# Veriopt Theories

February 9, 2022

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begin
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```

translations

```
n \le CONST Rep-intExp n

n \le CONST Rep-i32e n

n \le CONST Rep-i64e n
```

lemma  $\textit{vminusv} \colon \forall \, \textit{vv} \, \textit{v} \, . \, \textit{vv} = \textit{IntVal64} \, \textit{v} \longrightarrow \textit{v} - \textit{v} = \textit{0}$ 

**by** simp

thm-oracles vminusv

lemma redundant-sub:

 $\forall\, vv_1\ vv_2\ v_1\ v_2$  .  $vv_1=IntVal64\ v_1\wedge vv_2=IntVal64\ v_2\longrightarrow v_1-(v_1-v_2)=v_2$ 

by simp

 ${f thm ext{-}oracles}\ redundant ext{-}sub$ 

#### val-eq

 $\forall vv \ v. \ vv = IntVal64 \ v \longrightarrow v - v = 0$ 

 $\forall \, vv_1 \, vv_2 \, v_1 \, v_2. \, vv_1 = IntVal64 \, v_1 \wedge vv_2 = IntVal64 \, v_2 \longrightarrow v_1 - (v_1 - v_2) = v_2$ 

#### phase tmp

terminating size

begin

#### sub-same-32

optimization sub-same-32:  $(e::i32e) - e \longmapsto const (IntVal32 0)$ 

**apply** (unfold rewrite-preservation.simps, unfold rewrite-termination.simps, rule conjE, simp) **apply** auto[1] **using** Rep-i32e evalDet is-IntVal32-def **apply** (smt (verit, del-insts) eq-iff-diff-eq-0 evaltree.simps int-constants-valid int-val-sub.simps(1) is-int-val.simps(1) mem-Collect-eq)

unfolding size.simps

by (metis add-strict-increasing gr-implies-not0 less-one linorder-not-le size-gt-0)

## $\overline{sub}$ - $\overline{same}$ -64

optimization sub-same-64:  $(e::i64e) - e \longmapsto const (IntVal64 0)$ 

apply auto

 $\mathbf{apply} \ (metis \ (no\text{-}types, \ opaque\text{-}lifting) \ ConstantExpr \ bin\text{-}eval.simps} (3) \ bin\text{-}eval\text{-}preserves\text{-}validity } \\ cancel\text{-}comm\text{-}monoid\text{-}add\text{-}class.diff\text{-}cancel\ evalDet\ i64e\text{-}eval\ int\text{-}and\text{-}equal\text{-}bits.simps} (2) \\ intval\text{-}sub.simps} (2))$ 

by  $(simp\ add:\ Suc\ -le\ -eq\ add\ -strict\ -increasing\ size\ -gt\ -0)$  end

thm-oracles sub-same-32

## ast-example

 $BinaryExpr\ BinAdd\ (BinaryExpr\ BinMul\ x\ x)\ (BinaryExpr\ BinMul\ x\ x)$ 

# $abstract\hbox{-}syntax\hbox{-}tree$

# ${\bf datatype}\,\, \mathit{IRExpr} =$

 $UnaryExpr\ IRUnaryOp\ IRExpr$ 

 $BinaryExpr\ IRBinaryOp\ IRExpr\ IRExpr$ 

ConditionalExpr IRExpr IRExpr IRExpr

ParameterExpr nat Stamp

LeafExpr nat Stamp

| ConstantExpr Value

Constant Var (char list)

VariableExpr (char list) Stamp

## value

## ${\bf datatype}\ \mathit{Value} = \mathit{UndefVal}$

| IntVal32 (32 word)

IntVal64 (64 word)

ObjRef (nat option)

ObjStr (char list)

#### eval

 $unary\text{-}eval :: IRUnaryOp \Rightarrow Value \Rightarrow Value$ 

bin-eval ::  $IRBinaryOp \Rightarrow Value \Rightarrow Value \Rightarrow Value$ 

## tree-semantics

 $semantics: unary \quad semantics: binary \quad semantics: conditional \quad semantics: constant \quad semantics: parameter \quad semantics: leaf$ 

## tree-evaluation-deterministic

$$[m,p] \vdash e \mapsto v_1 \land [m,p] \vdash e \mapsto v_2 \Longrightarrow v_1 = v_2$$

thm-oracles evalDet

```
expression\hbox{-}refinement
```

begin

```
e_1 \sqsupseteq e_2 = (\forall \ m \ p \ v. \ [m,p] \vdash e_1 \mapsto v \longrightarrow [m,p] \vdash e_2 \mapsto v)
```

```
expression\mbox{-}refinement\mbox{-}monotone
     e \supseteq e'
                                            \implies UnaryExpr \ op \ e \supseteq UnaryExpr \ op \ e'
     x \sqsupseteq x' \land y \sqsupseteq y'
                                            \implies BinaryExpr op x \ y \supseteq BinaryExpr op x' \ y'
     ce \supseteq ce' \land te \supseteq te' \land fe \supseteq fe' \implies ConditionalExpr \ ce \ te \ fe \supseteq ConditionalExpr \ ce' \ te' \ fe'
\mathbf{ML} \leftarrow
(*fun get-list (phase: phase option) =
  case phase of
   NONE => []
   SOME \ p => (\#rewrites \ p)
fun\ get\text{-}rewrite\ name\ thy =
 let
   val\ (phases,\ lookup) = (case\ RWList.get\ thy\ of
     NoPhase\ store => store \mid
     InPhase (name, store, -) => store)
   val\ rewrites = (map\ (fn\ x => get\text{-}list\ (lookup\ x))\ phases)
  in
   rewrites
  end
fun \ rule-print \ name =
  Document-Output.antiquotation-pretty\ name\ (Args.term)
   (fn\ ctxt => fn\ (rule) => (*Pretty.str\ hello)*)
     Pretty.block (print-all-phases (Proof-Context.theory-of ctxt)));
(*
     Goal	ext{-}Display.pretty	ext{-}goal
       (Config.put Goal-Display.show-main-goal main ctxt)
       (#goal (Proof.goal (Toplevel.proof-of (Toplevel.presentation-state ctxt)))));
*)
val - = Theory.setup
(rule-print binding (rule));*)
phase SnipPhase
 terminating size
```

## Binary Fold Constant

**optimization** BinaryFoldConstant: BinaryExpr op (const v1) (const v2)  $\longmapsto$  ConstantExpr (bin-eval op v1 v2) when int-and-equal-bits v1 v2

**unfolding** rewrite-preservation.simps rewrite-termination.simps **apply** (rule conjE, simp, simp del: le-expr-def)

## Binary Fold Constant Obligation

- 1. int-and-equal-bits v1 v2 →
  BinaryExpr op (ConstantExpr v1) (ConstantExpr v2) □
  ConstantExpr (bin-eval op v1 v2)
- 2. int-and-equal-bits v1 v2  $\longrightarrow$   $trm(BinaryExpr\ op\ (ConstantExpr\ v1)$   $(ConstantExpr\ v2)) > trm(ConstantExpr\ (bin-eval\ op\ v1\ v2))$

using BinaryFoldConstant by auto

#### AddCommuteConstantRight

**optimization** AddCommuteConstantRight:  $((const\ v) + y) \longmapsto y + (const\ v)$  when  $\neg (is\text{-}ConstantExpr\ y)$ 

**unfolding** rewrite-preservation.simps rewrite-termination.simps **apply** (rule conjE, simp, simp del: le-expr-def)

# Add Commute Constant Right Obligation

¬ is-ConstantExpr y →
BinaryExpr BinAdd (ConstantExpr v) y □
BinaryExpr BinAdd y (ConstantExpr v)
 ¬ is-ConstantExpr y →
trm(BinaryExpr BinAdd (ConstantExpr v)
y) > trm(BinaryExpr BinAdd y (ConstantExpr v))

#### using AddShiftConstantRight by auto

#### AddNeutral

optimization  $AddNeutral: ((e::i32e) + (const (IntVal32 0))) \mapsto e$ 

unfolding rewrite-preservation.simps rewrite-termination.simps apply (rule conjE, simp, simp del: le-expr-def)

## Add Neutral Obligation

- 1.  $BinaryExpr\ BinAdd\ e\ (ConstantExpr\ (IntVal32\ 0)) \supseteq e$
- 2.  $trm(BinaryExpr\ BinAdd\ e\ (ConstantExpr\ (IntVal32\ 0))) > trm(e)$

using neutral-zero(1) rewrite-preservation.simps(1) apply blast by auto

## InverseLeftSub

**optimization** InverseLeftSub:  $((e_1::intExp) - (e_2::intExp)) + e_2 \longmapsto e_1$ 

unfolding rewrite-preservation.simps rewrite-termination.simps apply (rule conjE, simp, simp del: le-expr-def)

## InverseLeftSubObligation

- 1.  $BinaryExpr\ BinAdd\ (BinaryExpr\ BinSub\ e_1\ e_2)\ e_2\ \supseteq\ e_1$
- 2.  $trm(BinaryExpr\ BinAdd\ (BinaryExpr\ BinSub\ e_1\ e_2)\ e_2) > trm(e_1)$

using neutral-left-add-sub by auto

## InverseRightSub

**optimization**  $InverseRightSub: (e_2::intExp) + ((e_1::intExp) - e_2) \longmapsto e_1$ 

unfolding rewrite-preservation.simps rewrite-termination.simps apply (rule conjE, simp, simp del: le-expr-def)

# InverseRightSubObligation

- 1.  $BinaryExpr\ BinAdd\ e_2\ (BinaryExpr\ BinSub\ e_1\ e_2)\ \supseteq\ e_1$
- 2.  $trm(BinaryExpr\ BinAdd\ e_2\ (BinaryExpr\ BinSub\ e_1\ e_2)) > trm(e_1)$

using neutral-right-add-sub by auto

## AddToSub

**optimization**  $AddToSub: -e + y \longmapsto y - e$ 

unfolding rewrite-preservation.simps rewrite-termination.simps apply (rule conjE, simp, simp del: le-expr-def)

## Add To Sub Obligation

- 1.  $BinaryExpr\ BinAdd\ (UnaryExpr\ UnaryNeg\ e)\ y \supseteq BinaryExpr\ BinSub\ y\ e$
- 2.  $trm(BinaryExpr\ BinAdd\ (UnaryExpr\ UnaryNeg\ e)\ y) > trm(BinaryExpr\ BinSub\ y\ e)$

using AddLeftNegateToSub by auto

 $\mathbf{end}$ 

**definition** trm where trm = size

 $\overline{phase}$ 

phase AddCanonicalizations
 terminating trm
begin...end

hide-const (open) Form.wf-stamp

phase-example

phase Conditional terminating trm begin

phase-example-1

**optimization** negate-condition:  $(\neg e ? x : y) \longmapsto (e ? y : x)$ 

 ${\bf using} \ {\it Conditional Phase. negate-condition}$ 

**by** (auto simp: trm-def)

phase-example-2

**optimization** const-true:  $(true ? x : y) \longmapsto x$ 

 $\mathbf{by} \ (auto \ simp: \ trm-def)$ 

phase-example-3

**optimization** const-false: (false ? x : y)  $\longmapsto y$ 

**by** (auto simp: trm-def)

phase-example-4

**optimization** equal-branches:  $(e ? x : x) \longmapsto x$ 

**by** (auto simp: trm-def)

```
phase-example-5  \begin{array}{c} \textbf{optimization} \ \ condition\mbox{-}bounds\mbox{-}x\mbox{:}\ ((x < y) ? x : y) \longmapsto x \\  \  \  \  \  \  \mbox{when}\ (stamp\mbox{-}under\ (stamp\mbox{-}expr\ y) \\  \  \  \  \  \  \  \  \  \  \wedge\  \mbox{wf-}stamp\ x \wedge \mbox{wf-}stamp\ y) \end{array}
```

**using** ConditionalPhase.condition-bounds-x(1) **by** (blast, auto simp: trm-def)

```
phase-example-6

optimization condition-bounds-y: ((x < y) ? x : y) \mapsto y

when (stamp-under (stamp-expr y) (stamp-expr x) \land wf-stamp
x \land wf-stamp y)
```

**using** ConditionalPhase.condition-bounds-y(1) **by** (blast, auto simp: trm-def)

phase-example-7

end

## termination

```
\begin{array}{lll} trm(UnaryExpr\ op\ e) &=& trm(e)+1 \\ trm(BinaryExpr\ BinAdd\ x\ y) &=& trm(x)+2*trm(y) \\ trm(ConditionalExpr\ cond\ t\ f) &=& trm(cond)+trm(t)+trm(f)+2 \\ trm(ConstantExpr\ c) &=& 1 \\ trm(ParameterExpr\ ind\ s) &=& 2 \\ trm(LeafExpr\ nid\ s) &=& 2 \end{array}
```

## $graph\mbox{-}representation$

```
\mathbf{typedef}\ \mathit{IRGraph} = \{g :: \mathit{ID} \rightharpoonup (\mathit{IRNode} \times \mathit{Stamp}) \ . \ \mathit{finite}\ (\mathit{dom}\ g)\}
```

# graph2tree

rep:constant rep:parameter rep:conditional rep:unary rep:convert rep:binary rep:leaf rep:ref

```
preeval
is-preevaluated (InvokeNode\ n\ uu\ uv\ uw\ ux\ uy) = True
is-preevaluated (InvokeWithExceptionNode n uz va vb vc vd ve) =
True
is-preevaluated (NewInstanceNode n vf vg vh) = True
is-preevaluated (LoadFieldNode n vi vj vk) = True
is-preevaluated (SignedDivNode n vl vm vn vo vp) = True
is-preevaluated (SignedRemNode\ n\ vq\ vr\ vs\ vt\ vu) = True
is-preevaluated (ValuePhiNode n \ vv \ vw) = True
is-preevaluated (AbsNode\ v) = False
is-preevaluated (AddNode v va) = False
is-preevaluated (AndNode v va) = False
is-preevaluated (BeginNode v) = False
is-preevaluated (BytecodeExceptionNode v va vb) = False
is-preevaluated (ConditionalNode v va vb) = False
is-preevaluated (ConstantNode v) = False
is-preevaluated (DynamicNewArrayNode v va vb vc vd) = False
is-preevaluated EndNode = False
is-preevaluated (ExceptionObjectNode v va) = False
is-preevaluated (FrameState v va vb vc) = False
is-preevaluated (IfNode v va vb) = False
is-preevaluated (IntegerBelowNode v va) = False
is-preevaluated (IntegerEqualsNode v va) = False
is-preevaluated (IntegerLessThanNode v va) = False
is-preevaluated (IsNullNode\ v) = False
is-preevaluated (KillingBeginNode v) = False
is-preevaluated (LeftShiftNode v va) = False
is-preevaluated (LogicNegationNode v) = False
is-preevaluated (LoopBeginNode v va vb vc) = False
is-preevaluated (LoopEndNode\ v) = False
is-preevaluated (LoopExitNode\ v\ va\ vb) = False
is-preevaluated (MergeNode v va vb) = False
is-preevaluated (MethodCallTargetNode v va) = False
is-preevaluated (MulNode v va) = False
is-preevaluated (NarrowNode v va vb) = False
is-preevaluated (NegateNode v) = False
is-preevaluated (NewArrayNode v va vb) = False
is-preevaluated (NotNode v) = False
is-preevaluated (OrNode v va) = False
is-preevaluated (ParameterNode\ v) = False
is-preevaluated (PiNode\ v\ va) = False
is-preevaluated (ReturnNode v va) = False
is-preevaluated (RightShiftNode v va) = False
is-preevaluated (ShortCircuitOrNode v va) = False
```

is-preevaluated (SianExtendNode v va vb) = False

## $deterministic \hbox{-} representation$

$$g \vdash n \simeq e_1 \land g \vdash n \simeq e_2 \Longrightarrow e_1 = e_2$$

## thm-oracles repDet

## $well\mbox{-}formed\mbox{-}term\mbox{-}graph$

$$\exists\, e.\ g \vdash n \,\simeq\, e \,\wedge\, (\exists\, v.\ [m,p] \vdash e \mapsto v)$$

# graph-semantics

$$([g,m,p] \vdash n \mapsto v) = (\exists e. g \vdash n \simeq e \land [m,p] \vdash e \mapsto v)$$

## graph-semantics-deterministic

$$[g,m,p] \vdash n \mapsto v_1 \, \land \, [g,m,p] \vdash n \mapsto v_2 \Longrightarrow v_1 = v_2$$

 ${f thm ext{-}oracles}\ graphDet$ 

## **notation** (*latex*)

graph-refinement (term-graph-refinement -)

## graph-refinement

$$\begin{array}{l} \textit{term-graph-refinement} \ g_1 \ g_2 = \\ (\textit{ids} \ g_1 \subseteq \textit{ids} \ g_2 \land \\ (\forall \, n. \ n \in \textit{ids} \ g_1 \longrightarrow (\forall \, e. \ g_1 \vdash n \simeq e \longrightarrow g_2 \vdash n \unlhd e))) \end{array}$$

# translations

 $n <= CONST \ as ext{-}set \ n$ 

## graph-semantics-preservation

$$\begin{array}{l} {e_1}' \sqsupseteq {e_2}' \wedge \\ \{n\} \mathrel{\lessdot} g_1 \subseteq g_2 \wedge \\ g_1 \vdash n \simeq {e_1}' \wedge g_2 \vdash n \simeq {e_2}' \Longrightarrow \\ \mathit{term-graph-refinement} \ g_1 \ g_2 \end{array}$$

 ${\bf thm\text{-}oracles}\ \textit{graph-semantics-preservation-subscript}$ 

## $maximal\mbox{-}sharing$

```
maximal-sharing g = (\forall n_1 \ n_2.

n_1 \in true\text{-}ids \ g \land n_2 \in true\text{-}ids \ g \longrightarrow (\forall e. \ g \vdash n_1 \simeq e \land g \vdash n_2 \simeq e \longrightarrow n_1 = n_2))
```

## tree-to-graph-rewriting

$$\begin{array}{l} e_1 \mathrel{\sqsubseteq} e_2 \land \\ g_1 \vdash n \simeq e_1 \land \\ maximal\text{-}sharing \ g_1 \land \\ \{n\} \lessdot g_1 \mathrel{\subseteq} g_2 \land \\ g_2 \vdash n \simeq e_2 \land \\ maximal\text{-}sharing \ g_2 \Longrightarrow \\ term\text{-}graph\text{-}refinement \ g_1 \ g_2 \end{array}$$

## thm-oracles tree-to-graph-rewriting

## $term\hbox{-} graph\hbox{-} refines\hbox{-} term$

$$(g \vdash n \trianglelefteq e) = (\exists e'. g \vdash n \simeq e' \land e \sqsupseteq e')$$

## $term\mbox{-}graph\mbox{-}evaluation$

$$g \vdash n \trianglelefteq e \Longrightarrow \forall \ m \ p \ v. \ [m,p] \vdash e \mapsto v \longrightarrow [g,m,p] \vdash n \mapsto v$$

## graph-construction

$$\begin{array}{l} e_1 \sqsupseteq e_2 \wedge g_1 \subseteq g_2 \wedge g_2 \vdash n \simeq e_2 \Longrightarrow \\ g_2 \vdash n \trianglelefteq e_1 \wedge term\text{-}graph\text{-}refinement \ g_1 \ g_2 \end{array}$$

## $\mathbf{thm\text{-}oracles}\ \mathit{graph\text{-}construction}$

# $term\hbox{-} graph\hbox{-} reconstruction$

$$g \triangleleft e \leadsto (g', n) \Longrightarrow g' \vdash n \simeq e$$

#### end

theory SlideSnippets

#### imports

 $Semantics. Tree To Graph Thms \\ Snippets. Snipping$ 

#### begin

notation (latex)

 $kind \left( -\langle \langle -\rangle \rangle \right)$ 

notation (latex)

IRTreeEval.ord-IRExpr-inst.less-eq-IRExpr (-  $\longmapsto$  -)

## abstract-syntax-tree

datatype IRExpr =

UnaryExpr IRUnaryOp IRExpr

BinaryExpr IRBinaryOp IRExpr IRExpr

ConditionalExpr IRExpr IRExpr IRExpr

ParameterExpr nat Stamp

LeafExpr nat Stamp

 $Constant Expr\ Value$ 

Constant Var (char list)

| VariableExpr (char list) Stamp

#### tree-semantics

 $semantics: constant \quad semantics: parameter \quad semantics: unary \quad semantics: binary \quad semantics: leaf$ 

## $expression\hbox{-}refinement$

$$e_1 \supseteq e_2 = (\forall m \ p \ v. \ [m,p] \vdash e_1 \mapsto v \longrightarrow [m,p] \vdash e_2 \mapsto v)$$

## graph2tree

semantics:constant semantics:unary semantics:binary

## graph-semantics

$$([g,m,p] \vdash n \mapsto v) = (\exists e. g \vdash n \simeq e \land [m,p] \vdash e \mapsto v)$$

## graph-refinement

```
graph-refinement g_1 g_2 = (ids \ g_1 \subseteq ids \ g_2 \land (\forall n. \ n \in ids \ g_1 \longrightarrow (\forall e. \ g_1 \vdash n \simeq e \longrightarrow g_2 \vdash n \trianglelefteq e)))
```

## translations

 $n <= CONST \ as ext{-}set \ n$ 

## graph-semantics-preservation

## maximal-sharing

```
\begin{array}{l} \textit{maximal-sharing } g = \\ (\forall \, n_1 \ n_2. \\ \quad n_1 \in \textit{true-ids } g \land n_2 \in \textit{true-ids } g \longrightarrow \\ (\forall \, e. \ g \vdash n_1 \simeq e \land g \vdash n_2 \simeq e \longrightarrow n_1 = n_2)) \end{array}
```

## tree-to-graph-rewriting

```
\begin{array}{l} e_1 \sqsupseteq e_2 \wedge \\ g_1 \vdash n \simeq e_1 \wedge \\ maximal\text{-}sharing \ g_1 \wedge \\ \{n\} \lessdot g_1 \subseteq g_2 \wedge \\ g_2 \vdash n \simeq e_2 \wedge maximal\text{-}sharing \ g_2 \Longrightarrow \\ graph\text{-}refinement \ g_1 \ g_2 \end{array}
```

# graph-represents-expression

$$(g \vdash n \trianglelefteq e) = (\exists e'. g \vdash n \simeq e' \land e \sqsubseteq e')$$

## graph-construction

$$\begin{array}{l} e_1 \mathrel{\sqsubseteq} e_2 \land g_1 \mathrel{\subseteq} g_2 \land g_2 \vdash n \simeq e_2 \Longrightarrow \\ g_2 \vdash n \mathrel{\subseteq} e_1 \land \mathit{graph-refinement} \ g_1 \ g_2 \end{array}$$

 $\mathbf{end}$