Veriopt Theories

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1 Canonicalization Phase	
theory Common	
imports	
$Optimization DSL. \ Canonicalization$	
HOL-Eisbach.Eisbach	
begin	
fun $size :: IRExpr \Rightarrow nat$ where	
$size (UnaryExpr \ op \ e) = (size \ e) + 1$	
size (BinaryExpr BinAdd x y) = (size x) + ((size y) * 2)	
$size (BinaryExpr \ op \ x \ y) = (size \ x) + (size \ y) \mid$	
$size (Conditional Expr \ cond \ t \ f) = (size \ cond) + (size \ t) + (size \ f) + 2$	
size (ConstantExpr c) = 1	
$size (ParameterExpr ind s) = 2 \mid$	
$size (LeafExpr \ nid \ s) = 2 \mid$	
$size (Constant Var c) = 2 \mid$	
$size (Variable Expr \ x \ s) = 2$	
$egin{array}{c} \mathbf{method} \ unfold\mbox{-}optimization = \end{array}$	
$(unfold\ rewrite-preservation.simps,\ unfold\ rewrite-termination.simps,$	
unfold intval.simps,	
rule conjE, simp, simp del: le-expr-def)	
(unfold rewrite-preservation.simps, unfold rewrite-termination.simps,	
rule conjE, simp, simp del: le-expr-def)	
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end	

1.1 Conditional Expression

 $\begin{array}{c} \textbf{theory} \ \textit{ConditionalPhase} \\ \textbf{imports} \end{array}$

```
Common
   Proofs. Stamp Eval Thms
begin
phase Conditional
 terminating size
begin
lemma negates: is-IntVal32 e \lor is-IntVal64 e \Longrightarrow val-to-bool (val[e]) \equiv \neg (val-to-bool
  by (smt\ (verit,\ best)\ Value.disc(1)\ Value.disc(10)\ Value.disc(4)\ Value.disc(5)
Value.disc(6)\ Value.disc(9)\ intval-logic-negation.elims\ val-to-bool.simps(1)\ val-to-bool.simps(2)
zero-neq-one)
optimization negate-condition: ((\neg e) ? x : y) \mapsto (e ? y : x)
   apply unfold-optimization apply simp using negates
  \textbf{using } \textit{Conditional ExprE } \textit{Unary ExprE } \textit{intval-logic-negation. elims } \textit{unary-eval. simps} (4)
val-to-bool.simps(1) val-to-bool.simps(2) zero-neq-one
   apply (smt (verit) ConditionalExpr)
   unfolding size.simps by simp
optimization const-true: (true ? x : y) \mapsto x
  apply unfold-optimization
  apply force
  unfolding size.simps by simp
optimization const-false: (false ? x : y) \mapsto y
  apply unfold-optimization
  apply force
  unfolding size.simps by simp
optimization equal-branches: (e ? x : x) \mapsto x
  apply unfold-optimization
  apply force
 unfolding size.simps by auto
definition wff-stamps :: bool where
 wff-stamps = (\forall m \ p \ expr \ val \ . ([m,p] \vdash expr \mapsto val) \longrightarrow valid-value val \ (stamp-expr
expr))
optimization condition-bounds-x: ((x < y) ? x : y) \mapsto x when (stamp-under
(stamp-expr\ x)\ (stamp-expr\ y)\ \land\ wff-stamps)
  apply unfold-optimization
  \mathbf{using}\ stamp\text{-}under\text{-}semantics
 using wff-stamps-def apply fastforce
 unfolding size.simps by simp
optimization condition-bounds-y: ((x < y) ? x : y) \mapsto y when (stamp-under
```

```
(stamp-expr\ y)\ (stamp-expr\ x)\ \land\ wff-stamps)
  {\bf apply} \ {\it unfold-optimization}
  \mathbf{using}\ stamp\text{-}under\text{-}semantics\text{-}inversed
  using wff-stamps-def apply fastforce
  unfolding size.simps by simp
optimization b[intval]: ((x eq y) ? x : y) \mapsto y
  {\bf apply} \ unfold\text{-}optimization
     apply (smt\ (z3)\ bool-to-val.simps(2)\ intval-equals.elims\ val-to-bool.simps(1)
val-to-bool.simps(3))
   unfolding intval.simps
   \mathbf{apply} \; (smt \; (z3) \; BinaryExprE \; ConditionalExprE \; Value.inject(1) \; Value.inject(2)
bin-eval.simps(10)\ bool-to-val.simps(2)\ evalDet\ intval-equals.simps(1)\ intval-equals.simps(10)
intval-equals.simps(12) intval-equals.simps(15) intval-equals.simps(16) intval-equals.simps(2)
intval-equals.simps(5) intval-equals.simps(8) intval-equals.simps(9) le-expr-def val-to-bool.cases
val-to-bool. elims(2))
 unfolding size.simps by auto
end
end
```