# GraalVM Stamp Theory

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#### Abstract

The GraalVM compiler uses stamps to track type and range information during program analysis. Type information is recorded by using distinct subclasses of the abstract base class Stamp, i.e. IntegerStamp is used to represent an integer type. Each subclass introduces facilities for tracking range information. Every subclass of the Stamp class forms a lattice, together with an arbitrary top and bottom element each sublattice forms a lattice of all stamps. This Isabelle/HOL theory models stamps as instantiations of a lattice.

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### 1 Stamps: Type and Range Information

```
theory StampLattice
imports
Values
HOL.Lattices
begin
```

#### 1.1 Void Stamp

The VoidStamp represents a type with no associated values. The VoidStamp lattice is therefore a simple single element lattice.

```
datatype void =
  VoidStamp
instantiation \ void :: order
begin
definition less-eq\text{-}void :: void \Rightarrow void \Rightarrow bool where
 less-eq	ext{-}void\ a\ b=\ True
definition less\text{-}void :: void \Rightarrow void \Rightarrow bool where
  less-void\ a\ b=False
instance
 apply standard
    apply (simp add: less-eq-void-def less-void-def)
   apply (simp add: less-eq-void-def)
  apply (simp add: less-eq-void-def)
 by (metis (full-types) void.exhaust)
end
\mathbf{instantiation}\ \mathit{void} :: \mathit{semilattice-inf}
begin
definition inf-void :: void \Rightarrow void \Rightarrow void where
 inf-void\ a\ b = VoidStamp
instance
 apply standard
   apply (simp add: less-eq-void-def)
  apply (simp add: less-eq-void-def)
 by (metis (mono-tags) void.exhaust)
end
instantiation \ void :: semilattice-sup
begin
```

```
definition sup\text{-}void :: void \Rightarrow void \Rightarrow void where
  sup	ext{-}void\ a\ b=\ VoidStamp
instance
  apply standard
   apply (simp add: less-eq-void-def)
  apply (simp add: less-eq-void-def)
  by (metis (mono-tags) void.exhaust)
end
instantiation \ void :: bounded-lattice
begin
definition bot-void :: void where
  bot	ext{-}void = VoidStamp
definition top\text{-}void :: void \text{ where}
  top	ext{-}void = VoidStamp
instance
 apply standard
  apply (simp add: less-eq-void-def)
 by (simp add: less-eq-void-def)
end
Definition of the stamp type
datatype stamp =
```

### 1.2 Stamp Lattice



intstamp int64 int64 — Type: Integer; Range: Lower Bound & Upper Bound

#### 1.2.1 Stamp Order

Defines an ordering on the stamp type.

One stamp is less than another if the valid values for the stamp are a strict subset of the other stamp.

```
instantiation \ stamp :: order
begin
fun less-eq\text{-}stamp :: stamp \Rightarrow stamp \Rightarrow bool where
  less-eq-stamp (intstamp l1\ u1) (intstamp l2\ u2) = (\{l1..u1\}\subseteq\{l2..u2\})
fun less-stamp :: stamp \Rightarrow stamp \Rightarrow bool where
  less-stamp (intstamp l1\ u1) (intstamp l2\ u2) = (\{l1..u1\} \subset \{l2..u2\})
\mathbf{lemma}\ \mathit{less-le-not-le}:
  fixes x y :: stamp
  shows (x < y) = (x \le y \land \neg y \le x)
  using less-eq-stamp.simps less-stamp.simps
  using stamp.exhaust subset-not-subset-eq by metis
lemma order-refl:
  fixes x :: stamp
  shows x \leq x
  using less-eq-stamp.simps less-stamp.simps
  using dual-order.refl stamp.exhaust by metis
lemma order-trans:
  fixes x \ y \ z :: stamp
  shows x \le y \Longrightarrow y \le z \Longrightarrow x \le z
  fix x :: stamp and y :: stamp and z :: stamp
  assume x \leq y
  assume y \leq z
  obtain l1 u1 where xdef: x = intstamp l1 u1
   using stamp.exhaust
   by blast
  obtain l2\ u2 where ydef: y = intstamp\ l2\ u2
   using stamp.exhaust
  obtain l3\ u3 where zdef: z = intstamp\ l3\ u3
   using stamp.exhaust
   by blast
  have s1: \{l1..u1\} \le \{l2..u2\}
    \mathbf{using} \ \langle x \leq y \rangle \ less-eq\text{-}stamp.simps \ xdef \ ydef \ \mathbf{by} \ blast
  have s2: \{l2..u2\} \le \{l3..u3\}
   using \langle y \leq z \rangle less-eq-stamp.simps ydef zdef by blast
  from s1 \ s2 \ \text{have} \ \{l1..u1\} \le \{l3..u3\}
   by (meson dual-order.trans)
  then show x \leq z
```

```
using less-eq-stamp.simps
   using xdef zdef by presburger
qed
lemma antisym:
 fixes x y :: stamp
 shows x \le y \Longrightarrow y \le x \Longrightarrow x = y
proof -
 \mathbf{fix}\ x::stamp
 \mathbf{fix} \ y :: stamp
 assume xlessy: x \leq y
 assume ylessx: y \le x
 obtain l1 u1 where xdef: x = intstamp l1 u1
   using stamp.exhaust by blast
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by blast
 from xlessy have s1: \{l1..u1\} \subseteq \{l2..u2\}
   using less-eq-stamp.simps
   using xdef ydef by blast
  from ylessx have s2: \{l2..u2\} \subseteq \{l1..u1\}
   using less-eq-stamp.simps
   using xdef ydef by blast
 \mathbf{have}\ \{l1..u1\}\subseteq\{l2..u2\}\Longrightarrow\{l2..u2\}\subseteq\{l1..u1\}\Longrightarrow\{l1..u1\}=\{l2..u2\}
   by fastforce
  then have s3: \{l1..u1\} = \{l2..u2\} \Longrightarrow (l1 = l2) \land (u1 = u2)
 then have (l1 = l2) \land (u1 = u2) \Longrightarrow x = y
   using xdef ydef by fastforce
 then show x = y
   using s1 s2 s3 by fastforce
qed
instance
 apply standard
 using less-le-not-le apply simp
 using order-refl apply simp
 using order-trans apply simp
 using antisym by simp
end
```

#### 1.2.2 Stamp Join

Defines the *join* operation for stamps.

For any two stamps, the *join* is defined as the intersection of the valid values for the stamp.

```
instantiation \ stamp :: semilattice-inf
begin
notation inf (infix \sqcap 65)
fun inf-stamp :: stamp \Rightarrow stamp \Rightarrow stamp where
 inf-stamp (intstamp l1 u1) (intstamp l2 u2) = intstamp (max l1 l2) (min u1 u2)
lemma inf-le1:
 fixes x y :: stamp
 shows (x \sqcap y) \leq x
proof -
 \mathbf{fix} \ x :: stamp
 \mathbf{fix} \ y :: stamp
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by blast
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by blast
 have joindef: x \sqcap y = intstamp (max l1 l2) (min u1 u2)
   (is ?join = intstamp ?l3 ?u3)
   using inf-stamp.simps xdef ydef
   by force
 have leq: \{?l3..?u3\} \subseteq \{l1..u1\}
   by force
 have (x \sqcap y) \le x = (\{?l3..?u3\} \subseteq \{l1..u1\})
   using xdef joindef inf-stamp.simps
   by force
 then show (x \sqcap y) \leq x
   using leq
   by fastforce
qed
lemma inf-le2:
 \mathbf{fixes}\ x\ y::stamp
 shows (x \sqcap y) \leq y
proof -
 \mathbf{fix}\ x::stamp
 \mathbf{fix}\ y::stamp
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by blast
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by blast
 have joindef: x \sqcap y = intstamp (max l1 l2) (min u1 u2)
   (is ?join = intstamp ?l3 ?u3)
   using inf-stamp.simps xdef ydef
   by force
 have leg: \{?l3..?u3\} \subseteq \{l2..u2\}
   by force
 have (x \sqcap y) \le y = (\{?l3..?u3\} \subseteq \{l2..u2\})
```

```
using ydef joindef
   by force
 then show (x \sqcap y) \leq y
   using leq
   by fastforce
\mathbf{qed}
lemma inf-greatest:
 fixes x \ y \ z :: stamp
 shows x \leq y \Longrightarrow x \leq z \Longrightarrow x \leq (y \sqcap z)
proof -
 fix x y z :: stamp
 assume xlessy: x \leq y
 assume xlessz: x \leq z
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by blast
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by blast
 obtain l3\ u3 where zdef: z = intstamp\ l3\ u3
   using stamp.exhaust by blast
 obtain l4 u4 where yzdef: y \sqcap z = intstamp l4 u4
   by (meson inf-stamp.elims)
 have max4: l4 = max l2 l3
   using yzdef ydef zdef inf-stamp.simps by simp
 have min4: u4 = min u2 u3
   using yzdef ydef zdef inf-stamp.simps by simp
 have \{l1..u1\} \subseteq \{l2..u2\}
   using xlessy xdef ydef
   using less-eq-stamp.simps by blast
 have \{l1..u1\} \subseteq \{l3..u3\}
   using xlessz xdef zdef
   using less-eq-stamp.simps by blast
 have leq: \{l1..u1\} \subseteq \{l4..u4\}
   using \langle \{l1..u1\} \subseteq \{l2..u2\} \rangle \langle \{l1..u1\} \subseteq \{l3..u3\} \rangle max4 min4 by auto
 have x \le (y \sqcap z) = (\{l1..u1\} \subseteq \{l4..u4\})
   by (simp add: xdef yzdef)
 then show x \leq (y \sqcap z)
   using leq
   by fastforce
qed
instance
 apply standard
 using inf-le1 apply simp
 using inf-le2 apply simp
 using inf-greatest by simp
end
```

#### 1.2.3 Stamp Meet

Defines the *meet* operation for stamps.

For any two stamps, the *meet* is defined as the union of the valid values for the stamp.

```
instantiation \ stamp :: semilattice-sup
begin
notation sup (infix \sqcup 65)
fun sup\text{-}stamp :: stamp \Rightarrow stamp \Rightarrow stamp where
 sup-stamp (intstamp l1 u1) (intstamp l2 u2) = intstamp (min l1 l2) (max u1 u2)
lemma sup-ge1:
 fixes x y :: stamp
 shows x \leq x \sqcup y
proof -
 \mathbf{fix}\ x::stamp
 \mathbf{fix} \ y :: stamp
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by blast
  obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by blast
  have joindef: x \sqcup y = intstamp (min l1 l2) (max u1 u2)
   (is ?join = intstamp ?l3 ?u3)
   using inf-stamp.simps xdef ydef
   by force
 have leq: \{l1..u1\} \subseteq \{?l3..?u3\}
   by simp
 have x \le x \sqcup y = (\{l1..u1\} \subseteq \{?l3..?u3\})
   using xdef joindef inf-stamp.simps
   by force
 then show x \leq x \sqcup y
   using leq
   by fastforce
\mathbf{qed}
lemma sup-ge2:
 fixes x y :: stamp
 shows y \leq x \sqcup y
proof -
 \mathbf{fix} \ x :: stamp
 \mathbf{fix} \ y :: stamp
 obtain l1 u1 where xdef: x = intstamp \ l1 \ u1
   using stamp.exhaust by blast
 obtain l2\ u2 where ydef: y = intstamp\ l2\ u2
   using stamp.exhaust by blast
 have joindef: x \sqcup y = intstamp (min l1 l2) (max u1 u2)
   (is ?join = intstamp ?l3 ?u3)
```

```
using inf-stamp.simps xdef ydef
   by force
 have leq: \{l2..u2\} \subseteq \{?l3..?u3\} (is ?subset-thesis)
   by simp
 have ?thesis = (?subset-thesis)
   \mathbf{using}\ ydef\ joindef\ sup\text{-}stamp.simps\ less\text{-}eq\text{-}stamp.simps
   by (metis StampLattice.sup-ge1 max.commute min.commute sup-stamp.elims)
  then show ?thesis
   using leq
   by fastforce
qed
lemma sup-least:
 \mathbf{fixes}\ x\ y\ z::stamp
 shows y \le x \Longrightarrow z \le x \Longrightarrow ((y \sqcup z) \le x)
proof -
 \mathbf{fix} \ x \ y \ z :: stamp
 assume xlessy: y \leq x
 assume xlessz: z \leq x
 obtain l1 u1 where xdef: x = intstamp l1 u1
   using stamp.exhaust by blast
 obtain l2\ u2 where ydef:\ y=intstamp\ l2\ u2
   using stamp.exhaust by blast
 obtain l3\ u3 where zdef: z = intstamp\ l3\ u3
   using stamp.exhaust by blast
 \mathbf{have} \ yzdef: \ y \ \sqcup \ z = intstamp \ (min \ l2 \ l3) \ (max \ u2 \ u3)
   (is ?meet = intstamp ?l4 ?u4)
   using sup-stamp.simps
   by (simp add: ydef zdef)
 have s1: \{l2..u2\} \subseteq \{l1..u1\}
   using xlessy xdef ydef
   using less-eq-stamp.simps by blast
 have s2: \{l3..u3\} \subseteq \{l1..u1\}
   using xlessz xdef zdef
   using less-eq-stamp.simps by blast
 have leg: \{?l4...?u4\} \subset \{l1..u1\} (is ?subset-thesis)
   using s1 s2 unfolding atLeastatMost-subset-iff
   by (metis (no-types, opaque-lifting) inf.orderE inf-stamp.simps max.bounded-iff
max.cobounded2 min.bounded-iff min.cobounded2 stamp.inject xdef xlessy xlessz ydef
zdef
 have (y \sqcup z \leq x) = ?subset\text{-thesis}
   using yzdef xdef less-eq-stamp.simps
   by simp
 then show (y \sqcup z \leq x)
   using leq by fastforce
ged
instance
```

```
apply standard
using sup-ge1 apply simp
using sup-ge2 apply simp
using sup-least by simp
end
```

#### 1.2.4 Stamp Bounds

Defines the top and bottom elements of the stamp lattice.

This poses an interesting question as our stamp type is a union of the various Stamp subclasses, e.g. IntegerStamp, ObjectStamp, etc.

Each subclass should preferably have its own unique top and bottom element, i.e. An *IntegerStamp* would have the top element of the full range of integers allowed by the bit width and a bottom of a range with no integers. While the *ObjectStamp* should have *Object* as the top and *Void* as the bottom element.

```
instantiation stamp :: bounded-lattice
begin

notation bot \ (\bot 50)
notation top \ (\top 50)

definition width\text{-}min :: nat \Rightarrow int64 where
width\text{-}min \ bits = -(2 \cap bits - 1))

definition width\text{-}max :: nat \Rightarrow int64 where
width\text{-}max \ bits = (2 \cap bits - 1) - 1

value (sint \ (width\text{-}min \ 64), \ sint \ (width\text{-}max \ 64))
value max\text{-}word::int64

lemma
assumes x = width\text{-}min \ 64
assumes y = width\text{-}max \ 64
shows sint \ x < sint \ y
using assms unfolding width\text{-}min\text{-}def \ width\text{-}max\text{-}def} by simp
```

Note that this definition is valid for unsigned integers only.

The bottom and top element for signed integers would be (-9223372036854775808, 9223372036854775807).

For unsigned we have (0, 18446744073709551615).

For Java we are likely to be more concerned with signed integers. To use the appropriate bottom and top for signed integers we would need to change our definition of less\_eq from l1..u1 <= l2..u2 to sint l1..sint u1 <= sint l2..sint u2

We may still find an unsigned integer stamp useful. I plan to investigate

the Java code to see if this is useful and then apply the changes to switch to signed integers.

```
definition bot-stamp = intstamp (-1) \theta
definition top-stamp = intstamp \ \theta \ (-1)
lemma bot-least:
 fixes a :: stamp
 shows (\bot) \le a
proof -
 obtain min max where bot-def:\bot = intstamp max min
   using bot-stamp-def
   by force
 have min < max
   using bot-def
   unfolding bot-stamp-def width-min-def width-max-def
   using word-gt-\theta by fastforce
 then have \{max..min\} = \{\}
   using bot-def
   unfolding bot-stamp-def width-min-def width-max-def
   by auto
 then show ?thesis
   unfolding bot-stamp-def
   using less-eq-stamp.simps
   by (simp add: stamp.induct)
qed
lemma top-greatest:
 fixes a :: stamp
 shows a \leq (\top)
proof -
 obtain min max where top-def:\top = intstamp min max
   using top-stamp-def
   by force
 have max-is-max: \neg(\exists n. n > max)
   by (metis stamp.inject top-def top-stamp-def word-order.extremum-strict)
 have min-is-min: \neg(\exists n. n < min)
  by (metis not-less-iff-gr-or-eq stamp.inject top-def top-stamp-def word-coorder.not-eq-extremum)
 have \neg(\exists l u. \{min..max\} < \{l..u\})
   using max-is-max min-is-min
   by (metis atLeastatMost-psubset-iff not-less)
 then show ?thesis
   unfolding top-stamp-def
   using less-eq-stamp.simps
   using less-eq-stamp.elims(3) by fastforce
qed
instance
 apply standard
 using bot-least apply simp
```

```
using top-greatest by simp end
```

#### 1.3 Java Stamp Methods

The following are methods from the Java Stamp class, they are the methods primarily used for optimizations.

```
definition is-unrestricted :: stamp \Rightarrow bool where is-unrestricted s = (\top = s)

fun is-empty :: stamp \Rightarrow bool where is-empty s = (\bot = s)

fun as-constant :: stamp \Rightarrow Value option where as-constant (intstamp l u) = (if (card \{l..u\}) = 1 then Some (IntVal64 (SOME x. x \in \{l..u\})) else None)

definition always-distinct :: stamp \Rightarrow stamp \Rightarrow bool where always-distinct stamp1 stamp2 = (\bot = (stamp1 \sqcap stamp2))

definition never-distinct :: stamp \Rightarrow stamp \Rightarrow bool where never-distinct stamp1 stamp2 = (as-constant stamp1 = as-constant stamp2 \land as-constant stamp1 \neq None)
```

### 1.4 Mapping to Values

```
fun valid-value :: stamp => Value => bool where valid-value (intstamp l u) (IntVal64 v) = (v \in \{l..u\}) | valid-value (intstamp l u) -= False
```

The *valid-value* function is used to map a stamp instance to the values that are allowed by the stamp.

It would be nice if there was a slightly more integrated way to perform this mapping as it requires some infrastructure to prove some fairly simple properties.

```
lemma bottom-range-empty: \neg(valid\text{-}value\ (\bot)\ v) unfolding bot-stamp-def using valid-value.elims(2) by fastforce \text{lemma join-}values: assumes joined = x-stamp \sqcap y-stamp shows valid-value joined x \longleftrightarrow (valid\text{-}value\ x\text{-}stamp\ x \land valid\text{-}value\ y\text{-}stamp\ x) proof (cases\ x) case UndefVal then show ?thesis using valid\text{-}value.elims(2) by blast
```

```
next
 case (IntVal32 x2)
 then show ?thesis
   using valid-value. elims(2) by blast
next
  case (IntVal64 x3)
 obtain lx\ ux where xdef:\ x\text{-}stamp\ =\ intstamp\ lx\ ux
   using stamp.exhaust by blast
 obtain ly\ uy\ \mathbf{where}\ ydef:\ y\text{-}stamp\ =\ intstamp\ ly\ uy
   using stamp.exhaust by blast
 obtain v where x = IntVal64 v
   using IntVal64 by blast
 have joined = intstamp (max lx ly) (min ux uy)
   (is joined = intstamp ?lj ?uj)
   by (simp add: xdef ydef assms)
 then have valid-value joined (IntVal64 v) = (v \in \{?lj..?uj\})
   by simp
 then show ?thesis
   using \langle x = IntVal64 \ v \rangle \ xdef \ ydef \ \mathbf{by} \ force
\mathbf{next}
 case (ObjRef x5)
 then show ?thesis
   using valid-value. elims(2) by blast
next
 case (ObjStr\ x6)
 then show ?thesis
   using valid-value. elims(2) by blast
qed
lemma disjoint-empty:
 fixes x-stamp y-stamp :: stamp
 assumes \bot = x\text{-}stamp \sqcap y\text{-}stamp
 shows \neg(valid\text{-}value x\text{-}stamp x \land valid\text{-}value y\text{-}stamp x)
 {\bf using} \ assms \ bottom{-}range{-}empty \ join{-}values
 by blast
experiment begin
A possible equivalent alternative to the definition of less eq
fun less-eq-alt :: 'a::ord \times 'a \Rightarrow 'a \times 'a \Rightarrow bool where
 less-eq-alt (l1, u1) (l2, u2) = ((\neg l1 \le u1) \lor l2 \le l1 \land u1 \le u2)
Proof equivalence
lemma
 fixes 11 12 u1 u2 :: int
 assumes l1 \leq u1 \wedge l2 \leq u2
 shows \{l1..u1\} \subseteq \{l2..u2\} = ((l1 \ge l2) \land (u1 \le u2))
 by (simp add: assms)
```

```
lemma
  fixes 11 12 u1 u2 :: int
 shows \{l1..u1\} \subseteq \{l2..u2\} = less-eq-alt\ (l1, u1)\ (l2, u2)
  by simp
\mathbf{end}
```

#### 1.5 Generic Integer Stamp

Experimental definition of integer stamps generically, restricting the datatype to only allow valid ranges and the bottom integer element (max\_int..min\_int).

```
lemma
 assumes (x::int) > 0
 shows (2 \hat{x})/2 = (2 \hat{x} - 1)
 sorry
definition max-signed-int :: 'a::len word where
 max-signed-int = (2 \land (LENGTH('a) - 1)) - 1
definition min-signed-int :: 'a::len word where
 min-signed-int = -(2 \land (LENGTH('a) - 1))
definition int-bottom :: 'a::len word \times 'a word where
 int-bottom = (max-signed-int, min-signed-int)
definition int-top :: 'a::len word \times 'a word where
 int-top = (min-signed-int, max-signed-int)
lemma
 fixes x :: 'a :: len word
 shows sint \ x \leq sint \ (((2 \cap (LENGTH('a) - 1)) - 1)::'a \ word)
 using sint-greater-eq sorry
value sint (0::1 word)
value sint (1::1 word)
value sint (((2 \cap \theta) - 1)::1 \ word)
value sint (((2 \hat{\ } 31) - 1)::32 \ word)
lemma max-signed:
 fixes a :: 'a :: len word
 shows sint \ a \leq sint \ (max\text{-}signed\text{-}int::'a \ word)
proof (cases sint a = sint (max-signed-int::'a word))
 case True
 then show ?thesis by simp
next
```

```
{f case} False
 have sint\ a < sint\ (max-signed-int::'a\ word)
   using False unfolding max-signed-int-def sorry
  then show ?thesis by simp
qed
lemma min-signed:
 fixes a :: 'a::len word
 shows sint \ a \ge sint \ (min\text{-}signed\text{-}int::'a \ word)
 sorry
value max-signed-int :: 32 word
value int-bottom::(32 word \times 32 word)
value sint (2147483647::32 word)
value sint (2147483648::32 word)
typedef (overloaded) ('a::len) intstamp =
  \{bounds :: ('a \ word, 'a \ word) \ prod . ((fst \ bounds) \leq s \ (snd \ bounds) \lor bounds = \}
int-bottom)}
proof -
 show ?thesis
   by (smt (z3) mem-Collect-eq prod.sel(1) prod.sel(2) signed-minus-1 sint-0)
qed
setup-lifting type-definition-intstamp
lift-definition lower :: ('a::len) intstamp \Rightarrow 'a word
 is prod.fst \circ Rep-intstamp.
lift-definition upper :: ('a::len) intstamp \Rightarrow 'a word
 is prod.snd \circ Rep-intstamp.
lift-definition lower-int :: ('a::len) intstamp \Rightarrow int
 is sint \circ prod.fst.
lift-definition upper-int :: ('a::len) intstamp <math>\Rightarrow int
 is sint \circ prod.snd.
lift-definition range :: ('a::len) intstamp \Rightarrow int set
 is \lambda (l, u). \{sint \ l..sint \ u\}.
lift-definition bounds :: ('a::len) intstamp \Rightarrow ('a word \times 'a word)
 is Rep-intstamp.
lift-definition is-bottom :: ('a::len) intstamp \Rightarrow bool
 is \lambda x. x = int\text{-}bottom.
```

```
lift-definition from-bounds :: ('a::len word \times 'a word) \Rightarrow 'a intstamp
  is Abs-intstamp.
instantiation intstamp :: (len) order
begin
definition less-eq-intstamp :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow bool where
  less-eq-intstamp \ s1 \ s2 = (range \ s1 \subseteq range \ s2)
definition less-intstamp :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow bool where
  less-intstamp s1 s2 = (range \ s1 \subset range \ s2)
value int-bottom::(1 \ word \times 1 \ word)
value sint (0::1 word)
value sint (1::1 word)
value int-bottom::(2 word \times 2 word)
value sint (1::2 word)
value sint (2::2 word)
value sint((2 \hat{LENGTH}(32) - 1) - 1)::32 \ word) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word)) > sint((-(2 \hat{LENGTH}(32) + 1) - 1)::32 \ word))
- 1)))::32 word)
lemma bottom-is-bottom:
  assumes is-bottom s
  shows s \leq a
proof -
  have boundsdef: bounds s = int-bottom
    by (metis assms bounds.transfer is-bottom.rep-eq)
  obtain min max where bounds s = (max, min)
    by fastforce
  then have max \neq min
   \mathbf{by} \; (\textit{metis bounds def dual-order}. \textit{eq-iff fst-conv int-bottom-def less-minus-one-simps} (1)
max-signed min-signed not-less sint-0 sint-n1 snd-conv)
  then have sint min < sint max
    unfolding boundsdef int-bottom-def
    using max-signed
   by (metis \ (bounds \ s = (max, min)) \ bounds defint-bottom-deforder.not-eq-order-implies-strict
prod.sel(1) signed-word-eqI)
  then have range s = \{\}
    unfolding range-def bounds-def
    by (simp add: \langle bounds \ s = (max, min) \rangle bounds.transfer)
  then show ?thesis
    by (simp add: StampLattice.less-eq-intstamp-def)
qed
lemma bounds-has-value:
  fixes x y :: int
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assumes x < y
 shows card \{x..y\} > 0
 using assms by auto
lemma bounds-has-no-value:
 fixes x y :: int
 assumes x < y
 shows card \{y..x\} = 0
 using assms by auto
lemma bottom-unique:
 fixes a s :: 'a intstamp
 assumes is-bottom s
 shows a \leq s \longleftrightarrow is\text{-}bottom\ a
proof -
 have \forall x. \ sint \ (fst \ (bounds \ x)) \leq sint \ (snd \ (bounds \ x)) \lor is-bottom \ x
   unfolding bounds-def is-bottom-def
   using Rep-intstamp
   using word-sle-eq by auto
  then have \forall x. (card (range x)) > 0 \lor is\text{-bottom } x
   unfolding range-def using bounds-has-value
   by (simp add: bounds.transfer case-prod-beta)
  obtain min max where boundsdef: bounds s = (max, min)
   by fastforce
 have nooverlap: sint min < sint max
   using max-signed
  by (metis assms bounds.transfer boundsdef fst-conv int-bottom-def is-bottom.rep-eq
min-signed order.not-eq-order-implies-strict signed-word-eqI sint-0 snd-conv verit-la-disequality
zero-neg-one)
 have range s = \{sint \ max..sint \ min\}
   by (simp add: bounds.transfer boundsdef range.rep-eq)
 then have card (range s) = 0
   using nooverlap bounds-has-no-value by simp
  then have \forall x. (card (range x)) > 0 \longrightarrow s < x
     using \langle StampLattice.range\ s = \{sint\ max..sint\ min\} \rangle\ atLeastatMost-empty
less-intstamp-def by auto
 then show ?thesis
  \textbf{by} \; (\textit{meson} \; \forall \, \textit{x.} \; \textit{0} < \textit{card} \; (\textit{StampLattice.range} \; \textit{x}) \; \lor \; \textit{is-bottom} \; \textit{x} \lor \; \textit{bottom-is-bottom}
leD\ less-eq-intstamp-def\ less-intstamp-def)
qed
lemma bottom-antisym:
 assumes is-bottom x
 shows x \leq y \Longrightarrow y \leq x \Longrightarrow x = y
 using assms proof (cases is-bottom y)
case True
  then show ?thesis
   by (metis Rep-intstamp-inverse assms is-bottom.rep-eq)
```

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next
 {f case}\ {\it False}
 assume y \leq x
 have \neg (y \leq x)
   using bottom-unique False assms
   by simp
 then show ?thesis
   \mathbf{using} \ \langle y \leq x \rangle \ \mathbf{by} \ auto
qed
lemma int-antisym:
 fixes x y :: 'a intstamp
 \mathbf{shows}\ x \leq y \Longrightarrow y \leq x \Longrightarrow x = y
proof -
 \mathbf{fix} \ x :: 'a \ intstamp
 \mathbf{fix} \ y :: 'a \ intstamp
 assume xlessy: x \leq y
 assume ylessx: y \le x
 obtain l1 u1 where xdef: bounds x = (l1, u1)
   by fastforce
 obtain l2 u2 where ydef: bounds y = (l2, u2)
   by fastforce
 from xlessy have s1: \{sint \ l1..sint \ u1\} \subseteq \{sint \ l2..sint \ u2\} (is ?xlessy)
   using xdef ydef unfolding bounds-def range-def less-eq-intstamp-def
   by simp
  from ylessx have s2: \{sint \ l2..sint \ u2\} \subseteq \{sint \ l1..sint \ u1\} (is ?ylessx)
   using xdef ydef unfolding bounds-def range-def less-eq-intstamp-def
   by simp
 show x = y proof (cases is-bottom x)
   case True
   then show ?thesis using bottom-antisym xlessy ylessx
     \mathbf{by} \ simp
 next
   {f case} False
   then show ?thesis sorry
 qed
qed
instance
 apply standard
    apply (simp add: less-eq-intstamp-def less-intstamp-def less-le-not-le)
 apply blast
 using less-eq-intstamp-def apply force
 \mathbf{using}\ \mathit{less-eq-intstamp-def}\ \mathbf{apply}\ \mathit{force}
 by (simp add: int-antisym)
value take-bit LENGTH(63) 20::int
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```
value take-bit LENGTH(63) ((-20)::int)
value bit (20::int64) (63::nat)
value bit ((-20)::int64) (63::nat)
value ((-20)::int64) < (20::int64)
value take-bit LENGTH(63) ((-20)::int)
lift-definition smax :: 'a :: len \ word \Rightarrow 'a \ word \Rightarrow 'a \ word
 is \lambda a b. (if (sint a) \leq (sint b) then b else a).
lift-definition smin :: 'a::len \ word \Rightarrow 'a \ word \Rightarrow 'a \ word
 is \lambda a b. (if (sint a) \leq (sint b) then a else b).
instantiation intstamp :: (len) semilattice-inf
begin
notation inf (infix \sqcap 65)
definition join-bounds: 'a intstamp \Rightarrow 'a intstamp \Rightarrow ('a word \times 'a word) where
 join-bounds\ s1\ s2 = (smax\ (lower\ s1)\ (lower\ s2),\ smin\ (upper\ s1)\ (upper\ s2))
definition join-or-bottom :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow ('a word \times 'a word)
where
 join-or-bottom\ s1\ s2=(let\ bound=(join-bounds\ s1\ s2)\ in
   if sint (fst bound) \ge sint (snd bound) then int-bottom else bound)
definition inf-intstamp :: 'a intstamp \Rightarrow 'a intstamp \Rightarrow 'a intstamp where
  inf-intstamp s1 \ s2 = from-bounds (join-or-bottom s1 \ s2)
lemma always-valid:
 fixes s1 s2 :: 'a intstamp
 shows Rep-intstamp (from-bounds (join-or-bottom s1 s2)) = join-or-bottom s1 s2
 unfolding join-or-bottom-def join-bounds-def from-bounds-def
 using Abs-intstamp-inverse
 by (smt (z3) from-bounds.transfer from-bounds-def mem-Collect-eq word-sle-eq)
lemma invalid-join:
  fixes s1 s2 :: 'a intstamp
 assumes bound = join\text{-}bounds s1 s2
 assumes sint (fst bound) \ge sint (snd bound)
 shows from-bounds int-bottom = s1 \sqcap s2
 using assms(1) assms(2) inf-intstamp-def join-or-bottom-def by presburger
\mathbf{lemma} \ \mathit{unfold-bounds} :
  bounds \ x = (lower \ x, \ upper \ x)
 by (simp add: bounds.transfer lower.rep-eq upper.rep-eq)
```

```
lemma int-inf-le1:
    \mathbf{fixes}\ x\ y::\ 'a\ intstamp
    shows (x \sqcap y) \leq x
proof (cases is-bottom (x \sqcap y))
    case True
    then show ?thesis
       by (simp add: bottom-is-bottom)
\mathbf{next}
    case False
    then show ?thesis
    using False proof –
    obtain l1 u1 where xdef: lower x = l1 \land upper x = u1
       by fastforce
    obtain l2 u2 where ydef: lower y = l2 \land upper y = u2
       by fastforce
    have joindef: x \sqcap y = from\text{-bounds} ((smax l1 l2, smin u1 u2))
       (is x \sqcap y = from\text{-}bounds (?13, ?u3))
       using False
          by (smt (z3) StampLattice.inf-intstamp-def StampLattice.join-bounds-def al-
ways-valid is-bottom.rep-eq join-or-bottom-def xdef ydef)
    have leq: \{sint ?l3..sint ?u3\} \subseteq \{sint l1..sint u1\}
       by (smt (z3) atLeastatMost-subset-iff smax.transfer smin.transfer)
    have (x \sqcap y) \leq x = (\{sint ? l3..sint ? u3\} \subseteq \{sint l1..sint u1\})
       using xdef joindef range-def less-eq-intstamp-def
        by (smt (z3) False StampLattice.always-valid StampLattice.join-or-bottom-def
bounds. abs-eq\ case-prod-conv\ inf-intstamp-def\ is-bottom. rep-eq\ join-bounds-def\ range. rep-eq\ pointstamp-def\ is-bottom. The product of the product
unfold-bounds ydef)
    then show (x \sqcap y) \leq x
       using leq
       by fastforce
qed
qed
lemma int-inf-le2:
   fixes x y :: 'a intstamp
    shows (x \sqcap y) \leq y
proof (cases is-bottom (x \sqcap y))
    case True
    then show ?thesis
       by (simp add: bottom-is-bottom)
next
    case False
    then show ?thesis
    using False proof -
    obtain l1 u1 where xdef: lower x = l1 \land upper x = u1
       by fastforce
    obtain l2\ u2 where ydef:\ lower\ y=l2\ \land\ upper\ y=u2
       bv fastforce
    have joindef: x \sqcap y = from\text{-}bounds ((smax l1 l2, smin u1 u2))
```

```
(is x \sqcap y = from\text{-}bounds (?l3, ?u3))
        using False
           by (smt (z3) StampLattice.inf-intstamp-def StampLattice.join-bounds-def al-
ways-valid is-bottom.rep-eq join-or-bottom-def xdef ydef)
    have leq: \{sint ?l3..sint ?u3\} \subseteq \{sint l1..sint u1\}
        by (smt (z3) atLeastatMost-subset-iff smax.transfer smin.transfer)
    have (x \sqcap y) \leq y = (\{sint ?l3..sint ?u3\} \subseteq \{sint l2..sint u2\})
        using xdef joindef range-def less-eq-intstamp-def
         by (smt (z3) False StampLattice.always-valid StampLattice.join-or-bottom-def
bounds. abs-eq\ case-prod-conv\ inf-intstamp-def\ is-bottom. rep-eq\ join-bounds-def\ range. rep-eq\ pointstamp-def\ is-bottom. The product of the product
unfold-bounds ydef)
    then show (x \sqcap y) \leq y
        using leq
        by (smt (z3) atLeastatMost-subset-iff smax.transfer smin.transfer)
qed
qed
lemma
    assumes x \leq y
    assumes is-bottom y
    shows is-bottom x
    using bottom-is-bottom assms
    using bottom-unique by auto
\mathbf{lemma} \ \mathit{int-inf-greatest} :
    fixes x y :: 'a intstamp
    shows x \leq y \Longrightarrow x \leq z \Longrightarrow x \leq y \sqcap z
    sorry
instance
    apply standard
        apply (simp add: local.int-inf-le1)
     apply (simp add: local.int-inf-le2)
    by (simp add: local.int-inf-greatest)
end
instantiation intstamp :: (len) semilattice-sup
begin
notation sup (infix \sqcup 65)
instance sorry
end
instantiation intstamp :: (len) bounded-lattice
begin
```

```
notation bot (\perp 50)
notation top (\top 50)
definition bot-intstamp = int-bottom
definition top-intstamp = int-top
instance sorry
end
value sint (0::1 word)
value sint (1::1 word)
datatype Stamp =
 BottomStamp |
 TopStamp |
 VoidStamp \mid
 Int8Stamp 8 intstamp
 Int16Stamp 16 intstamp
 Int32Stamp 32 intstamp |
 Int64Stamp 64 intstamp
instantiation Stamp :: order
begin
fun less-eq-Stamp :: Stamp <math>\Rightarrow Stamp \Rightarrow bool where
 less-eq-Stamp\ BottomStamp\ -=\ True\ |
 less-eq-Stamp - TopStamp = True \mid
 less-eq-Stamp\ VoidStamp\ VoidStamp\ =\ True\ |
 less-eq-Stamp (Int8Stamp v1) (Int8Stamp v2) = (v1 \le v2) |
 less-eq-Stamp (Int16Stamp v1) (Int16Stamp v2) = (v1 \le v2)
 less-eq-Stamp (Int32Stamp v1) (Int32Stamp v2) = (v1 \le v2)
 less-eq-Stamp (Int64Stamp v1) (Int64Stamp v2) = (v1 \le v2) |
 less-eq-Stamp - - = False
fun less\text{-}Stamp :: Stamp \Rightarrow Stamp \Rightarrow bool where
 less-Stamp\ BottomStamp\ BottomStamp\ = False
 less-Stamp BottomStamp - = True |
 less-Stamp \ TopStamp \ TopStamp = False
 less-Stamp - TopStamp = True \mid
 less-Stamp\ VoidStamp\ VoidStamp\ = False
 less-Stamp (Int8Stamp v1) (Int8Stamp v2) = (v1 < v2)
 less-Stamp (Int16Stamp v1) (Int16Stamp v2) = (v1 < v2)
 less-Stamp (Int32Stamp v1) (Int32Stamp v2) = (v1 < v2)
 less-Stamp (Int64Stamp v1) (Int64Stamp v2) = (v1 < v2)
 less-Stamp - - = False
```

```
instance
 apply standard sorry
end
instantiation Stamp :: semilattice-inf
begin
notation inf (infix \sqcap 65)
fun inf-Stamp :: Stamp \Rightarrow Stamp \Rightarrow Stamp where
 inf-Stamp BottomStamp -= BottomStamp
 inf-Stamp - BottomStamp = BottomStamp
 inf-Stamp TopStamp - = TopStamp
 inf-Stamp - TopStamp = TopStamp
 inf-Stamp VoidStamp | VoidStamp |
 inf-Stamp (Int8Stamp v1) (Int8Stamp v2) = Int8Stamp (v1 \sqcap v2)
 inf-Stamp (Int16Stamp v1) (Int16Stamp v2) = Int16Stamp (v1 \sqcap v2)
 inf-Stamp (Int32Stamp v1) (Int32Stamp v2) = Int32Stamp (v1 \sqcap v2)
 inf-Stamp (Int64Stamp v1) (Int64Stamp v2) = Int64Stamp (v1 \sqcap v2)
instance
 apply standard sorry
end
instantiation Stamp :: semilattice-sup
begin
notation sup (infix \sqcup 65)
fun sup-Stamp :: Stamp \Rightarrow Stamp \Rightarrow Stamp where
 sup-Stamp BottomStamp - = BottomStamp
 sup-Stamp - BottomStamp = BottomStamp
 sup\text{-}Stamp \ TopStamp \ - = \ TopStamp \ |
 sup\text{-}Stamp - TopStamp = TopStamp
 sup-Stamp VoidStamp VoidStamp = VoidStamp
 sup-Stamp (Int8Stamp v1) (Int8Stamp v2) = Int8Stamp (v1 \sqcup v2) |
 sup\text{-}Stamp \ (Int16Stamp \ v1) \ (Int16Stamp \ v2) = Int16Stamp \ (v1 \sqcup v2) \mid
 sup-Stamp (Int32Stamp v1) (Int32Stamp v2) = Int32Stamp (v1 \sqcup v2)
 sup-Stamp (Int64Stamp v1) (Int64Stamp v2) = Int64Stamp (v1 \sqcup v2)
instance
 apply standard sorry
end
instantiation Stamp :: bounded-lattice
begin
```

```
notation bot \ (\bot 50)

notation top \ (\top 50)

definition top\text{-}Stamp :: Stamp \text{ where}

top\text{-}Stamp = TopStamp

definition bot\text{-}Stamp :: Stamp \text{ where}

bot\text{-}Stamp = BottomStamp

instance

apply standard \text{ sorry}

end

lemma [code]: Rep\text{-}intstamp \ (from\text{-}bounds \ (l, \ u)) = (l, \ u)

using Abs\text{-}intstamp\text{-}inverse \ from\text{-}bounds.rep\text{-}eq}

sorry

code-datatype Abs\text{-}intstamp
```