      by simp
    have inputs: {x} = inputs g n
      using kind unfolding inputs.simps by simp
    have x ∈ ids g
      using closed unfolding wf-closed-def
      using isin inputs by blast
    then have x ∉ new
      using new-def unchanged by blast

```

```

    then show ?case using NotNode
      using rep.NotNode by presburger
next
  case (NegateNode n x xe)
  then have kind: kind g n = NegateNode x
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using NegateNode
    using rep.NegateNode by presburger
next
  case (LogicNegationNode n x xe)
  then have kind: kind g n = LogicNegationNode x
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using LogicNegationNode
    using rep.LogicNegationNode by presburger
next
  case (AddNode n x y xe ye)
  then have kind: kind g n = AddNode x y
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x, y\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g \wedge y \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new} \wedge y \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using AddNode
    using rep.AddNode by presburger
next
  case (MulNode n x y xe ye)

```

```

then have kind: kind g n = MulNode x y
  by simp
then have isin: n ∈ ids g
  by simp
have inputs: {x, y} = inputs g n
  using kind unfolding inputs.simps by simp
have x ∈ ids g ∧ y ∈ ids g
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have x ∉ new ∧ y ∉ new
  using new-def unchanged by blast
then show ?case using MulNode
  using rep.MulNode by presburger
next
case (SubNode n x y xe ye)
then have kind: kind g n = SubNode x y
  by simp
then have isin: n ∈ ids g
  by simp
have inputs: {x, y} = inputs g n
  using kind unfolding inputs.simps by simp
have x ∈ ids g ∧ y ∈ ids g
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have x ∉ new ∧ y ∉ new
  using new-def unchanged by blast
then show ?case using SubNode
  using rep.SubNode by presburger
next
case (AndNode n x y xe ye)
then have kind: kind g n = AndNode x y
  by simp
then have isin: n ∈ ids g
  by simp
have inputs: {x, y} = inputs g n
  using kind unfolding inputs.simps by simp
have x ∈ ids g ∧ y ∈ ids g
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have x ∉ new ∧ y ∉ new
  using new-def unchanged by blast
then show ?case using AndNode
  using rep.AndNode by presburger
next
case (OrNode n x y xe ye)
then have kind: kind g n = OrNode x y
  by simp
then have isin: n ∈ ids g
  by simp

```



```

have inputs: {x, y} = inputs g n
  using kind unfolding inputs.simps by simp
have x ∈ ids g ∧ y ∈ ids g
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have x ∉ new ∧ y ∉ new
  using new-def unchanged by blast
then show ?case using OrNode
  using rep.OrNode by presburger
next
case (XorNode n x y xe ye)
then have kind: kind g n = XorNode x y
  by simp
then have isin: n ∈ ids g
  by simp
have inputs: {x, y} = inputs g n
  using kind unfolding inputs.simps by simp
have x ∈ ids g ∧ y ∈ ids g
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have x ∉ new ∧ y ∉ new
  using new-def unchanged by blast
then show ?case using XorNode
  using rep.XorNode by presburger
next
case (ShortCircuitOrNode n x y xe ye)
then have kind: kind g n = ShortCircuitOrNode x y
  by simp
then have isin: n ∈ ids g
  by simp
have inputs: {x, y} = inputs g n
  using kind unfolding inputs.simps by simp
have x ∈ ids g ∧ y ∈ ids g
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have x ∉ new ∧ y ∉ new
  using new-def unchanged by blast
then show ?case using ShortCircuitOrNode
  using rep.ShortCircuitOrNode by presburger
next
case (LeftShiftNode n x y xe ye)
then have kind: kind g n = LeftShiftNode x y
  by simp
then have isin: n ∈ ids g
  by simp
have inputs: {x, y} = inputs g n
  using kind unfolding inputs.simps by simp
have x ∈ ids g ∧ y ∈ ids g
  using closed unfolding wf-closed-def

```

```

    using isin inputs by blast
  then have  $x \notin \text{new} \wedge y \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using LeftShiftNode
    using rep.LeftShiftNode by presburger
next
case (RightShiftNode  $n\ x\ y\ xe\ ye$ )
then have kind:  $\text{kind } g\ n = \text{RightShiftNode } x\ y$ 
  by simp
then have isin:  $n \in \text{ids } g$ 
  by simp
have inputs:  $\{x, y\} = \text{inputs } g\ n$ 
  using kind unfolding inputs.simps by simp
have  $x \in \text{ids } g \wedge y \in \text{ids } g$ 
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have  $x \notin \text{new} \wedge y \notin \text{new}$ 
  using new-def unchanged by blast
then show ?case using RightShiftNode
  using rep.RightShiftNode by presburger
next
case (UnsignedRightShiftNode  $n\ x\ y\ xe\ ye$ )
then have kind:  $\text{kind } g\ n = \text{UnsignedRightShiftNode } x\ y$ 
  by simp
then have isin:  $n \in \text{ids } g$ 
  by simp
have inputs:  $\{x, y\} = \text{inputs } g\ n$ 
  using kind unfolding inputs.simps by simp
have  $x \in \text{ids } g \wedge y \in \text{ids } g$ 
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have  $x \notin \text{new} \wedge y \notin \text{new}$ 
  using new-def unchanged by blast
then show ?case using UnsignedRightShiftNode
  using rep.UnsignedRightShiftNode by presburger
next
case (IntegerBelowNode  $n\ x\ y\ xe\ ye$ )
then have kind:  $\text{kind } g\ n = \text{IntegerBelowNode } x\ y$ 
  by simp
then have isin:  $n \in \text{ids } g$ 
  by simp
have inputs:  $\{x, y\} = \text{inputs } g\ n$ 
  using kind unfolding inputs.simps by simp
have  $x \in \text{ids } g \wedge y \in \text{ids } g$ 
  using closed unfolding wf-closed-def
  using isin inputs by blast
then have  $x \notin \text{new} \wedge y \notin \text{new}$ 
  using new-def unchanged by blast
then show ?case using IntegerBelowNode

```

```

    using rep.IntegerBelowNode by presburger
next
  case (IntegerEqualsNode n x y xe ye)
  then have kind: kind g n = IntegerEqualsNode x y
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x, y\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g \wedge y \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new} \wedge y \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using IntegerEqualsNode
    using rep.IntegerEqualsNode by presburger
next
  case (IntegerLessThanNode n x y xe ye)
  then have kind: kind g n = IntegerLessThanNode x y
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x, y\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g \wedge y \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new} \wedge y \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using IntegerLessThanNode
    using rep.IntegerLessThanNode by presburger
next
  case (NarrowNode n inputBits resultBits x xe)
  then have kind: kind g n = NarrowNode inputBits resultBits x
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using NarrowNode
    using rep.NarrowNode by presburger
next
  case (SignExtendNode n inputBits resultBits x xe)
  then have kind: kind g n = SignExtendNode inputBits resultBits x

```

```

    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using SignExtendNode
    using rep.SignExtendNode by presburger
next
case (ZeroExtendNode n inputBits resultBits x xe)
  then have kind:  $\text{kind } g \ n = \text{ZeroExtendNode inputBits resultBits } x$ 
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{x\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $x \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $x \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case using ZeroExtendNode
    using rep.ZeroExtendNode by presburger
next
case (LeafNode n s)
  then show ?case
    by (metis no-encoding rep.LeafNode)
next
case (RefNode n n' e)
  then have kind:  $\text{kind } g \ n = \text{RefNode } n'$ 
    by simp
  then have isin:  $n \in \text{ids } g$ 
    by simp
  have inputs:  $\{n'\} = \text{inputs } g \ n$ 
    using kind unfolding inputs.simps by simp
  have  $n' \in \text{ids } g$ 
    using closed unfolding wf-closed-def
    using isin inputs by blast
  then have  $n' \notin \text{new}$ 
    using new-def unchanged by blast
  then show ?case
    using RefNode
    using rep.RefNode by presburger
qed
qed

```

```

lemma not-in-no-rep:
   $n \notin \text{ids } g \implies \forall e. \neg(g \vdash n \simeq e)$ 
  using eval-contains-id by blast

lemma unary-inputs:
  assumes  $\text{kind } g \ n = \text{unary-node } op \ x$ 
  shows  $\text{inputs } g \ n = \{x\}$ 
  using assms by (cases op; auto)

lemma unary-succ:
  assumes  $\text{kind } g \ n = \text{unary-node } op \ x$ 
  shows  $\text{succ } g \ n = \{\}$ 
  using assms by (cases op; auto)

lemma binary-inputs:
  assumes  $\text{kind } g \ n = \text{bin-node } op \ x \ y$ 
  shows  $\text{inputs } g \ n = \{x, y\}$ 
  using assms by (cases op; auto)

lemma binary-succ:
  assumes  $\text{kind } g \ n = \text{bin-node } op \ x \ y$ 
  shows  $\text{succ } g \ n = \{\}$ 
  using assms by (cases op; auto)

lemma unrep-contains:
  assumes  $g \oplus e \rightsquigarrow (g', n)$ 
  shows  $n \in \text{ids } g'$ 
  using assms
  using not-in-no-rep term-graph-reconstruction by blast

lemma unrep-preserves-contains:
  assumes  $n \in \text{ids } g$ 
  assumes  $g \oplus e \rightsquigarrow (g', n')$ 
  shows  $n \in \text{ids } g'$ 
  using assms
  by (meson subsetD unrep-ids-subset)

lemma unrep-preserves-closure:
  assumes wf-closed  $g$ 
  assumes  $g \oplus e \rightsquigarrow (g', n)$ 
  shows wf-closed  $g'$ 
  using assms(2,1) unfolding wf-closed-def
  proof (induction  $g \ e \ (g', n)$  arbitrary:  $g' \ n$ )
    case (ConstantNodeSame  $g \ c \ n$ )
    then show ?case
      by blast

```

```

next
  case (ConstantNodeNew g c n g')
  then have dom:  $ids\ g' = ids\ g \cup \{n\}$ 
    by (meson IRNode.distinct(683) add-node-ids-subset ids-add-update)
  have k:  $kind\ g'\ n = ConstantNode\ c$ 
    using ConstantNodeNew add-node-lookup by simp
  then have inp:  $\{\} = inputs\ g'\ n$ 
    unfolding inputs.simps by simp
  from k have suc:  $\{\} = succ\ g'\ n$ 
    unfolding succ.simps by simp
  have  $inputs\ g'\ n \subseteq ids\ g' \wedge succ\ g'\ n \subseteq ids\ g' \wedge kind\ g'\ n \neq NoNode$ 
    using inp suc k by simp
  then show ?case
    by (smt (verit) ConstantNodeNew.hyps(3) ConstantNodeNew.prem Un-insert-right
    add-changed changeonly.elims(2) dom inputs.simps insert-iff singleton-iff subset-insertI
    subset-trans succ.simps sup-bot-right)
  next
    case (ParameterNodeSame g i s n)
    then show ?case by blast
  next
    case (ParameterNodeNew g i s n g')
    then have dom:  $ids\ g' = ids\ g \cup \{n\}$ 
      using IRNode.distinct(2447) fresh-ids ids-add-update by presburger
    have k:  $kind\ g'\ n = ParameterNode\ i$ 
      using ParameterNodeNew add-node-lookup by simp
    then have inp:  $\{\} = inputs\ g'\ n$ 
      unfolding inputs.simps by simp
    from k have suc:  $\{\} = succ\ g'\ n$ 
      unfolding succ.simps by simp
    have  $inputs\ g'\ n \subseteq ids\ g' \wedge succ\ g'\ n \subseteq ids\ g' \wedge kind\ g'\ n \neq NoNode$ 
      using k inp suc by simp
    then show ?case
      by (smt (verit) ParameterNodeNew.hyps(3) ParameterNodeNew.prem Un-insert-right
      add-node-as-set dom inputs.elims insertE not-excluded-keep-type order-trans single-
      tonD subset-insertI succ.elims sup-bot-right)
    next
      case (ConditionalNodeSame g ce g2 c te g3 t fe g4 f s' n)
      then show ?case by blast
    next
      case (ConditionalNodeNew g4 c t f s' g ce g2 te g3 fe n g')
      then have dom:  $ids\ g' = ids\ g4 \cup \{n\}$ 
        by (meson IRNode.distinct(591) add-node-ids-subset ids-add-update)
      have k:  $kind\ g'\ n = ConditionalNode\ c\ t\ f$ 
        using ConditionalNodeNew add-node-lookup by simp
      then have inp:  $\{c, t, f\} = inputs\ g'\ n$ 
        unfolding inputs.simps by simp
      from k have suc:  $\{\} = succ\ g'\ n$ 
        unfolding succ.simps by simp
      have  $inputs\ g'\ n \subseteq ids\ g' \wedge succ\ g'\ n \subseteq ids\ g' \wedge kind\ g'\ n \neq NoNode$ 

```

```

    using k inp suc unrep-contains unrep-preserves-contains
    using ConditionalNodeNew
    by (smt (verit) IRNode.simps(643) Un-insert-right bot.extremum dom in-
sert-absorb insert-subset subset-insertI sup-bot-right)
    then show ?case using dom
    by (smt (z3) ConditionalNodeNew.hyps ConditionalNodeNew.premis Diff-eq-empty-iff
Diff-iff Un-insert-right Un-upper1 add-node-def inputs.simps insertE replace-node-def
replace-node-unchanged subset-trans succ.simps sup-bot-right)
  next
  case (UnaryNodeSame g xe g2 x s' op n)
  then show ?case by blast
next
  case (UnaryNodeNew g2 op x s' g xe n g')
  then have dom:  $ids\ g' = ids\ g2 \cup \{n\}$ 
    by (metis add-node-ids-subset add-node-lookup ids-add-update ids-some un-
rep.UnaryNodeNew unrep-contains)
  have k:  $kind\ g'\ n = unary-node\ op\ x$ 
    using UnaryNodeNew add-node-lookup
    by (metis fresh-ids ids-some)
  then have inp:  $\{x\} = inputs\ g'\ n$ 
    using unary-inputs by simp
  from k have suc:  $\{\} = succ\ g'\ n$ 
    using unary-succ by simp
  have inputs  $g'\ n \subseteq ids\ g' \wedge succ\ g'\ n \subseteq ids\ g' \wedge kind\ g'\ n \neq NoNode$ 
    using k inp suc unrep-contains unrep-preserves-contains
    using UnaryNodeNew
    by (metis Un-upper1 dom empty-subsetI ids-some insert-not-empty in-
sert-subsetI not-in-g-inputs subset-iff)
  then show ?case
    by (smt (verit) Un-insert-right UnaryNodeNew.hyps UnaryNodeNew.premis
add-changed changeonly.elims(2) dom inputs.simps insert-iff singleton-iff subset-insertI
subset-trans succ.simps sup-bot-right)
  next
  case (BinaryNodeSame g xe g2 x ye g3 y s' op n)
  then show ?case by blast
next
  case (BinaryNodeNew g3 op x y s' g xe g2 ye n g')
  then have dom:  $ids\ g' = ids\ g3 \cup \{n\}$ 
    by (metis binary-inputs fresh-ids ids-add-update ids-some insert-not-empty
not-in-g-inputs)
  have k:  $kind\ g'\ n = bin-node\ op\ x\ y$ 
    using BinaryNodeNew add-node-lookup
    by (metis fresh-ids ids-some)
  then have inp:  $\{x, y\} = inputs\ g'\ n$ 
    using binary-inputs by simp
  from k have suc:  $\{\} = succ\ g'\ n$ 
    using binary-succ by simp
  have inputs  $g'\ n \subseteq ids\ g' \wedge succ\ g'\ n \subseteq ids\ g' \wedge kind\ g'\ n \neq NoNode$ 
    using k inp suc unrep-contains unrep-preserves-contains

```

```

    using BinaryNodeNew
    by (metis Un-upper1 dom empty-subsetI ids-some insert-not-empty in-
sert-subsetI not-in-g-inputs subset-iff)
    then show ?case using dom BinaryNodeNew
    by (smt (verit, del-Insts) Diff-eq-empty-iff Diff-iff Un-insert-right Un-upper1
add-node-def inputs.simps insertE replace-node-def replace-node-unchanged subset-trans
succ.simps sup-bot-right)
  next
  case (AllLeafNodes g n s)
  then show ?case
  by blast
qed

```

inductive-cases *ConstUnrepE*: $g \oplus (\text{ConstantExpr } x) \rightsquigarrow (g', n)$

definition *constant-value* **where**

constant-value = (*IntVal* 32 0)

definition *bad-graph* **where**

```

bad-graph = irgraph [
  (0, AbsNode 1, constantAsStamp constant-value),
  (1, RefNode 2, constantAsStamp constant-value),
  (2, ConstantNode constant-value, constantAsStamp constant-value)
]

```

end

8 Control-flow Semantics

theory *IRStepObj*

imports

TreeToGraph

begin

8.1 Object Heap

The heap model we introduce maps field references to object instances to runtime values. We use the $H[f][p]$ heap representation. See \cite{heap-reps-2011}.

We also introduce the *DynamicHeap* type which allocates new object references sequentially storing the next free object reference as 'Free'.

heapdef

```

type-synonym ('a, 'b) Heap = 'a  $\Rightarrow$  'b  $\Rightarrow$  Value
type-synonym Free = nat
type-synonym ('a, 'b) DynamicHeap = ('a, 'b) Heap  $\times$  Free

fun h-load-field :: 'a  $\Rightarrow$  'b  $\Rightarrow$  ('a, 'b) DynamicHeap  $\Rightarrow$  Value where
  h-load-field f r (h, n) = h f r

fun h-store-field :: 'a  $\Rightarrow$  'b  $\Rightarrow$  Value  $\Rightarrow$  ('a, 'b) DynamicHeap  $\Rightarrow$  ('a, 'b)
  DynamicHeap where
  h-store-field f r v (h, n) = (h(f := ((h f)(r := v))), n)

fun h-new-inst :: ('a, 'b) DynamicHeap  $\Rightarrow$  ('a, 'b) DynamicHeap  $\times$  Value
where
  h-new-inst (h, n) = ((h, n+1), (ObjRef (Some n)))

type-synonym FieldRefHeap = (string, objref) DynamicHeap

```

```

definition new-heap :: ('a, 'b) DynamicHeap where
  new-heap = (( $\lambda$ f.  $\lambda$ p. UndefVal), 0)

```

8.2 Intraprocedural Semantics

```

fun find-index :: 'a  $\Rightarrow$  'a list  $\Rightarrow$  nat where
  find-index - [] = 0 |
  find-index v (x # xs) = (if (x=v) then 0 else find-index v xs + 1)

fun phi-list :: IRGraph  $\Rightarrow$  ID  $\Rightarrow$  ID list where
  phi-list g n =
    (filter ( $\lambda$ x.(is-PhiNode (kind g x)))
     (sorted-list-of-set (usages g n)))

fun input-index :: IRGraph  $\Rightarrow$  ID  $\Rightarrow$  ID  $\Rightarrow$  nat where
  input-index g n n' = find-index n' (inputs-of (kind g n))

fun phi-inputs :: IRGraph  $\Rightarrow$  nat  $\Rightarrow$  ID list  $\Rightarrow$  ID list where
  phi-inputs g i nodes = (map ( $\lambda$ n. (inputs-of (kind g n))!(i + 1)) nodes)

fun set-phis :: ID list  $\Rightarrow$  Value list  $\Rightarrow$  MapState  $\Rightarrow$  MapState where
  set-phis [] [] m = m |
  set-phis (n # xs) (v # vs) m = (set-phis xs vs (m(n := v))) |
  set-phis [] (v # vs) m = m |
  set-phis (x # xs) [] m = m

```

Intraprocedural semantics are given as a small-step semantics.

Within the context of a graph, the configuration triple, (ID, MethodState, Heap), is related to the subsequent configuration.

```

inductive step :: IRGraph  $\Rightarrow$  Params  $\Rightarrow$  (ID  $\times$  MapState  $\times$  FieldRefHeap)  $\Rightarrow$  (ID

```

$\times \text{MapState} \times \text{FieldRefHeap}) \Rightarrow \text{bool}$
 $(-, - \vdash - \rightarrow - \ 55)$ **for** $g \ p$ **where**

SequentialNode:

$\llbracket \text{is-sequential-node } (kind \ g \ nid);$
 $\quad nid' = (\text{successors-of } (kind \ g \ nid))!0 \rrbracket$
 $\implies g, p \vdash (nid, m, h) \rightarrow (nid', m, h) \mid$

IfNode:

$\llbracket kind \ g \ nid = (\text{IfNode } cond \ tb \ fb);$
 $\quad g \vdash cond \simeq condE;$
 $\quad [m, p] \vdash condE \mapsto val;$
 $\quad nid' = (\text{if } val\text{-to-bool } val \text{ then } tb \text{ else } fb) \rrbracket$
 $\implies g, p \vdash (nid, m, h) \rightarrow (nid', m, h) \mid$

EndNodes:

$\llbracket \text{is-AbstractEndNode } (kind \ g \ nid);$
 $\quad merge = \text{any-usage } g \ nid;$
 $\quad \text{is-AbstractMergeNode } (kind \ g \ merge);$

 $\quad i = \text{find-index } nid \ (\text{inputs-of } (kind \ g \ merge));$
 $\quad phis = (\text{phi-list } g \ merge);$
 $\quad inps = (\text{phi-inputs } g \ i \ phis);$
 $\quad g \vdash inps \simeq_L inpsE;$
 $\quad [m, p] \vdash inpsE \mapsto_L vs;$

 $\quad m' = \text{set-phis } phis \ vs \ m \rrbracket$
 $\implies g, p \vdash (nid, m, h) \rightarrow (merge, m', h) \mid$

NewInstanceNode:

$\llbracket kind \ g \ nid = (\text{NewInstanceNode } nid \ f \ obj \ nid');$
 $\quad (h', ref) = h\text{-new-inst } h;$
 $\quad m' = m(nid := ref) \rrbracket$
 $\implies g, p \vdash (nid, m, h) \rightarrow (nid', m', h') \mid$

LoadFieldNode:

$\llbracket kind \ g \ nid = (\text{LoadFieldNode } nid \ f \ (\text{Some } obj) \ nid');$
 $\quad g \vdash obj \simeq objE;$
 $\quad [m, p] \vdash objE \mapsto \text{ObjRef } ref;$
 $\quad h\text{-load-field } f \ ref \ h = v;$
 $\quad m' = m(nid := v) \rrbracket$
 $\implies g, p \vdash (nid, m, h) \rightarrow (nid', m', h) \mid$

SignedDivNode:

$\llbracket kind \ g \ nid = (\text{SignedDivNode } nid \ x \ y \ zero \ sb \ nxt);$
 $\quad g \vdash x \simeq xe;$
 $\quad g \vdash y \simeq ye;$
 $\quad [m, p] \vdash xe \mapsto v1;$

$$\begin{aligned}
& [m, p] \vdash ye \mapsto v2; \\
& v = (\text{intval-div } v1 \ v2); \\
& m' = m(\text{nid} := v) \\
\implies & g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nxt}, m', h) \mid
\end{aligned}$$

SignedRemNode:

$$\begin{aligned}
& \llbracket \text{kind } g \text{ nid} = (\text{SignedRemNode } \text{nid } x \ y \ \text{zero } sb \ \text{nxt}); \\
& \quad g \vdash x \simeq xe; \\
& \quad g \vdash y \simeq ye; \\
& \quad [m, p] \vdash xe \mapsto v1; \\
& \quad [m, p] \vdash ye \mapsto v2; \\
& \quad v = (\text{intval-mod } v1 \ v2); \\
& \quad m' = m(\text{nid} := v) \\
\implies & g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nxt}, m', h) \mid
\end{aligned}$$

StaticLoadFieldNode:

$$\begin{aligned}
& \llbracket \text{kind } g \text{ nid} = (\text{LoadFieldNode } \text{nid } f \ \text{None } \text{nid}'); \\
& \quad h\text{-load-field } f \ \text{None } h = v; \\
& \quad m' = m(\text{nid} := v) \\
\implies & g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m', h) \mid
\end{aligned}$$

StoreFieldNode:

$$\begin{aligned}
& \llbracket \text{kind } g \text{ nid} = (\text{StoreFieldNode } \text{nid } f \ \text{newval} - (\text{Some } \text{obj}) \ \text{nid}'); \\
& \quad g \vdash \text{newval} \simeq \text{newvalE}; \\
& \quad g \vdash \text{obj} \simeq \text{objE}; \\
& \quad [m, p] \vdash \text{newvalE} \mapsto \text{val}; \\
& \quad [m, p] \vdash \text{objE} \mapsto \text{ObjRef } \text{ref}; \\
& \quad h' = h\text{-store-field } f \ \text{ref } \text{val } h; \\
& \quad m' = m(\text{nid} := \text{val}) \\
\implies & g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m', h') \mid
\end{aligned}$$

StaticStoreFieldNode:

$$\begin{aligned}
& \llbracket \text{kind } g \text{ nid} = (\text{StoreFieldNode } \text{nid } f \ \text{newval} - \text{None } \text{nid}'); \\
& \quad g \vdash \text{newval} \simeq \text{newvalE}; \\
& \quad [m, p] \vdash \text{newvalE} \mapsto \text{val}; \\
& \quad h' = h\text{-store-field } f \ \text{None } \text{val } h; \\
& \quad m' = m(\text{nid} := \text{val}) \\
\implies & g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m', h')
\end{aligned}$$

code-pred (*modes*: $i \Rightarrow i \Rightarrow i * i * i \Rightarrow o * o * o \Rightarrow \text{bool}$) *step* .

8.3 Interprocedural Semantics

type-synonym *Signature* = *string*

type-synonym *Program* = *Signature* \rightarrow *IRGraph*

inductive *step-top* :: *Program* \Rightarrow (*IRGraph* \times *ID* \times *MapState* \times *Params*) *list* \times *FieldRefHeap* \Rightarrow (*IRGraph* \times *ID* \times *MapState* \times *Params*) *list* \times *FieldRefHeap* \Rightarrow

bool

$(- \vdash - \longrightarrow - \ 55)$

for P where

Lift:

$\llbracket g, p \vdash (nid, m, h) \rightarrow (nid', m', h') \rrbracket$
 $\implies P \vdash ((g, nid, m, p) \# stk, h) \longrightarrow ((g, nid', m', p) \# stk, h') \mid$

InvokeNodeStep:

$\llbracket is-Invoke \ (kind \ g \ nid);$

$callTarget = ir-callTarget \ (kind \ g \ nid);$

$kind \ g \ callTarget = (MethodCallTargetNode \ targetMethod \ arguments);$

$Some \ targetGraph = P \ targetMethod;$

$m' = new-map-state;$

$g \vdash arguments \simeq_L argsE;$

$[m, p] \vdash argsE \mapsto_L p \rrbracket$

$\implies P \vdash ((g, nid, m, p) \# stk, h) \longrightarrow ((targetGraph, 0, m', p') \# (g, nid, m, p) \# stk, h)$

\mid

ReturnNode:

$\llbracket kind \ g \ nid = (ReturnNode \ (Some \ expr) \ -);$

$g \vdash expr \simeq e;$

$[m, p] \vdash e \mapsto v;$

$cm' = cm(cnid := v);$

$cnid' = (successors-of \ (kind \ cg \ cnid))!0 \rrbracket$

$\implies P \vdash ((g, nid, m, p) \# (cg, cnid, cm, cp) \# stk, h) \longrightarrow ((cg, cnid', cm', cp) \# stk, h) \mid$

ReturnNodeVoid:

$\llbracket kind \ g \ nid = (ReturnNode \ None \ -);$

$cm' = cm(cnid := (ObjRef \ (Some \ (2048))));$

$cnid' = (successors-of \ (kind \ cg \ cnid))!0 \rrbracket$

$\implies P \vdash ((g, nid, m, p) \# (cg, cnid, cm, cp) \# stk, h) \longrightarrow ((cg, cnid', cm', cp) \# stk, h) \mid$

UnwindNode:

$\llbracket kind \ g \ nid = (UnwindNode \ exception);$

$g \vdash exception \simeq exceptionE;$

$[m, p] \vdash exceptionE \mapsto e;$

$kind \ cg \ cnid = (InvokeWithExceptionNode \ - \ - \ - \ - \ - \ exEdge);$

$cm' = cm(cnid := e) \rrbracket$

$\implies P \vdash ((g, nid, m, p) \# (cg, cnid, cm, cp) \# stk, h) \longrightarrow ((cg, exEdge, cm', cp) \# stk, h)$

code-pred (*modes*: $i \Rightarrow i \Rightarrow o \Rightarrow bool$) *step-top* .

8.4 Big-step Execution

type-synonym $Trace = (IRGraph \times ID \times MapState \times Params) \text{ list}$

fun $has\text{-}return :: MapState \Rightarrow bool$ **where**
 $has\text{-}return\ m = (m\ 0 \neq UndefinedVal)$

inductive $exec :: Program$
 $\Rightarrow (IRGraph \times ID \times MapState \times Params) \text{ list} \times FieldRefHeap$
 $\Rightarrow Trace$
 $\Rightarrow (IRGraph \times ID \times MapState \times Params) \text{ list} \times FieldRefHeap$
 $\Rightarrow Trace$
 $\Rightarrow bool$
 $(- \vdash - \mid - \longrightarrow * - \mid -)$
for P
where
 $\llbracket P \vdash (((g, nid, m, p) \# xs), h) \longrightarrow (((g', nid', m', p') \# ys), h') ;$
 $\neg(has\text{-}return\ m') ;$
 $l' = (l @ [(g, nid, m, p)]) ;$
 $exec\ P\ (((g', nid', m', p') \# ys), h')\ l'\ next\text{-}state\ l'' \rrbracket$
 $\implies exec\ P\ (((g, nid, m, p) \# xs), h)\ l\ next\text{-}state\ l''$
 \mid
 $\llbracket P \vdash (((g, nid, m, p) \# xs), h) \longrightarrow (((g', nid', m', p') \# ys), h') ;$
 $has\text{-}return\ m' ;$
 $l' = (l @ [(g, nid, m, p)]) \rrbracket$
 $\implies exec\ P\ (((g, nid, m, p) \# xs), h)\ l\ (((g', nid', m', p') \# ys), h')\ l'$
code-pred $(modes: i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool \text{ as } Exec)\ exec .$

inductive $exec\text{-}debug :: Program$
 $\Rightarrow (IRGraph \times ID \times MapState \times Params) \text{ list} \times FieldRefHeap$
 $\Rightarrow nat$
 $\Rightarrow (IRGraph \times ID \times MapState \times Params) \text{ list} \times FieldRefHeap$
 $\Rightarrow bool$
 $(\vdash \longrightarrow * - \mid -)$
where
 $\llbracket n > 0 ;$
 $p \vdash s \longrightarrow s' ;$
 $exec\text{-}debug\ p\ s'\ (n - 1)\ s' \rrbracket$
 $\implies exec\text{-}debug\ p\ s\ n\ s'' \mid$
 $\llbracket n = 0 \rrbracket$
 $\implies exec\text{-}debug\ p\ s\ n\ s$
code-pred $(modes: i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool)\ exec\text{-}debug .$

8.4.1 Heap Testing

definition $p3 :: Params$ **where**

$p3 = [IntVal\ 32\ 3]$

values $\{(prod.fst(prod.snd\ (prod.snd\ (hd\ (prod.fst\ res))))\ 0$
 $\mid res. (\lambda x. Some\ eg2-sq) \vdash [(eg2-sq, 0, new-map-state, p3), (eg2-sq, 0, new-map-state, p3)],$
 $new-heap) \rightarrow^* 2^* res\}$

definition $field-sq :: string$ **where**

$field-sq = "sq"$

definition $eg3-sq :: IRGraph$ **where**

$eg3-sq = irgraph\ [$
 $(0, StartNode\ None\ 4, VoidStamp),$
 $(1, ParameterNode\ 0, default-stamp),$
 $(3, MulNode\ 1\ 1, default-stamp),$
 $(4, StoreFieldNode\ 4\ field-sq\ 3\ None\ None\ 5, VoidStamp),$
 $(5, ReturnNode\ (Some\ 3)\ None, default-stamp)$
 $]$

values $\{h-load-field\ field-sq\ None\ (prod.snd\ res)$
 $\mid res. (\lambda x. Some\ eg3-sq) \vdash [(eg3-sq, 0, new-map-state, p3), (eg3-sq, 0,$
 $new-map-state, p3)], new-heap) \rightarrow^* 3^* res\}$

definition $eg4-sq :: IRGraph$ **where**

$eg4-sq = irgraph\ [$
 $(0, StartNode\ None\ 4, VoidStamp),$
 $(1, ParameterNode\ 0, default-stamp),$
 $(3, MulNode\ 1\ 1, default-stamp),$
 $(4, NewInstanceNode\ 4\ "obj-class"\ None\ 5, ObjectStamp\ "obj-class"\ True\ True$
 $True),$
 $(5, StoreFieldNode\ 5\ field-sq\ 3\ None\ (Some\ 4)\ 6, VoidStamp),$
 $(6, ReturnNode\ (Some\ 3)\ None, default-stamp)$
 $]$

values $\{h-load-field\ field-sq\ (Some\ 0)\ (prod.snd\ res) \mid res.$
 $(\lambda x. Some\ eg4-sq) \vdash [(eg4-sq, 0, new-map-state, p3), (eg4-sq, 0,$
 $new-map-state, p3)], new-heap) \rightarrow^* 3^* res\}$

end

8.5 Control-flow Semantics Theorems

theory $IRStepThms$

imports

$IRStepObj$

begin

We prove that within the same graph, a configuration triple will always transition to the same subsequent configuration. Therefore, our step semantics is deterministic.

8.5.1 Control-flow Step is Deterministic

theorem *stepDet*:

$(g, p \vdash (nid, m, h) \rightarrow next) \implies$
 $(\forall next'. ((g, p \vdash (nid, m, h) \rightarrow next') \longrightarrow next = next'))$

proof (*induction rule: step.induct*)

case (*SequentialNode nid next m h*)

have *notif*: $\neg(is_IfNode\ (kind\ g\ nid))$

using *SequentialNode.hyps(1) is-sequential-node.simps*

by (*metis is-IfNode-def*)

have *notend*: $\neg(is_AbstractEndNode\ (kind\ g\ nid))$

using *SequentialNode.hyps(1) is-sequential-node.simps*

by (*metis is-AbstractEndNode.simps is-EndNode.elims(2) is-LoopEndNode-def*)

have *notnew*: $\neg(is_NewInstanceNode\ (kind\ g\ nid))$

using *SequentialNode.hyps(1) is-sequential-node.simps*

by (*metis is-NewInstanceNode-def*)

have *notload*: $\neg(is_LoadFieldNode\ (kind\ g\ nid))$

using *SequentialNode.hyps(1) is-sequential-node.simps*

by (*metis is-LoadFieldNode-def*)

have *notstore*: $\neg(is_StoreFieldNode\ (kind\ g\ nid))$

using *SequentialNode.hyps(1) is-sequential-node.simps*

by (*metis is-StoreFieldNode-def*)

have *notdivrem*: $\neg(is_IntegerDivRemNode\ (kind\ g\ nid))$

using *SequentialNode.hyps(1) is-sequential-node.simps is-SignedDivNode-def*

is-SignedRemNode-def

by (*metis is-IntegerDivRemNode.simps*)

from *notif notend notnew notload notstore notdivrem*

show *?case using SequentialNode step.cases*

by (*smt (z3) IRNode.disc(1028) IRNode.disc(2270) IRNode.discI(31) Pair-inject*

is-sequential-node.simps(18) is-sequential-node.simps(43) is-sequential-node.simps(44))

next

case (*IfNode nid cond tb fb m val next h*)

then have *notseq*: $\neg(is_sequential-node\ (kind\ g\ nid))$

using *is-sequential-node.simps is-AbstractMergeNode.simps*

by (*simp add: IfNode.hyps(1)*)

have *notend*: $\neg(is_AbstractEndNode\ (kind\ g\ nid))$

using *is-AbstractEndNode.simps*

by (*simp add: IfNode.hyps(1)*)

have *notdivrem*: $\neg(is_IntegerDivRemNode\ (kind\ g\ nid))$

using *is-AbstractEndNode.simps*

by (*simp add: IfNode.hyps(1)*)

from *notseq notend notdivrem show ?case using IfNode repDet evalDet IRN-*

```

ode.distinct IRNode.inject(11) Pair-inject step.simps
  by (smt (z3) IRNode.distinct IRNode.inject(12) Pair-inject step.simps)
next
case (EndNodes nid merge i phis inputs m vs m' h)
have notseq: ¬(is-sequential-node (kind g nid))
  using EndNodes.hyps(1) is-AbstractEndNode.simps is-sequential-node.simps
  by (metis is-EndNode.elims(2) is-LoopEndNode-def)
have notif: ¬(is-IfNode (kind g nid))
  using EndNodes.hyps(1) is-IfNode-def is-AbstractEndNode.elims
  by (metis IRNode.distinct-disc(1058) is-EndNode.simps(12))
have notref: ¬(is-RefNode (kind g nid))
  using EndNodes.hyps(1) is-sequential-node.simps
  using IRNode.disc(1899) IRNode.distinct(1473) is-AbstractEndNode.simps
is-EndNode.elims(2) is-LoopEndNode-def is-RefNode-def
  by metis
have notnew: ¬(is-NewInstanceNode (kind g nid))
  using EndNodes.hyps(1) is-AbstractEndNode.simps
  using IRNode.distinct-disc(1442) is-EndNode.simps(29) is-NewInstanceNode-def
  by (metis IRNode.distinct-disc(1901) is-EndNode.simps(32))
have notload: ¬(is-LoadFieldNode (kind g nid))
  using EndNodes.hyps(1) is-AbstractEndNode.simps
  using is-LoadFieldNode-def
  by (metis IRNode.distinct-disc(1706) is-EndNode.simps(21))
have notstore: ¬(is-StoreFieldNode (kind g nid))
  using EndNodes.hyps(1) is-AbstractEndNode.simps is-StoreFieldNode-def
  by (metis IRNode.distinct-disc(1926) is-EndNode.simps(44))
have notdivrem: ¬(is-IntegerDivRemNode (kind g nid))
  using EndNodes.hyps(1) is-AbstractEndNode.simps is-SignedDivNode-def is-SignedRemNode-def
  using IRNode.distinct-disc(1498) IRNode.distinct-disc(1500) is-IntegerDivRemNode.simps
is-EndNode.simps(36) is-EndNode.simps(37)
  by auto
from notseq notif notref notnew notload notstore notdivrem
show ?case using EndNodes repAllDet evalAllDet
  by (smt (z3) is-IfNode-def is-LoadFieldNode-def is-NewInstanceNode-def is-RefNode-def
is-StoreFieldNode-def is-SignedDivNode-def is-SignedRemNode-def Pair-inject is-IntegerDivRemNode.elims(3)
step.cases)
next
case (NewInstanceNode nid f obj nxt h' ref h m' m)
then have notseq: ¬(is-sequential-node (kind g nid))
  using is-sequential-node.simps is-AbstractMergeNode.simps
  by (simp add: NewInstanceNode.hyps(1))
have notend: ¬(is-AbstractEndNode (kind g nid))
  using is-AbstractMergeNode.simps
  by (simp add: NewInstanceNode.hyps(1))
have notif: ¬(is-IfNode (kind g nid))
  using is-AbstractMergeNode.simps
  by (simp add: NewInstanceNode.hyps(1))
have notref: ¬(is-RefNode (kind g nid))
  using is-AbstractMergeNode.simps

```



```

    by (simp add: NewInstanceNode.hyps(1))
  have notload:  $\neg(\text{is-LoadFieldNode } (\text{kind } g \text{ nid}))$ 
    using is-AbstractMergeNode.simps
    by (simp add: NewInstanceNode.hyps(1))
  have notstore:  $\neg(\text{is-StoreFieldNode } (\text{kind } g \text{ nid}))$ 
    using is-AbstractMergeNode.simps
    by (simp add: NewInstanceNode.hyps(1))
  have notdivrem:  $\neg(\text{is-IntegerDivRemNode } (\text{kind } g \text{ nid}))$ 
    using is-AbstractMergeNode.simps
    by (simp add: NewInstanceNode.hyps(1))
  from notseq notend notif notref notload notstore notdivrem
  show ?case using NewInstanceNode.step.cases
    by (smt (z3) IRNode.disc(1028) IRNode.disc(2270) IRNode.discI(11) IRN-
ode.distinct(2311) IRNode.distinct(2313) IRNode.inject(31) Pair-inject)
next
  case (LoadFieldNode nid f obj nrt m ref h v m')
  then have notseq:  $\neg(\text{is-sequential-node } (\text{kind } g \text{ nid}))$ 
    using is-sequential-node.simps is-AbstractMergeNode.simps
    by (simp add: LoadFieldNode.hyps(1))
  have notend:  $\neg(\text{is-AbstractEndNode } (\text{kind } g \text{ nid}))$ 
    using is-AbstractEndNode.simps
    by (simp add: LoadFieldNode.hyps(1))
  have notdivrem:  $\neg(\text{is-IntegerDivRemNode } (\text{kind } g \text{ nid}))$ 
    using is-AbstractEndNode.simps
    by (simp add: LoadFieldNode.hyps(1))
  from notseq notend notdivrem
  show ?case using LoadFieldNode.step.cases repDet evalDet
    by (smt (z3) IRNode.distinct(1051) IRNode.distinct(1721) IRNode.distinct(1739)
IRNode.distinct(1741) IRNode.distinct(1745) IRNode.inject(20) Pair-inject Value.inject(2)
option.distinct(1) option.inject)
next
  case (StaticLoadFieldNode nid f nrt h v m' m)
  then have notseq:  $\neg(\text{is-sequential-node } (\text{kind } g \text{ nid}))$ 
    using is-sequential-node.simps is-AbstractMergeNode.simps
    by (simp add: StaticLoadFieldNode.hyps(1))
  have notend:  $\neg(\text{is-AbstractEndNode } (\text{kind } g \text{ nid}))$ 
    using is-AbstractEndNode.simps
    by (simp add: StaticLoadFieldNode.hyps(1))
  have notdivrem:  $\neg(\text{is-IntegerDivRemNode } (\text{kind } g \text{ nid}))$ 
    by (simp add: StaticLoadFieldNode.hyps(1))
  from notseq notend notdivrem
  show ?case using StaticLoadFieldNode.step.cases
    by (smt (z3) IRNode.distinct(1051) IRNode.distinct(1721) IRNode.distinct(1739)
IRNode.distinct(1741) IRNode.distinct(1745) IRNode.inject(20) Pair-inject option.distinct(1))
next
  case (StoreFieldNode nid f newval uu obj nrt m val ref h' h m')
  then have notseq:  $\neg(\text{is-sequential-node } (\text{kind } g \text{ nid}))$ 
    using is-sequential-node.simps is-AbstractMergeNode.simps
    by (simp add: StoreFieldNode.hyps(1))

```

```

have notend: ¬(is-AbstractEndNode (kind g nid))
  using is-AbstractEndNode.simps
  by (simp add: StoreFieldNode.hyps(1))
have notdivrem: ¬(is-IntegerDivRemNode (kind g nid))
  by (simp add: StoreFieldNode.hyps(1))
from notseq notend notdivrem
show ?case using StoreFieldNode step.cases repDet evalDet
  by (smt (z3) IRNode.distinct(1097) IRNode.distinct(1745) IRNode.distinct(2317)
IRNode.distinct(2605) IRNode.distinct(2627) IRNode.inject(43) Pair-inject Value.inject(2)
option.distinct(1) option.inject)
next
case (StaticStoreFieldNode nid f newval uv nxt m val h' h m')
then have notseq: ¬(is-sequential-node (kind g nid))
  using is-sequential-node.simps is-AbstractMergeNode.simps
  by (simp add: StaticStoreFieldNode.hyps(1))
have notend: ¬(is-AbstractEndNode (kind g nid))
  using is-AbstractEndNode.simps
  by (simp add: StaticStoreFieldNode.hyps(1))
have notdivrem: ¬(is-IntegerDivRemNode (kind g nid))
  by (simp add: StaticStoreFieldNode.hyps(1))
from notseq notend notdivrem
show ?case using StoreFieldNode step.cases repDet evalDet
  by (smt (z3) IRNode.distinct(1097) IRNode.distinct(1745) IRNode.distinct(2317)
IRNode.distinct(2605) IRNode.distinct(2627) IRNode.inject(43) Pair-inject Static-
StoreFieldNode.hyps(1) StaticStoreFieldNode.hyps(2) StaticStoreFieldNode.hyps(3)
StaticStoreFieldNode.hyps(4) StaticStoreFieldNode.hyps(5) option.distinct(1))
next
case (SignedDivNode nid x y zero sb nxt m v1 v2 v m' h)
then have notseq: ¬(is-sequential-node (kind g nid))
  using is-sequential-node.simps is-AbstractMergeNode.simps
  by (simp add: SignedDivNode.hyps(1))
have notend: ¬(is-AbstractEndNode (kind g nid))
  using is-AbstractEndNode.simps
  by (simp add: SignedDivNode.hyps(1))
from notseq notend
show ?case using SignedDivNode step.cases repDet evalDet
  by (smt (z3) IRNode.distinct(1091) IRNode.distinct(1739) IRNode.distinct(2311)
IRNode.distinct(2601) IRNode.distinct(2605) IRNode.inject(40) Pair-inject)
next
case (SignedRemNode nid x y zero sb nxt m v1 v2 v m' h)
then have notseq: ¬(is-sequential-node (kind g nid))
  using is-sequential-node.simps is-AbstractMergeNode.simps
  by (simp add: SignedRemNode.hyps(1))
have notend: ¬(is-AbstractEndNode (kind g nid))
  using is-AbstractEndNode.simps
  by (simp add: SignedRemNode.hyps(1))
from notseq notend
show ?case using SignedRemNode step.cases repDet evalDet
  by (smt (z3) IRNode.distinct(1093) IRNode.distinct(1741) IRNode.distinct(2313)

```

IRNode.distinct(2601) IRNode.distinct(2627) IRNode.inject(41) Pair-inject
qed

lemma *stepRefNode*:

$\llbracket \text{kind } g \text{ nid} = \text{RefNode nid}' \rrbracket \implies g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m, h)$

using *SequentialNode*

by (*metis IRNodes.successors-of-RefNode is-sequential-node.simps(7) nth-Cons-0*)

lemma *IfNodeStepCases*:

assumes *kind g nid = IfNode cond tb fb*

assumes $g \vdash \text{cond} \simeq \text{condE}$

assumes $[m, p] \vdash \text{condE} \mapsto v$

assumes $g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m, h)$

shows $\text{nid}' \in \{tb, fb\}$

using *step.IfNode repDet stepDet assms*

by (*metis insert-iff old.prod.inject*)

lemma *IfNodeSeq*:

shows *kind g nid = IfNode cond tb fb* $\longrightarrow \neg(\text{is-sequential-node } (\text{kind } g \text{ nid}))$

unfolding *is-sequential-node.simps*

using *is-sequential-node.simps(18)* **by** *presburger*

lemma *IfNodeCond*:

assumes *kind g nid = IfNode cond tb fb*

assumes $g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m, h)$

shows $\exists \text{ condE } v. ((g \vdash \text{cond} \simeq \text{condE}) \wedge ([m, p] \vdash \text{condE} \mapsto v))$

using *assms(2,1)* **by** (*induct (nid,m,h) (nid',m,h) rule: step.induct; auto*)

lemma *step-in-ids*:

assumes $g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m', h')$

shows $\text{nid} \in \text{ids } g$

using *assms* **apply** (*induct (nid, m, h) (nid', m', h') rule: step.induct*)

using *is-sequential-node.simps(45) not-in-g*

apply *simp*

apply (*metis is-sequential-node.simps(53)*)

using *ids-some*

using *IRNode.distinct(1113)* **apply** *presburger*

using *EndNodes(1) is-AbstractEndNode.simps is-EndNode.simps(45) ids-some*

apply (*metis IRNode.disc(1218) is-EndNode.simps(52)*)

by *simp+*

end

9 Proof Infrastructure

9.1 Bisimulation

theory *Bisimulation*

imports

Stuttering
begin

inductive *weak-bisimilar* :: $ID \Rightarrow IRGraph \Rightarrow IRGraph \Rightarrow bool$
 (· · · ~ ·) **for** *nid* **where**
 $\llbracket \forall P'. (g \ m \ p \ h \vdash \text{nid} \rightsquigarrow P') \longrightarrow (\exists Q'. (g' \ m \ p \ h \vdash \text{nid} \rightsquigarrow Q') \wedge P' = Q');$
 $\forall Q'. (g' \ m \ p \ h \vdash \text{nid} \rightsquigarrow Q') \longrightarrow (\exists P'. (g \ m \ p \ h \vdash \text{nid} \rightsquigarrow P') \wedge P' = Q') \rrbracket$
 $\implies \text{nid} \cdot g \sim g'$

A strong bisimilution between no-op transitions

inductive *strong-noop-bisimilar* :: $ID \Rightarrow IRGraph \Rightarrow IRGraph \Rightarrow bool$
 (· | · ~ ·) **for** *nid* **where**
 $\llbracket \forall P'. (g, p \vdash (\text{nid}, m, h) \rightarrow P') \longrightarrow (\exists Q'. (g', p \vdash (\text{nid}, m, h) \rightarrow Q') \wedge P' = Q');$
 $\forall Q'. (g', p \vdash (\text{nid}, m, h) \rightarrow Q') \longrightarrow (\exists P'. (g, p \vdash (\text{nid}, m, h) \rightarrow P') \wedge P' = Q') \rrbracket$
 $\implies \text{nid} \mid g \sim g'$

lemma *lockstep-strong-bisimilulation*:
assumes $g' = \text{replace-node } \text{nid} \ \text{node } g$
assumes $g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m, h)$
assumes $g', p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m, h)$
shows $\text{nid} \mid g \sim g'$
using *assms(2) assms(3) stepDet strong-noop-bisimilar.simps* **by** *metis*

lemma *no-step-bisimulation*:
assumes $\forall m \ p \ h \ \text{nid}' \ m' \ h'. \neg(g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m', h'))$
assumes $\forall m \ p \ h \ \text{nid}' \ m' \ h'. \neg(g', p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m', h'))$
shows $\text{nid} \mid g \sim g'$
using *assms*
by (*simp add: assms(1) assms(2) strong-noop-bisimilar.intros*)

end

9.2 Graph Rewriting

theory
Rewrites
imports
Stuttering
begin

fun *replace-usages* :: $ID \Rightarrow ID \Rightarrow IRGraph \Rightarrow IRGraph$ **where**
replace-usages *nid* *nid'* *g* = *replace-node* *nid* (*RefNode* *nid'*, *stamp* *g* *nid'*) *g*

lemma *replace-usages-effect*:
assumes $g' = \text{replace-usages } \text{nid} \ \text{nid}' \ g$

```

shows kind  $g'$   $nid = \text{RefNode } nid'$ 
using assms replace-node-lookup replace-usages.simps
by (metis IRNode.distinct(2755))

lemma replace-usages-changeonly:
  assumes  $nid \in \text{ids } g$ 
  assumes  $g' = \text{replace-usages } nid \text{ } nid' \text{ } g$ 
  shows changeonly  $\{nid\} \text{ } g \text{ } g'$ 
  using assms unfolding replace-usages.simps
  by (metis add-changed add-node-def replace-node-def)

lemma replace-usages-unchanged:
  assumes  $nid \in \text{ids } g$ 
  assumes  $g' = \text{replace-usages } nid \text{ } nid' \text{ } g$ 
  shows unchanged  $(\text{ids } g - \{nid\}) \text{ } g \text{ } g'$ 
  using assms unfolding replace-usages.simps
  using assms(2) disjoint-change replace-usages-changeonly by presburger

fun nextNid :: IRGraph  $\Rightarrow$  ID where
  nextNid  $g = (\text{Max } (\text{ids } g)) + 1$ 

lemma max-plus-one:
  fixes  $c :: \text{ID set}$ 
  shows  $\llbracket \text{finite } c; c \neq \{\} \rrbracket \Longrightarrow (\text{Max } c) + 1 \notin c$ 
  by (meson Max-gr-iff less-add-one less-irrefl)

lemma ids-finite:
  finite  $(\text{ids } g)$ 
by simp

lemma nextNidNotIn:
   $\text{ids } g \neq \{\} \longrightarrow \text{nextNid } g \notin \text{ids } g$ 
unfolding nextNid.simps
using ids-finite max-plus-one by blast

fun bool-to-val-width1 :: bool  $\Rightarrow$  Value where
  bool-to-val-width1 True = (IntVal 1 1) |
  bool-to-val-width1 False = (IntVal 1 0)

fun constantCondition :: bool  $\Rightarrow$  ID  $\Rightarrow$  IRNode  $\Rightarrow$  IRGraph  $\Rightarrow$  IRGraph where
  constantCondition val  $nid$  (IfNode cond t f)  $g =$ 
    replace-node  $nid$  (IfNode (nextNid  $g$ ) t f, stamp  $g$   $nid$ )
    (add-node (nextNid  $g$ ) ((ConstantNode (bool-to-val-width1 val)), constantA-
sStamp (bool-to-val-width1 val))  $g$ ) |
  constantCondition cond  $nid - g = g$ 

lemma constantConditionTrue:

```

assumes $\text{kind } g \text{ ifcond} = \text{IfNode cond } t \text{ } f$
assumes $g' = \text{constantCondition True ifcond (kind } g \text{ ifcond) } g$
shows $g', p \vdash (\text{ifcond}, m, h) \rightarrow (t, m, h)$
proof –
have $\text{ifn}: \bigwedge c \text{ } t \text{ } f. \text{IfNode } c \text{ } t \text{ } f \neq \text{NoNode}$
by *simp*
then have $\text{if}': \text{kind } g' \text{ ifcond} = \text{IfNode (nextNid } g) \text{ } t \text{ } f$
using $\text{assms}(1) \text{ assms}(2) \text{ constantCondition.simps}(1) \text{ replace-node-lookup}$
by *presburger*
have $\text{truedef}: \text{bool-to-val True} = (\text{IntVal } 32 \text{ } 1)$
by *auto*
from ifn **have** $\text{ifcond} \neq (\text{nextNid } g)$
by $(\text{metis assms}(1) \text{ emptyE ids-some nextNidNotIn})$
moreover have $\bigwedge c. \text{ConstantNode } c \neq \text{NoNode}$ **by** *simp*
ultimately have $\text{kind } g' (\text{nextNid } g) = \text{ConstantNode (bool-to-val-width1 True)}$
using $\text{add-changed add-node-def assms}(1) \text{ assms}(2) \text{ constantCondition.simps}(1)$
not-in-g other-node-unchanged replace-node-def replace-node-lookup singletonD
by $(\text{smt (z3) find-new-kind replace-node-unchanged})$
then have $c': \text{kind } g' (\text{nextNid } g) = \text{ConstantNode (IntVal } 1 \text{ } 1)$
using truedef **by** *simp*
have $\text{valid-value (IntVal } 1 \text{ } 1) (\text{constantAsStamp (IntVal } 1 \text{ } 1))$
unfolding $\text{constantAsStamp.simps valid-value.simps}$
using *nat-numeral* **by** *force*
then have $[g', m, p] \vdash \text{nextNid } g \mapsto \text{IntVal } 1 \text{ } 1$
using $\text{ConstantExpr ConstantNode Value.distinct}(1) \langle \text{kind } g' (\text{nextNid } g) = \text{ConstantNode (bool-to-val-width1 True)} \rangle \text{ encodeeval-def truedef}$
by $(\text{metis bool-to-val-width1.simps}(1) \text{ wf-value-def})$
from $\text{if}' \text{ } c'$ **show** $?thesis$ **using** IfNode
by $(\text{metis (no-types, opaque-lifting) val-to-bool.simps}(1) \langle [g', m, p] \vdash \text{nextNid } g \mapsto \text{IntVal } 1 \text{ } 1 \rangle \text{ encodeeval-def zero-neq-one})$
qed

lemma *constantConditionFalse*:

assumes $\text{kind } g \text{ ifcond} = \text{IfNode cond } t \text{ } f$
assumes $g' = \text{constantCondition False ifcond (kind } g \text{ ifcond) } g$
shows $g', p \vdash (\text{ifcond}, m, h) \rightarrow (f, m, h)$
proof –
have $\text{ifn}: \bigwedge c \text{ } t \text{ } f. \text{IfNode } c \text{ } t \text{ } f \neq \text{NoNode}$
by *simp*
then have $\text{if}': \text{kind } g' \text{ ifcond} = \text{IfNode (nextNid } g) \text{ } t \text{ } f$
by $(\text{metis assms}(1) \text{ assms}(2) \text{ constantCondition.simps}(1) \text{ replace-node-lookup})$
have $\text{falsedef}: \text{bool-to-val False} = (\text{IntVal } 32 \text{ } 0)$
by *auto*
from ifn **have** $\text{ifcond} \neq (\text{nextNid } g)$
by $(\text{metis assms}(1) \text{ equals0D ids-some nextNidNotIn})$
moreover have $\bigwedge c. \text{ConstantNode } c \neq \text{NoNode}$ **by** *simp*
ultimately have $\text{kind } g' (\text{nextNid } g) = \text{ConstantNode (bool-to-val-width1 False)}$
by $(\text{smt (z3) add-changed add-node-def assms}(1) \text{ assms}(2) \text{ constantCondition.simps}(1) \text{ find-new-kind not-in-g other-node-unchanged replace-node-def single-})$

```

tonD)
  then have c': kind g' (nextNid g) = ConstantNode (IntVal 1 0)
    using falsedef by simp
  have valid-value (IntVal 1 0) (constantAsStamp (IntVal 1 0))
    unfolding constantAsStamp.simps valid-value.simps
    using nat-numeral by force
  then have [g', m, p] ⊢ nextNid g ↦ IntVal 1 0
    by (meson ConstantExpr ConstantNode c' encodeeval-def wf-value-def)
  from if' c' show ?thesis using IfNode
    by (metis (no-types, opaque-lifting) val-to-bool.simps(1) ⟨[g',m,p] ⊢ nextNid g
    ↦ IntVal 1 0⟩ encodeeval-def)
qed

```

```

lemma diff-forall:
  assumes ∀ n ∈ ids g - {nid}. cond n
  shows ∀ n. n ∈ ids g ∧ n ∉ {nid} ⟶ cond n
  by (meson Diff-iff assms)

```

```

lemma replace-node-changeonly:
  assumes g' = replace-node nid node g
  shows changeonly {nid} g g'
  using assms replace-node-unchanged
  unfolding changeonly.simps using diff-forall
  by (metis add-changed add-node-def changeonly.simps replace-node-def)

```

```

lemma add-node-changeonly:
  assumes g' = add-node nid node g
  shows changeonly {nid} g g'
  by (metis Rep-IRGraph-inverse add-node.rep-eq assms replace-node.rep-eq re-
  place-node-changeonly)

```

```

lemma constantConditionNoEffect:
  assumes ¬(is-IfNode (kind g nid))
  shows g = constantCondition b nid (kind g nid) g
  using assms apply (cases kind g nid)
  using constantCondition.simps
  apply presburger+
  apply (metis is-IfNode-def)
  using constantCondition.simps
  by presburger+

```

```

lemma constantConditionIfNode:
  assumes kind g nid = IfNode cond t f
  shows constantCondition val nid (kind g nid) g =
    replace-node nid (IfNode (nextNid g) t f, stamp g nid)
    (add-node (nextNid g) ((ConstantNode (bool-to-val-width1 val)), constantA-
    sStamp (bool-to-val-width1 val)) g)
  using constantCondition.simps
  by (simp add: assms)

```

```

lemma constantCondition-changeonly:
  assumes  $nid \in ids\ g$ 
  assumes  $g' = constantCondition\ b\ nid\ (kind\ g\ nid)\ g$ 
  shows  $changeonly\ \{nid\}\ g\ g'$ 
proof (cases is-IfNode (kind g nid))
  case True
  have  $nextNid\ g \notin ids\ g$ 
    using nextNidNotIn by (metis emptyE)
  then show ?thesis using assms
  using replace-node-changeonly add-node-changeonly unfolding changeonly.simps
  using True constantCondition.simps(1) is-IfNode-def
  by (metis (no-types, lifting) insert-iff)
next
  case False
  have  $g = g'$ 
    using constantConditionNoEffect
    using False assms(2) by blast
  then show ?thesis by simp
qed

lemma constantConditionNoIf:
  assumes  $\forall\ cond\ t\ f.\ kind\ g\ ifcond \neq IfNode\ cond\ t\ f$ 
  assumes  $g' = constantCondition\ val\ ifcond\ (kind\ g\ ifcond)\ g$ 
  shows  $\exists\ nid'. (g\ m\ p\ h \vdash ifcond \rightsquigarrow nid') \longleftrightarrow (g'\ m\ p\ h \vdash ifcond \rightsquigarrow nid')$ 
proof –
  have  $g' = g$ 
    using assms(2) assms(1)
    using constantConditionNoEffect
    by (metis IRNode.collapse(11))
  then show ?thesis by simp
qed

lemma constantConditionValid:
  assumes  $kind\ g\ ifcond = IfNode\ cond\ t\ f$ 
  assumes  $[g, m, p] \vdash cond \mapsto v$ 
  assumes  $const = val\text{-}to\text{-}bool\ v$ 
  assumes  $g' = constantCondition\ const\ ifcond\ (kind\ g\ ifcond)\ g$ 
  shows  $\exists\ nid'. (g\ m\ p\ h \vdash ifcond \rightsquigarrow nid') \longleftrightarrow (g'\ m\ p\ h \vdash ifcond \rightsquigarrow nid')$ 
proof (cases const)
  case True
  have ifstep:  $g, p \vdash (ifcond, m, h) \rightarrow (t, m, h)$ 
    by (meson IfNode True assms(1) assms(2) assms(3) encodeeval-def)
  have ifstep':  $g', p \vdash (ifcond, m, h) \rightarrow (t, m, h)$ 
    using constantConditionTrue
    using True assms(1) assms(4) by presburger
  from ifstep ifstep' show ?thesis
    using StutterStep by blast

```



```

next
  case False
  have ifstep:  $g, p \vdash (\text{ifcond}, m, h) \rightarrow (f, m, h)$ 
    by (meson IfNode False assms(1) assms(2) assms(3) encodeeval-def)
  have ifstep':  $g', p \vdash (\text{ifcond}, m, h) \rightarrow (f, m, h)$ 
    using constantConditionFalse
  using False assms(1) assms(4) by presburger
  from ifstep ifstep' show ?thesis
  using StutterStep by blast
qed

end

```

9.3 Stuttering

```

theory Stuttering
  imports
    Semantics.IRStepThms
begin

```

```

inductive stutter:: IRGraph  $\Rightarrow$  MapState  $\Rightarrow$  Params  $\Rightarrow$  FieldRefHeap  $\Rightarrow$  ID  $\Rightarrow$ 
  ID  $\Rightarrow$  bool (- - -  $\vdash$  -  $\rightsquigarrow$  - 55)
  for g m p h where

```

```

  StutterStep:
     $\llbracket g, p \vdash (nid, m, h) \rightarrow (nid', m, h) \rrbracket$ 
     $\implies g \ m \ p \ h \vdash nid \rightsquigarrow nid' \mid$ 

```

```

  Transitive:
     $\llbracket g, p \vdash (nid, m, h) \rightarrow (nid'', m, h);$ 
     $g \ m \ p \ h \vdash nid'' \rightsquigarrow nid' \rrbracket$ 
     $\implies g \ m \ p \ h \vdash nid \rightsquigarrow nid'$ 

```

```

lemma stuttering-successor:
  assumes  $(g, p \vdash (nid, m, h) \rightarrow (nid', m, h))$ 
  shows  $\{P'. (g \ m \ p \ h \vdash nid \rightsquigarrow P')\} = \{nid'\} \cup \{nid''. (g \ m \ p \ h \vdash nid' \rightsquigarrow nid'')\}$ 
proof -
  have nextin:  $nid' \in \{P'. (g \ m \ p \ h \vdash nid \rightsquigarrow P')\}$ 
    using assms StutterStep by blast
  have nextsubset:  $\{nid''. (g \ m \ p \ h \vdash nid' \rightsquigarrow nid'')\} \subseteq \{P'. (g \ m \ p \ h \vdash nid \rightsquigarrow P')\}$ 
    by (metis Collect-mono assms stutter.Transitive)
  have  $\forall n \in \{P'. (g \ m \ p \ h \vdash nid \rightsquigarrow P')\} . n = nid' \vee n \in \{nid''. (g \ m \ p \ h \vdash nid' \rightsquigarrow nid'')\}$ 
    using stepDet
    by (metis (no-types, lifting) Pair-inject assms mem-Collect-eq stutter.simps)
  then show ?thesis
    using insert-absorb mk-disjoint-insert nextin nextsubset by auto
qed

```

end

9.4 Evaluation Stamp Theorems

```
theory StampEvalThms
  imports Graph.ValueThms
           Semantics.IRTreeEvalThms
begin
```

```
lemma
  assumes take-bit b v = v
  shows signed-take-bit b v = v
  using assms
  by (metis(full-types) eq-imp-le signed-take-bit-take-bit)
```

```
lemma unwrap-signed-take-bit:
  fixes v :: int64
  assumes 0 < b ∧ b ≤ 64
  assumes signed-take-bit (b - 1) v = v
  shows signed-take-bit 63 (Word.rep (signed-take-bit (b - Suc 0) v)) = sint v
  using assms using size64 unfolding signed-def by auto
```

```
lemma unrestricted-new-int-always-valid [simp]:
  assumes 0 < b ∧ b ≤ 64
  shows valid-value (new-int b v) (unrestricted-stamp (IntegerStamp b lo hi))
  unfolding unrestricted-stamp.simps new-int.simps valid-value.simps
  by (simp; metis One-nat-def assms int-power-div-base int-signed-value.simps
    int-signed-value-range linorder-not-le not-exp-less-eq-0-int zero-less-numeral)
```

```
lemma unary-undef: val = UndefVal ⇒ unary-eval op val = UndefVal
  by (cases op; auto)
```

```
lemma unary-obj: val = ObjRef x ⇒ unary-eval op val = UndefVal
  by (cases op; auto)
```

```
lemma unrestricted-stamp-valid:
  assumes s = unrestricted-stamp (IntegerStamp b lo hi)
  assumes 0 < b ∧ b ≤ 64
  shows valid-stamp s
  using assms
  by (smt (z3) Stamp.inject(1) bit-bounds.simps not-exp-less-eq-0-int prod.sel(1)
    prod.sel(2) unrestricted-stamp.simps(2) upper-bounds-equiv valid-stamp.elims(1))
```

```
lemma unrestricted-stamp-valid-value [simp]:
  assumes 1: result = IntVal b ival
  assumes take-bit b ival = ival
  assumes 0 < b ∧ b ≤ 64
```

```

  shows valid-value result (unrestricted-stamp (IntegerStamp b lo hi))
proof -
  have valid-stamp (unrestricted-stamp (IntegerStamp b lo hi))
    using assms unrestricted-stamp-valid by blast
  then show ?thesis
    unfolding 1 unrestricted-stamp.simps valid-value.simps
    using assms int-signed-value-bounds by presburger
qed

```

9.4.1 Support Lemmas for Integer Stamps and Associated IntVal values

Valid int implies some useful facts.

```

lemma valid-int-gives:
  assumes valid-value (IntVal b val) stamp
  obtains lo hi where stamp = IntegerStamp b lo hi ∧
    valid-stamp (IntegerStamp b lo hi) ∧
    take-bit b val = val ∧
    lo ≤ int-signed-value b val ∧ int-signed-value b val ≤ hi
  using assms
  by (smt (z3) Value.distinct(7) Value.inject(1) valid-value.elims(1))

```

And the corresponding lemma where we know the stamp rather than the value.

```

lemma valid-int-stamp-gives:
  assumes valid-value val (IntegerStamp b lo hi)
  obtains ival where val = IntVal b ival ∧
    valid-stamp (IntegerStamp b lo hi) ∧
    take-bit b ival = ival ∧
    lo ≤ int-signed-value b ival ∧ int-signed-value b ival ≤ hi
  by (metis assms valid-int valid-value.simps(1))

```

A valid int must have the expected number of bits.

```

lemma valid-int-same-bits:
  assumes valid-value (IntVal b val) (IntegerStamp bits lo hi)
  shows b = bits
  by (meson assms valid-value.simps(1))

```

A valid value means a valid stamp.

```

lemma valid-int-valid-stamp:
  assumes valid-value (IntVal b val) (IntegerStamp bits lo hi)
  shows valid-stamp (IntegerStamp bits lo hi)
  by (metis assms valid-value.simps(1))

```

A valid int means a valid non-empty stamp.

```

lemma valid-int-not-empty:
  assumes valid-value (IntVal b val) (IntegerStamp bits lo hi)

```

shows $lo \leq hi$
by (*metis* *assms* *order.trans* *valid-value.simps*(1))

A valid int fits into the given number of bits (and other bits are zero).

lemma *valid-int-fits*:
assumes *valid-value* (*IntVal* *b val*) (*IntegerStamp* *bits lo hi*)
shows *take-bit* *bits val* = *val*
by (*metis* *assms* *valid-value.simps*(1))

lemma *valid-int-is-zero-masked*:
assumes *valid-value* (*IntVal* *b val*) (*IntegerStamp* *bits lo hi*)
shows *and* *val* (*not* (*mask* *bits*)) = 0
by (*metis* (*no-types*, *lifting*) *assms* *bit.conj-cancel-right* *take-bit-eq-mask* *valid-int-fits*
word-bw-assocs(1) *word-log-esimps*(1))

Unsigned ints have bounds 0 up to 2^{bits} .

lemma *valid-int-unsigned-bounds*:
assumes *valid-value* (*IntVal* *b val*) (*IntegerStamp* *bits lo hi*)
shows *uint* *val* < 2^{bits}
by (*metis* *assms*(1) *mask-eq-iff* *take-bit-eq-mask* *valid-value.simps*(1))

Signed ints have the usual two-complement bounds.

lemma *valid-int-signed-upper-bound*:
assumes *valid-value* (*IntVal* *b val*) (*IntegerStamp* *bits lo hi*)
shows *int-signed-value* *bits val* < 2^{bits-1}
by (*metis* (*mono-tags*, *opaque-lifting*) *diff-le-mono* *int-signed-value.simps* *less-imp-diff-less*

linorder-not-le *one-le-numeral* *order-less-le-trans* *power-increasing* *signed-take-bit-int-less-exp-word* *sint-lt*)

lemma *valid-int-signed-lower-bound*:
assumes *valid-value* (*IntVal* *b val*) (*IntegerStamp* *bits lo hi*)
shows $-(2^{bits-1}) \leq \text{int-signed-value } bits \text{ } val$
by (*smt* (*verit*) *diff-le-self* *int-signed-value.simps* *linorder-not-less* *power-increasing-iff* *signed-take-bit-int-greater-eq-minus-exp-word* *sint-greater-eq*)

and *bit_bounds* versions of the above bounds.

lemma *valid-int-signed-upper-bit-bound*:
assumes *valid-value* (*IntVal* *b val*) (*IntegerStamp* *bits lo hi*)
shows *int-signed-value* *bits val* ≤ *snd* (*bit-bounds* *bits*)
proof –
have *b* = *bits* **using** *assms* *valid-int-same-bits* **by** *blast*
then show *?thesis*
using *assms* **by** *force*
qed

```

lemma valid-int-signed-lower-bit-bound:
  assumes valid-value (IntVal b val) (IntegerStamp bits lo hi)
  shows fst (bit-bounds bits)  $\leq$  int-signed-value bits val
proof –
  have b = bits using assms valid-int-same-bits by blast
  then show ?thesis
    using assms by force
qed

```

Valid values satisfy their stamp bounds.

```

lemma valid-int-signed-range:
  assumes valid-value (IntVal b val) (IntegerStamp bits lo hi)
  shows lo  $\leq$  int-signed-value bits val  $\wedge$  int-signed-value bits val  $\leq$  hi
  by (metis assms valid-value.simps(1))

```

9.4.2 Validity of all Unary Operators

We split the validity proof for unary operators into two lemmas, one for normal unary operators whose output bits equals their input bits, and the other case for the widen and narrow operators.

```

lemma eval-normal-unary-implies-valid-value:
  assumes [m,p]  $\vdash$  expr  $\mapsto$  val
  assumes result = unary-eval op val
  assumes op: op  $\in$  normal-unary
  assumes result  $\neq$  UndefVal
  assumes valid-value val (stamp-expr expr)
  shows valid-value result (stamp-expr (UnaryExpr op expr))
proof –
  obtain b1 v1 where v1: val = IntVal b1 v1
  by (metis Value.exhaust assms(1) assms(2) assms(4) assms(5) evaltree-not-undef
unary-obj valid-value.simps(11))
  then obtain b2 v2 where v2: result = IntVal b2 v2
  using assms(2) assms(4) is-IntVal-def unary-eval-int by presburger
  then have result = unary-eval op (IntVal b1 v1)
  using assms(2) v1 by blast
  then obtain vtmp where vtmp: result = new-int b2 vtmp
  using assms(3) v2 by auto
  obtain b' lo' hi' where stamp-expr expr = IntegerStamp b' lo' hi'
  by (metis assms(5) v1 valid-int-gives)
  then have stamp-unary op (stamp-expr expr) =
    unrestricted-stamp
      (IntegerStamp (if op  $\in$  normal-unary then b' else ir-resultBits op) lo' hi')
  using stamp-unary.simps(1) by presburger
  then obtain lo2 hi2 where s: (stamp-expr (UnaryExpr op expr)) = unrestricted-stamp
    (IntegerStamp b2 lo2 hi2)
  unfolding stamp-expr.simps
  using vtmp op
  by (smt (verit, best) Value.inject(1)  $\langle$ (result::Value) = unary-eval (op::IRUnaryOp)
```

```

(lo'::int) (hi'::int) › assms(5) insertE intval-abs.simps(1) intval-logic-negation.simps(1)
intval-negate.simps(1) intval-not.simps(1) new-int.elims singleton-iff unary-eval.simps(1)
unary-eval.simps(2) unary-eval.simps(3) unary-eval.simps(4) v1 valid-int-same-bits)
  then have 0 < b1 ∧ b1 ≤ 64
    using valid-int-gives
    by (metis assms(5) v1 valid-stamp.simps(1))
  then have fst (bit-bounds b2) ≤ int-signed-value b2 v2 ∧
    int-signed-value b2 v2 ≤ snd (bit-bounds b2)
    by (smt (verit, del-insts) Stamp.inject(1) assms(3) assms(5) int-signed-value-bounds
s stamp-expr.simps(1) stamp-unary.simps(1) unrestricted-stamp.simps(2) v1 valid-int-gives)
  then show ?thesis
    unfolding s v2 unrestricted-stamp.simps valid-value.simps
    by (smt (z3) assms(3) assms(5) is-stamp-empty.simps(1) new-int-take-bits s
stamp-expr.simps(1) stamp-unary.simps(1) unrestricted-stamp.simps(2) v1 v2 valid-int-gives
valid-stamp.simps(1) vtmp)
qed

```

lemma narrow-widen-output-bits:

```

assumes unary-eval op val ≠ UndefVal
assumes op ∉ normal-unary
shows 0 < (ir-resultBits op) ∧ (ir-resultBits op) ≤ 64
proof –
  consider ib ob where op = UnaryNarrow ib ob
    | ib ob where op = UnarySignExtend ib ob
    | ib ob where op = UnaryZeroExtend ib ob
  using IRUnaryOp.exhaust-sel assms(2) by blast
  then show ?thesis
  proof (cases)
    case 1
    then show ?thesis using assms intval-narrow-ok by force
  next
    case 2
    then show ?thesis using assms intval-sign-extend-ok by force
  next
    case 3
    then show ?thesis using assms intval-zero-extend-ok by force
  qed
qed

```

lemma eval-widen-narrow-unary-implies-valid-value:

```

assumes [m,p] ⊢ expr ↦ val
assumes result = unary-eval op val
assumes op: op ∉ normal-unary
assumes result ≠ UndefVal
assumes valid-value val (stamp-expr expr)
shows valid-value result (stamp-expr (UnaryExpr op expr))
proof –
  obtain b1 v1 where v1: val = IntVal b1 v1

```

```

  by (metis Value.exhaust assms(1) assms(2) assms(4) assms(5) evaltree-not-undef
    unary-obj valid-value.simps(11))
  then have result = unary-eval op (IntVal b1 v1)
    using assms(2) v1 by blast
  then obtain v2 where v2: result = new-int (ir-resultBits op) v2
    using assms by (cases op; simp; (meson new-int.simps)+)
  then obtain v3 where v3: result = IntVal (ir-resultBits op) v3
    using assms by (cases op; simp; (meson new-int.simps)+)
  then obtain lo2 hi2 where s: (stamp-expr (UnaryExpr op expr)) = unre-
    stricted-stamp (IntegerStamp (ir-resultBits op) lo2 hi2)
    unfolding stamp-expr.simps stamp-unary.simps
    using assms(3) assms(5) v1 valid-int-gives by fastforce
  then have outBits: 0 < (ir-resultBits op) ∧ (ir-resultBits op) ≤ 64
    using assms narrow-widen-output-bits
    by blast
  then have fst (bit-bounds (ir-resultBits op)) ≤ int-signed-value (ir-resultBits op)
    v3 ∧
    int-signed-value (ir-resultBits op) v3 ≤ snd (bit-bounds (ir-resultBits op))
    using int-signed-value-bounds
    by (smt (verit, del-insts) Stamp.inject(1) assms(3) assms(5) int-signed-value-bounds
      s stamp-expr.simps(1) stamp-unary.simps(1) unrestricted-stamp.simps(2) v1 valid-int-gives)
  then show ?thesis
    unfolding s v3 unrestricted-stamp.simps valid-value.simps
    using outBits v2 v3 by auto
qed

```

```

lemma eval-unary-implies-valid-value:
  assumes [m,p] ⊢ expr ↦ val
  assumes result = unary-eval op val
  assumes result ≠ UndefVal
  assumes valid-value val (stamp-expr expr)
  shows valid-value result (stamp-expr (UnaryExpr op expr))
  proof (cases op ∈ normal-unary)
    case True
    then show ?thesis by (metis assms eval-normal-unary-implies-valid-value)
  next
    case False
    then show ?thesis by (metis assms eval-widen-narrow-unary-implies-valid-value)
  qed

```

9.4.3 Support Lemmas for Binary Operators

```

lemma binary-undef: v1 = UndefVal ∨ v2 = UndefVal ⇒ bin-eval op v1 v2 =
  UndefVal
  by (cases op; auto)

```

```

lemma binary-obj: v1 = ObjRef x ∨ v2 = ObjRef y ⇒ bin-eval op v1 v2 =
  UndefVal
  by (cases op; auto)

```

Some lemmas about the three different output sizes for binary operators.

lemma *bin-eval-bits-binary-shift-ops*:

assumes $result = bin_eval\ op\ (IntVal\ b1\ v1)\ (IntVal\ b2\ v2)$
assumes $result \neq UndefinedVal$
assumes $op \in binary_shift_ops$
shows $\exists v. result = new_int\ b1\ v$
using *assms*
by (*cases op; simp; smt (verit, best) new-int.simps*)**+**

lemma *bin-eval-bits-fixed-32-ops*:

assumes $result = bin_eval\ op\ (IntVal\ b1\ v1)\ (IntVal\ b2\ v2)$
assumes $result \neq UndefinedVal$
assumes $op \in binary_fixed_32_ops$
shows $\exists v. result = new_int\ 32\ v$
using *assms*
apply (*cases op; simp*)
using *assms bool-to-val.simps bin-eval-new-int new-int.simps bin-eval-unused-bits-zero*
by *metis***+**

lemma *bin-eval-bits-normal-ops*:

assumes $result = bin_eval\ op\ (IntVal\ b1\ v1)\ (IntVal\ b2\ v2)$
assumes $result \neq UndefinedVal$
assumes $op \notin binary_shift_ops$
assumes $op \notin binary_fixed_32_ops$
shows $\exists v. result = new_int\ b1\ v$
using *assms* **apply** (*cases op; simp*)
using *assms* **apply** (*metis (mono-tags)*)**+**
using *take-bit-and* **apply** *metis*
using *take-bit-or* **apply** *metis*
using *take-bit-xor* **by** *metis*

lemma *bin-eval-input-bits-equal*:

assumes $result = bin_eval\ op\ (IntVal\ b1\ v1)\ (IntVal\ b2\ v2)$
assumes $result \neq UndefinedVal$
assumes $op \notin binary_shift_ops$
shows $b1 = b2$
using *assms* **apply** (*cases op; simp*)
by *presburger***+**

lemma *bin-eval-implies-valid-value*:

assumes $[m,p] \vdash expr1 \mapsto val1$
assumes $[m,p] \vdash expr2 \mapsto val2$
assumes $result = bin_eval\ op\ val1\ val2$
assumes $result \neq UndefinedVal$
assumes $valid_value\ val1\ (stamp_expr\ expr1)$
assumes $valid_value\ val2\ (stamp_expr\ expr2)$
shows $valid_value\ result\ (stamp_expr\ (BinaryExpr\ op\ expr1\ expr2))$

proof –


```

obtain b1 v1 where v1: val1 = IntVal b1 v1
  by (metis Value.collapse(1) assms(3) assms(4) bin-eval-inputs-are-ints bin-eval-int)
obtain b2 v2 where v2: val2 = IntVal b2 v2
  by (metis Value.collapse(1) assms(3) assms(4) bin-eval-inputs-are-ints bin-eval-int)
then obtain lo1 hi1 where s1: stamp-expr expr1 = IntegerStamp b1 lo1 hi1
  by (metis assms(5) v1 valid-int-gives)
then obtain lo2 hi2 where s2: stamp-expr expr2 = IntegerStamp b2 lo2 hi2
  by (metis assms(6) v2 valid-int-gives)
then have r: result = bin-eval op (IntVal b1 v1) (IntVal b2 v2)
  using assms(3) v1 v2 by blast
then obtain bres vtmp where vtmp: result = new-int bres vtmp
  using assms bin-eval-bits-binary-shift-ops
  by (meson bin-eval-new-int)
then obtain vres where vres: result = IntVal bres vres
  by force

then have sres: stamp-expr (BinaryExpr op expr1 expr2) =
  unrestricted-stamp (IntegerStamp bres lo1 hi1)
  ∧ 0 < bres ∧ bres ≤ 64
proof (cases op ∈ binary-shift-ops)
  case True
  then show ?thesis
    unfolding s1 s2 stamp-binary.simps stamp-expr.simps
    using assms bin-eval-bits-binary-shift-ops
    by (metis Value.inject(1) eval-bits-1-64 new-int.simps r v1 vres)
  next
  case False
  then have op ∉ binary-shift-ops
    by simp
  then have beq: b1 = b2
    using v1 v2 assms bin-eval-input-bits-equal by simp
  then show ?thesis
    proof (cases op ∈ binary-fixed-32-ops)
    case True
    then show ?thesis
      unfolding s1 s2 stamp-binary.simps stamp-expr.simps
      using assms bin-eval-bits-fixed-32-ops
      by (metis False Value.inject(1) beq bin-eval-new-int le-add-same-cancel1
new-int.simps numeral-Bit0 vres zero-le-numeral zero-less-numeral)
    next
    case False
    then show ?thesis
      unfolding s1 s2 stamp-binary.simps stamp-expr.simps
      using assms
      by (metis beq bin-eval-new-int eval-bits-1-64 intval-bits.simps unrestricted-new-int-always-valid
unrestricted-stamp.simps(2) v1 valid-int-same-bits vres)
    qed
  qed
then show ?thesis

```

```

    unfolding vres
    using unrestricted-new-int-always-valid vres vtmp by presburger
qed

```

9.4.4 Validity of Stamp Meet and Join Operators

```

lemma stamp-meet-integer-is-valid-stamp:
  assumes valid-stamp stamp1
  assumes valid-stamp stamp2
  assumes is-IntegerStamp stamp1
  assumes is-IntegerStamp stamp2
  shows valid-stamp (meet stamp1 stamp2)
  using assms unfolding is-IntegerStamp-def valid-stamp.simps meet.simps
  by (smt (verit, del-insts) meet.simps(2) valid-stamp.simps(1) valid-stamp.simps(8))

```

```

lemma stamp-meet-is-valid-stamp:
  assumes 1: valid-stamp stamp1
  assumes 2: valid-stamp stamp2
  shows valid-stamp (meet stamp1 stamp2)
  by (cases stamp1; cases stamp2; insert stamp-meet-integer-is-valid-stamp[OF 1 2]; auto)

```

```

lemma stamp-meet-commutes: meet stamp1 stamp2 = meet stamp2 stamp1
  by (cases stamp1; cases stamp2; auto)

```

```

lemma stamp-meet-is-valid-value1:
  assumes valid-value val stamp1
  assumes valid-stamp stamp2
  assumes stamp1 = IntegerStamp b1 lo1 hi1
  assumes stamp2 = IntegerStamp b2 lo2 hi2
  assumes meet stamp1 stamp2 ≠ IllegalStamp
  shows valid-value val (meet stamp1 stamp2)

```

proof –

```

  have m: meet stamp1 stamp2 = IntegerStamp b1 (min lo1 lo2) (max hi1 hi2)
    using assms by (metis meet.simps(2))
  obtain ival where val: val = IntVal b1 ival
    using assms valid-int by blast
  then have v: valid-stamp (IntegerStamp b1 lo1 hi1) ∧
    take-bit b1 ival = ival ∧
    lo1 ≤ int-signed-value b1 ival ∧ int-signed-value b1 ival ≤ hi1
    using assms by (metis valid-value.simps(1))
  then have mm: min lo1 lo2 ≤ int-signed-value b1 ival ∧ int-signed-value b1 ival
    ≤ max hi1 hi2
    by linarith
  then have valid-stamp (IntegerStamp b1 (min lo1 lo2) (max hi1 hi2))
    using assms v stamp-meet-is-valid-stamp
    by (metis meet.simps(2))
  then show ?thesis

```

```

    unfolding m val valid-value.simps
    using mm v by presburger
qed

```

and the symmetric lemma follows by the commutativity of meet.

```

lemma stamp-meet-is-valid-value:
  assumes valid-value val stamp2
  assumes valid-stamp stamp1
  assumes stamp1 = IntegerStamp b1 lo1 hi1
  assumes stamp2 = IntegerStamp b2 lo2 hi2
  assumes meet stamp1 stamp2 ≠ IllegalStamp
  shows valid-value val (meet stamp1 stamp2)
  using assms stamp-meet-commutes stamp-meet-is-valid-value1
  by metis

```

9.4.5 Validity of conditional expressions

```

lemma conditional-eval-implies-valid-value:
  assumes [m,p] ⊢ cond ↦ condv
  assumes expr = (if val-to-bool condv then expr1 else expr2)
  assumes [m,p] ⊢ expr ↦ val
  assumes val ≠ UndefinedVal
  assumes valid-value condv (stamp-expr cond)
  assumes valid-value val (stamp-expr expr)
  assumes compatible (stamp-expr expr1) (stamp-expr expr2)
  shows valid-value val (stamp-expr (ConditionalExpr cond expr1 expr2))
proof -
  have def: meet (stamp-expr expr1) (stamp-expr expr2) ≠ IllegalStamp
    using assms
    by (metis Stamp.distinct(13) Stamp.distinct(25) compatible.elims(2) meet.simps(1)
    meet.simps(2))
  then have valid-stamp (meet (stamp-expr expr1) (stamp-expr expr2))
    using assms
    by (smt (verit, best) compatible.elims(2) stamp-meet-is-valid-stamp valid-stamp.simps(2))

  then show ?thesis using stamp-meet-is-valid-value
    using assms def
    by (smt (verit, best) compatible.elims(2) never-void stamp-expr.simps(6) stamp-meet-commutes)
qed

```

9.4.6 Validity of Whole Expression Tree Evaluation

TODO: find a way to encode that conditional expressions must have compatible (and valid) stamps? One approach would be for all the stamp_expr operators to require that all input stamps are valid.

```

definition wf-stamp :: IRExp ⇒ bool where
  wf-stamp e = (∀ m p v. ([m, p] ⊢ e ↦ v) ⟶ valid-value v (stamp-expr e))

```

lemma *stamp-under-defn*:

```

assumes stamp-under (stamp-expr x) (stamp-expr y)
assumes wf-stamp x  $\wedge$  wf-stamp y
assumes ( $[m, p] \vdash x \mapsto xv$ )  $\wedge$  ( $[m, p] \vdash y \mapsto yv$ )
shows val-to-bool (bin-eval BinIntegerLessThan xv yv)  $\vee$  (bin-eval BinIntegerLessThan xv yv) = UndefVal
proof –
  have yval: valid-value yv (stamp-expr y)
    using assms wf-stamp-def by blast
  obtain b lx hi where xstamp: stamp-expr x = IntegerStamp b lx hi
    using assms(1)
    by (metis stamp-under.elims(2))
  then obtain b' lo hy where ystamp: stamp-expr y = IntegerStamp b' lo hy
    using assms(1)
    by (meson stamp-under.elims(2))
  obtain xvv where xvv: xv = IntVal b xvv
    by (metis assms(2) assms(3) valid-int wf-stamp-def xstamp)
  then have xval: valid-value (IntVal b xvv) (stamp-expr x)
    using assms(2) assms(3) wf-stamp-def by blast
  obtain yvv where yvv: yv = IntVal b' yvv
    by (metis valid-int ystamp yval)
  then have xval: valid-value (IntVal b' yvv) (stamp-expr y)
    using yval
    by blast
  have xunder: int-signed-value b xvv  $\leq$  hi
    using xvv xval valid-value.simps
    by (metis assms(2) assms(3) wf-stamp-def xstamp)
  have yunder: lo  $\leq$  int-signed-value b' yvv
    using yvv yval valid-value.simps
    by (metis ystamp)
  have unwrap:  $\forall$  cond. bool-to-val-bin b b cond = bool-to-val cond
    by simp
  from xunder yunder have int-signed-value b xvv  $<$  int-signed-value b' yvv
    using assms(1) xstamp ystamp by auto
  then have (intval-less-than xv yv) = IntVal 32 1  $\vee$  (intval-less-than xv yv) =
UndefVal
    using xvv yvv
    using intval-less-than.simps(1) unwrap
    using bool-to-val.simps(1)
    by simp
  then show ?thesis
    by force
qed

```

lemma *stamp-under-defn-inverse*:

```

assumes stamp-under (stamp-expr y) (stamp-expr x)
assumes wf-stamp x  $\wedge$  wf-stamp y
assumes ( $[m, p] \vdash x \mapsto xv$ )  $\wedge$  ( $[m, p] \vdash y \mapsto yv$ )

```

```

  shows  $\neg(\text{val-to-bool } (\text{bin-eval BinIntegerLessThan } xv \ yv)) \vee (\text{bin-eval BinIntegerLessThan } xv \ yv) = \text{UndefVal}$ 
proof -
  have  $yval$ :  $\text{valid-value } yv \ (\text{stamp-expr } y)$ 
    using  $\text{assms wf-stamp-def by blast}$ 
  obtain  $b \ lo \ hx$  where  $xstamp$ :  $\text{stamp-expr } x = \text{IntegerStamp } b \ lo \ hx$ 
    using  $\text{assms}(1)$ 
    by ( $\text{metis stamp-under.elims}(2)$ )
  then obtain  $b' \ ly \ hi$  where  $ystamp$ :  $\text{stamp-expr } y = \text{IntegerStamp } b' \ ly \ hi$ 
    using  $\text{assms}(1)$ 
    by ( $\text{meson stamp-under.elims}(2)$ )
  obtain  $xvv$  where  $xvv$ :  $xv = \text{IntVal } b \ xvv$ 
    by ( $\text{metis assms}(2) \ \text{assms}(3) \ \text{valid-int wf-stamp-def } xstamp$ )
  then have  $xval$ :  $\text{valid-value } (\text{IntVal } b \ xvv) \ (\text{stamp-expr } x)$ 
    using  $\text{assms}(2) \ \text{assms}(3) \ \text{wf-stamp-def by blast}$ 
  obtain  $yvv$  where  $yvv$ :  $yv = \text{IntVal } b' \ yvv$ 
    by ( $\text{metis valid-int } ystamp \ yval$ )
  then have  $xval$ :  $\text{valid-value } (\text{IntVal } b' \ yvv) \ (\text{stamp-expr } y)$ 
    using  $yval \ \text{by auto}$ 
  have  $yunder$ :  $\text{int-signed-value } b' \ yvv \leq hi$ 
    using  $yvv \ yval \ \text{valid-value.simps}$ 
    by ( $\text{metis } ystamp$ )
  have  $xover$ :  $lo \leq \text{int-signed-value } b \ xvv$ 
    using  $xvv \ xval \ \text{valid-value.simps}$ 
    by ( $\text{metis assms}(2) \ \text{assms}(3) \ \text{wf-stamp-def } xstamp$ )
  have  $unwrap$ :  $\forall \text{ cond. } \text{bool-to-val-bin } b \ b \ \text{cond} = \text{bool-to-val } \text{cond}$ 
    by  $\text{simp}$ 
  from  $xover \ yunder$  have  $\text{int-signed-value } b' \ yvv < \text{int-signed-value } b \ xvv$ 
    using  $\text{assms}(1) \ xstamp \ ystamp \ \text{by auto}$ 
  then have  $(\text{intval-less-than } xv \ yv) = \text{IntVal } 32 \ 0 \ \vee (\text{intval-less-than } xv \ yv) =$ 
 $\text{UndefVal}$ 
    using  $xvv \ yvv$ 
    using  $\text{intval-less-than.simps}(1) \ \text{unwrap by simp}$ 
  then show  $?thesis$ 
    by  $\text{force}$ 
qed

end

```

10 Optization DSL

10.1 Markup

```

theory Markup
  imports Semantics.IRTreeEval Snippets.Snipping
begin

datatype 'a Rewrite =
  Transform 'a 'a ( $- \mapsto - \ 10$ ) |

```

Conditional 'a 'a bool (- \mapsto - when - 11) |
Sequential 'a Rewrite 'a Rewrite |
Transitive 'a Rewrite

datatype 'a ExtraNotation =
ConditionalNotation 'a 'a 'a (- ? - : - 50) |
EqualsNotation 'a 'a (- eq -) |
ConstantNotation 'a (const - 120) |
TrueNotation (true) |
FalseNotation (false) |
ExclusiveOr 'a 'a (- \oplus -) |
LogicNegationNotation 'a (!-) |
ShortCircuitOr 'a 'a (- || -)

definition word :: ('a::len) word \Rightarrow 'a word **where**
 word x = x

ML-file \langle markup.ML \rangle

10.1.1 Expression Markup

ML \langle
 structure IRExprTranslator : DSL-TRANSLATION =
 struct
 fun markup DSL-Tokens.Add = @{term BinaryExpr} \$ @{term BinAdd}
 | markup DSL-Tokens.Sub = @{term BinaryExpr} \$ @{term BinSub}
 | markup DSL-Tokens.Mul = @{term BinaryExpr} \$ @{term BinMul}
 | markup DSL-Tokens.And = @{term BinaryExpr} \$ @{term BinAnd}
 | markup DSL-Tokens.Or = @{term BinaryExpr} \$ @{term BinOr}
 | markup DSL-Tokens.Xor = @{term BinaryExpr} \$ @{term BinXor}
 | markup DSL-Tokens.ShortCircuitOr = @{term BinaryExpr} \$ @{term Bin-ShortCircuitOr}
 | markup DSL-Tokens.Abs = @{term UnaryExpr} \$ @{term UnaryAbs}
 | markup DSL-Tokens.Less = @{term BinaryExpr} \$ @{term BinIntegerLessThan}
 | markup DSL-Tokens.Equals = @{term BinaryExpr} \$ @{term BinIntegerEquals}
 | markup DSL-Tokens.Not = @{term UnaryExpr} \$ @{term UnaryNot}
 | markup DSL-Tokens.Negate = @{term UnaryExpr} \$ @{term UnaryNeg}
 | markup DSL-Tokens.LogicNegate = @{term UnaryExpr} \$ @{term UnaryLogicNegation}
 | markup DSL-Tokens.LeftShift = @{term BinaryExpr} \$ @{term BinLeftShift}
 | markup DSL-Tokens.RightShift = @{term BinaryExpr} \$ @{term BinRightShift}
 | markup DSL-Tokens.UnsignedRightShift = @{term BinaryExpr} \$ @{term Bin-URightShift}
 | markup DSL-Tokens.Conditional = @{term ConditionalExpr}
 | markup DSL-Tokens.Constant = @{term ConstantExpr}
 | markup DSL-Tokens.TrueConstant = @{term ConstantExpr (IntVal 32 1)}
 | markup DSL-Tokens.FalseConstant = @{term ConstantExpr (IntVal 32 0)}
 end
 structure IRExprMarkup = DSL-Markup(IRExprTranslator);

>

ir expression translation

```
syntax -expandExpr :: term ⇒ term (exp[-])  
parse-translation < [( @{syntax-const -expandExpr} , IRExprMarkup.markup-expr []) ] >
```

ir expression example

```
value exp[(e1 < e2) ? e1 : e2]  
  
ConditionalExpr (BinaryExpr BinIntegerLessThan (e1::IRExpr)  
(e2::IRExpr)) e1 e2
```

10.1.2 Value Markup

ML <

```
structure IntValTranslator : DSL-TRANSLATION =  
struct  
  fun markup DSL-Tokens.Add = @{term intval-add}  
    | markup DSL-Tokens.Sub = @{term intval-sub}  
    | markup DSL-Tokens.Mul = @{term intval-mul}  
    | markup DSL-Tokens.And = @{term intval-and}  
    | markup DSL-Tokens.Or = @{term intval-or}  
    | markup DSL-Tokens.ShortCircuitOr = @{term intval-short-circuit-or}  
    | markup DSL-Tokens.Xor = @{term intval-xor}  
    | markup DSL-Tokens.Abs = @{term intval-abs}  
    | markup DSL-Tokens.Less = @{term intval-less-than}  
    | markup DSL-Tokens.Equals = @{term intval-equals}  
    | markup DSL-Tokens.Not = @{term intval-not}  
    | markup DSL-Tokens.Negate = @{term intval-negate}  
    | markup DSL-Tokens.LogicNegate = @{term intval-logic-negation}  
    | markup DSL-Tokens.LeftShift = @{term intval-left-shift}  
    | markup DSL-Tokens.RightShift = @{term intval-right-shift}  
    | markup DSL-Tokens.UnsignedRightShift = @{term intval-uright-shift}  
    | markup DSL-Tokens.Conditional = @{term intval-conditional}  
    | markup DSL-Tokens.Constant = @{term IntVal 32}  
    | markup DSL-Tokens.TrueConstant = @{term IntVal 32 1}  
    | markup DSL-Tokens.FalseConstant = @{term IntVal 32 0}  
end  
structure IntValMarkup = DSL-Markup(IntValTranslator);  
>
```

value expression translation

```
syntax -expandIntVal :: term ⇒ term (val[-])
parse-translation < [( @{syntax-const -expandIntVal} , IntVal-
Markup.markup-expr []) ] >
```

value expression example

```
value val[(e1 < e2) ? e1 : e2]

intval-conditional (intval-less-than (e1::Value) (e2::Value)) e1 e2
```

10.1.3 Word Markup

```
ML <
structure WordTranslator : DSL-TRANSLATION =
struct
fun markup DSL-Tokens.Add = @{term plus}
| markup DSL-Tokens.Sub = @{term minus}
| markup DSL-Tokens.Mul = @{term times}
| markup DSL-Tokens.And = @{term Bit-Operations.semiring-bit-operations-class.and}
| markup DSL-Tokens.Or = @{term or}
| markup DSL-Tokens.Xor = @{term xor}
| markup DSL-Tokens.Abs = @{term abs}
| markup DSL-Tokens.Less = @{term less}
| markup DSL-Tokens.Equals = @{term HOL.eq}
| markup DSL-Tokens.Not = @{term not}
| markup DSL-Tokens.Negate = @{term uminus}
| markup DSL-Tokens.LogicNegate = @{term logic-negate}
| markup DSL-Tokens.LeftShift = @{term shiftl}
| markup DSL-Tokens.RightShift = @{term signed-shiftr}
| markup DSL-Tokens.UnsignedRightShift = @{term shiftr}
| markup DSL-Tokens.Constant = @{term word}
| markup DSL-Tokens.TrueConstant = @{term 1}
| markup DSL-Tokens.FalseConstant = @{term 0}
end
structure WordMarkup = DSL-Markup(WordTranslator);
>
```

word expression translation

```
syntax -expandWord :: term ⇒ term (bin[-])
parse-translation < [( @{syntax-const -expandWord} , Word-
Markup.markup-expr []) ] >
```


word expression example

value *bin*[*x* & *y* | *z*]

intval-conditional (*intval-less-than* (*e*₁::*Value*) (*e*₂::*Value*)) *e*₁ *e*₂

value *bin*[$\neg x$]

value *val*[$\neg x$]

value *exp*[$\neg x$]

value *bin*[!*x*]

value *val*[!*x*]

value *exp*[!*x*]

value *bin*[$\neg x$]

value *val*[$\neg x$]

value *exp*[$\neg x$]

value *bin*[$\sim x$]

value *val*[$\sim x$]

value *exp*[$\sim x$]

value $\sim x$

end

10.2 Optimization Phases

theory *Phase*

imports *Main*

begin

ML-file *map.ML*

ML-file *phase.ML*

end

10.3 Canonicalization DSL

theory *Canonicalization*

imports

Markup

Phase

HOL-Eisbach.Eisbach

keywords

phase :: *thy-decl* **and**

terminating :: *quasi-command* **and**

print-phases :: *diag* **and**

export-phases :: *thy-decl* **and**

```

    optimization :: thy-goal-defn
begin

print-methods

ML <
datatype 'a Rewrite =
  Transform of 'a * 'a |
  Conditional of 'a * 'a * term |
  Sequential of 'a Rewrite * 'a Rewrite |
  Transitive of 'a Rewrite

type rewrite = {
  name: binding,
  rewrite: term Rewrite,
  proofs: thm list,
  code: thm list,
  source: term
}

structure RewriteRule : Rule =
struct
type T = rewrite;

(*
fun pretty-rewrite ctxt (Transform (from, to)) =
  Pretty.block [
    Syntax.pretty-term ctxt from,
    Pretty.str  $\mapsto$  ,
    Syntax.pretty-term ctxt to
  ]
| pretty-rewrite ctxt (Conditional (from, to, cond)) =
  Pretty.block [
    Syntax.pretty-term ctxt from,
    Pretty.str  $\mapsto$  ,
    Syntax.pretty-term ctxt to,
    Pretty.str when ,
    Syntax.pretty-term ctxt cond
  ]
| pretty-rewrite - - = Pretty.str not implemented*)

fun pretty-thm ctxt thm =
  (Proof-Context.pretty-fact ctxt (, [thm]))

fun pretty ctxt obligations t =
  let
    val is-skipped = Thm-Deps.has-skip-proof (#proofs t);

    val warning = (if is-skipped

```

```

    then [Pretty.str (proof skipped), Pretty.brk 0]
    else [];

val obligations = (if obligations
  then [Pretty.big-list
    obligations:
      (map (pretty-thm ctxt) (#proofs t)),
    Pretty.brk 0]
  else []);

fun pretty-bind binding =
  Pretty.markup
    (Position.markup (Binding.pos-of binding) Markup.position)
    [Pretty.str (Binding.name-of binding)];

in
  Pretty.block ([
    pretty-bind (#name t), Pretty.str : ,
    Syntax.pretty-term ctxt (#source t), Pretty.fbrk
  ] @ obligations @ warning)
end
end

structure RewritePhase = DSL-Phase(RewriteRule);

val - =
  Outer-Syntax.command command-keyword⟨phase⟩ enter an optimization phase
  (Parse.binding --| Parse.*** terminating -- Parse.const --| Parse.begin
  >> (Toplevel.begin-main-target true o RewritePhase.setup));

fun print-phases print-obligations ctxt =
  let
    val thy = Proof-Context.theory-of ctxt;
    fun print phase = RewritePhase.pretty print-obligations phase ctxt
  in
    map print (RewritePhase.phases thy)
  end

fun print-optimizations print-obligations thy =
  print-phases print-obligations thy |> Pretty.writeln-chunks

val - =
  Outer-Syntax.command command-keyword⟨print-phases⟩
  print debug information for optimizations
  (Parse.opt-bang >>
    (fn b => Toplevel.keep ((print-optimizations b) o Toplevel.context-of)));

fun export-phases thy name =
  let

```

```

val state = Toplevel.theory-tolevel thy;
val ctxt = Toplevel.context-of state;
val content = Pretty.string-of (Pretty.chunks (print-phases false ctxt));
val cleaned = YXML.content-of content;

val filename = Path.explode (name ^ ".rules");
val directory = Path.explode optimizations;
val path = Path.binding (
  Path.append directory filename,
  Position.none);
val thy' = thy |> Generated-Files.add-files (path, (Bytes.string content));

val - = Export.export thy' path [YXML.parse cleaned];

val - = writeln (Export.message thy' (Path.basic optimizations));
in
  thy'
end

val - =
  Outer-Syntax.command command-keyword ⟨export-phases⟩
  export information about encoded optimizations
  (Parse.path >>
    (fn name => Toplevel.theory (fn state => export-phases state name)))
  ,

```

ML-file *rewrites.ML*

10.3.1 Semantic Preservation Obligation

```

fun rewrite-preservation :: IRExp Rewrite ⇒ bool where
  rewrite-preservation (Transform x y) = (y ≤ x) |
  rewrite-preservation (Conditional x y cond) = (cond ⟶ (y ≤ x)) |
  rewrite-preservation (Sequential x y) = (rewrite-preservation x ∧ rewrite-preservation
y) |
  rewrite-preservation (Transitive x) = rewrite-preservation x

```

10.3.2 Termination Obligation

```

fun rewrite-termination :: IRExp Rewrite ⇒ (IRExp ⇒ nat) ⇒ bool where
  rewrite-termination (Transform x y) trm = (trm x > trm y) |
  rewrite-termination (Conditional x y cond) trm = (cond ⟶ (trm x > trm y)) |
  rewrite-termination (Sequential x y) trm = (rewrite-termination x trm ∧ rewrite-termination
y trm) |
  rewrite-termination (Transitive x) trm = rewrite-termination x trm

```

```

fun intval :: Value Rewrite ⇒ bool where
  intval (Transform x y) = (x ≠ UndefVal ∧ y ≠ UndefVal ⟶ x = y) |
  intval (Conditional x y cond) = (cond ⟶ (x = y)) |

```

$intval (Sequential\ x\ y) = (intval\ x \wedge intval\ y) \mid$
 $intval (Transitive\ x) = intval\ x$

10.3.3 Standard Termination Measure

```

fun size :: IRExpr  $\Rightarrow$  nat where
  unary-size:
    size (UnaryExpr op x) = (size x) + 2 |

  bin-const-size:
    size (BinaryExpr op x (ConstantExpr cy)) = (size x) + 2 |
  bin-size:
    size (BinaryExpr op x y) = (size x) + (size y) + 2 |
  cond-size:
    size (ConditionalExpr c t f) = (size c) + (size t) + (size f) + 2 |
  const-size:
    size (ConstantExpr c) = 1 |
  param-size:
    size (ParameterExpr ind s) = 2 |
  leaf-size:
    size (LeafExpr nid s) = 2 |
    size (ConstantVar c) = 2 |
    size (VariableExpr x s) = 2

```

10.3.4 Automated Tactics

named-theorems *size-simps size simplification rules*

```

method unfold-optimization =
  (unfold rewrite-preservation.simps, unfold rewrite-termination.simps,
   unfold intval.simps,
   rule conjE, simp, simp del: le-expr-def, force?)
| (unfold rewrite-preservation.simps, unfold rewrite-termination.simps,
   rule conjE, simp, simp del: le-expr-def, force?)

```

```

method unfold-size =
  (((unfold size.simps, simp add: size-simps del: le-expr-def)?
   ; (simp add: size-simps del: le-expr-def)?
   ; (auto simp: size-simps)?
   ; (unfold size.simps)?)[1])

```

print-methods

```

ML <
  structure System : RewriteSystem =
  struct
    val preservation = @{const rewrite-preservation};
    val termination = @{const rewrite-termination};
    val intval = @{const intval};

```

```

end

structure DSL = DSL-Rewrites(System);

val - =
  Outer-Syntax.local-theory-to-proof command-keyword ⟨optimization⟩
  define an optimization and open proof obligation
  (Parse-Spec.thm-name : -- Parse.term
   >> DSL.rewrite-cmd);
>

end

```

11 Canonicalization Optimizations

```

theory Common
  imports
    OptimizationDSL.Canonicalization
    Semantics.IRTreeEvalThms
begin

lemma size-pos[size-simps]: 0 < size y
  apply (induction y; auto?)
  by (smt (z3) add-2-eq-Suc' add-is-0 not-gr0 size.elims size.simps(12) size.simps(13)
      size.simps(14) size.simps(15) zero-neq-numeral zero-neq-one)

lemma size-non-add[size-simps]: size (BinaryExpr op a b) = size a + size b + 2
  ⟷ ¬(is-ConstantExpr b)
  by (induction b; induction op; auto simp: is-ConstantExpr-def)

lemma size-non-const[size-simps]:
  ¬ is-ConstantExpr y ⟹ 1 < size y
  using size-pos apply (induction y; auto)
  by (metis Suc-lessI add-is-1 is-ConstantExpr-def le-less linorder-not-le n-not-Suc-n
      numeral-2-eq-2 pos2 size.simps(2) size-non-add)

lemma size-binary-const[size-simps]:
  size (BinaryExpr op a b) = size a + 2 ⟷ (is-ConstantExpr b)
  by (induction b; auto simp: is-ConstantExpr-def size-pos)

lemma size-flip-binary[size-simps]:
  ¬(is-ConstantExpr y) ⟹ size (BinaryExpr op (ConstantExpr x) y) > size
  (BinaryExpr op y (ConstantExpr x))
  by (metis add-Suc not-less-eq order-less-asym plus-1-eq-Suc size.simps(11) size.simps(2)
      size-non-add)

lemma size-binary-lhs-a[size-simps]:
  size (BinaryExpr op (BinaryExpr op' a b) c) > size a
  by (metis add-lessD1 less-add-same-cancel1 pos2 size-binary-const size-non-add)

```

lemma *size-binary-lhs-b*[size-simps]:

size (BinaryExpr op (BinaryExpr op' a b) c) > size b
by (metis *IRExpr.disc(42) One-nat-def add.left-commute add.right-neutral is-ConstantExpr-def less-add-Suc2 numeral-2-eq-2 plus-1-eq-Suc size.simps(11) size-binary-const size-non-add size-non-const trans-less-add1*)

lemma *size-binary-lhs-c*[size-simps]:

size (BinaryExpr op (BinaryExpr op' a b) c) > size c
by (metis *IRExpr.disc(42) add.left-commute add.right-neutral is-ConstantExpr-def less-Suc-eq numeral-2-eq-2 plus-1-eq-Suc size.simps(11) size-non-add size-non-const trans-less-add2*)

lemma *size-binary-rhs-a*[size-simps]:

size (BinaryExpr op c (BinaryExpr op' a b)) > size a
by (smt (verit, best) *less-Suc-eq less-add-Suc2 less-add-same-cancel1 linorder-neqE-nat not-add-less1 order-less-trans pos2 size.simps(4) size-binary-const size-non-add*)

lemma *size-binary-rhs-b*[size-simps]:

size (BinaryExpr op c (BinaryExpr op' a b)) > size b
by (metis *add.left-commute add.right-neutral is-ConstantExpr-def lessI numeral-2-eq-2 plus-1-eq-Suc size.simps(11) size.simps(4) size-non-add trans-less-add2*)

lemma *size-binary-rhs-c*[size-simps]:

size (BinaryExpr op c (BinaryExpr op' a b)) > size c
by *simp*

lemma *size-binary-lhs*[size-simps]:

size (BinaryExpr op x y) > size x
by (metis *One-nat-def Suc-eq-plus1 add-Suc-right less-add-Suc1 numeral-2-eq-2 size-binary-const size-non-add*)

lemma *size-binary-rhs*[size-simps]:

size (BinaryExpr op x y) > size y
by (metis *IRExpr.disc(42) add-strict-increasing is-ConstantExpr-def linorder-not-le not-add-less1 size.simps(11) size-non-add size-non-const size-pos*)

lemmas *arith*[size-simps] = *Suc-leI add-strict-increasing order-less-trans trans-less-add2*

definition *well-formed-equal* :: *Value* \Rightarrow *Value* \Rightarrow *bool*

(*infix* \approx 50) **where**

well-formed-equal $v_1 v_2 = (v_1 \neq \text{UndefVal} \longrightarrow v_1 = v_2)$

lemma *well-formed-equal-defn* [*simp*]:

well-formed-equal $v_1 v_2 = (v_1 \neq \text{UndefVal} \longrightarrow v_1 = v_2)$

unfolding *well-formed-equal-def* **by** *simp*

end

11.1 AbsNode Phase

theory *AbsPhase*

imports

Common

begin

phase *AbsNode*

terminating *size*

begin

lemma *abs-pos*:

fixes $v :: ('a :: \text{len word})$

assumes $0 \leq_s v$

shows $(\text{if } v <_s 0 \text{ then } -v \text{ else } v) = v$

by $(\text{simp add: } \text{assms signed.leD})$

lemma *abs-neg*:

fixes $v :: ('a :: \text{len word})$

assumes $v <_s 0$

assumes $-(2^{\wedge}(\text{Nat.size } v - 1)) <_s v$

shows $(\text{if } v <_s 0 \text{ then } -v \text{ else } v) = -v \wedge 0 <_s -v$

by $(\text{smt } (\text{verit, ccfv-SIG}) \text{ assms(1) assms(2) signed-take-bit-int-greater-eq-minus-exp}$

$\text{signed-take-bit-int-greater-eq-self-iff sint-0 sint-word-ariths(4) word-sless-alt})$

lemma *abs-max-neg*:

fixes $v :: ('a :: \text{len word})$

assumes $v <_s 0$

assumes $-(2^{\wedge}(\text{Nat.size } v - 1)) = v$

shows $-v = v$

using *assms*

by $(\text{metis One-nat-def add.inverse-neutral double-eq-zero-iff mult-minus-right size-word.rep-eq})$

lemma *final-abs*:

fixes $v :: ('a :: \text{len word})$

assumes $\text{take-bit } (\text{Nat.size } v) v = v$

assumes $-(2^{\wedge}(\text{Nat.size } v - 1)) \neq v$

shows $0 \leq_s (\text{if } v <_s 0 \text{ then } -v \text{ else } v)$

proof $(\text{cases } v <_s 0)$

case *True*

then show *?thesis*

proof $(\text{cases } v = -(2^{\wedge}(\text{Nat.size } v - 1)))$

case *True*

then show *?thesis* **using** *abs-max-neg*


```

    using assms by presburger
  next
  case False
  then have  $-(2 \wedge (\text{Nat.size } v - 1)) < s \ v$ 
    unfolding word-sless-def using signed-take-bit-int-greater-self-iff
    by (smt (verit, best) One-nat-def diff-less double-eq-zero-iff len-gt-0 lessI less-irrefl
      mult-minus-right neg-equal-0-iff-equal signed.rep-eq signed-of-int
      signed-take-bit-int-greater-eq-self-iff signed-word-eqI sint-0 sint-range-size
      sint-sbintrunc' sint-word-ariths(4) size-word.rep-eq unsigned-0 word-2p-lem

      word-sless.rep-eq word-sless-def)
  then show ?thesis
    using abs-neg abs-pos signed.nless-le by auto
  qed
next
case False
then show ?thesis using abs-pos by auto
qed

```

```

lemma wf-abs: is-IntVal  $x \implies \text{intval-abs } x \neq \text{UndefVal}$ 
  using intval-abs.simps unfolding new-int.simps
  using is-IntVal-def by force

```

```

fun bin-abs :: 'a :: len word  $\Rightarrow$  'a :: len word where
  bin-abs  $v = (\text{if } (v < s \ 0) \text{ then } (- \ v) \text{ else } v)$ 

```

```

lemma val-abs-zero:
  intval-abs (new-int  $b \ 0$ ) = new-int  $b \ 0$ 
  by simp

```

```

lemma less-eq-zero:
  assumes val-to-bool (val[(IntVal  $b \ 0$ ) < (IntVal  $b \ v$ )])
  shows int-signed-value  $b \ v > 0$ 
  using assms unfolding intval-less-than.simps(1) apply simp
  by (metis bool-to-val.elims val-to-bool.simps(1))

```

```

lemma val-abs-pos:
  assumes val-to-bool(val[(new-int  $b \ 0$ ) < (new-int  $b \ v$ )])
  shows intval-abs (new-int  $b \ v$ ) = (new-int  $b \ v$ )
  using assms using less-eq-zero unfolding intval-abs.simps new-int.simps
  by force

```

```

lemma val-abs-neg:
  assumes val-to-bool(val[(new-int  $b \ v$ ) < (new-int  $b \ 0$ )])

```

```

shows intval-abs (new-int b v) = intval-negate (new-int b v)
using assms using less-eq-zero unfolding intval-abs.simps new-int.simps
by force

lemma val-bool-unwrap:
  val-to-bool (bool-to-val v) = v
by (metis bool-to-val.elims one-neq-zero val-to-bool.simps(1))

lemma take-bit-unwrap:
  b = 64  $\implies$  take-bit b (v1::64 word) = v1
by (metis size64 size-word.rep-eq take-bit-length-eq)

lemma bit-less-eq-def:
  fixes v1 v2 :: 64 word
  assumes b ≤ 64
  shows sint (signed-take-bit (b - Suc (0::nat)) (take-bit b v1))
    < sint (signed-take-bit (b - Suc (0::nat)) (take-bit b v2))  $\longleftrightarrow$ 
    signed-take-bit (63::nat) (Word.rep v1) < signed-take-bit (63::nat) (Word.rep
v2)
  using assms sorry

lemma less-eq-def:

  shows val-to-bool(val[(new-int b v1) < (new-int b v2)])  $\longleftrightarrow$  v1 < s v2
  unfolding new-int.simps intval-less-than.simps bool-to-val-bin.simps bool-to-val.simps

    int-signed-value.simps
  apply (simp add: val-bool-unwrap) apply auto[1]
  unfolding word-sless-def apply auto[1]
  unfolding signed-def apply auto[1] apply auto[2]
  using bit-less-eq-def apply (metis bot-nat-0.extremum take-bit-0)
  by (metis bit-less-eq-def bot-nat-0.extremum take-bit-0)

lemma val-abs-always-pos:
  assumes intval-abs (new-int b v) = (new-int b v')
  shows 0 ≤s v'
  using assms
proof (cases v = 0)
  case True
  then have v' = 0
  using val-abs-zero assms
  by (smt (verit, ccfv-threshold) Suc-diff-1 bit-less-eq-def bot-nat-0.extremum
diff-is-0-eq
    len-gt-0 len-of-numeral-defs(2) order-le-less signed-eq-0-iff take-bit-0
    take-bit-signed-take-bit take-bit-unwrap)
  then show ?thesis by simp
next
  case neg0: False
  then show ?thesis

```

```

proof (cases val-to-bool(val[(new-int b 0) < (new-int b v)]))
  case True
  then show ?thesis using less-eq-def
  using assms val-abs-pos
  by (smt (verit, ccfv-SIG) One-nat-def Suc-leI bit.compl-one bit-less-eq-def
    cancel-comm-monoid-add-class.diff-cancel diff-zero len-gt-0 len-of-numeral-defs(2)

    mask-0 mask-1 one-le-numeral one-neq-zero signed-word-eqI take-bit-dist-subL

    take-bit-minus-one-eq-mask take-bit-not-eq-mask-diff take-bit-signed-take-bit

    zero-le-numeral)
  next
  case False
  then have val-to-bool(val[(new-int b v) < (new-int b 0)])
  using neq0 less-eq-def
  by (metis signed.neqE)
  then show ?thesis using val-abs-neg less-eq-def unfolding new-int.simps
    intval-negate.simps
  by (metis signed.nless-le take-bit-0)
qed

qed

```

```

lemma intval-abs-elim:
  assumes intval-abs x  $\neq$  UndefVal
  shows  $\exists t v . x = \text{IntVal } t v \wedge \text{intval-abs } x = \text{new-int } t$  (if int-signed-value t v <
    0 then  $-v$  else v)
  using assms
  by (meson intval-abs.elims)

```

```

lemma wf-abs-new-int:
  assumes intval-abs (IntVal t v)  $\neq$  UndefVal
  shows intval-abs (IntVal t v) = new-int t v  $\vee$  intval-abs (IntVal t v) = new-int t
    ( $-v$ )
  using assms
  using intval-abs.simps(1) by presburger

```

```

lemma mono-undef-abs:
  assumes intval-abs (intval-abs x)  $\neq$  UndefVal
  shows intval-abs x  $\neq$  UndefVal
  using assms
  by force

```

```

lemma val-abs-idem:
  assumes intval-abs(intval-abs(x))  $\neq$  UndefVal
  shows intval-abs(intval-abs(x)) = intval-abs x

```

```

using assms
proof -
  obtain b v where in-def: intval-abs x = new-int b v
    using assms intval-abs-elim mono-undef-abs by blast
  then show ?thesis
  proof (cases val-to-bool(val[(new-int b v) < (new-int b 0)]))
    case True
    then have nested: (intval-abs (intval-abs x)) = new-int b (-v)
      using val-abs-neg intval-negate.simps in-def
      by simp
    then have x = new-int b (-v)
      using in-def True unfolding new-int.simps
    by (smt (verit, best) intval-abs.simps(1) less-eq-def less-eq-zero less-numeral-extra(1)

        mask-1 mask-eq-take-bit-minus-one neg-one.elims neg-one-signed new-int.simps

        one-le-numeral one-neq-zero signed.neqE signed.not-less take-bit-of-0
    val-abs-always-pos)
    then show ?thesis using val-abs-always-pos
      using True in-def less-eq-def signed.leD
      using signed.nless-le by blast
  next
    case False
    then show ?thesis
      using in-def by force
  qed
qed

lemma val-abs-negate:
  assumes intval-abs (intval-negate x) ≠ UndefVal
  shows intval-abs (intval-negate x) = intval-abs x
  using assms apply (cases x; auto)
  apply (metis less-eq-def new-int.simps signed.dual-order.strict-iff-not signed.less-linear

        take-bit-0)
  by (smt (verit, ccfv-threshold) add.inverse-neutral intval-abs.simps(1) less-eq-def
    less-eq-zero
    less-numeral-extra(1) mask-1 mask-eq-take-bit-minus-one neg-one.elims neg-one-signed

    new-int.simps one-le-numeral one-neq-zero signed.order.order-iff-strict take-bit-of-0

    val-abs-always-pos)

```

Optimisations

```

optimization AbsIdempotence: abs(abs(x)) ⟶ abs(x)
  apply auto[1]
  by (metis UnaryExpr unary-eval.simps(1) val-abs-idem)

optimization AbsNegate: (abs(-x)) ⟶ abs(x)

```

```

    apply auto[1] using val-abs-negate
  by (metis unary-eval.simps(1) unfold-unary)

end

end

```

11.2 AddNode Phase

```

theory AddPhase
  imports
    Common
begin

```

```

phase AddNode
  terminating size
begin

```

```

lemma binadd-commute:
  assumes bin-eval BinAdd x y  $\neq$ .UndefVal
  shows bin-eval BinAdd x y = bin-eval BinAdd y x
  using assms intval-add-sym by simp

```

```

optimization AddShiftConstantRight:  $((\text{const } v) + y) \mapsto y + (\text{const } v)$  when
 $\neg(\text{is-ConstantExpr } y)$ 
  using size-non-const
  apply (metis add-2-eq-Suc' less-Suc-eq plus-1-eq-Suc size.simps(11) size-non-add)
  unfolding le-expr-def
  apply (rule impI)
  subgoal premises 1
    apply (rule allI impI)+

    subgoal premises 2 for m p va
      apply (rule BinaryExprE[OF 2])
    subgoal premises 3 for x ya
      apply (rule BinaryExpr)
      using 3 apply simp
      using 3 apply simp
      using 3 binadd-commute apply auto
    done
  done
done
done

```

```

optimization AddShiftConstantRight2:  $((\text{const } v) + y) \mapsto y + (\text{const } v)$  when
 $\neg(\text{is-ConstantExpr } y)$ 

```

unfolding *le-expr-def*
apply (*auto simp: intval-add-sym*)

using *size-non-const*
by (*metis add-2-eq-Suc' lessI plus-1-eq-Suc size.simps(11) size-non-add*)

lemma *is-neutral-0 [simp]*:
assumes *1: intval-add (IntVal b x) (IntVal b 0) \neq UndefVal*
shows *intval-add (IntVal b x) (IntVal b 0) = (new-int b x)*
using *1 by auto*

optimization *AddNeutral*: $(e + (\text{const } (\text{IntVal } 32\ 0))) \mapsto e$
unfolding *le-expr-def* **apply** *auto*
using *is-neutral-0 eval-unused-bits-zero*
by (*smt (verit) add-cancel-left-right intval-add.elims val-to-bool.simps(1)*)

ML-val $\langle @\{term\ \langle x = y \rangle\} \rangle$

lemma *NeutralLeftSubVal*:
assumes *e1 = new-int b ival*
shows *val[(e1 - e2) + e2] \approx e1*
apply *simp using assms by (cases e1; cases e2; auto)*

optimization *RedundantSubAdd*: $((e_1 - e_2) + e_2) \mapsto e_1$
apply *auto using eval-unused-bits-zero NeutralLeftSubVal*
unfolding *well-formed-equal-defn*
by (*smt (verit) evalDet intval-sub.elims new-int.elims*)

lemma *allE2*: $(\forall x\ y. P\ x\ y) \implies (P\ a\ b \implies R) \implies R$
by *simp*

lemma *just-goal2*:
assumes *1: ($\forall\ a\ b. (\text{intval-add } (\text{intval-sub } a\ b)\ b \neq \text{UndefVal} \wedge a \neq \text{UndefVal}$*
 \longrightarrow
intval-add (intval-sub a b) b = a))
shows *(BinaryExpr BinAdd (BinaryExpr BinSub e1 e2) e2) \geq e1*
unfolding *le-expr-def unfold-binary bin-eval.simps*
by (*metis 1 evalDet evaltree-not-undef*)

optimization *RedundantSubAdd2*: $e_2 + (e_1 - e_2) \mapsto e_1$
apply (*metis add.commute add-less-cancel-right less-add-Suc2 plus-1-eq-Suc size-binary-const*)

size-non-add trans-less-add2)

by (*smt (verit, del-insts) BinaryExpr BinaryExprE RedundantSubAdd(1) bin-add-commute le-expr-def rewrite-preservation.simps(1)*)

lemma *AddToSubHelperLowLevel:*

shows *intval-add (intval-negate e) y = intval-sub y e (is ?x = ?y)*
by (*induction y; induction e; auto*)

print-phases

lemma *val-redundant-add-sub:*

assumes *a = new-int bb ival*
assumes *val[b + a] ≠ UndefVal*
shows *val[(b + a) - b] = a*
using *assms apply (cases a; cases b; auto)*
by *presburger*

lemma *val-add-right-negate-to-sub:*

assumes *val[x + e] ≠ UndefVal*
shows *val[x + (-e)] = val[x - e]*
using *assms by (cases x; cases e; auto)*

lemma *exp-add-left-negate-to-sub:*

exp[-e + y] ≥ exp[y - e]
apply (*cases e; cases y; auto*)
using *AddToSubHelperLowLevel by auto+*

Optimisations

optimization *RedundantAddSub: (b + a) - b ⟶ a*

apply *auto*
by (*smt (verit) evalDet intval-add.elims new-int.elims val-redundant-add-sub eval-unused-bits-zero*)

optimization *AddRightNegateToSub: x + -e ⟶ x - e*

apply (*metis Nat.add-0-right add-2-eq-Suc' add-less-mono1 add-mono-thms-linordered-field(2)*
less-SucI not-less-less-Suc-eq size-binary-const size-non-add size-pos)

```

using AddToSubHelperLowLevel intval-add-sym by auto

optimization AddLeftNegateToSub:  $-e + y \mapsto y - e$ 
  apply (smt (verit, best) One-nat-def add.commute add-Suc-right is-ConstantExpr-def
less-add-Suc2
    numeral-2-eq-2 plus-1-eq-Suc size.simps(1) size.simps(11) size-binary-const
size-non-add)
  using exp-add-left-negate-to-sub by blast

end

```

```

end

```

11.3 AndNode Phase

```

theory AndPhase
  imports
    Common
    Proofs.StampEvalThms
begin

context stamp-mask
begin

lemma AndRightFallthrough:  $((\text{and } (\text{not } (\downarrow x)) (\uparrow y)) = 0) \longrightarrow \text{exp}[x \ \& \ y] \geq$ 
 $\text{exp}[y]$ 
  apply simp apply (rule impI; (rule allI)+)
  apply (rule impI)
  subgoal premises p for m p v
  proof –
    obtain xv where xv:  $[m, p] \vdash x \mapsto xv$ 
      using p(2) by blast
    obtain yv where yv:  $[m, p] \vdash y \mapsto yv$ 
      using p(2) by blast
    have  $v = \text{val}[xv \ \& \ yv]$ 
      by (metis BinaryExprE bin-eval.simps(4) evalDet p(2) xv yv)
    then have  $v = yv$ 
      by (smt (verit) eval-unused-bits-zero intval-and.elims new-int.elims new-int-bin.elims
p(2)
    unfold-binary xv yv p(1) not-down-up-mask-and-zero-implies-zero)
    then show ?thesis using yv by simp
  qed
done

```



```

lemma AndLeftFallthrough: (((and (not ( $\downarrow$  y)) ( $\uparrow$  x)) = 0))  $\longrightarrow$  exp[x & y]  $\geq$ 
exp[x]
  apply simp apply (rule impI; (rule allI)+)
  apply (rule impI)
  subgoal premises p for m p v
  proof –
    obtain xv where xv: [m, p]  $\vdash$  x  $\mapsto$  xv
    using p(2) by blast
    obtain yv where yv: [m, p]  $\vdash$  y  $\mapsto$  yv
    using p(2) by blast
    have v = val[xv & yv]
    by (metis BinaryExprE bin-eval.simps(4) evalDet p(2) xv yv)
    then have v = xv
    by (smt (verit) and.commute eval-unused-bits-zero intval-and.elims new-int.simps

      new-int-bin.simps p(2) unfold-binary xv yv p(1) not-down-up-mask-and-zero-implies-zero)
    then show ?thesis using xv by simp
  qed
done
end

```

```

phase AndNode
  terminating size
begin

```

```

lemma bin-and-nots:
  ( $\sim$  x &  $\sim$  y) = ( $\sim$ (x | y))
  by simp

```

```

lemma bin-and-neutral:
  (x &  $\sim$  False) = x
  by simp

```

```

lemma val-and-equal:
  assumes x = new-int b v
  and val[x & x]  $\neq$  UndefVal
  shows val[x & x] = x
  using assms by (cases x; auto)

```

```

lemma val-and-nots:
  val[ $\sim$  x &  $\sim$  y] = val[ $\sim$ (x | y)]
  apply (cases x; cases y; auto) by (simp add: take-bit-not-take-bit)

```

```

lemma val-and-neutral:
  assumes x = new-int b v
  and val[x &  $\sim$ (new-int b' 0)]  $\neq$  UndefVal
  shows val[x &  $\sim$ (new-int b' 0)] = x

```

```

using assms apply (cases x; auto) apply (simp add: take-bit-eq-mask)
by presburger

```

```

lemma val-and-zero:
  assumes x = new-int b v
  shows val[x & (IntVal b 0)] = IntVal b 0
  using assms by (cases x; auto)

```

```

lemma exp-and-equal:
  exp[x & x] ≥ exp[x]
  apply auto
by (smt (verit) evalDet intval-and.elims new-int.elims val-and-equal eval-unused-bits-zero)

```

```

lemma exp-and-nots:
  exp[~x & ~y] ≥ exp[~(x | y)]
  apply (cases x; cases y; auto) using val-and-nots
by fastforce +

```

```

lemma exp-sign-extend:
  assumes e = (1 << In) - 1
  shows BinaryExpr BinAnd (UnaryExpr (UnarySignExtend In Out) x)
    (ConstantExpr (new-int b e))
    ≥ (UnaryExpr (UnaryZeroExtend In Out) x)

  apply auto
  subgoal premises p for m p va
  proof -
    obtain va where va: [m,p] ⊢ x ↦ va
    using p(2) by auto
    then have va ≠ UndefVal
    by (simp add: evaltree-not-undef)
    then have 1: intval-and (intval-sign-extend In Out va) (IntVal b (take-bit b
e)) ≠ UndefVal
    using evalDet p(1) p(2) va by blast
    then have 2: intval-sign-extend In Out va ≠ UndefVal
    by auto
    then have 21: (0::nat) < b
    using eval-bits-1-64 p(4) by blast
    then have 3: b ⊆ (64::nat)
    using eval-bits-1-64 p(4) by blast
    then have 4: - ((2::int) ^ b div (2::int)) ⊆ sint (signed-take-bit (b - Suc
(0::nat)) (take-bit b e))
    by (simp add: 21 int-power-div-base signed-take-bit-int-greater-eq-minus-exp-word)
    then have 5: sint (signed-take-bit (b - Suc (0::nat)) (take-bit b e)) < (2::int)
^ b div (2::int)

```

```

by (simp add: 21 3 Suc-le-lessD int-power-div-base signed-take-bit-int-less-exp-word)
then have 6:  $[m,p] \vdash \text{UnaryExpr } (\text{UnaryZeroExtend } In \ Out)$ 
       $x \mapsto \text{intval-and } (\text{intval-sign-extend } In \ Out \ va) \ (\text{IntVal } b \ (\text{take-bit } b \ e))$ 
apply (cases va; simp)
apply (simp add:  $\langle (va::\text{Value}) \neq \text{UndefVal} \rangle$ ) defer
subgoal premises p for x3
proof -
  have va = ObjRef x3
  using p(1) by auto
  then have sint (signed-take-bit (b - Suc (0::nat)) (take-bit b e)) <
(2::int) ^ b div (2::int)
  by (simp add: 5)
  then show ?thesis
  using 2 intval-sign-extend.simps(3) p(1) by blast
qed

subgoal premises p for x4
proof -
  have sg1: va = ObjStr x4
  using 2 p(1) by auto
  then have sint (signed-take-bit (b - Suc (0::nat)) (take-bit b e)) <
(2::int) ^ b div (2::int)
  by (simp add: 5)
  then show ?thesis
  using 1 sg1 by auto
qed

subgoal premises p for x21 x22
proof -
  have sgg1: va = IntVal x21 x22
  by (simp add: p(1))
  then have sgg2: sint (signed-take-bit (b - Suc (0::nat)) (take-bit b e))
< (2::int) ^ b div (2::int)
  by (simp add: 5)
  then show ?thesis
  sorry
qed
done
then show ?thesis
by (metis evalDet p(2) va)
qed
done

```

lemma *val-and-commute*[simp]:
 $val[x \& y] = val[y \& x]$
apply (*cases x; cases y; auto*)
by (*simp add: word-bw-comms(1)*)

Optimisations

optimization *AndEqual*: $x \& x \mapsto x$
using *exp-and-equal* **by** *blast*

optimization *AndShiftConstantRight*: $((const\ x) \& y) \mapsto y \& (const\ x)$
 $when\ \neg(is_ConstantExpr\ y)$
using *size-flip-binary* **by** *auto*

optimization *AndNots*: $(\sim x) \& (\sim y) \mapsto \sim(x \mid y)$
apply (*metis add-2-eq-Suc' less-SucI less-add-Suc1 not-less-eq size-binary-const size-non-add*)
using *exp-and-nots* **by** *auto*

optimization *AndSignExtend*: $BinaryExpr\ BinAnd\ (UnaryExpr\ (UnarySignExtend\ In\ Out)\ (x))$

$(const\ (new_int\ b\ e))$
 $\mapsto (UnaryExpr\ (UnaryZeroExtend\ In\ Out)\ (x))$
 $when\ (e = (1 << In) - 1)$

using *exp-sign-extend* **by** *simp*

optimization *AndNeutral*: $(x \& \sim(const\ (IntVal\ b\ 0))) \mapsto x$
 $when\ (wf_stamp\ x \wedge stamp_expr\ x = IntegerStamp\ b\ lo\ hi)$
apply *auto*
by (*smt (verit) Value.sel(1) eval-unused-bits-zero intval-and.elims intval-word.simps*
 $new_int.simps\ new_int_bin.simps\ take-bit-eq-mask$)

optimization *AndRightFallThrough*: $(x \& y) \mapsto y$
 $when\ (((and\ (not\ (IRExpr_down\ x))\ (IRExpr_up\ y)) = 0))$
by (*simp add: IRExpr-down-def IRExpr-up-def*)

optimization *AndLeftFallThrough*: $(x \& y) \mapsto x$
 $when\ (((and\ (not\ (IRExpr_down\ y))\ (IRExpr_up\ x)) = 0))$
by (*simp add: IRExpr-down-def IRExpr-up-def*)

end

end

11.4 BinaryNode Phase

```

theory BinaryNode
  imports
    Common
begin

phase BinaryNode
  terminating size
begin

optimization BinaryFoldConstant: BinaryExpr op (const v1) (const v2)  $\mapsto$  ConstantExpr (bin-eval op v1 v2)
  unfolding le-expr-def
  apply (rule allI impI)+
  subgoal premises bin for m p v
  print-facts
  apply (rule BinaryExprE[OF bin])
  subgoal premises prems for x y
  print-facts

  proof –
    have x: x = v1 using prems by auto
    have y: y = v2 using prems by auto
    have xy: v = bin-eval op x y using prems x y by simp
    have int:  $\exists b\ vv. v = \text{new-int } b\ vv$  using bin-eval-new-int prems by fast
    show ?thesis
      unfolding prems x y xy
      apply (rule ConstantExpr)
      using prems x y xy int sorry
    qed
  done
done

print-facts

end

end

```

11.5 ConditionalNode Phase

```

theory ConditionalPhase
  imports
    Common
    Proofs.StampEvalThms
begin

phase ConditionalNode

```

terminating *size*
begin

lemma *negates*: $\exists v\ b.\ e = \text{IntVal } b\ v \wedge b > 0 \implies \text{val-to-bool } (\text{val}[e]) \longleftrightarrow \neg(\text{val-to-bool } (\text{val}[!e]))$
unfolding *intval-logic-negation.simps*
by (*metis* (*mono-tags*, *lifting*) *intval-logic-negation.simps*(1) *logic-negate-def new-int.simps* *of-bool-eq*(2) *one-neq-zero take-bit-of-0 take-bit-of-1 val-to-bool.simps*(1))

lemma *negation-condition-intval*:
assumes $e = \text{IntVal } b\ ie$
assumes $0 < b$
shows $\text{val}[(!e)\ ?\ x : y] = \text{val}[e\ ?\ y : x]$
using *assms* **by** (*cases* *e*; *auto simp: negates logic-negate-def*)

lemma *negation-preserve-eval*:
assumes $[m, p] \vdash \text{exp}[!e] \mapsto v$
shows $\exists v'. ([m, p] \vdash \text{exp}[e] \mapsto v') \wedge v = \text{val}[!v']$
using *assms* **by** *auto*

lemma *negation-preserve-eval-intval*:
assumes $[m, p] \vdash \text{exp}[!e] \mapsto v$
shows $\exists v'\ b\ vv. ([m, p] \vdash \text{exp}[e] \mapsto v') \wedge v' = \text{IntVal } b\ vv \wedge b > 0$
using *assms*
by (*metis eval-bits-1-64 intval-logic-negation.elims negation-preserve-eval unfold-unary*)

optimization *NegateConditionFlipBranches*: $((!e)\ ?\ x : y) \mapsto (e\ ?\ y : x)$
apply *simp using negation-condition-intval negation-preserve-eval-intval*
by (*smt* (*verit*, *best*) *ConditionalExpr ConditionalExprE Value.distinct*(1) *evalDet negates negation-preserve-eval*)

optimization *DefaultTrueBranch*: $(\text{true}\ ?\ x : y) \mapsto x$.

optimization *DefaultFalseBranch*: $(\text{false}\ ?\ x : y) \mapsto y$.

optimization *ConditionalEqualBranches*: $(e\ ?\ x : x) \mapsto x$.

optimization *condition-bounds-x*: $((u < v)\ ?\ x : y) \mapsto x$
when (*stamp-under* (*stamp-expr* *u*) (*stamp-expr* *v*) \wedge *wf-stamp* *u* \wedge *wf-stamp* *v*)
using *stamp-under-defn* **by** *fastforce*

optimization *condition-bounds-y*: $((u < v)\ ?\ x : y) \mapsto y$
when (*stamp-under* (*stamp-expr* *v*) (*stamp-expr* *u*) \wedge *wf-stamp* *u* \wedge *wf-stamp* *v*)
using *stamp-under-defn-inverse* **by** *fastforce*

```

lemma val-optimise-integer-test:
  assumes  $\exists v. x = \text{IntVal } 32 \ v$ 
  shows  $\text{val}[(x \ \& \ (\text{IntVal } 32 \ 1)) \ \text{eq} \ (\text{IntVal } 32 \ 0)] \ ? \ (\text{IntVal } 32 \ 0) : (\text{IntVal } 32 \ 1)]$ 
  =
     $\text{val}[x \ \& \ \text{IntVal } 32 \ 1]$ 
  using assms apply auto
  apply (metis (full-types) bool-to-val.simps(2) val-to-bool.simps(1))
  by (metis (mono-tags, lifting) and-one-eq bool-to-val.simps(1) even-iff-mod-2-eq-zero
odd-iff-mod-2-eq-one val-to-bool.simps(1))

optimization ConditionalEliminateKnownLess:  $((x < y) \ ? \ x : y) \mapsto x$ 
   $\text{when } (\text{stamp-under } (\text{stamp-expr } x) \ (\text{stamp-expr } y))$ 
   $\wedge \text{wf-stamp } x \wedge \text{wf-stamp } y$ 
  using stamp-under-defn by fastforce

optimization ConditionalEqualIsRHS:  $((x \ \text{eq} \ y) \ ? \ x : y) \mapsto y$ 
  apply auto
  by (smt (verit) Value.inject(1) bool-to-val.simps(2) bool-to-val-bin.simps evalDet

    intval-equals.elims val-to-bool.elims(1))

optimization normalizeX:  $((x \ \text{eq} \ \text{const } (\text{IntVal } 32 \ 0)) \ ?$ 
   $(\text{const } (\text{IntVal } 32 \ 0)) : (\text{const } (\text{IntVal } 32 \ 1))) \mapsto x$ 
   $\text{when } (\text{IRExpr-up } x = 1) \wedge \text{stamp-expr } x = \text{IntegerStamp}$ 
  b 0 1
  apply auto
  subgoal premises p for m p v xa
  proof –
    obtain xa where xa:  $[m, p] \vdash x \mapsto xa$ 
    using p by blast
    have 3:  $[m, p] \vdash \text{if } \text{val-to-bool } (\text{intval-equals } xa \ (\text{IntVal } (32::\text{nat}) \ (0::64 \ \text{word})))$ 
       $\text{then } \text{ConstantExpr } (\text{IntVal } (32::\text{nat}) \ (0::64 \ \text{word}))$ 
       $\text{else } \text{ConstantExpr } (\text{IntVal } (32::\text{nat}) \ (1::64 \ \text{word})) \mapsto v$ 
    using evalDet p(3) p(5) xa
    using p(4) p(6) by blast
    then have 4:  $xa = \text{IntVal } 32 \ 0 \mid xa = \text{IntVal } 32 \ 1$ 
    sorry
    then have 6:  $v = xa$ 
    sorry
    then show ?thesis
    using xa by auto
  qed
done

optimization normalizeX2:  $((x \ \text{eq} \ (\text{const } (\text{IntVal } 32 \ 1))) \ ?$ 
   $(\text{const } (\text{IntVal } 32 \ 1)) : (\text{const } (\text{IntVal } 32 \ 0))) \mapsto x$ 

```

when ($x = \text{ConstantExpr } (\text{IntVal } 32\ 0) \mid (x = \text{ConstantExpr } (\text{IntVal } 32\ 1)))$).

optimization *flipX*: $((x \text{ eq } (\text{const } (\text{IntVal } 32\ 0))) \text{ ? } (\text{const } (\text{IntVal } 32\ 1)) : (\text{const } (\text{IntVal } 32\ 0))) \mapsto x \oplus (\text{const } (\text{IntVal } 32\ 1))$
when ($x = \text{ConstantExpr } (\text{IntVal } 32\ 0) \mid (x = \text{ConstantExpr } (\text{IntVal } 32\ 1)))$).

optimization *flipX2*: $((x \text{ eq } (\text{const } (\text{IntVal } 32\ 1))) \text{ ? } (\text{const } (\text{IntVal } 32\ 0)) : (\text{const } (\text{IntVal } 32\ 1))) \mapsto x \oplus (\text{const } (\text{IntVal } 32\ 1))$
when ($x = \text{ConstantExpr } (\text{IntVal } 32\ 0) \mid (x = \text{ConstantExpr } (\text{IntVal } 32\ 1)))$).

lemma *stamp-of-default*:
assumes *stamp-expr* $x = \text{default-stamp}$
assumes *wf-stamp* x
shows $([m, p] \vdash x \mapsto v) \longrightarrow (\exists vv. v = \text{IntVal } 32\ vv)$
using *assms*
by (*metis default-stamp valid-value-elim*(3) *wf-stamp-def*)

optimization *OptimiseIntegerTest*:
 $((x \ \& \ (\text{const } (\text{IntVal } 32\ 1))) \text{ eq } (\text{const } (\text{IntVal } 32\ 0))) \text{ ? } (\text{const } (\text{IntVal } 32\ 0)) : (\text{const } (\text{IntVal } 32\ 1))) \mapsto x \ \& \ (\text{const } (\text{IntVal } 32\ 1))$
when (*stamp-expr* $x = \text{default-stamp} \wedge \text{wf-stamp } x$)
apply *simp* **apply** (*rule impI*; (*rule allI*) $+$; *rule impI*)
subgoal **premises** *eval* **for** $m\ p\ v$
proof –
obtain xv **where** $[m, p] \vdash x \mapsto xv$
using *eval* **by** *fast*
then **have** $x32$: $\exists v. xv = \text{IntVal } 32\ v$
using *stamp-of-default eval* **by** *auto*
obtain lhs **where** $[m, p] \vdash \text{exp}[(((x \ \& \ (\text{const } (\text{IntVal } 32\ 1))) \text{ eq } (\text{const } (\text{IntVal } 32\ 0)))) \text{ ? } (\text{const } (\text{IntVal } 32\ 0)) : (\text{const } (\text{IntVal } 32\ 1)))] \mapsto lhs$
using *eval*(2) **by** *auto*
then **have** $lhsV$: $lhs = \text{val}[((xv \ \& \ (\text{IntVal } 32\ 1)) \text{ eq } (\text{IntVal } 32\ 0)) \text{ ? } (\text{IntVal } 32\ 0) : (\text{IntVal } 32\ 1)]$
using $xv \text{ evaltree.BinaryExpr evaltree.ConstantExpr evaltree.ConditionalExpr}$
by (*smt* (*verit*) *ConditionalExprE ConstantExprE bin-eval.simps*(11) *bin-eval.simps*(4) *evalDet intval-conditional.simps* *unfold-binary*)
obtain rhs **where** $[m, p] \vdash \text{exp}[x \ \& \ (\text{const } (\text{IntVal } 32\ 1))] \mapsto rhs$
using *eval*(2) **by** *blast*
then **have** $rhsV$: $rhs = \text{val}[xv \ \& \ \text{IntVal } 32\ 1]$
by (*metis BinaryExprE ConstantExprE bin-eval.simps*(4) *evalDet xv*)


```

have lhs = rhs using val-optimize-integer-test x32
using lhsV rhsV by presburger
then show ?thesis
by (metis eval(2) evalDet lhs rhs)
qed
done

```

```

optimization opt-optimize-integer-test-2:
  (((x & (const (IntVal 32 1))) eq (const (IntVal 32 0))) ?
    (const (IntVal 32 0)) : (const (IntVal 32 1)))  $\mapsto$ 
     $\begin{matrix} x \\ \text{when } (x = \text{ConstantExpr (IntVal 32 0)} \mid (x = \text{ConstantExpr (IntVal} \\ 32 \text{ 1)})) \end{matrix}$  .

```

end

end

11.6 MulNode Phase

```

theory MulPhase
imports
  Common
  Proofs.StampEvalThms
begin

```

```

fun mul-size :: IRExpr  $\Rightarrow$  nat where
  mul-size (UnaryExpr op e) = (mul-size e) + 2 |
  mul-size (BinaryExpr BinMul x y) = ((mul-size x) + (mul-size y) + 2) * 2 |
  mul-size (BinaryExpr op x y) = (mul-size x) + (mul-size y) + 2 |
  mul-size (ConditionalExpr cond t f) = (mul-size cond) + (mul-size t) + (mul-size
f) + 2 |
  mul-size (ConstantExpr c) = 1 |
  mul-size (ParameterExpr ind s) = 2 |
  mul-size (LeafExpr nid s) = 2 |
  mul-size (ConstantVar c) = 2 |
  mul-size (VariableExpr x s) = 2

```

```

phase MulNode
terminating mul-size
begin

```

```

lemma bin-eliminate-redundant-negative:
  uminus (x :: 'a::len word) * uminus (y :: 'a::len word) = x * y
by simp

lemma bin-multiply-identity:
  (x :: 'a::len word) * 1 = x
by simp

lemma bin-multiply-eliminate:
  (x :: 'a::len word) * 0 = 0
by simp

lemma bin-multiply-negative:
  (x :: 'a::len word) * uminus 1 = uminus x
by simp

lemma bin-multiply-power-2:
  (x :: 'a::len word) * (2j) = x << j
by simp

lemma take-bit64[simp]:
  fixes w :: int64
  shows take-bit 64 w = w
proof –
  have Nat.size w = 64
  by (simp add: size64)
  then show ?thesis
  by (metis lt2p-lem mask-eq-iff take-bit-eq-mask verit-comp-simplify1(2) wsst-TYs(3))
qed

lemma mergeTakeBit:
  fixes a :: nat
  fixes b c :: 64 word
  shows take-bit a (take-bit a (b) * take-bit a (c)) =
    take-bit a (b * c)
by (smt (verit, ccfv-SIG) take-bit-mult take-bit-of-int unsigned-take-bit-eq word-mult-def)

lemma val-eliminate-redundant-negative:
  assumes val[-x * -y] ≠ UndefVal
  shows val[-x * -y] = val[x * y]
  using assms apply (cases x; cases y; auto)
  using mergeTakeBit by auto

```

```

lemma val-multiply-neutral:
  assumes  $x = \text{new-int } b \ v$ 
  shows  $\text{val}[x * (\text{IntVal } b \ 1)] = \text{val}[x]$ 
  using assms by force

lemma val-multiply-zero:
  assumes  $x = \text{new-int } b \ v$ 
  shows  $\text{val}[x * (\text{IntVal } b \ 0)] = \text{IntVal } b \ 0$ 
  using assms by simp

lemma val-multiply-negative:
  assumes  $x = \text{new-int } b \ v$ 
  shows  $\text{val}[x * \text{intval-negate } (\text{IntVal } b \ 1)] = \text{intval-negate } x$ 
  by (smt (verit) Value.disc(1) Value.inject(1) add.inverse-neutral intval-negate.simps(1)

    is-IntVal-def mask-0 mask-eq-take-bit-minus-one new-int.elims of-bool-eq(2)
  take-bit-dist-neg
    take-bit-of-1 val-eliminate-redundant-negative val-multiply-neutral val-multiply-zero

    verit-minus-simplify(4) zero-neq-one assms)

lemma val-MulPower2:
  fixes  $i :: 64 \ \text{word}$ 
  assumes  $y = \text{IntVal } 64 \ (2 \wedge \text{unat}(i))$ 
  and  $0 < i$ 
  and  $i < 64$ 
  and  $\text{val}[x * y] \neq \text{UndefVal}$ 
  shows  $\text{val}[x * y] = \text{val}[x << \text{IntVal } 64 \ i]$ 
  using assms apply (cases  $x$ ; cases  $y$ ; auto)
  subgoal premises  $p$  for  $x2$ 
  proof –
    have  $63 :: \text{int}64 = \text{mask } 6$ 
    by eval
    then have  $(2 :: \text{int}) \wedge 6 = 64$ 
    by eval
    then have  $\text{uint } i < (2 :: \text{int}) \wedge 6$ 
    by (metis linorder-not-less lt2p-lem of-int-numeral  $p(4)$  size64 word-2p-lem
      word-of-int-2p wsst-TYs(3))
    then have  $\text{and } i \ (\text{mask } 6) = i$ 
    using mask-eq-iff by blast
    then show  $x2 << \text{unat } i = x2 << \text{unat } (\text{and } i \ (63 :: 64 \ \text{word}))$ 
    unfolding  $63$ 
    by force
  qed
by presburger

```

```

lemma val-MulPower2Add1:
  fixes  $i :: 64 \text{ word}$ 
  assumes  $y = \text{IntVal } 64 ((2 \wedge \text{unat}(i)) + 1)$ 
  and  $0 < i$ 
  and  $i < 64$ 
  and  $\text{val-to-bool}(\text{val}[\text{IntVal } 64 \ 0 < x])$ 
  and  $\text{val-to-bool}(\text{val}[\text{IntVal } 64 \ 0 < y])$ 
  shows  $\text{val}[x * y] = \text{val}[(x << \text{IntVal } 64 \ i) + x]$ 
  using assms apply (cases x; cases y; auto)
  subgoal premises  $p$  for  $x2$ 
  proof –
    have  $63 :: \text{int64} = \text{mask } 6$ 
    by eval
    then have  $(2 :: \text{int}) \wedge 6 = 64$ 
    by eval
    then have  $\text{and } i (\text{mask } 6) = i$ 
    using mask-eq-iff by (simp add: less-mask-eq p(6))
    then have  $x2 * ((2 :: 64 \text{ word}) \wedge \text{unat } i + (1 :: 64 \text{ word})) = (x2 * ((2 :: 64 \text{ word})$ 
     $\wedge \text{unat } i)) + x2$ 
    by (simp add: distrib-left)
    then show  $x2 * ((2 :: 64 \text{ word}) \wedge \text{unat } i + (1 :: 64 \text{ word})) = x2 << \text{unat } (\text{and } i$ 
     $(63 :: 64 \text{ word})) + x2$ 
    by (simp add: 63 and (i :: 64 word) (mask (6 :: nat)) = i)
  qed
  using val-to-bool.simps(2) by presburger

```

```

lemma val-MulPower2Sub1:
  fixes  $i :: 64 \text{ word}$ 
  assumes  $y = \text{IntVal } 64 ((2 \wedge \text{unat}(i)) - 1)$ 
  and  $0 < i$ 
  and  $i < 64$ 
  and  $\text{val-to-bool}(\text{val}[\text{IntVal } 64 \ 0 < x])$ 
  and  $\text{val-to-bool}(\text{val}[\text{IntVal } 64 \ 0 < y])$ 
  shows  $\text{val}[x * y] = \text{val}[(x << \text{IntVal } 64 \ i) - x]$ 
  using assms apply (cases x; cases y; auto)
  subgoal premises  $p$  for  $x2$ 
  proof –
    have  $63 :: \text{int64} = \text{mask } 6$ 
    by eval
    then have  $(2 :: \text{int}) \wedge 6 = 64$ 
    by eval
    then have  $\text{and } i (\text{mask } 6) = i$ 
    using mask-eq-iff by (simp add: less-mask-eq p(6))
    then have  $x2 * ((2 :: 64 \text{ word}) \wedge \text{unat } i - (1 :: 64 \text{ word})) = (x2 * ((2 :: 64 \text{ word})$ 
     $\wedge \text{unat } i)) - x2$ 
    by (simp add: right-diff-distrib')
    then show  $x2 * ((2 :: 64 \text{ word}) \wedge \text{unat } i - (1 :: 64 \text{ word})) = x2 << \text{unat } (\text{and } i$ 

```

```

(63::64 word)) - x2
  by (simp add: 63 and (i::64 word) (mask (6::nat)) = i)
qed
using val-to-bool.simps(2) by presburger

```

lemma *val-distribute-multiplication*:

```

assumes x = new-int 64 xx ∧ q = new-int 64 qq ∧ a = new-int 64 aa
shows val[x * (q + a)] = val[(x * q) + (x * a)]
apply (cases x; cases q; cases a; auto) using distrib-left assms by auto

```

lemma *val-MulPower2AddPower2*:

```

fixes i j :: 64 word
assumes y = IntVal 64 ((2 ^ unat(i)) + (2 ^ unat(j)))
and 0 < i
and 0 < j
and i < 64
and j < 64
and x = new-int 64 xx
shows val[x * y] = val[(x << IntVal 64 i) + (x << IntVal 64 j)]
using assms
proof -
  have 63: (63 :: int64) = mask 6
  by eval
  then have (2::int) ^ 6 = 64
  by eval
  then have n: IntVal 64 ((2 ^ unat(i)) + (2 ^ unat(j))) =
    val[(IntVal 64 (2 ^ unat(i))) + (IntVal 64 (2 ^ unat(j)))]

  using assms by (cases i; cases j; auto)
  then have 1: val[x * ((IntVal 64 (2 ^ unat(i))) + (IntVal 64 (2 ^ unat(j))))] =
    val[(x * IntVal 64 (2 ^ unat(i))) + (x * IntVal 64 (2 ^ unat(j)))]

  using assms val-distribute-multiplication val-MulPower2 by simp
  then have 2: val[(x * IntVal 64 (2 ^ unat(i)))] = val[x << IntVal 64 i]
  by (smt (verit) Value.distinct(1) intval-mul.simps(1) new-int.simps new-int-bin.simps
  assms
    val-MulPower2)
  then show ?thesis
  by (smt (verit, del-insts) 1 Value.distinct(1) assms(1) assms(3) assms(5)
  assms(6)
    intval-mul.simps(1) n new-int.simps new-int-bin.elims val-MulPower2)
qed

```

thm-oracles *val-MulPower2AddPower2*

lemma *exp-multiply-zero-64*:

```

exp[x * (const (IntVal 64 0))] ≥ ConstantExpr (IntVal 64 0)
  using val-multiply-zero apply auto
  by (smt (verit) Value.inject(1) constantAsStamp.simps(1) int-signed-value-bounds
      intval-mul.elims
      mult-zero-right new-int.simps new-int-bin.simps nle-le numeral-eq-Suc
      take-bit-of-0
      unfold-const valid-stamp.simps(1) valid-value.simps(1) zero-less-Suc wf-value-def)

```

```

lemma exp-multiply-neutral:
  exp[x * (const (IntVal b 1))] ≥ x
  using val-multiply-neutral apply auto
  by (smt (verit) Value.inject(1) eval-unused-bits-zero intval-mul.elims mult.right-neutral
      new-int.elims new-int-bin.elims)

```

thm-oracles exp-multiply-neutral

```

lemma exp-MulPower2:
  fixes i :: 64 word
  assumes y = ConstantExpr (IntVal 64 (2 ^ unat(i)))
  and 0 < i
  and i < 64
  and exp[x > (const IntVal b 0)]
  and exp[y > (const IntVal b 0)]
  shows exp[x * y] ≥ exp[x << ConstantExpr (IntVal 64 i)]
  using assms apply simp
  by (metis ConstantExprE equiv-exprs-def unfold-binary)

```

```

lemma exp-MulPower2Add1:
  fixes i :: 64 word
  assumes y = ConstantExpr (IntVal 64 ((2 ^ unat(i)) + 1))
  and 0 < i
  and i < 64
  and exp[x > (const IntVal b 0)]
  and exp[y > (const IntVal b 0)]
  shows exp[x * y] ≥ exp[(x << ConstantExpr (IntVal 64 i)) + x]
  using assms apply simp
  by (metis (no-types, lifting) ConstantExprE equiv-exprs-def unfold-binary)

```

```

lemma exp-MulPower2Sub1:
  fixes i :: 64 word
  assumes y = ConstantExpr (IntVal 64 ((2 ^ unat(i)) - 1))
  and 0 < i
  and i < 64
  and exp[x > (const IntVal b 0)]
  and exp[y > (const IntVal b 0)]
  shows exp[x * y] ≥ exp[(x << ConstantExpr (IntVal 64 i)) - x]
  using assms apply simp
  by (metis (no-types, lifting) ConstantExprE equiv-exprs-def unfold-binary)

```

```

lemma exp-MulPower2AddPower2:
  fixes i j :: 64 word
  assumes y = ConstantExpr (IntVal 64 ((2 ^ unat(i)) + (2 ^ unat(j))))
  and 0 < i
  and 0 < j
  and i < 64
  and j < 64
  and exp[x > (const IntVal b 0)]
  and exp[y > (const IntVal b 0)]
shows exp[x * y] ≥ exp[(x << ConstantExpr (IntVal 64 i)) + (x << ConstantExpr
(IntVal 64 j))]
  using assms apply simp
  by (metis (no-types, lifting) ConstantExprE equiv-exprs-def unfold-binary)

```

```

lemma greaterConstant:
  fixes a b :: 64 word
  assumes a > b
  and y = ConstantExpr (IntVal 64 a)
  and x = ConstantExpr (IntVal 64 b)
shows exp[y > x]
apply auto
sorry

```

```

lemma exp-distribute-multiplication:
  shows exp[(x * q) + (x * a)] ≥ exp[x * (q + a)]
sorry

```

Optimisations

```

optimization EliminateRedundantNegative:  $-x * -y \mapsto x * y$ 
  using mul-size.simps apply auto
  by (metis BinaryExpr val-eliminate-redundant-negative bin-eval.simps(2))

```

```

optimization MulNeutral:  $x * \text{ConstantExpr (IntVal b 1)} \mapsto x$ 
  using exp-multiply-neutral by blast

```

```

optimization MulEliminator:  $x * \text{ConstantExpr (IntVal b 0)} \mapsto \text{const (IntVal b 0)}$ 
apply auto
by (smt (verit) Value.inject(1) constantAsStamp.simps(1) int-signed-value-bounds
intval-mul.elims
  mult-zero-right new-int.simps new-int-bin.simps take-bit-of-0 unfold-const
  valid-stamp.simps(1) valid-value.simps(1) val-multiply-zero)

```

```

optimization MulNegate:  $x * -(\text{const (IntVal b 1)}) \mapsto -x$ 
apply auto

```

```

by (smt (verit) Value.distinct(1) Value.sel(1) add.inverse-inverse intval-mul.elims

    intval-negate.simps(1) mask-eq-take-bit-minus-one new-int.simps new-int-bin.simps

    take-bit-dist-neg unary-eval.simps(2) unfold-unary val-multiply-negative
    val-eliminate-redundant-negative val-multiply-negative wf-value-def)

fun isNonZero :: Stamp  $\Rightarrow$  bool where
  isNonZero (IntegerStamp b lo hi) = (lo > 0) |
  isNonZero - = False

lemma isNonZero-defn:
  assumes isNonZero (stamp-expr x)
  assumes wf-stamp x
  shows ([m, p]  $\vdash$  x  $\mapsto$  v)  $\longrightarrow$  ( $\exists$  vv b. (v = IntVal b vv  $\wedge$  val-to-bool val[(IntVal b
0) < v]))
  apply (rule impI) subgoal premises eval
proof -
  obtain b lo hi where xstamp: stamp-expr x = IntegerStamp b lo hi
  by (meson isNonZero.elims(2) assms)
  then obtain vv where vdef: v = IntVal b vv
  by (metis assms(2) eval valid-int wf-stamp-def)
  have lo > 0
  using assms(1) xstamp by force
  then have signed-above: int-signed-value b vv > 0
  using assms unfolding wf-stamp-def
  using eval vdef xstamp by fastforce
  have take-bit b vv = vv
  using eval eval-unused-bits-zero vdef by auto
  then have vv > 0
  by (metis bit-take-bit-iff int-signed-value.simps not-less-zero signed-eq-0-iff
    signed-take-bit-eq-if-positive take-bit-0 take-bit-of-0 verit-comp-simplify1(1)
word-gt-0
    signed-above)
  then show ?thesis
  using vdef signed-above
  by simp
qed
done

optimization MulPower2: x * y  $\mapsto$  x << const (IntVal 64 i)
  when (i > 0  $\wedge$ 
    64 > i  $\wedge$ 
    y = exp[const (IntVal 64 (2 ^ unat(i)))]])

  defer
  apply simp apply (rule impI; (rule allI)+; rule impI)
  subgoal premises eval for m p v
proof -
  obtain xv where xv: [m, p]  $\vdash$  x  $\mapsto$  xv

```



```

    using eval(2) by blast
  then obtain xv where xv: xv = IntVal 64 xv
    by (smt (verit) ConstantExprE bin-eval.simps(2) evalDet intval-bits.simps int-
    val-mul.elims
        new-int-bin.simps unfold-binary eval)
  obtain yv where yv: [m, p] ⊢ y ↦ yv
    using eval(1) eval(2) by blast
  then have lhs: [m, p] ⊢ exp[x * y] ↦ val[xv * yv]
    by (metis bin-eval.simps(2) eval(1) eval(2) evalDet unfold-binary xv)
  have [m, p] ⊢ exp[const (IntVal 64 i)] ↦ val[(IntVal 64 i)]
    by (smt (verit, ccfv-SIG) ConstantExpr constantAsStamp.simps(1) eval-bits-1-64
    take-bit64
        validStampIntConst wf-value-def valid-value.simps(1) xv xv)
  then have rhs: [m, p] ⊢ exp[x << const (IntVal 64 i)] ↦ val[xv << (IntVal 64
  i)]
    using xv xv using evaltree.BinaryExpr
  by (metis Value.simps(5) bin-eval.simps(8) intval-left-shift.simps(1) new-int.simps)
  have val[xv * yv] = val[xv << (IntVal 64 i)]
    by (metis ConstantExprE eval(1) evaltree-not-undef lhs yv val-MulPower2)
  then show ?thesis
    by (metis eval(1) eval(2) evalDet lhs rhs)
qed
done

```

```

optimization MulPower2Add1:  $x * y \mapsto (x \ll \text{const } (\text{IntVal } 64 \ i)) + x$ 
    when  $(i > 0 \wedge$ 
     $64 > i \wedge$ 
     $y = \text{ConstantExpr } (\text{IntVal } 64 \ ((2 \wedge \text{unat}(i)) + 1))$  )

  defer
  apply simp apply (rule impI; (rule allI)+; rule impI)
  subgoal premises p for m p v
  proof -
    obtain xv where xv: [m, p] ⊢ x ↦ xv
      using p by fast
    then obtain xv where xv: xv = IntVal 64 xv
      by (smt (verit) p ConstantExprE bin-eval.simps(2) evalDet intval-bits.simps
    intval-mul.elims
        new-int-bin.simps unfold-binary)
    obtain yv where yv: [m, p] ⊢ y ↦ yv
      using p by blast
    have ygezero: y > ConstantExpr (IntVal 64 0)
      using greaterConstant p wf-value-def by fastforce
    then have 1: 0 < i ∧
      i < 64 ∧
      y = ConstantExpr (IntVal 64 ((2 ∧ unat(i)) + 1))
      using p by blast
    then have lhs: [m, p] ⊢ exp[x * y] ↦ val[xv * yv]
      by (metis bin-eval.simps(2) evalDet p(1) p(2) xv yv unfold-binary)

```

```

then have  $[m, p] \vdash \text{exp}[\text{const } (\text{IntVal } 64 \ i)] \mapsto \text{val}[(\text{IntVal } 64 \ i)]$ 
by (metis wf-value-def verit-comp-simplify1(2) zero-less-numeral ConstantExpr

      constantAsStamp.simps(1) take-bit64 validStampIntConst valid-value.simps(1)
then have  $\text{rhs2}: [m, p] \vdash \text{exp}[x \ll \text{const } (\text{IntVal } 64 \ i)] \mapsto \text{val}[xv \ll (\text{IntVal } 64 \ i)]$ 
by (metis Value.simps(5) bin-eval.simps(8) intval-left-shift.simps(1) new-int.simps xv xv
      evaltree.BinaryExpr)
then have  $\text{rhs}: [m, p] \vdash \text{exp}[(x \ll \text{const } (\text{IntVal } 64 \ i)) + x] \mapsto \text{val}[(xv \ll (\text{IntVal } 64 \ i)) + xv]$ 
by (metis (no-types, lifting) intval-add.simps(1) rhs2 bin-eval.simps(1) Value.simps(5)
      evaltree.BinaryExpr intval-left-shift.simps(1) new-int.simps xv xv)
then have simple:  $\text{val}[xv * (\text{IntVal } 64 \ (2^{\text{unat}(i)}))] = \text{val}[xv \ll (\text{IntVal } 64 \ i)]$ 
using val-MulPower2 sorry
then have  $\text{val}[xv * yv] = \text{val}[(xv \ll (\text{IntVal } 64 \ i)) + xv]$ 
sorry
then show ?thesis
by (metis 1 evalDet lhs p(2) rhs)
qed
done

```

```

optimization MulPower2Sub1:  $x * y \mapsto (x \ll \text{const } (\text{IntVal } 64 \ i)) - x$ 
  when (i > 0 ∧
    64 > i ∧
    y = ConstantExpr (IntVal 64 ((2^unat(i)) - 1)) )

defer
apply simp apply (rule impI; (rule allI)+; rule impI)
subgoal premises p for m p v
proof –
  obtain xv where  $xv: [m, p] \vdash x \mapsto xv$ 
  using p by fast
  then obtain xv where  $xv: xv = \text{IntVal } 64 \ xv$ 
  by (smt (verit) p ConstantExprE bin-eval.simps(2) evalDet intval-bits.simps intval-mul.elims
    new-int-bin.simps unfold-binary)
  obtain yv where  $yv: [m, p] \vdash y \mapsto yv$ 
  using p by blast
  have ygezero:  $y > \text{ConstantExpr } (\text{IntVal } 64 \ 0)$ 
  by (smt (verit, del-Insts) eq-iff-diff-eq-0 mask-0 mask-eq-exp-minus-1 power-inject-exp
    uint-2p unat-eq-zero word-gt-0 zero-neq-one greaterConstant p)
  then have 1:  $0 < i \wedge$ 
     $i < 64 \wedge$ 
     $y = \text{ConstantExpr } (\text{IntVal } 64 \ ((2^{\text{unat}(i)} - 1)))$ 
  using p by blast

```

```

then have lhs: [m, p] ⊢ exp[x * y] ↦ val[xv * yv]
  by (metis bin-eval.simps(2) evalDet p(1) p(2) xv yv unfold-binary)
then have [m, p] ⊢ exp[const (IntVal 64 i)] ↦ val[(IntVal 64 i)]
  by (metis wf-value-def verit-comp-simplify1(2) zero-less-numeral ConstantExpr

      constantAsStamp.simps(1) take-bit64 validStampIntConst valid-value.simps(1))
then have rhs2: [m, p] ⊢ exp[x << const (IntVal 64 i)] ↦ val[xv << (IntVal
64 i)]
  by (metis Value.simps(5) bin-eval.simps(8) intval-left-shift.simps(1) new-int.simps
xv xvv
      evaltree.BinaryExpr)
then have rhs: [m, p] ⊢ exp[(x << const (IntVal 64 i)) - x] ↦ val[(xv <<
(IntVal 64 i)) - xv]
  by (smt (verit, ccfv-threshold) bin-eval.simps(3) new-int-bin.simps int-
val-sub.simps(1)
      rhs2 bin-eval.simps(1) Value.simps(5) evaltree.BinaryExpr intval-left-shift.simps(1)
      new-int.simps xv xvv )
then have val[xv * yv] = val[(xv << (IntVal 64 i)) - xv]
  using 1 exp-MulPower2Sub1 ygezzero sorry
then show ?thesis
  by (metis evalDet lhs p(1) p(2) rhs)
qed
done

end

end

```

11.7 Experimental AndNode Phase

```

theory NewAnd
imports
  Common
  Graph.JavaLong
begin

lemma bin-distribute-and-over-or:
  bin[z & (x | y)] = bin[(z & x) | (z & y)]
  by (smt (verit, best) bit-and-iff bit-eqI bit-or-iff)

lemma intval-distribute-and-over-or:
  val[z & (x | y)] = val[(z & x) | (z & y)]
  apply (cases x; cases y; cases z; auto)
  using bin-distribute-and-over-or by blast+

lemma exp-distribute-and-over-or:
  exp[z & (x | y)] ≥ exp[(z & x) | (z & y)]

```

```

apply simp using intval-distribute-and-over-or
using BinaryExpr bin-eval.simps(4,5)
using intval-or.simps(1) unfolding new-int-bin.simps new-int.simps apply auto
by (metis bin-eval.simps(4) bin-eval.simps(5) intval-or.simps(2) intval-or.simps(5))

```

```

lemma intval-and-commute:
   $val[x \& y] = val[y \& x]$ 
by (cases x; cases y; auto simp: and.commute)

```

```

lemma intval-or-commute:
   $val[x | y] = val[y | x]$ 
by (cases x; cases y; auto simp: or.commute)

```

```

lemma intval-xor-commute:
   $val[x \oplus y] = val[y \oplus x]$ 
by (cases x; cases y; auto simp: xor.commute)

```

```

lemma exp-and-commute:
   $exp[x \& z] \geq exp[z \& x]$ 
apply simp using intval-and-commute by auto

```

```

lemma exp-or-commute:
   $exp[x | y] \geq exp[y | x]$ 
apply simp using intval-or-commute by auto

```

```

lemma exp-xor-commute:
   $exp[x \oplus y] \geq exp[y \oplus x]$ 
apply simp using intval-xor-commute by auto

```

```

lemma bin-eliminate-y:
  assumes  $bin[y \& z] = 0$ 
  shows  $bin[(x | y) \& z] = bin[x \& z]$ 
  using assms
  by (simp add: and.commute bin-distribute-and-over-or)

```

```

lemma intval-eliminate-y:
  assumes  $val[y \& z] = IntVal\ b\ 0$ 
  shows  $val[(x | y) \& z] = val[x \& z]$ 
  using assms bin-eliminate-y by (cases x; cases y; cases z; auto)

```

```

lemma intval-and-associative:
   $val[(x \& y) \& z] = val[x \& (y \& z)]$ 
apply (cases x; cases y; cases z; auto)
by (simp add: and.assoc)+

```

```

lemma intval-or-associative:
   $val[(x | y) | z] = val[x | (y | z)]$ 

```

```

apply (cases x; cases y; cases z; auto)
by (simp add: or.assoc)+

lemma intval-xor-associative:
  val[(x ⊕ y) ⊕ z] = val[x ⊕ (y ⊕ z)]
apply (cases x; cases y; cases z; auto)
by (simp add: xor.assoc)+

lemma exp-and-associative:
  exp[(x & y) & z] ≥ exp[x & (y & z)]
apply simp using intval-and-associative by fastforce

lemma exp-or-associative:
  exp[(x | y) | z] ≥ exp[x | (y | z)]
apply simp using intval-or-associative by fastforce

lemma exp-xor-associative:
  exp[(x ⊕ y) ⊕ z] ≥ exp[x ⊕ (y ⊕ z)]
apply simp using intval-xor-associative by fastforce

lemma intval-and-absorb-or:
  assumes ∃ b v . x = new-int b v
  assumes val[x & (x | y)] ≠ UndefVal
  shows val[x & (x | y)] = val[x]
  using assms apply (cases x; cases y; auto)
  by (metis (mono-tags, lifting) intval-and.simps(5))

lemma intval-or-absorb-and:
  assumes ∃ b v . x = new-int b v
  assumes val[x | (x & y)] ≠ UndefVal
  shows val[x | (x & y)] = val[x]
  using assms apply (cases x; cases y; auto)
  by (metis (mono-tags, lifting) intval-or.simps(5))

lemma exp-and-absorb-or:
  exp[x & (x | y)] ≥ exp[x]
  apply auto using intval-and-absorb-or eval-unused-bits-zero
  by (smt (verit) evalDet intval-or.elims new-int.elims)

lemma exp-or-absorb-and:
  exp[x | (x & y)] ≥ exp[x]
  apply auto using intval-or-absorb-and eval-unused-bits-zero
  by (smt (verit) evalDet intval-or.elims new-int.elims)

lemma
  assumes y = 0
  shows x + y = or x y
  using assms

```

by *simp*

lemma *no-overlap-or*:

assumes *and* $x \ y = 0$

shows $x + y = \text{or } x \ y$

using *assms*

by (*metis bit-and-iff bit-xor-iff disjunctive-add xor-self-eq*)

context *stamp-mask*

begin

lemma *intval-up-and-zero-implies-zero*:

assumes *and* $(\uparrow x) (\uparrow y) = 0$

assumes $[m, p] \vdash x \mapsto xv$

assumes $[m, p] \vdash y \mapsto yv$

assumes $\text{val}[xv \ \& \ yv] \neq \text{UndefVal}$

shows $\exists \ b. \text{val}[xv \ \& \ yv] = \text{new-int } b \ 0$

using *assms* **apply** (*cases xv; cases yv; auto*)

using *up-mask-and-zero-implies-zero*

apply (*smt (verit, best) take-bit-and take-bit-of-0*)

by *presburger*

lemma *exp-eliminate-y*:

and $(\uparrow y) (\uparrow z) = 0 \longrightarrow \text{BinaryExpr BinAnd } (\text{BinaryExpr BinOr } x \ y) \ z \geq \text{BinaryExpr BinAnd } x \ z$

apply *simp* **apply** (*rule impI; rule allI; rule allI; rule allI*)

subgoal **premises** *p* **for** *m p v* **apply** (*rule impI*) **subgoal** **premises** *e*

proof –

obtain *xv* **where** *xv*: $[m, p] \vdash x \mapsto xv$

using *e* **by** *auto*

obtain *yv* **where** *yv*: $[m, p] \vdash y \mapsto yv$

using *e* **by** *auto*

obtain *zv* **where** *zv*: $[m, p] \vdash z \mapsto zv$

using *e* **by** *auto*

have *lhs*: $v = \text{val}[(xv \mid yv) \ \& \ zv]$

using *xv yv zv*

by (*smt (verit, best) BinaryExprE bin-eval.simps(4) bin-eval.simps(5) e*

evalDet)

then **have** $v = \text{val}[(xv \ \& \ zv) \mid (yv \ \& \ zv)]$

by (*simp add: intval-and-commute intval-distribute-and-over-or*)

```

also have  $\exists b. \text{val}[yv \ \& \ zv] = \text{new-int } b \ 0$ 
using intval-up-and-zero-implies-zero
by (metis calculation e intval-or.simps(5) p unfold-binary yv zv)
ultimately have rhs:  $v = \text{val}[xv \ \& \ zv]$ 
using intval-eliminate-y lhs by force
from lhs rhs show ?thesis
by (metis BinaryExpr BinaryExprE bin-eval.simps(4) e xv zv)
qed
done
done

```

```

lemma leadingZeroBounds:
  fixes  $x :: 'a::\text{len word}$ 
  assumes  $n = \text{numberOfLeadingZeros } x$ 
  shows  $0 \leq n \wedge n \leq \text{Nat.size } x$ 
  using assms unfolding numberOfLeadingZeros-def
  by (simp add: MaxOrNeg-def highestOneBit-def nat-le-iff)

```

```

lemma above-nth-not-set:
  fixes  $x :: \text{int64}$ 
  assumes  $n = 64 - \text{numberOfLeadingZeros } x$ 
  shows  $j > n \longrightarrow \neg(\text{bit } x \ j)$ 
  using assms unfolding numberOfLeadingZeros-def
  by (smt (verit, ccfv-SIG) highestOneBit-def int-nat-eq int-ops(6) less-imp-of-nat-less
max-set-bit size64 zerosAboveHighestOne)

```

no-notation *LogicNegationNotation* (!-)

```

lemma zero-horner:
  horner-sum of-bool 2 (map (\x. False) xs) = 0
  apply (induction xs) apply simp
  by force

```

```

lemma zero-map:
  assumes  $j \leq n$ 
  assumes  $\forall i. j \leq i \longrightarrow \neg(f \ i)$ 
  shows  $\text{map } f \ [0..<n] = \text{map } f \ [0..<j] @ \text{map } (\lambda x. \text{False}) \ [j..<n]$ 
  apply (insert assms)
  by (smt (verit, del-insts) add-diff-inverse-nat atLeastLessThan-iff bot-nat-0.extremum
leD map-append map-eq-conv set-upt upt-add-eq-append)

```

```

lemma map-join-horner:
  assumes  $\text{map } f \ [0..<n] = \text{map } f \ [0..<j] @ \text{map } (\lambda x. \text{False}) \ [j..<n]$ 
  shows  $\text{horner-sum of-bool } (2::'a::\text{len word}) \ (\text{map } f \ [0..<n]) = \text{horner-sum of-bool}$ 
 $2 \ (\text{map } f \ [0..<j])$ 
proof -
  have  $\text{horner-sum of-bool } (2::'a::\text{len word}) \ (\text{map } f \ [0..<n]) = \text{horner-sum of-bool}$ 
 $2 \ (\text{map } f \ [0..<j]) + 2 \wedge \text{length } [0..<j] * \text{horner-sum of-bool } 2 \ (\text{map } f \ [j..<n])$ 
  using horner-sum-append

```

```

    by (smt (verit) assms diff-le-self diff-zero le-add-same-cancel2 length-append
length-map length-upt map-append upt-add-eq-append)
    also have ... = horner-sum of-bool 2 (map f [0..<j]) + 2 ^ length [0..<j] *
horner-sum of-bool 2 (map (λx. False) [j..<n])
    using assms
    by (metis calculation horner-sum-append length-map)
    also have ... = horner-sum of-bool 2 (map f [0..<j])
    using zero-horner
    using mult-not-zero by auto
    finally show ?thesis by simp
qed

```

lemma *split-horner*:

```

    assumes j ≤ n
    assumes ∀ i. j ≤ i ⟶ ¬(f i)
    shows horner-sum of-bool (2::'a::len word) (map f [0..<n]) = horner-sum of-bool
2 (map f [0..<j])
    apply (rule map-join-horner)
    apply (rule zero-map)
    using assms by auto

```

lemma *transfer-map*:

```

    assumes ∀ i. i < n ⟶ f i = f' i
    shows (map f [0..<n]) = (map f' [0..<n])
    using assms by simp

```

lemma *transfer-horner*:

```

    assumes ∀ i. i < n ⟶ f i = f' i
    shows horner-sum of-bool (2::'a::len word) (map f [0..<n]) = horner-sum of-bool
2 (map f' [0..<n])
    using assms using transfer-map
    by (smt (verit, best))

```

lemma *L1*:

```

    assumes n = 64 - numberOfLeadingZeros (↑z)
    assumes [m, p] ⊢ z ↦ IntVal b zv
    shows and v zv = and (v mod 2^n) zv
proof -
    have nle: n ≤ 64
    using assms
    using diff-le-self by blast
    also have and v zv = horner-sum of-bool 2 (map (bit (and v zv)) [0..<64])
    using horner-sum-bit-eq-take-bit size64
    by (metis size-word.rep-eq take-bit-length-eq)
    also have ... = horner-sum of-bool 2 (map (λi. bit (and v zv) i) [0..<64])
    by blast
    also have ... = horner-sum of-bool 2 (map (λi. ((bit v i) ∧ (bit zv i))) [0..<64])
    using bit-and-iff by metis
    also have ... = horner-sum of-bool 2 (map (λi. ((bit v i) ∧ (bit zv i))) [0..<n])

```



```

proof –
  have  $\forall i. i \geq n \longrightarrow \neg(\text{bit } zv \ i)$ 
    using above-nth-not-set assms(1)
    using assms(2) not-may-implies-false
    by smt (verit, ccfv-SIG) One-nat-def diff-less int-ops(6) leadingZerosAddHighestOne linorder-not-le nat-int-comparison(2) not-numeral-le-zero size64 zero-less-Suc zerosAboveHighestOne)
  then have  $\forall i. i \geq n \longrightarrow \neg((\text{bit } v \ i) \wedge (\text{bit } zv \ i))$ 
    by auto
  then show ?thesis using nle split-horner
    by (metis (no-types, lifting))
qed
  also have  $\dots = \text{horner-sum of-bool } 2 \ (\text{map } (\lambda i. ((\text{bit } (v \bmod 2^n) \ i) \wedge (\text{bit } zv \ i))) \ [0..<n])$ 
proof –
  have  $\forall i. i < n \longrightarrow \text{bit } (v \bmod 2^n) \ i = \text{bit } v \ i$ 
    by (metis bit-take-bit-iff take-bit-eq-mod)
  then have  $\forall i. i < n \longrightarrow ((\text{bit } v \ i) \wedge (\text{bit } zv \ i)) = ((\text{bit } (v \bmod 2^n) \ i) \wedge (\text{bit } zv \ i))$ 
    by force
  then show ?thesis
    by (rule transfer-horner)
qed
  also have  $\dots = \text{horner-sum of-bool } 2 \ (\text{map } (\lambda i. ((\text{bit } (v \bmod 2^n) \ i) \wedge (\text{bit } zv \ i))) \ [0..<64])$ 
proof –
  have  $\forall i. i \geq n \longrightarrow \neg(\text{bit } zv \ i)$ 
    using above-nth-not-set assms(1)
    using assms(2) not-may-implies-false
    by smt (verit, ccfv-SIG) One-nat-def diff-less int-ops(6) leadingZerosAddHighestOne linorder-not-le nat-int-comparison(2) not-numeral-le-zero size64 zero-less-Suc zerosAboveHighestOne)
  then show ?thesis
    by (metis (no-types, lifting) assms(1) diff-le-self split-horner)
qed
  also have  $\dots = \text{horner-sum of-bool } 2 \ (\text{map } (\text{bit } (\text{and } (v \bmod 2^n) \ zv)) \ [0..<64])$ 
    by (meson bit-and-iff)
  also have  $\dots = \text{and } (v \bmod 2^n) \ zv$ 
    using horner-sum-bit-eq-take-bit size64
    by (metis size-word.rep-eq take-bit-length-eq)
  finally show ?thesis
    using  $\langle \text{and } (v::64 \ \text{word}) \ (zv::64 \ \text{word}) = \text{horner-sum of-bool } (2::64 \ \text{word}) \ (\text{map } (\text{bit } (\text{and } v \ zv)) \ [0::\text{nat}..<64::\text{nat}]) \rangle \langle \text{horner-sum of-bool } (2::64 \ \text{word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } ((v::64 \ \text{word}) \bmod (2::64 \ \text{word}) \wedge (n::\text{nat})) \ i \wedge \text{bit } (zv::64 \ \text{word}) \ i) \ [0::\text{nat}..<64::\text{nat}]) = \text{horner-sum of-bool } (2::64 \ \text{word}) \ (\text{map } (\text{bit } (\text{and } (v \bmod (2::64 \ \text{word}) \wedge n) \ zv)) \ [0::\text{nat}..<64::\text{nat}]) \rangle \langle \text{horner-sum of-bool } (2::64 \ \text{word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } ((v::64 \ \text{word}) \bmod (2::64 \ \text{word}) \wedge (n::\text{nat})) \ i \wedge \text{bit } (zv::64 \ \text{word}) \ i) \ [0::\text{nat}..<n]) = \text{horner-sum of-bool } (2::64 \ \text{word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } (v \bmod (2::64 \ \text{word}) \wedge n) \ i \wedge \text{bit } zv \ i) \ [0::\text{nat}..<64::\text{nat}]) \rangle \langle \text{horner-sum of-bool } (2::64 \ \text{word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } (v \bmod (2::64 \ \text{word}) \wedge n) \ i \wedge \text{bit } zv \ i) \ [0::\text{nat}..<64::\text{nat}]) \rangle$ 

```

$(v::64 \text{ word}) \ i \wedge \text{bit } (zv::64 \text{ word}) \ i \ [0::\text{nat}..<64::\text{nat}]] = \text{horner-sum of-bool } (2::64 \text{ word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } v \ i \wedge \text{bit } zv \ i) \ [0::\text{nat}..<n::\text{nat}])) \rangle \langle \text{horner-sum of-bool } (2::64 \text{ word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } (v::64 \text{ word}) \ i \wedge \text{bit } (zv::64 \text{ word}) \ i) \ [0::\text{nat}..<n::\text{nat}])) = \text{horner-sum of-bool } (2::64 \text{ word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } (v \bmod (2::64 \text{ word}) \wedge n) \ i \wedge \text{bit } zv \ i) \ [0::\text{nat}..<n])) \rangle \langle \text{horner-sum of-bool } (2::64 \text{ word}) \ (\text{map } (\text{bit } (\text{and } ((v::64 \text{ word}) \bmod (2::64 \text{ word}) \wedge (n::\text{nat})) \ (zv::64 \text{ word}))) \ [0::\text{nat}..<64::\text{nat}])) = \text{and } (v \bmod (2::64 \text{ word}) \wedge n) \ zv \rangle \langle \text{horner-sum of-bool } (2::64 \text{ word}) \ (\text{map } (\text{bit } (\text{and } (v::64 \text{ word}) \ (zv::64 \text{ word}))) \ [0::\text{nat}..<64::\text{nat}])) = \text{horner-sum of-bool } (2::64 \text{ word}) \ (\text{map } (\lambda i::\text{nat}. \text{bit } v \ i \wedge \text{bit } zv \ i) \ [0::\text{nat}..<64::\text{nat}])) \rangle$ **by** *presburger*

qed

lemma *up-mask-upper-bound*:

assumes $[m, p] \vdash x \mapsto \text{IntVal } b \ xv$

shows $xv \leq (\uparrow x)$

using *assms*

by (*metis* (*no-types*, *lifting*) *and*.*idem* *and*.*right-neutral bit.conj-cancel-left bit.conj-disj-distrib*(1) *bit.double-compl ucast-id up-spec word-and-le1 word-not-dist*(2))

lemma *L2*:

assumes $\text{numberOfLeadingZeros } (\uparrow z) + \text{numberOfTrailingZeros } (\uparrow y) \geq 64$

assumes $n = 64 - \text{numberOfLeadingZeros } (\uparrow z)$

assumes $[m, p] \vdash z \mapsto \text{IntVal } b \ zv$

assumes $[m, p] \vdash y \mapsto \text{IntVal } b \ yv$

shows $yv \bmod 2^{\wedge n} = 0$

proof –

have $yv \bmod 2^{\wedge n} = \text{horner-sum of-bool } 2 \ (\text{map } (\text{bit } yv) \ [0..<n])$

by (*simp add: horner-sum-bit-eq-take-bit take-bit-eq-mod*)

also have $\dots \leq \text{horner-sum of-bool } 2 \ (\text{map } (\text{bit } (\uparrow y)) \ [0..<n])$

using *up-mask-upper-bound assms*(4)

by (*metis* (*no-types*, *opaque-lifting*) *and*.*right-neutral bit.conj-cancel-right bit.conj-disj-distrib*(1) *bit.double-compl horner-sum-bit-eq-take-bit take-bit-and ucast-id up-spec word-and-le1 word-not-dist*(2))

also have $\text{horner-sum of-bool } 2 \ (\text{map } (\text{bit } (\uparrow y)) \ [0..<n]) = \text{horner-sum of-bool } 2 \ (\text{map } (\lambda x. \text{False}) \ [0..<n])$

proof –

have $\forall i < n. \neg(\text{bit } (\uparrow y) \ i)$

using *assms*(1,2) *zerosBelowLowestOne*

by (*metis* *add commute add-diff-inverse-nat add-lessD1 leD le-diff-conv numberOfTrailingZeros-def*)

then show *?thesis*

by (*metis* (*full-types*) *transfer-map*)

qed

also have $\text{horner-sum of-bool } 2 \ (\text{map } (\lambda x. \text{False}) \ [0..<n]) = 0$

using *zero-horner*

by *blast*

finally show *?thesis*

by *auto*

qed

thm-oracles *L1 L2*

lemma *unfold-binary-width-add:*

shows $([m,p] \vdash \text{BinaryExpr BinAdd } xe \ ye \mapsto \text{IntVal } b \ val) = (\exists \ x \ y. \\ (([m,p] \vdash xe \mapsto \text{IntVal } b \ x) \wedge \\ ([m,p] \vdash ye \mapsto \text{IntVal } b \ y) \wedge \\ (\text{IntVal } b \ val = \text{bin-eval BinAdd } (\text{IntVal } b \ x) (\text{IntVal } b \ y)) \wedge \\ (\text{IntVal } b \ val \neq \text{UndefVal})) \\)) \text{ (is } ?L = ?R)$

proof (*intro iffI*)
assume $\mathcal{I}: ?L$
show $?R$ **apply** (*rule evaltree.cases[OF \mathcal{I}]*)
apply *force+* **apply** *auto[1]*
apply (*smt (verit) intval-add.elims intval-bits.simps*)
by *blast*

next
assume $R: ?R$
then obtain $x \ y$ **where** $[m,p] \vdash xe \mapsto \text{IntVal } b \ x$
and $[m,p] \vdash ye \mapsto \text{IntVal } b \ y$
and $\text{new-int } b \ val = \text{bin-eval BinAdd } (\text{IntVal } b \ x) (\text{IntVal } b \ y)$
and $\text{new-int } b \ val \neq \text{UndefVal}$
by *auto*
then show $?L$
using R **by** *blast*

qed

lemma *unfold-binary-width-and:*

shows $([m,p] \vdash \text{BinaryExpr BinAnd } xe \ ye \mapsto \text{IntVal } b \ val) = (\exists \ x \ y. \\ (([m,p] \vdash xe \mapsto \text{IntVal } b \ x) \wedge \\ ([m,p] \vdash ye \mapsto \text{IntVal } b \ y) \wedge \\ (\text{IntVal } b \ val = \text{bin-eval BinAnd } (\text{IntVal } b \ x) (\text{IntVal } b \ y)) \wedge \\ (\text{IntVal } b \ val \neq \text{UndefVal})) \\)) \text{ (is } ?L = ?R)$

proof (*intro iffI*)
assume $\mathcal{I}: ?L$
show $?R$ **apply** (*rule evaltree.cases[OF \mathcal{I}]*)
apply *force+* **apply** *auto[1]* **using** *intval-and.elims intval-bits.simps*
apply (*smt (verit) new-int.simps new-int-bin.simps take-bit-and*)
by *blast*

next
assume $R: ?R$
then obtain $x \ y$ **where** $[m,p] \vdash xe \mapsto \text{IntVal } b \ x$
and $[m,p] \vdash ye \mapsto \text{IntVal } b \ y$
and $\text{new-int } b \ val = \text{bin-eval BinAnd } (\text{IntVal } b \ x) (\text{IntVal } b \ y)$
and $\text{new-int } b \ val \neq \text{UndefVal}$
by *auto*
then show $?L$
using R **by** *blast*

qed

lemma *mod-dist-over-add-right*:

fixes $a\ b\ c :: \text{int64}$
fixes $n :: \text{nat}$
assumes $1: 0 < n$
assumes $2: n < 64$
shows $(a + b \bmod 2^n) \bmod 2^n = (a + b) \bmod 2^n$
using *mod-dist-over-add*
by (*simp add: 1 2 add.commute*)

lemma *numberOfLeadingZeros-range*:

$0 \leq \text{numberOfLeadingZeros } n \wedge \text{numberOfLeadingZeros } n \leq \text{Nat.size } n$
unfolding *numberOfLeadingZeros-def highestOneBit-def* **using** *max-set-bit*
by (*simp add: highestOneBit-def leadingZeroBounds numberOfLeadingZeros-def*)

lemma *improved-opt*:

assumes $\text{numberOfLeadingZeros } (\uparrow z) + \text{numberOfTrailingZeros } (\uparrow y) \geq 64$
shows $\text{exp}[(x + y) \& z] \geq \text{exp}[x \& z]$
apply *simp* **apply** (*rule allI*)+; *rule impI*
subgoal **premises** *eval* **for** $m\ p\ v$

proof –

obtain n **where** $n: n = 64 - \text{numberOfLeadingZeros } (\uparrow z)$
by *simp*
obtain $b\ \text{val}$ **where** $\text{val}: [m, p] \vdash \text{exp}[(x + y) \& z] \mapsto \text{IntVal } b\ \text{val}$
by (*metis BinaryExprE bin-eval-new-int eval new-int.simps*)
then obtain $xv\ yv$ **where** $\text{addv}: [m, p] \vdash \text{exp}[x + y] \mapsto \text{IntVal } b\ (xv + yv)$
apply (*subst (asm) unfold-binary-width-and*) **by** (*metis add.right-neutral*)
then obtain yv **where** $yv: [m, p] \vdash y \mapsto \text{IntVal } b\ yv$
apply (*subst (asm) unfold-binary-width-add*) **by** *blast*
from addv **obtain** xv **where** $xv: [m, p] \vdash x \mapsto \text{IntVal } b\ xv$
apply (*subst (asm) unfold-binary-width-add*) **by** *blast*
from val **obtain** zv **where** $zv: [m, p] \vdash z \mapsto \text{IntVal } b\ zv$
apply (*subst (asm) unfold-binary-width-and*) **by** *blast*
have $\text{addv}: [m, p] \vdash \text{exp}[x + y] \mapsto \text{new-int } b\ (xv + yv)$
apply (*rule evaltree.BinaryExpr*)
using xv **apply** *simp*
using yv **apply** *simp*
by *simp*+
have $\text{lhs}: [m, p] \vdash \text{exp}[(x + y) \& z] \mapsto \text{new-int } b\ (\text{and } (xv + yv)\ zv)$
apply (*rule evaltree.BinaryExpr*)
using addv **apply** *simp*
using zv **apply** *simp*
using addv **apply** *auto*[1]
by *simp*
have $\text{rhs}: [m, p] \vdash \text{exp}[x \& z] \mapsto \text{new-int } b\ (\text{and } xv\ zv)$
apply (*rule evaltree.BinaryExpr*)
using xv **apply** *simp*
using zv **apply** *simp*
apply *force*

```

    by simp
  then show ?thesis
proof (cases numberOfLeadingZeros ( $\uparrow z$ ) > 0)
  case True
  have n-bounds:  $0 \leq n \wedge n < 64$ 
  using diff-le-self n numberOfLeadingZeros-range
  by (simp add: True)
  have and (xv + yv) zv = and ((xv + yv) mod  $2^n$ ) zv
  using L1 n zv by blast
  also have ... = and ((xv + (yv mod  $2^n$ )) mod  $2^n$ ) zv
  using mod-dist-over-add-right n-bounds
  by (metis take-bit-0 take-bit-eq-mod zero-less-iff-neq-zero)
  also have ... = and (((xv mod  $2^n$ ) + (yv mod  $2^n$ )) mod  $2^n$ ) zv
  by (metis bits-mod-by-1 mod-dist-over-add n-bounds order-le-imp-less-or-eq
power-0)
  also have ... = and ((xv mod  $2^n$ ) mod  $2^n$ ) zv
  using L2 n zv yv
  using assms by auto
  also have ... = and (xv mod  $2^n$ ) zv
  using mod-mod-trivial
  by (smt (verit, best) and.idem take-bit-eq-mask take-bit-eq-mod word-bw-assocs(1))
  also have ... = and xv zv
  using L1 n zv by metis
  finally show ?thesis
  using eval lhs rhs
  by (metis evalDet)
next
  case False
  then have numberOfLeadingZeros ( $\uparrow z$ ) = 0
  by simp
  then have numberOfTrailingZeros ( $\uparrow y$ )  $\geq 64$ 
  using assms(1)
  by fastforce
  then have yv = 0
  using yv
  by (metis (no-types, lifting) L1 L2 add-diff-cancel-left' and.comm-neutral
and.idem bit.compl-zero bit.conj-cancel-right bit.conj-disj-distrib(1) bit.double-compl
less-imp-diff-less linorder-not-le word-not-dist(2))
  then show ?thesis
  by (metis add.right-neutral eval evalDet lhs rhs)
qed
qed
done

thm-oracles improved-opt

```

```

phase NewAnd
  terminating size
begin

```

optimization *redundant-lhs-x-or:* $((x \mid y) \& z) \mapsto y \& z$
 when $((\text{and } (\text{IExpr-up } x) (\text{IExpr-up } z)) = 0)$
apply (*simp add: IExpr-up-def*)
using *simple-mask.exp-eliminate-y*
by (*meson exp-or-commute mono-binary order-refl order-trans*)

optimization *redundant-rhs-y-or:* $(z \& (x \mid y)) \mapsto z \& x$
 when $((\text{and} (\text{IExpr-up } y) (\text{IExpr-up } z)) = 0)$
apply (*simp add: IExpr-up-def*)
using *simple-mask.exp-eliminate-y*
by (*meson exp-and-commute order.trans*)

end

```
theory NotPhase
  imports
    Common
begin
```

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lemma *bin-not-cancel*:

$\text{bin}[\neg(\neg(e))] = \text{bin}[e]$

by *auto*

lemma *val-not-cancel*:

assumes $\text{val}[\sim(\text{new-int } b \ v)] \neq \text{UndefVal}$

shows $\text{val}[\sim(\sim(\text{new-int } b \ v))] = (\text{new-int } b \ v)$

by (*simp add: take-bit-not-take-bit*)

lemma *exp-not-cancel*:

$\text{exp}[\sim(\sim a)] \geq \text{exp}[a]$

using *val-not-cancel* **apply** *auto*

by (*metis eval-unused-bits-zero intval-logic-negation.cases new-int.simps intval-not.simps(1)*

intval-not.simps(2) intval-not.simps(3) intval-not.simps(4))

Optimisations

optimization *NotCancel*: $\text{exp}[\sim(\sim a)] \mapsto a$

by (*metis exp-not-cancel*)

end

end

11.9 OrNode Phase

theory *OrPhase*

imports

Common

begin

context *stamp-mask*

begin

Taking advantage of the truth table of or operations.

#	x	y	$x y$
1	0	0	0
2	0	1	1
3	1	0	1
4	1	1	1

If row 2 never applies, that is, $\text{canBeZero } x \ \& \ \text{canBeOne } y = 0$, then $(x|y) = x$.

Likewise, if row 3 never applies, $\text{canBeZero } y \ \& \ \text{canBeOne } x = 0$, then $(x|y) = y$.

```

lemma OrLeftFallthrough:
  assumes  $(\text{and } (\text{not } (\downarrow x)) (\uparrow y)) = 0$ 
  shows  $\text{exp}[x \mid y] \geq \text{exp}[x]$ 
  using assms
  apply simp apply  $((\text{rule } \text{allI})+; \text{rule } \text{impI})$ 
  subgoal premises eval for  $m \ p \ v$ 
  proof –
    obtain  $b \ vv$  where  $e: [m, p] \vdash \text{exp}[x \mid y] \mapsto \text{IntVal } b \ vv$ 
      by  $(\text{metis } \text{BinaryExprE } \text{bin-eval-new-int } \text{new-int.simps } \text{eval}(2))$ 
    from  $e$  obtain  $xv$  where  $xv: [m, p] \vdash x \mapsto \text{IntVal } b \ xv$ 
      apply  $(\text{subst } (\text{asm}) \text{ unfold-binary-width})$ 
      by force+
    from  $e$  obtain  $yv$  where  $yv: [m, p] \vdash y \mapsto \text{IntVal } b \ yv$ 
      apply  $(\text{subst } (\text{asm}) \text{ unfold-binary-width})$ 
      by force+
    have  $v\text{def}: v = \text{intval-or } (\text{IntVal } b \ xv) (\text{IntVal } b \ yv)$ 
      by  $(\text{metis } \text{bin-eval.simps}(5) \text{ eval}(2) \text{ evalDet } \text{unfold-binary } xv \ yv)$ 
    have  $\forall i. (\text{bit } xv \ i) \mid (\text{bit } yv \ i) = (\text{bit } xv \ i)$ 
      by  $(\text{metis } \text{assms } \text{bit-and-iff } \text{not-down-up-mask-and-zero-implies-zero } xv \ yv)$ 
    then have  $\text{IntVal } b \ xv = \text{intval-or } (\text{IntVal } b \ xv) (\text{IntVal } b \ yv)$ 
    by  $(\text{smt } (\text{verit}, \text{ccfv-threshold}) \text{ and.idem } \text{assms } \text{bit.conj-disj-distrib } \text{eval-unused-bits-zero}$ 
        
 $\text{intval-or.simps}(1) \text{ new-int.simps } \text{new-int-bin.simps } \text{not-down-up-mask-and-zero-implies-zero}$ 
        
 $\text{word-ao-absorbs}(3) \ xv \ yv)$ 
    then show ?thesis
      using  $xv \ v\text{def}$  by presburger
  qed
done

```

```

lemma OrRightFallthrough:
  assumes  $(\text{and } (\text{not } (\downarrow y)) (\uparrow x)) = 0$ 
  shows  $\text{exp}[x \mid y] \geq \text{exp}[y]$ 
  using assms
  apply simp apply  $((\text{rule } \text{allI})+; \text{rule } \text{impI})$ 
  subgoal premises eval for  $m \ p \ v$ 
  proof –
    obtain  $b \ vv$  where  $e: [m, p] \vdash \text{exp}[x \mid y] \mapsto \text{IntVal } b \ vv$ 
      by  $(\text{metis } \text{BinaryExprE } \text{bin-eval-new-int } \text{new-int.simps } \text{eval}(2))$ 
    from  $e$  obtain  $xv$  where  $xv: [m, p] \vdash x \mapsto \text{IntVal } b \ xv$ 
      apply  $(\text{subst } (\text{asm}) \text{ unfold-binary-width})$ 
      by force+
    from  $e$  obtain  $yv$  where  $yv: [m, p] \vdash y \mapsto \text{IntVal } b \ yv$ 
      apply  $(\text{subst } (\text{asm}) \text{ unfold-binary-width})$ 
      by force+
    have  $v\text{def}: v = \text{intval-or } (\text{IntVal } b \ xv) (\text{IntVal } b \ yv)$ 
      by  $(\text{metis } \text{bin-eval.simps}(5) \text{ eval}(2) \text{ evalDet } \text{unfold-binary } xv \ yv)$ 
    have  $\forall i. (\text{bit } xv \ i) \mid (\text{bit } yv \ i) = (\text{bit } yv \ i)$ 
      by  $(\text{metis } \text{assms } \text{bit-and-iff } \text{not-down-up-mask-and-zero-implies-zero } xv \ yv)$ 

```



```

    then have IntVal b yv = intval-or (IntVal b xv) (IntVal b yv)
    by (metis (no-types, lifting) assms eval-unused-bits-zero intval-or.simps(1)
new-int.elims
      new-int-bin.elims stamp-mask.not-down-up-mask-and-zero-implies-zero
stamp-mask-axioms
      word-ao-absorbs(8) xv yv)
    then show ?thesis
    using vdef yv by presburger
  qed
done

end

phase OrNode
  terminating size
begin

lemma bin-or-equal:
  bin[x | x] = bin[x]
  by simp

lemma bin-shift-const-right-helper:
  x | y = y | x
  by simp

lemma bin-or-not-operands:
  (¬x | ¬y) = (¬(x & y))
  by simp

lemma val-or-equal:
  assumes x = new-int b v
  and     (val[x | x] ≠ UndefVal)
  shows   val[x | x] = val[x]
  apply (cases x; auto) using bin-or-equal assms
  by auto+

lemma val-elim-redundant-false:
  assumes x = new-int b v
  and     val[x | false] ≠ UndefVal
  shows   val[x | false] = val[x]
  using assms apply (cases x; auto) by presburger

lemma val-shift-const-right-helper:
  val[x | y] = val[y | x]
  apply (cases x; cases y; auto)
  by (simp add: or.commute)+

```

lemma *val-or-not-operands*:
 $val[\sim x \mid \sim y] = val[\sim(x \& y)]$
apply (*cases x; cases y; auto*)
by (*simp add: take-bit-not-take-bit*)

lemma *exp-or-equal*:
 $exp[x \mid x] \geq exp[x]$
using *val-or-equal* **apply** *auto[1]*
by (*smt (verit, ccfv-SIG) evalDet eval-unused-bits-zero intval-negate.elims int-*
val-or.simps(2)
intval-or.simps(6) intval-or.simps(7) new-int.simps val-or-equal)

lemma *exp-elim-redundant-false*:
 $exp[x \mid false] \geq exp[x]$
using *val-elim-redundant-false* **apply** *auto[1]*
by (*smt (verit) Value.sel(1) eval-unused-bits-zero intval-or.elims new-int.simps*
new-int-bin.simps val-elim-redundant-false)

Optimisations

optimization *OrEqual*: $x \mid x \mapsto x$
by (*meson exp-or-equal*)

optimization *OrShiftConstantRight*: $((const\ x) \mid y) \mapsto y \mid (const\ x)$ when $\neg(is-ConstantExpr\ y)$
using *size-flip-binary* **apply** *force*
apply *auto[1]*
by (*simp add: BinaryExpr unfold-const val-shift-const-right-helper*)

optimization *EliminateRedundantFalse*: $x \mid false \mapsto x$
by (*meson exp-elim-redundant-false*)

optimization *OrNotOperands*: $(\sim x \mid \sim y) \mapsto \sim(x \& y)$
apply (*metis add-2-eq-Suc' less-SucI not-add-less1 not-less-eq size-binary-const*
size-non-add)
apply *auto[1]*
by (*metis BinaryExpr UnaryExpr bin-eval.simps(4) intval-not.simps(2) unary-eval.simps(3)*
val-or-not-operands)

optimization *OrLeftFallthrough*:
 $x \mid y \mapsto x$ when $((and\ (not\ (IExpr-down\ x))\ (IExpr-up\ y)) = 0)$
using *simple-mask.OrLeftFallthrough* **by** *blast*

optimization *OrRightFallthrough*:
 $x \mid y \mapsto y$ when $((and\ (not\ (IExpr-down\ y))\ (IExpr-up\ x)) = 0)$
using *simple-mask.OrRightFallthrough* **by** *blast*

end

end

11.10 ShiftNode Phase

theory *ShiftPhase*

imports

Common

begin

phase *ShiftNode*

terminating *size*

begin

fun *intval-log2* :: *Value* \Rightarrow *Value* **where**

intval-log2 (*IntVal* *b v*) = *IntVal* *b* (*word-of-int* (*SOME* *e. v=2^e*)) |

intval-log2 - = *UndefVal*

fun *in-bounds* :: *Value* \Rightarrow *int* \Rightarrow *int* \Rightarrow *bool* **where**

in-bounds (*IntVal* *b v*) *l h* = (*l* < *sint* *v* \wedge *sint* *v* < *h*) |

in-bounds - *l h* = *False*

lemma

assumes *in-bounds* (*intval-log2* *val-c*) 0 32

shows *intval-left-shift* *x* (*intval-log2* *val-c*) = *intval-mul* *x* *val-c*

apply (*cases* *val-c*; *auto*) **using** *intval-left-shift.simps*(1) *intval-mul.simps*(1) *intval-log2.simps*(1)

sorry

lemma *e-intval*:

n = *intval-log2* *val-c* \wedge *in-bounds* *n* 0 32 \longrightarrow

intval-left-shift *x* (*intval-log2* *val-c*) =

intval-mul *x* *val-c*

proof (*rule* *impI*)

assume *n* = *intval-log2* *val-c* \wedge *in-bounds* *n* 0 32

show *intval-left-shift* *x* (*intval-log2* *val-c*) =

intval-mul *x* *val-c*

proof (*cases* $\exists v . \text{val-c} = \text{IntVal } 32 \ v$)

case *True*

obtain *vc* **where** *val-c* = *IntVal* 32 *vc*

using *True* **by** *blast*

then have *n* = *IntVal* 32 (*word-of-int* (*SOME* *e. vc=2^e*))

using $\langle n = \text{intval-log2 } \text{val-c} \wedge \text{in-bounds } n \ 0 \ 32 \rangle$ *intval-log2.simps*(1) **by**

presburger

then show *?thesis* **sorry**

next

case *False*

then have $\exists v . \text{val-c} = \text{IntVal } 64 \ v$

```

      sorry
    then obtain  $vc$  where  $val-c = IntVal\ 64\ vc$ 
      by auto
    then have  $n = IntVal\ 64\ (word-of-int\ (SOME\ e.\ vc=2^e))$ 
      using  $\langle n = intval-log2\ val-c \wedge in-bounds\ n\ 0\ 32 \rangle\ intval-log2.simps(1)$  by
    presburger
    then show ?thesis sorry
  qed
qed

```

```

optimization  $e$ :
   $x * (const\ c) \mapsto x << (const\ n)$  when  $(n = intval-log2\ c \wedge in-bounds\ n\ 0\ 32)$ 
  using e-intval
  using BinaryExprE ConstantExprE bin-eval.simps(2,7) sorry
end
end

```

11.11 SignedDivNode Phase

```

theory SignedDivPhase
  imports
    Common
begin

phase SignedDivNode
  terminating size
begin

```

```

lemma val-division-by-one-is-self-32:
  assumes  $x = new-int\ 32\ v$ 
  shows  $intval-div\ x\ (IntVal\ 32\ 1) = x$ 
  using assms apply (cases x; auto)
  by (simp add: take-bit-signed-take-bit)

```

end

end

11.12 SignedRemNode Phase

```

theory SignedRemPhase
  imports

```

```

    Common
begin

phase SignedRemNode
  terminating size
begin

lemma val-remainder-one:
  assumes intval-mod  $x$  (IntVal 32 1)  $\neq$  UndefVal
  shows intval-mod  $x$  (IntVal 32 1) = IntVal 32 0
  using assms apply (cases  $x$ ; auto) sorry

value word-of-int (sint ( $x2::32$  word) smod 1)

end

end

```

11.13 SubNode Phase

```

theory SubPhase
  imports
    Common
    Proofs.StampEvalThms
begin

phase SubNode
  terminating size
begin

lemma bin-sub-after-right-add:
  shows  $((x::('a::len)$  word) +  $(y::('a::len)$  word)) -  $y = x$ 
  by simp

lemma sub-self-is-zero:
  shows  $(x::('a::len)$  word) -  $x = 0$ 
  by simp

lemma bin-sub-then-left-add:
  shows  $(x::('a::len)$  word) -  $(x + (y::('a::len)$  word)) =  $-y$ 
  by simp

lemma bin-sub-then-left-sub:
  shows  $(x::('a::len)$  word) -  $(x - (y::('a::len)$  word)) =  $y$ 
  by simp

lemma bin-subtract-zero:

```

```

shows (x :: 'a::len word) - (0 :: 'a::len word) = x
by simp

lemma bin-sub-negative-value:
  (x :: ('a::len) word) - (-(y :: ('a::len) word)) = x + y
by simp

lemma bin-sub-self-is-zero:
  (x :: ('a::len) word) - x = 0
by simp

lemma bin-sub-negative-const:
  (x :: 'a::len word) - (-(y :: 'a::len word)) = x + y
by simp

lemma val-sub-after-right-add-2:
  assumes x = new-int b v
  assumes val[(x + y) - y] ≠ UndefVal
  shows val[(x + y) - y] = val[x]
  using bin-sub-after-right-add
  using assms apply (cases x; cases y; auto)
  by (metis (full-types) intval-sub.simps(2))

lemma val-sub-after-left-sub:
  assumes val[(x - y) - x] ≠ UndefVal
  shows val[(x - y) - x] = val[-y]
  using assms apply (cases x; cases y; auto)
  using intval-sub.elims by fastforce

lemma val-sub-then-left-sub:
  assumes y = new-int b v
  assumes val[x - (x - y)] ≠ UndefVal
  shows val[x - (x - y)] = val[y]
  using assms apply (cases x; cases y; auto)
  by (metis (mono-tags) intval-sub.simps(5))

lemma val-subtract-zero:
  assumes x = new-int b v
  assumes intval-sub x (IntVal b 0) ≠ UndefVal
  shows intval-sub x (IntVal b 0) = val[x]
  using assms by (induction x; simp)

lemma val-zero-subtract-value:
  assumes x = new-int b v
  assumes intval-sub (IntVal b 0) x ≠ UndefVal
  shows intval-sub (IntVal b 0) x = val[-x]
  using assms by (induction x; simp)

```

```

lemma val-sub-then-left-add:
  assumes  $\text{val}[x - (x + y)] \neq \text{UndefVal}$ 
  shows  $\text{val}[x - (x + y)] = \text{val}[-y]$ 
  using assms apply (cases x; cases y; auto)
  by (metis (mono-tags, lifting) intval-sub.simps(5))

lemma val-sub-negative-value:
  assumes  $\text{val}[x - (-y)] \neq \text{UndefVal}$ 
  shows  $\text{val}[x - (-y)] = \text{val}[x + y]$ 
  using assms by (cases x; cases y; auto)

lemma val-sub-self-is-zero:
  assumes  $x = \text{new-int } b \ v \wedge \text{val}[x - x] \neq \text{UndefVal}$ 
  shows  $\text{val}[x - x] = \text{new-int } b \ 0$ 
  using assms by (cases x; auto)

lemma val-sub-negative-const:
  assumes  $y = \text{new-int } b \ v \wedge \text{val}[x - (-y)] \neq \text{UndefVal}$ 
  shows  $\text{val}[x - (-y)] = \text{val}[x + y]$ 
  using assms by (cases x; cases y; auto)

lemma exp-sub-after-right-add:
  shows  $\text{exp}[(x + y) - y] \geq \text{exp}[x]$ 
  apply auto
  by (smt (verit) evalDet eval-unused-bits-zero intval-add.elims new-int.simps
    val-sub-after-right-add-2)

lemma exp-sub-after-right-add2:
  shows  $\text{exp}[(x + y) - x] \geq \text{exp}[y]$ 
  using exp-sub-after-right-add apply auto
  by (smt (z3) Value.inject(1) diff-eq-eq evalDet eval-unused-bits-zero intval-add.elims
    intval-sub.elims new-int.simps new-int-bin.simps take-bit-dist-subL bin-eval.simps(1)
    bin-eval.simps(3) intval-add-sym unfold-binary)

lemma exp-sub-negative-value:
   $\text{exp}[x - (-y)] \geq \text{exp}[x + y]$ 
  apply simp
  by (smt (verit) bin-eval.simps(1) bin-eval.simps(3) evaltree-not-undef unary-eval.simps(2)
    unfold-binary unfold-unary val-sub-negative-value)

lemma exp-sub-then-left-sub:
   $\text{exp}[x - (x - y)] \geq \text{exp}[y]$ 
  using val-sub-then-left-sub apply auto
  subgoal premises p for m p xa xaa ya
  proof—

```

```

obtain xa where xa:  $[m, p] \vdash x \mapsto xa$ 
  using p(2) by blast
obtain ya where ya:  $[m, p] \vdash y \mapsto ya$ 
  using p(5) by auto
obtain xaa where xaa:  $[m, p] \vdash x \mapsto xaa$ 
  using p(2) by blast
have 1:  $val[xa - (xaa - ya)] \neq UndefinedVal$ 
  by (metis evalDet p(2) p(3) p(4) p(5) xa xaa ya)
then have  $val[xaa - ya] \neq UndefinedVal$ 
  by auto
then have  $[m, p] \vdash y \mapsto val[xa - (xaa - ya)]$ 
  by (metis 1 Value.exhaust evalDet eval-unused-bits-zero evaltree-not-undef
    intval-sub.simps(6) intval-sub.simps(7) new-int.simps p(5) val-sub-then-left-sub
    xa xaa ya)
  then show ?thesis
    by (metis evalDet p(2) p(4) p(5) xa xaa ya)
qed
done

```

thm-oracles *exp-sub-then-left-sub*

Optimisations

optimization *SubAfterAddRight*: $((x + y) - y) \mapsto x$
using *exp-sub-after-right-add* **by** *blast*

optimization *SubAfterAddLeft*: $((x + y) - x) \mapsto y$
using *exp-sub-after-right-add2* **by** *blast*

optimization *SubAfterSubLeft*: $((x - y) - x) \mapsto -y$
apply (*metis Suc-lessI add-2-eq-Suc' add-less-cancel-right less-trans-Suc not-add-less1*

size-binary-const size-binary-lhs size-binary-rhs size-non-add)
apply *auto*
by (*metis evalDet unary-eval.simps(2) unfold-unary val-sub-after-left-sub*)

optimization *SubThenAddLeft*: $(x - (x + y)) \mapsto -y$
apply *auto*
by (*metis evalDet unary-eval.simps(2) unfold-unary val-sub-then-left-add*)

optimization *SubThenAddRight*: $(y - (x + y)) \mapsto -x$
apply *auto*
by (*metis evalDet intval-add-sym unary-eval.simps(2) unfold-unary val-sub-then-left-add*)

optimization *SubThenSubLeft*: $(x - (x - y)) \mapsto y$
using *size-simps* **apply** *simp*
using *exp-sub-then-left-sub* **by** *blast*

optimization *SubtractZero*: $(x - (const IntVal b 0)) \mapsto x$

apply *auto*
by (*smt* (*verit*) *add.right-neutral diff-add-cancel eval-unused-bits-zero intval-sub.elims*
intval-word.simps new-int.simps new-int-bin.simps)

thm-oracles *SubtractZero*

optimization *SubNegativeValue*: $(x - (-y)) \mapsto x + y$
apply (*metis add-2-eq-Suc' less-SucI less-add-Suc1 not-less-eq size-binary-const*
size-non-add)
using *exp-sub-negative-value* **by** *simp*

thm-oracles *SubNegativeValue*

lemma *negate-idempotent*:
assumes $x = \text{IntVal } b \ v \wedge \text{take-bit } b \ v = v$
shows $x = \text{val}[-(-x)]$
using *assms*
using *is-IntVal-def* **by** *force*

optimization *ZeroSubtractValue*: $((\text{const IntVal } b \ 0) - x) \mapsto (-x)$
when (*wf-stamp* $x \wedge \text{stamp-expr } x = \text{IntegerStamp } b \ \text{lo}$
 $hi \wedge \neg(\text{is-ConstantExpr } x)$)
defer
apply *auto unfolding wf-stamp-def*
apply (*smt* (*verit*) *diff-0 intval-negate.simps(1) intval-sub.elims intval-word.simps*
new-int-bin.simps unary-eval.simps(2) unfold-unary)
using *add-2-eq-Suc' size.simps(2) size-flip-binary* **by** *presburger*

optimization *SubSelfIsZero*: $(x - x) \mapsto \text{const IntVal } b \ 0$ *when*
(wf-stamp $x \wedge \text{stamp-expr } x = \text{IntegerStamp } b \ \text{lo } hi)$
apply *simp-all*
apply *auto*
using *IRExpr.disc(42) One-nat-def size-non-const* **apply** *presburger*
by (*smt* (*verit*, *best*) *wf-value-def ConstantExpr evalDet eval-bits-1-64 eval-unused-bits-zero*
new-int.simps take-bit-of-0 val-sub-self-is-zero validDefIntConst valid-int
wf-stamp-def)

end

end

11.14 XorNode Phase

theory *XorPhase*

imports

Common

Proofs.StampEvalThms

begin

phase *XorNode*

terminating *size*

begin

lemma *bin-xor-self-is-false:*

$\text{bin}[x \oplus x] = 0$

by *simp*

lemma *bin-xor-commute:*

$\text{bin}[x \oplus y] = \text{bin}[y \oplus x]$

by (*simp add: xor.commute*)

lemma *bin-eliminate-redundant-false:*

$\text{bin}[x \oplus 0] = \text{bin}[x]$

by *simp*

lemma *val-xor-self-is-false:*

assumes $\text{val}[x \oplus x] \neq \text{UndefVal}$

shows $\text{val-to-bool}(\text{val}[x \oplus x]) = \text{False}$

using *assms* **by** (*cases x; auto*)

lemma *val-xor-self-is-false-2:*

assumes $(\text{val}[x \oplus x]) \neq \text{UndefVal}$

and $x = \text{IntVal } 32 \ v$

shows $\text{val}[x \oplus x] = \text{bool-to-val } \text{False}$

using *assms* **by** (*cases x; auto*)

lemma *val-xor-self-is-false-3:*

assumes $\text{val}[x \oplus x] \neq \text{UndefVal} \wedge x = \text{IntVal } 64 \ v$

shows $\text{val}[x \oplus x] = \text{IntVal } 64 \ 0$

using *assms* **by** (*cases x; auto*)

lemma *val-xor-commute:*

$\text{val}[x \oplus y] = \text{val}[y \oplus x]$

```

apply (cases x; cases y; auto)
by (simp add: xor.commute)+

lemma val-eliminate-redundant-false:
  assumes x = new-int b v
  assumes val[x  $\oplus$  (bool-to-val False)]  $\neq$  UndefVal
  shows val[x  $\oplus$  (bool-to-val False)] = x
  using assms apply (cases x; auto)
  by meson

lemma exp-xor-self-is-false:
  assumes wf-stamp x  $\wedge$  stamp-expr x = default-stamp
  shows exp[x  $\oplus$  x]  $\geq$  exp[false]
  using assms apply auto unfolding wf-stamp-def
  by (smt (z3) validDefIntConst IntVal0 Value.inject(1) bool-to-val.simps(2)
    constantAsStamp.simps(1) evalDet int-signed-value-bounds new-int.simps un-
    fold-const
    val-xor-self-is-false-2 valid-int valid-stamp.simps(1) valid-value.simps(1) wf-value-def)

lemma exp-eliminate-redundant-false:
  shows exp[x  $\oplus$  false]  $\geq$  exp[x]
  using val-eliminate-redundant-false apply auto
  subgoal premises p for m p xa
  proof –
    obtain xa where xa: [m, p]  $\vdash$  x  $\mapsto$  xa
    using p(2) by blast
    then have val[xa  $\oplus$  (IntVal 32 0)]  $\neq$  UndefVal
    using evalDet p(2) p(3) by blast
    then have [m, p]  $\vdash$  x  $\mapsto$  val[xa  $\oplus$  (IntVal 32 0)]
    apply (cases xa; auto) using eval-unused-bits-zero xa by auto
    then show ?thesis
    using evalDet p(2) xa by blast
  qed
done

Optimisations

optimization XorSelfIsFalse: (x  $\oplus$  x)  $\mapsto$  false when
  (wf-stamp x  $\wedge$  stamp-expr x = default-stamp)
  using size-non-const apply force
  using exp-xor-self-is-false by auto

optimization XorShiftConstantRight: ((const x)  $\oplus$  y)  $\mapsto$  y  $\oplus$  (const x) when
 $\neg$ (is-ConstantExpr y)
  using size-flip-binary apply force
  unfolding le-expr-def using val-xor-commute
  by auto

optimization EliminateRedundantFalse: (x  $\oplus$  false)  $\mapsto$  x

```

using *exp-eliminate-redundant-false* **by** *blast*

end

end

12 Conditional Elimination Phase

theory *ConditionalElimination*

imports

Semantics.IRTreeEvalThms

Proofs.Rewrites

Proofs.Bisimulation

begin

12.1 Individual Elimination Rules

The set of rules used for determining whether a condition $q1::'a$ implies another condition $q2::'a$ or its negation. These rules are used for conditional elimination.

inductive *impliesx* :: *IRExpr* \Rightarrow *IRExpr* \Rightarrow *bool* ($- \Rightarrow -$) **and**

impliesnot :: *IRExpr* \Rightarrow *IRExpr* \Rightarrow *bool* ($- \Rightarrow \neg -$) **where**

q-imp-q:

$q \Rightarrow q$ |

eq-impliesnot-less:

$(\text{BinaryExpr BinIntegerEquals } x \ y) \Rightarrow \neg (\text{BinaryExpr BinIntegerLessThan } x \ y)$ |

eq-impliesnot-less-rev:

$(\text{BinaryExpr BinIntegerEquals } x \ y) \Rightarrow \neg (\text{BinaryExpr BinIntegerLessThan } y \ x)$ |

less-impliesnot-rev-less:

$(\text{BinaryExpr BinIntegerLessThan } x \ y) \Rightarrow \neg (\text{BinaryExpr BinIntegerLessThan } y \ x)$

|

less-impliesnot-eq:

$(\text{BinaryExpr BinIntegerLessThan } x \ y) \Rightarrow \neg (\text{BinaryExpr BinIntegerEquals } x \ y)$ |

less-impliesnot-eq-rev:

$(\text{BinaryExpr BinIntegerLessThan } x \ y) \Rightarrow \neg (\text{BinaryExpr BinIntegerEquals } y \ x)$ |

negate-true:

$\llbracket x \Rightarrow \neg y \rrbracket \Longrightarrow x \Rightarrow (\text{UnaryExpr UnaryLogicNegation } y)$ |

negate-false:

$\llbracket x \Rightarrow y \rrbracket \Longrightarrow x \Rightarrow \neg (\text{UnaryExpr UnaryLogicNegation } y)$

The relation $q1::IRExpr \Rightarrow q2::IRExpr$ indicates that the implication $(q1::bool) \longrightarrow (q2::bool)$ is known true (i.e. universally valid), and the relation $q1::IRExpr \Rightarrow \neg q2::IRExpr$ indicates that the implication $(q1::bool) \longrightarrow (q2::bool)$ is known false (i.e. $(q1::bool) \longrightarrow \neg (q2::bool)$ is universally valid). If neither

$q1::IRExpr \Rightarrow q2::IRExpr$ nor $q1::IRExpr \Rightarrow \neg q2::IRExpr$ then the status is unknown. Only the known true and known false cases can be used for conditional elimination.

```
fun implies-valid :: IRExpr  $\Rightarrow$  IRExpr  $\Rightarrow$  bool (infix  $\rightarrow$  50) where
  implies-valid q1 q2 =
    ( $\forall m\ p\ v1\ v2. ([m, p] \vdash q1 \mapsto v1) \wedge ([m, p] \vdash q2 \mapsto v2) \longrightarrow$ 
      (val-to-bool v1  $\longrightarrow$  val-to-bool v2))
```

```
fun impliesnot-valid :: IRExpr  $\Rightarrow$  IRExpr  $\Rightarrow$  bool (infix  $\nrightarrow$  50) where
  impliesnot-valid q1 q2 =
    ( $\forall m\ p\ v1\ v2. ([m, p] \vdash q1 \mapsto v1) \wedge ([m, p] \vdash q2 \mapsto v2) \longrightarrow$ 
      (val-to-bool v1  $\longrightarrow$   $\neg$ val-to-bool v2))
```

The relation $(q1::IRExpr) \rightarrow (q2::IRExpr)$ means $(q1::bool) \longrightarrow (q2::bool)$ is universally valid, and the relation $(q1::IRExpr) \nrightarrow (q2::IRExpr)$ means $(q1::bool) \longrightarrow \neg (q2::bool)$ is universally valid.

```
lemma eq-impliesnot-less-helper:
   $v1 = v2 \longrightarrow \neg(int\text{-}signed\text{-}value\ b\ v1 < int\text{-}signed\text{-}value\ b\ v2)$ 
by force
```

```
lemma eq-impliesnot-less-val:
   $val\text{-}to\text{-}bool(intval\text{-}equals\ v1\ v2) \longrightarrow \neg val\text{-}to\text{-}bool(intval\text{-}less\text{-}than\ v1\ v2)$ 
using eq-impliesnot-less-helper bool-to-val.simps bool-to-val-bin.simps intval-equals.simps
  intval-less-than.elims val-to-bool.elims val-to-bool.simps
by (smt (verit))
```

```
lemma eq-impliesnot-less-rev-val:
   $val\text{-}to\text{-}bool(intval\text{-}equals\ v1\ v2) \longrightarrow \neg val\text{-}to\text{-}bool(intval\text{-}less\text{-}than\ v2\ v1)$ 
proof –
  have a:  $intval\text{-}equals\ v1\ v2 = intval\text{-}equals\ v2\ v1$ 
    using bool-to-val-bin.simps intval-equals.simps intval-equals.elims
    by (smt (verit))
  show ?thesis using a eq-impliesnot-less-val by presburger
qed
```

```
lemma less-impliesnot-rev-less-val:
   $val\text{-}to\text{-}bool(intval\text{-}less\text{-}than\ v1\ v2) \longrightarrow \neg val\text{-}to\text{-}bool(intval\text{-}less\text{-}than\ v2\ v1)$ 
by (smt (verit, del-insts) Value.exhaust Value.inject(1) bool-to-val.simps(2)
  bool-to-val-bin.simps intval-less-than.simps(1) intval-less-than.simps(5)
  intval-less-than.simps(6) intval-less-than.simps(7) val-to-bool.elims(2))
```

```
lemma less-impliesnot-eq-val:
   $val\text{-}to\text{-}bool(intval\text{-}less\text{-}than\ v1\ v2) \longrightarrow \neg val\text{-}to\text{-}bool(intval\text{-}equals\ v1\ v2)$ 
using eq-impliesnot-less-val by blast
```

```
lemma logic-negate-type:
  assumes  $[m, p] \vdash UnaryExpr\ UnaryLogicNegation\ x \mapsto v$ 
  shows  $\exists b\ v2. [m, p] \vdash x \mapsto IntVal\ b\ v2$ 
```

```

using assms
by (metis UnaryExprE intval-logic-negation.elims unary-eval.simps(4))

lemma intval-logic-negation-inverse:
  assumes  $b > 0$ 
  assumes  $x = \text{IntVal } b \ v$ 
  shows  $\text{val-to-bool } (\text{intval-logic-negation } x) \longleftrightarrow \neg(\text{val-to-bool } x)$ 
  using assms by (cases  $x$ ; auto simp: logic-negate-def)

lemma logic-negation-relation-tree:
  assumes  $[m, p] \vdash y \mapsto \text{val}$ 
  assumes  $[m, p] \vdash \text{UnaryExpr } \text{UnaryLogicNegation } y \mapsto \text{invval}$ 
  shows  $\text{val-to-bool } \text{val} \longleftrightarrow \neg(\text{val-to-bool } \text{invval})$ 
  using assms using intval-logic-negation-inverse
  by (metis UnaryExprE evalDet eval-bits-1-64 logic-negate-type unary-eval.simps(4))

The following theorem shows that the known true/false rules are valid.

theorem implies-impliesnot-valid:
  shows  $((q1 \Rightarrow q2) \longrightarrow (q1 \mapsto q2)) \wedge$ 
     $((q1 \Rightarrow \neg q2) \longrightarrow (q1 \mapsto \neg q2))$ 
     $(\text{is } (?imp \longrightarrow ?val) \wedge (?notimp \longrightarrow ?notval))$ 
proof (induct  $q1 \ q2$  rule: impliesx-impliesnot.induct)
  case ( $q\text{-imp-}q \ q$ )
  then show  $?case$ 
    using evalDet by fastforce
next
  case ( $eq\text{-impliesnot-less } x \ y$ )
  then show  $?case$  apply auto using eq-impliesnot-less-val evalDet by blast
next
  case ( $eq\text{-impliesnot-less-rev } x \ y$ )
  then show  $?case$  apply auto using eq-impliesnot-less-rev-val evalDet by blast
next
  case ( $less\text{-impliesnot-rev-less } x \ y$ )
  then show  $?case$  apply auto using less-impliesnot-rev-less-val evalDet by blast
next
  case ( $less\text{-impliesnot-eq } x \ y$ )
  then show  $?case$  apply auto using less-impliesnot-eq-val evalDet by blast
next
  case ( $less\text{-impliesnot-eq-rev } x \ y$ )
  then show  $?case$  apply auto using eq-impliesnot-less-rev-val evalDet by metis
next
  case ( $negate\text{-true } x \ y$ )
  then show  $?case$  apply auto
    by (metis logic-negation-relation-tree unary-eval.simps(4) unfold-unary)
next
  case ( $negate\text{-false } x \ y$ )
  then show  $?case$  apply auto
    by (metis UnaryExpr logic-negation-relation-tree unary-eval.simps(4))
qed

```

We introduce a type *TriState::'a* (as in the GraalVM compiler) to represent when static analysis can tell us information about the value of a Boolean expression. If *Unknown::'a* then no information can be inferred and if *KnownTrue::'a*/*KnownFalse::'a* one can infer the expression is always true/false.

datatype *TriState* = *Unknown* | *KnownTrue* | *KnownFalse*

The implies relation corresponds to the *LogicNode.implies* method from the compiler which attempts to infer when one logic nodes value can be inferred from a known logic node.

inductive *implies* :: *IRGraph* \Rightarrow *IRNode* \Rightarrow *IRNode* \Rightarrow *TriState* \Rightarrow *bool*

(- \vdash - & - \hookrightarrow -) **for** *g* **where**

eq-imp-less:

g \vdash (*IntegerEqualsNode* *x y*) & (*IntegerLessThanNode* *x y*) \hookrightarrow *KnownFalse* |

eq-imp-less-rev:

g \vdash (*IntegerEqualsNode* *x y*) & (*IntegerLessThanNode* *y x*) \hookrightarrow *KnownFalse* |

less-imp-rev-less:

g \vdash (*IntegerLessThanNode* *x y*) & (*IntegerLessThanNode* *y x*) \hookrightarrow *KnownFalse* |

less-imp-not-eq:

g \vdash (*IntegerLessThanNode* *x y*) & (*IntegerEqualsNode* *x y*) \hookrightarrow *KnownFalse* |

less-imp-not-eq-rev:

g \vdash (*IntegerLessThanNode* *x y*) & (*IntegerEqualsNode* *y x*) \hookrightarrow *KnownFalse* |

x-imp-x:

g \vdash *x* & *x* \hookrightarrow *KnownTrue* |

negate-false:

$\llbracket g \vdash x \text{ \& } (kind\ g\ y) \hookrightarrow KnownTrue \rrbracket \implies g \vdash x \text{ \& } (LogicNegationNode\ y) \hookrightarrow KnownFalse$ |

negate-true:

$\llbracket g \vdash x \text{ \& } (kind\ g\ y) \hookrightarrow KnownFalse \rrbracket \implies g \vdash x \text{ \& } (LogicNegationNode\ y) \hookrightarrow KnownTrue$

Total relation over partial implies relation

inductive *condition-implies* :: *IRGraph* \Rightarrow *IRNode* \Rightarrow *IRNode* \Rightarrow *TriState* \Rightarrow *bool*

(- \vdash - & - \rightarrow -) **for** *g* **where**

$\llbracket \neg(g \vdash a \text{ \& } b \hookrightarrow imp) \rrbracket \implies (g \vdash a \text{ \& } b \rightarrow Unknown)$ |

$\llbracket (g \vdash a \text{ \& } b \hookrightarrow imp) \rrbracket \implies (g \vdash a \text{ \& } b \rightarrow imp)$

inductive *implies-tree* :: *IRExpr* \Rightarrow *IRExpr* \Rightarrow *bool* \Rightarrow *bool*

(- & - \hookrightarrow -) **where**

eq-imp-less:

(*BinaryExpr* *BinIntegerEquals* *x y*) & (*BinaryExpr* *BinIntegerLessThan* *x y*) \hookrightarrow *False* |

eq-imp-less-rev:

(*BinaryExpr* *BinIntegerEquals* *x y*) & (*BinaryExpr* *BinIntegerLessThan* *y x*) \hookrightarrow *False* |

less-imp-rev-less:

```

  (BinaryExpr BinIntegerLessThan x y) & (BinaryExpr BinIntegerLessThan y x)
  ↪ False |
  less-imp-not-eq:
  (BinaryExpr BinIntegerLessThan x y) & (BinaryExpr BinIntegerEquals x y) ↪
  False |
  less-imp-not-eq-rev:
  (BinaryExpr BinIntegerLessThan x y) & (BinaryExpr BinIntegerEquals y x) ↪
  False |
  x-imp-x:
  x & x ↪ True |
  negate-false:
   $\llbracket x \& y \hookrightarrow \text{True} \rrbracket \implies x \& (\text{UnaryExpr UnaryLogicNegation } y) \hookrightarrow \text{False} \mid$ 
  negate-true:
   $\llbracket x \& y \hookrightarrow \text{False} \rrbracket \implies x \& (\text{UnaryExpr UnaryLogicNegation } y) \hookrightarrow \text{True}$ 

```

Proofs that the implies relation is correct with respect to the existing evaluation semantics.

lemma *logic-negation-relation*:

```

  assumes [g, m, p] ⊢ y ↦ val
  assumes kind g neg = LogicNegationNode y
  assumes [g, m, p] ⊢ neg ↦ invval
  assumes invval ≠ UndefVal
  shows val-to-bool val ⟷ ¬(val-to-bool invval)
  using assms
  by (metis LogicNegationNode encodeeval-def logic-negation-relation-tree repDet)

```

lemma *implies-valid*:

```

  assumes x & y ↪ imp
  assumes [m, p] ⊢ x ↦ v1
  assumes [m, p] ⊢ y ↦ v2
  shows (imp ⟶ (val-to-bool v1 ⟶ val-to-bool v2)) ∧
        (¬imp ⟶ (val-to-bool v1 ⟶ ¬(val-to-bool v2)))
  (is (?TP ⟶ ?TC) ∧ (?FP ⟶ ?FC))
  apply (intro conjI; rule impI)
proof -
  assume KnownTrue: ?TP
  show ?TC
  using assms(1) KnownTrue assms(2-) proof (induct x y imp rule: implies-tree.induct)
  case (eq-imp-less x y)
  then show ?case by simp
  next
  case (eq-imp-less-rev x y)
  then show ?case by simp
  next
  case (less-imp-rev-less x y)
  then show ?case by simp
  next
  case (less-imp-not-eq x y)

```



```

    then show ?case by simp
next
  case (less-imp-not-eq-rev x y)
  then show ?case by simp
next
  case (x-imp-x)
  then show ?case
    by (metis evalDet)
next
  case (negate-false x1)
  then show ?case using evalDet
    using assms(2,3) by blast
next
  case (negate-true x y)
  then show ?case
    using logic-negation-relation-tree sorry
qed
next
  assume KnownFalse: ?FP
  show ?FC using assms KnownFalse proof (induct x y imp rule: implies-tree.induct)
    case (eq-imp-less x y)
    obtain xval where xval: [m, p] ⊢ x ↦ xval
      using eq-imp-less(1) eq-imp-less.prem(3)
      by blast
    then obtain yval where yval: [m, p] ⊢ y ↦ yval
      using eq-imp-less.prem(3)
      using eq-imp-less.prem(2) by blast
    have egeval: [m, p] ⊢ (BinaryExpr BinIntegerEquals x y) ↦ intval-equals xval
      yval
      using xval yval evaltree.BinaryExpr
      by (metis BinaryExprE bin-eval.simp(11) eq-imp-less.prem(1) evalDet)
    have lesseval: [m, p] ⊢ (BinaryExpr BinIntegerLessThan x y) ↦ intval-less-than
      xval yval
      using xval yval evaltree.BinaryExpr
      by (metis BinaryExprE bin-eval.simp(12) eq-imp-less.prem(2) evalDet)
    have val-to-bool (intval-equals xval yval) ⟶ ¬(val-to-bool (intval-less-than xval
      yval))
      apply (cases xval; cases yval; auto)
      by (smt (verit, best) bool-to-val.simp(2) val-to-bool.simp(1))
    then show ?case
      using egeval lesseval
      by (metis eq-imp-less.prem(1) eq-imp-less.prem(2) evalDet)
  next
    case (eq-imp-less-rev x y)
    obtain xval where xval: [m, p] ⊢ x ↦ xval
      using eq-imp-less-rev.prem(3)
      using eq-imp-less-rev.prem(2) by blast
    obtain yval where yval: [m, p] ⊢ y ↦ yval
      using eq-imp-less-rev.prem(3)

```

```

    using eq-imp-less-rev.premis(2) by blast
    have egeval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerEquals } x \ y) \mapsto \text{intval-equals } xval$ 
    yval
    using xval yval evaltree.BinaryExpr
    by (metis BinaryExprE bin-eval.simps(11) eq-imp-less-rev.premis(1) evalDet)
    have lesseval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerLessThan } y \ x) \mapsto \text{intval-less-than}$ 
    yval xval
    using xval yval evaltree.BinaryExpr
    by (metis BinaryExprE bin-eval.simps(12) eq-imp-less-rev.premis(2) evalDet)
    have val-to-bool (intval-equals xval yval)  $\longrightarrow \neg(\text{val-to-bool (intval-less-than yval xval)})$ 
    apply (cases xval; cases yval; auto)
    by (metis (full-types) bool-to-val.simps(2) less-irrefl val-to-bool.simps(1))
    then show ?case
    using egeval lesseval
    by (metis eq-imp-less-rev.premis(1) eq-imp-less-rev.premis(2) evalDet)
  next
  case (less-imp-rev-less x y)
  obtain xval where xval:  $[m, p] \vdash x \mapsto xval$ 
  using less-imp-rev-less.premis(3)
  using less-imp-rev-less.premis(2) by blast
  obtain yval where yval:  $[m, p] \vdash y \mapsto yval$ 
  using less-imp-rev-less.premis(3)
  using less-imp-rev-less.premis(2) by blast
  have lesseval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerLessThan } x \ y) \mapsto \text{intval-less-than}$ 
  xval yval
  using xval yval evaltree.BinaryExpr
  by (metis BinaryExprE bin-eval.simps(12) evalDet less-imp-rev-less.premis(1))
  have revlesseval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerLessThan } y \ x) \mapsto \text{int-}$ 
  val-less-than yval xval
  using xval yval evaltree.BinaryExpr
  by (metis BinaryExprE bin-eval.simps(12) evalDet less-imp-rev-less.premis(2))
  have val-to-bool (intval-less-than xval yval)  $\longrightarrow \neg(\text{val-to-bool (intval-less-than yval xval)})$ 
  apply (cases xval; cases yval; auto)
  by (smt (verit) bool-to-val.simps(2) val-to-bool.simps(1))
  then show ?case
  by (metis evalDet less-imp-rev-less.premis(1) less-imp-rev-less.premis(2) lesseval revlesseval)
  next
  case (less-imp-not-eq x y)
  obtain xval where xval:  $[m, p] \vdash x \mapsto xval$ 
  using less-imp-not-eq.premis(3)
  using less-imp-not-eq.premis(1) by blast
  obtain yval where yval:  $[m, p] \vdash y \mapsto yval$ 
  using less-imp-not-eq.premis(3)
  using less-imp-not-eq.premis(1) by blast
  have egeval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerEquals } x \ y) \mapsto \text{intval-equals } xval$ 
  yval

```

```

    using xval yval evaltree.BinaryExpr
    by (metis BinaryExprE bin-eval.simps(11) evalDet less-imp-not-eq.prem(2))
    have lesseval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerLessThan } x \ y) \mapsto \text{intval-less-than}$ 
xval yval
    using xval yval evaltree.BinaryExpr
    by (metis BinaryExprE bin-eval.simps(12) evalDet less-imp-not-eq.prem(1))
    have val-to-bool (intval-less-than xval yval)  $\longrightarrow \neg(\text{val-to-bool (intval-equals xval yval)})$ 
    apply (cases xval; cases yval; auto)
    by (smt (verit, best) bool-to-val.simps(2) val-to-bool.simps(1))
    then show ?case
    by (metis egeval evalDet less-imp-not-eq.prem(1) less-imp-not-eq.prem(2)
lesseval)
  next
  case (less-imp-not-eq-rev x y)
  obtain xval where xval:  $[m, p] \vdash x \mapsto \text{xval}$ 
  using less-imp-not-eq-rev.prem(3)
  using less-imp-not-eq-rev.prem(1) by blast
  obtain yval where yval:  $[m, p] \vdash y \mapsto \text{yval}$ 
  using less-imp-not-eq-rev.prem(3)
  using less-imp-not-eq-rev.prem(1) by blast
  have egeval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerEquals } y \ x) \mapsto \text{intval-equals yval xval}$ 
  using xval yval evaltree.BinaryExpr
  by (metis BinaryExprE bin-eval.simps(11) evalDet less-imp-not-eq-rev.prem(2))
  have lesseval:  $[m, p] \vdash (\text{BinaryExpr BinIntegerLessThan } x \ y) \mapsto \text{intval-less-than}$ 
xval yval
  using xval yval evaltree.BinaryExpr
  by (metis BinaryExprE bin-eval.simps(12) evalDet less-imp-not-eq-rev.prem(1))
  have val-to-bool (intval-less-than xval yval)  $\longrightarrow \neg(\text{val-to-bool (intval-equals yval xval)})$ 
  apply (cases xval; cases yval; auto)
  by (smt (verit, best) bool-to-val.simps(2) val-to-bool.simps(1))
  then show ?case
  by (metis egeval evalDet less-imp-not-eq-rev.prem(1) less-imp-not-eq-rev.prem(2)
lesseval)
  next
  case (x-imp-x x1)
  then show ?case by simp
  next
  case (negate-false x y)
  then show ?case sorry
  next
  case (negate-true x1)
  then show ?case by simp
qed
qed

```

lemma *implies-true-valid*:

```

assumes  $x \ \& \ y \hookrightarrow imp$ 
assumes  $imp$ 
assumes  $[m, p] \vdash x \mapsto v1$ 
assumes  $[m, p] \vdash y \mapsto v2$ 
shows  $val\text{-}to\text{-}bool \ v1 \longrightarrow val\text{-}to\text{-}bool \ v2$ 
using  $assms \ implies\text{-}valid$ 
by  $blast$ 

```

```

lemma  $implies\text{-}false\text{-}valid$ :
assumes  $x \ \& \ y \hookrightarrow imp$ 
assumes  $\neg imp$ 
assumes  $[m, p] \vdash x \mapsto v1$ 
assumes  $[m, p] \vdash y \mapsto v2$ 
shows  $val\text{-}to\text{-}bool \ v1 \longrightarrow \neg(val\text{-}to\text{-}bool \ v2)$ 
using  $assms \ implies\text{-}valid \ \mathbf{by} \ blast$ 

```

The following relation corresponds to the `UnaryOpLogicNode.tryFold` and `BinaryOpLogicNode.tryFold` methods and their associated concrete implementations.

The relation determines if a logic operation can be shown true or false through the stamp typing information.

```

inductive  $tryFold :: IRNode \Rightarrow (ID \Rightarrow Stamp) \Rightarrow bool \Rightarrow bool$ 
where
   $\llbracket alwaysDistinct \ (stamps \ x) \ (stamps \ y) \rrbracket$ 
     $\Longrightarrow tryFold \ (IntegerEqualsNode \ x \ y) \ stamps \ False \mid$ 
   $\llbracket neverDistinct \ (stamps \ x) \ (stamps \ y) \rrbracket$ 
     $\Longrightarrow tryFold \ (IntegerEqualsNode \ x \ y) \ stamps \ True \mid$ 
   $\llbracket is\text{-}IntegerStamp \ (stamps \ x);$ 
     $is\text{-}IntegerStamp \ (stamps \ y);$ 
     $stpi\text{-}upper \ (stamps \ x) < stpi\text{-}lower \ (stamps \ y) \rrbracket$ 
     $\Longrightarrow tryFold \ (IntegerLessThanNode \ x \ y) \ stamps \ True \mid$ 
   $\llbracket is\text{-}IntegerStamp \ (stamps \ x);$ 
     $is\text{-}IntegerStamp \ (stamps \ y);$ 
     $stpi\text{-}lower \ (stamps \ x) \geq stpi\text{-}upper \ (stamps \ y) \rrbracket$ 
     $\Longrightarrow tryFold \ (IntegerLessThanNode \ x \ y) \ stamps \ False$ 

```

Proofs that show that when the stamp lookup function is well-formed, the `tryFold` relation correctly predicts the output value with respect to our evaluation semantics.

```

lemma
assumes  $kind \ g \ nid = IntegerEqualsNode \ x \ y$ 
assumes  $[g, m, p] \vdash nid \mapsto v$ 
assumes  $([g, m, p] \vdash x \mapsto xval) \wedge ([g, m, p] \vdash y \mapsto yval)$ 
shows  $val\text{-}to\text{-}bool \ (intval\text{-}equals \ xval \ yval) \longleftrightarrow v = IntVal \ 32 \ 1$ 
proof –
have  $v = intval\text{-}equals \ xval \ yval$ 
using  $assms(1, 2, 3) \ BinaryExprE \ IntegerEqualsNode \ bin\text{-}eval.simps(7)$ 
by  $(smt \ (verit) \ bin\text{-}eval.simps(11) \ encodeeval\text{-}def \ evalDet \ repDet)$ 

```

```

then show ?thesis using intval-equals.simps val-to-bool.simps
by (smt (verit) bool-to-val.simps(1) bool-to-val.simps(2) bool-to-val-bin.simps
    intval-equals.elims one-neq-zero)
qed

lemma tryFoldIntegerEqualsAlwaysDistinct:
  assumes wf-stamp g stamps
  assumes kind g nid = (IntegerEqualsNode x y)
  assumes [g, m, p] ⊢ nid ↦ v
  assumes alwaysDistinct (stamps x) (stamps y)
  shows v = IntVal 32 0
proof –
  have ∀ val. ¬(valid-value val (join (stamps x) (stamps y)))
    using assms(1,4) unfolding alwaysDistinct.simps
    by (smt (verit, best) is-stamp-empty.elims(2) valid-int valid-value.simps(1))
  obtain xv where [g, m, p] ⊢ x ↦ xv
    using assms using assms(2,3) unfolding encodeeval-def sorry
  have ¬(∃ val . ([g, m, p] ⊢ x ↦ val) ∧ ([g, m, p] ⊢ y ↦ val))
    using assms(1,4) unfolding alwaysDistinct.simps wf-stamp.simps encodee-
    val-def sorry
  then show ?thesis sorry
qed

lemma tryFoldIntegerEqualsNeverDistinct:
  assumes wf-stamp g stamps
  assumes kind g nid = (IntegerEqualsNode x y)
  assumes [g, m, p] ⊢ nid ↦ v
  assumes neverDistinct (stamps x) (stamps y)
  shows v = IntVal 32 1
  using assms IntegerEqualsNodeE sorry

lemma tryFoldIntegerLessThanTrue:
  assumes wf-stamp g stamps
  assumes kind g nid = (IntegerLessThanNode x y)
  assumes [g, m, p] ⊢ nid ↦ v
  assumes stpi-upper (stamps x) < stpi-lower (stamps y)
  shows v = IntVal 32 1
proof –
  have stamp-type: is-IntegerStamp (stamps x)
    using assms
    sorry
  obtain xval where xval: [g, m, p] ⊢ x ↦ xval
    using assms(2,3) sorry
  obtain yval where yval: [g, m, p] ⊢ y ↦ yval
    using assms(2,3) sorry
  have is-IntegerStamp (stamps x) ∧ is-IntegerStamp (stamps y)
    using assms(4)
    sorry
  then have val-to-bool (intval-less-than xval yval)

```

```

    sorry
  then show ?thesis
    sorry
qed

```

```

lemma tryFoldIntegerLessThanFalse:
  assumes wf-stamp g stamps
  assumes kind g nid = (IntegerLessThanNode x y)
  assumes [g, m, p] ⊢ nid ↦ v
  assumes stpi-lower (stamps x) ≥ stpi-upper (stamps y)
  shows v = IntVal 32 0
  proof -
    have stamp-type: is-IntegerStamp (stamps x)
      using assms
    sorry
    obtain xval where xval: [g, m, p] ⊢ x ↦ xval
      using assms(2,3) sorry
    obtain yval where yval: [g, m, p] ⊢ y ↦ yval
      using assms(2,3) sorry
    have is-IntegerStamp (stamps x) ∧ is-IntegerStamp (stamps y)
      using assms(4)
    sorry
    then have ¬(val-to-bool (intval-less-than xval yval))
      sorry
    then show ?thesis
      sorry
  qed

```

```

theorem tryFoldProofTrue:
  assumes wf-stamp g stamps
  assumes tryFold (kind g nid) stamps True
  assumes [g, m, p] ⊢ nid ↦ v
  shows val-to-bool v
  using assms(2) proof (induction kind g nid stamps True rule: tryFold.induct)
  case (1 stamps x y)
    then show ?case using tryFoldIntegerEqualsAlwaysDistinct assms
      by force
  next
    case (2 stamps x y)
    then show ?case using tryFoldIntegerEqualsAlwaysDistinct assms
      by (smt (verit, best) one-neq-zero tryFold.cases tryFoldIntegerEqualsNeverDis-
        tinct tryFoldIntegerLessThanTrue val-to-bool.simps(1))
  next
    case (3 stamps x y)
    then show ?case using tryFoldIntegerLessThanTrue assms
      by (smt (verit, best) one-neq-zero tryFold.cases tryFoldIntegerEqualsNeverDis-
        tinct val-to-bool.simps(1))
  next
    case (4 stamps x y)

```

```

    then show ?case using tryFoldIntegerLessThanFalse assms sorry
qed

theorem tryFoldProofFalse:
  assumes wf-stamp g stamps
  assumes tryFold (kind g nid) stamps False
  assumes  $[g, m, p] \vdash \text{nid} \mapsto v$ 
  shows  $\neg(\text{val-to-bool } v)$ 
using assms(2) proof (induction kind g nid stamps False rule: tryFold.induct)
case (1 stamps x y)
  then show ?case using tryFoldIntegerEqualsAlwaysDistinct assms sorry
next
case (2 stamps x y)
  then show ?case using tryFoldIntegerEqualsNeverDistinct assms sorry
next
case (3 stamps x y)
  then show ?case using tryFoldIntegerLessThanTrue assms sorry
next
case (4 stamps x y)
  then show ?case using tryFoldIntegerLessThanFalse assms sorry

qed

```

inductive-cases *StepE*:

$$g, p \vdash (\text{nid}, m, h) \rightarrow (\text{nid}', m', h)$$

Perform conditional elimination rewrites on the graph for a particular node. In order to determine conditional eliminations appropriately the rule needs two data structures produced by static analysis. The first parameter is the set of IRNodes that we know result in a true value when evaluated. The second parameter is a mapping from node identifiers to the flow-sensitive stamp.

The relation transforms the third parameter to the fifth parameter for a node identifier which represents the fourth parameter.

inductive *ConditionalEliminationStep* ::
 $\text{IRExpr set} \Rightarrow (\text{ID} \Rightarrow \text{Stamp}) \Rightarrow \text{IRGraph} \Rightarrow \text{ID} \Rightarrow \text{IRGraph} \Rightarrow \text{bool}$ **where**
impliesTrue:
 $\llbracket \text{kind } g \text{ ifcond} = (\text{IfNode } \text{cid } t \text{ } f);$
 $g \vdash \text{cid} \simeq \text{cond};$
 $\exists \text{ ce} \in \text{conds} . (\text{ce} \Rightarrow \text{cond});$
 $g' = \text{constantCondition True ifcond (kind } g \text{ ifcond) } g$
 $\rrbracket \Rightarrow \text{ConditionalEliminationStep conds stamps } g \text{ ifcond } g' \mid$
impliesFalse:
 $\llbracket \text{kind } g \text{ ifcond} = (\text{IfNode } \text{cid } t \text{ } f);$
 $g \vdash \text{cid} \simeq \text{cond};$

```

 $\exists$   $ce \in \text{conds} . (ce \Rightarrow \neg \text{cond});$ 
 $g' = \text{constantCondition False ifcond (kind g ifcond) g}$ 
 $\parallel \Rightarrow \text{ConditionalEliminationStep conds stamps g ifcond g'}$ 

```

```

tryFoldTrue:
 $\parallel \text{kind g ifcond} = (\text{IfNode cid t f});$ 
 $\text{cond} = \text{kind g cid};$ 
 $\text{tryFold (kind g cid) stamps True};$ 
 $g' = \text{constantCondition True ifcond (kind g ifcond) g}$ 
 $\parallel \Rightarrow \text{ConditionalEliminationStep conds stamps g ifcond g'}$ 

```

```

tryFoldFalse:
 $\parallel \text{kind g ifcond} = (\text{IfNode cid t f});$ 
 $\text{cond} = \text{kind g cid};$ 
 $\text{tryFold (kind g cid) stamps False};$ 
 $g' = \text{constantCondition False ifcond (kind g ifcond) g}$ 
 $\parallel \Rightarrow \text{ConditionalEliminationStep conds stamps g ifcond g'}$ 

```

code-pred (*modes*: $i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow \text{bool}$) *ConditionalEliminationStep* .

thm *ConditionalEliminationStep.equation*

12.2 Control-flow Graph Traversal

```

type-synonym Seen = ID set
type-synonym Condition = IRExpr
type-synonym Conditions = Condition list
type-synonym StampFlow = (ID  $\Rightarrow$  Stamp) list

```

`nextEdge` helps determine which node to traverse next by returning the first successor edge that isn't in the set of already visited nodes. If there is not an appropriate successor, `None` is returned instead.

```

fun nextEdge :: Seen  $\Rightarrow$  ID  $\Rightarrow$  IRGraph  $\Rightarrow$  ID option where
  nextEdge seen nid g =
    (let nids = (filter ( $\lambda \text{nid}' . \text{nid}' \notin \text{seen}$ ) (successors-of (kind g nid)))) in
    (if length nids > 0 then Some (hd nids) else None)

```

`pred` determines which node, if any, acts as the predecessor of another.

Merge nodes represent a special case where-in the predecessor exists as an input edge of the merge node, to simplify the traversal we treat only the first input end node as the predecessor, ignoring that multiple nodes may act as a successor.

For all other nodes, the predecessor is the first element of the predecessors set. Note that in a well-formed graph there should only be one element in the predecessor set.

```

fun pred :: IRGraph  $\Rightarrow$  ID  $\Rightarrow$  ID option where
  pred g nid = (case kind g nid of

```



```

    (MergeNode ends -) ⇒ Some (hd ends) |
  - ⇒
    (if IRGraph.predecessors g nid = {}
     then None else
     Some (hd (sorted-list-of-set (IRGraph.predecessors g nid)))
    )
)

```

When the basic block of an if statement is entered, we know that the condition of the preceding if statement must be true. As in the GraalVM compiler, we introduce the registerNewCondition function which roughly corresponds to the ConditionalEliminationPhase.registerNewCondition. This method updates the flow-sensitive stamp information based on the condition which we know must be true.

```

fun clip-upper :: Stamp ⇒ int ⇒ Stamp where
  clip-upper (IntegerStamp b l h) c = (IntegerStamp b l c) |
  clip-upper s c = s
fun clip-lower :: Stamp ⇒ int ⇒ Stamp where
  clip-lower (IntegerStamp b l h) c = (IntegerStamp b c h) |
  clip-lower s c = s

```

```

fun registerNewCondition :: IRGraph ⇒ IRNode ⇒ (ID ⇒ Stamp) ⇒ (ID ⇒ Stamp) where

```

```

  registerNewCondition g (IntegerEqualsNode x y) stamps =
    (stamps
     (x := join (stamps x) (stamps y)))
    (y := join (stamps x) (stamps y)) |

```

```

  registerNewCondition g (IntegerLessThanNode x y) stamps =
    (stamps
     (x := clip-upper (stamps x) (stpi-lower (stamps y))))
    (y := clip-lower (stamps y) (stpi-upper (stamps x))) |
  registerNewCondition g - stamps = stamps

```

```

fun hdOr :: 'a list ⇒ 'a ⇒ 'a where
  hdOr (x # xs) de = x |
  hdOr [] de = de

```

The Step relation is a small-step traversal of the graph which handles transitions between individual nodes of the graph.

It relates a pairs of tuple of the current node, the set of seen nodes, the always true stack of IfNode conditions, and the flow-sensitive stamp information.

inductive Step

```

  :: IRGraph ⇒ (ID × Seen × Conditions × StampFlow) ⇒ (ID × Seen ×
  Conditions × StampFlow) option ⇒ bool

```

for g **where**

— Hit a BeginNode with an IfNode predecessor which represents the start of a

basic block for the IfNode. 1. nid' will be the successor of the begin node. 2. Find the first and only predecessor. 3. Extract condition from the preceding IfNode. 4. Negate condition if the begin node is second branch (we've taken the else branch of the condition) 5. Add the condition or the negated condition to stack 6. Perform any stamp updates based on the condition using the `registerNewCondition` function and place them on the top of the stack of stamp information

$\llbracket kind\ g\ nid = BeginNode\ nid';$

$nid \notin seen;$
 $seen' = \{nid\} \cup seen;$

$Some\ ifcond = pred\ g\ nid;$
 $kind\ g\ ifcond = IfNode\ cond\ t\ f;$

$i = find-index\ nid\ (successors-of\ (kind\ g\ ifcond));$
 $c = (if\ i = 0\ then\ kind\ g\ cond\ else\ LogicNegationNode\ cond);$
 $rep\ g\ cond\ ce;$
 $ce' = (if\ i = 0\ then\ ce\ else\ UnaryExpr\ UnaryLogicNegation\ ce);$
 $conds' = ce' \# conds;$

$flow' = registerNewCondition\ g\ c\ (hdOr\ flow\ (stamp\ g));$
 $\implies Step\ g\ (nid,\ seen,\ conds,\ flow)\ (Some\ (nid',\ seen',\ conds',\ flow' \# flow)) \mid$

— Hit an EndNode 1. nid' will be the usage of EndNode 2. pop the conditions and stamp stack

$\llbracket kind\ g\ nid = EndNode;$

$nid \notin seen;$
 $seen' = \{nid\} \cup seen;$

$nid' = any-usage\ g\ nid;$

$conds' = tl\ conds;$
 $flow' = tl\ flow$
 $\implies Step\ g\ (nid,\ seen,\ conds,\ flow)\ (Some\ (nid',\ seen',\ conds',\ flow')) \mid$

— We can find a successor edge that is not in seen, go there

$\llbracket \neg(is-EndNode\ (kind\ g\ nid));$
 $\neg(is-BeginNode\ (kind\ g\ nid));$

$nid \notin seen;$
 $seen' = \{nid\} \cup seen;$

$Some\ nid' = nextEdge\ seen'\ nid\ g$
 $\implies Step\ g\ (nid,\ seen,\ conds,\ flow)\ (Some\ (nid',\ seen',\ conds,\ flow)) \mid$

— We cannot find a successor edge that is not in seen, give back None

$\llbracket \neg(is-EndNode\ (kind\ g\ nid));$
 $\neg(is-BeginNode\ (kind\ g\ nid));$

$nid \notin seen;$
 $seen' = \{nid\} \cup seen;$

$None = nextEdge\ seen'\ nid\ g$
 $\implies Step\ g\ (nid, seen, conds, flow)\ None \mid$

— We've already seen this node, give back None
 $\llbracket nid \in seen \rrbracket \implies Step\ g\ (nid, seen, conds, flow)\ None$

code-pred ($modes: i \Rightarrow i \Rightarrow o \Rightarrow bool$) $Step$.

The ConditionalEliminationPhase relation is responsible for combining the individual traversal steps from the Step relation and the optimizations from the ConditionalEliminationStep relation to perform a transformation of the whole graph.

inductive ConditionalEliminationPhase

$:: IRGraph \Rightarrow (ID \times Seen \times Conditions \times StampFlow) \Rightarrow IRGraph \Rightarrow bool$

where

— Can do a step and optimise for the current node
 $\llbracket Step\ g\ (nid, seen, conds, flow)\ (Some\ (nid', seen', conds', flow')) \rrbracket;$
 $ConditionalEliminationStep\ (set\ conds)\ (hdOr\ flow\ (stamp\ g))\ g\ nid\ g';$

$ConditionalEliminationPhase\ g'\ (nid', seen', conds', flow')\ g'$
 $\implies ConditionalEliminationPhase\ g\ (nid, seen, conds, flow)\ g'' \mid$

— Can do a step, matches whether optimised or not causing non-determinism We need to find a way to negate ConditionalEliminationStep
 $\llbracket Step\ g\ (nid, seen, conds, flow)\ (Some\ (nid', seen', conds', flow')) \rrbracket;$

$ConditionalEliminationPhase\ g\ (nid', seen', conds', flow')\ g'$
 $\implies ConditionalEliminationPhase\ g\ (nid, seen, conds, flow)\ g' \mid$

— Can't do a step but there is a predecessor we can backtrace to
 $\llbracket Step\ g\ (nid, seen, conds, flow)\ None \rrbracket;$
 $Some\ nid' = pred\ g\ nid;$
 $seen' = \{nid\} \cup seen;$
 $ConditionalEliminationPhase\ g\ (nid', seen', conds, flow)\ g'$
 $\implies ConditionalEliminationPhase\ g\ (nid, seen, conds, flow)\ g' \mid$

— Can't do a step and have no predecessors so terminate
 $\llbracket Step\ g\ (nid, seen, conds, flow)\ None \rrbracket;$
 $None = pred\ g\ nid$
 $\implies ConditionalEliminationPhase\ g\ (nid, seen, conds, flow)\ g$

code-pred ($modes: i \Rightarrow i \Rightarrow o \Rightarrow bool$) ConditionalEliminationPhase .

definition runConditionalElimination $:: IRGraph \Rightarrow IRGraph$ **where**

```

runConditionalElimination g =
  (Predicate.the (ConditionalEliminationPhase-i-i-o g (0, {}, ([], []))))

```

```

end

```